



SPARK 2014: Ghost Code

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What is ghost code?

ghost code is part of the program that is added for the purpose of specification
Why3 team, "The Spirit of Ghost Code"

... or verification

addition by SPARK team

Examples of ghost code:

- contracts (Pre, Post, Contract_Cases, etc.)
- assertions (pragma Assert, loop (in)variants, etc.)
- special values Func'Result, Var'Old, Var'Loop_Entry

Is it enough?

Ghost code – A trivial example

```
Data : Data_Array;
Free : Natural;

procedure Alloc is
begin
    -- some computations here
    assert that Free "increases"
end Alloc;
```

how to express it?

Ghost variables – aka auxiliary variables

Variables declared with aspect Ghost

declaration is discarded by compiler when ghost code ignored

Ghost assignments to ghost variables

assignment is discarded by compiler when ghost code ignored

```
Data : Data_Array;
Free : Natural;

procedure Alloc is
    Free_Init : Natural with Ghost;

begin
    Free_Init := Free;
    -- some computations here
    pragma Assert (Free > Free_Init);
end Alloc;
```

Ghost variables – non-interference rules

Ghost variable cannot be assigned to non-ghost one

```
Free := Free_Init;
```

 Ghost variable cannot indirectly influence assignment to non-ghost one

```
if Free_Init < Max then
   Free := Free + 1;
end if;</pre>
```

```
procedure Assign (From : Natural; To : out Natural) is
begin
   To := From;
end Assign;

Assign (From => Free_Init, To => Free);
```

Ghost statements

Ghost variables can only appear in ghost statements

- assignments to ghost variables
- assertions and contracts
- calls to ghost procedures

```
procedure Assign (From : Natural; To : out Natural)
    with Ghost
is
begin
    To := From;
end Assign;

Assign (From => Free, To => Free_Init);

Assign (From => Free_Init, To => Free);
```

Ghost procedures

Ghost procedures cannot write into non-ghost variables

```
procedure Assign (Value : Natural) with Ghost is
begin
   Free := Value;
end Assign;
```

- Used to group statements on ghost variables
 - in particular statements not allowed in non-ghost procedures

```
procedure Assign_Cond (Value : Natural) with Ghost is
begin
    if Condition then
        Free_Init := Value;
    end if;
end Assign_Cond;
```

Can have Global (including Proof_In) & Depends contracts

Ghost functions

Functions for queries used only in contracts

```
procedure Alloc with
    Pre => Free_Memory > 0,
    Post => Free_Memory < Free_Memory'Old;

function Free_Memory return Natural with Ghost;</pre>
```

Typically implemented as expression functions

- in private part proof of client code can use expression
- or in body only proof of unit can use expression

```
function Free_Memory return Natural is (...);
-- if completion of ghost function declaration

function Free_Memory return Natural is (...) with Ghost;
-- if function body as declaration
```

Imported ghost functions

Ghost functions without a body

cannot be executed

```
function Free Memory return Natural with Ghost, Import;
```

Typically used with abstract ghost private types

definition in SPARK_Mode(Off) → type is abstract for GNATprove

```
type Memory_Chunks is private;
function Free_Memory return Memory_Chunks
    with Ghost, Import;
private
    pragma SPARK_Mode (Off);
    type Memory_Chunks is null record;
```

Definition of ghost types/functions given in proof

- either in Why3 using External_Axiomatization
- or in an interactive prover (Coq, Isabelle, etc.)

Ghost packages and ghost abstract state

- Every entity in a ghost package is ghost
 - local ghost package can group all ghost entities
 - library-level ghost package can be withed/used in regular units
- Ghost abstract state can only represent ghost variables

```
package Mem with
   Abstract_State => (State with Ghost)
is

package body Mem with
   Refined_State => (State => (Data, Free, Free_Init))
is
```

 Non-ghost abstract state can contain both ghost and nonghost variables

Executing ghost code

- Ghost code can be enabled globally
 - using compilation switch --gnata (for all assertions)
- Ghost code can be enabled selectively
 - using pragma Assertion_Policy (Ghost => Check)
 - SPARK rules enforce consistency in particular no write disabled

```
pragma Assertion_Policy (Ghost => Ignore, Pre => Check);

procedure Alloc with
    Pre => Free_Memory > 0;

function Free_Memory return Natural with Ghost;
```

GNATprove analyzes all ghost code and assertions

Example of use – encoding a state automaton

• Tetris in SPARK

at http://blog.adacore.com/tetris-in-spark-on-arm-cortex-m4

Global state encoded in global ghost variable

updated at the end of procedures of the API

```
type State is (Piece_Falling, ...) with Ghost;
Cur_State : State with Ghost;
```

