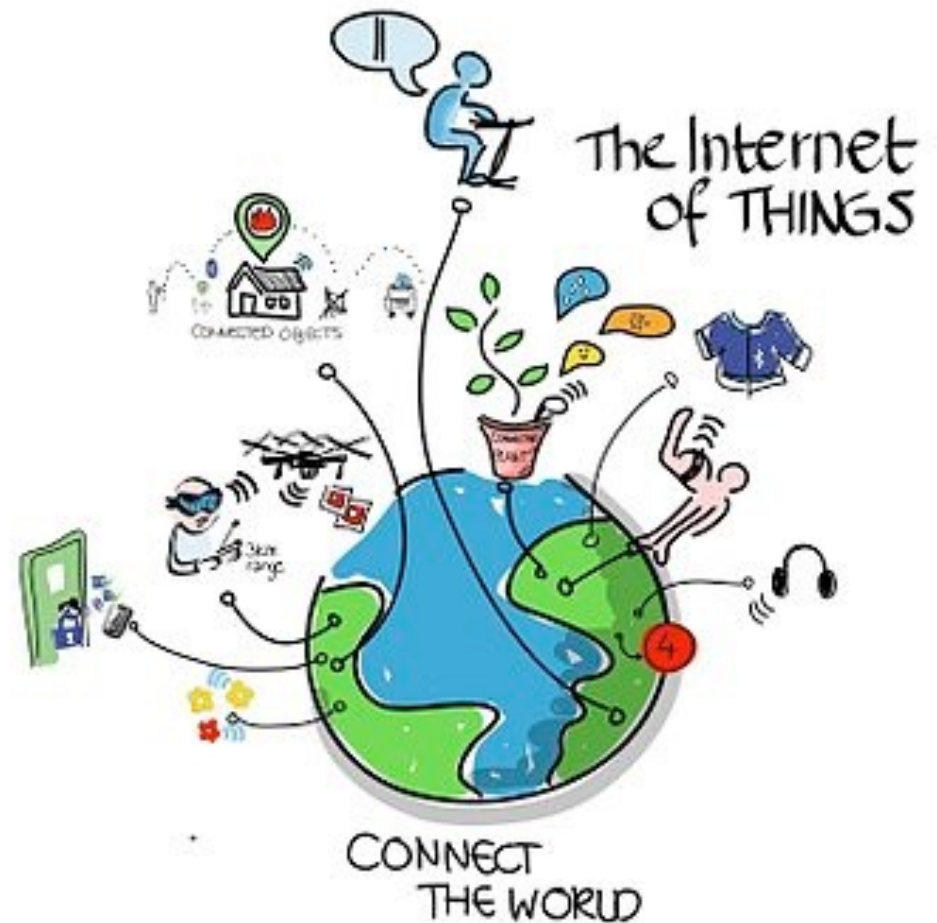
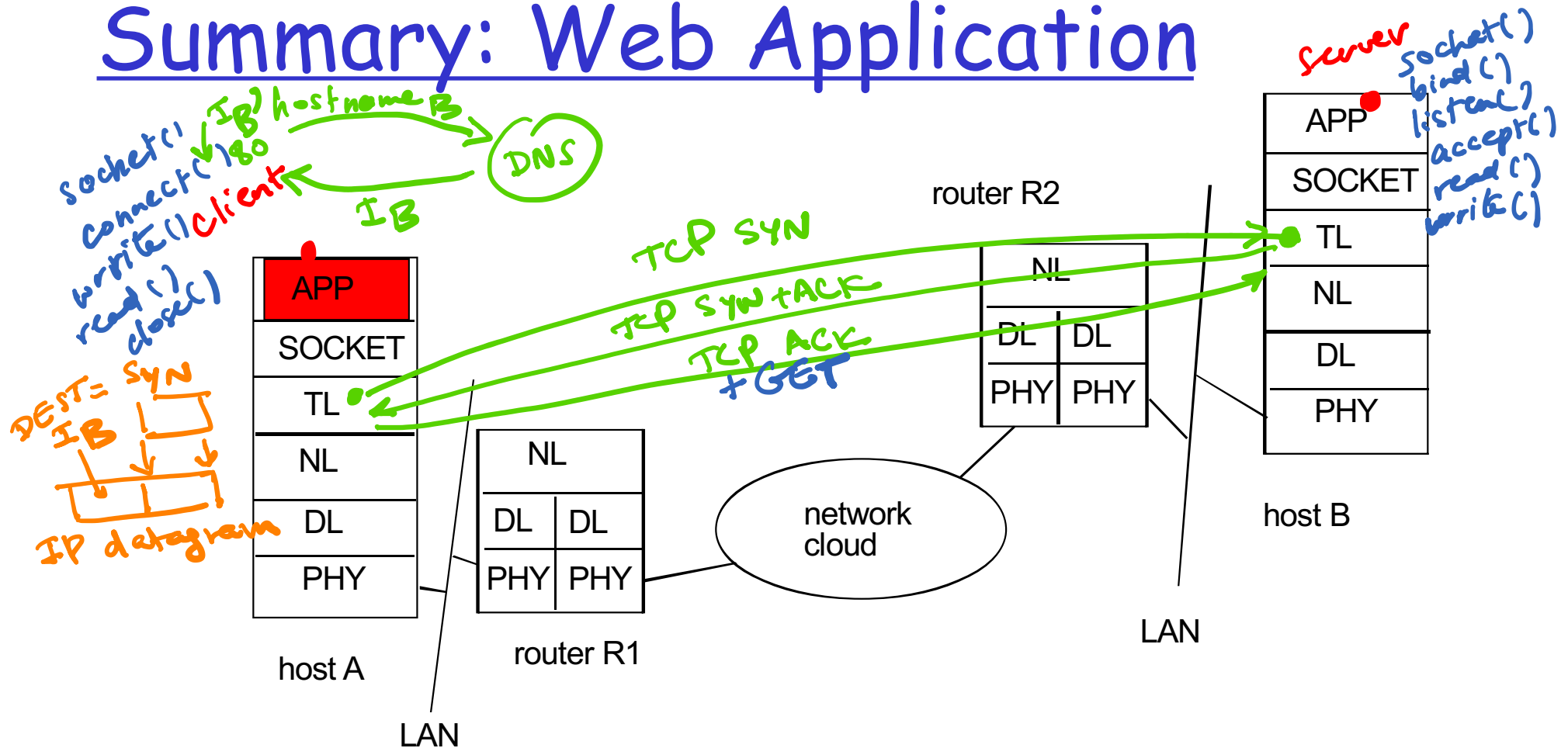


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Computer Science
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Putting it Together & Closing Remarks

