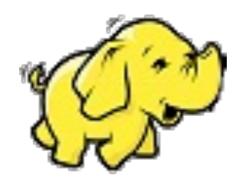
[544] Hadoop Ecosystem

Tyler Caraza-Harter



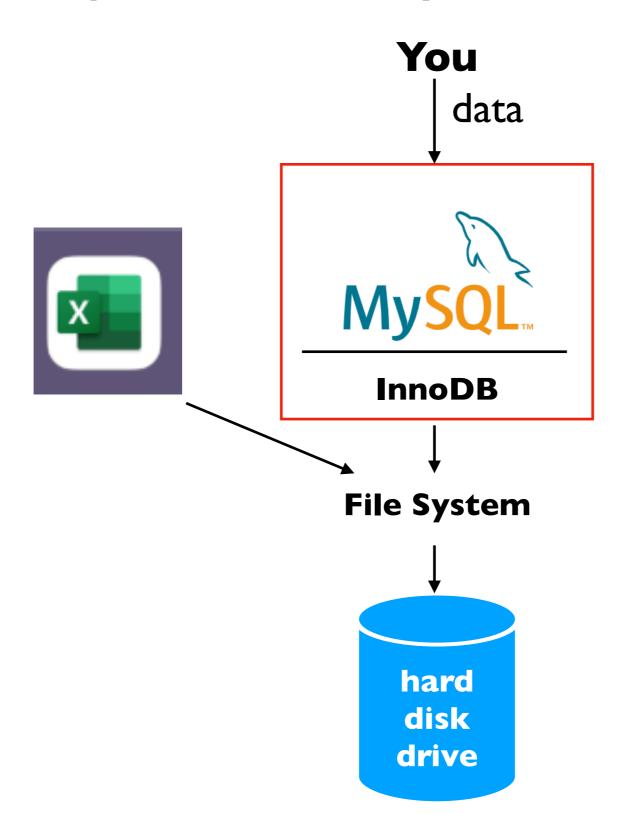
Outline: Hadoop Ecosystem

Architecture: Layered Data Systems

Hadoop File System

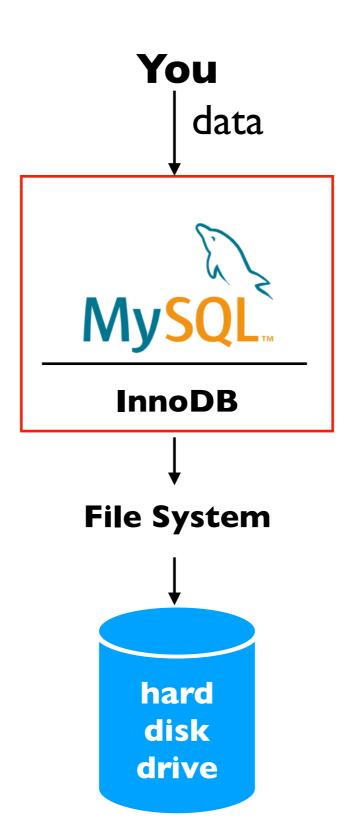
Hadoop MapReduce

Design: storage systems are generally built as a composition of layered subsystems

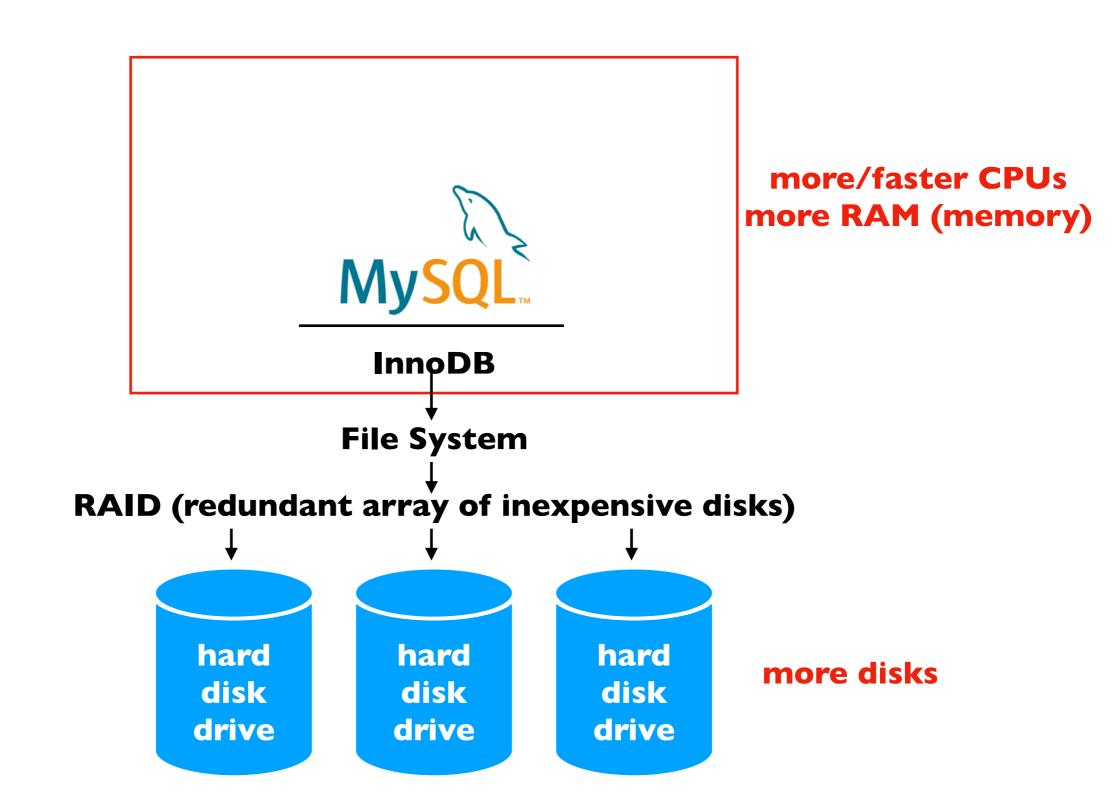


Today: 3 layered systems in the Hadoop Ecosystem

What if your data is too big for your server?

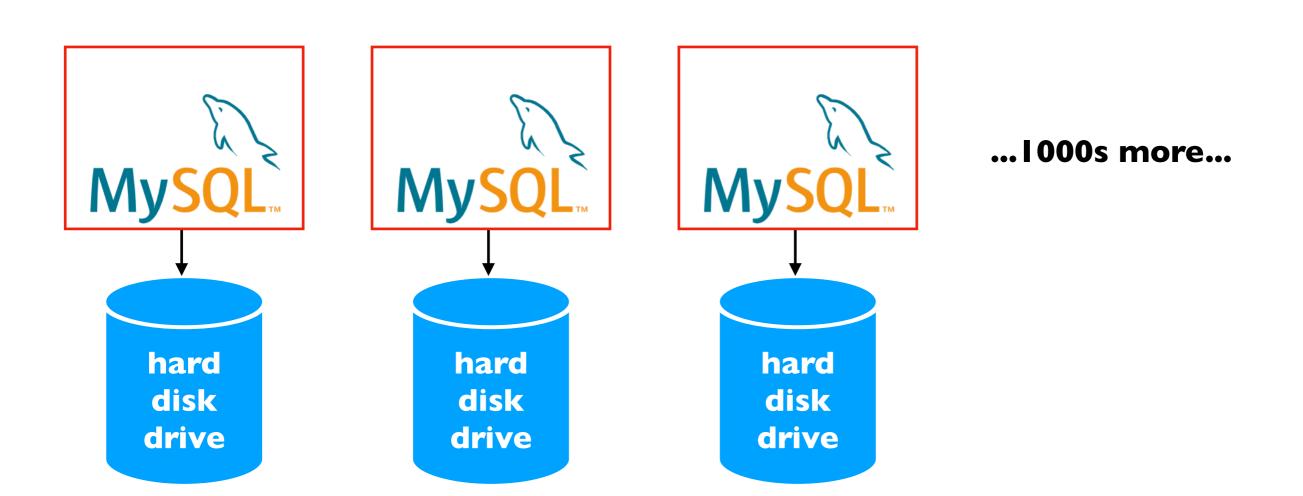


What if your data is too big for your server? Option I: scale up (buy better hardware)



What if your data is too big for your server? Option 2: scale out (more machines)

where does the data actually go?



Approach: partition the tables

tbl users

user ID name
1 "Yiyin"
2 "Ivan"
3 "Poulami"

tbl_purchases

user ID amt
2 \$10
2 \$15
3 \$20

• • •



tbl users

user ID name
1 "Yiyin"
2 "Ivan"

tbl_purchases

user ID amt 3 \$20

• • •

tbl users

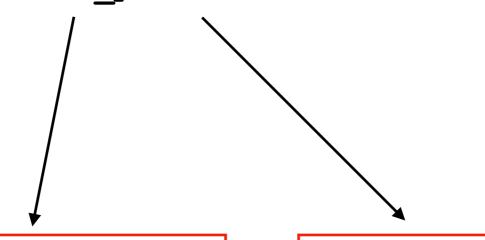
user ID name 3 "Poulami"

tbl_purchases

user ID amt
2 \$10
2 \$15

Approach: send queries to multiple DBs...

SELECT * FROM tbl_purchase WHERE amt > 12



tbl users

user ID name
1 "Yiyin"
2 "Ivan"

tbl_purchases

user ID amt 3 \$20

. . .

tbl users

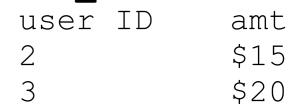
user ID name 3 "Poulami"

tbl_purchases

...combine results

SELECT * FROM tbl_purchase WHERE amt > 12

tbl_purchases



tbl users

user ID name
1 "Yiyin"
2 "Ivan"

tbl_purchases

user ID amt 3 \$20

• • •

tbl users

user ID name 3 "Poulami"

tbl_purchases

What is a query that would break things?

