

# Differential Cryptanalysis

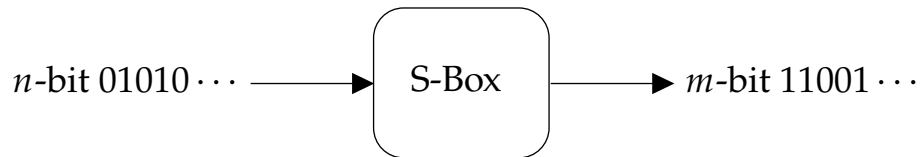
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Ji, Yong-hyeon

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## 1 Definitions

### 1.1 S-Box



#### S-Box

**Definition 1.** Let  $n, m \in \mathbb{Z}^+$ . A function

$$S : \mathbb{F}_2^n \rightarrow \mathbb{F}_2^m$$

is a **S-Box**.

### 1.2 The XOR operation

$$\begin{aligned} \oplus & : \mathbb{F}_2 \times \mathbb{F}_2 \longrightarrow \mathbb{F}_2 \\ (x, y) & \longmapsto z = x + y \bmod 2 \end{aligned}$$

$x$	$y$	$z = x \oplus y$
0	0	0
0	1	1
1	0	1
1	1	0

**Note** (Thinking).

$$\begin{aligned} \oplus & : \mathbb{F}_2 \longrightarrow [\mathbb{F}_2 \rightarrow \mathbb{F}_2] \\ x & \longmapsto \oplus_x = \begin{cases} \text{Id}(y) & : x = 0, \\ \neg(y) & : x = 1. \end{cases} \end{aligned}$$

$$\begin{aligned} \oplus_n : \quad \mathbb{F}_2^n \times \mathbb{F}_2^n &\longrightarrow \mathbb{F}_2^n \\ \left( \{x_i\}_{i=1}^n, \{y_i\}_{i=1}^n \right) &\longmapsto \{x_i \oplus y_i\}_{i=1}^n \end{aligned}$$

**Note.** We use the notation  $0_n$  to denote  $0_n = (0, \dots, 0)$ .

**Note (Thinking).**

$$\begin{aligned} \oplus_n : \quad \mathbb{F}_2^n &\longrightarrow [\mathbb{F}_2^n \rightarrow \mathbb{F}_2^n] \\ \{x_i\}_{i=1}^n &\longmapsto (\oplus_n)_{\{x_i\}_{i=1}^n} = \{z_i\}_{i=1}^n, \text{ where } z_i = \begin{cases} \text{Id}(y_i) & : x_i = 0, \\ \neg(y_i) & : x_i = 1. \end{cases} \end{aligned}$$

**Proposition 1.** Let  $X, Y, Z \in \mathbb{F}_2^n$ . Then

- (1)  $X \oplus_n Y = Y \oplus_n X$
- (2)  $(X \oplus_n Y) \oplus Z = X \oplus_n (Y \oplus Z)$
- (3)  $X \oplus_n 0_n = X = 0_n \oplus_n X$
- (4)  $X \oplus_n X = 0_n$
- (5)  $X \oplus_n Y = 0_n \implies X = Y$
- (6)  $A \oplus_n X = B \implies X = A \oplus_n B$

**Note.** By (4) and (5), we have  $X \oplus_n Y = 0 \iff X = Y$ .

*Proof.* PASS □

**Remark 1.**

- The binary operation  $\oplus$  provides the structure of an **abelian group** on the set  $\mathbb{F}_2^n$  with identity element  $0_n$ .
- Because of the property (4)  $X \oplus_n X = 0_n$ , we see that *the inverse of any element is itself* with respect to the operation  $\oplus$ .

**Definition 2.** The **difference** of  $X \in \mathbb{F}_2^n$  and  $Y \in \mathbb{F}_2^n$  is defined as  $X \oplus_n Y \in \mathbb{F}_2^n$ .

