



Logic Programming and Prolog

Finish reading: Scott, Chapter 12

Lecture Outline

- Prolog
 - Imperative control flow
 - Negation by failure
 - Generate and test paradigm

Imperative Control Flow

- Programmer has **explicit control** on backtracking process

cut (!)

- **!** is a subgoal
- As a goal it succeeds, but with a **side effect**:
 - Commits interpreter to **all bindings** made since unifying **left-hand side of current rule** with **parent goal**

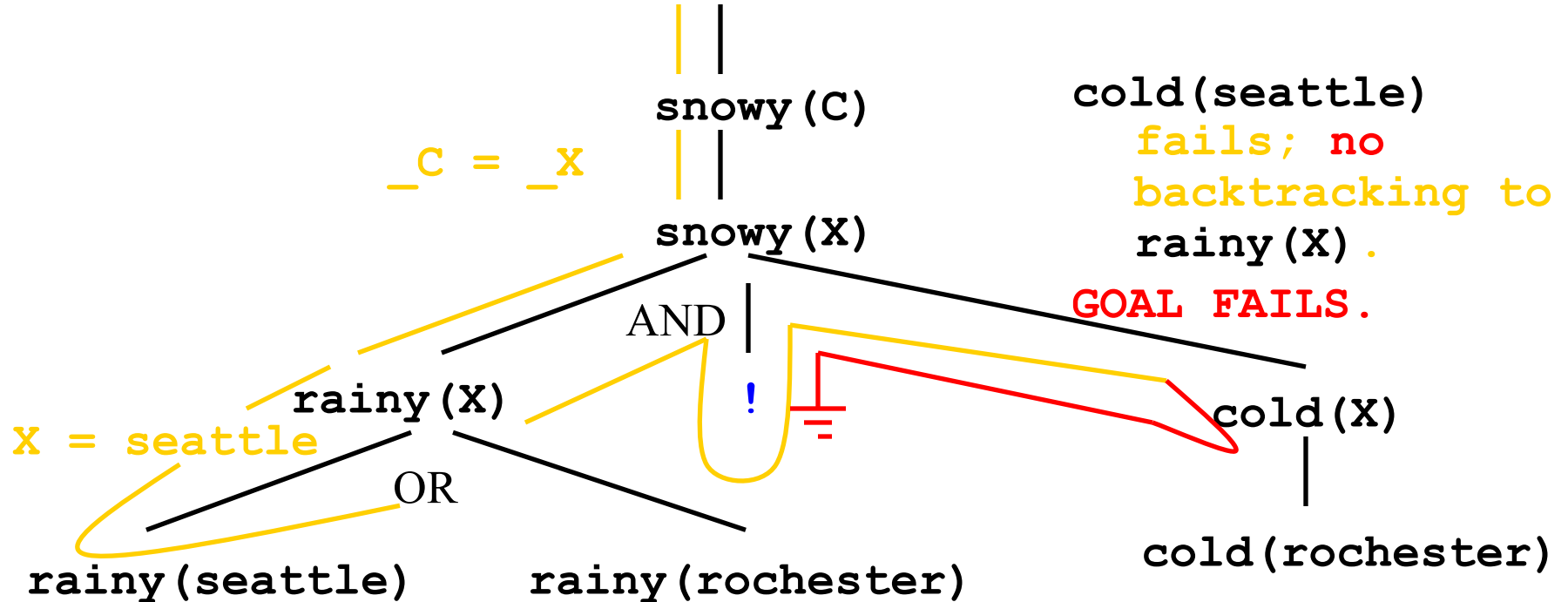
Cut (!) Example

```
rainy(seattle) .  
rainy(rochester) .  
cold(rochester) .  
snowy(X) :- rainy(X) , ! , cold(X) .
```

```
?- snowy(C) .
```

Cut (!) Example

```
rainy(seattle).  
rainy(rochester).  
cold(rochester).  
snowy(X) :- rainy(X), !, cold(X).
```



