



## COMPUTER NETWORKS

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## Computer Networks and the Internet

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Unlike other delays ( $d_{proc}$ ,  $d_{trans}$ ,  $d_{prop}$ ),  $d_{queue}$  is interesting.

- Can vary from packet to packet.
- Characterize  $d_{queue}$  -> average, variance, probability that it exceeds some specified value.

**When is the queuing delay large and when is it insignificant?**

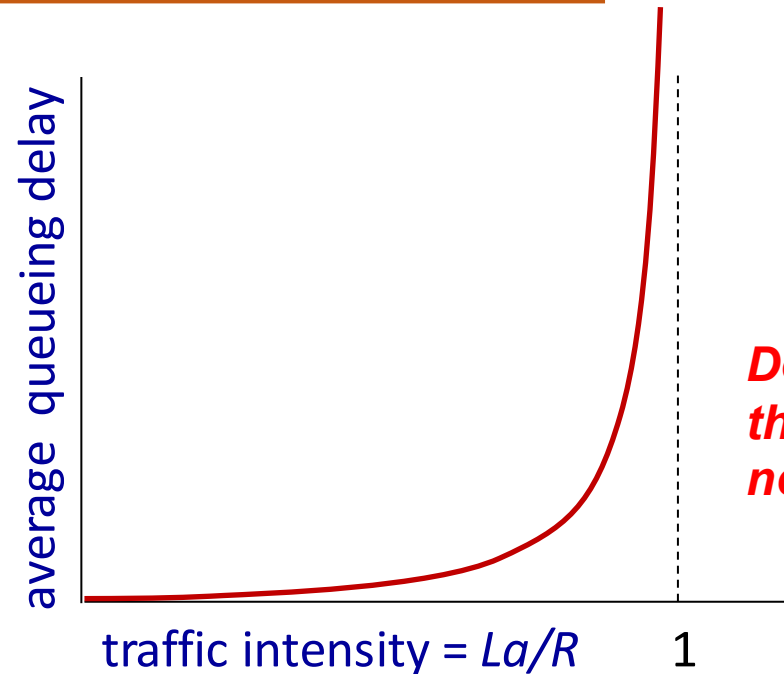
- Rate at which traffic arrives at the queue,
- Transmission rate of the link,
- Nature of the arriving traffic – periodically or in bursts

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## Performance: Packet Queueing Delay revisited

- $R$ : link bandwidth (bps)
- $L$ : packet length (bits)
- $a$ : average packet arrival rate (pps)
- $La$ : avg. rate at which bits arrive at the queue
- $La/R > 1$ : more “work” arriving is more than can be serviced - average delay infinite!
- $La/R \leq 1$ : nature of arriving traffic
- $La/R \sim 0$ : avg. queueing delay small

$La/R > 1$ : Average rate at which bits arrive at the queue exceeds the rate at which the bits can be transmitted from the queue.



*Design your system so that the traffic intensity is no greater than 1.*

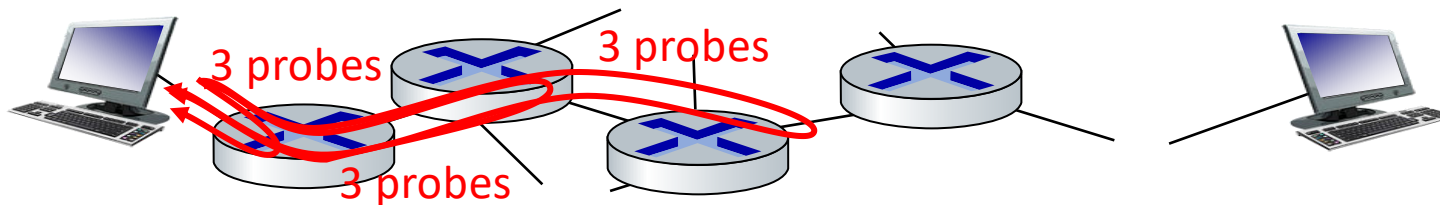


$La/R \sim 0$



$La/R \rightarrow 1$

- what do “real” Internet delay & loss look like?
- **tracert** program: provides delay measurement from source to router along end-end Internet path towards destination.  
For all  $i$ :
  - sends three packets that will reach router  $i$  on path towards destination (with time-to-live field value of  $i$ )
  - router  $i$  will return packets to sender
  - sender measures time interval between transmission and reply



**traceroute:** gaia.cs.umass.edu to www.eurecom.fr

3 delay measurements from  
gaia.cs.umass.edu to cs-gw.cs.umass.edu

3 delay measurements  
to border1-rt-fa5-1-0.gw.umass.edu

trans-oceanic link

looks like delays  
decrease! Why?

