

go A S

...while reviewing our postal blueprints and debating  
 latest fashion trend of "Leg Warmers." F asked me an  
 question. He said that the plans in these blueprints were  
 obviously complex, and he wondered if anyone else had  
 come up with this idea.

I internally debated whether I should tell him  
 about my Muse. F is a very superstitious  
 man—he crosses himself when he  
 walks over graves, and chastises  
 me for saying, "What the  
 Devil!" Although I have  
 always wanted to tell someone  
 about my divine experience,  
 I worry that he might think  
 I've gone mad all these years  
 in seclusion, or worse—that I'm  
 tangled in some kind of  
 unsavory black magic.

No matter. I told him that  
 with hard work anything is



# CODES:



CAESAR

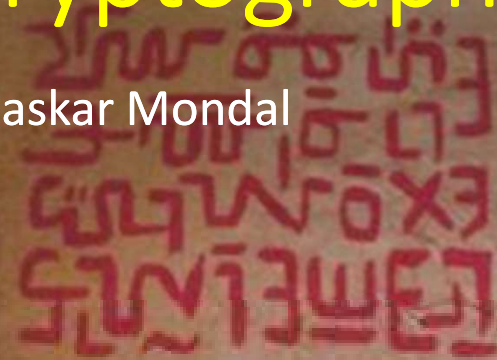


ATBASH

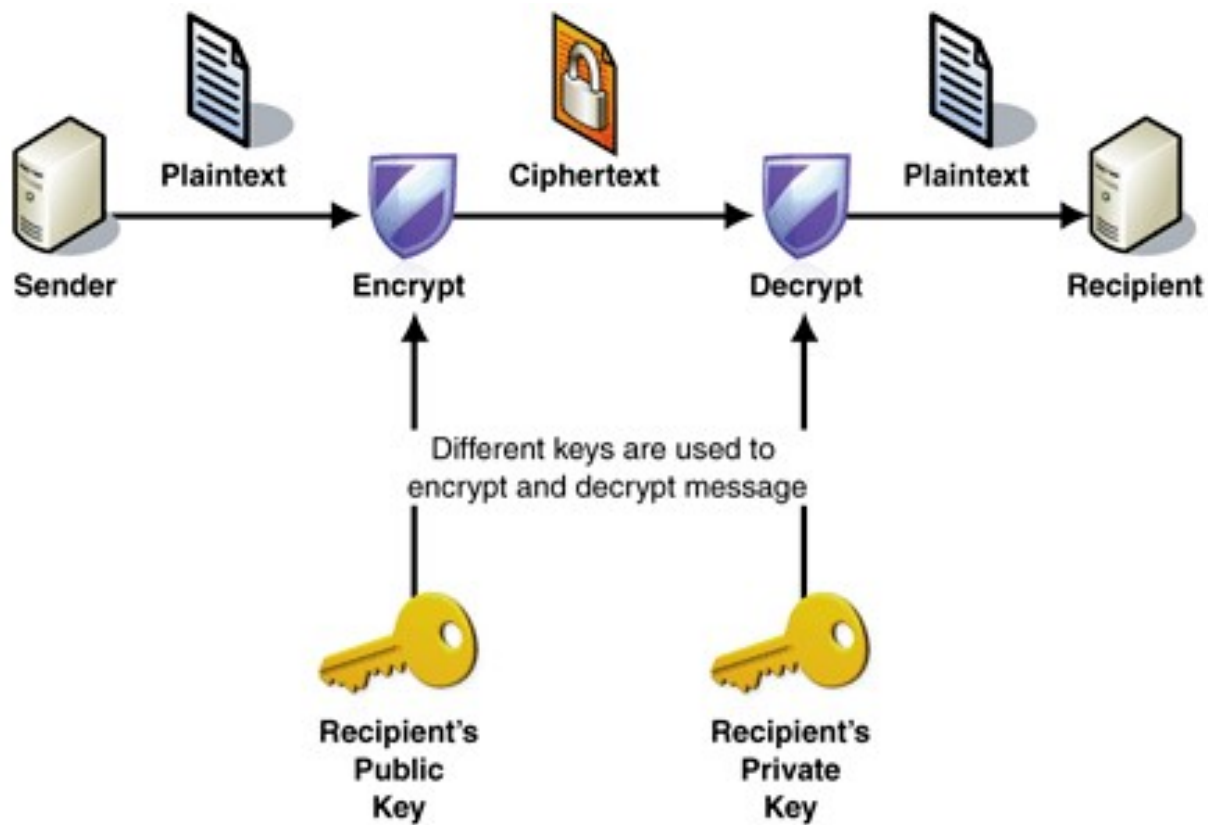


## Classic Cryptography

Dr. Bhaskar Mondal



It occurs to me that if I must keep secrets from F, I might  
 as well begin writing certain passages of this book in code.  
 I used Cryptology in college, so this will be fun! (Aft



Components  
of a  
Cryptosystem

# Ancient Ciphers

- 3100 B.C ancient Egyptians enciphered some of their *hieroglyphic* writing on monuments
- 100 BC years ago Julius Ceasar used a simple substitution cipher, now known as the **Caesar cipher**
- 1200: Roger Bacon described several methods