Cut (!) Example 2

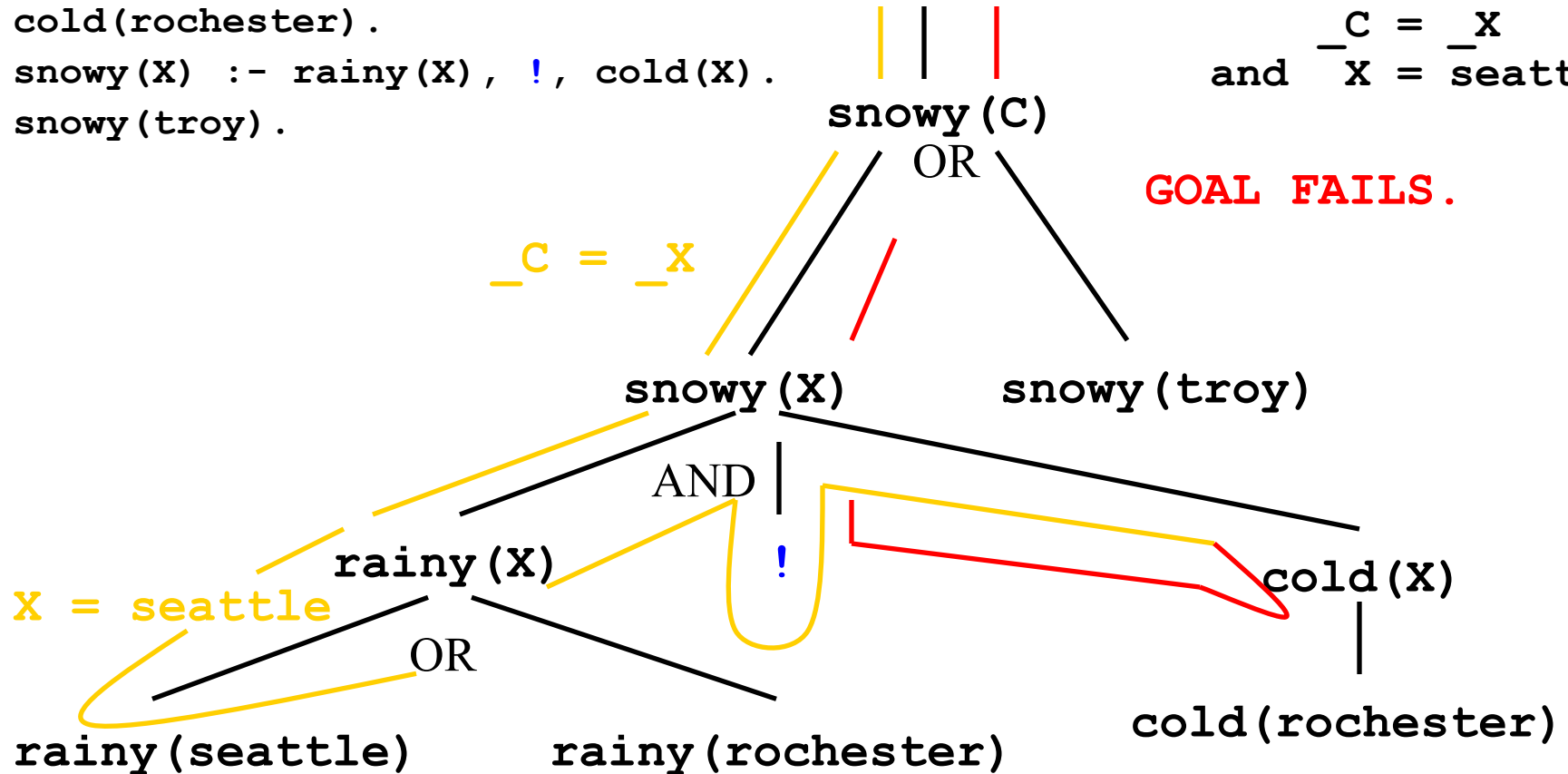
```
rainy(seattle) .  
rainy(rochester) .  
cold(rochester) .  
snowy(X) :- rainy(X) , ! , cold(X) .  
snowy(troy) .
```

```
?- snowy(C) .
```

