

演化計算 → 由基因演算法而來

evolutionary algorithms, genetic algorithms, evolution strategy, evolutionary programming, artificial life

Optimization: → 最佳化 question !

1. specific optimization (differential method
gradient method)

↳ 用在特定 problem

2. general ----- (random search, GA)

↳ 用在通用 problem

基因演算法

P. S. O

Ant C.

△. G. A (1975, Holland)

operators: 1. reproduction

複製

2. crossover

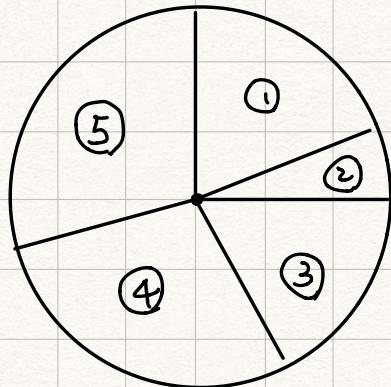
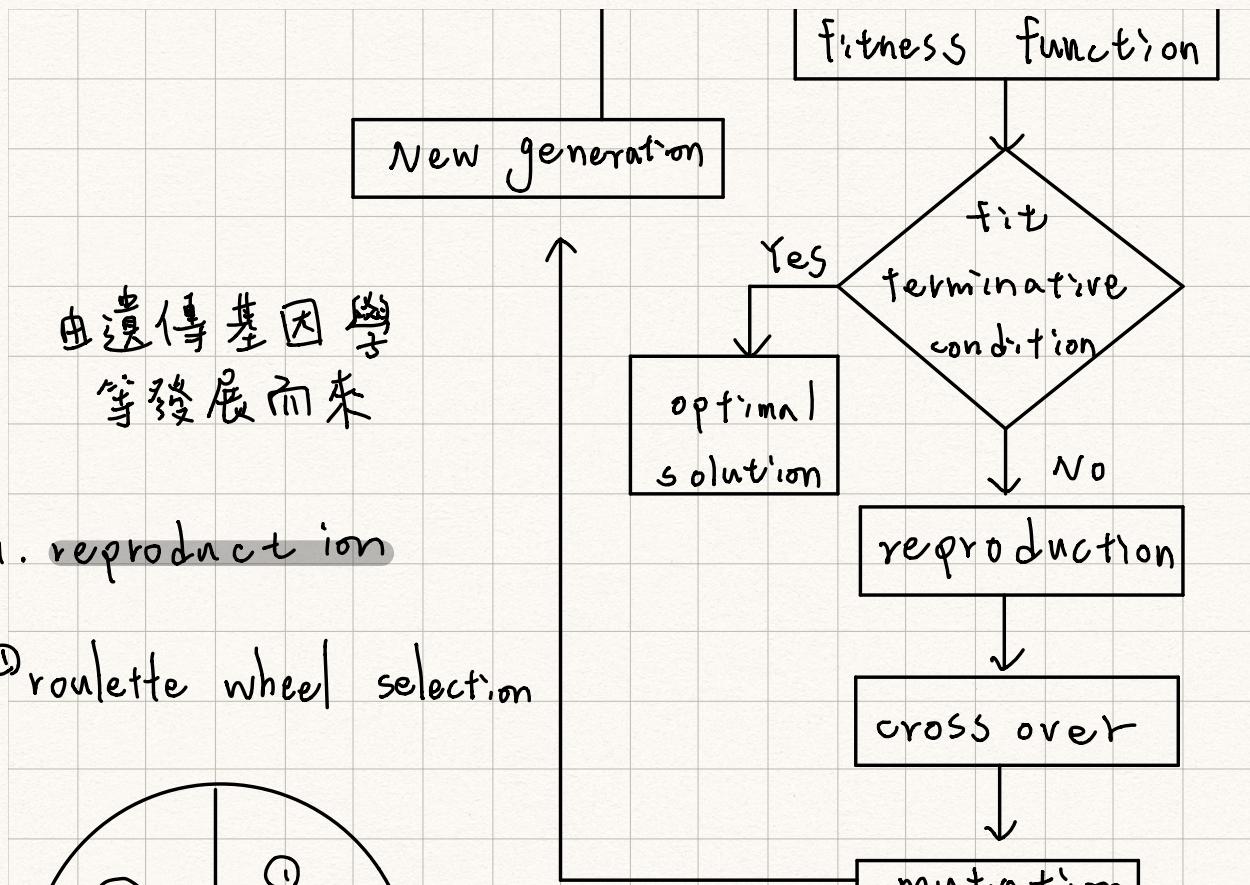
交配