SELECT ...

tbl users

user ID name
1 "Yiyin"
2 "Ivan"

tbl_purchases

user ID amt 3 \$20

• • •

tbl users

user ID name 3 "Poulami"

tbl_purchases

What is a query that would break things?

```
SELECT * FROM tbl_users
INNER JOIN tbl_purchases
ON tbl_users.user_id = tbl_purchases.user_id
```

```
tbl_users
```

user ID name
1 "Yiyin"
2 "Ivan"

tbl_purchases

user ID amt \$20

• • •

tbl users

tbl_purchases

Why use a traditional/relational DB if basic things like JOIN don't work right at scale?

example: Cassandra documentation

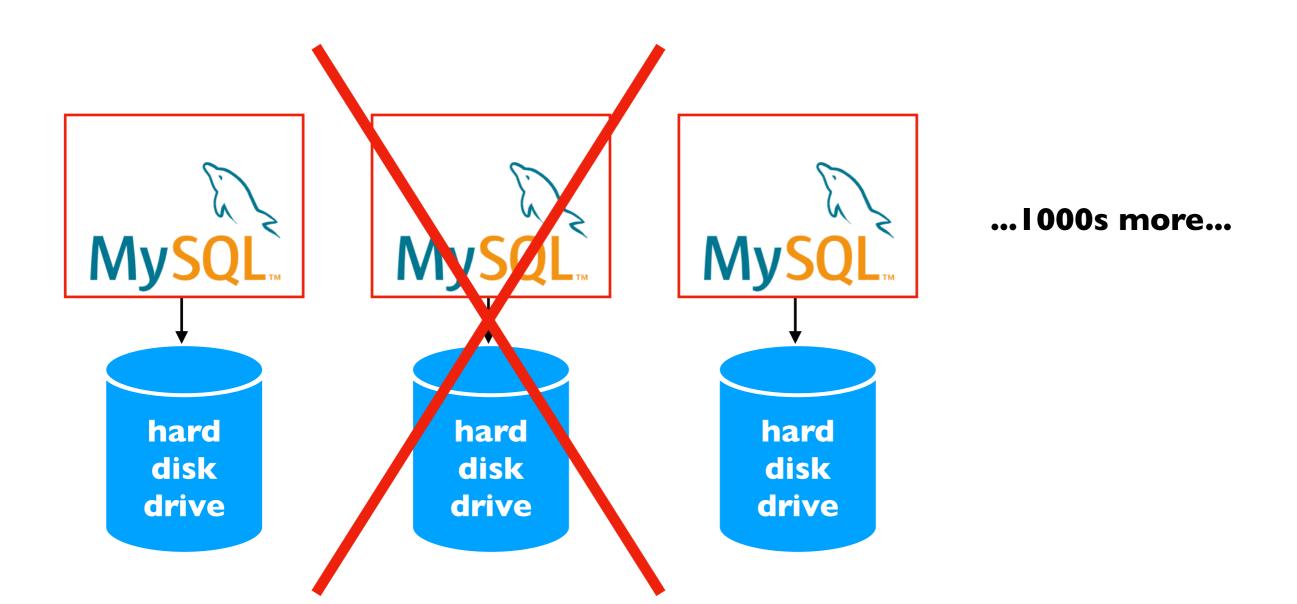
STEP 3: CREATE FILES

The Cassandra Query Language (CQL) is very similar to SQL but suited for the JOINless structure of Cassandra.

https://cassandra.apache.org/_/quickstart.html

What if a server dies?

happens all the time when you have 1000s of machines

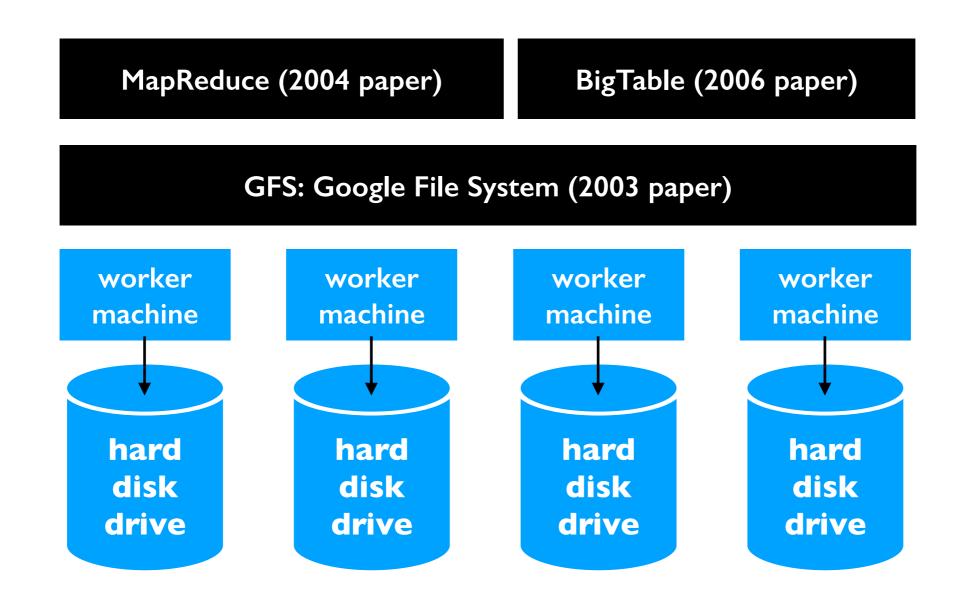


Motivation for System Redesign

Features

- some classic features (like JOINS) may not be essential
- scaling to many machines is essential
- fault tolerance is essential

Google Architecture



radical idea: base everything on lots of cheap, commodity hardware

Hadoop Ecosystem

Yahoo, Facebook, Cloudera, and others developed opensource Hadoop ecosystem, mirroring Google's systems

	Google (paper only)	Hadoop, 1st gen (open source)	Modern Hadoop
Distributed File System	GFS	HDFS covered	
Distributed Analytics	MapReduce	Hadoop MapReduce	Spark
Distributed Database	BigTable	HBase	Cassandra