Cut (!) Example 2

```
rainy(seattle).
rainy(rochester).
cold(rochester).
snowy(X) :- rainy(X), !, cold(X).
snowy(troy).
```

2 committed OR
bindings:
 $_C = _X$
and $_X = \text{seattle}$



How about query ?- `snowy(troy)` ?

Cut (!) Example 3

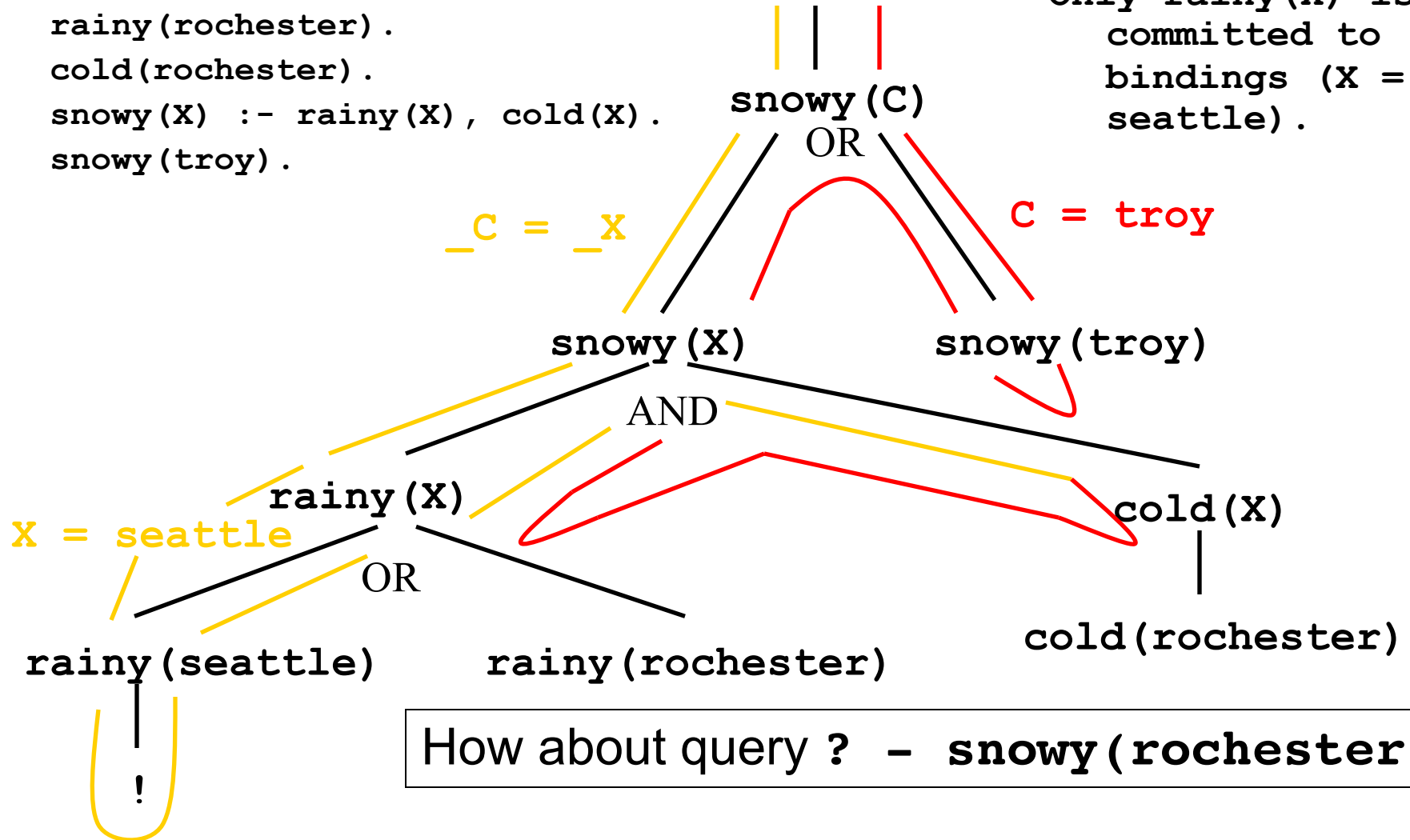
```
rainy(seattle) :- !.  
rainy(rochester).  
cold(rochester).  
snowy(X) :- rainy(X), cold(X).  
snowy(troy).  
  
?- snowy(C).
```


