



SPARK 2014: Overview

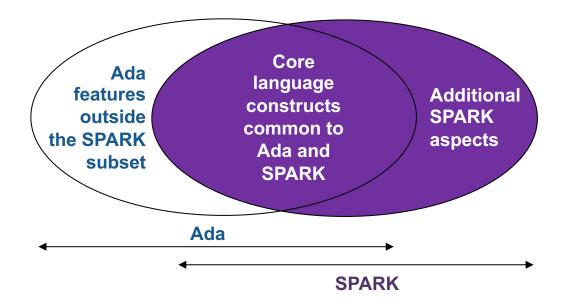
Claire Dross and Martyn Pike

University.adacore.com

SPARK 2014 – What is it?

A programming language

- A subset of Ada, designed for static verification
- Additional features to enhance program specification



A set of program verification tools

SPARK 2014 – What do the tools do?

Performs static verification of Ada source code. It can check that:

- a program is "well-formed"
- it performs its intended functionality

Static verification can take several forms:

- static-semantic checking (strong typing, visibility...)
- flow analysis (initialization of variables, data dependencies...)
- integrity and functional verification (absence of runtime error, functional correctness...)

SPARK 2014 – Key Tools

GNAT compiler

- checks conformance of source with Ada 2012 and SPARK 2014
- compiles source into executable

GNATprove

- performs additional SPARK 2014 checks
- performs flow analysis
- performs integrity and functional verification using formal proof
- GNAT Programming Studio (GPS)
- GNATbench plugin of Eclipse

SPARK 2014 – A trivial example

```
procedure Increment
           (X : in out Integer)
                                            data dependencies 🚫
  with Global => null,
                                            flow dependencies 🗸
        Depends \Rightarrow (X \Rightarrow X),
       Pre => X < Integer'Last,</pre>
                                            functionality 🗸
        Post \Rightarrow X = X'Old + 1;
procedure Increment
           (X : in out Integer)
is
begin
                                            absence of runtime error 🚫
   X := X + 1;
end Increment;
```

SPARK 2014 – The Programming Language

SPARK design principles:

- Exclude language features difficult to specify/verify
- Eliminate sources of ambiguity

Excluded from SPARK:

- access types and allocators
- expressions (including function calls) with side-effects
- aliasing of names
- goto statements
- controlled types
- handling of exceptions (raise statements can be used in a program but proof will attempt to show that they cannot be executed)

SPARK 2014 – Limitations – no side-effects in expressions

- The analysis of SPARK 2014 programs relies on the property that expressions are free from side-effects
- An expression is side-effect free if its evaluation does not update any object
- Side effects can introduce non-determinism in expression evaluation

```
G : Integer;
function F (X : in out Integer) return Integer;
G := F (G) + F (G); -- ??
```

SPARK 2014 – Limitations – no side-effects in expressions

 Function calls can appear in expressions, so function calls must be free from side-effects

 The objects updated by a function call are any parameters of mode out (or in out), plus global variables updated by the function body

```
function F (X : in out Integer) return Integer; -- Illegal
function Incr (X : Integer) return Integer; -- OK?
function Incr_And_Log (X : Integer) return Integer; -- OK?
```

SPARK 2014 – Limitations – no side-effects in expressions

• This is OK...

```
function Incr (X : Integer) return Integer; -- OK?
function Incr_And_Log (X : Integer) return Integer; -- OK?
```

 ... provided that the analysis of the bodies doesn't detect any side-effects

```
function Incr (X : in Integer) return Integer
is (X + 1); -- OK

Call_Count : Natural := 0;

function Incr_And_Log (X : in Integer) return Integer is
begin
    Call_Count := Call_Count + 1; -- Illegal
    return X + 1;
end Incr_And_Log;
```

SPARK 2014 – Limitations – no aliasing of names

- Aliasing occurs when two names refer to the same object
- SPARK forbids aliasing of outputs of subprograms (parameters of mode out or in out and global variables modified by the subprogram)
- This can make the subprogram give different results depending on parameter passing mechanism (by reference or by value)
- Results can seem unexpected, even for the user

SPARK 2014 – Limitations – no aliasing of names

```
Total : Natural := 0;

procedure Move_To_Total (Source : in out Natural) is
begin
   Total := Total + Source;
   Source := 0;
end Move_To_Total;
```

- The above subprogram is OK
- Absence of aliasing will be verified whenever a call is made:

```
X : Natural := 3;

Move_To_Total (X); -- OK
Move_To_Total (Total); -- Error
```

SPARK 2014 – Identifying SPARK Code

- SPARK and full Ada code can be mixed
 - To verify as much as possible of a code base
 - While retaining usage of excluded features when necessary
- SPARK_Mode (aspect or pragma) specifies whether code is intended to be SPARK
 - It is then checked by the GNATprove tool
- Defaults to 'Off', should be added to units to be analyzed
 - Or can be set globally using a configuration pragma
 - SPARK_Mode on 'withed' units defaults to 'auto' (up to the tool to determine whether or not a given construct is in SPARK)

