

DIGITAL DESIGN & COMPUTER ORGANISATION

Division

Sudarshan T S B., Ph.D.

Department of Computer Science & Engineering



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Course Outline



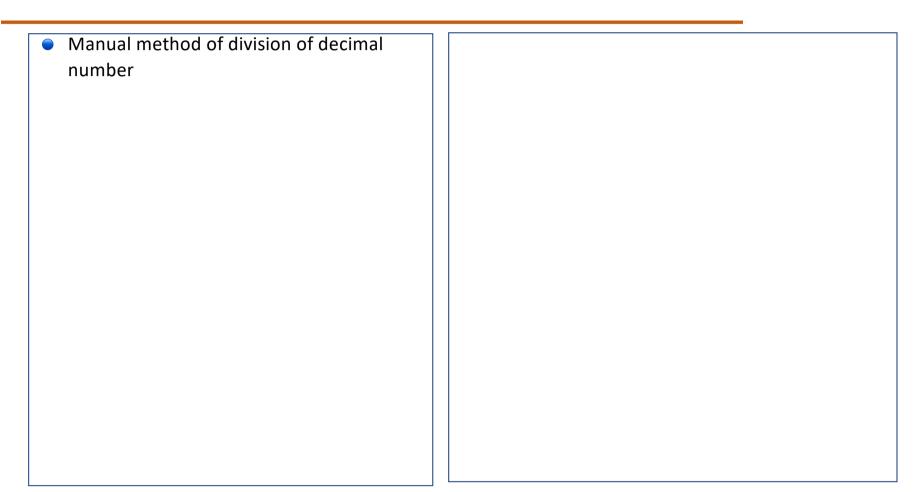
- Digital Design
 - Combinational logic design
 - Sequential logic design
 - ★ Divider
- Computer Organisation
 - Architecture (microprocessor instruction set)
 - Microarchitecture (microprocessor operation)

Concepts covered

Binary Integer Division



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1 Г	
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1 1	





- Manual method of division of decimal number
- Binary division also can be performed in a similar way



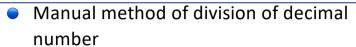
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```
(13) Divisor, M. => 1101
(274) Dividend, Q => 100010010
```



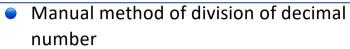


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```
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```

```
021
13) 274
26
14
13
```



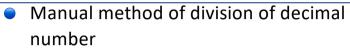


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```
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(274) Dividend, Q => 100010010
```

021	
13) 274	1101) 100010010
26	,
<u> 14</u>	
13	
1	



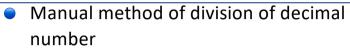


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```
(13) Divisor, M. => 1101
(274) Dividend, Q => 100010010
```

021	0
13) 274	1101) 100010010
26	,
<u> 14</u>	
13_	
1	



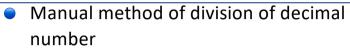


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021	00
13) 274	1101) 100010010
26	,
14	
13	
<u> </u>	



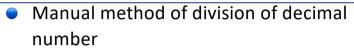


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021	000
13) 274	1101) 100010010
26	,
14	
13	



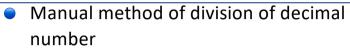


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021	0000
13) 274	1101) 100010010
26	,
14	
13	
<u> </u>	



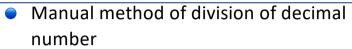


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```
(13) Divisor, M. => 1101
(274) Dividend, Q => 100010010
```

021	00001
13) 274	1101) 100010010
26	,
14	
13	
<u> </u>	

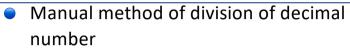




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021	00001
13) 274	1101) 100010010
26	1101
14	
13	
<u> </u>	

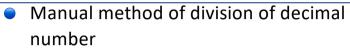




- Binary division also can be performed in a similar way
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021	00001
13) 274	1101) 100010010
_26	1101
14	100
13	
1	

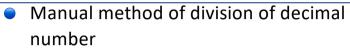




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021	00001
13) 274	1101) 100010010
_26	1101
14	100
13	
1	

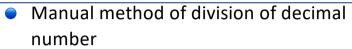




- Binary division also can be performed in a similar way
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021	00001
13) 274	1101) 100010010
26	1101
14	1000
13	
1	

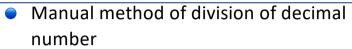




- Binary division also can be performed in a similar way
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021	000010
13) 274	1101) 100010010
_26	1101
14	1000
13_	

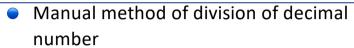




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021	000010
13) 274	1101) 100010010
_26	1101
14	1000
13	
1	

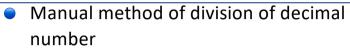




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021	000010
13) 274	1101) 100010010
26	1101
14	10000
13	
1	

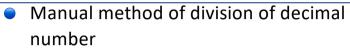




- Binary division also can be performed in a similar way
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021	0000101
13) 274	1101) 100010010
26	1101
14	10000
13	
1	

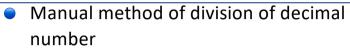




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021	0000101
13) 274	1101) 100010010
26	1101↓↓
14	10000
13_	1101
1	

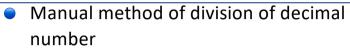




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021	0000101
13) 274	1101) 100010010
_26	1101 📗
14	10000
13_	1101
1	11

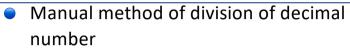




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021	0000101
13) 274	1101) 100010010
_26	1101 📗
14	10000
13_	1101
1	11

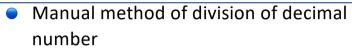




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021	0000101
13) 274	1101) 100010010
_26	1101 📗
14	10000
13_	1101
1	111

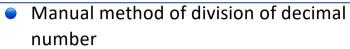




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021	00001010
13) 274	1101) 100010010
26	1101 📗
14	10000
13_	1101
1	111

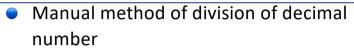




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13) 274	1101) 100010010
26	1101
14	10000
13	1101
1	111

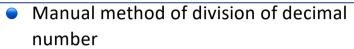




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021	00001010
13) 274	1101) 100010010
26	1101
14	10000
13_	1101
1	1110

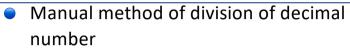




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021	000010101
13) 274	1101) 100010010
_26	1101
14	10000
13_	1101
1	1110

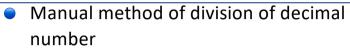




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14	10000
13_	1101
1	1110
	1101



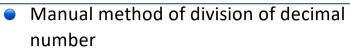


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021_	000010101
13) 274	1101) 100010010
_26	1101
14	10000
13_	1101
1	1110
	1101
	1



Long Hand Method



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021	000010101
13) 274	1101) 100010010
_26	1101
14	10000
13_	1101∐
1	1110
	1101
	1

Quotient = 21 => 10101; Remainder = 1



Long Hand Method

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 A subtraction of divisor is performed with dividend



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021	000010101
13) 274	1101) 100010010
_26	1101
14	10000
13_	1101
1	1110
	1101
	<u></u>

Quotient = 21 => 10101; Remainder = 1

- A subtraction of divisor is performed with dividend
- If the remainder is zero or positive quotient bit is 1, the remainder is extended by another bit of the dividend to repeat the process



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021	000010101
13) 274	1101) 100010010
26	1101
14	10000
13	1101
1	1110
	1101
	1

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- Repeat the process until all the digits in the dividend are considered



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- Repeat the process until all the digits in the dividend are considered
- This is called Restoring Division algorithm



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- A subtraction of divisor is performed with dividend
- If the remainder is zero or positive quotient bit is 1, the remainder is extended by another bit of the dividend to repeat the process
- If the remainder is negative, quotient is 0 the dividend is restored by adding back the divisor.
- Repeat the process until all the digits in the dividend are considered
- This is called Restoring Division algorithm
- The other method to be familiar with is Non-Restoring Division



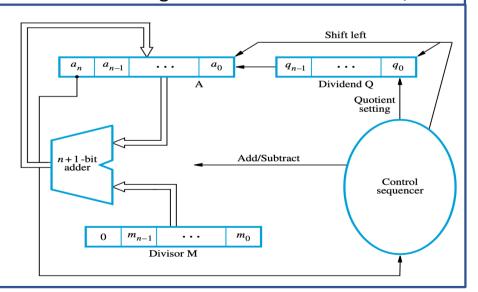
DIVISION Restoring Division



- We Require:
 - Accumulator register A (m+1) bits (Remainder)
 - ▶ Dividend register Q (Quotient) n-bits
 - Divisor register M (m+1 bit => Sign)
 - ► n+1-bit Adder
 - Control signals for Shift left and Add/Sub

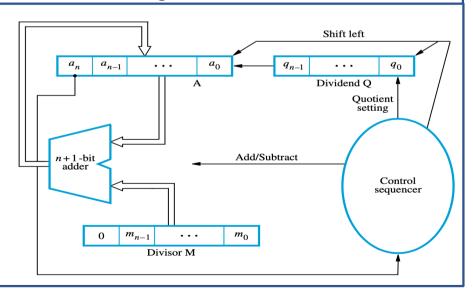


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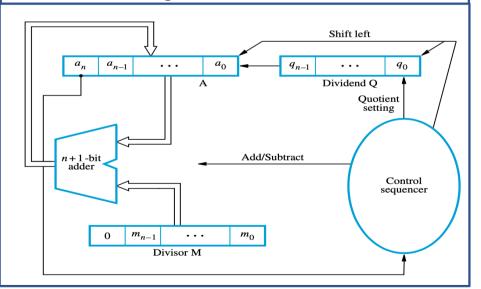


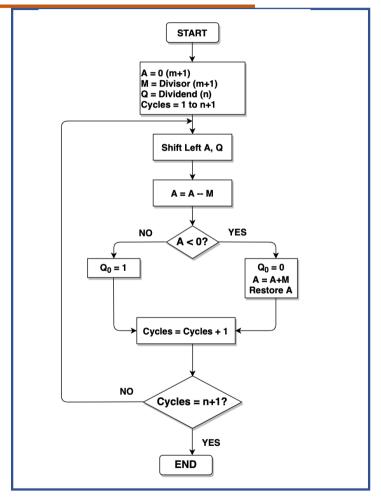
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Restoring Division

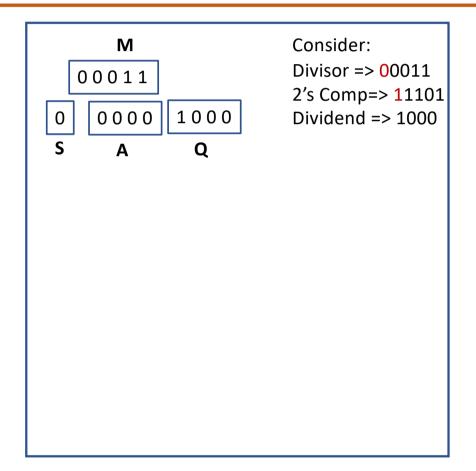


Consider:

