

— LAMBDA CALCULUS — REVIEW

1. Reducible expression (REDEx)

$(\lambda \underline{x}. \underline{E}) \underline{M}$ is called a redex

$$(\lambda x. E) M \rightarrow_{\beta} \underline{E_1}$$

$$\lambda \underline{x}. \underline{E}$$

- $(\lambda x. x) y \rightarrow_{\beta} y$

- $(\lambda x. x) (\lambda z. z) \rightarrow_{\beta}$

- $\frac{(\lambda x. \lambda y. x) \underline{z} \underline{w}}{\underline{x} \quad \underline{E} \quad \underline{M}} \rightarrow_{\beta} \frac{(\lambda y. \underline{z}) \underline{w}}{\underline{-} \quad \underline{-}} \rightarrow_{\beta} \underline{z}$

2. Substitution

Each β -reduction triggers a substitution

$$(\lambda x. E) M \rightarrow_{\beta} \underline{E_1}$$

$$E[M/x]$$

- Substitution may require renaming of formal parameters to avoid making free variables bound
- In our examples, we will be performing renaming only when necessary

$$(\lambda x. \lambda y. x) z w \rightarrow_{\beta} (\lambda y. z) w \rightarrow z$$

- Automated algorithm does aggressive renaming (Lec 15, slide 20)

$$\underline{(\lambda y. E)} [M/x]$$

3. Normal Form

3.1 Canonical normal form NF

- $\lambda z. z$ is in NF

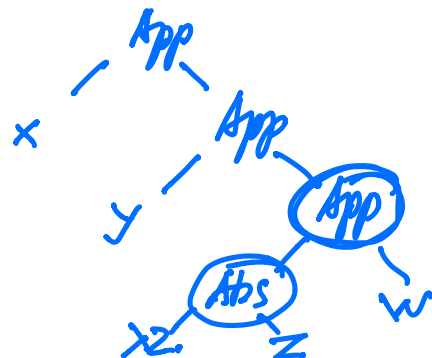
- $z v w$ is in NF

- $z (v w)$ is in NF

- $(\lambda x. x) z$ is not in NF

- x $((\lambda x. x) (\lambda y. y))$ is not in NF
 $E_1 \quad E_2$

- x $(y ((\lambda z. z) w))$ is not in NF
 $E_1 \quad E_2$
 $E_3 \quad E_4$



3.2 Weak normal forms (lazy normal forms)

HNF:

- x is in HNF
- • $x \underline{E_1} \underline{E_2} \underline{E_3} \dots \underline{E_n}$ is in HNF
- $\lambda x. E$ is in HNF if E is in HNF
- $\lambda z. z$ is in HNF
- $x ((\lambda x. x) (\lambda y. y))$ is in HNF

WHNF:

- x is in WHNF
- x $E_1 E_2 \dots E_n$ is in WHNF
- $\lambda x. \underline{E}$ is WHNF
- $\lambda x. \underline{(\lambda y. y) z}$ is in WHNF

4. Combinators are terms with no free variables

$$I = \lambda x. x$$

$$S = \lambda x. \lambda y. \lambda z. x z (y z)$$

$$K = \lambda x. \lambda y. x$$

$$\begin{aligned} \rightarrow S' K S' K &= ((S' K) S') K \\ &\xrightarrow{\beta} (\lambda x. \lambda y. \lambda z. x z (y z)) K S' K \\ &\xrightarrow{\beta} (\lambda y. \lambda z. K z (y z)) S' K \\ &\xrightarrow{\beta} (\lambda z. K z (S' z)) K \end{aligned}$$

$$\begin{aligned} &\underline{K K} (S' K) = \\ &\underline{(\lambda x. \lambda y. x)} \underline{K} \underline{(S' K)} \xrightarrow{\beta} K \end{aligned}$$