

Space and Congruence Compression of Proofs

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European Master in Computational Logic

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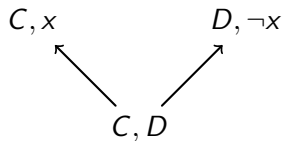
Space and Congruence Compression of Proofs

Space and Congruence Compression of **Proofs**

Resolution Rule

$$\frac{C \vee x \quad D \vee \neg x}{C \vee D}$$

Resolution Rule



Example

$$(x_1 \vee x_2 \vee x_3) \wedge (x_1 \vee \neg x_2) \wedge (x_1 \vee \neg x_3) \wedge (\neg x_1)$$

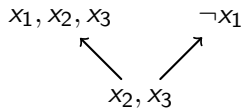
Example

$$(x_1 \vee x_2 \vee x_3) \wedge (x_1 \vee \neg x_2) \wedge (x_1 \vee \neg x_3) \wedge (\neg x_1)$$

$$\neg x_1$$

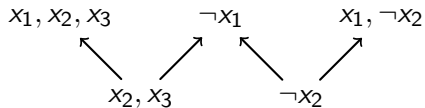
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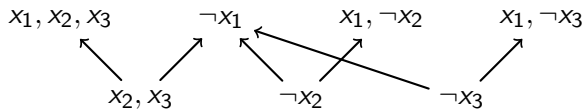
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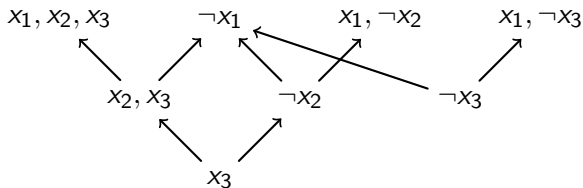
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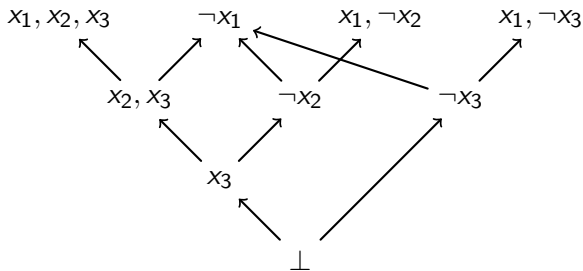
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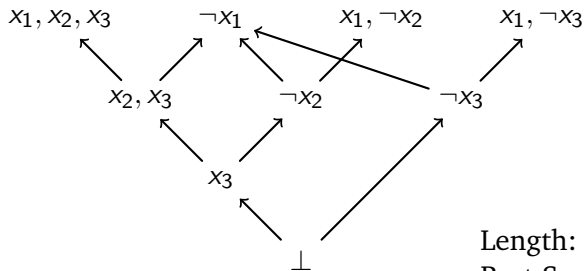
Example

$$(x_1 \vee x_2 \vee x_3) \wedge (x_1 \vee \neg x_2) \wedge (x_1 \vee \neg x_3) \wedge (\neg x_1)$$



Example

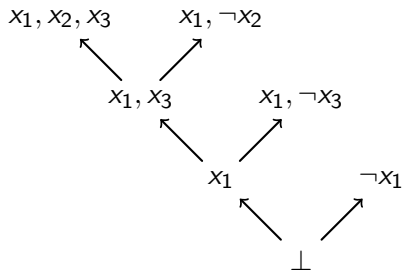
$$(x_1 \vee x_2 \vee x_3) \wedge (x_1 \vee \neg x_2) \wedge (x_1 \vee \neg x_3) \wedge (\neg x_1)$$



Length: 9
Best Space: 4

Example

$$(x_1 \vee x_2 \vee x_3) \wedge (x_1 \vee \neg x_2) \wedge (x_1 \vee \neg x_3) \wedge (\neg x_1)$$



Length: 7
Best Space: 3

Space and Congruence **Compression** of Proofs

- Smaller unsat cores, interpolants
- Easier proof processing
- Smaller proofs libraries
- Easier trusted interaction of deductive systems
- Proof generalization
- Proof carrying code

Synthesizing Multiple Boolean Functions using Interpolation on a Single Proof

- Georg Hofferek, et al, 2013, TU Graz

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Method

- ① Formulate problem in SMT theory of uninterpreted functions
- ② Obtain proof of unsatisfiability from SMT solver
- ③ Transform proof (local first, colorable)
- ④ Extract a single n -interpolant from the proof
- ⑤ Extract multiple interpolants from the single interpolant

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- Worst case exponential runtime in proof size

Proof Compression using Skeptik

- Input proof: 1,870,407 nodes
- Output proof: 868,760 nodes (53,6% compression)

Space and **Congruence** Compression of Proofs

Example

Knowledge

① $f(a) = a$

② $a = b$

③ $b = f(b)$

④ $f(a) \neq f(b)$

Example

Knowledge

- ① $f(a) = a$
- ② $a = b$
- ③ $b = f(b)$
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Unsatisfiable!

Example

Knowledge

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- ④ $f(a) \neq f(b)$

Unsatisfiable!

Proof

Equality is transitive, therefore from $f(a) = a$, $a = b$ and $b = f(b)$ follows $f(a) = f(b)$, which contradicts $f(a) \neq f(b)$

Example

Knowledge

- ① $f(a) = a$
- ② $a = b$
- ③ $b = f(b)$
- ④ $f(a) \neq f(b)$

Unsatisfiable!

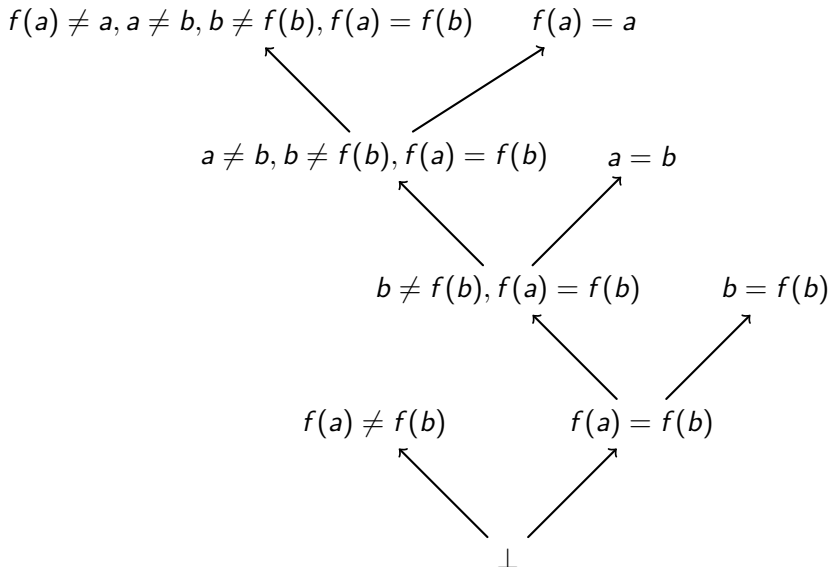
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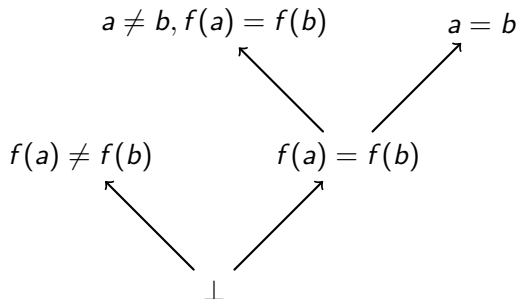
A different Proof

$f(\cdot)$ is a function, therefore from $a = b$ follows $f(a) = f(b)$, which contradicts $f(a) \neq f(b)$

A proof



A different proof



Congruence Closure

Congruence Closure

Ground Terms

- Constants a, b, c, \dots
- Compound Terms $f(t_1, \dots, t_n)$

Congruence Closure

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- Constants a, b, c, \dots
- Compound Terms $f(t_1, \dots, t_n)$

Congruence Relation R

- Reflexive: $\forall t (t, t) \in R$
- Symmetric: $(s, t) \in R \Rightarrow (t, s) \in R$
- Transitive: $(t_1, t_2) \in R \dots (t_{m-1}, t_m) \in R \Rightarrow (t_1, t_m) \in R$
- Compatible: $\forall i (t_i, s_i) \in R \Rightarrow (f(t_1, \dots, t_n), f(s_1, \dots, s_n)) \in R$

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Congruence Closure E^* of set of equations E

- Smallest Congruence Relation containing E
- Computable in $O(n \log(n))$
- E is explanation for $(s, t) \in E^*$

Example revisited

Knowledge

- ① $f(a) = a$
- ② $a = b$
- ③ $b = f(b)$
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Example revisited

Knowledge

- ① $f(a) = a$
- ② $a = b$
- ③ $b = f(b)$
- ④ $f(a) \neq f(b)$

Explanation for $f(a) = f(b)$

$\{ f(a) = a, a = b, b = f(b) \}$

Example revisited

Knowledge

- ① $f(a) = a$
- ② $a = b$
- ③ $b = f(b)$
- ④ $f(a) \neq f(b)$

Explanation for $f(a) = f(b)$

{ $a = b, b = f(b)$ }

Example revisited

Knowledge

- ① $f(a) = a$
- ② $a = b$
- ③ $b = f(b)$
- ④ $f(a) \neq f(b)$

Explanation for $f(a) = f(b)$

{ $a = b$ }

Example revisited

Knowledge

- ❶ $f(a) = a$
- ❷ $a = b$
- ❸ $b = f(b)$
- ❹ $f(a) \neq f(b)$

Explanation for $f(a) = f(b)$

{ $a = b$ }

Short explanation \rightsquigarrow short (sub)proof

Short Explanation Decision Problem

Given a set of input equations E , a target equation $s = t$ and $k \in \mathbb{N}$, does there exist an explanation $E' \subseteq E$ of $s = t$ with $|E'| \leq k$?

