# **SpinalHDL Documentation**

**SpinalHDL contributors** 

# **CONTENTS**

1	What is the overhead of SpinalHDL generated RTL compared to human written VHDL/Verilog?  What if SpinalHDL becomes unsupported in the future?  Does SpinalHDL keep comments in generated VHDL/verilog?  Could SpinalHDL scale up to big projects?  How SpinalHDL came to be  Why develop a new language when there is VHDL/Verilog/SystemVerilog?	1 1 1 1 2 2 2
2	2.1 Communication channels	<b>3</b> 3
3	3.1 Companies	<b>5</b> 5 5
4	4.1 Getting Started          4.2 Motivation          4.3 Presentation          4.4 Scala Guide          4.5 Help for VHDL people	7 10 10 10 20 30
5	5.1       Bool       3         5.2       Bits       3         5.3       UInt/SInt       4         5.4       SpinalEnum       4         5.5       Bundle       5         5.6       Vec       5         5.7       UFix/SFix       5         5.8       Floating       6	33 33 36 40 48 51 54 56 60
6	6.1       Component and hierarchy       6         6.2       Area       6         6.3       Function       6         6.4       Clock domains       7         6.5       Instantiate VHDL and Verilog IP       7	65 68 69 70 78
7	Semantic 9	93

	7.1 7.2 7.3	Assignments
8	<b>Seque</b> 8.1	ential logic Registers
	8.2	RAM/ROM
9	Desig	n errors 111
	9.1	Assignment overlap
	9.2	Clock crossing violation
	9.3	Combinatorial loop
	9.4	Hierarchy violation
	9.5	Io bundle
	9.6	Latch detected
	9.7	No driver on
	9.8 9.9	NullPointerException    118      Register defined as component input    119
	9.9	Scope violation
	9.10	Spinal can't clone class
	9.12	Unassigned register
	9.13	Unreachable is statement
		Width mismatch
		Introduction
10		language features 127
		Utils
		Assertions
		Report
		Formal
		Analog and inout          VHDL and Verilog generation
	10.7	Introduction
	10.7	introduction
11	Libra	
		Utils
		Stream
		Flow
		Fragment
	11.5	State machine
	11.6 11.7	Bus Slave Factory
		Fiber framework
		Bus
		Com
		IO
	11.12	Graphics
	11.13	EDĀ
		Misc
	11.15	Introduction
10	G• 1	105
12	Simul	
	12.1 12.2	Installation instructions
	12.2	Accessing signals of the simulation
	12.3	Clock domains
	12.5	Thread-full API
		Thread-less API
	12.7	Sensitive API

	12.9 12.10 12.11 12.12	Simulation engine  Examples  Introduction  How SpinalHDL simulates the hardware with Verilator backend  How SpinalHDL simulates the hardware with GHDL/Icarus Verilog backend  Performance	198 204 204 204
13	Exam 13.1 13.2 13.3 13.4		<b>207</b> 207 214 232 244
14		RiscV	
15	15.1 15.2	Bus Slave Factory Implementation	268
	16.1 16.2 16.3	ome to SpinalHDL's documentation!  Site purpose and structure	281 282

ONE

**FAQ** 

# 1.1 What is the overhead of SpinalHDL generated RTL compared to human written VHDL/Verilog?

The overhead is null, SpinalHDL is not an HLS approach. Its goal is not to translate any arbitrary code into RTL, but to provide a powerful language to describe RTL and raise the abstraction level.

# 1.2 What if SpinalHDL becomes unsupported in the future?

This question has two sides:

- 1. SpinalHDL generates VHDL/Verilog files, which means that SpinalHDL will be supported by all EDA tools for many decades.
- 2. If there is a bug in SpinalHDL and there is no longer support to fix it, it's not a deadly situation, because the SpinalHDL compiler is fully open source. For simple issues, you may be able to fix the issue yourself in few hours. Remember how much time it takes to EDA companies to fix issues or to add new features in their closed tools.

# 1.3 Does SpinalHDL keep comments in generated VHDL/verilog?

No, it doesn't. Generated files should be considered as a netlist. For example, when you compile C code, do you care about your comments in the generated assembly code?

# 1.4 Could SpinalHDL scale up to big projects?

Yes, some experiments were done, and it appears that generating hundreds of 3KLUT CPUs with caches takes around 12 seconds, which is a ridiculously short time compared to the time required to simulate or synthesize this kind of design.

# 1.5 How SpinalHDL came to be

Between December 2014 and April 2016, it was as a personal hobby project. But since April 2016 one person is working full time on it. Some people are also regularly contributing to the project.

# 1.6 Why develop a new language when there is VHDL/Verilog/SystemVerilog?

This page is dedicated to this topic.

# 1.7 How to use an unreleased version of SpinalHDL (but committed on git)?

For instance, if you want to try the dev branch, do the following in a dummy folder:

git clone https://github.com/SpinalHDL/SpinalHDL.git -b dev
cd SpinalHDL
sbt clean publishLocal

Then in your project, don't forget to update the SpinalHDL version specified in the build.sbt file, see https://github.com/SpinalHDL/SpinalTemplateSbt/blob/master/build.sbt#L10

To know which version you have to set, look in https://github.com/SpinalHDL/SpinalHDL/blob/dev/project/Version.scala#L7

2 Chapter 1. FAQ

**CHAPTER** 

**TWO** 

# **SUPPORT**

# 2.1 Communication channels

For bug reporting and feature requests, do not hesitate to create github issues: https://github.com/SpinalHDL/SpinalHDL/issues

For questions about SpinalHDL syntax and live talks, a Gitter channel is available: https://gitter.im/SpinalHDL/SpinalHDL

For questions, you can also use the forum StackOverflow with the tag SpinalHDL : https://stackoverflow.com/

A Google group is also available. Feel free to post whatever subject you want related to SpinalHDL: https://groups.google.com/forum/#!forum/spinalhdl-hardware-description-language

# 2.2 Commercial support

If you are interested in a presentation, a workshop, or consulting, do not hesitate to contact us by email: spinalhdl@gmail.com

# **CHAPTER**

# **THREE**

# **USERS**

# 3.1 Companies

• QsPin, Belgium

# 3.2 Repositories

- J1Sc Stack CPU
- VexRiscv CPU and SoC

6 Chapter 3. Users

### **GETTING STARTED**

# 4.1 Getting Started

*SpinalHDL* is a hardware description language written in Scala, a statically-typed functional language using the Java virtual machine (JVM). In order to start programming with *SpinalHDL*, you must have a JVM as well as the Scala compiler. In the next section, we will explain how to download those tools if you don't have them already.

## 4.1.1 Requirements / Things to download to get started

Before you download the SpinalHDL tools, you need to install:

- A Java JDK, which can be downloaded from here for instance.
- A Scala 2.11.X distribution, which can be downloaded here (not required if you use SBT).
- The SBT build tool, which can be downloaded here.

#### Optionally:

- An IDE (which is not compulsory). We advise you to get IntelliJ with its Scala plugin.
- Git, which is a tool for version control.

#### 4.1.2 How to start programming with SpinalHDL

Once you have downloaded all the requirements, there are two ways to get started with SpinalHDL programming.

- 1. The SBT way: If you already are familiar with the SBT build system and/or if you don't need an IDE.
- 2. The IDE way: Get a project already set up for you in an IDE and start programming right away.

## The SBT way

We have prepared a ready-to-go project for you on Github.

- Either clone or download the "getting started" repository.
- Open a terminal in the root of it and run sbt run. When you execute it for the first time, the process could take some time as it will download all the dependencies required to run *SpinalHDL*.

Normally, this command must generate an output file MyTopLevel.vhd, which corresponds to the top level *Spinal-HDL* code defined in src\main\scala\MyCode.scala, which corresponds to the *most simple SpinalHDL example* 

From a clean Debian distribution you can type the following commands into the shell. The commands will install Java, Scala, SBT, download the base project, and generate the corresponding VHDL file. Don't worry if it takes some time the first time that you run it.

```
sudo apt-get install openjdk-8-jdk
sudo apt-get install scala
echo "deb https://repo.scala-sbt.org/scalasbt/debian all main" | sudo tee /etc/apt/
sources.list.d/sbt.list
echo "deb https://repo.scala-sbt.org/scalasbt/debian /" | sudo tee /etc/apt/sources.
slist.d/sbt_old.list
curl -sL "https://keyserver.ubuntu.com/pks/lookup?op=get&
search=0x2EE0EA64E40A89B84B2DF73499E82A75642AC823" | sudo apt-key add
sudo apt-get update
sudo apt-get install sbt
git clone https://github.com/SpinalHDL/SpinalTemplateSbt.git SpinalTemplateSbt
cd SpinalTemplateSbt
sbt run # select "mylib.MyTopLevelVhdl" in the menu
ls MyTopLevel.vhd
```

#### SBT in a environnement isolated from internet

Normally, SBT uses online repositories to download and cache your projects dependencies, this cache is located in your home/.ivy2 folder. The way to set up an internet-free environnement is to copy this cache from an internet-full environnement where the cache was already filled once, and copy it over to your internet-less environnement.

You can get a portable SBT setup here:

https://www.scala-sbt.org/download.html

#### The IDE way, with IntelliJ IDEA and its Scala plugin

In addition to the aforementioned *requirements*, you also need to download the IntelliJ IDEA (the free *Community edition* is enough). When you have installed IntelliJ, also check that you have enabled its Scala plugin (install information can be found here).

And do the following:

- Either clone or download the "getting started" repository.
- In *Intellij IDEA*, "import project" with the root of this repository, the choose the *Import project from external model SBT* and be sure to check all boxes.
- In addition, you might need to specify some path like where you installed the JDK to IntelliJ.
- In the project (Intellij project GUI), right click on src/main/scala/mylib/MyTopLevel.scala and select "Run MyTopLevel".

This should generate the output file MyTopLevel.vhd in the project directory, which implements a simple 8-bit counter.

## 4.1.3 A very simple SpinalHDL example

The following code generates an and gate between two one-bit inputs.

```
import spinal.core._
class AND_Gate extends Component {
    /**
    * This is the component definition that corresponds to
```

(continues on next page)

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```
* the VHDL entity of the component

*/
val io = new Bundle {
  val a = in Bool()
  val b = in Bool()
  val c = out Bool()
}

// Here we define some asynchronous logic
io.c := io.a & io.b
}

object AND_Gate {
  // Let's go
  def main(args: Array[String]) {
    SpinalVhdl(new AND_Gate)
  }
}
```

As you can see, the first line you have to write in SpinalHDL is import spinal.core.\_ which indicates that we are using the *Spinal* components in the file.

#### **Generated code**

Once you have successfully compiled your code, the compiler should have emitted the following VHDL code:

```
package pkg_enum is
end pkg_enum;
package pkg_scala2hdl is
end pkg_scala2hdl;
library ieee;
use ieee.std_logic_1164.all;
use ieee.numeric_std.all;
library work;
use work.pkg_scala2hdl.all;
use work.all;
use work.pkg_enum.all;
entity AND_Gate is
 port(
    io_a : in std_logic;
    io_b : in std_logic;
    io_c : out std_logic
  );
end AND_Gate;
architecture arch of AND_Gate is
begin
```

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```
io_c <= (io_a and io_b);
end arch;</pre>
```

#### 4.1.4 What to do next?

It's up to you, but why not have a look at what the *types* are in SpinalHDL or discover what primitives the language provides to describe hardware components? You could also have a look at our *examples* to see some samples of what you could do next.

# 4.2 Motivation

Redirection to https://github.com/SpinalHDL/SpinalDoc/blob/master/presentation/en/workshop/taste.pdf

## 4.3 Presentation

Redirection to https://github.com/SpinalHDL/SpinalDoc/blob/master/presentation/en/presentation.pdf

## 4.4 Scala Guide

**Important:** Variables and functions should be defined into object, class, function. You can't define them on the root of a Scala file.

#### **4.4.1 Basics**

## **Types**

In Scala, there are 5 major types:

Туре	Literal	Description
Boolean	true, false	
Int	3, 0x32	32 bits integer
Float	3.14f	32 bits floating point
Double	3.14	64 bits floating point
String	"Hello world"	UTF-16 string

#### **Variables**

In Scala, you can define a variable by using the var keyword:

```
var number : Int = 0
number = 6
number += 4
println(number) // 10
```

Scala is able to infer the type automatically. You don't need to specify it if the variable is assigned at declaration:

```
var number = 0 //The type of 'number' is inferred as an Int during compilation.
```

However, it's not very common to use var in Scala. Instead, constant values defined by val are often used:

```
val two = 2
val three = 3
val six = two * three
```

#### **Functions**

For example, if you want to define a function which returns true if the sum of its two arguments is bigger than zero, you can do as follows:

```
def sumBiggerThanZero(a: Float, b: Float): Boolean = {
  return (a + b) > 0
}
```

Then, to call this function, you can write:

```
sumBiggerThanZero(2.3f, 5.4f)
```

You can also specify arguments by name, which is useful if you have many arguments:

```
sumBiggerThanZero(
   a = 2.3f,
   b = 5.4f
)
```

#### Return

The return keyword is not necessary. In absence of it, Scala takes the last statement of your function as the returned value.

```
def sumBiggerThanZero(a: Float, b: Float): Boolean = {
  (a + b) > 0
}
```

#### Return type inferation

Scala is able to automatically infer the return type. You don't need to specify it:

```
def sumBiggerThanZero(a: Float, b: Float) = {
   (a + b) > 0
}
```

4.4. Scala Guide 11

#### **Curly braces**

Scala functions don't require curly braces if your function contains only one statement:

```
def sumBiggerThanZero(a: Float, b: Float) = (a + b) > 0
```

#### **Function that returns nothing**

If you want a function to return nothing, the return type should be set to Unit. It's equivalent to the C/C++ void type.

```
def printer(): Unit = {
  println("1234")
  println("5678")
}
```

#### **Argument default values**

You can specify a default value for each argument of a function:

```
def sumBiggerThanZero(a: Float, b: Float = 0.0f) = {
   (a + b) > 0
}
```

#### **Apply**

Functions named apply are special because you can call them without having to type their name:

```
class Array() {
  def apply(index: Int): Int = index + 3
}

val array = new Array()
val value = array(4)  //array(4) is interpreted as array.apply(4) and will return 7
```

This concept is also applicable for Scala object (static)

```
object MajorityVote {
  def apply(value: Int): Int = ...
}
val value = MajorityVote(4) // Will call MajorityVote.apply(4)
```

#### **Object**

In Scala, there is no static keyword. In place of that, there is object. Everything defined inside an object definition is static.

The following example defines a static function named pow2 which takes a floating point value as parameter and returns a floating point value as well.

```
object MathUtils {
  def pow2(value: Float): Float = value * value
}
```

Then you can call it by writing:

```
MathUtils.pow2(42.0f)
```

#### **Entry point (main)**

The entry point of a Scala program (the main function) should be defined inside an object as a function named main.

```
object MyTopLevelMain{
  def main(args: Array[String]) {
    println("Hello world")
  }
}
```

#### **Class**

The class syntax is very similar to Java. Imagine that you want to define a Color class which takes as construction parameters three Float values (r,g,b):

```
class Color(r: Float, g: Float, b: Float) {
  def getGrayLevel(): Float = r * 0.3f + g * 0.4f + b * 0.4f
}
```

Then, to instantiate the class from the previous example and use its getGrayLevel function:

```
val blue = new Color(0, 0, 1)
val grayLevelOfBlue = blue.getGrayLevel()
```

Be careful, if you want to access a construction parameter of the class from the outside, this construction parameter should be defined as a val:

```
class Color(val r: Float, val g: Float, val b: Float) { ... }
...
val blue = new Color(0, 0, 1)
val redLevelOfBlue = blue.r
```

4.4. Scala Guide 13

#### Inheritance

As an example, suppose that you want to define two classes, Rectangle and Square, which extend the class Shape:

```
class Shape {
  def getArea(): Float
}

class Square(sideLength: Float) extends Shape {
  override def getArea() = sideLength * sideLength
}

class Rectangle(width: Float, height: Float) extends Shape {
  override def getArea() = width * height
}
```

#### Case class

Case class is an alternative way of declaring classes.

```
case class Rectangle(width: Float, height: Float) extends Shape {
  override def getArea() = width * height
}
```

Then there are some differences between case class and class:

- case classes don't need the new keyword to be instantiated.
- construction parameters are accessible from outside; you don't need to define them as val.

In SpinalHDL, this explains the reasoning behind the coding conventions: it's in general recommended to use case class instead of class in order to have less typing and more coherency.

#### Templates / Type parameterization

Imagine you want to design a class which is a queue of a given datatype, in that case you need to provide a type parameter to the class:

```
class Queue[T](){
  def push(that: T) : Unit = ...
  def pop(): T = ...
}
```

If you want to restrict the T type to be a sub class of a given type (for example Shape), you can use the <: Shape syntax :

```
class Shape() {
    def getArea(): Float
}
class Rectangle() extends Shape { ... }

class Queue[T <: Shape]() {
    def push(that: T): Unit = ...
    def pop(): T = ...
}</pre>
```

The same is possible for functions:

```
def doSomething[T <: Shape](shape: T): Something = { shape.getArea() }</pre>
```

# 4.4.2 Coding conventions

#### Introduction

The coding conventions used in SpinalHDL are the same as the ones documented in the Scala Style Guide. Some additional practical details and cases are explained in next pages.

#### class vs case class

When you define a Bundle or a Component, it is preferable to declare it as a case class.

The reasons are:

- It avoids the use of new keywords. Never having to use it is better than sometimes, under some conditions.
- A case class provides a clone function. This is useful in SpinalHDL when there is a need to clone a Bundle, for example, when you define a new Reg or a new Stream of some kind.
- Construction parameters are directly visible from outside.

#### [case] class

All classes names should start with a uppercase letter

```
class Fifo extends Component {
}
class Counter extends Area {
}
case class Color extends Bundle {
}
```

#### companion object

A companion object should start with an uppercase letter.

```
object Fifo {
  def apply(that: Stream[Bits]): Stream[Bits] = {...}
}
object MajorityVote {
  def apply(that: Bits): UInt = {...}
}
```

An exception to this rule is when the companion object is used as a function (only apply inside), and these apply functions don't generate hardware:

4.4. Scala Guide 15

```
object log2 {
  def apply(value: Int): Int = {...}
}
```

#### **function**

A function should always start with a lowercase letter:

```
def sinTable = (0 until sampleCount).map(sampleIndex => {
   val sinValue = Math.sin(2 * Math.PI * sampleIndex / sampleCount)
   S((sinValue * ((1 << resolutionWidth) / 2 - 1)).toInt, resolutionWidth bits)
})

val rom = Mem(SInt(resolutionWidth bit), initialContent = sinTable)</pre>
```

#### instances

Instances of classes should always start with a lowercase letter:

```
val fifo = new Fifo()
val buffer = Reg(Bits(8 bits))
```

#### if / when

Scala if and SpinalHDL when should normally be written in the following way:

```
if(cond) {
    ...
} else if(cond) {
    ...
} else {
    ...
}
when(cond) {
    ...
}.elseWhen(cond) {
    ...
}.otherwise {
    ...
}
```

Exceptions could be:

- It's fine to omit the dot before otherwise.
- It's fine to compress an if/when statement onto a single line if it makes the code more readable.

#### switch

SpinalHDL switch should normally be written in the following way:

```
switch(value) {
  is(key) {
  }
  is(key) {
  }
  default {
  }
}
```

It's fine to compress an is/default statement onto a single line if it makes the code more readable.

#### **Parameters**

Grouping parameters of a Component/Bundle inside a case class is generally welcome because:

- Easier to carry/manipulate to configure the design
- · Better maintainability

```
case class RgbConfig(rWidth: Int, gWidth: Int, bWidth: Int) {
  def getWidth = rWidth + gWidth + bWidth
}

case class Rgb(c: RgbConfig) extends Bundle {
  val r = UInt(c.rWidth bit)
  val g = UInt(c.gWidth bit)
  val b = UInt(c.bWidth bit)
}
```

But this should not be applied in all cases. For example: in a FIFO, it doesn't make sense to group the dataType parameter with the depth parameter of the fifo because, in general, the dataType is something related to the design, while the depth is something related to the configuration of the design.

```
class Fifo[T <: Data](dataType: T, depth: Int) extends Component {
}</pre>
```

## 4.4.3 Interaction

#### Introduction

SpinalHDL is, in fact, not an language: it's a regular Scala library. This could seem strange at first glance, but it is a very powerful combination.

You can use the whole Scala world to help you in the description of your hardware via the SpinalHDL library, but to do that properly, it's important to understand how SpinalHDL interacts with Scala.

4.4. Scala Guide

#### How SpinalHDL works behind the API

When you execute your SpinalHDL hardware description, each time you use SpinalHDL functions, operators, or classes, it will build an in-memory graph that represents the netlist of your design.

Then, when the elaboration is done (instantiation of your top-level Component classes), SpinalHDL will do some passes on the graph that was constructed, and if everything is fine, it will flush that graph into a VHDL or Verilog file.

#### **Everything is a reference**

For example, if you define a Scala function which takes a parameter of type Bits, when you call it, it will be passed as a reference. As consequence of that, if you assign that argument inside the function, it has the same effect on the underlying Bits object as if you had assigned to it outside the function.

#### **Hardware types**

Hardware data types in SpinalHDL are the combination of two things:

- An instance of a given Scala type
- The configuration of that instance

For example Bits(8 bits) is the combination of the Scala type Bits and its 8 bits configuration (as a construction parameter).

#### **RGB** example

Let's take an Rgb bundle class as example:

```
case class Rgb(rWidth: Int, gWidth: Int, bWidth: Int) extends Bundle {
  val r = UInt(rWidth bits)
  val g = UInt(gWidth bits)
  val b = UInt(bWidth bits)
}
```

The hardware data type here is the combination of the Scala Rgb class and its rWidth, gWidth, and bWidth parameterization.

Here is an example of usage:

```
// Define an Rgb signal
val myRgbSignal = Rgb(5, 6, 5)

// Define another Rgb signal of the same data type as the preceding one
val myRgbCloned = cloneOf(myRgbSignal)
```

You can also use functions to define various kinds of type factories (typedef):

```
// Define a type factory function
def myRgbTypeDef = Rgb(5, 6, 5)

// Use that type factory to create an Rgb signal
val myRgbFromTypeDef = myRgbTypeDef
```

#### Names of signals in the generated RTL

To name signals in the generated RTL, SpinalHDL uses Java reflections to walk through your entire component hierarchy, collecting all references stored inside the class attributes, and naming them with their attribute name.

