

ERIC BAILEY

GRAPH THEORY EXERCISES

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Graphs

Exercises

1. $\mathcal{P}(\{1, 2, 3\}) := \{\emptyset, \{1\}, \{2\}, \{3\}, \{1, 2\}, \{1, 3\}, \{2, 3\}, \{1, 2, 3\}\}$

2.

$$\frac{\frac{\emptyset \not\subseteq A}{\exists x \in \emptyset : x \notin A} \quad \overline{\neg \exists x \in \emptyset}}{\frac{\perp}{\emptyset \subseteq A}}$$

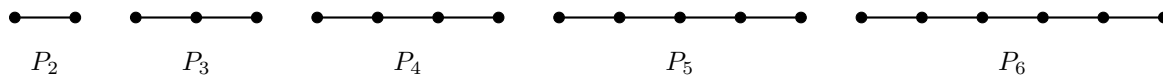
□

3.

$$\frac{\overline{S(b) := \{m \in V : m \notin S(m)\}} \quad \overline{b \in V}}{b \in S(b) \iff b \notin S(b)} \equiv \frac{\overline{R = \{x \notin x\}}}{R \in R \iff R \notin R}$$

4. Let S be the collection of all sets that can be described in an English sentence of twenty-five words or less. S is not a set, because S can be described in fewer than twenty-five words, and if it were a set, then S would have to be a member of itself, which violates the axiomatic definition of a set.

5. The first five *path graphs*.



$$V(P_v) = \{1, 2, \dots, v\}$$

$$E(P_v) = \{\{n-1, n\} : n \in \{2, \dots, v\}\}$$

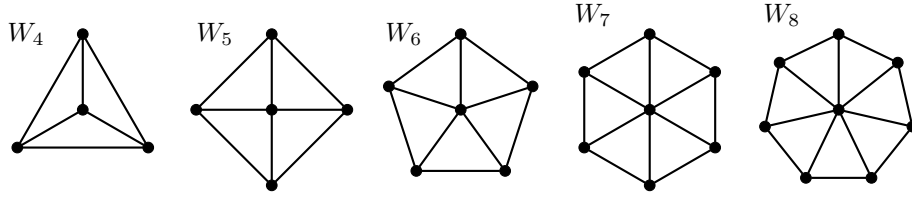
$$|E(P_v)| = v - 1$$

$$e = v - 1$$

The number of edges in a *path graph* P_v , where $v \geq 2$, is given by the formula $e = v - 1$.

□

6.

Figure 1: The first five *wheel graphs*.

$$\begin{aligned}
 V(W_v) &= \{1, 2, 3, \dots, v\} \\
 E(W_v) &= \{ \{1, 2\}, \{1, 3\}, \dots, \{1, v\}, \\
 &\quad \{2, 3\}, \{3, 4\}, \dots, \{v-1, v\}, \\
 &\quad \{v, 2\} \} \\
 &= \{ \{1, n\} : n \in \{2, \dots, v\} \}, \\
 &\quad \{ \{n-1, n\} : n \in \{3, \dots, v\} \}, \\
 &\quad \{v, 2\} \} \\
 |E(W_v)| &= (v-1) + (v-2) + 1 \\
 &= (v-1) + (v-1) \\
 e &= 2(v-1)
 \end{aligned}$$

The number of edges in the *wheel graph* on v vertices W_v , where $v \geq 4$, is given by the formula $e = 2(v-1)$.

□

7.

$$\begin{aligned}
 1 + 2 + \dots + (v-1) &= (1/2)v(v-1) \\
 &= E(K_v) \\
 &= (v-1) + (v-2) + \dots + (v-(v-1)) \\
 &= 1 + 2 + \dots + (v-1)
 \end{aligned}
 \tag{T2}$$

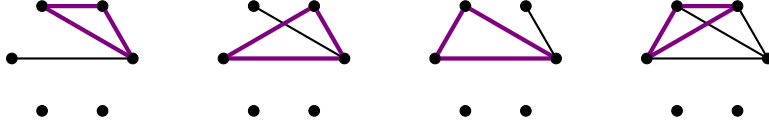
Imagine drawing K_v by joining vertex 1 to vertices 2 through v , creating $v-1$ edges; then joining vertex 2 to vertices 3 through v , creating $v-2$ edges; and so on, i.e. $(v-1) + (v-2) + \dots + (v-(v-1))$, or equivalently, $1 + 2 + \dots + (v-1)$.

