

# **Memory Topics in MIPS: Addressing, Global Vars & Arrays**

**CS 64: Computer Organization and Design Logic  
Lecture #7  
Winter 2019**

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# *This Week on “Didja Know Dat?!”*



Steve Wozniak and Steve Job's first commercial venture was the **Apple 1** in 1976 using an **8-bit MOS 6502 CPU**. It was built for \$500 and initially **sold for \$666.66** because Wozniak “*liked repeating digits*” (about \$2900 in today’s dollars). Keyboard and TV not included. They sold about 200 of them in 10 months, thus assuring the continuation of their company.

Previously, the only other popular “personal” computer was the Altair 8800, which you had to operate with switches!



# Administrative

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- Lab 4 starts Thursday (Due Monday)
- Midterm Exam on Feb. 5<sup>th</sup> (Next Week Tue.)

# What's on the Midterm?

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## What's on It?

- Everything we've done so far from start to end of this week

## What Should I Bring?

- Your pencil(s), eraser, MIPS Ref. Card
- THAT'S ALL!

## What Else Should I Do?

- Come to the classroom 5-10 minutes EARLY
- I will have some of you re-seated
- Bring your UCSB ID

# Lecture Outline

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- Addressing MIPS Memory
- Global Variables
- Arrays

# Any Questions From Last Lecture?

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# Pop Quiz!

- You have 5 minutes to fill in the missing code. You can use your MIPS Reference Card.
- Fill in the 4 blank spaces :

**main:** # assume \$t0 has been declared earlier (not here)

li \$t1, 0

li \_\_\_\_\_

blt \_\_\_\_\_

li \$t1, 1

**exit:** \_\_\_\_\_

\_\_\_\_\_

**In C++, the code would be:**

```
if (t0 >= 5)
    t1 = 1;
else
    t1 = 0;
```

# Pop Quiz Answers!

- You have 5 minutes to fill in the missing code. You can use your MIPS Reference Card.
- Fill in the 4 blank spaces :

**main:** # assume \$t0 has been declared earlier (not here)

li \$t1, 0

li \$t2, 5 # something to compare!

blt \$t0, \$t2, exit

li \$t1, 1;

**exit:** li \$v0, 10

syscall

**In C++, the code would be:**

```
if (t0 >= 5)
    t1 = 1;
else
    t1 = 0;
```



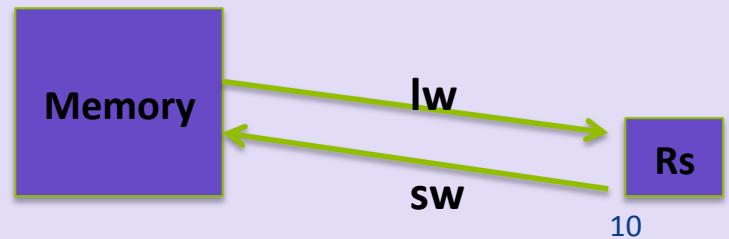
## Example 4

```
.data
num1: .word 42      # define 32b w/ value = 42
num2: .word 7       # define 32b w/ value = 7
num3: .space 1      # define one (1) 32b space

.text
main:
    lw $t0, num1      # load what's in num1 (42) into $t0
    lw $t1, num2      # load what's in num2 (7) into $t1
    add $t2, $t0, $t1 # ($t0 + $t1) → $t2
    sw $t2, num3      # load what's in $t2 (49) into num3 space

    li $v0, 1
    lw $a0, num3      # put the number you want to print in $a0
    syscall            # print integer

    li $v0, 10          # exit
    syscall
```



# Addressing Memory

- If you're not using the `.data` declarations, then you need *starting addresses* of the data in memory with `lw` and `sw` instructions

Example:    `lw $t0, 0x0000400A`    ← not a real address, just looks like one...

Example:    `lw $t0, 16($s0)`

- 1 word = 32 bits (in MIPS)
  - So, in a 32-bit unit of memory, that's 4 bytes
  - Represented with 8 hexadecimals                 $8 \times 4 \text{ bits} = 32 \text{ bits}$ ... checks out...
- MIPS addresses sequential memory addresses, but not in "words"
  - Addresses are in Bytes instead
  - MIPS words *must* start at addresses that are multiples of 4
  - Called an ***alignment restriction***

