

## CSC 225: Lab 6

### In-place Heapsort

You are to implement heapsort in an array given as a max-heap. You may use the programming language of your choice. The heapsort algorithm takes a max-heap input array  $A[1 \dots n]$ , where  $n = A.length$ . Since the maximum element of the array is stored at the root  $A[1]$ , we can put it into its correct final position by exchanging it with  $A[n]$ . If we now discard node  $n$  from the heap—and we can do so by simply decrementing  $A.heap\_size$ —we observe that the children of the root remain max-heaps, but the new root element might violate the max-heap property. All we need to do to restore the max-heap property, however, is call  $MaxHeapify(A, 1)$ , which leaves a max-heap in  $A[1 \dots n - 1]$ . The heapsort algorithm then repeats this process for the max-heap of size  $n - 1$  down to a heap of size 2.

In order to maintain the max-heap property, we call the procedure  $MaxHeapify$ . Its inputs are an array  $A$  and an index  $i$  into the array. When it is called,  $MaxHeapify$  assumes that the binary trees rooted at  $2i$  and  $2i + 1$  are max-heaps, but that  $A[i]$  might be smaller than its children, thus violating the max-heap property.  $MaxHeapify$  lets the value at  $A[i]$  “bubble down” in the max-heap so that the subtree rooted at index  $i$  obeys the max-heap property.

**Testers:** Show TA your results of sorting these two arrays in ascending order:  $A1 = [6, 5, 4, 1, 2, 3]$  and  $A2 = [99, 19, 9, 7, 11, 3, 3, 2, 2, 1]$ .