

Balanced Trees

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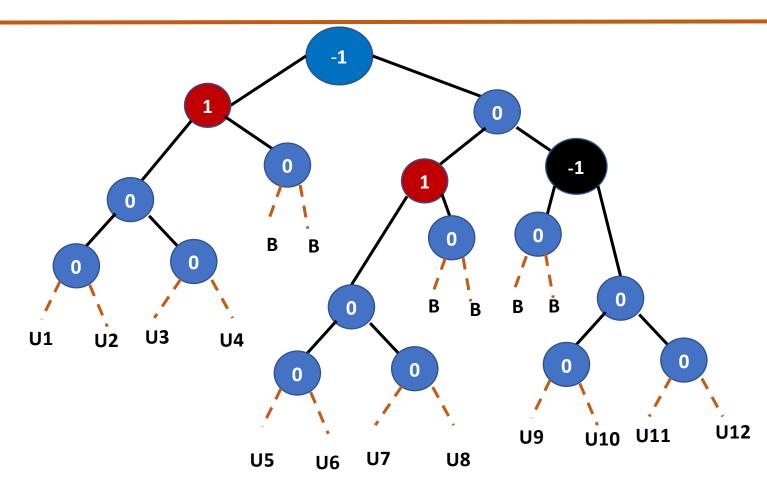
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Possible insertion into AVL tree

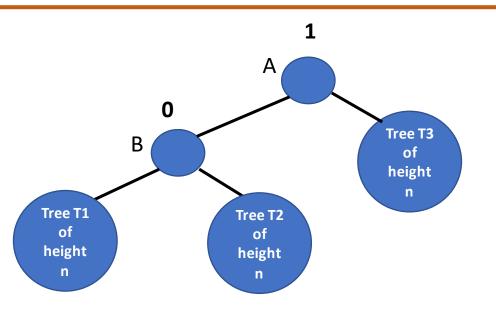




- Unbalanced insertions are indicated by U
- Balanced insertions are indicated by B

Unbalanced Tree after inserting a node to left subtree





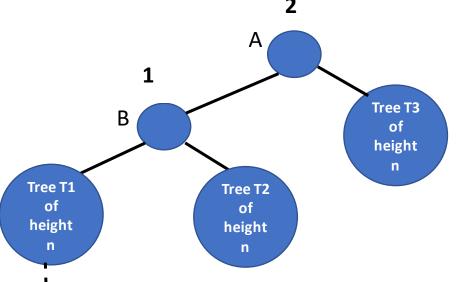
Balance factor(A) =
$$(n+1) - n = 1$$

Balance factor(B) = $n - n = 0$

- Let us consider A is the youngest ancestor which becomes unbalanced
- Balance factor of A should 1 before insertion
- A should have a left child B with the balance factor of 0

Unbalanced Tree after inserting a node to left subtree

- Newly created node in inserted into left subtree of node B
- Changing the balance B to 1 and A to 2
- A is the youngest ancestor of the new node to become unbalanced



Newly inserted

node

- To maintain the balance: Tree needs to be transformed
 - ✓ Transformed tree is balanced
 - ✓ Binary search tree property is maintained after

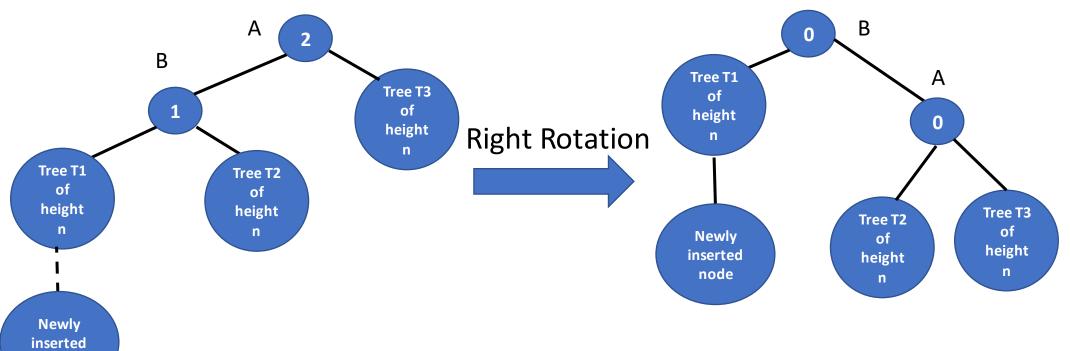


Transformed Balanced Tree after Rotations

node



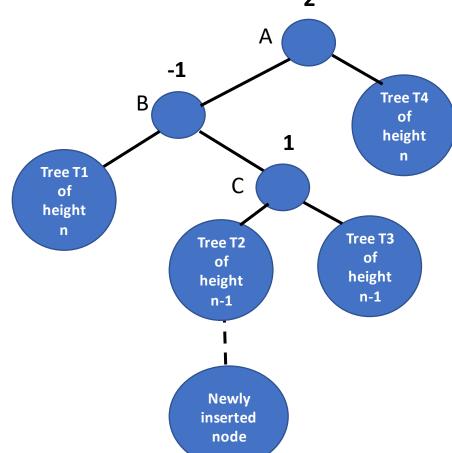
• To maintain a balance we need to rotate sub tree B rooted at node A