## 2 Difference Set

### 2.1 Definition and Property

#### Difference Set of Bit-Sequence

**Definition 3.** Given  $\alpha \in \mathbb{F}_2^n$ , we define the **difference set** of  $\alpha$  as follow:

$$\Delta_\alpha = \{(x_1, x_2) \in \mathbb{F}_2^n \times \mathbb{F}_2^n : x_1 \oplus x_2 = \alpha \in \mathbb{F}_2^n\} \subseteq \mathbb{F}_2^n \times \mathbb{F}_2^n.$$

**Proposition 2.** For any  $\alpha \in \mathbb{F}_2^n$  the set  $\Delta_\alpha$  contains  $2^n$  elements and can be expressed as

$$\Delta_\alpha = \{(x, x \oplus \alpha) : x \in \mathbb{F}_2^n\}.$$

*Proof.* Let

$$S := \{(x_1, x_2) \in \mathbb{F}_2^n \times \mathbb{F}_2^n : x_1 \oplus x_2 = \alpha \in \mathbb{F}_2^n\},$$

$$T := \{(x, x \oplus \alpha) : x \in \mathbb{F}_2^n\}.$$

We must show that  $S = T$ :

$(S \subseteq T)$  Let  $(x, y) \in S$  then by definition  $x \oplus y = \alpha$ . Since  $(x \oplus y = \alpha) \Rightarrow (y = x \oplus \alpha)$ ,

$$(x, y) = (x, x \oplus \alpha) \in T.$$

$(T \subseteq S)$  Let  $(x, x \oplus \alpha) \in T$ . Since

$$x \oplus (x \oplus \alpha) = \alpha,$$

$$(x, x \oplus \alpha) \in S.$$

□

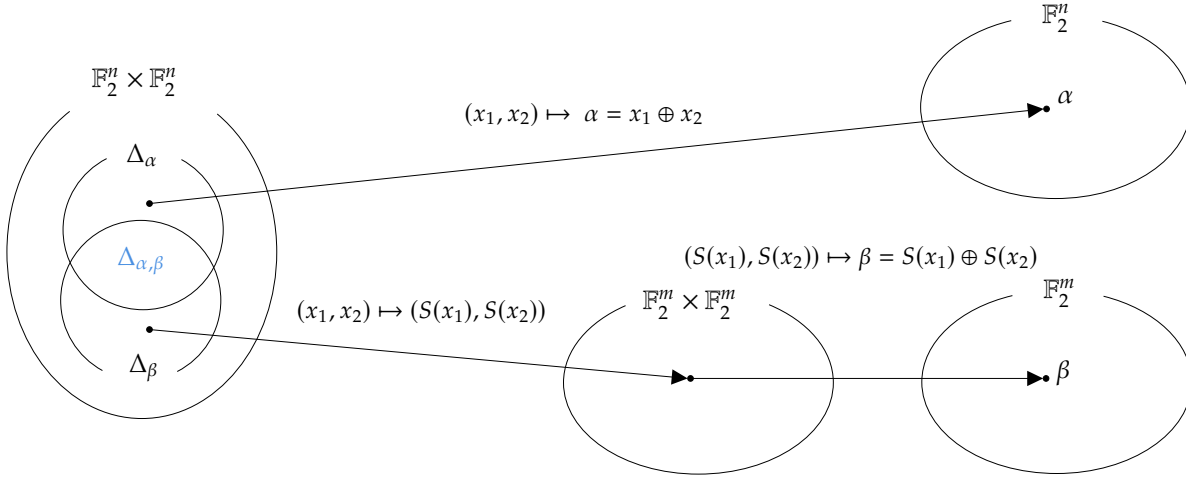
**Corollary 2.1.** For any  $\alpha \in \mathbb{F}_2^n$ , we have  $\Delta_\alpha \simeq \mathbb{F}_2^n$ .

**Remark 2.** Let us consider the case  $\alpha = 0$  for the set  $\Delta_\alpha$ . When  $\alpha = 0$  the difference set is

$$\Delta_0 = \{(x, x) : x \in \mathbb{F}_2^n\}$$

This set is often called the **diagonal** of  $\mathbb{F}_2^n \times \mathbb{F}_2^n$ .

