# FB tags

### 10 regular expression matching

recursion is very short
 base case = both string are empty
 the first char is matching
 or the other is matching

```
class Solution {
    public boolean isMatch(String s, String p) {
        if (p.length() == 0) {
            return s.length() == 0;
        }
        boolean first = s.length() > 0 && (p.charAt(0) == s.ch
arAt(0) || p.charAt(0) == '.');
        if (p.length() >= 2 && p.charAt(1) == '*') {
            return isMatch(s, p.substring(2)) || (first && isM
atch(s.substring(1), p));
        } else {
            return first && isMatch(s.substring(1), p.substrin
g(1));
        }
    }
}
```

1. 2d dp problem, understand when to remove 2 and when to remove only 1.

```
class Solution {
   public boolean isMatch(String s, String p) {
     if (s == null || p == null) {
        return false;
   }
```

```
boolean[][] dp = new boolean[s.length() + 1][p.length
() + 1];
        char[] sc = s.toCharArray();
        char[] pc = p.toCharArray();
        dp[0][0] = true;
        for (int i = 1; i <= p.length(); i++) {
            if (pc[i - 1] == '*') {
                dp[0][i] = dp[0][i - 2];
            }
        }
        for (int i = 1; i <= s.length(); i++) {
            for (int j = 1; j <= p.length(); j++) {
                if (pc[j - 1] == sc[i - 1] || pc[j - 1] ==
'.') {
                    dp[i][j] = dp[i - 1][j - 1];
                \} else if (pc[j-1] == '*') {
                    if (pc[j - 2] == sc[i - 1] \mid | pc[j - 2] ==
'.') {
                         dp[i][j] = dp[i - 1][j] || dp[i][j -
2];
                    } else {
                        dp[i][j] = dp[i][j - 2];
                    }
                }
            }
        }
        return dp[s.length()][p.length()];
    }
}
```

\*

\*\*\*\*

### 11 Whildcard matching

1. 2d dp, 2 cases when the p is having the \*, either up or right, cause you can either remove the \* or use the loop

```
class Solution {
    public boolean isMatch(String s, String p) {
        if (s == null || p == null) {
            return false;
        }
        boolean[][] dp = new boolean[s.length() + 1][p.length
() + 1];
        char[] sc = s.toCharArray();
        char[] pc = p.toCharArray();
        dp[0][0] = true;
        for (int i = 1; i <= p.length(); i++) {
            if (pc[i - 1] == '*') {
                dp[0][i] = dp[0][i - 1];
            }
        }
        for (int i = 1; i <= s.length(); i++) {
            for (int j = 1; j <= p.length(); j++) {
                if (sc[i - 1] == pc[j - 1] \mid\mid pc[j - 1] ==
'?') {
                    dp[i][j] = dp[i - 1][j - 1];
                } else {
                    if (pc[j - 1] == '*') {
                         dp[i][j] = dp[i - 1][j] || dp[i][j -
1];
                    }
```

```
}
}
return dp[s.length()][p.length()];
}
```

# 102. Binary Tree Level Order Traversal

- 1. bfs
- 2. dfs, cool and simple using the relationship of level and size of the size

```
class Solution {
    public List<List<Integer>> levelOrder(TreeNode root) {
        List<List<Integer>> res = new ArrayList<>();
        dfs(root, res, 0);
        return res;
    }
    void dfs(TreeNode root, List<List<Integer>> res, int l) {
        if (root == null) {
            return;
        }
        if (l >= res.size()) {
            res.add(new ArrayList<Integer>());
        }
        res.get(l).add(root.val);
        dfs(root.left, res, l + 1);
        dfs(root.right, res, l + 1);
    }
}
```

# 121. Best Time to Buy and Sell Stock

#### dp solution

```
class Solution {
   public int maxProfit(int[] prices) {
      int buy = Integer.MIN_VALUE;
      int sell = 0;
      for (int p : prices) {
           buy = Math.max(-p, buy);
           sell = Math.max(p + buy, sell);
      }
      return sell;
   }
}
```

# 122. Best Time to Buy and Sell Stock II

```
class Solution {
   public int maxProfit(int[] prices) {
      if (prices == null || prices.length == 0) {
          return 0;
      }
      int sum = 0;
      int prev = prices[0];
      for (int p : prices) {
          sum += Math.max(0, p - prev);
          prev = p;
      }
      return sum;
   }
}
```

