

# Sparse Optimization

## Lecture: Basic Sparse Optimization Models

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online discussions on piazza.com

Those who complete this lecture will know

- basic  $\ell_1$ ,  $\ell_{2,1}$ , and nuclear-norm models
- some applications of these models
- how to reformulate them into standard conic programs
- which conic programming solvers to use

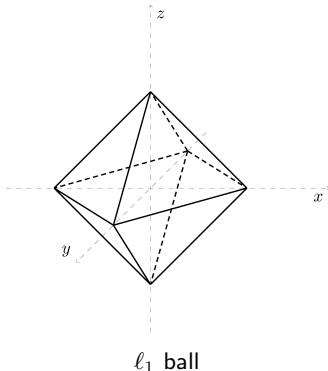
# Examples of Sparse Optimization Applications

See online seminar at [piazza.com](http://piazza.com)

# Basis pursuit

$$\min\{\|\mathbf{x}\|_1 : \mathbf{Ax} = \mathbf{b}\}$$

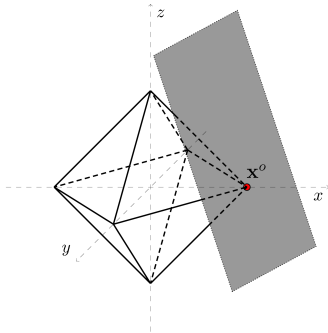
- find least  $\ell_1$ -norm point on the affine plane  $\{\mathbf{x} : \mathbf{Ax} = \mathbf{b}\}$
- tends to return a sparse point (sometimes, the sparsest)



# Basis pursuit

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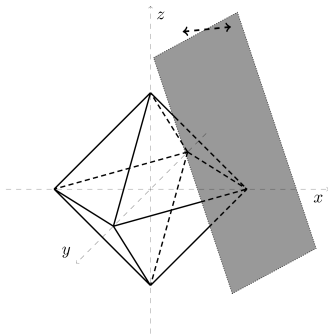
$\ell_1$  ball touches the affine plane

## Basis pursuit denoising, LASSO

$$\min_{\mathbf{x}} \{ \|\mathbf{Ax} - \mathbf{b}\|_2 : \|\mathbf{x}\|_1 \leq \tau \}, \quad (1a)$$

$$\min_{\mathbf{x}} \|\mathbf{x}\|_1 + \frac{\mu}{2} \|\mathbf{Ax} - \mathbf{b}\|_2^2, \quad (1b)$$

$$\min_{\mathbf{x}} \{ \|\mathbf{x}\|_1 : \|\mathbf{Ax} - \mathbf{b}\|_2 \leq \sigma \}. \quad (1c)$$



all models allow  $\mathbf{Ax}^* \neq \mathbf{b}$

## Basis pursuit denoising, LASSO

$$\min_{\mathbf{x}} \{ \|\mathbf{Ax} - \mathbf{b}\|_2 : \|\mathbf{x}\|_1 \leq \tau \}, \quad (2a)$$

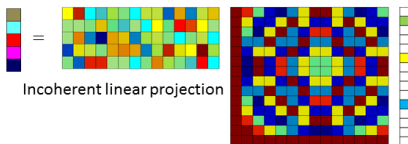
$$\min_{\mathbf{x}} \|\mathbf{x}\|_1 + \frac{\mu}{2} \|\mathbf{Ax} - \mathbf{b}\|_2^2, \quad (2b)$$

$$\min_{\mathbf{x}} \{ \|\mathbf{x}\|_1 : \|\mathbf{Ax} - \mathbf{b}\|_2 \leq \sigma \}. \quad (2c)$$

- $\|\cdot\|_2$  is most common for error but can be generalized to loss function  $\mathcal{L}$
- (2a) seeks for a least-squares solution with “bounded sparsity”
- (2b) is known as LASSO (least absolute shrinkage and selection operator). it seeks for a balance between sparsity and fitting
- (2c) is referred to as BPDN (basis pursuit denoising), seeking for a sparse solution from tube-like set  $\{\mathbf{x} : \|\mathbf{Ax} - \mathbf{b}\|_2 \leq \sigma\}$
- they are equivalent (see later slides)
- in terms of regression, they select a (sparse) set of features (i.e., columns of  $\mathbf{A}$ ) to linearly express the observation  $\mathbf{b}$