Outline: Hadoop Ecosystem

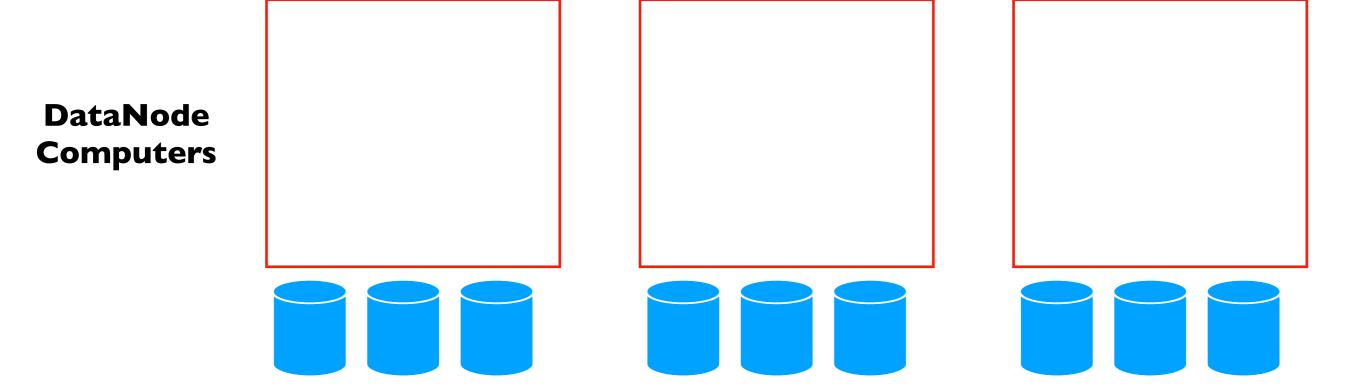
Architecture: Layered Data Systems

Hadoop File System

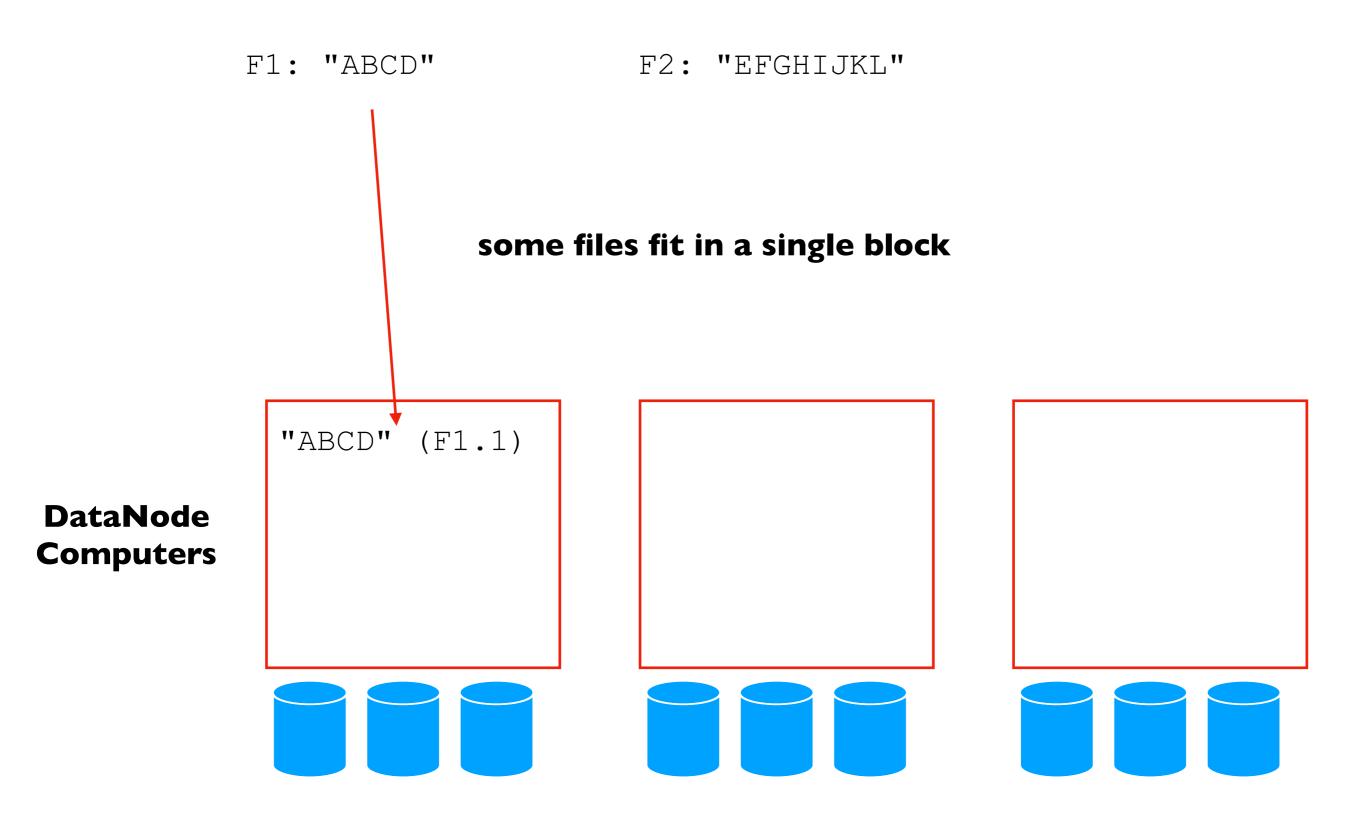
Hadoop MapReduce

HDFS: DataNodes store File Blocks

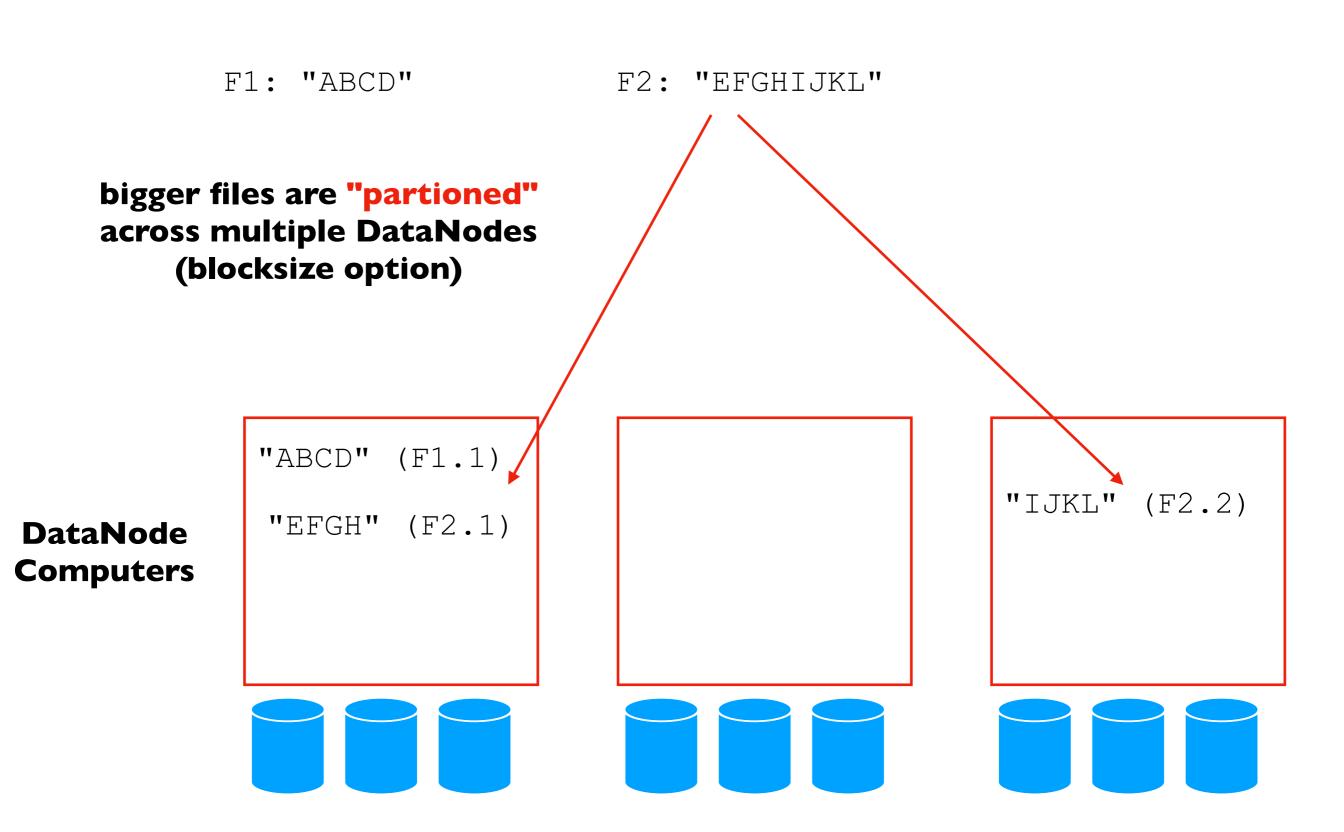
F1: "ABCD" F2: "EFGHIJKL"



HDFS: DataNodes store File Blocks



Partitioning Across DataNodes



Replication Across DataNodes

F1: "ABCD"

F2: "EFGHIJKL"

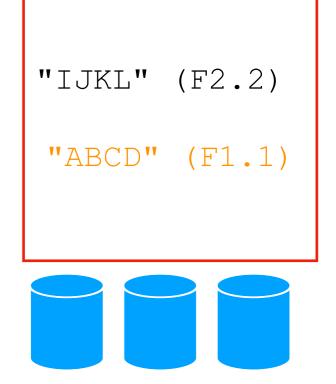
3x replication

2x relpication

DataNode Computers

```
"ABCD" (F1.1)
"EFGH" (F2.1)
```

```
"ABCD" (F1.1)
"EFGH" (F2.1)
"IJKL" (F2.2)
```



Replication Across DataNodes

F1: "ABCD"

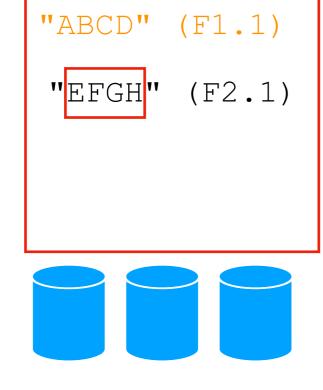
3x replication

F2: "EFGHIJKL"

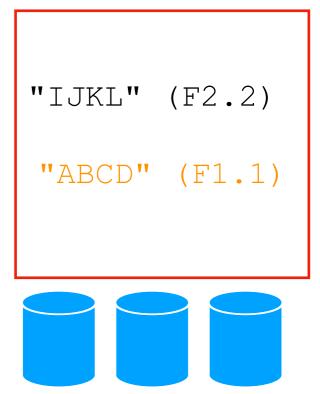
2x relpication

logical vs. physical blocks

DataNode Computers



```
"ABCD" (F1.1)
"EFGH" (F2.1)
"IJKL" (F2.2)
```



Replication Across DataNodes

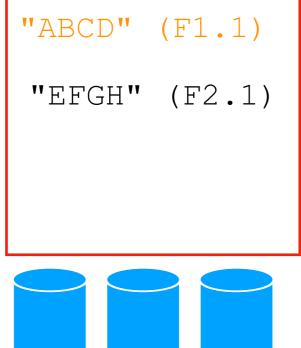
F1: "ABCD" F2: "EFGHIJKL"

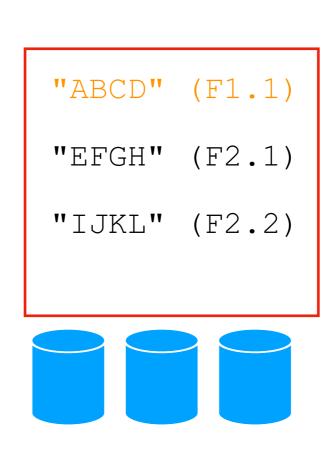
3x replication

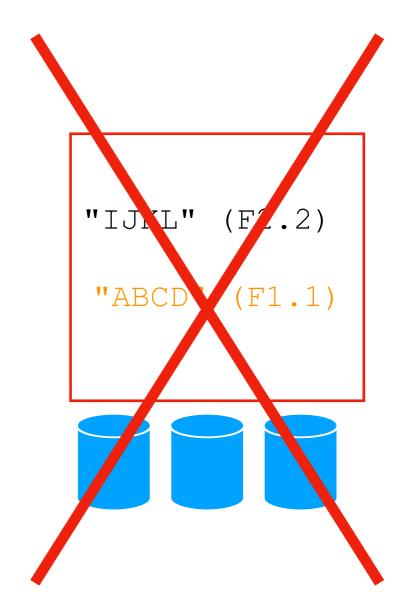
2x relpication

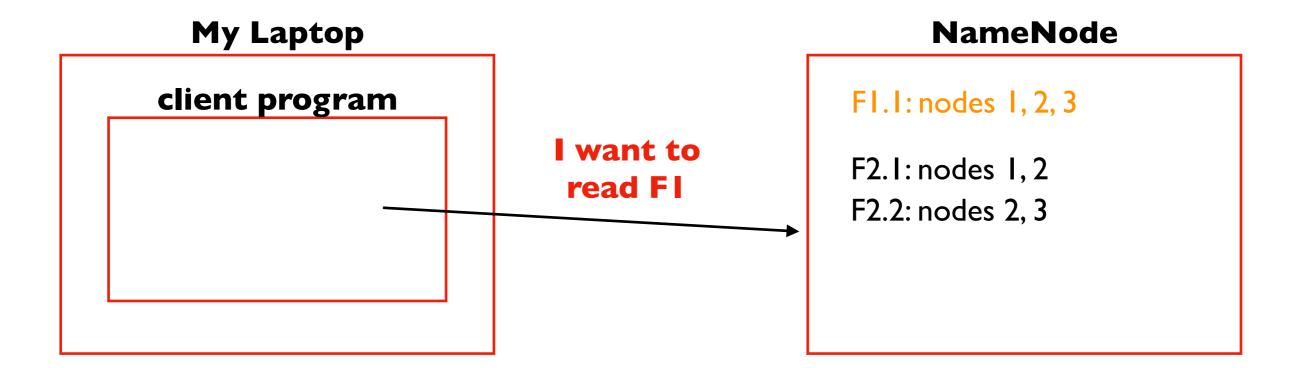
if a DataNode dies, we still have all the data. Which file (FI or F2) is safer in general?

DataNode Computers









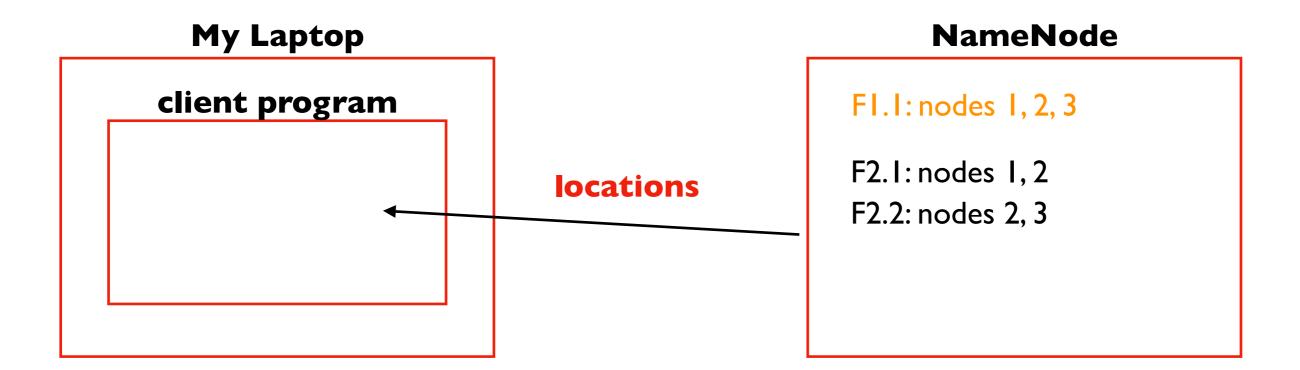
DataNode Computers

```
"ABCD" (F1.1)
"EFGH" (F2.1)
```

```
"ABCD" (F1.1)
"EFGH" (F2.1)
"IJKL" (F2.2)
```

"IJKL" (F2.2)
"ABCD" (F1.1)

