

Q1

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Time Complexity of finding the middle of the linked list size  $n$  is  $O(n)$ . This is because the slower pointer is at the middle node once we pass through the entire linked list of  $n$  elements. Another way to think of it is the linked list is traversed for the entire linked list and second till the middle of the linked list. Big  $O$  of traversing the entire list is  $O(n)$  and for the middle is  $O(n/2)$  adding  $O(n) + O(n/2)$  gives  $O(n)$ .

Q2

→ 2.1)

The time complexity of the given code is  $O(n \log n)$ . The outer loop has  $\log(n)$  iterations as  $i$  is doubled in each iteration. In the inner loop,  $j$  starts from 1 and goes up to  $i$  so the number of iterations of the inner loop depends on the value of  $i$  in each iteration of the outer loop.

Inner loop iterations is  $1 + 2 + 4 + \dots + n$  which represents a geometric series and hence can be calculated as  $n$ .

2.2) Time complexity of outer loop is  $\log_2 n$ .

Each iteration of  $i$  is doubled in this loop. For the inner loop, in each iteration, the value of  $j$  is doubled, and when  $j$  becomes greater than  $i$ , the loop will terminate. This gives time complexity of this code:  $O(\log^2 n)$

2.3) The time complexity of outer loop is  $O(\log n)$  as the value of  $i$  increases with a factor of 3 in each iteration. For the inner loop, time complexity is  $O(\log n)$  as the value of  $j$  increases with a factor of 2 in each iteration. Overall time complexity is  $O(\log^2 n)$ .

2.4) In the outer loop, the time complexity of this loop is  $O(n)$ .

In the inner loop, time complexity is  $O(\log n)$ .

Total time complexity of the code is  $O(n \log n)$ .

Q3

→ 3.1]

$$1) f_A(n) = 4n^3 + 3n^2 + 2n - 5$$

$$4n^3 \leq 4n^3 + 3n^2 + 2n - 5 \leq 4n^3 + 3n^3 + 2n^3 - 5n^3$$

$$4n^3 \leq f_A(n) \leq (4+3+2-5)n^3$$

$$4n^3 \leq f_A(n) \leq 4n^3$$

$$4f(n) \leq f_A(n) \leq 4f(n)$$

$$c_1 = 4$$

$$c_2 = 4$$

$$n_0 = 1$$

$$= f(n) + \Theta(n^3)$$

$$= O(n^3)$$

$$2) n \log n + 3n - 15$$

$$f(n) = \Theta(n \log n)$$

$$g(n) = n \log n$$

$$c_1 g(n) \leq f(n) \leq c_2 g(n)$$

$$c_1 n \log n \leq n \log n + 3n - 15 \leq c_2 n \log n$$

$$n \log n \leq n \log n + 3n - 15 \leq 11 n \log n$$

$$c_1 = 1$$

$$c_2 = 11$$

$$n_0 = 1$$

$$f(n) \in \Theta(n \log n)$$

$$\text{big } \Theta = \Theta(n \log n)$$

3.2]

$$1) n^2 \log(3n) + \underline{2n^4} + 3n^2 + 12$$

$$= O(n^4)$$

$$2) \underline{\sqrt{2n}} + 30 \log(4n)^2 + 27n - 3$$

$$= O(\sqrt{n})$$

$$3) \underline{(n+1)!} + 2^n$$

$$= O(n!)$$

$$4) \sqrt[3]{n^2} + 3n \log^2(2n) + \underline{4n^{\frac{6}{5}}}$$

$$= O(n^{\frac{6}{5}})$$