Properties encoded in ghost functions

```
function Valid_Configuration return Boolean is
   (case Cur_State is
     when Piece_Falling => ...,
     when ...)
with Ghost;
```

Example of use – expressing useful lemmas

GCD in SPARK

at http://www.spark-2014.org/entries/detail/gnatprove-tips-and-tricks-proving-the-ghost-common-denominator-gcd

Lemmas expressed as ghost procedures

```
procedure Lemma_Not_Divisor (Arg1, Arg2 : Positive) with
  Ghost,
  Global => null,
  Pre => Arg1 in Arg2 / 2 + 1 .. Arg2 - 1,
  Post => not Divides (Arg1, Arg2);
```

Most complex lemmas further refined into other lemmas

code in procedure body used to guide proof (e.g. for induction)

Example of use – specifying an API through a model

Red black trees in SPARK

at http://www.spark-2014.org/entries/detail/research-corner-auto-active-verification-in-spark

Invariants of data structures expressed as ghost functions

using Type_Invariant on private types

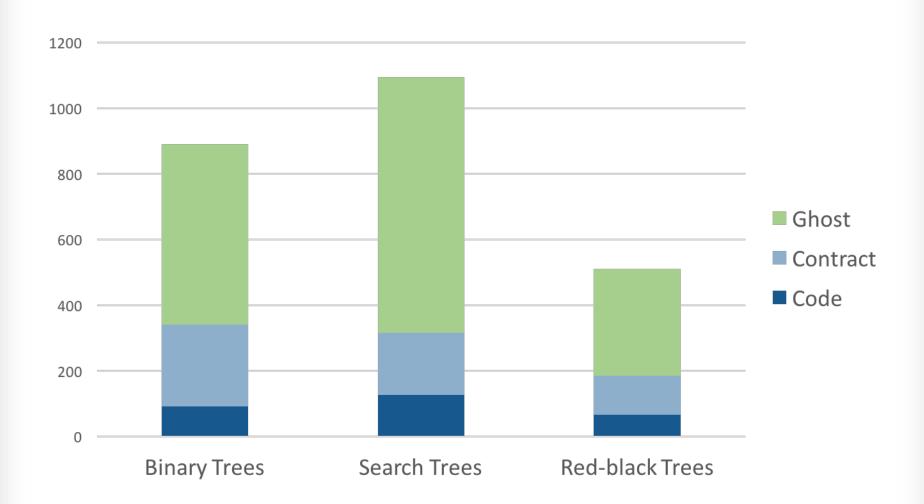
Model of data structures expressed as ghost functions

called from Pre/Post of subprograms from the API

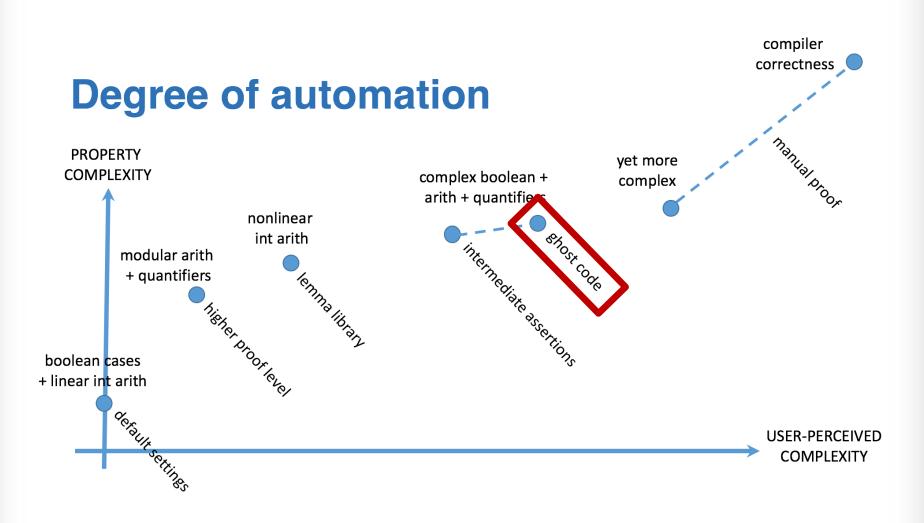
Lemmas expressed as ghost procedures

sometimes without contracts to benefit from inlining in proof

Extreme proving with ghost code – red black trees in SPARK



Positioning ghost code in proof techniques









Is this correct? 1/10



```
Data : Data_Array;
Free : Natural;

procedure Alloc is
    Free_Init : Natural with Ghost;
begin
    Free_Init := Free;
    -- some computations here
    if Free <= Free_Init then
        raise Program_Error;
    end if;
end Alloc;</pre>
```

