

**NAME**

getpid, getppid – get process identification

**SYNOPSIS**

```
#include <sys/types.h>
```

```
#include <unistd.h>
```

```
pid_t getpid(void);
```

```
pid_t getppid(void);
```

**DESCRIPTION**

**getpid()** returns the process ID (PID) of the calling process. (This is often used by routines that generate unique temporary filenames.)

**getppid()** returns the process ID of the parent of the calling process. This will be either the ID of the process that created this process using **fork()**, or, if that process has already terminated, the ID of the process to which this process has been reparented (either **init(1)** or a "subreaper" process defined via the **prctl(2)** **PR\_SET\_CHILD\_SUBREAPER** operation).

**ERRORS**

These functions are always successful.

**CONFORMING TO**

POSIX.1-2001, POSIX.1-2008, 4.3BSD, SVr4.

**NOTES**

If the caller's parent is in a different PID namespace (see **pid\_namespaces(7)**), **getppid()** returns 0.

From a kernel perspective, the PID (which is shared by all of the threads in a multithreaded process) is sometimes also known as the thread group ID (TGID). This contrasts with the kernel thread ID (TID), which is unique for each thread. For further details, see **gettid(2)** and the discussion of the **CLONE\_THREAD** flag in **clone(2)**.

**C library/kernel differences**

From glibc version 2.3.4 up to and including version 2.24, the glibc wrapper function for **getpid()** cached PIDs, with the goal of avoiding additional system calls when a process calls **getpid()** repeatedly. Normally this caching was invisible, but its correct operation relied on support in the wrapper functions for **fork(2)**, **vfork(2)**, and **clone(2)**: if an application bypassed the glibc wrappers for these system calls by using **syscall(2)**, then a call to **getpid()** in the child would return the wrong value (to be precise: it would return the PID of the parent process). In addition, there were cases where **getpid()** could return the wrong value even when invoking **clone(2)** via the glibc wrapper function. (For a discussion of one such case, see **BUGS** in **clone(2)**.) Furthermore, the complexity of the caching code had been the source of a few bugs within glibc over the years.

Because of the aforementioned problems, since glibc version 2.25, the PID cache is removed: calls to **getpid()** always invoke the actual system call, rather than returning a cached value.

On Alpha, instead of a pair of **getpid()** and **getppid()** system calls, a single **getxpid()** system call is provided, which returns a pair of PID and parent PID. The glibc **getpid()** and **getppid()** wrapper functions transparently deal with this. See **syscall(2)** for details regarding register mapping.

**SEE ALSO**

**clone(2)**, **fork(2)**, **gettid(2)**, **kill(2)**, **exec(3)**, **mkstemp(3)**, **tempnam(3)**, **tmpfile(3)**, **tmpnam(3)**, **credentials(7)**, **pid\_namespaces(7)**

**COLOPHON**

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