- 1392: Geoffrey Chaucer included several ciphers in his works
- 1460s: Leon Alberti devised a **cipher wheel**, and described the principles of frequency analysis
- 1553: Vigenère invented by Giovan Batista Belaso
- 1585: Blaise de Vigenère published a book on cryptology, and described the **polyalphabetic substitution cipher**

# Classical Cryptographic Techniques



have two basic components of classical ciphers: substitution and transposition



**substitution** ciphers: letters are replaced by other letters



**transposition** ciphers: the letters are arranged in a different order these ciphers may be:

**monoalphabetic** - only one substitution/ transposition is used, or  
**polyalphabetic** - where several substitutions/ transpositions are used



several such ciphers may be concatenated together to form a **product cipher**



## Substitution Cipher: Atbash

- Alphabets are reversed
- Plain Text: This is a Cipher
- Cipher Text: Gsrh Rh Z Xrksvi

# Shift Cipher: Caesar Cipher

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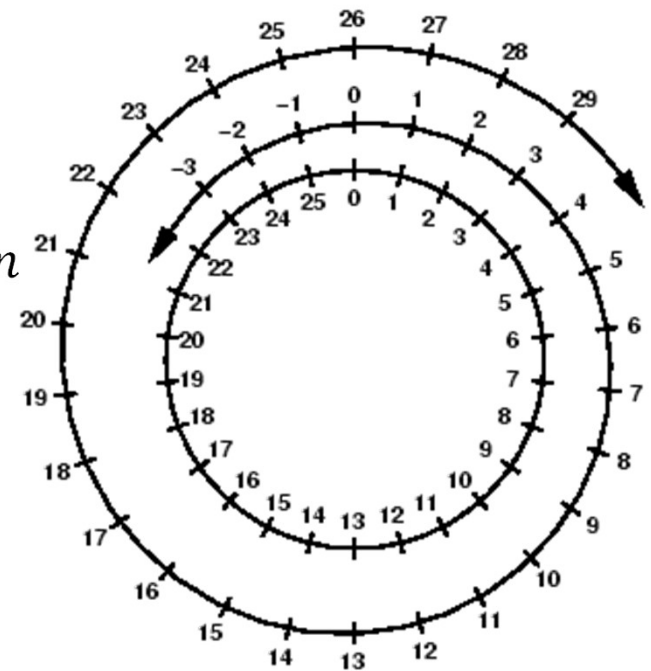
- a monoalphabetic Cipher
- reputedly used by Julius Caesar
- replace each letter of message by a letter a fixed distance away e.g. use the 3rd letter on
- $E_n(x) = (x + n) \bmod 26$
- $D_n(x) = (x - n) \bmod 26$
- Text : ATTACKATONCE
- Shift: 4
- Cipher: EXXEGOEXSRGI



