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OLLIR (Object-Oriented Low-Level Intermediate Representation) and the OLLIRTool

Compilers course

Masters in Informatics and Computing Engineering (MIEIC), 3rd Year

João M. P. Cardoso

Email: jmpc@fe.up.pt

Outline

- About three-address code representations
- About OLLIR
- OLLIR Examples
- About the OLLIRTool
- From OLLIR code to JVM code

About three-address code representations

- Intermediate representations mainly consisting of instructions with three operands
 - typical instructions include an assignment and a binary operation
 - closer to the register file based machines, such as RISC machines
 - includes conditional and unconditional branches
 - includes memory accesses
- Example of a three-address code representations
 - GCC RTL code

About OLLIR

- Inspired by three-address code representations
- Includes calls to functions following a multiple arguments, single result, format
- An OLLIR file represents a class
- Includes put and get primitives to write/read fields
- Supports classes and objects
- Elements of the OLLIR statements include attributes explicitly identifying the type of data
- It can be seen an intermediate low-level representation to represent Java classes

OLLIR Examples (1)

ex1.lir

```
myClass {
  .construct myClass().V {
    invokespecial(this, "<init>").V;
  }
  .method public sum(A.array.i32).i32 {
    sum.i32 :=.i32 0.i32;
    i.i32 :=.i32 0.i32;
  Loop:
    t1.i32 :=.i32 arraylength($1.A.array.i32).i32;
    if (i.i32 >=.i32 t1.i32) goto End;
    t2.i32 :=.i32 $1.A[i.i32].i32;
    sum.i32 :=.i32 sum.i32 +.i32 t2.i32;
    i.i32 :=.i32 i.i32 +.i32 1.i32;
    goto Loop;
  End:
    ret.i32 sum.i32;
  }
}
```

```
class myClass {
  public int sum(int[] A){
    int sum = 0;
    for(int i=0; i<A.length; i++) {
      sum += A[i];
    }
    return sum;
  }
}
```

OLLIR Examples (2)

ex3.lir

```
myClass {  
  .method public sum(A.array.i32, B.array.i32) .array {  
    t1.i32 :=.i32 arraylength($1.A.array.i32).i32;  
    C.array :=.array new(array, t1.i32).array;  
    i.i32 :=.i32 0.i32;  
    Loop:  
    t1.i32 :=.i32 arraylength($1.A.array.i32).i32;  
    if (i.i32 >=.i32 t1.i32) goto End;  
    t2.i32 :=.i32 $1.A[i.i32].i32;  
    t3.i32 :=.i32 $2.B[i.i32].i32;  
    t4.i32 :=.i32 t2.i32 +.i32 t3.i32;  
    C[i.i32].i32 :=.i32 t4.i32;  
    i.i32 :=.i32 i.i32 +.i32 1.i32;  
    goto Loop;  
    End:  
    ret.array.i32 C.array.i32;  
  }  
}
```

```
class myClass {  
  public int[] sum(int[] A, int[] B){  
    int[] C = new int[A.length];  
    for(int i=0; i<A.length; i++) {  
      C[i] = A[i] + B[i];  
    }  
    return C;  
  }  
}
```

