

March 2021: IN SEMESTER ASSESSMENT B Tech IV SEMESTER

TEST – 1

Scheme and Solution

UE19CS254: Operating Systems

Time: 2Hr

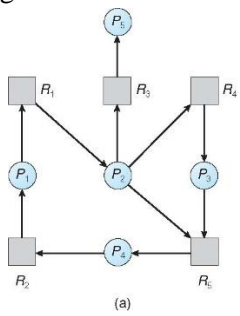
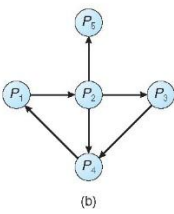
Answer All Questions

Max Marks: 60

1.	a)	<p>What is peer to peer computing? Name the applications of the peer-to-peer computing environment.</p> <p>Ans:</p> <ul style="list-style-type: none">P2P does not distinguish clients and servers. Instead, all nodes are considered peers Where each peer acts as client, server or both-----1MExamples include Napster and Gnutella, Voice over IP (VoIP) such as Skype -----1M	2																		
	b)	<p>What are the services provided by the operating systems?</p> <p>Ans</p> <p>Services provided by OS are,</p> <ul style="list-style-type: none">User interface - Almost all operating systems have a user interface (UI). Varies between Command-Line (CLI), Graphics User Interface (GUI), Batch Program execution - The system must be able to load a program into memory and to run that program, end execution, either normally or abnormally (indicating error)I/O operations - A running program may require I/O, which may involve a file or an I/O device.File-system manipulation - The file system is of particular interest. Programs need to read and write files and directories, create and delete them, search them, list file Information, permission management.Communications – Processes may exchange information, on the same computer or between computers over a networkCommunications may be via shared memory or through message passing (packets moved by the OS)Error detection – OS needs to be constantly aware of possible errors.Resource allocation - When multiple users or multiple jobs running concurrently, resources must be allocated to each of them, types of resources are- CPU cycles, main memory, file storage, I/O devices.Accounting - To keep track of which users use how much and what kinds of computer resources <p>Any 3 services must be specified 1M each</p>	3																		
	c)	<p>Consider the set of 5 processes whose arrival time and burst time are given below-</p> <table><tr><th>Process ID</th><th>Arrival time</th><th>Burst Time(ms)</th></tr><tr><td>P1</td><td>4</td><td>6</td></tr><tr><td>P2</td><td>2</td><td>4</td></tr><tr><td>P3</td><td>0</td><td>2</td></tr><tr><td>P4</td><td>3</td><td>3</td></tr><tr><td>P5</td><td>1</td><td>4</td></tr></table>	Process ID	Arrival time	Burst Time(ms)	P1	4	6	P2	2	4	P3	0	2	P4	3	3	P5	1	4	5
Process ID	Arrival time	Burst Time(ms)																			
P1	4	6																			
P2	2	4																			
P3	0	2																			
P4	3	3																			
P5	1	4																			

	<p>i. Calculate the waiting time and average waiting for round Robin Scheduling algorithm (given time quantum= 3)</p> <table><tr><td>P3</td><td>P5</td><td>P2</td><td>P4</td><td>P1</td><td>P5</td><td>P2</td><td>P1</td></tr><tr><td>0</td><td>2</td><td>5</td><td>8</td><td>11</td><td>14</td><td>15</td><td>16</td><td>19</td></tr></table> <p>Waiting time for P1=16-3-4=9 Waiting time for P2=15-3-2=10 Waiting Time for P3=0 Waiting Time for P4=8-3=5 Waiting Time for P5=14-3-1=10 Average waiting time=(9+10+0+5+10)/5=6.8ms-----1/2M</p> <p>ii. What will be the number of context switches for process P5? (Exclude the context switches at the start of the process and at termination) No.of context switches=2-----1/2M</p> <p>iii. Does an increase in time quantum decrease the number of context switches? Yes context switches are reduced-----1/2M</p> <p>iv. List the processes in the ready queue after 11ms. Process in the ready queue: P5, P2, P1-----1M</p>	P3	P5	P2	P4	P1	P5	P2	P1	0	2	5	8	11	14	15	16	19	
P3	P5	P2	P4	P1	P5	P2	P1												
0	2	5	8	11	14	15	16	19											
2.	<p>a) What are the possible conditions for the process to move from running state to ready state? A diagram depicting the same is required. Soln Diagram—2M</p> <p>The process moves from running state to ready state when,-----2M</p> <p>I process needs I/o II child process is created using fork() III time slice has expired (in case of RR scheduling) IV Process is interrupted</p>	4																	
	<p>b) Why short-term scheduler executes at a higher frequency than long-term scheduler? The process executes for the few milliseconds before it goes for I/O. Context switch times is very less</p>	2																	
	<p>c) An operating system uses Shortest Remaining Time First (SRTF) process scheduling algorithm. Consider the arrival times and execution times for the following processes:</p> <table><tr><td>Process</td><td>Execution time</td><td>Arrival time</td></tr><tr><td>P1</td><td>20</td><td>0</td></tr><tr><td>P2</td><td>25</td><td>15</td></tr><tr><td>P3</td><td>10</td><td>30</td></tr><tr><td>P4</td><td>15</td><td>45</td></tr></table> <p>i. What is the average waiting time for these processes?</p>	Process	Execution time	Arrival time	P1	20	0	P2	25	15	P3	10	30	P4	15	45	4 (2+1+1)		
Process	Execution time	Arrival time																	
P1	20	0																	
P2	25	15																	
P3	10	30																	
P4	15	45																	