# 122. Best Time to Buy and Sell Stock III

```
class Solution {
   public int maxProfit(int[] prices) {
      int buy1 = Integer.MIN_VALUE;
      int buy2 = Integer.MIN_VALUE;
      int sell1 = 0;
      int sell2 = 0;
      for (int p : prices) {
            buy1 = Math.max(-p, buy1);
            sell1 = Math.max(p + buy1, sell1);
            buy2 = Math.max(sell1 - p, buy2);
            sell2 = Math.max(p + buy2, sell2);
      }
      return sell2;
   }
}
```

2. left and right side max, final total max

# 188. Best Time to Buy and Sell Stock IV

1. dp similar to previous ones, only build new int array to fill buy and sell

```
class Solution {
  public int maxProfit(int k, int[] prices) {
    if (prices == null || prices.length == 0) {
      return 0;
    }
  int n = prices.length;
  int[] buy = new int[k + 1];
  int[] sell = new int[k + 1];
  if (k > n / 2) {
      return maxall(prices);
  }
```

```
Arrays.fill(buy, Integer.MIN_VALUE);
        for (int i = 0; i < n; i++) {
            for (int j = 1; j \le k; j++) {
                buy[j] = Math.max(sell[j - 1] - prices[i], buy
[j]);
                sell[j] = Math.max(buy[j] + prices[i], sell
[j]);
            }
        }
        return sell[k];
    }
    int maxall(int[] prices) {
        int sum = 0;
        for (int i = 1; i < prices.length; i++) {</pre>
            if (prices[i] - prices[i - 1] > 0) {
                sum += prices[i] - prices[i - 1];
            }
        }
        return sum;
    }
}
```

# 309. Best Time to Buy and Sell Stock with Cooldown

dp solution only with slight changes

```
class Solution {
   public int maxProfit(int[] prices) {
     if (prices == null || prices.length <= 1) {
       return 0;
   }
}</pre>
```

```
}
        int n = prices.length;
        int[] buy = new int[n];
        int[] sell = new int[n];
        Arrays.fill(buy, Integer.MIN_VALUE);
        buy[0] = -prices[0];
        sell[0] = 0;
        buy[1] = Math.max(-prices[0], -prices[1]);
        sell[1] = Math.max(0, prices[1] - prices[0]);
        for (int i = 2; i < n; i++) {
            buy[i] = Math.max(sell[i - 2] - prices[i], buy[i -
1]);
            sell[i] = Math.max(buy[i - 1] + prices[i], sell[i
- 1]);
        }
        return sell[n - 1];
    }
}
```

### 133. Clone Graph

#### 1. dfs

```
public class Solution {
    public UndirectedGraphNode cloneGraph(UndirectedGraphNode
    node) {
        Map<UndirectedGraphNode, UndirectedGraphNode> map = ne
    w HashMap<>();
        return dfs(node, map);
    }
    UndirectedGraphNode dfs(UndirectedGraphNode node, Map<UndirectedGraphNode, UndirectedGraphNode> map) {
        if (node == null) {
```

```
return null;
}
if (map.containsKey(node)) {
    return map.get(node);
}
UndirectedGraphNode temp = new UndirectedGraphNode(nod e.label);
    map.put(node, temp);
    for (UndirectedGraphNode n : node.neighbors) {
        temp.neighbors.add(dfs(n, map));
    }
    return temp;
}
```

1. bfs

# 138. Copy List with Random Pointer

1. Use a hashmap. First round, make a copy of list aside, set up the next pointer. Then do the random copy.

```
public class Solution {
    public RandomListNode copyRandomList(RandomListNode head)
{
        RandomListNode dummy = new RandomListNode(0);
        RandomListNode trav = dummy;
        Map<RandomListNode, RandomListNode> map = new HashMap<
>();
        while (head != null) {
            RandomListNode temp = new RandomListNode(head.labe);
            map.put(head, temp);
            temp.random = head.random; // copy the random here trav.next = temp;
```

```
trav = trav.next;
head = head.next;
}
trav = dummy.next;
while (trav != null) {
    if (trav.random != null) {
        trav.random = map.get(trav.random);
    }
    trav = trav.next;
}
return dummy.next;
}
```