# Modular Arithmetic monoalphabetic cipher

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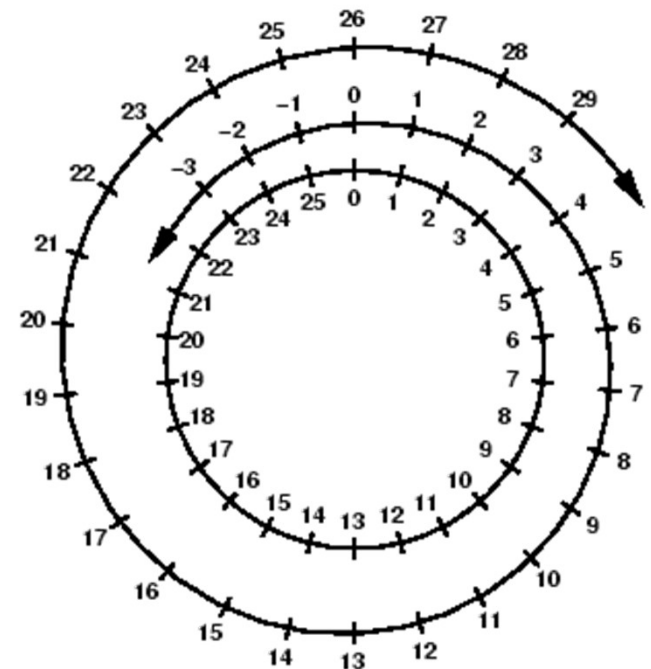
- If  $a \bmod n = b \bmod n$  then there is an integer  $k$ 
  - such that  $a - b = k \cdot n$
- $(a + b) \bmod n = ((a \bmod n) + (b \bmod n)) \bmod n$
- $(a \cdot b) \bmod n = ((a \bmod n) \cdot (b \bmod n)) \bmod n$
- $-a \bmod n = n - (a \bmod n)$





# Modular Arithmetic monoalphabetic cipher

- If  $a \bmod n = b \bmod n$  then there is an integer  $k$ 
  - such that  $a - b = k \cdot n$
- Say  $a = 17, n = 5$
- $17 \bmod 5 = 2$
- $(17 - 2) = 3 \times 5$  where  $b = 2$  and  $k = 3$