		<p>Solution:</p> <p>Gantt Chart:</p> <table><tr><td>P1</td><td>P2</td><td>P3</td><td>P2</td><td>P4</td></tr><tr><td>20</td><td>30</td><td>40</td><td>55</td><td>70</td></tr></table> <p>Waiting Time = Start Time - Arrival Time-execution time</p> <p>i. Waiting time for P1=0</p> <p>waiting time for P2=40-15-10=15</p> <p>Waiting time for P3=30-30-0=0</p> <p>Waiting time for P4=55-45-0=10</p> <p>Average waiting time=(0+15+0+10)/4=6.25 time units-----1M</p> <p>ii. No-----1M</p> <p>iii. Round Robin Scheduling algorithm-----1M</p>	P1	P2	P3	P2	P4	20	30	40	55	70	
P1	P2	P3	P2	P4									
20	30	40	55	70									
3.	a)	<p>Why is it that threads are faster to create than processes? What are the benefits of threads?</p> <p>Threads are faster than process as they share parent processes address space.—1M</p> <p>Benefits of threads</p> <ul style="list-style-type: none">• Responsiveness – may allow continued execution if part of process is blocked or performing lengthy operations.• Resource Sharing – threads share resources of process, easier than shared memory or message passing.• Economy – cheaper than process creation, thread switching lower overhead than context switching.• Scalability – process can take advantage of multiprocessor architectures. <p>Any 3 benefits specified—3M</p>	4										
	b)	<p>Write Peterson's solution to the critical section problem.</p> <p>Ans</p> <pre>do { flag[i] = true;-----1m turn = j;-----1m while (flag[j] && turn == j);-----1m /* critical section */ flag[i] = false;-----1m /* remainder section */ }while (true);</pre>	4										
	c)	<p>List two programming examples of multithreading giving improved performance over a single-threaded solution.</p> <p>Ans</p> <ul style="list-style-type: none">• A Web server that services each request in a separate thread.• A parallelized application such as matrix multiplication where different parts of the matrix may be worked on in parallel.• An interactive GUI program such as a debugger where a thread is used to monitor user input, another thread represents the running application, and a third thread monitors performance.	2										
4	a)	<p>Does the following code satisfy the requirements for solution to critical section problem?</p> <pre>while(true) { while(lock!=0); lock=1; Critical section</pre>	4										

	<pre>lock=0; Remainder section }</pre> <p>If not, suggest the solution to critical section problem.</p> <p>Ans:</p> <p>Does not satisfy the requirements for critical section problem,-----1M</p> <p>Suggested solution</p> <p>Test and set must be atomic operation</p> <p>boolean test_and_set (boolean *lock)-----3M</p> <pre>{ boolean rv = *lock; *lock = TRUE; return rv; }</pre>											
b)	<p>What is the significance of wait-for graph and resource allocation graph? Give an example for each.</p> <p>Solution</p> <p>A wait-for graph is a directed graph used for deadlock detection in operating systems. IT is applicable for single instance resource type.</p> <p>Resource allocation graph: The resource allocation graph is a directed graph. the resource allocation graph gives the complete information about all the processes which are holding some resources or waiting for some resources</p> <div><div><p>(a)</p><p>Resource-Allocation Graph</p></div><div><p>(b)</p><p>wait for graph</p></div></div> <p>2M each</p>	4										
c)	<p>A counting semaphore S is initialized to 20. Then, 10 P operations and 6 V operations are performed on S. What is the final value of S?</p> <p>Ans:</p> <p>Final value of semaphore variable S</p> $= 20 - (10 \times 1) + (6 \times 1)$ $= 20 - 10 + 6$ <p>S=16-----2M</p>	2										
5	<p>a) Consider a program of size 26 bytes and a page size of 4 bytes. The physical address consists of 200 frames. The user program consists of 26 instructions a, b, c, . . . z. Assume the size of each instruction to be 1byte. Given a page table entry as shown below</p> <table><tr><th>Page number</th><th>Frame number</th></tr><tr><td>0</td><td>7</td></tr><tr><td>1</td><td>26</td></tr><tr><td>2</td><td>52</td></tr><tr><td>3</td><td>20</td></tr></table>	Page number	Frame number	0	7	1	26	2	52	3	20	5
Page number	Frame number											
0	7											
1	26											
2	52											
3	20											

4	55
5	6

Page table

i. What will be the number of pages for a given process?