Unbalanced Tree after inserting a node to right subtree

Newly created node is inserted into right subtree changing the balance:

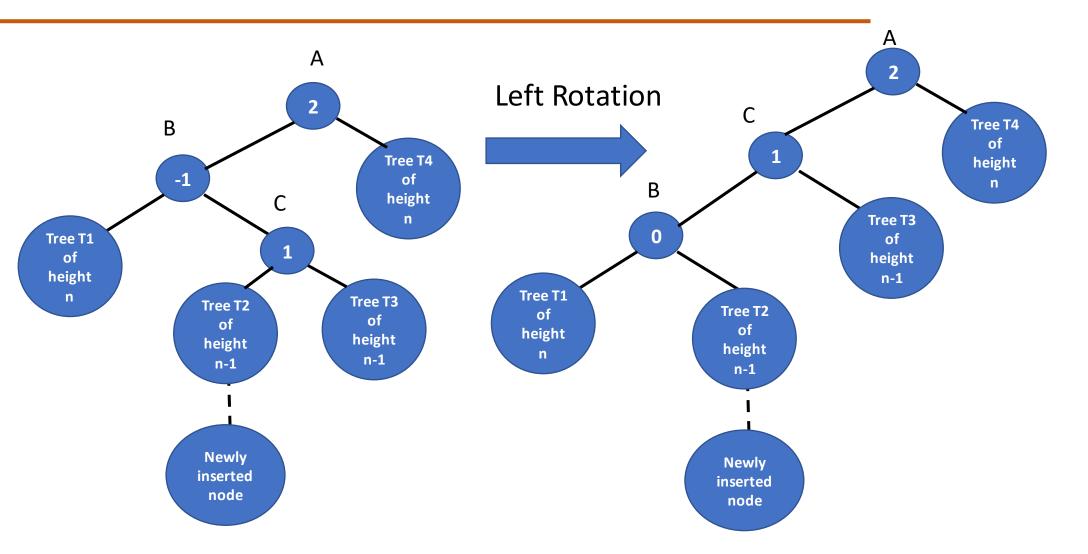
Balance factor(C)= n-(n-1) = 1Balance factor(B)= n-(n+1) = -1Balance factor(A)= n+2-n=2





Transformed Balanced tree after Rotations

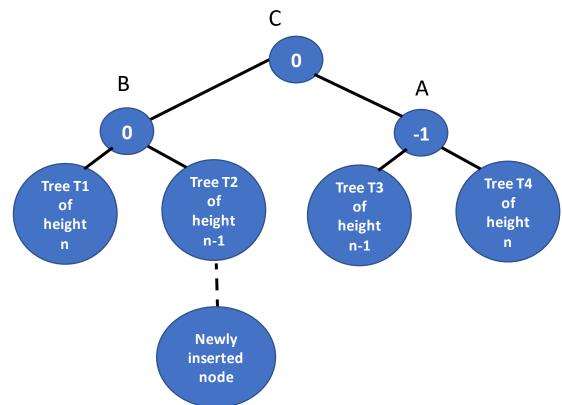




Transformed Balanced tree after Rotations







AVL tree Insertion



- Insertion in AVL tree is performed using standard BST Insertion
- If tree becomes unbalanced, we rebalance the tree using left or right rotation
- If node X is inserted into BST
- we need to find the youngest ancestor which becomes unbalanced

Four cases:

- Balance factor of node is greater than 1 unbalanced node(U)
 - ✓ Left-Left case
 - if newly inserted key is less than the key in the left subtree' root
 - ✓ Left-Right case
 - if newly inserted key is greater than the key in the left subtree's root

