# [544] Cassandra Replication

Tyler Caraza-Harter

### Outline

Replication

Quorum Reads/Writes

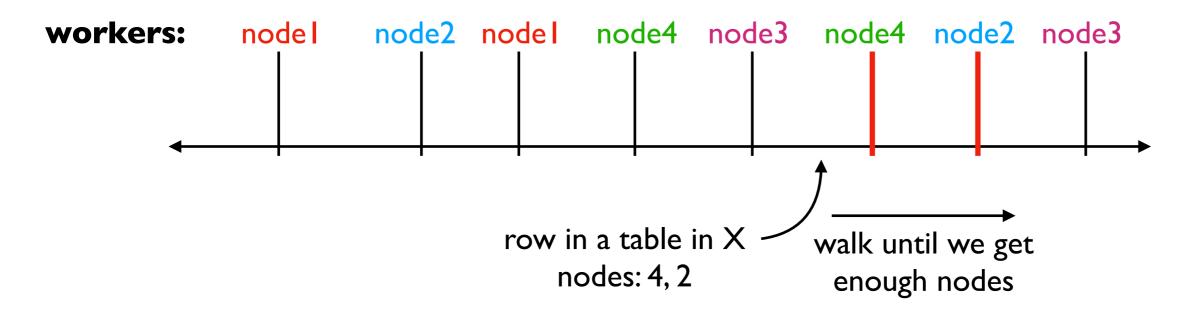
**Conflict Resolution** 

Cassandra Demos

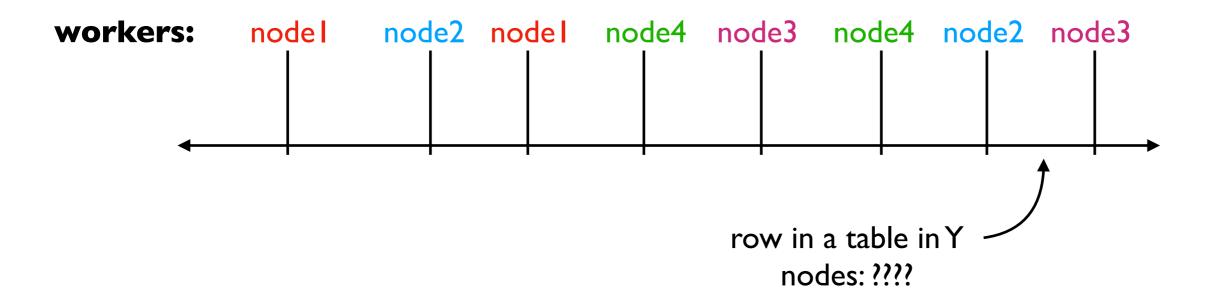
We replicate (create multiple copies on different nodes) to improve durability - meaning we don't want data to be lost when nodes die.

Cassandra lets us choose a different RF (replication factor) for each keyspace:

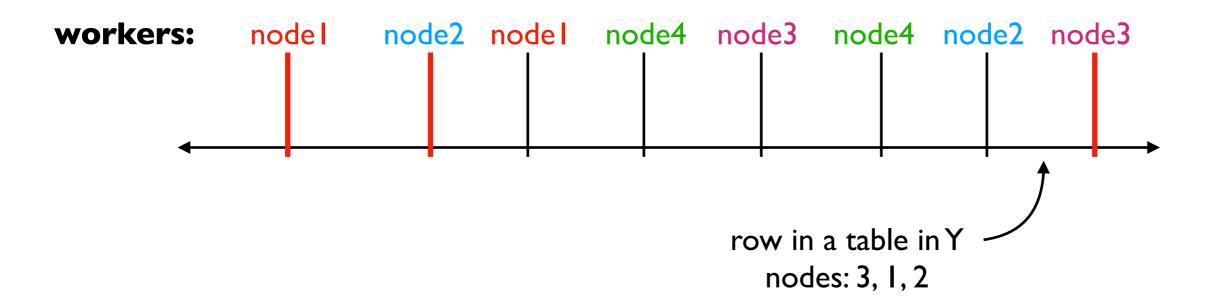
```
token(node1) = {t1, t2}
token(node2) = {t3, t4}
token(node3) = {t5, t6}
token(node4) = {t7, t8}
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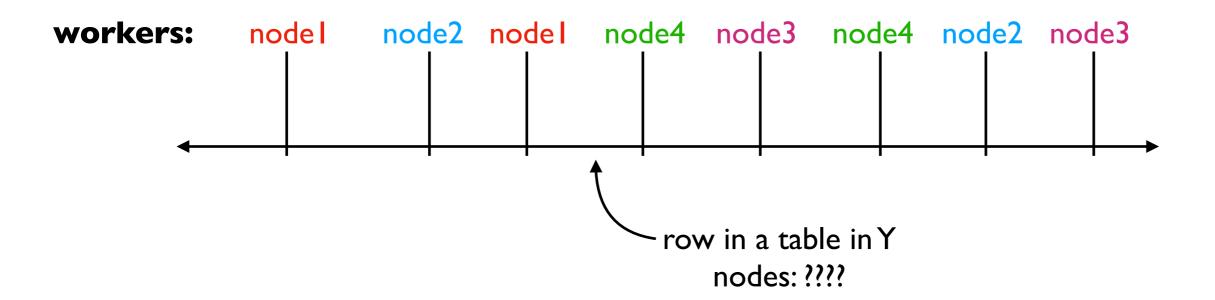
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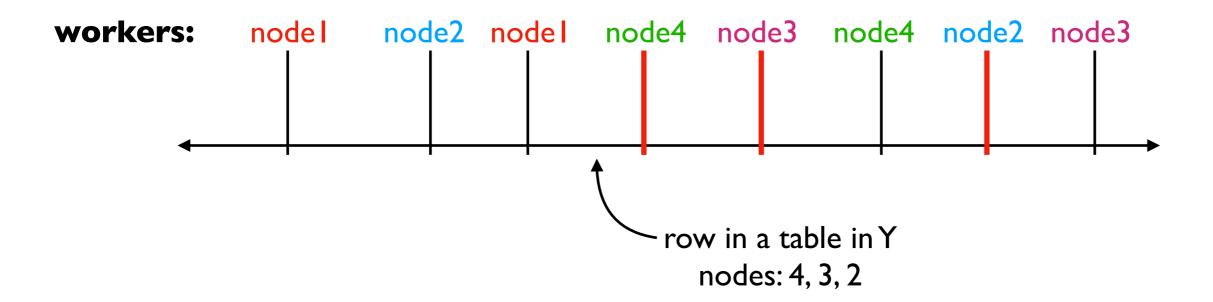
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#### **Token Map:**



Important! Keeping multiple copies on vnodes on the same node provides little safety (when a node dies, all its vnodes die). Same "failure domain".

Cassandra can skip nodes as it "walks the ring".