This is why the names of every signal defined inside a function are lost:

```
def myFunction(arg: UInt) {
   val temp = arg + 1 // You will not retrieve the `temp` signal in the generated RTL
   return temp
}
val value = myFunction(U"000001") + 42
```

One solution if you want preserve the names of the internal variables in the generated RTL, is to use Area:

```
def myFunction(arg: UInt) new Area {
  val temp = arg + 1 // You will not retrieve the temp signal in the generated RTL
}

val myFunctionCall = myFunction(U"000001") // Will generate `temp` with_
  `myFunctionCall_temp` as the name
val value = myFunctionCall.temp + 42
```

#### Scala is for elaboration, SpinalHDL for hardware description

For example, if you write a Scala for loop to generate some hardware, it will generate the unrolled result in VHDL/Verilog.

Also, if you want a constant, you should not use SpinalHDL hardware literals but the Scala ones. For example:

#### Scala elaboration capabilities (if, for, functional programming)

All of Scala's syntax can be used to elaborate hardware designs, for instance, a Scala if statement could be used to enable or disable the generation of hardware:

The same is true for Scala for loops:

4.4. Scala Guide 19

Also, functional programming techniques can be used with many SpinalHDL types:

## 4.4.4 Scala guide

#### Introduction

Scala is a very capable programming language that was influenced by a unique set of languages, but often, this set of languages doesn't cross the ones that most programmers use. That can hinder newcomers' understanding of the concepts and design choices behind Scala.

The following pages will present Scala, and try to provide enough information about it for newcomers to be comfortable with SpinalHDL.

# 4.5 Help for VHDL people

## 4.5.1 VHDL comparison

#### Introduction

This page will show the main differences between VHDL and SpinalHDL. Things will not be explained in depth.

#### **Process**

Processes are often needed when you write RTL, however, their semantics can be clunky to work with. Due to how they work in VHDL, they can force you to split your code and duplicate things.

To produce the following RTL:



You will have to write the following VHDL:

```
signal mySignal : std_logic;
 signal myRegister : std_logic_vector(3 downto 0);
  signal myRegisterWithReset : std_logic_vector(3 downto 0);
begin
 process(cond)
 begin
    mySignal <= '0';</pre>
    if cond = '1' then
      mySignal <= '1';</pre>
    end if;
 end process;
 process(clk)
 begin
    if rising_edge(clk) then
      if cond = '1' then
        myRegister <= myRegister + 1;</pre>
      end if;
    end if:
 end process;
 process(clk,reset)
 begin
    if reset = '1' then
      myRegisterWithReset <= (others => '0');
    elsif rising_edge(clk) then
      if cond = '1' then
        myRegisterWithReset <= myRegisterWithReset + 1;</pre>
      end if;
    end if:
 end process;
```

While in SpinalHDL, it's:

```
val mySignal = Bool()
val myRegister = Reg(UInt(4 bits))
val myRegisterWithReset = Reg(UInt(4 bits)) init(0)

mySignal := False
when(cond) {
  mySignal := True
  myRegister := myRegister + 1
  myRegisterWithReset := myRegisterWithReset + 1
}
```

#### Implicit vs explicit definitions

In VHDL, when you declare a signal, you don't specify if it is a combinatorial signal or a register. Where and how you assign to it decides whether it is combinatorial or registered.

In SpinalHDL these kinds of things are explicit. Registers are defined as registers directly in their declaration.

#### **Clock domains**

In VHDL, every time you want to define a bunch of registers, you need the carry the clock and the reset wire to them. In addition, you have to hardcode everywhere how those clock and reset signals should be used (clock edge, reset polarity, reset nature (async, sync)).

In SpinalHDL you can define a ClockDomain, and then define the area of your hardware that uses it.

For example:

#### Component's internal organization

In VHDL, there is a block feature that allows you to define sub-areas of logic inside your component. However, almost no one uses this feature, because most people don't know about them, and also because all signals defined inside these regions are not readable from the outside.

In SpinalHDL you have an Area feature that does this concept much more nicely:

```
val timeout = new Area {
  val counter = Reg(UInt(8 bits)) init(0)
  val overflow = False
  when(counter =/= 100) {
    counter := counter + 1
  } otherwise {
    overflow := True
  }
}

val core = new Area {
  when(timeout.overflow) {
    timeout.counter := 0
  }
}
```

Variables and signals defined inside of an Area are accessible elsewhere in the component, including in other Area regions.

#### Safety

In VHDL as in SpinalHDL, it's easy to write combinatorial loops, or to infer a latch by forgetting to drive a signal in the path of a process.

Then, to detect those issues, you can use some lint tools that will analyze your VHDL, but those tools aren't free. In SpinalHDL the lint process in integrated inside the compiler, and it won't generate the RTL code until everything is fine. It also checks clock domain crossing.

#### **Functions and procedures**

Functions and procedures are not used very often in VHDL, probably because they are very limited:

- You can only define a chunk of combinational hardware, or only a chunk of registers (if you call the function/procedure inside a clocked process).
- You can't define a process inside them.
- You can't instantiate a component inside them.
- The scope of what you can read/write inside them is limited.

In SpinalHDL, all those limitations are removed.

An example that mixes combinational logic and a register in a single function:

```
def simpleAluPipeline(op: Bits, a: UInt, b: UInt): UInt = {
  val result = UInt(8 bits)

switch(op) {
  is(0){ result := a + b }
  is(1){ result := a - b }
  is(2){ result := a * b }
}

return RegNext(result)
}
```

An example with the queue function inside the Stream Bundle (handshake). This function instantiates a FIFO component:

```
class Stream[T <: Data](dataType: T) extends Bundle with IMasterSlave with_
    DataCarrier[T] {
    val valid = Bool()
    val ready = Bool()
    val payload = cloneOf(dataType)

    def queue(size: Int): Stream[T] = {
        val fifo = new StreamFifo(dataType, size)
        fifo.io.push <> this
        fifo.io.pop
    }
}
```

An example where a function assigns a signal defined outside of itself:

```
val counter = Reg(UInt(8 bits)) init(0)
counter := counter + 1

def clear() : Unit = {
   counter := 0
}
when(counter > 42) {
   clear()
}
```

#### **Buses and Interfaces**

VHDL is very boring when it comes to buses and interfaces. You have two options:

1) Define buses and interfaces wire-by-wire, each time and everywhere:

```
PADDR : in unsigned(addressWidth-1 downto 0);

PSEL : in std_logic;

PENABLE : in std_logic;

PREADY : out std_logic;

PWRITE : in std_logic;

PWDATA : in std_logic_vector(dataWidth-1 downto 0);

PRDATA : out std_logic_vector(dataWidth-1 downto 0);
```

2) Use records but lose parameterization (statically fixed in the package), and you have to define one for each directions:

```
P_m : in APB_M;
P_s : out APB_S;
```

SpinalHDL has very strong support for bus and interface declarations with limitless parameterizations:

```
val P = slave(Apb3(addressWidth, dataWidth))
```

You can also use object oriented programming to define configuration objects:

```
val coreConfig = CoreConfig(
 pcWidth = 32,
 addrWidth = 32,
  startAddress = 0x000000000,
 regFileReadyKind = sync,
 branchPrediction = dynamic,
 bypassExecute0 = true,
 bypassExecute1 = true,
 bypassWriteBack = true,
 bypassWriteBackBuffer = true,
 collapseBubble = false,
 fastFetchCmdPcCalculation = true,
 dynamicBranchPredictorCacheSizeLog2 = 7
)
// The CPU has a system of plugins which allows adding new features into the core.
// Those extensions are not directly implemented in the core, but are kind of an.
→additive logic patch defined in a separate area.
coreConfig.add(new MulExtension)
coreConfig.add(new DivExtension)
```

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#### Signal declaration

VHDL forces you to define all signals at the top of your architecture description, which is annoying.

```
..
.. (many signal declarations)
..
signal a : std_logic;
..
.. (many signal declarations)
..
begin
..
.. (many logic definitions)
..
a <= x & y
..
.. (many logic definitions)
..</pre>
```

SpinalHDL is flexible when it comes to signal declarations.

```
val a = Bool
a := x & y
```

It also allows you to define and assign signals in a single line.

```
val a = x & y
```

#### **Component instantiation**

VHDL is very verbose about this, as you have to redefine all signals of your sub-component entity, and then bind them one-by-one when you instantiate your component.

```
divider_cmd_valid : in std_logic;
divider_cmd_ready : out std_logic;
divider_cmd_numerator : in unsigned(31 downto 0);
divider_cmd_denominator : in unsigned(31 downto 0);
divider_rsp_valid : out std_logic;
divider_rsp_ready : in std_logic;
divider_rsp_quotient : out unsigned(31 downto 0);
divider_rsp_remainder : out unsigned(31 downto 0);
divider : entity work.UnsignedDivider
 port map (
    clk
                    => clk,
   reset
                   => reset,
    cmd_valid => divider_cmd_valid,
cmd_ready => divider_cmd_ready,
    cmd_numerator => divider_cmd_numerator,
    cmd_denominator => divider_cmd_denominator,
   rsp_valid => divider_rsp_valid,
rsp_ready => divider_rsp_ready,
    rsp_quotient => divider_rsp_quotient,
    rsp_remainder => divider_rsp_remainder
 );
```

SpinalHDL removes that, and allows you to access the IO of sub-components in an object-oriented way.

```
val divider = new UnsignedDivider()

// And then if you want to access IO signals of that divider:
divider.io.cmd.valid := True
divider.io.cmd.numerator := 42
```

#### Casting

There are two annoying casting methods in VHDL:

- boolean <> std\_logic (ex: To assign a signal using a condition such as mySignal <= myValue < 10 is not legal)
- unsigned <> integer (ex: To access an array)

SpinalHDL removes these casts by unifying things.

boolean/std\_logic:

```
val value = UInt(8 bits)
val valueBiggerThanTwo = Bool
valueBiggerThanTwo := value > 2 // value > 2 return a Bool
```

unsigned/integer:

```
val array = Vec(UInt(4 bits),8)
val sel = UInt(3 bits)
val arraySel = array(sel) // Arrays are indexed directly by using UInt
```

#### Resizing

The fact that VHDL is strict about bit size is probably a good thing.

```
my8BitsSignal <= resize(my4BitsSignal, 8);</pre>
```

In SpinalHDL you have two ways to do the same:

```
// The traditional way
my8BitsSignal := my4BitsSignal.resize(8)

// The smart way
my8BitsSignal := my4BitsSignal.resized
```

#### **Parameterization**

VHDL prior to the 2008 revision has many issues with generics. For example, you can't parameterize records, you can't parameterize arrays in the entity, and you can't have type parameters.

Then VHDL 2008 came and fixed those issues. But RTL tool support for VHDL 2008 is really weak depending on the vendor.

SpinalHDL has full support for generics integrated natively in its compiler, and it doesn't rely on VHDL generics.

Here is an example of parameterized data structures:

```
val colorStream = Stream(Color(5, 6, 5)))
val colorFifo = StreamFifo(Color(5, 6, 5), depth = 128)
colorFifo.io.push <> colorStream
```

Here is an example of a parameterized component:

```
class Arbiter[T <: Data](payloadType: T, portCount: Int) extends Component {
  val io = new Bundle {
    val sources = Vec(slave(Stream(payloadType)), portCount)
    val sink = master(Stream(payloadType))
  }
  // ...
}</pre>
```

#### Meta hardware description

VHDL has kind of a closed syntax. You can't add abstraction layers on top of it.

SpinalHDL, because it's built on top of Scala, is very flexible, and allows you to define new abstraction layers very easily.

Some examples of this flexibility are the FSM library, the BusSlaveFactory library, and also the JTAG library.

## 4.5.2 VHDL equivalences

### **Entity and architecture**

In SpinalHDL, a VHDL entity and architecture are both defined inside a Component.

Here is an example of a component which has 3 inputs (a, b, c) and an output (result). This component also has an offset construction parameter (like a VHDL generic).

```
case class MyComponent(offset: Int) extends Component {
  val io = new Bundle{
    val a, b, c = in UInt(8 bits)
    val result = out UInt(8 bits)
}
io.result := a + b + c + offset
}
```

Then to instantiate that component, you don't need to bind it:

```
case class TopLevel extends Component {
    ...
    val mySubComponent = MyComponent(offset = 5)
    ...
    mySubComponent.io.a := 1
    mySubComponent.io.b := 2
    mySubComponent.io.c := 3
    ???? := mySubComponent.io.result
    ...
}
```

#### **Data types**

SpinalHDL data types are similar to the VHDL ones:

VHDL	SpinalHDL
std_logic	Bool
std_logic_vector	Bits
unsigned	UInt
signed	SInt

In VHDL, to define an 8 bit unsigned you have to give the range of bits unsigned(7 downto 0), whereas in SpinalHDL you simply supply the number of bits UInt(8 bits).

VHDL	SpinalHDL
records	Bundle
array	Vec
enum	SpinalEnum

Here is an example of the SpinalHDL Bundle definition. channelWidth is a construction parameter, like VHDL generics, but for data structures:

```
case class RGB(channelWidth: Int) extends Bundle {
  val r, g, b = UInt(channelWidth bits)
}
```

Then for example, to instantiate a Bundle, you need to write val myColor = RGB(channelWidth=8).

#### **Signal**

Here is an example about signal instantiations:

```
case class MyComponent(offset: Int) extends Component {
  val io = new Bundle {
    val a, b, c = UInt(8 bits)
    val result = UInt(8 bits)
}

val ab = UInt(8 bits)
ab := a + b

val abc = ab + c  // You can define a signal directly with its value
io.result := abc + offset
}
```

#### **Assignments**

In SpinalHDL, the := assignment operator is equivalent to the VHDL signal assignment (<=):

```
val myUInt = UInt(8 bits)
myUInt := 6
```

Conditional assignments are done like in VHDL by using if/case statements:

```
val clear
           = Bool()
val counter = Reg(UInt(8 bits))
when(clear) {
 counter := 0
}.elsewhen(counter === 76) {
  counter := 79
}.otherwise {
  counter(7) := ! counter(7)
switch(counter) {
  is(42) {
    counter := 65
  default {
    counter := counter + 1
  }
}
```

#### Literals

Literals are a little bit different than in VHDL:

```
val myBool = Bool()
myBool := False
myBool := True
myBool := Bool(4 > 7)

val myUInt = UInt(8 bits)
myUInt := "0001_1100"
myUInt := "xEE"
myUInt := 42
myUInt := U(54,8 bits)
myUInt := ((3 downto 0) -> myBool, default -> true)
when(myUInt === U(myUInt.range -> true)) {
    myUInt(3) := False
}
```

#### **Registers**

In SpinalHDL, registers are explicitly specified while in VHDL registers are inferred. Here is an example of SpinalHDL registers:

```
val counter = Reg(UInt(8 bits)) init(0)
counter := counter + 1  // Count up each cycle

// init(0) means that the register should be initialized to zero when a reset occurs
```

#### **Process blocks**

Process blocks are a simulation feature that is unnecessary to design RTL. It's why SpinalHDL doesn't contain any feature analogous to process blocks, and you can assign what you want, where you want.

```
val cond = Bool()
val myCombinatorial = Bool()
val myRegister = UInt(8 bits)

myCombinatorial := False
when(cond) {
   myCombinatorial := True
   myRegister = myRegister + 1
}
```

### 4.6 Cheatsheets

#### 4.6.1 Core

Redirection to https://github.com/SpinalHDL/SpinalDoc/blob/master/cheatSheet\_core\_oo.pdf

## 4.6.2 Lib

 $Redirection\ to\ https://github.com/SpinalHDL/SpinalDoc/blob/master/cheatSheet\_lib\_oo.pdf$ 

# 4.6.3 Symbolic

 $Redirection\ to\ https://github.com/SpinalHDL/SpinalDoc/blob/master/cheatSheet\_symbolic.pdf$ 

4.6. Cheatsheets 31

**CHAPTER** 

**FIVE** 

## **DATA TYPES**

# **5.1 Bool**

# 5.1.1 Description

The Bool type corresponds to a boolean value (True or False).

## 5.1.2 Declaration

The syntax to declare a boolean value is as follows: (everything between [] is optional)

Syntax	Description	Return
Bool[()]	Create a Bool	Bool
True	Create a Bool assigned with true	Bool
False	Create a Bool assigned with false	Bool
Bool(value:	Create a Bool assigned with a Scala Boolean(true, false)	Bool
Boolean)		

## 5.1.3 Operators

The following operators are available for the Bool type:

## Logic

Operator	Description	Return type
!x	Logical NOT	Bool
	Logical AND	Bool
x && y		
x & y		
	Logical OR	Bool
x    y		
x y		
x ^ y	Logical XOR	Bool
x.set[()]	Set x to True	
x.clear[()]	Set x to False	
x.setWhen(cond)	Set x when cond is True	Bool
x.clearWhen(cond)	Clear x when cond is True	Bool
x.riseWhen(cond)	Set x when x is False and cond is	Bool
	True	
x.fallWhen(cond)	Clear x when x is True and cond is	Bool
	True	

```
val a, b, c = Bool()
val res = (!a \& b) \land c // ((NOT a) AND b) XOR c
val d = False
when(cond) {
  d.set() // equivalent to d := True
}
val e = False
e.setWhen(cond) // equivalent to when(cond) { d := True }
val f = RegInit(False) fallWhen(ack) setWhen(req)
 /** equivalent to
 * when(f && ack) { f := False }
 * when(req) { f := True }
 * or
 * f := req || (f && !ack)
// mind the order of assignments!
val g = RegInit(False) setWhen(req) fallWhen(ack)
// equivalent to g := ((!g) && req) || (g && !ack)
```

## **Edge detection**

Operator	Description	Return
		type
x.edge[()]	Return True when x changes state	Bool
x.edge(initAt: Bool)	Same as x.edge but with a reset value	Bool
x.rise[()]	Return True when x was low at the last cycle and is now high	Bool
x.rise(initAt: Bool)	Same as x.rise but with a reset value	Bool
x.fall[()]	Return True when x was high at the last cycle and is now low	Bool
x.fall(initAt: Bool)	Same as x.fall but with a reset value	Bool
x.edges[()]	Return a bundle (rise, fall, toggle)	BoolEdges
x.edges(initAt: Bool)	Same as x.edges but with a reset value	BoolEdges

```
when(myBool_1.rise(False)) {
    // do something when a rising edge is detected
}

val edgeBundle = myBool_2.edges(False)
when(edgeBundle.rise) {
    // do something when a rising edge is detected
}
when(edgeBundle.fall) {
    // do something when a falling edge is detected
}
when(edgeBundle.toggle) {
    // do something at each edge
}
```

## Comparison

Operator	Description	Return type
x === y	Equality	Bool
x =/= y	Inequality	Bool

```
when(myBool) { // Equivalent to when(myBool === True)
     // do something when myBool is True
}
when(!myBool) { // Equivalent to when(myBool === False)
     // do something when myBool is False
}
```

5.1. Bool 35

#### Type cast

Operator	Description	Return
x.asBits	Binary cast to Bits	Bits(w(x) bits)
x.asUInt	Binary cast to UInt	UInt(w(x) bits)
x.asSInt	Binary cast to SInt	SInt(w(x) bits)
x.asUInt(bitCount)	Binary cast to UInt and resize	UInt(bitCount bits)
x.asBits(bitCount)	Binary cast to Bits and resize	Bits(bitCount bits)

```
// Add the carry to an SInt value
val carry = Bool()
val res = mySInt + carry.asSInt
```

#### Misc

Operator	Description	Return
x ## y	Concatenate, x->high, y->low	Bits(w(x) + w(y) bits)

```
val a, b, c = Bool

// Concatenation of three Bool into a Bits
val myBits = a ## b ## c
```

## **5.2 Bits**

## 5.2.1 Description

The Bits type corresponds to a vector of bits that does not convey any arithmetic meaning.

## 5.2.2 Declaration

The syntax to declare a bit vector is as follows: (everything between [] is optional)

Syntax	Description	Return
Bits [()]	Create a BitVector, bits count is inferred	Bits
Bits(x bits)	Create a BitVector with x bits	Bits
	Create a BitVector with x bits assigned with 'value'	Bits
B(value: Int[, x bits])		
B(value: BigInt[, x bits])		
B"[[size']base]value"	Create a BitVector assigned with 'value' (Base: 'h', 'd', 'o', 'b')	Bits
B([x bits,] element,)	Create a BitVector assigned with the value specified by elements	Bits

```
// Declaration
val myBits = Bits()  // the size is inferred
val myBits1 = Bits(32 bits)
val myBits2 = B(25, 8 bits)
val myBits3 = B"8'xFF"  // Base could be x,h (base 16)
```

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## 5.2.3 Operators

The following operators are available for the Bits type:

#### Logic

Operator	Description	Return type
~X	Bitwise NOT	Bits(w(x) bits)
x & y	Bitwise AND	Bits(w(xy) bits)
x   y	Bitwise OR	Bits(w(xy) bits)
x ^ y	Bitwise XOR	Bits(w(xy) bits)
x.xorR	XOR all bits of x	Bool
x.orR	OR all bits of x	Bool
x.andR	AND all bits of x	Bool
x >> y	Logical shift right, y: Int	Bits(w(x) - y bits)
x >> y	Logical shift right, y: UInt	Bits(w(x) bits)
x << y	Logical shift left, y: Int	Bits(w(x) + y bits)
x << y	Logical shift left, y: UInt	Bits(w(x) + max(y) bits)
x  >> y	Logical shift right, y: Int/UInt	Bits(w(x) bits)
x  << y	Logical shift left, y: Int/UInt	Bits(w(x) bits)
x.rotateLeft(y)	Logical left rotation, y: UInt/Int	Bits(w(x) bits)
x.rotateRight(y)	Logical right rotation, y: UInt/Int	Bits(w(x) bits)
x.clearAll[()]	Clear all bits	
x.setAll[()]	Set all bits	
x.setAllTo(value: Boolean)	Set all bits to the given Boolean value	
x.setAllTo(value: Bool)	Set all bits to the given Bool value	

```
// Bitwise operator
val a, b, c = Bits(32 bits)
c := ~(a & b) // Inverse(a AND b)

val all_1 = a.andR // Check that all bits are equal to 1

// Logical shift
val bits_10bits = bits_8bits << 2 // shift left (results in 10 bits)
val shift_8bits = bits_8bits |<< 2 // shift left (results in 8 bits)

// Logical rotation
val myBits = bits_8bits.rotateLeft(3) // left bit rotation</pre>
```

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5.2. Bits 37

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```
// Set/clear
val a = B"8'x42"
when(cond) {
  a.setAll() // set all bits to True when cond is True
}
```

## Comparison

Operator	Description	Return type
x === y	Equality	Bool
x =/= y	Inequality	Bool

```
when(myBits === 3) {
}
when(myBits_32 =/= B"32'x44332211") {
}
```

## Type cast

Operator	Description	Return
x.asBits	Binary cast to Bits	Bits(w(x) bits)
x.asUInt	Binary cast to UInt	UInt(w(x) bits)
x.asSInt	Binary cast to SInt	SInt(w(x) bits)
x.asBools	Cast to an array of Bools	Vec(Bool, w(x))
B(x: T)	Cast Data to Bits	Bits(w(x) bits)

To cast a Bool, UInt or an SInt into a Bits, you can use B(something):