Is this correct? 1/10



```
Data : Data_Array;
Free : Natural;

procedure Alloc is
   Free_Init : Natural with Ghost;
begin
   Free_Init := Free;
   -- some computations here
   if Free <= Free_Init then
        raise Program_Error;
   end if;
end Alloc;</pre>
```

ghost entity cannot appear in this context



Is this correct? 2/10



```
Data: Data Array;
Free: Natural;
procedure Alloc is
   Free Init : Natural with Ghost;
   procedure Check with Ghost is
   begin
      if Free <= Free Init then</pre>
         raise Program Error;
      end if:
   end Check;
begin
   Free Init := Free;
   -- some computations here
   Check;
end Alloc;
```



Is this correct? 2/10



```
Data: Data Array;
Free: Natural;
procedure Alloc is
   Free Init: Natural with Ghost;
   procedure Check with Ghost is
   begin
      if Free <= Free Init then</pre>
         raise Program Error;
      end if:
   end Check;
begin
   Free Init := Free;
   -- some computations here
   Check;
end Alloc;
```

Note that procedure Check is inlined for proof (no contract).



Is this correct? 3/10



```
pragma Assertion_Policy (Pre => Check);

procedure Alloc with
    Pre => Free_Memory > 0;

function Free_Memory return Natural with Ghost;
```



```
pragma Assertion Policy (Pre => Check);
procedure Alloc with
   Pre = Free Memory > 0;
function Free Memory return Natural with Ghost;
```

Incompatible ghost policies in effect during compilation, as ghost code is ignored by default.

Note that GNATprove accepts this code as it enables all ghost code and assertions.



Is this correct? 4/10





```
procedure Alloc with
   Post => Free Memory < Free Memory'Old;
function Free Memory return Natural with Ghost;
Max : constant := 1000;
function Free Memory return Natural is
begin
   return Max - Free + 1;
end Free Memory;
procedure Alloc is
begin
   Free := Free + 10;
end Alloc;
```

Is this correct? 4/10





```
procedure Alloc with
   Post => Free Memory < Free Memory'Old;
function Free Memory return Natural with Ghost;
Max : constant := 1000;
function Fre∉ Memory return Natural is
begin
   return Max - Free + 1;
end Free Memory;
procedure Alloc is
begin
   Free = Free + 10;
end Alloc;
```

No postcondition on Free_Memory that would allow proving the postcondition on Alloc.



Is this correct? 5/10



```
procedure Alloc with
    Post => Free_Memory < Free_Memory'Old;

function Free_Memory return Natural with Ghost;

Max : constant := 1000;

function Free_Memory return Natural is (Max - Free + 1);

procedure Alloc is
begin
    Free := Free + 10;
end Alloc;</pre>
```



Is this correct? 5/10





```
procedure Alloc with
    Post => Free_Memory < Free_Memory'Old;

function Free_Memory return Natural with Ghost;

Max : constant := 1000;

function Free_Memory return Natural is (Max - Free + 1);

procedure Alloc is begin
    Free := Free + 10;
end Alloc;</pre>
```

Free_Memory has an implicit postcondition as an expression function.



Is this correct? 6/10



```
subtype Resource is Natural range 0 .. 1000;
subtype Num is Natural range 0 .. 6;
subtype Index is Num range 1 .. 6;
type Data is array (Index) of Resource;
function Sum (D : Data; To : Num) return Natural is
  (if To = 0 then 0 else D(To) + Sum(D, To-1))
with Ghost;
procedure Create (D : out Data) with
   Post \Rightarrow Sum (D, D'Last) < 42
is
begin
  for J in D'Range loop
      D(J) := J;
     pragma Loop_Invariant (2 * Sum(D,J) \le J * (J+1));
   end loop;
end Create;
```

Is this correct? 6/10



```
subtype Resource is Natural range 0 .. 1000;
subtype Num is Natural range 0 .. 6;
subtype Index is Num range 1 .. 6;
type Data is array (Index) of Resource;
function Sum (D : Data; To : Num) return Natural is
  (if To = 0 then 0 else D(To) + Sum(D, To-1))
with Ghost;
procedure Create (D : out Data) with
   Post \Rightarrow Sum (D, D'Last) < 42
is
begin
  for J in D'Range loop
      D(J) := J;
     pragma Loop Invariant (2 * Sum(D, J) \le J * (J+1));
   end loop;
end Create;
```

3

info: expression function body not available for proof
("Sum" may not return)



Is this correct? 7/10



```
subtype Resource is Natural range 0 .. 1000;
subtype Num is Natural range 0 .. 6;
subtype Index is Num range 1 .. 6;
type Data is array (Index) of Resource;
function Sum (D : Data; To : Num) return Natural is
  (if To = 0 then 0 else D(To) + Sum(D, To-1))
with Ghost, Annotate => (GNATprove, Terminating);
procedure Create (D : out Data) with
   Post \Rightarrow Sum (D, D'Last) < 42
is
begin
  for J in D'Range loop
      D(J) := J;
     pragma Loop_Invariant (2 * Sum(D,J) \le J * (J+1));
   end loop;
end Create;
```



