Knowledge Discovery in Databases

Data Warehousing and OLAP Technology



School of Software, Nanjing University

Data Warehousing and OLAP Technology

- **■** What is a data warehouse?
- A multi-dimensional data model
- Data warehouse architecture
- Data warehouse implementation
- Further development of data cube technology
- From data warehousing to data mining

What is Data Warehouse?

- Defined in many different ways, but not rigorously.
 - A decision support database that is maintained separately from the organization's operational database
 - Support information processing by providing a solid platform of consolidated, historical data for analysis.
- "A data warehouse is a <u>subject-oriented</u>, <u>integrated</u>, <u>time-variant</u>, and <u>nonvolatile</u> collection of data in support of management's decision-making process."—W. H. Inmon
- Data warehousing:
 - The process of constructing and using data warehouses

Data Warehouse—Subject-Oriented

- Organized around major subjects, such as customer, product, sales.
- Focusing on the modeling and analysis of data for decision makers, not on daily operations or transaction processing.
- Provide a simple and concise view around particular subject issues by excluding data that are not useful in the decision support process.

Data Warehouse—Integrated

- Constructed by integrating multiple, heterogeneous data sources
 - Relational databases, flat files, on-line transaction records
- Data cleaning and data integration techniques are applied.
 - Ensure consistency in naming conventions, encoding structures, attribute measures, etc. among different data sources
 - ✓ E.g., Hotel price: currency, tax, breakfast covered, etc.
 - When data is moved to the warehouse, it is converted.

Data Warehouse—Time Variant

- The time horizon for the data warehouse is significantly longer than that of operational systems.
 - Operational database: current value data.
 - Data warehouse data: provide information from a historical perspective (e.g., past 5-10 years)
- Every key structure in the data warehouse
 - Contains an element of time, explicitly or implicitly
 - But the key of operational data may or may not contain "time element".

Data Warehouse—Non-Volatile

- A physically separate store of data transformed from the operational environment.
- Operational update of data does not occur in the data warehouse environment.
 - Does not require transaction processing, recovery, and concurrency control mechanisms
 - Requires only two operations in data accessing:
 - ✓ initial loading of data and access of data.

Data Warehouse vs. Heterogeneous DBMS

■ Traditional heterogeneous DB integration:

- Build wrappers/mediators on top of heterogeneous databases
- Query driven approach
 - ✓ When a query is posed to a client site, a meta-dictionary is used to translate the query into queries appropriate for individual heterogeneous sites involved, and the results are integrated into a global answer set
 - ✓ Complex information filtering, compete for resources

Data warehouse: update-driven, high performance

 Information from heterogeneous sources is integrated in advance and stored in warehouses for direct query and analysis

Data Warehouse vs. Operational DBMS

■ OLTP (On-Line Transaction Processing)

- Major task of traditional relational DBMS
- Day-to-day operations: purchasing, inventory, banking, manufacturing, payroll, registration, accounting, etc.

OLAP (On-Line Analytical Processing)

- Major task of data warehouse system
- Data analysis and decision making

■ Distinct features (OLTP vs. OLAP):

- User and system orientation: customer vs. market
- Data contents: current, detailed vs. historical, consolidated
- Database design: ER + application vs. star + subject
- View: current, local vs. evolutionary, integrated
- Access patterns: update vs. read-only but complex queries

OLTP vs. OLAP

| | OLTP | OLAP |
|--------------------|---------------------------|------------------------------|
| users | clerk, IT professional | knowledge worker |
| function | day to day operations | decision support |
| DB design | application-oriented | subject-oriented |
| data | current, up-to-date | historical, |
| | detailed, flat relational | summarized, multidimensional |
| | isolated | integrated, consolidated |
| usage | repetitive | ad-hoc |
| access | read/write | lots of scans |
| | index/hash on prim. key | |
| unit of work | short, simple transaction | complex query |
| # records accessed | tens | millions |
| #users | thousands | hundreds |
| DB size | 100MB-GB | 100GB-TB |
| metric | transaction throughput | query throughput, response |

Why Separate Data Warehouse?

■ High performance for both systems

- DBMS— tuned for OLTP: access methods, indexing, concurrency control, recovery
- Warehouse—tuned for OLAP: complex OLAP queries, multidimensional view, consolidation.

Different functions and different data:

- missing data: Decision support requires historical data which operational DBs do not typically maintain
- data consolidation: DS requires consolidation (aggregation, summarization) of data from heterogeneous sources
- data quality: different sources typically use inconsistent data representations, codes and formats which have to be reconciled

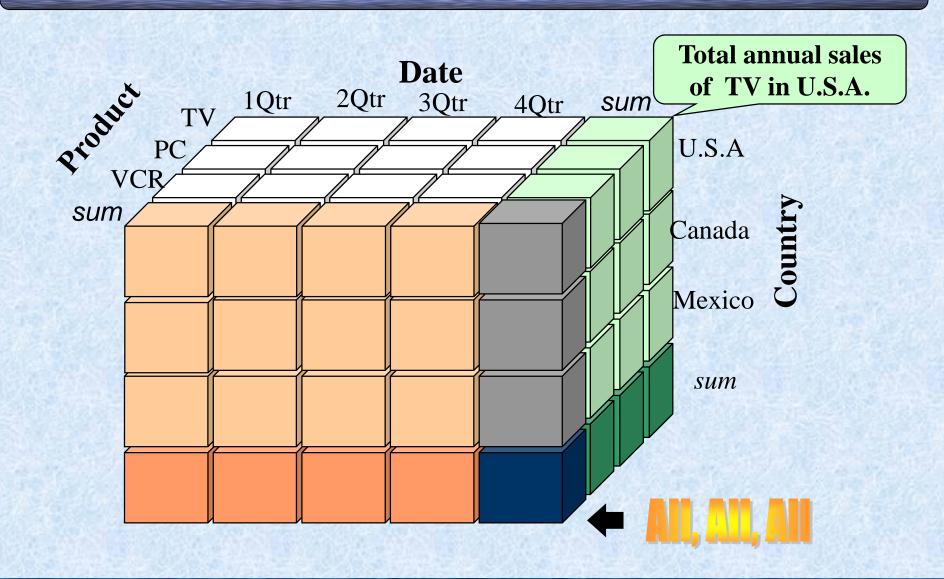
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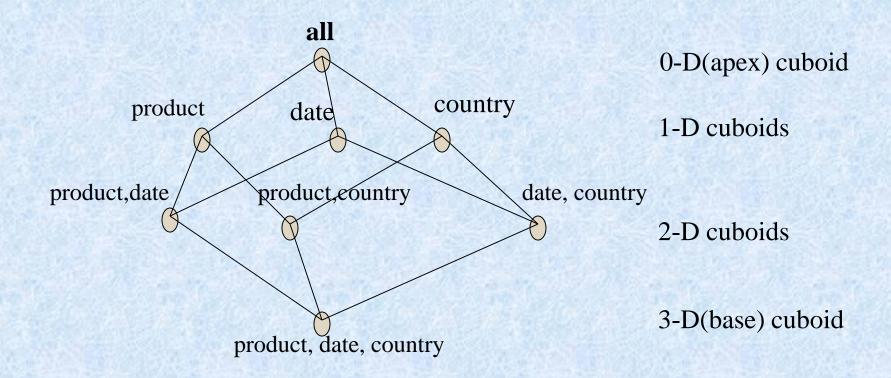
From Tables and Spreadsheets to Data Cubes

- A data warehouse is based on a multidimensional data model which views data in the form of a data cube
- A data cube, such as sales, allows data to be modeled and viewed in multiple dimensions
 - Dimension tables, such as item (item_name, brand, type), or time(day, week, month, quarter, year)
 - Fact table contains measures (such as dollars_sold) and keys to each
 of the related dimension tables

Multi-Dimensional Data Model(Cube)

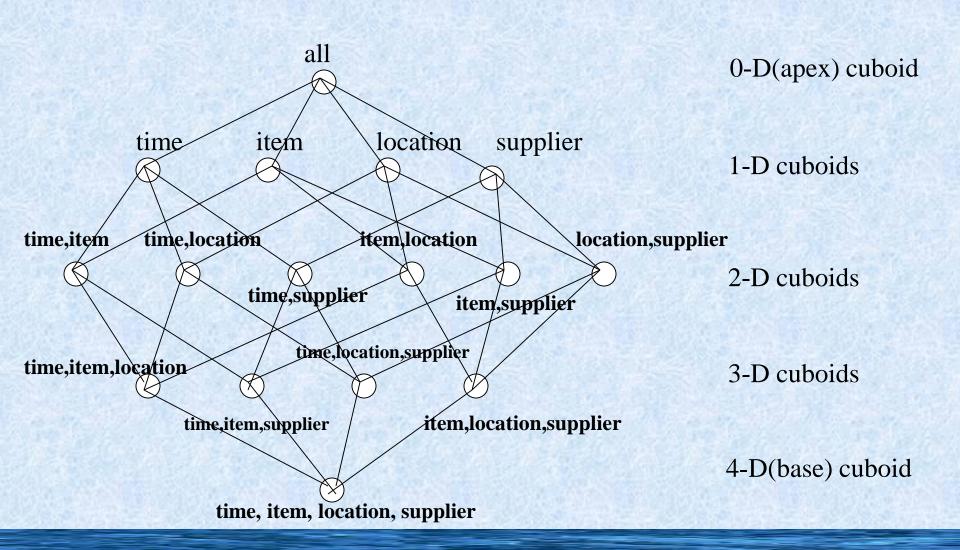


Cube: A Lattice of Cuboids(I)



In data warehousing literature, an n-D base cube is called a base cuboid. The top most 0-D cuboid, which holds the highest-level of summarization, is called the apex cuboid. The lattice of cuboids forms a data cube.