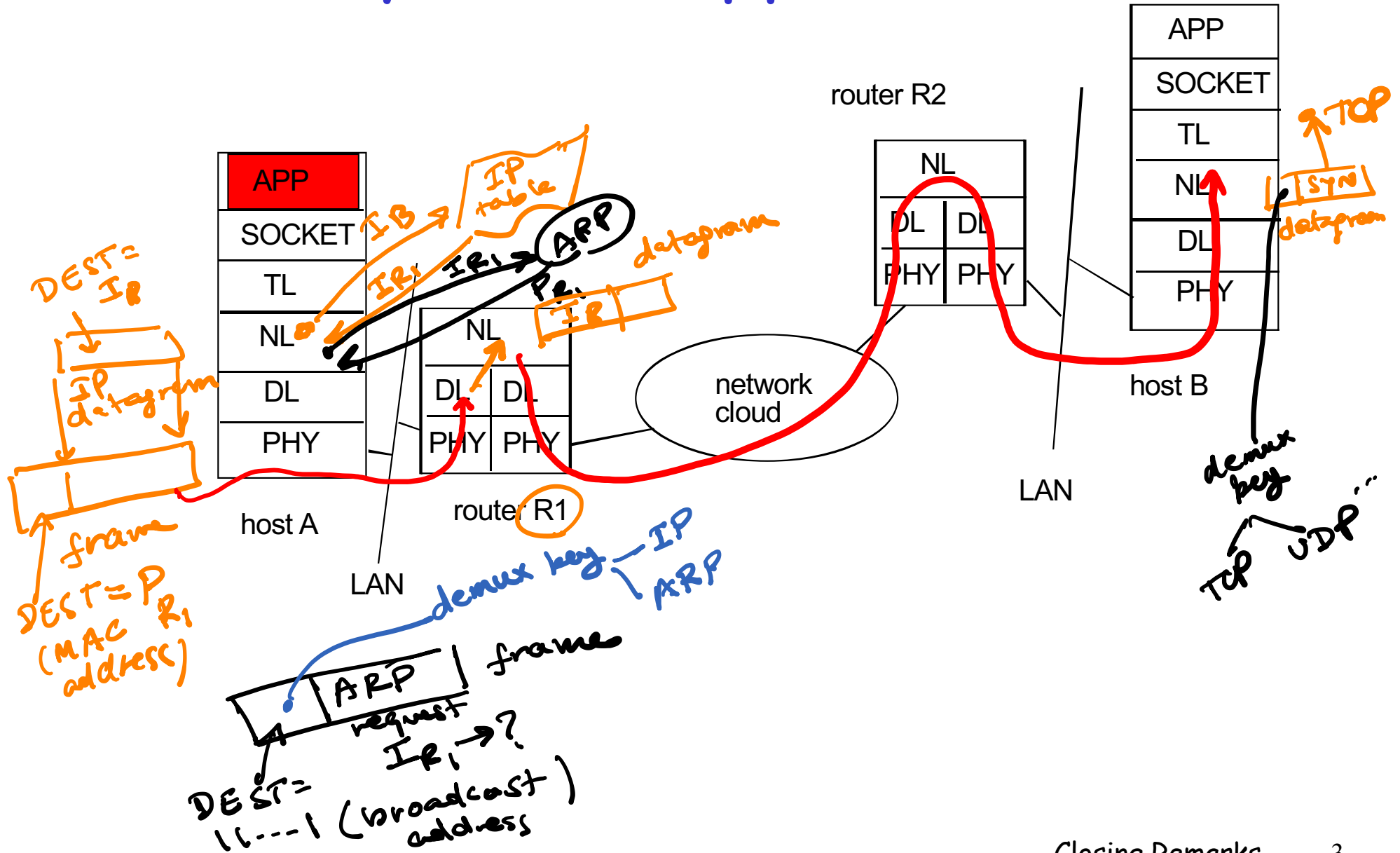


Summary: Web Application

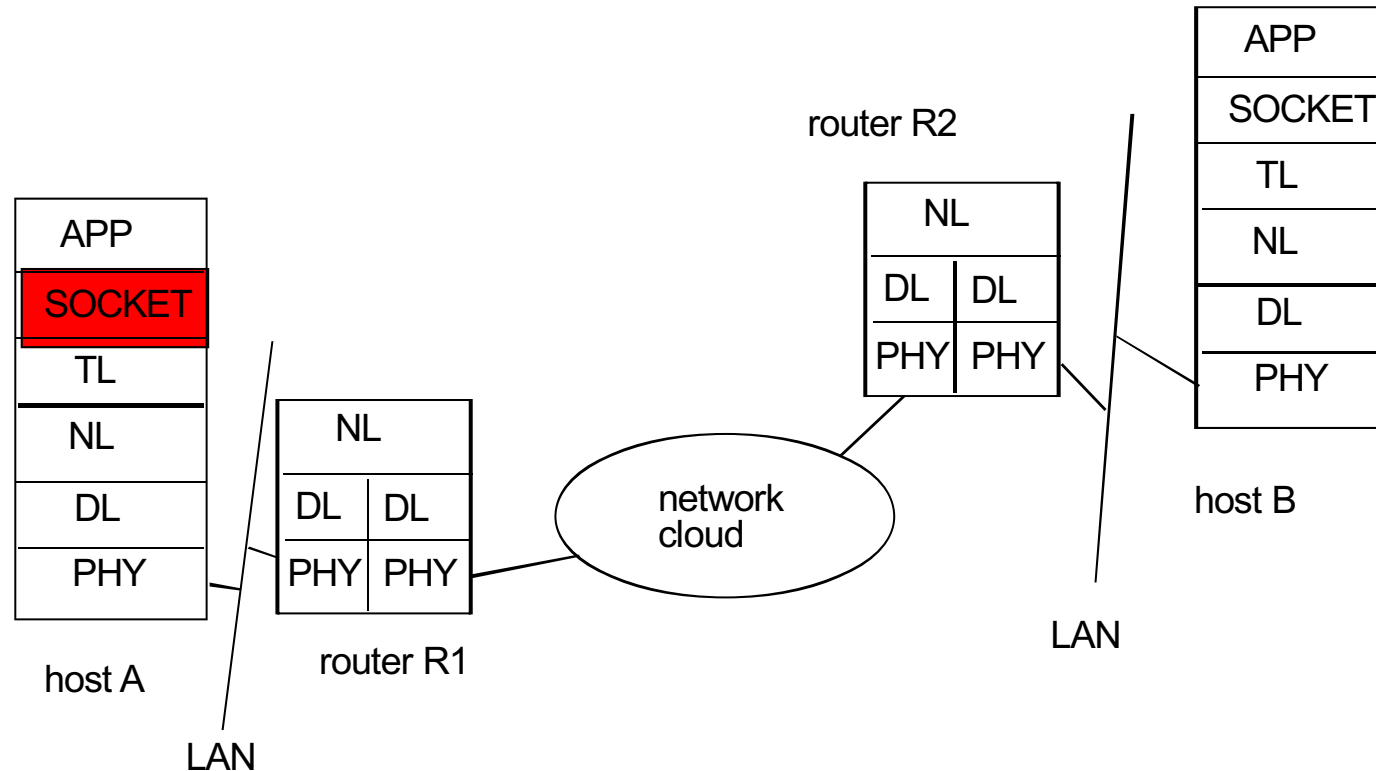


r user enters URL (e.g., www.wireshark.org/download.html) into WWW browser

Summary: Web Application



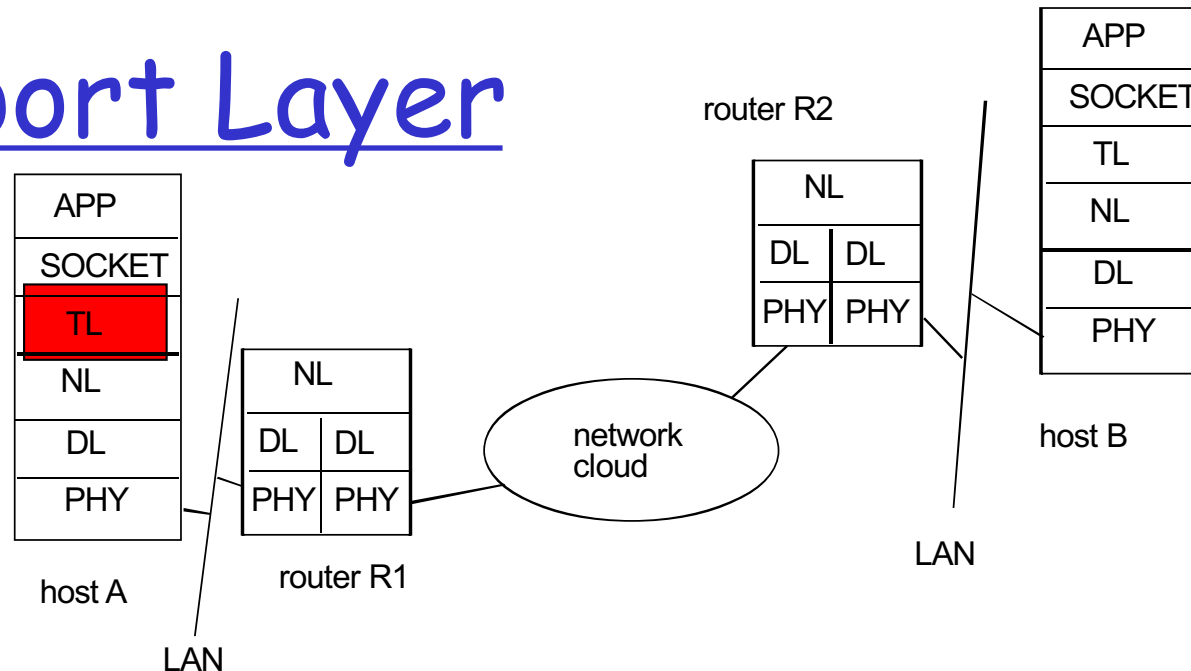
API



- r browser (client) determines host name, uses DNS to get server's IP address, `gethostbyname()`
- r client creates stream socket, `socket()`
- r client calls `connect()`, server port 80

