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## **DESIGN AND ANALYSIS OF ALGORITHMS**

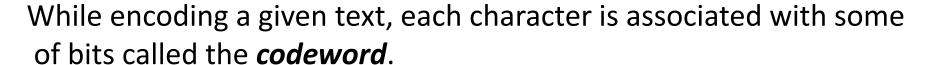
# **Unit 4: Greedy Technique- Huffman Trees**

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#### **Huffman Trees**

The Huffman trees are constructed for encoding a given text of n characters.



- **Fixed length encoding**: Assigns to each character a bit string of the same length.
- Variable length encoding: Assigns codewords of different lengths to different characters.
- **Prefix free code**: In Prefix free code, no codeword is a prefix of a codeword of another character.



#### **Binary prefix code:**

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- The characters are associated with the leaves of a binary tree.
- All left edges are labeled as 0, and right edges as 1.
- Codeword of a character is obtained by recording the labels on the simple path from the root to the character's leaf.
- Since, there is no simple path to a leaf that continues to another leaf, no codeword can be a prefix of another codeword.

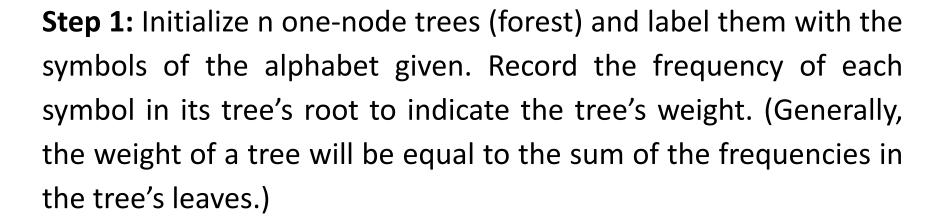
# **About Huffman Algorithm:**

- Invented by David A Huffman in 1951.
- Constructs binary prefix code tree.
- Huffman's algorithm achieves data compression by finding the best variable length binary encoding scheme for the symbols that occur in the file to be compressed. Huffman coding uses frequencies of the symbols in the string to build a variable rate prefix code
  - O Each symbol is mapped to a binary string
  - More frequent symbols have shorter codes
  - O No code is a prefix of another code (prefix free code)
- Huffman Codes for data compression achieves 20-90% Compression.



#### **Huffman Algorithm:**

**Input:** Alphabet and frequency of each symbol in the text.



**Step 2:** Repeat the following operation until a single tree is obtained. Find two trees with the smallest weight. Make them the left and right subtree of a new tree and record the sum of their weights in the root of the new tree as its weight.



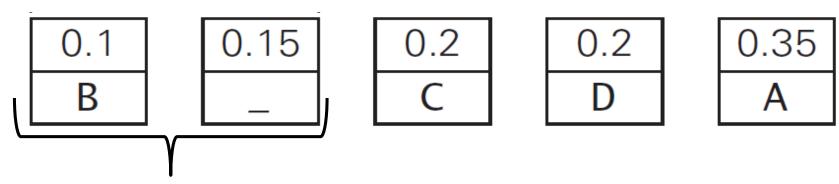
#### **Huffman Algorithm:**



#### **EXAMPLE**

symbol	Α	В	C	D	_
frequency	0.35	0.1	0.2	0.2	0.15

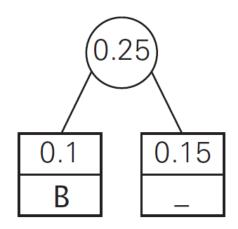
#### 1. First arrange the characters in ascending order of their probabilities



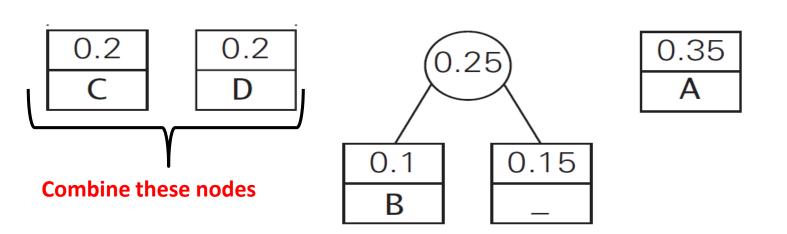
Combine these two nodes and form a single node

