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# 377 Operating Systems

Introduction to Concurrency





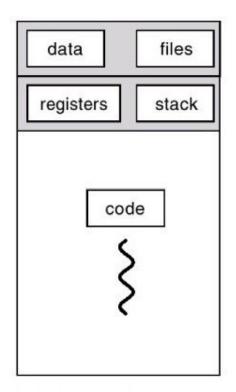
Virtualization: Processes

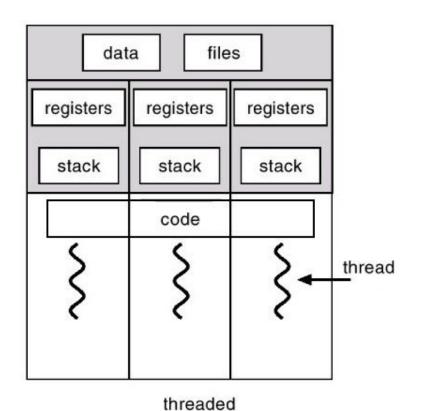
Reality: Single CPU

Virtualization: Multiple Processes



#### **Another Abstraction: Threads**

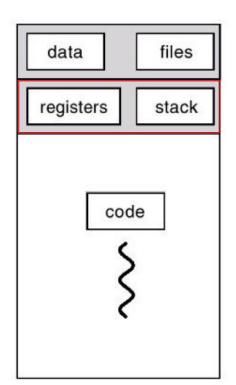


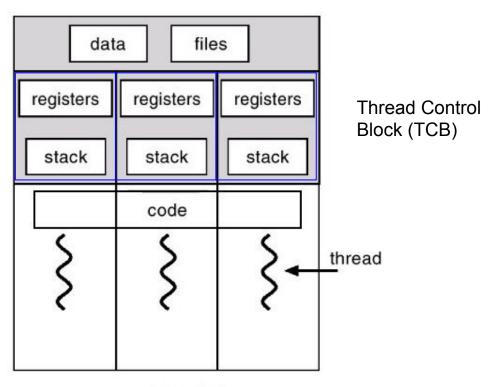


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#### **Another Abstraction: Threads**

Process Control Block (PCB)





## Single-Threaded and Multi-Threaded Address Space

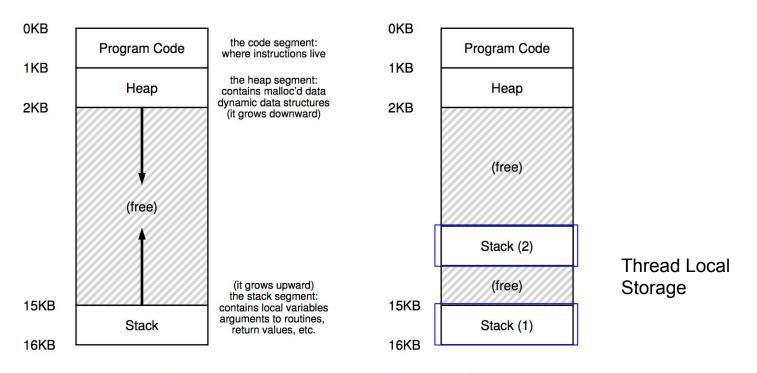


Figure 26.1: Single-Threaded And Multi-Threaded Address Spaces

The first one is simple: parallelism

Imagine you are writing a program that performs operations on very large arrays, for example, adding two large arrays together, or incrementing the value of each element in the array by some amount.

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If you are running on just a single processor, the task is straightforward: just perform each operation and be done.

The first one is simple: parallelism

Imagine you are writing a program that performs operations on very large arrays, for example, adding two large arrays together, or incrementing the value of each element in the array by some amount.

If you are running on just a single processor, the task is straightforward: just perform each operation and be done.

However, if you are executing the program on a system with multiple processors, you have the potential of speeding up this process considerably by using the processors to each perform a portion of the work.

