Spectral Clustering: A Graph Cut Formulation



Objective



Define the graph partition formulation



Learn how to solve simple graph partition

Clustering as Graph Partition/Cut

Find a partition of a graph such that the edges between different groups have a very low weight while the edges within a group have high weight.

E.g., minimum cut

CutSize = 2

More general, consider weighted edges.

2-way Spectral Graph Partitioning

Weighted adjacency matrix W

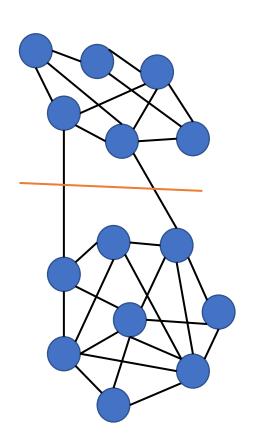
 $| w_{i,j} :$ the weight between two vertices *i* and *j*

(Cluster) Membership vector q

$$q_i = \begin{cases} 1 & i \in Cluster \ A \\ -1 & i \in Cluster \ B \end{cases}$$

$$\mathbf{q} = \underset{\mathbf{q} \in [-1, 1]^n}{\operatorname{argmin}} CutSize$$
,

CutSize =
$$J = \frac{1}{4} \sum_{i,j} (q_i - q_j)^2 w_{i,j}$$



Solving the Optimization Problem

$$\mathbf{q} = \underset{\mathbf{q} \in [-1, 1]^n}{\operatorname{argmin}} \frac{1}{4} \sum_{i,j} (q_i - q_j)^2 w_{i,j},$$

Directly solving the above problem requires combinatorial search \rightarrow exponential complexity

How to reduce the computational complexity?

Relaxation Approach

Key difficulty: q_i has to be either -1,1.

- | Relax q_i to be any real number.
- | Impose Constraint: $\sum_{i=1}^{n} q_i^2 = n$

$$J = \frac{1}{4} \sum_{i,j} (q_i - q_j)^2 w_{i,j} = \frac{1}{4} \sum_{i,j} (q_i^2 - 2q_i q_j + q_j^2) w_{i,j}$$

$$= \frac{1}{4} \sum_i 2q_i^2 (\sum_j w_{i,j}) - \frac{1}{4} \sum_{i,j} 2q_i q_j w_{i,j}$$

$$= \frac{1}{2} \sum_i q_i^2 d_i - \frac{1}{2} \sum_{i,j} q_i (d_i \delta_{i,j} - w_{i,j}) q_j$$

where $d_i = \sum_j w_{i,j}$ and $D \equiv [d_i \delta_{i,j}]$

Relaxation Approach (cont'd)

The final problem formulation:

$$\mathbf{q} = \underset{\mathbf{q}}{\operatorname{argmin}} J = \underset{\mathbf{q}}{\operatorname{argmin}} \mathbf{q}^T (\mathbf{D} - \mathbf{W}) \mathbf{q}$$
, subject to $\sum_{i=1}^n q_i^2 = n$

Solution: the second minimum eigenvector for **D-W**

$$(\mathbf{D} - \mathbf{W})\mathbf{q} = \lambda_2 \mathbf{q}$$

Graph Laplacian

$$L = D - W$$

- L is semi-positive definitive matrix.
 - For any x, we have $\mathbf{x}^T \mathbf{L} \mathbf{x} \geq 0$. (Why?)
- Minimum eigenvalue $\lambda_1 = 0$ (what is the eigenvector?)

$$0 = \lambda_1 < \lambda_2 < \lambda_3 \dots < \lambda_k$$

The eigenvector that corresponds to the second minimum eigenvalue λ_2 gives the best bipartite graph partition.

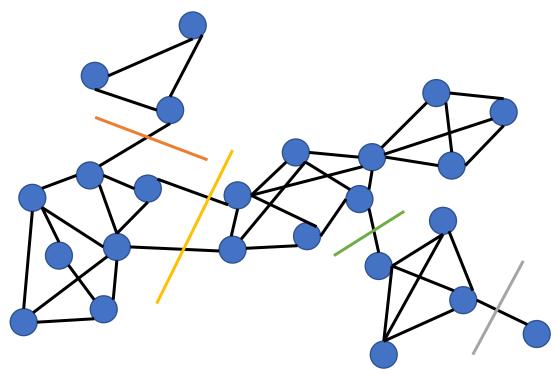
Recovering the Partitions

- Due to the relaxation, **q** can be any number (not just -1 and 1)
- How to construct the partition based on the eigenvector?
- A simple strategy:

$$A = \{i | q_i < 0\}, \quad B = \{i | q_i \ge 0\}$$

One Obvious Drawback

Minimum cut does not balance the size of bipartite graphs.



How should we consider other factors like the sizes of the partitions?