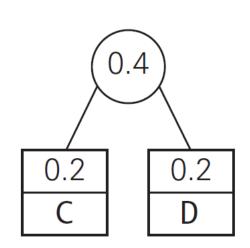
#### **Huffman Algorithm:**





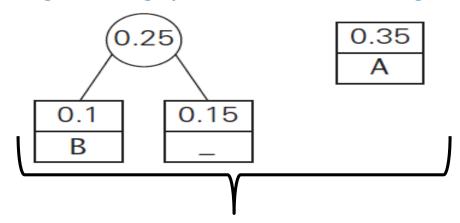
#### 2. Again arrange probabilities in ascending order

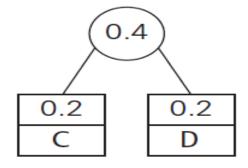




#### **Huffman Algorithm:**

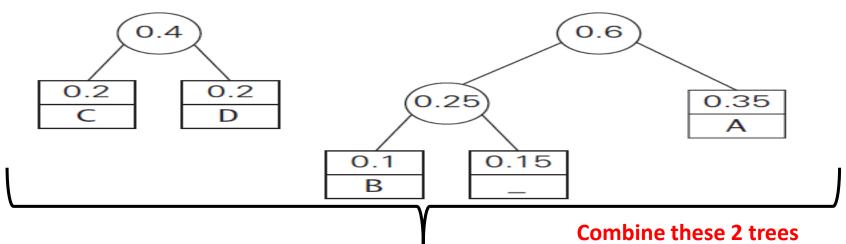
#### 3. Again arrange probabilities in ascending order





#### **Combine these 2 trees**

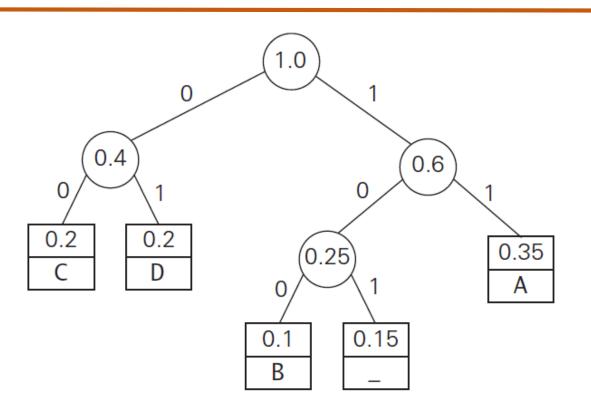
#### 4. Arrange probabilities in ascending order





#### **Huffman Tree**





DAD is encoded as 011101, and 10011011011101 is decoded as BAD\_AD

#### The resulting codewords are as follows

symbol	Α	В	C	D	_
frequency	0.35	0.1	0.2	0.2	0.15
codeword	11	100	00	01	101

#### **Compression Ratio**



Compression Ratio: A standard measure of a compression algorithm's efficiency.

Character	A	В	C	D	-
probability	0.4	0.1	0.2	0.15	0.15
Character	A	В	C	D	-
probability	0.4	0.1	0.2	0.15	0.15
Code word	0	100	111	101	110

## Compute compression ratio:

Bits per character = Codeword length \* Frequency

$$= (1*0.4) + (3*0.1) + (3*0.2) + (3*0.15) + (3*0.15)$$

= 2.20

Compression ratio is = (3 - 2.20)/3.100% = 26.6%

#### **Features of Huffman's Encoding**

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- Huffman's encoding is one of the most important filecompression methods.
- In addition to its simplicity and versatility, it yields an optimal, i.e., minimal-length, encoding (provided the frequencies of symbol occurrences are independent and known in advance).
- The simplest version of Huffman compression calls, in fact, for a preliminary scanning of a given text to count the frequencies of symbol occurrences in it. Then these frequencies are used to construct a Huffman coding tree and encode the text.

#### **Text Books**

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Chapter-9 Greedy Technique Introduction to the Design & Analysis of Algorithms- Anany Levitin



# **THANK YOU**

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