Cube: A Lattice of Cuboids(II)

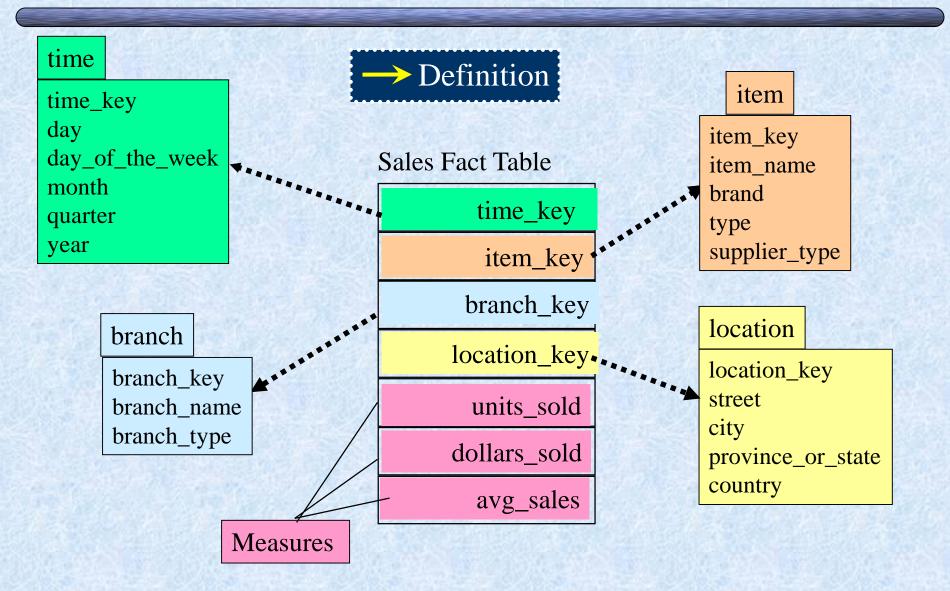


Conceptual Modeling of Data Warehouses

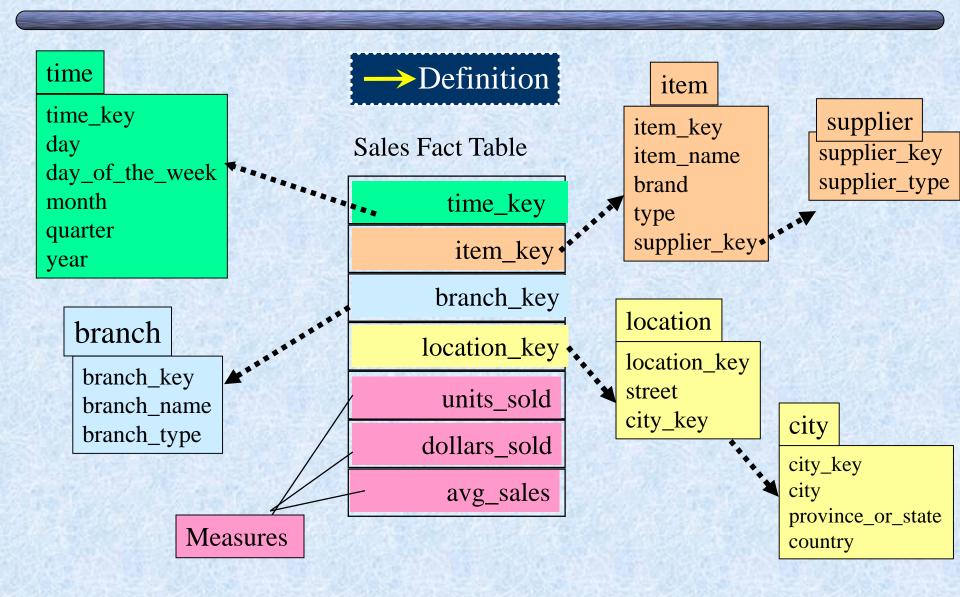
■ Modeling data warehouses: dimensions & measures

- Star schema: A fact table in the middle connected to a set of dimension tables
- Snowflake schema: A refinement of star schema where some dimensional hierarchy is normalized into a set of smaller dimension tables, forming a shape similar to snowflake
- <u>Fact constellations</u>: Multiple fact tables share dimension tables,
 viewed as a collection of stars, therefore called galaxy schema or fact
 constellation

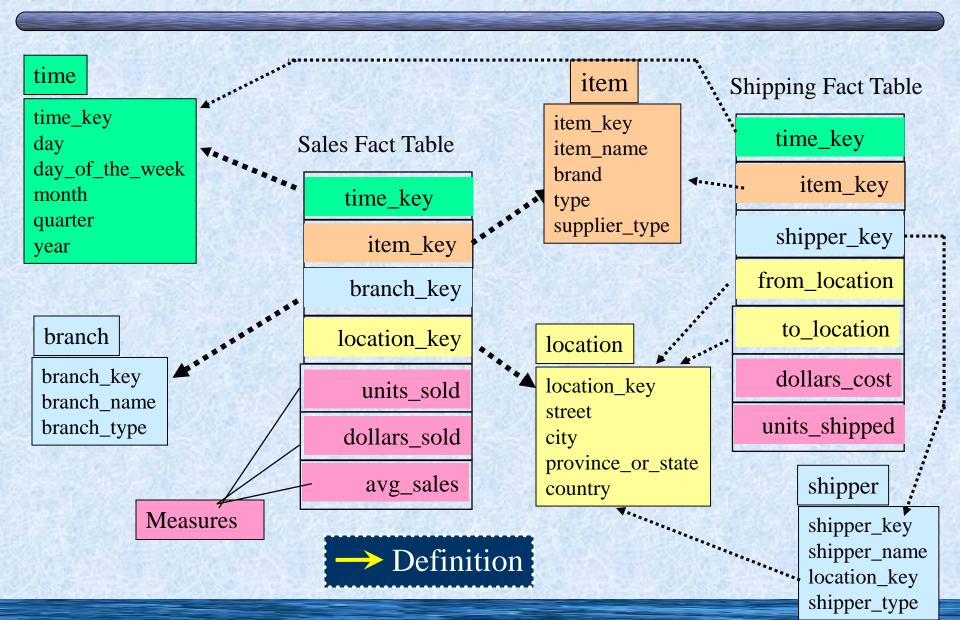
Example of Star Schema



Example of Snowflake Schema



Example of Fact Constellation



A Data Mining Query Language, DMQL

Cube Definition (Fact Table)

```
define cube <cube_name> [<dimension_list>]:
    <measure_list>
```

■ Dimension Definition (Dimension Table)

```
define dimension <dimension_name> as
  (<attribute_or_subdimension_list>)
```

- **■** Special Case (Shared Dimension Tables)
 - First time as "cube definition"
 - define dimension <dimension_name> as
 <dimension_name_first_time> in cube <cube_name_first_time>

Defining a Star Schema in DMQL

define cube sales_star [time, item, branch, location]: Figure

```
dollars_sold = sum(sales_in_dollars), avg_sales =
  avg(sales_in_dollars), units_sold = count(*)
```

- define dimension item as (item_key, item_name, brand, type, supplier_type)
- define dimension branch as (branch_key,
 branch_name, branch_type)
- define dimension location as (location_key, street, city, province_or_state, country)

Defining a Snowflake Schema in DMQL

define cube sales_snowflake [time, item, branch, location]:

Figure

```
dollars_sold = sum(sales_in_dollars), avg_sales =
  avg(sales_in_dollars), units_sold = count(*)
```

- define dimension time as (time_key, day, day_of_week, month, quarter, year)
- define dimension item as (item_key, item_name,
 brand, type, supplier(supplier_key, supplier_type))
- define dimension branch as (branch_key,
 branch_name, branch_type)
- define dimension location as (location_key, street,
 city(city_key, province_or_state, country))

Defining a Fact Constellation in DMQL

define cube sales [time, item, branch, location]:

dollars_sold = sum(sales_in_dollars), avg_sales = avg(sales_in_dollars),
 units_sold = count(*)

- define dimension time as (time_key, day, day_of_week, month, quarter, year)
- define dimension item as (item_key, item_name, brand,
 type, supplier_type)
- define dimension location as (location_key, street, city,
 province_or_state, country)
- define cube shipping [time, item, shipper, from_location,
 to_location]:

dollar_cost = sum(cost_in_dollars), unit_shipped = count(*)



define dimension time as time in cube sales define dimension item as item in cube sales

define dimension shipper as (shipper_key, shipper_name, location as location in cube sales, shipper_type)

define dimension from_location as location in cube sales define dimension to_location as location in cube sales

Measures: Three Categories

■ <u>distributive</u>: if the result derived by applying the function to *n* aggregate values is the same as that derived by applying the function on all the data without partitioning.

✓ E.g., count(), sum(), min(), max().

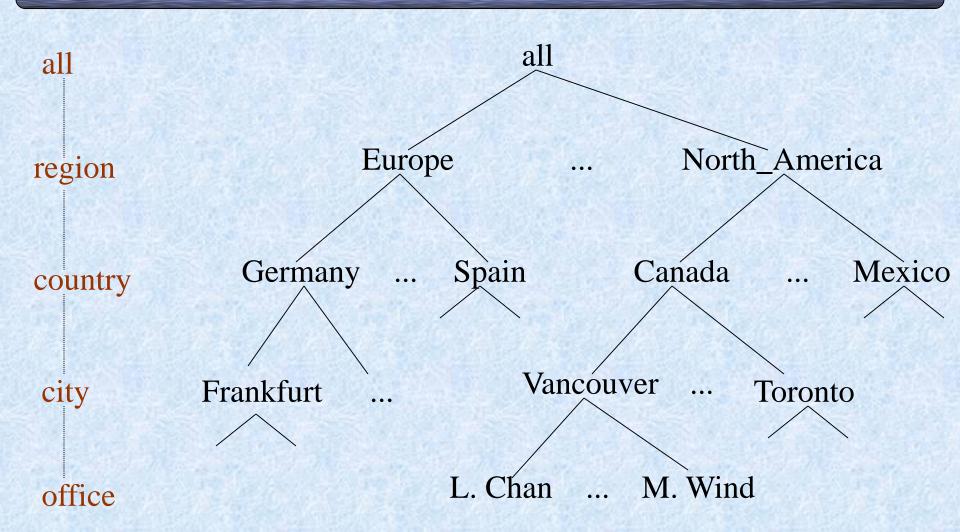
■ <u>algebraic</u>: if it can be computed by an algebraic function with *M* arguments (where *M* is a bounded integer), each of which is obtained by applying a distributive aggregate function.

✓ E.g., avg(), standard_deviation().

■ <u>holistic</u>: if there is no constant bound on the storage size needed to describe a subaggregate.

✓ E.g., median(), mode(), rank().