3. mutation

突變

generate
initial generation
randomly

decode (binary)
string and calculate



圓弧長(面積)大小由 fitness function 來決定 (適應值)

ex: 5 · 3 · 4 · 1 · 5

② tournament selection

Randomly select two or more chromosomes from **parent generation** and put the chromosome with highest fitness value in the mating pool.

Repeat N times, where N is the population size.
似人口，要維持不變

擇 K 個，選一個最大的，擇至 N 個數為止
可重複。

ex: 5 · 1 · 4 · 3 · 4

2. crossover

① one-point crossover

$$\begin{array}{cccccc} A = & 0 & 0 & 0 & 0 & 0 \\ B = & 1 & 1 & 1 & 1 & 1 \end{array} \longrightarrow \begin{array}{ccccc} A' = & 0 & 0 & 1 & 1 & 1 \\ B' = & 1 & 1 & 0 & 0 & 0 \end{array}$$

五個位置，在#2後相互交換

② or #2開始交換

固定做法即可，but 做法不限定

②

two-point crossover

$$\begin{array}{ccccc} 0 & 0 & 0 & 1 \\ 0 & 1 & 1 & 0 \\ 1 & 0 & 1 & 0 \\ 1 & 1 & 0 & 1 \end{array}$$

	※2	※4	
$A = 0 \ 0 \ 0 \ 0 \ 0$			$A' = 0 \ 1 \ 1 \ 1 \ 0$
$B = 1 \ 1 \ 1 \ 1 \ 1$			$B' = 1 \ 0 \ 0 \ 0 \ 1$

介於 2~4 之間的點進行交換

③ uniform crossover (mask crossover)

	$\begin{matrix} \text{mask} \\ 00101 \end{matrix}$	
$A = 0 \ 0 \ 0 \ 0 \ 0$		$A' = 0 \ 0 \ 1 \ 0 \ 1$
$B = 1 \ 1 \ 1 \ 1 \ 1$		$B' = 1 \ 1 \ 0 \ 1 \ 0$

mask 由亂數產生，看其為 1 的地方
交換數字

以上的※2, ※4 等等皆由亂數產生

3. mutation

① bit mutation

	$\begin{matrix} \text{bit} \\ \text{mutation} \\ \text{※4} \end{matrix}$	
$A = 0 \ 1 \ 0 \ 1 \ 0$		$A' = 0 \ 1 \ 0 \ 0 \ 0$

②

string mutation

$$A = 0 \ 1 \ 0 \ 1 \ 0 \rightarrow A' = 1 \ 0 \ 1 \ 0 \ 1$$

全部倒反

③

uniform mutation

mask

$$A = 0 \ 1 \ 0 \ 1 \ 0 \xrightarrow{\begin{matrix} 0 & 0 & 1 & 0 & 1 \\ *3 & *5 \end{matrix}} A' = 0 \ 1 \ 1 \ 1 \ 1$$

example: $f(x) = x^2$, $0 \leq x \leq 31$ (32個可能)

find x such that $f(x)$ is maximum

(求極值，一次微分等於 0)