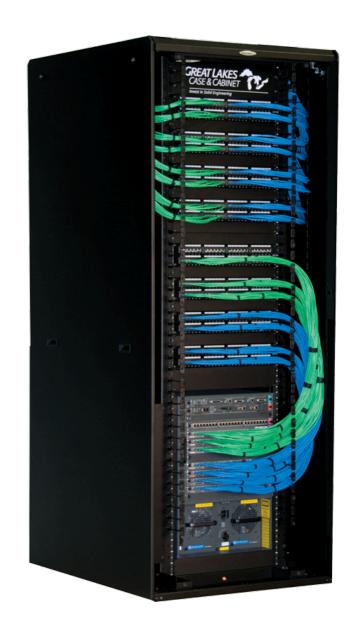
### Network Infrastructure



Server



Data Center



Rack

https://www.dotmagazine.online/issues/digital-infrastructure-and-transforming-markets/data-center-models

 $https://buy.hpe.com/us/en/servers/proliant-dl-servers/proliant-dl10-servers/proliant-dl20-server/hpe-proliant-dl20-gen10-plus-e-2336-2-9ghz-6-core-1p-16gb-u-4sff-500w-rps-server/p/p44115-b21?ef_id=Cj0KCQiAt66eBhCnARlsAKf3ZNFJsg49UV6Zm33R7lkRqi-XOd_JECmdyqNMAm2CKLSm_F-z6JTYDTQaAgMTEALw_wcB:G:s&s_kwcid=AL!13472!3!331628972784!!!g!318267171339!!1707918369!67076417419&gclsrc=aw.ds&gclid=Cj0KCQiAt66eBhCnARlsAKf3ZNFJsg49UV6Zm33R7lkRqi-XOd_JECmdyqNMAm2CKLSm_F-z6JTYDTQaAgMTEALw_wcB$ 

### Correlated Failures

One server goes down, all of its vnodes are gone.



Server



Data Center



#### Whole-rack problems:

- top-of-rack switch fails
- rack's power supply fails



Rack

https://www.dotmagazine.online/issues/digital-infrastructure-and-transforming-markets/data-center-models

https://buy.hpe.com/us/en/servers/proliant-dl-servers/proliant-dl10-servers/proliant-dl10-servers/proliant-dl20-server/hpe-proliant-dl20-gen10-plus-e-2336-2-9ghz-6-core-1p-16gb-u-4sff-500w-rps-server/p/p44115-b21?ef\_id=Cj0KCQiAt66eBhCnARlsAKf3ZNFJsg49UV6Zm33R7lkRqi-XOd\_JECmdyqNMAm2CKLSm\_F-z6JTYDTQaAgMTEALw\_wcB:G:s&s\_kwcid=AL!13472!3!331628972784!!!g!318267171339!!1707918369!67076417419&gclsrc=aw.ds&gclid=Cj0KCQiAt66eBhCnARlsAKf3ZNFJsg49UV6Zm33R7lkRqi-XOd\_JECmdyqNMAm2CKLSm\_F-z6JTYDTQaAgMTEALw\_wcB

## Replication Policy

Cassandra replication strategies are "pluggable", with a couple built-in options.

#### SimpleStrategy

- all nodes are considered equal
- skips vnodes on same machine
- ignores rack and data center placement
- used in CS 544

#### NetworkTopologyStrategy

- considers data centers and racks
- when walking the ring, some vnodes may be skipped to protect against multiple kinds of correlated failure

### Worksheet

### Outline

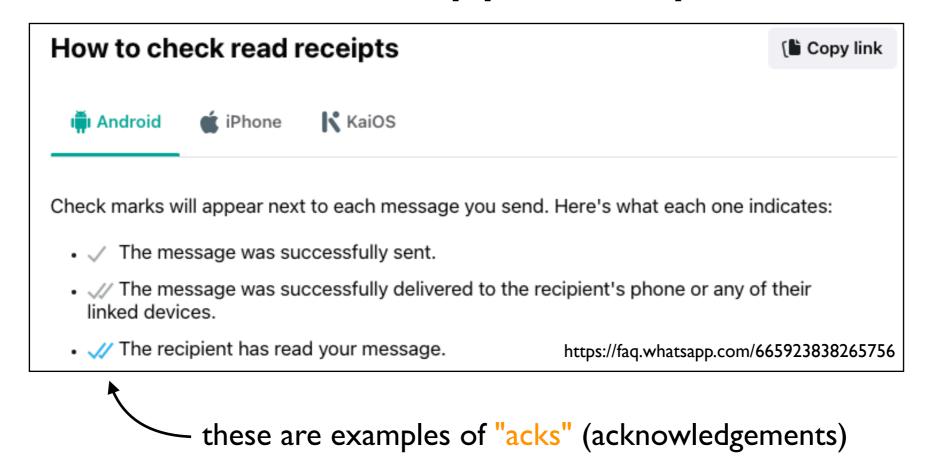
Replication

Quorum Reads/Writes

**Conflict Resolution** 

Cassandra Demos

## Write Acks: Whats App Example



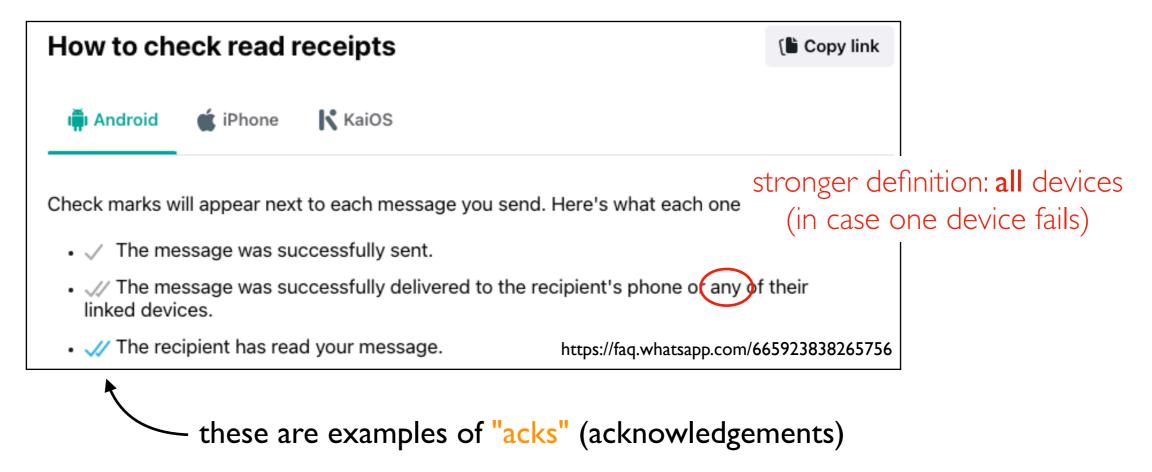
In distributed storage systems/databases, an ack means our data is committed.

"Committed" means our data is "safe", even if bad things happen. The definition varies system to system, based on what bad things are considered. For example:

- a node could hang until rebooted; a node's disk could permanently fail
- a rack could lose power; a data center could be destroyed

Obviously, no data is ever completely safe against any circumstance (e.g., comet strikes earth, leading to destruction of humankind and all our data centers).