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NP-complete

Short Explanation Decision Problem

Given a set of input equations E , a target equation $s = t$ and $k \in \mathbb{N}$, does there exist an explanation $E' \subseteq E$ of $s = t$ with $|E'| \leq k$?

NP-complete

Reduction of SAT to the short explanation decision problem

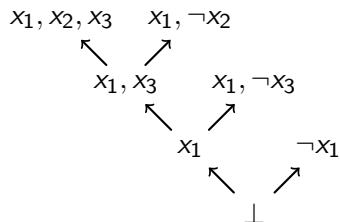
Space and Congruence Compression of Proofs

Is it possible to remove parts of the proof from memory during proof processing?

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Which parts?

Space Measure Example



Space Measure Example

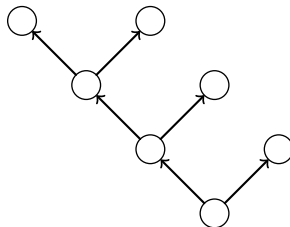
Maximum number of nodes in memory: 0



Not in memory

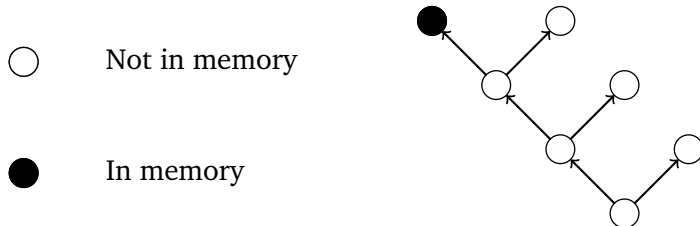


In memory



Space Measure Example

Maximum number of nodes in memory: 1



Space Measure Example

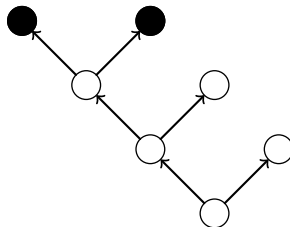
Maximum number of nodes in memory: 2



Not in memory



In memory



Space Measure Example

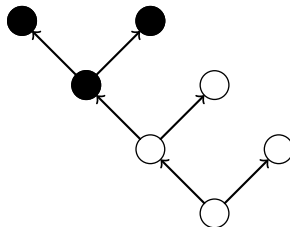
Maximum number of nodes in memory: 3



Not in memory



In memory



Space Measure Example

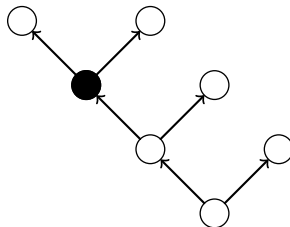
Maximum number of nodes in memory: 3



Not in memory

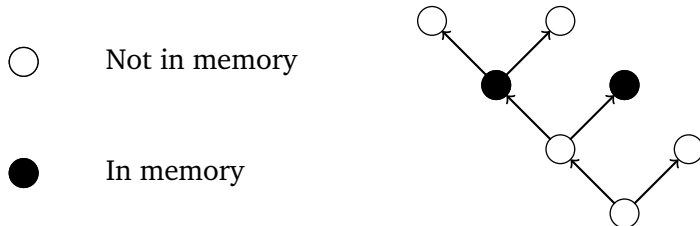


In memory



Space Measure Example

Maximum number of nodes in memory: 3



Space Measure Example

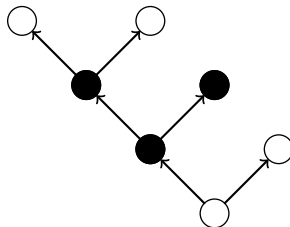
Maximum number of nodes in memory: 3



Not in memory



In memory



Space Measure Example

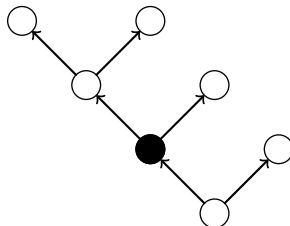
Maximum number of nodes in memory: 3



Not in memory



In memory



Space Measure Example

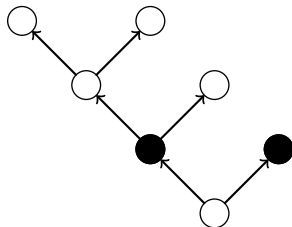
Maximum number of nodes in memory: 3



Not in memory



In memory



Space Measure Example

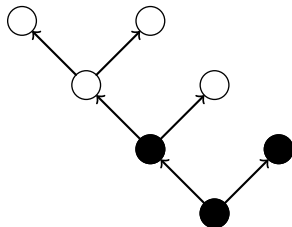
Maximum number of nodes in memory: 3



Not in memory



In memory



Space Measure Example

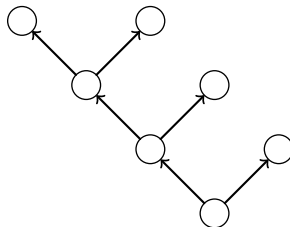
Maximum number of nodes in memory: 0



Not in memory



In memory



Space Measure Example

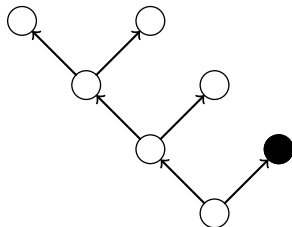
Maximum number of nodes in memory: 1



Not in memory



In memory



Space Measure Example

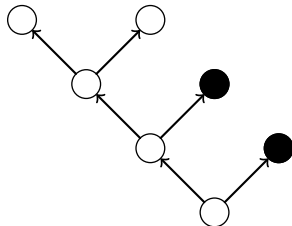
Maximum number of nodes in memory: 2



Not in memory



In memory



Space Measure Example

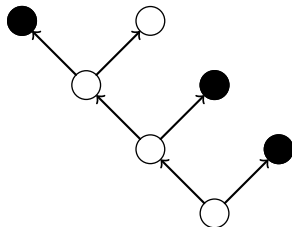
Maximum number of nodes in memory: 3



Not in memory



In memory



Space Measure Example

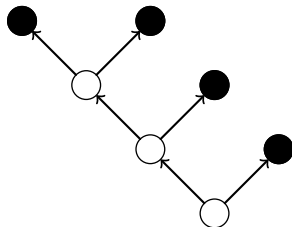
Maximum number of nodes in memory: 4



Not in memory



In memory



Space Measure Example

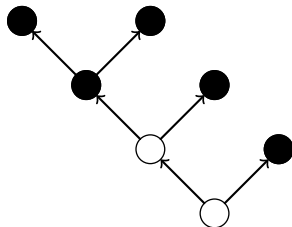
Maximum number of nodes in memory: 5



Not in memory

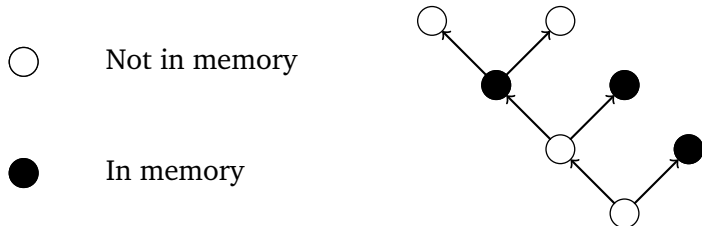


In memory



Space Measure Example

Maximum number of nodes in memory: 5



Space Measure Example

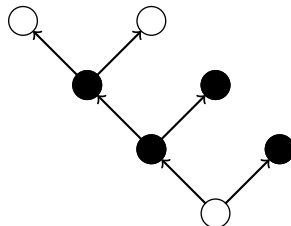
Maximum number of nodes in memory: 5



Not in memory

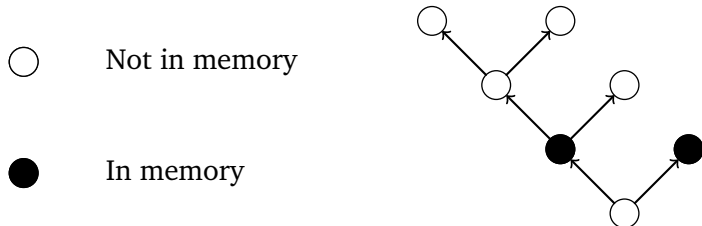


In memory



Space Measure Example

Maximum number of nodes in memory: 5



Space Measure Example

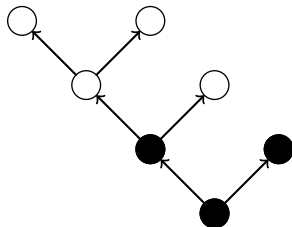
Maximum number of nodes in memory: 5



Not in memory



In memory



Space Measure Example

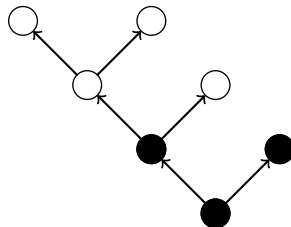
Maximum number of nodes in memory: 5



Not in memory



In memory



Good traversal orders are essential!