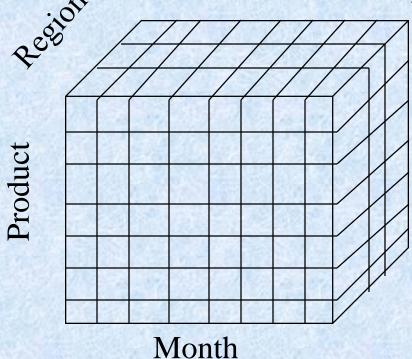
A Concept Hierarchy: Dimension (location)

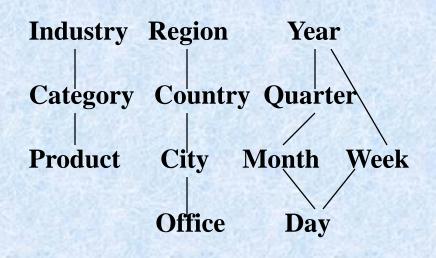


Multidimensional Data

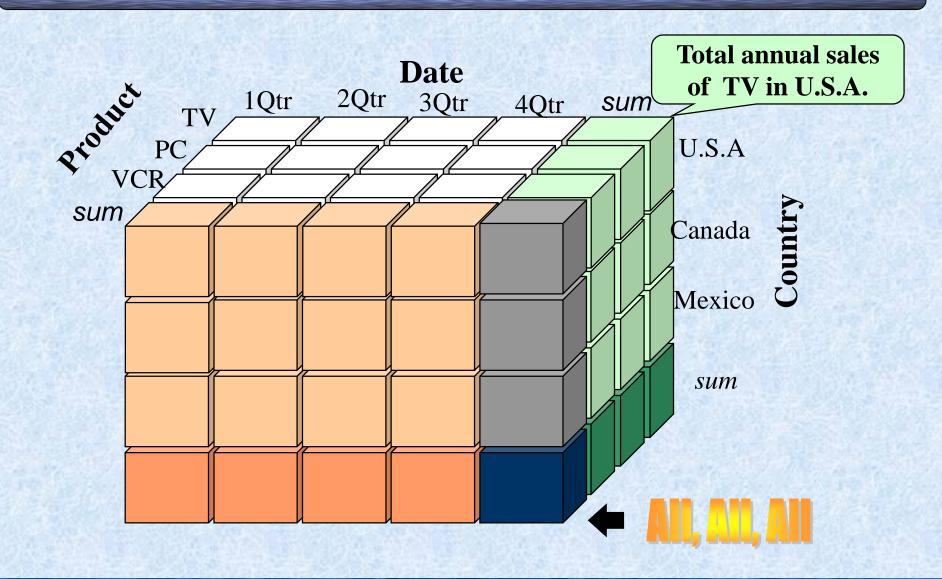
Sales volume as a function of product, month, and region

Dimensions: Product, Location, Time Hierarchical summarization paths

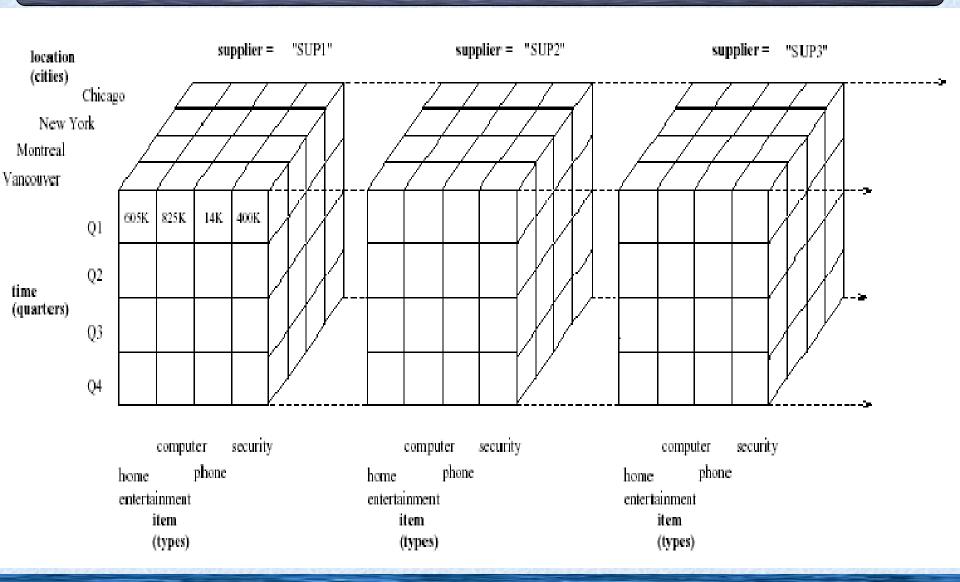




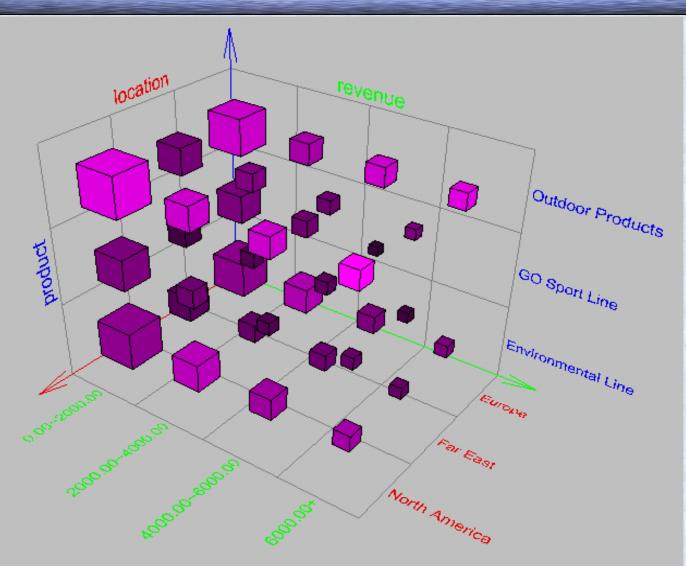
A Sample Data Cube(I)



A Sample Data Cube(II)



Browsing a Data Cube

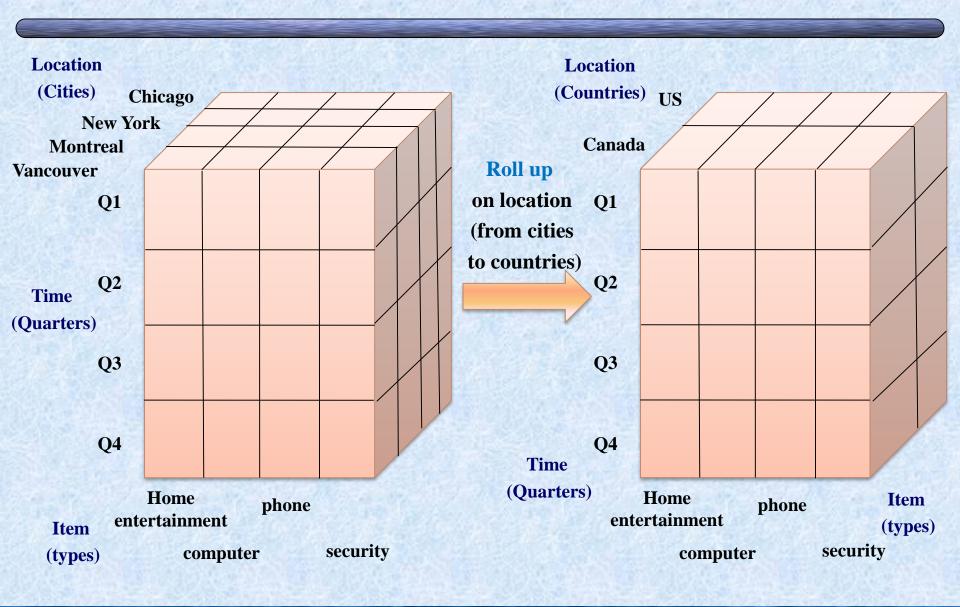


- Visualization
- OLAP capabilities
- Interactive manipulation

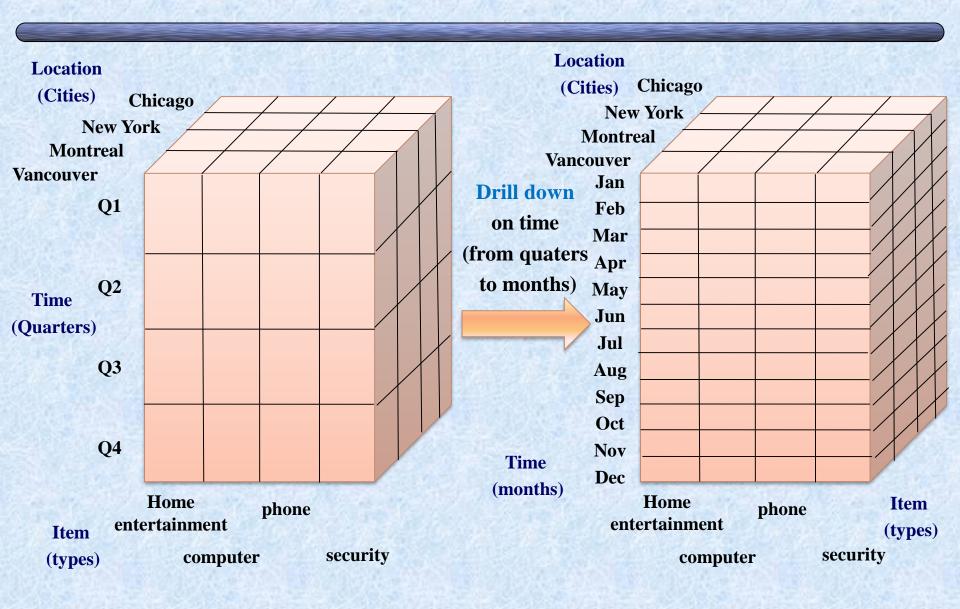
Typical OLAP Operations

- Roll up (drill-up): summarize data
 - by climbing up hierarchy or by dimension reduction
- Drill down (roll down): reverse of roll-up
 - from higher level summary to lower level summary or detailed data, or introducing new dimensions
- **■** Slice and dice:
 - project and select
- Pivot (rotate):
 - reorient the cube, visualization, 3D to series of 2D planes.
- Other operations
 - drill across: involving (across) more than one fact table
 - drill through: through the bottom level of the cube to its back-end relational tables (using SQL)