## Write Acks: WhatsApp Example



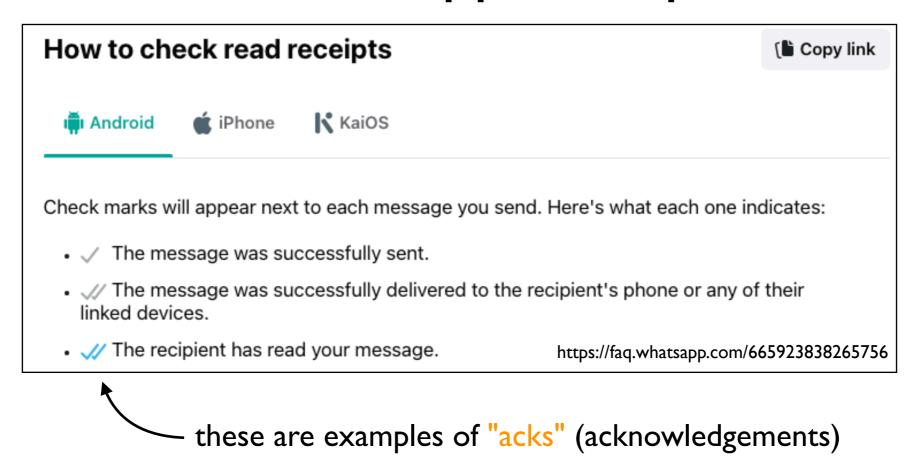
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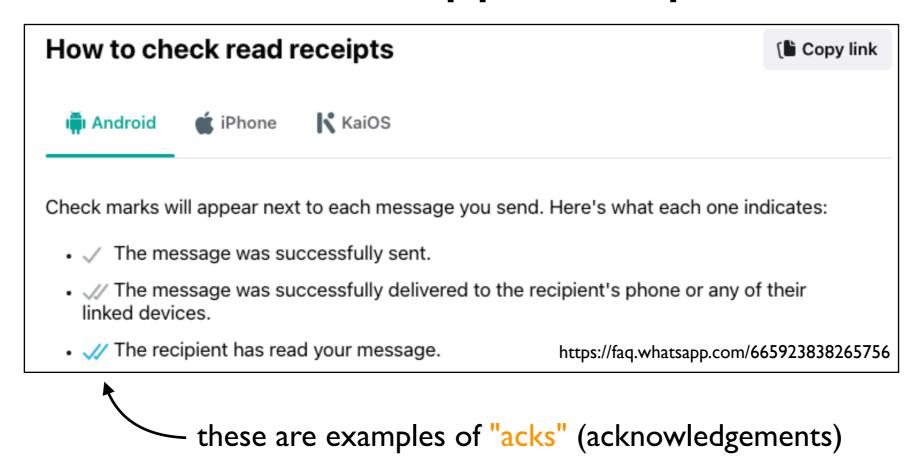
## Write Acks: WhatsApp Example



two checks (in WhatsApp) mean the message reached the destination.

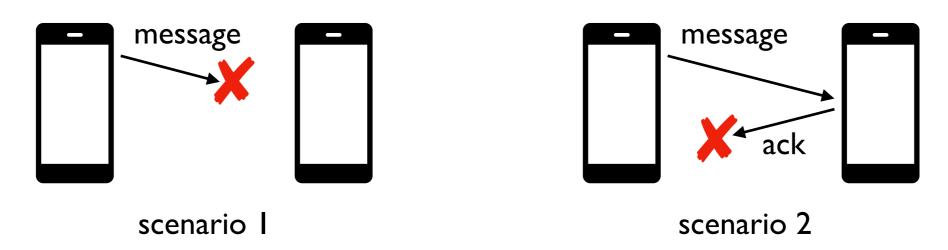
Does only one check mean the message has NOT reached the destination?

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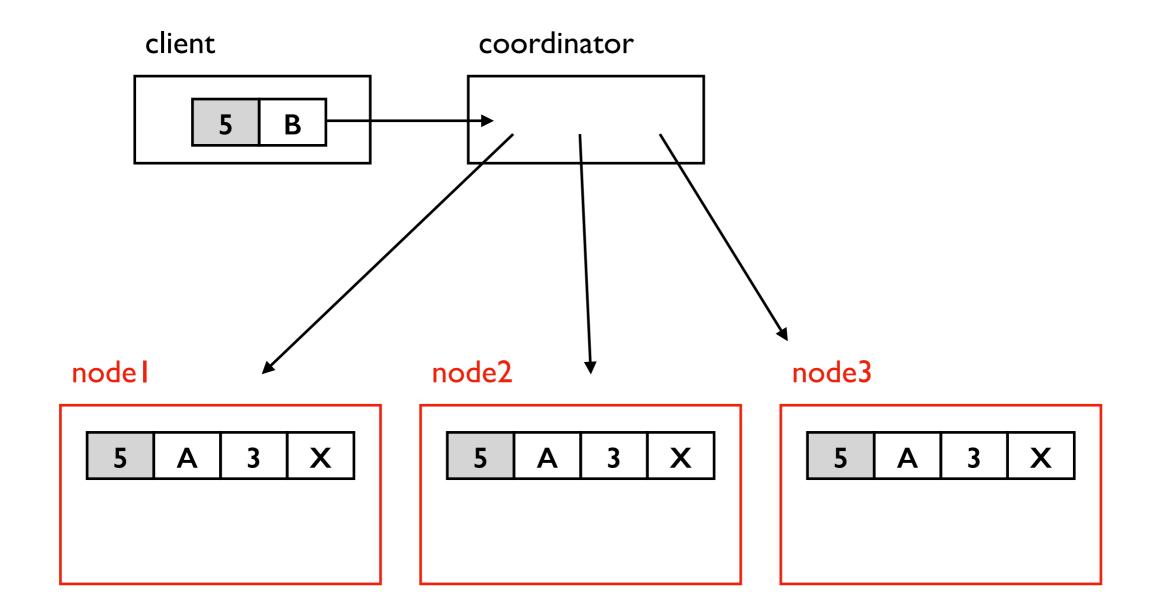


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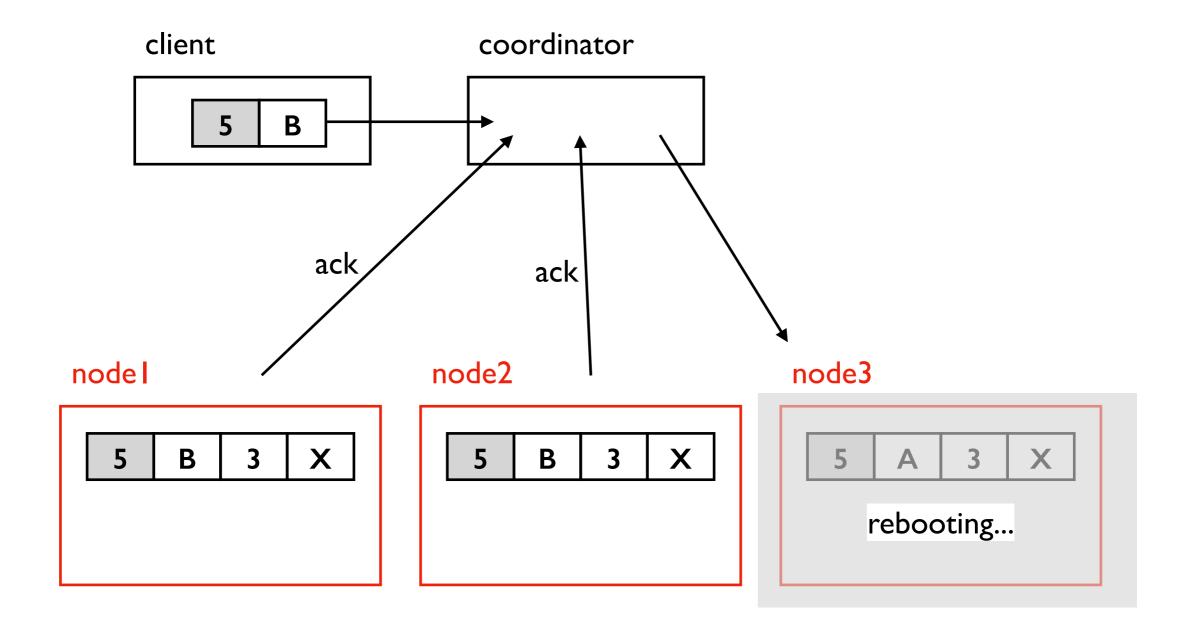


### Cassandra Writes



Say RF=3. Coordinator will attempt to write data to all 3 replicas.