Space measure of a proof and a traversal order

Maximal amount of nodes that have to be kept in memory at once while processing the proof following the traversal order

Construct Traversal Orders

Construct Optimal Order

- NP-complete
- Optimal strategy in some pebbling game

Construct Good Order

- Greedy Algorithms
- Heuristic choices
- Top-Down
- Bottom-Up

Experiments, Unsung Heroes & Conclusion

Experimental Results

Congruence Compression

- 3965 proofs from problems in QF_UF logic
- 2% average effective compression in proof length
- 28% compression in explanation length

Space Compression

- 7555 SAT and SMT proofs
- Bottom-Up outperforms Top-Down
- Average space measure is 44 times smaller than proof length

- Explanation producing congruence closure algorithm
 - Using immutable data structures
 - Modified version of Dijkstra's shortest path algorithm
- Proof producing algorithm
- Resolution calculus extended with equality
- SAT translation of optimal traversal order
- Correctness & soundness proofs
- Implementation of all presented methods

- Proofs can be compressed in length and space
- Finding the shortest explanation is NP-complete
- Proof production is tricky
- Construct traversal orders Bottom-Up

Thank you for your attention!

NP-completeness proof sketch

From a propositional logic formula Φ obtain ...

- a set of equations E_Φ
- a target equation $s_\Phi = t_\Phi$
- $k_\Phi \in \mathbb{N}$

such that ...

Φ is satisfiable if and only if there is an explanation $E' \subseteq E_\Phi$ of $s_\Phi = t_\Phi$ with $|E'| \leq k_\Phi$

NP-completeness proof sketch example

Formula

$$(x_1 \vee x_2 \vee \neg x_3) \wedge (\neg x_2 \vee x_3) \wedge (\neg x_1 \vee \neg x_2)$$

Translation to equations

NP-completeness proof sketch example

Formula

$$(x_1 \vee x_2 \vee \neg x_3) \wedge (\neg x_2 \vee x_3) \wedge (\neg x_1 \vee \neg x_2)$$

Translation to equations

$$\perp_1 \multimap \hat{x}_1 \multimap \top_1 \qquad \perp_2 \multimap \hat{x}_2 \multimap \top_2 \qquad \perp_3 \multimap \hat{x}_3 \multimap \top_3$$

NP-completeness proof sketch example

Formula

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Translation to equations

$$\perp_1 \text{ --- } \hat{x}_1 \text{ --- } \top_1 \qquad \perp_2 \text{ --- } \hat{x}_2 \text{ --- } \top_2 \qquad \perp_3 \text{ --- } \hat{x}_3 \text{ --- } \top_3$$

$$\begin{array}{c} t_1(\hat{x}_1) \\ \diagdown \\ \hat{c}_1 - t_1(\hat{x}_2) \\ \diagup \\ f_1(\hat{x}_3) \end{array}$$

NP-completeness proof sketch example

Formula

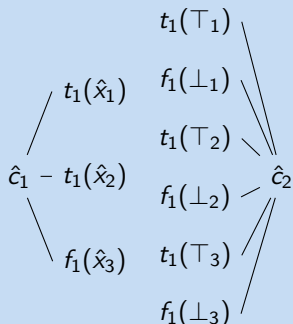
$$(x_1 \vee x_2 \vee \neg x_3) \wedge (\neg x_2 \vee x_3) \wedge (\neg x_1 \vee \neg x_2)$$

Translation to equations

$$\perp_1 \text{ --- } \hat{X}_1 \text{ --- } \top_1$$

$$\perp_2 \text{ --- } \hat{x}_2 \text{ --- } \top_2$$

$$\perp_3 \text{ --- } \hat{X}_3 \text{ --- } \top_3$$

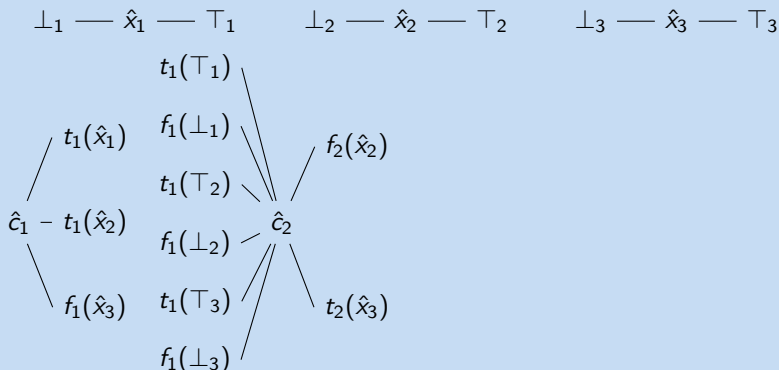


NP-completeness proof sketch example

Formula

$$(x_1 \vee x_2 \vee \neg x_3) \wedge (\neg x_2 \vee x_3) \wedge (\neg x_1 \vee \neg x_2)$$

Translation to equations

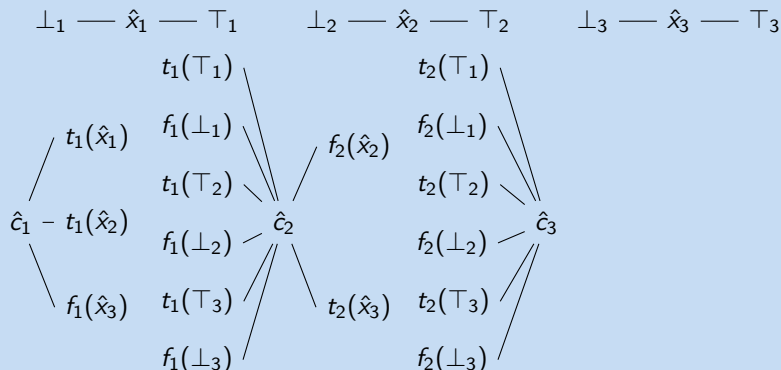


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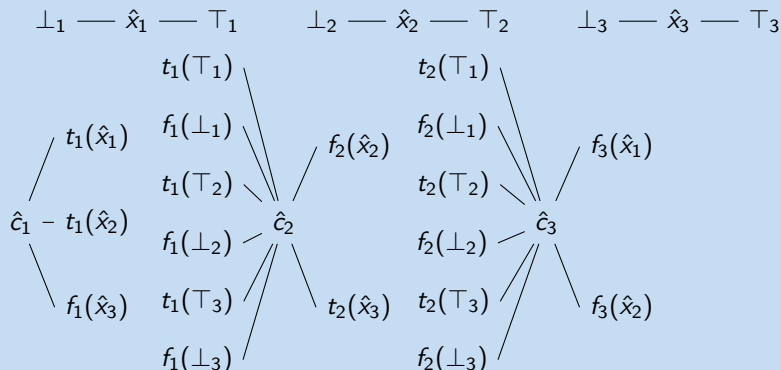


NP-completeness proof sketch example

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Translation to equations

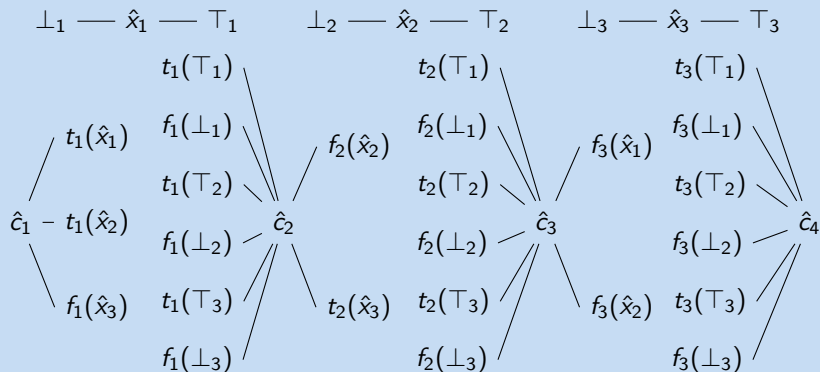


NP-completeness proof sketch example

Formula

$$(x_1 \vee x_2 \vee \neg x_3) \wedge (\neg x_2 \vee x_3) \wedge (\neg x_1 \vee \neg x_2)$$

Translation to equations



NP-completeness proof sketch example

Formula

$$(x_1 \vee x_2 \vee \neg x_3) \wedge (\neg x_2 \vee x_3) \wedge (\neg x_1 \vee \neg x_2)$$

Small subset corresponding to satisfying assignment

$$\hat{x}_1 \text{ --- } \top_1 \qquad \perp_2 \text{ --- } \hat{x}_2 \qquad \hat{x}_3 \text{ --- } \top_3$$

$$\hat{c}_1 = t_1(\hat{x}_1) \quad t_1(\top_1) = \hat{c}_2 = f_2(\hat{x}_2) \quad f_2(\perp_2) = \hat{c}_3 = f_3(\hat{x}_2) \quad f_3(\perp_2) = \hat{c}_4$$

Short explanation, long proof

$$t_a \leftarrow f(f(a, b), f(a, a))$$

$$t_b \leftarrow f(f(b, a), f(b, b))$$

$$f(a, b) \neq f(b, a), f(a, a) \neq f(b, b), t_a = t_b$$

$$a \neq b, f(a, b) = f(b, a)$$

$$a \neq b, f(a, a) \neq f(b, b), t_a = t_b$$

$$a \neq b, f(a, a) = f(b, b)$$

$$a \neq b, t_a = t_b$$

Compatibility proof

$$t_a \neq a, a \neq b, b \neq t_b, t_a = t_b$$

Transitivity proof

Congruence Experimental Results

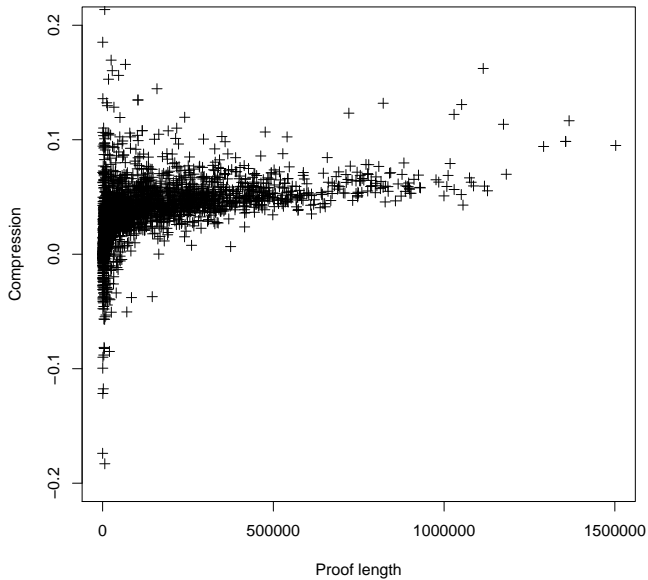
- 3965 proofs of problems of the SMT-LIB benchmark in the QF_UF logic