AVL tree Insertion

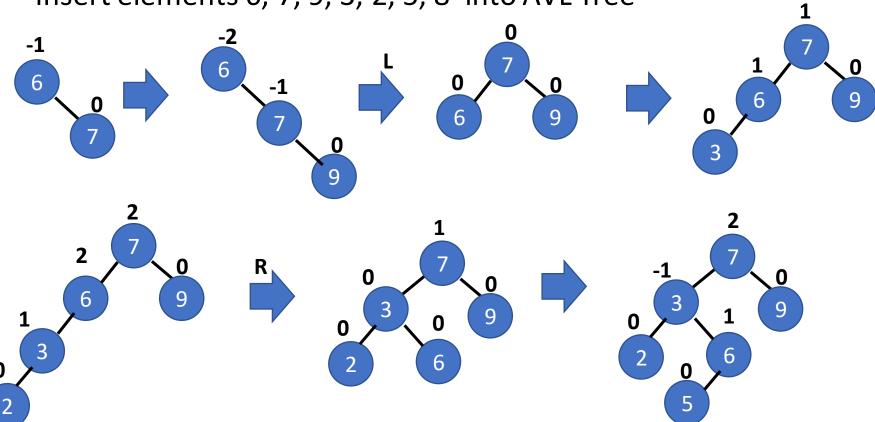


Four cases:

- Balance factor of node is lesser than -1 unbalanced node(U)
 - ✓ Right-Right case
 - if newly inserted key is greater than the key in the right subtree' root.
 - ✓ Right-Left case
 - if newly inserted key is lesser than the key in the right subtree's root.

Examples – AVL Tree Insertions

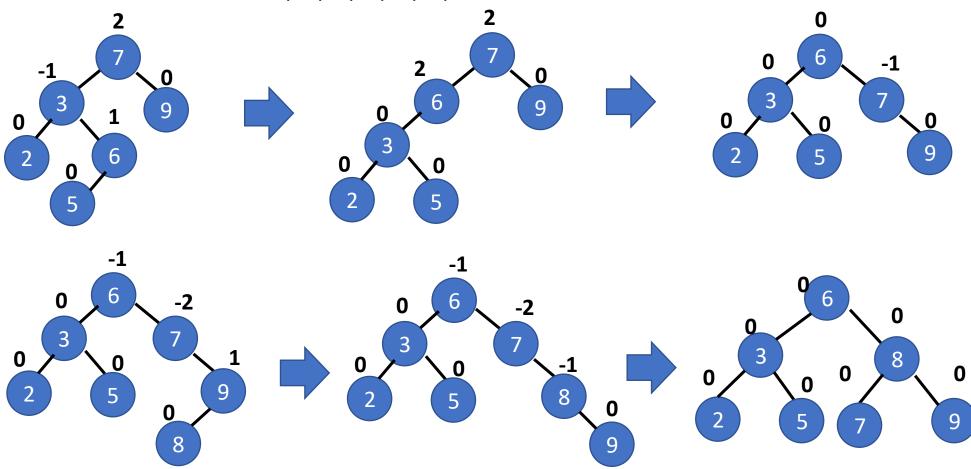
• Insert elements 6, 7, 9, 3, 2, 5, 8 into AVL Tree





Example – AVL Tree Insertions

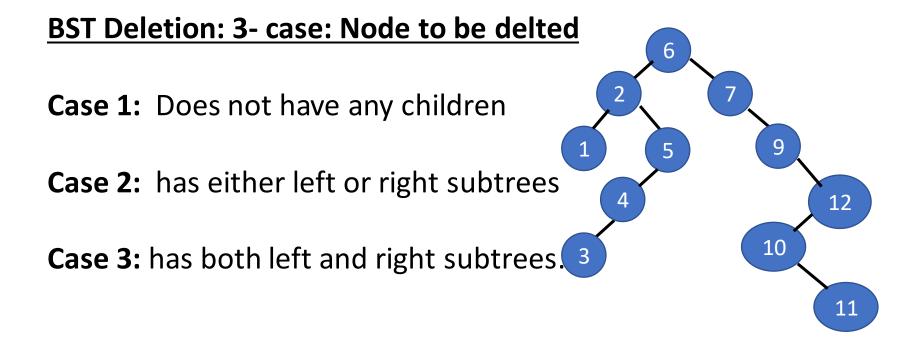
• Insert elements 6, 7, 9, 3, 2, 5, 8 into AVL Tree





Deletions in AVL tree

- Deletion in AVL tree is performed using standard BST Deletion
- If tree becomes unbalanced, we rebalance the tree using left or right rotation





