

# Divide-and-Conquer

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# Divide and Conquer

## Divide and Conquer

- Break up problems into parts
- Solve each part recursively
- Combine solutions of each part

## Examples

- Merge sort
- Exponentiation by squaring ( $a^n$ )
- [BOJ#1725 히스토그램](#)

# Merge Sort (1)

## Pseudo-code

```
MERGE-SORT(array):  
    MERGE_SORT(left half of array)  
    MERGE_SORT(right half of array)  
    array = MERGE(left, right)  
  
MERGE(left, right):  
    result := empty array  
    while left, right are not empty:  
        add min(left[0], right[0]) to result  
        and remove it from left/right  
    if left or right is not empty:  
        add it to result  
    return result
```

# Merge Sort (2)

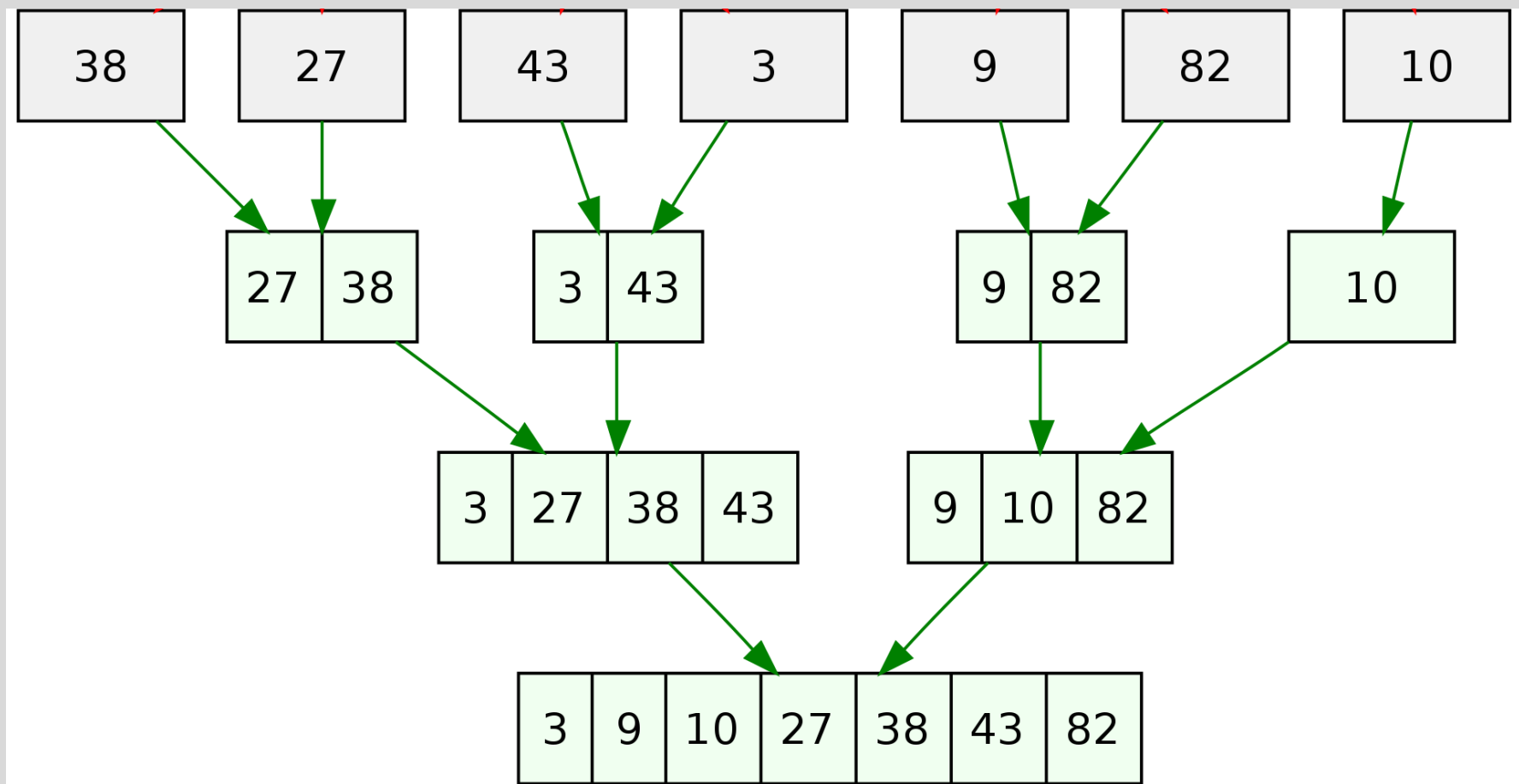
## Time Complexity

```
MERGE-SORT(list m):  
    MERGE_SORT(left)  
    MERGE_SORT(right)  
    MERGE(left, right)
```