SPARK 2014 – Identifying SPARK Code

```
package P
 with SPARK Mode => On
is
   -- package spec is SPARK, so can be used
   -- by SPARK clients
end P;
package body P
  with SPARK Mode => Off
is
   -- body is NOT SPARK, so assumed to
   -- be full Ada
end P;
```

SPARK 2014 – Identifying SPARK Code

- SPARK_Mode can be 'On' or 'Off' for
 - visible part of package specification
 - private part of package specification
 - declarative part of package body
 - statement part of package body
 - subprogram specification
 - subprogram body
- If SPARK mode is 'Off' for a package or subprogram then it cannot be switched 'On' for a later part of that package or subprogram







Is this correct? 1/10



```
package Stack_Package
  with SPARK_Mode => On

is
  type Element is new Natural;
  type Stack is private;

function Empty return Stack;
  procedure Push (S : in out Stack; E : Element);
  function Pop (S : in out Stack) return Element;

private
  ...
end Stack_Package;
```

Is this correct? 1/10



```
package Stack_Package
  with SPARK_Mode => On
is
  type Element is new Natural;
  type Stack is private;

function Empty return Stack;
  procedure Push (S : in out Stack; E : Element);
  function Pop (S : in out Stack) return Element;

private
  ...
end Stack_Package;
```

Function with "in out"
parameter is not allowed in
SPARK



Is this correct? 2/10





```
package body Global Stack
 with SPARK Mode => On
is
  Max : constant Natural := 100;
  type Element Array is array (1 .. Max) of Element;
  Content : Element Array;
  Top : Natural;
  function Pop return Element is
    E : constant Element := Content (Top);
 begin
    Top := Top -1;
    return E;
  end Pop;
end Global Stack;
```

Is this correct? 2/10



```
package body Global Stack
 with SPARK Mode => On
is
 Max : constant Natural := 100;
  type Element Array is array (1 .. Max) of Element;
  Content : Element Array;
  Top : Natural;
  function Pop return Element is
   E : constant Element := Content (Top);
 begin
    Top := Top -1;
   return \s
  end Pop;
end Global Stack;
```

Function updating a global variable is not allowed in SPARK



Is this correct? 3/10



```
package body P
  with SPARK Mode => On
is
  procedure Permute (X, Y, Z : in out Positive) is
    Tmp : constant Positive := X;
 begin
   X := Y;
   Y := Z;
    Z := Tmp;
  end Permute;
  procedure Swap (X, Y : in out Positive) is
 begin
    Permute (X, Y, Y);
  end Swap;
end P;
```

Is this correct? 3/10





```
package body P
  with SPARK Mode => On
is
 procedure Permute (X, Y, Z : in out Positive) is
    Tmp : constant Positive := X;
 begin
   X := Y;
   Y := Z;
    Z := Tmp;
  end Permute;
 procedure Swap (X, Y : in out Positive) is
 begin
    Permute (X, Y, Y);
  end Swap;
end P;
```

Aliasing between "in out" parameters is not allowed in SPARK



Is this correct? 4/10



```
package body P
  with SPARK Mode => On
is
  procedure Swap (X, Y : in out Positive);
  type Rec is record
   F1 : Positive;
   F2 : Positive;
  end record;
  procedure Swap Fields (R : in out Rec) is
  begin
    Swap (R.F1, R.F2);
  end Swap Fields;
end P;
```



Is this correct? 4/10



```
package body P
  with SPARK Mode => On
is
 procedure Swap (X, Y : in out Positive);
  type Rec is record
   F1 : Positive;
   F2 : Positive;
  end record;
 procedure Swap Fields (R : in out Rec) is
 begin
    Swap (R.F1, R.F2);
  end Swap Fields;
end P;
```

Two different fields of the same record always refer to distinct objects.



Is this correct? 5/10



```
package body P
  with SPARK_Mode => On
is
  procedure Swap (X, Y : in out Positive);

type P_Array is array (Natural range <>) of Positive;

procedure Swap_Indexes (A : in out P_Array, I, J : Natural) is
begin
    Swap (P (I), P (J));
end Swap_Indexes;
...
end P;
```

Is this correct? 5/10



```
package body P
  with SPARK_Mode => On
is
  procedure Swap (X, Y : in out Positive);

type P_Array is array (Natural range <>) of Positive;

procedure Swap_Indexes (A : in out P_Array, I, J : Natural) is
begin
  Swap (A (I), A (J));
end Swap_Indexes;
...
end P;
```

If I = J, the two
components of the array can
be aliased.