Deletions in AVL tree

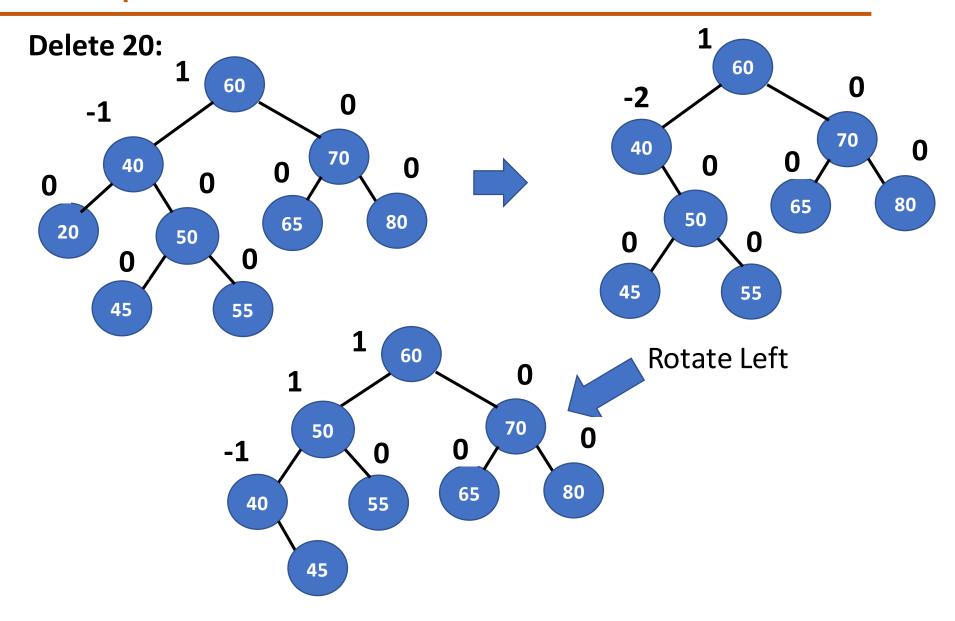
- If node X is deleted from the BST
- we need to find the youngest ancestor which becomes unbalanced

Four cases:

- Balance factor of a node is greater than 1 unbalanced node(U)
 - ✓ Left-Left case
 - Balance factor of left subtree's root is greater than or equal to 0
 - ✓ Left-Right case
 - Balance factor of left subtree's root is less than 0
- Balance factor of unbalanced node is less than -1
 - ✓ Right-Right case
 - Balance factor of Right subtree's root is less than or equal to 0
 - √ Right-Left case
 - ■Balance factor of sub tree is grater than 0



