Security and Protection in xv6

CS 344 Assignment 4

Contents

- Intro to Security and Protection
- Comparison between features of xv6 and Linux
- Drawing out a strategy
- Pseudocode for implementation

Protection

Internal Threats

- Controls the access to the system resources
- Uses principle of least privilege and is implemented using authorisation/permission policies

Security

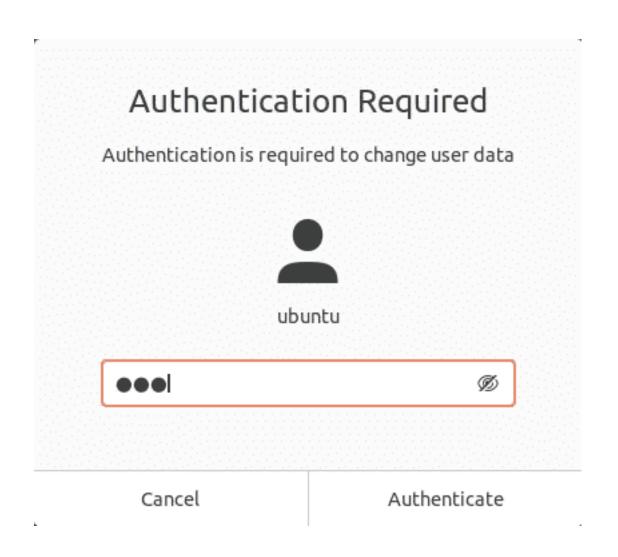
External Threats

- Allows only legitimate users to use the system
- Implemented through authentication, encryption, firewalls, etc.

Linux OS

Security and Protection Features

- User Authentication Usually a username password mechanism to identify user.
 - Authentication is done right after booting and before privileged operations.
- File System Protection Defines which users/groups can read, write and execute given files.
 - Ensures that users can't perform unauthorised operations on files.
- Firewalls Linux uses the Netfilter framework to implement firewalls



XV6

Security and Protection Features

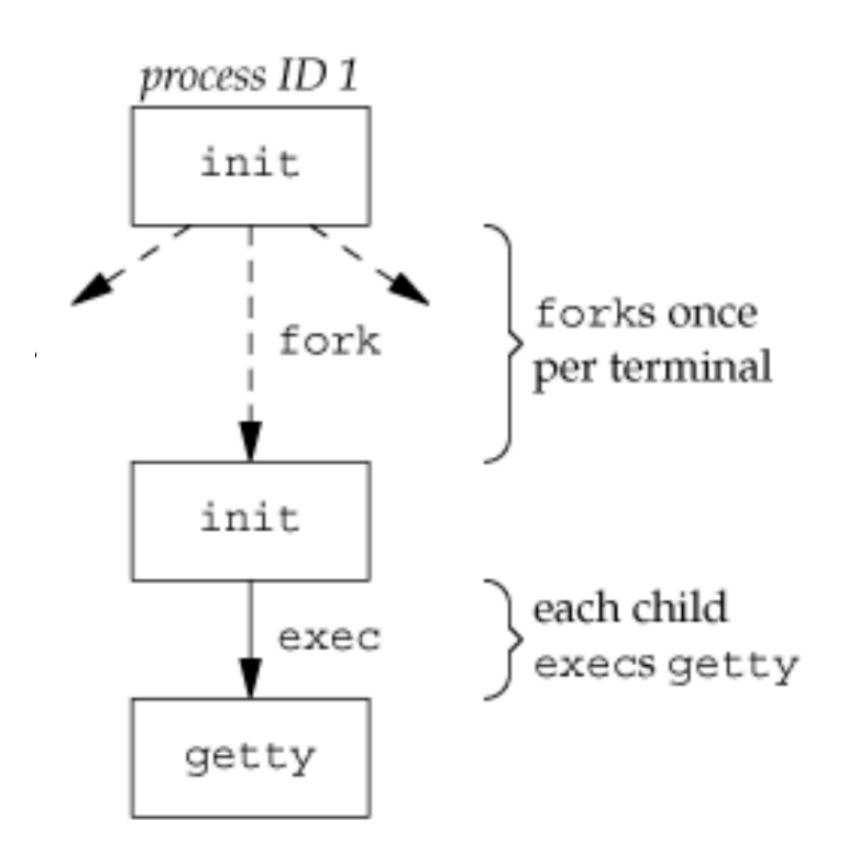
- Protection limited to memory access.
- Uses segmentation to protect parts of a program. Direct access to code and stack segment is not allowed.
- Implements protection of physical memory pages from improper access using PTE_U bit in page table entry.
- No notion of users (not multiuser OS). Hence, no form of security and authentication in place.
- User Authentication
- File System Access Control