\* means no response (probe lost, router not replying)

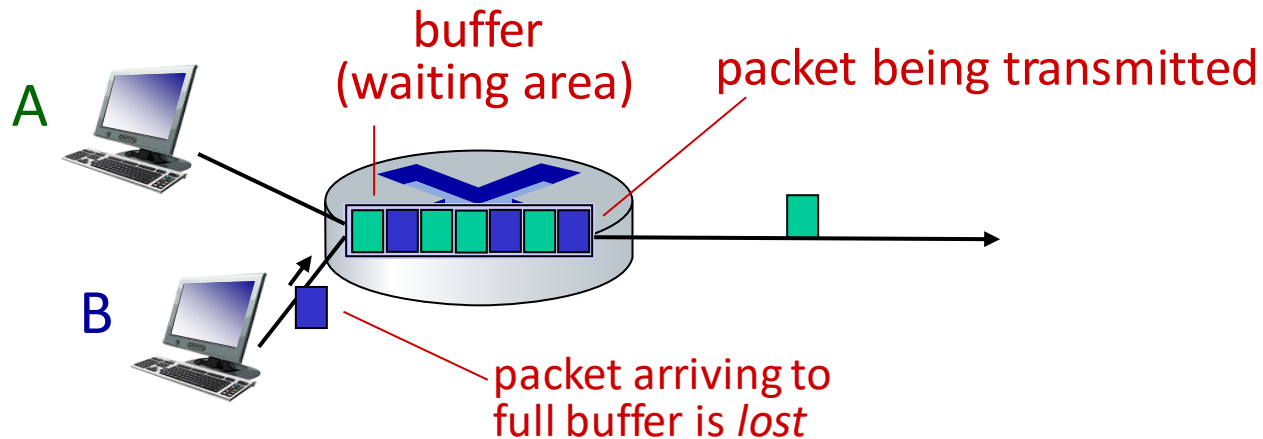
```
1 cs-gw (128.119.240.254) 1 ms 1 ms 2 ms
2 border1-rt-fa5-1-0.gw.umass.edu (128.119.3.145) 1 ms 1 ms 2 ms
3 cht-vbns.gw.umass.edu (128.119.3.130) 6 ms 5 ms 5 ms
4 jn1-at1-0-0-19.wor.vbns.net (204.147.132.129) 16 ms 11 ms 13 ms
5 jn1-so7-0-0-0.wae.vbns.net (204.147.136.136) 21 ms 18 ms 18 ms
6 abilene-vbns.abilene.ucaid.edu (198.32.11.9) 22 ms 18 ms 22 ms
7 nycm-wash.abilene.ucaid.edu (198.32.8.46) 22 ms 22 ms 22 ms
8 62.40.103.253 (62.40.103.253) 104 ms 109 ms 106 ms
9 de2-1.de1.de.geant.net (62.40.96.129) 109 ms 102 ms 104 ms
10 de.fr1.fr.geant.net (62.40.96.50) 113 ms 121 ms 114 ms
11 renater-gw.fr1.fr.geant.net (62.40.103.54) 112 ms 114 ms 112 ms
12 nio-n2.cssi.renater.fr (193.51.206.13) 111 ms 114 ms 116 ms
13 nice.cssi.renater.fr (195.220.98.102) 123 ms 125 ms 124 ms
14 r3t2-nice.cssi.renater.fr (195.220.98.110) 126 ms 126 ms 124 ms
15 eurecom-valbonne.r3t2.ft.net (193.48.50.54) 135 ms 128 ms 133 ms
16 194.214.211.25 (194.214.211.25) 126 ms 128 ms 126 ms
17 * * *
18 * * *
19 fantasia.eurecom.fr (193.55.113.142) 132 ms 128 ms 136 ms
```

\* Do some traceroutes from exotic countries at [www.traceroute.org](http://www.traceroute.org)

