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11:23  
14 min  
14\*4 = 56 min  
⇒ ok length

University of Victoria  
Faculty of Engineering  
Department of Computer Science  
CSc 230 Computer Architecture and Assembly Language

**MIDTERM EXAM October 19, 2015**

NAME: LillAnne (Print)

ID NUMBER: \_\_\_\_\_

SIGNATURE: \_\_\_\_\_

TIME: 75 minutes

INSTRUCTOR: LillAnne Jackson

| Question     | Value               | Mark |
|--------------|---------------------|------|
| 1            | 1                   |      |
| 2            | 84                  |      |
| 3            | 4                   |      |
| 4            | 4                   |      |
| 5            | 2                   |      |
| 6            | 1                   |      |
| 7            | 2                   |      |
| 8            | 6                   |      |
| 9            | 6                   |      |
| 10           | 8                   |      |
| <b>TOTAL</b> | <del>40</del><br>38 |      |

109 in class  
1 excused late  
6 RCSD

116.

121 = Registered.

STUDENTS MUST CHECK THE NUMBER OF PAGES IN THIS EXAMINATION PAPER BEFORE BEGINNING TO WRITE, AND REPORT ANY DISCREPANCY IMMEDIATELY TO THE INVIGILATOR

- The questions are to be answered on the examination paper.
- The exam is **NOT** open book. Calculators are **NOT** permitted.
- A formula for a calculation is acceptable anywhere a calculation might be required.
- A copy of the *AVR Instruction Set Summary* will be provided with this exam.
- The marks assigned to each question are printed within square brackets.
- There are 8 pages in this document, including this cover page.
- It is strongly recommended that you read the entire exam through from beginning to end before beginning to answer the questions.

1. [1 mark] What is the major distinction between the von Neumann and Harvard architectures?

von Neuman - One bank of memory addresses.  
- shared by code + data.

Harvard - Separate addresses for code memory + data memory.

2. [4 marks] Given two hexadecimal numbers, state their corresponding values :  
a. in binary  
b. as a decimal, if the number is an unsigned integer

| Hex              | Binary (8 bits):       | Decimal Integer: unsigned     |
|------------------|------------------------|-------------------------------|
| 3E <sub>16</sub> | 0011 1110 <sub>2</sub> | 3*16 + 14 = 62 <sub>10</sub>  |
| B8 <sub>16</sub> | 1011 1000 <sub>2</sub> | 11*16 + 8 = 184 <sub>10</sub> |

3. [4 marks] Given two positive decimal numbers which are in the range of 8-bit binary numbers:  
a. provide the binary equivalent of each number  
b. provide the 2's Complement

| Decimal           | Binary (8 bits):     | 2's Complement |
|-------------------|----------------------|----------------|
| 41 <sub>10</sub>  | 0010 1001<br>32 16 8 |                |
| 103 <sub>10</sub> | 0110 0111<br>64 32   |                |

positive,  
so  
same as  
binary

positive,  
so same as  
binary

4. [4 marks] Given the two signed 8-bit binary numbers and their sum below, what would be the value that would be in the C (carry), V (overflow), Z (zero), and N (negative) bits of a typical (i.e., AVR's) Status Register immediately after the sum is created?

|  | Values in Status Register (SR)              |
|--|---|
| $\begin{array}{r} 0110\ 1110 \\ 0101\ 0111 \\ \hline 1100\ 0101 \end{array}$ | C <u>0</u> V <u>1</u> Z <u>0</u> N <u>1</u> |

5. [2 marks] Given the two 8-bit binary numbers below, perform bitwise AND and OR operations (on each bit in the numbers).

|  |   |
|--|---|
| $\begin{array}{r} 0110\ 1110 \\ \text{AND } 0101\ 0111 \\ \hline 0100\ 0110 \end{array}$ | $\begin{array}{r} 0110\ 1110 \\ \text{OR } 0101\ 0111 \\ \hline 0111\ 1111 \end{array}$ |
|--|---|

6. [1 mark] In a system with 24 address lines, what is the address space? (I.e., What is the maximum number of locations that can be addressed in this system?)

