

GaussSynth.m -- The Two-Qubit Clifford + CS Circuit Synthesis Package

Written and Maintained by Andrew Glaudell

The GaussSynth.m package is a package for quantum compiling on two qubits using the Clifford group and the Controlled-Phase gate CS. The circuits which are exactly expressible over this gate set constitute every 4×4 unitary matrix which can be written as a matrix of Gaussian integers divided by some non-negative integer power of $2^{1/2}$ -- hence the package name. In this package, we supply a number of functions for performing quantum circuit synthesis on this gate set, both in the exact and approximate case. The algorithms in this package are based off of the work of Andrew Glaudell, Julien Ross, Matthew Amy, and Jake Taylor, and for details related to how these algorithms were developed, I suggest reading the articles [1-3] in the sources section below.

Package Details

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History:

1.0 - Initial version, completed 11/4/2019

1.1 - Added further functionality to the code, 10/2/2020

Keywords: Quantum Compiling, Quantum Circuit Synthesis, Clifford Group, Controlled Phase Gate, Normal Forms, Exact Synthesis, Approximate Synthesis

Sources:

[1] Amy, Matthew, Andrew N. Glaudell, and Neil J. Ross. "Number-theoretic characterizations of some restricted clifford+ t circuits." *Quantum* 4 (2020): 252.

[2] Glaudell, Andrew N., Neil J. Ross, and Jacob M. Taylor. "Optimal Two-Qubit Circuits for Universal Fault-Tolerant Quantum Computation." *arXiv preprint arXiv:2001.05997* (2020). (under review)

[3] Glaudell, Andrew Noble. *Quantum Compiling Methods for Fault-Tolerant Gate Sets of Dimension Greater than Two*. Diss. 2019.

Warnings:

I have used a fair amount of input checking so that functions only accept inputs of the appropriate form. This comes at the cost of some speed -- that being said, these checks cause constant overhead, and so their performance impact is worth it to prevent some erroneous calculation from being carried out. If you don't care about this input checking, one could relatively easily define their own functions from my own internal ones to slightly speed up their performance.

Limitations:

This package is only intended for usage on two-qubit circuits. To perform circuit synthesis on larger circuits, I suggest loading this package and using these functions as subroutines.

Discussion:

Rather than describe these algorithms in detail here, I defer to the sources [1-3] listed above or the function descriptions.

Requirements:

None

```
BeginPackage["GaussSynth`"];
```

Function Usage

```
GaussSynth::usage = "GaussSynth is a package for quantum compiling for two-qubit c";

Id::usage = "Id is the 4x4 Identity Matrix.";
X1::usage = "X1 is the unitary representation of the  $X \otimes I$  gate.";
X2::usage = "X2 is the unitary representation of the  $I \otimes X$  gate.";
Z1::usage = "Z1 is the unitary representation of the  $Z \otimes I$  gate.";
Z2::usage = "Z2 is the unitary representation of the  $I \otimes Z$  gate.";
W::usage = "W is the unitary representation of the primitive 8th root of unity  $\omega$ .";
H1::usage = "H1 is the unitary representation of the  $H \otimes I$  gate.";
H2::usage = "H2 is the unitary representation of the  $I \otimes H$  gate.";
S1::usage = "S1 is the unitary representation of the  $S \otimes I$  gate.";
S2::usage = "S2 is the unitary representation of the  $I \otimes S$  gate.";
CZ::usage = "CZ is the unitary representation of the CZ gate.";
CNOT12::usage = "CNOT12 is the unitary representation of the CNOT gate with control";
CNOT21::usage = "CNOT21 is the unitary representation of the CNOT gate with control";
EX::usage = "EX is the unitary representation of the SWAP (Exchange) gate.";
CS::usage = "CS is the unitary representation of the CS gate.";

R::usage = "R[P,Q] takes as input two strings P and Q of the form \"AB\" or \"-AB\".
  These correspond to two-qubit Hermitian Pauli operators, and if P and Q are defined,
  computes the corresponding unitary gate as defined in our work. These compositions
  to study the Clifford + CS gate set.";

U4ToS06::usage = "U4ToS06[U] maps the 4x4 unitary U to the equivalent S0(6) representation
  for some phase  $\varphi$  and U' an element of SU(4).";

IdS06::usage = "Id is the 6x6 Identity Matrix.";
X1S06::usage = "X1S06 is the S0(6) representation of the  $X \otimes I$  gate.";
X2S06::usage = "X2S06 is the S0(6) representation of the  $I \otimes X$  gate.";
Z1S06::usage = "Z1S06 is the S0(6) representation of the  $Z \otimes I$  gate.";
Z2S06::usage = "Z2S06 is the S0(6) representation of the  $I \otimes Z$  gate.";
IS06::usage = "IS06 is the S0(6) representation of the complex phase I. As SU(4) element, it is the 4x4 unitary representation of the phase I.";
```

```

we have  $\omega^2 = \text{Id}_{S(6)}$ .";
H1S06::usage = "H1S06 is the  $S(6)$  representation of the  $H \otimes I$  gate.";
H2S06::usage = "H2S06 is the  $S(6)$  representation of the  $I \otimes H$  gate.";
S1S06::usage = "S1S06 is the  $S(6)$  representation of the  $S \otimes I$  gate.";
S2S06::usage = "S2S06 is the  $S(6)$  representation of the  $I \otimes S$  gate.";
CZS06::usage = "CZS06 is the  $S(6)$  representation of the CZ gate. Note that this is
    by a primitive 8th root of unity  $\omega$  before
    performing the transformation.";
CNOT12S06::usage = "CNOT12S06 is the  $S(6)$  representation of the CNOT gate with control 1 and target 2.
    Note that this gate is not special unitary, and so we multiply by a primitive 8th root of unity  $\omega$ 
    before performing the transformation.";
CNOT21S06::usage = "CNOT21S06 is the  $S(6)$  representation of the CNOT gate with control 2 and target 1.
    Note that this gate is not special unitary, and so we multiply by a primitive 8th root of unity  $\omega$ 
    before performing the transformation.";
EXS06::usage = "EXS06 is the  $S(6)$  representation of the SWAP (Exchange) gate. Note that this is not special unitary,
    and so we multiply by a primitive 8th root of unity  $\omega$  before performing the transformation.";
CSS06::usage = "CSS06 is the  $S(6)$  representation of the CS gate. Note that this is not special unitary,
    a primitive 16th root of unity  $\omega^2$  before performing the transformation.";

RS06::usage = "R[P,Q] takes as input two strings P and Q of the form \"AB\" or \"C\". These correspond to two-qubit Hermitian Pauli operators, and if P and Q are different, R computes the corresponding  $S(6)$  gate as defined in our work. These composite gates are used to study the Clifford + CS gate set. Note that this gate is not special unitary, and so we multiply by a primitive 8th root of unity  $\omega$  before performing the transformation.";

FromSequence::usage = "FromSequence[str] reads in the string str and interprets the characters as gates. The string str should be a valid sequence of gates, and the output is a list of gates."

FromHexDec::usage = "FromHexDec[str] attempts to read in a string of a signed hexadecimal number. The sign indicates whether the operator corresponds to using symmetric or asymmetric gates. The string str should be a valid signed hexadecimal number, and the output is a list of gates. After the syllables comes the flag 00000 which indicates the end of the string."

CliffordQ::usage = "CliffordQ[U] returns True if U is a Clifford and False otherwise."
CliffordSynth::usage = "CliffordSynth[U] gives the index number (in Hexadecimal) of the Clifford gate U. The input U should be a string, an element of U, or a list of gates."
RightCliffordSimilar::usage = "CliffordSimilarRight[U,V] Returns True if there is a Clifford gate C such that C.U.C^-1 = V."
LeftCliffordSimilar::usage = "CliffordSimilarLeft[U,V] Returns True if there is a Clifford gate C such that C.U.C^-1 = V."

GaussianQ::usage = "GaussianQ[U] is a Boolean function which checks if U corresponds to a Gaussian gate."

LDE::usage = "LDE[U] takes as input either a  $U(4)$  or  $S(6)$  Clifford + CS operator and returns a list of gates."
OptimalCSCCount::usage = "OptimalCSCCount[U] takes as input any representation of a Clifford + CS operator and returns the minimum number of CS gates required to implement U."

SyllableList::usage = "A list of 15 syllables of CS-count one which are not right-Clifford C which conjugates CS as C.CS.C^-1. The input U should be a string, an element of U, or a list of gates. The output is a list of syllables."