# Memory Allocation Map

## MEMORY ALLOCATION

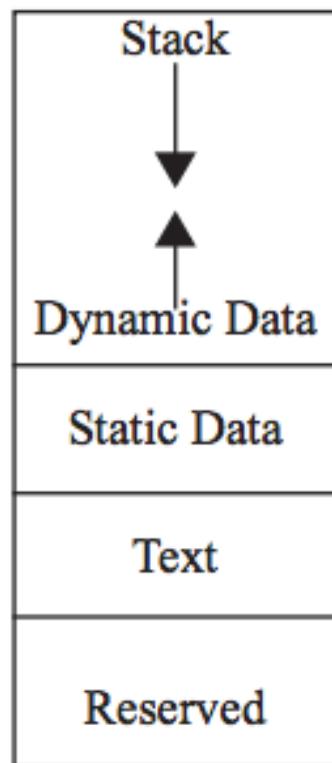
\$sp → 7fff fffc<sub>hex</sub>

\$gp → 1000 8000<sub>hex</sub>

1000 0000<sub>hex</sub>

pc → 0040 0000<sub>hex</sub>

0<sub>hex</sub>



*This is found on your  
MIPS Reference Card*

*How much memory does a  
programmer get to directly  
use in MIPS?*

### NOTE:

Not all memory addresses can be accessed by the programmer.

Although the address space is 32 bits, the top addresses from **0x80000000** to **0xFFFFFFFF** are not available to user programs. They are used mostly by the OS.

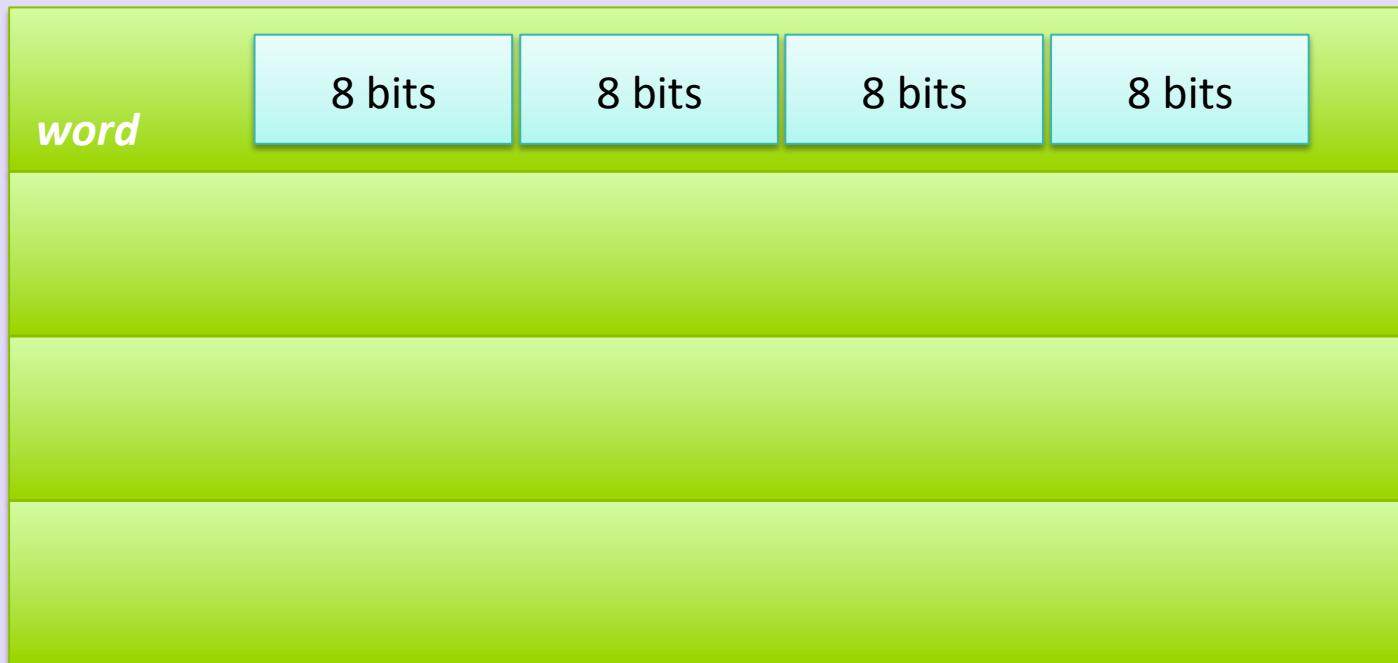
# Mapping MIPS Memory

(say that 10 times fast!)