No.of pages=26/4=6.5(7 pages)-----1M

ii. Find the physical addresses for the instructions a, e , i, o, u? ---2M

PA for a=7x4+0=28

PA for e=26x4+0=104

PA for i=52x4+0=208

PA for o=20x4+2=82

PA for u=6x4+0=24

iii. Calculate the fragmentation if exists? ----1m

Yes, internal fragment exists.

Size of fragment=4-2=2bytes

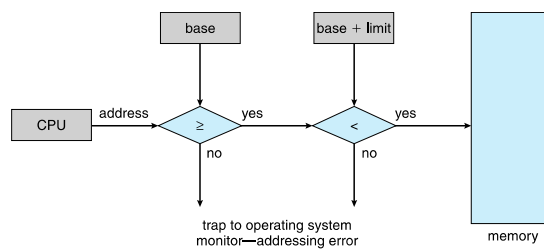
iv. Reference to which instruction causes page fault?-----1m

Reference to instruction y or z causes page fault

b) How are the processes protected from each other during their execution?

Ans:

Protection is provided by using two registers base and limit.---1M



The logical address generated must be greater than or equal to base register content and it must be less than base+limit. ----2M

c) Consider the following segment table:

Segment #	Base address	Length
0	219	600
1	2300	14
2	90	100
3	1327	580
4	1952	96

What are the physical addresses for the following logical addresses?

a. 0,430

Physical address for segment 0, byte 430=219+430=649

b. 1,10

Physical address for segment 1, byte 10=2300+10=2310

c. 2,500

Physical address for segment 2, byte 500 is invalid address

d. 3,400

Physical address for segment 3, byte 400=1327+400=1727

½ mark for each correct answer

6	<p>a) Consider the reference stream 1,4,3,4,2,1,5,6,2,1,2,3,7,6 for a particular process execution. How many page faults occur while using LRU page replacement algorithms with 4 frames?</p> <p>I. Calculate the number of page faults-----2M</p> <table><tr><td>1</td><td>4</td><td>3</td><td>4</td><td>2</td><td>1</td><td>5</td><td>6</td><td>2</td><td>1</td><td>2</td><td>3</td><td>7</td><td>6</td></tr><tr><td>1</td><td>1</td><td>1</td><td rowspan="4">hit</td><td>1</td><td rowspan="4">hit</td><td>1</td><td>1</td><td rowspan="4">hit</td><td rowspan="4">hit</td><td rowspan="4">hit</td><td>1</td><td>1</td><td>6</td></tr><tr><td></td><td>4</td><td>4</td><td>4</td><td>6</td><td>6</td><td>7</td><td>7</td></tr><tr><td></td><td></td><td>3</td><td>3</td><td>5</td><td>5</td><td>3</td><td>3</td><td>3</td></tr><tr><td></td><td></td><td></td><td>2</td><td>2</td><td>2</td><td>2</td><td>2</td></tr></table> <p>No.of page faults=9-----1M</p> <p>II. What are the data structures used to implement this algorithm?----1M Counter and stack is used in implementing LRU page replacement algorithm</p> <p>III. Does it suffer from Belady's anomaly? It does not suffer from Belady's anomaly --1M</p>	1	4	3	4	2	1	5	6	2	1	2	3	7	6	1	1	1	hit	1	hit	1	1	hit	hit	hit	1	1	6		4	4	4	6	6	7	7			3	3	5	5	3	3	3				2	2	2	2	2	5
1	4	3	4	2	1	5	6	2	1	2	3	7	6																																										
1	1	1	hit	1	hit	1	1	hit	hit	hit	1	1	6																																										
	4	4		4		6	6				7	7																																											
		3		3		5	5				3	3	3																																										
				2		2	2				2	2																																											
	<p>b) What is the significance of valid-invalid bits of page table in paging and in demand paging? Paging: Valid indicates that the associated page is in the logical address space. Invalid indicates that the associated page is not in logical address space. Demand Paging: a bit 1 in valid-invalid bit signifies that the page is in memory and 0 signifies that the page may be invalid or haven't brought into the memory just yet. 1M each</p>	2																																																					
	<p>c) What is thrashing? What are the methods to prevent thrashing? Ans: Thrashing is a condition or a situation when the system is spending a major portion of its time in servicing the page faults, but the actual processing done is very negligible --1M Methods to prevent thrashing----2M 1. Working set model 2. Page-fault frequency</p>	3																																																					