Method	Compression	Min	Max	Speed
EqGraph	5.350 %	-18.302 %	81.347 %	0.343
Proof Forest	5.196 %	-43.985 %	77.202 %	0.611
DAGify	3.368 %	0.0 %	14.433 %	1.655

Table: Compression Results

Congruence Graph	Compressed	Compression
Equation Graph	12.42 %	28.34 %
Proof Forest	11.459 %	28.69 %

Table: Explanation Size Results



Space Compression Results

- VeriT: Problems from SMT-lib
- TraceCheck: Problems from SATLIB, computed with PicoSAT

Name	Number of proofs	Maximum length	Average length
TraceCheck ₁	2239	90756	5423
TraceCheck ₂	215	1768249	268863
veriT ₁	4187	2241042	103162
veriT ₂	914	120075	5391

Table: Proof Benchmark Sets

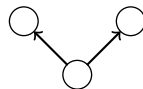
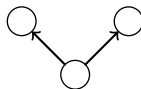
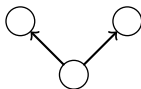
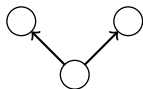
Space Compression Results

Algorithm Heuristic	Relative Performance (%)	Speed (nodes/ms)
Bottom-Up		
Children	17.52	88.6
LastChild	26.31	84.5
Distance(1)	9.46	21.2
Distance(3)	-0.40	0.5
Top-Down		
Children	-27.47	0.3
LastChild	-31.98	1.9
Distance(1)	-70.14	0.6
Distance(3)	-74.33	0.1

Table: Experimental Results

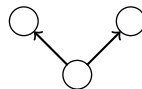
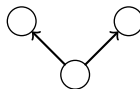
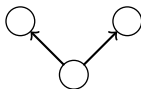
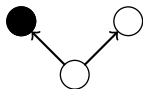
Construct Order Top-Down