User Authentication

- Create a /passwd file in the root directory to store username, uid, private group and password
 - username = root and password = 123456 (arbitrary)
 - Format: <username>:<password>:<uid>:<private-gid>
- New user programs:
 - useradd Create a new user
 - Format: useradd <username>
 - passwd Change/add a password for a user
 - Format: passwd <username> <passwd>

User Authentication

- Create getty.c for login prompt.
 - This process should access /passwd to verify the identity of the user.



File System Protection

- We will add groups into xv6. Create a similar /group file in root to store group related details.
 - Format: <group_name>:<gid>:<members>
- Store access mode in struct inode and uid, gid and euid in struct proc.
- System calls:
 - int setuid(int uid) allows root to change uid of process
 - int setgid(int gid)
 - int chown(char * path, int uid) Changes owner uid of a file

File System Protection

- System calls:
 - int chmod(char* path, int mode) changes the access permission mode of a file
 - int chgrp (char* path, int gid) changes group owner of a file
- User programs:
 - groupadd Adds a group entry in /group file
 - usermod Adds user into a group
 - Format: usermod <username> <groupname>
 - chmod, chgrp, chown Allows user to call resp. syscalls

User Authentication- Add /passwd file into xv6 file system

- mkfs.c is responsible for mounting the file system in xv6
- Create a passwd file in the xv6 source code folder
- Within this file, add the root user credentials in the specified format. Note that root user must have a uid = 0
- Update Makefile to reflect this change:

User Authentication- struct user

- Create struct user to store user details.
- This struct can be used a return type while parsing the /passwd file later on

```
struct user {
  int uid;
  int gid;
  char username[25];
  char password[25];
```

User Authentication- useradd.c

- User program to add a new user.
- Command: useradd <username>

useradd.c:

```
Command Line Arguments: username

passwdFile = OPEN("/passwd", READ|APPEND_MODE)

groupFile = OPEN("/group", READ_ONLY_MODE)

uid = minimum unused UID in passwdFile

gid = minimum unused GID in groupFile

passwdFile.WRITE("%username::%uid,:%gid")

passwdFile.CLOSE()
```

User Authentication- passwd.c

- Updates password of a user
- Command: passwd <username> <password>

passwd.c: Command Line Arguments : uname, pass passwdFile = OPEN("/passwd",READ_WRITE_MODE) FOR EACH line IN passwdFile IF line[username] == uname line[password] = pass passwdFile.WRITE(line) // Update password field and rewrite ENDIF ENDFOR passwdFile.CLOSE()

User Authentication- getty.c

This is login prompt process that is forked by the init process.