$2^{24}$ , from 0 to  $2^{24} - 1$

(An expression or formula, rather than the calculated result is acceptable.)

7. [2 marks]

a. What is the output of an assembler program?

binary numbers that represent machine language

b. What does a two-pass assembler produce on each of its two passes?

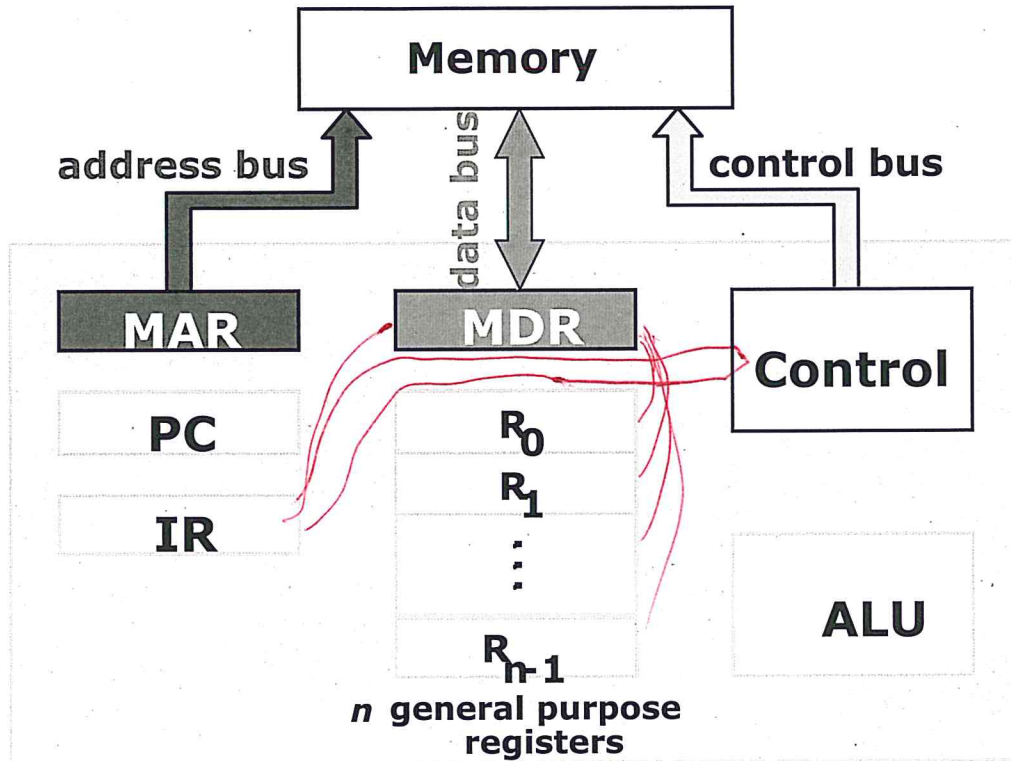
pass 1: a table of the labels + their corresponding addresses

pass 2: The machine language, including the resolution of labels to numbers

8. [6 marks] Consider the diagram below of a (generic) CPU with its internal datapath connections. Use the diagram to give the detailed steps of the Fetch and Decode/Execute cycles of the instruction: LD R0, 0x021F

Note: The instruction uses direct addressing.

→ The address of the data is in the instruction



Fetch:

MAR ← PC  
 address bus ← MAR  
 Control bus ← read signal  
 <<wait for Data>>.  
 data bus ← memory (instruction)  
 MDR ← data bus  
 IR ← MDR

PC update:

ALU ← PC  
 ALU (add).

<The next page is provided to give additional space for you to write this answer>

PC ← ALU



## Decode

Control  $\leftarrow$  IR

$\langle\langle$  Control unit determines opcode + operands  $\rangle\rangle$ .