## Sparse under basis $\Psi$ / $\ell_1$ -synthesis model

$$\min_{\mathbf{s}} \{\|\mathbf{s}\|_1 : \mathbf{A}\Psi\mathbf{s} = \mathbf{b}\} \quad (3)$$



- signal  $\mathbf{x}$  is *sparsely synthesized* by atoms from  $\Psi$ , so vector  $\mathbf{s}$  is sparse
- $\Psi$  is referred to as the *dictionary*
- commonly used dictionaries include both analytic and trained ones
- analytic examples: Id, DCT, wavelets, curvelets, gabor, etc., also their combinations; they have analytic properties, often easy to compute (for example, multiplying a vector takes  $O(n \log n)$  instead of  $O(n^2)$ )
- $\Psi$  can also be numerically learned from *training data* or *partial signal*
- they can be orthogonal, frame, or general

## Sparse under basis $\Psi$ / $\ell_1$ -synthesis model

If  $\Psi$  is **orthogonal**, problem (3) is equivalent to

$$\min_{\mathbf{x}} \{ \|\Psi^* \mathbf{x}\|_1 : \mathbf{A} \mathbf{x} = \mathbf{b} \} \quad (4)$$

by change of variable  $\mathbf{x} = \Psi \mathbf{s}$ , equivalently  $\mathbf{s} = \Psi^* \mathbf{x}$ .

Related models for noise and approximate sparsity:

$$\begin{aligned} & \min_{\mathbf{x}} \{ \|\mathbf{A} \mathbf{x} - \mathbf{b}\|_2 : \|\Psi^* \mathbf{x}\|_1 \leq \tau \}, \\ & \min_{\mathbf{x}} \|\Psi^* \mathbf{x}\|_1 + \frac{\mu}{2} \|\mathbf{A} \mathbf{x} - \mathbf{b}\|_2^2, \\ & \min_{\mathbf{x}} \{ \|\Psi^* \mathbf{x}\|_1 : \|\mathbf{A} \mathbf{x} - \mathbf{b}\|_2 \leq \sigma \}. \end{aligned}$$



## Sparse after transform / $\ell_1$ -analysis model

$$\min_{\mathbf{x}} \{ \|\Psi^* \mathbf{x}\|_1 : \mathbf{A}\mathbf{x} = \mathbf{b} \} \quad (5)$$

Signal  $\mathbf{x}$  becomes sparse under the transform  $\Psi$  (may not be orthogonal)

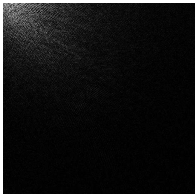
**Examples** of  $\Psi$ :

- DCT, wavelets, curvelets, ridgelets, ....
- tight frames, Gabor, ...
- (weighted) total variation

When  $\Psi$  is not orthogonal, the analysis is more difficult

## Example: sparsify an image





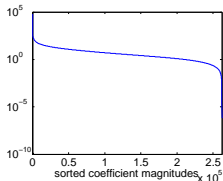
(a) DCT coefficients



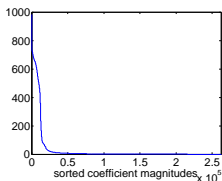
(b) Haar wavelets



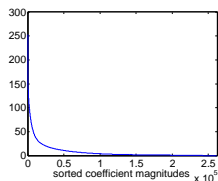
(c) Local variation



(d) DCT coeff's decay



(e) Haar wavelets



(f) Local variation

**Figure:** the DCT and wavelet coefficients are scaled for better visibility.

# Questions

1. Can we trust these models to return intended sparse solutions?
2. When will the solution be unique?
3. Will the solution be robust to noise in  $\mathbf{b}$ ?
4. Are constrained and unconstrained models equivalent? in what sense?

Questions 1–4 will be addressed in *next lecture*.