```

SyllableListAsymmetric::usage = "An alternative list of 15 syllables of CS-count (Each syllable is equivalent to C.CS for C a Clifford. For each syllable, we show the operator's 6x6 SO(6) representation, and its name according to the generalization of the Pauli matrices.";

NormIt::usage = "NormIt[U,options] takes as input a Gaussian Clifford + T operator. This normal form is output as a string of generators using the standard syllable notation. The options for the "OutputType" are "String" or "HexDec", the options for "UpToPhase" are the booleans True and False. When either of the following characters: "W", "S1", "S2", "H1", "H2", "CZ", "EX", "S1" is the gate S1, and so on.";

FrobeniusDistance::usage = "FrobeniusDistance[U,V] computes the distance between U and V.";

PauliRotation::usage = "PauliRotation[φ, ϵ , Pauli] finds a unitary Gaussian Clifford decomposition of the Pauli rotation $e^{i\varphi P}$ for the Pauli P. The Pauli can be one of the fifteen strings "XI", "YI", "ZI", "IX", "XZ", "YZ", or "ZZ".";

PauliRotationSequence::usage = "PauliRotationSequence[φ, ϵ , Pauli, options] finds a unitary Gaussian Clifford decomposition of the Pauli rotation $e^{i\varphi P}$ for the Pauli P. The Pauli can be one of the fifteen strings "XI", "YI", "ZI", "IX", "XZ", "YZ", or "ZZ". It then outputs a normalized sequence of Clifford gates. The options for the "OutputType" are "String" or "HexDec" and the options for the "SyllableType" are "Normal" or "Asymmetric".";

PauliDecomposition::usage = "PauliDecomp[U] finds a list of 15 angle parameters θ_j for the decomposition of the form $e^{i\sum_{j=1}^{15} \theta_j P_j}$ for the Paulis P_j is (ZI, XI, ZI, IZ, IX, IZ, XX, YY, ZZ, ...).";

ApproximateOp::usage = "Approximate[U, ϵ] finds an approximation within Frobenius distance ϵ of U. If U is an element of U(4), the result is a U(4) representation of a Clifford SO(6) representation of a Clifford + CS circuit. Note that the Frobenius distance is a representation.";

ApproximateSequence::usage = "ApproximateSequence[U, ϵ , options] finds a normalized sequence of Clifford gates (in the Unitary representation and up to an irrelevant phase) for the input U string unless otherwise specified in the options. The options for the "OutputType" are "String" or "HexDec" and the options for the "SyllableType" are "Normal" or "Asymmetric", and the options for "IfGate" are "String" or "HexDec".";

RandomCliffCS::usage = "RandomCliffCS[n] pseudorandomly samples from the uniform distribution over the Clifford SO(6) representations of a Clifford + CS circuit of size n.";

Possible Errors

```

General::invldopt = "Option `2` for function `1` received invalid value `3`";
U4ToS06::notunitary = "The argument must be a 4x4 unitary matrix.";
R::invalidpaulis = "The arguments `1` and `2` must be strings of the form \"AB\"
    Furthermore, the strings must correspond to commuting non-identity two-qubit I
FromSequence::notstring = "You have not entered a string.";
FromList::notagate = "The string `1` is not one of \"W\", \"S1\", \"S2\", \"H1\",
    You may have forgotten a space between gate names or used a name for a gate w
FromHexDec::invalidnumber = "The string `1` is not a valid hexadecimal representa
    seperating the syllables from the Clifford. Otherwise, ensure the Clifford ha
    before being written in hexadecimal representation, and that your syllables o
CliffordQ::notacircuit = "You have not entered a string of operators, a valid hexa
CliffordSynth::notaclifford = "Your input is not a Clifford operator in string fo
GaussianQ::notacircuit = "You have not entered a string of operators, a valid hexa
LDE::notamatrix = "You have not entered an element of SO(6) or U(4) which is a C
OptimalCSCCount::notacircuit = "You have not entered a string of operators, a vali
NormIt::badopt = "The option `1` is not valid for `2`.";
NormIt::notacircuit = "You have not entered a string of operators, a valid hexade
FrobeniusDistance::notequidimensionalmatrices = "Your inputs are not two matrices
CandidateFinder::notreals = "Your inputs are not two real numbers";
PauliRotation::notreals = "Your input does not include two real numbers";
PauliRotation::invldstring = "Your input does not include one string from the set
    \"YX\", \"ZX\", \"XY\", \"YY\", \"ZY\", \"XZ\", \"YZ\", or \"ZZ\".";
PauliDecomposition::notanoperator = "Your input is neither an element of U(4) or
ApproximateOp::notreal = "Your error tolerance is not a real number.";
ApproximateOp::notanoperator = "Your input is neither an element of U(4) or SO(6)
RandomCliffCS::notanatural = "Your input is not a natural number.";

```

```
Begin["`Private`"];
```

Function Definitions

Functions for option checking

These functions will be used to check options for functions which accept them. For each such function, we must define a `test[f,op]` function for a particular option type `op` of function `f`. Credit for this code snippet goes to Mr. Wizard in the Stack Overflow post <https://mathematica.stackexchange.com/questions/116623>.

```

optsMsg[f_][op_, val_] :=
  test[f, op][val] || Message[General::invldopt, f, op, val];

Attributes[optsCheck] = {HoldFirst};

optsCheck @ head_[___, opts : OptionsPattern[]] :=
  And @@ optsMsg[head] @@@ FilterRules[{opts}, Options @ head];

```

Constants and Single-Qubit operators

For internal use only.

```

 $\omega$  = (1+I)/Sqrt[2];
 $\zeta$  = Exp[I*Pi/8];
s = DiagonalMatrix[{1,I}];
h = 1/Sqrt[2]*{{1,1},{1,-1}};
id = PauliMatrix[0];
x = PauliMatrix[1];
y = PauliMatrix[2];
z = PauliMatrix[3];

```

Unitary Representations of Two-qubit Clifford + CS operators

These operators are exported to the user as 4x4 matrices in the standard Mathematica format.

Fixed Two-Qubit Gates

These are the traditional two-qubit Clifford + CS gates.

```

Id = IdentityMatrix[4];
W =  $\omega$ *Id;
S1 = KroneckerProduct[s,id];
S2 = KroneckerProduct[id,s];
H1 = KroneckerProduct[h,id];
H2 = KroneckerProduct[id,h];
CZ = DiagonalMatrix[{1,1,1,-1}];
CNOT12 = {
    {1,0,0,0},
    {0,1,0,0},
    {0,0,0,1},
    {0,0,1,0}
};
CNOT21 = {
    {1,0,0,0},
    {0,0,0,1},
    {0,0,1,0},
    {0,1,0,0}
};
EX = {
    {1,0,0,0},
    {0,0,1,0},
    {0,1,0,0},
    {0,0,0,1}
};
CS = DiagonalMatrix[{1,1,1,I}];
X1 = KroneckerProduct[x,id];
X2 = KroneckerProduct[id,x];
Z1 = KroneckerProduct[z,id];
Z2 = KroneckerProduct[id,z];

```

R[P,Q] Operators

This code provides the R[P,Q] gate constructor, as in [2].


```

NoPhaseHermitianPaulis = {"IX","IY","IZ","XI","XX","XY","XZ","YI","YX","YY","YZ",
HermitianPaulis = Join[NoPhaseHermitianPaulis,Map["-"<>#&,NoPhaseHermitianPaulis]

SinglePauliTranslator["I"] = id;
SinglePauliTranslator["X"] = x;
SinglePauliTranslator["Y"] = y;
SinglePauliTranslator["Z"] = z;

PauliTranslator[str_] := PauliTranslator[str] = Module[{ablist,a,b,phase},
  ablist = Characters @ StringTake[str,-2];
  {a,b} = Map[SinglePauliTranslator,ablist];
  phase = If[StringStartsQ[str,"-"],-1,1];
  phase * KroneckerProduct[a,b]
];

HermitianPauliQ[P_] := MemberQ[HermitianPaulis,P];
DistCommuteHermitePauliQ[P_?HermitianPauliQ,Q_?HermitianPauliQ] := Module[{p,q,cor
  p = PauliTranslator[P];
  q = PauliTranslator[Q];
  commutebool = p.q == q.p;
  distinctbool = Not[MemberQ[{Id,-Id},p.q]];
  commutebool && distinctbool
];
DistCommuteHermitePauliQ[P_,Q_] = False;

R[P_,Q_] /; DistCommuteHermitePauliQ[P,Q] := R[P,Q] = Module[{p,q},
  p = PauliTranslator[P];
  q = PauliTranslator[Q];
  Id + (I-1)/4*(Id - p).(Id - q)
];
R[_,_] := (
  Message[R::invalidpaulis];
  $Failed
);

```

Checks For U(4) and SO(6)

These Boolean functions determine whether an operator is an element of U(4) or SO(6), respectively.

```

U4Q[U_] := UnitaryMatrixQ[U] && (Dimensions[U] == {4,4});
SO6Q[O_] := OrthogonalMatrixQ[O] && (Dimensions[O] == {6,6}) && (Det[O] == 1);

```

The $SU(4) \cong SO(6)$ Isomorphism

These definitions and functions allow one to compute the $SO(6)$ representation of an element of $U(4)$ (up to a phase).

Rules for Inner Products of Wedge Products

Defined using the unassigned Mathematica symbols of \wedge , \langle , and \rangle .

```

<a_,0_,b_*c_>:=b*<a,0,c>;
<a_*b_,0_,c_>:=Conjugate[a]*<b,0,c>;
<a_+b_,0_,c_>:=<a,0,c>+<b,0,c>;
<a_,0_,b_+c_>:=<a,0,b>+<a,0,c>;
<x_∧y_,0_,u_∧v_>:= (Conjugate[x].0.u)*(Conjugate[y].0.v) - (Conjugate[x].0.v)*(C

```

Orthonormal Basis for Subscript[\mathbb{C} , 6]

This basis is such that computing the above inner products for an element of $U(4, \mathbb{C})$ will produce an element of $SO(6, \mathbb{R})$. Moreover, the representations for Clifford + CS operators are easy to work with in this basis.

```

Basis6 = {
  (I/Sqrt[2])* (UnitVector[4,1] ^UnitVector[4,2] - UnitVector[4,3] ^UnitVector[4,4]
  (1/Sqrt[2])* (UnitVector[4,1] ^UnitVector[4,2] + UnitVector[4,3] ^UnitVector[4,4]
  (I/Sqrt[2])* (UnitVector[4,2] ^UnitVector[4,3] - UnitVector[4,1] ^UnitVector[4,4]
  (1/Sqrt[2])* (UnitVector[4,2] ^UnitVector[4,4] + UnitVector[4,3] ^UnitVector[4,1]
  (I/Sqrt[2])* (UnitVector[4,2] ^UnitVector[4,4] - UnitVector[4,3] ^UnitVector[4,1]
  (1/Sqrt[2])* (UnitVector[4,2] ^UnitVector[4,3] + UnitVector[4,1] ^UnitVector[4,4]
};

```

Calculating the $SO(6)$ Representation for an element of $SU(4)$ (and from $U(4)$ up to a phase)

