



OPERATING SYSTEMS

Computer-System Architecture, OS Structure and Operations

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OPERATING SYSTEMS

Slides Credits for all PPTs of this course



- The slides/diagrams in this course are an **adaptation**, **combination**, and **enhancement** of material from the following resources and persons:
 1. Slides of Operating System Concepts, Abraham Silberschatz, Peter Baer Galvin, Greg Gagne - 9th edition 2013 and some slides from 10th edition 2018
 2. Some conceptual text and diagram from Operating Systems - Internals and Design Principles, William Stallings, 9th edition 2018
 3. Some presentation transcripts from A. Frank – P. Weisberg
 4. Some conceptual text from Operating Systems: Three Easy Pieces, Remzi Arpaci-Dusseau, Andrea Arpaci Dusseau

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Computer Architecture and Computer Organization



Computer Architecture	Computer Organization
Computer Architecture is concerned with the way hardware components are connected together to form a computer system.	Computer Organization is concerned with the structure and behaviour of a computer system as seen by the user.
It acts as the interface between hardware and software.	It deals with the components of a connection in a system.
Computer Architecture helps us to understand the functionalities of a system.	Computer Organization tells us how exactly all the units in the system are arranged and interconnected.
A programmer can view architecture in terms of instructions, addressing modes and registers.	Whereas Organization expresses the realization of architecture.
While designing a computer system architecture is considered first.	An organization is done on the basis of architecture.
Computer Architecture deals with high-level design issues.	Computer Organization deals with low-level design issues.
Architecture involves Logic (Instruction sets, Addressing modes, Data types, Cache optimization)	Organization involves Physical Components (Circuit design, Adders, Signals, Peripherals)

- Most systems use a single general-purpose processor
 - Most systems have other special-purpose processors as well.
 - Device specific processors like disk, keyboard, graphic controller
 - Special-purpose processors run a limited number of instructions
 - Managed by OS.
 - OS monitors the status.
- Example Disk controller microprocessor
 - Receives sequence of requests from CPU.
 - implements its own disk queue and scheduling algorithm
 - relieves the main CPU of the overhead of disk scheduling.
 - special-purpose processors are low-level components built into the hardware

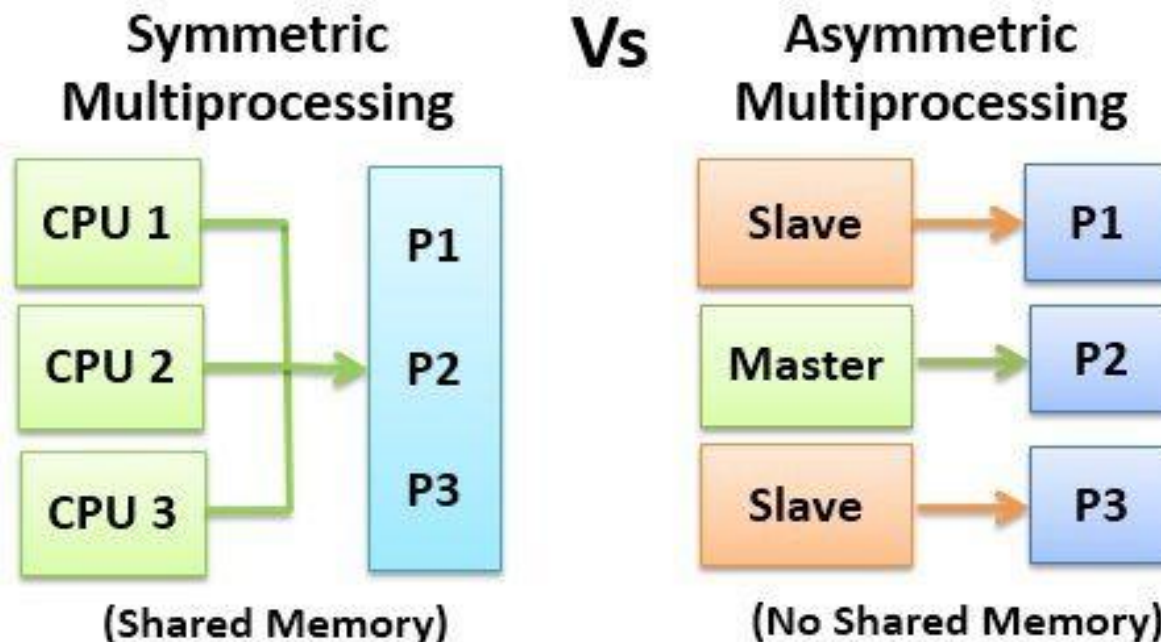
- **Multiprocessors** systems growing in use and importance
 - Also known as **parallel systems**, **tightly-coupled systems**
 - Advantages include:
 1. **Increased throughput**
 2. **Economy of scale**
 3. **Increased reliability** – graceful degradation or fault tolerance

