SAPIENZA
UNIVERSITÀ DI ROMA

IDK something about proofs, search problems, circuits, protocols and P $\neq$ NP?

Faculty of Information Engineering, Computer Science and Statistics
Bachelor's Degree in Computer Science

Simone Bianco
ID number 1986936

Advisor
Prof. Nicola Galesi
Prof. Massimo Lauria

Academic Year 2023/2024

Thesis not yet defended

IDK something about proofs, search problems, circuits, protocols and P $\neq$ NP?
Bachelor's Thesis. Sapienza University of Rome

This thesis has been typeset by LaTeX and the Sapthesis class.

Author's email: bianco.simone@outlook.it

Please help me.

Abstract

Uhhh, idk

Contents

Chapter 1

Notes

The following definitions and proofs are a reformulation of the results shown in [BFI23; GHJ+22]

Definition 1. *A total query search problem is a sequence of relations* $R = \{R_i : R_i \subseteq \{0,1\}^n \times O_i, i \in [n]\}$*, where* O_i *are finite sets called outcome sets, such that* $\forall x \in \{0,1\}^n$ *there is a* $o \in O_i$ *for which* $(x, o) \in R_n$*.*

A total search problem R *is in* TFNP^{dt} *if its solutions are verifiable through decision trees: for each* $o \in O_n$ *there is a decision tree* T_o *with* $\mathrm{poly}(\log n)$*-depth such that* $T_o(x) = 1 \iff (x, i) \in R_n$

While total search problems are formally defined as sequences $R = \{R_1, \ldots, R_n\}$, it will often make sense to speak of an individual search problem R_i in the sequence. Therefore, we will slightly abuse the notation and also call R_i a total search problem.

Definition 2. *Given* $R \subseteq \{0,1\}^n \times O$ *and* $S \subseteq \{0,1\}^m \times O'$*, a decision tree reduction from* R *to* S *is a set of decision trees* $T_i : \{0,1\}^n \to \{0,1\}$ *and* $T'_o : \{0,1\}^n \to O$*, for each* $i \in [m]$ *and each* $o \in O'$*, such that*

$$\forall x \in 0,1^n \;\; ((T_1(x), \ldots, T_m(x)), o) \in S \implies (x, T'_o(x)) \in R$$

To give an easier intuition of a decision tree reduction, the decision trees T_i map inputs from Q to R, while the decision trees T'_o map solutions to R back into solutions of Q.

The *size* s of the reduction is the number of input bits to S, meaning that $s = m$, while the *depth* d of the reduction is the maximum depth of any tree involved in the reduction

$$d := \max(\{D(T_i) : i \in [m]\} \cup \{D(T'_o) : o \in O_m\})$$

Finally, we define the *complexity* of the reduction as $\log s + d$. Moreover, the denote as $R^{dt}(S)$ the minimum complexity of any decision tree reduction from R to S.

Using the previous definition, we can define complexity classes of total query search problems via decision tree reductions. Given a total query search problem $S = (S_n)$, we define the subclass of problems reducible to S as:

$$S^{dt} := \{R : S^{dt}(R) = \mathrm{poly}(\log n)\}$$

where $R = (R_n)$.

Any total query search problem with solution verifiers T_o for each $o \in O$ can be *encoded* into a canonical unsatisfiable CNF formula.

Proposition 1. *Given a total query search problem $R \subseteq \{0,1\}^n \times O$, there exists an unsatisfiable* CNF *formula F defined on $|O|$-many variables such that $R = \mathsf{Search}(F)$.*

This formula is called the canonical CNF *encoding of R.*

Proof. Since $R \in \mathsf{TFNP}^{dt}$, for each $o \in O$ there is a $\mathrm{poly}(\log n)$-depth decision tree T_o that verifies R. Then, for each T_o, let C_o be the clause obtained by taking the disjunction over the conjunction of the literals along each of the accepting paths in T_o, meaning that C_o is a DNF and, by De Morgan's theorem, that $\overline{C_o}$ is a CNF.