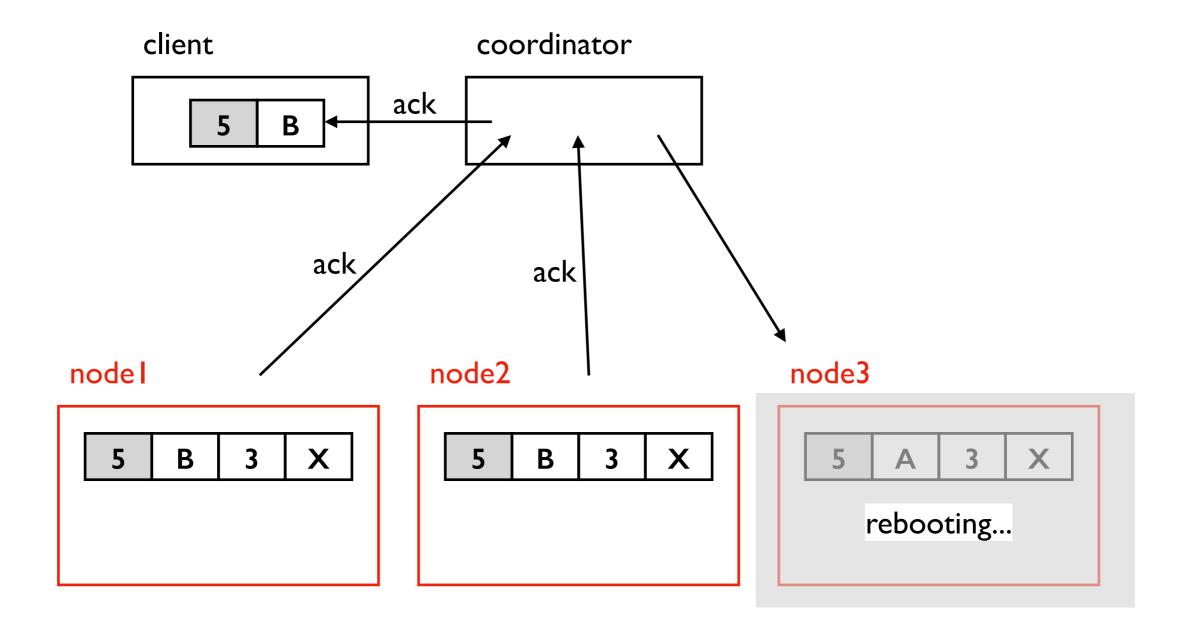
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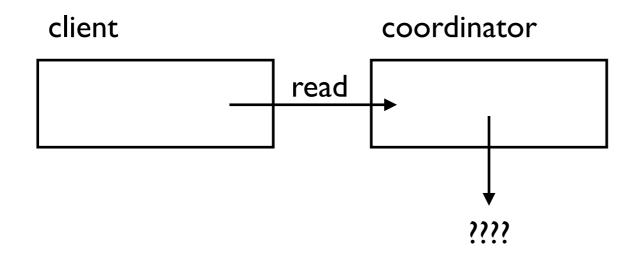
At what point should we send an ack to the client?

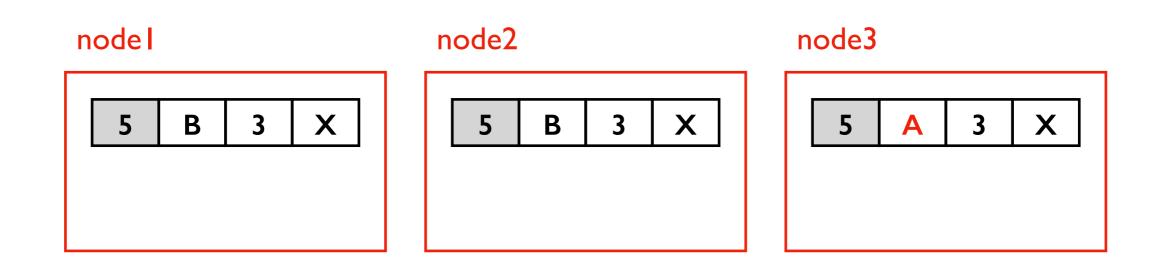
### Cassandra Writes



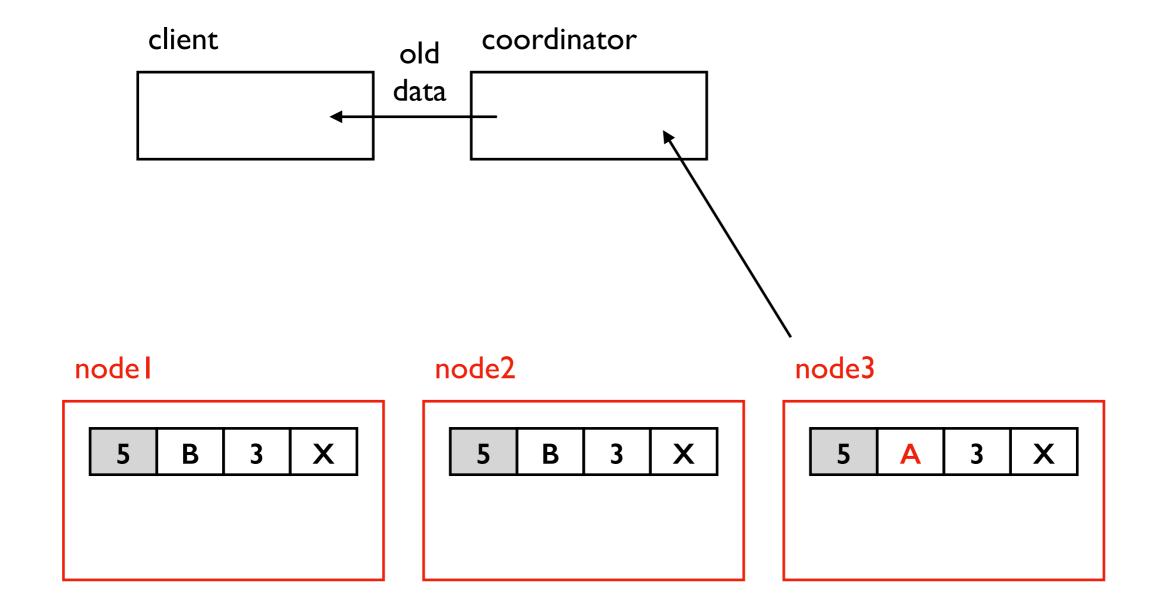
Say RF=3. Coordinator will attempt to write data to all 3 replicas.