- Imagine computer memory like a big array of words
- Size of computer memory is:

$$2^{32} = 4 \text{ Gbits, or } 512 \text{ MBytes (MB)}$$

- We only get to use 2 Gbits, or 256 MB
- That's  $(256 \text{ MB}) / (\text{groups of } 4 \text{ B}) = 64 \text{ million words}$



# MIPS Computer Memory Addressing Conventions

A  
→

1A	80	C5	29
0x0000	0x0001	0x0002	0x0003
52	00	37	EE
0x0004	0x0005	0x0006	0x0007
B1	11	1A	A5
0x0008	0x0009	0x000A	0x000B

# MIPS Computer Memory Addressing Conventions

*or...*

B  
←

1A	80	C5	29
0x0003	0x0002	0x0001	0x0000
52	00	37	EE
0x0007	0x0006	0x0005	0x0004
B1	11	1A	A5
0x000B	0x000A	0x0009	0x0008

# A Tale of 2 Conventions...

**BIG END (MSByte)**  
gets addressed first

1A	80	C5	29
0x0000	0x0001	0x0002	0x0003
52	00	37	EE
0x0004	0x0005	0x0006	0x0007
B1	11	1A	A5
0x0008	0x0009	0x000A	1A
80	C5	29	

← BIG ENDIAN

**LITTLE END (LSByte)**  
gets addressed first

1A	80	C5	29
0x0003	0x0002	0x0001	0x0000
52	00	37	EE
0x0007	0x0006	0x0005	0x0004
B1	11	1A	A5
0x000B	0x000A	0x0009	0x0008

LITTLE ENDIAN →

# The Use of Big Endian vs. Little Endian

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*Origin: Jonathan Swift (author) in “Gulliver's Travels”.  
Some people preferred to eat their hard boiled eggs from the “little end” first (thus, little endians), while others prefer to eat from the “big end” (i.e. big endians).*

- MIPS users typically go with Big Endian convention
  - MIPS allows you to program “endian-ness”
- Most Intel processors go with Little Endian...
- It's just a convention – it makes no difference to a CPU!

# Global Variables

## Recall:

- Typically, global variables are placed directly in memory, not registers
- **lw** and **sw** for **load word** and **save word**
  - **lw ≠ la ≠ move !!!**
  - Syntax:  
 $lw \text{ } register\_destination, N(register\_with\_address)$   
Where **N** = **offset** of address in bytes
- Let's take a look at: [\*\*access\\_global.asm\*\*](#)

# access\_global.asm

## Load Address (la) and Load Word (lw)

```
.data
myVariable: .word 42
.text
main:
    $t0 = &myVariable
    la $t0, myVariable      ← WHAT'S IN $t0??
    lw $t1, 0($t0)          ← WHAT DID WE DO HERE??
    li $v0, 1
    move $a0, $t1
    syscall                 ← WHAT SHOULD WE SEE HERE??
```

# access\_global.asm

**Store Word (sw) (...continuing from last page...)**

```
li $t1, 5
sw $t1, 0($t0)      ← WHAT'S IN $t0 AGAIN??
```

```
li $t1, 0
lw $t1, 0($t0)      ← WHAT DID WE DO HERE??
```

```
li $v0, 1
move $a0, $t1
syscall               ← WHAT SHOULD WE SEE HERE??
```

# Arrays

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- Question:

As far as memory is concerned, what is the *major difference* between an **array** and a **global variable**?

- Arrays contain multiple elements

- Let's take a look at:

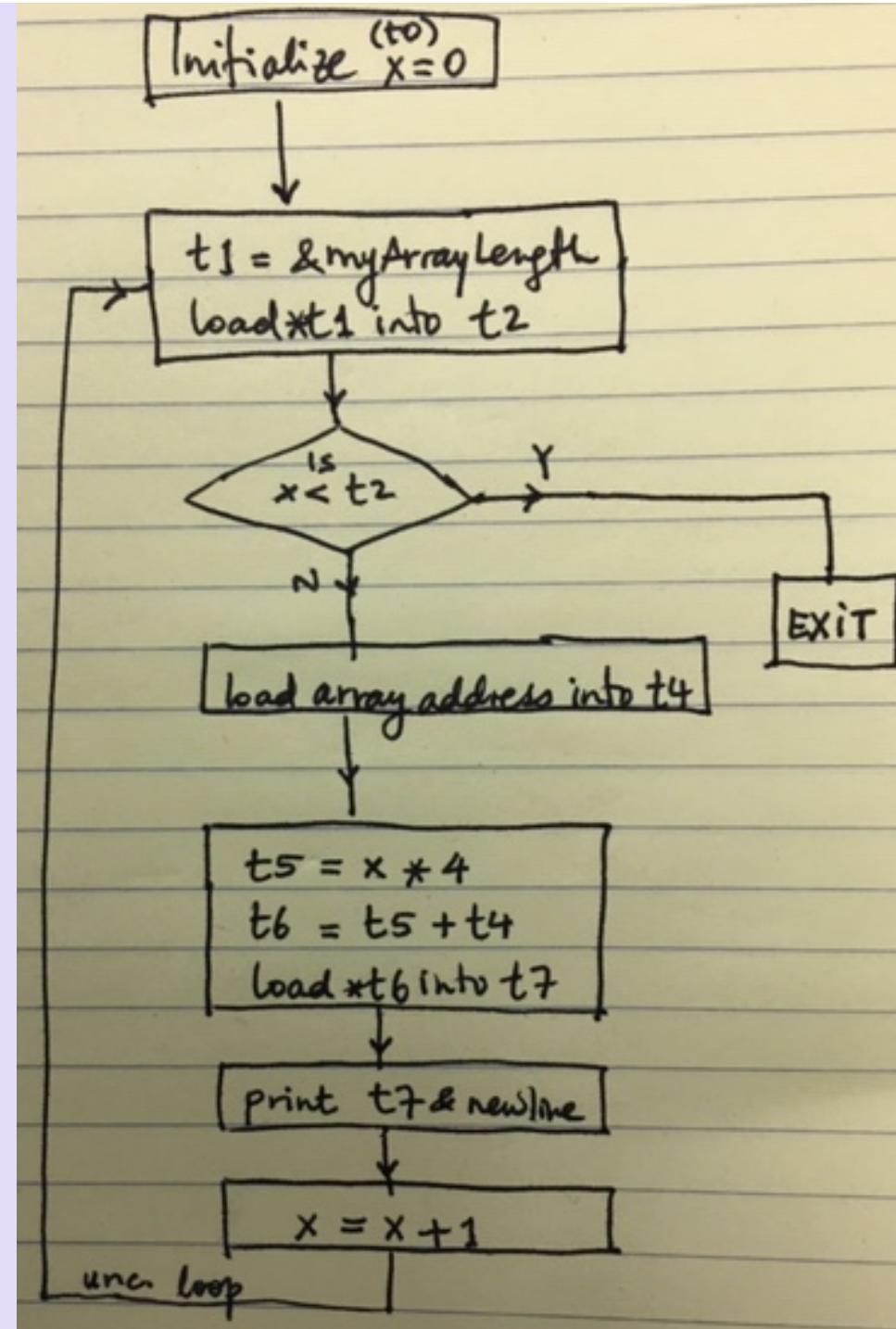
- print\_array1.asm
  - print\_array2.asm
  - print\_array3.asm

# print\_array1.asm

---

```
int myArray[]  
= {5, 32, 87, 95, 286, 386};  
int myArrayLength = 6;  
int x;  
  
for (x = 0; x < myArrayLength; x++)  
{  
    print(myArray[x]);  
    print("\n");  
}
```

# Flow Chart for print\_array1



```

# C code:
# int myArray[] =
#     {5, 32, 87, 95, 286, 386}
# int myArrayLength = 6
# for (x = 0; x < myArrayLength; x++) {
#   print(myArray[x])
#   print("\n")
.data
newline: .asciiz "\n"
myArray: .word 5 32 87 95 286 386
myArrayLength: .word 6

.text
main:
    # t0: x
    # initialize x
    li $t0, 0
loop:
    # get myArrayLength, put result in $t2
    # $t1 = &myArrayLength
    la $t1, myArrayLength
    lw $t2, 0($t1)

    # see if x < myArrayLength
    # put result in $t3
    slt $t3, $t0, $t2
    # jump out if not true
    beq $t3, $zero, end_main

```

```

# get the base of myArray
la $t4, myArray

# figure out where in the array we need
# to read from. This is going to be the array
# address + (index << 2). The shift is a
# multiplication by four to index bytes
# as opposed to words.
# Ultimately, the result is put in $t7
sll $t5, $t0, 2
add $t6, $t5, $t4
lw $t7, 0($t6)

# print it out, with a newline
li $v0, 1
move $a0, $t7
syscall
li $v0, 4
la $a0, newline
syscall

# increment index
addi $t0, $t0, 1

# restart loop
j loop

end_main:
    # exit the program
    li $v0, 10
    syscall

```

# print\_array2.asm

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- Same as print\_array1.asm, *except that* in the assembly code, we lift redundant computation out of the loop.
- This is the sort of thing a decent compiler (**clang** or **gcc** or **g++**, for example) will do with a HLL program
- Your homework: **Go through this assembly code!**

# print\_array3.asm

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```
int myArray[]  
    = {5, 32, 87, 95, 286, 386};  
int myArrayLength = 6;  
int* p;  
  
for ( p = myArray; p < myArray + myArrayLength; p++)  
{  
    print(*p);  
    print("\n");  
}
```

Your homework: Go through this assembly code!

# YOUR TO-DOs

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- Review ALL the demo codes
  - Available via the class website
- Lab 4 on Thursday!



</LECTURE>