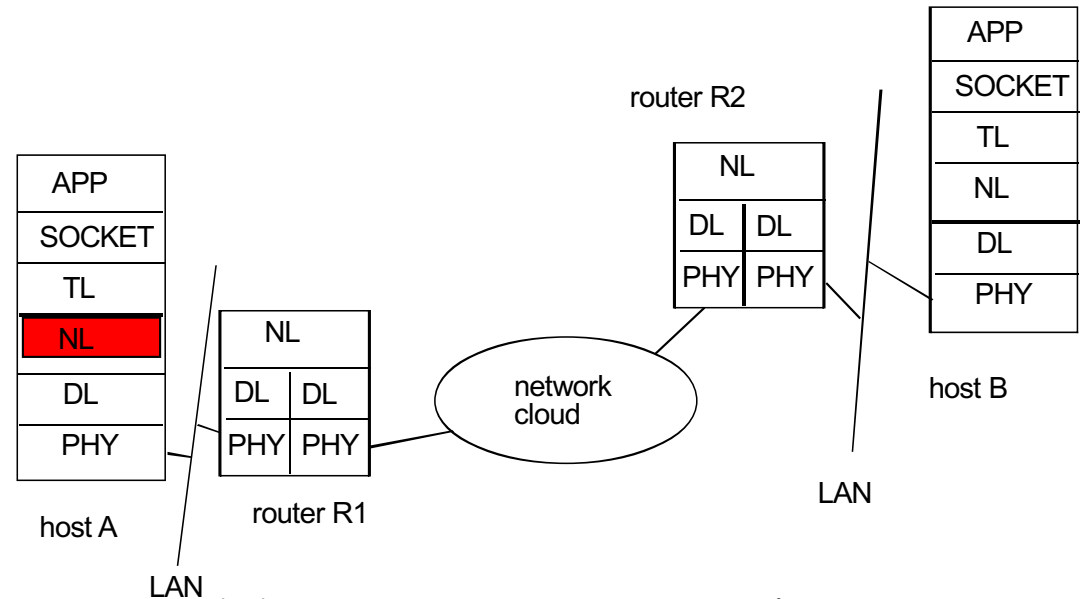
A DNS packet walks into a
liquor store - where do I find
beer "ABC"? Clerk: aisle 4, top
row on the right.
@fsmontenegro

Transport Layer



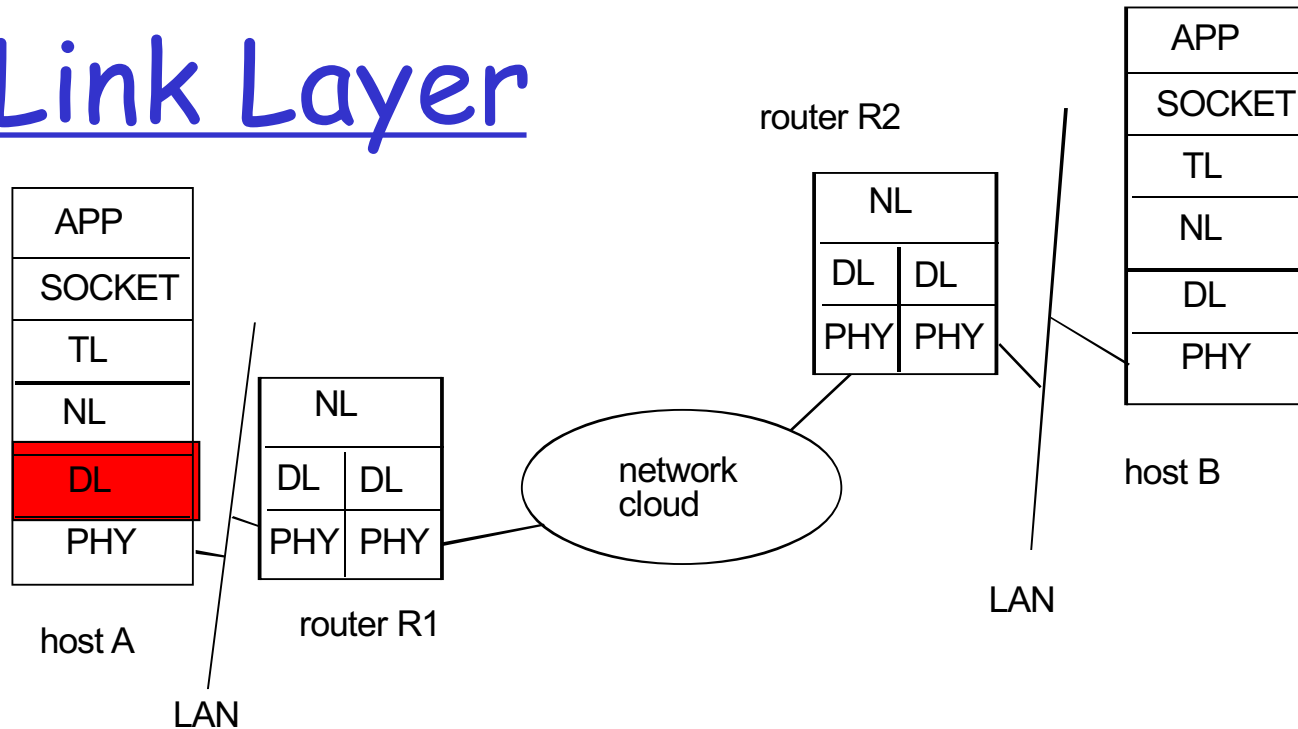
- r connect() call causes TCP connection establishment
- r choose initial sequence number
- r generate SYN segment, server IP address, port 80
- r TCP forms segment, computes checksum
- r TCP calls IP, passing SYN segment and IP address information

Network Layer



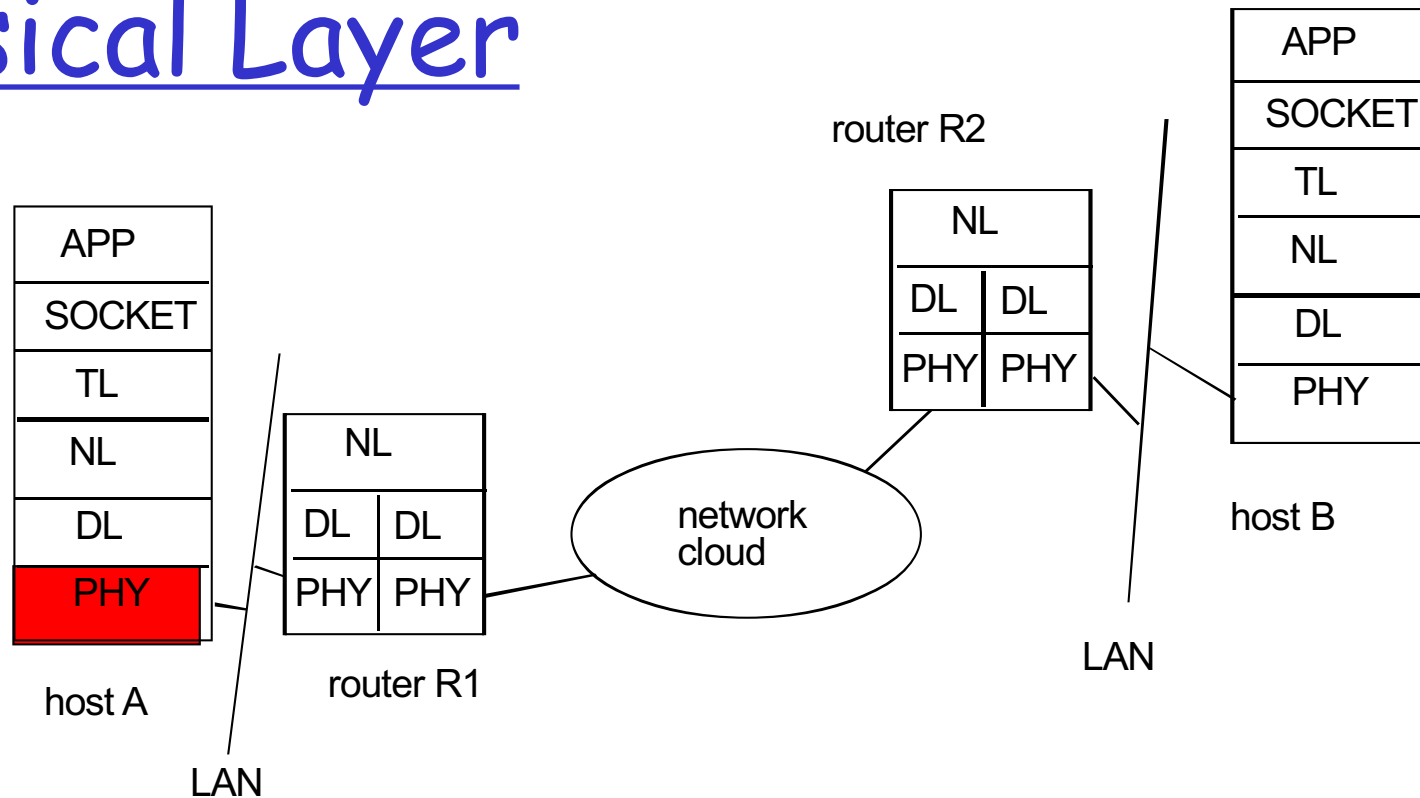
- r adds IP source, destination addresses in IP packet
- r IP forwarding consults routing table
 - m routing table gives IP address of, and local interface to get to next router (i.e., on its LAN), R, on route to destination
- r runs ARP to get 802.3 physical address corresponding to R's IP address
- r ARP will generate Ethernet broadcast frame on LAN, requesting R to reply with its physical address
- r R replies with physical address