Let $F := \bigwedge_{o \in O} \overline{C_o}$. Since each R is a total search problem, for each input there is a valid output, implying that at least one tree T_o will have an accepting path. Hence, by definition of C_o, we get that $\overline{C_o} = 0$, implying that there will always be a false clause in F and thus that F is an unsatisfiable CNF, concluding that:

$$(x, o) \in R \iff T_o(x) = 1 \iff \overline{C_o} = 0 \iff (x, o) \in \mathsf{Search}(F)$$

□

This result implies that black-box TFNP is *exactly* the study of the false clause search problem. Then, instead of studying a total search problem R, it's sufficient to study the search problem $\mathsf{Search}(F)$ associated with the canonical CNF encoding F of R.

Given a proof system P and an unsatisfiable CNF formula F, we define the *complexity required by P to prove F*, denoted with $P(F)$, as:

$$P(F) := \min(\{\deg(\Pi) + \log \mathrm{size}(\Pi) : \Pi \text{ is a } P\text{-proof of } F\})$$

where deg is the *degree* of the proof, which is a measure defined by the proof system itself. For example, for the Resolution proof system the degree is defined as the *width* of the proof, which is the maximum number of literals in any clause in Π.

Definition 3. *Given $R \in \mathsf{TFNP}^{dt}$, we say that R characterizes and a proof system P (and that P characterizes R) if it holds that $R^{dt} = \{\mathsf{Search}(F) : P(F) = \mathrm{poly}(\log n)\}$.*

Tree-like resolutions for an unsatisfiable CNF formula are strictly connected to the decision trees that solve its associated search problem. In particular, it can be proven that the smallest tree-like refutation has the exact same structure of the smallest decision tree.

Lemma 1. *[BGL13] Let F be an unsatisfiable* CNF *formula. If there is a tree-like refutation of F with structure T, there also exists a decision tree with structure T that solves* Search(F)

Proof. We procede by induction on the size s of the refutation of F.

Let $F = C_1 \wedge \ldots \wedge C_m$. If $s = 1$, then the refutation is made up of only one step that ends with the empty clause, implying that $\exists i \in [m]$ such that $F = C_i = \bot$. Hence, Search(F) can be solved by the decision tree made of only one vertex labeled with i.

We now assume that every formula with a tree-like refutation with a structure of size s there exists a decision tree with the same structure that solves the search problem associated with the formula.

Suppose now that the size s of the refutation is bigger than 1. Let x be the last variable resolved by the refutation and let T_0 and T_1 be the subtrees of T such that x is the root of T_0 and $\overline{x}$ is the root of T_1.

Consider now the formulas $F\restriction_{x=0}$ and $F\restriction_{x=1}$, respectively corresponding the formula F with the value 0 or 1 assigned to x. It's easy to see that the subtrees T_0 and T_1 are valid refutations of the formulas $F\restriction_{x=0}$ and $F\restriction_{x=1}$: if $b = 0$, then x evaluates to 0, otherwise if $b = 1$ then $\overline{x}$ evaluates to 0.

Since T_0 and T_1 have size $s - 1$, by inductive hypothesis there exist two decision tree with structure T_0 and T_1 that solve Search$(F\restriction_{x=0})$ and Search$(F\restriction_{x=1})$.

Finally, the search problem Search(F) can be solved by the decision tree that queries x and proceeds with the decision tree T_b based on the value $b \in \{0, 1\}$ such that $x = b$.

□

Definition 4. *Given two rooted trees T and T', we say that T is embeddable in T' if there exists a mapping $f : V(T) \to V(T')$ such that, for any vertices $u, v \in V(T)$, if u is a parent of v in T then $f(u)$ is an ancestor of $f(v)$ in T'.*

Lemma 2. *[BGL13; LNN+95] Let F be an unsatisfiable* CNF *formula. If there is a decision tree with structure T that solves* Search(F)*, there also exists a tree-like refutation of F with structure T' such that T' is embeddable in T.*

Proof. The main idea is to associate inductively, starting from the leaves, a clause to each vertex of T in order to transform T in a tree-like refutation of F. In particular, each vertex v gets associated to a clause $C(v)$ such that every input of the decision tree that reaches v falsifies $C(v)$.

Let $F = C_1 \wedge \ldots \wedge C_m$. For all $i \in [m]$, we associate the clause C_i to the leaf of T labeled with i. This constitutes our base case.

Consider now a vertex v that isn't a leaf. Let x be the variable that labels v and let u_0, u_1 be the vertices such that the edge (v, u_0) is taken if $x = 0$ and the edge (v, u_1) is taken if $x = 1$. By induction, assume that u_0 and u_1 have already been associated with the clauses C_0 and C_1.