基因演算法解之，先進行編碼

$$0 \ 0 \ 0 \ 0 \ 0 \rightarrow 0$$

$$0 \ 0 \ 0 \ 0 \ 1 \rightarrow 1$$

$$\begin{array}{ccccc} & & \vdots & & \\ 1 & 1 & 1 & 1 & 1 \end{array} \rightarrow 31$$

let
population size = 4

step 1 : initial generation

*1 : 01110 ← 14

*2 : 11000 ← 24

*3 : 10001 ← 17

*4 : 00011 ← 7

encode

fitness function
(fitness value)

$$f_{fit}(x) = f(x) = x^2$$

= 2f(x) 可有多種 function

越符合答案的值，它的
fitness value 便越大

找出 fitness value

14: 196

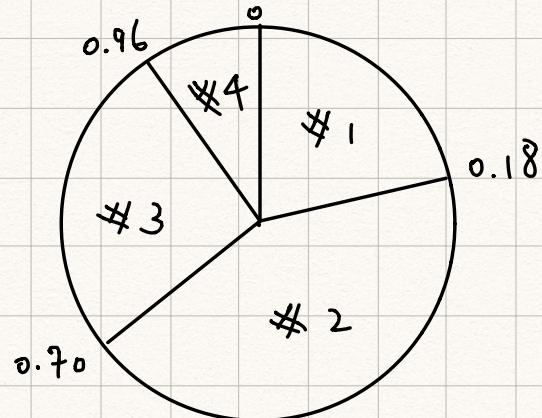
24: 576

17: 289

7: 49

統計 + %, 將各數所佔比
例得出

step 2 : reproduction (RWS)



按百分比將※N 畫入
- 圓餅圖中

照百分比
取和 population 一樣的數量

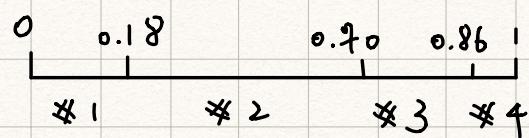
① 01110 (*1)

② 11000 (*2)

③ 11000 (*2)

④ 10001 (*3)

程式的寫法



0~1 隨意取值，取到何

值便是選到它惹、

Step 3: crossover

Let

$$\text{crossover rate} = 1 (100\%)$$

$0 < c.r. \leq 1$, 表 population 中有多少比例才能進行交配。有些人會和下一代交配

for large population(100)

$$c.r. = 0.6$$

(60) + 40, 60 丟入交配池進行計算

for small population(30)

$$c.r. = 0.9$$

一定都是兩兩交配

bit #2

② 1 1 0 0 0

① 0 1 1 1 0

②' 1 1 1 1 0 30 900

①' 0 1 0 0 0 8 64

取到 0.18 為 #1 依此
0.70 為 #2 類推

Step 4: mutation

Let

$$\text{mutation rate} = 0.01$$

$$5 \times 4 = 20 \text{ bit/gen}$$

每-generation 有 20 個 bit

$$20 \times 0.01 = 0.2 \text{ bit/gen}$$

不足 1. 怎辦？，使之等於 1 時再突變即可
(5th generation)