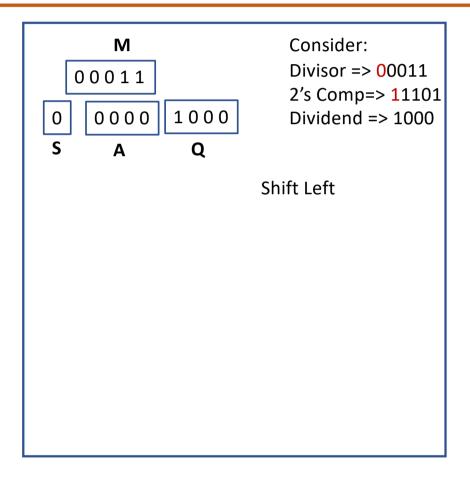
Divisor => **0**0011

2's Comp=> 11101

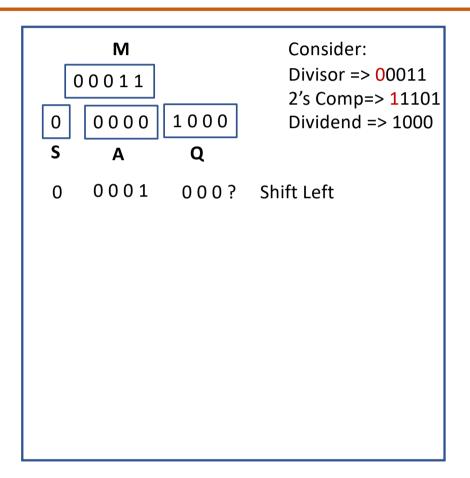
Dividend => 1000



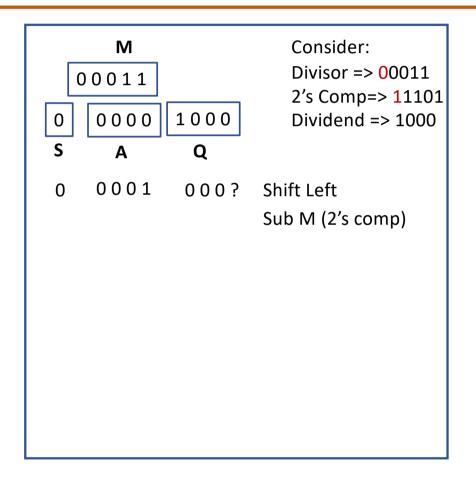




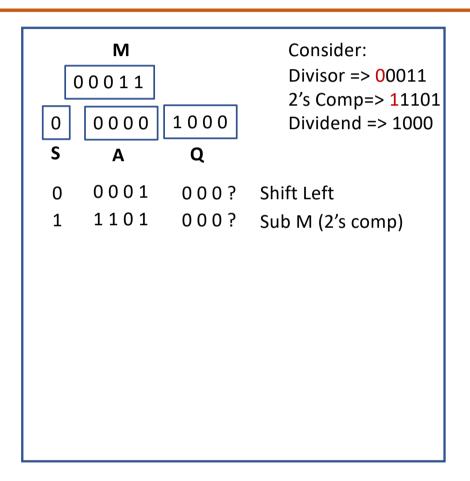




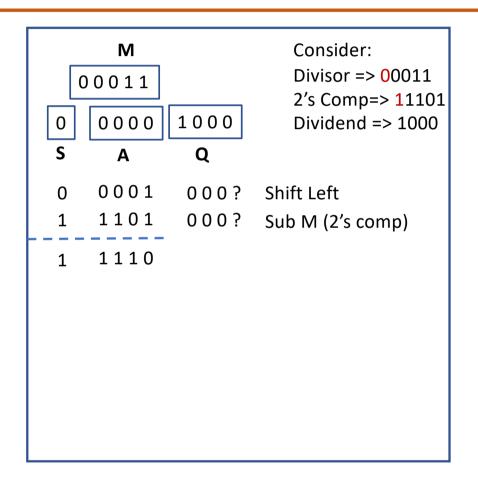




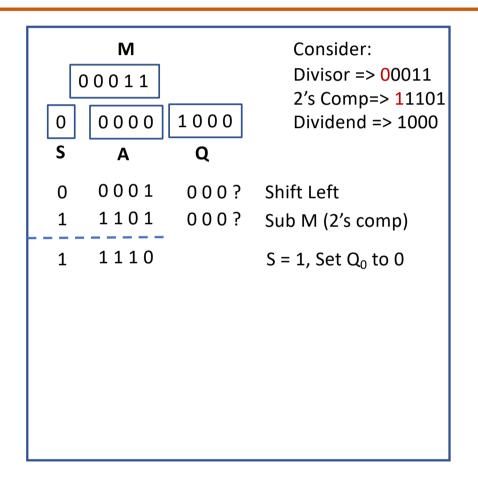




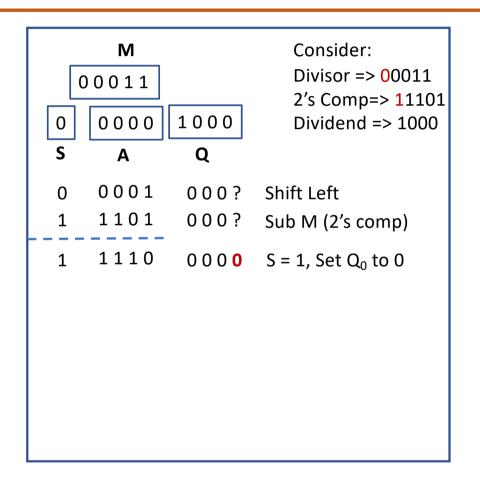




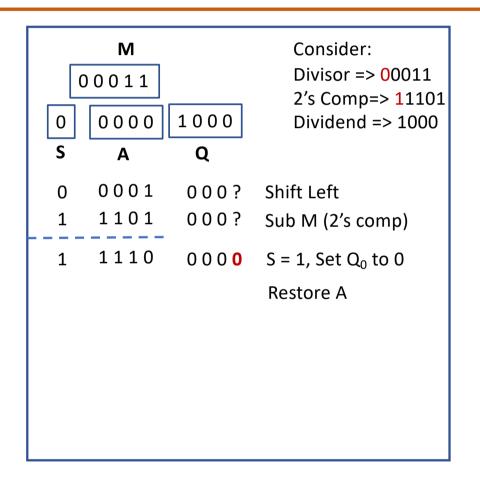




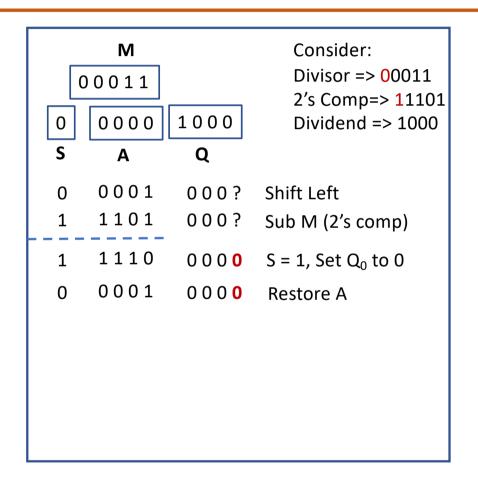




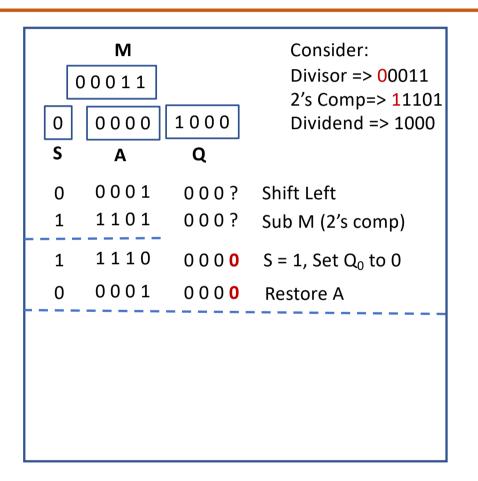




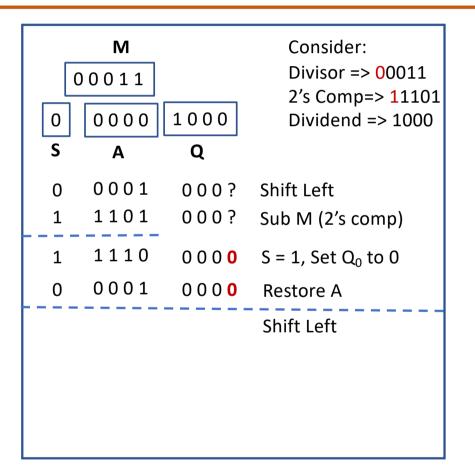




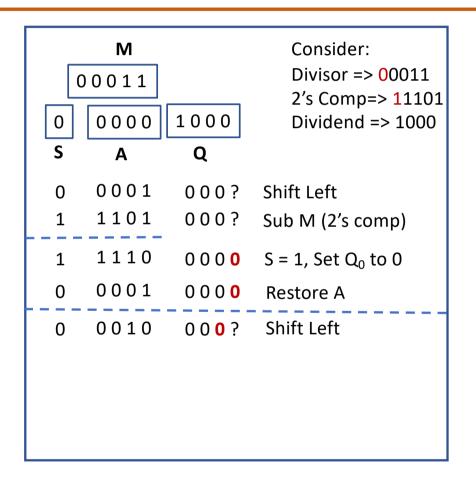




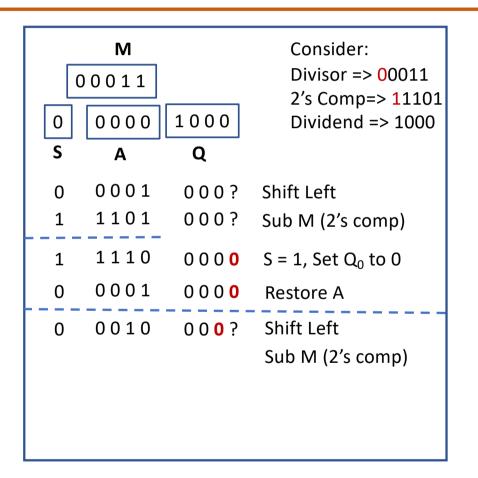




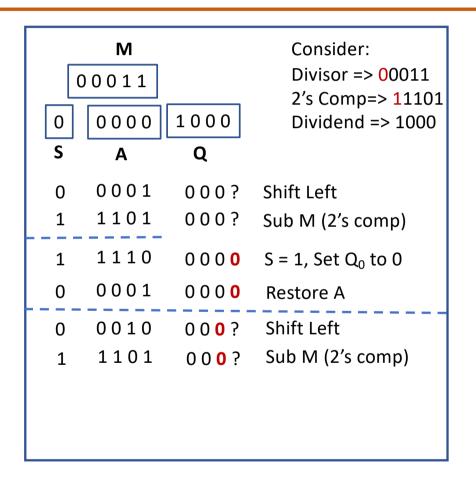




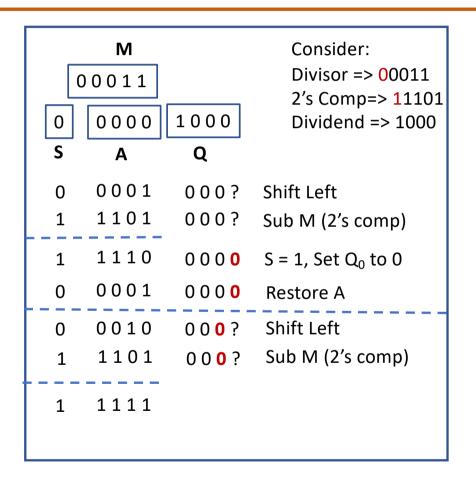




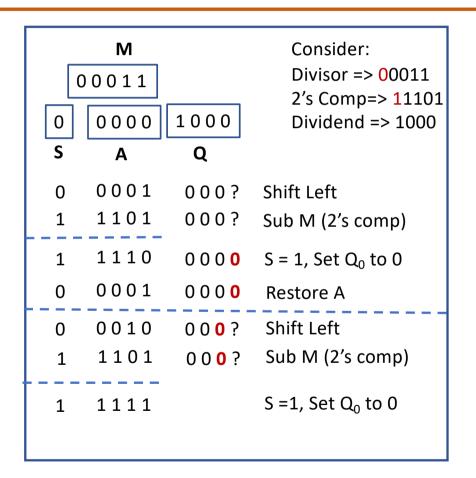




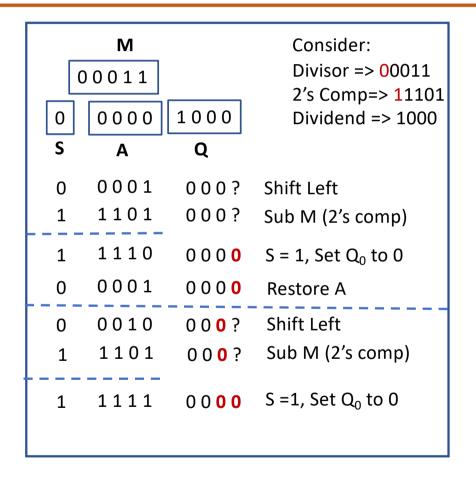








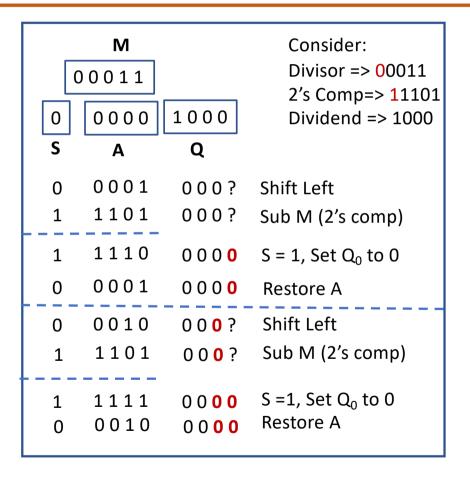




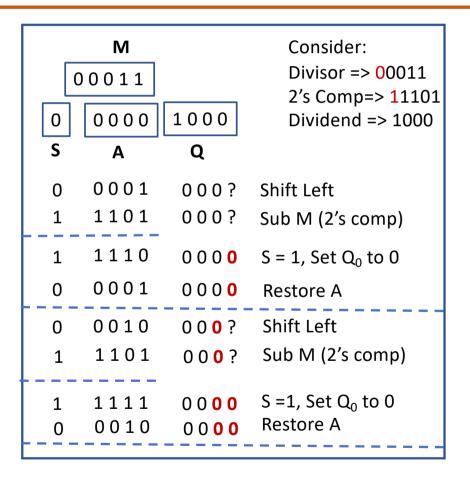


	М		Consider:
Ĺ	00011		Divisor => 00011 2's Comp=> 11101
0	0000	1000	Dividend => 1000
S	Α	Q	
0	0001	000?	Shift Left
1	1101	000?	Sub M (2's comp)
1	1110	0000	S = 1, Set Q ₀ to 0
0	0001	0000	Restore A
0	0010	000?	Shift Left
1	1101	000?	Sub M (2's comp)
1	1111	0000	S =1, Set Q ₀ to 0 Restore A







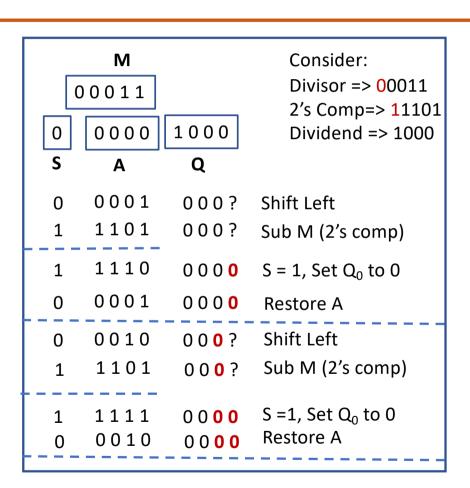


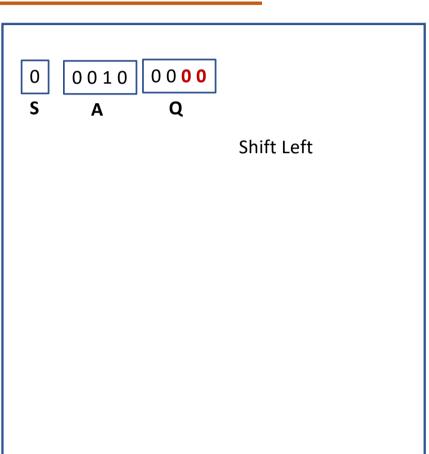




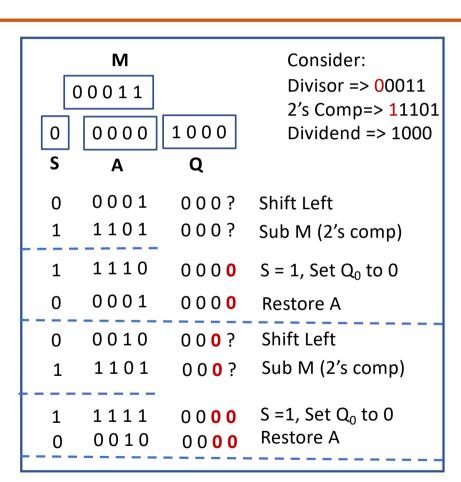
```
Consider:
      M
                        Divisor => 00011
  00011
                        2's Comp=> 11101
            1000
0
    0000
                        Dividend => 1000
S
              Q
      Α
    0001
             000?
                     Shift Left
0
    1101
             000?
                     Sub M (2's comp)
1
    1110
             0000
                     S = 1, Set Q_0 to 0
1
    0001
             0000
0
                     Restore A
    0010
             000?
                     Shift Left
0
    1101
                     Sub M (2's comp)
             000?
                     S = 1, Set Q_0 to 0
1
    1111