DNI DN2 DN3

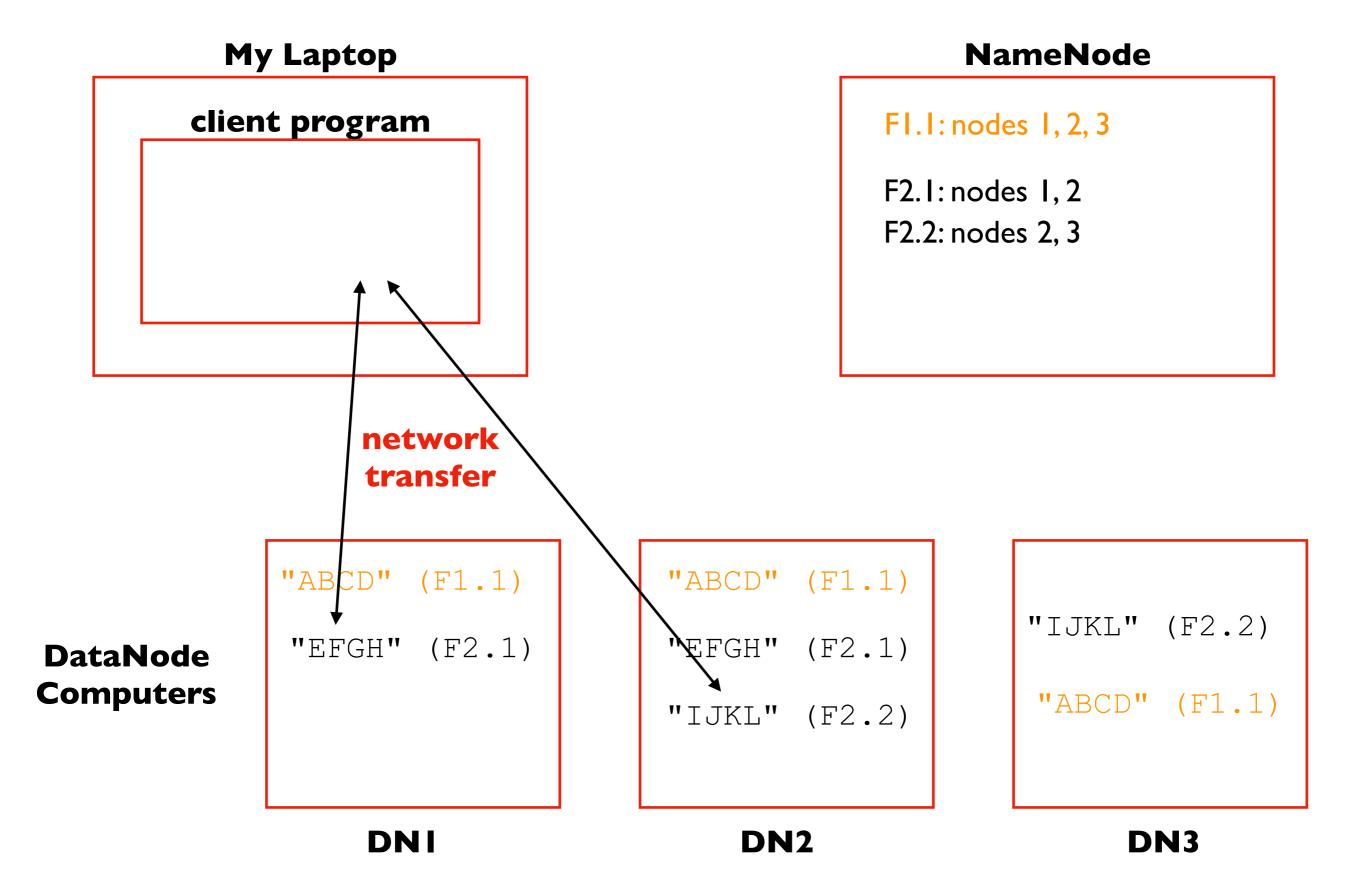


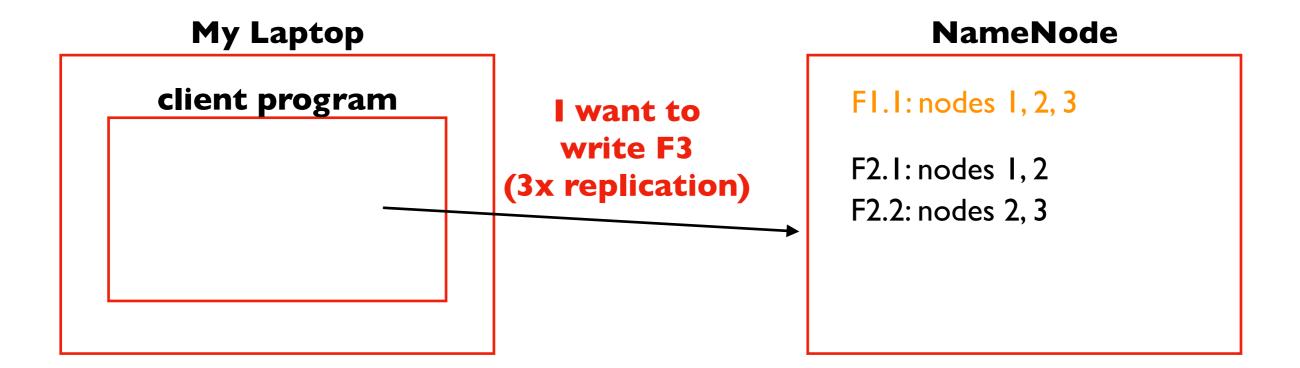
DataNode Computers

```
"ABCD" (F1.1)
"EFGH" (F2.1)
```

```
"ABCD" (F1.1)
"EFGH" (F2.1)
"IJKL" (F2.2)
```

DNI DN2 DN3

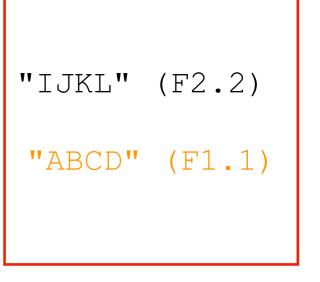




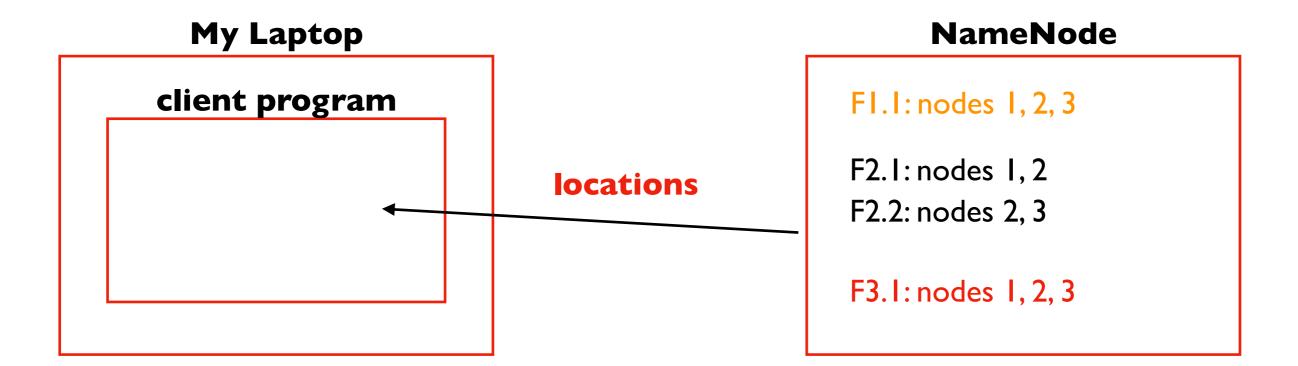
DataNode Computers

```
"ABCD" (F1.1)
"EFGH" (F2.1)
```

```
"ABCD" (F1.1)
"EFGH" (F2.1)
"IJKL" (F2.2)
```