Data Link Layer



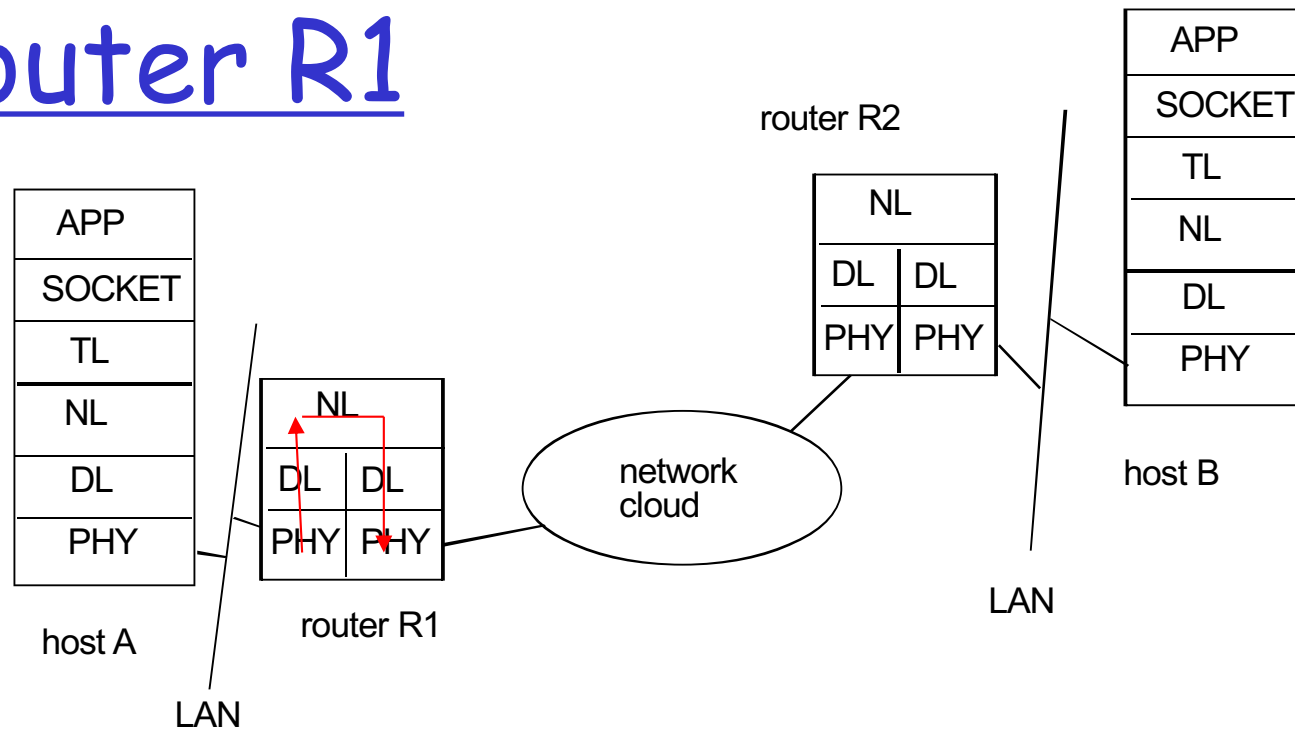
- r TCP SYN segment (inside IP packet), as payload in Ethernet frame sent onto LAN using Ethernet protocol
- r transparent bridge may be involved (not shown)

Physical Layer



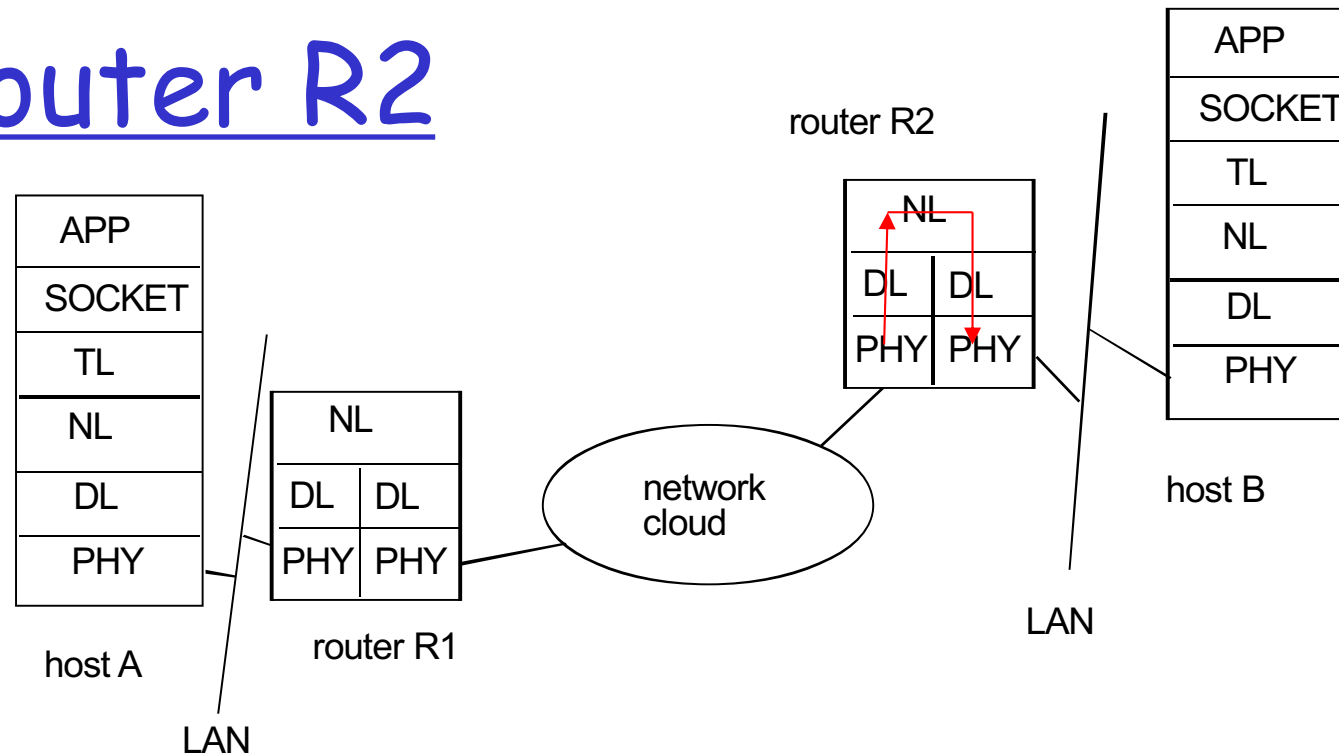
- r Ethernet frame transmitted at 100Mbps

At Router R1



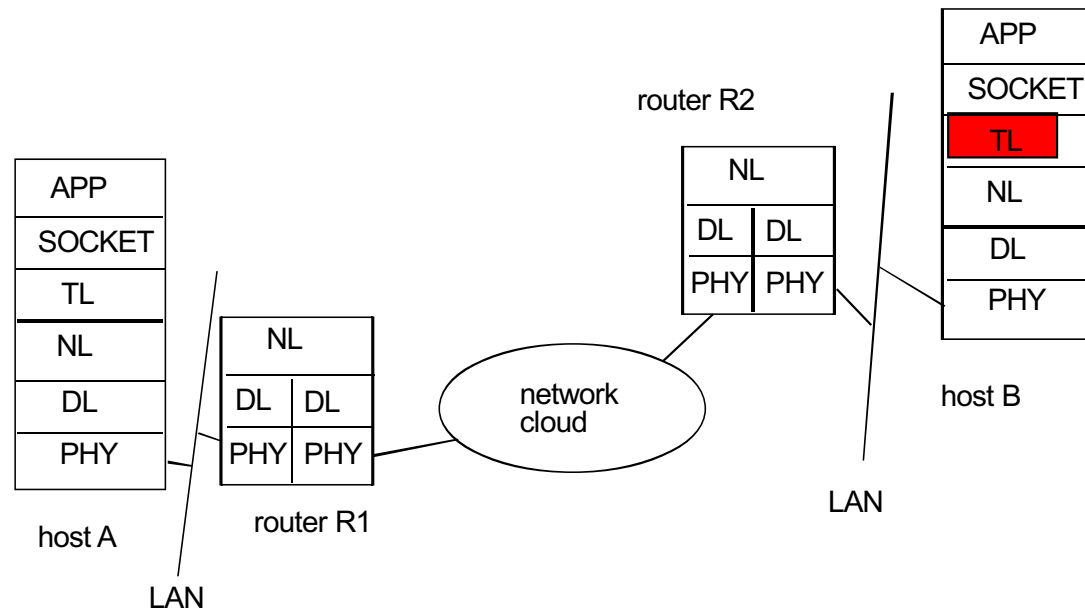
- r physical layer receives Ethernet frame bits, passes up
- r data link layer recognizes frame, computes OK checksum, removes IP packet, passes up
- r network layer consults routing table
 - m passes IP packet down to data link layer