Cut (!) Example 3

```
rainy(seattle) :- !.  
rainy(rochester).  
cold(rochester).  
snowy(X) :- rainy(X), cold(X).  
snowy(troy).
```

C = troy
SUCCEEDS

Only rainy(X) is
committed to
bindings (X =
seattle).



How about query ? - **snowy(rochester)**?

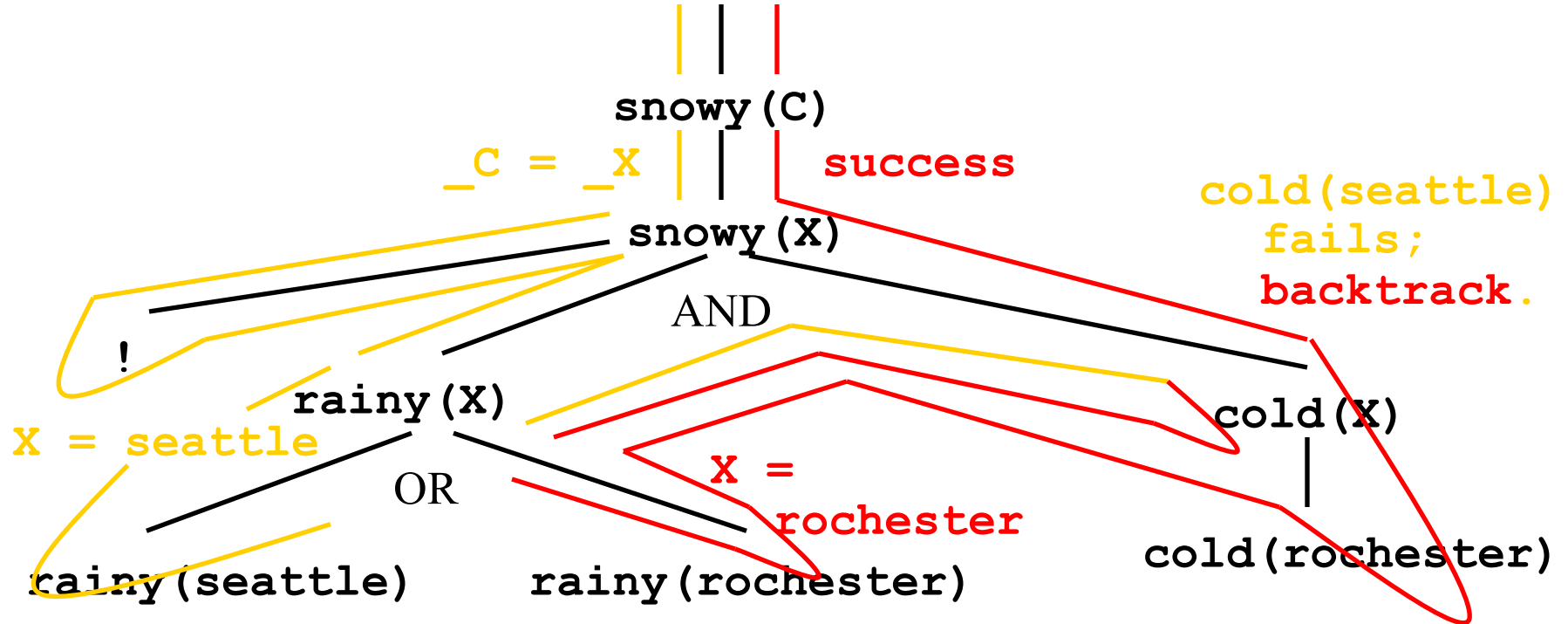
Cut (!) Example 4

```
rainy(seattle) .  
rainy(rochester) .  
cold(rochester) .  
snowy(X) :- !, rainy(X), cold(X) .
```

```
?- snowy(C) .
```