At what point should we send an ack to the client? Configurable. W=2 lets coordinator ack now, and data is fairly safe.

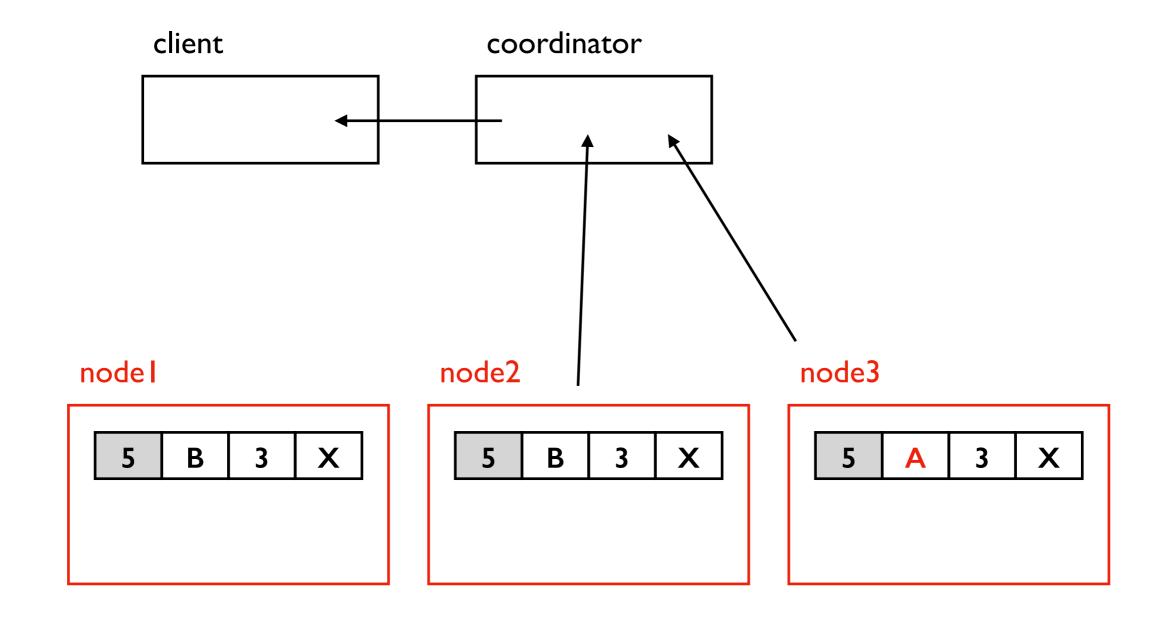




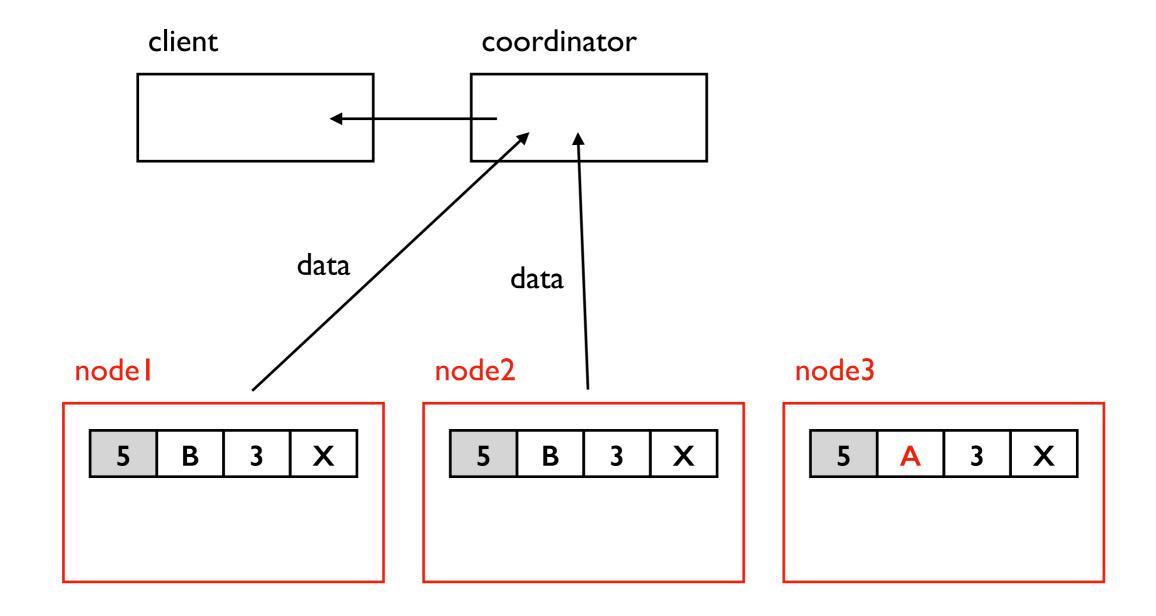
HDFS reads go to one replica. What if Cassandra tries that?



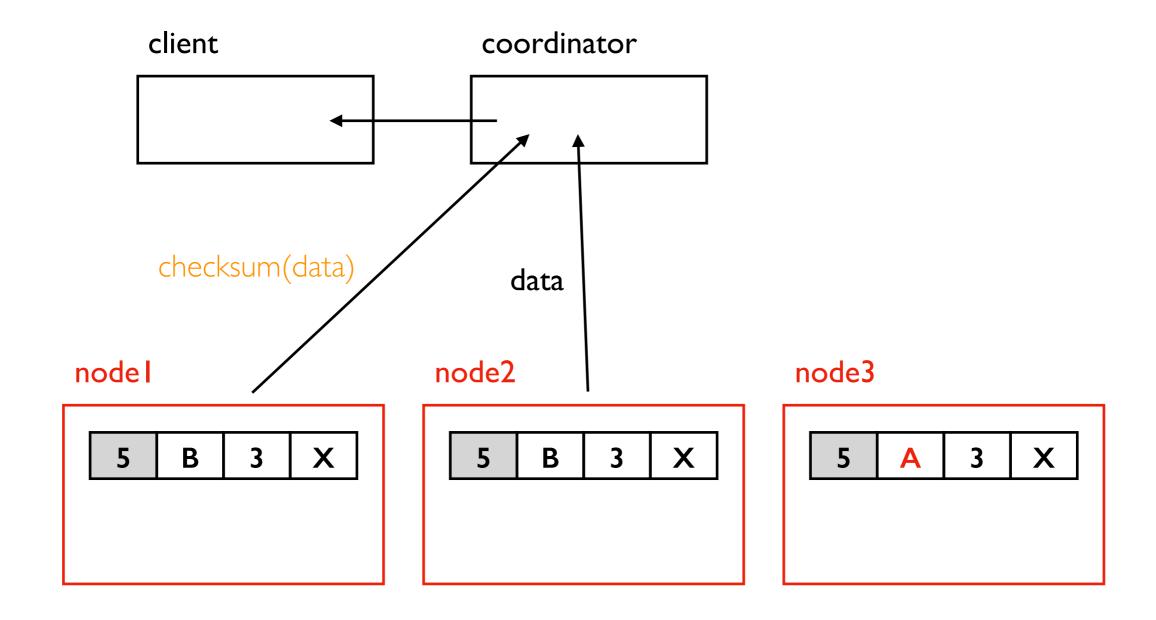
HDFS reads go to one replica. What if Cassandra tries that?



