

DESIGN AND ANALYSIS OF ALGORITHMS

All Pairs Shortest Path (Floyd's Algorithm

Reetinder Sidhu

Department of Computer Science and Engineering



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UNIT 5: Limitations of Algorithmic Power and Coping with the Limitations

PES UNIVERSITY ONLINE

- Dynamic Programming
 - Computing a Binomial Coefficient
 - The Knapsack Problem
 - Memory Functions
 - Warshall's and Floyd's Algorithms
- Limitations of Algorithmic Power
 - Lower-Bound Arguments
 - Decision Trees
 - P, NP, and NP-Complete, NP-Hard Problems
- Coping with the Limitations
 - Backtracking
 - Branch-and-Bound. Architecture (microprocessor instruction set)

Concepts covered

- All Pairs Shortes Path (Floyd's Algorithm)
 - Definition
 - Algorithm
 - Example



Problem Definition

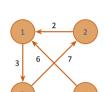


- Given a undirected or directed graph, with weighted edges, find the shortest path between every pair of vertices
 - ▶ Dijkstra's algorithm found shortest paths from given vertex to remaining n-1 vertices ($\Theta(n)$ paths)
 - ▶ Current problem is to find the shortest path between every pair of vertices ($\Theta(n^2)$ paths)
- Solution approach is similar to the transitive closure approach: Compute transitive closure via sequence of $n \times n$ matrices $R^{(0)}, \ldots, R^{(k)}, \ldots, R^{(n)}$ where $R^{(k)}[i,j] = 1$ iff there is nontrivial path from i to j with only first k vertices allowed as intermediate vertices
- Compute all pairs shortest paths via sequence of $n \times n$ matrices $D^{(0)}, \ldots, D^{(k)}, \ldots, D^{(n)}$ where $D^{(k)}[i,j]$ is the shortest path from i to j with only first k vertices allowed as intermediate vertices

Example



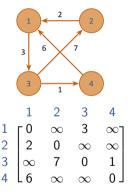
Example





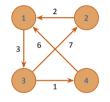
Example



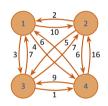


Example



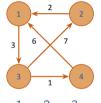


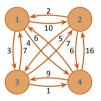




Example







Algorithm



Transitive Closure (Floyd's Algorithm)

```
1: procedure FLOYD(()A[1 \dots n, 1 \dots n])
2: \triangleright Input: Weight matrix A of a graph with no negative length cycles
3: \triangleright Output: Distance matrix of shortest paths
4: D \leftarrow W
5: for k \leftarrow 1 to n do
6: for i \leftarrow 1 to n do
7: for j \leftarrow 1 to n do
8: D[i,j] \leftarrow \min(D[i,j], D[i,k] + D[k,j])
9: return D
```

Algorithm



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• Complexity: $\Theta(n^3)$









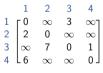




	1	2	3	4
1	0	∞	3	∞
2	2	0	∞	∞
3	∞	7	0	1
4	6	∞	∞	0]









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3	∞	7	0	1	
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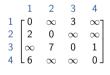
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2	2	0	5	6
3	9	7	0	1
4	L6	16	9	0_









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1	0	∞	3	∞
2	2	0	5	∞
3	∞	7	0	1
4	6	∞	9	0]



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2	2	0	5	∞
3	9	7	0	1
4	L6	∞	9	0]



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1	L0	10	3	47
2	2	0	5	6
3	9	7	0	1
4	L6	16	9	0_



	1	2	3	4
1	L0	10	3	47
2	2	0	5	6
3	9	7	0	1
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	1	2	3	4
1	L0	10	3	47
2	2	0	5	6
3	9	7	0	1
4	L6	16	9	0_



	1	2	3	4
1	L0	10	3	47
2	2	0	5	6
3	7	7	0	1
4	-6	16	9	0_

Think About It



- Give an example of a graph with negative weights for which Floyd's algorithm does not yield the correct result
- Enhance Floyd's algorithm so that shortest paths themselves, not just their lengths, can be found