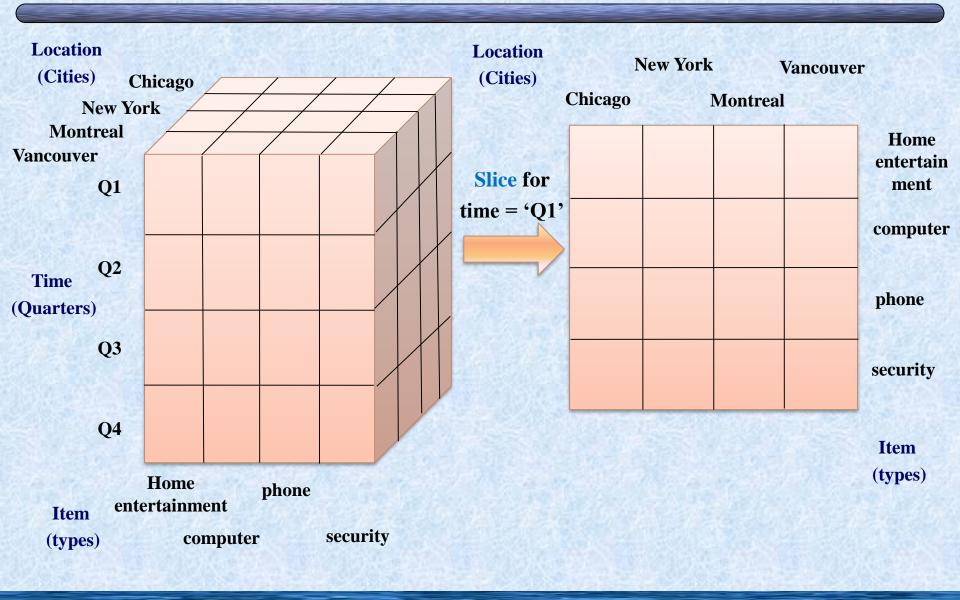
Typical OLAP Operations—— Roll-up



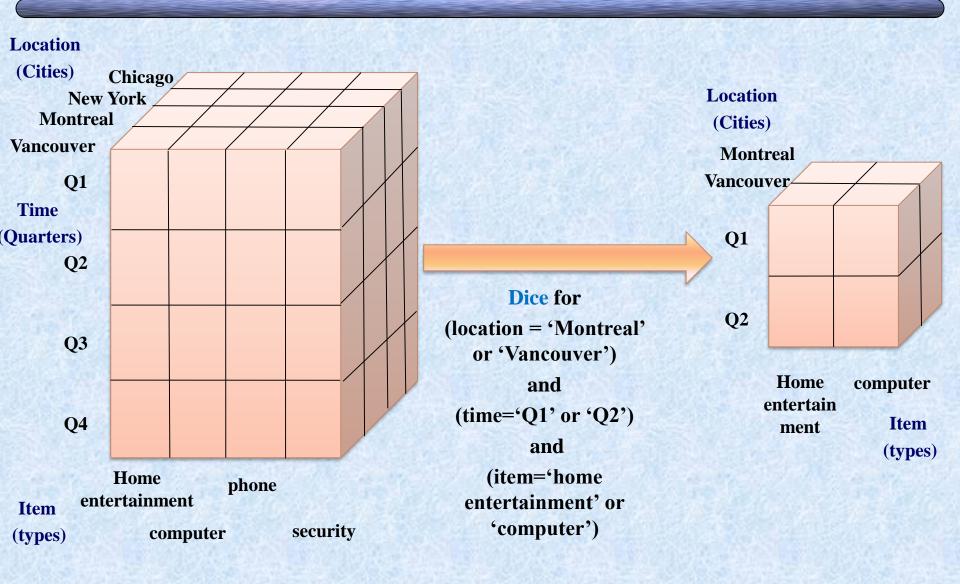
Typical OLAP Operations—— Drill-down



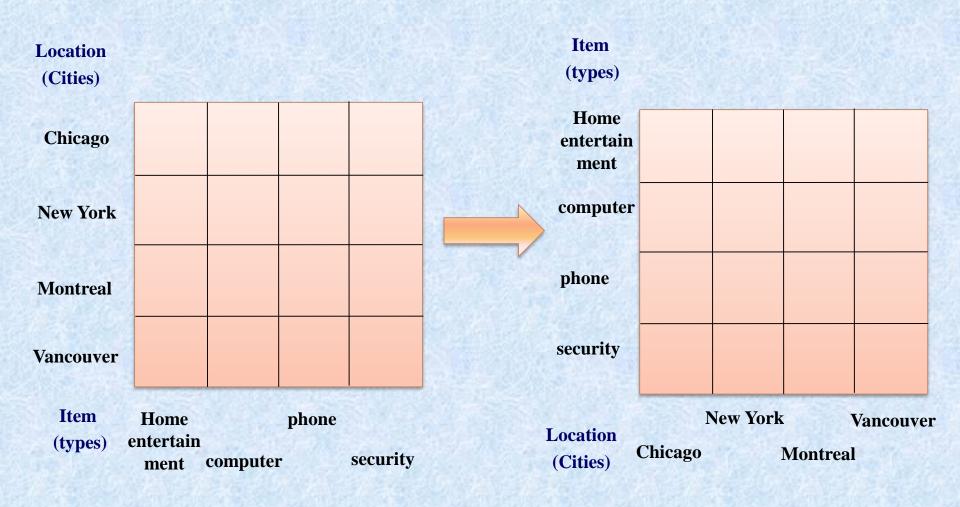
Typical OLAP Operations—— Slice



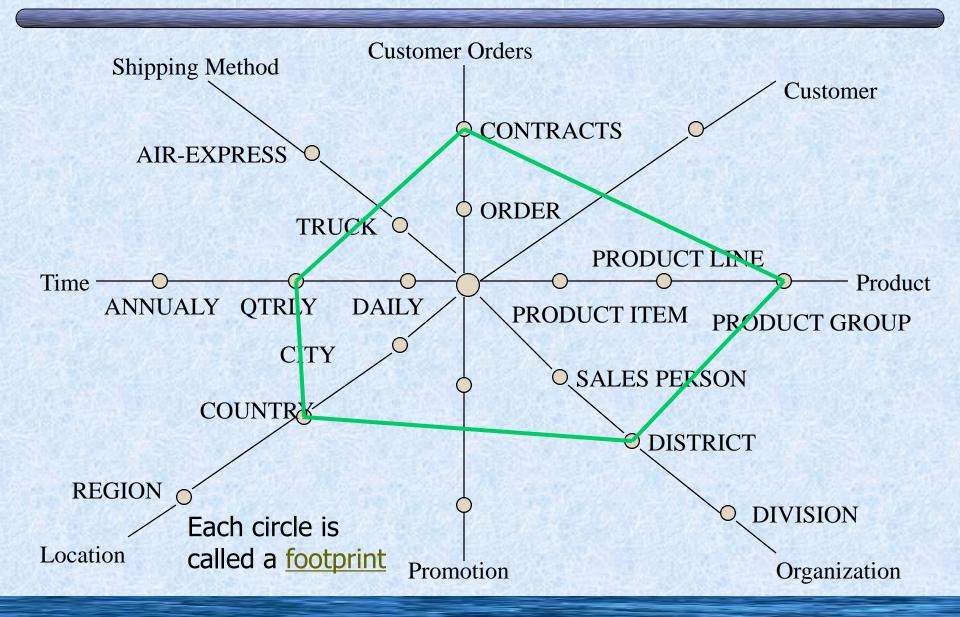
Typical OLAP Operations — Dice



Typical OLAP Operations — Pivot(Rotate)



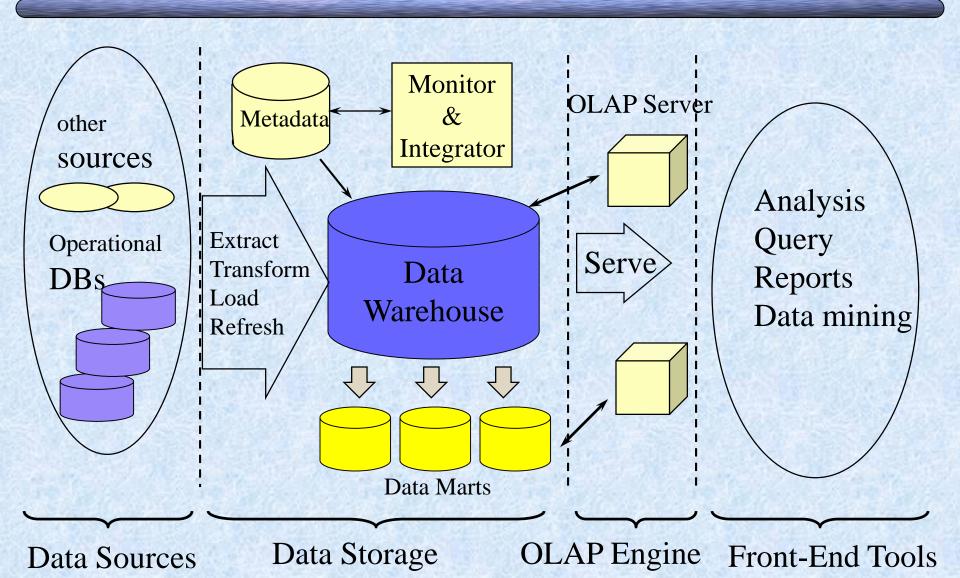
A Star-Net Query Model



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Multi-Tiered Architecture



Three Data Warehouse Models

■ Enterprise warehouse

 collects all of the information about subjects spanning the entire organization

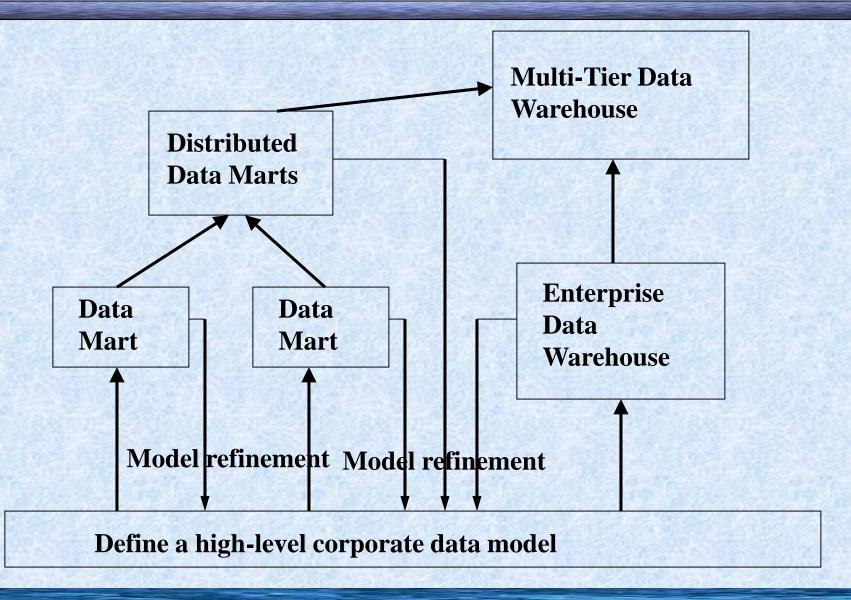
Data Mart

- a subset of corporate-wide data that is of value to a specific groups of users. Its scope is confined to specific, selected groups, such as marketing data mart
 - ✓ Independent vs. dependent (directly from warehouse) data mart

Virtual warehouse

- A set of views over operational databases
- Only some of the possible summary views may be materialized

Data Warehouse Development: A Recommended Approach



OLAP Server Architectures

Relational OLAP (ROLAP)

- Use relational or extended-relational DBMS to store and manage warehouse data and OLAP middleware to support missing pieces
- Include optimization of DBMS backend, implementation of aggregation navigation logic, and additional tools and services
- greater scalability

Multidimensional OLAP (MOLAP)

- Array-based multidimensional storage engine (sparse matrix techniques)
- fast indexing to pre-computed summarized data

Hybrid OLAP (HOLAP)

User flexibility, e.g., low level: relational, high-level: array

Specialized SQL servers

specialized support for SQL queries over star/snowflake schemas

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数据仓库设计

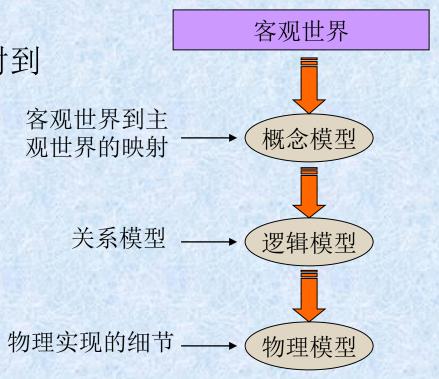
- ≥ 数据仓库设计和数据库设计的差异
 - ◆ 面向需求不同:数据库面向具体应用,需求一开始就很明确,而数据仓库是一个渐进的过程
 - ◆ 设计目标不同: OLTP vs. OLAP
 - ◆ 处理类型不同: 面向操作型应用 vs. 面向分析型应用
 - ◆ 数据来源不同: 业务员输入 vs. 已存在的业务系统数据
 - ◆ 系统设计的方法不同:
 - ✓数据仓库可以采用"数据驱动"的设计方法
 - ✓数据仓库设计可以分为数据仓库模型设计和数据装 载接口设计两部分

