

# **PES University, Bangalore**

## **Department of Computer Science and Engineering**

## **Automata Formal Languages & Logic**

# Q&A for Propositional Theorem Proving- Inference Rule AND Resolution Algorithm

### **Problem:**

1. Given axioms P  $\Lambda$  ((P  $\Lambda$  Q) -> R)  $\Lambda$  ((S V T)-> Q)  $\Lambda$  T Prove R using Resolution algorithm.

**Solution:-** We have

R1: P

R2:  $(P \land Q) \rightarrow R \equiv (P \land Q) \lor R \equiv P \lor Q \lor R$ 

Next (S V T)-> Q  $\equiv$   $\sim$  (S V T) V Q  $\equiv$  ( $\sim$ S  $\wedge$   $\sim$ T) V Q  $\equiv$  ( $\sim$ S V Q )  $\wedge$  ( $\sim$ T V Q)

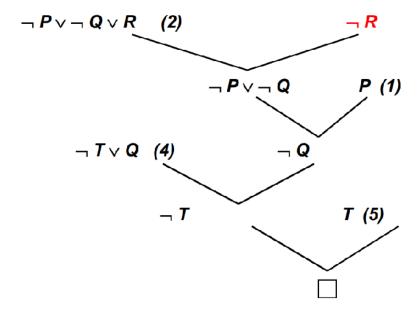
R3: (~S V Q)

R4: (~T V Q)

R5: T

R6: ~R





Received an empty clause, which is equivalent to False. Empty clause represents a contradiction and hence "R" is true.

### 2. Given

- a. John is in Paris or John is in Australia
- b. If John is in Paris, then It is raining
- c. If John is in Australia, then It is raining DO the following
  - 1. Write the clauses
  - 2. Convert it in CNF form
  - 3. **Prove** It is raining using Resolution.

#### Solution:-

Let

P: John is in Paris

Q: John is in Australia

R: It is raining

To prove R: It is raining

- 1. Write all the statements in clausal form
  - a. John is in Paris or John is in Australia. R1: (P V Q)



b. If John is in Paris, then It is raining. R2: (P-> R)

c. If John is in Australia, then It is raining. R3: (Q-> R)
 To prove R include R4: ~R

2. In CNF Form

R1: (P V Q)

R2:  $(P-> R) \equiv P VQ$ 

R3:  $(Q \rightarrow R) \equiv Q \vee R$ 

R4: ~R

3.

Step	Formula	Derivation
1	PvQ	Given
2	¬PvR	Given
3	¬Q∨R	Given
4	¬ R	Negated conclusion
5	QvR	1,2
6	¬ P	2,4
7	¬ Q	3,4
8	R	5,7
9	•	4,8

Received an empty clause, which is equivalent to False. Empty clause represents a contradiction and hence "R" is true.