             0000
                     Restore A
    0010
             0000
0
```

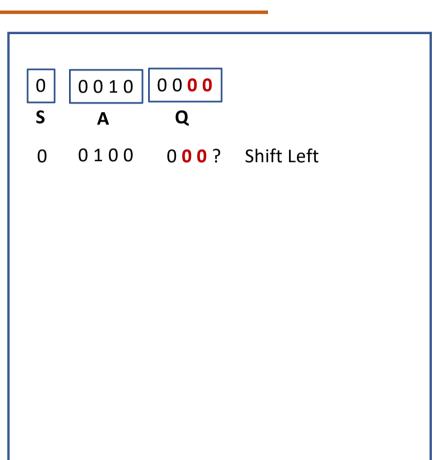
```
0000
   0010
0
           Q
    Α
```









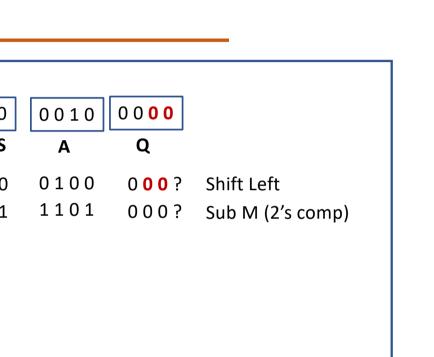






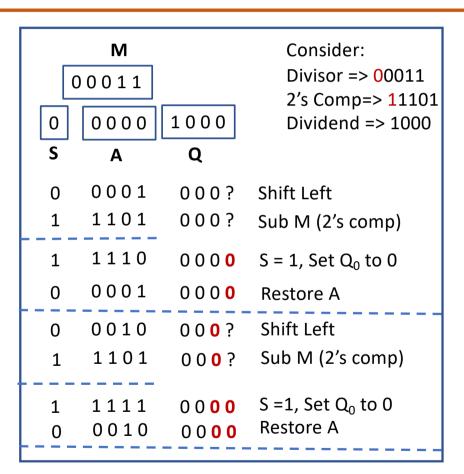
```
Consider:
      M
                        Divisor => 00011
  00011
                        2's Comp=> 11101
0
    0000
            1000
                        Dividend => 1000
S
              Q
      Α
    0001
             000?
                     Shift Left
0
    1101
             000?
                     Sub M (2's comp)
1
    1110
             0000
1
                     S = 1, Set Q_0 to 0
    0001
             0000
0
                     Restore A
    0010
             000?
                     Shift Left
0
    1101
                     Sub M (2's comp)
             000?
                     S = 1, Set Q_0 to 0
    1111
             0000
1
                     Restore A
    0010
             0000
0
```

```
0000
    0010
0
             Q
     Α
    0100
            000?
                    Shift Left
0
                    Sub M (2's comp)
```



	M		Consider: Divisor => 00011 2's Comp=> 11101
0	0000	1000	Dividend => 1000
S	A	Q	
0	0001	000?	Shift Left
1	1101	000?	Sub M (2's comp)
1	1110	0000	$S = 1$, Set Q_0 to 0
0	0001	0000	Restore A
0	0010	000?	Shift Left
1	1101	000?	Sub M (2's comp)
1 0	1111	0000	S =1, Set Q ₀ to 0 Restore A





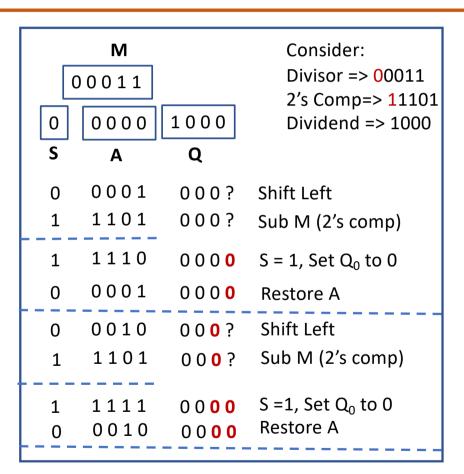
```
0 0010 0000

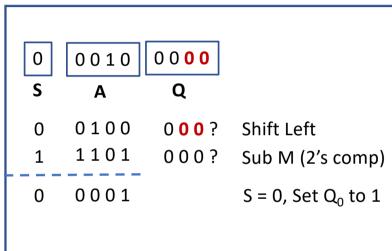
S A Q

0 0100 000? Shift Left
1 1101 000? Sub M (2's comp)

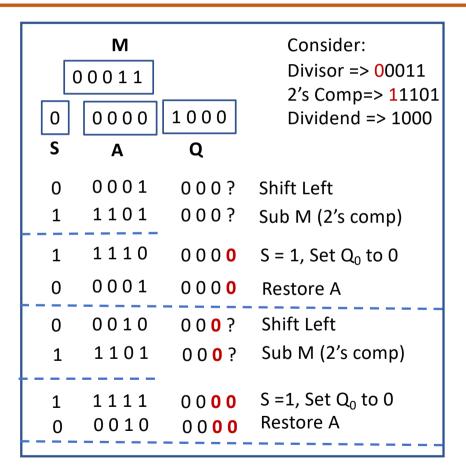
0 0001
```

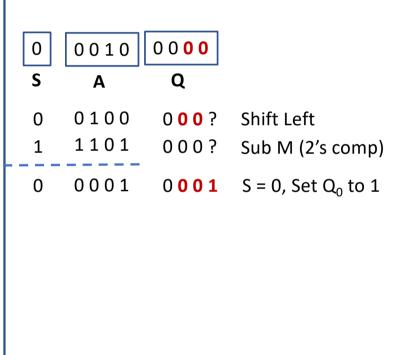
















M			Consider: Divisor => 00011	
0	0000	1000	2's Comp=> 1 1101 Dividend => 1000	
S	Α	Q		
0	0001	000?	Shift Left	
1	1101	000?	Sub M (2's comp)	
1	1110	0000	S = 1, Set Q ₀ to 0	
0	0001	0000	Restore A	
0	0010	000?	Shift Left	
1	1101	000?	Sub M (2's comp)	
1 0	1111 0010	0000	S =1, Set Q ₀ to 0 Restore A	

