

ERIC BAILEY

INTRODUCTION TO GRAPH THEORY: EXERCISES

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Graphs

Exercises

1. $\mathcal{P}(\{1, 2, 3\}) := \{\emptyset, \{1\}, \{2\}, \{3\}, \{1, 2\}, \{1, 3\}, \{2, 3\}, \{1, 2, 3\}\}$

2.

$$\frac{\frac{\emptyset \not\subseteq A}{\exists x \in \emptyset : x \notin A} \quad \overline{\neg \exists x \in \emptyset}}{\frac{\perp}{\emptyset \subseteq A}}$$

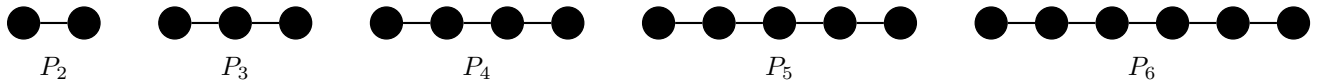
□

3.

$$\frac{\overline{S(b) := \{m \in V : m \notin S(m)\}} \quad \overline{b \in V}}{b \in S(b) \iff b \notin S(b)} \equiv \frac{\overline{R = \{x \notin x\}}}{R \in R \iff R \notin R}$$

4. Let S be the collection of all sets that can be described in an English sentence of twenty-five words or less. S is not a set, because S can be described in fewer than twenty-five words, and if it were a set, then S would have to be a member of itself, which violates the axiomatic definition of a set.

5. The first five *path graphs*.

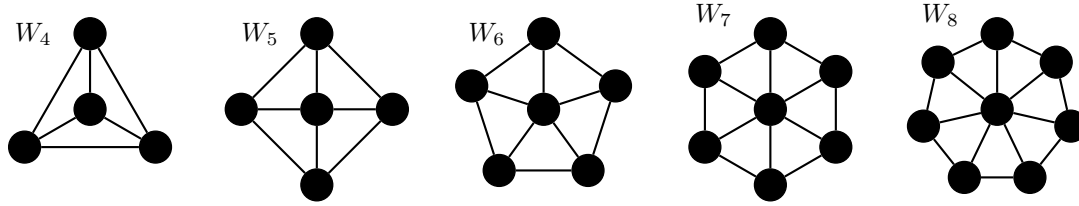


$$\begin{aligned} V(P_v) &= \{1, 2, \dots, v\} \\ E(P_v) &= \{\{n-1, n\} : n \in \{2, \dots, v\}\} \\ |E(P_v)| &= v-1 \\ e &= v-1 \end{aligned}$$

The number of edges in a *path graph* P_v , where $v \geq 2$, is given by the formula $e = v - 1$.

□

6. The first five *wheel graphs*.



$$\begin{aligned}
 V(W_v) &= \{1, 2, 3, \dots, v\} \\
 E(W_v) &= \{ \{1, 2\}, \{1, 3\}, \dots, \{1, v\}, \\
 &\quad \{2, 3\}, \{3, 4\}, \dots, \{v-1, v\}, \\
 &\quad \{v, 2\} \} \\
 &= \{ \{ \{1, n\} : n \in \{2, \dots, v\} \}, \\
 &\quad \{ \{n-1, n\} : n \in \{3, \dots, v\} \}, \\
 &\quad \{v, 2\} \} \\
 |E(W_v)| &= (v-1) + (v-2) + 1 \\
 &= (v-1) + (v-1) \\
 e &= 2(v-1)
 \end{aligned}$$

The number of edges in the *wheel graph* on v vertices W_v , where $v \geq 4$, is given by the formula $e = 2(v-1)$.

□

7.

$$\begin{aligned}
 1 + 2 + \dots + (v-1) &= (1/2)v(v-1) \\
 &= E(K_v) \\
 &= (v-1) + (v-2) + \dots + (v-(v-1)) \\
 &= 1 + 2 + \dots + (v-1)
 \end{aligned} \tag{T2}$$

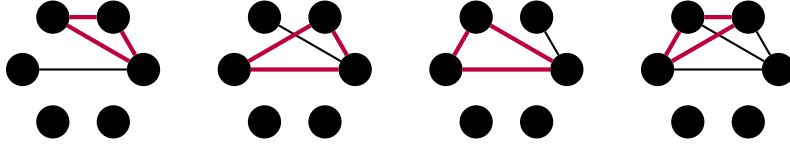
Imagine drawing K_v by joining vertex 1 to vertices 2 through v , creating $v-1$ edges; then joining vertex 2 to vertices 3 through v , creating $v-2$ edges; and so on, i.e. $(v-1) + (v-2) + \dots + (v-(v-1))$, or equivalently, $1 + 2 + \dots + (v-1)$.