## 2.2 Difference Sets of a S-BOX



### Difference Set of a S-BOX

Let  $S : \mathbb{F}_2^n \rightarrow \mathbb{F}_2^m$  is a S-box. Let  $\alpha \in \mathbb{F}_2^n$  and  $\beta \in \mathbb{F}_2^m$ . Consider

$$\Delta_\alpha = \{(x_1, x_2) : x_1 \oplus x_2 = \alpha \in \mathbb{F}_2^n\} \subseteq \mathbb{F}_2^n \times \mathbb{F}_2^n \quad \text{and}$$

$$\Delta_\beta = \{(x_1, x_2) : S(x_1) \oplus S(x_2) = \beta \in \mathbb{F}_2^m\} \subseteq \mathbb{F}_2^n \times \mathbb{F}_2^n.$$

We define the **difference set** of  $S$  with respect to  $\alpha$  and  $\beta$  by

$$\Delta_{\alpha,\beta} = \Delta_\alpha \cap \Delta_\beta = \{(x_1, x_2) \in \mathbb{F}_2^n \times \mathbb{F}_2^n : x_1 \oplus x_2 = \alpha \in \mathbb{F}_2^n \text{ and } S(x_1) \oplus S(x_2) = \beta \in \mathbb{F}_2^m\}.$$

That is,  $\Delta_{\alpha,\beta}$  is the set of ordered pairs of elements from  $\mathbb{F}_2^n$  which have a difference of  $\alpha$  and such that their images under  $S$  have a difference of  $\beta$ .

### Remark 3.

- This can also written as

$$\Delta_{\alpha,\beta} = \{(x_1, x_2) \in \Delta_\alpha : (S(x_1), S(x_2)) \in \Delta_\beta\}.$$

- $\Delta_{\alpha,\beta}$  is always defined w.r.t. a given S-Box  $S$ . If we want to make this dependence explicit we can write  $\Delta_{\alpha,\beta}^S$ .

### Cardinality of a Difference Set

We define  $\delta_{\alpha,\beta}$  to be the cardinality of the finite set  $\Delta_{\alpha,\beta}$ , namely

$$\delta_{\alpha,\beta} := |\Delta_{\alpha,\beta}| \in \mathbb{Z}_{\geq 0}.$$

### Proposition 3.

$$\delta_{\alpha,\beta} = \# \{x \in \mathbb{F}_2^n : S(x) \oplus S(x \oplus \alpha) = \beta \in \mathbb{F}_2^m\}.$$

*Proof.*

$$\begin{aligned} \Delta_{\alpha,\beta} &= \Delta_\alpha \cap \Delta_\beta = \{(x_1, x_2) : x_1 \oplus x_2 = \alpha\} \cap \{(x_1, x_2) : S(x_1) \oplus S(x_2) = \beta\} \\ &= \{(x, x \oplus \alpha) : x \in \mathbb{F}_2^n\} \cap \{(x_1, x_2) : S(x_1) \oplus S(x_2) = \beta\} \\ &= \{x \in \mathbb{F}_2^n : S(x) \oplus S(x \oplus \alpha) = \beta \in \mathbb{F}_2^m\}. \end{aligned}$$

□

### Remark 4.

- When  $\alpha = 0$  and  $\beta = 0$  we have  $\Delta_{0,0} = \Delta_0 = \{(x, x) : x \in \mathbb{F}_2^n\}$ .
- In general when  $\alpha = 0$  we find that  $\Delta_{0,\beta} = \begin{cases} \Delta_0 & : \beta = 0 \\ \emptyset & : \beta \neq 0 \end{cases}$
- Since  $|\Delta_0| = 2^n$  and  $|\emptyset| = 0$ ,  $\delta_{0,\beta} = \begin{cases} 2^n & : \beta = 0 \\ 0 & : \beta \neq 0 \end{cases}$ .

**Proposition 4.** *The integer  $\delta_{\alpha,\beta} \in \mathbb{Z}_{\geq 0}$  is always even.*

*Proof.* Recall that 0 is even.