Maximum number of nodes in memory: 0



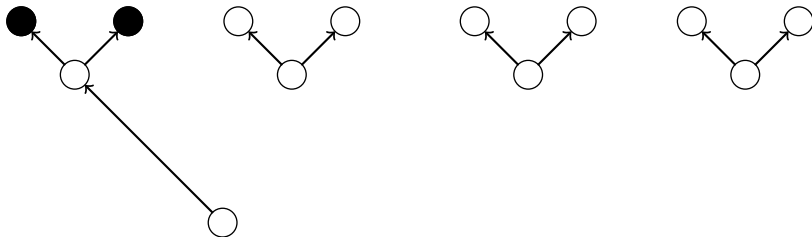
Construct Order Top-Down

Maximum number of nodes in memory: 1



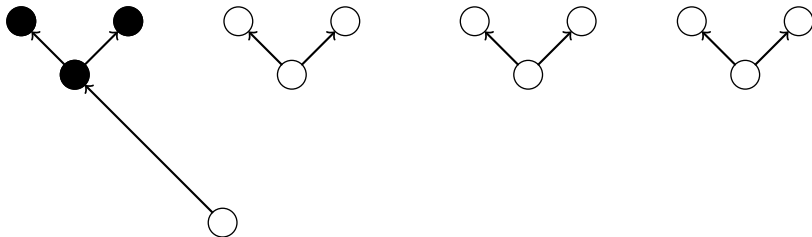
Construct Order Top-Down

Maximum number of nodes in memory: 2



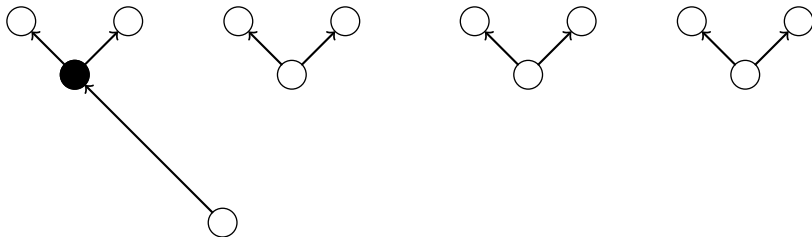
Construct Order Top-Down

Maximum number of nodes in memory: 3



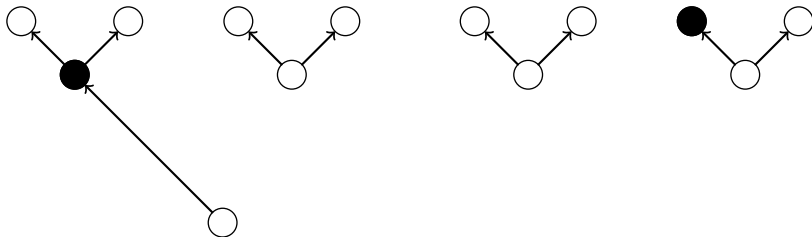
Construct Order Top-Down

Maximum number of nodes in memory: 3



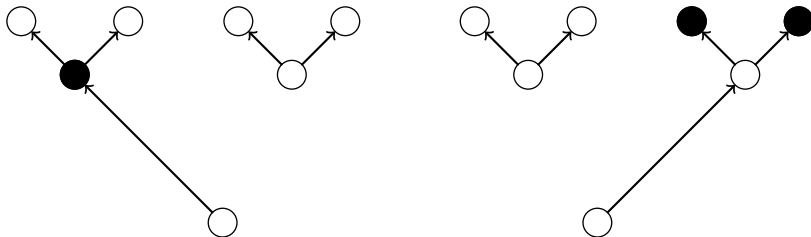
Construct Order Top-Down

Maximum number of nodes in memory: 3



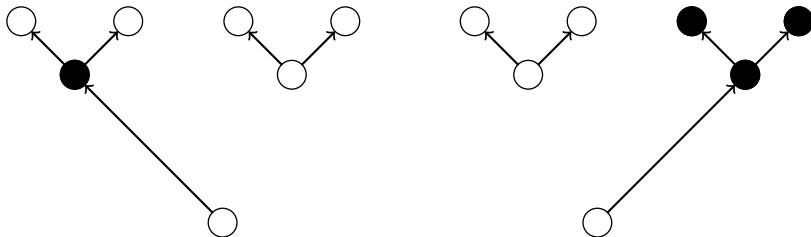
Construct Order Top-Down

Maximum number of nodes in memory: 3



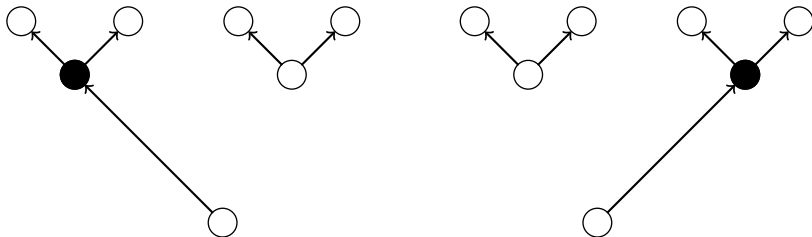
Construct Order Top-Down

Maximum number of nodes in memory: 4



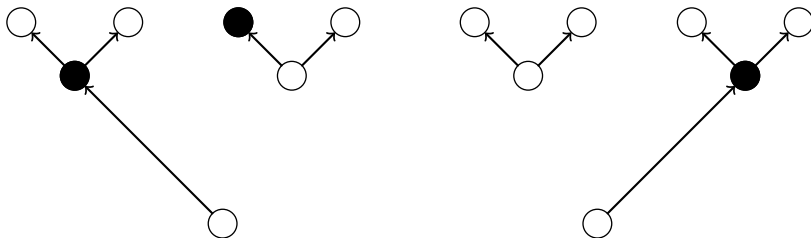
Construct Order Top-Down

Maximum number of nodes in memory: 4



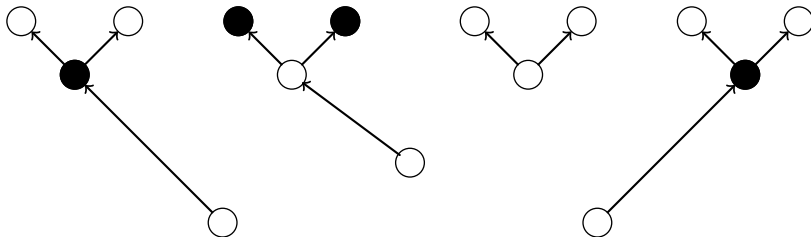
Construct Order Top-Down

Maximum number of nodes in memory: 4



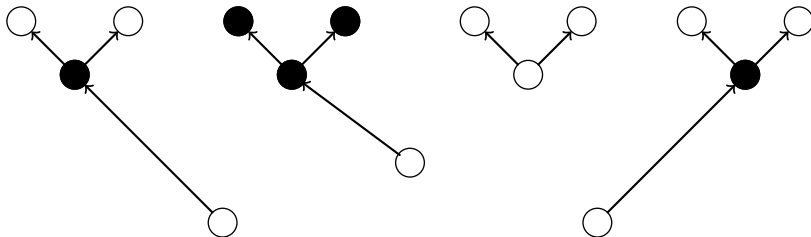
Construct Order Top-Down

Maximum number of nodes in memory: 4



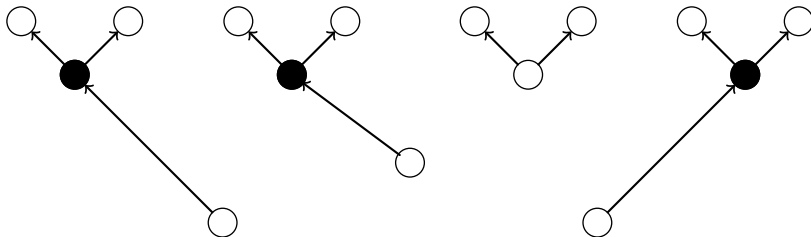
Construct Order Top-Down

Maximum number of nodes in memory: 5



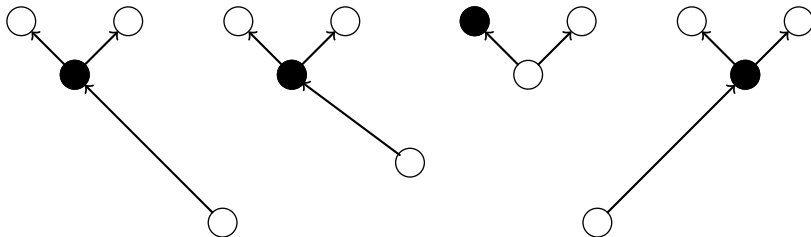
Construct Order Top-Down

Maximum number of nodes in memory: 5



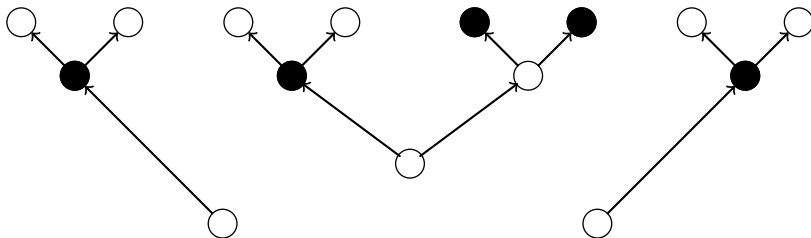
Construct Order Top-Down

Maximum number of nodes in memory: 5



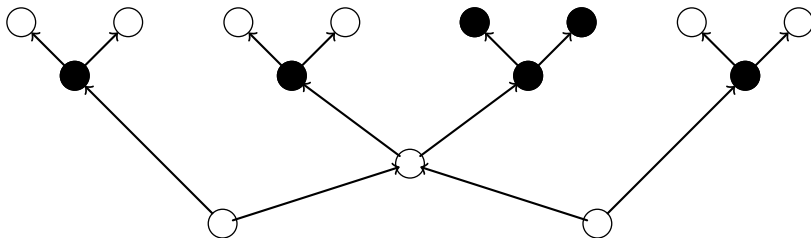