Functions for mapping elements of $SU(4)$ to $SO(6)$ and $U(4)$ to $SO(6)$ (by converting that element of $U(4)$ to an element of $SU(4)$).

```

SU4ToS06[U_] := Table[Simplify[<Basis6[[i]],U,Basis6[[j]]>],{i,1,6},{j,1,6}];

U4ToS06[U_;/;U4Q[U]] := SU4ToS06[1/(Det[U])^(1/4)*U];
U4ToS06[U_] := (
  Message[U4ToS06::notunitary];
  $Failed
);

```

SO(6) Representations of Two-qubit Clifford + CS operators

Calculated using our transformations. These operators are exported to the user as 6x6 matrices in the standard Mathematica format. Note that we have to multiply by overall phases to ensure that the transformation uses an element of SU(4).

```
IdS06 = SU4ToS06[Id];
IS06 = SU4ToS06[W.W];
H1S06 = SU4ToS06[H1];
H2S06 = SU4ToS06[H2];
S1S06 = SU4ToS06[ $\omega^{(-1)}$ *S1];
S2S06 = SU4ToS06[ $\omega^{(-1)}$ *S2];
CZS06 = SU4ToS06[ $\omega^{(-1)}$ *CZ];
CNOT12S06 = SU4ToS06[ $\omega^{(-1)}$ *CNOT12];
CNOT21S06 = SU4ToS06[ $\omega^{(-1)}$ *CNOT21];
EXS06 = SU4ToS06[ $\omega^{(-3)}$ *EX];
CSS06 = SU4ToS06[ $\xi^{(-1)}$ *CS];
X1S06 = SU4ToS06[X1];
X2S06 = SU4ToS06[X2];
Z1S06 = SU4ToS06[Z1];
Z2S06 = SU4ToS06[Z2];
RS06[P_,Q_] := RS06[P,Q] = SU4ToS06[ $\xi^{(-1)}$ *R[P,Q]];
```

Custom Representations of SO(6) Clifford + CS operators

Our synthesis algorithms will use a custom data type for the SO(6) representation of a Clifford + CS operator. The basic data structure of this special representation is as follows:

$$\{k,M\} := 2^{(-k/2)} \cdot M$$

This allows easy tracking of the lde. They are packed in SparseArrays to help make things even a little faster, as every Clifford is just a permutation matrix. These representations are for internal use only.

Switching between the standard representation for a 6x6 matrix and a representation specifically for Clifford + CS operators.

Functions for converting to and from the special representation.

```

S06ToSpecialRep[M_] := Module[{LDE, IntegerMat},
  LDE = FullSimplify[Max[Map[Simplify[Log[Sqrt[2], Denominator[#]]] &, Flatten[Sim
  IntegerMat = FullSimplify[(Sqrt[2]^LDE)*M];
  {LDE, SparseArray[IntegerMat]}
];
SpecialRepToS06[{k_, M_}] := Simplify[1/Sqrt[2]^k*Normal[M]];

```

Special Representations for SO(6) Clifford + CS operators

```

IdSp = S06ToSpecialRep[IdS06];
ISp = S06ToSpecialRep[IS06];
H1Sp = S06ToSpecialRep[H1S06];
H2Sp = S06ToSpecialRep[H2S06];
S1Sp = S06ToSpecialRep[S1S06];
S2Sp = S06ToSpecialRep[S2S06];
CZSp = S06ToSpecialRep[CZS06];
CNOT12Sp = S06ToSpecialRep[CNOT12S06];
CNOT21Sp = S06ToSpecialRep[CNOT21S06];
EXSp = S06ToSpecialRep[EXS06];
CSSp = S06ToSpecialRep[CSS06];
X1Sp = S06ToSpecialRep[X1S06];
X2Sp = S06ToSpecialRep[X2S06];
Z1Sp = S06ToSpecialRep[Z1S06];
Z2Sp = S06ToSpecialRep[Z2S06];
RSp[P_, Q_] := RSp[P, Q] = S06ToSpecialRep[RS06[P, Q]];

```

Basic Matrix Operations for the Special Representation

These internal functions are used to reduce the denominator exponent to the lde, multiply operators in the special representation, and invert the special representation.

```

KReduceOnce[{k_, a_}] := If[AllTrue[a, EvenQ, 2] && k > 1, {k - 2, a/2}, {k, a}];
KReduce[o_] := FixedPoint[KReduceOnce, o];

Dot2Sp[{k1_, a1_}, {k2_, a2_}] := KReduce[{k1 + k2, a1.a2}];
DotSp[x_] := Fold[Dot2Sp, IdSp, {x}];

InvSp[{k_, a_}] := {k, Transpose[a]};

```

Functions for Residues Modulo 2 and Finding Paired Matrix Rows

These functions are used for finding paired rows in operators in the special representation.

```

PatternMats[{k_,a_}] := Mod[a,2];

MatrixRowPairs[list_] := GatherBy[
  Range@Length[Normal[list]],
  Normal[list][[#]]&
];

RowPairs[x_] := Sort@MatrixRowPairs@PatternMats@x

```

String Reading

Functions for reading in a string of operators. The string is always read in as an element of $U(4)$.

```

FromSequence[str_String] := Module[{strlist},
  strlist = StringSplit[str];
  FromList[strlist]
];
FromSequence[x_] := (
  Message[FromSequence::notstring];
  $Failed
);

FromList[{}] = Id;
FromList[{"W"}] = W;
FromList[{"S1"}] = S1;
FromList[{"S2"}] = S2;
FromList[{"H1"}] = H1;
FromList[{"H2"}] = H2;
FromList[{"CZ"}] = CZ;
FromList[{"CS"}] = CS;
FromList[{"EX"}] = EX;
FromList[{"X1"}] = X1;
FromList[{"X2"}] = X2;
FromList[{"Z1"}] = Z1;
FromList[{"Z2"}] = Z2;
FromList[{str_/(StringMatchQ[str,"R[*,*]" && StringCount[str,","] == 1)}] := Module[
  ops = StringSplit[str,{"R[","",""]"}];
  R @@ ops
];
FromList[{str_}] := (
  Message[FromList::notagate,str];
  $Failed
);
FromList[{h_,t_}] := FromList[{h}] . FromList[{t}];

```

Hexadecimal Representations

We shall use strings of signed hexadecimal integers to represent Clifford + CS operators. This form is much more compact than, for example, the string form of an operator. The string must be of the following form:

`\"(| -)(1-9 | a-f)* 00000 C\"`

where C is the hexadecimal representation of an integer from 1 to 92160 and the first option is simply an optional "-" character.

```

ValidHexDecQ[num_String] := Module[{separator, typenum, syllabletype, syllablescliff},
  separator = StringCount[num,"00000"] == 1;

```

```

typenum = StringCount[num,"-"];
syllabletype = (typenum == 0) || ((typenum == 1) && StringStartsQ[num,"-"]);
syllablescliffordok = If[
  separator && syllabletype,
    split = StringSplit[num,{"00000","-"}];
    {syllables,clifford} = If[Length[split] > 1,split,{"",split[[1]]}];
    syllablescharacterlist = Union[Characters[syllables]];
    syllablesok = SubsetQ[{"1","2","3","4","5","6","7","8","9","a","b","c"}];
    cliffordcharacterlist = Union[Characters[clifford]];
    cliffordpossible = (StringLength[clifford]) <= 5 && SubsetQ[{"0","1",
    cliffordok = If[
      cliffordpossible,
        1 <= FromDigits[clifford,16] <= 92160,
        False
      ];
    syllablesok && cliffordok,
    False
  ];
syllablescliffordok
];
ValidHexDecQ[U_] := False

PhaseSet = {Id,MatrixPower[W,4],W,MatrixPower[W,5]};
L1Set = {Id,H1,S1.H1.MatrixPower[W,7]};
L2Set = {Id,H2,S2.H2.MatrixPower[W,7]};
CZSet = {Id,CZ.MatrixPower[W,7]};
S1Set = {Id,S1.MatrixPower[W,7]};
P1Set = {Id,X1,X1.Z1,Z1};
S2Set = {Id,S2.MatrixPower[W,7]};
P2Set = {Id,X2,X2.Z2,Z2};
SWAPSet = {Id,EX.MatrixPower[W,5]};
ISet = {Id,MatrixPower[W,2]};

CliffordFromNumber[cliffnum_] := Module[{digits,cz,l1,l2,index},
  digits = PadLeft[IntegerDigits[cliffnum-1,MixedRadix[{4,10,3,3,2,4,2,4,2,2}]]
  cz = Unitize[digits[[2]]];
  l2 = Mod[digits[[2]]-cz,3];
  l1 = (digits[[2]]-l2-cz)/3;
  index = Join[{digits[[1]],l1,l2,cz},digits[[3;;-1]]] + 1;
  PhaseSet[[index[[1]]] . L1Set[[index[[2]]] . L2Set[[index[[3]]] . CZSet[[i
  L1Set[[index[[5]]] . L2Set[[index[[6]]] . S1Set[[index[[7]]] . P1Set[[
  S2Set[[index[[9]]] . P2Set[[index[[10]]] . SWAPSet[[index[[11]]] . ISe
];

```

```

FromHexDec[str_String/;ValidHexDecQ[str]] := Module[{syllablelist,split,syllables,
  syllablelist = If[StringStartsQ[str,"-"],SyllableListAsymmetric[{All,1}],SyllableListAsymmetric[{All,1}]];
  split = StringSplit[str,{"00000","-"}];
  {syllablestr,cliffstr} = If[Length[split] > 1,split,{"",split[[1]]}];
  cliffnum = FromDigits[cliffstr,16];
  cliff = CliffordFromNumber[cliffnum];
  syllables = Map[syllablelist[[FromDigits[#,16]]]&,Characters[syllablestr]];
  Simplify[Dot @@ (Append[syllables,cliff])]
];
FromHexDec[str_] := (
  Message[FromHexDec::invalidnumber,str];
  $Failed
);

```

The Two-Qubit Clifford Group

Here we develop some basic functions for the two-qubit Clifford group.

