Ontology Languages Beyond ALC



Objectives



Objective

Explain additional constructs in description logics and how they add expressivity

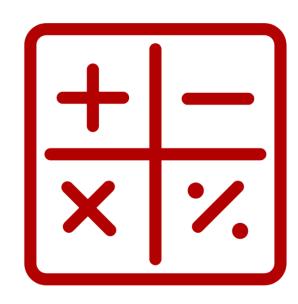
Number Restrictions

Role restriction cannot express that a teacher teaches at least 3 (or at most 5) courses

Number restrictions can express arithmetic constraints on the number of fillers of a role

Examples:

- BusyTeacher ≡ Teacher □ (≥ 3 teaches)
- ConsciousTeacher ≡ Teacher □ (≤5 teaches)



Number Restrictions, cont'd

Semantics:

$$-(\ge n R)^{I} = \{x \in \Delta^{I} \mid |\{y \mid (x, y) \in R^{I}\}| \ge n\}$$

$$-(\le n R)^{I} = \{x \in \Delta^{I} \mid |\{y \mid (x, y) \in R^{I}\}| \le n\}$$

Observation: $(\geq 1 R) \equiv \exists R. \top$

$$\pi_{x}(\geq 2R)$$

$$= \exists y_{1}y_{2}(R(x_{1})\wedge R(x_{1})) \wedge R(x_{1}) \wedge y_{1} \neq y_{2})$$

$$\pi_{x}(\leq 2R)$$

$$= \forall y_{1}y_{2}y_{3}(R(x_{1})\wedge R(x_{1})) \wedge R(x_{1}) \wedge R(x_{2})$$

$$\wedge R(x_{1},y_{3})$$

y== 43 V

4>=41)

The above number restrictions are called unqualified. \rightarrow ($y_i = y_i \lor$

$$-\pi_{x}(\geq n R) = \exists y_{1} \dots y_{n}(\bigwedge_{k=1}^{n} R(x, y_{k}) \land \bigwedge_{i < j} y_{j} \neq y_{i})$$

$$-\pi_{x}(\leq n R) = \forall y_{1} \dots y_{n+1}(\bigwedge_{k=1}^{n+1} R(x, y_{k}) \rightarrow \bigvee_{i < j} y_{i} = y_{i})$$

Notation: The description logic which extends ALC with number restrictions is denoted by ALCN.

Qualified Number Restrictions

Examples:

- BusyTeacher \equiv Teacher \sqcap (≥ 3 teaches. Course)
- Conscious Teacher □ (\le 5 teaches. Course)

Semantics:

- $(\ge n R.C)^I = \{x \in \Delta^I \mid \text{at least } n R^I \text{ successors of } x \text{ are in } C^I \}$
- $(\le n R.C)^I = \{x \in \Delta^I \mid \text{at most } n R^I \text{ successors of } x \text{ are in } C^I \}$

Notation: The description logic which extends ALC with qualified number restrictions is denoted by ALCQ.

Enumerations - Nominals

Sometimes it is useful to define a concept that contains exactly the individuals I_1, \ldots, I_m . This concept is written as $\{I_1, \ldots, I_m\}$.

Examples:

- Weekday ≡ {MON, TUE, WED, THU, FRI, SAT, SUN}
- Citizen ≡ Person □ ∃hasCountry.Country
- German ≡ Citizen □ ∃hasCountry.{Germany}

Nominals

A nominal is a concept that contains exactly one individual.

If we have the ability to define nominals, then using ⊔, we can define concepts containing more than one individual.

Example:

- Weekend ≡ {SAT, SUN}

Nominals, cont'd

If we have nominals and we do not want to make the UNA, then we can explicitly state whether two individuals are the same or different.

Examples:

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- {US} \equiv {USA}, {US} \sqcap {Mexico} \sqsubseteq \bot
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Notation: If we add nominals to DL ALCQ, we get the DL ALCQO.

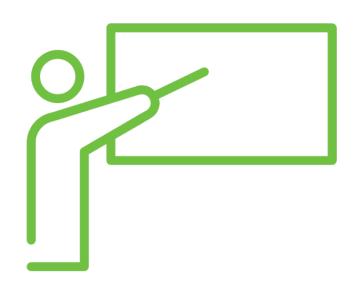
Example

Let a knowledge base K be:

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BusyTeacher ≡ Teacher □ (≥3 teaches)
ConsciousTeacher ≡ Teacher □ (≤ 5 teaches)
Teacher(MARY),
teaches(MARY, AI), teaches(MARY, KR),
teaches(MARY, DB)
```

Questions:

- K ⊨ BusyTeacher(MARY) ?
- K ⊨ ConsciousTeacher(MARY) ?



Riddle: Two Fathers and Two Sons

Two fathers and two sons go fishing together in the same boat. They all catch a fish but the total catch for the day is three fish. How is this possible?

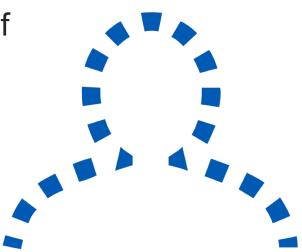
OWS and **CWA**

Contrary to databases, DLs make the Open World Assumption.

 Absence of information is not interpreted as presence of negative information but simply as lack of knowledge.

Thus in the previous example:

K ⊭ ConsciousTeacher(MARY)



Wrap-Up