Construct Order Top-Down

Maximum number of nodes in memory: 5



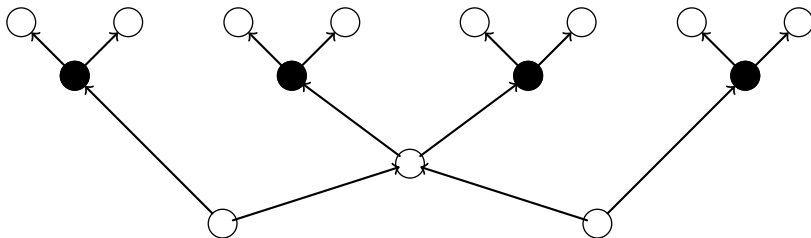
Construct Order Top-Down

Maximum number of nodes in memory: 6



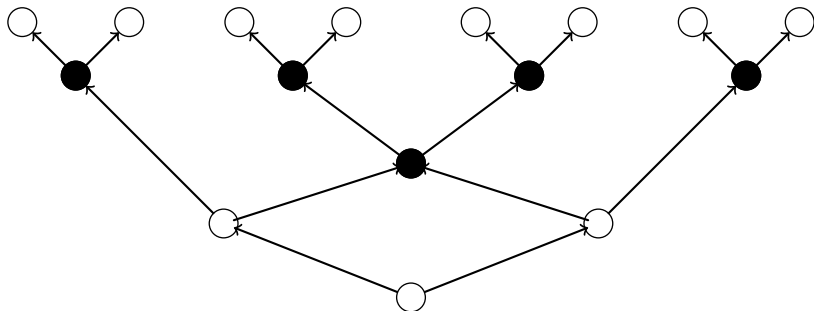
Construct Order Top-Down

Maximum number of nodes in memory: 6



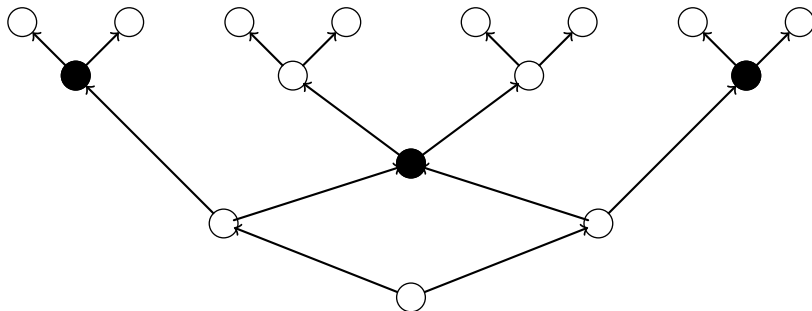
Construct Order Top-Down

Maximum number of nodes in memory: 6



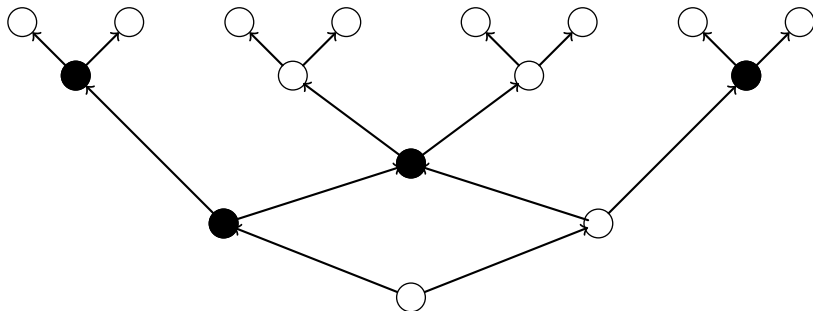
Construct Order Top-Down

Maximum number of nodes in memory: 6



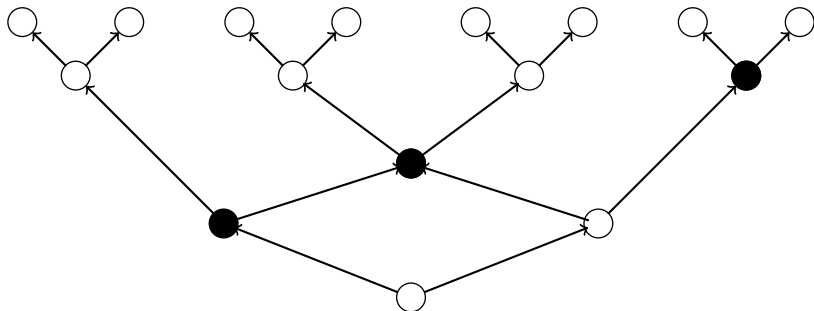
Construct Order Top-Down

Maximum number of nodes in memory: 6



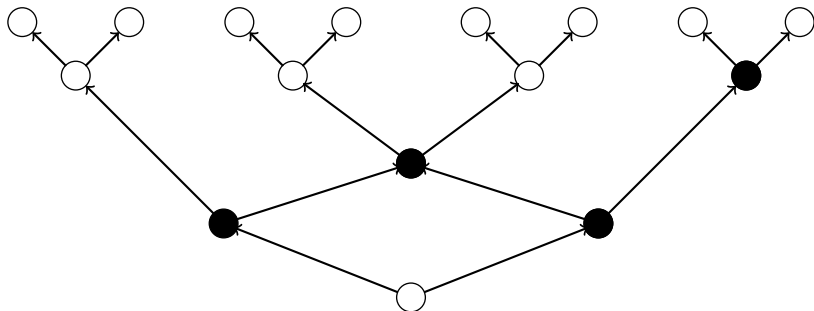
Construct Order Top-Down

Maximum number of nodes in memory: 6



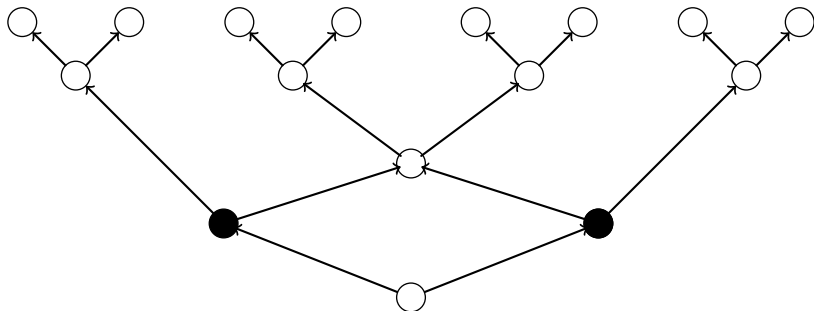
Construct Order Top-Down

Maximum number of nodes in memory: 6



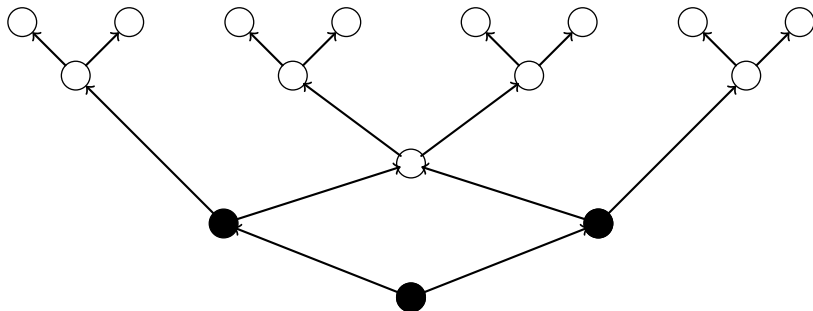
Construct Order Top-Down

Maximum number of nodes in memory: 6



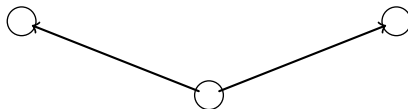
Construct Order Top-Down

Maximum number of nodes in memory: 6



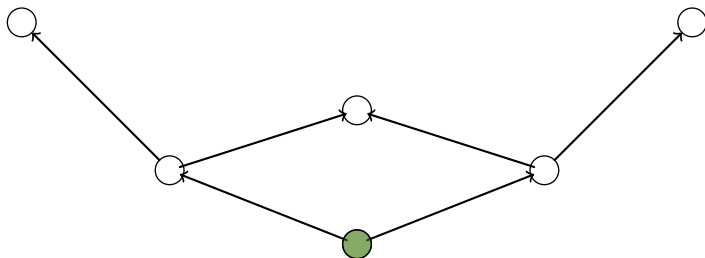
Construct Order Bottom-Up

Maximum number of nodes in memory: 0



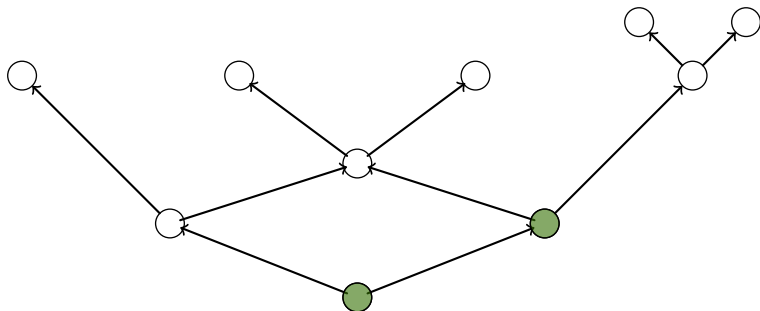
Construct Order Bottom-Up

Maximum number of nodes in memory: 0



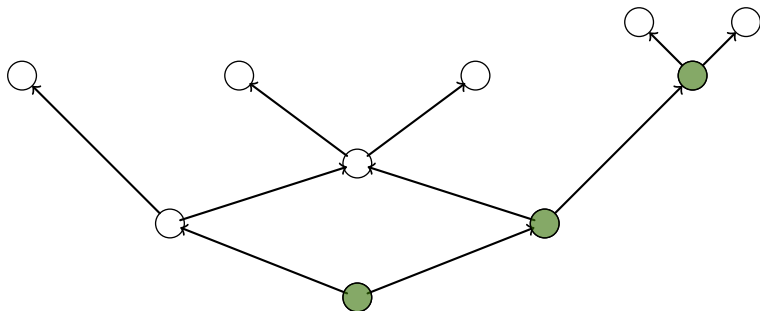
Construct Order Bottom-Up

Maximum number of nodes in memory: 0



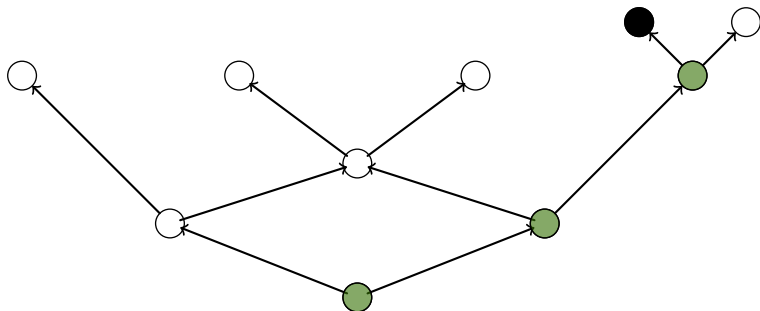
Construct Order Bottom-Up

Maximum number of nodes in memory: 0