5. How to choose parameters?
  - $\tau$  (sparsity),  $\mu$  (weight), and  $\sigma$  (noise level) have different meanings
  - applications determine which one is easier to set
  - generality: use a test data set, then scale parameters for real data
  - cross validation: reserve a subset of data to test the solution

## Joint/group sparsity

Joint sparse recovery model:

$$\min_{\mathbf{X}} \{ \|\mathbf{X}\|_{2,1} : \mathcal{A}(\mathbf{X}) = \mathbf{b} \} \quad (6)$$

where

$$\|\mathbf{X}\|_{2,1} := \sum_{i=1}^m \|[x_{i1} \ x_{i,2} \cdots x_{in}]\|_2.$$

- $\ell_2$ -norm is applied to each row of  $\mathbf{X}$
- $\ell_{2,1}$ -norm ball has sharp boundaries “across different rows”, which tend to be touched by  $\{\mathbf{X} : \mathcal{A}(\mathbf{X}) = \mathbf{b}\}$ , so the solution tends to be *row-sparse*
- also  $\|\mathbf{X}\|_{p,q}$  for  $1 < p \leq \infty$ , affects magnitudes of entries on the same row
- complex-valued signals are a special case

## Joint/group sparsity

Decompose  $\{1, \dots, n\} = \mathcal{G}_1 \cup \mathcal{G}_2 \cup \dots \cup \mathcal{G}_S$ .

- non-overlapping groups:  $\mathcal{G}_i \cap \mathcal{G}_j = \emptyset, \forall i \neq j$ .
- otherwise, groups may overlap (modeling many interesting structures).

Group-sparse recovery model:

$$\min_{\mathbf{x}} \{\|\mathbf{x}\|_{\mathcal{G},2,1} : \mathbf{A}\mathbf{x} = \mathbf{b}\} \quad (7)$$

where

$$\|\mathbf{x}\|_{\mathcal{G},2,1} = \sum_{s=1}^S w_s \|\mathbf{x}_{\mathcal{G}_s}\|_2.$$

## Auxiliary constraints

Auxiliary constraints introduce additional structures of the underlying signal into its recovery, which sometimes *significantly* improve recovery quality

- nonnegativity:  $\mathbf{x} \geq \mathbf{0}$
- bound (box) constraints:  $\mathbf{l} \leq \mathbf{x} \leq \mathbf{u}$
- general inequalities:  $\mathbf{Q}\mathbf{x} \leq \mathbf{q}$

They can be very effective in practice. They also generate “corners.”

## Reduce to conic programs

Sparse optimization often has *nonsmooth* objectives.

Classic conic programming solvers do not handle nonsmooth functions.

Basic idea: *model nonsmoothness by inequality constraints*.

Example: for given  $\mathbf{x}$ , we have

$$\|\mathbf{x}\|_1 = \min_{\mathbf{x}_1, \mathbf{x}_2} \{\mathbf{1}^T(\mathbf{x}_1 + \mathbf{x}_2) : \mathbf{x}_1 - \mathbf{x}_2 = \mathbf{x}, \mathbf{x}_1 \geq \mathbf{0}, \mathbf{x}_2 \geq \mathbf{0}\}. \quad (8)$$

Therefore,

- $\min\{\|\mathbf{x}\|_1 : \mathbf{Ax} = \mathbf{b}\}$  reduces to a linear program (LP)
- $\min_{\mathbf{x}} \|\mathbf{x}\|_1 + \frac{\mu}{2} \|\mathbf{Ax} - \mathbf{b}\|_2^2$  reduces to a bound constrained quadratic program (QP)
- $\min_{\mathbf{x}} \{\|\mathbf{Ax} - \mathbf{b}\|_2 : \|\mathbf{x}\|_1 \leq \tau\}$  reduces to a bound constrained QP
- $\min_{\mathbf{x}} \{\|\mathbf{x}\|_1 : \|\mathbf{Ax} - \mathbf{b}\|_2 \leq \sigma\}$  reduces to a second-order cone program (SOCP)



# Conic programming

Basic form:

$$\min_{\mathbf{x}} \{ \mathbf{c}^T \mathbf{x} : \mathbf{F} \mathbf{x} + \mathbf{g} \succeq_{\mathcal{K}} \mathbf{0}, \mathbf{A} \mathbf{x} = \mathbf{b}. \}$$

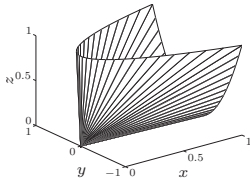
" $\mathbf{a} \succeq_{\mathcal{K}} \mathbf{b}$ " stands for  $\mathbf{a} - \mathbf{b} \in \mathcal{K}$ , which is a convex, closed, pointed cone.