Read from R replicas (configurable). Here R=2. Hopefully at least one of the replicas has new data.



R=2 means we'll often read identical data from two replicas (wasteful!)

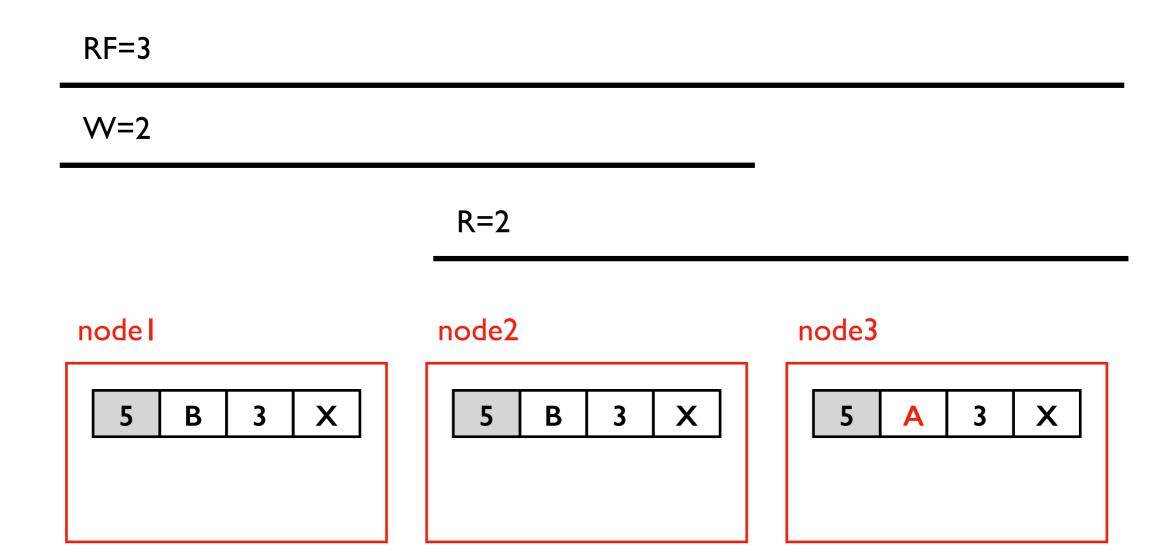


R=2 means we'll often read identical data from two replicas (wasteful!)

Improvement: read one copy, and only request checksum from others.

A checksum (like md5) is a hash function where collisions are extremely rare and hard to find.

### When R+W > RF



When R+W > RF, the replicas read+written will overlap.

There are some caveats (related to ring membership and something called "hinted handoff") not covered in 544.

## Tuning R and W

#### Say RF=3

#### W=3, R=I

- reads are highly available and fast -- only need one replica to respond before we can get back to the client!
- writes will not succeed (from the clients perspective) if even one node is down. But the data may still get recorded on some nodes.

#### W=1, R=3

- writes are highly available and fast -- only need one replica to respond before we can get back to the client!