Identifying if an operator is a Clifford and if two operators are Clifford-similar

Boolean functions for checking if an operator is a Clifford or if two operators are Clifford-similar.

```

CliffordQ[U_;/;U4Q[U]] := (Det[U]^2 == 1) && AllTrue[Flatten[U4ToS06[U]],IntegerQ];
CliffordQ[U_;/;S06Q[U]] := AllTrue[Flatten[U],IntegerQ];
CliffordQ[U_String/;ValidHexDecQ[U]] := CliffordQ[FromHexDec[U]];
CliffordQ[U_String] := CliffordQ[FromSequence[U]];
CliffordQ[U_] := (
  Message[CliffordQ::notacircuit];
  $Failed
);

RightCliffordSimilar[U_String/;ValidHexDecQ[U],V_String/;ValidHexDecQ[V]] := RightCliffordSimilar[U_String,V_String] := RightCliffordSimilar[FromSequence[U],FromSequence[V]];
RightCliffordSimilar[U_,V_] := CliffordQ[ConjugateTranspose[U].V];

LeftCliffordSimilar[U_String/;ValidHexDecQ[U],V_String/;ValidHexDecQ[V]] := LeftCliffordSimilar[U_String,V_String] := LeftCliffordSimilar[FromSequence[U],FromSequence[V]];
LeftCliffordSimilar[U_,V_] := CliffordQ[V.ConjugateTranspose[U]];

```

Synthesis of Clifford Circuits (with at most 1 CZ gate and 1 SWAP gate)

Rather than carry around an explicit lookup table with 92160 elements in the $U(4)$ case and 23040 elements in the $SO(6)$ case, we instead will use a synthesis algorithm based on the special representa-

tion; this takes advantage of the sparsity of Cliffords in this representation. We implicitly define our regular Clifford synthesis algorithm, CliffordSynth, in terms of this other algorithm, to be defined below.

```

PhaseFixU4[{str_,power_,SUcliffnum_},U_] := Module[{unitary,phase,Wpower,diff,Wco:
  unitary = FromSequence[str];
  phase = Simplify[(U. ConjugateTranspose[unitary])[[1,1]]];
  Wpower = Mod[Log[(1+I)/Sqrt[2],phase],8];
  diff = Mod[Wpower - power,8];
  Wcosetnumber = Which[
    (diff == 0) || (diff == 2),0,
    (diff == 4) || (diff == 6),1,
    (diff == 1) || (diff == 3),2,
    True,3
  ];
  {IntegerString[23040*Wcosetnumber + FromDigits[SUcliffnum,16],16],str<>StringI
];
PhaseFixS06[{str_,power_,SUcliffnum_}] := {SUcliffnum,str<>StringRepeat[" W",power

CliffordSynth[U_String;/;ValidHexDecQ[U]] := Module[{unitary},
  unitary = FromSequence[U];
  CliffordSynth[unitary]
];
CliffordSynth[U_String] := Module[{unitary},
  unitary = FromSequence[U];
  CliffordSynth[unitary]
];
CliffordSynth[U_/(U4Q[U] && CliffordQ[U])] := Module[{orthogonal,synth},
  orthogonal = U4ToS06[U];
  synth = CliffordSynthSp[{0,SparseArray[orthogonal]}];
  PhaseFixU4[synth,U]
];
CliffordSynth[U_/(CliffordQ[U])] := PhaseFixS06[CliffordSynthSp[{0,SparseArray[U]}]
CliffordSynth[U_] := (
  Message[CliffordSynth::notaclifford];
  $Failed
);

```

Clifford Group in the Special representation

Explicit CliffordSynthSp algorithm. We perform the synthesis by decomposing an element of the signed permutation group in terms of a generating set which is constructed from basic Clifford operators in the SO(6) representation.

```

L1Cosets = {
  {IdSp,"",0,0},

```

```

    {H1Sp,"H1 ",0,1},
    {DotSp[S1Sp,H1Sp],"S1 H1 ",7,2}
};
L2Cosets = {
    {IdSp,"",0,0},
    {H2Sp,"H2 ",0,1},
    {DotSp[S2Sp,H2Sp],"S2 H2 ",7,2}
};
CZSets = {
    {IdSp,"",0,0},
    {CZSp,"CZ ",7,1}
};
S1Sets = {
    {IdSp,"",0,0},
    {S1Sp,"S1 ",7,1},
};
Pauli1Sets = {
    {IdSp,"",0,0},
    {X1Sp,"X1 ",0,1},
    {DotSp[X1Sp,Z1Sp],"X1 Z1 ",0,2},
    {Z1Sp,"Z1 ",0,3}
};
S2Sets = {
    {IdSp,"",0,0},
    {S2Sp,"S2 ",7,1},
};
Pauli2Sets = {
    {IdSp,"",0,0},
    {X2Sp,"X2 ",0,1},
    {DotSp[X2Sp,Z2Sp],"X2 Z2 ",0,2},
    {Z2Sp,"Z2 ",0,3}
};
SwapSets = {
    {IdSp,"",0,0},
    {EXSp,"EX ",5,1}
};
ISets = {
    {IdSp,"",0,0},
    {ISp,"",2,1}
};

CliffordQSp[U_] := Module[{k,M},
    {k,M} = KReduce[U];

```

```

(k == 0) && OrthogonalMatrixQ[M] && (Det[M] == 1)
];

CliffordSynthSp[{_,M_}] := Module[
{
    CliffordList,SwapTest,MUnSwapped,UpperLeft,LowerRight,RemovedLeftCosets,C
    UpperLeftCol,LowerRightCol,RemovedLeftCosetsNew,S1Test,S2Test,RemovedSSet
    RemovedPauli1Sets,d3new,d4,d5,d6,RemovedPauli2Sets,ITest,str,phase,list,f
},

CliffordList = ConstantArray[0,11];

SwapTest = Total[Abs[M[[1;;3,1;;3]]],2];
CliffordList[[10]] = If[SwapTest > 1,SwapSets[[1]],SwapSets[[2]]];
MUnSwapped = M.Transpose[CliffordList[[10,1,2]]];

UpperLeft = Total[Abs[MUnSwapped[[1;;3,1;;3]]],{2}];
LowerRight = Total[Abs[MUnSwapped[[4;;6,4;;6]]],{2}];
CliffordList[[1]] = Which[
    UpperLeft == {0,1,1},L1Cosets[[2]],
    UpperLeft == {1,0,1},L1Cosets[[3]],
    True,L1Cosets[[1]]
];
CliffordList[[2]] = Which[
    LowerRight == {0,1,1},L2Cosets[[2]],
    LowerRight == {1,0,1},L2Cosets[[3]],
    True,L2Cosets[[1]]
];
RemovedLeftCosets = Transpose[CliffordList[[1,1,2]].CliffordList[[2,1,2]]].MUn

CZRows = Total[Abs[RemovedLeftCosets[[3,1;;3]]],2];
CliffordList[[3]] = If[CZRows == 1,CZSets[[1]],CZSets[[2]]];
RemovedCZSet = Transpose[CliffordList[[3,1,2]].RemovedLeftCosets;

UpperLeftCol = Abs[RemovedCZSet[[1;;3,3]]];
LowerRightCol = Abs[RemovedCZSet[[4;;6,6]]];
CliffordList[[4]] = Which[
    UpperLeftCol == {0,0,1},L1Cosets[[1]],
    UpperLeftCol == {0,1,0},L1Cosets[[3]],
    True,L1Cosets[[2]]
];
CliffordList[[5]] = Which[
    LowerRightCol == {0,0,1},L2Cosets[[1]],

```

```

LowerRightCol == {0,1,0},L2Cosets[[3]],
True,L2Cosets[[2]]
];
RemovedLeftCosetsNew = Transpose[CliffordList[[4,1,2]].CliffordList[[5,1,2]]]

S1Test = Abs[RemovedLeftCosetsNew[[1,1]]];
S2Test = Abs[RemovedLeftCosetsNew[[4,4]]];
CliffordList[[6]] = If[S1Test == 1,S1Sets[[1]],S1Sets[[2]]];
CliffordList[[8]] = If[S2Test == 1,S2Sets[[1]],S2Sets[[2]]];
RemovedSSets = Transpose[CliffordList[[6,1,2]].CliffordList[[8,1,2]].Removed

{d1,d2,d3} = Diagonal[RemovedSSets[[1;;3,1;;3]]];
CliffordList[[7]] = Which[
  (d1 == d2) && (d2 == d3),Pauli1Sets[[1]],
  (d1 != d2) && (d2 == d3),Pauli1Sets[[2]],
  (d1 != d2) && (d2 != d3),Pauli1Sets[[3]],
  True,Pauli1Sets[[4]]
];
RemovedPauli1Sets = Transpose[CliffordList[[7,1,2]].RemovedSSets;

{d3new,d4,d5,d6} = Diagonal[RemovedPauli1Sets[[3;;6,3;;6]]];
CliffordList[[9]] = Which[
  (d4 == d5) && (d5 == d6) && (d3new*d4 == 1),Pauli2Sets[[1]],
  (d4 != d5) && (d5 == d6),Pauli2Sets[[2]],
  (d4 != d5) && (d5 != d6),Pauli2Sets[[3]],
  True,Pauli2Sets[[4]]
];
RemovedPauli2Sets = Transpose[CliffordList[[9,1,2]].RemovedPauli1Sets;

ITest = RemovedPauli2Sets[[1,1]];
CliffordList[[11]] = If[ITest == 1,ISets[[1]],ISets[[2]]];

{str,phase,list} = Fold[{
  #1[[1]]<>#2[[2]],
  Mod[#1[[2]]+#2[[3]],8],
  Append[#1[[3]],#2[[4]]]
}&,
{"",0,{}},
CliffordList
];
firstdigit = list[[3]] * (list[[1]]*3 + list[[2]] + 1);
cliffordnumber = FromDigits[Prepend[list[[4;;-1]],firstdigit],MixedRadix[{10,

```

```
{str,phase,IntegerString[cliffordnumber,16]}
];
```

The Two-Qubit Clifford + CS Group

Here we develop some basic constructors for Clifford + CS circuits. We also provide a function for checking if an operator corresponds to a Clifford + CS circuit.