At Router R2



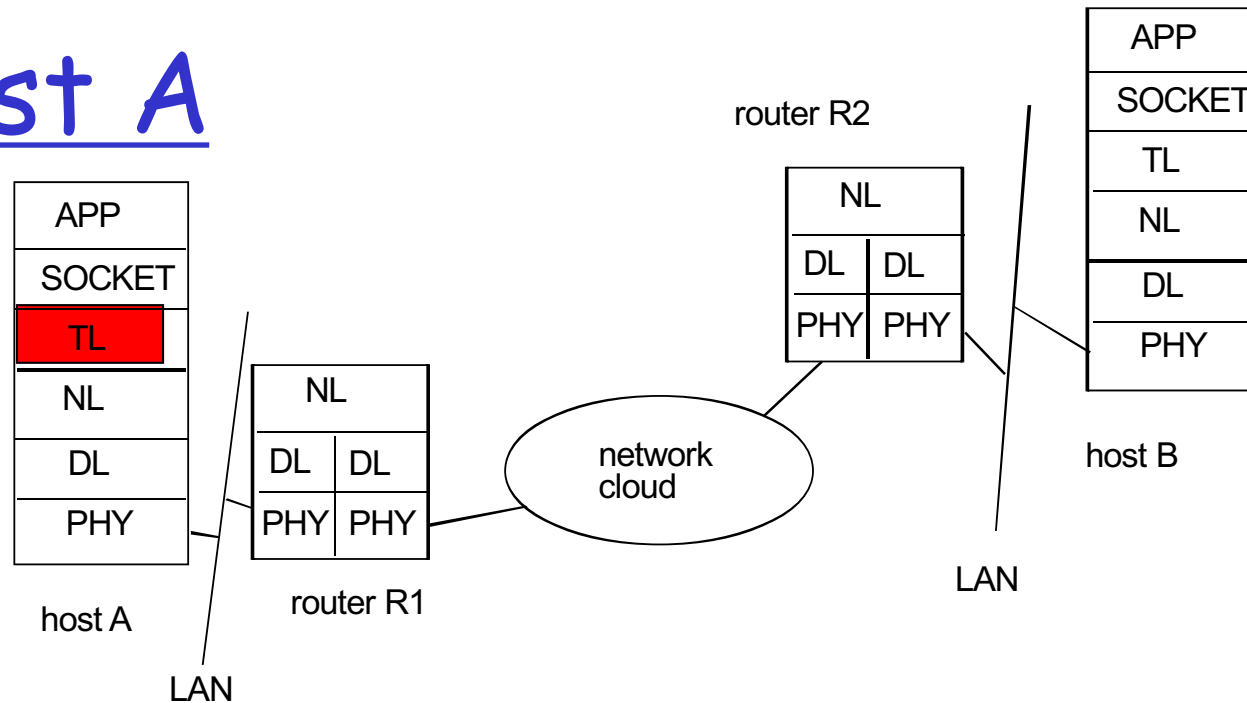
- r frame arrives, passed up to network layer
- r network layer determines outgoing interface to get to host B
- r Ethernet frame sent (R2 knew/learned B's physical address)

At Host B



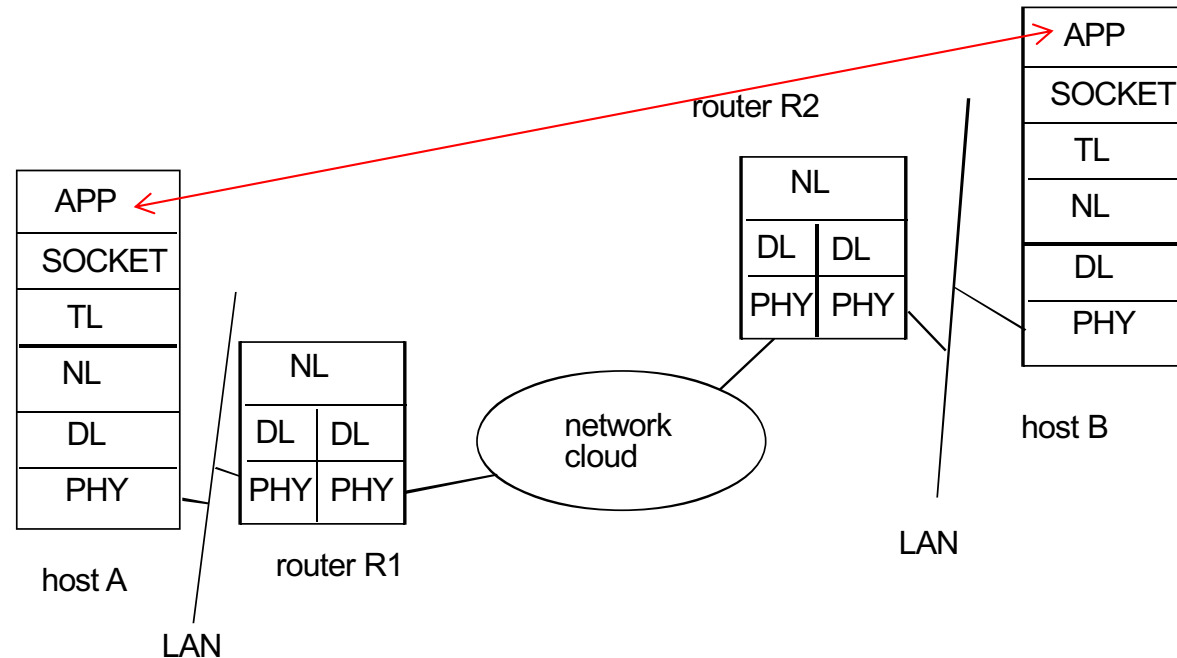
- r Ethernet frame arrives, checksum OK, pass up to IP
- r IP layer extracts TCP segment, demultiplexes up to TCP (note: not UDP message)
- r TCP sees SYN segment
 - m server must have previously opened socket and made accept(), else SYN dropped
 - m TCP determines flow control window, chooses initial sequence no
 - m sends SYN+ACK back

At Host A



- r SYN+ACK eventually received
- r send transport-level ACK to B
- r move to established state
- r return from connect() system call

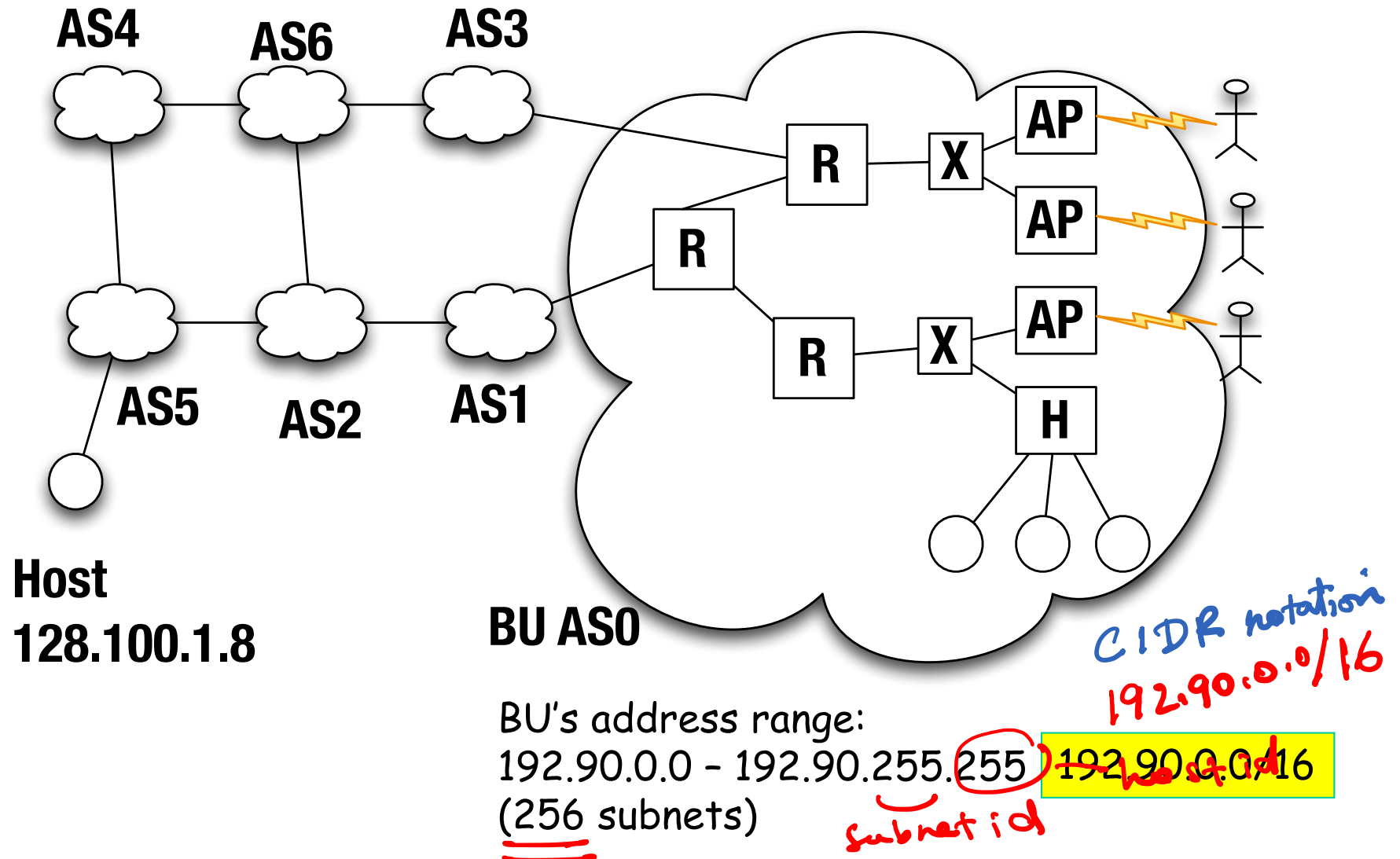
Finally!