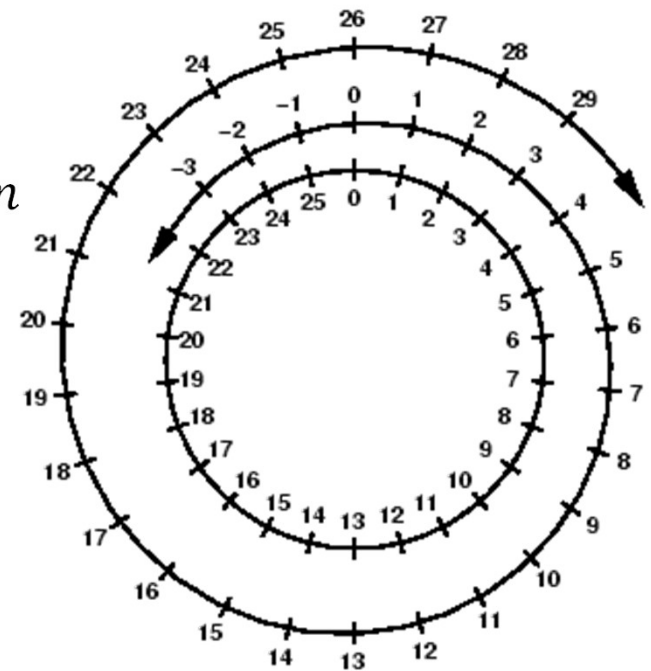




# Modular Arithmetic monoalphabetic cipher

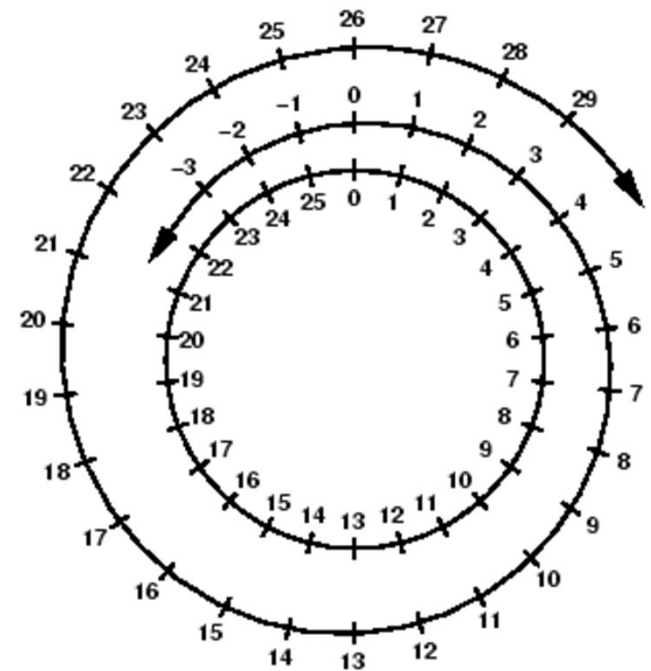
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- $(a + b) \bmod n = ((a \bmod n) + (b \bmod n)) \bmod n$
- Say  $a = 17, b = 3, n = 5$
- $((17 \bmod 5) + (3 \bmod 5)) \bmod 5$
- $= (2 + 3) \bmod 5 = 0$



# Modular Arithmetic monoalphabetic cipher

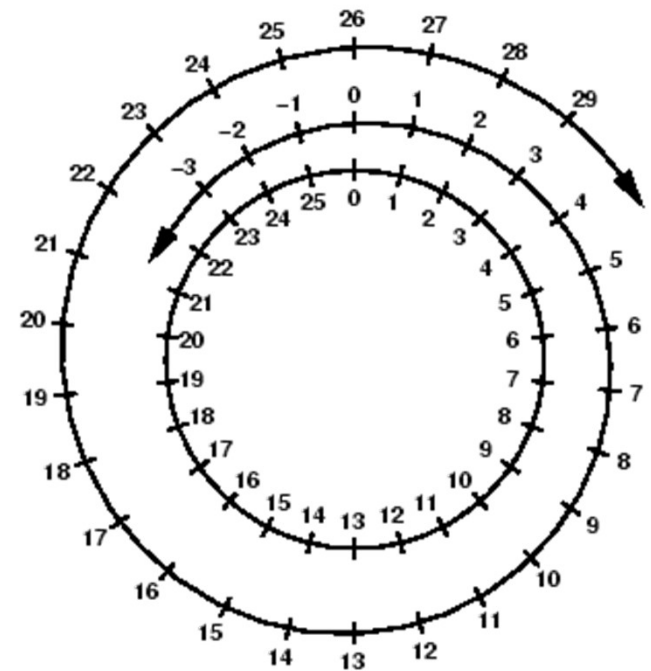
- $(a \cdot b) \bmod n = ((a \bmod n) \cdot (b \bmod n)) \bmod n$
- Say  $a = 17, b = 3, n = 5$
- $((17 \bmod 5) \cdot (3 \bmod 5)) \bmod 5 = (2 \times 3) \bmod 5 = 1$



# Modular Arithmetic monoalphabetic cipher

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- $-a \bmod n = n - (a \bmod n)$
- Say  $a = 17, n = 5$
- $5 - (17 \bmod 5) = 5 - 2 = 3$