```
// cast a Bits to SInt
val mySInt = myBits.asSInt

// create a Vector of bool
val myVec = myBits.asBools

// Cast a SInt to Bits
val myBits = B(mySInt)
```

## Bit extraction

Operator	Description	Return
x(y)	Readbit, y: Int/UInt	Bool
x(offset,width bits)	Read bitfield, offset: UInt, width: Int	Bits(width bits)
x(range)	Read a range of bit. Ex: myBits(4 downto 2)	Bits(range bits)
x(y) := z	Assign bits, y: Int/UInt	Bool
x(offset, width bits) :=	Assign bitfield, offset: UInt, width: Int	Bits(width bits)
Z		
x(range) := z	Assign a range of bit. Ex: myBits(4 downto 2) := B"010"	Bits(range bits)

```
// get the element at the index 4
val myBool = myBits(4)

// assign
myBits(1) := True

// Range
val myBits_8bits = myBits_16bits(7 downto 0)
val myBits_7bits = myBits_16bits(0 to 6)
val myBits_6bits = myBits_16Bits(0 until 6)

myBits_8bits(3 downto 0) := myBits_4bits
```

#### **Misc**

```
println(myBits_32bits.getWidth) // 32
myBool := myBits.lsb // Equivalent to myBits(0)
// Concatenation
myBits_24bits := bits_8bits_1 ## bits_8bits_2 ## bits_8bits_3
// Subdivide
val sel = UInt(2 bits)
val myBitsWord = myBits_128bits.subdivideIn(32 bits)(sel)
   // sel = 0 => myBitsWord = myBits_128bits(127 downto 96)
   // sel = 1 => myBitsWord = myBits_128bits( 95 downto 64)
   // sel = 2 => myBitsWord = myBits_128bits( 63 downto 32)
   // sel = 3 => myBitsWord = myBits_128bits( 31 downto 0)
// If you want to access in reverse order you can do:
val myVector = myBits_128bits.subdivideIn(32 bits).reverse
val myBitsWord = myVector(sel)
// Resize
myBits_32bits := B"32'x112233344"
myBits_8bits := myBits_32bits.resizeLeft(8) // resize to 8 bits (myBits_8bits = 0x11)
```

5.2. Bits 39

## 5.3 UInt/SInt

## 5.3.1 Description

The UInt/SInt type corresponds to a vector of bits that can be used for signed/unsigned integer arithmetic.

## 5.3.2 Declaration

The syntax to declare an integer is as follows: (everything between [] is optional)

Syntax	Description	Return
	Create an unsigned/signed integer, bits count is inferred	
UInt[()]		UInt
SInt[()]		SInt
	Create an unsigned/signed integer with x bits	
UInt(x bits)		UInt
SInt(x bits)		SInt
	Create an unsigned/signed integer assigned with 'value'	
U(value: Int[,x bits])		UInt
U(value: BigInt[,x bits])		SInt
S(value: Int[,x bits])		
S(value: BigInt[,x bits])		
	Create an unsigned/signed integer assigned with 'value' (Base:	
U"[[size']base]value"	'h', 'd', 'o', 'b')	UInt
S"[[size']base]value"		SInt
	Create an unsigned integer assigned with the value specified by	
$U([x \text{ bits,}] \text{ element,} \dots)$	elements	UInt
S([x bits,] element,)		SInt

```
val myUInt = UInt(8 bits)
myUInt := U(2,8 bits)
myUInt := U(2)
myUInt := U"0000_0101" // Base per default is binary => 5
                        // Base could be x (base 16)
myUInt := U"h1A"
                                        h (base 16)
                        //
                        //
                                          d (base 10)
                                          o (base 8)
                                          b (base 2)
myUInt := U"8'h1A"
myUInt := 2
                        // You can use a Scala Int as a literal value
val myBool := myUInt === U(7 \rightarrow true, (6 downto 0) \rightarrow false)
val myBool := myUInt === U(myUInt.range -> true)
// For assignment purposes, you can omit the U/S, which also allows the use of the
→[default -> ???] feature
```

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## 5.3.3 Operators

The following operators are available for the UInt and SInt types:

## Logic

Operator	Description	Return type
x ^ y	Logical XOR	Bool
~X	Bitwise NOT	T(w(x) bits)
x & y	Bitwise AND	T(max(w(xy) bits)
x   y	Bitwise OR	T(max(w(xy) bits)
x ^ y	Bitwise XOR	T(max(w(xy) bits)
x.xorR	XOR all bits of x	Bool
x.orR	OR all bits of x	Bool
x.andR	AND all bits of x	Bool
x >> y	Arithmetic shift right, y: Int	T(w(x) - y  bits)
x >> y	Arithmetic shift right, y: UInt	T(w(x) bits)
x << y	Arithmetic shift left, y: Int	T(w(x) + y  bits)
x << y	Arithmetic shift left, y: UInt	T(w(x) + max(y) bits)
x  >> y	Logical shift right, y: Int/UInt	T(w(x) bits)
x  << y	Logical shift left, y: Int/UInt	T(w(x) bits)
x.rotateLeft(y)	Logical left rotation, y: UInt/Int	T(w(x) bits)
x.rotateRight(y)	Logical right rotation, y: UInt/Int	T(w(x) bits)
x.clearAll[()]	Clear all bits	
x.setAll[()]	Set all bits	
x.setAllTo(value :	Set all bits to the given Boolean value	
Boolean)		
x.setAllTo(value : Bool)	Set all bits to the given Bool value	

```
// Bitwise operator
val a, b, c = SInt(32 bits)
c := ~(a & b) // Inverse(a AND b)

val all_1 = a.andR // Check that all bits are equal to 1

// Logical shift
val uint_10bits = uint_8bits << 2 // shift left (resulting in 10 bits)
val shift_8bits = uint_8bits |<< 2 // shift left (resulting in 8 bits)

// Logical rotation
val myBits = uint_8bits.rotateLeft(3) // left bit rotation

// Set/clear
val a = B"8'x42"
when(cond) {
   a.setAll() // set all bits to True when cond is True
}</pre>
```

5.3. UInt/SInt 41

#### **Arithmetic**

Operator	Description	Return
x + y	Addition	$T(\max(w(x), w(y)), \text{ bits})$
x +^ y	Addition with carry	$T(\max(w(x), w(y) + 1), \text{ bits})$
x +  y	Addition by sat carry bit	$T(\max(w(x), w(y)), \text{ bits})$
x - y	Subtraction	$T(\max(w(x), w(y)), \text{ bits})$
x -^ y	Subtraction with carry	$T(\max(w(x), w(y) + 1), \text{ bits})$
x -  y	Subtraction by sat carry bit	$T(\max(w(x), w(y)), \text{ bits})$
x * y	Multiplication	T(w(x) + w(y)), bits)
x / y	Division	T(w(x), bits)
x % y	Modulo	T(w(x), bits)

```
// Addition
val res = mySInt_1 + mySInt_2
```

## Comparison

Operator	Description	Return type
x === y	Equality	Bool
x =/= y	Inequality	Bool
x > y	Greater than	Bool
x >= y	Greater than or equal	Bool
x < y	Less than	Bool
x <= y	Less than or equal	Bool

```
// Comparison between two SInts
myBool := mySInt_1 > mySInt_2

// Comparison between a UInt and a literal
myBool := myUInt_8bits >= U(3, 8 bits)

when(myUInt_8bits === 3) {
    ..
}
```

## Type cast

Operator	Description	Return
x.asBits	Binary cast to Bits	Bits(w(x), bits)
x.asUInt	Binary cast to UInt	UInt(w(x), bits)
x.asSInt	Binary cast to SInt	SInt(w(x), bits)
x.asBools	Cast into a array of Bool	Vec(Bool, w(x))
S(x: T)	Cast a Data into a SInt	SInt(w(x), bits)
U(x: T)	Cast a Data into an UInt	UInt(w(x), bits)
x.intoSInt	convert to SInt expand signbit	SInt(w(x) + 1, bits)

To cast a Bool, a Bits, or an SInt into a UInt, you can use U(something). To cast things into an SInt, you can use S(something).

```
// Cast an SInt to Bits
val myBits = mySInt.asBits

// Create a Vector of Bool
val myVec = myUInt.asBools

// Cast a Bits to SInt
val mySInt = S(myBits)
```

#### Bit extraction

Operator	Description	Return
x(y)	Readbit, y: Int/UInt	Bool
x(offset, width)	Read bitfield, offset: UInt, width: Int	T(width bits)
x(range)	Read a range of bits. Ex: myBits(4 downto 2)	T(range bits)
x(y) := z	Assign bits, y: Int/UInt	Bool
x(offset, width) := z	Assign bitfield, offset: UInt, width: Int	T(width bits)
x(range) := z	Assign a range of bit. Ex: myBits(4 downto 2) := U"010"	T(range bits)

```
// get the bit at index 4
val myBool = myUInt(4)

// assign bit 1 to True
mySInt(1) := True

// Range
val myUInt_8bits = myUInt_16bits(7 downto 0)
val myUInt_7bits = myUInt_16bits(0 to 6)
val myUInt_6bits = myUInt_16Bits(0 until 6)

mySInt_8bits(3 downto 0) := mySInt_4bits
```

5.3. UInt/Sint 43

#### Misc

Operator	Description	Return
x.getWidth	Return bitcount	Int
x.msb	Return the most significant bit	Bool
x.lsb	Return the least significant bit	Bool
x.range	Return the range (x.high downto 0)	Range
x.high	Return the upper bound of the type x	Int
x ## y	Concatenate, x->high, y->low	Bits(w(x)
		+ w(y)
		bits)
x @ @ y	Concatenate x:T with y:Bool/SInt/UInt	T(w(x) +
•		w(y) bits)
x.subdivideIn(y slices)	Subdivide x into y slices, y: Int	Vec(T, y)
x.subdivideIn(y bits)	Subdivide x into multiple slices of y bits, y: Int	Vec(T,
•		w(x)/y)
x.resize(y)		T(y bits)
	Determine a socied converted if coloured it is filled with some	
	Return a resized copy of x, if enlarged, it is filled with zero	
	for UInt or filled with the sign for SInt, y: Int	
x.resized		T(w(x)
		bits)
	Return a version of x which is allowed to be automatically	
	resized where needed	
mvUInt.twoComplement(	enUse the two's complement to transform an UInt into an SInt	SInt(w(myUIn
Bool)	che se die two s component to dansform air can time air sant	+ 1, bits)
mySInt.abs	Return the absolute value of the UInt value	UInt(w(mySIn
my ome.aos	retain the assorate value of the offic value	bits)
mySInt.abs(en: Bool)	Return the absolute value of the UInt value when en is True	UInt(w(mySIn
<b>3</b>		bits)
mySInt.sign	Return most significant bit	Bool
x.expand	Return x with 1 bit expand	T(w(x)+1
•		
		bit)
mySInt.absWithSym	Return the absolute value of the UInt value with symmetric, shrink	bit) UInt(w(mySInt

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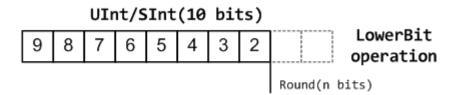
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## 5.3.4 FixPoint operations

For fixpoint, we can divide it into two parts:

- Lower bit operations (rounding methods)
- High bit operations (saturation operations)

## Lower bit operations



About Rounding: https://en.wikipedia.org/wiki/Rounding

SpinalHDL-	Wikipedia-	API	Mathematic Algo-	re-	Sup-
Name	Name		rithm	turn(align=false)	ported
FLOOR	RoundDown	floor	floor(x)	w(x)-n bits	Yes
FLOOR-	RoundToZero	floor-	sign*floor(abs(x))	w(x)-n bits	Yes
TOZERO		ToZero			
CEIL	RoundUp	ceil	ceil(x)	w(x)-n+1 bits	Yes
CEILTOINF	RoundToInf	ceilToInf	sign*ceil(abs(x))	w(x)-n+1 bits	Yes
ROUNDUP	RoundHalfUp	roundUp	floor(x+0.5)	w(x)-n+1 bits	Yes
ROUNDDOWN	RoundHalf-	roundDown	ceil(x-0.5)	w(x)-n+1 bits	Yes
	Down				
ROUND-	Round-	round-	sign*ceil(abs(x)-0.5)	w(x)-n+1 bits	Yes
TOZERO	HalfToZero	ToZero			
ROUNDTOINF	Round-	roundToInf	sign*floor(abs(x)+0.5)	w(x)-n+1 bits	Yes
	HalfToInf				
ROUNDTO-	RoundHalfTo-	roundTo-			No
EVEN	Even	Even			
ROUND-	Round-	round-			No
TOODD	HalfToOdd	ToOdd			

**Note:** The **RoundToEven** and **RoundToOdd** modes are very special, and are used in some big data statistical fields with high accuracy concerns, SpinalHDL doesn't support them yet.

5.3. UInt/SInt 45

You will find *ROUNDUP*, *ROUNDDOWN*, *ROUNDTOZERO*, *ROUNDTOINF*, *ROUNDTOEVEN*, *ROUNTOODD* are very close in behavior, *ROUNDTOINF* is the most common. The behavior of rounding in different programming languages may be different.

Programming lan-	default-	Example	comments
guage	RoundType		
Matlab	ROUNDTOINF	round(1.5)=2,round(2.5)=3;round(-1.5)=-	round to
		2,round(-2.5)=-3	$\pm Infinity$
python2	ROUNDTOINF	round(1.5)=2,round(2.5)=3;round(-1.5)=-	round to
		2,round(-2.5)=-3	$\pm Infinity$
python3	ROUNDTO-	round(1.5)=round(2.5)=2; round(-1.5)=round(-	close to Even
	EVEN	2.5)=-2	
Scala.math	ROUNDTOUP	round(1.5)=2,round(2.5)=3;round(-1.5)=-	always to +In-
		1,round(-2.5)=-2	finity
SpinalHDL	ROUNDTOINF	round(1.5)=2,round(2.5)=3;round(-1.5)=-	round to
		2,round(-2.5)=-3	$\pm Infinity$

**Note:** In SpinalHDL *ROUNDTOINF* is the default RoundType (round = roundToInf)

```
val A = SInt(16 bit)
val B = A.roundToInf(6 bits) // default 'align = false' with carry, got 11 bit
val B = A.roundToInf(6 bits, align = true) // sat 1 carry bit, got 10 bit
val B = A.floor(6 bits) // return 10 bit
val B = A.floorToZero(6 bits)
                                 // return 10 bit
val B = A.ceil(6 bits)
                                  // ceil with carry so return 11 bit
val B = A.ceil(6 bits, align = true) // ceil with carry then sat 1 bit return 10 bit
val B = A.ceilToInf(6 bits)
val B = A.roundUp(6 bits)
val B = A.roundDown(6 bits)
val B = A.roundToInf(6 bits)
val B = A.roundToZero(6 bits)
val B = A.round(6 bits)
                                  // SpinalHDL uses roundToInf as the default.
→rounding mode
val B0 = A.roundToInf(6 bits, align = true)
                                                     /--> equal
val B1 = A.roundToInf(6 bits, align = false).sat(1) // ---+
```

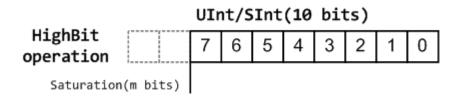
**Note:** Only floor and floorToZero work without the align option; they do not need a carry bit. Other rounding operations default to using a carry bit.

## round Api

API	UInt/SInt	description	Return(align=false)	Return(align=true)
floor	Both		w(x)-n bits	w(x)-n bits
floorToZero	SInt	equal to floor in UInt	w(x)-n bits	w(x)-n bits
ceil	Both		w(x)-n+1 bits	w(x)-n bits
ceilToInf	SInt	equal to ceil in UInt	w(x)-n+1 bits	w(x)-n bits
roundUp	Both	simple for HW	w(x)-n+1 bits	w(x)-n bits
roundDown	Both		w(x)-n+1 bits	w(x)-n bits
roundToInf	SInt	most Common	w(x)-n+1 bits	w(x)-n bits
roundToZero	SInt	equal to roundDown in UInt	w(x)-n+1 bits	w(x)-n bits
round	Both	SpinalHDL chose roundToInf	w(x)-n+1 bits	w(x)-n bits

**Note:** Although roundToInf is very common, roundUp has the least cost and good timing, with almost no performance loss. As a result, roundUp is strongly recommended for production use.

## **High bit operations**



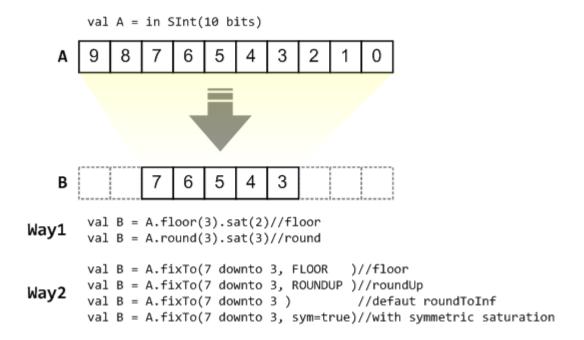
function	Opera-	Positive-Op	Negative-Op
	tion		
sat	Saturation	when(Top[w-1, w-n].orR) set max-	When(Top[w-1, w-n].andR) set min-
		Value	Value
trim	Discard	N/A	N/A
symme-	Symmet-	N/A	minValue = -maxValue
try	ric		

Symmetric is only valid for SInt.

#### fixTo function

Two ways are provided in UInt/SInt to do fixpoint:

5.3. UInt/SInt 47



fixTo is strongly recommended in your RTL work, you don't need to handle carry bit alignment and bit width calculations manually like **Way1** in the above diagram.

Factory Fix function with Auto Saturation:

Function	Description	Return
fixTo(section, roundType, symmetric)	Factory FixFunction	section.size bits

```
val A = SInt(16 bit)
val B = A.fixTo(10 downto 3) // default RoundType.ROUNDTOINF, sym = false
val B = A.fixTo( 8 downto 0, RoundType.ROUNDUP)
val B = A.fixTo( 9 downto 3, RoundType.CEIL, sym = false)
val B = A.fixTo(16 downto 1, RoundType.ROUNDTOINF, sym = true )
val B = A.fixTo(10 downto 3, RoundType.FLOOR) // floor 3 bit, sat 5 bit @ highest
val B = A.fixTo(20 downto 3, RoundType.FLOOR) // floor 3 bit, expand 2 bit @ highest
```

# 5.4 SpinalEnum

## 5.4.1 Description

The Enumeration type corresponds to a list of named values.

## 5.4.2 Declaration

The declaration of an enumerated data type is as follows:

```
object Enumeration extends SpinalEnum {
  val element0, element1, ..., elementN = newElement()
}
```

For the example above, the default encoding is used. The native enumeration type is used for VHDL and a binary encoding is used for Verilog.

The enumeration encoding can be forced by defining the enumeration as follows:

```
object Enumeration extends SpinalEnum(defaultEncoding=encodingOfYourChoice) {
  val element0, element1, ..., elementN = newElement()
}
```

**Note:** If you want to define an enumeration as in/out for a given component, you have to do as following: in(MyEnum()) or out(MyEnum())

## **Encoding**

The following enumeration encodings are supported:

Encod-	Bit	Description
ing	width	
native		Use the VHDL enumeration system, this is the default encoding
binary-	log2Up(st	ate Scothits to store states in declaration order (value from 0 to n-1)
Sequen-		
tial		
binary-	state-	Use Bits to store state. Each bit corresponds to one state
OneHot	Count	

Custom encodings can be performed in two different ways: static or dynamic.