*3
隨機得到 0 1 0 0 0



0 1 1 0 0

以單點突變之

突變後的人要取代原

b.i.t * 4

④ 1 0 0 0 1

③ 1 1 0 0 0

④' 1 0 0 0 0 16 256

③' 1 1 0 0 1 25 625

本人的位置

突變用意在於突破
區域或最佳解的狀況

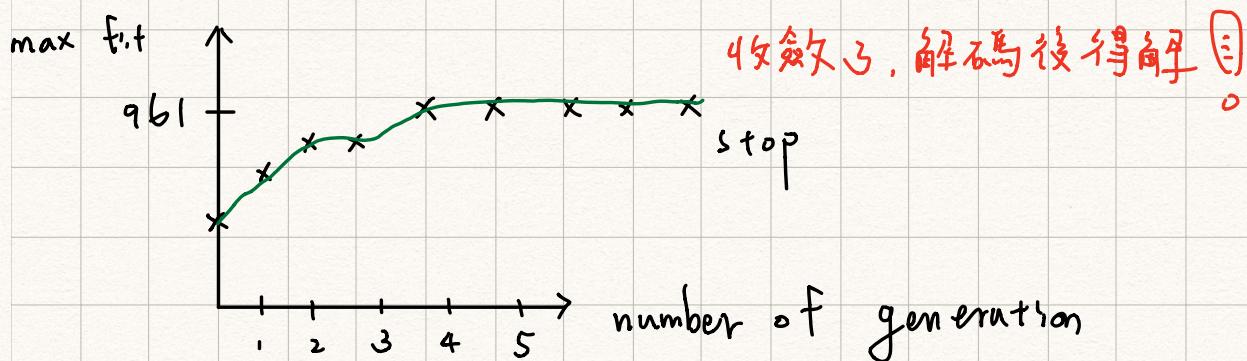
→ 基因演算化
具備全土或最佳解能力



做題目時要畫出 fitness value 的趨勢圖

譬如找出每個 generation 中的 Max 值

隨世代變換的結果 or 平均



不一定要找出全土或最佳解，若能允許
誤差，則不必尋找最佳解

以實際測試來判斷此解是否OK
(若客戶、老闆皆可，那此解也不錯)

Δ. Characteristics of GAs :

1. search encoded parameters , not the parameters .
2. avoid local optimum .
3. use fitness function to determine the goodness .
4. use probability method instead of crisp rules .

Δ. Affecting factors of GAs:

1. parameter setting
 - ① coding range
 - ② string length
 - ③ population size
 - ④ crossover rate
 - ⑤ mutation rate
 - ⑥ fitness function

2. encoding & decoding

$$X_1: 00011 \quad \frac{3}{31}$$

更多位元

$$X_2: 01100 \quad \frac{12}{31}$$

→ 解析度增加
→ 答案越精確

$$X_3: 00001 \quad \frac{1}{31}$$

$$f(x_1, x_2, x_3) = (x_1 - x_2)^2 + (x_3 + x_2)^2$$

find x_1, x_2, x_3 such that f is maximum, $0 \leq x \leq 1$

3. string length

① large length has higher accuracy, but takes more computation time

② length depends on the required accuracy

4. crossover rate

① higher rate may cause better

chromosomes to be changed via crossover process (好 x 上壠 → 變壞了)

② lower rate may show down search process

5. mutation rate

mutation rate must be small, otherwise it will become random search

1 1 1 0 1

1 1 1 0 1
1 1 1 0 1

1 1 1 0 1

→ 交配後極為定值

6. avoidance of (local) optimum

① annealing crossover rate and mutation rate (BP: crossover rate 和 mutation rate 不維持一致)

② increasing c.r. and m.r. while evolution process stop

7. Fitness Function

- ① problem dependent (= 題目敘述)
- ② increasing the difference between best chromosome and better chromosome

$31 \rightarrow 961$

$30 \rightarrow 900$

$[f(x)]^2$, 加大 range 使之明顯

停止條件：population 長得一樣

or 誤差在範圍內 (ex: 0.05%)

精度要高，編碼數要變長 → 用實數來做？

△. Real - Valued GAs

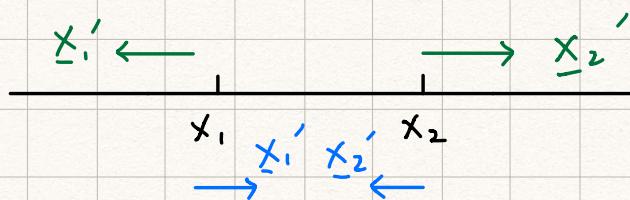
Chromosome type: $\underline{x} = \begin{bmatrix} x_1 \\ x_2 \\ \vdots \\ x_p \end{bmatrix}$

$x_1 \sim x_p$ 是不同的變量

Reproduction (same as binary GAs)

Crossover: for $x_1 < x_2$, $0 < \sigma < 1$

$$\left\{ \begin{array}{l} \underline{x}_1' = \underline{x}_1 + \sigma(\underline{x}_2 - \underline{x}_1) \\ \underline{x}_2' = \underline{x}_2 - \sigma(\underline{x}_2 - \underline{x}_1) \end{array} \right. \quad \begin{array}{l} \text{↑ 異數越少，每次} \\ \text{不一樣} \end{array}$$