DNI DN2 DN3



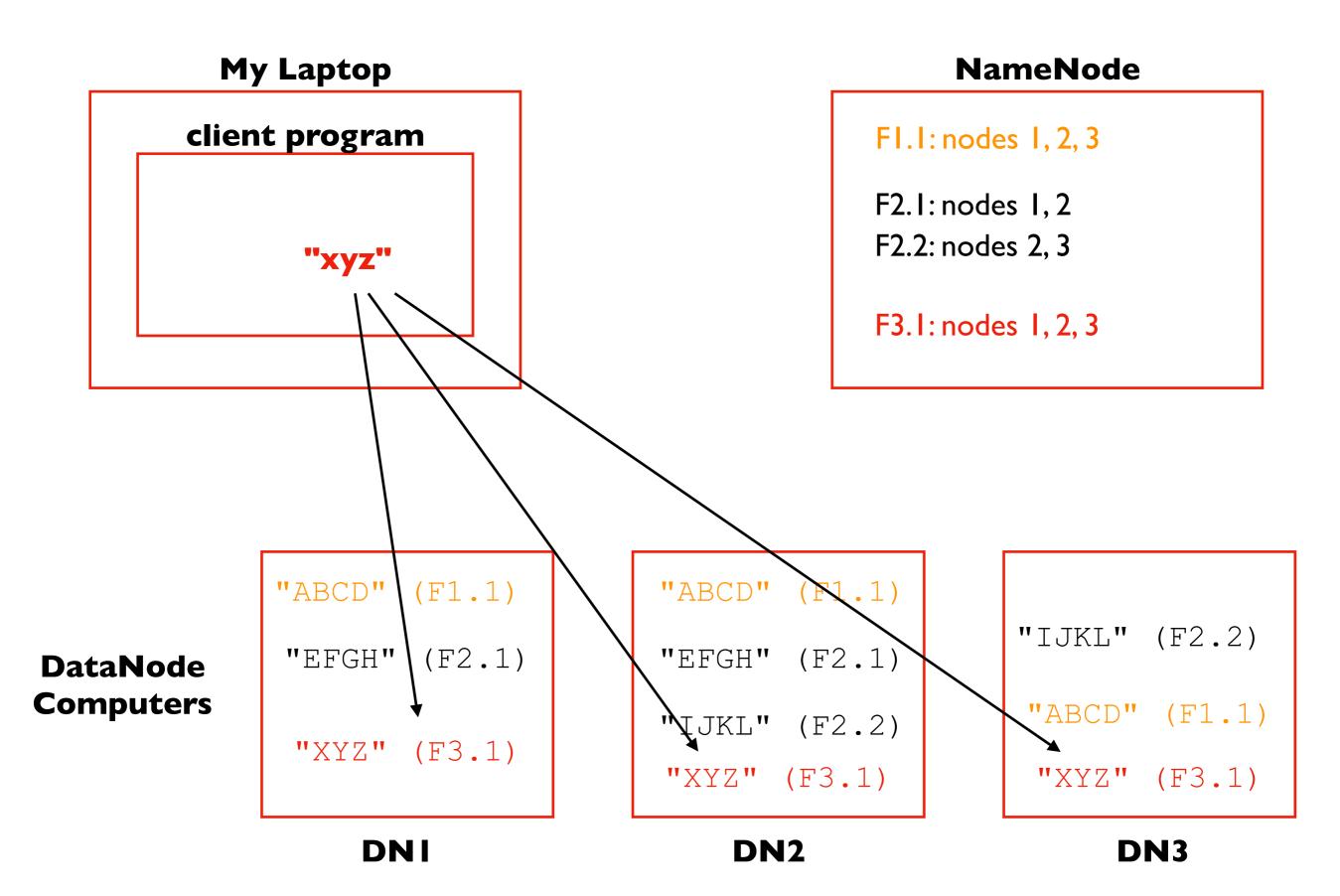
DataNode Computers

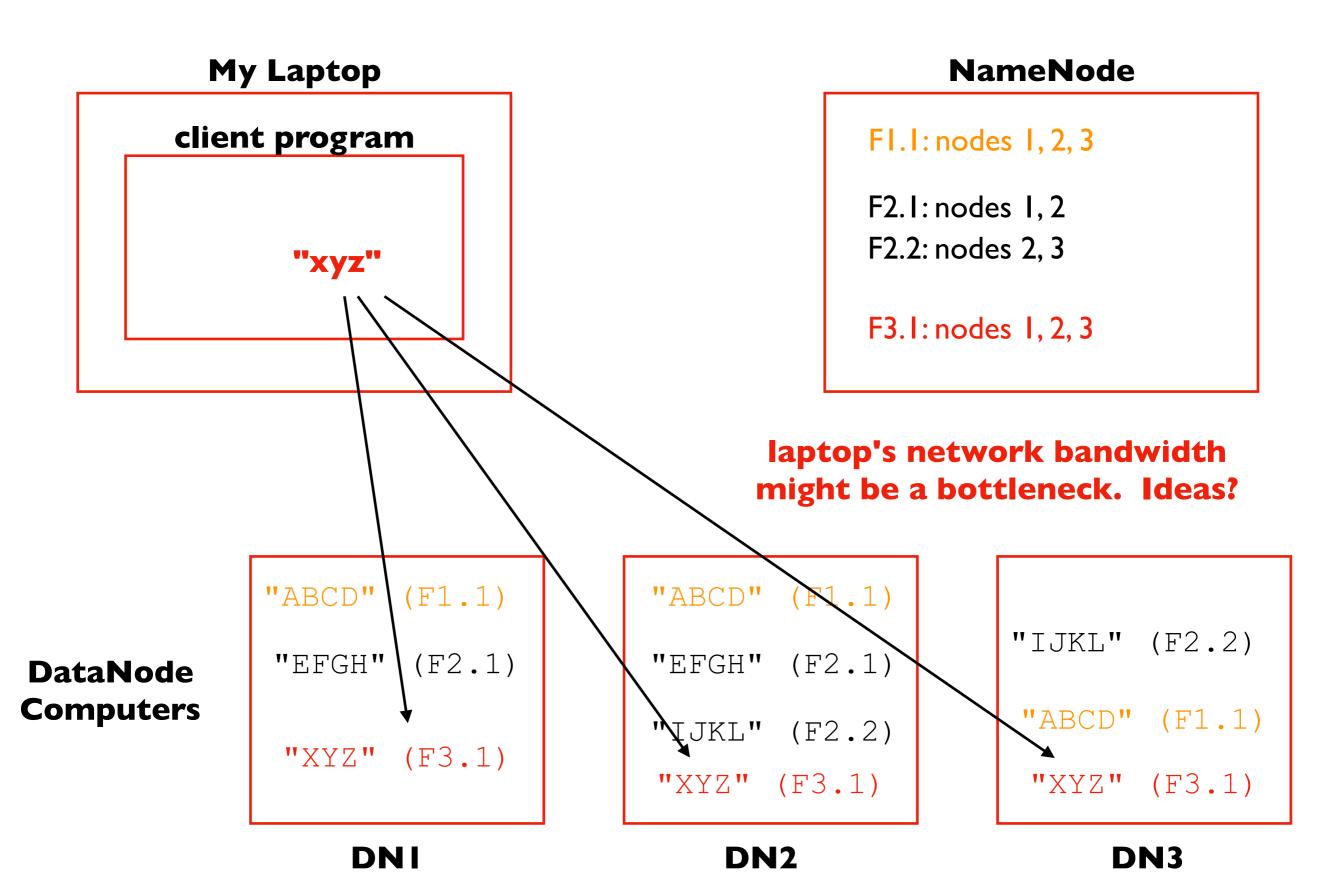
```
"ABCD" (F1.1)
"EFGH" (F2.1)
```

"ABCD" (F1.1)
"EFGH" (F2.1)
"IJKL" (F2.2)

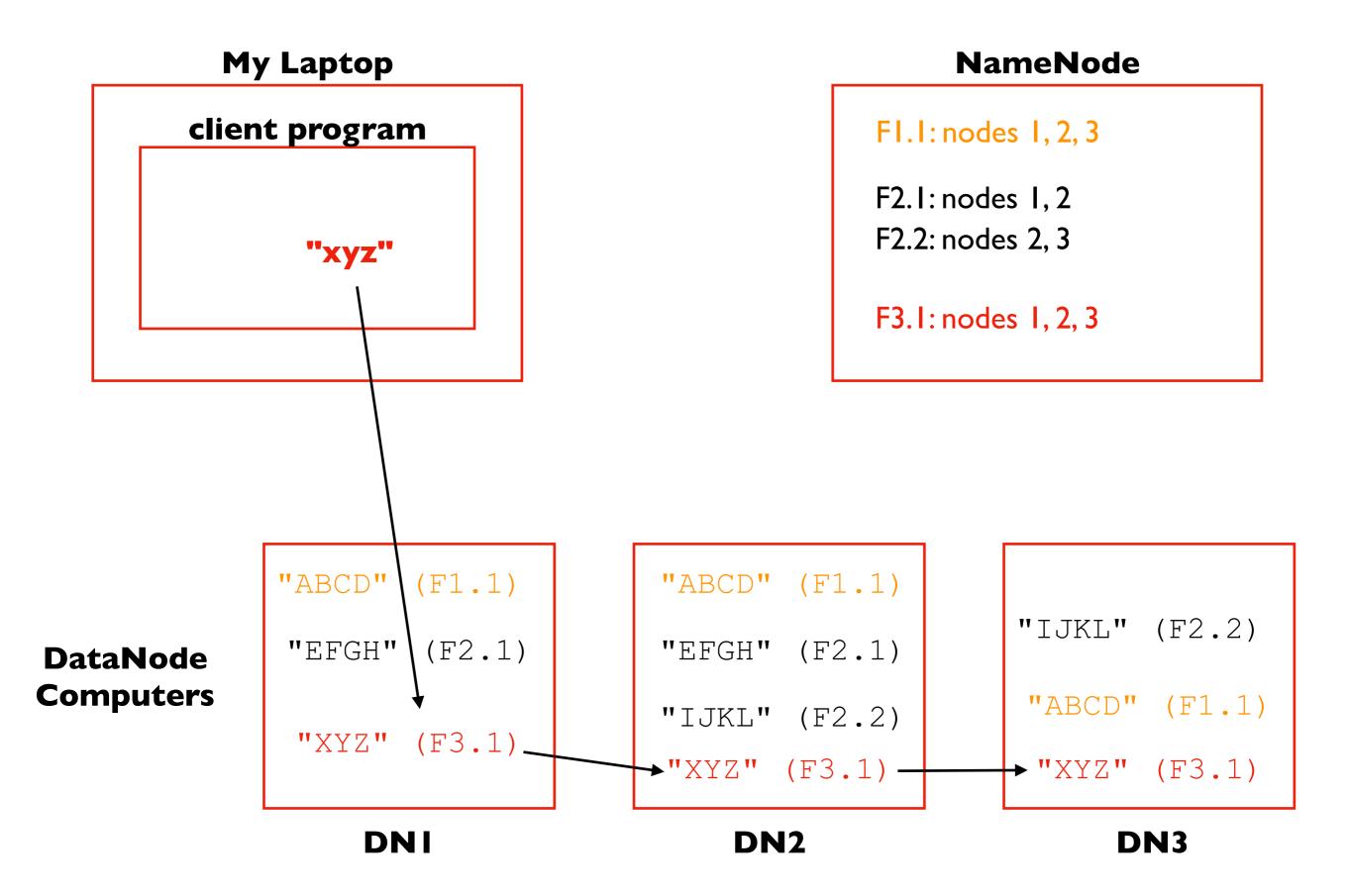
"IJKL" (F2.2)
"ABCD" (F1.1)

DNI DN2 DN3





Pipeline Writes



How are reads/writes amplified at disk level?

if a client **writes** 4 MB to a 2x replicated file, how much data do we **write** to hard drives?

if a client **reads** 2 MB to a 3x replicated file, how much data do we **read** from hard drives?

NameNode

FI.I: nodes 1, 2, 3

F2.1: nodes 1, 2

F2.2: nodes 2, 3

F3.1: nodes 1, 2, 3

DataNode Computers

```
"ABCD" (F1.1)
"EFGH" (F2.1)
"XYZ" (F3.1)
```

```
"ABCD" (F1.1)

"EFGH" (F2.1)

"IJKL" (F2.2)

"XYZ" (F3.1)
```

```
"IJKL" (F2.2)

"ABCD" (F1.1)

"XYZ" (F3.1)
```

DNI DN2 DN3

What are the tradeoffs of replication factor and block size?

NameNode

benefits of high replication?

benefits of low replication?

benefits of large block size?

benefits of small block size?

```
FI.1: nodes 1, 2, 3
```

F2.1: nodes 1, 2

F2.2: nodes 2, 3

F3.1: nodes 1, 2, 3

DataNode Computers

```
"ABCD" (F1.1)

"EFGH" (F2.1)

"XYZ" (F3.1)
```

```
"ABCD" (F1.1)

"EFGH" (F2.1)

"IJKL" (F2.2)

"XYZ" (F3.1)
```

DNI DN2 DN3

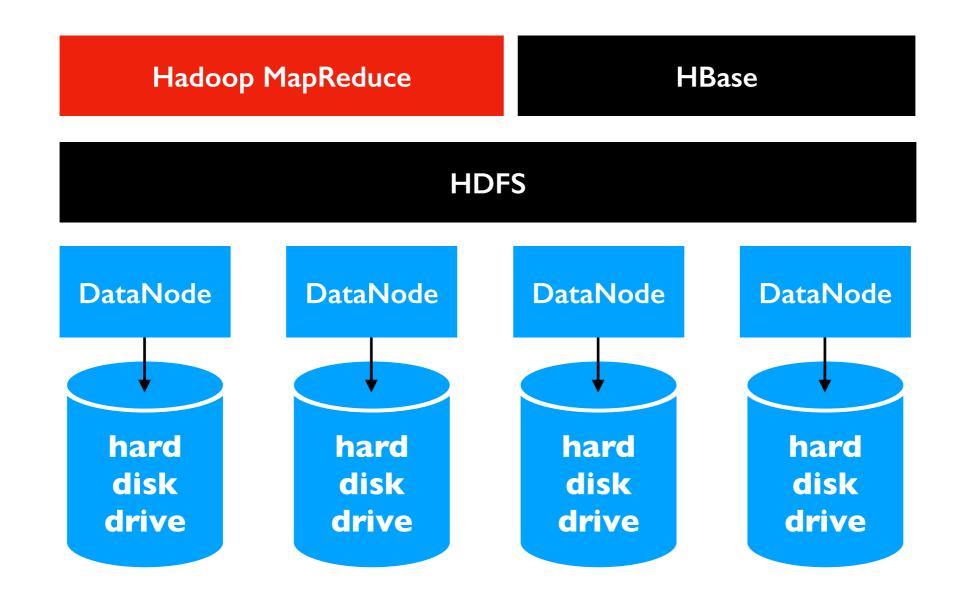
Outline: Hadoop Ecosystem

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Hadoop File System

Hadoop MapReduce

MapReduce

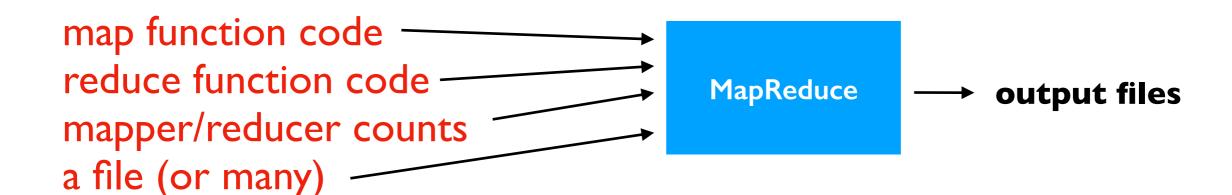


How do we answer questions?

SQL:

a query, "SELECT * FROM ..." → Database → results

MapReduce



How do we answer questions?

