

DESIGN AND ANALYSIS OF ALGORITHMS

The Knapsack Problem

Reetinder Sidhu

Department of Computer Science and Engineering



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UNIT 5: Limitations of Algorithmic Power and Coping with the Limitations

- Dynamic Programming
 - Computing a Binomial Coefficient
 - The Knapsack Problem
 - Memory Functions
 - Warshall's and Floyd's Algorithms
- Limitations of Algorithmic Power
 - Lower-Bound Arguments
 - Decision Trees
 - P, NP, and NP-Complete, NP-Hard Problems
- Coping with the Limitations
 - Backtracking
 - Branch-and-Bound. Architecture (microprocessor instruction set)

Concepts covered

- The Knapsack Problem
 - Introduction
 - Recurrence
 - Example



Problem Definition



- Given
 - integer weights: w_1 w_2 ... w_n values: v_1 v_2 ... v_n
 - knapsack of capacity W (integer W > 0)
- ullet Find the most valuable subset of items such that sum of their weights does not exceed W

Kanpsack Recurrence



- To design a dynamic programming algorithm, we need to derive a recurrence relation that expresses a solution to an instance of the knapsack problem in terms of solutions to its smaller subinstances
- Consider the smaller knapsack problem where number of items is i ($i \le n$) and the knapsack capacity id i (i < W)
- Then

$$F(i,j) = \begin{cases} \max(F(i-1,j), & v_i + F(i-1,j-w_i)) & \text{if } j - w_i \ge 0 \\ F(i-1,j) & \text{if } j - w_i < 0 \end{cases}$$



$$F(i,j) = egin{cases} ext{max}(F(i-1,j), & v_i + F(i-1,j-w_i)) & ext{if } j-w_i \geq 0 \ F(i-1,j) & ext{if } j-w_i < 0 \end{cases}$$

Dynamic Programming Example

item i	weight w _i	value <i>v_i</i>
1	2	12
2	1	10
3	3	20
4	2	15



$$F(i,j) = egin{cases} \max(F(i-1,j), & v_i + F(i-1,j-w_i)) & \text{if } j-w_i \geq 0 \ F(i-1,j) & \text{if } j-w_i < 0 \end{cases}$$

Dynamic Programming Example

item i	weight w _i	value <i>v_i</i>
1	2	12

	capacity <i>j</i>				
i	1	2	3	4	5
1					
2					
3					
4					



$$F(i,j) = egin{cases} \max(F(i-1,j), & v_i + F(i-1,j-w_i)) & \text{if } j - w_i \geq 0 \ F(i-1,j) & \text{if } j - w_i < 0 \end{cases}$$

Dynamic Programming Example

item i	weight w _i	value <i>v_i</i>
1	2	12

		ca	pacit	y <i>j</i>	
i	1	2	3	4	5
1	0				
2					
3					
4					



$$F(i,j) = egin{cases} \max(F(i-1,j), & v_i + F(i-1,j-w_i)) & \text{if } j - w_i \geq 0 \ F(i-1,j) & \text{if } j - w_i < 0 \end{cases}$$

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i	1	2	3	4	5
1	0	12			
2					
3					
4					



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1	2	12

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i	1	2	3	4	5
1	0	12	12		
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3					
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$$F(i,j) = \begin{cases} \max(F(i-1,j), & v_i + F(i-1,j-w_i)) & \text{if } j - w_i \ge 0 \\ F(i-1,j) & \text{if } j - w_i < 0 \end{cases}$$

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1	0	12	12	12	
2					
3					
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$$F(i,j) = \begin{cases} \max(F(i-1,j), & v_i + F(i-1,j-w_i)) & \text{if } j - w_i \ge 0 \\ F(i-1,j) & \text{if } j - w_i < 0 \end{cases}$$

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What	is	the	maximum	value	that	can	be
stored	l in	a kn	apsack of c	apacit	y 6?		

		ca	pacit	y j	
i	1	2	3	4	5
1	0	12	12	12	12
2					
3					
4					



$$F(i,j) = \begin{cases} \max(F(i-1,j), & v_i + F(i-1,j-w_i)) & \text{if } j - w_i \ge 0 \\ F(i-1,j) & \text{if } j - w_i < 0 \end{cases}$$

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2	10				
3					
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1	0	12	12	12	12
2	10	12			
3					
4					



$$F(i,j) = \begin{cases} \max(F(i-1,j), & v_i + F(i-1,j-w_i)) & \text{if } j - w_i \ge 0 \\ F(i-1,j) & \text{if } j - w_i < 0 \end{cases}$$