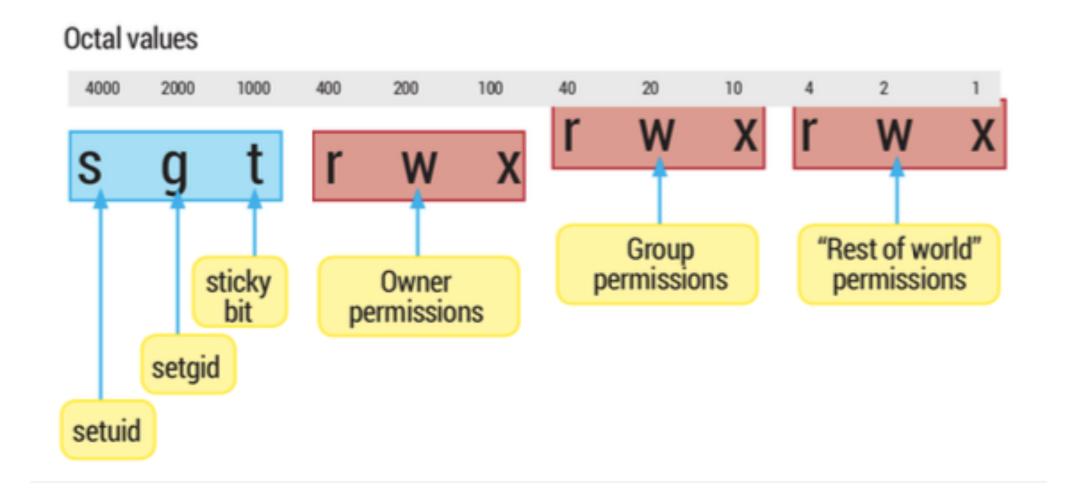
In init.c, remove exec("sh") and replace with exec("getty")

User Authentication- proc.c

- We need to make sure every process stores the logged in user's uid, gid and has an effective uid.
- Effective uid may not be equal to uid if setuid bit is used for a file.
- allocproc(): Initialise uid, gid and euid to 0.
- fork(): Make child process inherit uid, gid, euid from parent.

File System Protection - Initial changes

- Create a /group file in a similar manner as /password.
- Add uid, gid field for ownership in struct inode and field mode for permissions
- In sysfile.c, update the create() function.
 - This function is used to allocate inodes for new files.
 - Make sure that a new inode, borrows the running process' uid and gid



File System Protection - setuid() / setgid() system call

- System call to change uid/gid of current running process.
- Can be called only by root user.
- Hence, the pseudo code is:

```
sys_setuid()
  newUID = ARGINT(0) // Fetch argument from user stack
  IF myproc()->uid == ROOT_UID
        myproc()->uid = newUID
        RETURN 0
  ENDIF
  RETURN -1
```

• setgid() is almost identical but uses gid instead of uid.

File System Protection - Enforcing permissions

```
int allowRead(char* path)
    fileInode = nameiparent(path) // Get inode
    ownUID = fileInode->uid
    ownGID = fileInode->gid
   mode = fileInode->mode
   IF fileInode->mode & SETUID
       myproc()->euid = ownUID
    ENDIF
    fMembers = list of UIDs belonging to ownGID group in /group
   currUID = myproc()->euid
    currGID = myproc()->gid
    IF currUID == ROOT_UID // Root can read all files
        RETURN 1
    ELSE IF currUID == ownUID AND mode & U_READ // Owner can read ?
        RETURN 1
    ELSE IF currGID == ownGID AND mode & G_READ // Group can read ?
        RETURN 1
    ELSE IF currUID IN fMembers AND mode & G_READ // Group can read ?
        RETURN 1
    ELSE if mode & O_READ // Others can read ?
       RETURN 1
    ENDIF
    RETURN 0
```

- Call allowRead() in sys_open() of sysfile.c to ensure that read permission is allowed for current user.
- Similarly, create allowWrite() and use it in sys_write().
- Similarly, create
 allowExec() function and
 check for permission in
 exec.c

File System Protection - chown(), chgrp() and chmod() system call

- System calls to change ownership
- Implementation details: int chown(char* path, int uid)
 - 1. Open inode of file at path.
 - 2. Check if current user is root or owner of file. If not, return -1
 - 3. Update the uid file of inode
- chgrp() and chmod() implemented in the same fashion.
- Also add user programs, chown.c, chgrp.c and chmod.c to allow user to call these from shell process.

File System Protection - groupadd and usermod

- Similar to useradd user program. Implementation of groupadd.c user program:
 - Open /group file. (Will be automatically disallowed if not root)
 - Find minimum unused gid.
 - Create new entry in /group file with given groupname and gid.
- usermod.c user program to add user to a group:
 - Open /group file.
 - Seek to the line having entered group details.
 - In members field of line, add a comma separated entry for new user

Thank you!