ex4.lir

```
myClass {
  .field private a.i32;
  .construct myClass(n.i32).V {
    invokespecial(this, "<init>").V;
    putfield(this, a.i32, $1.n.i32).V;
  }
  .construct myClass().V {
    invokespecial(this, "<init>").V;
  }
  .method public get().i32 {
    t1.i32 :=.i32 getfield(this, a.i32).i32;
    ret.i32 t1.i32;
  }
  .method public put(n.i32).V {
    putfield(this, a.i32, $1.n.i32).V;
  }
  .method public m1().V {
    putfield(this, a.i32, 2.i32).V; // this.a = 2;
    t2.String :=.String ldc("val = ").String;
    t1.i32 :=.i32 invokevirtual(this, "get").i32;
    invokestatic(io, "println", t2.String, t1.i32).V; // io.println("val = ", this.get());
    c1.myClass :=.myClass new(myClass, 3.i32).myClass;
    invokespecial(c1.myClass, "<init>").V; // myClass c1 = new myClass(3);
    t3.i32 :=.i32 invokevirtual(c1.myClass, "get").i32;
    invokestatic(io, "println", t2.String, t3.i32).V; // io.println("val = ", c1.get());
    invokevirtual(c1.myClass, "put", 2.i32).V; // c1.put(2);
    t4.i32 :=.i32 invokevirtual(c1.myClass, "get").i32;
    invokestatic(io, "println", t2.String, t4.i32).V; // io.println("val = ", c1.get());
  }
  .method public static main(args.array.String).V {
    A.myClass :=.myClass new(myClass).myClass;
    invokespecial(A.myClass, "<init>").V;
    invokevirtual(A.myClass, "m1").V;
  }
}
```

Note: since the Java-- version used this year does not accept Strings, some of the instructions in this exemple are not needed (see next slide with a similar exemple, but without strings)

ex4b.lir

```
myClass {
    .field private a.i32;
    .construct myClass(n.i32).V {
        invokespecial(this, "<init>").V;
        putfield(this, a.i32, $1.n.i32).V;
    }
    .construct myClass().V {
        invokespecial(this, "<init>").V;
    }
    .method public get().i32 {
        t1.i32 :=.i32 getfield(this, a.i32).i32;
        ret.i32 t1.i32;
    }
    .method public put(n.i32).V {
        putfield(this, a.i32, $1.n.i32).V;
    }
    .method public m1().V {
        putfield(this, a.i32, 2.i32).V; // this.a = 2;
        t1.i32 :=.i32 invokevirtual(this, "get").i32;
        invokestatic(io, "println", t1.i32).V; // io.println(this.get());
        c1.myClass :=.myClass new(myClass, 3.i32).myClass;
        invokespecial(c1.myClass, "<init>").V; // myClass c1 = new myClass(3);
        t3.i32 :=.i32 invokevirtual(c1.myClass, "get").i32;
        invokestatic(io, "println", t3.i32).V; // io.println(c1.get());
        invokevirtual(c1.myClass, "put", 2.i32).V; // c1.put(2);
        t4.i32 :=.i32 invokevirtual(c1.myClass, "get").i32;
        invokestatic(io, "println", t4.i32).V; // io.println(c1.get());
    }
    .method public static main(args.array.String).V {
        A.myClass :=.myClass new(myClass).myClass;
        invokespecial(A.myClass, "<init>").V;
        invokevirtual(A.myClass, "m1").V;
    }
}
```

```
class myClass {
    int a;

    myClass(int n){
        this.a = n;
    }
    myClass(int n){
        this.a = n;
    }
    public int get(){
        return this.a;
    }
    public void put(int n){
        this.a = n;
    }
    public void m1(){
        this.a = 2;
        io.println(this.get());

        myClass c1 = new myClass(3);
        io.println(c1.get());

        c1.put(2);
        io.println(c1.get());
    }
    public static void main(String[] args){
        myClass A = new myClass();
        A.m1();
    }
}
```


OLLIR Examples (4)

ex5.lir

```
myClass {  
  .construct myClass().V {  
    invokespecial(this, "<init>").V;  
  }  
  .method public check(A.array.i32, N.i32, T.i32).bool {  
    i.i32 :=.i32 0.i32;  
    all.bool :=.bool 0.bool;  
    Loop:  
    t1.bool :=.bool i.i32 <.i32 $2.N.i32;  
    t2.i32 :=.i32 $1.A[i.i32].i32;  
    t3.bool :=.bool t2.i32 <.i32 $3.T;  
    if (t1.bool & .bool t3.bool) goto Body;  
    goto EndLoop;  
    Body:  
    i.i32 :=.i32 i.i32 +.i32 1.i32;  
    goto Loop;  
    EndLoop:  
    if (i.i32 ==.i32 $2.N.i32) goto Then;  
    goto End;  
    Then:  
    all.bool :=.bool 1.bool;  
    End:  
    ret.bool all.bool;  
  }  
}
```

```
class myClass {  
  public boolean check(int[] A, int N, int T){  
    int i = 0;  
    boolean all = false;  
    while((i < N) && (A[i] < T)) {  
      i++;  
    }  
    if(i == N) all = true;  
    return all;  
  }  
}
```