By way of contradiction, suppose that C_0 contains the literal $\overline{x}$. Then, since in a decision tree each variable can be queried only once in every path, there will always be an input with $x = 0$ that reaches v. Since $x = 0$ and since C_0 contains $\overline{x}$, this input would satisfy C_0, contradicting the fact that C_0 was associated to u_0 in a way that it is falsified by every input.

Thus, the only possibility is that C_0 can't contain the literal $\overline{x}$. Similarly, we can show that C_1 cant' contain the literal x. This leaves us with only two possibilities: either $C_0 = x \vee \alpha$ and $C_1 = \overline{x} \vee \beta$ or one of C_0, C_1 doesn't contain $x, \overline{x}$.

In the first case, we can simply associate to v the clause $C = \alpha \vee \beta$. In the second case, we associate to v the clause that doesn't contain $x, \overline{x}$ (chose any of them if both clauses do not contain $x, \overline{x}$).

In particular, we notice that the first case directly emulates the resolution rule, while the second case essentially represent "redundant steps". By "skipping" these redundant steps, we can obtain a tree T' that is embeddable in T and that contains only nodes on which the first case was applied. Finally, it's easy to deduce that the root node of T' will always be associated with the empty clause $\perp$, concluding that T' is the structure of a tree-like refutation of F.

□

Theorem 1. *Let F be an unsatisfiable* CNF *formula. The smallest tree-like refutation of F has size s and depth d if and only if the smallest decision tree solving* Search(F) *has size s and depth d.*

Proof. Let s and d be the size and depth of the smallest tree-like refutation of F. Likewise, let x and y be the size and depth of the smallest decision tree solving Search(F).

Then, by Lemma 1, we know that there exists a decision tree that solved Search(F) with the same structure of the smallest refutation. Let α and β be the size and depth of this decision tree. It's easy to see that $s = \alpha \geq x$ and $d = \beta \geq y$.

Viceversa, by the Lemma 2, we know that there exists a tree-like refutation of F such that its structure is embeddable in the one of the smallest decision tree. Let γ and β be the size and depth of this tree-like refutation. Since the latter is embedded in the smallest decision tree, it's structure must be smaller or equal. Hence, it's easy to see that $x \geq \gamma \geq s$ and $y \geq \delta \geq d$. Thus, we can conclude that $s = x$ and $d = y$.

□

Note: the theorem should be generalizable to each tree and not only for the smallest trees.