数据仓库的设计步骤

将企业模型映射到数据仓库系统的过程

≥逻辑模型设计

≥ 物理模型设计



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OLAP建模方法

- ≥ 维表设计
 - ◆ 维的变化
 - ◆ 维表的共享
 - ◆ 层次信息和分类信息的位置

- ≥ 事实表设计
 - ◆ 事实表的特性
 - ◆ 通用数据和专用数据事实表

维表的变化(一)

- 维表通过记录因素的属性描述事件中包含的诸多因素
- 维表的本质是多维分析空间在某个角度上的投影
- 由于维表描述的是事物的属性,因此随着事物本身的变化,其属性也会产生改变
 - ◆ 如果该属性与决策没有太大关系,例如电话号码属性 对于分析顾客购买行为没有什么作用,则此属性的变 化可以忽略不计
 - ◆ 如果该属性与决策有关,例如某位顾客搬家后离超市 更远了,我们试图分析其购买行为与家里距离变远有 何关系,则不能将之忽略

维表的变化(二)

| | 对于需要记录其改变的维, | 有若干方法可以进行处理 |
|---|--------------|-------------|
| _ | | |

◆ 当属性进行变化时, 创建一个新记录

✓例如:✓缺点:345 顾客A 东城区 → 368 顾客A 西城区

- 由于ID产生变化,被认为是两条记录
- ◆ 创建一个新的字段,将新地址填入

✓例如:✓缺点:✓缺点:
345 顾客A 东城区 西城区

- ■可扩展性不佳
- ◆ 增加一个修订号码字段和当前标记字段

✓例如: 345 0 过往 东城区 345 1 当前 西城区

✔缺点:

■ 维表和事实表连接时需要采用"主关键字"+"修订号码",增加了事实表的复杂性

维表的变化(三)

- ■最为理想的解决方案
 - ◆ 新建一个关键字客户ID,通过主关键字与之相连,使 用时间字段标志当前的值
 - ✔缺点: 相对较为复杂

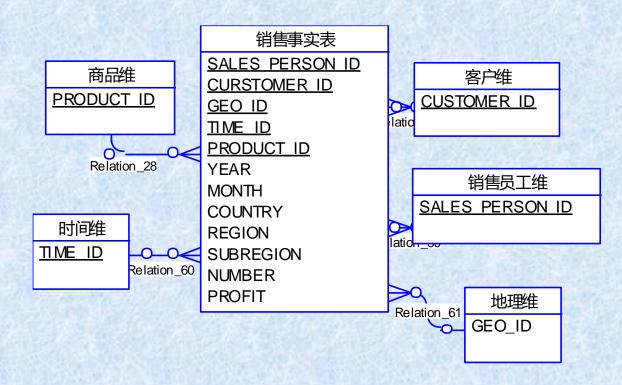
| 345 | 2003 / 1 | 当前 | 西城区 |
|-----|----------|----|-----|
| 345 | 1999 / 3 | 过往 | 东城区 |
| 345 | 1998 / 7 | 过往 | 海淀区 |

维表的共享

- ≥ 多个维表中可能包含相同的属性:
 - ◆ 供货商维中包含地址维,而销售商维中可能也包含地址维,因而可能共享维表
 - ◆ 由于数据仓库中时间维的重要性,各个维中都有可能 包含时间维,因而可能共享时间维表
- ■可以采用扩展的星座结构来描述共享维表
- 具体内容请参见上节

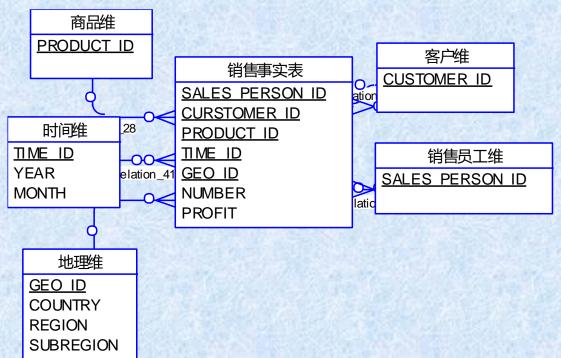
层次信息的位置(一)

- ≥ 将维层次信息放入事实表
 - ◆ 优点: 计算极为方便
 - ◆ 缺点: 事实表会因此变得极为庞大



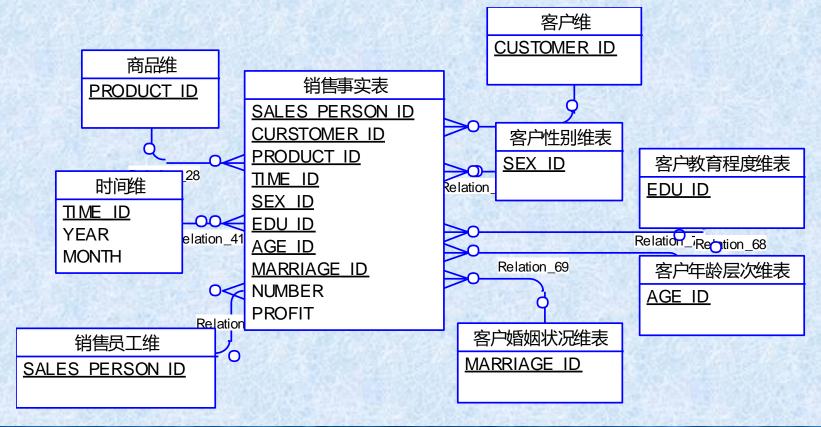
层次信息的位置(二)

- ▶ 将维层次放在各自的维表中,通过主关键字与事实表相连
 - ◆ 优点: 减少了事实表的大小
 - ◆ 缺点: OLAP性能比上一种方式差



分类信息的位置(一)

- ▲ 在事实表中体现维分类
 - ◆ 优点: OLAP性能好
 - ◆ 缺点: 事实表的字段数和记录数增加



分类信息的位置(二)

≥ 在维表中体现维分类 客户性别维表 Relation 88 SEX ID 商品维 PRODUCT ID 客户教育程度维表 客户维 销售事实表 EDU ID **CUSTOMER ID** SALES PERSON ID AGE ID **CURSTOMER ID** 时间维 28 EDU ID PRODUCT ID TIME ID Relation Relation 89 MARRIAGE ID TIME ID YEAR elation 41 客户年龄层次维表 SEX ID NUMBER **MONTH** AGE ID **PROFIT** Relation 90 Relation_30 销售员工维 客户婚姻状况维表 SALES PERSON ID MARRIAGE ID Relation 91

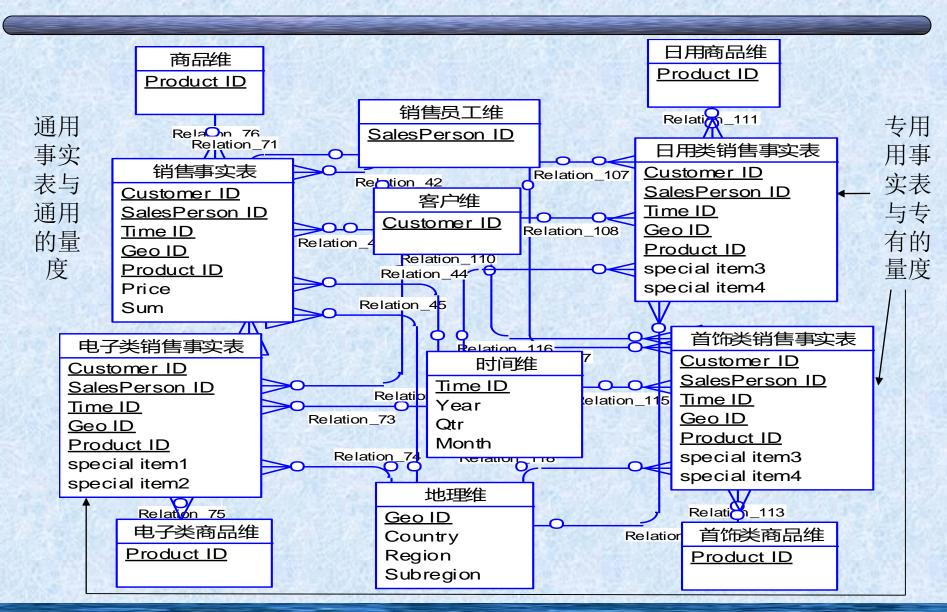
事实表的特征

- ≥ 与维表相比,事实表具有以下特征:
 - ◆ 记录数量非常多
 - ◆ 除了度量外,其他字段都是维表或者中间维表(雪花模型)的关键字
 - ◆ 如果事实相关的维度很多,则事实表的字段数也会比较多
- 因此应当尽量减小一条记录的长度,才能避免事实表过 大而难于管理
- 数据的粒度是影响事实表大小的关键因素,因而必须认 真设计

通用数据和专用数据事实表

- ☑ 对应一个问题通常采用一个事实表,但在特殊情况下,也允许采用多个事实表
- 例如:超市里出售多种商品,由于商品本身分类不同,因此 所采用的量度可能也不相同
 - ◆ 如果将这些量度全部置于一个事实表中,由于某种类型的 商品的量度其他商品可能不具备,因此则不可避免将在事 实表中造成大量的数据空缺
 - ◆解决办法:采用多个事实表,分为通用数据事实表与专用数据事实表加以管理

通用数据和专用数据事实表



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数据仓库的逻辑模型设计

- ■系统数据量估算
- ≥数据粒度的选择
- ≥数据的分割
- ≥表的合理划分
- ≥去除纯操作数据
- ≥增加导出字段
- ≥定义关系模式
- ≥定义元数据存储
- ≥定义记录系统

数据仓库的逻辑模型设计

■ 系统数据量估算

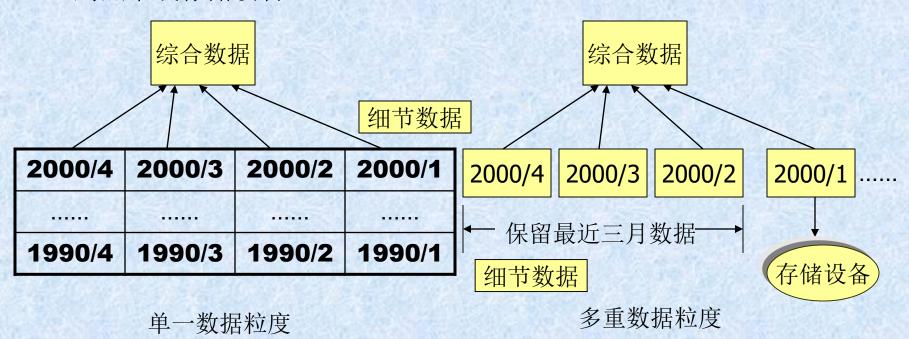
◆ 设在概念模型中出现的表的个数为N(这些表中应不包含不会放进数据仓库的表),对于每个表i(0≤i≤N), 计算表的大小S_i和表的主关键字大小K_i,然后估计每张表i 在单位时间内最大记录数 L_{max}和最少记录数L_{min},则数据仓库的粗略数据量在以下范围:

$$\alpha \times \left(\sum_{i=1}^{N} (S_i + K_i) \times L_{\max}^i \times T\right) \sim \alpha \times \left(\sum_{i=1}^{N} (S_i + K_i) \times L_{\min}^i \times T\right)$$