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## Performance: Packet loss

- queue (aka buffer) preceding link in buffer has finite capacity
- packet arriving to full queue dropped (aka lost)
- lost packet may be retransmitted by previous node, by source end system, or not at all

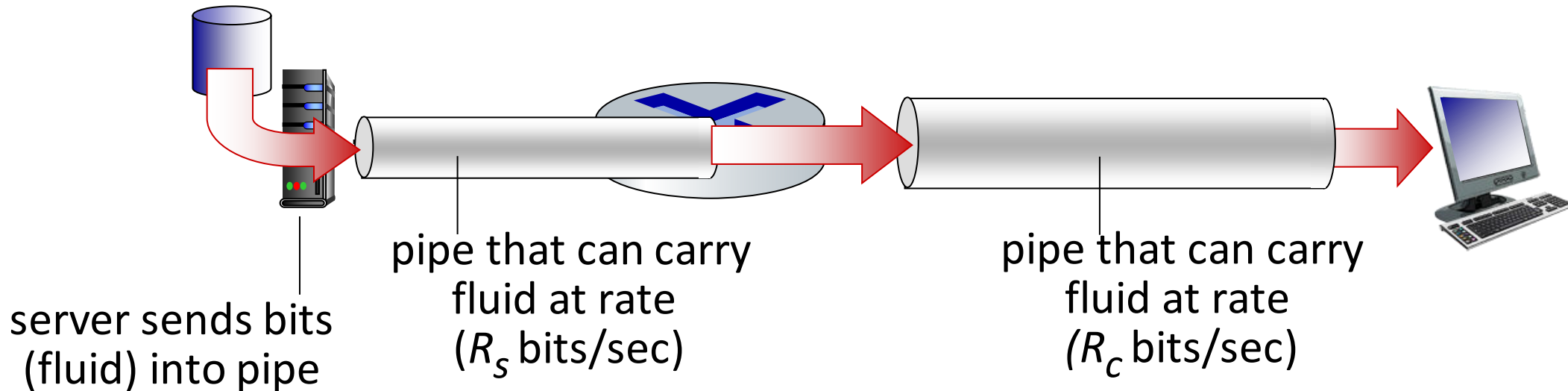


\* Check out the Java applet for an interactive animation on queuing and loss

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## Performance: Throughput

- *throughput*: rate (bits/time unit) at which bits are being sent from sender to receiver
  - *instantaneous*: rate at given point in time
  - *average*: rate over longer period of time

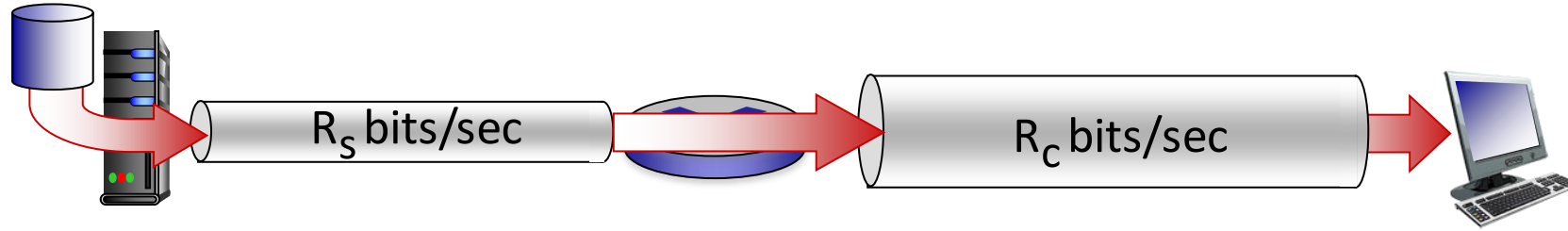




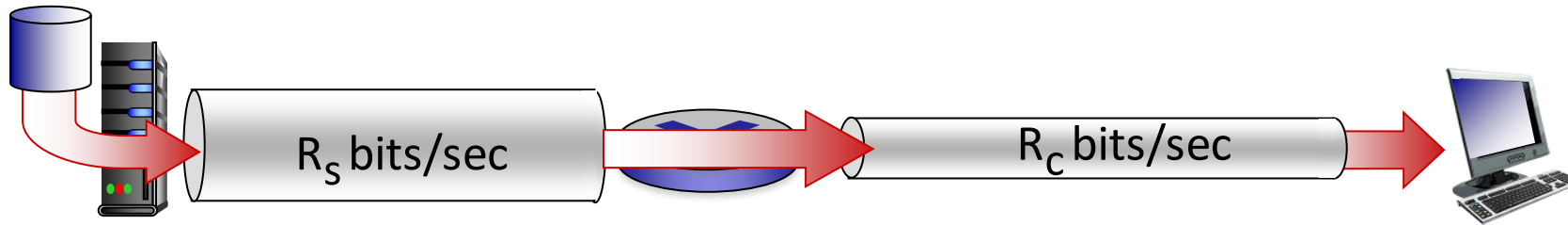
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## Performance: Throughput (more)