- reads will not return data when even one node is down.
- risky: if the one node that took the write fails permanently, we'll lose committed data

#### W=2, R=2

relatively balanced approach

#### W=I, R=I

speed+availability more important that correct data

### Worksheet

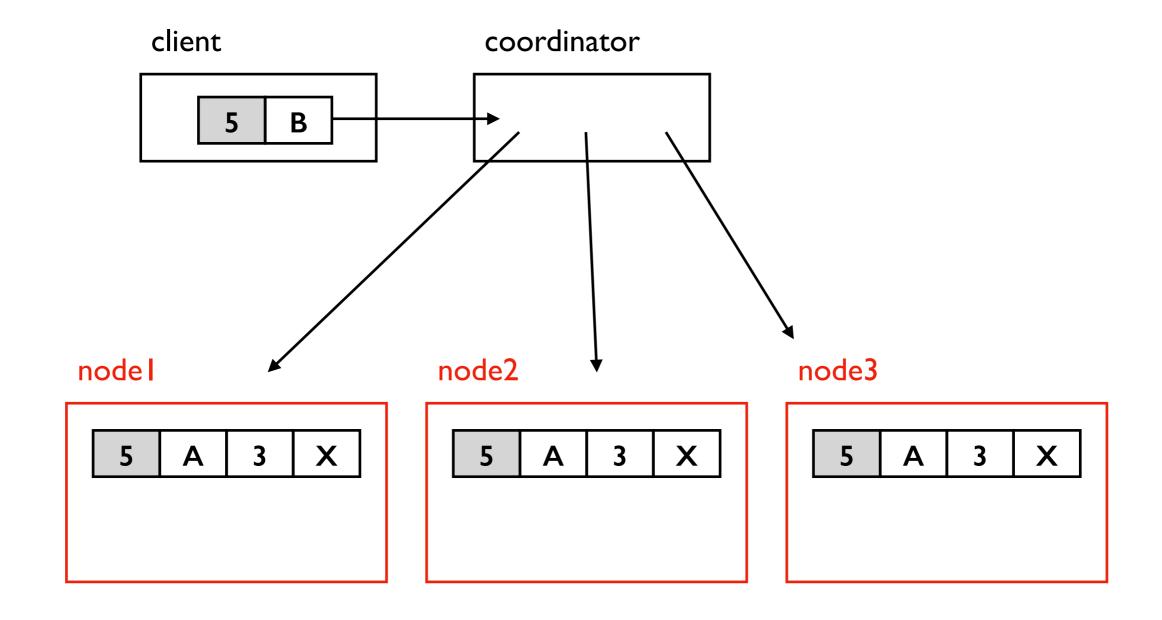
### Outline

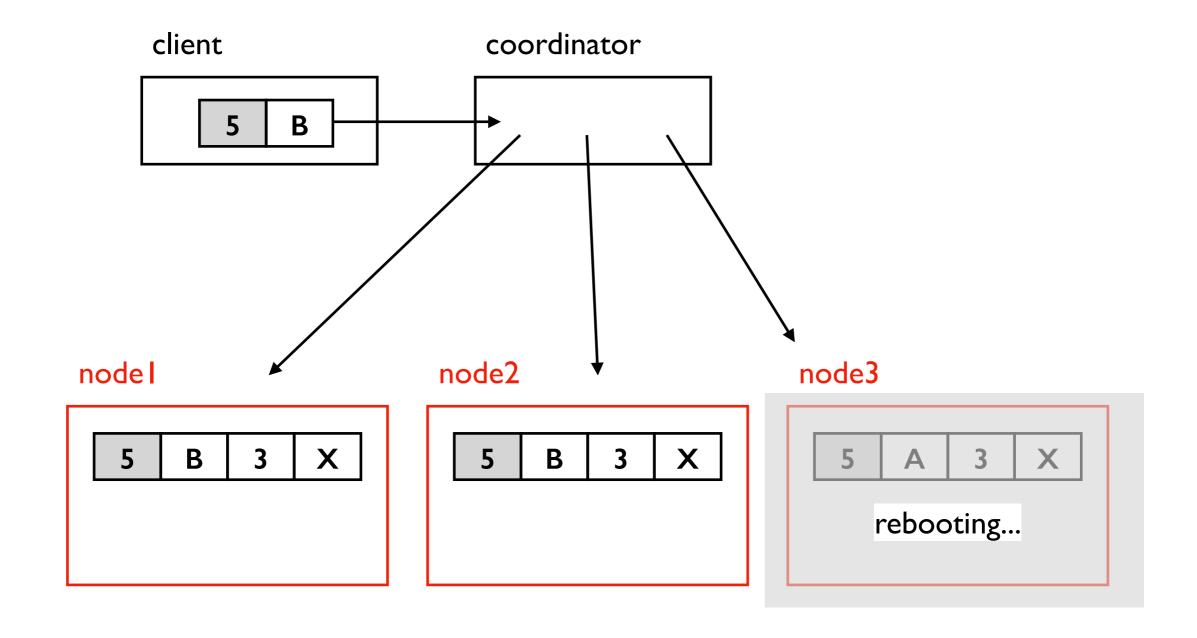
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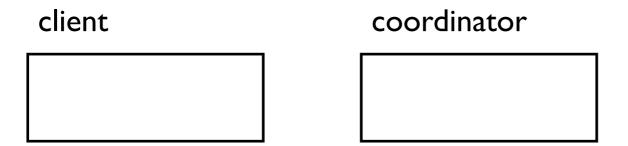
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**Conflict Resolution** 

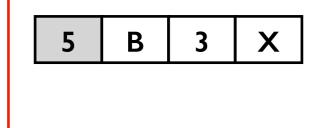
Cassandra Demos



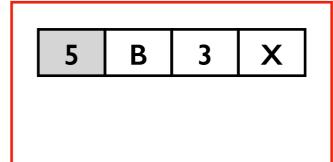




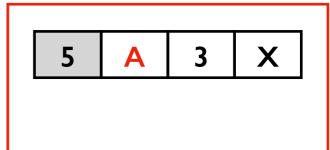
nodel

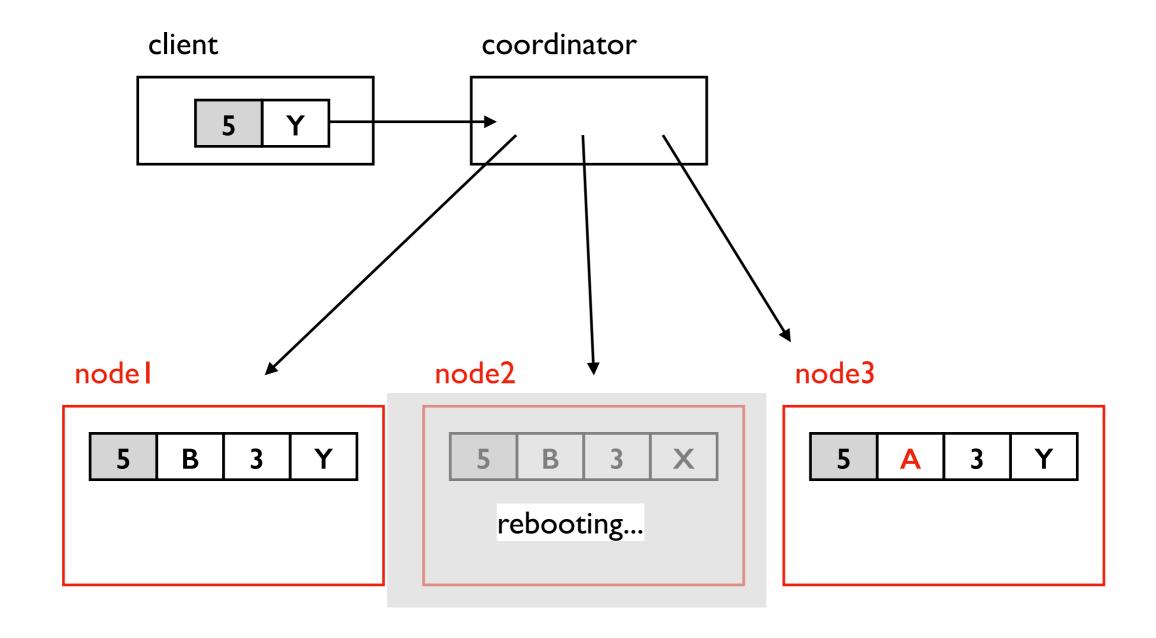


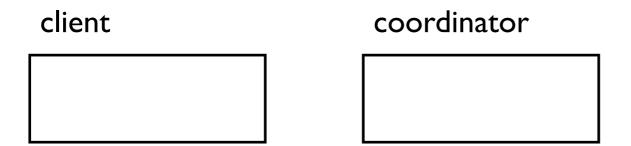
node2



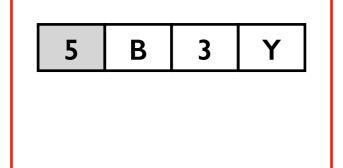
node3



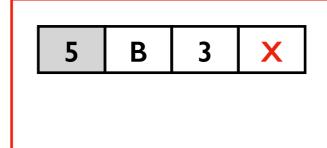




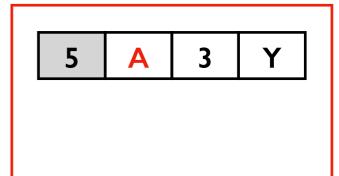
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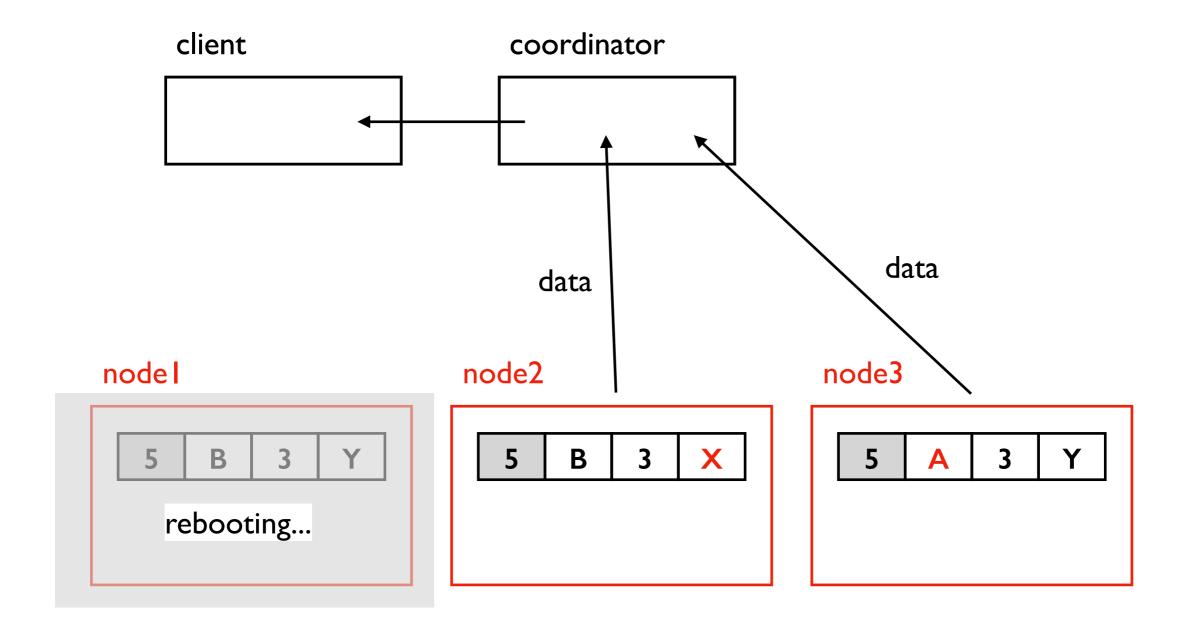


