

IOB-CACHE

User Guide, 0.1 , Build 4d26003



May 4, 2022



Contents

1	Introduction	5
1.1	Features	5
1.2	Benefits	5
1.3	Deliverables	5
2	Description	6
2.1	Block Diagram	6
2.2	Configuration	6
2.2.1	Macros	6
2.2.2	Parameters	7
2.3	Interface Signals	9
3	Usage	9
3.1	Simulation	11
3.2	Synthesis	12

List of Tables

1	Synthesis Macros.	7
2	Synthesis Parameters.	8
3	General Interface Signals	9
4	IObundle Master Interface Signals	9
5	Control-Status Interface Signals	9
6	IObundle Slave Interface Signals	9

List of Figures

1	IP Core Symbol.	5
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2	High-Level Block Diagram.	6
3	Core Instance and Required Surrounding Blocks	10
4	Testbench Block Diagram	11

1 Introduction

The IObundle CACHE is an open-source pipelined-memory cache. It is a performance-wise and highly configurable IP core. The cache core is isolated from the processor and memory interfaces in order to make it easy to adopt new processors or memory controllers while keeping the core functionality intact. It implements a simple front-end native interface. It also implements an AXI4 interface with configurable data width which allows maximum use of the available memory bandwidth. The IObundle CACHE can be implemented as a Direct-Mapped cache or K-Way Set-Associative cache. It supports both fixed write-through not-allocate policy and write-back policy.

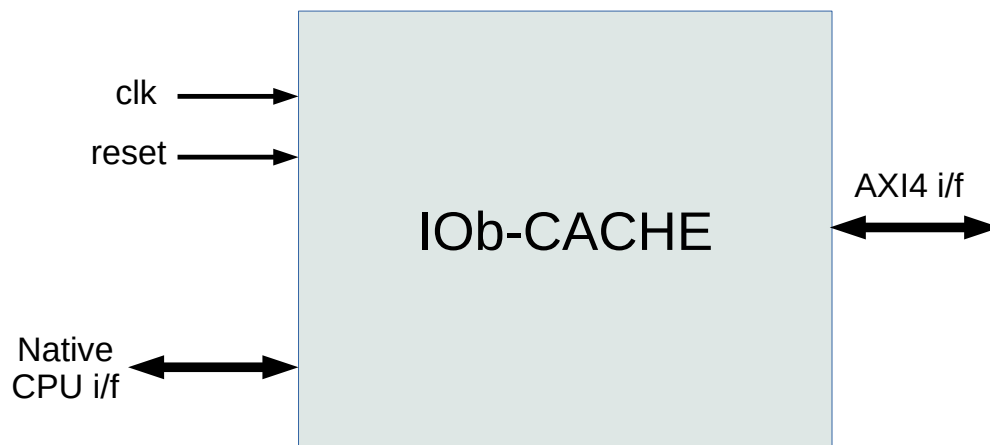


Figure 1: IP Core Symbol.

1.1 Features

- Pipelined-memory (1 request/clock-cycle)
- AXI4 interface with configurable data width
- Simple front-end native interface
- Direct-Mapped or K-Way Set-Associative
- Fixed write-through not-allocate policy
- Write-back policy

1.2 Benefits

- Compact and easy to integrate hardware and software implementation
- Can fit many instances in low cost FPGAs and ASICs
- Low power consumption

1.3 Deliverables

- ASIC or FPGA synthesized netlist or Verilog source code, and respective synthesis and implementation scripts

- ASIC or FPGA verification environment by simulation and emulation
- Bare-metal software driver and example user software
- User documentation for easy system integration
- Example integration in IOb-SoC (optional)

2 Description

2.1 Block Diagram

Figure 2 presents a high-level block diagram of the core, followed by a brief description of each block.

