

AUTOMATA FORMAL LANGUAGES AND LOGIC

First Order Logic – The Electronic Circuits Domain

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Electronic Circuits Domain

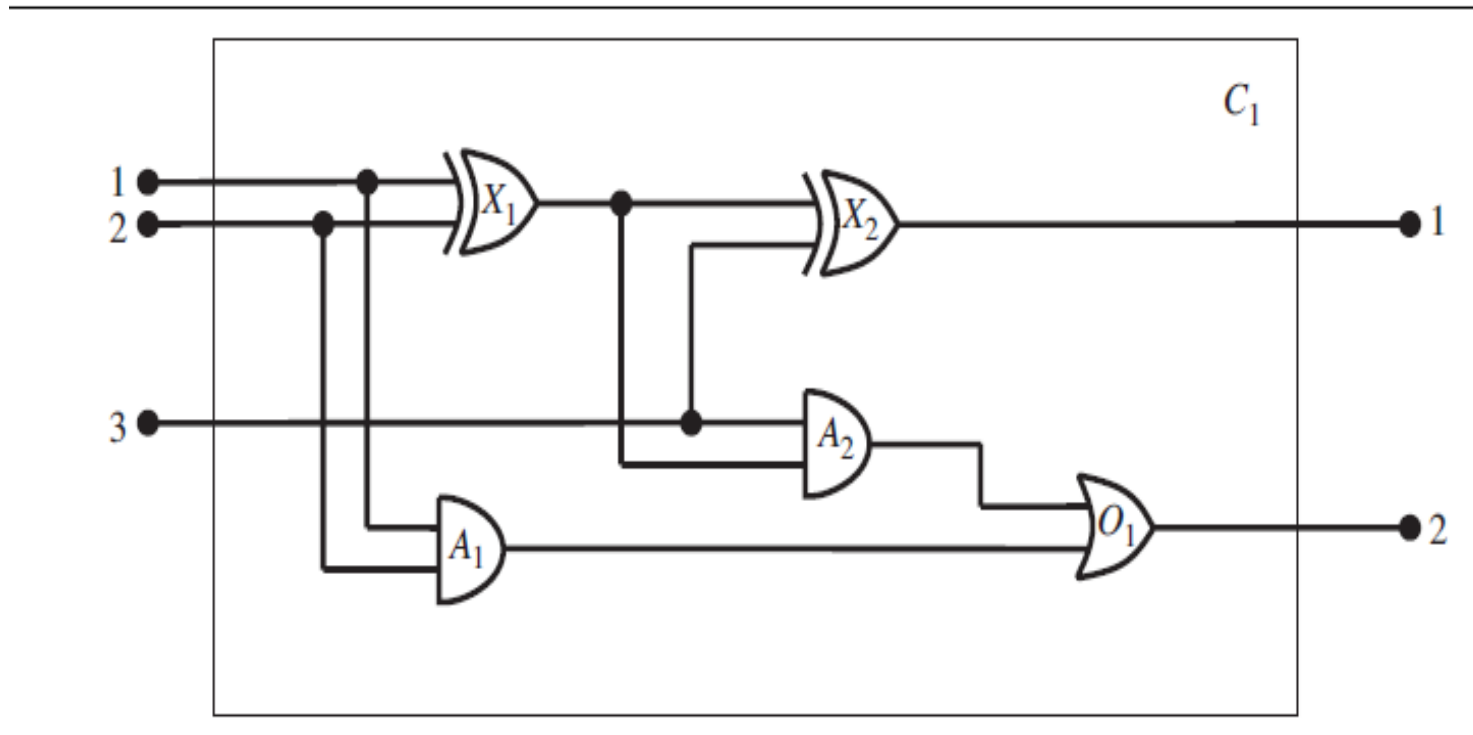


Seven-step process for Knowledge Engineering

1. Identify the task
2. Assemble the relevant knowledge
3. Decide on a vocabulary
4. Encode general knowledge of the domain
5. Encode the specific problem instance
6. Pose queries to the inference procedure
7. Debug the knowledge base

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Digital circuits



Decide on a vocabulary

- Next we consider **terminals**, which are identified by the predicate **Terminal (x)**.
- A gate or circuit can have one or more input terminals and one or more output terminals.
- We use the function **In(1,X1)** to denote the first input terminal for gate X1.
- A similar function Out is used for output terminals.
- The function **Arity(c, i, j)** says that circuit c has i input and j output terminals.
- The connectivity between gates can be represented by a predicate, Connected, which takes two terminals as arguments, as in **Connected(Out(1,X1), In(1,X2))**.