□

8. Let G be a graph with v vertices and e edges. In terms of v and e , \overline{G} has $(1/2)v(v-1) - e$ edges.

9. If a graph G has $v = 6$ then G or \overline{G} (possibly both) has a subgraph isomorphic to K_3 .

In the graph G or \overline{G} there exists a vertex a of degree three or more. Let there be three adjacent vertices b, c , and d . If any of the edges $\{b, c\}$, $\{b, d\}$, or $\{c, d\}$ are present in the graph, then the graph contains a subgraph isomorphic to K_3 . If none of those edges is present, then they are all present in the other graph, which thus contains a subgraph isomorphic to K_3 .



10. From the proof in [Exercise 9](#), it follows that in any graph of degree 6 there exists a subgraph H , such that either H or \overline{H} is isomorphic to K_3 . Equivalently, in any gathering of six people (a graph G with $v = 6$) there are either three people who are mutually acquainted (a subgraph isomorphic to K_3) or three people who are mutually unacquainted (a subgraph whose complement is isomorphic to K_3).
- 11.

$$2|\emptyset| = \sum_{v \in V} \deg(v) = 0 \quad (\text{BC})$$

$$G(V, E) \rightarrow 2|E| = \sum_{v \in V} \deg(v) \quad (\text{IH})$$

$$H(V', E'), |E'| = |E| + 1 \quad (1)$$

$$\exists H_0(V', E'_0), |E'_0| = |E| \quad (2)$$

$$2|E'_0| = \sum_{v \in V'} \deg(v) \quad (\text{by IH})$$

$$H \cong (V', E'_0 \cup \{e\}), e \notin E'_0 \quad (3)$$

$$\sum_{v \in V'} \deg(v) + 2 = 2|E| \quad (4)$$

$$2|E| + 2 = 2(|E| + 1) = \sum_{v \in V'} \deg(v) \quad (5)$$

□

12. a)

$$(4 \times 3) + (2 \times 5) + (2 \times 6) + (1 \times 8) = 2e$$

$$12 + 10 + 12 + 8 = 2e$$

$$21 = e$$

b) $(7 \times 3) = 2e \rightarrow e \notin \mathbb{Q}$

13. If $C_n \cong \overline{C_n}$, then it must be the case that $|E(C_n)| + |E(\overline{C_n})| = |E(K_n)|$, i.e. $n + n = \frac{n(n-1)}{2}$.

Since $n = 5$ is a solution, $|E(C_5)| + |E(\overline{C_5})| = |E(K_5)|$, and thus $C_5 \cong \overline{C_5}$.

□

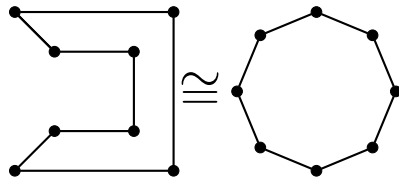
A graph with no edges ($n = 0$) has no cycles. Therefore the only self-complementary cycle graph is C_5 .

□

14. If $G \cong \overline{G}$ for a graph G of order v , then $|E(G)| + |E(\overline{G})| = |E(K_v)|$, i.e. $e + e = 2e = \frac{v(v-1)}{2}$. So the number of edges e in G is $\frac{v(v-1)}{4}$, and $4 \mid v(v-1)$. Therefore, for any self-complementary graph G of order v , $4 \mid v$ or $4 \mid v-1$.

□

15. _____
16. _____
17. P_4 is the self-complementary graph with $v = 4$.
18. Between the two set of 10 elements, $\{a, b, \dots, j\}$ and $\{1, 2, \dots, 10\}$, there exist $10!$, or 3,628,800, one-to-one correspondences.
19. $f : \mathbb{Z}^+ \rightarrow 2\mathbb{Z}^+$, $f(x) = 2n$ is bijective.
20. The first graph in Figure 35 contains a subgraph



which is isomorphic to C_8 , whereas the second graph contains no subgraph isomorphic to C_8 . Therefore, the two graphs are not isomorphic.