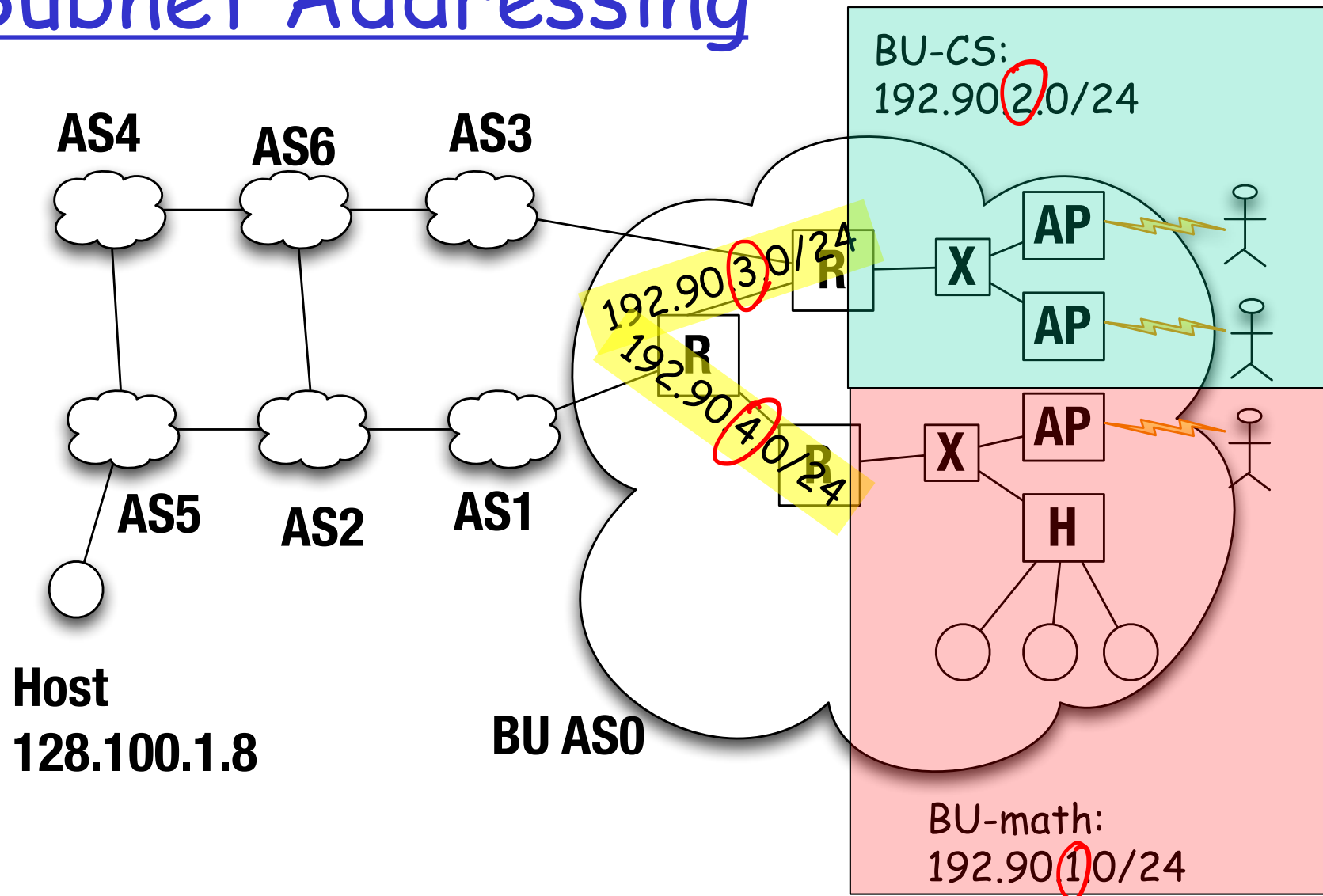
- r Host A can now request .html file, write()
- r Host B reads request, read()
- r Host B sends requested file, write()
- r Host A reads requested file, read()
- r Data is transferred using TCP
 - m Error, flow and congestion control
 - m R1, R2 and A,B (network layer and below) act exactly the same with data as with SYN/ACK segments

A TCP packet walks in to a bar
and says "I want a beer",
barman says "you want a beer?"
and TCP packet says "yes, a
beer" @stevie_chambers