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Decide on a vocabulary

- Finally, we need to know whether a *signal* is on or off.
 - One possibility is to use a unary predicate, *On(t)*, which is true when the signal at a terminal is on.
 - This makes it a little difficult, however, to pose questions such as “**What are all the possible values of the signals at the output terminals of circuit C1 ?**”
- We therefore introduce as objects two signal values, 1 and 0, and a function *Signal (t)* that denotes the signal value for the terminal t.

- One sign that we have a good ontology is that we require only a few general rules, which can be stated clearly and concisely.
- **These are all the axioms we will need:**
- ***1. If two terminals are connected, then they have the same signal:***
$$\forall t1, t2 \text{ Terminal } (t1) \wedge \text{Terminal } (t2) \wedge \text{Connected}(t1, t2) \Rightarrow \text{Signal } (t1) = \text{Signal } (t2) .$$
- ***2. The signal at every terminal is either 1 or 0:***
$$\forall t \text{ Terminal } (t) \Rightarrow \text{Signal } (t) = 1 \vee \text{Signal } (t) = 0 .$$
- ***3. Connected is commutative:***
$$\forall t1, t2 \text{ Connected}(t1, t2) \Leftrightarrow \text{Connected}(t2, t1) .$$

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Encode general knowledge of the domain



4. There are four types of gates:

$\forall g \text{ Gate}(g) \wedge k = \text{Type}(g) \Rightarrow k = \text{AND} \vee k = \text{OR} \vee k = \text{XOR} \vee k = \text{NOT} .$

5. An AND gate's output is 0 if and only if any of its inputs is 0:

$\forall g \text{ Gate}(g) \wedge \text{Type}(g)=\text{AND} \Rightarrow \text{Signal}(\text{Out}(1, g))=0 \Leftrightarrow \exists n \text{ Signal}(\text{In}(n, g))=0 .$

6. An OR gate's output is 1 if and only if any of its inputs is 1:

$\forall g \text{ Gate}(g) \wedge \text{Type}(g)=\text{OR} \Rightarrow \text{Signal}(\text{Out}(1, g))=1 \Leftrightarrow \exists n \text{ Signal}(\text{In}(n, g))=1 .$

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Encode general knowledge of the domain

8. A NOT gate's output is different from its input:

$$\forall g \text{ Gate}(g) \wedge (\text{Type}(g)=\text{NOT}) \Rightarrow \text{Signal}(\text{Out}(1, g)) = \text{Signal}(\text{In}(1, g)) .$$

9. The gates (except for NOT) have two inputs and one output.

$$\forall g \text{ Gate}(g) \wedge \text{Type}(g) = \text{NOT} \Rightarrow \text{Arity}(g, 1, 1) .$$

$$\forall g \text{ Gate}(g) \wedge k = \text{Type}(g) \wedge (k = \text{AND} \vee k = \text{OR} \vee k = \text{XOR}) \Rightarrow \text{Arity}(g, 2, 1)$$

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Encode general knowledge of the domain



10. A circuit has terminals, up to its input and output arity, and nothing beyond its arity:

$$\forall c, i, j \text{ Circuit}(c) \wedge \text{Arity}(c, i, j) \Rightarrow \forall n (n \leq i \Rightarrow \text{Terminal}(\text{In}(c, n))) \wedge (n > i \Rightarrow \text{In}(c, n) = \text{Nothing}) \wedge$$
$$\forall n (n \leq j \Rightarrow \text{Terminal}(\text{Out}(c, n))) \wedge (n > j \Rightarrow \text{Out}(c, n) = \text{Nothing})$$