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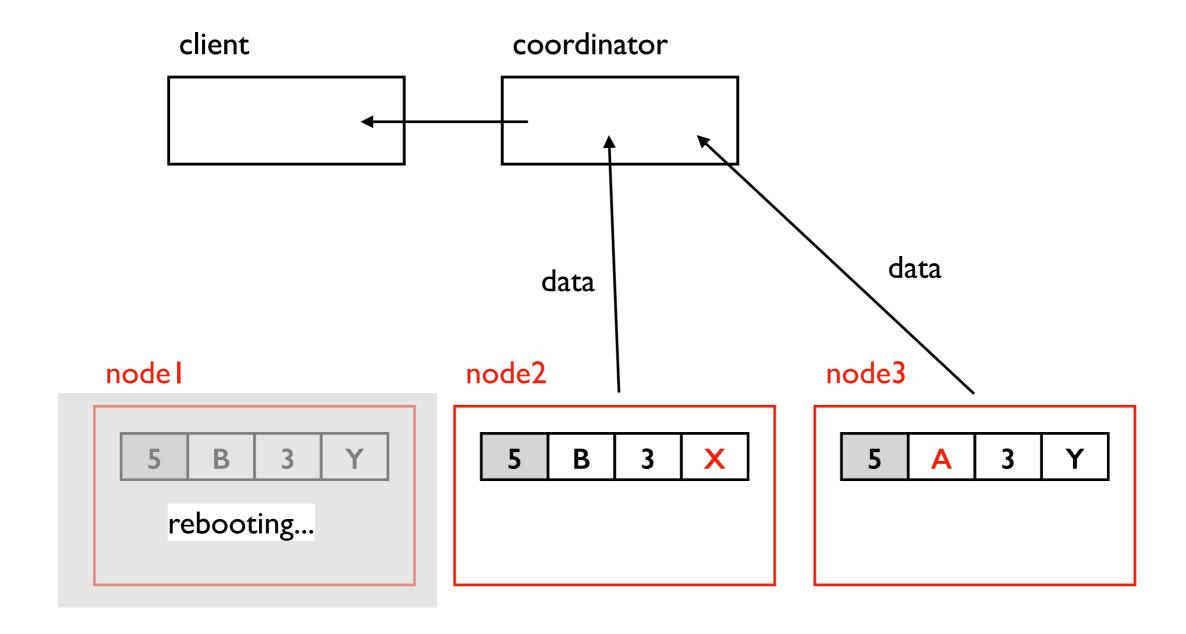
node3





Which version of row 5 should be sent back? Both contain some new data not contained by other.

Systems that allow conflicting versions to co-exist, fixing it up later are "eventually consistent"

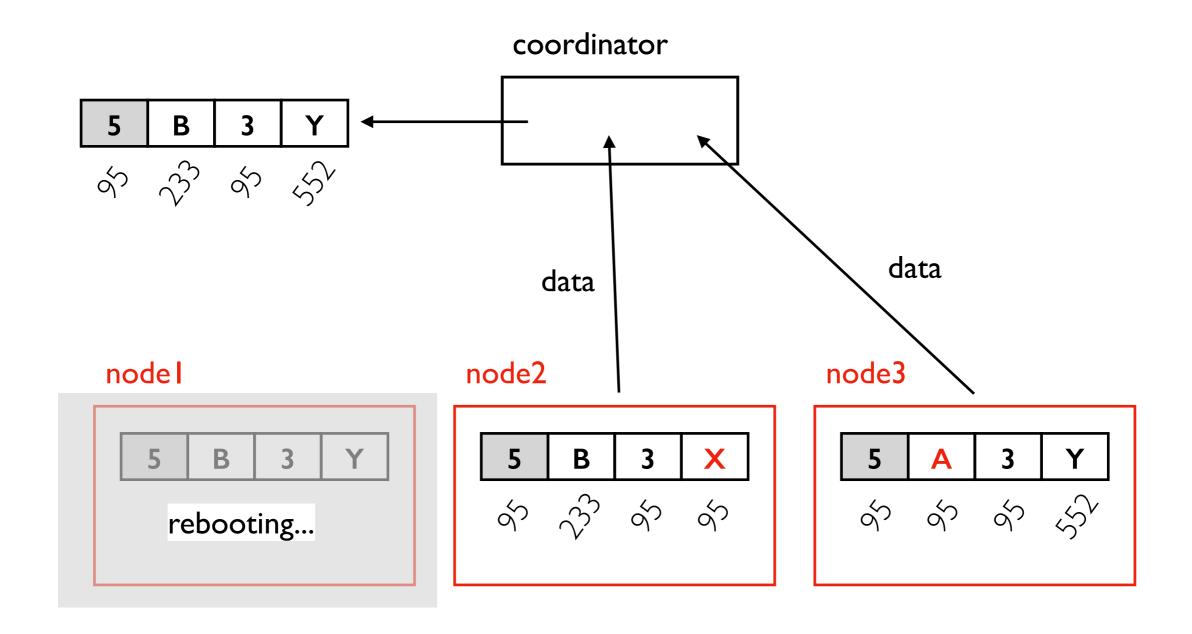


#### Approaches:

- send all version back to the client, which will need specialized conflict resolution code
- automatically combine them into a new row, and write that (if possible to all replicas)

Dynamo supports both. Cassandra uses second approach.

### Timestamps



#### Every cell of every table has a timestamp:

- approximate (since clocks of nodes in a cluster are never perfectly in sync)
- policy is LWW (last writer wins), meaning prefer newer data
- Cassandra lets you query the timestamp of each cell

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Conflict Resolution

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