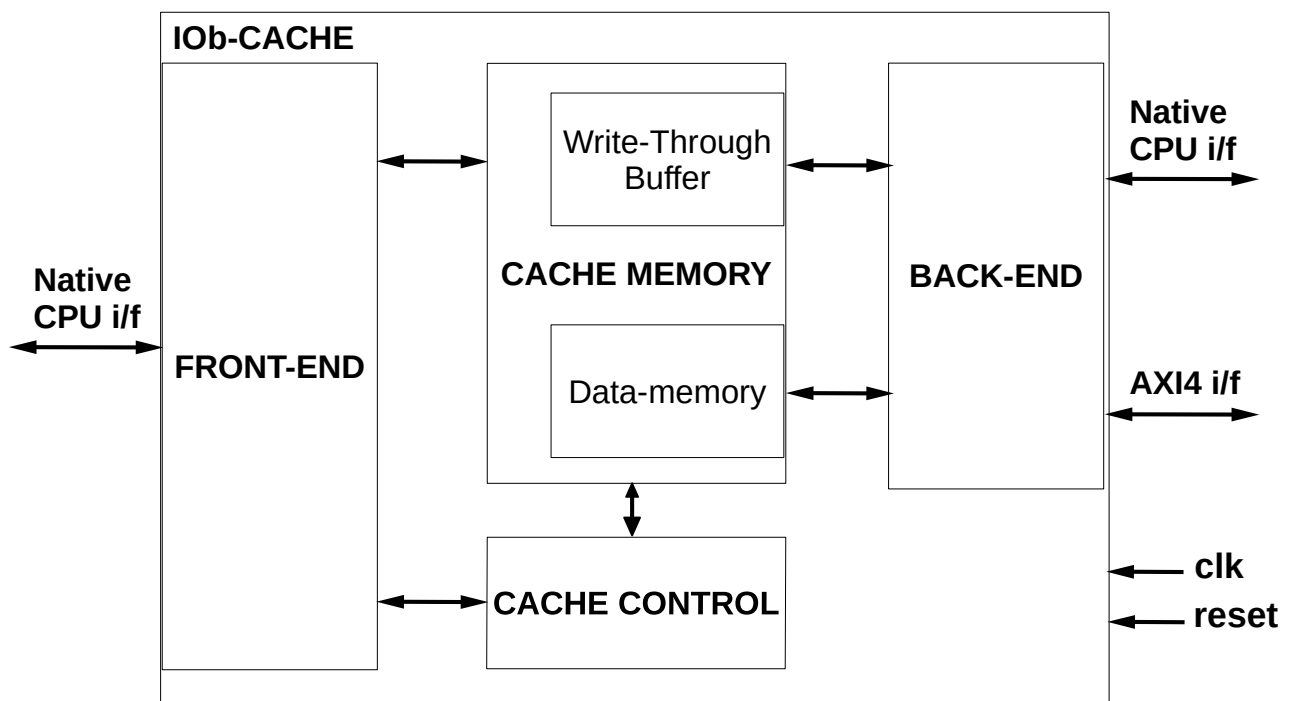


Figure 2: High-Level Block Diagram.

FRONT-END Front-end block.

CACHE MEMORY Cache memory block.

BACK-END Back-end block.

CACHE CONTROL Cache control block.

2.2 Configuration

2.2.1 Macros

The configuration macros apply to all instances of the core, and are listed in Table 1.

Parameter	Min	Typ	Max	Description
REP_POLICY	?	PLRU_mru	?	LRU - Least Recently Used; PLRU_mru (1) - mru-based pseudoLRU; PLRU_tree (3) - tree-based pseudoLRU
WRITE_POL	?	WRITE_THROUGH	?	write policy: write-through (0), write-back (1)

Table 1: Synthesis Macros.

2.2.2 Parameters

The configuration parameters of the core are presented in Table 2. Configuration parameters can vary from instance to instance.

Parameter	Min	Typ	Max	Description
FE_ADDR_W	?	32	?	Address width - width of the Master's entire access address (including the LSBs that are discarded, but discarding the Controller's)
FE_DATA_W	?	32	?	Data width - word size used for the cache
N_WAYS	?	2	?	Number of Cache Ways (Needs to be Potency of 2: 1, 2, 4, 8, ..)
LINE_OFF_W	?	7	?	Line-Offset Width - 2**NLINE_W total cache lines
WORD_OFF_W	?	3	?	Word-Offset Width - 2**OFFSET_W total FE_DATA_W words per line - WARNING about LINE2MEM_W (can cause word_counter [-1:0])
WTBUF_DEPTH_W	?	5	?	Depth Width of Write-Through Buffer
NWAY_W	?	clog2(N_WAYS)	?	Cache Ways Width
FE_NBYTES	?	FE_DATA_W/8	?	Number of Bytes per Word
FE_BYTE_W	?	clog2(FE_NBYTES)	?	Byte Offset
BE_ADDR_W	?	FE_ADDR_W	?	Address width of the higher hierarchy memory
BE_DATA_W	?	FE_DATA_W	?	Data width of the memory
BE_NBYTES	?	BE_DATA_W/8	?	Number of bytes
BE_BYTE_W	?	clog2(BE_NBYTES)	?	Offset of Number of Bytes
LINE2MEM_W	?	WORD_OFF_W-clog2(BE_DATA_W/FE_DATA_W)	?	Logarithm Ratio between the size of the cache-line and the BE's data width
CTRL_CACHE	?	0	?	Adds a Controller to the cache, to use functions sent by the master or count the hits and misses
CTRL_CNT	?	0	?	Counters for Cache Hits and Misses - Disabling this and previous, the Controller only store the buffer states and allows cache invalidation

Table 2: Synthesis Parameters.

