

# Dependent Types and Theorem Proving: Introduction to Dependent Types

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## 1 Inductive types

# Inductive types are discriminated unions

- Recall how ordinary inductive types work in F# (where they are called discriminated unions).
- To define an inductive type  $I : \text{Type}$ , we list its constructors.
- The constructors are ordinary functions which take some arguments and return an element of the type.
- To create an element of  $I$ , we use one of the constructors and provide it with the arguments it requires.
- To use an element of  $I$ , we pattern match on it and for each case we provide an expression which will be computed if that case matches.

# Inductive families

- Now watch the analogy unfold...
- To define an **inductive family**  $I : (\#a : \text{Type}) \rightarrow a \rightarrow \text{Type}$ , we list its constructors.
- **The constructors are dependent functions which take some arguments and return an element of the type  $I\ x$ , where  $x : a$  was constructed from the arguments.**
- To create an element of  $I\ x$ , we use one of the constructors and provide it with the arguments it requires.
- To use an element of  $I\ x$ , we pattern match on it and for each case we provide an expression which will be computed if that case matches.

## Code snippet no 5 - inductive families in $F^*$

- Let's see how inductive families work in  $F^*$ .
- See the code snippet `Lecture1/InductiveFamilies.fst`

## The running summary 5

- Dependent types are types that can depend on values.
- In dependently typed languages:
- There is a universe – a type whose elements are themselves types.
- There is a type of dependent functions which are just like ordinary functions, but their output TYPE can depend on the VALUE of their input.
- Dependent record types are just like ordinary records, but the TYPES of later fields can depend on the VALUE of earlier fields.
- Inductive families are just like ordinary inductive types, but the TYPES in the family can depend on the VALUE of the index.

# Inductive types and polynomials 1/2

- An inductive type is **EITHER** constructor 1 applied to arguments  $x_1$  **and**  $x_2 \dots$  **and**  $x_N$  **OR** constructor 2 applied to arguments  $\dots$  **OR** constructor  $M$  applied to arguments  $\dots$
- In math, OR means **addition**, whereas AND means **multiplication**.
- So, an inductive type boils down to a **Sum of Products**.
- These products are made of two kinds of arguments: recursive arguments (whose type is the inductive type that is being defined) and non-recursive ones.
- If you think about it long enough, **inductive types correspond to polynomials!**

# Inductive types and polynomials 2/2

- This could be hard to swallow, so let's see examples.
- Lists satisfy the equation  $\text{List}(A) = 1 + A \times \text{List}(A)$ .
- Here 1 corresponds to the `nil` constructor, whereas the  $A$  and  $\text{List}(A)$  on the right correspond to the arguments of the `cons` constructor.
- This corresponds to the polynomial  $F(X) = 1 + A \times X$ .
- $\text{List}(A)$  is the least fixed point of this polynomial, i.e. the smallest type  $X$  that satisfies  $F(X) = X$ .
- Here “fixed point” corresponds to the fact that we create lists using constructors (`nil` and `cons`), whereas “least” corresponds to the fact that all lists are made of finitely many constructors.