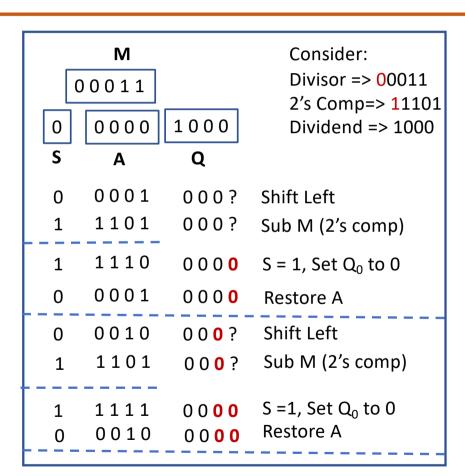
```
0 0010 0000

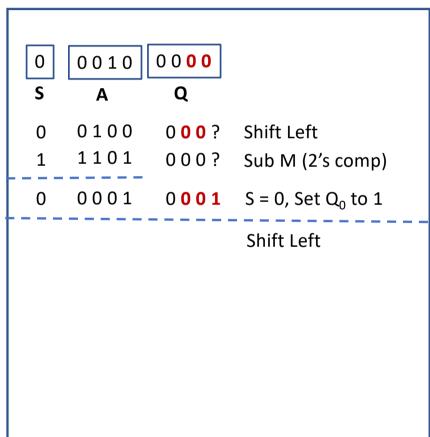
S A Q

0 0100 000? Shift Left

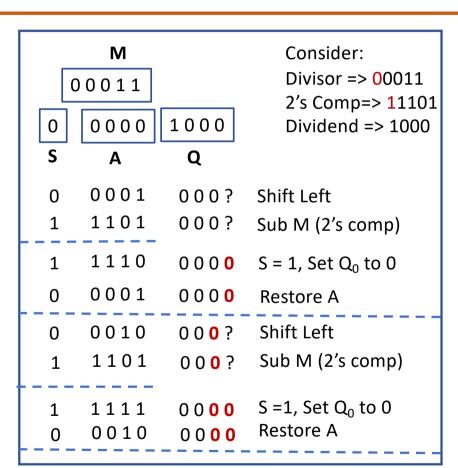
1 1101 000? Sub M (2's comp)

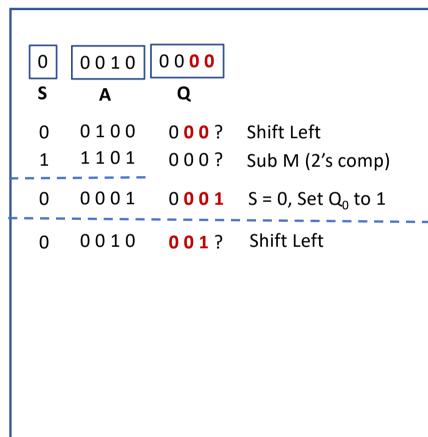
0 0001 0001 S = 0, Set Q<sub>0</sub> to 1
```



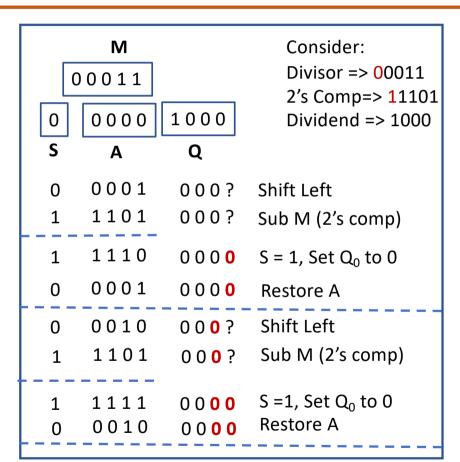


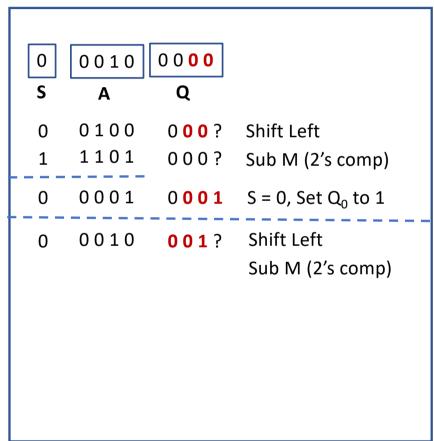




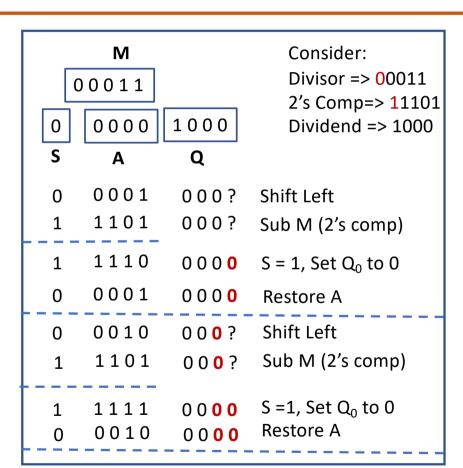






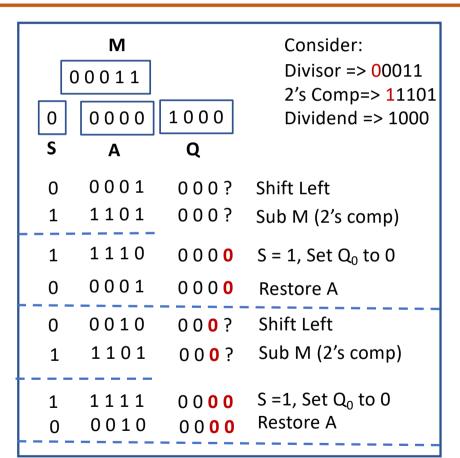


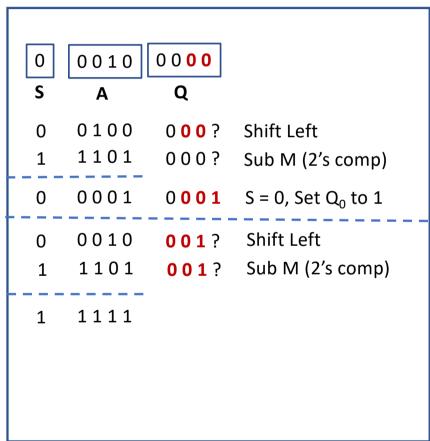




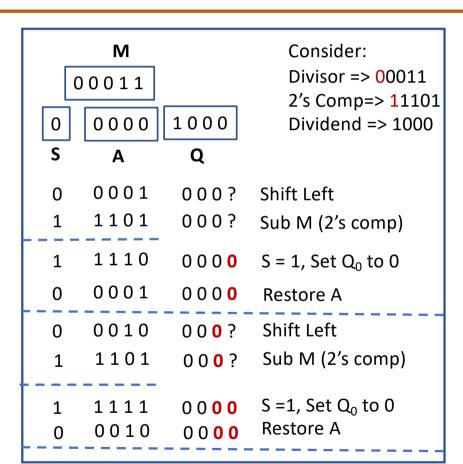
0	0010	0000	
S	Α	Q	
0	0100	000?	Shift Left
1	1101	000?	Sub M (2's comp)
0	0001	0001	S = 0, Set Q ₀ to 1
0	0010	001?	Shift Left
1	1101	001?	Sub M (2's comp)





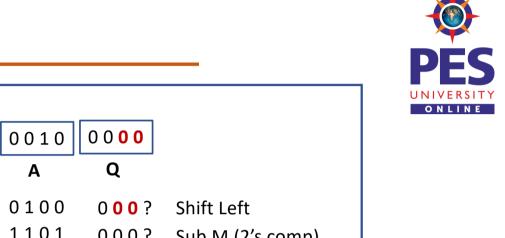


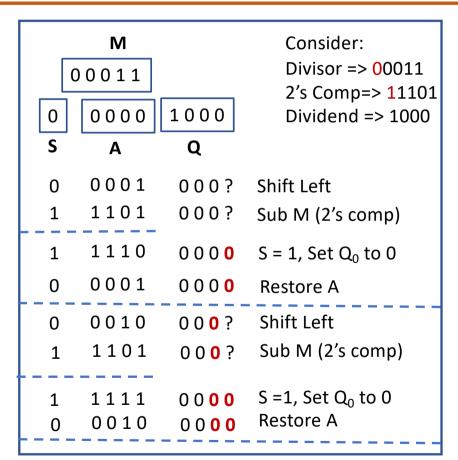




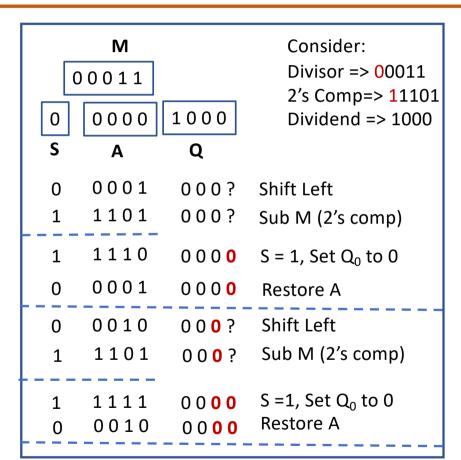
0	0010	0000	
S	Α	Q	
0	0100	000?	Shift Left
1	1101	000?	Sub M (2's comp)
0	0001	0001	S = 0, Set Q ₀ to 1
0	0010	001?	Shift Left
1	1101	001?	Sub M (2's comp)
1	1111		S =1, Set Q ₀ to 0

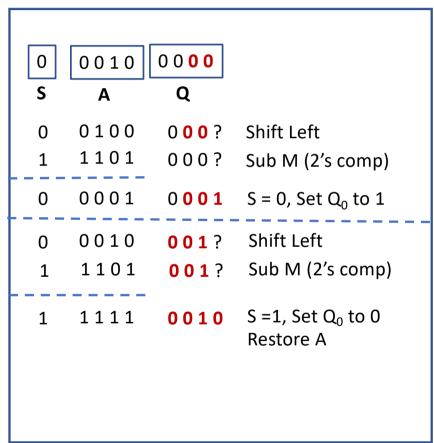




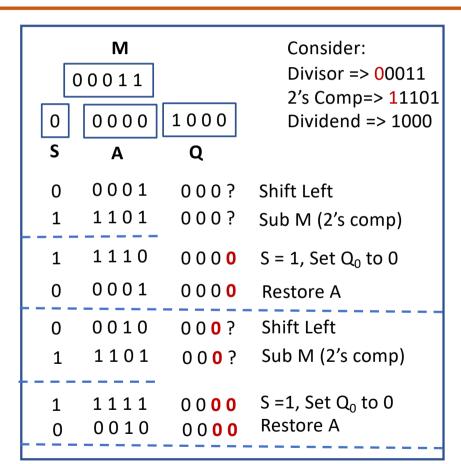


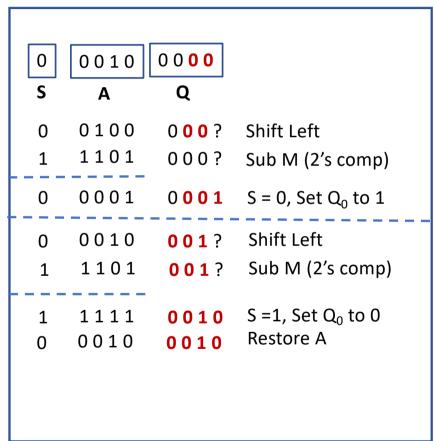
0	0010	0000	
S	Α	Q	
0	0100	000?	Shift Left
1	1101	000?	Sub M (2's comp)
0	0001	0001	S = 0, Set Q ₀ to 1
0	0010	001?	Shift Left
1	1101	001?	Sub M (2's comp)
1	1111	0010	S = 1, Set Q_0 to 0