The second reason is more subtle:

to avoid blocking program progress due to slow I/O

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Imagine that you are writing a program that performs different types of I/O: either waiting to send or receive a message, for an explicit disk I/O to complete, or even (implicitly) for a trap/fault to finish.

The second reason is more subtle:

to avoid blocking program progress due to slow I/O

Imagine that you are writing a program that performs different types of I/O: either waiting to send or receive a message, for an explicit disk I/O to complete, or even (implicitly) for a trap/fault to finish.

Instead of waiting, your program may wish to do something else, including utilizing the CPU to perform computation, or even issuing further I/O requests. Using threads is a natural way to avoid getting stuck.

The second reason is more subtle:

to avoid blocking program progress due to slow I/O

while one thread in your program waits (i.e., is blocked waiting for I/O), the CPU scheduler can switch to other threads, which are ready to run and do something useful.

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Threading enables overlap of I/O with other activities within a single program, much like multiprogramming did for processes across programs.

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while one thread in your program waits (i.e., is blocked waiting for I/O), the CPU scheduler can switch to other threads, which are ready to run and do something useful.

Threading enables overlap of I/O with other activities within a single program, much like multiprogramming did for processes across programs.

Examples: many modern server-based applications (web servers, database management systems, and the like) make use of threads in their implementations.

#### Why Not Use Processes?

Of course, in either of the cases mentioned above, you could use multiple processes instead of threads.

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However, threads **share an address space** and thus make it easy to share data.

### Why Not Use Processes?

Of course, in either of the cases mentioned above, you could use multiple processes instead of threads.

However, threads **share an address space** and thus make it easy to share data.

Thus, threads are a natural choice when constructing these types of programs.

Processes are a more sound choice for logically separate tasks where little sharing of data structures in memory is needed.

```
#include <stdio.h>
    #include <assert.h>
    #include <pthread.h>
4
    void *mythread(void *arg) {
        printf("%s\n", (char *) arg);
        return NULL;
8
9
    int.
10
    main(int argc, char *argv[]) {
11
        pthread_t p1, p2;
12
        int rc:
13
       printf("main: begin\n");
14
        rc = pthread_create(&p1, NULL, mythread, "A"); assert(rc == 0);
15
        rc = pthread_create(&p2, NULL, mythread, "B"); assert(rc == 0);
16
        // join waits for the threads to finish
17
        rc = pthread_join(p1, NULL); assert(rc == 0);
18
        rc = pthread_join(p2, NULL); assert(rc == 0);
19
        printf("main: end\n");
20
        return 0;
21
22
```

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        rc = pthread_join(p2, NULL); assert(rc == 0);
19
        printf("main: end\n");
20
        return 0;
21
22
```

#### Question

#### Which prints first A or B?