(Case I) When  $\alpha = 0$ , we saw either  $\delta_{0,\beta} \in \{0, 2^n\}$  and these are even in either case.

(Case II) Suppose that  $\alpha \neq 0$  and  $\Delta_{\alpha,\beta} \neq \emptyset$ . Let  $(x_1, x_2) \in \Delta_{\alpha,\beta}$  then

$$(x_2, x_1) \in \Delta_{\alpha,\beta} \quad \text{and} \quad x_1 \neq x_2.$$

Therefore  $(x_1, x_2) \neq (x_2, x_1)$ . So if we pair  $(x_1, x_2)$  and  $(x_2, x_1)$ , we can partition  $\Delta_{\alpha,\beta}$  into subsets, each subset having cardinality of 2.

□

### 3 The DDT of a S-Box

#### 3.1 Definition and Property

##### Differential Distribution Table

Let  $S : \mathbb{F}_2^n \rightarrow \mathbb{F}_2^m$  be a S-Box. The **differential distribution table** (abbreviated DDT) of  $S$  is a table (or matrix) with  $2^n$ -rows and  $2^m$ -columns. We denote it by  $\mathcal{D}_S$  or just by  $\mathcal{D}$ .

- The rows are indexed by the elements  $\alpha \in \mathbb{F}_2^n = \{0, \dots, 2^n - 1\}$ .
- The columns are indexed by the elements  $\beta \in \mathbb{F}_2^m = \{0, \dots, 2^m - 1\}$ .
- The entry at row index  $\alpha$  and column index  $\beta$  is given by  $\delta_{\alpha,\beta} = |\Delta_{\alpha,\beta}|$ . That is,

$$\mathcal{D} = (\delta_{\alpha,\beta})_{2^n \times 2^m}.$$

**Remark 5.** The DDT of a S-Box is just table of all the possible integer values  $\delta_{\alpha,\beta}$ .

$\mathcal{D}$	0	1	...	$\beta$	...	$2^m - 1$
0	$2^n$	0		...		0
1						
$\vdots$						
$\alpha$				$\delta_{\alpha,\beta}$		
$\vdots$						
$2^n - 1$						

**Proposition 5.** Let  $S : \mathbb{F}_2^n \rightarrow \mathbb{F}_2^m$  be a S-Box with differential distribution table  $\mathcal{D}$ . The following properties hold for  $\mathcal{D}$ .

- (1) Every entry in  $\mathcal{D}$  is a non-negative even integer between 0 and  $2^n$ .
- (2) The top-left entry of  $\mathcal{D}$  is  $2^n$ .
- (3) The first row of  $\mathcal{D}$  consists of all zeros except the first entry which is  $2^n$ .
- (4) If  $S$  is one-to-one then the first column of  $\mathcal{D}$  consists all zeros except the first entry which is  $2^n$ .
- (5) The sum of the entries of each row is  $2^n$ .
- (6) If  $S$  is bijective then every row and column of the DDT add up to  $2^n$ .

### 3.2 The DDT and Probabilities

When designing and analyzing attacks by differential cryptanalysis, we often want to translate the integer values from the DDT into probabilities.

#### Probability of DDT

**Definition 4.** Let the universe is  $\Omega = \Delta_\alpha$  and the event  $E = \Delta_{\alpha,\beta}$ . We define probability  $P[\alpha \rightarrow \beta] = p_{\alpha,\beta}$  as follows:

$$P[\alpha \rightarrow \beta] = p_{\alpha,\beta} = \frac{|E|}{|\Omega|} = \frac{|\Delta_{\alpha,\beta}|}{|\Delta_\alpha|} = \frac{\delta_{\alpha,\beta}}{2^n}.$$

**Remark 6.** Equivalently we can regard the universe  $\Omega = \mathbb{F}_2^n$  and the event as the set

$$E = \{x \in \mathbb{F}_2^n : S(x) \oplus S(x \oplus \alpha) = \beta\}.$$

### 3.3 The DDT of a Linear S-Box

#### Linear S-Box

**Definition 5.** A S-Box  $L : \mathbb{F}_2^n \rightarrow \mathbb{F}_2^m$  is said to be **linear** if

$$x_1, x_2 \in \mathbb{F}_2^n \implies L(x_1 \oplus x_2) = L(x_1) \oplus L(x_2).$$

