## Tutorial 4 11.02.20

## 1 Applications of Gaussian elimination

For this exercise, K is a field, and we consider an ambient linear space  $K^n$  for some  $n \geq 2$ . All vectors will be row vectors.

- 1. Let  $V = \text{Span}\{\mathbf{v}_1, \dots, \mathbf{v}_m\}$  be a linear subspace. Give an algorithm to compute a basis of V.
- 2. Let  $W = \text{Span}\{\mathbf{w}_1, \dots, \mathbf{w}_e\}$  be another linear subspace. Give an algorithm to compute a basis of V + W.
- 3. Give an algorithm to compute a basis of  $V \cap W$ .

## 2 An algorithm for computing the characteristic polynomial

Let  $A \in \mathcal{M}_n(\mathbb{K})$ , the goal of the following method is to compute the characteristic polynomial of A in  $O(n^3)$  time.

- 1. Let T be the transformation which acts on the left of a matrix A through  $L_i \leftarrow L_i + \alpha L_j$ , i.e.,  $T = I_n + \alpha E_{i,j}$ . Here  $E_{i,j}$  denotes an  $n \times n$  matrix with 1 on the (i,j) position and 0s everywhere else. Describe the action of  $T^{-1}$  on the right of A in terms of column operations.
- 2. Using Question 1, show that one can find a matrix R such that

$$RAR^{-1} = \begin{bmatrix} a_{1,1} & a'_{1,2} & \cdots & a'_{1,n} \\ l_2 & a'_{2,2} & \ddots & a'_{2,n} \\ 0 & a'_{3,2} & \ddots & a'_{3,n} \\ \vdots & \vdots & \ddots & \vdots \\ 0 & a'_{n,2} & \cdots & a'_{n,n} \end{bmatrix}.$$

(Hint: perform row operations by multiplying on the left by some transformation matrices  $T_i$  and see what happens on the columns when you multiply on the right by  $T_i^{-1}$ ).

3. Give an algorithm to compute the matrices  $R_n$  and M such that

$$R_n A R_n^{-1} = M = \begin{bmatrix} m_{1,1} & m_{1,2} & m_{1,3} & \cdots & m_{1,n} \\ \ell_2 & m_{2,2} & m_{2,3} & \ddots & m_{2,n} \\ 0 & \ell_3 & m_{3,3} & \ddots & m_{3,n} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & \cdots & 0 & \ell_n & m_{n,n} \end{bmatrix}$$

using  $O(n^3)$  operations in  $\mathbb{K}$ .

Remark: such an "almost triangular" shape matrix is called an upper Hessenberg matrix, i.e., a matrix that has zero entries below the first subdiagonal. We have shown how to reduce any matrix into (upper) Hessenberg form.

- 4. Deduce an algorithm to compute the characteristic polynomial of A, with a complexity bound  $O(n^3)$ . Use the fact that two similar matrices have the same characteristic polynomial.
- 5. Could it be possible to find R such that  $R^{-1}AR = M$  is upper triangular by (arbitrarily many) elementary operations in  $\mathbb{K}$ ? If yes, explain how. If not, explain why.

## 3 A faster algorithm for characteristic polynomial

Let A be an  $n \times n$  matrix. In this exercise, we will denote by  $n^{\omega}$  the number of operations in K needed to multiply two n by n matrices with coefficients in K. You have seen in class that given a n by n matrix  $M \in \mathcal{M}_n(K)$  (if you have not, take it as a fact), we can compute  $M^{-1}$  using  $O(n^{\omega} \log n)$  operations in K (computing the inverse is asymptotically the same as multiplying).

1. Assume that v is a vector such that v, Av,  $A^2v$ , ...,  $A^{n-1}v$  is a basis of  $K^n$ ; then if B is the matrix with columns v, Av,  $A^2v$ , ...,  $A^{n-1}v$ , prove that  $B^{-1}AB$  is a companion matrix, that is, a matrix of the following form

$$C = \begin{bmatrix} 0 & 0 & \cdots & 0 & c_0 \\ 1 & 0 & \cdots & 0 & c_1 \\ 0 & 1 & \cdots & 0 & c_2 \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & \cdots & 1 & c_{n-1} \end{bmatrix}.$$

- 2. If B is given, what is the cost of computing the characteristic polynomial of A using the previous question.
- 3. Explain how from an  $n \times n$  matrix multiplication in time  $O(n^{\omega})$  we can deduce a  $n \times m$  by  $m \times k$  matrix multiplication algorithm in time  $O(\max(n, m, k)^{\omega})$
- 4. Define  $w_0 = v, w_1 = (v, Av), w_2 = (v, Av, A^2v, A^3v), \dots, w_k = (v, Av, A^2v, \dots, A^{2^k-1}v)$ Prove that  $w_k$  can be computed in time  $O(kn^{\omega})$  for  $k < \log n$ .
- 5. Under the assumption that v exists and that you know it, give a  $O(n^{\omega} \log n)$  algorithm for computing the characteristic polynomial of a square matrix.
- 6. Does there always exist a v as in Question 1?