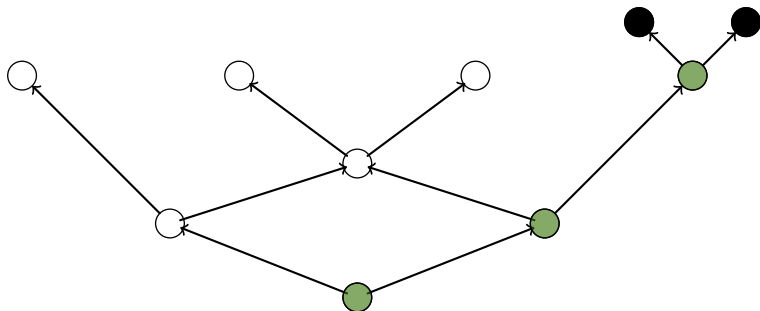
Construct Order Bottom-Up

Maximum number of nodes in memory: 1



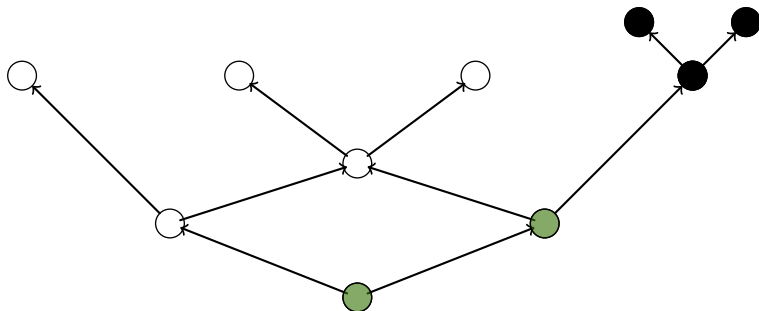
Construct Order Bottom-Up

Maximum number of nodes in memory: 2



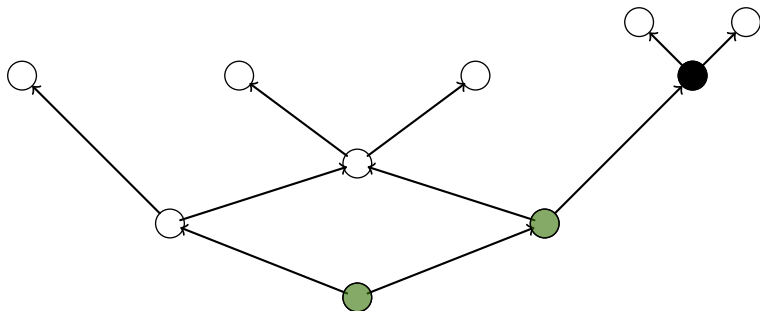
Construct Order Bottom-Up

Maximum number of nodes in memory: 3



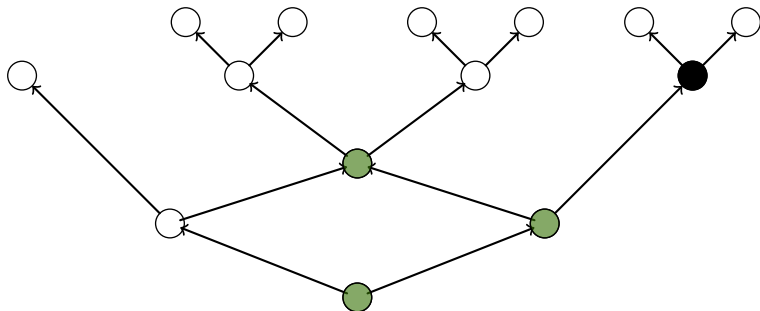
Construct Order Bottom-Up

Maximum number of nodes in memory: 3



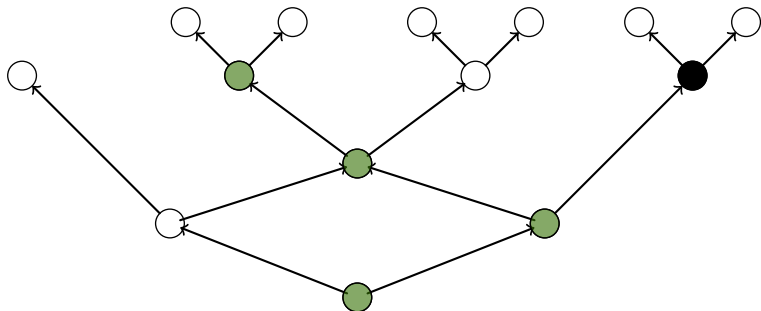
Construct Order Bottom-Up

Maximum number of nodes in memory: 3



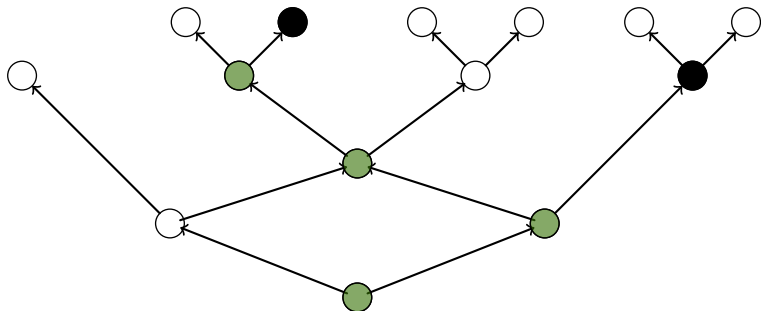
Construct Order Bottom-Up

Maximum number of nodes in memory: 3



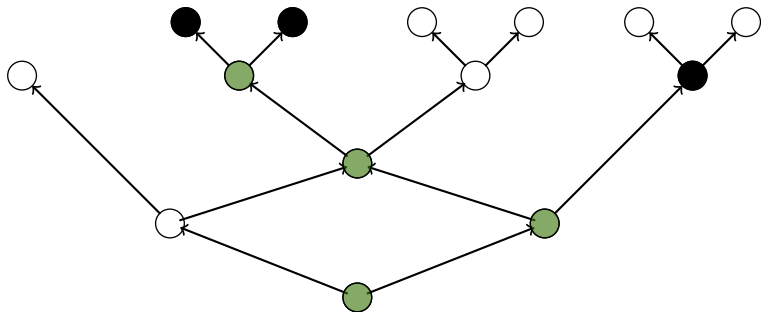
Construct Order Bottom-Up

Maximum number of nodes in memory: 3



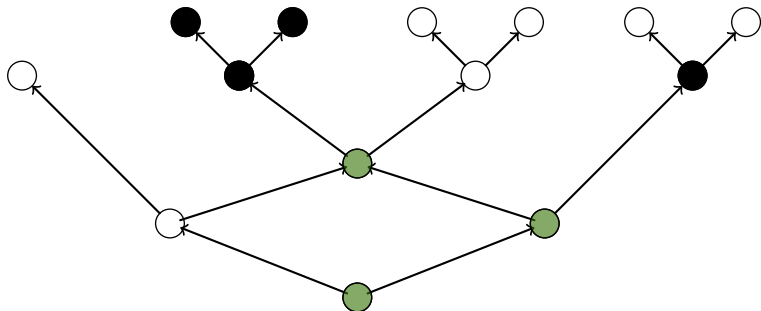
Construct Order Bottom-Up

Maximum number of nodes in memory: 3



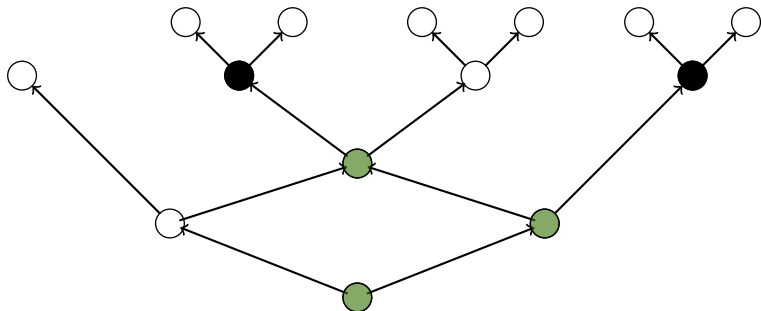
Construct Order Bottom-Up

Maximum number of nodes in memory: 4



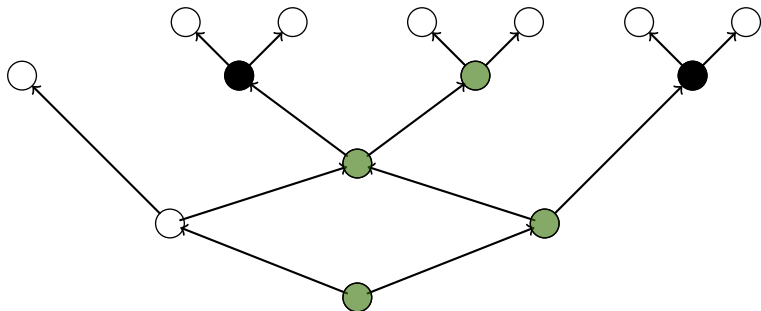
Construct Order Bottom-Up

Maximum number of nodes in memory: 4



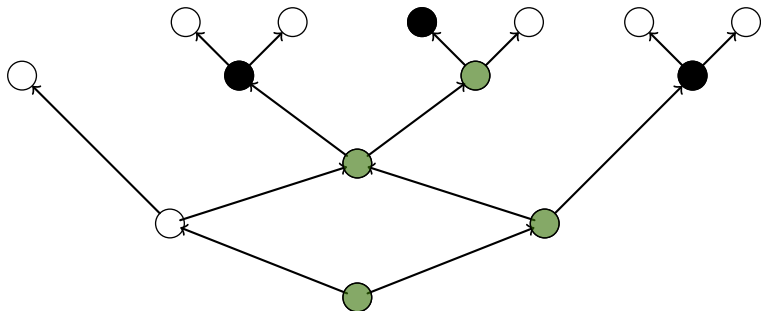
Construct Order Bottom-Up

Maximum number of nodes in memory: 4



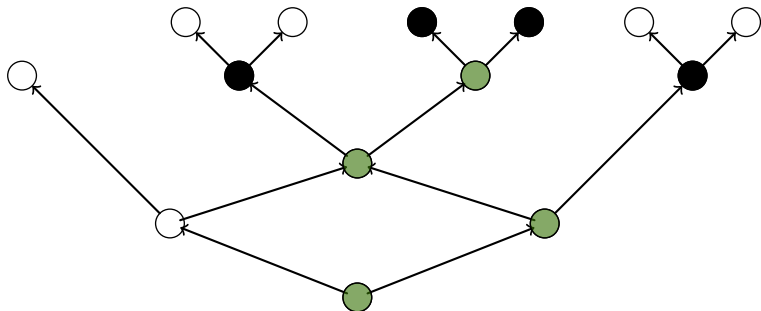
Construct Order Bottom-Up

Maximum number of nodes in memory: 4



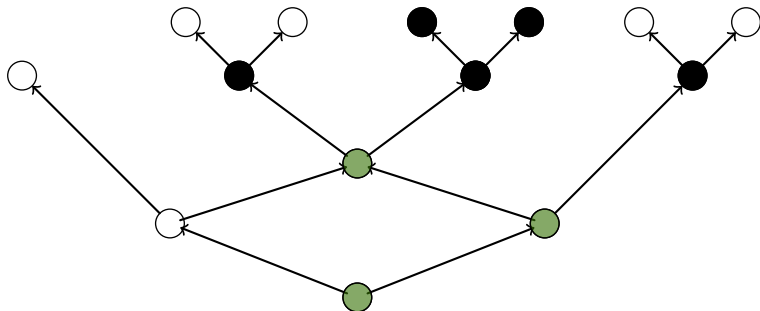
Construct Order Bottom-Up

Maximum number of nodes in memory: 4



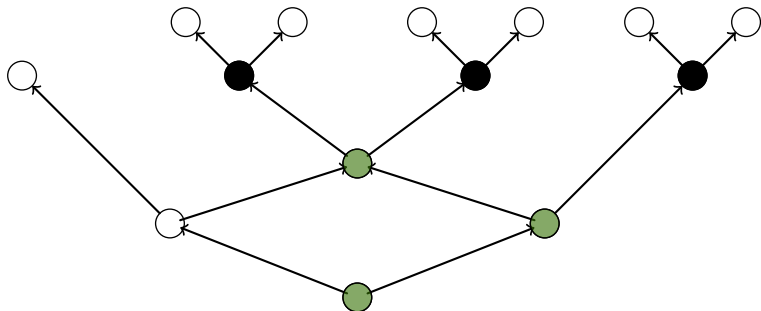
Construct Order Bottom-Up