#### 139. Word Break

1. return true or false, use dp

```
}
return dp[s.length()];
}
```

#### 140. Word Break II

#### 1. dfs + memo

```
class Solution {
    Map<String, List<String>> map = new HashMap<>();
    public List<String> wordBreak(String s, List<String> wordD
ict) {
        List<String> res = new ArrayList<>();
        if (s == null || s.length() == 0) {
            return res;
        }
        if (map.containsKey(s)) {
            return map.get(s);
        }
        if (wordDict.contains(s)) res.add(s);
        for (int i = 1; i < s.length(); i++) {
            String temp = s.substring(i);
            if (wordDict.contains(temp)) {
                List<String> next = wordBreak(s.substring(0,
i), wordDict);
                if (next.size() > 0) {
                    for (String n : next) {
                        res.add(n + " " + temp);
                    }
                }
            }
```

```
map.put(s, res);
return res;
}
```

### 151. Reverse Words in a String

此题不难,注意clarification里的细节,要去除头尾空格,单词中间多个空格reverse后只能有一个。

### 186. Reverse Words in a String II

### 152. Maximum Product Subarray

max positive, min negative, 分情况讨论

```
class Solution {
    public int maxProduct(int[] nums) {
        if (nums == null || nums.length == 0) {
            return 0;
        }
        int maxPos = nums[0];
        int maxNeg = nums[0];
        int res = nums[0];
        for (int i = 1; i < nums.length; i++) {</pre>
            if (nums[i] >= 0) {
                maxPos = Math.max(maxPos * nums[i], nums[i]);
                maxNeg = Math.min(maxNeg * nums[i], nums[i]);
            } else {
                int temp = maxNeg;
                maxNeg = Math.min(maxPos * nums[i], nums[i]);
                maxPos = Math.max(temp * nums[i], nums[i]);
```

```
}
res = Math.max(res, maxPos);
}
return res;
}
```

#### 161. One Edit Distance

think about the 3 cases. Either use equals substring or compare one by one.

```
public boolean isOneEditDistance(String s, String t) {
    int len = Math.min(s.length(), t.length());
    for (int i = 0; i < len; i++) {
        if (s.charAt(i) != t.charAt(i)) {
            if (s.length() == t.length()) return s.substring
(i + 1).equals(t.substring(i + 1)); // replace
            else if (s.length() < t.length()) return s.subst</pre>
ring(i).equals(t.substring(i + 1)); // delete t
                    return s.substring(i + 1).equals(t.substri
            else
ng(i)); // delete s
        }
    }
    return Math.abs(s.length() - t.length()) == 1; // corner c
ase: ""
}
```

Not using equals or substring, use recursion to pass one second boolean

```
class Solution {
  public boolean isOneEditDistance(String s, String t) {
    if(s == null && t == null){
      return false;
    }
  if(s == null && t.length() != 1){
```

```
return false;
        }
        if(t == null && s.length() != 1){
            return false;
        }
        if(Math.abs(s.length() - t.length()) > 1){
            return false;
        }
        return helper(s, 0, t, 0, false);
    }
    boolean helper(String s, int sstart, String t, int tstart,
boolean second){
        while(sstart < s.length() && tstart < t.length()){</pre>
            if(s.charAt(sstart) == t.charAt(tstart)){
                sstart++;
                tstart++;
            }else{
                if(second){
                     return false;
                }
                if(s.length() > t.length()){
                     return helper(s, sstart + 1, t, tstart, tr
ue);
                }else if(t.length() > s.length()){
                     return helper(s, sstart, t, tstart + 1, tr
ue);
                }else{
                     return helper(s, sstart + 1, t, tstart +
1, true);
                }
            }
```

```
if(!second && Math.abs(s.length() - t.length()) == 1){
    return true;
}
if(second && sstart == s.length() && tstart == t.lengt
h()){
    return true;
}
return false;
}
```

#### 17. Letter Combinations of a Phone Number

classic dfs

#### 173. Binary Search Tree Iterator

push all function, use next to have all the logic

Tree preorder, inorder, postorder traversal.

#### preorder

```
public List<Integer> preorderTraversal(TreeNode root) {
   List<Integer> res = new ArrayList<>();
   if (root == null) return res; // corner check
   Stack<TreeNode> stack = new Stack<>();
   stack.push(root);
   while (!stack.empty()) {
      res.add(stack.pop().val);
      if (root.right != null) stack.push(root.right);
      if (root.left != null) stack.push(root.left);
```

```
}
return res;
}
```

#### inorder

```
public List<Integer> inorderTraversal(TreeNode root) {
   List<Integer> res = new ArrayList<>();
   Stack<TreeNode> stack = new Stack<>();
   while (root != null || !stack.empty()) {
      while (root != null) {
         stack.push(root);
         root = root.left;
      }
      res.add(stack.pop().val);
      root = root.right;
   }
   return res;
}
```

