

Homework assignment 4 – Symbolic Systems I – UvA, June 2020

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Question 1

TBox

Person

Country

City

Republic \sqsubseteq *Country*

ConstitutionalMonarchy \sqsubseteq *Country*

Republic \sqcap *ConstitutionalMonarchy* $\sqsubseteq \perp$

Person \sqcap *Country* $\sqsubseteq \perp$

City \sqcap *Country* $\sqsubseteq \perp$

Person \sqcap *City* $\sqsubseteq \perp$

Country $\sqsubseteq \exists \text{hasCapital.City}$

Country $\sqsubseteq \forall \text{hasCapital.City}$

Country $\sqsubseteq \exists \text{hasHead.Person}$

City $\sqsubseteq \exists \text{hasHead.Person}$

President \sqsubseteq *Person*

Monarch \sqsubseteq *Person*

Mayor \sqsubseteq *Person*

ConstitutionalMonarchy $\sqsubseteq \forall \text{hasHead.Monarch}$

City $\sqsubseteq \forall \text{hasHead.Mayor}$

ABox

NL : *Country*

NL : *ConstitutionalMonarchy*

(*NL*, *WillemAlexander*) : *hasHead*

(*Amsterdam*, *Femke*) : *hasHead*

Continued:

GCI's:

1. (*Republic* \sqcap *Country*) $\sqsubseteq \forall \text{hasHead.}\neg \text{Mayor}$

It may be noted that republic is already a subset of country. So the following is correct as well:

Republic $\sqsubseteq \forall \text{hasHead.}\neg \text{Mayor}$

No, this does not follow from the knowledge base.

Republic \sqsubseteq *Country*

Country $\sqsubseteq \forall \text{hasHead.Person}$

President \sqsubseteq *Person*

Monarch \sqsubseteq *Person*

Mayor \sqsubseteq *Person*

From this KB it follows that a republic country always has a head which is a person. This person could in this knowledge base be of the set President, Monarch and/or Mayor. To make this statement follow from the knowledge base the following statements could be added:

Republic $\sqsubseteq \forall \text{hasHead.President} \sqcup \text{Monarch}$

President \sqcap *Mayor* $\sqsubseteq \perp$

$Monarch \sqcap Mayor \sqsubseteq \perp$

2. *Femke* : *Mayor*

Yes, this follows from the knowledge base.

$(Amsterdam, Femke) : hasHead$

$City \sqsubseteq \forall hasHead.Mayor$

Therefore, Femke is in the set Mayor.

3. *NL* : $\neg Republic$

Yes, this follows from the knowledge base.

$NL : ConstitutionalMonarchy$

$Republic \sqcap ConstitutionalMonarchy \sqsubseteq \perp$

So, NL cannot be Republic.

4. *WillemAlexander* : $\neg Mayor$

No, this does not follow from the knowledge base.

$(NL, WillemAlexander) : hasHead$

$NL : ConstitutionalMonarchy$

$ConstitutionalMonarchy \sqsubseteq \forall hasHead.Monarch$

But, there is never mentioned that a Monarch cannot also be a Mayor.

Therefore, the following statement needs to be added:

$Monarch \sqcap Mayor \sqsubseteq \perp$

For consistency, the following statements could be added but this does not influence the statement *WillemAlexander* : $\neg Mayor$.

$Monarch \sqcap President \sqsubseteq \perp$

$Mayor \sqcap President \sqsubseteq \perp$

Question 2

Constructing a complex concept of a 3-CNF formula can be done by replacing every propositional variable x_1, \dots, x_n in ϕ to the concept names A_1, \dots, A_n .

So for instance if ϕ is :

$$\phi = (x_1 \vee x_2 \vee \neg x_3) \wedge (x_3 \vee \neg x_4 \vee x_5)$$

Then constructing this into a complex concept C_ϕ :

$$C_\phi \equiv (A_1 \sqcup A_2 \sqcup \neg A_3) \sqcap (A_3 \sqcup \neg A_4 \sqcup A_5)$$

Furthermore, we could set the first clause concept $C_1 \equiv A_1 \sqcup A_2 \sqcup \neg A_3$.

A 3-CNF formula is satisfiable if every clause in the formula is true. A clause in a 3-CNF formula is true if at least one literal, being x_i or $\neg x_i$, in the clause is true.

If for an interpretation $I = (\Delta^I, \cdot^I)$ x_i is true, then $A_i^I \neq \emptyset$. Furthermore, if in a clause concept one concept is non empty, then we can conclude that the entire clause is non empty: $C_1 \equiv A_1 \sqcup A_2 \sqcup \neg A_3$ and $A_1^I \neq \emptyset$ leads to $C_1^I \neq \emptyset$.

$$C_\phi \equiv (A_1 \sqcup A_2 \sqcup \neg A_3) \sqcap (A_3 \sqcup \neg A_4 \sqcup A_5) \equiv C_1 \sqcap C_2$$

For this example, $C_\phi \neq \emptyset$ iff both $C_1 \neq \emptyset$ and $C_2 \neq \emptyset$. $C_1 \neq \emptyset$ and $C_2 \neq \emptyset$ iff for both clauses one literal is true. So therefore, if the 3-CNF formula ϕ is satisfiable there exists an interpretation I for which $C_\phi^I \neq \emptyset$

Question 3

For all $n \in \{1, \dots, m\}$

KB_n

TBox

Add the TBox of all the KB_i , with $0 < i < n$

$$A_{n-1} \sqsubseteq \forall r. ((B_{n,a} \sqcap A_n) \sqcup (B_{n,b} \sqcap A_n))$$

$$\exists r. (B_{n,a} \sqcap A_n)$$

$$\exists r. (B_{n,b} \sqcap A_n)$$

$$B_{n,a} \sqcap B_{n,b} \sqsubseteq \perp$$

$$\forall r. B_{n,a}$$

$$\forall r. B_{n,b}$$

$$\underline{\text{ABox}}$$

$$x : A_0$$

Question 4