Check If an Operator Is a Gaussian Clifford + T Matrix

We define a function for determining if an operator is a representation for a Gaussian Clifford + T matrix, i.e. a Clifford + CS operator.

```
DyadicQ[n_] := Module[{numerator,denominator},
  {numerator,denominator} = NumeratorDenominator[n];
  IntegerQ[numerator] && IntegerQ[Log2[denominator]]
];
GaussianDyadicQ[n_] := AllTrue[ReIm[n],DyadicQ];

GaussianQ[U_;/;U4Q[U]] := Module[{UGaussDyadicQ,UWGuassDyadicQ},
  UGaussDyadicQ = AllTrue[Flatten[U],GaussianDyadicQ];
  UWGuassDyadicQ = AllTrue[Flatten[W.U],GaussianDyadicQ];
  UGaussDyadicQ || UWGuassDyadicQ
];
GaussianQ[U_;/;S06Q[U]] := Module[{sp},
  sp = S06ToSpecialRep[U];
  (Det[U] == 1) && IntegerQ[sp[[1]]] && AllTrue[Flatten[sp[[2]]],IntegerQ]
];
GaussianQ[U_;/;ValidHexDecQ[U]] := True;
GaussianQ[U_String] := GaussianQ[FromSequence[U]];
GaussianQ[U_] := (Message[GaussianQ::notacircuit];
  $Failed
);
```

Least Denominator Exponents and CS-counts

These functions can be used to find denominator exponents and optimal-CS counts.

```

LDE[U_/(GaussianQ[U] && U4Q[U])] := Module[{reimlist},
  reimdenomlist = Denominator @ Flatten @ ReIm[U];
  FullSimplify[Max[Map[Simplify[Log[Sqrt[2],#]]&,reimdenomlist]]]
];
LDE[U_/(GaussianQ[U] && S06Q[U])] := Module[{sp},
  sp = S06ToSpecialRep[U];
  sp[[1]]
];
LDE[_] := (
  Message[LDE::notamatrix];
  $Failed
);

OptimalCSCCount[U_/(GaussianQ[U] && U4Q[U])] := LDE[U4ToS06[U]];
OptimalCSCCount[0_/(GaussianQ[0] && S06Q[0])] := LDE[0]
OptimalCSCCount[hexdec_String;ValidHexDecQ[hexdec]] := OptimalCSCCount[FromHexDec[hexdec]];
OptimalCSCCount[str_String] := OptimalCSCCount[FromSequence[str]];
OptimalCSCCount[U_] := (
  Message[OptimalCSCCount::notacircuit,U];
  $Failed
);

```

Syllable Lists

For the exported versions:

```

SyllableList = {
  {R["XI","IX"],RS06["XI","IX"],"R[XI,IX] "},
  {R["YI","IY"],RS06["YI","IY"],"R[YI,IY] "},
  {R["ZI","IZ"],RS06["ZI","IZ"],"R[ZI,IZ] "},
  {R["YI","IZ"],RS06["YI","IZ"],"R[YI,IZ] "},
  {R["ZI","IY"],RS06["ZI","IY"],"R[ZI,IY] "},
  {R["ZI","IX"],RS06["ZI","IX"],"R[ZI,IX] "},
  {R["XI","IZ"],RS06["XI","IZ"],"R[XI,IZ] "},
  {R["XI","IY"],RS06["XI","IY"],"R[XI,IY] "},
  {R["YI","IX"],RS06["YI","IX"],"R[YI,IX] "},
  {R["XX","YY"],RS06["XX","YY"],"R[XX,YY] "},
  {R["XX","ZY"],RS06["XX","ZY"],"R[XX,ZY] "},
  {R["ZX","YY"],RS06["ZX","YY"],"R[ZX,YY] "},
  {R["YX","XY"],RS06["YX","XY"],"R[YX,XY] "},
  {R["ZX","XY"],RS06["ZX","XY"],"R[ZX,XY] "},
  {R["YX","ZY"],RS06["YX","ZY"],"R[YX,ZY] "}
};

SyllableListAsymmetric = {
  {H1.H2.CS,H1S06.H2S06.CSS06,"H1 H2 CS "},
  {S1.H1.S2.H2.CS,S1S06.H1S06.S2S06.H2S06.CSS06,"S1 H1 S2 H2 CS "},
  {CS,CSS06,"CS "},
  {S1.H1.CS,S1S06.H1S06.CSS06,"S1 H1 CS "},
  {S2.H2.CS,S2S06.H2S06.CSS06,"S2 H2 CS "},
  {H2.CS,H2S06.CSS06,"H2 CS "},
  {H1.CS,H1S06.CSS06,"H1 CS "},
  {H1.S2.H2.CS,H1S06.S2S06.H2S06.CSS06,"H1 S2 H2 CS "},
  {S1.H1.H2.CS,S1S06.H1S06.H2S06.CSS06,"S1 H1 H2 CS "},
  {CN0T12.H1.CS,CN0T12S06.H1S06.CSS06,"H2 CZ H1 H2 CS "},
  {CZ.S1.H1.S2.H2.CS,CZS06.S1S06.H1S06.S2S06.H2S06.CSS06,"CZ S1 H1 S2 H2 CS "},
  {CZ.H1.H2.CS,CZS06.H1S06.H2S06.CSS06,"CZ H1 H2 CS "},
  {CN0T12.S1.H1.CS,CN0T12S06.S1S06.H1S06.CSS06,"H2 CZ S1 H1 H2 CS "},
  {CZ.S1.H1.H2.CS,CZS06.S1S06.H1S06.H2S06.CSS06,"CZ S1 H1 H2 CS "},
  {CZ.H1.S2.H2.CS,CZS06.H1S06.S2S06.H2S06.CSS06,"CZ H1 S2 H2 CS "}
};

```

And the internal data structure:

```

SyllableListSp = MapIndexed[
  {S06ToSpecialRep[#1[[2]]], #1[[3]], IntegerString[#2, 16]} &,
  SyllableList
];
SyllableListAsymmetricSp = MapIndexed[
  {S06ToSpecialRep[#1[[2]]], #1[[3]], IntegerString[#2, 16]} &,
  SyllableListAsymmetric
];

```

Normalization

In this section we develop a function for normalizing a circuit in any of our representations. This algorithm is described in detail in [2].

Earliest Generator Ordering and associated row pairings

We develop here an association which matches up each syllable to their row pairings under EGO.

```

TwoFourPaired = Map[
  {#, Complement[Range[1, 6], #]} &,
  Subsets[Range[1, 6], {4}]
];
TwoTwoTwoPaired = Flatten[Map[
  Table[{#[[1]], {#[[2, 1]], #[[2, k]]}}, Complement[{#[[2]], {#[[2, 1]], #[[2, k]]}}], {k, 1, 4}],
  Table[{1, j}, Complement[Range[1, 6], {1, j}]], {j, 2, 6}]]
];
PairList = Map[Sort, Join[TwoTwoTwoPaired, TwoFourPaired]];

MatchingPairings[list_] := {
  list,
  {list[[3]], Union[list[[1]], list[[2]]]},
  {list[[2]], Union[list[[1]], list[[3]]]},
  {list[[1]], Union[list[[2]], list[[3]]]}
};
PossibleSyllableListPairings = Map[
  MatchingPairings[RowPairs[First[#]]] &,
  SyllableListSp
];
SyllableListPairings = Map[
  # -> First[FirstPosition[PossibleSyllableListPairings, #]] &,
  PairList
];
PairingKey = Association @@ SyllableListPairings;

```


Finding a leftmost syllable and normalizing in the Special representation

We develop separate synthesis algorithms here based on whether we want to synthesize a circuit using the symmetric or asymmetric syllables.

```

LeftmostSyllable[U_] := SyllableListSp[PairingKey[RowPairs[U]]];
RemoveLeftmost[{k_,M_},str_,numberstr_] := Module[{leftmost,clifford},
  Which[
    k > 0,
      leftmost = LeftmostSyllable[{k,M}];
      {Dot2Sp[InvSp[leftmost[[1]]],{k,M}},str<>leftmost[[2]],numberstr<>lef
    k == 0,
      clifford = CliffordSynthSp[{k,M}];
      {{-1,SparseArray[ConstantArray[0,{6,6}]]},str<>clifford[[1]]<>StringR
    True,
      {{k,M},str,numberstr}
  ]
];

LeftmostSyllableAsymmetric[U_] := SyllableListAsymmetricSp[PairingKey[RowPairs[U]]];
RemoveLeftmostAsymmetric[{k_,M_},str_,numberstr_] := Module[{leftmost,clifford},
  Which[
    k > 0,
      leftmost = LeftmostSyllableAsymmetric[{k,M}];
      {Dot2Sp[InvSp[leftmost[[1]]],{k,M}},str<>leftmost[[2]],numberstr<>lef
    k == 0,
      clifford = CliffordSynthSp[{k,M}];
      {{-1,SparseArray[ConstantArray[0,{6,6}]]},str<>clifford[[1]]<>StringR
    True,
      {{k,M},str,numberstr}
  ]
];

test[NormItSp,"SyllableType"] := MemberQ[{"Normal","Asymmetric"},#]&;

Options[NormItSp] = {"SyllableType" -> "Normal"};
NormItSp[U_,OptionsPattern[]]?optsCheck := Module[{type},
  type = If[OptionValue["SyllableType"] == "Normal",RemoveLeftmost,RemoveLeftmo:
  FixedPoint[type,{U,"",If[OptionValue["SyllableType"] == "Normal","", "-"]}] [2
];

```

Normalizing from a string, a hexadecimal, U(4), or SO(6).

