

# **Flow Control & Memory Use in Assembly**

**CS 64: Computer Organization and Design Logic  
Lecture #6  
Winter 2019**

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Legend: Adm. Grace Hopper coined the term "debugging" when a moth was removed from the computer she was working on (see below)

Reality: The term "bug" was used in engineering in the 19<sup>th</sup> century. As seen independently from various scientists, including Ada Lovelace and Thomas Edison.

# This Week on "Didja Know Dat?!"

1300 (033) MP - MC 2.130476415 (23) 4.615925059(-2)  
(033) PRO 2 2.130476415  
conch 2.130676415  
Relays 6-2 in 033 failed special speed test  
in relay " 11.00 test .

1100 Started Cosine Tape (Sine check)  
1525 Started Multi-Adder Test.

1545

Relay #70 Panel F  
(moth) in relay.

First actual case of bug being found.

1630 Antangut started.  
1700 closed down.

Relay 2145  
Relay 3376

# Lecture Outline

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- **.data** Directives and Basic Memory Use
- Branching (Conditionals)
- Loops
- Accessing Data in Memory
- Memory Addressing

# Any Questions From Last Lecture?

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# MIPS Peculiarity: NOR used a NOT

- How to make a NOT function using **NOR** instead
- Recall: NOR = NOT OR
- Truth-Table:

A	B	A NOR B
0	0	1
0	1	0
1	0	0
1	1	0

Note that:  
 $0 \text{ NOR } x = \text{NOT } x$

- So, in the absence of a NOT function,  
use a NOR with a 0 as one of the inputs!

# **.data Declaration Types**

## *w/ Examples*

var1:	.byte 9	# declare a single byte with value 9
var2:	.half 63	# declare a 16-bit half-word w/ val. 63
<b>var3:</b>	<b>.word 9433</b>	<b># declare a 32-bit word w/ val. 9433</b>
num1:	.float 3.14	# declare 32-bit floating point number
num2:	.double 6.28	# declare 64-bit floating pointer number
str1:	.ascii "Text"	# declare a string of chars
<b>str3:</b>	<b>.asciiz "Text"</b>	<b># declare a null-terminated string</b>
<b>str2:</b>	<b>.space 5</b>	<b># reserve 5 bytes of space (useful for arrays)</b>

*These are now reserved in memory and we can call them up by loading their memory address into the appropriate registers.*

***Highlighted ones are the ones most commonly used in this class.***

**.data**

```
name: .asciiz "Jimbo Jones is "  
rtn: .asciiz " years old.\n"
```

**.text**

main:

```
    li $v0, 4  
    la $a0, name    # la = load memory address  
    syscall
```

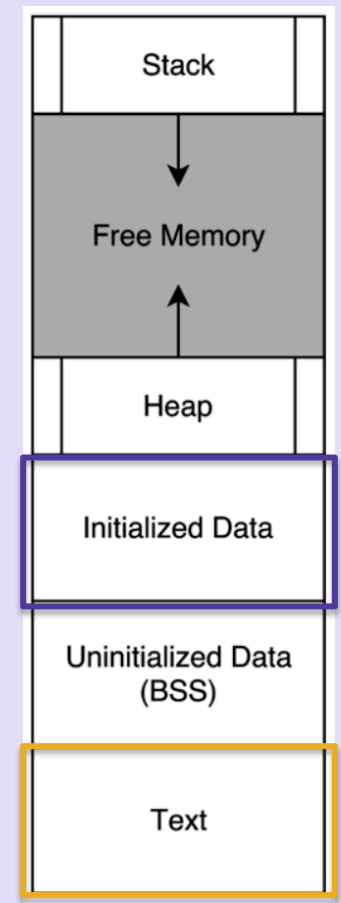
```
    li $v0, 1  
    li $a0, 15  
    syscall
```

```
    li $v0, 4  
    la $a0, rtn  
    syscall
```

```
    li $v0, 10  
    syscall
```

# Example

*What does this do?*



*What goes in here? →*

*What goes in here? →*

# Conditionals

- What if we wanted to do:

```
if (x == 0) { printf("x is zero"); }
```

- Can we write this in assembly with what we know?

- No... we haven't covered **if-else** (aka *branching*)

- What do we need to implement this?

- A way to *compare* numbers

- A way to *conditionally execute* code



# Relevant Instructions in MIPS

*for use with branching conditionals*

- Comparing numbers:
  - set-less-than (slt)**
    - Set some register (i.e. make it “1”) if a less-than comparison of some other registers is true
- Conditional execution:
  - branch-on-equal (beq)**
  - branch-on-not-equal (bne)**
    - “Go to” some other place in the code (i.e. jump)

```
if (x == 0) { printf("x is zero"); }
```

```
.data
```

```
x_is_zero: .asciiz "x is zero"
```

Create a constant string called "x\_is\_zero"

```
.text
```

```
bne $t0, $zero, after_print
```

If  $\$t0 \neq 0$  go to the block labeled as "**after\_print**"

```
li $v0, 4
```

```
la $a0, x_is_zero
```

(otherwise) prepare to print a string...

```
syscall
```

...and that string is inside of "x\_is\_zero"

```
after_print:
```

```
li $v0, 10
```

```
syscall
```