Cut (!) Example 4

```
rainy(seattle).  
rainy(rochester).  
cold(rochester).  
snowy(X) :- !, rainy(X), cold(X).
```



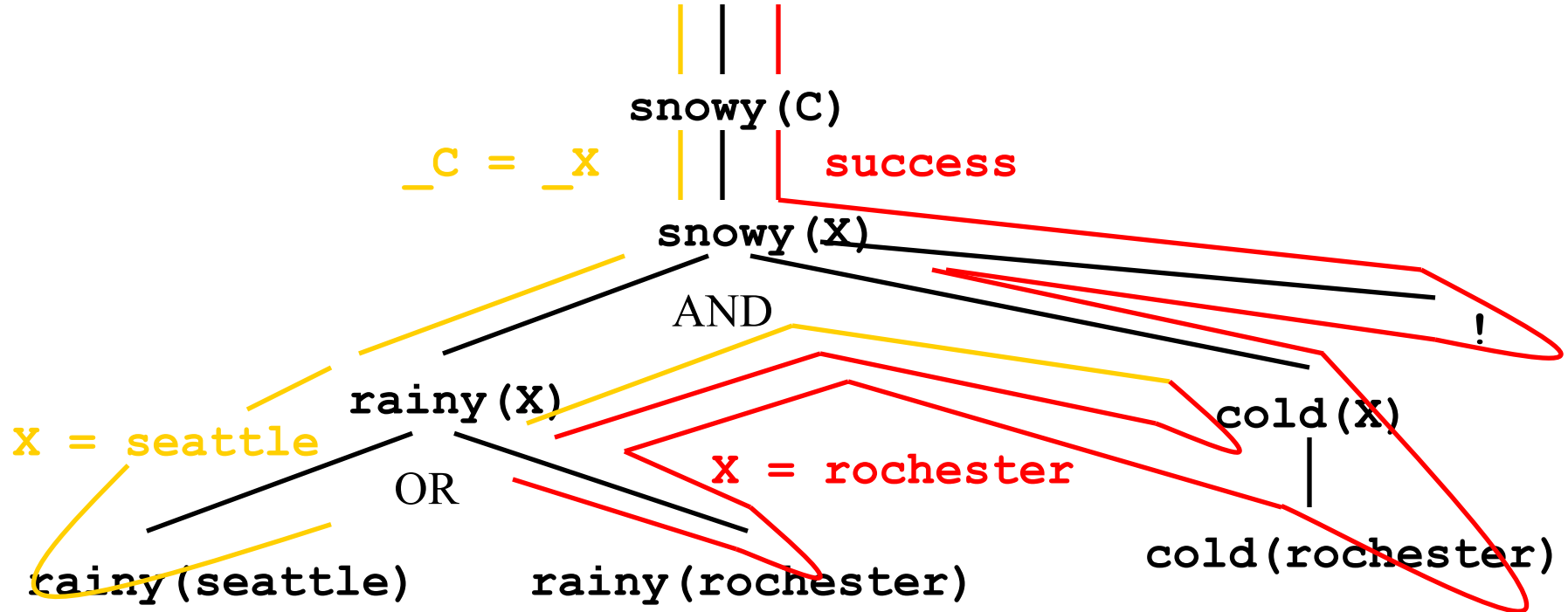
Cut (!) Example 5

```
rainy(seattle) .  
rainy(rochester) .  
cold(rochester) .  
snowy(X) :- rainy(X) , cold(X) , ! .
```

```
?- snowy(C) .
```