□

8. Let G be a graph with v vertices and e edges. In terms of v and e , \overline{G} has $(1/2)v(v-1) - e$ edges.

9. If a graph G has $v = 6$ then G or \overline{G} (possibly both) has a subgraph isomorphic to K_3 .

In the graph G or \overline{G} there exists a vertex a of degree three or more. Let there be three adjacent vertices b, c , and d . If any of the edges $\{b, c\}$, $\{b, d\}$, or $\{c, d\}$ are present in the graph, then the graph contains a subgraph isomorphic to K_3 . If none of those edges is present, then they are all present in the other graph, which thus contains a subgraph isomorphic to K_3 .



10. From the proof in [Exercise 9](#), it follows that in any graph of degree 6 there exists a subgraph H , such that either H or \overline{H} is isomorphic to K_3 . Equivalently, in any gathering of six people (a graph G with $v = 6$) there are either three people who are mutually acquainted (a subgraph isomorphic to K_3) or three people who are mutually unacquainted (a subgraph whose complement is isomorphic to K_3).

11.

$$2|\emptyset| = \sum_{v \in V} \deg(v) = 0 \quad (\text{BC})$$

$$G(V, E) \rightarrow 2|E| = \sum_{v \in V} \deg(v) \quad (\text{IH})$$

$$H(V', E'), |E'| = |E| + 1 \quad (1)$$

$$\exists H_0(V', E'_0), |E'_0| = |E| \quad (2)$$

$$2|E'_0| = \sum_{v \in V'} \deg(v) \quad (\text{by IH})$$

$$H \cong (V', E'_0 \cup \{e\}), e \notin E'_0 \quad (3)$$

$$\sum_{v \in V'} \deg(v) + 2 = 2|E| \quad (4)$$

$$2|E| + 2 = 2(|E| + 1) = \sum_{v \in V'} \deg(v) \quad (5)$$

□

12. a)

$$(4 \times 3) + (2 \times 5) + (2 \times 6) + (1 \times 8) = 2e$$

$$12 + 10 + 12 + 8 = 2e$$

$$21 = e$$

b) $(7 \times 3) = 2e \rightarrow e \notin \mathbb{Q}$

13. If $C_n \cong \overline{C_n}$, then it must be the case that

$$|E(C_n)| + |E(\overline{C_n})| = |E(K_n)|, \text{ i.e. } n + n = \frac{n(n-1)}{2}.$$

Since $n = 5$ is a solution, $|E(C_5)| + |E(\overline{C_5})| = |E(K_5)|$, and thus $C_5 \cong \overline{C_5}$.

□

A graph with no edges ($n = 0$) has no cycles. Therefore the only self-complementary cycle graph is C_5 .

□

14. If $G \cong \overline{G}$ for a graph G of order v , then $|E(G)| + |E(\overline{G})| = |E(K_v)|$, i.e. $e + e = 2e = \frac{v(v-1)}{2}$. So the number of edges e in G is $\frac{v(v-1)}{4}$, and $4 \mid v(v-1)$. Therefore, for any self-complementary graph G of order v , $4 \mid v$ or $4 \mid v-1$.

□

15. .

Find two other self-complementary graphs
(Trudeau, *Graph Theory* 57)

$$n + n = \frac{n(n-1)}{2}$$

$$2n = \frac{n(n-1)}{2}$$

$$4n = n(n-1)$$

$$4n = n^2 - n$$

$$0 = n^2 - 5n$$

$$n = \{0, 5\}$$

Glossary

path graph If v is a an integer greater than or equal to 2, the *path graph* on v vertices, denoted “ P_v ”, is the graph having the vertex set $\{1, 2, 3, \dots, v\}$ and edge set $\{\{1, 2\}, \{2, 3\}, \{3, 4\}, \dots, \{v-1, v\}\}$.
5

wheel graph . 6

define