CS558 Network Security

Lecture 4: ARP and BGP



Current Events

02.16.2022

Blumenthal & Blackburn Introduce Comprehensive Kids' Online Safety Legislation

The introduction of the Kids Online Safety Act follows bombshell reporting & a series of watershed bipartisan hearings revealing Big Tech's repeated failures to protect children & teens from serious dangers on their platforms



Current Events



About Issues Our Work Take Action Tools Donate Q

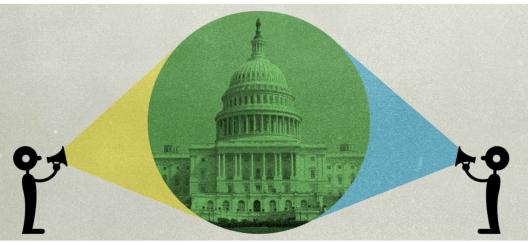
02.16.2022

Blumenthal Introduce Co Online Safet

The introduction of the Kids Online Sa bipartisan hearings revealing Big Tech' dangers on their platforms

Dangerous "Kids Online Safety Act" Does Not Belong in Must-Pass Legislation

BY JASON KELLEY AND AARON MACKEY | DECEMBER 15, 2022



Review: Encrypted Email



Review: Encrypt Then Sign vs Sign Then Encrypt



Ad-Hoc, Fast, Decentralized Routing is Hard

(That's it. That's the tweet.)



Updated by: <u>5227</u>, <u>5494</u>
Network Working Group
Request For Comments: 826

Internet Standard
David C. Plummer
(DCP@MIT-MC)
November 1982

An Ethernet Address Resolution Protocol

The Problem:

The world is a jungle in general, and the networking game contributes many animals. At nearly every layer of a network architecture there are several potential protocols that could be used. For example, at a high level, there is TELNET and SUPDUP for remote login. Somewhere below that there is a reliable byte stream protocol, which might be CHAOS protocol, DOD TCP, Xerox BSP or DECnet. Even closer to the hardware is the logical transport layer, which might be CHAOS, DOD Internet, Xerox PUP, or DECnet. The 10Mbit Ethernet allows all of these protocols (and more) to coexist on a single cable by means of a type field in the Ethernet packet header. However, the 10Mbit Ethernet requires 48.bit addresses on the physical cable, yet most protocol addresses are not 48.bits long, nor do they necessarily have any relationship to the 48.bit Ethernet address of the hardware. For example, CHAOS addresses are 16.bits, DOD Internet addresses are 32.bits, and Xerox PUP addresses are 8.bits. A protocol is needed to dynamically distribute the correspondences between a protocol, address pair and a 48.bit Ethernet address.



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This format allows the packet buffer to be reused if a reply is generated; a reply has the same length as a request, and several of the fields are the same.

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Link Layer Reminder



A(dress) R(esolution) P(rotocol) Reminder



A(dress) R(esolution) P(rotocol) Reminder

- Request
- Response
- Gratuitous ARP
 - Request/Announcement: IP of sender in both
 - Response: Normal response to no request



```
648 10,124459
                        Apple f0:71:99
                                                   Broadcast
                                                                              ARP
                                                                                            42 Who has 192.168.86.1? Tell 192.168.86.48
     653 10.214837
                        Google c9:6c:39
                                                   Apple_f0:71:99
                                                                              ARP
                                                                                            42 192,168,86,1 is at 70:3a:cb:c9:6c:39
▶ Frame 648: 42 bytes on wire (336 bits), 42 bytes captured (336 bits) on interface 0
▼ Ethernet II, Src: Apple_f0:71:99 (a4:5e:60:f0:71:99), Dst: Broadcast (ff:ff:ff:ff:ff:ff)
  ▶ Destination: Broadcast (ff:ff:ff:ff:ff)
  Source: Apple f0:71:99 (a4:5e:60:f0:71:99)
    Type: ARP (0x0806)
▼ Address Resolution Protocol (request)
    Hardware type: Ethernet (1)
    Protocol type: IPv4 (0x0800)
    Hardware size: 6
    Protocol size: 4
    Opcode: request (1)
    Sender MAC address: Apple_f0:71:99 (a4:5e:60:f0:71:99)
    Sender IP address: 192.168.86.48
    Target MAC address: 00:00:00 00:00:00 (00:00:00:00:00:00)
    Target IP address: 192.168.86.1
0000
     ff ff ff ff ff a4 5e 60 f0 71 99 08 06 00 01
                                                         .......^ `.q...V0
0010 08 00 06 04 00 01 a4 5e 60 f0 71 99 c0 a8 56 30
                                                         . . . . . . . . V .
```