GNATprove will detect this possibility



Is this correct? 6/10



```
package P
  with SPARK Mode => On
is
  subtype Letter is Character range 'a' .. 'z';
  type String Access is access all String;
  type Dictionary is array (Letter) of String Access;
 procedure Store (D : in out Dictionary; W : String);
end P;
package body P
 with SPARK Mode => On
is
 procedure Store (D : in out Dictionary; W : String) is
    First Letter : constant Letter := W (W'First);
 begin
   D (First Letter) := new String'(W);
  end Store;
end P;
```



Is this correct? 6/10



```
package P
  with SPARK Mode => On
is
  subtype Letter is Character range 'a' .. 'z';
type String Access is access all String;
  type Dictionary is array (Letter) of String Access;
 procedure Store (D : in out Dictionary; W : String);
end P;
package body P
  with SPARK Mode => On
is
  procedure Store (D : in out Diationary; W : String) is
    First Letter : constant Letter := W (W'First);
 begin
    D (First Letter) := new String'(N);
  end Store;
end P;
```

General access types are not allowed in SPARK



Is this correct? 7/10



```
package P
  with SPARK Mode => On
is
  subtype Letter is Character range 'a' .. 'z';
  type String Access is private;
  type Dictionary is array (Letter) of String Access;
  function New String Access (W : String) return String Access;
 procedure Store (D : in out Dictionary; W : String);
private
 pragma SPARK Mode (Off);
  type String Access is access all String;
  function New String Access (W : String) return String Access is
    (new String'(W));
end P;
```



Is this correct? 7/10



```
package P
  with SPARK Mode => On
is
  subtype Letter is Character range 'a' .. 'z';
  type String Access is private;
  type Dictionary is array (Letter) of String Access;
  function New String Access (W : String) return String Access;
 procedure Store (D : in out Dictionary; W : String);
private
 pragma SPARK Mode (Off);
  type String Access is access all String;
  function New String Access (W : String) return String Access is
    (new String'(W));
end P:
```

SPARK_Mode has been turned off in the private section

The type String Access is declared in full Ada.



Is this correct? 8/10



```
package P with SPARK Mode => On is
  subtype Letter is Character range 'a' .. 'z';
  type String Access is private;
  type Dictionary is array (Letter) of String Access;
  function New String Access (W : String) return String Access;
  procedure Store (D : in out Dictionary; W : String);
private
 pragma SPARK Mode (Off);
end P;
package body P with SPARK Mode => On is
  procedure Store (D : in out Dictionary; W : String) is
    First Letter : constant Letter := W (W'First);
  begin
    D (First Letter) := New String Access (W);
  end Store;
end P;
```



Is this correct? 8/10



```
package P with SPARK Mode => On is
   subtype Letter is Character range 'a' .. 'z';
   type String Access is private;
   type Dictionary is array (Letter) of String Access;
   function New String Access (W : String) return String Access;
  procedure Store (D : in out Dictionary; W : String);
private
  pragma SPARK Mode (Off);
 end P;
package body P with SPARK Mode => On is
  procedure Store ( : in out Dictionary; W : String) is
     First Letter : constant Letter := W (W'First);
  begin
    D (First Letter) := New String Access (W);
  end Store;
 end P;
```

SPARK mode is 'Off' for P's private part. It cannot be switched back 'On' for P's body.



Is this correct? 9/10



```
package P with SPARK Mode => On is
  subtype Letter is Character range 'a' .. 'z';
  type String Access is private;
  type Dictionary is array (Letter) of String Access;
  function New String Access (W : String) return String Access;
private
 pragma SPARK Mode (Off);
end P;
with P; use P;
package Q with SPARK Mode => On is
 procedure Store (D : in out Dictionary; W : String);
end 0;
package body Q with SPARK Mode => On is
  procedure Store (D : in out Dictionary; W : String) is
    First Letter : constant Letter := W (W'First);
 begin
    D (First Letter) := New String Access (W);
  end Store;
end 0;
```



Is this correct? 9/10



```
package P with SPARK Mode => On is
  subtype Letter is Character range 'a' .. 'z';
  type String Access is private;
  type Dictionary is array (Letter) of String Access;
  function New String Access (W : String) return String Access;
private
 pragma SPARK Mode (Off);
end P;
with P; use P;
package Q with SPARK Mode => On is
 procedure Store (D : in out Dictionary; W : String);
end 0;
package body Q with SPARK Mode => On is
  procedure Store (D : in out Dictionary; W : String) is
    First Letter : constant Letter := W (W'First);
 begin
    D (First Letter) := New String Access (W);
  end Store;
end 0;
```



Is this correct? 10/10



```
package body P with SPARK Mode => On is
  type N Array is array (Positive range <>) of Natural;
  Not Found : exception;
  function Search Zero P (A : N Array) return Positive is
 begin
    for I in A'Range loop
      if A(I) = 0 then
        return I;
      end if:
    end loop;
    raise Not Found;
  end Search Zero P;
  function Search Zero N (A : N Array) return Natural
    with SPARK Mode => Off is
 begin
    return Search Zero P (A);
  exception
    when Not Found => return 0;
  end Search Zero N;
end P;
```



Is this correct? 10/10



```
package body P with SPARK Mode => On is
  type N Array is array (Positive range <>) of Natural;
  Not Found : exception;
  function Search Zero P (A : N Array) return Positive is
 begin
    for I in A'Range loop
      if A (I) = 0 then
        return I;
      end if;
    end loop;
    raise Not Found;
  end Search Zero P;
  function Search Zero N (A : N Array) return Natural
   with SPARK Mode => Off is
  begin
    return Search Zero P (A);
  exception
    when Not Found => return 0;
  end Search Zero N;
end P;
```

Raising an exception is allowed. GNATprove will try to demonstrate that it will never happen at runtime.

Handlers must be in full Ada parts though.





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