2.3 Interface Signals

Name	Direction	Width	Description
clk	INPUT	1	System clock input
reset	INPUT	1	System reset, asynchronous and active high

Table 3: General Interface Signals

Name	Direction	Width	Description
valid	INPUT	1	Native CPU interface valid signal
addr	INPUT	CTRL_CACHE + FE_ADDR_W - FE_BYTE_W	Native CPU interface address signal
addr	INPUT	CTRL_CACHE + FE_ADDR_W	Native CPU interface address signal
wdata	INPUT	FE_DATA_W	Native CPU interface data write signal
wstrb	INPUT	FE_NBYTES	Native CPU interface write strobe signal
rdata	OUTPUT	FE_DATA_W	Native CPU interface read data signal
ready	OUTPUT	1	Native CPU interface ready signal

Table 4: IObundle Master Interface Signals

Name	Direction	Width	Description
force_inv_in	INPUT	1	force 1'b0 if unused
force_inv_out	OUTPUT	1	cache invalidate signal
wtb_empty_in	INPUT	1	force 1'b1 if unused
wtb_empty_out	OUTPUT	1	write-through buffer empty signal

Table 5: Control-Status Interface Signals

Name	Direction	Width	Description
mem_valid	OUTPUT	1	Native CPU interface valid signal
mem_addr	OUTPUT	BE_ADDR_W	Native CPU interface address signal
mem_wdata	OUTPUT	BE_DATA_W	Native CPU interface data write signal
mem_wstrb	OUTPUT	BE_NBYTES	Native CPU interface write strobe signal
mem_rdata	INPUT	BE_DATA_W	Native CPU interface read data signal
mem_ready	INPUT	1	Native CPU interface ready signal

Table 6: IObundle Slave Interface Signals

3 Usage

Figure 4 illustrates how to instantiate the IP core and, if applicable, the required external blocks.

bla bla bla...