### 2<sup>nd</sup> Fetch (Required!)

MAR  $\leftarrow$  0x021F (from operands of instruction)

address bus  $\leftarrow$  MAR

Control bus  $\leftarrow$  read signal

$\langle\langle$  wait for Data  $\rangle\rangle$ .

data bus  $\leftarrow$  memory (Data from address 012F)

MDR  $\leftarrow$  data bus

## Execute

RO  $\leftarrow$  MDR.

Example

9. [6 marks] Given the picture of the system's registers, including some initial values in those registers, and some of its memory (below) determine the values in the registers after the execution of the following assembly language code.

```

→ LDI R26, 00 ; X = R27:R26 (High):(Low)
→ LDI R27, 02
→ LD R0, X
→ INC R26
→ CP R0, R1
→ BREQ skip
→ ADD R2, R1
skip: → LSL R3

```

|     |  | Memory Addresses | Memory Content |
|-----|--|------------------|----------------|
| R0  | <div>10<br/>0x00</div>   | 0x0200           | 10             |
| R1  | <div>2F<br/>0x2f</div>   | 0x0201           | 2f             |
| R2  | <div>3F<br/>0x10</div>   | 0x0202           | 10             |
| R3  | <div>3E<br/>0x1f<br/><small>0011 110 = 3E<br/>0001 111</small></div> | 0x0203           | fd             |
| R26 | <div>01<br/>0x10</div>   | 0x0204           | 06             |
| R27 | <div>02<br/>0x10</div>   | 0x0205           | 00             |
|     |  | 0x0206           | 0f             |
|     |  | 0x0207           | de             |

Final Values in Registers:

|    |                 |    |                 |     |                 |
|----|-----------------|----|-----------------|-----|-----------------|
| R0 | <div>0x10</div> | R1 | <div>0x2F</div> | R26 | <div>0x01</div> |
| R2 | <div>0x3F</div> | R3 | <div>0x3E</div> | R27 | <div>0x02</div> |

10. [8 marks] Complete the AVR assembly language program below. It must calculate the sum of the finite geometric series:

$$\sum_{k=0}^7 2^k = 2^0 + 2^1 + 2^2 + 2^3 + 2^4 + 2^5 + 2^6 + 2^7$$

And store the result in the memory (RAM) location called `sum`. Pseudo-code is provided as an algorithm that describes how to complete the code.

```
.include "m2560def.inc"
;This program calculates the sum of a finite geometric series.
;Pseudo- code:
;    sum = 0
;    term = 1
;    for( index = 0; index < size; index++) {
;        sum = sum + term
;        term = term * 2
;    }
; Registers use:
;    r16 <-> sum
;    r17 <-> term
;    r18 <-> loop index
;    r19 <-> size, i.e., number of terms
;    r__ <-> _____ (fill in if other register used)
;    r__ <-> _____ (fill in if other register used)
;    r__ <-> _____ (fill in if other register used)
```

```
.cseg
```

```
.def sum=r16
.def term=r17
.def index=r18
.def terms=r19
size
```

;Continued on next page

```
;initializing the registers
clr sum          ; sum = 0
ldi term, 0x01   ; term = 1
```

```
;for( index = 0; index < size; index++) {
;    sum = sum + term
;    term = term * 2
;}
```

```
** Insert your code here: perform the for loop above **
```

```
ldi size, 0x08
```

```
clr index
```

```
for: cp index, size
     breq end do the store done
```

```
add sum, term
```

```
lsl term
```

```
inc index
```

```
rjmp for
```

0, 1, 2, 3, 4, 5, 6, 7  
1, 1, 1, 1, 1, 1, 1  
8

add term, term

do the store

```
;storing the result in dseg location called sum
ldi r27, HIGH(sum)
ldi r26, LOW(sum)
st X, sum
```

```
done: jmp done
;End of code segment
```

```
; Data Storage area
.dseg
.org 0x200
sum .byte 2
```

The End