- MERGE\_SORT: only function call overhead
- MERGE:  $O(\text{right}-\text{left})$

# Merge Sort (3)

## Time Complexity



- Height of tree:  $O(\log n)$
- Cost of MERGE on each level:  $O(n)$
- Time complexity:  $O(n \log n)$

# Exponentiation by squaring (1)

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## Recursion

$$x^n = \begin{cases} x (x^2)^{\frac{n-1}{2}}, & \text{if } n \text{ is odd} \\ (x^2)^{\frac{n}{2}}, & \text{if } n \text{ is even.} \end{cases}$$

## Pseudo-code

```
POWER(x, n):  
    if n == 0: return 1  
    else if n == 1: return x  
    else if n is odd:  
        return x*POWER(x*x, (n-1)/2)  
    else:  
        return POWER(x*x, n/2)
```

# Exponentiation by squaring (2)

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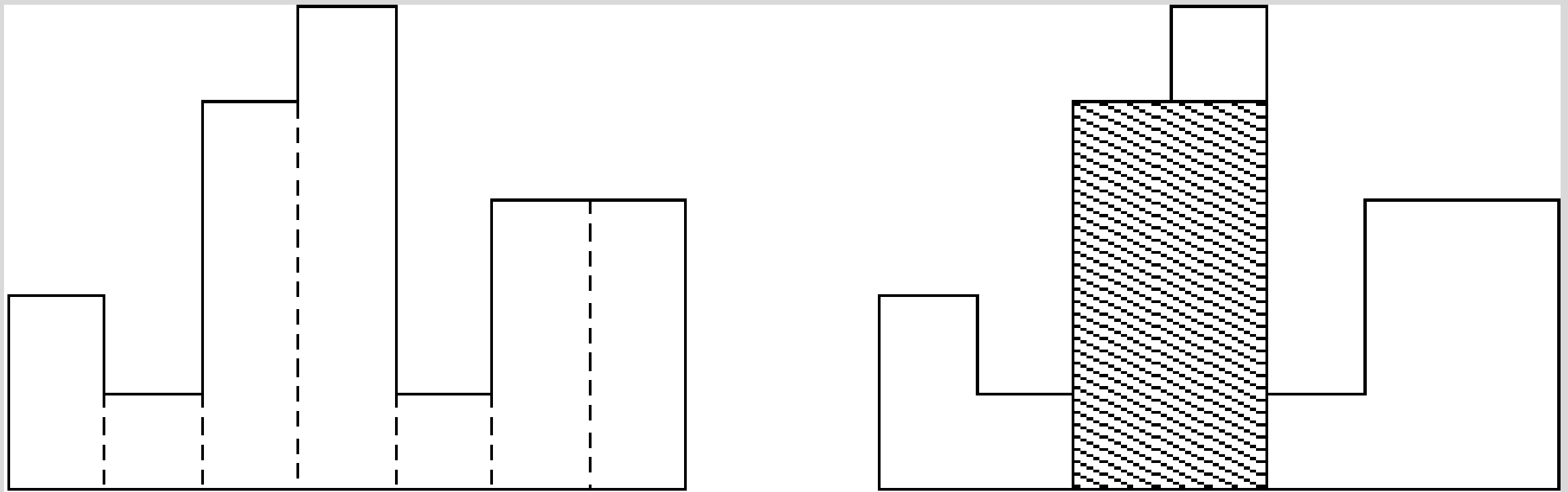
## Time Complexity

```
POWER(x, n):  
    if n == 0: return 1  
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```

- Number of multiplication:  
 $(\log n) + (\# \text{ of 1-bit in } n) = O(2 \log n)$   
 $= O(\log n)$

# Largest Rectangle in a Histogram

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- Area of largest rectangle?
- Number of rectangles:  $1 \leq n \leq 100,000$
- Height of each rectangle:  $0 \leq h_i \leq 1,000,000,000$
- Width of each rectangle: 1



## BOJ

- [10827  \$a^b\$](#)
- [10830 행렬 제곱](#)
- [1780 종이의 개수](#)
- [1725 히스토그램](#)
- [2339 석판 자르기](#)