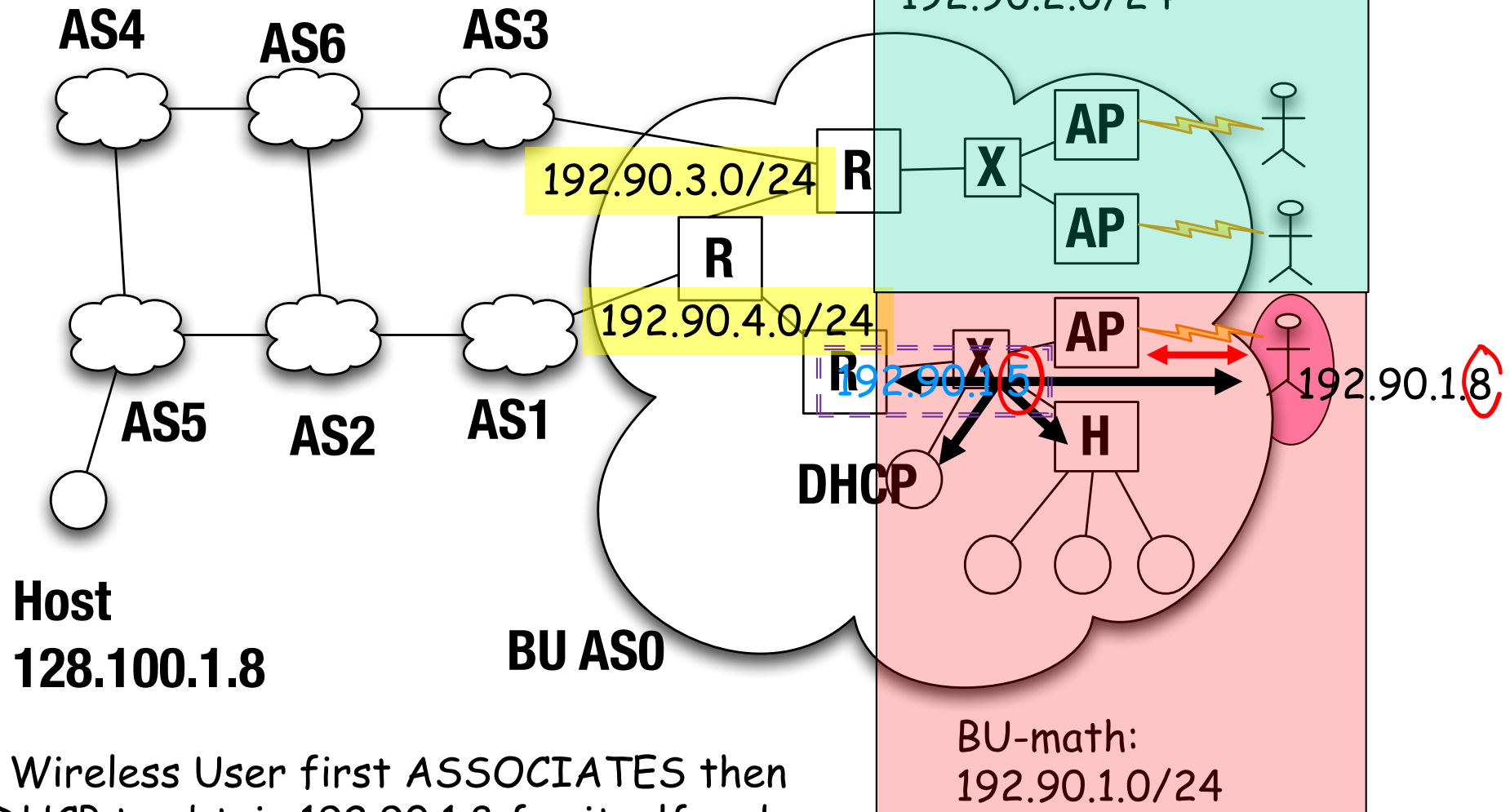
Summary: Addressing & Routing



Subnet Addressing



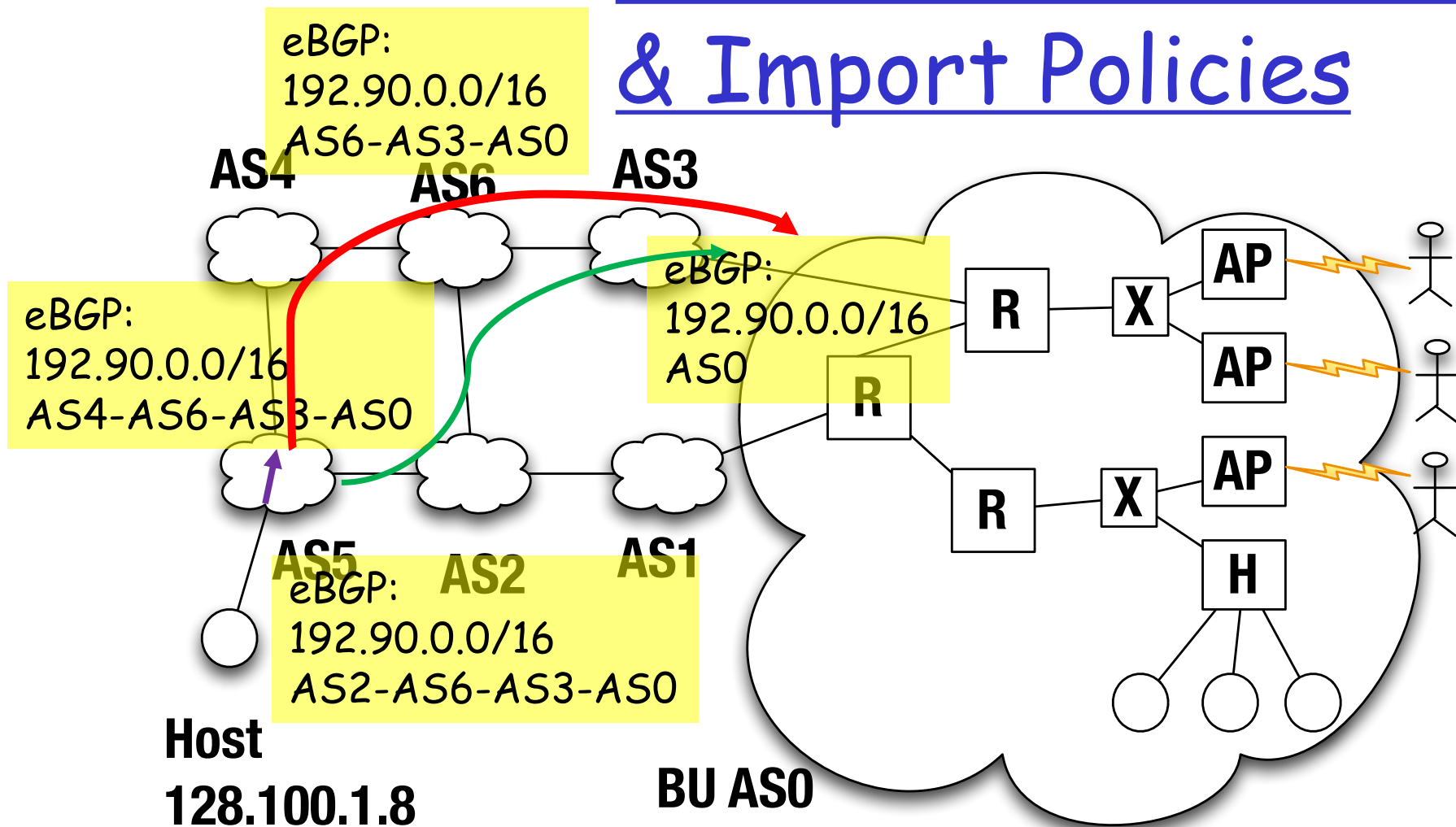
Dynamic Host Config



Math Wireless User first ASSOCIATES then runs DHCP to obtain 192.90.1.8 for itself and learns of 192.90.1.5 (gateway router)

A DHCP packet walks into a bar
and asks for a beer. Bartender
says , "here, but I'll need that
back in an hour!"
@brandoncarroll

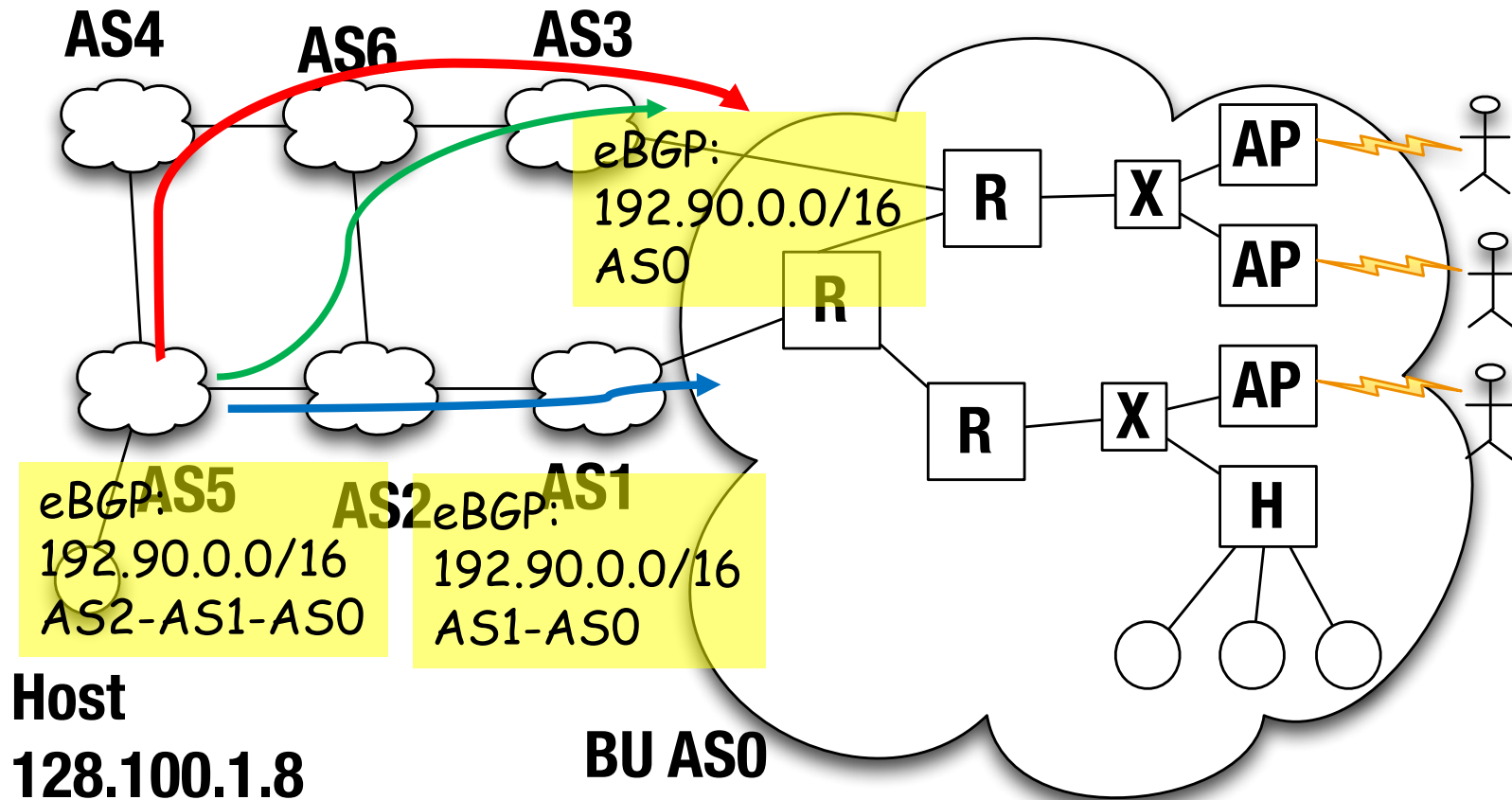
Inter-domain Routing & Import Policies



- AS5 learns of AS4-AS6-AS3-AS0 (LOC_PREF=X), AS2-AS6-AS3-AS0 (LOC_PREF=Y).