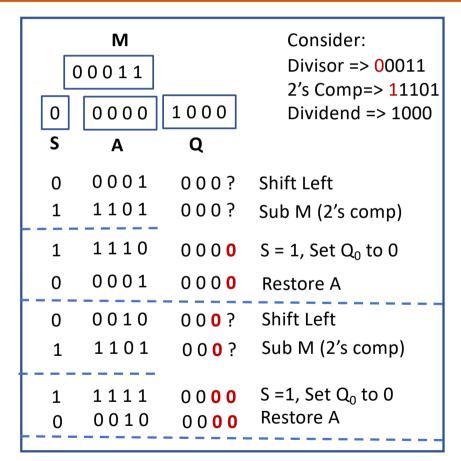


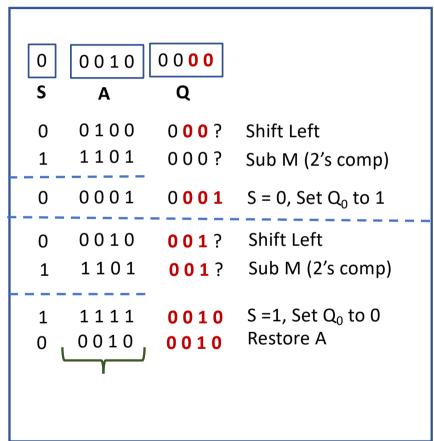




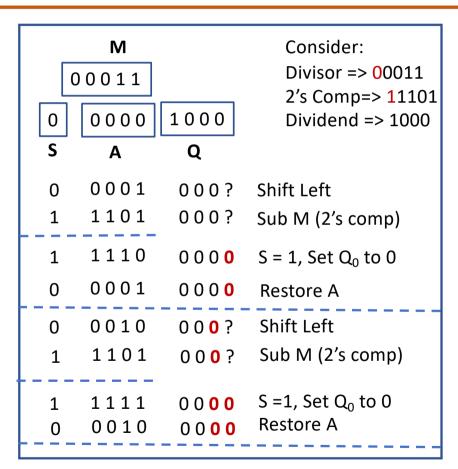


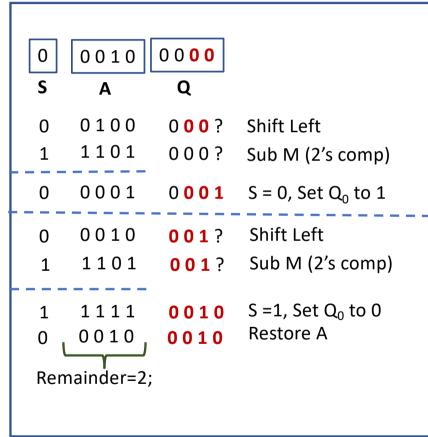




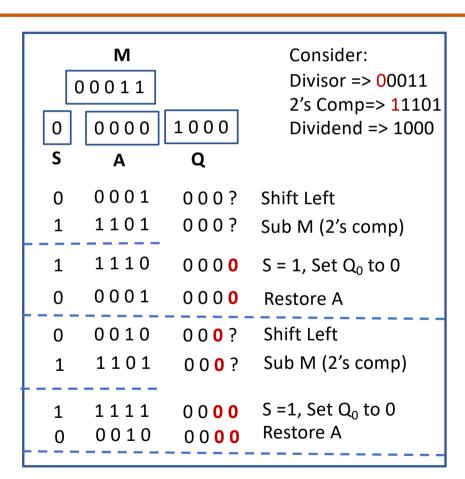


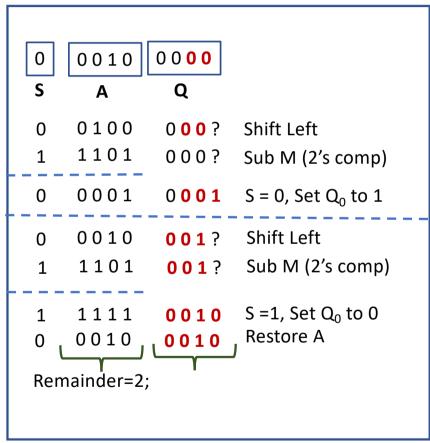




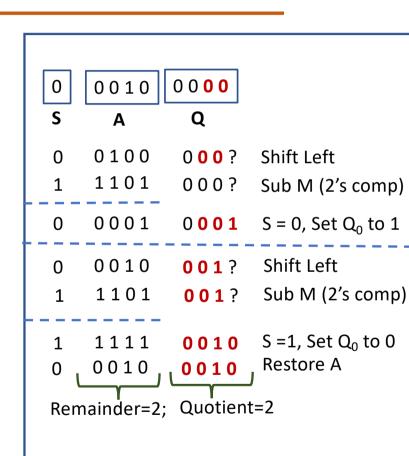






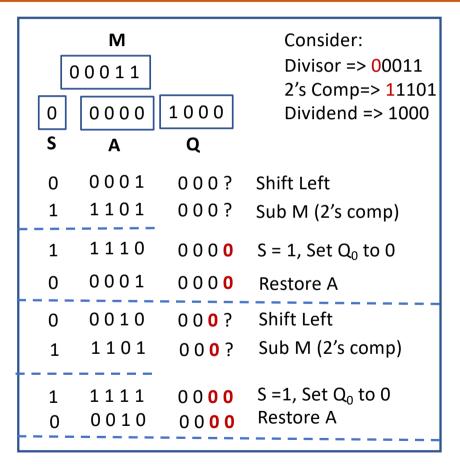


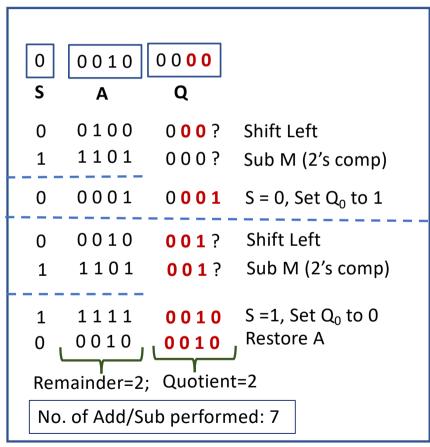






```
M
                        Consider:
                        Divisor => 00011
  00011
                        2's Comp=> 11101
    0000
            1000
0
                        Dividend => 1000
S
      Α
              Q
    0001
             000?
                     Shift Left
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    1101
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    0001
             0000
                      Restore A
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    0010
             000?
                      Shift Left
0
                      Sub M (2's comp)
    1101
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    1111
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                     S = 1, Set Q_0 to 0
1
                     Restore A
    0010
             0000
0
```







DIVISION Restoring Division





Observations:		

Restoring Division



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• This method is used for unsigned binary numbers division or positive signed binary numbers division

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- 2A+M, represents shift left and add M, when A is negative.

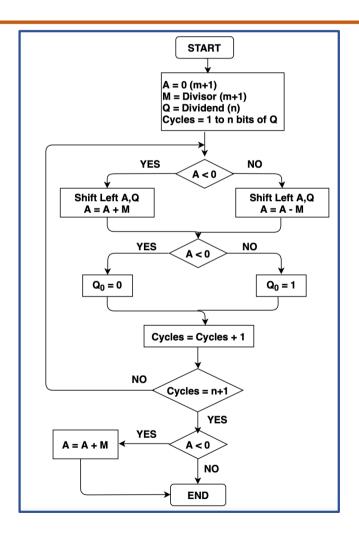


Restoring Division

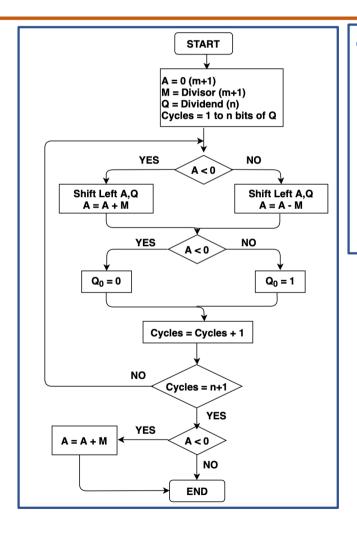
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- 2A+M, represents shift left and add M, when A is negative.
- This leads to Non-Restoring division algorithm









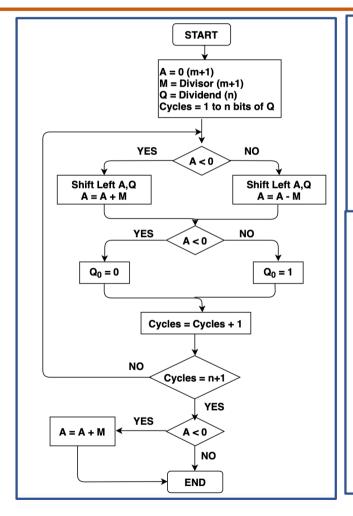




- Accumulator register A (m+1) bits (Remainder)
- ▶ Dividend register Q (Quotient) n-bits
- Divisor register M (m+1 bit => Sign)
- ▶ n+1-bit Adder
- Control signals for Shift left and Add/Sub

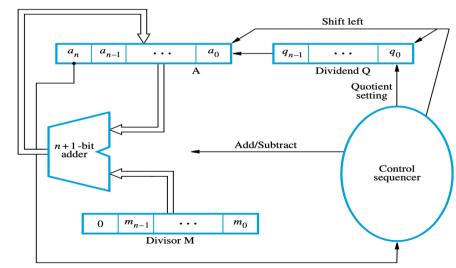


Non-Restoring Division



We Require:

- Accumulator register A (m+1) bits (Remainder)
- ▶ Dividend register Q (Quotient) n-bits
- Divisor register M (m+1 bit => Sign)
- ▶ n+1-bit Adder
- Control signals for Shift left and Add/Sub





Non-Restoring Division

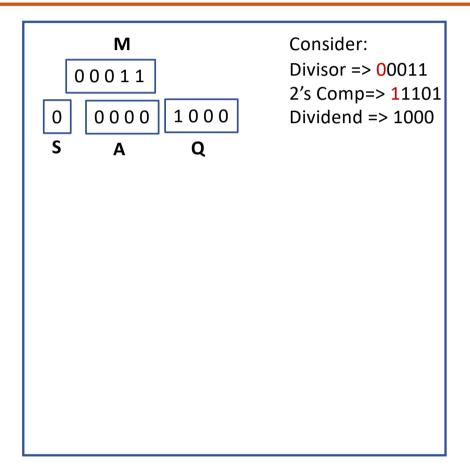


Consider:

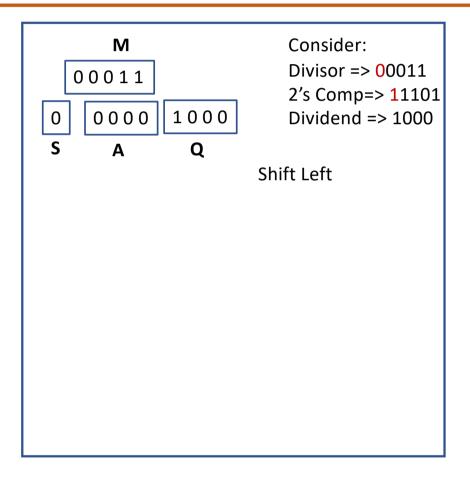
Divisor => **0**0011

2's Comp=> 11101

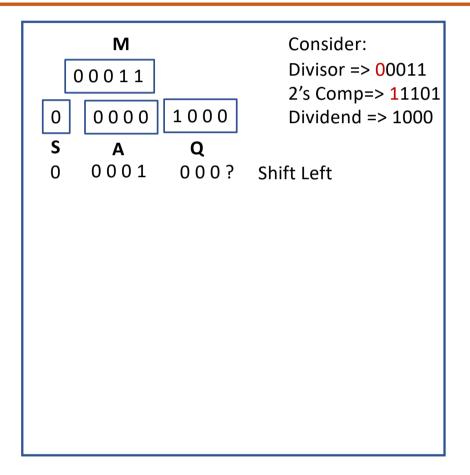
Dividend => 1000



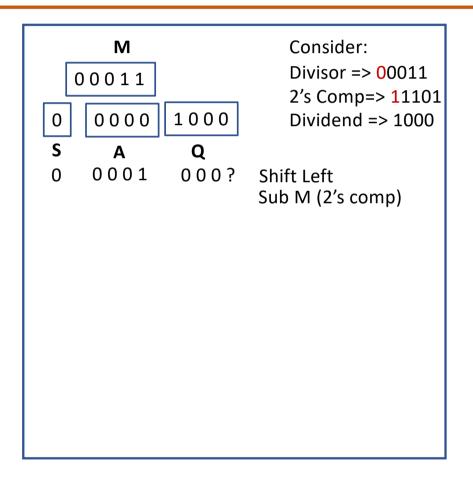




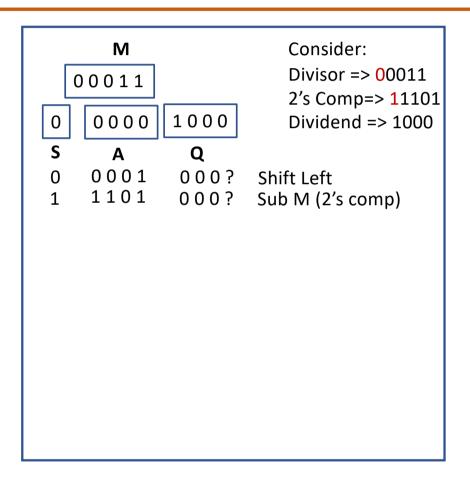




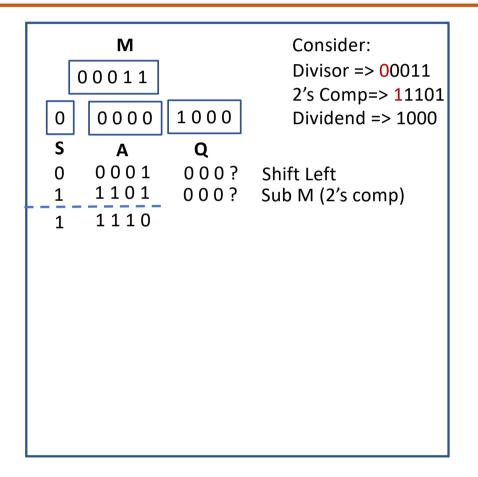




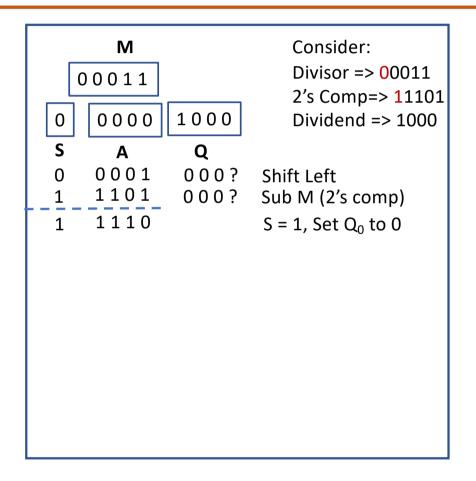




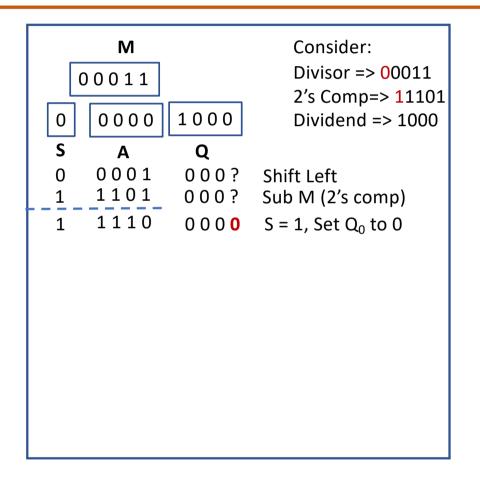




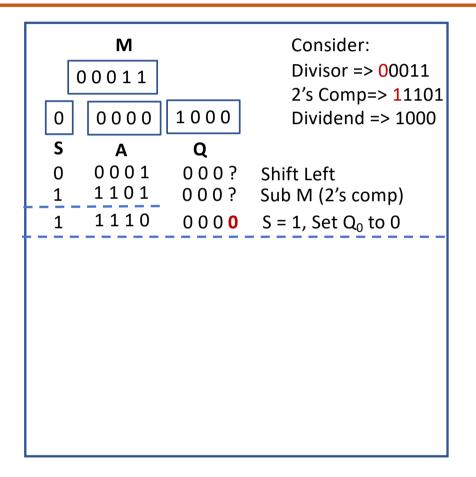




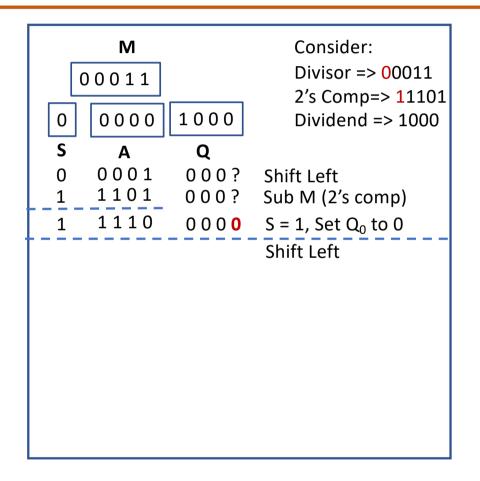




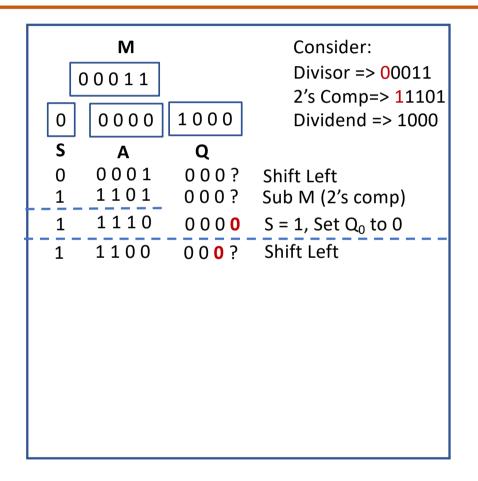




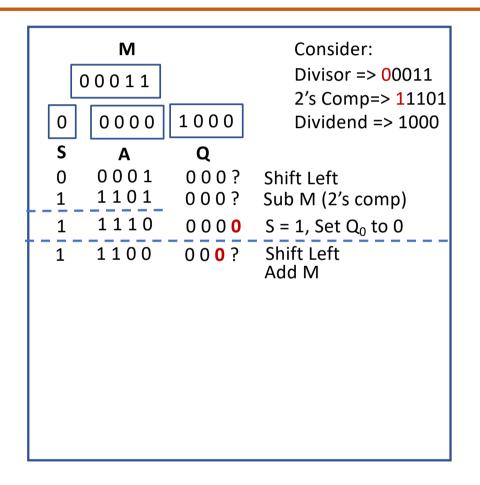




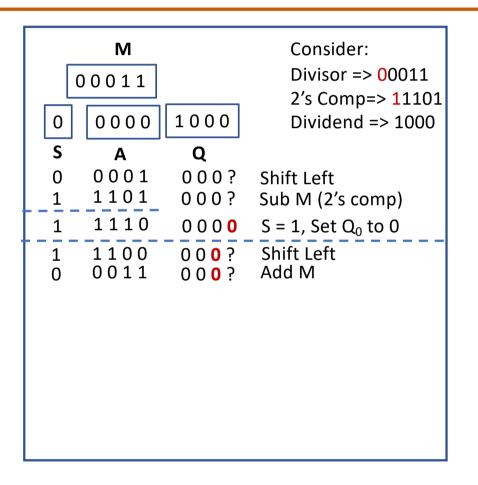




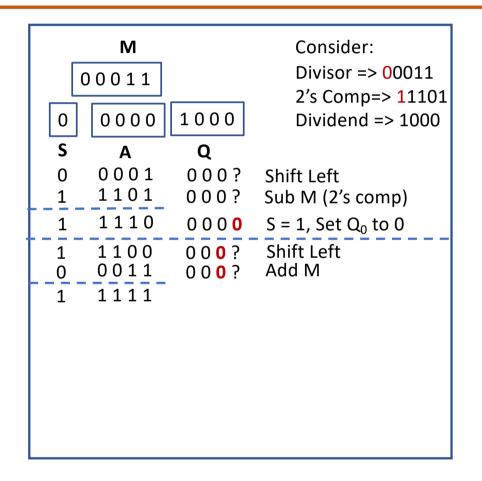




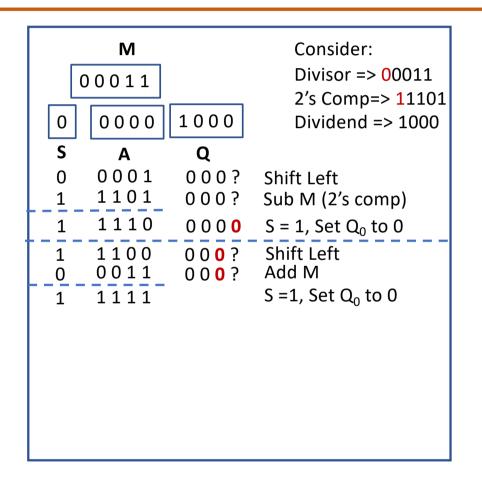




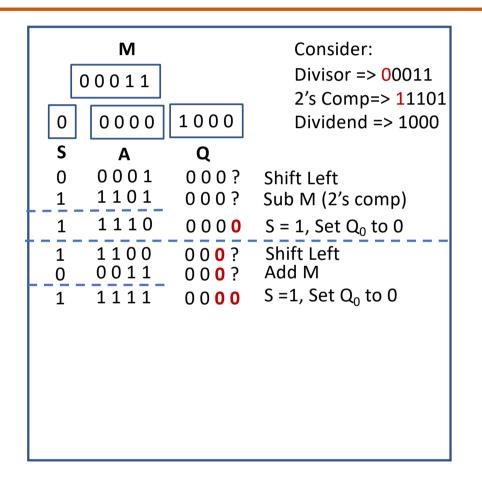




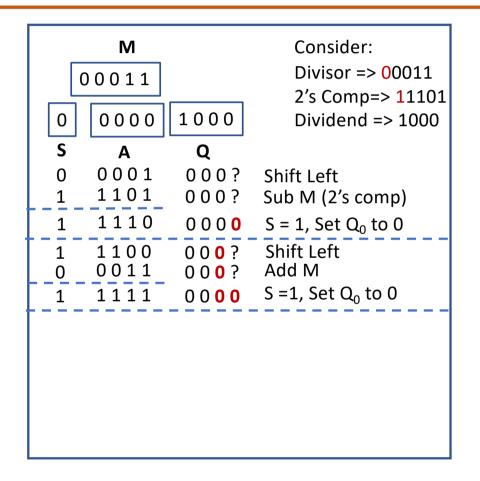




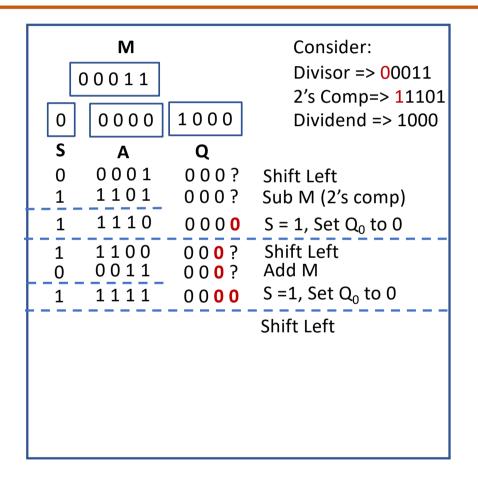




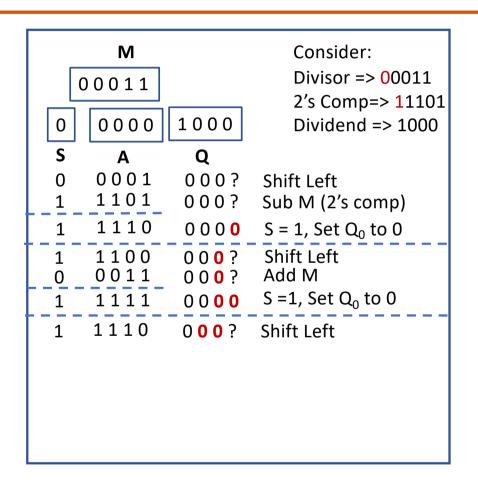




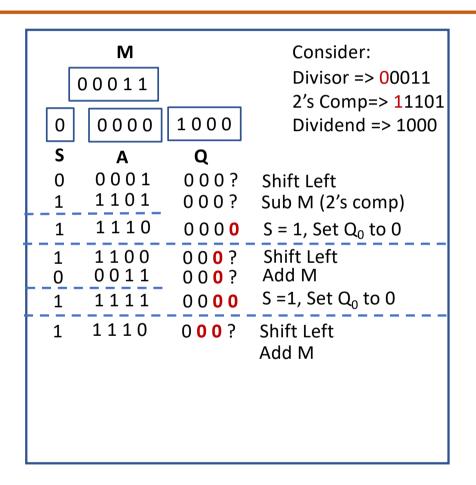




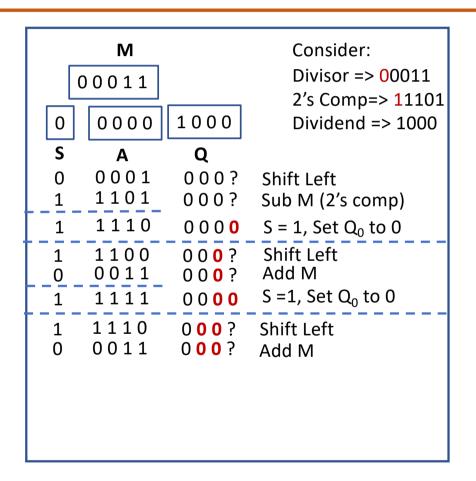




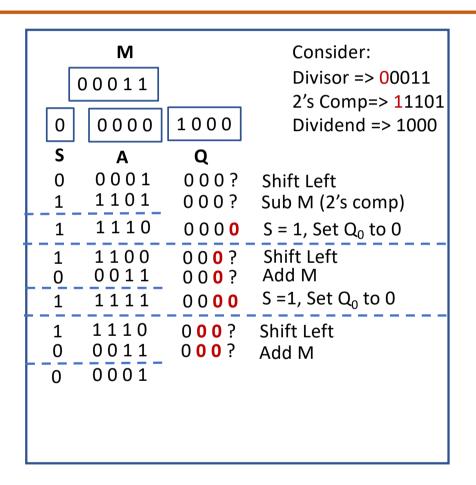




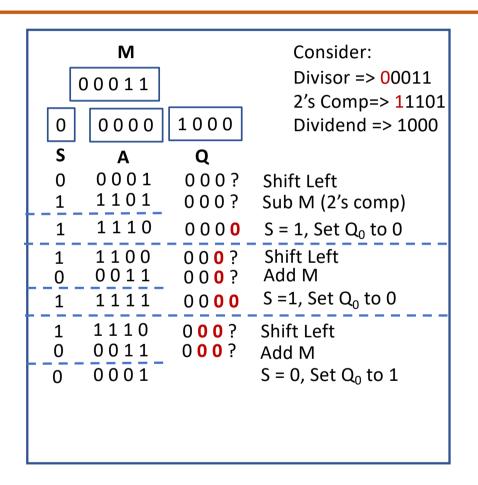