# Cryptanalysis of the Caesar Cipher

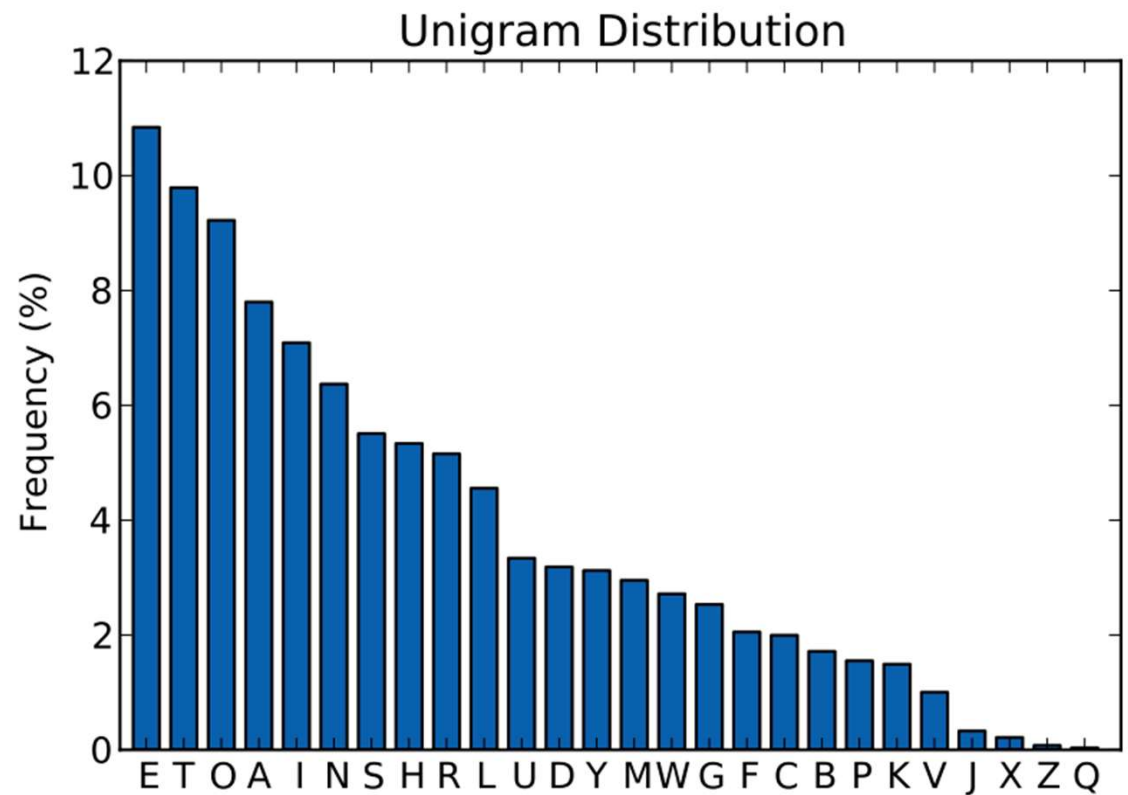
- only have 26 possible ciphers
- could simply try each in turn - exhaustive key search

# Cryptanalysis of the Caesar Cipher

Single Letter	Double Letter	Triple Letter
E	TH	THE
T	HE	AND
R	IN	TIO
N	ER	ATI
I	RE	FOR
O	ON	THA
A	AN	TER
S	EN	RES

# Character Frequencies

- in most languages' letters are not equally common
- in English e is by far the most common letter
- have tables of single double & triple letter frequencies



# Affine Cipher

- transform each of the letter in plaintext to the corresponding integer in the range 0 to  $m-1$  ( $m=26$ )
- Encryption  $E(x) = (ax + b) \bmod m$ 
  - $a$  and  $b$  are the key
- Say key  $a = 5, b = 8$

Plaintext Alphabet	a	b	c	d	e	f	g	h	i	j	k	l	m	n	o	p	q	r	s	t	u	v	w	x	y	z
Plaintext Value	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25



# Affine Cipher

- Encryption  $E(x) = (ax + b) \bmod m$
- key  $a = 5, b = 8$

Plaintext	a	f	f	i	n	e		c	i	p	h	e	r
x	0	5	5	8	13	4		2	8	15	7	4	17
5x+8	8	33	33	48	73	28		18	48	83	43	28	93
(5x+8) mod 26	8	7	7	22	21	2		18	22	5	17	2	15
Ciphertext	I	H	H	W	V	C		S	W	F	R	C	P