```
void *mythread(void *arg) {
    printf("%s\n", (char *) arg);
    return NULL;
int
main(int argc, char *argv[]) {
    pthread_t p1, p2;
    int rc;
    printf("main: begin\n");
    rc = pthread_create(&p1, NULL, mythread, "A"); assert(rc == 0);
    rc = pthread_create(&p2, NULL, mythread, "B"); assert(rc == 0);
    // join waits for the threads to finish
    rc = pthread_join(pl, NULL); assert(rc == 0);
    rc = pthread_join(p2, NULL); assert(rc == 0);
    printf("main: end\n");
    return 0;
```

main	Thread 1	Thread2
starts running		
prints "main: begin"		
creates Thread 1		
creates Thread 2		
waits for T1		
	runs	
	prints "A"	
	returns	
waits for T2		
		runs
		prints "B"
		returns
prints "main: end"		

main	Thread 1	Thread2
starts running		
prints "main: begin"		
creates Thread 1		
creates Thread 2		
waits for T1		
	runs	
	prints "A"	
	returns	
waits for T2		
		runs
		prints "B"
		returns
prints "main: end"		

main	Thread 1	Thread2
starts running		
prints "main: begin"		
creates Thread 1		
creates Thread 2		
waits for T1		
	runs	
	prints "A"	
	returns	
waits for T2		
		runs
		prints "B"
		returns
prints "main: end"		

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main	Thread 1	Thread2
starts running		
prints "main: begin"		
creates Thread 1		
creates Thread 2		
waits for T1		
and to for TO	runs prints "A" returns	
waits for T2		
prints "main, and"		runs prints "B" returns

main	Thread 1	Thread2
starts running		
prints "main: begin"		
creates Thread 1		
creates Thread 2		
waits for T1		
	runs	
	prints "A"	
	returns	
waits for T2		
		runs
		prints "B"
		returns
prints "main: end"		

main	Thread 1	Thread2
starts running		
prints "main: begin"		
creates Thread 1		
creates Thread 2		
waits for T1		
	runs	
	prints "A"	
	returns	
waits for T2		
		runs
		prints "B"
		returns
prints "main: end"		

main	Thread 1	Thread2
starts running		
prints "main: begin"		
creates Thread 1		
creates Thread 2		
waits for T1		
	runs	
	prints "A"	
	returns	
waits for T2		
		runs
		prints "B"
		returns
prints "main: end"		

main	Thread 1	Thread2
starts running		
prints "main: begin"		
creates Thread 1		
creates Thread 2		
waits for T1		
	runs	
	prints "A"	
	returns	
waits for T2		
		runs
		prints "B"
		returns
prints "main: end"		

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main	Thread 1	Thread2
starts running		
prints "main: begin"		
creates Thread 1		
creates Thread 2		
waits for T1		
	runs	
	prints "A"	
	returns	
waits for T2		
		runs
		prints "B"
		returns
prints "main: end"		

main	Thread 1	Thread2
starts running		
prints "main: begin"		
creates Thread 1		
	runs prints "A" returns	
creates Thread 2		
		runs prints "B" returns
waits for T1 returns immediately; T1 is done		
waits for T2  returns immediately; T2 is done prints "main: end"		

main	Thread 1	Thread2
starts running		
prints "main: begin"		
creates Thread 1		
	runs prints "A" returns	
creates Thread 2		
		runs prints "B" returns
waits for T1 returns immediately; T1 is done		
waits for T2  returns immediately; T2 is done prints "main: end"		

main	Thread 1	Thread2
starts running		
prints "main: begin"		
creates Thread 1		
	runs	
	prints "A"	
	returns	
creates Thread 2		
		runs
		prints "B"
		returns
waits for T1		
returns immediately; T1 is done		
waits for T2		
returns immediately; T2 is done		
prints "main: and"		

main	Thread 1	Thread2
starts running		=
prints "main: begin"		
creates Thread 1		
	runs	
	prints "A"	
	returns	
creates Thread 2		
		runs
		prints "B"
		returns
waits for T1		
returns immediately; T1 is done		
waits for T2		
returns immediately; T2 is done		
prints "main: end"		

main	Thread 1	Thread2
starts running		-
prints "main: begin"		
creates Thread 1		
	runs	
	prints "A"	
	returns	
creates Thread 2		
		runs
		prints "B"
		returns
waits for T1		
returns immediately; T1 is done		
waits for T2		
returns immediately; T2 is done		
prints "main: end"		

main	Thread 1	Thread2
starts running		
prints "main: begin"		
creates Thread 1		
	runs	
	prints "A"	
	returns	
creates Thread 2		
		runs
		prints "B"
		returns
waits for T1		
returns immediately; T1 is done		
waits for T2		
returns immediately; T2 is done		
prints "main: end"		

main	Thread 1	Thread2
starts running		
prints "main: begin"		
creates Thread 1		
	runs	
	prints "A"	
	returns	
creates Thread 2		
		runs
		prints "B"
		returns
waits for T1		TOTALLO
returns immediately; T1 is done		
waits for T2		
returns immediately; T2 is done		
prints "main: end"		

main	Thread 1	Thread2
starts running		
prints "main: begin"		
creates Thread 1		
	runs prints "A" returns	
creates Thread 2		
		runs prints "B" returns
waits for T1 returns immediately; T1 is done		
waits for T2 returns immediately; T2 is done		
muinta "maine and"		

main	Thread 1	Thread2
starts running		-
prints "main: begin" creates Thread 1		
	runs prints "A" returns	
creates Thread 2		
		runs prints "B" returns
waits for T1		
returns immediately; T1 is done		
waits for T2		
returns immediately; T2 is done		
prints "main: end"		

main	Thread 1	Thread2
starts running		
prints "main: begin"		
creates Thread 1		
creates Thread 2		
waits for T1		runs prints "B" returns
	runs	
	prints "A" returns	
waits for T2  returns immediately; T2 is done prints "main: end"	retarrio	

main	Thread 1	Thread2
starts running		
prints "main: begin"		
creates Thread 1		
creates Thread 2		
waits for T1		runs prints "B" returns
	runs prints "A" returns	
waits for T2  returns immediately; T2 is done prints "main: end"		

main	Thread 1	Thread2
starts running		
prints "main: begin"		
creates Thread 1		
creates Thread 2		
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main	Thread 1	Thread2
starts running		
prints "main: begin"		
creates Thread 1		
creates Thread 2		
		runs prints "B" returns
waits for T1		
	runs	
	prints "A"	
	returns	
waits for T2  returns immediately; T2 is done prints "main: end"		

main	Thread 1	Thread2
starts running		
prints "main: begin"		
creates Thread 1		
creates Thread 2		
		runs
		prints "B"
		returns
waits for T1		
	runs	
	prints "A"	
	returns	
waits for T2  returns immediately; T2 is done prints "main: end"		

main	Thread 1	Thread2
starts running		
prints "main: begin"		
creates Thread 1		
creates Thread 2		
		runs prints "B" returns
waits for T1		
	runs	
	prints "A"	
	returns	
waits for T2  returns immediately; T2 is done prints "main: end"		

returns immediately; T2 is done

prints "main: end"

main	Thread 1	Thread2
starts running		
prints "main: begin"		
creates Thread 1		
creates Thread 2		
		runs prints "B" returns
waits for T1		
	runs	
	prints "A"	
	returns	
waits for T2		

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main	Thread 1	Thread2
starts running		
prints "main: begin"		
creates Thread 1		
creates Thread 2		
		runs
		prints "B"
		returns
waits for T1		
	runs	
	prints "A"	
	returns	
waits for T2		
returns immediately; T2 is done		
prints "main: and"		

main	Thread 1	Thread2
starts running		
prints "main: begin"		
creates Thread 1		
creates Thread 2		
waits for T1	runs prints "A" returns	runs prints "B" returns
waits for T2		
returns immediately; T2 is done		
prints "main: end"		