```
* Static encoding
object MyEnumStatic extends SpinalEnum {
  val e0, e1, e2, e3 = newElement()
  defaultEncoding = SpinalEnumEncoding("staticEncoding")(
    e0 -> 0,
    e1 \rightarrow 2
    e2 \rightarrow 3,
    e3 -> 7)
}
 * Dynamic encoding with the function : \_ * 2 + 1
     e.g. : e0 \Rightarrow 0 * 2 + 1 = 1
             e1 \Rightarrow 1 * 2 + 1 = 3
 *
             e2 \implies 2 * 2 + 1 = 5
 *
             e3 \implies 3 * 2 + 1 = 7
val encoding = SpinalEnumEncoding("dynamicEncoding", _ * 2 + 1)
```

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5.4. SpinalEnum 49

(continued from previous page)

```
object MyEnumDynamic extends SpinalEnum(encoding) {
  val e0, e1, e2, e3 = newElement()
}
```

#### **Example**

Instantiate an enumerated signal and assign a value to it:

```
object UartCtrlTxState extends SpinalEnum {
  val sIdle, sStart, sData, sParity, sStop = newElement()
}

val stateNext = UartCtrlTxState()
  stateNext := UartCtrlTxState.sIdle

// You can also import the enumeration to have visibility of its elements
import UartCtrlTxState._
  stateNext := sIdle
```

## 5.4.3 Operators

The following operators are available for the Enumeration type:

## Comparison

Operator	Description	Return type
x === y	Equality	Bool
x =/= y	Inequality	Bool

```
import UartCtrlTxState._
val stateNext = UartCtrlTxState()
stateNext := sIdle
when(stateNext === sStart) {
    ...
}
switch(stateNext) {
    is(sIdle) {
     ...
}
    is(sStart) {
    ...
}
...
}
```

#### Type cast

Operator	Description	Return
x.asBits	Binary cast to Bits	Bits(w(x) bits)
x.asUInt	Binary cast to UInt	UInt(w(x) bits)
x.asSInt	Binary cast to SInt	SInt(w(x) bits)

```
import UartCtrlTxState._
val stateNext = UartCtrlTxState()
myBits := sIdle.asBits
```

## 5.5 Bundle

## 5.5.1 Description

The Bundle is a composite type that defines a group of named signals (of any SpinalHDL basic type) under a single name.

A Bundle can be used to model data structures, buses, and interfaces.

## 5.5.2 Declaration

The syntax to declare a bundle is as follows:

```
case class myBundle extends Bundle {
  val bundleItem0 = AnyType
  val bundleItem1 = AnyType
  val bundleItemN = AnyType
}
```

For example, a bundle holding a color could be defined as:

```
case class Color(channelWidth: Int) extends Bundle {
  val r, g, b = UInt(channelWidth bits)
}
```

You can find an APB3 definition among the Spinal HDL examples.

## 5.5.3 Operators

The following operators are available for the Bundle type:

5.5. Bundle 51

## Comparison

Operator	Description	Return type
x === y	Equality	Bool
x =/= y	Inequality	Bool

```
val color1 = Color(8)
color1.r := 0
color1.g := 0
color1.b := 0

val color2 = Color(8)
color2.r := 0
color2.g := 0
color2.b := 0
myBool := color1 === color2
```

## Type cast

Operator	Description	Return	
x.asBits	Binary cast to Bits	Bits(w(x) bits)	

```
val color1 = Color(8)
val myBits := color1.asBits
```

#### **Convert Bits back to Bundle**

The .assignFromBits operator can be viewed as the reverse of .asBits.

Operator	Description	
x.assignFromBits(y)	Convert Bits (y) to Bundle(x)	Unit
x.assignFromBits(y, hi, lo)	Convert Bits (y) to Bundle(x) with high/low boundary	Unit

The following example saves a Bundle called CommonDataBus into a circular buffer (3rd party memory), reads the Bits out later and converts them back to CommonDataBus format.

# CommonDataBus Input

# CommonDataBus Output

.assignFromBits

```
case class TestBundle () extends Component {
 val io = new Bundle {
   val we
            = in
                        Bool()
   val addrWr = in
                      UInt (7 bits)
   val dataIn = slave (CommonDataBus())
   val addrRd = in
                      UInt (7 bits)
   val dataOut = master (CommonDataBus())
 }
 val mm = Ram3rdParty_1w_1rs (G_DATA_WIDTH = io.dataIn.getBitsWidth,
                             G_ADDR_WIDTH = io.addrWr.getBitsWidth,
                                         = "Intel_Arria10_M20K")
                             G_VENDOR
 mm.io.clk_in := clockDomain.readClockWire
               := clockDomain.readClockWire
 mm.io.clk_out
 mm.io.we
                := io.we
 mm.io.addr_wr := io.addrWr.asBits
 mm.io.d
                := io.dataIn.asBits
 mm.io.addr_rd := io.addrRd.asBits
 io.dataOut.assignFromBits(mm.io.q)
```

5.5. Bundle 53

#### 5.5.4 IO Element direction

When you define a Bundle inside the IO definition of your component, you need to specify its direction.

#### in/out

If all elements of your bundle go in the same direction you can use in(MyBundle()) or out(MyBundle()).

For example:

```
val io = new Bundle {
  val input = in (Color(8))
  val output = out(Color(8))
}
```

#### master/slave

If your interface obeys to a master/slave topology, you can use the IMasterSlave trait. Then you have to implement the function def asMaster(): Unit to set the direction of each element from the master's perspective. Then you can use the master(MyBundle()) and slave(MyBundle()) syntax in the IO definition.

For example:

```
case class HandShake(payloadWidth: Int) extends Bundle with IMasterSlave {
 val valid = Bool()
 val ready = Bool()
 val payload = Bits(payloadWidth bits)
 // You have to implement this asMaster function.
 // This function should set the direction of each signals from an master point of
\hookrightarrow view
 override def asMaster(): Unit = {
    out(valid, payload)
    in(ready)
 }
}
val io = new Bundle {
 val input = slave(HandShake(8))
 val output = master(HandShake(8))
}
```

## 5.6 Vec

## 5.6.1 Description

A Vec is a composite type that defines a group of indexed signals (of any SpinalHDL basic type) under a single name.

## 5.6.2 Declaration

The syntax to declare a vector is as follows:

Declaration	Description		
Vec(type: Data, size: Int)	Create a vector capable of holding size elements of type Data		
Vec(x, y,)			
	Create a vector where indexes point to the provided elements.  This constructor supports mixed element width.		

#### **Examples**

```
// Create a vector of 2 signed integers
val myVecOfSInt = Vec(SInt(8 bits), 2)
myVecOfSInt(0) := 2
myVecOfSInt(1) := myVecOfSInt(0) + 3
// Create a vector of 3 different type elements
val myVecOfMixedUInt = Vec(UInt(3 bits), UInt(5 bits), UInt(8 bits))
val x, y, z = UInt(8 bits)
val myVecOf_xyz_ref = Vec(x, y, z)
// Iterate on a vector
for(element <- myVecOf_xyz_ref) {</pre>
  element := 0 // Assign x, y, z with the value 0
}
// Map on vector
myVecOfMixedUInt.map(\_ := 0) // Assign all elements with value 0
// Assign 3 to the first element of the vector
myVecOf_xyz_ref(1) := 3
```

## 5.6.3 Operators

The following operators are available for the Vec type:

## Comparison

Operator	Description	Return type
x === y	Equality	Bool
x =/= y	Inequality	Bool

```
// Create a vector of 2 signed integers
val vec2 = Vec(SInt(8 bits), 2)
val vec1 = Vec(SInt(8 bits), 2)
myBool := vec2 === vec1 // Compare all elements
```

5.6. Vec 55

#### Type cast

Operator	Description	Return	
x.asBits	Binary cast to Bits	Bits(w(x) bits)	

```
// Create a vector of 2 signed integers
val vec1 = Vec(SInt(8 bits), 2)

myBits_16bits := vec1.asBits
```

#### Misc

Operator	Description	Return
x.getBitsWidth	Return the full size of the Vec	Int

```
// Create a vector of 2 signed integers
val vec1 = Vec(SInt(8 bits), 2)
println(vec1.getBitsWidth) // 16
```

**Warning:** SpinalHDL fixed-point support is only partially used/tested, if you find any bugs with it, or you think that some functionality is missing, please create a Github issue. Also, please do not use undocumented features in your code.

## 5.7 UFix/SFix

# 5.7.1 Description

The UFix and SFix types correspond to a vector of bits that can be used for fixed-point arithmetic.

## 5.7.2 Declaration

The syntax to declare a fixed-point number is as follows:

## **Unsigned Fixed-Point**

Syntax		bit	resolution	max	min
		width			
UFix(peak: Exp	Number, reso-	peak-	2^resolution	2^peak-2^resolution	0
lution: ExpNun	nber)	resolution			
UFix(peak:	ExpNumber,	width	2^(peak-width)	2^peak-2^(peak-width)	0
width: BitCoun	t)				

#### **Signed Fixed-Point**

Syntax		bit	resolution	max	min
		width			
SFix(peak: Exp	pNumber, reso-	peak-	2^resolution	2^peak-2^resolution	-
lution: ExpNun	nber)	resolution	+1		(2^peak)
SFix(peak:	ExpNumber,	width	2^(peak-width-1)	2^peak-2^(peak-width-1)	-
width: BitCoun	nt)				(2^peak)

#### **Format**

The chosen format follows the usual way of defining fixed-point number format using Q notation. More information can be found on the Wikipedia page about the Q number format.

For example Q8.2 will mean a fixed-point number of 8+2 bits, where 8 bits are used for the natural part and 2 bits for the fractional part. If the fixed-point number is signed, one more bit is used for the sign.

The resolution is defined as being the smallest power of two that can be represented in this number.

**Note:** To make representing power-of-two numbers less error prone, there is a numeric type in spinal.core called ExpNumber, which is used for the fixed-point type constructors. A convenience wrapper exists for this type, in the form of the exp function (used in the code samples on this page).

#### **Examples**

## 5.7.3 Assignments

#### **Valid Assignments**

An assignment to a fixed-point value is valid when there is no bit loss. Any bit loss will result in an error.

If the source fixed-point value is too big, the .truncated function will allow you to resize the source number to match the destination size.

5.7. UFix/SFix 57

#### **Example**

#### From a Scala constant

Scala BigInt or Double types can be used as constants when assigning to UFix or SFix signals.

#### **Example**

## 5.7.4 Raw value

The integer representation of the fixed-point number can be read or written by using the raw property.

#### **Example**

## 5.7.5 Operators

The following operators are available for the UFix type:

## **Arithmetic**

Ор-	Description	Returned resolu-	Returned amplitude
er-		tion	
a-			
tor			
X	Addition	Min(x.resolution,	Max(x.amplitude, y.amplitude)
+		y.resolution)	
У			
X -	Subtraction	Min(x.resolution,	Max(x.amplitude, y.amplitude)
У		y.resolution)	
X	Multiplication	x.resolution *	x.amplitude * y.amplitude
*		y.resolution)	
У			
X	Arithmetic shift right, y: Int	x.amplitude >> y	x.resolution >> y
>>			
У			
X	Arithmetic shift left, y: Int	x.amplitude << y	x.resolution << y
<<			
У			
X	Arithmetic shift right, y: Int	x.amplitude >> y	x.resolution
>>			
у			
X	Arithmetic shift left, y: Int	x.amplitude << y	x.resolution
<<			
у			

# Comparison

Operator	Description	Return type
x === y	Equality	Bool
x =/= y	Inequality	Bool
x > y	Greater than	Bool
x >= y	Greater than or equal	Bool
x > y	Less than	Bool
x >= y	Less than or equal	Bool

# Type cast

Operator	Description	Return
x.asBits	Binary cast to Bits	Bits(w(x) bits)
x.asUInt	Binary cast to UInt	UInt(w(x) bits)
x.asSInt	Binary cast to SInt	SInt(w(x) bits)
x.asBools	Cast into a array of Bool	Vec(Bool,width(x))
x.toUInt	Return the corresponding UInt (with truncation)	UInt
x.toSInt	Return the corresponding SInt (with truncation)	SInt
x.toUFix	Return the corresponding UFix	UFix
x.toSFix	Return the corresponding SFix	SFix

5.7. UFix/SFix 59

#### **Misc**

Name	Return	Description
x.maxValue	Return the maximum value storable	Double
x.minValue	Return the minimum value storable	Double
x.resolution	x.amplitude * y.amplitude	Double

**Warning:** SpinalHDL floating-point support is under development and only partially used/tested, if you have any bugs with it, or you think that some functionality is missing, please create a Github issue. Also, please do not use undocumented features in your code.

# 5.8 Floating

## 5.8.1 Description

The Floating type corresponds to IEEE-754 encoded numbers. A second type called RecFloating helps in simplifying your design by recoding the floating-point value simplify some edge cases in IEEE-754 floating-point.

It's composed of a sign bit, an exponent field and a mantissa field. The widths of the different fields are defined in the IEEE-754 or de-facto standards.

This type can be used with the following import:

import spinal.lib.experimental.math.\_

### **IEEE-754 floating format**

The numbers are encoded into IEEE-754 floating-point format.

#### **Recoded floating format**

Since IEEE-754 has some quirks about denormalized numbers and special values, Berkeley proposed another way of recoding floating-point values.

The mantissa is modified so that denormalized values can be treated the same as the normalized ones.

The exponent field is one bit larger that one of the IEEE-754 number.

The sign bit is kept unchanged between the two encodings.

Examples can be found here

## Zero

The zero is encoded with the three leading zeros of the exponent field being set to zero.

#### **Denormalized values**

Denormalized values are encoded in the same way as a normal floating-point number. The mantissa is shifted so that the first one becomes implicit. The exponent is encoded as 107 (decimal) plus the index of the highest bit set to 1.

#### **Normalized values**

The recoded mantissa for normalized values is exactly the same as the original IEEE-754 mantissa. The recoded exponent is encoded as 130 (decimal) plus the original exponent value.

## Infinity

The recoded mantissa value is treated as don't care. The recoded exponent three highest bits is 6 (decimal), the rest of the exponent can be treated as don't care.

#### NaN

The recoded mantissa for normalized values is exactly the same as the original IEEE-754 mantissa. The recoded exponent three highest bits is 7 (decimal), the rest of the exponent can be treated as don't care.

#### 5.8.2 Declaration

The syntax to declare a floating-point number is as follows:

#### **IEEE-754 Number**

Syntax	Description
Floating(exponentSize: Int, mantissaS-	IEEE-754 Floating-point value with a custom exponent and man-
ize: Int)	tissa size
Floating16()	IEEE-754 Half precision floating-point number
Floating32()	IEEE-754 Single precision floating-point number
Floating64()	IEEE-754 Double precision floating-point number
Floating128()	IEEE-754 Quad precision floating-point number

## **Recoded floating-point number**

Syntax	Description
RecFloating(exponentSize: Int, mantissaS-	Recoded Floating-point value with a custom exponent and man-
ize: Int)	tissa size
RecFloating16()	Recoded Half precision floating-point number
RecFloating32()	Recoded Single precision floating-point number
RecFloating64()	Recoded Double precision floating-point number
RecFloating128()	Recoded Quad precision floating-point number

5.8. Floating 61

## 5.8.3 Operators

The following operators are available for the Floating and RecFloating types:

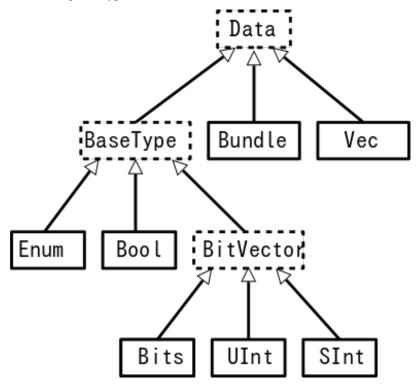
#### Type cast

Operator	Description	Return
x.asBits	Binary cast to Bits	Bits(w(x)
		bits)
x.asBools	Cast into a array of Bool	Vec(Bool,width(x))
x.toUInt(size: Int)	Return the corresponding UInt (with truncation)	UInt
x.toSInt(size: Int)	Return the corresponding SInt (with truncation)	SInt
x.fromUInt	Return the corresponding Floating (with truncation)	UInt
x.fromSInt	Return the corresponding Floating (with truncation)	SInt

## 5.9 Introduction

The language provides 5 base types, and 2 composite types that can be used.

- Base types: Bool, Bits, UInt for unsigned integers, SInt for signed integers and Enum.
- Composite types: Bundle and Vec.



In addition to the base types, Spinal has support under development for:

- Fixed-point numbers (partial support)
- *Floating-point* numbers (experimental support)

Finally, a special type is available for checking equality between a BitVector and a bits constant that contains holes (don't care values). An example is shown below:

```
val myBits = Bits(8 bits)
val itMatch = myBits === M"00--10--" // - don't care value
```

5.9. Introduction 63

## **STRUCTURING**

# 6.1 Component and hierarchy

## 6.1.1 Introduction

Like in VHDL and Verilog, you can define components that can be used to build a design hierarchy. However, in SpinalHDL, you don't need to bind their ports at instantiation:

```
class AdderCell extends Component {
 // Declaring external ports in a Bundle called `io` is recommended
 val io = new Bundle {
   val a, b, cin = in Bool()
   val sum, cout = out Bool()
 // Do some logic
 io.sum := io.a ^ io.b ^ io.cin
 io.cout := (io.a & io.b) | (io.a & io.cin) | (io.b & io.cin)
class Adder(width: Int) extends Component {
 // Create 2 AdderCell instances
 val cell0 = new AdderCell
 val cell1 = new AdderCell
 cell1.io.cin := cell0.io.cout // Connect cout of cell0 to cin of cell1
 // Another example which creates an array of ArrayCell instances
 val cellArray = Array.fill(width)(new AdderCell)
 cellArray(1).io.cin := cellArray(0).io.cout // Connect cout of cell(0) to cin of_
\hookrightarrow cell(1)
```

#### Tip:

```
val io = new Bundle { ... }
```

Declaring external ports in a Bundle called io is recommended. If you name your bundle io, SpinalHDL will check that all of its elements are defined as inputs or outputs.

## 6.1.2 Input / output definition

The syntax to define inputs and outputs is as follows:

Syntax	Description	Return
in Bool()/out Bool()	Create an input Bool/output Bool	Bool
in/out Bits/UInt/SInt[(x bit)]	Create an input/output of the corresponding type	Bits/UInt/SInt
in/out(T)	For all other data types, you may have to add some	T
	brackets around it. Sorry, this is a Scala limitation.	
master/slave(T)	This syntax is provided by the spinal.lib library (If	T
	you annotate your object with the slave syntax, then	
	import spinal.lib.slave instead). T should ex-	
	tend IMasterSlave – Some documentation is avail-	
	able <i>here</i> . You may not actually need the brackets, so	
	master T is fine as well.	

There are some rules to follow with component interconnection:

- Components can only **read** output and input signals of child components.
- Components can read their own output port values (unlike in VHDL).

**Tip:** If for some reason you need to read signals from far away in the hierarchy (such as for debugging or temporal patches), you can do it by using the value returned by some.where.else.theSignal.pull()

## 6.1.3 Pruned signals

SpinalHDL only generates things which are directly or indirectly required to drive the outputs of your top-level entity.

All other signals (the useless ones) are removed from the RTL generation and are inserted into a list of pruned signals. You can get this list via the printPruned and the printPrunedIo functions on the generated SpinalReport object:

```
class TopLevel extends Component {
 val io = new Bundle {
   val a,b = in UInt(8 bits)
   val result = out UInt(8 bits)
 }
 io.result := io.a + io.b
 val unusedSignal = UInt(8 bits)
 val unusedSignal2 = UInt(8 bits)
 unusedSignal2 := unusedSignal
object Main {
 def main(args: Array[String]) {
   SpinalVhdl(new TopLevel).printPruned()
   //This will report :
   // [Warning] Unused wire detected : toplevel/unusedSignal : UInt[8 bits]
       [Warning] Unused wire detected : toplevel/unusedSignal2 : UInt[8 bits]
 }
```

If you want to keep a pruned signal in the generated RTL for debugging reasons, you can use the keep function of that signal:

```
class TopLevel extends Component {
 val io = new Bundle {
   val a, b = in UInt(8 bits)
   val result = out UInt(8 bits)
 }
 io.result := io.a + io.b
 val unusedSignal = UInt(8 bits)
 val unusedSignal2 = UInt(8 bits).keep()
 unusedSignal := 0
 unusedSignal2 := unusedSignal
object Main {
 def main(args: Array[String]) {
   SpinalVhdl(new TopLevel).printPruned()
    // This will report nothing
 }
}
```

# 6.1.4 Parametrized Hardware ("Generic" in VHDL, "Parameter" in Verilog)

If you want to parameterize your component, you can give parameters to the constructor of the component as follows:

```
class MyAdder(width: BitCount) extends Component {
  val io = new Bundle {
    val a, b = in UInt(width)
    val result = out UInt(width)
}
  io.result := io.a + io.b
}

object Main {
  def main(args: Array[String]) {
    SpinalVhdl(new MyAdder(32 bits))
  }
}
```

If you have several parameters, it is a good practice to give a specific configuration class as follows:

### 6.1.5 Synthesized component names

Within a module, each component has a name, called a "partial name". The "full" name is built by joining every component's parent name with "\_", for example: io\_clockDomain\_reset. You can use setName to replace this convention with a custom name. This is especially useful when interfacing with external components. The other methods are called getName, setPartialName, and getPartialName respectively.

When synthesized, each module gets the name of the Scala class defining it. You can override this as well with setDefinitionName.

### 6.2 Area

#### 6.2.1 Introduction

Sometimes, creating a Component to define some logic is overkill because you:

- Need to define all construction parameters and IO (verbosity, duplication)
- Split your code (more than needed)

For this kind of case you can use an Area to define a group of signals/logic:

```
class UartCtrl extends Component {
 val timer = new Area {
   val counter = Reg(UInt(8 bit))
   val tick = counter === 0
    counter := counter - 1
   when(tick) {
      counter := 100
    }
 }
 val tickCounter = new Area {
   val value = Reg(UInt(3 bit))
   val reset = False
   when(timer.tick) {
                                // Refer to the tick from timer area
      value := value + 1
   when(reset) {
      value := 0
 val stateMachine = new Area {
 }
}
```

#### Tip:

In VHDL and Verilog, sometimes prefixes are used to separate variables into logical sections. It is suggested that you use Area instead of this in SpinalHDL.

**Note:** ClockingArea is a special kind of Area that allows you to define chunks of hardware which use a given

### 6.3 Function

#### 6.3.1 Introduction

The ways you can use Scala functions to generate hardware are radically different than VHDL/Verilog for many reasons:

- You can instantiate registers, combinational logic, and components inside them.
- You don't have to play with process/@always blocks that limit the scope of assignment of signals.
- Everything is passed by reference, which allows easy manipulation.

  For example, you can give a bus to a function as an argument, then the function can internally read/write to it. You can also return a Component, a Bus, or anything else from Scala and the Scala world.

### 6.3.2 RGB to gray

For example, if you want to convert a Red/Green/Blue color into greyscale by using coefficients, you can use functions to apply them:

```
// Input RGB color
val r, g, b = UInt(8 bits)

// Define a function to multiply a UInt by a Scala Float value.
def coef(value: UInt, by: Float): UInt = (value * U((255 * by).toInt, 8 bits) >> 8)

// Calculate the gray level
val gray = coef(r, 0.3f) + coef(g, 0.4f) + coef(b, 0.3f)
```

# 6.3.3 Valid Ready Payload bus

For instance, if you define a simple bus with valid, ready, and payload signals, you can then define some useful functions inside of it.