## Affine Cipher: Decryption

- Decryption  $D(x) = c(x - b) \bmod m$
- $c$  is the modular multiplicative inverse of  $a$ :  $a \times c = 1 \bmod m$
- key  $a = 5, b = 8$
- $C = 21$
- since  $21 \times 5 = 105 = 1 \bmod 26$ 
  - As  $26 \times 4 = 104$ , and  $105 - 104 = 1$

# Affine Cipher: Decryption

- Decryption  $D(x) = c(x - b) \bmod m$
- $c$  is the modular multiplicative inverse of  $a$ :  
 $a \times c = 1 \bmod m$
- key  $a = 5, b = 8$
- $C = 21$

Ciphertext	I	H	H	W	V	C		S	W	F	R	C	P
$y$	8	7	7	22	21	2		18	22	5	17	2	15
$21(y - 8)$	0	-21	-21	294	273	-126		210	294	-63	189	-126	147
$21(y - 8) \bmod 26$	0	5	5	8	13	4		2	8	15	7	4	17
Plaintext	a	f	f	i	n	e		c	i	p	h	e	r

# Affine Cipher: Drawbacks

- key creates the situation where more than one plaintext letter is encrypted to the same ciphertext letter (for example, above both "e" and "r" encipher to "V")
- there is more than one number that can be multiplied by 4 to get 1 modulo 26.
- Solution: The inverse of  $a$  modulo  $m$  exists if and only if  $a$  and  $m$  are coprime  $\gcd(a, m) = 1$

## Affine Cipher: Drawbacks keys

- there are 12 numbers less than 26 which are coprime to 26
- for each of these there are 26 possibilities for the value of  $b$
- total of  $12 \times 26 = 312$  possible keys. Brute Force Attack!!!

# Polyalphabetic Substitution



in general use more than one substitution alphabet



makes cryptanalysis harder since have more alphabets to guess



and because flattens frequency distribution



(since same plaintext letter gets replaced by several ciphertext letter, depending on which alphabet is used)

# Vigenère Cipher



basically multiple caesar ciphers



key is multiple letters long  $K = k_{(1)} k_{(2)} \dots k_{(d)}$



ith letter specifies ith alphabet to use



use each alphabet in turn, repeating from start after  $d$  letters in message



## Many Caesars: the Vigenère CIPHER

- Giovan Batista Belaso in 1553
- Key=jasmine (9, 0, 18, 12, 8, 13, 4)
- Plain text= we are getting hungry

	w	e	a	r	e	g	e	t	t	i	n	g	h	u	n	g	r	y
	22	4	0	17	4	6	4	19	19	8	13	6	7	20	13	6	17	24
+	9	0	18	12	8	13	4	9	0	18	12	8	13	8	4	9	0	18
=	5	4	18	3	12	19	8	2	19	0	25	14	20	2	17	15	17	16
	F	E	A	D	M	T	I	C	T	A	Z	O	U	C	R	P	R	Q

# Vigenère Cipher: A tabula recta

- how people in the sixteenth century used the cipher
- 26 x 26 table listing all the shifts
- plaintext letter gives the row, and the key letter gives the column
- Apply numerical conversion and the modular addition

	a	b	c	d	e	f	g	h	i	j	k	l	m	n	o	p	q	r	s	t
a	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T
b	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U
c	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V
d	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W
e	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X
f	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y
g	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z
h	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A
i	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B
j	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C
k	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D
l	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E
m	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F
n	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G
o	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H
p	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I
q	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J
r	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K
s	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L
t	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M
u	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N
v	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O
w	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P
x	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q
y	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R
z	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S

# Kasiski examination

- For about 300 years, it was believed to be unbreakable until the 1920s
- Charles Babbage and Friedrich Kasiski independently determined a method of breaking it in the middle of the nineteenth century.
- repeated patterns in the text to determine the length of the key