```
static volatile int counter = 0;
void *
mythread(void *arg)
{
    printf("%s: begin\n", (char *) arg);
    int i;
    for (i = 0; i < 1e7; i++) {
        counter = counter + 1;
    }
    printf("%s: done\n", (char *) arg);
    return NULL;
}</pre>
```

Threads are useful, but there are problems.

```
int
main(int argc, char *argv[])
{
    pthread_t p1, p2;
    printf("main: begin (counter = %d)\n", counter);
    Pthread_create(&p1, NULL, mythread, "A");
    Pthread_create(&p2, NULL, mythread, "B");

    // join waits for the threads to finish
    Pthread_join(p1, NULL);
    Pthread_join(p2, NULL);
    printf("main: done with both (counter = %d)\n", counter);
    return 0;
}
```

```
static volatile int counter = 0;
void *
mythread(void *arg)
   printf("%s: begin\n", (char *) arg);
    int i:
    for (i = 0; i < 1e7; i++) {
        counter = counter + 1;
   printf("%s: done\n", (char *) arg);
   return NULL;
int
main(int argc, char *argv[])
   pthread_t p1, p2;
```

Pthread create (&pl, NULL, mythread, "A"); Pthread create (&p2, NULL, mythread, "B");

// join waits for the threads to finish

printf("main: done with both (counter = %d)\n", counter);

Pthread\_join(p1, NULL); Pthread join(p2, NULL);

return 0;

# Threads are useful, but there are problems.