```
case class MyBus(payloadWidth: Int) extends Bundle with IMasterSlave {
  val valid = Bool()
  val ready = Bool()
  val payload = Bits(payloadWidth bits)

// Define the direction of the data in a master mode
  override def asMaster(): Unit = {
    out(valid, payload)
    in(ready)
  }

// Connect that to this
  def <<(that: MyBus): Unit = {
    this.valid := that.valid
    that.ready := this.ready
    this.payload := that.payload
  }
}</pre>
```

(continues on next page)

6.3. Function 69

```
// Connect this to the FIFO input, return the fifo output
def queue(size: Int): MyBus = {
    val fifo = new MyBusFifo(payloadWidth, size)
    fifo.io.push << this
    return fifo.io.pop
}
}
class MyBusFifo(payloadWidth: Int, depth: Int) extends Component {
    val io = new Bundle {
        val push = slave(MyBus(payloadWidth))
        val pop = master(MyBus(payloadWidth))
    }

    val mem = Mem(Bits(payloadWidth bits), depth)

// ...
}</pre>
```

# 6.4 Clock domains

### 6.4.1 Introduction

In SpinalHDL, clock and reset signals can be combined to create a **clock domain**. Clock domains can be applied to some areas of the design and then all synchronous elements instantiated into those areas will then **implicitly** use this clock domain.

Clock domain application works like a stack, which means that if you are in a given clock domain you can still apply another clock domain locally.

#### 6.4.2 Instantiation

The syntax to define a clock domain is as follows (using EBNF syntax):

```
ClockDomain(
  clock: Bool
  [,reset: Bool]
  [,softReset: Bool]
  [,clockEnable: Bool]
  [,frequency: IClockDomainFrequency]
  [,config: ClockDomainConfig]
)
```

This definition takes five parameters:

Argu-	Description	De-
ment		fault
clock	Clock signal that defines the domain	
reset	Reset signal. If a register exists which needs a reset and the clock domain doesn't provide	null
	one, an error message will be displayed	
softRe	sæteset which infers an additional synchronous reset	null
clockE	nable goal of this signal is to disable the clock on the whole clock domain without having to	null
	manually implement that on each synchronous element	
freque	n <b>Ay</b> llows you to specify the frequency of the given clock domain and later read it in your design	Un-
		known-
		Fre-
		quency
config	Specify the polarity of signals and the nature of the reset	Cur-
		rent
		config

An applied example to define a specific clock domain within the design is as follows:

```
val coreClock = Bool()
val coreReset = Bool()

// Define a new clock domain
val coreClockDomain = ClockDomain(coreClock, coreReset)

// Use this domain in an area of the design
val coreArea = new ClockingArea(coreClockDomain) {
   val coreClockedRegister = Reg(UInt(4 bit))
}
```

### Configuration

In addition to *constructor parameters*, the following elements of each clock domain are configurable via a ClockDomainConfigclass:

Property	Valid values	
clockEdge	RISING, FALLING	
resetKind	ASYNC, SYNC, and BOOT which is supported by some FPGAs (where FF values are loaded by	
	the bitstream)	
resetActiveLevenIGH, LOW		
softResetActi	v <b>ěli@i</b> i;LOW	
clockEnableAc	t <del>indi</del> ķvēw	

```
class CustomClockExample extends Component {
  val io = new Bundle {
    val clk = in Bool()
    val resetn = in Bool()
    val result = out UInt (4 bits)
  }

// Configure the clock domain
val myClockDomain = ClockDomain(
    clock = io.clk,
    reset = io.resetn,
    config = ClockDomainConfig(
```

(continues on next page)

6.4. Clock domains 71

```
clockEdge = RISING,
  resetKind = ASYNC,
  resetActiveLevel = LOW
)

// Define an Area which use myClockDomain
val myArea = new ClockingArea(myClockDomain) {
  val myReg = Reg(UInt(4 bits)) init(7)

  myReg := myReg + 1

  io.result := myReg
}
```

By default, a ClockDomain is applied to the whole design. The configuration of this default domain is:

- Clock: rising edge
- Reset: asynchronous, active high
- · No clock enable

This corresponds to the following ClockDomainConfig:

#### Internal clock

An alternative syntax to create a clock domain is the following:

```
ClockDomain.internal(
  name: String,
  [config: ClockDomainConfig,]
  [withReset: Boolean,]
  [withSoftReset: Boolean,]
  [withClockEnable: Boolean,]
  [frequency: IClockDomainFrequency]
)
```

This definition takes six parameters:

Argument	Description Default		
name	Name of <i>clk</i> and <i>reset</i> signal		
config	Specify polarity of signals and the nature of the reset	Current con-	
		fig	
withReset	Add a reset signal true		
withSoftResetAdd a soft reset signal false			
withClockEnabAded a clock enable		false	
frequency	Frequency of the clock domain Un		
		Frequency	

The advantage of this approach is to create clock and reset signals with a known/specified name instead of an inherited one.

Once created, you have to assign the ClockDomain's signals, as shown in the example below:

```
class InternalClockWithPllExample extends Component {
 val io = new Bundle {
   val clk100M = in Bool()
   val aReset = in Bool()
   val result = out UInt (4 bits)
 }
  // myClockDomain.clock will be named myClockName_clk
 // myClockDomain.reset will be named myClockName_reset
 val myClockDomain = ClockDomain.internal("myClockName")
 // Instantiate a PLL (probably a BlackBox)
 val pll = new Pll()
 pll.io.clkIn := io.clk100M
 // Assign myClockDomain signals with something
 myClockDomain.clock := pll.io.clockOut
 myClockDomain.reset := io.aReset || !pll.io.
 // Do whatever you want with myClockDomain
 val myArea = new ClockingArea(myClockDomain) {
   val myReg = Reg(UInt(4 bits)) init(7)
   myReg := myReg + 1
    io.result := myReg
 }
}
```

### **External clock**

You can define a clock domain which is driven by the outside anywhere in your source. It will then automatically add clock and reset wires from the top level inputs to all synchronous elements.

```
ClockDomain.external(
  name: String,
  [config: ClockDomainConfig,]
  [withReset: Boolean,]
  [withSoftReset: Boolean,]
  [withClockEnable: Boolean,]
  [frequency: IClockDomainFrequency]
)
```

The arguments to the ClockDomain.external function are exactly the same as in the ClockDomain.internal function. Below is an example of a design using ClockDomain.external:

```
class ExternalClockExample extends Component {
  val io = new Bundle {
    val result = out UInt (4 bits)
  }

  // On the top level you have two signals :
    // myClockName_clk and myClockName_reset
  val myClockDomain = ClockDomain.external("myClockName")
```

(continues on next page)

6.4. Clock domains 73

```
val myArea = new ClockingArea(myClockDomain) {
   val myReg = Reg(UInt(4 bits)) init(7)
   myReg := myReg + 1

   io.result := myReg
}
```

### Context

You can retrieve in which clock domain you are by calling ClockDomain.current anywhere.

The returned ClockDomain instance has the following functions that can be called:

name	Description	Return
fre-	Return the frequency of the clock domain	Double
quency.getValu	ae	
hasReset	Return if the clock domain has a reset signal	Boolean
hasSoftReset	Return if the clock domain has a soft reset signal	Boolean
hasClockEn-	Return if the clock domain has a clock enable signal	Boolean
able		
readClock-	Return a signal derived from the clock signal	Bool
Wire		
read-	Return a signal derived from the soft reset signal	Bool
ResetWire		
readSoftRe-	Return a signal derived from the reset signal	Bool
setWire		
readClock-	Return a signal derived from the clock enable signal	Bool
EnableWire		
isResetAc-	Return True when the reset is active	Bool
tive		
isSoftRese-	Return True when the soft reset is active	Bool
tActive		
isClockEn-	Return True when the clock enable is active	Bool
ableActive		

An example is included below where a UART controller uses the frequency specification to set its clock divider:

### 6.4.3 Clock domain crossing

SpinalHDL checks at compile time that there are no unwanted/unspecified cross clock domain signal reads. If you want to read a signal that is emitted by another ClockDomain area, you should add the crossClockDomain tag to the destination signal as depicted in the following example:

```
//
// clkA
// rstA
// Implementation where clock and reset pins are given by components' IO
class CrossingExample extends Component {
 val io = new Bundle {
   val clkA = in Bool()
   val rstA = in Bool()
   val clkB = in Bool()
   val rstB = in Bool()
   val dataIn = in Bool()
   val dataOut = out Bool()
 // sample dataIn with clkA
 val area_clkA = new ClockingArea(ClockDomain(io.clkA,io.rstA)) {
   val reg = RegNext(io.dataIn) init(False)
 }
 // 2 register stages to avoid metastability issues
 val area_clkB = new ClockingArea(ClockDomain(io.clkB,io.rstB)) {
   val buf0 = RegNext(area_clkA.reg) init(False) addTag(crossClockDomain)
                                  init(False)
   val buf1 = RegNext(buf0)
 }
 io.dataOut := area_clkB.buf1
}
// Alternative implementation where clock domains are given as parameters
class CrossingExample(clkA : ClockDomain,clkB : ClockDomain) extends Component {
 val io = new Bundle {
   val dataIn = in Bool()
   val dataOut = out Bool()
 }
 // sample dataIn with clkA
 val area_clkA = new ClockingArea(clkA) {
   val reg = RegNext(io.dataIn) init(False)
 // 2 register stages to avoid metastability issues
 val area_clkB = new ClockingArea(clkB) {
```

(continues on next page)

6.4. Clock domains 75

```
val buf0 = RegNext(area_clkA.reg) init(False) addTag(crossClockDomain)
val buf1 = RegNext(buf0) init(False)
}
io.dataOut := area_clkB.buf1
}
```

In general, you can use 2 or more flip-flop driven by the destination clock domain to prevent metastability. The BufferCC(input: T, init: T = null, bufferDepth: Int = 2) function provided in spinal.lib. \_ will instantiate the necessary flip-flops (the number of flip-flops will depends on the bufferDepth parameter) to mitigate the phenomena.

```
class CrossingExample(clkA : ClockDomain,clkB : ClockDomain) extends Component {
  val io = new Bundle {
    val dataIn = in Bool()
    val dataOut = out Bool()
}

// sample dataIn with clkA
val area_clkA = new ClockingArea(clkA) {
    val reg = RegNext(io.dataIn) init(False)
}

// BufferCC to avoid metastability issues
val area_clkB = new ClockingArea(clkB) {
    val buf1 = BufferCC(area_clkA.reg, False)
}

io.dataOut := area_clkB.buf1
}
```

**Warning:** The BufferCC function is only for signals of type Bit, or Bits operating as Gray-coded counters (only 1 bit-flip per clock cycle), and can not used for multi-bit cross-domain processes.

### 6.4.4 Special clocking Areas

#### **Slow Area**

A SlowArea is used to create a new clock domain area which is slower than the current one:

```
class TopLevel extends Component {

   // Use the current clock domain : 100MHz
   val areaStd = new Area {
     val counter = out(CounterFreeRun(16).value)
   }

   // Slow the current clockDomain by 4 : 25 MHz
   val areaDiv4 = new SlowArea(4) {
     val counter = out(CounterFreeRun(16).value)
   }

   // Slow the current clockDomain to 50MHz
```

```
val area50Mhz = new SlowArea(50 MHz) {
   val counter = out(CounterFreeRun(16).value)
}

def main(args: Array[String]) {
   new SpinalConfig(
     defaultClockDomainFrequency = FixedFrequency(100 MHz)
   ).generateVhdl(new TopLevel)
}
```

#### **ResetArea**

A ResetArea is used to create a new clock domain area where a special reset signal is combined with the current clock domain reset:

### **ClockEnableArea**

A ClockEnableArea is used to add an additional clock enable in the current clock domain:

```
class TopLevel extends Component {
  val clockEnable = Bool()

  // Add a clock enable for this area
  val area_1 = new ClockEnableArea(clockEnable) {
    val counter = out(CounterFreeRun(16).value)
  }
}
```

6.4. Clock domains 77

# 6.5 Instantiate VHDL and Verilog IP

# 6.5.1 Description

A blackbox allows the user to integrate an existing VHDL/Verilog component into the design by just specifying its interfaces. It's up to the simulator or synthesizer to do the elaboration correctly.

# 6.5.2 Defining an blackbox

An example of how to define a blackbox is shown below:

```
// Define a Ram as a BlackBox
class Ram_1w_1r(wordWidth: Int, wordCount: Int) extends BlackBox {
 // Add VHDL Generics / Verilog parameters to the blackbox
 // You can use String, Int, Double, Boolean, and all SpinalHDL base
 // types as generic values
 addGeneric("wordCount", wordCount)
 addGeneric("wordWidth", wordWidth)
 // Define IO of the VHDL entity / Verilog module
 val io = new Bundle {
   val clk = in Bool()
   val wr = new Bundle {
     val en = in Bool()
     val addr = in UInt (log2Up(wordCount) bit)
     val data = in Bits (wordWidth bit)
   val rd = new Bundle {
     val en = in Bool()
     val addr = in UInt (log2Up(wordCount) bit)
     val data = out Bits (wordWidth bit)
   }
 }
 // Map the current clock domain to the io.clk pin
 mapClockDomain(clock=io.clk)
```

In VHDL, signals of type Bool will be translated into std\_logic and Bits into std\_logic\_vector. If you want to get std\_ulogic, you have to use a BlackBoxULogic instead of BlackBox.

In Verilog, BlackBoxUlogic has no effect.

```
class Ram_1w_1r(wordWidth: Int, wordCount: Int) extends BlackBoxULogic {
    ...
}
```

### 6.5.3 Generics

There are two different ways to declare generics:

```
class Ram(wordWidth: Int, wordCount: Int) extends BlackBox {
   addGeneric("wordCount", wordCount)
   addGeneric("wordWidth", wordWidth)

// OR

val generic = new Generic {
   val wordCount = Ram.this.wordCount
   val wordWidth = Ram.this.wordWidth
  }
}
```

# 6.5.4 Instantiating a blackbox

Instantiating a BlackBox is just like instantiating a Component:

```
// Create the top level and instantiate the Ram
class TopLevel extends Component {
 val io = new Bundle {
   val wr = new Bundle {
     val en = in Bool()
     val addr = in UInt (log2Up(16) bit)
     val data = in Bits (8 bit)
   }
   val rd = new Bundle {
     val en = in Bool()
     val addr = in UInt (log2Up(16) bit)
     val data = out Bits (8 bit)
   }
 }
 // Instantiate the blackbox
 val ram = new Ram_1w_1r(8,16)
 // Connect all the signals
 io.wr.en <> ram.io.wr.en
 io.wr.addr <> ram.io.wr.addr
 io.wr.data <> ram.io.wr.data
 io.rd.en <> ram.io.rd.en
 io.rd.addr <> ram.io.rd.addr
 io.rd.data <> ram.io.rd.data
object Main {
 def main(args: Array[String]): Unit = {
   SpinalVhdl(new TopLevel)
 }
}
```

### 6.5.5 Clock and reset mapping

In your blackbox definition you have to explicitly define clock and reset wires. To map signals of a ClockDomain to corresponding inputs of the blackbox you can use the mapClockDomain or mapCurrentClockDomain function. mapClockDomain has the following parameters:

name	type	default	description
clockDo-	ClockDo-	ClockDo-	Specify the clockDomain which provides the signals
main	main	main.current	
clock	Bool	Nothing	Blackbox input which should be connected to the clockDomain clock
reset	Bool	Nothing	Blackbox input which should be connected to the clockDomain reset
enable	Bool	Nothing	Blackbox input which should be connected to the clockDomain en-
			able

mapCurrentClockDomain has almost the same parameters as mapClockDomain but without the clockDomain.

For example:

```
class MyRam(clkDomain: ClockDomain) extends BlackBox {
  val io = new Bundle {
    val clkA = in Bool()
    ...
  val clkB = in Bool()
    ...
}

// Clock A is map on a specific clock Domain
mapClockDomain(clkDomain, io.clkA)
// Clock B is map on the current clock domain
mapCurrentClockDomain(io.clkB)
}
```

# 6.5.6 io prefix

In order to avoid the prefix "io\_" on each of the IOs of the blackbox, you can use the function noIoPrefix() as shown below:

```
// Define the Ram as a BlackBox
class Ram_1w_1r(wordWidth: Int, wordCount: Int) extends BlackBox {
  val generic = new Generic {
    val wordCount = Ram_1w_1r.this.wordCount
    val wordWidth = Ram_1w_1r.this.wordWidth
  }

  val io = new Bundle {
    val clk = in Bool()

    val wr = new Bundle {
      val en = in Bool()
      val addr = in UInt (log2Up(_wordCount) bit)
      val data = in Bits (_wordWidth bit)
    }
  val rd = new Bundle {
      val en = in Bool()
    }
```

```
val addr = in UInt (log2Up(_wordCount) bit)
  val data = out Bits (_wordWidth bit)
}
noIoPrefix()
mapCurrentClockDomain(clock=io.clk)
}
```

#### 6.5.7 Rename all io of a blackbox

IOs of a BlackBox or Component can be renamed at compile-time using the addPrePopTask function. This function takes a no-argument function to be applied during compilation, and is useful for adding renaming passes, as shown in the following example:

```
class MyRam() extends Blackbox {
 val io = new Bundle {
   val clk = in Bool()
   val portA = new Bundle{
      val cs = in Bool()
      val rwn = in Bool()
      val dIn = in Bits(32 bits)
     val dOut = out Bits(32 bits)
   val portB = new Bundle{
      val cs = in Bool()
      val rwn = in Bool()
      val dIn = in Bits(32 bits)
      val dOut = out Bits(32 bits)
   }
 }
 // Map the clk
 mapCurrentClockDomain(io.clk)
  // Remove io_ prefix
 noIoPrefix()
 // Function used to rename all signals of the blackbox
 private def renameIO(): Unit = {
    io.flatten.foreach(bt => {
      if(bt.getName().contains("portA")) bt.setName(bt.getName().repalce("portA_", "
→") + "_A")
      if(bt.getName().contains("portB")) bt.setName(bt.getName().repalce("portB_", "
\hookrightarrow") + "_B")
   })
 }
 // Execute the function renameIO after the creation of the component
 addPrePopTask(() => renameIO())
// This code generate these names:
```

```
// clk
// cs_A, rwn_A, dIn_A, dOut_A
// cs_B, rwn_B, dIn_B, dOut_B
```

### 6.5.8 Add RTL source

With the function addRTLPath() you can associate your RTL sources with the blackbox. After the generation of your SpinalHDL code you can call the function mergeRTLSource to merge all of the sources together.

```
class MyBlackBox() extends Blackbox {
 val io = new Bundle {
   val clk = in Bool()
   val start = in Bool()
   val dIn = in Bits(32 bits)
   val dOut = out Bits(32 bits)
   val ready = out Bool()
 // Map the clk
 mapCurrentClockDomain(io.clk)
 // Remove io_ prefix
 noIoPrefix()
 // Add all rtl dependencies
 addRTLPath("./rtl/RegisterBank.v")
                                                            // Add a verilog file
                                                            // Add a vhdl file
 addRTLPath(s"./rtl/myDesign.vhd")
 addRTLPath(s"${sys.env("MY_PROJECT")}/myTopLevel.vhd")
                                                            // Use an environement_
→variable MY_PROJECT (System.getenv("MY_PROJECT"))
}
val report = SpinalVhdl(new MyBlackBox)
report.mergeRTLSource("mergeRTL") // Merge all rtl sources into mergeRTL.vhd and.
→mergeRTL.v files
```

### 6.5.9 VHDL - No numeric type

If you want to use only std\_logic\_vector in your blackbox component, you can add the tag noNumericType to the blackbox.

```
class MyBlackBox() extends BlackBox{
  val io = new Bundle {
    val clk = in Bool()
    val increment = in Bool()
    val initValue = in UInt(8 bits)
    val counter = out UInt(8 bits)
}

mapCurrentClockDomain(io.clk)
noIoPrefix()
```

```
addTag(noNumericType) // Only std_logic_vector
}
```

The code above will generate the following VHDL:

```
component MyBlackBox is
  port(
    clk : in std_logic;
    increment : in std_logic;
    initValue : in std_logic_vector(7 downto 0);
    counter : out std_logic_vector(7 downto 0)
);
end component;
```

# 6.6 Preserving names

### 6.6.1 Introduction

This page will describe how SpinalHDL propagate names from the scala code to the generated hardware. Knowing them should enable you to preserve those names as much as possible to generate understandable netlists.

### 6.6.2 Nameable base class

All the things which can be named in SpinalHDL extends the Nameable base class which.

So in practice, the following classes extends Nameable:

- Component
- Area
- Data (UInt, SInt, Bundle, ...)

There is a few example of that Nameable API

Will generation:

In general, you don't realy need to access that API, unless you want to do tricky stuff for debug reasons or for elaboration purposes.

#### 6.6.3 Name extraction from Scala

First, since version 1.4.0, SpinalHDL use a scala compiler plugin which can provide a call back each time a new val is defined during the construction of an class.

There is a example showing more or less how SpinalHDL itself is implemented:

```
//spinal.idslplugin.ValCallback is the Scala compiler plugin feature which will.
→provide the callbacks
class Component extends spinal.idslplugin.ValCallback{
 override def valCallback[T](ref: T, name: String) : T = {
   println(s"Got $ref named $name") // Here we just print what we got as a demo.
   ref
 }
}
class UInt
class Bits
class MyComponent extends Component{
 val two = 2
 val wuff = "miaou"
 val toto = new UInt
 val rawrr = new Bits
object Debug3 extends App{
 new MyComponent()
 // ^ This will print :
 // Got 2 named two
 // Got miaou named wuff
 // Got spinal.tester.code.sandbox.UInt@691a7f8f named toto
 // Got spinal.tester.code.sandbox.Bits@161b062a named rawrr
}
```

Using that ValCallback "introspection" feature, SpinalHDL's Component classes are able to be aware of their content and its name.

But this also mean that if you want something to get a name, and you only rely on this automatic naming feature, the reference to your Data (UInt, SInt, ...) instances should be stored somewhere in a Component val.

For instance:

```
class MyComponent extends Component {
  val a,b = in UInt(8 bits) // Will be properly named
  val toto = out UInt(8 bits) // same

  def doStuff(): Unit = {
    val tmp = UInt(8 bits) // This will not be named, as it isn't stored anywhere in_
    a component val (but there is a solution explained later)
    tmp := 0x20
    toto := tmp
  }
  doStuff()
}
```

Will generate:

```
module MyComponent (
input [7:0] a,

(continues on next page)
```

```
input [7:0] b,
output [7:0] toto
);
//Note that the tmp signal defined in scala was "shortcuted" by SpinalHDL, as it_
was unamed and technicaly "shortcutable"
assign toto = 8'h20;
endmodule
```

# 6.6.4 Area in a Component

One important aspect in the naming system is that you can define new namespaces inside components and manipulate

For instance via Area:

Will generate

### 6.6.5 Area in a function

You can also define function which will create new Area which will provide a namespace for all its content :

```
class MyComponent extends Component {
  def isZero(value: UInt) = new Area {
    val comparator = value === 0
  }

  val value = in UInt (8 bits)
  val someLogic = isZero(value)

  val result = out Bool()
  result := someLogic.comparator
}
```

Which will generate:

```
module MyComponent (
  input [7:0] value,
  output result
```

# 6.6.6 Composite in a function

Added in SpinalHDL 1.5.0, Composite which allow you to create a scope which will use as prefix another Nameable:

```
class MyComponent extends Component {
    //Basicaly, a Composite is an Area that use its construction parameter as namespace...
    prefix
    def isZero(value: UInt) = new Composite(value) {
        val comparator = value === 0
        }.comparator //Note we don't return the Composite, but the element of the...
        composite that we are interested in

    val value = in UInt (8 bits)
    val result = out Bool()
    result := isZero(value)
}
```

Will generate:

```
module MyComponent (
  input [7:0] value,
  output result
);
  wire value_comparator;

  assign value_comparator = (value == 8'h0);
  assign result = value_comparator;
endmodule
```

### 6.6.7 Composite chains

You can also chain composites:

```
class MyComponent extends Component {
  def isZero(value: UInt) = new Composite(value) {
    val comparator = value === 0
  }.comparator

def inverted(value: Bool) = new Composite(value) {
    val inverter = !value
  }.inverter

val value = in UInt(8 bits)
```

```
val result = out Bool()
  result := inverted(isZero(value))
}
```

Will generate:

```
module MyComponent (
  input [7:0] value,
  output result
);
  wire value_comparator;
  wire value_comparator_inverter;

  assign value_comparator = (value == 8'h0);
  assign value_comparator_inverter = (! value_comparator);
  assign result = value_comparator_inverter;

endmodule
```

### 6.6.8 Composite in a Bundle's function

This behaviour can be very useful when implementing Bundles utilities. For instance in the spinal.lib.Stream class is defined the following:

```
class Stream[T <: Data](val payloadType : HardType[T]) extends Bundle {</pre>
 val valid = Bool()
 val ready
            = Bool()
 val payload = payloadType()
 def queue(size: Int): Stream[T] = new Composite(this){
   val fifo = new StreamFifo(payloadType, size)
    fifo.io.push << self</pre>
                          // 'self' refer to the Composite construction argument_
→(this in that example). It avoid having to do a boring 'Stream.this'
 }.fifo.io.pop
 def m2sPipe(): Stream[T] = new Composite(this) {
   val m2sPipe = Stream(payloadType)
   val rValid = RegInit(False)
   val rData = Reg(payloadType)
   self.ready := (!m2sPipe.valid) || m2sPipe.ready
   when(self.ready) {
      rValid := self.valid
      rData := self.payload
   }
   m2sPipe.valid := rValid
   m2sPipe.payload := rData
  }.m2sPipe
}
```

Which allow nested calls while preserving the names:

```
class MyComponent extends Component {
  val source = slave(Stream(UInt(8 bits)))
  val sink = master(Stream(UInt(8 bits)))
  sink << source.queue(size = 16).m2sPipe()
}</pre>
```

Will generate

```
module MyComponent (
  input
                      source_valid,
 output
                      source_ready,
 input
             [7:0]
                      source_payload,
 output
                      sink_valid,
 input
                      sink_ready,
 output
             [7:0]
                      sink_payload,
 input
                      clk.
 input
                      reset
 wire
                      source_fifo_io_pop_ready;
 wire
                      source_fifo_io_push_ready;
 wire
                      source_fifo_io_pop_valid;
 wire
             [7:0]
                      source_fifo_io_pop_payload;
 wire
             [4:0]
                      source_fifo_io_occupancy;
 wire
             [4:0]
                      source_fifo_io_availability;
 wire
                      source_fifo_io_pop_m2sPipe_valid;
 wire
                      source_fifo_io_pop_m2sPipe_ready;
 wire
             [7:0]
                      source_fifo_io_pop_m2sPipe_payload;
 reg
                      source_fifo_io_pop_rValid;
             [7:0]
 reg
                      source_fifo_io_pop_rData;
 StreamFifo source_fifo (
    .io_push_valid
                         (source_valid
                                                       ), //i
    .io_push_ready
                         (source_fifo_io_push_ready
                                                       ), //o
    .io_push_payload
                         (source_payload
                                                       ), //i
    .io_pop_valid
                         (source_fifo_io_pop_valid
                                                       ), //o
    .io_pop_ready
                         (source_fifo_io_pop_ready
                                                       ), //i
    .io_pop_payload
                         (source_fifo_io_pop_payload
                                                       ), //o
    .io_flush
                         (1'b0
                                                       ), //i
                                                       ), //o
    .io_occupancy
                         (source_fifo_io_occupancy
                                                       ), //o
    .io_availability
                         (source_fifo_io_availability
    .clk
                         (clk
                                                       ), //i
    .reset
                        (reset
                                                       ) //i
 );
 assign source_ready = source_fifo_io_push_ready;
 assign source_fifo_io_pop_ready = ((1'b1 && (! source_fifo_io_pop_m2sPipe_valid))_
→ | | source_fifo_io_pop_m2sPipe_ready);
 assign source_fifo_io_pop_m2sPipe_valid = source_fifo_io_pop_rValid;
 assign source_fifo_io_pop_m2sPipe_payload = source_fifo_io_pop_rData;
 assign sink_valid = source_fifo_io_pop_m2sPipe_valid;
 assign source_fifo_io_pop_m2sPipe_ready = sink_ready;
 assign sink_payload = source_fifo_io_pop_m2sPipe_payload;
 always @ (posedge clk or posedge reset) begin
    if (reset) begin
      source_fifo_io_pop_rValid <= 1'b0;</pre>
    end else begin
      if(source_fifo_io_pop_ready)begin
        source_fifo_io_pop_rValid <= source_fifo_io_pop_valid;</pre>
```

```
end
end
end
always @ (posedge clk) begin
if(source_fifo_io_pop_ready)begin
    source_fifo_io_pop_rData <= source_fifo_io_pop_payload;
    end
end
end
endmodule</pre>
```

# 6.6.9 Unamed signal handling

Since 1.5.0, for signal which end up without name, SpinalHDL will find a signal which is driven by that unamed signal and propagate its name. This can produce useful results as long you don't have too large island of unamed stuff.

The name attributed to such unamed signal is : \_zz\_ + drivenSignal.getName()

Note that this naming pattern is also used by the generation backend when they need to breakup some specific expressions or long chain of expression into multiple signals.

### Verilog expression splitting

There is an instance of expressions (ex: the + operator) that SpinalHDL need to express in dedicated signals to match the behaviour with the Scala API:

```
class MyComponent extends Component {
  val a,b,c,d = in UInt(8 bits)
  val result = a + b + c + d
}
```

Will generate

```
module MyComponent (
 input
             [7:0]
                       a,
 input
             [7:0]
                       b,
 input
             [7:0]
                       С,
 input
             [7:0]
);
                      _zz_result;
 wire
             [7:0]
 wire
             [7:0]
                      _zz_result_1;
 wire
             [7:0]
                      result;
 assign _zz_result = (_zz_result_1 + c);
 assign _zz_result_1 = (a + b);
 assign result = (_zz_result + d);
endmodule
```

#### Verilog long expression splitting

There is a instance of how a very long expression chain will be splited up by SpinalHDL:

Will generate

```
module MyComponent (
 input
                   conditions_0,
 input
                   conditions_1,
 input
                   conditions_2,
 input
                   conditions_3,
 . . .
                   conditions_58,
 input
                   conditions_59,
 input
                   conditions_60,
 input
 input
                   conditions_61,
 input
                   conditions_62,
                   conditions_63
 input
);
                   _zz_result;
 wire
 wire
                   _zz_result_1;
 wire
                   _zz_result_2;
 wire
                   result;
 →33) || conditions_34) || conditions_35) || conditions_36) || conditions_37) || ⊔
→conditions_38) || conditions_39) || conditions_40) || conditions_41) || conditions_
→42) || conditions_43) || conditions_44) || conditions_45) || conditions_46) || □
_{\hookrightarrow}17) || conditions_18) || conditions_19) || conditions_20) || conditions_21) || _{\square}
→conditions_22) || conditions_23) || conditions_24) || conditions_25) || conditions_
→26) || conditions_27) || conditions_28) || conditions_29) || conditions_30) ||
assign _zz_result_2 = ((((((((((((((conditions_0 | | conditions_1) | | conditions_2)_
\rightarrow || conditions_3) || conditions_4) || conditions_5) || conditions_6) || conditions_
→7) || conditions_8) || conditions_9) || conditions_10) || conditions_11) || □
→conditions_12) || conditions_13) || conditions_14) || conditions_15);
 →conditions_50) || conditions_51) || conditions_52) || conditions_53) || conditions_
→54) || conditions_55) || conditions_56) || conditions_57) || conditions_58) || □
→conditions_59) || conditions_60) || conditions_61) || conditions_62) || conditions_
\hookrightarrow63);
endmodule
```

#### When statement condition

The *when(cond) { }* statements condition are generated into separated signals named *when\_* + fileName + line. A similar thing will also be done for switch statements.

```
//In file Test.scala
class MyComponent extends Component {
  val value = in UInt(8 bits)
  val isZero = out(Bool())
  val counter = out(Reg(UInt(8 bits)))

  isZero := False
  when(value === 0){ //At line 117
    isZero := True
    counter := counter + 1
  }
}
```

Will generate

```
module MyComponent (
 input
                      value,
             [7:0]
 output reg
                       isZero,
 output reg [7:0]
                       counter,
 input
                       clk,
 input
                       reset
);
 wire
                       when_Test_l117;
 always @ (*) begin
    isZero = 1'b0;
    if(when_Test_1117)begin
      isZero = 1'b1;
    end
 end
 assign when_Test_l117 = (value == 8'h0);
 always @ (posedge clk) begin
    if(when_Test_l117)begin
      counter <= (counter + 8'h01);</pre>
    end
 end
endmodule
```

#### In last resort

In last resort, if a signal has no name (anonymous signal), SpinalHDL will seek for a named signal which is driven by the anonymous signal, and use it as a name postfix :

```
ret := ret + 1
}
return ret
}
value := count(enable)
}
```

Will generate

```
module MyComponent (
                       enable,
  input
  output
              [7:0]
                       value,
  input
                       clk,
  input
                       reset
);
              [7:0]
                       _zz_value; //Name given to the register in last resort by_
 reg
→looking what was driven by it
  assign value = _zz_value;
  always @ (posedge clk) begin
    if(enable)begin
      _{zz\_value} \leftarrow (_{zz\_value} + 8'h01);
    end
  end
endmodule
```

This last resort naming skim isn't ideal in all cases, but can help out.

Note that signal starting with a underscore aren't stored in the Verilator waves (on purpose)

**CHAPTER** 

# **SEVEN**

# **SEMANTIC**

# 7.1 Assignments

# 7.1.1 Assignments

There are multiple assignment operators:

Symbol	Description	
:=	Standard assignment, equivalent to <= in VHDL/Verilog. The last assignment to a variable	
	wins; the value is not updated until the next simulation delta cycle.	
\=	Equivalent to := in VHDL and = in Verilog. The value is updated instantly in-place.	
<>	Automatic connection between 2 signals or two bundles of the same type. Direction is in-	
	ferred by using signal direction (in/out). (Similar behavior to :=)	

```
// Because of hardware concurrency, `a` is always read as '1' by b and c
val a, b, c = UInt(4 bits)
a := 0
b := a
a := 1 // a := 1 "wins"
c := a
var x = UInt(4 bits)
val y, z = UInt(4 bits)
x := 0
y := x
            // y read x with the value 0
x = x + 1
            // z read x with the value 1
z := x
// Automatic connection between two UART interfaces.
uartCtrl.io.uart <> io.uart
```

It is important to understand that in SpinalHDL, the nature of a signal (combinational/sequential) is defined in its declaration, not by the way it is assigned. All datatype instances will define a combinational signal, while a datatype instance wrapped with Reg(...) will define a sequential (registered) signal.

# 7.1.2 Width checking

SpinalHDL checks that the bit count of the left side and the right side of an assignment matches. There are multiple ways to adapt the width of a given BitVector (Bits, UInt, SInt):

Resizing techniques	Description
x := y.resized Assign x with a resized copy of y, resize value is automatically infer	
	match x
x := y.resize(newWidth)	Assign x with a resized copy of y, size is manually calculated

There is one case where Spinal automatically resizes a value:

Assignment	Problem	SpinalHDL action
$myUIntOf_8bit := U(3)$	U(3) creates an UInt of 2 bits, which	Because U(3) is a "weak" bit count in-
	doesn't match the left side (8 bits)	ferred signal, SpinalHDL resizes it au-
		tomatically

# 7.1.3 Combinatorial loops

SpinalHDL checks that there are no combinatorial loops (latches) in your design. If one is detected, it raises an error and SpinalHDL will print the path of the loop.

# 7.2 When/Switch/Mux

#### 7.2.1 When

As in VHDL and Verilog, signals can be conditionally assigned when a specified condition is met:

```
when(cond1) {
    // Execute when cond1 is true
}.elsewhen(cond2) {
    // Execute when (not cond1) and cond2
}.otherwise {
    // Execute when (not cond1) and (not cond2)
}
```

### 7.2.2 Switch

As in VHDL and Verilog, signals can be conditionally assigned when a signal has a defined value:

```
switch(x) {
  is(value1) {
    // Execute when x === value1
  }
  is(value2) {
    // Execute when x === value2
  }
  default {
    // Execute if none of precedent conditions met
  }
}
```

### 7.2.3 Local declaration

It is possible to define new signals inside a when/switch statement:

```
val x, y = UInt(4 bits)
val a, b = UInt(4 bits)

when(cond) {
    val tmp = a + b
    x := tmp
    y := tmp + 1
} otherwise {
    x := 0
    y := 0
}
```

**Note:** SpinalHDL checks that signals defined inside a scope are only assigned inside that scope.

### 7.2.4 Mux

If you just need a Mux with a Bool selection signal, there are two equivalent syntaxes:

Syntax	Return	Description	
Mux(cond, whenTrue, whenFalse)  T Return whenTrue when		Return whenTrue when cond is True,	
		whenFalse otherwise	
cond ? whenTrue   whenFalse	Т	Return whenTrue when cond is True,	
		whenFalse otherwise	

```
val cond = Bool
val whenTrue, whenFalse = UInt(8 bits)
val muxOutput = Mux(cond, whenTrue, whenFalse)
val muxOutput2 = cond ? whenTrue | whenFalse
```

### 7.2.5 Bitwise selection

A bitwise selection looks like the VHDL when syntax.

### **Example**

```
val bitwiseSelect = UInt(2 bits)
val bitwiseResult = bitwiseSelect.mux(
    0 -> (io.src0 & io.src1),
    1 -> (io.src0 | io.src1),
    2 -> (io.src0 ^ io.src1),
    default -> (io.src0)
)
```

Also, if all possible values are covered in your mux, you can omit the default value:

```
val bitwiseSelect = UInt(2 bits)
val bitwiseResult = bitwiseSelect.mux(
0 -> (io.src0 & io.src1),
```

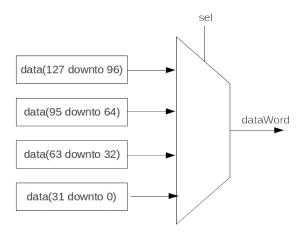
```
1 -> (io.src0 | io.src1),

2 -> (io.src0 ^ io.src1),

3 -> (io.src0)

)
```

muxLists(...) is another bitwise selection which takes a sequence of tuples as input. Below is an example of dividing a Bits of 128 bits into 32 bits:



# 7.3 Rules

### 7.3.1 Introduction

The semantics behind SpinalHDL are important to learn, so that you understand what is really happening behind the scenes, and how to control it.

These semantics are defined by multiple rules:

- Signals and registers are operating concurrently with each other (parallel behavioral, as in VHDL and Verilog)
- An assignment to a combinational signal is like expressing a rule which is always true
- An assignment to a register is like expressing a rule which is applied on each cycle of its clock domain
- For each signal, the last valid assignment wins
- Each signal and register can be manipulated as an object during hardware elaboration in a OOP manner

### 7.3.2 Concurrency

The order in which you assign each combinational or registered signal has no behavioral impact.

For example, both of the following pieces of code are equivalent:

```
val a, b, c = UInt(8 bits) // Define 3 combinational signals
c := a + b // c will be set to 7
b := 2 // b will be set to 2
a := b + 3 // a will be set to 5
```

This is equivalent to:

More generally, when you use the := assignment operator, it's like specifying a new rule for the left side signal/register.

### 7.3.3 Last valid assignment wins

If a combinational signal or register is assigned multiple times, the last valid one wins.

As an example:

This will produce the following truth table:

Χ	У	=>	result
False	False		1
False	True		1
True	False		2
True	True		3

# 7.3.4 Signal and register interactions with Scala (OOP reference + Functions)

In SpinalHDL, each hardware element is modeled by a class instance. This means you can manipulate instances by using their references, such as passing them as arguments to a function.

As an example, the following code implements a register which is incremented when inc is True and cleared when clear is True (clear has priority over inc):

7.3. Rules 97

```
when(inc) {
  counter := counter + 1
}
when(clear) {
  counter := 0  // If inc and clear are True, then this assignment wins (Lasture valid assignment rule)
}
```

You can implement exactly the same functionality by mixing the previous example with a function that assigns to counter:

```
val inc, clear = Bool()
val counter = Reg(UInt(8 bits))

def setCounter(value : UInt): Unit = {
    counter := value
}

when(inc) {
    setCounter(counter + 1) // Set counter with counter + 1
}
when(clear) {
    counter := 0
}
```

You can also integrate the conditional check inside the function:

```
val inc, clear = Bool()
val counter = Reg(UInt(8 bits))

def setCounterWhen(cond : Bool, value : UInt): Unit = {
    when(cond) {
        counter := value
    }
}

setCounterWhen(cond = inc, value = counter + 1)
setCounterWhen(cond = clear, value = 0)
```

And also specify what should be assigned to the function:

```
val inc, clear = Bool()
val counter = Reg(UInt(8 bits))

def setSomethingWhen(something : UInt, cond : Bool, value : UInt): Unit = {
   when(cond) {
      something := value
   }
}

setSomethingWhen(something = counter, cond = inc, value = counter + 1)
setSomethingWhen(something = counter, cond = clear, value = 0)
```

All of the previous examples are strictly equivalent both in their generated RTL and also in the SpinalHDL compiler's perspective. This is because SpinalHDL only cares about the Scala runtime and the objects instantiated there, it doesn't care about the Scala syntax itself.

In other words, from a generated RTL generation / SpinalHDL perspective, when you use functions in Scala which

generate hardware, it is like the function was inlined. This is also true case for Scala loops, as they will appear in unrolled form in the generated RTL.

7.3. Rules 99

# **SEQUENTIAL LOGIC**

# 8.1 Registers

### 8.1.1 Introduction

Creating registers in SpinalHDL is very different than in VHDL or Verilog.

In Spinal, there are no process/always blocks. Registers are explicitly defined at declaration. This difference from traditional event-driven HDL has a big impact:

- You can assign registers and wires in the same scope, meaning the code doesn't need to be split between process/always blocks
- It make things much more flexible (see *Functions*)

Clocks and resets are handled separately, see the *Clock domain* chapter for details.

### 8.1.2 Instantiation

There are 4 ways to instantiate a register:

Syntax	Description	
Reg(type : Data)	Register of the given type	
RegInit(resetValue : Data)	Register loaded with the given resetValue when a reset	
	occurs	
RegNext(nextValue : Data)	Register that samples the given nextValue each cycle	
RegNextWhen(nextValue : Data, cond :	Register that samples the given nextValue when a con-	
Bool)	dition occurs	

Here is an example declaring some registers:

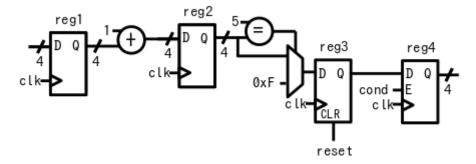
```
// UInt register of 4 bits
val reg1 = Reg(UInt(4 bit))

// Register that samples reg1 each cycle
val reg2 = RegNext(reg1 + 1)

// UInt register of 4 bits initialized with 0 when the reset occurs
val reg3 = RegInit(U"0000")
reg3 := reg2
when(reg2 === 5) {
    reg3 := 0xF
}
```

```
// Register that samples reg3 when cond is True
val reg4 = RegNextWhen(reg3, cond)
```

The code above will infer the following logic:



**Note:** The reg3 example above shows how you can assign the value of a RegInit register. It's possible to use the same syntax to assign to the other register types as well (Reg, RegNext, RegNextWhen). Just like in combinational assignments, the rule is 'Last assignment wins', but if no assignment is done, the register keeps its value.

Also, RegNext is an abstraction which is built over the Reg syntax. The two following sequences of code are strictly equivalent:

```
// Standard way
val something = Bool()
val value = Reg(Bool())
value := something

// Short way
val something = Bool()
val value = RegNext(something)
```

### 8.1.3 Reset value

In addition to the RegInit(value: Data) syntax which directly creates the register with a reset value, you can also set the reset value by calling the init(value: Data) function on the register.

```
// UInt register of 4 bits initialized with 0 when the reset occurs
val reg1 = Reg(UInt(4 bit)) init(0)
```

If you have a register containing a Bundle, you can use the init function on each element of the Bundle.

## 8.1.4 Initialization value for simulation purposes

For registers that don't need a reset value in RTL, but need an initialization value for simulation (to avoid x-propagation), you can ask for a random initialization value by calling the randBoot() function.

```
// UInt register of 4 bits initialized with a random value
val reg1 = Reg(UInt(4 bit)) randBoot()
```

## 8.2 RAM/ROM

# **8.2.1 Syntax**

To create a memory in SpinalHDL, the Mem class should be used. It allows you to define a memory and add read and write ports to it.

The following table shows how to instantiate a memory:

Syntax	Description
Mem(type : Data, size : Int)	Create a RAM
Mem(type : Data, initialContent :	Create a ROM. If your target is an FPGA, because the
Array[Data])	memory can be inferred as a block ram, you can still
	create write ports on it.

**Note:** If you want to define a ROM, elements of the initialContent array should only be literal values (no operator, no resize functions). There is an example *here*.

**Note:** To give a RAM initial values, you can also use the init function.

The following table show how to add access ports on a memory:

8.2. RAM/ROM 103

tax men <b>Siyanddhresso</b> ) us write :=	tı
data merA(sy)nchronous read	Γ
Synchronous write with an optional mask. If no enable is specified, it's automatically inferred from the conditional scope where this function is called mem.write(	1
address	
data	
[enable]	
[mask]	
Asynchronous read with an optional read-under-write policy	Т
mem.readAsync(	
address	
[readUnderWrite]	
Synchronous read with an optional enable, read-under-write policy, and clockCrossing mode	Γ
mem.readSync(	
address	
[enable]	
[readUnderWrite]	
[clockCrossing]	
	Γ
Infer a read/write port. mem.readWriteSync( data is written when enable && write.	
Return the read data, the read occurs when enable is true	
address and dam, are read seems when chapte is a de	
data	
enable	
write	
04 [mask] Chapter 8. Sequential Ic	gi
[readUnderWrite]	

**Note:** If for some reason you need a specific memory port which is not implemented in Spinal, you can always abstract over your memory by specifying a BlackBox for it.

**Important:** Memory ports in SpinalHDL are not inferred, but are explicitly defined. You should not use coding templates like in VHDL/Verilog to help the synthesis tool to infer memory.

Here is a example which infers a simple dual port ram (32 bits \* 256):

```
val mem = Mem(Bits(32 bits), wordCount = 256)
mem.write(
  enable = io.writeValid,
  address = io.writeAddress,
  data = io.writeData
)

io.readData := mem.readSync(
  enable = io.readValid,
  address = io.readAddress
)
```

# 8.2.2 Read-under-write policy

This policy specifies how a read is affected when a write occurs in the same cycle to the same address.

Kinds	Description
dontCare	Don't care about the read value when the case occurs
readFirst	The read will get the old value (before the write)
writeFirst	The read will get the new value (provided by the write)

**Important:** The generated VHDL/Verilog is always in the readFirst mode, which is compatible with dontCare but not with writeFirst. To generate a design that contains this kind of feature, you need to enable *automatic memory blackboxing*.

8.2. RAM/ROM 105

# 8.2.3 Mixed-width ram

You can specify ports that access the memory with a width that is a power of two fraction of the memory width using these functions:

Syntax	Description
- Cymax	Similar to mem.write
mem.writeMixed	Width(
address	
data	
[readUnde	rWrite]
)	
	Similar to mem.readAsync, but in place of returning the read value, it drives the sig-
	nal/object given as the data argument
mem.readAsync	whited whith(
address	
data	
[readUnde	rWrita]
	write]
)	
	Similar to mem.readSync, but in place of returning the read value, it drives the signal/object
	given as the data argument
mem.readSyncN	
address	
data	
[enable]	
[enacie]	
[readUnde	rWrite]
<b>L</b>	
[clockCros	sing]
)	
,	
	Equivalent to mem.readWriteSync
	PyncMixedWidth(
address	
data	
enable	
write	
[mask]	
[readUnde	rWrite]
[clockCros	sking]
)	

**Important:** As for read-under-write policy, to use this feature you need to enable *automatic memory blackboxing*,

because there is no universal VHDL/Verilog language template to infer mixed-width ram.

#### 8.2.4 Automatic blackboxing

Because it's impossible to infer all ram kinds by using regular VHDL/Verilog, SpinalHDL integrates an optional automatic blackboxing system. This system looks at all memories present in your RTL netlist and replaces them with blackboxes. Then the generated code will rely on third party IP to provide the memory features, such as the read-during-write policy and mixed-width ports.

Here is an example of how to enable blackboxing of memories by default:

```
def main(args: Array[String]) {
    SpinalConfig()
        .addStandardMemBlackboxing(blackboxAll)
        .generateVhdl(new TopLevel)
}
```

If the standard blackboxing tools don't do enough for your design, do not hesitate to create a Github issue. There is also a way to create your own blackboxing tool.

#### **Blackboxing policy**

There are multiple policies that you can use to select which memory you want to blackbox and also what to do when the blackboxing is not feasible:

Description	
Blackbox all memory.	
Throw an error on unblackboxable memory	
Blackbox all memory that is blackboxable	
blackboxRequestedAndUninferable	
Blackbox memory specified by the user and memory that is known to be uninferable (mixed-width,).	
Throw an error on unblackboxable memory	
Blackbox memory specified by the user	
Throw an error on unblackboxable memory	

To explicitly set a memory to be blackboxed, you can use its generateAsBlackBox function.