0020 00 00 00 00 00 00 c0 a8 56 01



```
653 10.214837
                        Google c9:6c:39
                                                    Apple_f0:71:99
                                                                               ARP
                                                                                              42 192.168.86.1 is at 70:3a:cb:c9:6c:39
▶ Frame 653: 42 bytes on wire (336 bits), 42 bytes captured (336 bits) on interface 0
Ethernet II, Src: Google_c9:6c:39 (70:3a:cb:c9:6c:39), Dst: Apple_f0:71:99 (a4:5e:60:f0:71:99)
  Destination: Apple f0:71:99 (a4:5e:60:f0:71:99)
  Source: Google_c9:6c:39 (70:3a:cb:c9:6c:39)
    Type: ARP (0x0806)
▼ Address Resolution Protocol (reply)
    Hardware type: Ethernet (1)
    Protocol type: IPv4 (0x0800)
    Hardware size: 6
    Protocol size: 4
    Opcode: reply (2)
    Sender MAC address: Google c9:6c:39 (70:3a:cb:c9:6c:39)
    Sender IP address: 192,168,86,1
    Target MAC address: Apple_f0:71:99 (a4:5e:60:f0:71:99)
    Target IP address: 192.168.86.48
      a4 5e 60 f0 71 99 70 3a cb c9 6c 39 08 06 00 01 ^`.g.p: ..l9...
                                                          · · · · · p: · · l9 · · V ·
0010 08 00 06 04 00 02 70 3a cb c9 6c 39 c0 a8 56 01
0020 a4 5e 60 f0 71 99 c0 a8 56 30
                                                           ·^` · q · · · V0
```



ARP Security?





ARP Spoofing Mitigations

Detection: Watching for potential mismatches

Prevention: Centralization of some kind

- Cryptography: Secure-ARP



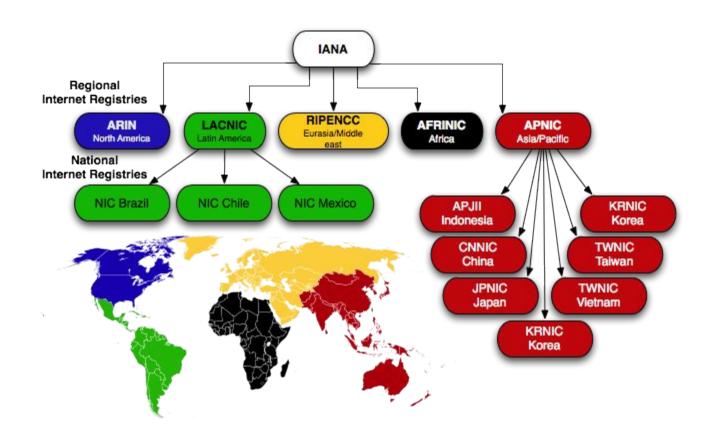
Prepping for BGP



Internet Organization/s

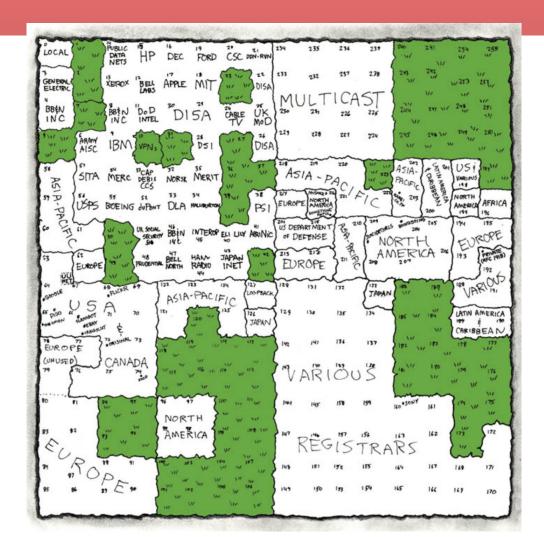
- CIDR (Class-less internet domain routing) notation
- Internet Assigned Numbers Authority
- Regional Internet Registrars
- ISPs and Backbone Networks (Level3)
- Autonomous Systems

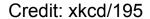














Autonomous Systems Reminder



111 search

AS number 111 **BOSTONU-AS** AS name organization **Boston University** United States country **AS** rank 11752 10 165888 customer cone prefix address 5 global 3 4 1 0 **AS** degree transit provider peer customer Spoofer 01/2022-01/2023 **Tested IP Blocks** 17 0 0 (0.0%) IPv4 /24s 0 (0.0%) IPv6 /40s **Spoofing IP Blocks** see more spoofer data >

AS Rank A	AS neighbors ▼	Organization	AS customer cone ∇	number of paths	relationship
1	3356	Level 3 Parent, LLC	48838	137	provider
3	174	Cogent Communications	34689	54	provider
68	32787	Akamai Technologies, Inc.	658	206	provider
578	10578	Harvard University	66	40	provider
24263	10961	Boston GigaPoP	1	12	peer