```
prompt> gcc -o main main.c -Wall -pthread
                                    prompt> ./main
                                   main: begin (counter = 0)
                                   A: begin
                                    B: begin
                                    A: done
                                    B: done
                                   main: done with both (counter = 20000000)
printf("main: begin (counter = %d) \n", counter);
```

```
static volatile int counter = 0;
void *
mythread(void *arg)
   printf("%s: begin\n", (char *) arg);
   int i:
   for (i = 0; i < 1e7; i++) {
       counter = counter + 1;
   printf("%s: done\n", (char *) arg);
```

Pthread create (&pl, NULL, mythread, "A"); Pthread create (&p2, NULL, mythread, "B");

// join waits for the threads to finish

return NULL;

return 0;

main(int argc, char \*argv[])

Pthread\_join(p1, NULL); Pthread join(p2, NULL);

pthread\_t p1, p2;

int

# Threads are useful, but there are problems.

```
B: begin
                                             A: done
                                             B: done
                                             main: done with both (counter = 19345221)
printf("main: begin (counter = %d) \n", counter);
printf("main: done with both (counter = %d)\n", counter);
```

A: begin

prompt> ./main

main: begin (counter = 0)

```
53
```

```
static volatile int counter = 0;
void *
mythread(void *arg)
   printf("%s: begin\n", (char *) arg);
    int i:
    for (i = 0; i < 1e7; i++) {
       counter = counter + 1;
   printf("%s: done\n", (char *) arg);
   return NULL;
int
main(int argc, char *argv[])
```

Pthread create (&pl, NULL, mythread, "A"); Pthread create (&p2, NULL, mythread, "B");

// join waits for the threads to finish

printf("main: done with both (counter = %d)\n", counter);

pthread\_t p1, p2;

return 0;

Pthread\_join(p1, NULL); Pthread join(p2, NULL);

# Threads are useful, but there are problems.

```
prompt> ./main
                                         main: begin (counter = 0)
                                         A: begin
                                         B: begin
                                         A: done
                                         B: done
                                         main: done with both (counter = 19221041
printf("main: begin (counter = %d) \n", counter);
```

```
static volatile int counter = 0;
void *
mythread(void *arg)
                                                 nroblome
   printf("%s: be Aren't computers supposed to produce
   int i:
   for (i = 0; i deterministic results, as you have been
                                                                    = 0)
       counter =
                  taught?!
   printf("%s: do:
   return NULL;
                  Not only is each run wrong, but also
int
                  yields a different result!
main(int argc, char
   pthread_t p1, p
   printf("main: be
                  A big question remains:
   Pthread create
                      why does this happen?
   Pthread create (
   // join waits for the threads to finish
   Pthread_join(p1, NULL);
   Pthread join(p2, NULL);
   printf("main: done with both (counter = %d)\n", counter);
   return 0;
```

Threads are useful, but there are

counter = 19221041

To understand why this happens, we must understand the code sequence that the compiler generates for the update to counter. In this case, we wish to simply add a number (1) to counter.

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Thus, the code sequence for doing so might look something like this (in x86):

```
mov 0x8049a1c, %eax add $0x1, %eax mov %eax, 0x8049a1c
```

To understand why this happens, we must understand the code sequence that the compiler generates for the update to counter. In this case, we wish to simply add a number (1) to counter.

Thus, the code sequence for doing so might look something like this (in x86):