#### postorder

```
public List<Integer> postorderTraversal(TreeNode root) {
   List<Integer> res = new ArrayList<>();
   Stack<TreeNode> stack = new Stack<>();
   TreeNode prev = null;
   while (root != null || !stack.empty()) {
      if (root != null) {
        stack.push(root);
        root = root.left;
      } else {
        TreeNode tmp = stack.peek();
        if (tmp.right != null && tmp.right != prev)
```

```
root = tmp.right;
else {
          stack.pop();
          res.add(tmp.val);
          prev = tmp;
        }
}
return res;
}
```

### 200 Number of Islands

- 1. DFS
- 2. BFS
- 3. Union Find

follow up perimeter of islands. Make sure to understand what is the perimeter

#### 206 Reverse LinkedList

- 2. iterative
- 3. recursive

# 211. Add and Search Word - Data structure design.java

1. trie, for the . , use dfs with the len to count the level

# 215. Kth Largest Element in an Array

### 221. Maximal Square

1. dp 滚动数组 int[] prev, int[] cur, 最后 prev = cur

#### 230. Kth Smallest Element in a BST

- 1. keep a stack, count the number of prev you have, return the correct prev
- 2. dfs can also work

### 235. Lowest Common Ancestor of Binary Search Tree

# 236. Lowest Common Ancestor of Binary Tree

similar

### 253. Meeting Rooms II

1. find the relationship between start and end

```
class Solution {
    public int minMeetingRooms(Interval[] intervals) {
        int idx = 0;
        int[] start = new int[intervals.length];
        int[] end = new int[intervals.length];
        for (Interval i : intervals) {
            start[idx] = i.start;
            end[idx] = i.end;
            idx++;
        }
        Arrays.sort(start);
        Arrays.sort(end);
        int start_idx = 0;
        int end idx = 0;
        int res = 0;
        while (start_idx < intervals.length) {</pre>
            if (start[start_idx] < end[end_idx]) {</pre>
                 res++;
                 start_idx++;
```

2. priority queue to get the earliest meeting end time.

# 257. Binary Tree Paths

# 261. Graph Valid Tree

1. union find

```
class Solution {
   public boolean validTree(int n, int[][] edges) {
      int[] father = new int[n];
      for (int i = 0; i < n; i++) {
          father[i] = i;
      }
      for (int[] e : edges) {
          int f1 = find(father, e[0]);
          int f2 = find(father, e[1]);
          if (f1 == f2) {
               return false;
          }
          father[f1] = father[f2];
    }
    return edges.length == n - 1;</pre>
```

```
}
int find(int[] nums, int i) {
    while (i != nums[i]) {
        nums[i] = nums[nums[i]];
        i = nums[i];
    }
    return i;
}
```

2. dfs or bfs

# 273. Integer to English Words

1. recursion with hardcoded

### 277. Find the Celebrity

1. O(n) time 2 passes

#### 278. First Bad Version

# 28. Implement strStr()

### 282. Expression Add Operators

1. DFS, remember to carry a number to reduce from when doing multiplying

#### 283 Move Zeros

#### 285 Inorder Sucessor in BST

### 29 Divide Two Integers

use bits manipulation, watch out for overflow

```
class Solution {
    public int divide(int dividend, int divisor) {
        if (divisor == 0 || (dividend == Integer.MIN_VALUE &&
divisor == -1)) {
            return Integer.MAX_VALUE;
        }
        int sign = 1;
        if ((dividend < 0 && divisor > 0) || (dividend > 0 &&
divisor < 0)) {
            sign = -1;
        }
        long ddd = Math.abs((long)dividend);
        long dvs = Math.abs((long)divisor);
        int res = 0;
        while (dvs <= ddd) {</pre>
            long temp = dvs;
            long mul = 1;
            while (ddd >= temp << 1) {</pre>
                temp <<= 1;
                mul <<= 1;
            }
            ddd -= temp;
            res += mul;
        }
        return sign * res;
    }
}
```

#### 295. Find Median from Data Stream

# 297. Serialize and Deserialize Binary Tree

use preorder and when deserialize, use a queue comparing to 449. Serialize and Deserialize BST, all elements on the left can be smaller than the root

# 300. Longest Increasing Subsequence

A binary search trick, use binary search to update the length of the array, time complexity is O(n logn)