End the program

Note the flow

# Loops

- How might we translate the following C++ to assembly?

```
n = 3;
sum = 0;
while (n != 0)
{
    sum += n;
    n--;
}
cout << sum;
```

*n = 3; sum = 0;*  
*while (n != 0) { sum += n; n--; }*

`.text`

`main:`

`li $t0, 3    # n`  
`li $t1, 0    # running sum`

`loop:`

`beq $t0, $zero, loop_exit`  
`addu $t1, $t1, $t0`  
`addi $t0, $t0, -1`  
`j loop`

`loop_exit:`

`li $v0, 1`  
`move $a0, $t1`  
`syscall`

`li $v0, 10`  
`syscall`

Set up the variables in \$t0, \$t1

If \$t0 == 0 go to "loop\_exit"

(otherwise) make \$t1 the (unsigned) sum of \$t1 and \$t0 (i.e. **sum += n**)

decrement \$t0 (i.e. **n--**)

jump to the code labeled "loop"  
(i.e. **repeat loop**)

prepare to print out an integer,  
which is inside the \$t1 reg. (i.e. **print sum**)

end the program

# Let's Run More Programs!!

## Using SPIM

- More!!
- This time exploring conditional logic and loops



These assembly code programs are made available to you via the class webpage

# More Branching Examples

```
int y;
if (x == 5)
{
    y = 8;
}
else if (x < 7)
{
    y = x + x;
}
else
{
    y = -1;
}
print(y)
```

```
.text
main:    # t0: x and t1: y
        li $t0, 5      # example
        li $t2, 5      # what's this?
        beq $t0, $t2, equal_5

        # check if less than 7
        li $t2, 7
        slt $t3, $t0, $t2
        bne $t3, $zero, less_than_7

        # fall through to final else
        li $t1, -1
        j after_branches

equal_5:
        li $t1, 8
        j after_branches
```

```
less_than_7:
        add $t1, $t0, $t0
        # could jump to after_branches,
        # but this is what we will fall
        # through to anyways

after_branches:
# print out the value in y ($t1)
        li $v0, 1
        move $a0, $t1
        syscall

        # exit the program
        li $v0, 10
        syscall
```

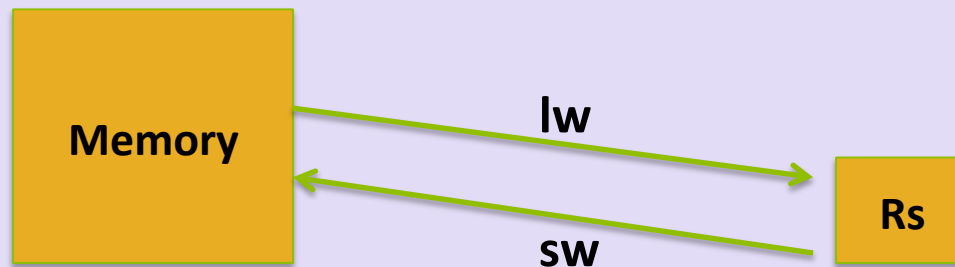
# Larger Data Structures

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- Recall: registers vs. memory
  - Where would data structures, arrays, etc. go?
  - Which is faster to access? Why?
- Some data structures have to be stored in memory
  - So we need instructions that “shuttle” data to/from the CPU and computer memory (RAM)

# Accessing Memory

- Two base instructions:
  - load-word (**lw**) from memory to registers
  - store-word (**sw**) from registers to memory



- MIPS lacks instructions that do more with memory than access it (e.g., retrieve something from memory and then add)
  - Operations are done step-by-step
  - Mark of RISC architecture



**.data**

num1: .word 42

num2: .word 7

num3: .space 1

**.text**

main:

lw \$t0, num1

lw \$t1, num2

add \$t2, \$t0, \$t1

sw \$t2, num3

li \$v0, 1

lw \$a0, num3

syscall

li \$v0, 10

syscall

## Example 4

*What does this do?*



## Example 4

**.data**

```
num1: .word 42    # define 32b w/ value = 42
num2: .word 7     # define 32b w/ value = 7
num3: .space 1    # define one (1) 32b space
```

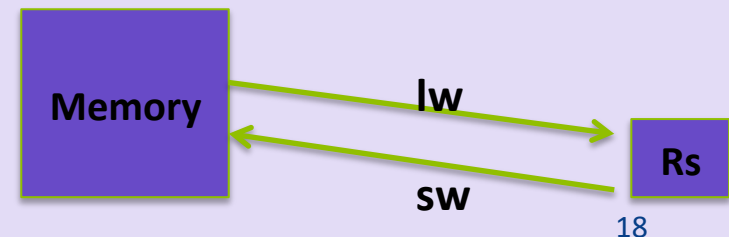
**.text**

**main:**

```
    lw $t0, num1      # load what's in num1 (42) into $t0
    lw $t1, num2      # load what's in num2 (7) into $t1
    add $t2, $t0, $t1 # ($t0 + $t1) → $t2
    sw $t2, num3      # load what's in $t2 (49) into num3 space
```

```
    li $v0, 1
    lw $a0, num3      # put the number you want to print in $a0
    syscall           # print integer
```

```
    li $v0, 10        # exit
    syscall
```



# YOUR TO-DOs

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- Review ALL the demo codes
  - Available via the class website
- Assignment #3
  - Due Monday!

**</LECTURE>**