Homework 3 - Integer Programming

Adv. Analytics and Metaheuristics

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1 - Problem 1

1.1 Mathematical Formulation

1.1.1 Sets

Set Name	Description			
GENERATORS PERIODS	Set of generators i that can be used (A,B,C) 2 possible periods p (1, 2) in the production day			

1.1.2 Parameters

Parameter Name	Description		
$\overline{S_i}$	Fixed cost to start a generator		
	$(i \in GENERATORS)$ in the entire day		
F_i	Fixed cost to operate a generator		
	$(i \in GENERATORS)$ in any period		
C_{i}	Variable cost per megawatt to operator a		
	generator $(i \in GENERATORS)$ in any		
	period		
U_i	Max. megawatts generated for generator		
	$(i \in GENERATORS)$ in any period		
$demand_p$	Total demanded megawatts for period		
1	$(p \in PERIODS)$		
M	Large constant to map watts used by each		
	generator $(i \in GENERATORS)$		

1.1.3 Decision Variables

Variable Name	Description		
$\overline{watts_{i,p}}$	Integer variable: Number of watts to		
	produce per generator		
	$(i \in GENERATORS)$ per period		
	$(p \in PERIODS)$		
$x_{i,p}$	Binary variable: 1 if a generator		
	$(i \in GENERATORS)$ is in period p		
	$(p \in PERIODS)$, 0 if not turned on at all		
y_i	Binary variable: 1 if a generator		
	$(i \in GENERATORS)$ is used, 0 if not		
	turned on at all		

1.1.4 Objective Function

$$minimize\ cost: \sum_{i \in GENERATORS} \left(\left(\sum_{p \in PERIODS} (watts_{i,p}) \times C_i \right) + \left(F_i \times \sum_{p \in PERIODS} x_{i,p} \right) + \left(S_i \times y_i \right) \right)$$

1.1.5 Constraints

C1: For each period, meet the demanded megawatts

$$requiredWatts: \sum_{i \in GENERATORS} (watts_{i,p}) = demand_p, \forall \ p \in PERIODS$$

C2: For each generator, don't surpass the allowable megawatts

$$upperBound: \sum_{p \in PERIODS} (watts_{i,p}) \leq U_i, \forall i \in GENERATORS$$

C3: For each generator, map decision variables together to account for the fixed costs in a given day S_i

$$mapVars: \sum_{p \in PERIODS} (watts_{i,p}) \leq M_i \times y_i, \forall i \in GENERATORS$$

C4: For each generator and period, map decision variables y and watts together to account for the fixed costs in a per period p

$$mapVars2: watts_{i,p} \leq M_i \times x_{i,p}, \forall i \in GENERATORS, p \in PERIODS$$

C5 Non-negativity or Binary restraints of decision variables

$$watts_{i,p} \ge 0$$

$$x_{i,p}, y_i \in (0,1)$$

1.2 Code and Output

1.2.1 Code

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| Particle | Control | Cont
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1.2.2 Output

```
ampl: model 'C:\Users\daniel.carpenter\OneDrive - the Chickasaw
CPLEX 20.1.0.0: optimal integer solution; objective 46100
7 MIP simplex iterations
0 branch-and-bound nodes

Which generators are used?
y [*] :=
A 1
B 1
C 1;

Which periods were the generators used?
x :=
A 1 0
A 2 1
B 1 0
B 2 1
C 1 1
C 2 0;

Optimal Amount of Megawatts for each generator and period:
watts :=
A 1 0
A 2 2100
B 1 0
B 2 1800
C 1 2900
C 2 0;
```

1.2.2.1 Analysis of the Output

- The minimized cost is \$46, 100
- Generator A, B, and C run
- Generator C runs in period 1. Generator A and B run in period 2
- Generator A produces 2, 100 megawatts in total
- Generator B produces 1,800 megawatts in total
- Generator C produces 2,900 megawatts in total

2 - Problem 2

2.1 Mathematical Formulation (Part a)

2.1.1 Sets

Set Name	Description
PRODUCTS	5 types of landscaping and construction products (e.g., cement, sand, etc.) labeled product (p) A, B, C, D , and E
SILOS	8 different silos s that each product must be stored in $(1, 2,, 8)$

2.1.2 Parameters

Parameter Name	Description		
$\overline{cost_{s,p}}$	Cost of storing one ton of product $p \in PRODUCTS$ in silo $s \in SILOS$		
$supply_p$	Total supply in tons available of product $p \in PRODUCTS$		
$capacity_s$	Total capacity in tons of silo $s \in SILOS$. Can store products.		
M	Variable to map $decision \ variable$ $tonsOfProduct_{p,s}$ to $isStored_{p,s}$. Uses big M method.		

2.1.3 Decision Variables

Variable Name	Description		
$tonsOfProduct_{p,s}$ $isStored_{p,s}$	Tons of product $p \in PRODUCTS$ to store in silo $s \in SILOS$. Non-negative. Binary variable indicating if product $p \in PRODUCTS$ is stored in silo $s \in SILOS$.		