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item i	weight w _i	value <i>v_i</i>
1	2	12
2	1	10

		ca	pacit	y j	
i	1	2	3	4	5
1	0	12	12	12	12
2	10	12	22		
3					
4					



$$F(i,j) = \begin{cases} \max(F(i-1,j), & v_i + F(i-1,j-w_i)) & \text{if } j - w_i \ge 0 \\ F(i-1,j) & \text{if } j - w_i < 0 \end{cases}$$

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1	0	12	12	12	12
2	10	12	22	22	22
3					
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		ca	pacit	y J	
i	1	2	3	4	5
1	0	12	12	12	12
2	10	12	22	22	22
3					
4					



$$F(i,j) = \begin{cases} \max(F(i-1,j), & v_i + F(i-1,j-w_i)) & \text{if } j - w_i \ge 0 \\ F(i-1,j) & \text{if } j - w_i < 0 \end{cases}$$

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		capacity <i>j</i>							
i	1	2	3	4	5				
1	0	12	12	12	12				
2	10	12	22	22	22				
3	10								
1									



$$F(i,j) = egin{cases} \max(F(i-1,j), & v_i + F(i-1,j-w_i)) & \text{if } j-w_i \geq 0 \ F(i-1,j) & \text{if } j-w_i < 0 \end{cases}$$

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1	0	12	12	12	12			
2	10	12	22	22	22			
3	10	12						
1								



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3	10	12	22					



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1	0	12	12	12	12		
2	10	12	22	22	22		
3	10	12	22	30			
4							



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1	0	12	12	12	12	
2	10	12	22	22	22	
3	10	12	22	30	32	
1						



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4	2	15

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i	1	2	3	4	5	
1	0	12	12	12	12	
2	10	12	22	22	22	
3	10	12	22	30	32	
1						



$$F(i,j) = \begin{cases} \max(F(i-1,j), & v_i + F(i-1,j-w_i)) & \text{if } j - w_i \ge 0 \\ F(i-1,j) & \text{if } j - w_i < 0 \end{cases}$$

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		ca	pacit	y <i>j</i>	
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1	0	12	12	12	12
2	10	12	22	22	22
3	10	12	22	30	32
4	10				



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1	0	12	12	12	12
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3	10	12	22	30	32
4	10	15			



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1	2	12
2	1	10
3	3	20
4	2	15

	capacity <i>j</i>				
i	1	2	3	4	5
1	0	12	12	12	12
2	10	12	22	22	22
3	10	12	22	30	32
4	10	15	25		



$$F(i,j) = \begin{cases} \max(F(i-1,j), & v_i + F(i-1,j-w_i)) & \text{if } j - w_i \ge 0 \\ F(i-1,j) & \text{if } j - w_i < 0 \end{cases}$$

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1	2	12
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3	3	20
4	2	15

	capacity <i>j</i>				
i	1	2	3	4	5
1	0	12	12	12	12
2	10	12	22	22	22
3	10	12	22	30	32
4	10	15	25	30	



$$F(i,j) = \begin{cases} \max(F(i-1,j), & v_i + F(i-1,j-w_i)) & \text{if } j - w_i \ge 0 \\ F(i-1,j) & \text{if } j - w_i < 0 \end{cases}$$

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1	2	12
2	1	10
3	3	20
4	2	15

	capacity <i>j</i>				
i	1	2	3	4	5
1	0	12	12	12	12
2	10	12	22	22	22
3	10	12	22	30	32
4	10	15	25	30	37



$$F(i,j) = \begin{cases} \max(F(i-1,j), & v_i + F(i-1,j-w_i)) & \text{if } j - w_i \ge 0 \\ F(i-1,j) & \text{if } j - w_i < 0 \end{cases}$$

Dynamic Programming Example

item <i>i</i>	weight w _i	value <i>v_i</i>
1	2	12
2	1	10
3	3	20
4	2	15

What is the maximum value that can be stored in a knapsack of capacity 6?

	capacity <i>j</i>				
i	1	2	3	4	5
1	0	12	12	12	12
2	10	12	22	22	22
3	10	12	22	30	32
4	10	15	25	30	37

Given above 6 items, maximum value that can be stored in a knapsack of capacity 5 is **37**

THE KNAPSACK PROBLEM Complexity



- Space complexity: $\Theta(nW)$
- Time complexity: $\Theta(nW)$
- Time to compose optimal solution: O(n)

Think About It



Think About It



 Write pseudocode of the bottom-up dynamic programming algorithm for the knapsack problem

Think About It



- Write pseudocode of the bottom-up dynamic programming algorithm for the knapsack problem
- True or False:
 - A sequence of values in a row of the dynamic programming table for the knapsack problem is always nondecreasing?

Think About It



- Write pseudocode of the bottom-up dynamic programming algorithm for the knapsack problem
- True or False:
 - A sequence of values in a row of the dynamic programming table for the knapsack problem is always nondecreasing?
 - A sequence of values in a column of the dynamic programming table for the knapsack problem is always nondecreasing?