

Traveling salesman

Saturday, October 28, 2023

6:31 AM

Given a complete graph with weighted edges. What is the min cost Hamiltonian cycle.

Path that visits every node once ↗

It is an NP-complete problem.

Brute force $O(n!)$ Dynamic programming $O(2^n \cdot n^2)$ Selling on ebay: $O(1)$

Dynamic programming

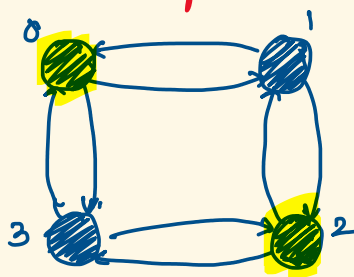
- for computing paths of len = N use data for N-
 - Store the optimal direct path from a node to every other node.
- ↗ optimization.

State:

1. The nodes visited so far.
2. The index of the last visited node.

In total, there are $2^N \cdot N$ states.
space complexity.

We use a single 32-bit integer & use its bits to represent selected nodes



State
(0101, 2)
(4) ↗ last visited

Code:

```
1 def TSP_recursive(prev, mask):
2     if (prev, mask) in dp: return dp[(prev, mask)]
3     if mask == (1 << n) - 1:
4         return 0 ## all nodes visited
5
6     cost = inf
7     for i in range(n):
8         if mask & (1 << i): continue
9         cost = min(cost, dist[prev][i] + TSP(i, mask | (1 << i)))
10    dp[(prev, mask)] = cost # memoize
11    return cost
12
13 def TSP_iterative():
14     dp = [[inf] * (1 << n) for _ in range(n)]
15
16     for i in range(n): dp[i][(1 << i)] = 0
17
18     for mask in range(1, 1 << n):
19         for i in range(n):
20             if mask & (1 << i):
21                 for j in range(n):
22                     if i != j and mask & (1 << j): continue
23                     dp[j][mask | (1 << j)] = min(
24                         dp[j][mask | (1 << j)],
25                         dp[i][mask] + dist[i][j])
26
27
28     return min([
29         x[(1 << n) - 1] for x in dp[i]
30     ], default = inf)
```

Notice that we use a push dp in the iterative code. This is because, for doing bottom up there is no way for us to know in advance, what the best cost is for mask $(1 << n) - 1$.

Return the min cost for mask that includes all nodes ending at one of the nodes.