# Mixed Alphabets Cipher

- most generally we could use an arbitrary mixed (jumbled) alphabet
- each plaintext letter is given a different random ciphertext letter, hence key is 26 letters long
- Plain: ABCDEFGHIJKLMNOPQRSTUVWXYZ
- Cipher: DKVQFIBJWPESCXHTMYAUOLRGZN
- Plaintext: IFWEWISHTOREPLACELETTERS
- Ciphertext: WIRFRWAJUHYFTSDVFSFUUFYA
- now have a total of  $26!$  keys

26! is secure !!

- brute force attack NOT possible!!
- problem is not the number of keys, rather:
  - there is lots of statistical information in message
  - can solve the problem piece by piece

(ie. can get key nearly right, and nearly get message, near enough MUST NOT be good enough!)

# Cryptanalysis



use frequency counts to guess letter by letter



also have frequencies for digraphs & trigraphs

# General Monoalphabetic: A special case of Mix Alphabets

- special form of mixed alphabet
- use key as follows:
  - write key (with repeated letters deleted)
  - then write all remaining letters in columns underneath
  - then read off by columns to get ciphertext equivalents
- Keyword: monoalphabetic
- second "o" is skipped as it has already appeared

Plaintext Alphabet	a	b	c	d	e	f	g	h	i	j	k	l	m	n	o	p	q	r	s	t	u	v	w	x	y	z
Ciphertext Alphabet	M	O	N	A	L	P	H	B	E	T	I	C	D	F	G	J	K	Q	R	S	U	V	W	X	Y	Z

- "u" encrypts to "U", "v" to "V" and so on to "z". This problem occurs if the keyword does not contain any letters from near the end of the plaintext alphabet. To combat this problem, we can choose a keyword with a letter from near the end of the alphabet.



# General Monoalphabetic: A special case of Mix Alphabets

- First we realise that there are 26 possible choices for the first letter in the ciphertext alphabet. Now, for the second letter, we can use any letter *APART* from the letter we have already selected for the first position, so there are 25 choices for the second position. For the third position we can choose any letter, apart from either the letter in the first position or the second position, and hence there are 24 choices here. Thus, for the first 3 places, there are  $26 \times 25 \times 24$  possible choices.
- Continuing in this way, we quickly find that there are 26!



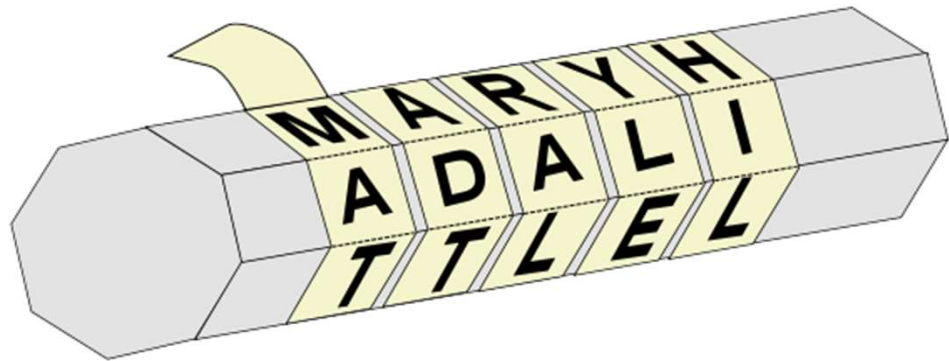
# Transposition Ciphers

# Transposition Ciphers

- transposition or permutation ciphers hide the message contents by rearranging the order of the letters
- Scytale cipher
  - an early Greek transposition cipher
  - a strip of paper was wound round a staff
  - message written along staff in rows, then paper removed
  - leaving a strip of seemingly random letters
  - not very secure as key was width of paper & staff

# Scytale

- used by the Ancient Greeks and Spartans
- consisted of a polygonal rod or cylinder, around which was wrapped a piece of parchment.