We provide the function NormIt for normalizing operators in any of our forms.

```

Options[FixPhase] = {"UpToPhase" -> True};
FixPhase[{str_, numberstr_}, U_, OptionsPattern[]] := Module[{syllablesplit, syllablenum, syllablepart, clifford, {cliffnum, cliffstr}, numberstrfixed, nophasestr, phasecount, strfixed},
  If[
    OptionValue["UpToPhase"] == True,
    {str, numberstr},
    syllablesplit = StringSplit[numberstr, "00000"];
    syllablenum = If[Length[syllablesplit] > 1, syllablesplit[[1]], ""];
    syllablepart = FromHexDec[syllablenum<>"000001"];
    clifford = Simplify[ConjugateTranspose[syllablepart] . U];
    {cliffnum, cliffstr} = CliffordSynth[clifford];
    numberstrfixed = syllablenum<>"00000"<>cliffnum;
    nophasestr = StringDelete[str, "W "];
    phasecount = StringCount[cliffstr, "W "];
    strfixed = nophasestr<>StringRepeat["W ", phasecount];
    {strfixed, numberstrfixed}
  ]
];

test[NormIt, "SyllableType"] := MemberQ[{"Normal", "Asymmetric"}, #]&;
test[NormIt, "OutputType"] := MemberQ[{"String", "HexDec"}, #]&;
test[NormIt, "UpToPhase"] := BooleanQ;

Options[NormIt] = {"OutputType" -> "String", "SyllableType" -> "Normal", "UpToPhase" -> True};
NormIt[U_ /; (GaussianQ[U] && U4Q[U]), OptionsPattern[]]?optsCheck := Module[{index},
  index = If[OptionValue["OutputType"] == "String", 1, 2];
  FixPhase[NormItSp[S06ToSpecialRep[U4ToS06[U]], "SyllableType" -> OptionValue["SyllableType"]], U, {index, "UpToPhase"}];
];
NormIt[U_ /; (GaussianQ[U] && S06Q[U]), OptionsPattern[]]?optsCheck := Module[{index},
  index = If[OptionValue["OutputType"] == "String", 1, 2];
  NormItSp[S06ToSpecialRep[U], "SyllableType" -> OptionValue["SyllableType"]][[index]];
];
NormIt[hexdec_String /; ValidHexDecQ[hexdec], opts:OptionsPattern[]] := NormIt[FromHexDec[hexdec], opts];
NormIt[str_String, opts:OptionsPattern[]] := NormIt[FromSequence[str], opts];
NormIt[U_] := (
  Message[NormIt::notacircuit, U];
  $Failed
);

```

ϵ -Approximating Pauli Rotations

This section describes algorithms used for finding approximations to Pauli Rotations

Frobenius Distance of Matrices

Computes the Frobenius distance between two matrices A and B (i.e. the Frobenius norm of the difference between A and B)

```
FrobeniusDistance[A_;/;MatrixQ[A],B_;/;MatrixQ[B]]/;(Dimensions[A]==Dimensions[B])
FrobeniusDistance[_,_]:= (
  Message[FrobeniusDistance::notequidimensionalmatrices];
  $Failed
);
```

Continued fractions and affine transformation

We describe a scheme for finding an appropriate affine transformation based off of continued fractions in the rationals.

```
NextIterate[{aN_,rN_,pN_,qN_,pNm1_,qNm1_}] := Module[{aNp1},
  aNp1 = IntegerPart[1/rN];
  {aNp1,1/rN - aNp1,aNp1*pN + pNm1,aNp1*qN + qNm1,pN,qN}
];
ContinuedFractionToPrecision[α_,tol_] := Module[{a0,iterate},
  a0 = IntegerPart[α];
  iterate = NestWhile[NextIterate,{a0,N[α - a0,2*Log10[1/tol]],a0,1,1,0},#[[4]]
  If[iterate[[2]] == 0 && iterate[[4]] <= 1/tol,
    iterate[[3;;4]],
    iterate[[5;;6]]
  ]
];
AffineMatrix[φ_,ε_] := Module[{α,q,r,s,t},
  α = Tan[φ/2];
  {r,q} = ContinuedFractionToPrecision[α,Sqrt[ε/(8+ε*α)]]*Sec[φ/2];
  {t,s} = ExtendedGCD[q,-r][[2]];
  {{q,r},{s,t}}
];
```

Finding angle candidates for bounded and unbounded angles

For calculating a candidate solution, we first map every angle into the interval $[0,\pi/2]$. We then use our affine transformation to find solutions to the integer programming problem.

```
RealQ[x_] := Element[x,Reals];
CandidateFinder[φ_?RealQ,ε_?RealQ] := Module[{modφ},
  modφ=Mod[φ,4*Pi];
  Which[
```

```

0 <= mod  $\varphi$  <= Pi/2,
  CandidateFinderBounded[mod  $\varphi$ ,  $\epsilon$ ],
Pi/2 < mod  $\varphi$  <= Pi,
  MapAt[Reverse, CandidateFinderBounded[Pi - mod  $\varphi$ ,  $\epsilon$ ], 2],
Pi < mod  $\varphi$  <= 2*Pi,
  MapAt[{-#[[2]], #[[1]]}&, CandidateFinder[mod  $\varphi$  - Pi,  $\epsilon$ ], 2],
True,
  MapAt[{-#[[1]], -#[[2]]}&, CandidateFinder[mod  $\varphi$  - 2*Pi,  $\epsilon$ ], 2]
];
CandidateFinder[_ , _] := (
  Message[CandidateFinder::notreals];
  $Failed
);

CandidateFinderBounded[ $\varphi$ _,  $\epsilon$ _] := Module[{root2k, c, s,  $\alpha$ , p,  $\Delta$ ,  $\delta$ , A, invA, scaledp, scaled $\Delta$ , scaled $\delta$ ,  $\theta_{\Delta}$ },
  root2k = 1;
  c = Cos[ $\varphi/2$ ];
  s = Sin[ $\varphi/2$ ];
   $\alpha$  = 1 -  $\epsilon^2/8$ ;
  p = {c* $\alpha$  + s*Sqrt[1 -  $\alpha^2$ ], s* $\alpha$  - c*Sqrt[1 -  $\alpha^2$ ]};
   $\Delta$  = 2*Sqrt[1 -  $\alpha^2$ ]*{-s, c};
   $\delta$  =  $\epsilon^2/8$ *{c, s};
  A = AffineMatrix[ $\varphi$ ,  $\epsilon$ ];
  invA = Inverse[A];
  scaledp = A.p;
  scaled $\Delta$  = A. $\Delta$ ;
  scaled $\delta$  = A. $\delta$ ;
   $\theta_{\Delta}$  = VectorAngle[scaled $\Delta$ , {1, 0}];
  Which[
     $\theta_{\Delta}$  > Pi/2,
    m1 = scaled $\Delta$ [[2]] / scaled $\Delta$ [[1]];
    m2 = scaled $\delta$ [[2]] / scaled $\delta$ [[1]];
    valid = Catch[Do[
      {x1, x2, x3, x4} = {
        Ceiling[scaledp[[1]] + scaled $\Delta$ [[1]]],
        Floor[scaledp[[1]]],
        Floor[scaledp[[1]] + scaled $\Delta$ [[1]] + scaled $\delta$ [[1]]],
        Floor[scaledp[[1]] + scaled $\delta$ [[1]]]
      };
    Do[
      y0 = If[x <= x2,

```

```

        Ceiling[m1*(x - scaledp[[1]]) + scaledp[[2]]],
        Ceiling[m2*(x - scaledp[[1]]) + scaledp[[2]]]
    ];
    y1 = If[x <= x3,
        Floor[m2*(x - scaledp[[1]]-scaledΔ[[1]]) + scaledp[[2]]
        Floor[m1*(x - scaledp[[1]]-scaledδ[[1]]) + scaledp[[2]]
    ];
    Do[
        transformed = invA.{x,y};
        If[Norm[transformed] <= Sqrt[2]^k,Throw[{k,transformed}]]
        {y,y0,y1}
    ];,
    {x,x1,x4}
];
{scaledp,scaledΔ,scaledδ} *= Sqrt[2];,
{k,0,Ceiling[4*Log2[1/ε]+6]}
];];,
θΔ' == Pi/2,
m1 = scaledδ[[2]] / scaledδ[[1]];
valid = Catch[Do[
    {x1,x2} = {
        Ceiling[scaledp[[1]]],
        Floor[scaledp[[1]]+scaledδ[[1]]]
    };
    Do[
        y0 = Ceiling[m1*(x - scaledp[[1]]) + scaledp[[2]]];
        y1 = Floor[m1*(x - scaledp[[1]]-scaledΔ[[1]]) + scaledp[[2]]
        Do[
            transformed = invA.{x,y};
            If[Norm[transformed] <= Sqrt[2]^k,Throw[{k,transformed}]]
            {y,Floor[(y0+y1)/2],y1}
        ];,
        {x,x1,x2}
    ];
    {scaledp,scaledΔ,scaledδ} *= Sqrt[2];,
    {k,0,Ceiling[4*Log2[1/ε]+6]}
];];,
True,
m1 = scaledδ[[2]] / scaledδ[[1]];
m2 = scaledΔ[[2]] / scaledΔ[[1]];
valid = Catch[Do[
    {x1,x2,x3,x4} = {
        Ceiling[scaledp[[1]]],

```

```

Floor[scaledp[[1]] + scaledδ[[1]]],
Floor[scaledp[[1]] + scaledΔ[[1]]],
Floor[scaledp[[1]] + scaledδ[[1]] + scaledΔ[[1]]]
};
Do[
  y0 = If[x <= x2,
    Ceiling[m1*(x - scaledp[[1]]) + scaledp[[2]]],
    Ceiling[m2*(x - scaledp[[1]] - scaledδ[[1]]) + scaledp[[2]]];
  y1 = If[x <= x3,
    Floor[m2*(x - scaledp[[1]]) + scaledp[[2]]],
    Floor[m1*(x - scaledp[[1]] - scaledΔ[[1]]) + scaledp[[2]]];
  Do[
    transformed = invA.{x,y};
    If[Norm[transformed] <= Sqrt[2]^k, Throw[{k,transformed}]]
    {y,y0,y1}
  ];,
  {x,x1,x4}
];
{scaledp,scaledΔ,scaledδ} *= Sqrt[2];,
{k,0,Ceiling[4*Log2[1/ε]+6]}
];];
];
valid
];