- ◆ 其中,T是数据在数据仓库中存在的周期。α是考虑由于数据冗余和数据索引而使数据量增大的冗余因子,通常取1.2~2
- ◆ 本公式的含义是:
 - ✓数据仓库数据量={累加[(每张表记录大小+每张表主关键字 大小)*每张表单位时间内记录的数量]*存储时间}*冗余因子
- ◆ 公式估算出的结果仅能作为参考

数据粒度的选择

- 数据量较小的情况下使用单一的数据粒度,即直接存储细节数据并 定期在细节数据基础上进行数据综合
- 对于大数据量需要采用双重粒度,对于细节数据只保存近期的数据 在数据仓库之中,当保留周期到达时,将距离当前较远的数据导出 到磁带或存储设备上



数据粒度策略

| 一年内数据量(行) | 五年内数据量(行) | 位(行) 数据粒度策略 | | |
|------------|------------|----------------------|--|--|
| 10,000 | 100,000 | 简单的单一粒度策略 | | |
| 100,000 | 1,000,000 | 如果选择单一粒度, 则需要认真设计 | | |
| 1,000,000 | 10,000,000 | 最好使用双重粒度 | | |
| 10,000,000 | 20,000,000 | 必须采用双重粒度且 认真设计 | | |

数据的分割

- ▶ 为何要进行数据分割?
 - ◆ 数据仓库中数据量过大时,检索速度很慢
- 数据分割是指将数据分散到各自的物理单元里以便能够独立处理,以提高数据处理的效率
- ≥ 数据分割没有固定的标准,分割的方法和粒度应当根据实际情况确定
 - 通常选择时间、地点、业务等划分
 - ◆ 一般按照时间分割数据分布比较均匀,因此按照时间分割最为常见



合理的表划分

- 直接存储字段数目很大的表,会造成以下问题:
 - ◆ 各个字段的更新频率不一,放在一张表里造成数据追加工作的浪费
 - ◆ 各个字段的访问频率不一,放在一张表里影响访问效率

- 因此需要对表中的内容进行合理的划分
 - ◆ 按数据的稳定性划分
 - ◆ 按业务规则进行表划分(略)

按数据稳定性进行表划分

原始表

几乎不变的数据

有时改变的数据

经常改变的数据

商品ID 数商品加载时间 高据和称 高品类型格 高品质量,以为 最后,则为 最后,则为 是后,以为 是后,以为 是一,成本 商品ID 商品名称 商品类型 商品规格 商品质量

商品**ID** 主要原料供应商 生产成本 商品ID 数据加载时间 最后订购时间 最后订购数量

避免了整张表记录快速增长的现象, 节约了存储空间

删除纯操作数据及增加导出字段

- - ◆例: "收款人"字段

- 导出数据本身是冗余的,但是增加导出字段有利 于数据以后的使用
 - ◆例:在按月综合表中,可以加入"平均价格", "供货总价","供货总数量"等导出字段

Meta Data

- What is meta data?
 Meta data are data describing data
- In data warehouse, there are several data levels:
 - meta data
 - current detailed data
 - older detailed data
 - lightly summarized data
 - highly summarized data

Meta Data Storage

- Meta data is the data defining warehouse objects. It has the following kinds
 - Description of the structure of the warehouse
 - ✓ schema, view, dimensions, hierarchies, derived data definition, data mart locations and contents
 - Operational meta-data
 - ✓ data lineage (history of migrated data and transformation path), currency of data (active, archived, or purged), monitoring information (warehouse usage statistics, error reports, audit trails)
 - The algorithms used for summarization
 - The mapping from operational environment to the data warehouse
 - Business data
 - ✓ business terms and definitions, ownership of data, charging policies

记录系统的定义

- ≥ 记录系统是操作型元数据的一部分
- 记录系统指明数据仓库中关系表的各个字段来源于业务数据库何处

◆ 例:

| 主题 | 表名 | 属性字段 | 源数据库 | 源表名 | 源字段 |
|----|-----|-------|------|-----|------|
| 商品 | 供应表 | 交易ID | 采购 | 采购表 | 交易号 |
| 商品 | 供应表 | 商品ID | 采购 | 采购表 | 商品号 |
| 商品 | 供应表 | 供应商ID | 采购 | 采购表 | 供应商号 |
| 商品 | 供应表 | 供货数量 | 采购 | 采购表 | 供货数量 |
| 商品 | 供应表 | 供货价格 | 采购 | 采购表 | 供货价 |
| | | | | | |

数据仓库物理模型设计

- 确定数据的存储结构
- ≥ 索引策略
- ≥ 数据存储策略与性能优化
 - ◆ 多路聚集优化
 - ◆ 表的归并
 - ◆ 分割表的存放
 - ◆ 按列存储
 - ◆ 存储分配优化
- ≥ 数据装载接口设计
- ≥ 并行优化设计

Define the Storage Structure

- 一般的数据库数据量相对较小,除非业务要求必须保证数据的安全性和可恢复性,否则可以不采用并行存储结构
- ≥ 数据仓库由于数据的巨量存储,必须采用并行存储结构,例如RAID

Indexing Technology: why not B-tree?

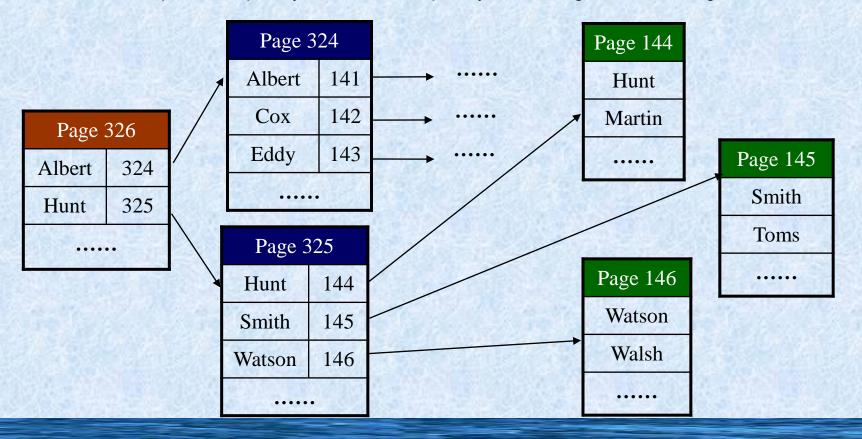
- B-tree is a widely-used technology in database
 - B-tree indexing has high performance when used to find some records in the database

■ While B-tree is not a good technology for data warehouse, Why?

B-tree Indexing Technology

Because:

- B-tree demands that the attribute must have many different values, such as itemID, customer ID, etc.
- B-tree demands that the query should have simpler conditions and less results
- The space complexity and time complexity of creating B-tree are huge



Indexing OLAP Data: Bitmap Index

■ Index on a particular column

Pogo toblo

- Each value in the column has a bit vector: bit-op is fast
- The length of the bit vector: # of records in the base table
- The +th bit is set if the +th row of the base table has the value for the indexed column

Indox on Dogion

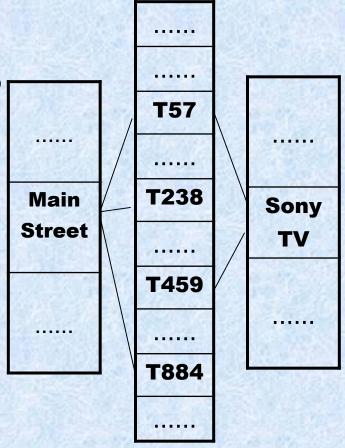
not suitable for high cardinality domains

| D | ase table | | mu | ex on | region | muex on Type | | | |
|------|-----------|--------|-------|-------|---------------|----------------|-------|--------|--------|
| Cust | Region | Туре | RecID | Asia | Europe | America | RecID | Retail | Dealer |
| C1 | Asia | Retail | 1 | 1 | 0 | 0 | 1 | 1 | 0 |
| C2 | Europe | Dealer | 2 | 0 | 1 | 0 | 2 | 0 | 1 |
| C3 | Asia | Dealer | 3 | 1 | 0 | 0 | 3 | 0 | 1 |
| C4 | America | Retail | 4 | 0 | 0 | 1 | 4 | 1 | 0 |
| C5 | Europe | Dealer | 5 | 0 | 1 | 0 | 5 | 0 | 1 |

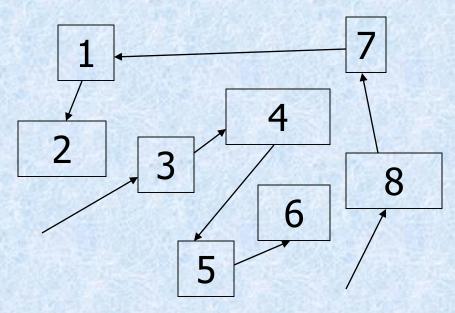
Inday on Type

Indexing OLAP Data: Join Index

- Join index: JI(R-id, S-id) where R (R-id, ...)
 ▷▷□ S (S-id, ...)
- Traditional indices map the values to a list of record ids
- In data warehouses, join index relates the values of the <u>dimensions</u> of a star schema to <u>rows</u> in the fact table.
 - E.g. fact table: Sales and two dimensions city and product
 - ✓ A join index on city maintains for each distinct city a list of R-IDs of the tuples recording the Sales in the city
 - Join indices can span multiple dimensions



■ Table Mergence(I)



Common access orders

3 4 5 6

8 7 1 2

Optimized access orders

■ Table Mergence(II)

| ID | NAME | TYPE |
|----|------------------------|----------|
| 1 | Dell PC | Computer |
| 2 | Cisco 2600 | Router |
| 3 | IBM PC | Computer |
| 4 | HP Laserjet 1000 | Printer |

| ID | Storage ID | Amount |
|-----|---------------|--------|
| 1\ | 5 | 300 |
| 1 \ | 3 | 35 |
| 2 | 3 | 67 |
| 2 | 2 | 348 |
| 3 | 6 | 20 |
| 4 | 1 | 908 |

Add Redundancy

ID Name Type Model Weight

ID Storage ID Amount ID Name Type Model Weight

ID Storage ID Amount Name

Original tables

Add redundancy

Dividing Table

- In logical design, big tables can be divided into small tables.
 When a big table is visited, the table can be substituted with small ones.
- In physical design, distributed storage methods can be used.
- Tables can be stored into a disk array

Data Warehousing and OLAP Technology

- What is a data warehouse?
- A multi-dimensional data model
- Data warehouse architecture
- Data warehouse implementation
 - Data warehouse Design
 - OLAP Modeling Methods
 - Optimization of Logical Model/Physical Model
 - Efficient Methods for Data Cube Computation
 - Plan and Implementation of Data Warehouse
- Further development of data cube technology
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Efficient Methods for Data Cube Computation

■ Data Cube: High Efficiency Computing

MutiWay Array Aggregation

• BUC

Data Cube: High Efficiency Computing

Data cube can be viewed as a lattice of cuboids

- The bottom-most cuboid is the base cuboid
- The top-most cuboid (apex) contains only one cell
- How many cuboids in an n-dimensional cube with L levels?

$$T = \prod_{i=1}^{n} (L_i + 1)$$

$T = \prod_{i=1}^{n} (L_i + 1)$ Materialization of data cube

- Materialize <u>every</u> (cuboid) (full materialization), <u>none</u> (no materialization), or some (partial materialization)
- Selection of which cuboids to materialize
 - ✓ Based on size, sharing, access frequency, etc.