11. Gates, terminals, signals, gate types, and Nothing are all distinct.

$$\forall g, t \text{ Gate}(g) \wedge \text{Terminal}(t) \Rightarrow g = t = 1 = 0 = \text{OR} = \text{AND} = \text{XOR} = \text{NOT} = \text{Nothing} .$$

12. Gates are circuits. $\forall g \text{ Gate}(g) \Rightarrow \text{Circuit}(g)$

The circuit is encoded as circuit C1 with the following description.

First, we categorize the circuit and its component gates:

Circuit(C1) \wedge Arity(C1, 3, 2)

Gate(X1) \wedge Type(X1)=XOR

Gate(X2) \wedge Type(X2)=XOR

Gate(A1) \wedge Type(A1)=AND

Gate(A2) \wedge Type(A2)=AND

Gate(O1) \wedge Type(O1)=OR .

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Encode the specific problem instance

Then, we show the connections between them:

Connected(Out(1,X1), In(1,X2)) Connected(In(1,C1), In(1,X1))
Connected(Out(1,X1), In(2,A2)) Connected(In(1,C1), In(1,A1))
Connected(Out(1,A2), In(1,O1)) Connected(In(2,C1), In(2,X1))
Connected(Out(1,A1), In(2,O1)) Connected(In(2,C1), In(2,A1))
Connected(Out(1,X2), Out(1,C1)) Connected(In(3,C1), In(2,X2))
Connected(Out(1,O1), Out(2,C1)) Connected(In(3,C1), In(1,A2))

What combinations of inputs would cause the first output of C1 (the sum bit) to be 0 and the second output of C1 (the carry bit) to be 1?

$\exists i1, i2, i3 \text{ Signal (In(1, C1))}=i1 \wedge \text{Signal (In(2, C1))}=i2 \wedge \text{Signal (In(3, C1))}=i3 \wedge \text{Signal (Out(1, C1))}=0 \wedge \text{Signal (Out(2, C1))}=1 .$

The answers are substitutions for the variables $i1$, $i2$, and $i3$ such that the resulting sentence is entailed by the knowledge base.

ASKVARS will give us three such substitutions:

$\{i1/1, i2/1, i3/0\} \{i1/1, i2/0, i3/1\} \{i1/0, i2/1, i3/1\} .$

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Pose queries to the inference procedure

What are the possible sets of values of all the terminals for the adder circuit?

$\exists i1, i2, i3, o1, o2 \text{ Signal (In(1, C1))=i1} \wedge \text{Signal (In(2, C1))=i2}$
 $\wedge \text{Signal (In(3, C1))=i3} \wedge \text{Signal (Out(1, C1))=o1} \wedge \text{Signal (Out(2, C1))=o2} .$

- We can perturb the knowledge base in various ways to see what kinds of erroneous behaviors emerge.
- For example, we can ask $\exists i1, i2, o \text{ Signal (In(1,C1))=i1} \wedge \text{Signal (In(2,C1))=i2} \wedge \text{Signal (Out(1,X1))}$, which reveals that no outputs are known at X1 for the input cases 10 and 01.
- Then, we look at the axiom for XOR gates, as applied to X1:

$\text{Signal (Out(1,X1))=1} \Leftrightarrow \text{Signal (In(1,X1))} = \text{Signal (In(2,X1))} .$

- If the inputs are known to be, say, 1 and 0, then this reduces to $\text{Signal (Out(1,X1))=1} \Leftrightarrow 1 = 0 .$
- Now the problem is apparent: the system is unable to infer that $\text{Signal (Out(1,X1))=1}$, so we need to tell it that $1 = 0$.

Reference

- Some of the definitions/diagrams/examples are from the book “Artificial Intelligence – A Modern Approach”, Stuart Russell and Peter Norvig, Pearson, 3rd Edition (Paperback), 2016



THANK YOU

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