SQL: a query, "SELECT * FROM ..." → **Database** results MapReduce map function code reduce function code -**MapReduce** output files mapper/reducer counts a file (or many)

input/output files are generally in HDFS

How do we answer questions?

SQL: a query, "SELECT * FROM ..." → **Database** results MapReduce map function code reduce function code **MapReduce** output files mapper/reducer counts a file (or many)

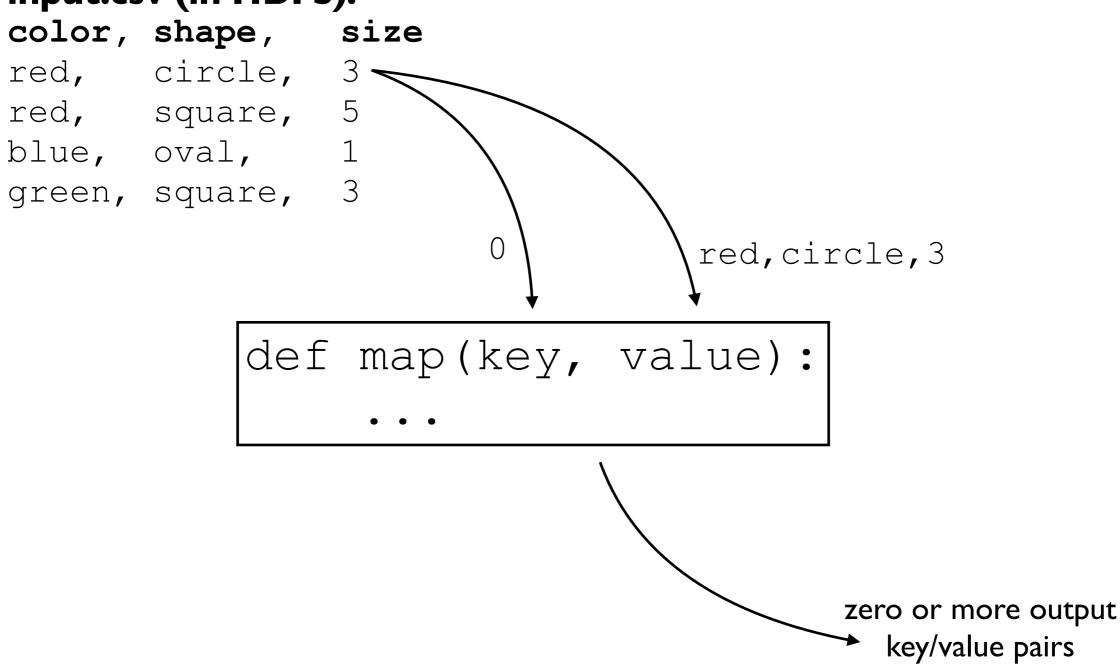
input.csv (in HDFS):

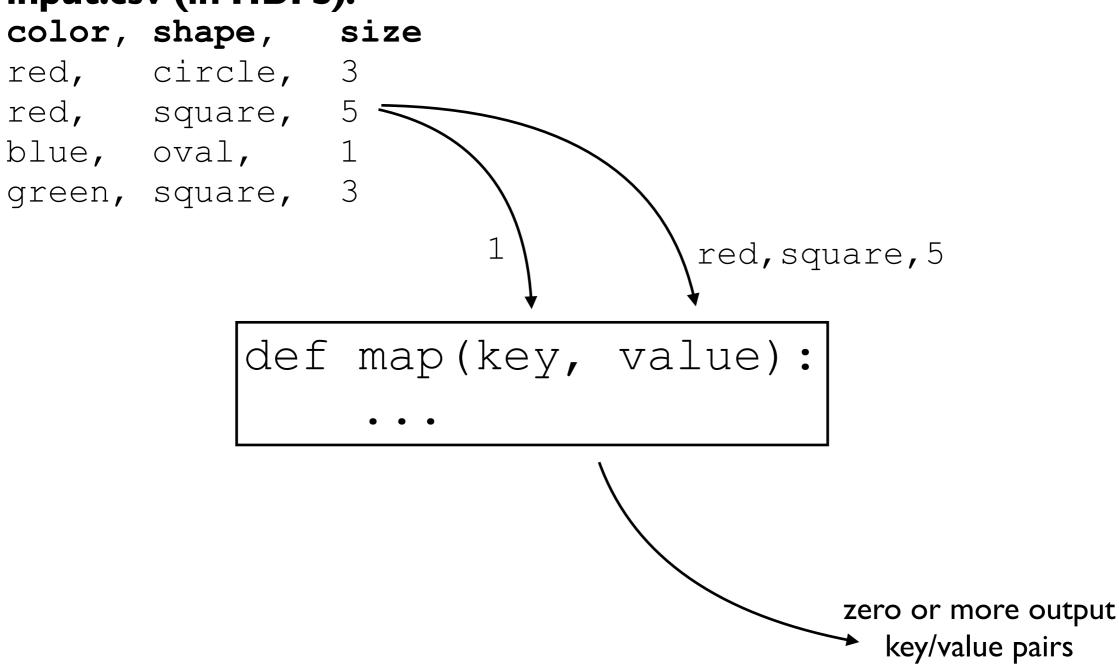
```
color, shape, size
red, circle, 3
red, square, 5
blue, oval, 1
green, square, 3
```

```
def map(key, value):
    ...
```

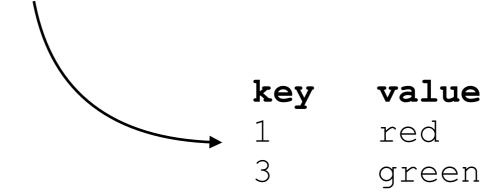
In SQL:

SELECT color FROM table WHERE shape = "square";



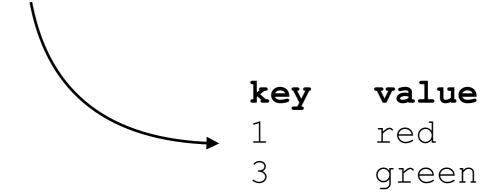


```
def map(key, value):
   if value.shape = square:
       emit(key, value.color)
```

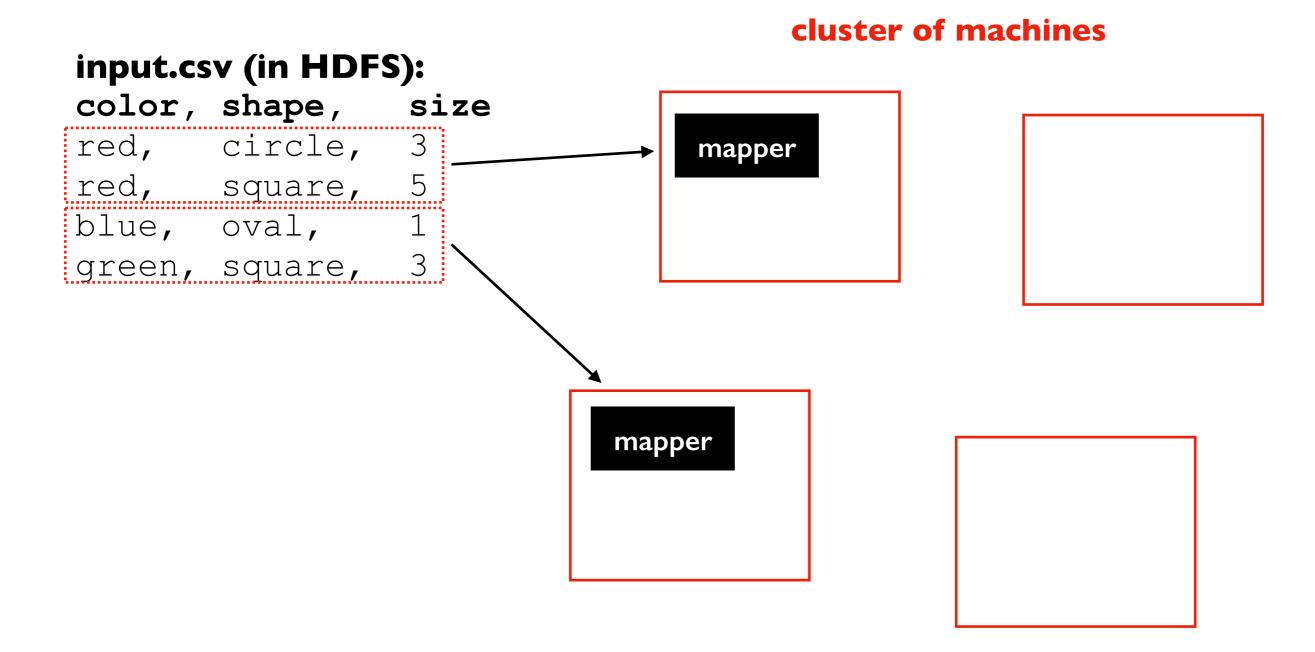


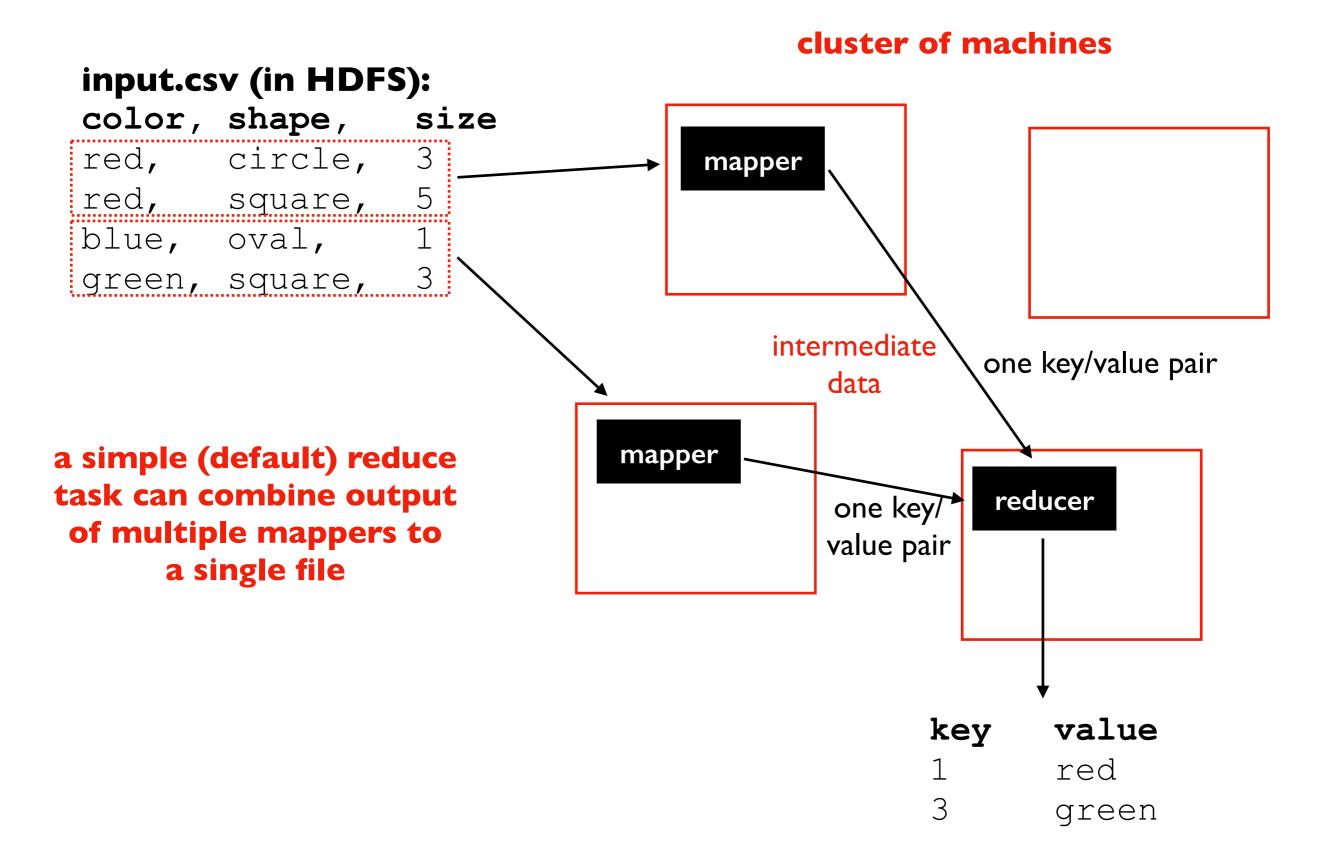
what if the data is huge?

```
def map(key, value):
   if value.shape = square:
       emit(key, value.color)
```