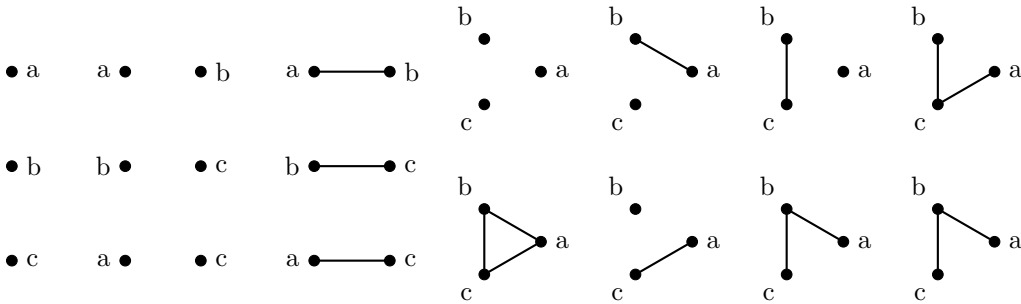
$$\begin{aligned} n + n &= \frac{n(n-1)}{2} \\ 2n &= \frac{n(n-1)}{2} \\ 4n &= n(n-1) \\ 4n &= n^2 - n \\ 0 &= n^2 - 5n \\ n &= \{0, 5\} \end{aligned}$$

Figure 2: Solving for n in [Exercise 13](#)

Find two other self-complementary graphs (Trudeau, *Graph Theory* 57).

Find a self-complementary graph with $v = 8$.

21.

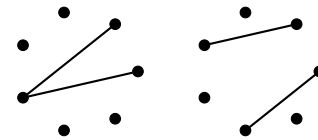
Figure 3: The 17 unequal subgraphs of K_3

22.

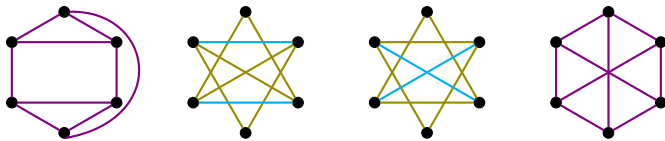
Figure 4: The seven *nonisomorphic* subgraphs of K_3