2.1.4 Objective Function

$$minimize\ costOfStorage: \sum_{p \in PRODUCTS} \sum_{s \in SILOS} tonsOfProduct_{p,s} \times cost_{p,s}$$

2.1.5 Constraints

C1: For each silo s, the tons of the supplied product p must be less than or equal to the capacity limit of silo s

$$meetCapacity: \sum_{p \in PRODUCTS} tonsOfProduct_{p,s} \leq capacity_s, \ \forall \ s \in SILOS$$

C2: For each product p, must use all of the total product that is available

$$useAllProduct: \sum_{s \in SILOS} tonsOfProduct_{p,s} = supply_p, \ \forall \ p \in PRODUCTS$$

C3: For each silo s and product p,

$$oneProductInSilo: \sum_{pinPRODUCTS} isStored_{p,s} = 1, \ \forall \ s \in SILOS$$

C4: Map the decision variables together using the Big M method

$$mapVars: tonsOfProduct_{p,s} \leq M \times isStored_{p,s}, \ \forall \ p \in PRODUCTS, \ \forall \ s \in SILOS$$

C5 Non-negativity or Binary restraints of decision variables

$$tonsOfProduct_{p,s} \geq 0$$

$$isStored_{p,s} \in (0,1)$$

2.2 Code and Output (Part a)

2.2.1 Code

```
# gracinal part of the control of th
```

2.2.2 Output (Part a)

```
ampl: model 'C:\Users\daniel.carpenter\OneDrive - the Chickasaw Nation\D CPLEX 20.1.0.0: optimal integer solution; objective 320 48 MIP simplex iterations 0 branch-and-bound nodes
0
0
0
               0
0
0
0
12345678
                              1
0
               1
0
0
0
                              1
1
                                      0
                       0
Optimal tons of product allocated to each silo: tonsOfProduct [*,*] (tr)
: A B C D E :=
                 0
0
0
0
50
                                      0
                                               0000
                                    25 0 0 0
                                             20
```

2.2.2.1 Analysis of the Output

- Minimized loading cost for 250 tons of 5 products over the 8 silos is 320 (problem does not state cost units).
- Product A stores 25 tons in silo 1 and 50 tons in silo 4
- Product B stores 50 tons in silo 5
- Product C stores 25 tons in silo 3
- Product D stores 25 tons in silo 2, 5tons in silo 7, and and 50 tons in silo 8
- Product E stores 20 tons in silo 6

2.3 Problem 2 b

- Create a new objective that also minimizes the distance between capacity and stored tons of product
- For each silo, minimize the variance between the total capacity and the tons of product

 $minimize\ capacity Actual Variance: capacity_s - \sum_{p \in PRODUCTS} tonsOfProduct_{p,s},\ \forall s \in SILOS$

2.4 Code and Output (Part b)

2.4.1 Code

```
Common interactions of Colonia in Colonia State Colonia (Colonia) (Colonia)
```

2.4.2 Output (Part b)

```
ampl: model 'C:\Users\daniel.carpenter\OneDrive - the Chickasav CPLEX 20.1.0.0: optimal integer solution; objective 320 48 MIP simplex iterations 0 branch-and-boundes
Objective = costOfStorage
which silo(s) stores what product?
isStored [*,*] (tr)
: A B C D E :=
        A
1
0
                                         0
                0
                         0
                                 0
                         ŏ
                0
                                 1
                                 0
        0
                0
                                         0
                                         0
        1
                0
                                 0
        ō
                                 ŏ
                                         1
0
0
        0
Optimal tons of product allocated to each silo: tonsOfProduct [*,*] (tr)
: A B C D E :=
                                       0
25
                             0
0
25
0
0
0
12345678
                                                    0
          0
                    000
          0
                                       0
0
5
50
                                                  0
20
0
0
                   50
0
0
          0
```

2.4.2.1 Analysis of the Output

- The optimal cost actually stays the same, but the amount of iterations to get to that solution is much more.
- The values of the decision variables are the same.

3 - Problem 3

3.1 Mathematical Formulation

3.1.1 Sets

Set Name Description

3.1.2 Parameters

Parameter Name Description

3.1.3 Decision Variables

Variable Name Description

3.1.4 Objective Function

3.1.5 Constraints

C1:

3.2 Code and Output

3.2.1 Code

```
options solver cplex; # Using cplex for simplex alg
 11 param theDemand := 55000; # The demanded amount of products
        param mcWII := 4.95; # Marginal cost compnent of WII
        param availWII := 18000; # Amount of WII that is available
        s.t. upperBoundWII: WII <= availWII;</pre>
         param mcWRS := 2.30; # Marginal cost compnent of WRS
        param fixWRS := 20000; # Fixed Cost component of WRS
        param availWRS := 14000; # Amount of WRS that is available
         var WRS >= 0; # amt of product WRS to produce
        var yWRS1 binary; # Binary used for fixed cost if used
         s.t. map_yWRS1: WRS <= availWRS * yWRS1; # Upper bound and map</pre>
        param mcWE1 := 3.95; # If buy from WRS, m. cost for WE
        param mcWE2 := 4.10; # Else m. cost for WE
        param availWE := 7000; # Amount of WE that is available
           var yWRS binary; # If WRS is selected
            var yWII binary; # If WII is selected
            var yWE1 binary; # If WE is selected
             var yWE2 binary;
                      binary;
```

```
s.t. mapWE1: WE1 <= M * yWE1; # Map the W vars to the y binary
s.t. mapWE2: WE2 <= M * yWE2; # ""
s.t. mapWRS: WRS <= M * yWRS; # ""
s.t. mapWII: WII <= M * yWII; # ""</pre>
                    s.t. ifWRS_ThenWE1:    yWRS <= yWE1 + M*z;</pre>
                     # If WE2, cannot do WII. (Use of Mz to choose one constraint)
s.t. ifWRS_thenNotWII: yWE2 + yWII <= 1 + M*(1-z);</pre>
                     s.t. only1WE: yWE1 + yWE2 <= 1;
                     s.t. setWE: WE == WE1 + WE2;
                 s.t. upperBoundWE: WE <= availWE; # Meet the upper bound limit
          param mcWU
          param availWU := 22000; # Amount of WU that is available
          param minBuyAmt := 15000; # Must buy at Least 15k
           var WU >= 0; # amt of product WU to produce
           var yWU binary; # Binary used for fixed cost if used
            s.t. buyAtLeastMin: WU <= availWU * yWU; # Buy at least min amount
s.t. map_yWU: WU >= minBuyAmt * yWU; # Under the upper bound
            param mcWOW1 := 9.50; param mcWOW1Upper := 3000; # 3000 upper bound
            param mcWOW2 := 4.90; param mcWOW2Upper := 6000; # 3000 + 6000 = 9000 upper bound
            param mcWOW3 := 2.75; param mcWOW3Upper := 25000; # Cannot exceed 25000 due to supply
            var WOW >= 0; #amt of product WOW to produce
            var d1WOW >=0; # piecewise component 1 of var WOW
            var d2WOW >=0; # piecewise component 2 of var WOW
var d3WOW >=0; # piecewise component 3 of var WOW
            var y1WOW binary; #to model piecewise cost for var WOW
            var y2WOW binary; #to model piecewise cost for var WOW
```

```
group23_HW3_p3.mod ×

□

group23_HW3_p3.mod ×
              s.t. X_WOW: WOW = d1WOW + d2WOW + d3WOW;
              s.t. piecela: mcWOW1Upper*y1WOW <= d1WOW;</pre>
              s.t. piece1b: d1WOW <= mcWOW1Upper;</pre>
             s.t. piece2a: mcWOW2Upper*y2WOW <= d2WOW;</pre>
              s.t. piece2b: d2WOW <= mcWOW2Upper*y1WOW;</pre>
              s.t. piece3: d3WOW <= mcWOW3Upper*y2WOW;</pre>
              s.t. upperBoundWOW: WOW <= mcWOW3Upper;</pre>
      s.t. meetTheDemand: WII + WRS + WE + WU + WOW >= theDemand;
163 minimize cost: mcWII*WII
                      + fixWRS*yWRS1 + mcWRS*WRS # WRS: Fixed plus variable
                      + mcWE1*WE1 + mcWE2*WE2 # WE: Continguint mc based on scenario
                      + mcWU*WU
                      + mcWOW1*d1WOW + mcWOW2*d2WOW + mcWOW3*d3WOW # WOW: Piecewise
      printf "Demand\t| WII\t| WRS\t| WE\t| WU\t| WOW\t| Total Cost";
      printf "\n%s\t %s\t %s\t %s\t %s\t %f", theDemand, WII, WRS, WE, WU, WOW, cost;
```

3.2.2 Output
Summary table of Output

Demand	WII	WRS	WE	WU	WOW	Total Cost
5000	0	0	5000	0	0	19750.000000
10000	3000	0	7000	0	0	42500.000000
25000	4000	14000	7000	0	0	99650.000000
35000	0	14000	6000	15000	0	139650.000000
45000	0	14000	6000	0	25000	177800.000000
50000	4000	14000	7000	0	25000	201550.000000
55000	0	14000	1000	15000	25000	221800.000000

Snapshots of Compilation

ampl: model 'C:\Users\daniel.carpenter\OneDrive - the Chickasaw ! CPLEX 20.1.0.0: optimal integer solution; objective 19750 3 MIP simplex iterations 0 branch-and-bound nodes Demand WII WRS WE WU WOW | Total Cost 5000 19750.000000 5000 0 0 0 | Total Cost Demand | WII l WRS l WE WOW WU 10000 3000 7000 0 42500.000000 ampl. WOW | Total Cost Demand WII WRS WE WU 25000 4000 14000 7000 0 99650.000000 WRS | Total Cost Demand WII WE WU WOW 139650.000000 35000 0 14000 6000 15000 0 Demand WII l WRS l WE WU WOW | Total Cost 45000 25000 177800.000000 14000 6000 0 -mn1. Demand WII WRS WE WU WOW | Total Cost 50000 4000 14000 7000 0 25000 201550.000000 Demand WII WRS WE WU WOW | Total Cost 55000 0 14000 1000 15000 25000 221800.000000