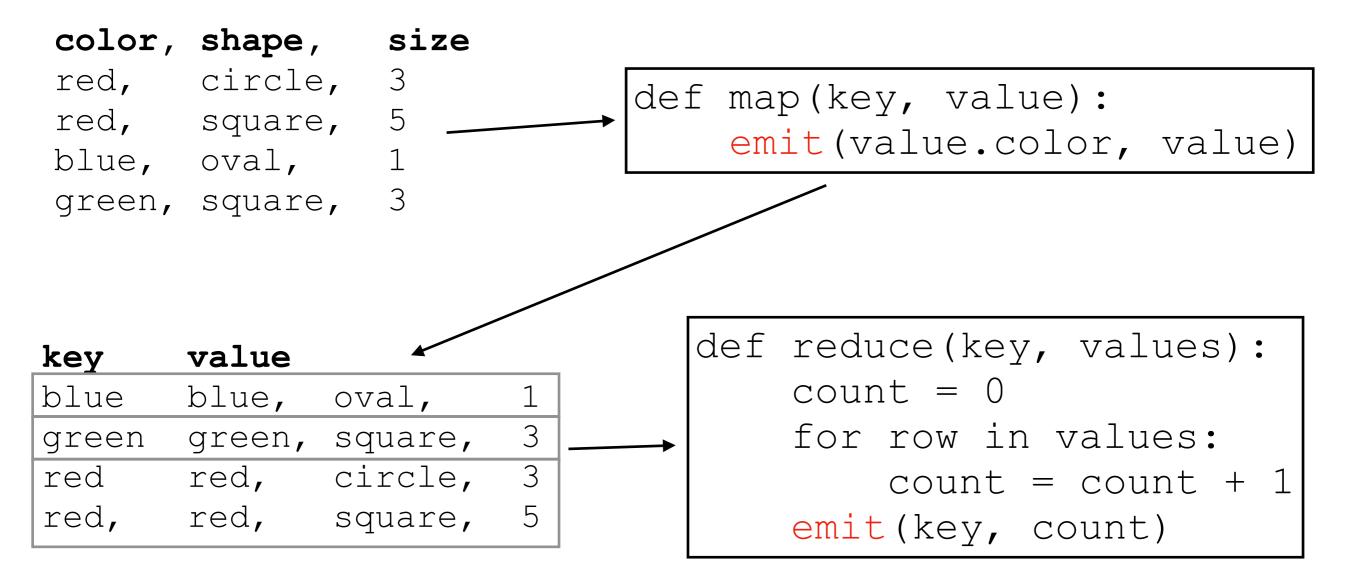
Mappers Run on Multiple Machines at Once



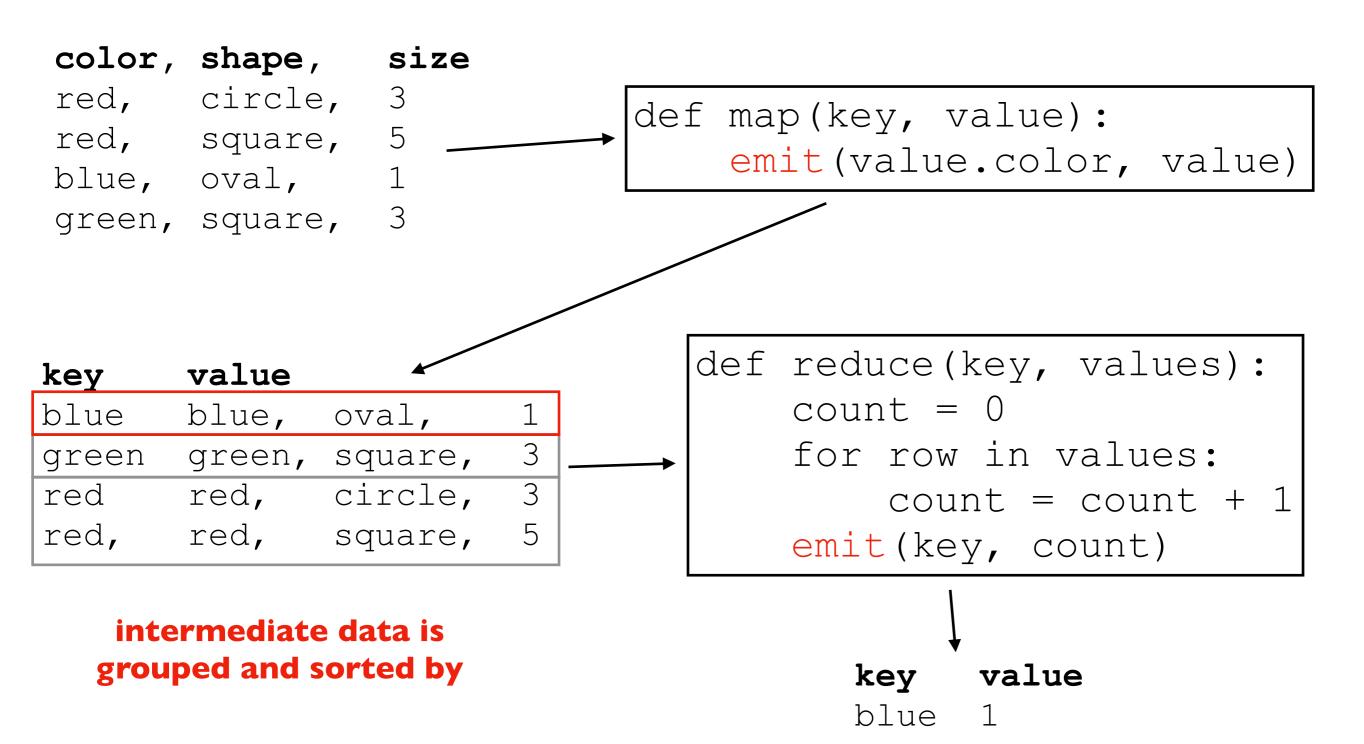


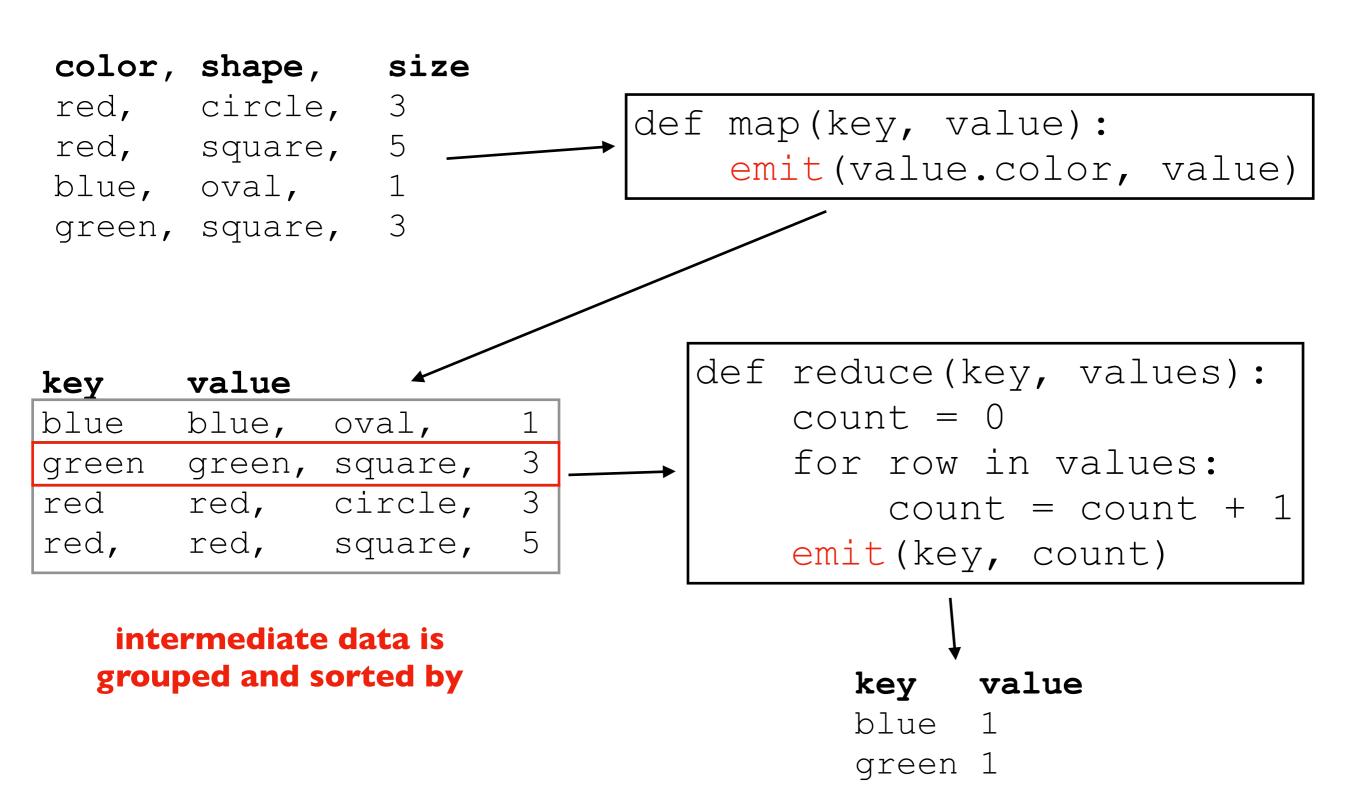
reducers can output exactly their input, OR have further computation