**Proposition 6.** Let  $L : \mathbb{F}_2^n \rightarrow \mathbb{F}_2^m$  be a linear S-Box. Let  $\alpha \in \mathbb{F}_2^n$  and  $\beta \in \mathbb{F}_2^m$ . Then the difference sets of  $L$  and their cardinalities are given by

$$\Delta_{\alpha,\beta} = \begin{cases} \Delta_\alpha & : \beta = L(\alpha) \\ \emptyset & : \beta \neq L(\alpha) \end{cases}, \quad \delta_{\alpha,\beta} = \begin{cases} 2^n & : \beta = L(\alpha) \\ 0 & : \beta \neq L(\alpha) \end{cases}.$$

That is, every row of the DDT of  $L$  consists of all zeros except one entry with a value of  $2^n$ .

*Proof.* Let  $\alpha \in \mathbb{F}_2^n$ . Assume that  $(x_1, x_2) \in \Delta_\alpha$ , namely  $x_1 \oplus x_2 = \alpha$  then

$$L(x_1) \oplus L(x_2) = L(x_1 \oplus x_2) = L(\alpha).$$

So  $(x_1, x_2) \in \Delta_{\alpha,\beta} \iff \beta = L(\alpha).$  □

**Remark 7.**

- In the context of probabilities this states

$$p_{\alpha,\beta} = \begin{cases} 1 & : \beta = L(\alpha) \\ 0 & : \beta \neq L(\alpha) \end{cases}$$

- So the DDT of a linear S-Box is not interesting since every entry is either 0 or  $2^n$ .

**Note.**

- Note the DDT of a linear S-Box contains large values (namely,  $2^n$ , which is as large as an entry can be).
- This shows that any block cipher that uses a linear S-Box for its “confusion” layer would have a catastrophic vulnerability. It would likely be vulnerable to an attack by differential cryptanalysis.
- This is why the “confusion” layer is required to be non-linear, so we would never do this anyway.



### 3.4 The DDT of the XOR S-Box

- Let  $k \in \mathbb{F}_2^n$  be fixed and define the bijective S-Box:

$$\begin{aligned} S &: \mathbb{F}_2^n \longrightarrow \mathbb{F}_2^n \\ x &\longmapsto x \oplus k \end{aligned}$$

- Note that  $S(x_1 \oplus x_2) \neq S(x_1) \oplus S(x_2)$  when  $k \neq 0_n$
- Therefore when  $k \neq 0$  the S-Box  $S$  is non-linear
- Consider the equation for  $S$

$$S(x) \oplus S(x \oplus \alpha) = \beta$$

- By the definition of  $S$  this is equivalent to

$$(x \oplus k) \oplus ((x \oplus \alpha) \oplus k) = \beta$$

- Which is equivalent to  $\alpha = \beta$ .
- We find the following properties hold for the XOR S-Box

$$\Delta_{\alpha,\beta} = \begin{cases} \Delta_\alpha & : \alpha = \beta \\ \emptyset & : \alpha \neq \beta \end{cases} \quad \delta_{\alpha,\beta} = \begin{cases} 2^n & : \alpha = \beta \\ 0 & : \alpha \neq \beta \end{cases}$$



Department of Information Security, Cryptology and Mathematics  
College of Science and Technology  
Kookmin University

## References

- [1] "Cryptanalysis - L6 Differential Cryptanalysis" YouTube, uploaded by Maria Eichlseder, 22 April 2021, <https://www.youtube.com/watch?v=GQX8W8zKf2Q>
- [2] "The Differential Distribution Table (DDT) of a S-Box" YouTube, uploaded by JacksonInfoSec, 5 August 2021, [https://www.youtube.com/watch?v=di90\\_oL4lxx](https://www.youtube.com/watch?v=di90_oL4lxx)