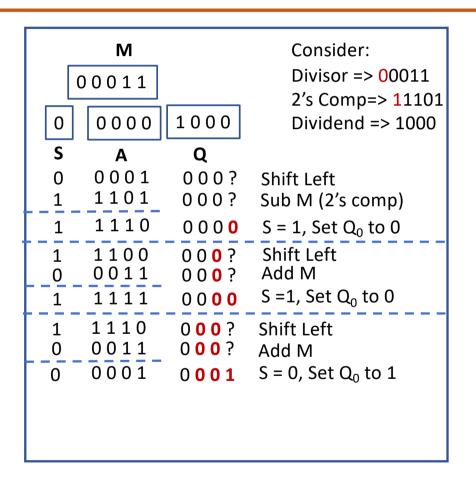




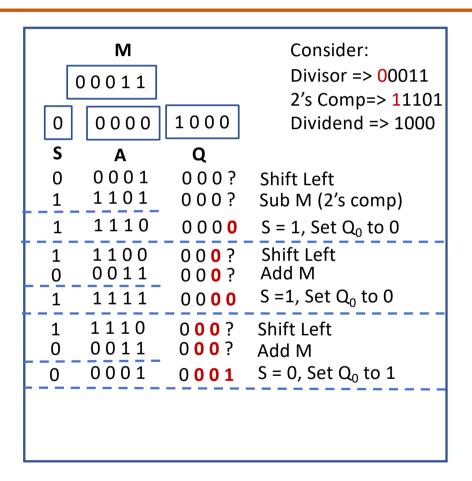




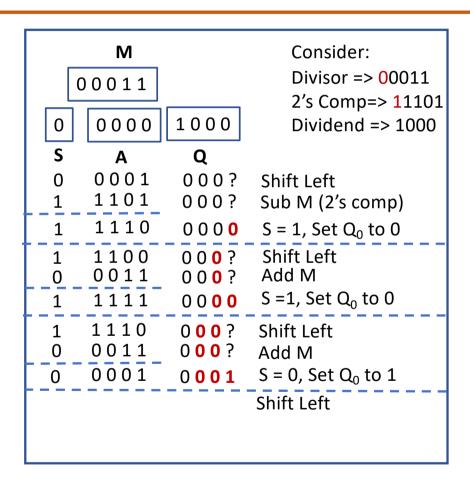




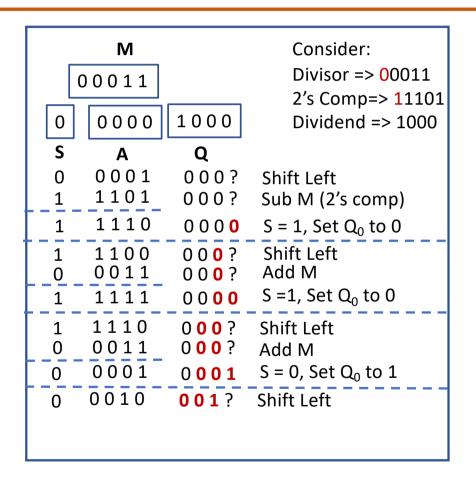












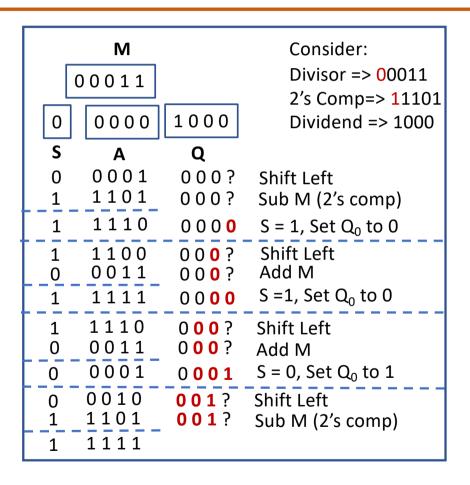


	М		Consider:
	00011		Divisor => 00011
<u></u> '			2's Comp=> 1 1101
0	0000	1000	Dividend => 1000
S	Α	Q	
0	0001	000?	Shift Left
_ 1	1101	000?	Sub M (2's comp)
_ 1	1110	0000	$S = 1$, Set Q_0 to 0
1	1100	000?	Shift Left
_ 0	0011	000?	Add M
1	1111	0000	$S = 1$, $Set Q_0 to 0$
1	1110	000?	Shift Left
0	0011	0003	Add M
0	0001	0001	$S = 0$, Set Q_0 to 1
0	0010	001?	Shift Left
			Sub M (2's comp)

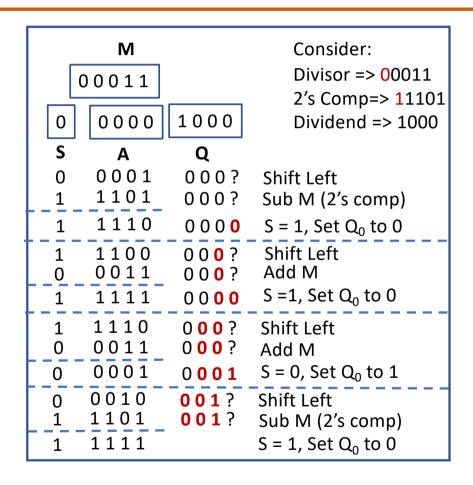


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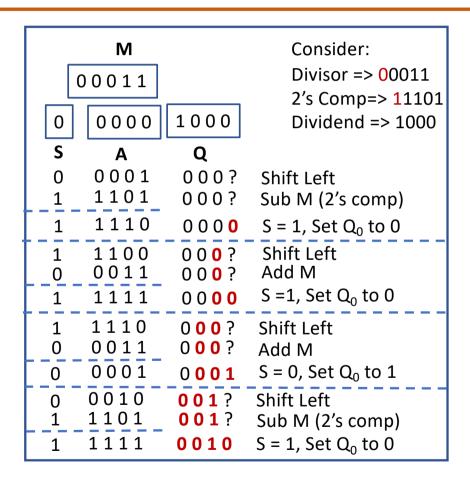








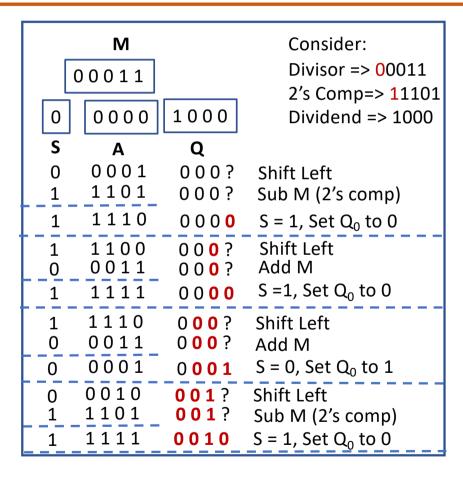


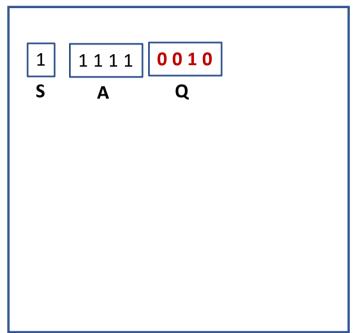




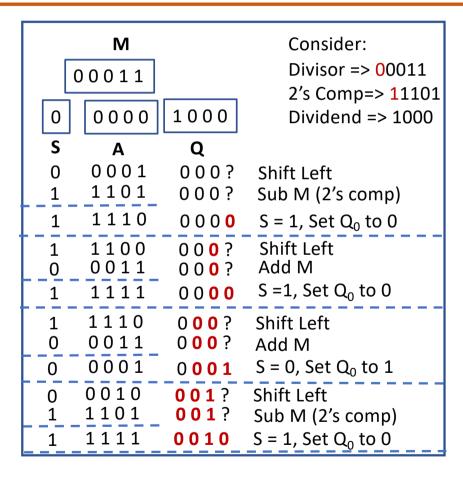
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0	0010	001?	Shift Left
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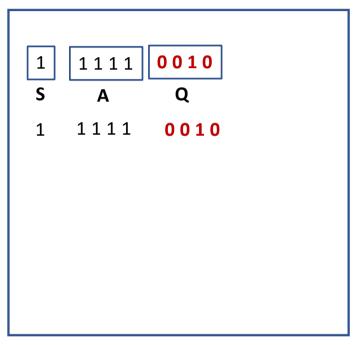