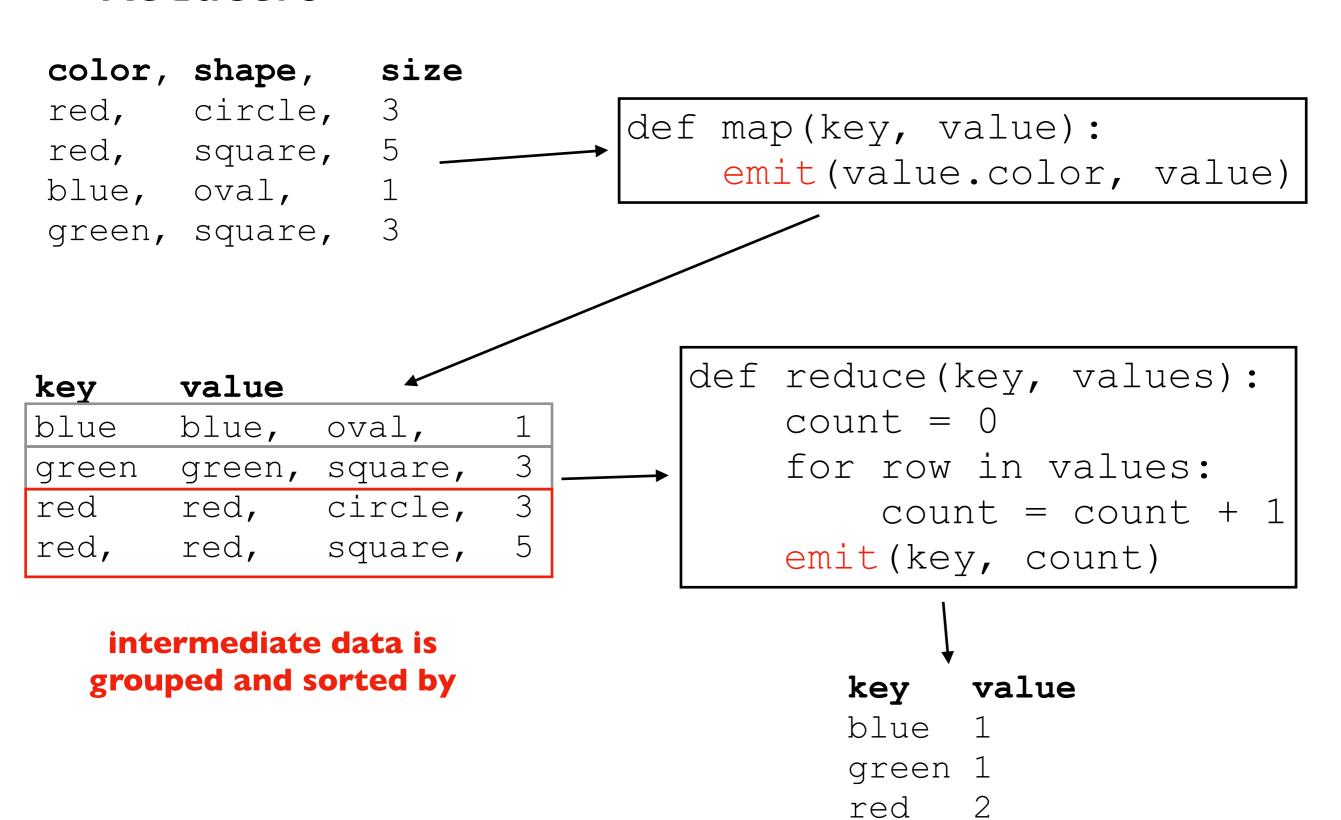
```
def reduce(key, values):
   for row in values:
       emit(key, row)
```



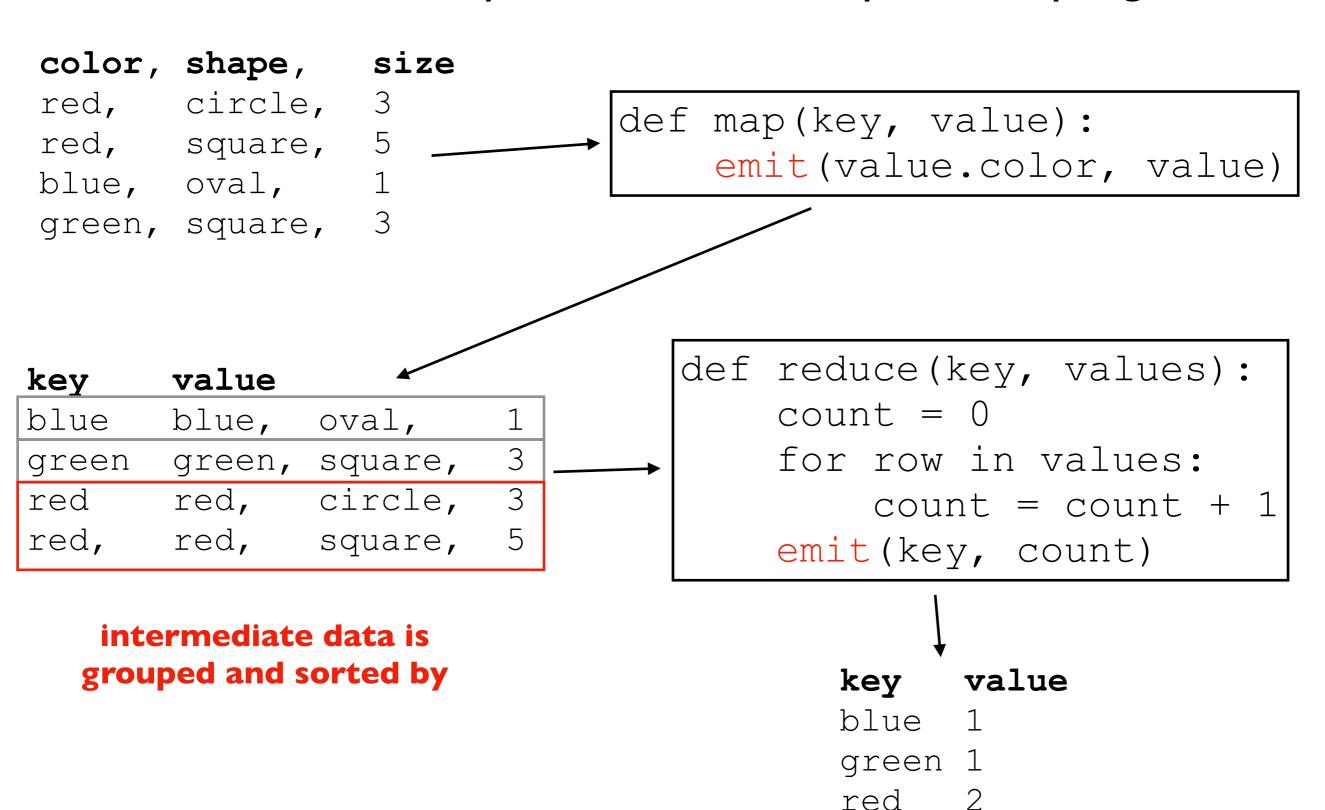
intermediate data is grouped and sorted by





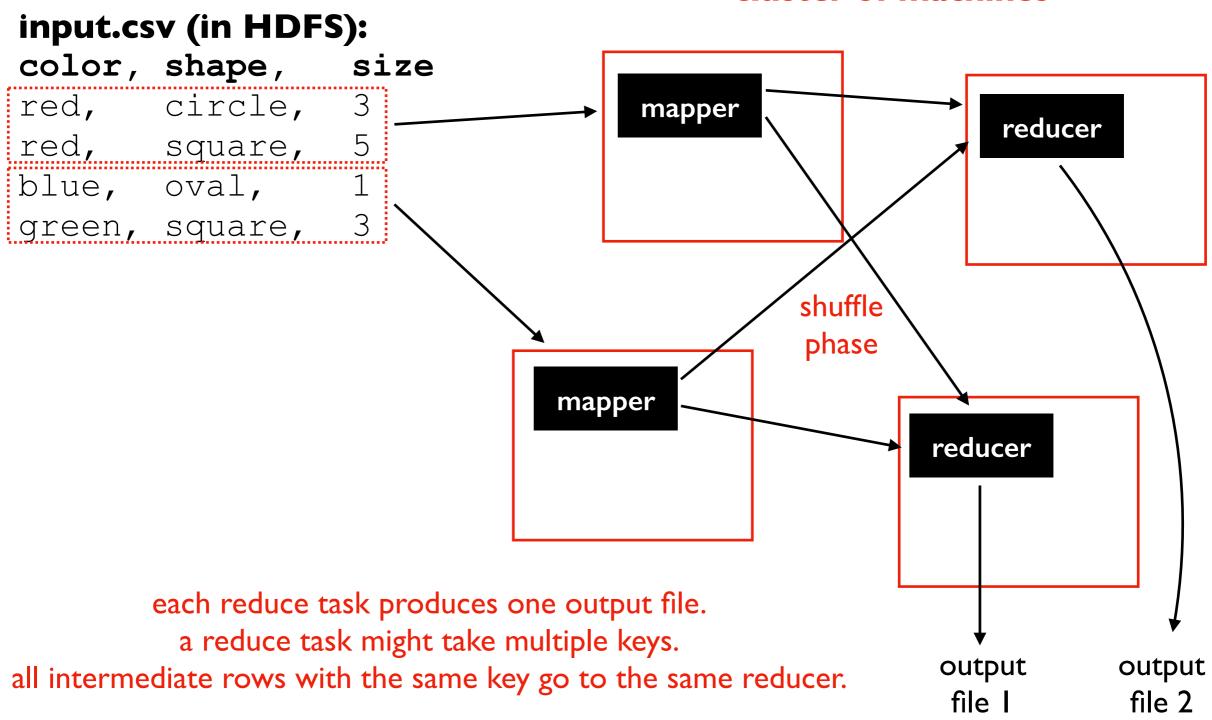


What is the SQL equivalent of this MapReduce program?



Multiple Reducers (for big intermediate data)

cluster of machines



SQL => MapReduce

Map Phase

SELECT, WHERE, JOIN

Shuffle Phase

ORDER BY

Reduce Phase

GROUPBY/AGGREGATE, HAVING, JOIN

MapReduce is more flexible. (for example, how to do a GROUP BY where one row goes to mutliple groups in SQL?)

Projects like **HiveQL** try to make MapReduce more accessible.

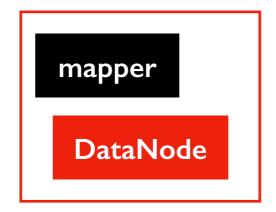
Data Locality: Avoid Network Transfers

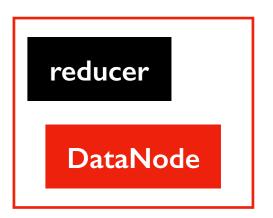
Run on same machines

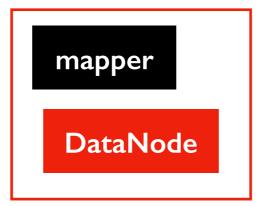
- HDFS DataNodes
- MapReduce executor

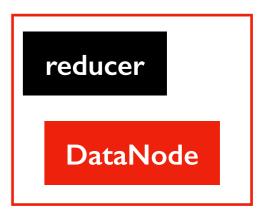
Try to run mappers on machine where DataNode has needed data. Uses disk but not network.

cluster of machines









Summary: Some Key Ideas

To build complex systems...

compose layers of subsystems

To scale out...

partition your data

To handle faults...

replicate your data

To optimize I/O...

- pipeline writes
- co-locate computation (MapReduce) with data (HDFS)