```
val mem = Mem(Rgb(rgbConfig), 1 << 16)
mem.generateAsBlackBox()</pre>
```

You can also define your own blackboxing policy by extending the MemBlackboxingPolicy class.

8.2. RAM/ROM 107

#### Standard memory blackboxes

Shown below are the VHDL definitions of the standard blackboxes used in SpinalHDL:

```
-- Simple asynchronous dual port (1 write port, 1 read port)
component Ram_1w_1ra is
 generic(
   wordCount : integer;
   wordWidth : integer;
   technology : string;
   readUnderWrite : string;
   wrAddressWidth : integer;
   wrDataWidth : integer;
   wrMaskWidth : integer;
   wrMaskEnable : boolean;
   rdAddressWidth : integer;
   rdDataWidth : integer
 );
 port(
   clk : in std_logic;
   wr_en : in std_logic;
   wr_mask : in std_logic_vector;
   wr_addr : in unsigned;
   wr_data : in std_logic_vector;
   rd_addr : in unsigned;
   rd_data : out std_logic_vector
 );
end component;
-- Simple synchronous dual port (1 write port, 1 read port)
component Ram_1w_1rs is
 generic(
   wordCount : integer;
   wordWidth : integer;
   clockCrossing : boolean;
   technology : string;
   readUnderWrite : string;
   wrAddressWidth : integer;
   wrDataWidth : integer;
   wrMaskWidth : integer;
   wrMaskEnable : boolean;
   rdAddressWidth : integer;
   rdDataWidth : integer;
   rdEnEnable : boolean
 );
 port(
   wr_clk : in std_logic;
   wr_en : in std_logic;
   wr_mask : in std_logic_vector;
   wr_addr : in unsigned;
   wr_data : in std_logic_vector;
   rd_clk : in std_logic;
   rd_en : in std_logic;
   rd_addr : in unsigned;
   rd_data : out std_logic_vector
 );
end component;
```

(continued from previous page)

```
-- Single port (1 readWrite port)
component Ram_1wrs is
 generic(
   wordCount : integer;
   wordWidth : integer;
   readUnderWrite : string;
    technology: string
 );
 port(
   clk : in std_logic;
    en : in std_logic;
   wr : in std_logic;
   addr : in unsigned;
   wrData : in std_logic_vector;
   rdData : out std_logic_vector
 );
end component;
--True dual port (2 readWrite port)
component Ram_2wrs is
 generic(
   wordCount : integer;
   wordWidth : integer;
    clockCrossing : boolean;
    technology : string;
   portA_readUnderWrite : string;
   portA_addressWidth : integer;
   portA_dataWidth : integer;
   portA_maskWidth : integer;
   portA_maskEnable : boolean;
   portB_readUnderWrite : string;
   portB_addressWidth : integer;
   portB_dataWidth : integer;
   portB_maskWidth : integer;
   portB maskEnable : boolean
 );
 port(
   portA_clk : in std_logic;
   portA_en : in std_logic;
   portA_wr : in std_logic;
   portA_mask : in std_logic_vector;
   portA_addr : in unsigned;
   portA_wrData : in std_logic_vector;
   portA_rdData : out std_logic_vector;
   portB_clk : in std_logic;
   portB_en : in std_logic;
   portB_wr : in std_logic;
   portB_mask : in std_logic_vector;
   portB_addr : in unsigned;
   portB_wrData : in std_logic_vector;
   portB_rdData : out std_logic_vector
 );
end component;
```

As you can see, blackboxes have a technology parameter. To set it, you can use the setTechnology function on the corresponding memory. There are currently 4 kinds of technologies possible:

8.2. RAM/ROM 109

# **SpinalHDL Documentation**

- auto
- ramBlock
- distributedLut
- registerFile

**CHAPTER** 

NINE

# **DESIGN ERRORS**

# 9.1 Assignment overlap

#### 9.1.1 Introduction

SpinalHDL will check that no signal assignment completely erases a previous one.

# 9.1.2 Example

The following code

```
class TopLevel extends Component {
  val a = UInt(8 bits)
  a := 42
  a := 66 // Erase the a := 42 assignment
}
```

will throw the following error:

```
ASSIGNMENT OVERLAP completely the previous one of (toplevel/a : UInt[8 bits])

***

Source file location of the a := 66 assignment via the stack trace

***
```

A fix could be:

```
class TopLevel extends Component {
  val a = UInt(8 bits)
  a := 42
  when(something) {
    a := 66
  }
}
```

But in the case when you really want to override the previous assignment (as there are times when overriding makes sense), you can do the following:

```
class TopLevel extends Component {
  val a = UInt(8 bits)
  a := 42
  a.allowOverride
  a := 66
}
```

# 9.2 Clock crossing violation

#### 9.2.1 Introduction

SpinalHDL will check that every register of your design only depends (through combinational logic paths) on registers which use the same or a synchronous clock domain.

## 9.2.2 Example

The following code:

```
class TopLevel extends Component {
  val clkA = ClockDomain.external("clkA")
  val clkB = ClockDomain.external("clkB")

val regA = clkA(Reg(UInt(8 bits))) // PlayDev.scala:834
  val regB = clkB(Reg(UInt(8 bits))) // PlayDev.scala:835

val tmp = regA + regA // PlayDev.scala:838
  regB := tmp
}
```

will throw:

```
CLOCK CROSSING VIOLATION from (toplevel/regA : UInt[8 bits]) to (toplevel/regB : __
UInt[8 bits]).
- Register declaration at
    ***
    Source file location of the toplevel/regA definition via the stack trace
    ***
- through
    >>> (toplevel/regA : UInt[8 bits]) at ***(PlayDev.scala:834) >>>
    >>> (toplevel/tmp : UInt[8 bits]) at ***(PlayDev.scala:838) >>>
    >>> (toplevel/regB : UInt[8 bits]) at ***(PlayDev.scala:835) >>>
```

There are multiple possible fixes, listed below:

- crossClockDomain tags
- setSyncronousWith method
- BufferCC type

#### crossClockDomain tag

The crossClockDomain tag can be used to communicate "It's alright, don't panic about this specific clock crossing" to the SpinalHDL compiler.

```
class TopLevel extends Component {
  val clkA = ClockDomain.external("clkA")
  val clkB = ClockDomain.external("clkB")

  val regA = clkA(Reg(UInt(8 bits)))
  val regB = clkB(Reg(UInt(8 bits))).addTag(crossClockDomain)

  val tmp = regA + regA
```

(continued from previous page)

```
regB := tmp
}
```

#### setSyncronousWith

You can also specify that two clock domains are synchronous together by using the setSynchronousWith method of one of the ClockDomain objects.

```
class TopLevel extends Component {
  val clkA = ClockDomain.external("clkA")
  val clkB = ClockDomain.external("clkB")
  clkB.setSyncronousWith(clkA)

  val regA = clkA(Reg(UInt(8 bits)))
  val regB = clkB(Reg(UInt(8 bits)))

  val tmp = regA + regA
  regB := tmp
}
```

#### **BufferCC**

When exchanging single-bit signals (such as Bool types), or Gray-coded values, you can use BufferCC to safely cross different ClockDomain regions.

**Warning:** Do not use BufferCC with multi-bit signals, as there is a risk of corrupted reads on the receiving side if the clocks are asynchronous. See the *Clock Domains* page for more details.

```
class AsyncFifo extends Component {
  val popToPushGray = Bits(ptrWidth bits)
   val pushToPopGray = Bits(ptrWidth bits)
  val pushCC = new ClockingArea(pushClock) {
                   = Counter(depth << 1)</pre>
    val pushPtr
     val pushPtrGray = RegNext(toGray(pushPtr.valueNext)) init(0)
    val popPtrGray = BufferCC(popToPushGray, B(0, ptrWidth bits))
    val full
                   = isFull(pushPtrGray, popPtrGray)
   }
  val popCC = new ClockingArea(popClock) {
                 = Counter(depth << 1)</pre>
    val popPtr
     val popPtrGray = RegNext(toGray(popPtr.valueNext)) init(0)
     val pushPtrGray = BufferCC(pushToPopGray, B(0, ptrWidth bit))
                   = isEmpty(popPtrGray, pushPtrGray)
     val empty
   }
}
```

# 9.3 Combinatorial loop

#### 9.3.1 Introduction

SpinalHDL will check that there are no combinatorial loops in the design.

#### 9.3.2 Example

The following code:

```
class TopLevel extends Component {
  val a = UInt(8 bits) // PlayDev.scala line 831
  val b = UInt(8 bits) // PlayDev.scala line 832
  val c = UInt(8 bits)
  val d = UInt(8 bits)

  a := b
  b := c | d
  d := a
  c := 0
}
```

will throw:

```
COMBINATORIAL LOOP :
   Partial chain :
    >>> (toplevel/a : UInt[8 bits]) at ***(PlayDev.scala:831) >>>
    >>> (toplevel/d : UInt[8 bits]) at ***(PlayDev.scala:834) >>>
    >>> (toplevel/b : UInt[8 bits]) at ***(PlayDev.scala:832) >>>
    >>> (toplevel/a : UInt[8 bits]) at ***(PlayDev.scala:831) >>>

Full chain :
    (toplevel/a : UInt[8 bits])
    (toplevel/d : UInt[8 bits])
    (UInt | UInt)[8 bits]
    (toplevel/b : UInt[8 bits])
    (toplevel/a : UInt[8 bits])
```

A possible fix could be:

```
class TopLevel extends Component {
  val a = UInt(8 bits) // PlayDev.scala line 831
  val b = UInt(8 bits) // PlayDev.scala line 832
  val c = UInt(8 bits)
  val d = UInt(8 bits)

  a := b
  b := c | d
  d := 42
  c := 0
}
```

# 9.3.3 False-positives

It should be said that SpinalHDL's algorithm to detect combinatorial loops can be pessimistic, and it may give false positives. If it is giving a false positive, you can manually disable loop checking on one signal of the loop like so:

```
class TopLevel extends Component {
  val a = UInt(8 bits)
  a := 0
  a(1) := a(0) // False positive because of this line
}
```

could be fixed by:

```
class TopLevel extends Component {
  val a = UInt(8 bits).noCombLoopCheck
  a := 0
  a(1) := a(0)
}
```

It should also be said that assignments such as (a(1) := a(0)) can make some tools like Verilator unhappy. It may be better to use a Vec(Bool, 8) in this case.

# 9.4 Hierarchy violation

#### 9.4.1 Introduction

SpinalHDL will check that signals are never accessed outside of the current component's scope.

The following signals can be read inside a component:

- All directionless signals defined in the current component
- All in/out/inout signals of the current component
- All in/out/inout signals of child components

In addition, the following signals can be assigned to inside of a component:

- All directionless signals defined in the current component
- All out/inout signals of the current component
- All in/inout signals of child components

If a HIERARCHY VIOLATION error appears, it means that one of the above rules was violated.

#### 9.4.2 Example

The following code:

```
class TopLevel extends Component {
  val io = new Bundle {
    val a = in UInt(8 bits)
  }
  val tmp = U"x42"
  io.a := tmp
}
```

will throw:

```
HIERARCHY VIOLATION : (toplevel/io_a : in UInt[8 bits]) is driven by (toplevel/tmp : __
UInt[8 bits]), but isn't accessible in the toplevel component.

***

Source file location of the `io.a := tmp` via the stack trace

***
```

A fix could be:

```
class TopLevel extends Component {
  val io = new Bundle {
    val a = out UInt(8 bits) // changed from in to out
  }
  val tmp = U"x42"
  io.a := tmp
}
```

## 9.5 lo bundle

#### 9.5.1 Introduction

SpinalHDL will check that each io bundle contains only in/out/inout signals.

#### 9.5.2 Example

The following code:

```
class TopLevel extends Component {
  val io = new Bundle {
    val a = UInt(8 bits)
  }
}
```

will throw:

```
IO BUNDLE ERROR : A direction less (toplevel/io_a : UInt[8 bits]) signal was defined_
into toplevel component's io bundle

***

Source file location of the toplevel/io_a definition via the stack trace

***
```

A fix could be:

```
class TopLevel extends Component {
  val io = new Bundle {
    val a = in UInt(8 bits)
  }
}
```

But if for meta hardware description reasons you really want io.a to be directionless, you can do:

```
class TopLevel extends Component {
  val io = new Bundle {
    val a = UInt(8 bits)
  }
```

(continued from previous page)

```
a.allowDirectionLessIo
}
```

# 9.6 Latch detected

# 9.6.1 Introduction

SpinalHDL will check that no combinational signals will infer a latch during synthesis. In other words, this is a check that no combinational signals are partially assigned.

## 9.6.2 Example

The following code:

```
class TopLevel extends Component {
  val cond = in(Bool)
  val a = UInt(8 bits)

  when(cond) {
    a := 42
  }
}
```

will throw:

```
LATCH DETECTED from the combinatorial signal (toplevel/a : UInt[8 bits]), defined at

***

Source file location of the toplevel/io_a definition via the stack trace

***
```

A fix could be:

```
class TopLevel extends Component {
  val cond = in(Bool)
  val a = UInt(8 bits)

a := 0
  when(cond) {
    a := 42
  }
}
```

## 9.7 No driver on

#### 9.7.1 Introduction

SpinalHDL will check that all combinational signals which have an impact on the design are assigned by something.

9.6. Latch detected 117

#### 9.7.2 Example

The following code:

```
class TopLevel extends Component {
  val result = out(UInt(8 bits))
  val a = UInt(8 bits)
  result := a
}
```

will throw:

```
NO DRIVER ON (toplevel/a : UInt[8 bits]), defined at

***

Source file location of the toplevel/a definition via the stack trace

***
```

A fix could be:

```
class TopLevel extends Component {
  val result = out(UInt(8 bits))
  val a = UInt(8 bits)
  a := 42
  result := a
}
```

# 9.8 NullPointerException

#### 9.8.1 Introduction

NullPointerException is a Scala runtime reported error which can happen when a variable is accessed before it has been initialized.

# 9.8.2 Example

The following code:

```
class TopLevel extends Component {
  a := 42
  val a = UInt(8 bits)
}
```

will throw:

```
Exception in thread "main" java.lang.NullPointerException

***

Source file location of the a := 42 assignment via the stack trace

***
```

A fix could be:

```
class TopLevel extends Component {
  val a = UInt(8 bits)
  a := 42
}
```

#### Issue explanation

SpinalHDL is not a language, it is a Scala library, which means that it obeys the same rules as the Scala general purpose programming language.

When running the above SpinalHDL hardware description to generate the corresponding VHDL/Verilog RTL, the SpinalHDL hardware description will be executed as a Scala program, and a will be a null reference until the program executes val a = UInt(8 bits), so trying to assign to it before then will result in a NullPointerException.

# 9.9 Register defined as component input

#### 9.9.1 Introduction

In SpinalHDL, you are not allowed to define a component that has a register as an input. The reasoning behind this is to prevent surprises when the user tries to drive the inputs of child components with the registered signal. If a registered input is desired, you will need to declare the unregistered input in the io bundle, and register the signal in the body of the component.

#### 9.9.2 Example

The following code:

```
class TopLevel extends Component {
  val io = new Bundle {
    val a = in(Reg(UInt(8 bits)))
  }
}
```

will throw:

```
REGISTER DEFINED AS COMPONENT INPUT : (toplevel/io_a : in UInt[8 bits]) is defined as_
a registered input of the toplevel component, but isn't allowed.

***

Source file location of the toplevel/io_a definition via the stack trace

***
```

A fix could be:

```
class TopLevel extends Component {
  val io = new Bundle {
    val a = in UInt(8 bits)
  }
}
```

If a registered a is required, it can be done like so:

```
class TopLevel extends Component {
  val io = new Bundle {
    val a = in UInt(8 bits)
  }
  val a = RegNext(io.a)
}
```

# 9.10 Scope violation

#### 9.10.1 Introduction

SpinalHDL will check that there are no signals assigned outside the scope they are defined in. This error isn't easy to trigger as it requires some specific meta hardware description tricks.

# **9.10.2 Example**

The following code:

```
class TopLevel extends Component {
  val cond = Bool()

  var tmp : UInt = null
  when(cond) {
    tmp = UInt(8 bits)
  }
  tmp := U"x42"
}
```

will throw:

```
SCOPE VIOLATION : (toplevel/tmp : UInt[8 bits]) is assigned outside its declaration... scope at
***

Source file location of the tmp := U"x42" via the stack trace
***
```

A fix could be:

```
class TopLevel extends Component {
  val cond = Bool()

  var tmp : UInt = UInt(8 bits)
  when(cond) {

  }
  tmp := U"x42"
}
```

# 9.11 Spinal can't clone class

#### 9.11.1 Introduction

This error happens when SpinalHDL wants to create a new datatype instance via the cloneOf function but isn't able to do it. The reason for this is nearly always because it can't retrieve the construction parameters of a Bundle.

#### **9.11.2 Example**

The following code:

will throw:

A fix could be:

```
case class RGB(width : Int) extends Bundle {
  val r, g, b = UInt(width bits)
}
class TopLevel extends Component {
  val tmp = Stream(RGB(8))
}
```

# 9.12 Unassigned register

#### 9.12.1 Introduction

SpinalHDL will check that all registers which impact the design have been assigned somewhere.

#### **9.12.2 Example**

The following code:

```
class TopLevel extends Component {
  val result = out(UInt(8 bits))
  val a = Reg(UInt(8 bits))
  result := a
}
```

will throw:

```
UNASSIGNED REGISTER (toplevel/a : UInt[8 bits]), defined at
   ***
```

(continued from previous page)

```
Source file location of the toplevel/a definition via the stack trace ***
```

A fix could be:

```
class TopLevel extends Component {
  val result = out(UInt(8 bits))
  val a = Reg(UInt(8 bits))
  a := 42
  result := a
}
```

#### 9.12.3 Register with only init

In some cases, because of the design parameterization, it could make sense to generate a register which has no assignment but only an init statement.

```
class TopLevel extends Component {
  val result = out(UInt(8 bits))
  val a = Reg(UInt(8 bits)) init(42)

if(something)
  a := somethingElse
  result := a
}
```

will throw:

```
UNASSIGNED REGISTER (toplevel/a : UInt[8 bits]), defined at

***

Source file location of the toplevel/a definition via the stack trace

***
```

To fix it, you can ask SpinalHDL to transform the register into a combinational one if no assignment is present but it has an init statement:

```
class TopLevel extends Component {
  val result = out(UInt(8 bits))
  val a = Reg(UInt(8 bits)).init(42).allowUnsetRegToAvoidLatch

  if(something)
    a := somethingElse
  result := a
}
```

#### 9.13 Unreachable is statement

#### 9.13.1 Introduction

SpinalHDL will check to ensure that all is statements in a switch are reachable.

#### **9.13.2 Example**

The following code:

```
class TopLevel extends Component {
  val sel = UInt(2 bits)
  val result = UInt(4 bits)
  switch(sel) {
    is(0){ result := 4 }
    is(1){ result := 6 }
    is(2){ result := 8 }
    is(3){ result := 9 }
    is(0){ result := 2 } // Duplicated is statement!
  }
}
```

will throw:

```
UNREACHABLE IS STATEMENT in the switch statement at

***

Source file location of the is statement definition via the stack trace

***
```

A fix could be:

```
class TopLevel extends Component {
  val sel = UInt(2 bits)
  val result = UInt(4 bits)
  switch(sel) {
    is(0){ result := 4 }
    is(1){ result := 6 }
    is(2){ result := 8 }
    is(3){ result := 9 }
}
```

# 9.14 Width mismatch

#### 9.14.1 Introduction

SpinalHDL will check that operators and signals on the left and right side of assignments have the same widths.

#### 9.14.2 Assignment example

The following code:

```
class TopLevel extends Component {
  val a = UInt(8 bits)
  val b = UInt(4 bits)
  b := a
}
```

will throw:

```
WIDTH MISMATCH on (toplevel/b : UInt[4 bits]) := (toplevel/a : UInt[8 bits]) at
    ***
Source file location of the OR operator via the stack trace
    ***
```

A fix could be:

```
class TopLevel extends Component {
  val a = UInt(8 bits)
  val b = UInt(4 bits)
  b := a.resized
}
```

#### 9.14.3 Operator example

The following code:

```
class TopLevel extends Component {
  val a = UInt(8 bits)
  val b = UInt(4 bits)
  val result = a | b
}
```

will throw:

```
WIDTH MISMATCH on (UInt | UInt)[8 bits]
- Left operand : (toplevel/a : UInt[8 bits])
- Right operand : (toplevel/b : UInt[4 bits])
at
***
Source file location of the OR operator via the stack trace
***
```

A fix could be:

```
class TopLevel extends Component {
  val a = UInt(8 bits)
  val b = UInt(4 bits)
  val result = a | (b.resized)
}
```

# 9.15 Introduction

The SpinalHDL compiler will perform many checks on your design to be sure that the generated VHDL/Verilog will be safe for simulation and synthesis. Basically, it should not be possible to generate a broken VHDL/Verilog design. Below is a non-exhaustive list of SpinalHDL checks:

- Assignment overlapping
- · Clock crossing
- · Hierarchy violation
- · Combinatorial loops
- Latches
- Undriven signals
- · Width mismatch
- Unreachable switch statements

On each SpinalHDL error report, you will find a stack trace, which can be useful to accurately find out where the design error is. These design checks may look like overkill at first glance, but they becomes invaluable as soon as you start to move away from the traditional way of doing hardware description.

9.15. Introduction 125

**CHAPTER** 

TEN

# **OTHER LANGUAGE FEATURES**

## **10.1 Utils**

#### 10.1.1 General

Many tools and utilities are present in *spinal.lib* but some are already present in the SpinalHDL Core.

Syntax	Return	Description	
widthOf(x : BitVector)	Int	Return the width of a Bits/UInt/SInt signal	
log2Up(x : BigInt)	Int	Return the number of bits needed to represent <b>x</b> states	
isPow2(x : BigInt)	Boolean	Return true if x is a power of two	
roundUp(that : BigInt, by :	BigInt	Return the first by multiply from that (included)	
BigInt)			
Cat(x : Data*)	Bits	Concatenate all arguments, the first in MSB, the last	
		in LSB	

# 10.1.2 Cloning hardware datatypes

You can clone a given hardware data type by using the cloneOf(x) function. It will return a new instance of the same Scala type and parameters.

For example:

```
def plusOne(value : UInt) : UInt = {
    // Will recreate a UInt with the same width than ``value``
    val temp = cloneOf(value)
    temp := value + 1
    return temp
}

// treePlusOne will become a 8 bits value
val treePlusOne = plusOne(U(3, 8 bits))
```

You can get more information about how hardware data types are managed on the Hardware types page.

**Note:** If you use the cloneOf function on a Bundle, this Bundle should be a case class or should override the clone function internally.

#### 10.1.3 Passing a datatype as construction parameter

Many pieces of reusable hardware need to be parameterized by some data type. For example if you want to define a FIFO or a shift register, you need a parameter to specify which kind of payload you want for the component.

There are two similar ways to do this.

#### The old way

A good example of the old way to do this is in this definition of a ShiftRegister component:

```
case class ShiftRegister[T <: Data](dataType: T, depth: Int) extends Component {
  val io = new Bundle {
    val input = in (cloneOf(dataType))
    val output = out(cloneOf(dataType))
  }
  // ...
}</pre>
```

And here is how you can instantiate the component:

```
val shiftReg = ShiftRegister(Bits(32 bits), depth = 8)
```

As you can see, the raw hardware type is directly passed as a construction parameter. Then each time you want to create an new instance of that kind of hardware data type, you need to use the cloneOf(...) function. Doing things this way is not super safe as it's easy to forget to use cloneOf.

#### The safe way

An example of the safe way to pass a data type parameter is as follows:

And here is how you instantiate the component (exactly the same as before):

```
val shiftReg = ShiftRegister(Bits(32 bits), depth = 8)
```

Notice how the example above uses a HardType wrapper around the raw data type T, which is a "blueprint" definition of a hardware data type. This way of doing things is easier to use than the "old way", because to create a new instance of the hardware data type you only need to call the apply function of that HardType (or in other words, just add parentheses after the parameter).

Additionally, this mechanism is completely transparent from the point of view of the user, as a hardware data type can be implicitly converted into a HardType.

## 10.1.4 Frequency and time

SpinalHDL has a dedicated syntax to define frequency and time values:

```
val frequency = 100 MHz
val timeoutLimit = 3 ms
val period = 100 us

val periodCycles = frequency * period
val timeoutCycles = frequency * timeoutLimit
```

For time definitions you can use following postfixes to get a TimeNumber:

```
fs, ps, ns, us, ms, sec, mn, hr
```

For time definitions you can use following postfixes to get a HertzNumber:

```
Hz, KHz, MHz, GHz, THz
```

TimeNumber and HertzNumber are based on the PhysicalNumber class which use scala BigDecimal to store numbers.

# 10.1.5 Binary prefix

SpinalHDL allows the definition of integer numbers using binary prefix notation according to IEC.