```

Solving Lagrange Four-Squares

Rather than implement our own solver based off of well known algorithms, we instead use a basic Mathematica function as the inputs for this problem never get too big.

```
Lagrange4[n_] := Module[{x1,x2,x3,x4}, {x1,x2,x3,x4} /. FindInstance[x1^2+x2^2+x3^2+x4^2 == n, {x1,x2,x3,x4}, Integers]
```

SU(4) Approximations Using a Candidate Solution

We provide an algorithm for finding a rotation by angle φ up to error ϵ for any of the 15 Pauli matrices.

```

SU4Z1Finder[φ_,ε_] := Module[{k,x,y,a,b,c,d,Invroot2k,α,β,χ},
  {k,{x,y}} = CandidateFinder[φ,ε];
  {a,b,c,d} = Lagrange4[2^k-x^2-y^2];
  Invroot2k = 1/Sqrt[2]^k;
  α = Invroot2k*(x+I*y);
  β = Invroot2k*(a+I*b);
  χ = Invroot2k*(c+I*d);

```

```

 $\chi = \text{Invroot2k}*(c + I*d);$ 
{
  { $\alpha, 0, -\text{Conjugate}[\beta], -\text{Conjugate}[\chi]$ },
  {0,  $\alpha, \chi, -\beta$ },
  { $\beta, -\text{Conjugate}[\chi], \text{Conjugate}[\alpha], 0$ },
  { $\chi, \text{Conjugate}[\beta], 0, \text{Conjugate}[\alpha]$ }
}
];

PauliRotation[ $\varphi_/, \text{Not}[\text{RealQ}[\varphi]], \epsilon_/, \text{Not}[\text{RealQ}[\epsilon]], \_ ] := (
  Message[PauliRotation::notreals];
  $Failed
);

PauliRotation[ $\varphi_/, \epsilon_/, \text{"ZI"}$ ] := SU4Z1Finder[ $\varphi, \epsilon$ ];
PauliRotation[ $\varphi_/, \epsilon_/, \text{"XI"}$ ] := H1 . SU4Z1Finder[ $\varphi, \epsilon$ ] . H1;
PauliRotation[ $\varphi_/, \epsilon_/, \text{"YI"}$ ] := S1 . H1 . SU4Z1Finder[ $\varphi, \epsilon$ ] . H1 . ConjugateTranspose
PauliRotation[ $\varphi_/, \epsilon_/, \text{"IZ"}$ ] := EX . SU4Z1Finder[ $\varphi, \epsilon$ ] . EX;
PauliRotation[ $\varphi_/, \epsilon_/, \text{"IX"}$ ] := EX . H1 . SU4Z1Finder[ $\varphi, \epsilon$ ] . H1 . EX;
PauliRotation[ $\varphi_/, \epsilon_/, \text{"IY"}$ ] := EX . S1 . H1 . SU4Z1Finder[ $\varphi, \epsilon$ ] . H1 . ConjugateTrans
PauliRotation[ $\varphi_/, \epsilon_/, \text{"ZZ"}$ ] := CNOT21 . SU4Z1Finder[ $\varphi, \epsilon$ ] . CNOT21;
PauliRotation[ $\varphi_/, \epsilon_/, \text{"XZ"}$ ] := H1 . CNOT21 . SU4Z1Finder[ $\varphi, \epsilon$ ] . CNOT21 . H1;
PauliRotation[ $\varphi_/, \epsilon_/, \text{"YZ"}$ ] := S1 . H1 . CNOT21 . SU4Z1Finder[ $\varphi, \epsilon$ ] . CNOT21 . H1 . (
PauliRotation[ $\varphi_/, \epsilon_/, \text{"ZX"}$ ] := H2 . CNOT21 . SU4Z1Finder[ $\varphi, \epsilon$ ] . CNOT21 . H2;
PauliRotation[ $\varphi_/, \epsilon_/, \text{"XX"}$ ] := H1 . H2 . CNOT21 . SU4Z1Finder[ $\varphi, \epsilon$ ] . CNOT21 . H2 . I
PauliRotation[ $\varphi_/, \epsilon_/, \text{"YX"}$ ] := S1 . H1 . H2 . CNOT21 . SU4Z1Finder[ $\varphi, \epsilon$ ] . CNOT21 . I
PauliRotation[ $\varphi_/, \epsilon_/, \text{"ZY"}$ ] := S2 . H2 . CNOT21 . SU4Z1Finder[ $\varphi, \epsilon$ ] . CNOT21 . H2 . (
PauliRotation[ $\varphi_/, \epsilon_/, \text{"XY"}$ ] := H1 . S2 . H2 . CNOT21 . SU4Z1Finder[ $\varphi, \epsilon$ ] . CNOT21 . I
PauliRotation[ $\varphi_/, \epsilon_/, \text{"YY"}$ ] := S1 . H1 . S2 . H2 . CNOT21 . SU4Z1Finder[ $\varphi, \epsilon$ ] . CNOT
PauliRotation[_, _, _] := (
  Message[PauliRotation::invldstring];
  $Failed
);

test[PauliRotationSequence, "SyllableType"] := MemberQ[{"Normal", "Asymmetric"}, #] &
test[PauliRotationSequence, "OutputType"] := MemberQ[{"String", "HexDec"}, #] &;

Options[PauliRotationSequence] = {"OutputType" -> "String", "SyllableType" -> "Nor

PauliRotationSequence[ $\varphi_/, \epsilon_/, \text{pauli}_/, \text{opts:OptionsPattern[]} ]? \text{optsCheck} := \text{NormIt}[\text{Pau}$$ 
```

General Unitary Decomposition

Here we develop an algorithm for the approximation of any $U(4)$ matrix using the Clifford + CS gate set.

Angles of rotation in the Pauli decomposition

First, we develop a method for finding the Pauli decomposition of an operator in terms of its 15 Pauli-rotation angles.

```
SafeArcTan[a_,b_] := If[(a == 0) && (b == 0),0,ArcTan[a,b]];
AngleFind[{{a_,_},{b_,_}}] := SafeArcTan[a,b];
SmallPauliDecomposition[o_] := Module[{v1,v2,v3, $\varphi$ 1, $\varphi$ 2, $\varphi$ 3,M1,M2,o',o''},
  v1 = o[[1;;2,3]];
   $\varphi$ 1 = -SafeArcTan @@ Reverse[v1];
  M1 = {{Cos[ $\varphi$ 1],Sin[ $\varphi$ 1],0},{-Sin[ $\varphi$ 1],Cos[ $\varphi$ 1],0},{0,0,1}};
  o' = Simplify[M1 . o];
  v2 = o'[[2;;3,3]];
   $\varphi$ 2 = -SafeArcTan @@ Reverse[v2];
  M2 = {{1,0,0},{0,Cos[ $\varphi$ 2],Sin[ $\varphi$ 2]},{0,-Sin[ $\varphi$ 2],Cos[ $\varphi$ 2]}};
  o'' = Simplify[M2 . o'];
  v3 = o''[[1;;2,1]];
   $\varphi$ 3 = SafeArcTan @@ v3;
  { $\varphi$ 1, $\varphi$ 2, $\varphi$ 3}
];

PauliDecomposition[U_?U4Q] := PauliDecomposition @ U4ToS06 @ U
PauliDecomposition[U_?S06Q] := Module[{a,b,c,d,U1a',Da,U2a',U1d',Dd,U2d',U1a,U2a,U1
  a = U[[1;;3,1;;3]];
  b = U[[4;;6,1;;3]];
  c = U[[1;;3,4;;6]];
  d = U[[4;;6,4;;6]];
  {U1a',Da,U2a'} = SingularValueDecomposition[a];
  {U1d',Dd,U2d'} = SingularValueDecomposition[d];
  {U1a,U2a,U1d,U2d} = {
    FullSimplify[Det[U1a']] * U1a',
    FullSimplify[Det[U2a']] * U2a',
    FullSimplify[Det[U1d']] * U1d',
    FullSimplify[Det[U2d']] * U2d'
  };
  Db = Simplify[Transpose[U1d] . b . U2a];
  Dc = Simplify[Transpose[U1a] . c . U2d];
  blocks = Map[{{Da[[#,#]],Db[[#,#]]},{Dc[[#,#]],Dd[[#,#]]}}&,Range[1,3]];
  { $\varphi$ XX, $\varphi$ YY, $\varphi$ ZZ} = -Map[AngleFind,blocks];
  { $\varphi$ ZI1, $\varphi$ XI1, $\varphi$ ZI2} = -SmallPauliDecomposition[U1a];
  { $\varphi$ ZI3, $\varphi$ XI2, $\varphi$ ZI4} = -SmallPauliDecomposition[Transpose[U2a]];
  { $\varphi$ IZ1, $\varphi$ IX1, $\varphi$ IZ2} = -SmallPauliDecomposition[U1d];
  { $\varphi$ IZ3, $\varphi$ IX2, $\varphi$ IZ4} = -SmallPauliDecomposition[Transpose[U2d]];
  {
```



```

      { $\varphi_{ZI1}$ , "ZI"}, { $\varphi_{XI1}$ , "XI"}, { $\varphi_{ZI2}$ , "ZI"},
      { $\varphi_{IZ1}$ , "IZ"}, { $\varphi_{IX1}$ , "IX"}, { $\varphi_{IZ2}$ , "IZ"},
      { $\varphi_{XX}$ , "XX"}, { $\varphi_{YY}$ , "YY"}, { $\varphi_{ZZ}$ , "ZZ"},
      { $\varphi_{ZI3}$ , "ZI"}, { $\varphi_{XI2}$ , "XI"}, { $\varphi_{ZI4}$ , "ZI"},
      { $\varphi_{IZ3}$ , "IZ"}, { $\varphi_{IX2}$ , "IX"}, { $\varphi_{IZ4}$ , "IZ"}
    }
  ];
  PauliDecomposition[_] := (
    Message[PauliDecomposition::notanoperator];
    $Failed
  );