Examples:

- first orthant (cone):  $\mathbb{R}_+^n = \{\mathbf{x} \in \mathbb{R}^n : \mathbf{x} \geq \mathbf{0}\}$ .
- norm cone (2nd order cone):  $\mathcal{Q} = \{(\mathbf{x}, t) : \|\mathbf{x}\| \leq t\}$
- polyhedral cone:  $\mathcal{P} = \{\mathbf{x} : \mathbf{A} \mathbf{x} \geq \mathbf{0}\}$
- positive semidefinite cone:  $\mathbf{S}_+ = \{\mathbf{X} : \mathbf{X} \succeq \mathbf{0}, \mathbf{X}^T = \mathbf{X}\}$

**Example:**

$$\left\{ (x, y, z) : \begin{bmatrix} x & y \\ y & z \end{bmatrix} \in \mathbf{S}_+ \right\}$$



# Linear program

Model

$$\min\{\mathbf{c}^T \mathbf{x} : \mathbf{Ax} = \mathbf{b}, \mathbf{x} \succeq_{\mathcal{K}} \mathbf{0}\}$$

where  $\mathcal{K}$  is the nonnegative cone (first orthant).

$$\mathbf{x} \succeq_{\mathcal{K}} \mathbf{0} \iff \mathbf{x} \geq \mathbf{0}.$$

Algorithms

- the Simplex method (move between vertices)
- interior-point methods (IPMs) (move inside the polyhedron)
- decomposition approaches (divide and conquer)

In primal IPM,  $\mathbf{x} \geq 0$  is replaced by its logarithmic barrier:

$$\psi(\mathbf{y}) = \sum_i \log(y_i)$$

log-barrier formulation:

$$\min\{\mathbf{c}^T \mathbf{x} - (1/t) \sum_i \log(x_i) : \mathbf{Ax} = \mathbf{b}\}$$

## Second-order cone program

Model

$$\min\{\mathbf{c}^T \mathbf{x} : \mathbf{Ax} = \mathbf{b}, \mathbf{x} \succeq_{\mathcal{K}} \mathbf{0}\}$$

where  $\mathcal{K} = \mathcal{K}_1 \times \cdots \times \mathcal{K}_K$ ; each  $\mathcal{K}_k$  is the second-order cone

$$\mathcal{K}_k = \left\{ \mathbf{y} \in \mathbb{R}^{n_k} : y_{n_k} \geq \sqrt{y_1^2 + \cdots + y_{n_k-1}^2} \right\}.$$

IPM is the standard solver (though other options also exist)

Log-barrier of  $\mathcal{K}_k$ :

$$\psi(\mathbf{y}) = \log(y_{n_k}^2 - (y_1^2 + \cdots + y_{n_k-1}^2))$$

# Semi-definite program

Model

$$\min\{\mathbf{C} \bullet \mathbf{X} : \mathcal{A}(\mathbf{X}) = \mathbf{b}, \mathbf{X} \succeq_{\mathcal{K}} \mathbf{0}\}$$

where  $\mathcal{K} = \mathcal{K}_1 \times \cdots \times \mathcal{K}_K$ ; each  $\mathcal{K}_k = \mathbf{S}_+^{n_k}$ .

IPM is the standard solver (though other options also exist)

Log-barrier of  $\mathbf{S}_+^{n_k}$  (still a concave function):

$$\psi(\mathbf{Y}) = \log \det(\mathbf{Y}).$$

(from Boyd & Vandenberghe, *Convex Optimization*)

**properties** (without proof): for  $y \succ_K 0$ ,

$$\nabla\psi(y) \succeq_{K^*} 0, \quad y^T \nabla\psi(y) = \theta$$

- nonnegative orthant  $\mathbf{R}_+^n$ :  $\psi(y) = \sum_{i=1}^n \log y_i$

$$\nabla\psi(y) = (1/y_1, \dots, 1/y_n), \quad y^T \nabla\psi(y) = n$$

- positive semidefinite cone  $\mathbf{S}_+^n$ :  $\psi(Y) = \log \det Y$

$$\nabla\psi(Y) = Y^{-1}, \quad \text{tr}(Y \nabla\psi(Y)) = n$$

- second-order cone  $K = \{y \in \mathbf{R}^{n+1} \mid (y_1^2 + \dots + y_n^2)^{1/2} \leq y_{n+1}\}$ :

$$\nabla\psi(y) = \frac{2}{y_{n+1}^2 - y_1^2 - \dots - y_n^2} \begin{bmatrix} -y_1 \\ \vdots \\ -y_n \\ y_{n+1} \end{bmatrix}, \quad y^T \nabla\psi(y) = 2$$

