

## Introduction and Background

- P4 is a domain specific language for programming network switches in SDNs. P4 enables rapid experimentation and deployment of network protocols.
- Networking requirements are rapidly changing and rising in complexity. Software Defined Networking (SDN) introduces opportunities for advancing network protocols.
- P4 programs can have bugs and safety issues which cause undesirable behavior in SDN systems.
- p4-analyzer is introduced to detect these issues in P4 programs.

## Undesirable P4 Behaviors

- P4 programs can have many undesirable behaviors:
  - Use-before-initialization:** Using a variable before it has been initialized with a value.
  - Header-use-before-valid:** Reading from a header before it has been read from the input packet.
  - Undefined-forwarding:** Failing to forward or drop a packet.
  - ...and more
- Most of these issues can be detected statically.

## p4-analyzer

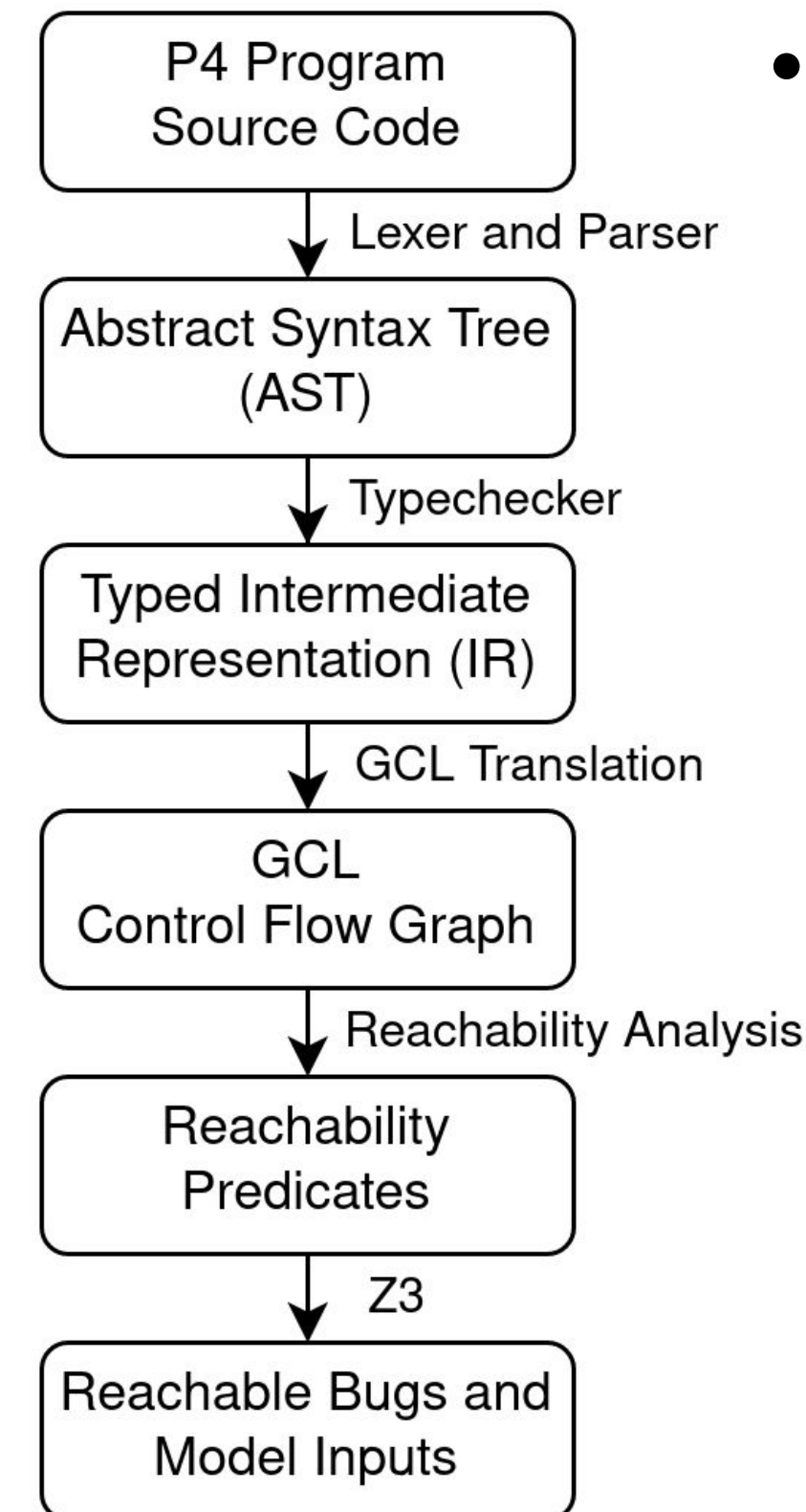
```
struct metadata_t { bool value; }

control my_control(inout metadata_t meta) {
    bool var;

    apply {
        if (meta.value) {
            // Bug: 'var' is uninitialized
            meta.value = var;
        }
    }
}
```

Figure 1. A P4 program with a bug.

## p4-analyzer



- p4-analyzer merges good ideas from previous works:
  - From p4v [3]: Translation to **GCL**.
    - Allows for a more semantic analysis of P4 programs.
  - From bf4 [2]: Translation to a **Control Flow Graph (CFG)** and **bug node injection**.
    - During GCL translation, bug nodes are inserted whenever undesirable behavior could occur.
  - From Petr4 [1]: The **type checker** and some of the **GCL translation** are based on the static and dynamic P4 semantics introduced by Petr4
    - The formal foundations from Petr4 give strong guarantees about the correctness of p4-analyzer.

Figure 2. The flow of data in p4-analyzer.