```
mov 0x8049a1c, %eax add $0x1, %eax mov %eax, 0x8049a1c
```

counter

			(afte	er instr	uction)
OS	Thread 1	Thread 2	PC	%eax	counter
	before critical sect	ion	100	0	50
	mov 0x8049a1c,	%eax	105	50	50
	add \$0x1, %eax		108	51	50
interrupt					
save T1's state					
restore T2's sta	te		100	0	50
		mov 0x8049a1c, %eax	105	50	50
		add \$0x1, %eax	108	51	50
		mov %eax, 0x8049a1c	113	51	51
interrupt		1701 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1			
save T2's state					
restore T1's sta	te		108	51	51
	mov %eax, 0x8049a1c				

			(after instruction)		action)
OS	Thread 1	Thread 2	PC	%eax	counter
. <del></del>	before critical sec	ction	100	0	50
	mov 0x8049a1c	c, %eax	105	50	50
	add \$0x1, %eax		108	51	50
interrupt					
save T1's state					
restore T2's sta	te		100	0	50
		mov 0x8049a1c, %eax	105	50	50
		add \$0x1, %eax	108	<b>51</b>	50
		mov %eax, 0x8049a1c	113	51	51
interrupt					
save T2's state					
restore T1's sta	te		108	51	51
mov %eax, 0x8049a1c			113	51	51

	(after instruction)				
OS	Thread 1	Thread 2	PC	%eax	counter
8 <del>.</del>	before critical sec	ction	100	0	50
	mov 0x8049a1c	c, %eax	105	50	50
	add \$0x1, %eax	x	108	<b>51</b>	50
interrupt					
save T1's state	N Comment				
restore T2's sta	ate		100	0	50
		mov 0x8049a1c, %eax	105	50	50
		add \$0x1, %eax	108	<b>51</b>	50
	mov %eax, 0x8049a1c			51	51
interrupt					
save T2's state					
restore T1's sta	ate		108	51	51
mov %eax, 0x8049a1c			113	51	51

			(after instruction)		
OS	Thread 1	Thread 2	PC	%eax	counter
	before critical sec	tion	100	0	50
	mov 0x8049a1c, %eax			50	50
	add \$0x1, %eax		108	51	50
interrupt					
save T1's state					
restore T2's sta	ite		100	0	50
		mov 0x8049a1c, %eax	105	50	50
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	mov %eax, 0x8049a1c				51

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<i></i>	before critical section			0	50
	mov 0x8049a1c, %eax		105	50	50
	add \$0x1, %eax	(	108	<b>51</b>	50
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1000	mov %eax, 0x8049a1c			51	51

#### The Wish for Atomicity

The problem is that we must execute 3 instructions to perform the increment:

```
mov 0x8049a1c, %eax add $0x1, %eax mov %eax, 0x8049a1c
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```
mov 0x8049a1c, %eax add $0x1, %eax mov %eax, 0x8049a1c
```

What if we could do this with only a single <u>atomic</u> instruction:

```
memory-add 0x8049a1c, $0x1
```

#### The Wish for Atomicity

Thus, what we will instead do is ask the hardware for a few useful instructions upon which we can build a general set of what we call synchronization primitives.

By using these hardware synchronization primitives, in combination with some help from the operating system, we will be able to build multi-threaded code that accesses <u>critical sections</u> in a synchronized and controlled manner.

#### THE CRUX:

HOW TO PROVIDE SUPPORT FOR SYNCHRONIZATION

What support do we need from the hardware in order to build useful synchronization primitives? What support do we need from the OS? How can we build these primitives correctly and efficiently? How can programs use them to get the desired results?

#### One more problem: waiting for another

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accessing shared variables and the need to support atomicity for critical sections.

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One thread must wait for another to complete some action before it continues.

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We will be looking at how <u>synchronization primitives</u> and <u>conditional variables</u> help tame **uncontrolled scheduling**.

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#### Why study this in an OS class?

Before wrapping up, one question that you might have is: why are we studying this in OS class?

"History" is the one-word answer!

The OS was the **first concurrent program**, and many techniques were created for use within the OS. Later, with multi-threaded processes, application programmers also had to consider such things.

# Ehe End