(from Boyd & Vandenberghe, *Convex Optimization*)

## Central path

- for  $t > 0$ , define  $x^*(t)$  as the solution of

$$\begin{array}{ll}\text{minimize} & t f_0(x) + \phi(x) \\ \text{subject to} & Ax = b\end{array}$$

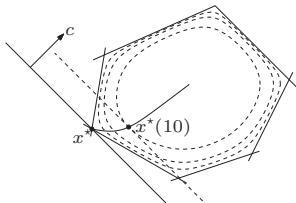
(for now, assume  $x^*(t)$  exists and is unique for each  $t > 0$ )

- central path is  $\{x^*(t) \mid t > 0\}$

**example:** central path for an LP

$$\begin{array}{ll}\text{minimize} & c^T x \\ \text{subject to} & a_i^T x \leq b_i, \quad i = 1, \dots, 6\end{array}$$

hyperplane  $c^T x = c^T x^*(t)$  is tangent to level curve of  $\phi$  through  $x^*(t)$



Log-barrier formulation:

$$\min\{tf_0(\mathbf{x}) + \phi(\mathbf{x}) : \mathbf{Ax} = \mathbf{b}\}$$

Complexity of log-barrier interior-point method:

$$k \sim \left\lceil \frac{\log((\sum_i \theta_i)/(\varepsilon t^{(0)}))}{\log \mu} \right\rceil$$

## Primal-dual interior-point methods

more efficient than barrier method when high accuracy is needed

- update primal and dual variables at each iteration; no distinction between inner and outer iterations
- often exhibit superlinear asymptotic convergence
- search directions can be interpreted as Newton directions for modified KKT conditions
- can start at infeasible points
- cost per iteration same as barrier method



## $\ell_1$ minimization by interior-point method

Model

$$\min_{\mathbf{x}} \{ \|\mathbf{x}\|_1 : \|\mathbf{Ax} - \mathbf{b}\|_2 \leq \sigma \}$$

$\Leftrightarrow$

$$\min_{\mathbf{x}} \min_{\mathbf{x}_1, \mathbf{x}_2} \{ \mathbf{1}^T (\mathbf{x}_1 + \mathbf{x}_2) : \mathbf{x}_1 \geq \mathbf{0}, \mathbf{x}_2 \geq \mathbf{0}, \mathbf{x}_1 - \mathbf{x}_2 = \mathbf{x}, \|\mathbf{Ax} - \mathbf{b}\|_2 \leq \sigma \}$$

$\Leftrightarrow$

$$\min_{\mathbf{x}_1, \mathbf{x}_2} \{ \mathbf{1}^T (\mathbf{x}_1 + \mathbf{x}_2) : \mathbf{x}_1 \geq \mathbf{0}, \mathbf{x}_2 \geq \mathbf{0}, \|\mathbf{A}(\mathbf{x}_1 - \mathbf{x}_2) - \mathbf{b}\|_2 \leq \sigma \}$$

$\Leftrightarrow$

$$\min_{\mathbf{x}_1, \mathbf{x}_2} \{ \mathbf{1}^T \mathbf{x}_1 + \mathbf{1}^T \mathbf{x}_2 : \mathbf{Ax}_1 - \mathbf{Ax}_2 + \mathbf{y} = \mathbf{b}, z = \sigma, (\mathbf{x}_1, \mathbf{x}_2, z, \mathbf{y}) \succeq_{\mathcal{K}} \mathbf{0} \}$$

where  $(\mathbf{x}_1, \mathbf{x}_2, z, \mathbf{y}) \succeq_{\mathcal{K}} \mathbf{0}$  means

- $\mathbf{x}_1, \mathbf{x}_2 \in \mathbb{R}_+^n$ ,
- $(t, \mathbf{y}) \in \mathcal{Q}^{m+1}$ .

Solver: Mosek, SDPT3, Gurobi.

Also, modeling language CVX and YALMIP.