Maximum number of nodes in memory: 5



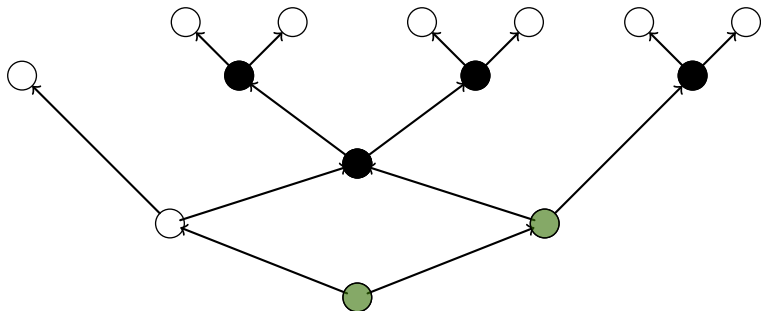
Construct Order Bottom-Up

Maximum number of nodes in memory: 5



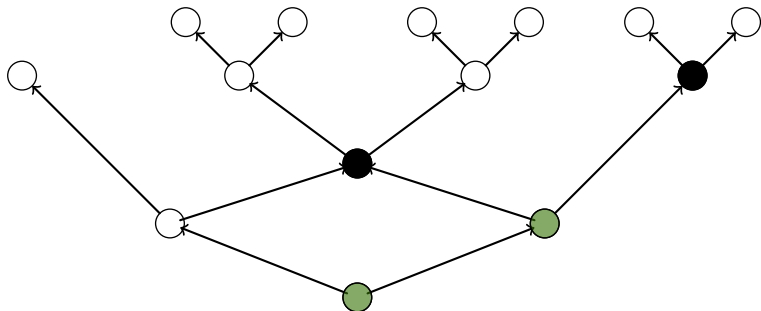
Construct Order Bottom-Up

Maximum number of nodes in memory: 5



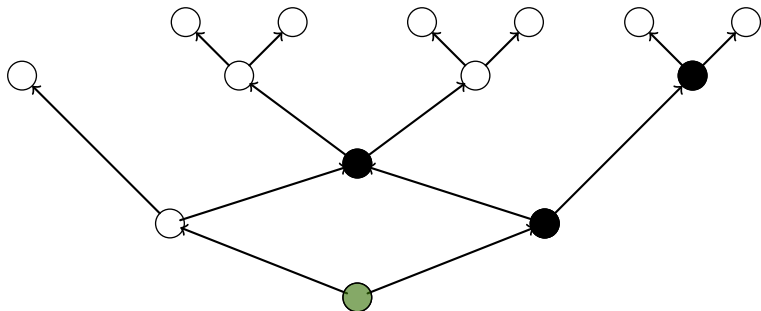
Construct Order Bottom-Up

Maximum number of nodes in memory: 5



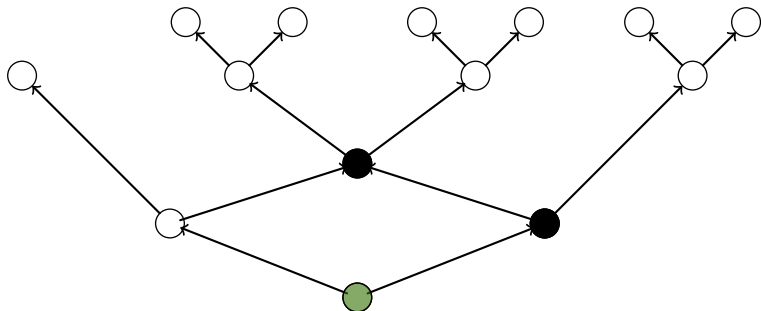
Construct Order Bottom-Up

Maximum number of nodes in memory: 5



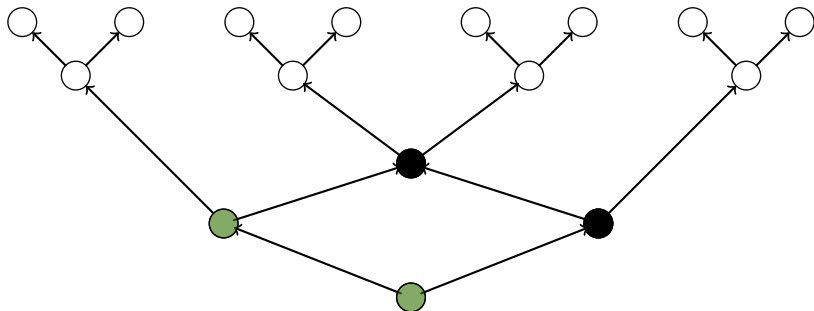
Construct Order Bottom-Up

Maximum number of nodes in memory: 5



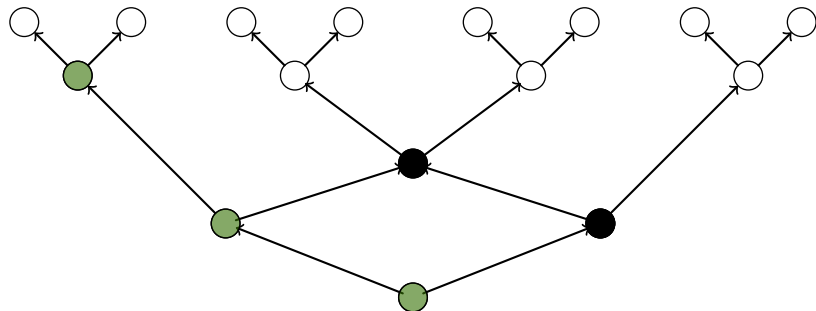
Construct Order Bottom-Up

Maximum number of nodes in memory: 5



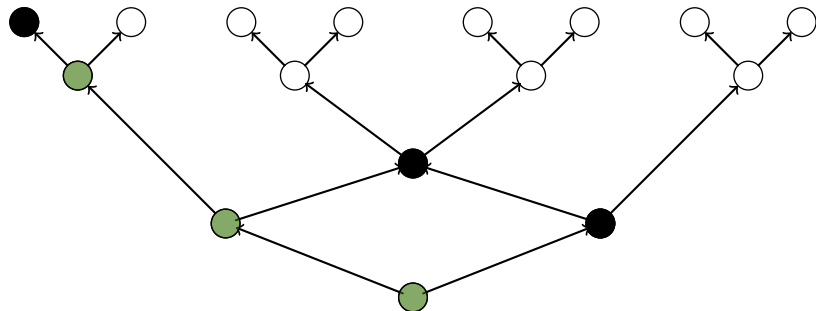
Construct Order Bottom-Up

Maximum number of nodes in memory: 5



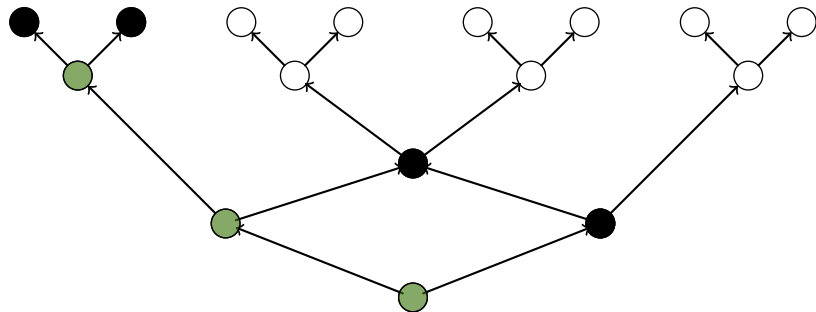
Construct Order Bottom-Up

Maximum number of nodes in memory: 5



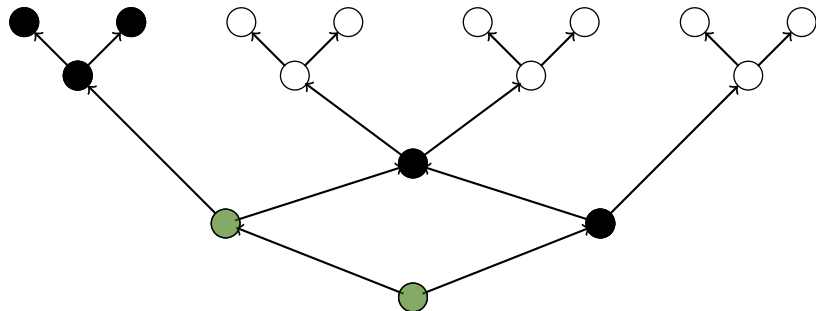
Construct Order Bottom-Up

Maximum number of nodes in memory: 5



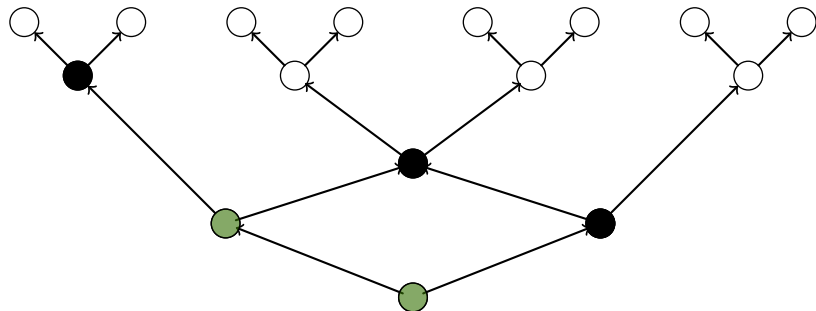
Construct Order Bottom-Up

Maximum number of nodes in memory: 5



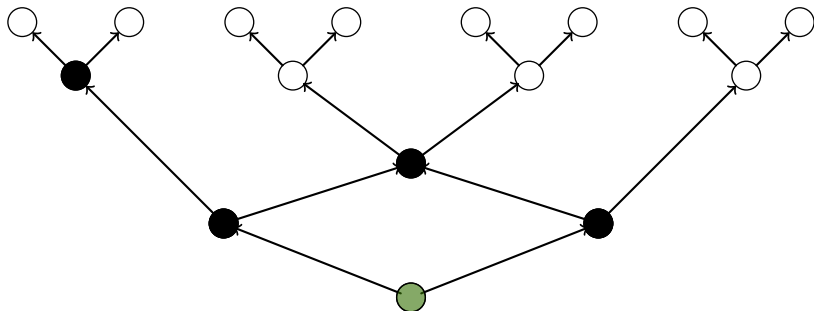
Construct Order Bottom-Up

Maximum number of nodes in memory: 5



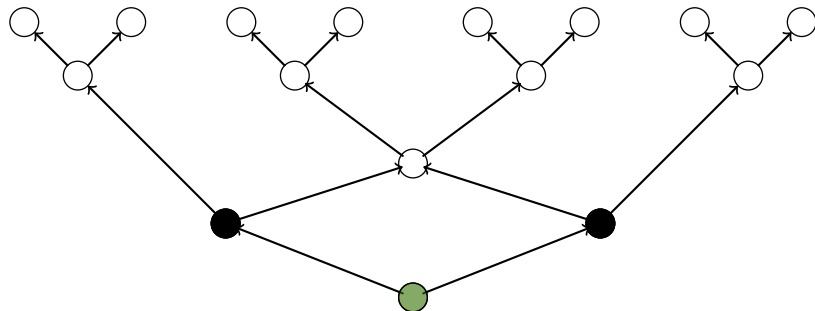
Construct Order Bottom-Up

Maximum number of nodes in memory: 5



Construct Order Bottom-Up

Maximum number of nodes in memory: 5



Construct Order Bottom-Up

Maximum number of nodes in memory: 5