Data Cube Operation

■ Cube definition and computation in DMQL

define cube sales[item, city, year]: sum(sales_in_dollars)
compute cube sales

■ Transform it into a SQL-like language (with a new operator cube by, introduced by Gray et al. '96)

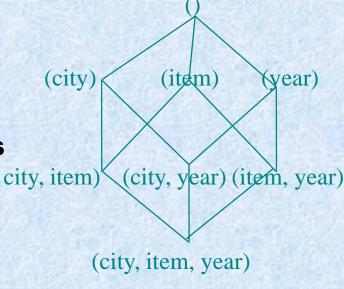
SELECT item, city, year, SUM (amount)

FROM SALES

CUBE BY item, city, year

№ Need compute the following Group-Bys

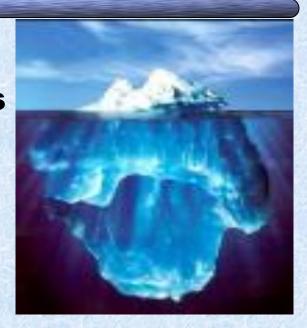
(date, product, customer), (date,product),(date, customer), (product, customer), (date), (product), (customer) ()



Iceberg Cube

■ Computing only the cuboid cells whose count or other aggregates satisfying the condition like

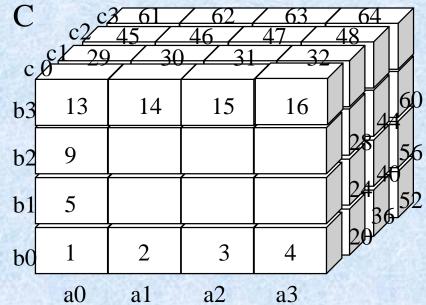
HAVING COUNT(*) >= minsup



Motivation

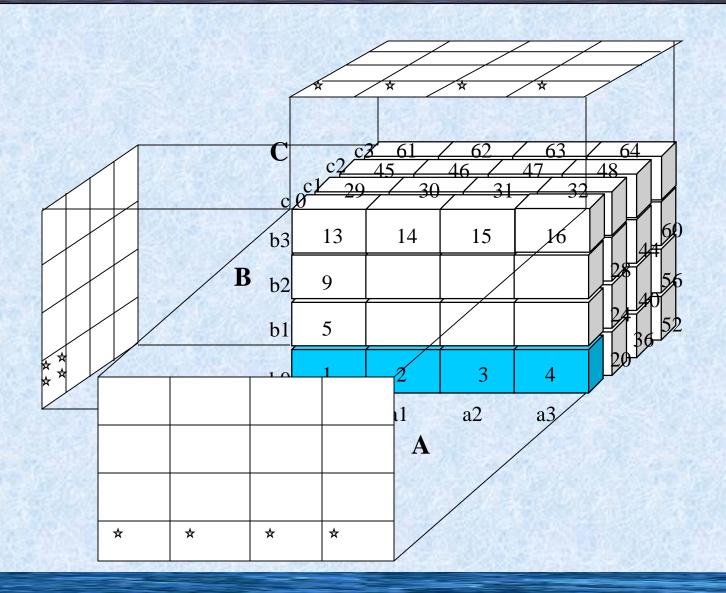
- Only a small portion of cube cells may be "above the water" in a sparse cube
- Only calculate "interesting" cells—data above certain threshold
- Avoid explosive growth of the cube

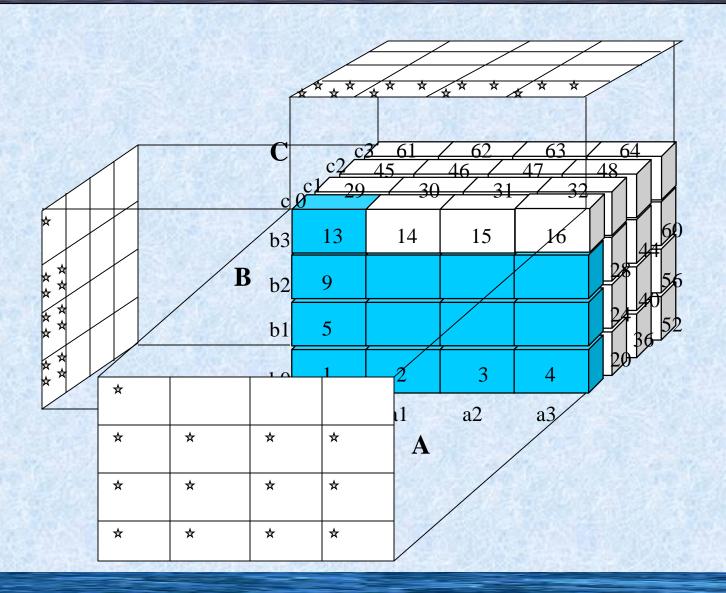
- Partition arrays into chunks (a small subcube which fits in memory).
- Compressed sparse array addressing: (chunk_id, offset)
- Compute aggregates in "multiway" by visiting cube cells in the order which minimizes the # of times to visit each cell, and reduces memory access and storage cost.

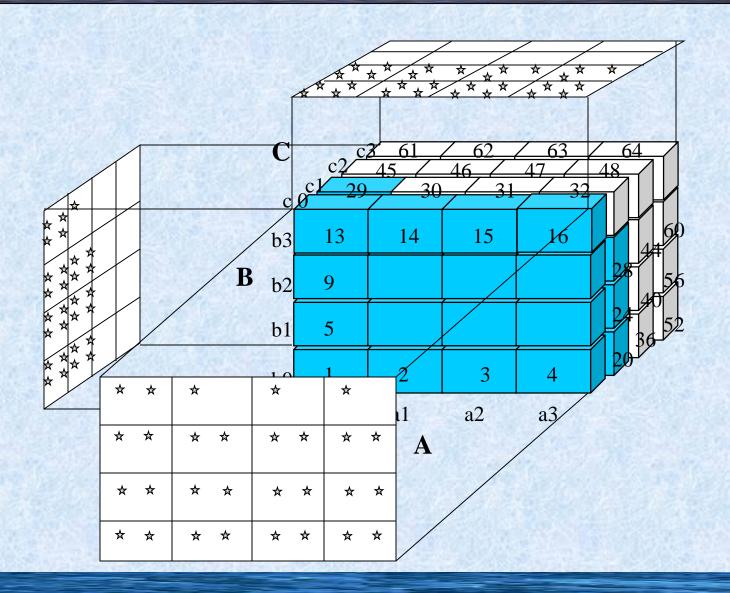


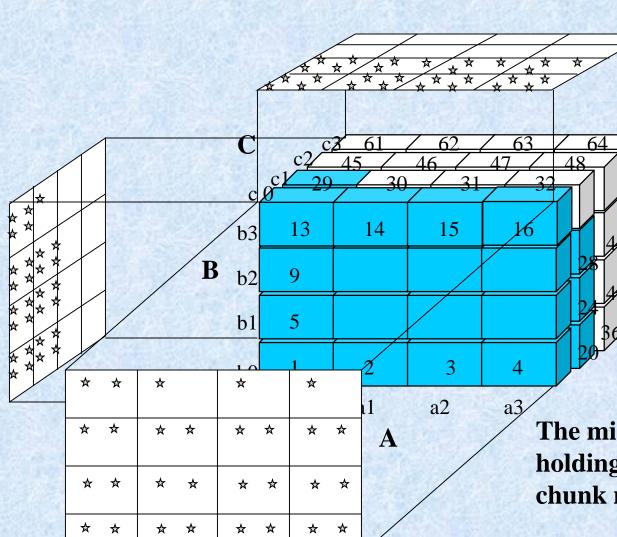
B

What is the best traversing order to do multi-way aggregation?









If the dimensions are sorted as:

A: 10, B:100, C: 1000

For BC plane:

100*1000=100,000

For AC plane:

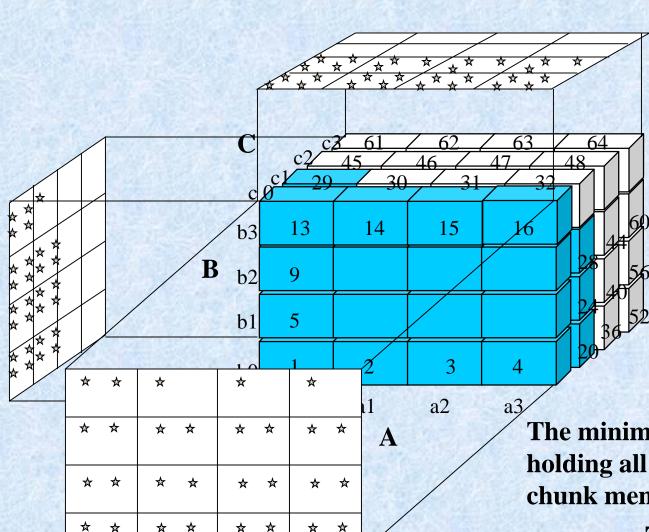
40*1000=40,000

For AB plane:

40*400=16,000

The minimal memory required for holding all relevant 2-D planes in chunk memory:

Total: 156,000



If the dimensions are sorted as:

A: 1000, B:100, C: 10

For BC plane:

10*100=1,000

For AC plane:

10*4000=40,000

For AB plane:

400*4000=1,600,000

The minimal memory required for holding all relevant 2-D planes in chunk memory:

Total: 1,641,000

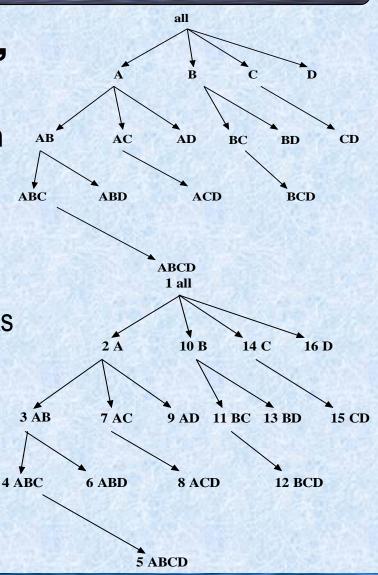
- Method: the planes should be sorted and computed according to their size in ascending order.
 - Idea: keep the smallest plane in the main memory, fetch and compute only one chunk at a time for the largest plane
- Limitation of the method: computing well only for a small number of dimensions
 - If there are a large number of dimensions, iceberg cube computation methods can be explored