# Nuclear-norm minimization by interior-point method

If we can model

$$\min_{\mathbf{X}} \{\|\mathbf{X}\|_* : \mathcal{A}(\mathbf{X}) = \mathbf{b}\} \quad (9)$$

as an SDP ... (how? see next slide) ...

then, we can also model

- $\min_{\mathbf{X}} \{\|\mathbf{X}\|_* : \|\mathcal{A}(\mathbf{X}) - \mathbf{b}\|_F \leq \sigma\}$
- $\min_{\mathbf{X}} \{\|\mathcal{A}(\mathbf{X}) - \mathbf{b}\|_F : \|\mathbf{X}\|_* \leq \tau\}$
- $\min_{\mathbf{X}} \mu \|\mathbf{X}\|_* + \frac{1}{2} \|\mathcal{A}(\mathbf{X}) - \mathbf{b}\|_F^2$

as well as problems involving  $\text{tr}(\mathbf{X})$  and spectral norm  $\|\mathbf{X}\|$ .

$$\|\mathbf{X}\| \leq \alpha \iff \alpha I - \mathbf{X} \succeq \mathbf{0}.$$

## Sparse calculus for $\ell_1$

- inspect  $|x|$  to get some ideas:

$$y, z \geq 0 \text{ and } \sqrt{yz} \geq |x| \implies \frac{1}{2}(y+z) \geq \sqrt{yz} \geq |x|.$$

$$\text{moreover, } \frac{1}{2}(y+z) = \sqrt{yz} = |x| \text{ if } y = z = |x|.$$

- observe

$$y, z \geq 0 \text{ and } \sqrt{yz} \geq |x| \iff \begin{bmatrix} y & x \\ x & z \end{bmatrix} \succeq \mathbf{0}.$$

So,

$$\begin{bmatrix} y & x \\ x & z \end{bmatrix} \succeq \mathbf{0} \implies \frac{1}{2}(y+z) \geq |x|.$$

- we attain  $\frac{1}{2}(y+z) = |x|$  if  $y = z = |x|$ .

Therefore, given  $x$ , we have

$$|x| = \min_{\mathbf{M}} \left\{ \frac{1}{2} \text{tr}(\mathbf{M}) : \begin{bmatrix} 0 & 1 \\ 0 & 0 \end{bmatrix} \bullet \mathbf{M} = x, \mathbf{M} = \mathbf{M}^T, \mathbf{M} \succeq \mathbf{0} \right\}.$$

## Generalization to nuclear norm

- Consider  $\mathbf{X} \in \mathbb{R}^{m \times n}$  (w.o.l.g., assume  $m \leq n$ ) and let's try imposing

$$\begin{bmatrix} \mathbf{Y} & \mathbf{X} \\ \mathbf{X}^T & \mathbf{Z} \end{bmatrix} \succeq \mathbf{0}$$

- Diagonalize  $\mathbf{X} = \mathbf{U}\Sigma\mathbf{V}^T$ ,  $\Sigma = \text{diag}(\sigma_1, \dots, \sigma_m)$ ,  $\|\mathbf{X}\|_* = \sum_i \sigma_i$ .

$$\begin{aligned} [\mathbf{U}^T, -\mathbf{V}^T] \begin{bmatrix} \mathbf{Y} & \mathbf{X} \\ \mathbf{X}^T & \mathbf{Z} \end{bmatrix} \begin{bmatrix} \mathbf{U} \\ -\mathbf{V} \end{bmatrix} &= \mathbf{U}^T \mathbf{Y} \mathbf{U} + \mathbf{V}^T \mathbf{Z} \mathbf{V} - \mathbf{U}^T \mathbf{X} \mathbf{V} - \mathbf{V}^T \mathbf{X}^T \mathbf{U} \\ &= \mathbf{U}^T \mathbf{Y} \mathbf{U} + \mathbf{V}^T \mathbf{Z} \mathbf{V} - 2\Sigma \succeq \mathbf{0}. \end{aligned}$$

So,  $\text{tr}(\mathbf{U}^T \mathbf{Y} \mathbf{U} + \mathbf{V}^T \mathbf{Z} \mathbf{V} - 2\Sigma) = \text{tr}(\mathbf{Y}) + \text{tr}(\mathbf{Z}) - 2\|\mathbf{X}\|_* \geq 0$ .

- To attain "=", we can let  $\mathbf{Y} = \mathbf{U}\Sigma\mathbf{U}^T$  and  $\mathbf{Z} = \mathbf{V}\Sigma_{n \times n}\mathbf{V}^T$ .