OLLIR Examples (5)

Fac.lir

```
Fac {
  .method public compFac(num.i32).i32 {
    if ($1.num.i32 >=.i32 1.i32) goto else;
    num_aux.i32 :=.i32 1.i32;
    goto endif;
  else:
    aux1.i32 :=.i32 $1.num.i32 -.i32 1.i32;
    aux2.i32 :=.i32 invokevirtual(this, "compFac", aux1.i32).i32;
    num_aux.i32 :=.i32 $1.num.i32 *.i32 aux2.i32;
  endif:
    ret.i32 num_aux.i32;
  }
  .method public static main(args.array.String).V {
    aux1.Fac :=.Fac new(Fac).Fac;
    invokespecial(aux1.Fac, "<init>").V;
    aux2.i32 :=.i32 invokevirtual(aux1.Fac, "compFac", 10.i32).i32;
    invokestatic(io, "println", aux2.i3).V;
  }
}
```

```
class Fac {
  public int compFac(int num){
    int num_aux;
    if (num < 1)
      num_aux = 1;
    else
      num_aux = num * (this.compFac(num-1));
    return num_aux;
  }
  public static void main(String[] args){
    io.println(new Fac().compFac(10));
  }
}
```

About the OLLIRTool

- Parses OLLIR code and represents the code as Java classes
 - Each method includes a CFG with the OLLIR instructions
 - Each method includes a VarTable (a symbol table of parameters and local variables)
- Current parser does not include error recovery and grammar is permissive
- The parser considers that the OLLIR code has been generated from a compiler and is correct
 - so, it is possible that some OLLIR incorrect code is not identified as incorrect!

From OLLIR code to JVM code (1)

- The JVM is a stack-based machine
 - The instructions of stack-based machines are also known as zero-address code
- The translation of OLLIR code to JVM code is almost direct, with exception of the instructions of each method
 - A translation between three-address code to zero-address code is needed!

From OLLIR code to JVM code (2)

- Non-optimized JVM code – each three-address OLLIR instruction is translated to a sequence of JVM instructions always loading/storing values from/to JVM local variables
 - i.e., without considering that the stack can be very useful to communicate intermediate results

source code:

```
int a, b, c;  
...  
a = b+c;
```

OLLIR code:

```
a.i32 :=.i32 b.i32 +.i32 c.i32;
```

Symbol Table:

OLLIR var	JVM local var
...	...
b	1
c	2
a	3
...	...

JVM code:

```
iload_1  
iload_2  
iadd  
istore_3
```

From OLLIR code to JVM code (3)

- Non-optimized JVM code – each three-address OLLIR instruction is translated to a sequence of JVM instructions always loading/storing values from/to JVM local variables
 - i.e., without considering that the stack can be very useful to communicate intermediate results

source code:

```
int a, b, c;  
...  
a = b+c+1;
```

OLLIR code:

```
t1.i32 := .i32 b.i32 +.i32 b.i32;  
a.i32 := t1.i32 +.i32 1.i32;
```

Symbol Table:

OLLIR var	JVM local var
...	...
b	1
c	2
t1	3
a	4
...	...

JVM code:

```
iload_1  
iload_2  
iadd  
istore_3  
iload_3  
iconst_1  
iadd  
istore_4
```

JVM code:

```
iload_1  
iload_2  
iadd  
iconst_1  
iadd  
istore_4
```