## Analysis of the Figure 1 P4 Program

● - Reachable ● - Reachable bug ● - Unreachable

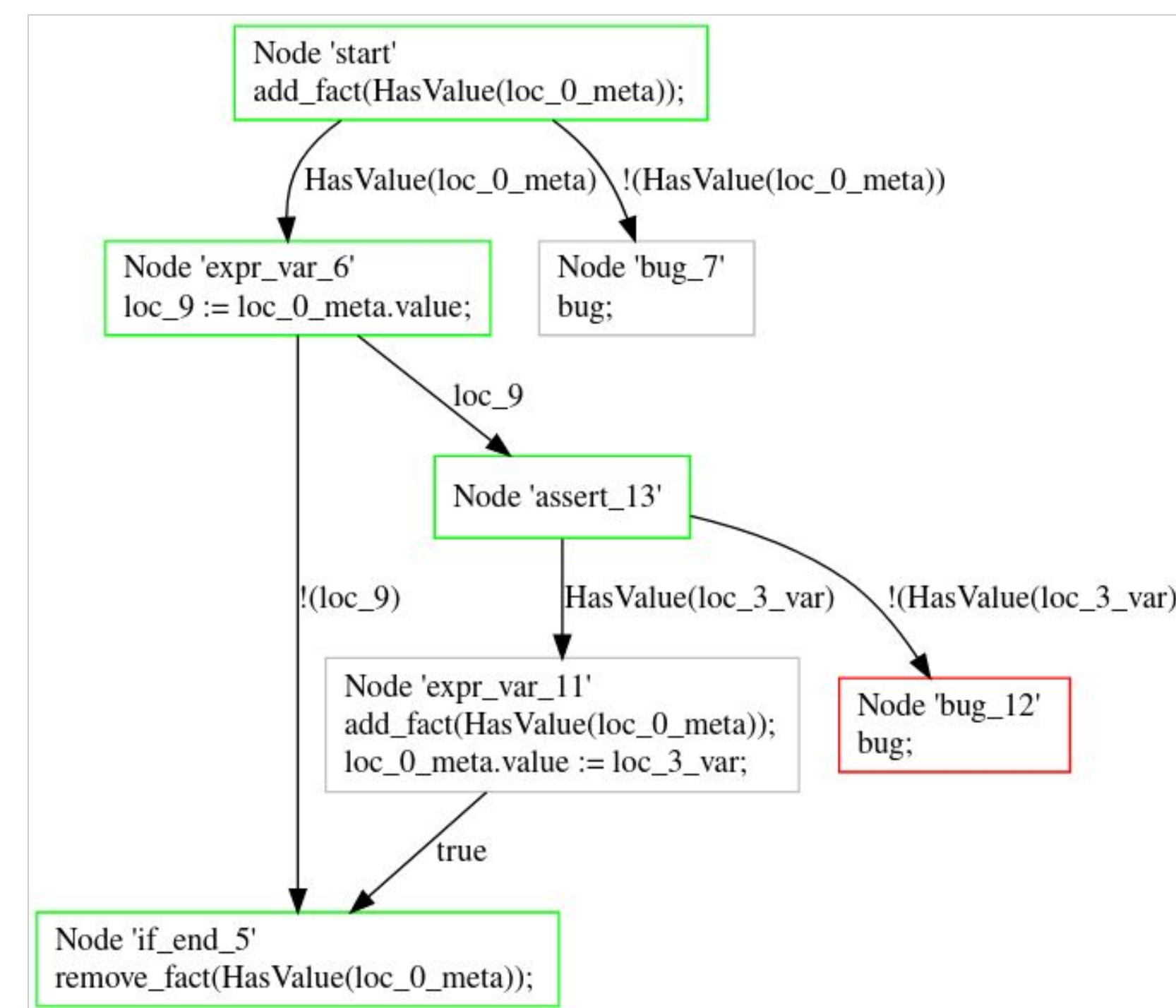


Figure 3. Computed GCL CFG, reachability predicates, and model of buggy input.

## Contributions

- p4-analyzer can detect non-trivial bugs in P4 programs.
  - Its foundation is **theoretically stronger** than previous work due to the combination of semantic analysis and use of Petr4's formal type system calculus.
- The translation of reachability predicates to Z3 includes a **novel translation** of advanced P4 types into Z3 types.
  - Struct types are translated into Z3 data types. This simplifies translation to and from Z3.
  - Example:
 

```

struct metadata_t {
    bool value;
}
Z3: meta = (metadata_t true)
P4: meta = (metadata_t) { value = true }
          
```
- Advanced uninitialized variable bug detection
  - The p4c reference compiler does simple uninitialized variable usage detection, but fails in more advanced cases.
  - bf4, being based off of p4c, inherits this simple bug detection.
  - p4-analyzer introduces an analysis of P4 programs which detects non-trivial instances of uninitialized variable usage.

## Conclusions and Future Work

- p4-analyzer performs advanced static analysis of P4.
- Support more of P4, including integers, bit strings, and headers.
  - p4-analyzer currently only supports a subset of P4.
    - Control blocks which use booleans and structs.
- Analyze P4 match-action tables and infer a control plane interface.
  - This would enable detecting control-plane-interface bugs.
- Investigate syntactic vs semantic analysis of P4.

## References

- Doenges et al. (2021). Petr4: Formal Foundations for P4 Data Planes. *Proc. ACM Program. Lang.*, 5(POPL). <https://doi.org/10.1145/3434322>
- Dumitrescu et al. (2020). Bf4: Towards Bug-Free P4 Programs. *Sigcomm '20*, 571–585. <https://doi.org/10.1145/3387514.3405888>
- Liu et al. (2018). P4v: Practical Verification for Programmable Data Planes. *Sigcomm '18*, 490–503. <https://doi.org/10.1145/3230543.3230582>