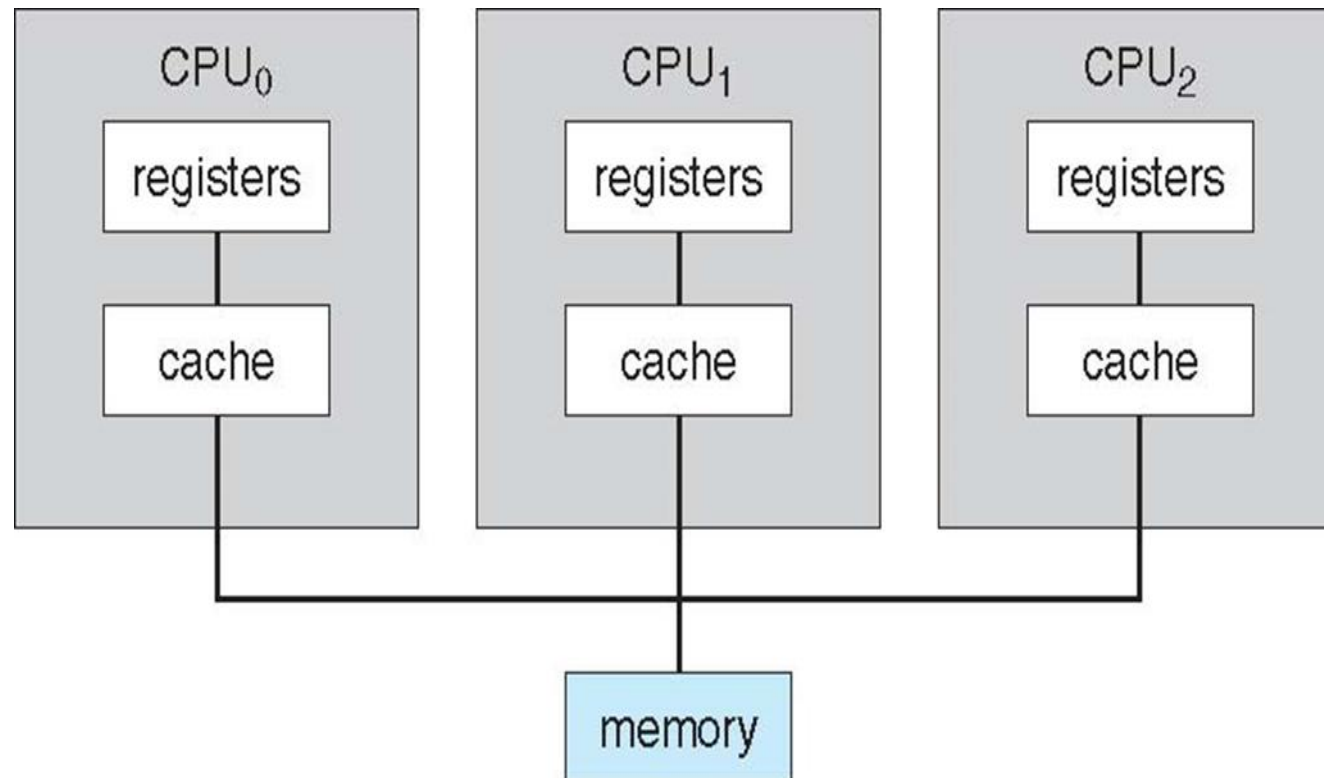
Advantages include

- Increased reliability is crucial

□ Two types of Multiprocessor Systems

1. **Asymmetric Multiprocessing** – each processor is assigned a specific task.
2. **Symmetric Multiprocessing** – each processor performs all tasks