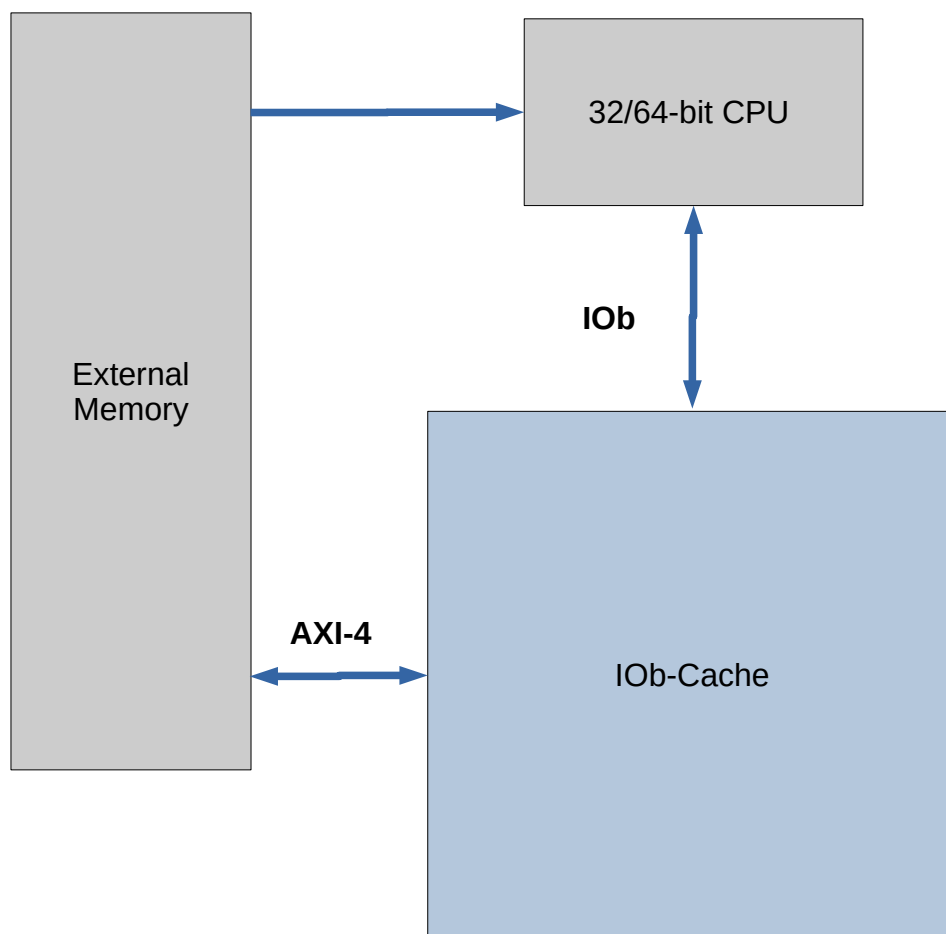


Figure 3: Core Instance and Required Surrounding Blocks

3.1 Simulation

The provided testbench uses the core instance described in Section 3. A high-level block diagram of the testbench is shown in Figure 4. The testbench is organized in a modular fashion, with each test described in a separate file. The test suite consists of all the test case files to make adding, modifying, or removing tests easy.

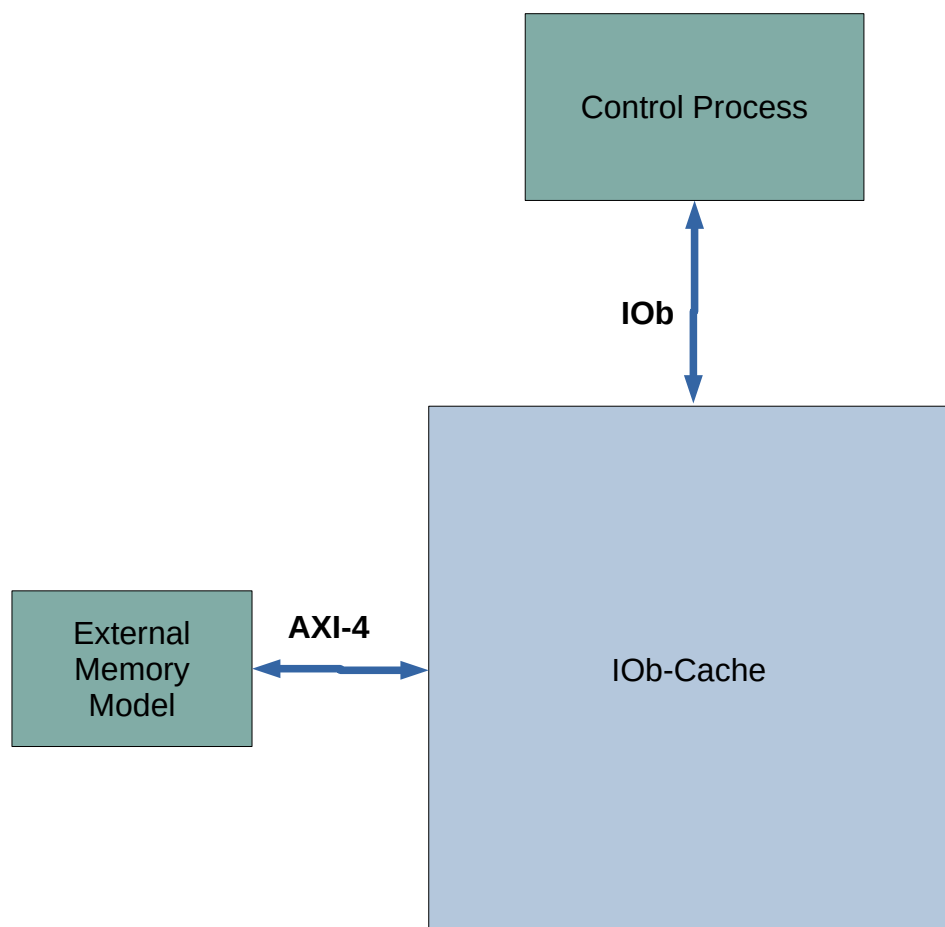


Figure 4: Testbench Block Diagram

In this preliminary version, simulation is not yet fully functional. The provided testbench merely allows compilation for simulation, and drives the clock and reset signals. Behavioural memory models to allow pre-synthesis simulation are already included. In the case of ROMs, their programming data is also included in the form of .hex files.

3.2 Synthesis

A simple `.tcl` script is provided for the Cadence Genus synthesis tool. The script reads the technology files, compiles and elaborates the design, and proceeds to synthesise it. The timing constraints are contained within the constraints file provided, or provided in a separate file.

After synthesis, reports on silicon area usage, power consumption, and timing closure are generated. A post-synthesis Verilog file is created, to be used in post-synthesis simulation.

In this preliminary version, synthesis of the IP core without the memories is functional. The memories are for now treated as black boxes.

It is important not to include the memory models provided in the simulation directory in synthesis, unless they are to be synthesised as logic. In a next version, the memories will be generated for the target technology and included.