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PES University, Bangalore (Established under Karnataka Act No. 16 of 2013)

UE19CS205

END SEMESTER ASSESSMENT (ESA) - B.TECH III SEMESTER - December, 2020

Automata Formal Languages & Logic Answer All Questions

Time: 3 Hrs

Max Marks: 100

1	а	Answer the following:						
		a) Construct a minimal DFA(DFA with minimum states) to accept binary strin	gs (5+3)					
		that contain at most one occurrence of the string "00".						
		b) Describe the language (i.e., set of all strings) accepted by the follow	wing					
		automaton:						
		State Input = a Input = b						
		$ ightarrow q_0$ q_2 q_1						
		$q_1 \hspace{0.1cm} q_1 \hspace{0.1cm} q_1$						
		q_2 q_3 q_2						
		$*q_3$ q_3 q_2	-					
	b	Construct an NFA/ λ -NFA that accepts the language over the alphabet {a, b} where words contain either "baa" as substring or where any 'a' is immediately followed at least two b's.	by 6					
	С		6					
		a a	[3+3]					
		$\frac{\lambda}{q0}$ $\frac{\lambda}{q1}$	[2.5]					
		d d d d						
		1) What is the lambda closure of q0, q1 and q2. [3 marks]						
		2) Convert the NFA to DFA using Subset Construction and provide a Transition	1 1					
		diagram of converted DFA, clearly mention the start and the final state(s).	1 1					
		[3 marks]						
2	а	Construct a regular expression for the regular language of strings over $\{a, b\}$ in who some number of a 's is followed by some number of b 's with the total length of	hich 4 the					
		string being divisible by 3.						

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	b		4
		a) What is the reversal of the given language L defined by regular expression	(2+2)
		01*+10*	
		b) True or False: Regular expressions that do not contain the star operator can	
		represent only finite languages.	
	С	Consider the following language over $\Sigma = \{ 0, E \}$:	6
		PARITY = $\{ w \mid w \text{ has even length and has the form } E^n \text{ or } w \text{ has odd length and has the } \}$	(3+3)
		form O ⁿ }	
		For example, EE \in PARITY, 00000 \in PARITY, EEEE \in PARITY, and $\varepsilon \in$ PARITY,	
		but	
		EEE ∉ PARITY, EO ∉ PARITY, and OOOO ∉ PARITY.	
		Write a regular expression and regular grammar for PARITY.	
	d	Convert the finite automata to regular expression using state elimination method:	6
		Eliminate the states in the order : q1, q0	
	-	C	
		(q1)	
	-	\ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \	
		9/	
		q_0	
	a	Construct a PDA for the language $L = \{a^nb^m \text{ where } m = n * 3 \text{ and } m,n >= 1\}$. Make sure	5
		that the PDA is deterministic .	
	,		
	b	Show that the grammar G with productions	5
		$S \rightarrow a \mid aAb \mid abSb$	
		A → aAAb bS is ambiguous.	
	-	Consider a language called ADD, which consists of all strings describing unary	10
	С	encodings of two sums that equal one another. For example:	10
ı		1 + 3 = 4 would be encoded as 1+111=1111	
		4 = 1 + 3 would be encoded as 1111=1111	
		2 + 2 = 1 + 3 would be encoded as $1111 = 1 + 111$	
		2+0+2+0=0+4+0 would be encoded as 11+11+111+	
1		0=0 would be encoded as =	
I		Notice that there can be any number of summands on each side of the =, but there	
ı		should be exactly one = in the string; thus 1=1=1 \iff ADD.	
		Answer the following:	
		I. Write a CFG that generates ADD. [5 marks]	
1		II. Show a parse tree for 1+1=11+ and +=+. [5 marks]	
	_	Prove using pumping lemma that the language L={ $a^nb^mc^m$, where $n \neq m$ }, is not a	-
+	a	110VC using pullipling lending that the language Let in it where it 2 mt is not a	>< 1
+	a	context-free language. [Note: list all the cases]	8

_		100			j.					
	b	Typ Typ Typ Fill t	h of the languages below in the table e DEC It is Turing-decidable. e TMR It is Turing-recognizable, be NTR It is not Turing-recognizable the column "TYPE" with an appropri C or TMR or NTR).	ut not decidable.		4				
		#	Langu	age	Туре					
		1	$EQ_{TM} = \{ < M_1, M_2 > M_1, M_2 \text{ are } T$	Ms with $L(M_1) = L(M_2)$.						
		2	HALT $_{TM} = \{ < M, w > M \text{ is a TM that}$	at halts on input w }.						
		3	$EQ_{DFA} = \{ < M_1, M_2 > M_1, M_2 \text{ are } D_1$	FAs with $L(M_1) = L(M_2)$ }.	-					
		4	$\overline{A_{TM}}$ [Complement of A_{TM}] wh $A_{TM} = \{ < M, w > M \text{ is a TM that acc} \}$			-				
	С	What	hat is the function computed by the following Turing Machines:							
		Turi	ing Machine 1	Turing Machine 2		(4+4)				
			q1	do di	□ , R					
5	а	predic (i) No	P(x) be "x is a professor" Q(x) be "x is ignorant" R(x) be "x is vain" ss the following statements using quates given above, where the domain professors are ignorant. I ignorant people are vain.	antifiers, logical connectives, an consists of all people.	d the	4 (2+2)				
	b	Use Re P1 : "II P2 : "II P3 : "II P4 : "II P5 : "S	esolution Refutation to show that the f Superman were unable to prevent f Superman were unwilling to preve f Superman were able and willing to f Superman exists, then he is neither uperman does not prevent evil" the conclusion Q : "Superman does not prevent evil"	evil, then he would be powerless nt evil, then he would be evil-mi prevent evil, then he would prev powerless nor evil-minded", and	nded", vent evil".	10				

	Pr	(int: Convert all the statements P1, P2, opositional logic. Assume the following p nversion:											
		nversion: ev: Superman prevents evil											
		le: Superman is able to prevent evil											
	willing: Superman is willing to prevent evil												
	pless: Superman is powerless												
	evilm: Superman is evil-minded exists: Superman exists												
	Ma	arks distribution:											
			e co	nclusion	to sen	tences	in Pro	nositi	ional				
	5 marks : To convert the statements and the conclusion to sentences in Propositional logic using given propositional symbols.												
		narks: To prove the conclusion Q using	reso	lution re	futatio	n.]							
С		nsider a group of men and women. Let,											
	M(x): x is a man												
	W(x): x is a woman												
	L(x, y) : x likes y												
	Match the following assertions with their equivalent logical expressions using the												
	abo	ve predicates.											
A A	1	Some men like all the women	a)M)xE	x) /\ \	≠y(W(y) → I	(y, x)))				
	2	For every man, there is at least one woman who likes him	b	∃x(M(:	() /\ \	/y(W(y) → L	(x, y)))				
	_					1 2-3							
	3	Some women does not like any man	С	∃x(M(x	() A E	y(W(y	y) /\ L	(x, y)))				
	4	There is a man whom all women like	d	∀x(M(x	() → Ξ	y(W(y	y) / L	(y, x)))				
	5	There is a man who likes a woman	е	∀x(W(x	() → <i>t</i>	/ y(M()	/) → L	(x, y)))				
	6	All women like all men	f	∃x(W(x	x) /\ \	y(M(y	₇) → −	L(x, y	7)))				
	Provide your answer in the following format only:												
	A	ssertions		1 2	3	4	5	6					
	ar	quivalent logical expression(write your iswer as a,b,c,d,e,f only - one of them atches each assertion).											

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