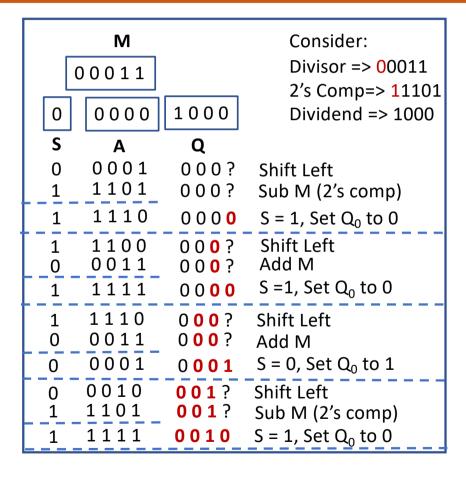


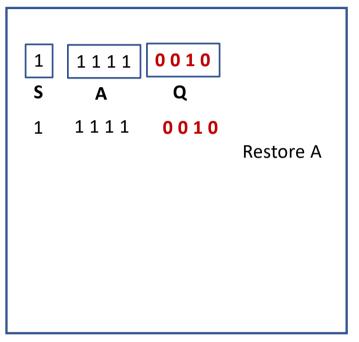




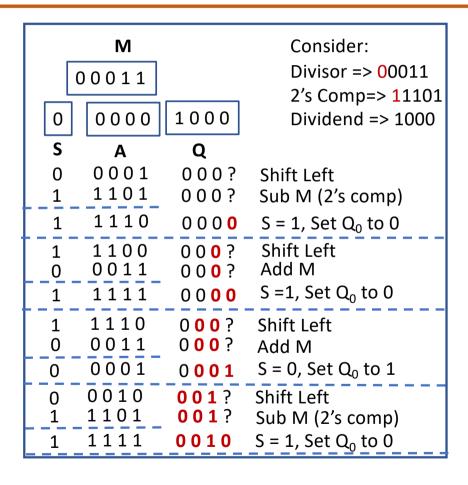


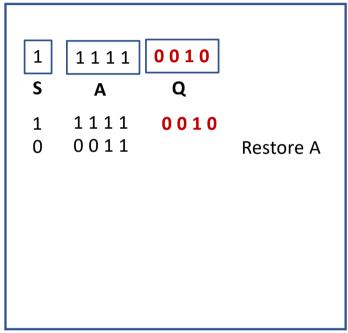




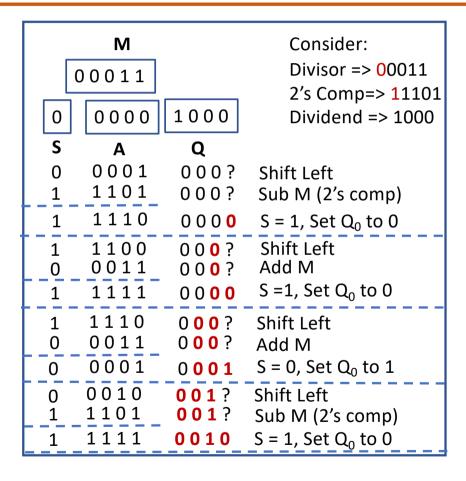


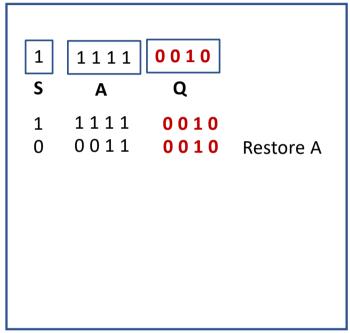




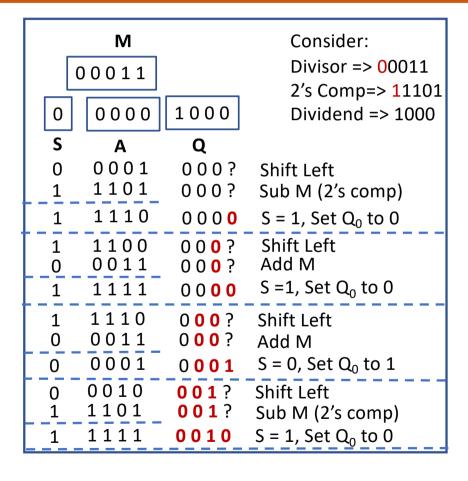


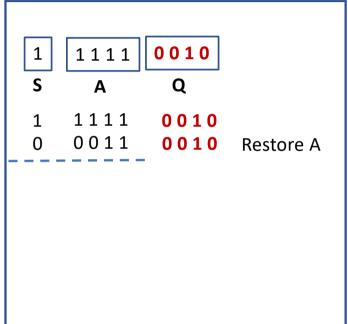




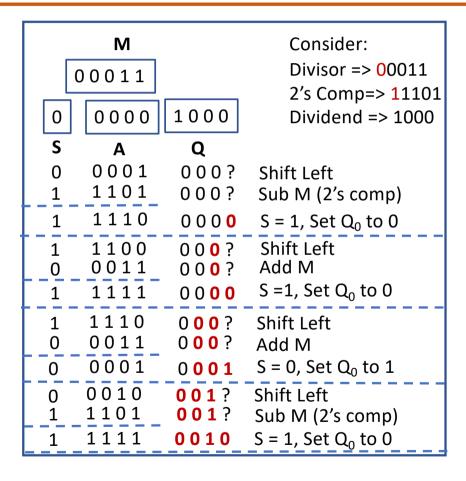


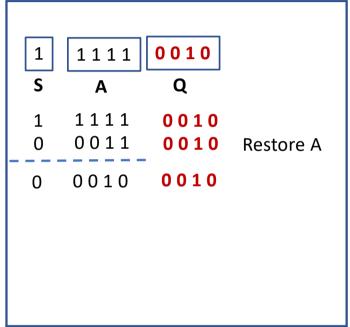




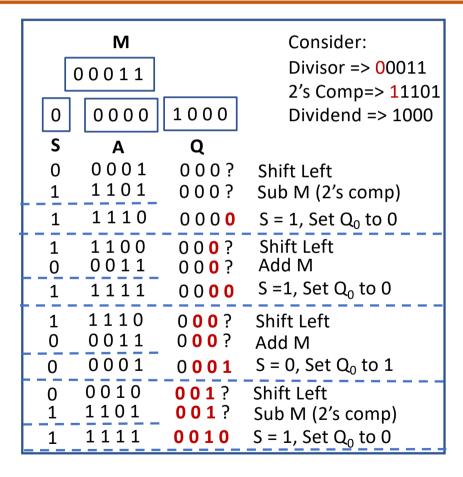


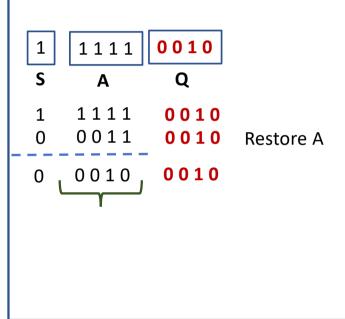




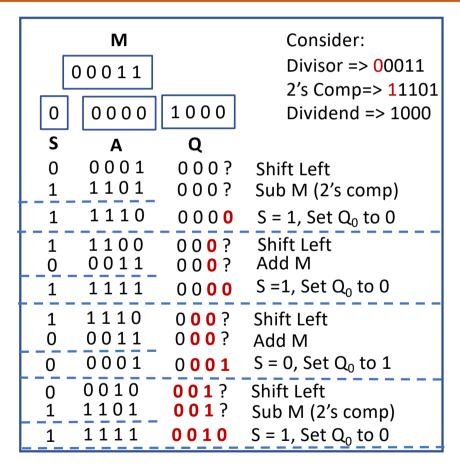


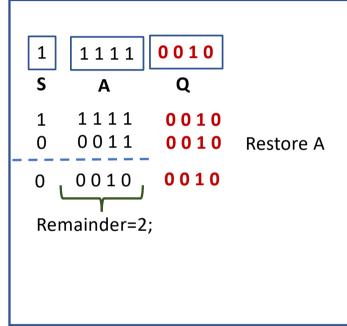




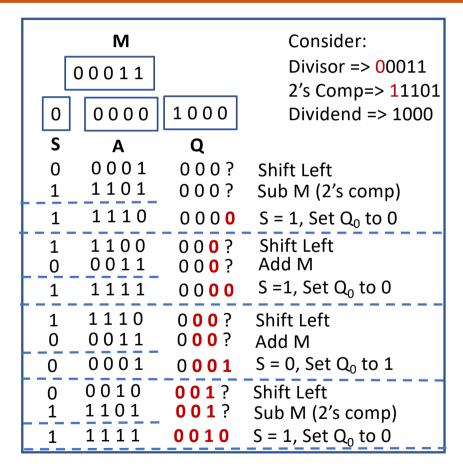


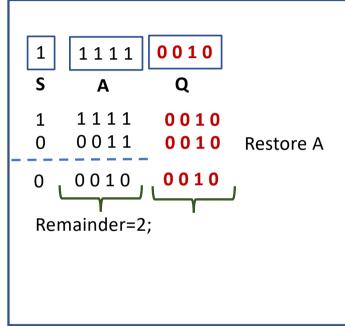




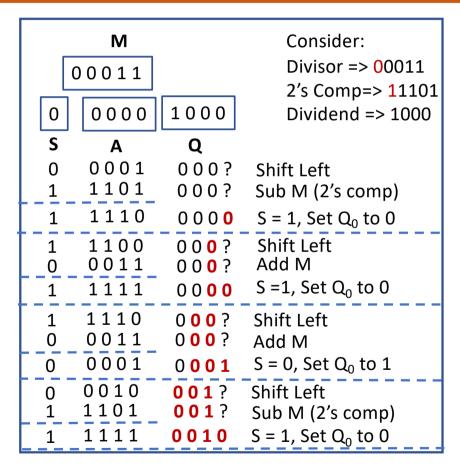


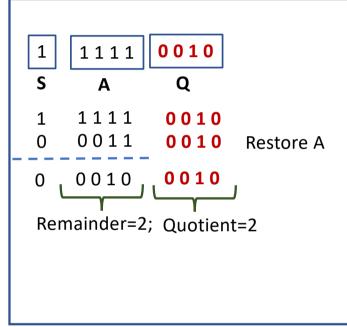




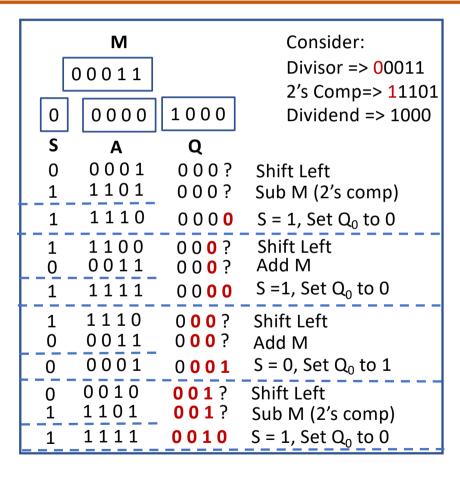


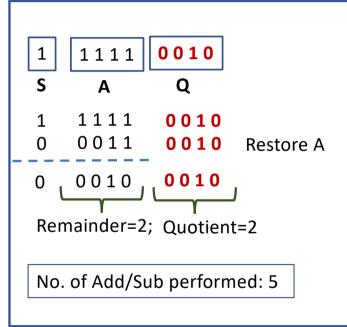














Non-Restoration Division



Points to be noted:

- Both Restoration and Non Restoration Division works for unisigned and positive signed numbers
- For signed numbers, convert the numbers to positive and adjust the sign later
- Simple solution Sign of the result will be XOR function of signs of Divisor and Dividend
- Ex: $7 \div 2 = 3$ and the remainder is 1 $-7 \div 2 = -3$ and the remainder is -1

Think about it



 Compare the number of addition / subtraction for the following operation using restoration and non-restoration methods:

For n-bits division what will be the worst case in terms of number of additions and subtractions for restoration and non-restoration division?



THANK YOU

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