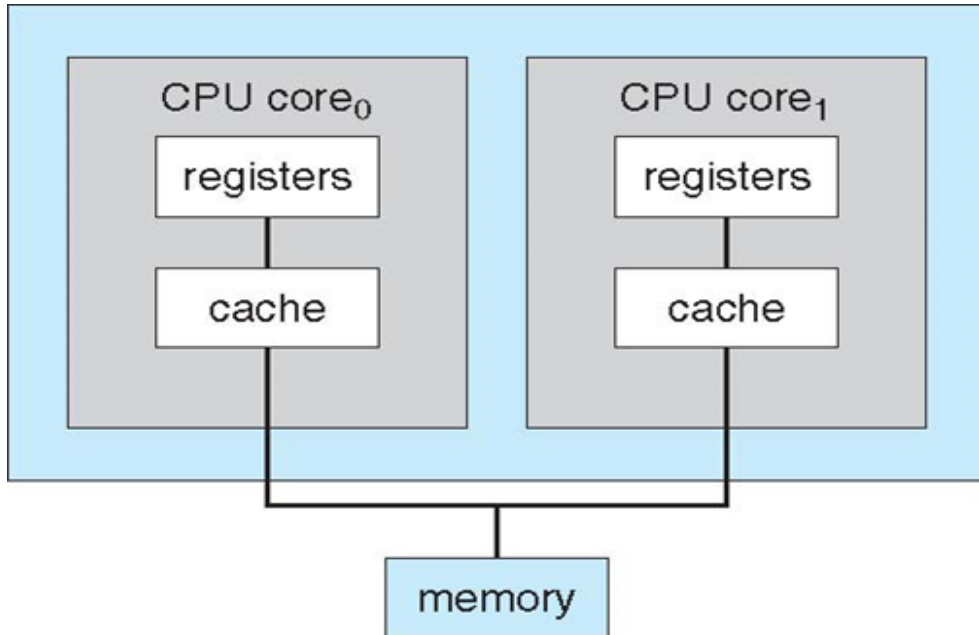




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A Dual-Core Design

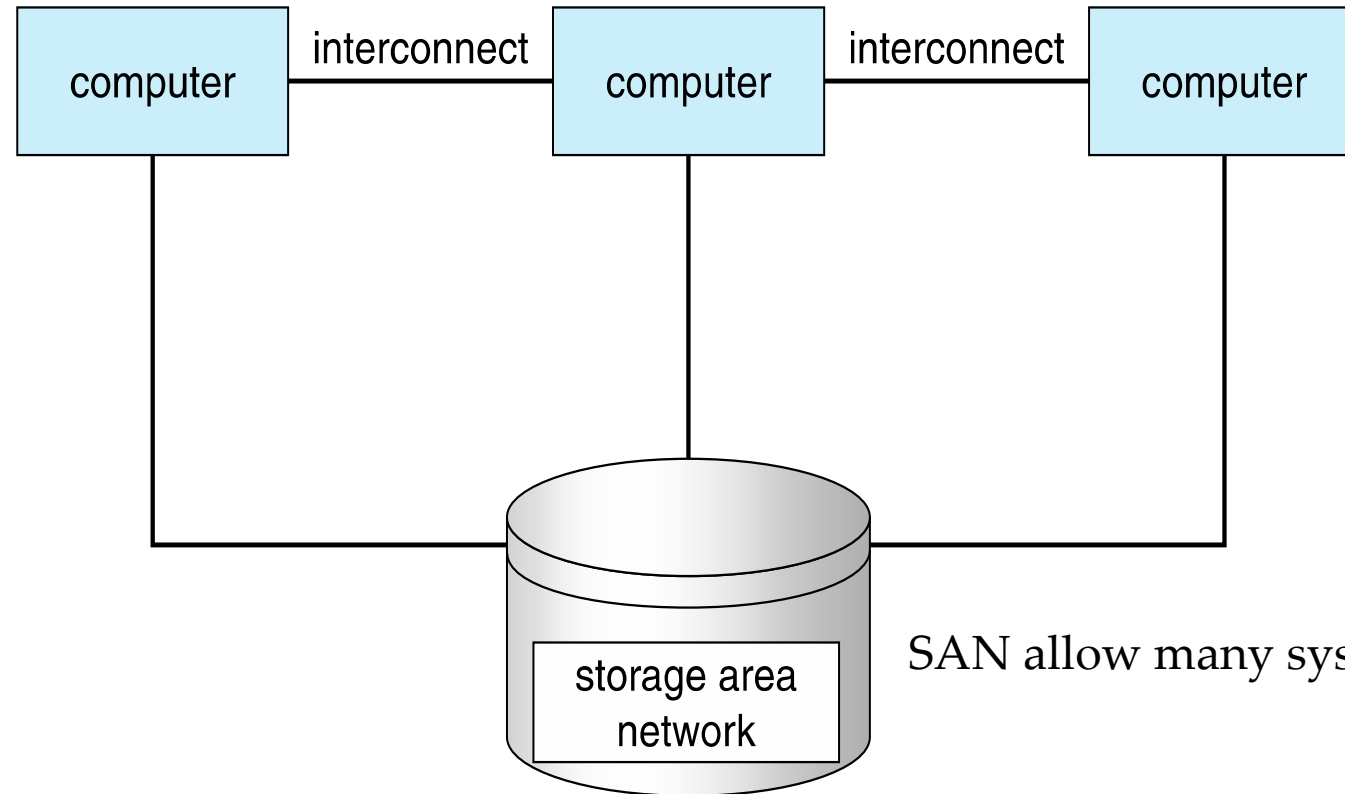
- ❑ Multi-chip and **multicore**
- ❑ Systems containing all chips
 - ❑ Chassis containing multiple separate systems