```

Reconstruction using Pauli-rotation angles

```

ApproximateOp[_ ,  $\epsilon$ _ / ; Not[RealQ[ $\epsilon$ ]]] := (
  Message[Approximate::notreal];
  $Failed
);
ApproximateOp[U_?U4Q,  $\epsilon$ _] := Module[{anglesaxes, pauliapproximations},
  anglesaxes = PauliDecomposition[U];
  pauliapproximations = Map[PauliRotation[#[[1]],  $\epsilon$ /15, #[[2]]]&, anglesaxes];
  Simplify[Dot @@ pauliapproximations]
];
ApproximateOp[U_?S06Q,  $\epsilon$ _] := Module[{anglesaxes, pauliapproximations},
  anglesaxes = PauliDecomposition[U];
  pauliapproximations = Map[PauliRotation[#[[1]],  $\epsilon$ /15, #[[2]]]&, anglesaxes];
  U4ToS06 @ Simplify @ Dot @@ pauliapproximations
];
ApproximateOp[_ , _] := (
  Message[Approximate::notanoperator];
  $Failed
);

test[ApproximateSequence, "IfGaussianDoExact"] := BooleanQ;
test[ApproximateSequence, "SyllableType"] := MemberQ[{"Normal", "Asymmetric"}, #]&;
test[ApproximateSequence, "OutputType"] := MemberQ[{"String", "HexDec"}, #]&;

Options[ApproximateSequence] = {"IfGaussianDoExact" -> True, "OutputType" -> "String"};

ApproximateSequence[_ ,  $\epsilon$ _ / ; Not[RealQ[ $\epsilon$ ]], OptionsPattern[]] := (
  Message[ApproximateSequence::notreal];
  $Failed
);
ApproximateSequence[U_ ,  $\epsilon$ _ , opts:OptionsPattern[]] := Module[{normitops},
  normitops = FilterRules[{opts}, Options[NormIt]];
  If[
    OptionValue["IfGaussianDoExact"] && GaussianQ[U],
    NormIt[U, Append[normitops, "UpToPhase" -> False]],
    NormIt[ApproximateOp[U,  $\epsilon$ ], normitops]
  ]
];

```

Pseudorandom Circuit Generator

In this section, we develop a pseudo-random Clifford + CS circuit generator using our unique normal form. These circuits are random in the sense that given some integer n , we pseudorandomly sample uniformly from the set of all Clifford + CS operators whose optimal CS-count is at most n . In the limit

that n goes to infinity, this is uniformly random. Note that the overwhelming majority of operators in this set have an optimal CS-count around n .

Adjacency Matrices of Automata Graphs

The adjacency matrices of the graphs corresponding to the automata in [2].

```
adjacency3 = SparseArray[{
  {0,1,1},
  {1,0,1},
  {1,1,0}
}];
adjacency9 = SparseArray[{
  {0,1,1,1,1,0,0,0,0},
  {1,0,1,0,0,1,1,0,0},
  {1,1,0,0,0,0,0,1,1},
  {1,0,0,0,1,1,0,1,0},
  {1,0,0,1,0,0,1,0,1},
  {0,1,0,1,0,0,1,1,0},
  {0,1,0,0,1,1,0,0,1},
  {0,0,1,1,0,1,0,0,1},
  {0,0,1,0,1,0,1,1,0}
}];
adjacency15 = SparseArray[{
  {0,1,1,1,1,0,0,0,0,0,0,1,1,1,1},
  {1,0,1,0,0,1,1,0,0,0,1,0,1,1,1},
  {1,1,0,0,0,0,0,1,1,0,1,1,0,1,1},
  {1,0,0,0,1,1,0,1,0,1,0,1,1,0,1},
  {1,0,0,1,0,0,1,0,1,1,0,1,1,1,0},
  {0,1,0,1,0,0,1,1,0,1,1,0,1,0,1},
  {0,1,0,0,1,1,0,0,1,1,1,0,1,1,0},
  {0,0,1,1,0,1,0,0,1,1,1,1,0,0,1},
  {0,0,1,0,1,0,1,1,0,1,1,1,0,1,0},
  {0,0,0,1,1,1,1,1,1,0,0,0,0,1,1},
  {0,1,1,0,0,1,1,1,1,0,0,1,1,0,0},
  {1,0,1,1,1,0,0,1,1,0,1,0,1,0,0},
  {1,1,0,1,1,1,1,0,0,0,1,1,0,0,0},
  {1,1,1,0,1,0,1,0,1,1,0,0,0,0,1},
  {1,1,1,1,0,1,0,1,0,1,0,0,0,1,0}
```

Next Vertex Keys

Association keys which make random walks on the automata graphs easy to compute.

```

S3NextVertexList = GatherBy[adjacency3["NonzeroPositions"],First][[All,All,2]];
S3VertexKey = Association @@ MapIndexed[First[#2] -> #1&,S3NextVertexList];

S9NextVertexList = GatherBy[adjacency9["NonzeroPositions"],First][[All,All,2]];
S9VertexKey = Association @@ MapIndexed[First[#2] -> #1&,S9NextVertexList];

S15NextVertexList = GatherBy[adjacency15["NonzeroPositions"],First][[All,All,2]];
S15VertexKey = Association @@ MapIndexed[First[#2] -> #1&,S15NextVertexList];

```

Random Walk Vertex Lists

Functions to generate a list of vertices for random walks of a given length on our automata.

```

S13Random[0] = {};
S13Random[n_] := Module[{seed,tail,sylnumlist},
  seed = RandomInteger[{1,3},1];
  tail = RandomInteger[{1,2},n-1];
  sylnumlist = Fold[Append[#1,S3VertexKey[#1[[-1]]][[#2]]]&,seed,tail];
  sylnumlist
];

S49Random[0] = {};
S49Random[n_] := Module[{seed,tail,sylnumlist},
  seed = RandomInteger[{4,9},1];
  tail = RandomInteger[{1,4},n-1];
  sylnumlist = Fold[Append[#1,S9VertexKey[#1[[-1]]][[#2]]]&,seed,tail];
  sylnumlist
];

S1015Random[0] = {};
S1015Random[n_] := Module[{seed,tail,sylnumlist},
  seed = RandomInteger[{10,15},1];
  tail = RandomInteger[{1,8},n-1];
  sylnumlist = Fold[Append[#1,S15VertexKey[#1[[-1]]][[#2]]]&,seed,tail];
  sylnumlist
];

```

Sampling from the probability distribution of sub-walk lengths given a total walk length

Because Mathematica doesn't support sampling from most custom multivariate probability density functions, we use the marginal probability for sub-walks having a certain length given that the total walk length is at most n .

```

marginalpdfk[n_,k_] := (2/3)^(DiscreteDelta[k])*3*2^k*(15*8^(n-k)-7*4^(n-k)-1)/(4!
marginalpdfl[n_,k_,l_] := (2/3)^(DiscreteDelta[l])*3/2*4^l*(6*8^(n-k-l)+1)/(15*8^
marginalpdfm[n_,k_,l_,m_] := (4/3)^(DiscreteDelta[m])*21/4*8^m/(6*8^(n-k-l)+1);
distributionk[n_] := ProbabilityDistribution[marginalpdfk[n,k],{k,0,n,1}];
distributionl[n_,k_] := ProbabilityDistribution[marginalpdfl[n,k,l],{l,0,n-k,1}];
distributionm[n_,k_,l_] := ProbabilityDistribution[marginalpdfm[n,k,l,m],{m,0,n-k-
l}];

randklm[n_] := Module[{k,l,m},
  k = RandomVariate[distributionk[n]];
  l = RandomVariate[distributionl[n,k]];
  m = RandomVariate[distributionm[n,k,l]];
  {k,l,m}
];

```

Pseudorandom Circuits

Using our pseudorandom walk-length generator, our pseudorandom walk generator, our syllable list, a random Clifford generator, and our unique normal form from [2], we develop a function to generate uniformly pseudorandom Clifford + CS unitaries whose optimal CS-count is at most some integer n .

```

randsylnumlist[n_] := Module[{k,l,m},
  {k,l,m} = randklm[n];
  Join[S13Random[k],S49Random[l],S1015Random[m]]
];

RandomCliffCS[n_/(IntegerQ[n] && n > 0)] := Module[{sylnumlist,cliffnum,Rmult,clifford},
  sylnumlist = randsylnumlist[n];
  cliffnum = RandomInteger[{1,92160}];
  Rmult = Fold[Dot[#1,SyllableList[#2,1]]&,Id,sylnumlist];
  clifford = CliffordFromNumber[cliffnum];
  Rmult.clifford
];

RandomCliffCS[_] := (
  Message[RandomCliffCS::notanatural];
  $Failed
);

```

End of functions in the private context

```
End[];
```

Exported Functions

```
Protect[Id,X1,X2,Z1,Z2,W,H1,H2,S1,S2,CZ,CNOT12,CNOT21,EX,CS,R,U4ToS06,
      IdS06,X1S06,X2S06,Z1S06,Z2S06,I0S06,H1S06,H2S06,S1S06,S2S06,CZS06,CNOT12S06,CNOT21S06,EXS06,CS06,RS06,U4ToS06S06,
      FromSequence,FromHexDec,CliffordQ,CliffordSynth,RightCliffordSimilar,LeftCliffordSimilar,CandidateFinder,PauliRotation,PauliRotationSequence,PauliDecomposition,Approx
```

```
EndPackage[];
```