

OPERATING SYSTEMS

I/O Management, System Protection and Security

Arya S S

Department of Computer Science

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System Protection- Access Matrix

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Slides Credits for all PPTs of this course



- The slides/diagrams in this course are an **adaptation, combination,** and **enhancement** of material from the following resources and persons:
1. Slides of Operating System Concepts, Abraham Silberschatz, Peter Baer Galvin, Greg Gagne - 9th edition 2013 and some slides from 10th edition 2018
 2. Some conceptual text and diagram from Operating Systems - Internals and Design Principles, William Stallings, 9th edition 2018
 3. Some presentation transcripts from A. Frank – P. Weisberg
 4. Some conceptual text from Operating Systems: Three Easy Pieces, Remzi Arpaci-Dusseau, Andrea Arpaci Dusseau

- View protection abstractly as a matrix (**access matrix**)
- Rows represent domains
- Columns represent objects
- **Access(i, j)** is the set of operations that a process executing in Domain_i can invoke on Object_j

object domain	F_1	F_2	F_3	printer
D_1	read		read	
D_2				print
D_3		read	execute	
D_4	read write		read write	

- If a process in Domain D_i tries to do “op” on object O_j , then “op” must be in the access matrix
- User who creates object can define access column for that object
- Can be expanded to dynamic protection
 - Operations to add, delete access rights
 - Special access rights:
 - 4 *owner of O_i*
 - 4 *copy op from O_i to O_j (denoted by “*”)*
 - 4 *control – D_i can modify D_j access rights*
 - 4 *transfer – switch from domain D_i to D_j*
 - *Copy and Owner* applicable to an object
 - *Control* applicable to domain object

- **Access matrix** design separates mechanism from policy
 - Mechanism
 - 4 Operating system provides access-matrix + rules
 - 4 It ensures that the matrix is only manipulated by authorized agents and that rules are strictly enforced
 - Policy
 - 4 User dictates policy
 - 4 Who can access what object and in what mode
- But doesn't solve the general confinement problem i.e. preventing a process from taking disallowed actions
 - Ex: In a client/server situation, preventing a server from leaking information that the client considers confidential

Processes should be able to switch from one domain to another.

Switching from domain D_i to domain D_j is allowed if and only if the access right switch $\in \text{access}(i, j)$.

domain \ object	F_1	F_2	F_3	laser printer	D_1	D_2	D_3	D_4
D_1	read		read			switch		
D_2				print			switch	switch
D_3		read	execute					
D_4	read write		read write		switch			

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Access Matrix with Copy Rights

- ❑ The ability to copy an access right from one domain (or row) of the access matrix to another is denoted by an asterisk (*) appended to the access right.

domain \ object	F_1	F_2	F_3
D_1	execute		write*
D_2	execute	read*	execute
D_3	execute		

(a)

domain \ object	F_1	F_2	F_3
D_1	execute		write*
D_2	execute	read*	execute
D_3	execute	read	

(b)

The copy scheme has 3 variants:

1. A right is copied from $\text{access}(i,j)$ to $\text{access}(k,j)$ is not limited: This action is called **copy**.
 - When the right Read^* is copied from $\text{access}(i,j)$ to $\text{access}(k,j)$, the Read^* is created.
 - So, a process executing in D_k can further copy the right Read^* .
2. Propagation of the copy right may be limited: This action is called **limited copy**.
 - When the right Read^* is copied from $\text{access}(i,j)$ to $\text{access}(k,j)$, only the Read (not Read^*) is created.
 - So, a process executing in D_k cannot further copy the right Read .
3. A right is copied from $\text{access}(i,j)$ to $\text{access}(k,j)$; it is then removed from $\text{access}(i,j)$.
 - This action is called a **transfer of a right**, rather than a copy

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Access Matrix with Owner Rights

object domain	F_1	F_2	F_3
D_1	owner execute		write
D_2		read* owner	read* owner write
D_3	execute		

(a)

object domain	F_1	F_2	F_3
D_1	owner execute		write
D_2		owner read* write*	read* owner write
D_3		write	write

(b)

- ❑ **Owner** right controls addition of new rights and removal of some rights.
- ❑ Domain D_1 is the owner of F_1 and can add /delete any valid right in column F_1
- ❑ Owner rights allow a process to change the entries in a column

- The copy and owner rights allow a process to change the entries in a column.
- A mechanism is now needed to change the entries in a row.
- The control right is applicable only to domain objects (rows).
- If $\text{access}(i,j)$ includes the control right, then a process executing in D_i can remove any access right from row j .
- If we include the *control* right in $\text{access}(D_2, D_4)$, then, a process executing in domain D_2 could modify domain D_4 .

object \ domain	F_1	F_2	F_3	laser printer	D_1	D_2	D_3	D_4
D_1	read		read			switch		
D_2				print			switch	switch
D_3		read	execute					
D_4	read write		read write		switch			

Fig A

object \ domain	F_1	F_2	F_3	laser printer	D_1	D_2	D_3	D_4
D_1	read		read			switch		
D_2				print			switch	switch control
D_3		read	execute					
D_4	write		write		switch			

Fig B: Modified access matrix of fig A

Add **control** right in access(D_2 , D_4).

Then, a process executing in D_2 (row) could modify D_4 (row).



THANK YOU

Arya S S

Department of Computer Science Engineering

aryadeep@pes.edu