

# An Algebraic Introduction to Mathematical Logic

## Chapter 3 Properties of the Propositional Calculus

### Section 1 Introduction

### Exercises

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**Definition 1.** A logic  $\mathcal{L}$  is a system consisting of a set  $P$  of elements (called propositions), a set  $\mathcal{V}$  of functions (called valuations) from  $P$  to some value set  $W$ , and, for each subset  $A$  of  $P$ , a set of finite sequences of elements of  $P$  (called proofs from the assumptions  $A$ ).

**Example 1.** The logic called The Propositional Calculus on the set  $X$ , denoted  $Prop(X)$ , consists of the set  $\mathcal{V}$  of all homomorphisms of  $P(X)$  onto  $\mathbb{Z}_2$ , and the set of proofs defined as in section 4 of chapter 2.

**Definition 2.** A logic  $\mathcal{L}$  is *sound* if  $A \vdash p$  implies that  $A \models p$ .

**Definition 3.** A logic  $\mathcal{L}$  is *consistent* if  $F$  is not a theorem.

**Definition 4.** A logic  $\mathcal{L}$  is *adequate* if  $A \models p$  implies that  $A \implies p$ .

**Definition 5.** A proposition is *valid* or *tautological* in a logic if for every valuation  $v \in \mathcal{V}$ ,  $v(p) = 1$  where  $W = \mathbb{Z}_2$  and 1 captures our intuitive notion of truth.

**Definition 6.** A logic  $\mathcal{L}$  is *decidable for validity* if there exists an algorithm which determines for every proposition  $p$ , in a finite number of steps, whether or not  $p$  is valid.

**Definition 7.** A logic  $\mathcal{L}$  is *decidable for provability* if there exists an algorithm which determines for every proposition  $p$ , in a finite number of steps, whether or not  $p$  is a theorem.