Cut (!) Example 5

```
rainy(seattle).  
rainy(rochester).  
cold(rochester).  
snowy(X) :- rainy(X), cold(X), !.
```



Negation by Failure: `not (X)` , `\+ (X)`

- **`not (C)`** succeeds when `C` fails
 - Called **negation by failure**, defined:
`not (X) :- X, !, fail.`
`not (_).`
- Not the same as negation in logic $\neg X$!
- In Prolog, we can assert that something is true, but we **cannot assert that something is false**

Exercise

takes(jane, his).

takes(jane, cs).

takes(ajit, art).

takes(ajit, cs).

classmates(X,Y) :- takes(X,Z), takes(Y,Z).

?- classmates(jane,Y).

What are the bindings of **Y**?

How can we change rule **classmates(X,Y)** to prevent binding **Y = jane**?

Exercise

- $p(X) \text{ :- } q(X), \text{ not}(r(X)) .$
 $r(X) \text{ :- } w(X), \text{ not}(s(X)) .$
 $q(a) . \quad q(b) . \quad q(c) .$
 $s(a) . \quad s(c) .$
 $w(a) . \quad w(b) .$

- Evaluate:
 - $?- p(a) .$
 - $?- p(b) .$
 - $?- p(c) .$

Lecture Outline

- Prolog
 - Imperative control flow
 - Negation by failure
 - Generate and test paradigm

Generate and Test Paradigm

- Search in space
- Prolog rules to **generate** potential solutions
- Prolog rules to **test** potential solutions for desired properties
- Easy prototyping of search
`solve(P) :- generate(P), test(P).`

A Classical Example: n Queens

- Given an n by n chessboard, place each of n queens on the board so that no queen can attack another in one move
 - Queens can move either vertically,
 - horizontally, or
 - diagonally.
- A classical generate and test problem

n Queens

```
my_not(X):- X, !, fail.    %same as not
my_not(_).
in(H, [H|_]).              %same as member
in(H, [_|T]):- in(H,T).
```

```
nums(H,H,[H]).
```

```
nums(L,H,[L|R]):- L<H, N is L+1, nums(N,H,R).
```

%%%nums generates a list of integers between two other numbers, L,H by putting the first number at the front of the list returned by a recursive call with a number 1 greater than the first. It only works when the first argument is bound to an integer. It stops when it gets to the higher number

```
queen_no(4).
```

%%%The number of queens/size of board - use 4

n Queens (ii)

ranks(L):- queen_no(N), nums(1,N,L).

files(L):- queen_no(N), nums(1,N,L).

%%%ranks and files generate the x and y axes of the chess board. Both are lists of numbers up to the number of queens; that is, ranks(L) binds L to the list [1,2,3,...,#queens].

rank(R):- ranks(L), in(R,L).

%%% R is a rank on the board; selects a particular rank R from the list of all ranks L.

file(F):- files(L), in(F,L).

%%% F is a file on the board; selects a particular file F from the list of all files L.

n Queens (iii)

*%%% Squares on the board are (rank,file) coordinates.
attacks decides if a queen on the square at rank R1,
file F1 attacks the square at rank R2, file F2 or
vice versa. A queen attacks every square on the same
rank, the same file, or the same diagonal.*

attacks((R,_),(R,_)).