Command to know the number of cores, cache details
`$cat /proc/cpuinfo|more`

- **blade servers** are a recent development in which multiple processor boards, I/Oboards, and networking boards are placed in the same chassis.
 - blade-processor board boots independently and runs its own operating system.
 - Some blade-server boards are multiprocessor as well, which blurs the lines between types of computers.
 - In essence, these servers consist of multiple independent multiprocessor systems.

- ❑ Like multiprocessor systems, but multiple systems working together
 - ❑ Usually sharing storage via a **storage-area network (SAN)**
 - ❑ Provides a **high-availability** service which survives failures
 - ▶ **Asymmetric clustering** has one machine in hot-standby mode
 - ▶ **Symmetric clustering** has multiple nodes running applications, monitoring each other
 - ❑ Some clusters are for **high-performance computing (HPC)**
 - ▶ Applications must be written to use **parallelization**
 - ❑ Some have **distributed lock manager (DLM)** to avoid conflicting operations (Ex: when multiple hosts access the same data on shared storage)



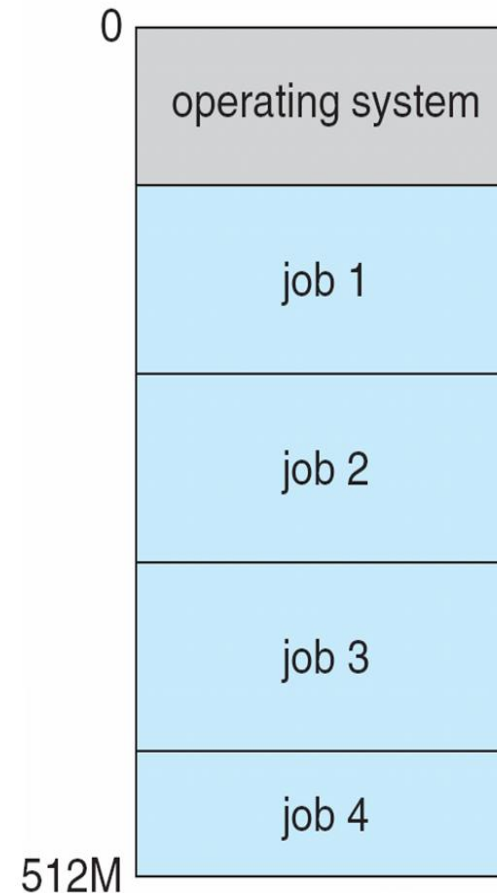
SAN allow many systems to attach to a pool of storage

General structure of a clustered system.

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Operating-System Structure - Multiprogramming

- ❑ **Multiprogramming** (**Batch system**) needed for efficiency
 - ❑ Single user cannot keep CPU and I/O devices busy at all times
 - ❑ Multiprogramming organizes jobs (code and data) so CPU always has one to execute
 - ❑ A subset of total jobs in system is kept in memory
 - ❑ One job selected and run via **job scheduling**
 - ❑ When it has to wait (for I/O for example), OS switches to another job



- **Timesharing** (**multitasking**) is logical extension in which CPU switches jobs so frequently that users can interact with each job while it is running, creating **interactive** computing
 - **Response time** should be < 1 second
 - Each user has at least one program executing in memory
⇒ **process**
 - If several jobs ready to run at the same time ⇒ **CPU scheduling**
 - If processes don't fit in memory, **swapping** moves them in and out to run
 - **Virtual memory** allows execution of processes not completely in memory