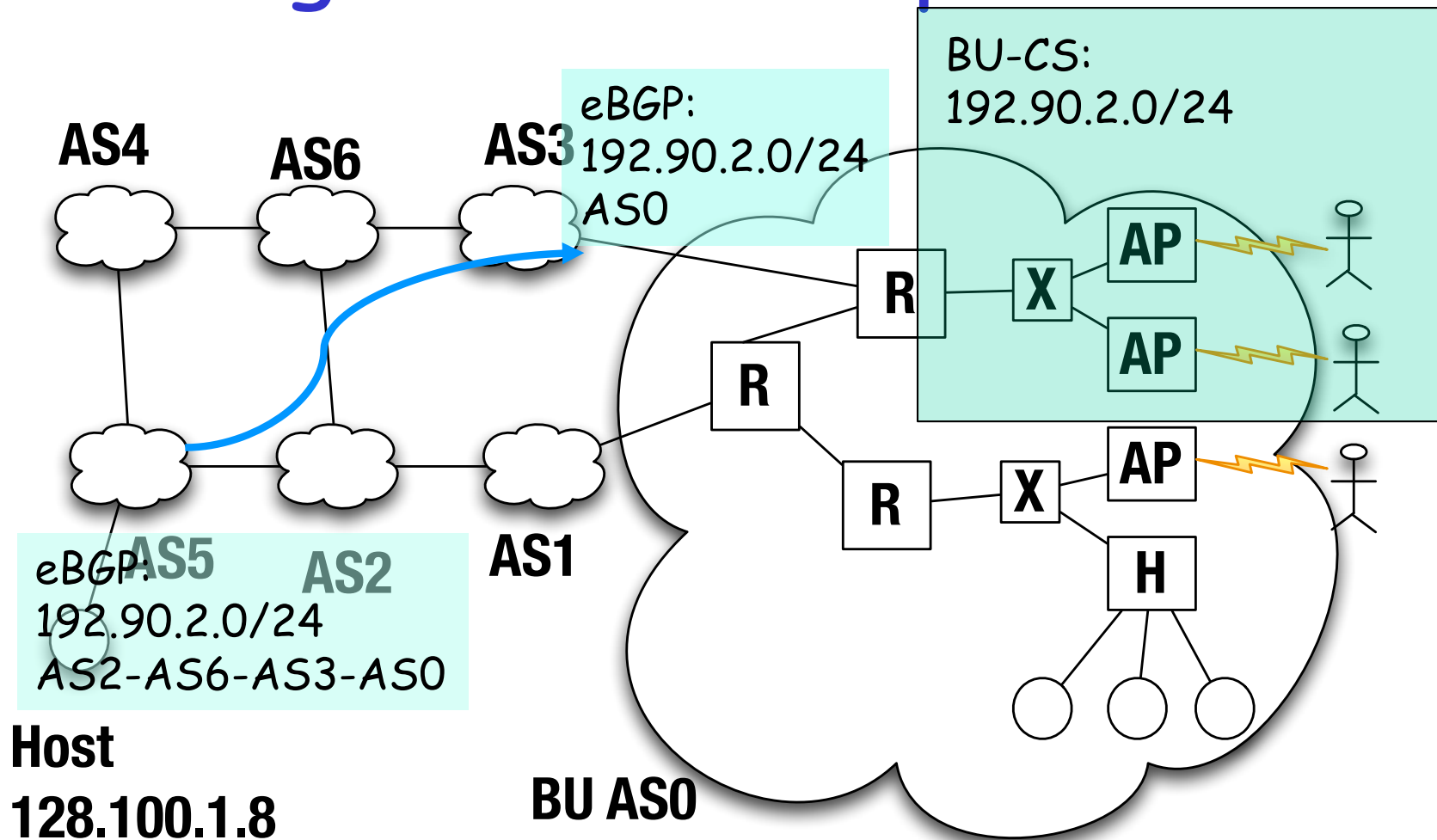
- Pick AS path with higher LOC_PREF. E.g., AS2-AS6-AS3-AS0 if $Y > X$.
- Otherwise ($X=Y$), router inside AS5 picks closest border router (**hot potato routing**)

AS Path Selection



- ❑ AS5 learns of 192.90.0.0/16, AS2-AS1-AS0.
- ❑ Under equal LOC_PREF, AS5 picks it since it is the **shortest AS path**

IP Routing Table Lookup



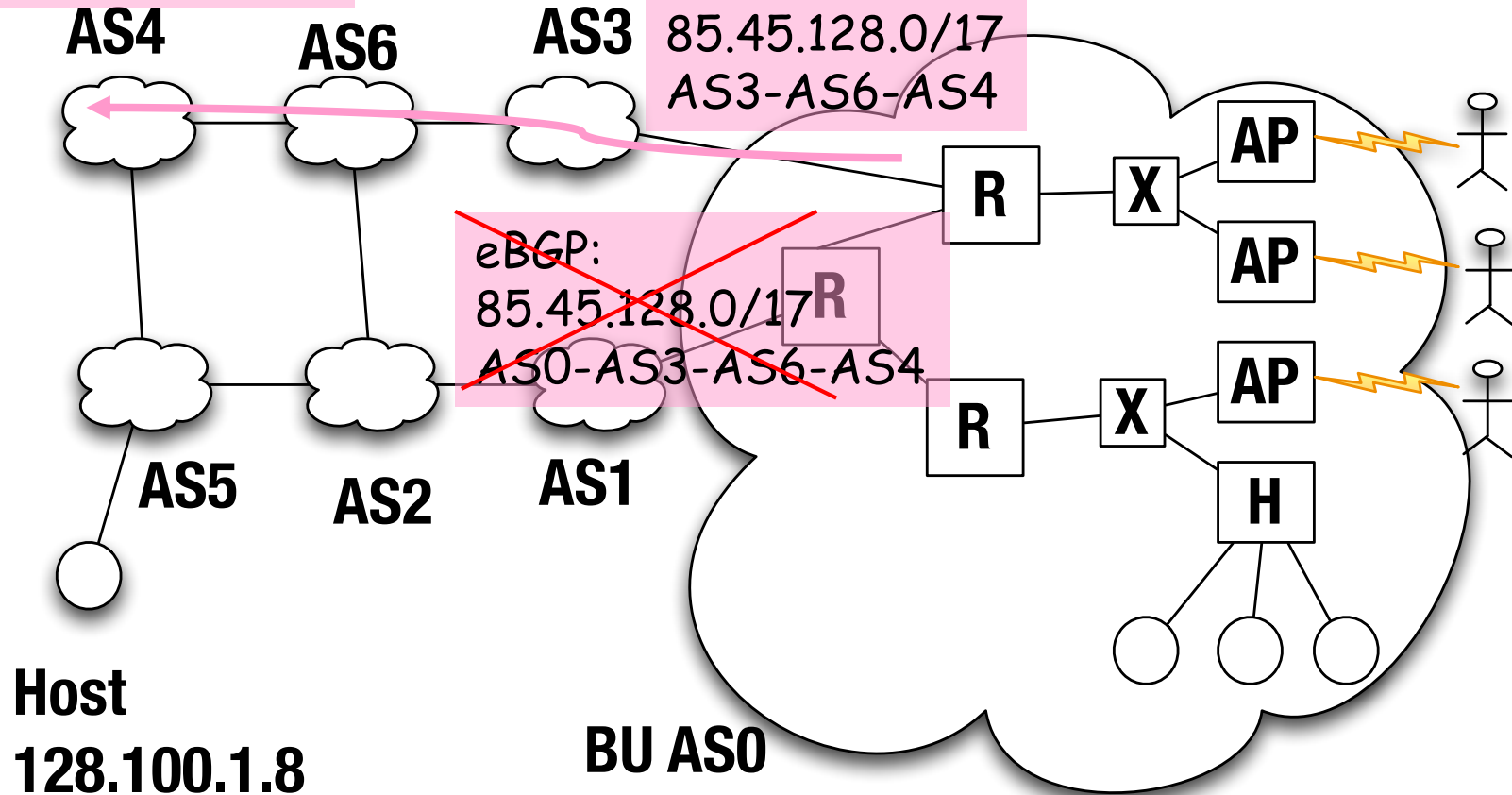
- AS5 learns of 192.90.2.0/24, AS2-AS6-AS3-AS0.
- AS5 uses this path to route to BU-CS (**longest prefix matching**)

Export Policies

eBGP:
85.45.128.0/17
AS4

eBGP:
85.45.128.0/17
AS3-AS6-AS4

~~eBGP:
85.45.128.0/17
AS0-AS3-AS6-AS4~~



- AS0 does NOT advertise 85.45.128.0/17 to AS1 since BU is a **multi-homed stub** AS and not a "transit" AS. This is set by BU's "export policies"

An IPv6 packet walks into a
bar. Nobody talks to him.
@fsmontenegro

So what's more / next?

- ❑ Further study: multimedia, mobility, high-speed, traffic engineering, virtual networks, ... (CS 556)
- ❑ Security (CS 558, CS 511 & 512, ...)
 - CS 511 & 512 teaches formal methods to verify protocol correctness
- ❑ Data center networks (CS 528, ...)
- ❑ Application management (CS 451/651, ...)
- ❑ Internet of Things (CS 552, CS 654 , ...)
- ❑
- ❑ Certification programs!
 - Network Programmability (e.g. Cisco), Network Virtualization (e.g. VMware), Risk & Information Systems Control, Cloud Networking (e.g. Amazon), Wireless Networking, Storage Networking, ...