轉成 $-1 < \sigma < 1$
便可用一公式做即可

$$\left\{ \begin{array}{l} \underline{x}_1' = \underline{x}_1 + \sigma(\underline{x}_2 - \underline{x}_1) \\ \underline{x}_2' = \underline{x}_2 - \sigma(\underline{x}_2 - \underline{x}_1) \end{array} \right. \quad [\text{get closer}]$$

$$\begin{array}{c} \downarrow \quad \downarrow \\ A_1: 10101010 (170) \quad A_1': 100000010 (130) \\ \rightarrow \\ A_2: 00000000 (0) \quad A_2': 00101000 (40) \end{array}$$

$$\begin{array}{c} \downarrow \quad \downarrow \\ A_3: 10111110 (190) \quad A_3': 10101010 (170) \\ \rightarrow \\ A_4: 11101011 (235) \quad A_4': 11111111 (225) \end{array}$$

觀察 binary 的 crossover, 出來的結果不是靠近就是拉遠, 因此實數公式也是如此做的

Mutation:

$$A_1: 101\underline{0}1010 \text{ (170)} \rightarrow A'_1: 101\underline{1}1010 \text{ (186)}$$

$$A_2: \underline{10101010} \text{ (170)} \rightarrow A'_2: \underline{01010101} \text{ (85)}$$

突變後, 不是變大就是變小

$$\underline{x}' = \underline{x} + s \times \text{random_noise(a vector)}$$

noise $\in \{-1, +1\}$, s is positive value
隨 x 值大小選擇
problem dependent

A. Evolutionary Algorithm (EA)

$$\underline{x} = \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ x_p \end{pmatrix}$$

Reproduction (same as binary GAs)

crossover: \underline{x}_j and \underline{x}_k are parents

① average crossover ② convex combination

$$\underline{x} = \frac{1}{2} [\underline{x}_j + \underline{x}_k]$$

$$\underline{x} = r \underline{x}_j + (1-r) \underline{x}_k$$

$$r \in (0, 1)$$

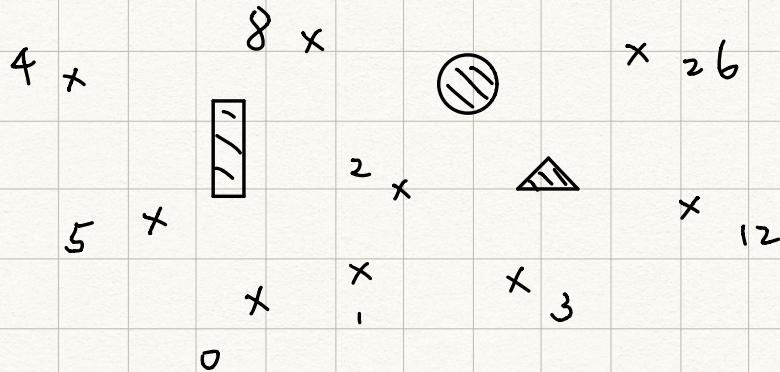
random

mutation:

$$\underline{x}' = \underline{x} + r \underline{d}$$

random

<Ex> Shortest Path



從 0 走到 6，中間有許多障礙物
每個節點為一種走法

兩者需有重複 (15)

crossover

parent 1: 0 - 1 - 2 - 11 - 15 - 16 - 21 - 24 - 26

parent 2: 0 - 4 - 8 - 10 - 15 - 20 - 23 - 26



... 1': 0 - 1 - 2 - 11 - 15 - 20 - 23 - 26

... 2': 0 - 4 - 8 - 10 - 15 - 16 - 21 - 24 - 26

如果發生 0 - 1 - 3 - 9 - 11 - 15 - 10 - 7 - 11 - 26

11 ~ 11 間的 15, 10, 7
要刪掉, 11, 11 合併

mutation



0 - 1 - 2 - 3 - 9 - 11 - 15 - 20 - 23 - 26

0 - 1 - 2 - 3 - 9 - 11 - 14 - 22 - 16

隨機抓取一個, 其後 random

fitness function 設定: $D = d + d + \dots$ (路徑和)

fitness = D^{-1} (越好的值越大)

<Ex> 最佳空制序列

$$x(k+1) = Ax(k) + Bu(k)$$

$$Y = C\underline{x}(k) + Du(k)$$

$u(0), u(1), u(2), \dots, u(k)$?

$$e = Y_d - Y$$

蝶馬主義演算法 · PSO