# Rail Fence cipher



write message with letters on alternate rows



read off cipher row by row

D		F		N		T		E		A		T		A		L
	E		E		D		H		E		S		W		L	

Plain: defend the east wall

Cipher: DFNTEATALEEDHESWL

Key=1

# Rail Fence cipher

- Plain: defend the east wall
- Key=3
- Cipher: DNETLEEDHESWLXFTAAX

D				N				E				T				L		
	E		E		D		H		E		S		W		L		X	
		F				T				A				A				X

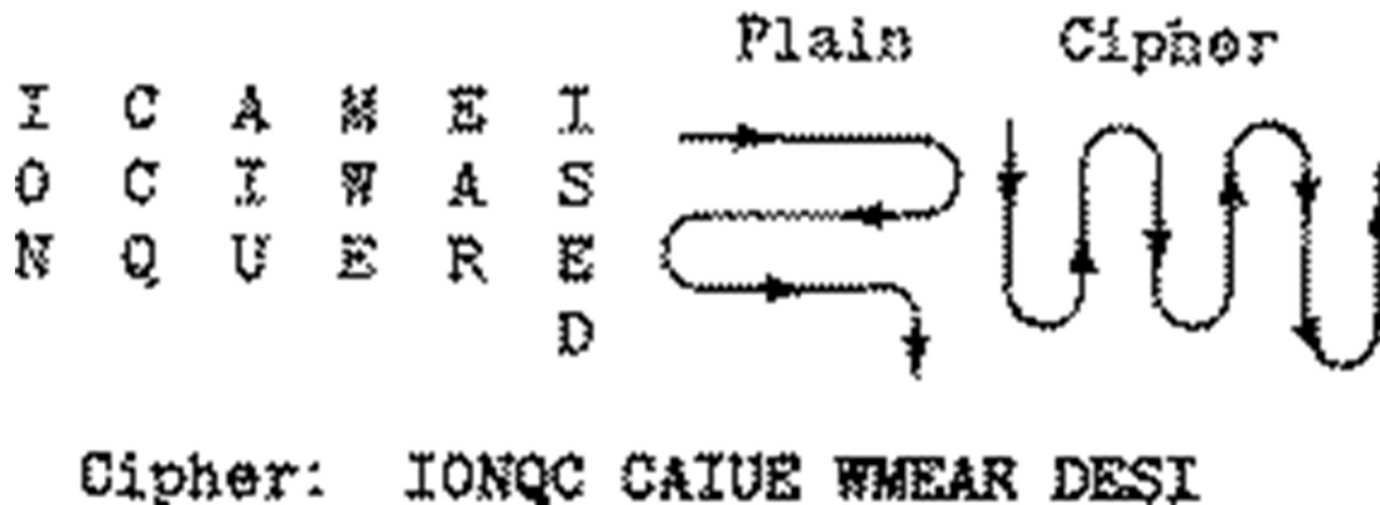
## Rail Fence cipher

- limited number of usable keys, especially for short messages
- for there to be enough movement of letters, the length of the message needs to be at least twice the key, but preferably 3 times the key
- The use of nulls can also have a detrimental effect on the security of the cipher, as an interceptor can use them to identify where the end of the line is, and so have a sensible guess at the key.
- Frequency Analysis



# Geometric Figure/ Route Cipher

- write message following one pattern and read out with another



# Row Transposition ciphers

- in general write message in a number of columns and then use some rule to read off from these columns
- key could be a series of number being the order to: read off the cipher; or write in the plaintext

Plain: THESIMPLESTPOSSIBLETRANSPOSITIONSXX

Key (R): 2 5 4 1 3

Key (W): 4 1 5 3 2

THE S I S T I E H M P L E S E M S L P T P O S S S T S O P I B L E T E I T L B R A  
N S P S R P N A O S I T I T O I I S O N S X X X O X S N

Cipher: S T I E H E M S L P S T S O P E I T L B S R P N A T O I I S X O X S N