$R_s < R_c$  What is average end-end throughput?



$R_s > R_c$  What is average end-end throughput?

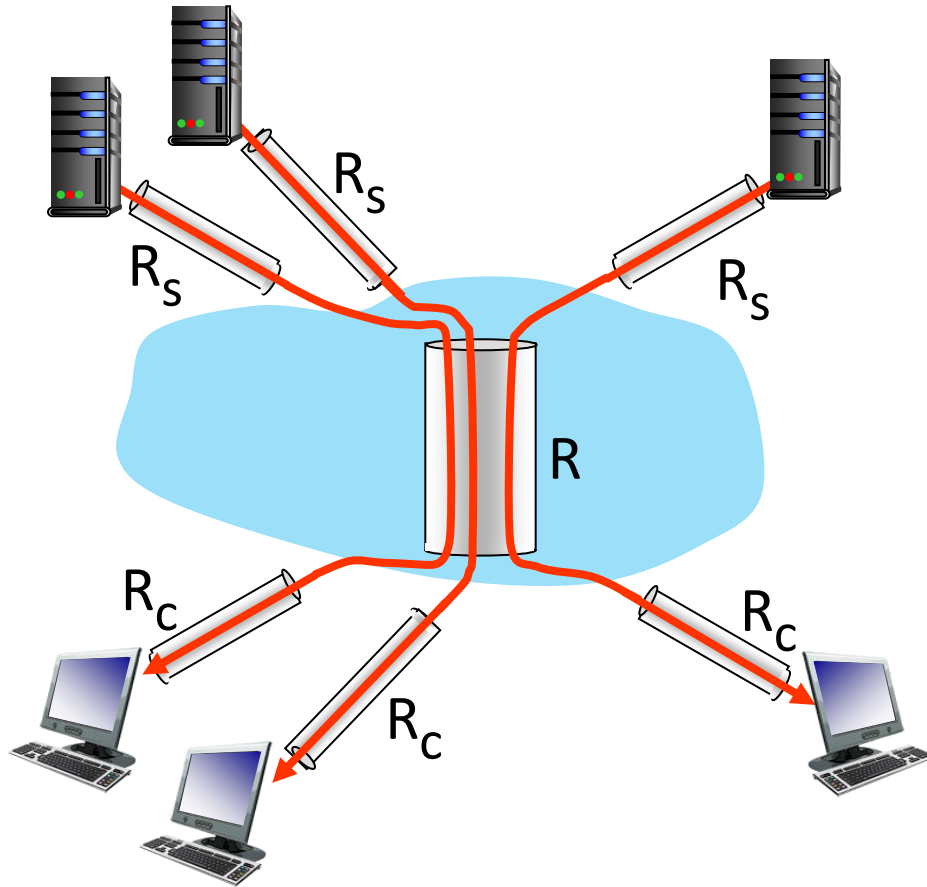


*bottleneck link*

link on end-end path that constrains end-end throughput

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## Performance: Throughput – Network Scenario



10 connections (fairly) share  
backbone bottleneck link  $R$  bits/sec

- per-connection end-end throughput:  $\min(R_c, R_s, R/10)$
- in practice:  $R_c$  or  $R_s$  is often bottleneck

\* Check out the online interactive exercises for more examples:  
[http://gaia.cs.umass.edu/kurose\\_ross/](http://gaia.cs.umass.edu/kurose_ross/)

- Suppose  $R_s = 2$  Mbps,  $R_c = 1$  Mbps,  $R = 5$  Mbps
- 10 clients from 10 servers = 10 downloads

End-to-end throughput for each  
download is now reduced to 500 kbps.



Thank You  
For Your Attention



**THANK YOU**

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