Processing OLAP Query Effectively

- Determine which operations should be performed on the available cuboids:
 - transform drill, roll, etc. into corresponding SQL and/or OLAP operations, e.g, dice = selection + projection
- Determine to which materialized cuboid(s) the relevant operations should be applied.
- Exploring indexing structures and compressed vs. dense array structures in MOLAP

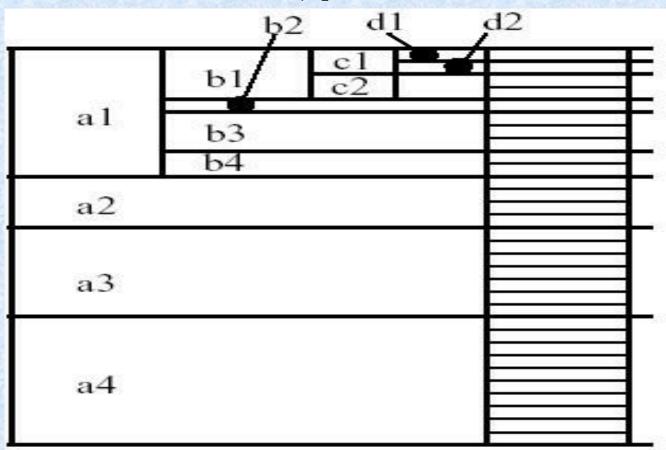
Bottom-Up Computation (BUC)

- BUC (Beyer & Ramakrishnan, SIGMOD'99)
- Bottom-up cube computation
- Divides dimensions into partitions and facilitates iceberg pruning
 - If a partition does not satisfy min_sup, its descendants can be pruned
 - ◆ If minsup = 1 ⇒ compute full CUBE!



BUC: Partitioning

- Usually, entire data set can't fit in main memory
- Sort distinct values, partition into blocks that fit



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数据仓库的投资分析

- ≥ 数据仓库的应用目标
 - ◆ 企业的核心业务
 - ◆ 优化企业内部管理控制
 - ◆ 为企业增加商业机会
- ≥ 建设数据仓库的必要性



可通过计算ROI(Return of Investment)来衡量投资回报的价值

数据仓库主题的选择和阶段规划

- 数据仓库的实施是一个极为复杂的长期过程,因此,应 选择当前最急需、能在短期内产生效益、业务模型清晰 的任务首期实现
- ≥ 选择首期实现主题的参考原则:
 - ◆ 优先实现管理者目前需求最迫切和最关心的主题
 - ◆ 优先选择能在短期内产生效益的主题
 - ◆ 推后选择业务逻辑准备不充分的主题
 - ◆ 推后实施技术难度较大、可实现性较低、投资风险大的主题
- ≥ 维护阶段
 - ◆ 数据仓库的维护极为重要,一般数据仓库在建立完成 之后,都需要一至两年的维护
 - ◆数据仓库的维护过程就是DSS逐步产生效益的过程

数据仓库后端工具

- 数据抽取(Data extraction):
 - get data from multiple, heterogeneous, and external sources
- 数据清洗(Data cleaning):
 - detect errors in the data and rectify them when possible
- 数据转换(Data transformation):
 - convert data from legacy or host format to warehouse format
- ≥ 数据装载(Load):
 - sort, summarize, consolidate, compute views, check integrity, and build indicies and partitions
- ☑ 刷新(Refresh):
 - propagate the updates from the data sources to the warehouse

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Discovery-Driven Exploration of Data Cubes

- Hypothesis-driven: exploration by user, huge search space
- Discovery-driven (Sarawagi et al. '98)
 - pre-compute measures indicating exceptions, guide user in the data analysis,
 at all levels of aggregation
 - Exception: significantly different from the value anticipated, based on a statistical model
 - Visual cues such as background color are used to reflect the degree of exception of each cell
 - Computation of exception indicator (modeling fitting and computing SelfExp, InExp, and PathExp values) can be overlapped with cube construction

Examples: Discovery-Driven Data Cubes

| item | all |
|--------|-----|
| region | all |

Selfexp: Background color; Inexp: Margin; Pathexp: Which path will lead to an exception?

| Sum of sales | mont | month | | | | | | | | | | | |
|--------------|------|-------|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|--|
| | Jan | Feb | Mar | Apr | May | Jun | Jul | Aug | Sep | Oct | Nov | Dec | |
| Total | | 1% | -1% | 0% | 1% | 3% | -1 | -9% | -1% | 2% | -4% | 3% | |

| Avg sales | топ | ıtlı | | | | | | | | | | |
|-------------------------|-----|------|-----|-----|-----|-----|------|---------------|-----|------|------|------|
| item | Jan | Feb | Mac | Apr | May | Jun | Jul | Aug | Sep | Oct | Nov | Dec |
| Sony b/w printer | | 9% | -8% | 2% | -5% | 14% | 4% | 0% | 41% | -13% | -15% | -11% |
| Sony color printer | | 0% | 0% | 3% | 2% | 4% | -10% | -13% | 0% | 4% | -6% | 4% |
| HP b/w printer | | -2% | 1% | 2% | 3% | 8% | 0% | -12% | -9% | 3% | -3% | 6% |
| HP color printer | | 0% | 0% | -2% | 1% | 0% | -1% | -7% | -2% | 1% | -5% | 1% |
| IBM home computer | | 1% | -2% | -1% | -1% | 3% | 3% | -10% | 4% | 1% | 4% | -1% |
| IBM laptop computer | | 0% | 0% | -1% | 3% | 4% | 2% | -1 0 % | -2% | 0% | -9% | 3% |
| Toshiba home computer | | -2% | -5% | 1% | 1% | -1% | 1% | 5% | -3% | -5% | -1% | -1% |
| Toshiba laptop computer | | 1% | 0% | 3% | 0% | -2% | -2% | -5% | 3% | 2% | -1% | 0% |
| Logitech mouse | | 3% | -2% | -1% | 0% | 4% | 6% | -11% | 2% | 1% | 4% | 0% |
| Ergo-way mouse | | 0% | 0% | 2% | 3% | 1% | -2% | -2% | -5% | 0% | -5% | 8% |

| item | IB1 | M home | comput | ter | | | | | | | | |
|--------------------------------|-----|-------------------------|--------------------------------|-------------------------|------------------------|-----------------------|-------------------------|--------------------------|-------------------------|-----------------------|-------------------------|------------------------|
| Avg sales | топ | nth | | | | | | | | | | |
| region | Jan | Feb | Mar | Арг | May | Jun | Jul | Aug | Sep | Oct | Nov | Dec |
| North South East West | | -1% -1% -1% 4% | -3% 1% -2% 0 % | -1% -9% 2% -1% | 0% 6% -3% -3% | 3% -1% 1% 5% | 4% -39% 18% 1% | -7% 9% -2% -18% | 1% -34% 11% 8% | 0% 4% -3% 5% | -3% 1% -2% -8% | -3% 7% -1% 1% |

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Data Warehouse Usage

Three kinds of data warehouse applications

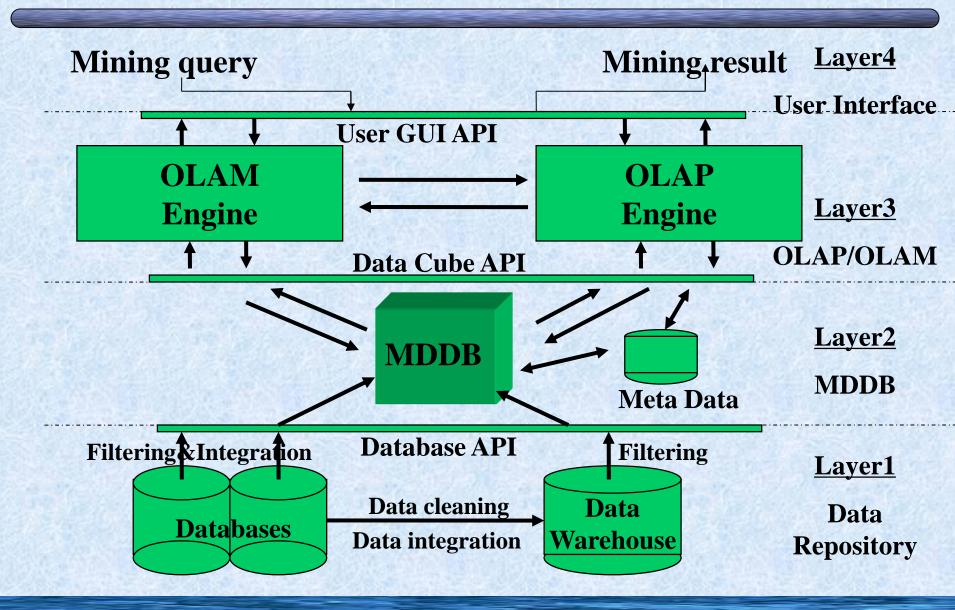
- Information processing
 - ✓ supports querying, basic statistical analysis, and reporting using crosstabs, tables, charts and graphs
- Analytical processing
 - ✓ multidimensional analysis of data warehouse data
 - ✓ supports basic OLAP operations, slice-dice, drilling, pivoting
- Data mining
 - ✓ knowledge discovery from hidden patterns
 - ✓ supports associations, constructing analytical models, performing classification and prediction, and presenting the mining results using visualization tools.
- **■** Differences among the three tasks

From On-Line Analytical Processing to On Line Analytical Mining (OLAM)

■ Why online analytical mining?

- High quality of data in data warehouses
 - ✓ DW contains integrated, consistent, cleaned data
- Available information processing structure surrounding data warehouses
 - ✓ODBC, OLEDB, Web accessing, service facilities, reporting and OLAP tools
- OLAP-based exploratory data analysis
 - ✓ mining with drilling, dicing, pivoting, etc.
- On-line selection of data mining functions
 - ✓ integration and swapping of multiple mining functions, algorithms, and tasks

An OLAM Architecture



Summary

Data warehouse

- A <u>subject-oriented</u>, <u>integrated</u>, <u>time-variant</u>, and <u>nonvolatile</u> collection of data in support of management's decision-making process
- A multi-dimensional model of a data warehouse
 - Star schema, snowflake schema, fact constellations
 - A data cube consists of dimensions & measures
- OLAP operations: drilling, rolling, slicing, dicing and pivoting
- **OLAP servers: ROLAP, MOLAP, HOLAP**
- Data warehouse implementation
- **■** Efficient computation of data cubes
 - Partial vs. full vs. no materialization
 - Multiway array aggregation
 - Bitmap index and join index implementations
- Further development of data cube technology
 - Discovery-drive and multi-feature cubes
 - From OLAP to OLAM (on-line analytical mining)

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Thank you !!!