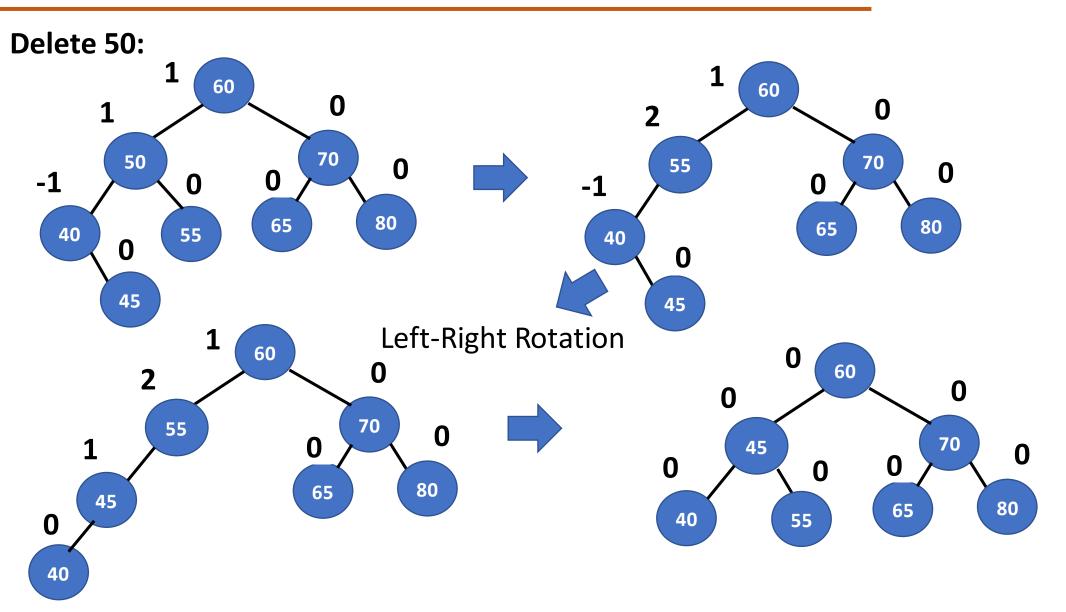
Example – Deletions in AVL tree





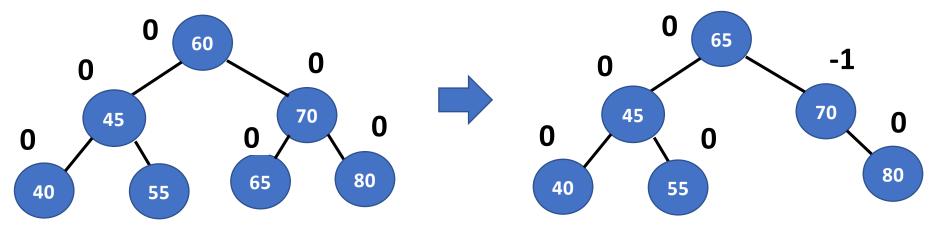
Example – Deletions in AVL tree



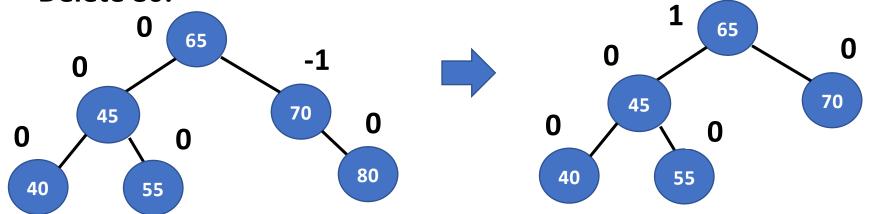


Example – Deletions in AVL tree

Delete 60:



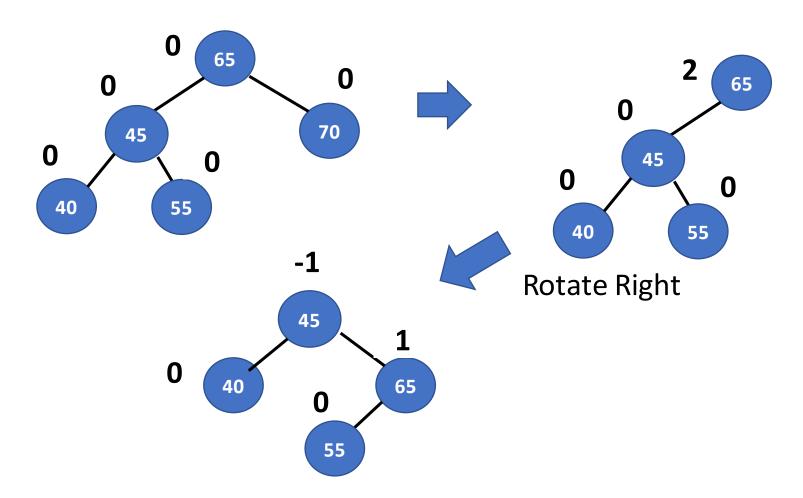
Delete 80:





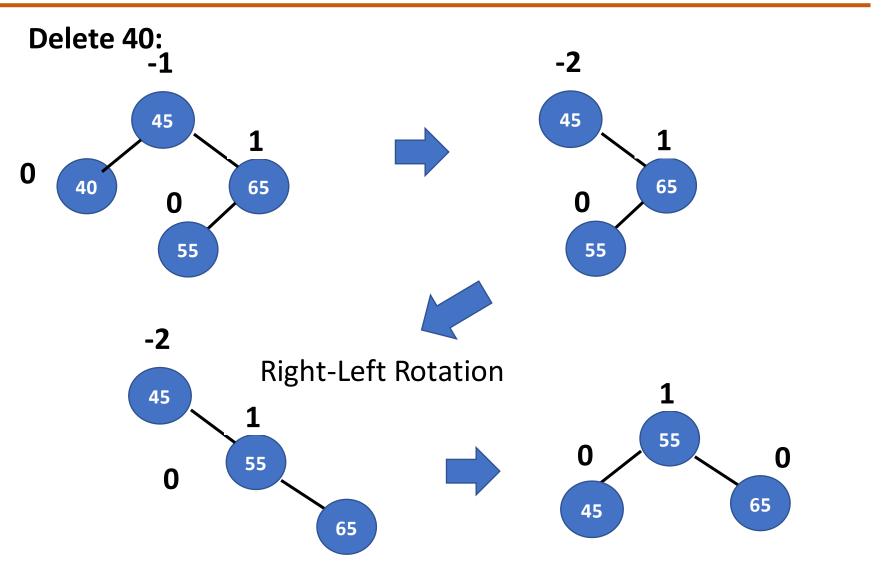
Example – Deletions in AVL tree

Delete 70:





Example – Deletions in AVL tree







THANK YOU

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