Is this correct? 7/10



```
subtype Resource is Natural range 0 .. 1000;
subtype Num is Natural range 0 .. 6;
subtype Index is Num range 1 .. 6;
type Data is array (Index) of Resource;
function Sum (D : Data; To : Num) return Natural is
  (if To = 0 then 0 else D(To) + Sum(D, To-1))
with Ghost, Annotate => (GNATprove, Terminating);
procedure Create (D : out Data) with
   Post \Rightarrow Sum (D, D'Last) < 42
is
begin
  for J in D'Range loop
      D(J) := J;
     pragma Loop_Invariant (2 * Sum(D, J) \le J * (J+1));
   end loop;
end Create;
```

Note that GNATprove does not prove the termination of Sum here.



Is this correct? 8/10



```
subtype Resource is Natural range 0 .. 1000;
subtype Num is Natural range 0 .. 6;
subtype Index is Num range 1 .. 6;
type Data is array (Index) of Resource;
function Sum (D : Data; To : Num) return Natural is
  (if To = 0 then 0 else D(To) + Sum(D, To-1))
with Ghost, Annotate => (GNATprove, Terminating);
procedure Create (D : out Data) with
   Post \Rightarrow Sum (D, D'Last) < 42
is
begin
  for J in D'Range loop
      D(J) := J;
   end loop;
end Create;
```



Is this correct? 8/10



```
subtype Resource is Natural range 0 .. 1000;
subtype Num is Natural range 0 .. 6;
subtype Index is Num range 1 .. 6;
type Data is array (Index) of Resource;
function Sum (D : Data; To : Num) return Natural is
  (if To = 0 then 0 else D(To) + Sum(D, To-1))
with Ghost, Annotate => (GNATprove, Terminating);
procedure Create (D : out Data) with
   Post \Rightarrow Sum (D, D'Last) < 42
is
begin
  for J in D'Range loop
      D(J) := J;
   end loop;
end Create;
```

The loop is unrolled by GNATprove here, as D'Range is 0..6. The automatic prover unrolls the recursive definition of Sum.



Is this correct? 9/10



```
subtype Resource is Natural range 0 .. 1000;
subtype Index is Natural range 1 .. 42;
package Segs is new
  Ada. Containers. Functional Vectors (Index, Resource);
use Seqs;
function Create return Sequence with
   Post => (for all K in 1 .. Last (Create'Result) =>
              Get (Create'Result, K) = K)
is
   S : Sequence;
begin
  for K in 1...42 loop
     S := Add (S, K);
   end loop;
  return S;
end Create;
```

Is this correct? 9/10



```
subtype Resource is Natural range 0 .. 1000;
subtype Index is Natural range 1 .. 42;
package Segs is new
  Ada. Containers. Functional Vectors (Index, Resource);
use Seqs;
function Create return Sequence with
   Post => (for all K in 1 .. Last (Create'Result) =>
              Get (Create'Result, K) = K)
is
   S : Sequence;
begin
  for K in 1 .. 42 loop
   S := Add (S, K);
   end loop;
  return S;
end Create;
```

Loop requires a loop invariant to prove the postcondition.



Is this correct? 10/10



```
subtype Resource is Natural range 0 .. 1000;
subtype Index is Natural range 1 .. 42;
package Segs is new
  Ada. Containers. Functional Vectors (Index, Resource);
use Seqs;
function Create return Sequence with
   Post => (for all K in 1 .. Last (Create'Result) =>
              Get (Create'Result, K) = K)
is
   S : Sequence;
begin
  for K in 1...42 loop
     S := Add (S, K);
     pragma Loop_Invariant (Integer (Length (S)) = K);
     pragma Loop Invariant
       (for all J in 1 ... K \Rightarrow Get(S, J) = J);
   end loop;
  return S;
end Create;
```



Is this correct? 10/10



```
subtype Resource is Natural range 0 .. 1000;
subtype Index is Natural range 1 .. 42;
package Segs is new
  Ada. Containers. Functional Vectors (Index, Resource);
use Seqs;
function Create return Sequence with
   Post => (for all K in 1 .. Last (Create'Result) =>
              Get (Create'Result, K) = K)
is
   S : Sequence;
begin
  for K in 1...42 loop
     S := Add (S, K);
     pragma Loop_Invariant (Integer (Length (S)) = K);
    pragma Loop Invariant
       (for all J in 1 ... K \Rightarrow Get(S, J) = J);
   end loop;
  return S;
end Create;
```





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