Therefore,

$$\|\mathbf{X}\|_* = \min_{\mathbf{Y}, \mathbf{Z}} \left\{ \frac{1}{2}(\text{tr}(\mathbf{Y}) + \text{tr}(\mathbf{Z})) : \begin{bmatrix} \mathbf{Y} & \mathbf{X} \\ \mathbf{X}^T & \mathbf{Z} \end{bmatrix} \succeq \mathbf{0} \right\} \quad (10)$$

$$= \min_{\mathbf{M}} \left\{ \frac{1}{2} \text{tr}(\mathbf{M}) : \begin{bmatrix} \mathbf{0} & \mathbf{I} \\ \mathbf{0} & \mathbf{0} \end{bmatrix} \bullet \mathbf{M} = \mathbf{X}, \mathbf{M} = \mathbf{M}^T, \mathbf{M} \succeq \mathbf{0} \right\}. \quad (11)$$

Exercise: express the following problems as SDPs

- $\min_{\mathbf{X}} \{\|\mathbf{X}\|_* : \mathcal{A}(\mathbf{X}) = \mathbf{b}\}$
- $\min_{\mathbf{X}} \mu \|\mathbf{X}\|_* + \frac{1}{2} \|\mathcal{A}(\mathbf{X}) - \mathbf{b}\|_F$
- $\min_{\mathbf{L}, \mathbf{S}} \{\|\mathbf{L}\|_* + \lambda \|\mathbf{S}\|_1 : \mathcal{A}(\mathbf{L} + \mathbf{S}) = \mathbf{b}\}$

## Practice of interior-point methods (IPMs)

- LP, SOCP, SDP are well known and have reliable (commercial, off-the-shelf) solvers
- Yet, the most reliable solvers cannot handle large-scale problems (e.g., images, video, manifold learning, distributed stuff, ...)
  - Example: to recover a still image, there can be 10M variables and 1M constraints. Even worse, the constraint coefficients are dense. Result: Out of memory.
- Simplex and active-set methods: matrix containing  $\mathbf{A}$  must be inverted or factorized to compute the next point (unless  $\mathbf{A}$  is very sparse).
- IPMs approximately solve a Newton system and thus also factorize a matrix involving  $\mathbf{A}$ .
- Even large and dense matrices can be handled, for sparse optimization, one should take advantages of the solution sparsity.
- Some compressive sensing problems have  $\mathbf{A}$  with structures friendly for operations like  $\mathbf{A}\mathbf{x}$  and  $\mathbf{A}^T\mathbf{y}$ .

## Practice of interior-point methods (IPMs)

- The Simplex, active-set, and IPMs have *reliable* solvers; good to be the benchmark.
- They have nice interfaces (including *CVX* and *YALMIP*, which save you time.)  
*CVX* and *YALMIP* are not solvers; they translate problems and then call solvers; see <http://goo.gl/zU1MK> and <http://goo.gl/1u0xP>.
- They can return *highly accurate* solutions; some first-order algorithms (coming later in this course) do not always.
- There are other remedies; see next slide.

# Papers of large-scale SDPs

- Low-rank factorizations:
  - S. Burer and R. D. C. Monteiro, A nonlinear programming algorithm for solving semidefinite programs via low-rank factorization, Math. Program., 95:329–357, 2003.
  - LMaFit, <http://lmafit.blogs.rice.edu/>
- First-order methods for conic programming:
  - Z. Wen, D. Goldfarb, and W. Yin. Alternating direction augmented Lagrangian methods for semidefinite programming. Math. Program. Comput., 2(3-4):203–230, 2010.
- Matrix-free IPMs:
  - K. Fountoulakis, J. Gondzio, P. Zhlobich. Matrix-free interior point method for compressed sensing problems, 2012. <http://www.maths.ed.ac.uk/~gondzio/reports/mfCS.html>



## Subgradient methods

Sparse optimization is typically nonsmooth, so it is natural to consider subgradient methods.

- apply subgradient descent to, say,  $\min_{\mathbf{x}} \|\mathbf{x}\|_1 + \frac{\mu}{2} \|\mathbf{Ax} - \mathbf{b}\|_2^2$ .
- apply projected subgradient descent to, say,  $\min_{\mathbf{x}} \{\|\mathbf{x}\|_1 : \mathbf{Ax} = \mathbf{b}\}$ .

Good: subgradients are easy to derive, methods are simple to implement.

Bad: convergence requires carefully chosen step sizes (classical ones require diminishing step sizes). Convergence rate is weak on paper (and in practice, too?)

Further readings: <http://arxiv.org/pdf/1001.4387.pdf>,  
<http://goo.gl/qFVA6>, <http://goo.gl/vC21a>.