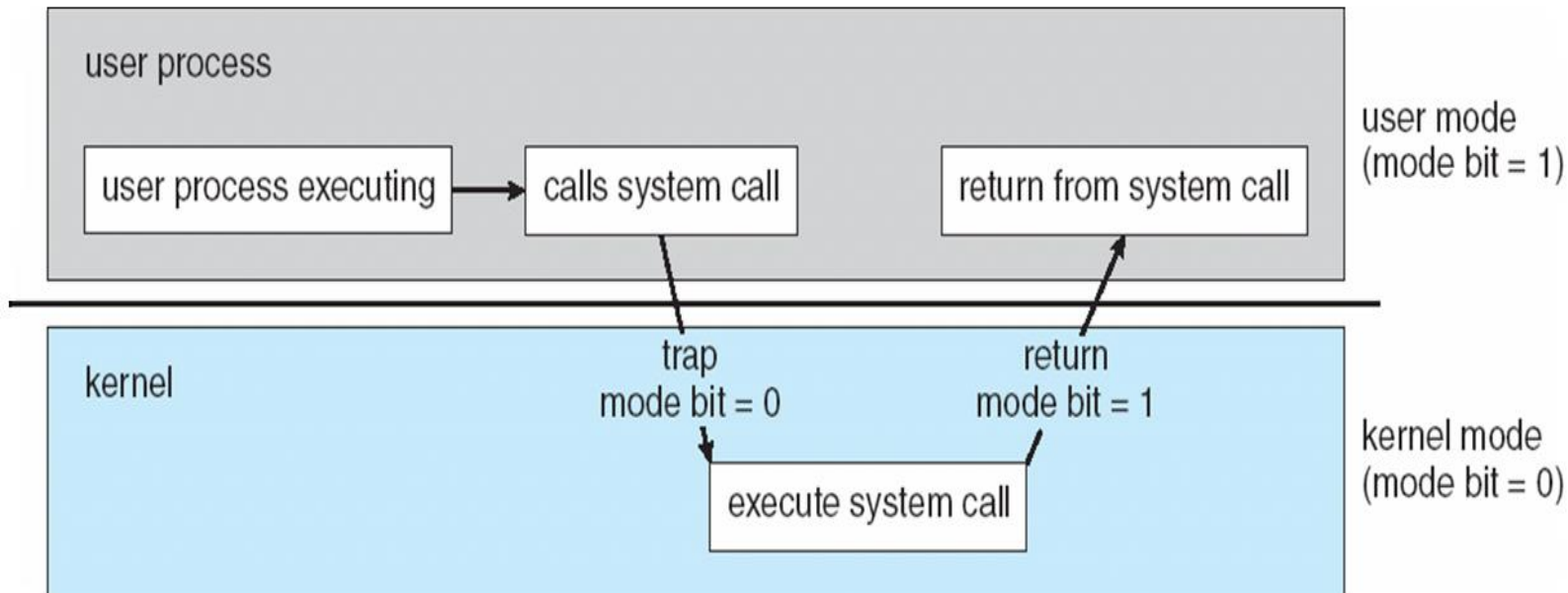
- **Interrupt driven** (hardware and software)
 - Hardware interrupt by one of the devices
 - Software interrupt (**exception** or **trap**):
 - ▶ Software error (e.g., division by zero)
 - ▶ Request for operating system service
 - ▶ Other process problems include infinite loop, processes modifying each other or the operating system

- **Dual-mode** operation allows OS to protect itself and other system components
 - **User mode** and **kernel mode**
 - **Mode bit** provided by hardware
 - ▶ Provides ability to distinguish when system is running user code or kernel code
 - ▶ Some instructions designated as **privileged**, only executable in kernel mode
 - ▶ System call changes mode to kernel, return from call resets it to user
- Increasingly CPUs support multi-mode operations
 - i.e. **virtual machine manager (VMM)** mode for guest **VMs**

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Transition from user to kernel mode

- When a trap or interrupt occurs, hardware switches from user mode to kernel mode (changes the state of the mode bit to 0).
- When the request is fulfilled, the system always switches to user mode (by setting the mode bit to 1) before passing control to a user program.



- ❑ Timer to prevent infinite loop / process hogging resources
 - ❑ Timer is set to interrupt the computer after a specified period (fixed 1/60 sec or variable 1 msec to 1 sec)
 - ❑ A **variable timer** is generally implemented by a fixed-rate clock and a counter.
 - ❑ Operating system sets the counter (privileged instruction)
 - ❑ Every time the clock ticks, the counter is decremented.
 - ❑ When counter reaches zero, an interrupt occurs
 - ❑ Timer can be used to prevent a user program from running too long (terminate the program)



THANK YOU

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