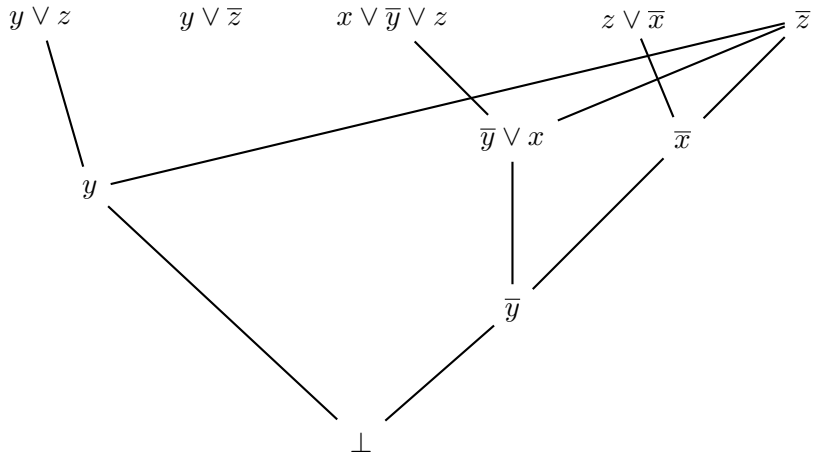

Figure 1.1. Dag-like refutation of the previous formula

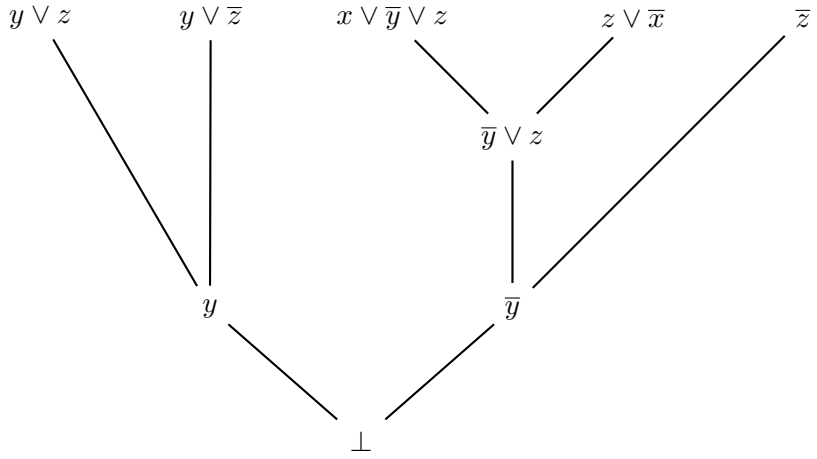

Figure 1.2. Tree-like refutation of the previous formula

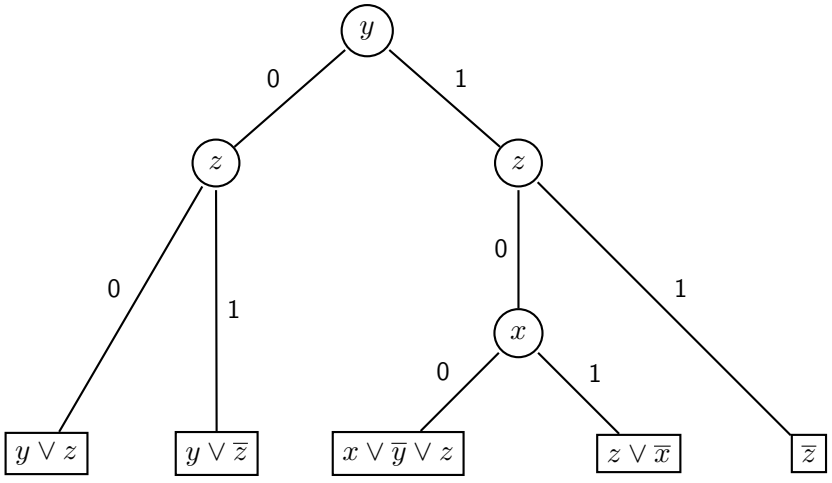

Figure 1.3. Decision tree for the previous formula

Acknowledgements

this is sad, alexa play despacito

Bibliography

[BFI23] Sam Buss, Noah Fleming, and Russell Impagliazzo. "TFNP Characterizations of Proof Systems and Monotone Circuits". In: *14th Innovations in Theoretical Computer Science Conference (ITCS 2023)*. Ed. by Yael Tauman Kalai. Vol. 251. Leibniz International Proceedings in Informatics (LIPIcs). Dagstuhl, Germany: Schloss Dagstuhl – Leibniz-Zentrum für Informatik, 2023, 30:1–30:40. ISBN: 978-3-95977-263-1. DOI: 10.4230/LIPIcs.ITCS.2023.30. URL: https://drops-dev.dagstuhl.de/entities/document/10.4230/LIPIcs.ITCS.2023.30.

[BGL13] Olaf Beyersdorff, Nicola Galesi, and Massimo Lauria. "A characterization of tree-like Resolution size". In: *Information Processing Letters* 113.18 (2013), pp. 666–671. ISSN: 0020-0190. DOI: https://doi.org/10.1016/j.ipl.2013.06.002. URL: https://www.sciencedirect.com/science/article/pii/S0020019013001592.

[GHJ+22] Mika Göös, Alexandros Hollender, Siddhartha Jain, et al. "Separations in Proof Complexity and TFNP". In: *2022 IEEE 63rd Annual Symposium on Foundations of Computer Science (FOCS)*. 2022, pp. 1150–1161. DOI: 10.1109/FOCS54457.2022.00111.

[LNN+95] László Lovász, Moni Naor, Ilan Newman, et al. "Search Problems in the Decision Tree Model". In: *SIAM J. Discret. Math.* 8.1 (Feb. 1995), pp. 119–132. ISSN: 0895-4801. DOI: 10.1137/S0895480192233867. URL: https://doi.org/10.1137/S0895480192233867.