23. C_3 is a subgraph of the first graph, but not the second. Therefore, the two graphs are not isomorphic.

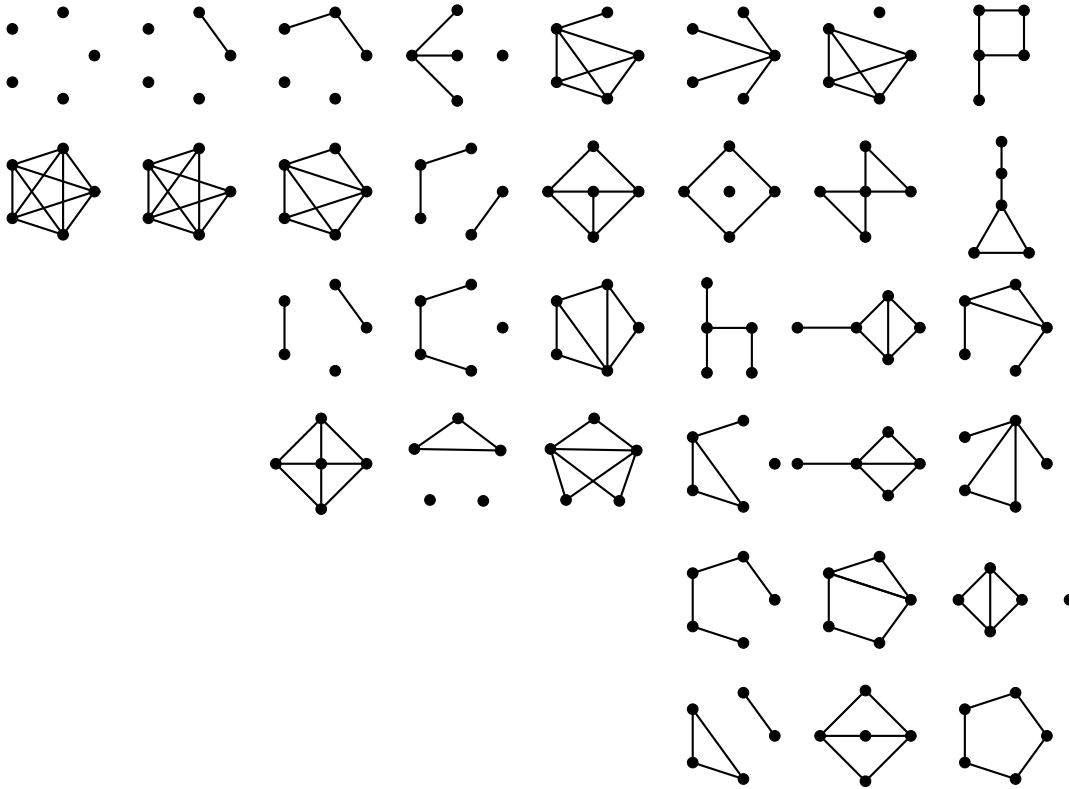
24. The graphs of Figure 31 are not isomorphic because their complements have different degree distributions: the first has one vertex of degree 2, two of degree 1, and four of degree 0, whereas the second has four of degree 1, and three of degree 0.



25. The graphs of Figure 46 are not isomorphic because their complements are nonisomorphic, as shown below. Notably, the cyan edges in one complement are missing in the other.



26.

Figure 5: All 34 graphs having $v = 5$

There are an odd number of nonisomorphic graphs with $v = 4$.

Is this special? Are there other values of v for which the number of nonisomorphic graphs is also odd? (Trudeau, *Graph Theory*

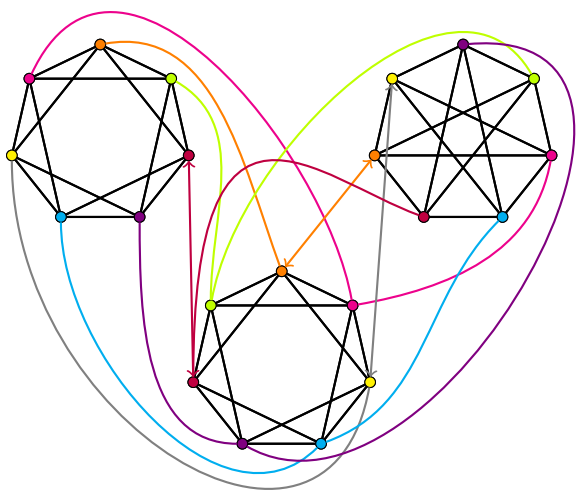
59). <https://math.stackexchange.com/questions/86790/>

[is-there-a-reason-why-the-number-of-non-isomorphic-graphs-with-v-4-is-odd](https://math.stackexchange.com/questions/86790/is-there-a-reason-why-the-number-of-non-isomorphic-graphs-with-v-4-is-odd)

27.

28.

Figure 6: A pair of isomorphic graphs
and their shared canonization



Glossary

nonisomorphic . 9

Define nonisomorphic

path graph If v is a an integer greater than or equal to 2, the *path graph* on v vertices, denoted " P_v ", is the graph having the vertex set $\{1, 2, 3, \dots, v\}$ and edge set $\{\{1, 2\}, \{2, 3\}, \{3, 4\}, \dots, \{v-1, v\}\}$.

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wheel graph Given a vertex set of $\{1, 2, 3, \dots, v\}$, where $v \geq 4$, the edge set of the *wheel graph* W_v is $\{\{1, 2\}, \{1, 3\}, \dots, \{1, v\}, \{2, 3\}, \{3, 4\}, \dots, \{v-1, v\}, \{v, 2\}\}$.

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