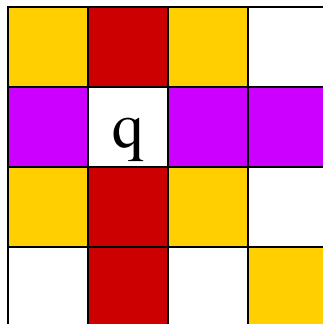
attacks((_,F), (_,F)). *%a Prolog tuple*

attacks((R1,F1), (R2,F2)):-

diagonal((R1,F1),(R2,F2)).

%%%can decompose a Prolog tuple by unification

*(X,Y)=(1,2) results in X=1,Y=2; tuples have fixed
size and there is not head-tail type construct for
tuples*



same rank

same file

same diagonal

**What is safe placement
for next queen on board?**

n Queens (iv)

%%% Two squares are on the same diagonal if the slope of the line between them is 1 or -1. Since / is used, real number values for 1 and -1 are needed.

diagonal((X,Y),(X,Y)). *%degenerate case*

**diagonal((X1,Y1),(X2,Y2)):-N is Y2-Y1,D is X2-X1,
Q is N/D, Q is 1 .** *%diagonal needs bound
arguments!*

**diagonal((X1,Y1),(X2,Y2)):-N is Y2-Y1,D is X2-X1,
Q is N/D, Q is -1 .**

%%%because of use of "is", diagonal is NOT invertible.

n Queens (v)

```
%%% This solution works by generating every list of  
squares, such that the length of the list is the same  
as the number of queens, and then checks every list  
generated to see if it represents a valid placement of  
queens to solve the N queens problem;  
assume list length function
```

```
queens(P) :- queen_no(N), length(P,N),  
    placement(P), ok_place(P).
```

“generate” code given first

“test” code follows

n Queens (vi)

%%%placement can be used as a generator. If placement is called with a free variable, it will construct every possible list of squares on a chess board.

The first predicate will allow it to establish the empty list as a list of squares on the board. The second predicate will allow it to add any (R,F) pair onto the front of a list of squares if R is a rank of the board and F is a file of the board.

placement first generates all 1 element lists, then all 2 element lists, etc. Switching the order of predicates in the second clause will cause it to try varying the length of the list before it varies the squares added to the list

placement([]).

placement([(R,F)|P]) :- placement(P), rank(R), file(F).

n Queens (vii)

```
%%%these two routines check the placement of the next
queen
%%%Checks a list of squares to see that no queen on
any of them would attack any other. does by checking
that position j doesn't conflict with positions
(j+1),(j+2) etc.
ok_place([]).
ok_place([(R,F)|P]):- no_attacks((R,F),P),ok_place(P).
%%% Checks that a queen at square (R,F) doesn't attack
any square (rank,file pair) in list L; uses attacks
predicate defined previously
no_attacks(_,[]).
no_attacks((R,F),[(R2,F2)|P]):-
    my_not(attacks((R,F),(R2,F2))), no_attacks((R,F),P).
```

Solution Structure

- Typical Prolog homework: search in space (e.g., paths in a maze, paths in graph, parsing sequences, various puzzles)
- Typical solution:

```
search(F,Partial,Total) :-
```

```
    final(F), ... % get Total from Partial
```

```
search(C,Partial,Total) :-
```

```
    generate(C,N), % generate next position
```

```
    valid(N),... % test if N is a valid position
```

```
    augment(Partial,New_partial),
```

```
% augment Partial solution with N, typically  
    we would need not(member(N,Partial)) too.
```

```
search(N,New_partial,Total).
```

A Harder Exercise

- Remember the grammar...
 1. $S \rightarrow aSbS$
 2. $S \rightarrow bSaS$
 3. $S \rightarrow \varepsilon$
- Write a **top-down depth-first** parser in Prolog:
?- parse([a,b,a,b],R) .
R = [1, 2, 3, 3, 3] ; // seq. of productions
R = [1, 3, 1, 3, 3] ; // different seq
false. // no more seqs
- Hint: break list into constituent parts

The End