```
val memSize = 512 MiB
val dpRamSize = 4 KiB
```

The following binary prefix notations are available:

Binary Prefix	Value
Byte, Bytes	1
KiB	1024 == 1 << 10
MiB	$1024^2 == 1 << 20$
GiB	$1024^3 == 1 << 30$
TiB	$1024^4 == 1 << 40$
PiB	$1024^5 == 1 << 50$
EiB	$1024^6 == 1 << 60$
ZiB	$1024^7 == 1 << 70$
YiB	$1024^8 == 1 << 80$

## 10.2 Assertions

In addition to Scala run-time assertions, you can add hardware assertions using the following syntax:

```
assert(assertion : Bool, message : String = null, severity: AssertNodeSeverity =
Error)
```

Severity levels are:

10.2. Assertions 129

Name	Description
NOTE	Used to report an informative message
WARNING	Used to report an unusual case
ERROR	Used to report an situation that should not happen
FAILURE	Used to report a fatal situation and close the simulation

One practical example could be to check that the valid signal of a handshake protocol never drops when ready is low:

```
class TopLevel extends Component {
  val valid = RegInit(False)
  val ready = in Bool

  when(ready) {
    valid := False
  }
  // some logic

assert(
    assertion = !(valid.fall && !ready),
    message = "Valid dropped when ready was low",
    severity = ERROR
  )
}
```

# 10.3 Report

You can add debugging in RTL for simulation, using the following syntax:

```
object Enum extends SpinalEnum{
    val MIAOU, RAWRR = newElement()
}

class TopLevel extends Component {
    val a = Enum.RAWRR()
    val b = U(0x42)
    val c = out(Enum.RAWRR())
    val d = out (U(0x42))
    report(Seq("miaou", a, b, c, d))
}
```

It will generate the following Verilog code for example:

```
$display("NOTE miaou %s%x%s%x", a_string, b, c_string, d);
```

Since SpinalHDL 1.4.4, the following syntax is also supported:

```
report(L'miaou $a $b $c $d")
```

## 10.4 Formal

#### **10.4.1 General**

There is limited support for SystemVerilog Assertions (SVA).

You can add formal statements (assume, assert, etc.) in the Component definition, like in the example below:

```
class TopLevel extends Component {
 val io = new Bundle {
   val ready = in Bool()
   val valid = out Bool()
 val valid = RegInit(False)
when(io.ready) {
  valid := False
 io.valid <> valid
 // some logic
 import spinal.core.GenerationFlags._
 import spinal.core.Formal._
GenerationFlags.formal {
  when(initstate()) {
     assume(clockDomain.isResetActive)
     assume(io.ready === False)
   }.otherwise {
     assert(!(valid.fall && !io.ready))
}
}
```

To generate a design which includes the formal statements you can use  ${\tt includeFormal}$ :

```
object MyToplevelSystemVerilogWithFormal {
  def main(args: Array[String]) {
    val config = SpinalConfig(defaultConfigForClockDomains = ClockDomainConfig(resetKind=SYNC, resetActiveLevel=HIGH))
    config.includeFormal.generateSystemVerilog(new TopLevel())
  }
}
```

10.4. Formal 131

## 10.4.2 Clocked assertion without reset

```
ClockDomain.current.withoutReset(){
  assert(wuff === 0)
}
```

# 10.4.3 Specifying the initial value of a signal

For instance, for the reset signal of the current clockdomain (usefull at the top)

```
ClockDomain.current.readResetWire initial(False)
```

# 10.4.4 Specifying a initial assumption

```
assumeInitial(!clockDomain.isResetActive)
```

# 10.4.5 Specifying assertion in the reset scope

```
ClockDomain.current.duringReset {
  assume(rawrrr === 0)
  assume(wuff === 3)
}
```

# 10.4.6 Supported features

Syntax	Returns	Creates in SystemVerilog
assert()		assert()
cover()		cover()
	T	past(that)
<pre>past(that : T, delay : Int) past(that : T)</pre>		
rose(that : Bool)	Bool	rose(that)
fell(that : Bool)	Bool	fell(that)
changed(that : Bool)	Bool	changed(that)
stable(that : Bool)	Bool	stable(that)
initstate()	Bool	<pre>\$initstate()</pre>

#### 10.4.7 Limitations

No support for unclocked assertions.

# 10.5 Analog and inout

#### 10.5.1 Introduction

You can define native tristate signals by using the Analog/inout features. These features were added for the following reasons:

- Being able to add native tristate signals to the toplevel (it avoids having to manually wrap them with some hand-written VHDL/Verilog).
- Allowing the definition of blackboxes which contain inout pins.
- Being able to connect a blackbox's inout pin through the hierarchy to a toplevel inout pin.

As those features were only added for convenience, please do not try other fancy stuff with tristate logic just yet.

If you want to model a component like a memory-mapped GPIO peripheral, please use the *TriState/TriStateArray* bundles from the Spinal standard library, which abstract over the true nature of tristate drivers.

## 10.5.2 **Analog**

Analog is the keyword which allows a signal to be defined as something analog, which in the digital world could mean 0, 1, or Z (the disconnected, high-impedance state).

For instance:

#### 10.5.3 inout

inout is the keyword which allows you to set an Analog signal as a bidirectional (both "in" and "out") signal.

For instance:

#### 10.5.4 InOutWrapper

InOutWrapper is a tool which allows you to transform all master TriState/TriStateArray/ReadableOpenDrain bundles of a component into native inout(Analog(...)) signals. It allows you to keep your hardware description free of any Analog/inout things, and then transform the toplevel to make it synthesis ready.

For instance:

```
case class Apb3Gpio(gpioWidth : Int) extends Component {
  val io = new Bundle{
    val gpio = master(TriStateArray(gpioWidth bits))
    val apb = slave(Apb3(Apb3Gpio.getApb3Config()))
  }
  ...
}
SpinalVhdl(InOutWrapper(Apb3Gpio(32)))
```

Will generate:

```
entity Apb3Gpio is
 port(
    io_gpio : inout std_logic_vector(31 downto 0); -- This io_gpio was originally a_
→ TriStateArray Bundle
    io_apb_PADDR : in unsigned(3 downto 0);
    io_apb_PSEL : in std_logic_vector(0 downto 0);
    io_apb_PENABLE : in std_logic;
    io_apb_PREADY : out std_logic;
    io_apb_PWRITE : in std_logic;
    io_apb_PWDATA : in std_logic_vector(31 downto 0);
    io_apb_PRDATA : out std_logic_vector(31 downto 0);
    io_apb_PSLVERROR : out std_logic;
   clk : in std_logic;
   reset : in std_logic
 );
end Apb3Gpio;
```

Instead of:

```
entity Apb3Gpio is
 port(
    io_gpio_read : in std_logic_vector(31 downto 0);
    io_gpio_write : out std_logic_vector(31 downto 0);
    io_gpio_writeEnable : out std_logic_vector(31 downto 0);
    io_apb_PADDR : in unsigned(3 downto 0);
    io_apb_PSEL : in std_logic_vector(0 downto 0);
    io_apb_PENABLE : in std_logic;
    io_apb_PREADY : out std_logic;
    io_apb_PWRITE : in std_logic;
    io_apb_PWDATA : in std_logic_vector(31 downto 0);
    io_apb_PRDATA : out std_logic_vector(31 downto 0);
    io_apb_PSLVERROR : out std_logic;
   clk : in std_logic;
   reset : in std_logic
 );
end Apb3Gpio;
```

#### 10.5.5 Manually driving Analog bundles

If an Analog bundle is not driven, it will default to being high-Z. Therefore to manually implement a tristate driver (in case the InOutWrapper type can't be used for some reason) you have to conditionally drive the signal.

To manually connect a TriState signal to an Analog bundle:

```
case class Example extends Component {
  val io = new Bundle {
    val tri = slave(TriState(Bits(16 bit)))
    val analog = inout Analog(Bits(16 bit))
  }
  tri.read := analog
  when(tri.writeEnable) { analog := tri.write }
}
```

# 10.6 VHDL and Verilog generation

# 10.6.1 Generate VHDL and Verilog from a SpinalHDL Component

To generate the VHDL from a SpinalHDL component you just need to call SpinalVhdl(new YourComponent) in a Scala main.

Generating Verilog is exactly the same, but with SpinalVerilog in place of SpinalVHDL

```
import spinal.core._
// A simple component definition.
class MyTopLevel extends Component {
  // Define some input/output signals. Bundle like a VHDL record or a Verilog struct.
 val io = new Bundle {
   val a = in Bool()
   val b = in Bool()
   val c = out Bool()
 }
 // Define some asynchronous logic.
 io.c := io.a & io.b
}
// This is the main function that generates the VHDL and the Verilog corresponding to
→ MyTopLevel.
object MyMain {
 def main(args: Array[String]) {
    SpinalVhdl(new MyTopLevel)
    SpinalVerilog(new MyTopLevel)
 }
}
```

**Important:** SpinalVhdl and SpinalVerilog may need to create multiple instances of your component class, therefore the first argument is not a Component reference, but a function that returns a new component.

**Important:** The SpinalVerilog implementation began the 5th of June, 2016. This backend successfully passes the same regression tests as the VHDL one (RISCV CPU, Multicore and pipelined mandelbrot, UART RX/TX, Single clock fifo, Dual clock fifo, Gray counter, ...).

If you have any issues with this new backend, please make a Github issue describing the problem.

#### **Parametrization from Scala**

Argument name	Туре	Default	Description
mode	SpinalMode	null	
			Set the SpinalHDL hdl generation mode.
			Can be set to VHDL or Verilog
defaultCor	nfChrohRd-ocl mainCon-		Set the clock configuration that will be used as the default value for all new ClockDomain.
	fig	RisingEdgeC	Clock
		Asynchronou	usReset
		ResetActivel	High
		ClockEnable	ActiveHigh
	ji <b>BVeletoo</b> rAt	_	Change all unsigned/signed toplevel io into std_logic_vector.
defaultClo	c <b>kDomkDo</b> Fre	. ,	Default clock frequency.
	mainFre-	Frequency	
	quency		
targetDire	c <b>Stoin</b> g	Current di-	Directory where files are generated.
		rectory	

And this is the syntax to specify them:

```
SpinalConfig(mode=VHDL, targetDirectory="temp/myDesign").generate(new UartCtrl)

// Or for Verilog in a more scalable formatting:
SpinalConfig(
  mode=Verilog,
  targetDirectory="temp/myDesign"
).generate(new UartCtrl)
```

#### Parametrization from shell

You can also specify generation parameters by using command line arguments.

```
def main(args: Array[String]): Unit = {
    SpinalConfig.shell(args)(new UartCtrl)
}
```

The syntax for command line arguments is:

```
Usage: SpinalCore [options]

--vhdl

Select the VHDL mode

--verilog

Select the Verilog mode
```

(continued from previous page)

```
-d | --debug
Enter in debug mode directly
-o <value> | --targetDirectory <value>
Set the target directory
```

# 10.6.2 Generated VHDL and Verilog

How a SpinalHDL RTL description is translated into VHDL and Verilog is important:

- Names in Scala are preserved in VHDL and Verilog.
- Component hierarchy in Scala is preserved in VHDL and Verilog.
- when statements in Scala are emitted as if statements in VHDL and Verilog.
- switch statements in Scala are emitted as case statements in VHDL and Verilog in all standard cases.

#### Organization

When you use the VHDL generator, all modules are generated into a single file which contain three sections:

- 1. A package that contains the definition of all Enums
- 2. A package that contains functions used by the architectural elements
- 3. All components needed by your design

When you use the Verilog generation, all modules are generated into a single file which contains two sections:

- 1. All enumeration definitions used
- 2. All modules needed by your design

#### **Combinational logic**

Scala:

```
class TopLevel extends Component {
 val io = new Bundle {
   val cond
                    = in Bool()
   val value = in UInt(4 bits)
   val withoutProcess = out UInt(4 bits)
   val withProcess = out UInt(4 bits)
 io.withoutProcess := io.value
 io.withProcess := 0
 when(io.cond) {
   switch(io.value) {
     is(U"0000") {
       io.withProcess := 8
     is(U"0001") {
       io.withProcess := 9
     }
     default {
       io.withProcess := io.value+1
     }
   }
```

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```
}
```

VHDL:

```
entity TopLevel is
  port(
    io_cond : in std_logic;
    io_value : in unsigned(3 downto 0);
    io_withoutProcess : out unsigned(3 downto 0);
    io_withProcess : out unsigned(3 downto 0)
  );
end TopLevel;
architecture arch of TopLevel is
begin
  io_withoutProcess <= io_value;</pre>
  process(io_cond,io_value)
    io_withProcess <= pkg_unsigned("0000");</pre>
    if io_cond = '1' then
      case io_value is
        when pkg_unsigned("0000") =>
          io_withProcess <= pkg_unsigned("1000");</pre>
        when pkg_unsigned("0001") =>
          io_withProcess <= pkg_unsigned("1001");</pre>
        when others =>
          io_withProcess <= (io_value + pkg_unsigned("0001"));</pre>
      end case;
    end if;
  end process;
end arch;
```

#### **Sequential logic**

Scala:

```
class TopLevel extends Component {
  val io = new Bundle {
    val cond = in Bool()
    val value = in UInt (4 bit)
    val resultA = out UInt(4 bit)
    val resultB = out UInt(4 bit)
}

val regWithReset = Reg(UInt(4 bits)) init(0)
val regWithoutReset = Reg(UInt(4 bits))

regWithReset := io.value
regWithoutReset := 0
when(io.cond) {
    regWithoutReset := io.value
}

io.resultA := regWithReset
```

(continued from previous page)

```
io.resultB := regWithoutReset
}
```

VHDL:

```
entity TopLevel is
 port(
    io_cond : in std_logic;
    io_value : in unsigned(3 downto 0);
    io_resultA : out unsigned(3 downto 0);
    io_resultB : out unsigned(3 downto 0);
    clk : in std_logic;
    reset : in std_logic
 );
end TopLevel;
architecture arch of TopLevel is
 signal regWithReset : unsigned(3 downto 0);
 signal regWithoutReset : unsigned(3 downto 0);
begin
 io_resultA <= regWithReset;</pre>
 io_resultB <= regWithoutReset;</pre>
 process(clk,reset)
 begin
    if reset = '1' then
      regWithReset <= pkg_unsigned("0000");</pre>
    elsif rising_edge(clk) then
      regWithReset <= io_value;</pre>
    end if;
 end process;
 process(clk)
 begin
    if rising_edge(clk) then
      regWithoutReset <= pkg_unsigned("0000");</pre>
      if io_cond = '1' then
        regWithoutReset <= io_value;</pre>
      end if;
    end if;
 end process;
end arch;
```

## 10.6.3 VHDL and Verilog attributes

In some situations, it is useful to give attributes for some signals in a design to modify how they are synthesized. To do that, you can call the following functions on any signals or memories in the design:

Syntax	Description
addAttribute(name)	Add a boolean attribute with the given name set to true
addAttribute(name, value)	Add a string attribute with the given name set to value

Example:

```
val pcPlus4 = pc + 4
pcPlus4.addAttribute("keep")
```

Produced declaration in VHDL:

```
attribute keep : boolean;
signal pcPlus4 : unsigned(31 downto 0);
attribute keep of pcPlus4: signal is true;
```

Produced declaration in Verilog:

```
(* keep *) wire [31:0] pcPlus4;
```

## 10.7 Introduction

### 10.7.1 Introduction

The core of the language defines the syntax for many features:

- Types / Literals
- Register / Clock domains
- · Component / Area
- RAM / ROM
- When / Switch / Mux
- BlackBox (to integrate VHDL or Verilog IPs inside Spinal)
- SpinalHDL to VHDL converter

Then, by using these features, you can define digital hardware, and also build powerful libraries and abstractions. It's one of the major advantages of SpinalHDL over other commonly used HDLs, because you can extend the language without having knowledge about the compiler.

One good example of this is the SpinalHDL lib which adds many utilities, tools, buses, and methodologies.

To use features introduced in the following chapter you need to  $import\ spinal.core.\_in\ your\ sources.$ 

## **CHAPTER**

# **ELEVEN**

# **LIBRARIES**

## 11.1 Utils

Some utils are also present in *spinal.core* 

## 11.1.1 State less utilities

Syntax	Return	Description
toGray(x : UInt)	Bits	Return the gray value converted from <b>x</b> (UInt)
fromGray(x : Bits)	UInt	Return the UInt value converted value from <b>x</b> (gray)
Reverse(x : T)	T	Flip all bits $(lsb + n \rightarrow msb - n)$
	UInt	Return the index of the single bit set (one hot) in x
OHToUInt(x:		
Seq[Bool])		
OHToUInt(x:		
BitVector)		
	UInt	Return the number of bit set in x
CountOne(x:		
Seq[Bool])		
CountOne(x : BitVector)		
	Bool	Return True if the number of bit set is > x.size / 2
MajarityVata(v.		
MajorityVote(x : Seq[Bool])		
MajorityVote(x:		
BitVector)		
Dit vector)		
EndiannessSwap(that:	T	Big-Endian <-> Little-Endian
T[, base:BitCount])		
OHMasking.first(x :	Bits	Apply a mask on x to only keep the first bit set
Bits)	-	
OHMasking.last(x :	Bits	Apply a mask on x to only keep the last bit set
Bits)	Bits	
	DIIS	
OHMasking.roundRobin(		Apply a mask on x to only keep the bit set from requests.
requests: Bits,		it start looking in requests from the ohPriority position.
ohPriority : Bits		For example if requests is "1001" and ohPriority is "0010",
		the roundRobin function will start looking in <i>requests</i> from its
)		second bit and will return "1000".
	T	December 17 formation 1 1 1 1 2
	T	Returns the muxed T from the inputs based on the oneHot vector.
MuxOH (		
oneHot:		
IndexedSeq[Bool],		
inputs:		
Iterable[T]		
)		

## 11.1.2 State full utilities

Syntax	Return	Description
Delay(that: T, cycleCount: Int)	Т	Return that delayed by cycleCount cycles
History(that: T, length: Int[,when:	List[T]	
Bool])		Return a Vec of length elements
		The first element is that, the last one is that delayed by
		length-1
		The internal shift register sample when when is asserted
BufferCC(input : T)	T	Return the input signal synchronized with the current clock
		domain by using 2 flip flop

#### Counter

The Counter tool can be used to easily instanciate an hardware counter.

Instanciation syntax	Notes
Counter(start: BigInt, end: BigInt[, inc : Bool])	
Counter(range : Range[, inc : Bool])	Compatible with the x to y x until y syntaxes
Counter(stateCount: BigInt[, inc : Bool])	Start at zero and finish at stateCount - 1
Counter(bitCount: BitCount[, inc : Bool])	Start at zero and finish at (1 << bitCount) - 1

There is an example of different syntaxes which could be used with the Counter tool

When a Counter overflow its end value, it restart to its start value.

Note: Currently, only up counter are supported.

#### **Timeout**

The Timeout tool can be used to easily instanciate an hardware timeout.

Instanciation syntax	Notes
Timeout(cycles : BigInt)	Tick after cycles clocks
Timeout(time : TimeNumber)	Tick after a time duration
Timeout(frequency : HertzNumber)	Tick at an frequency rate

There is an example of different syntaxes which could be used with the Counter tool

11.1. Utils 143

**Note:** If you instanciate an Timeout with an time or frequency setup, the implicit ClockDomain should have an frequency setting.

### **ResetCtrl**

The ResetCtrl provide some utilities to manage resets.

## asyncAssertSyncDeassert

You can filter an asynchronous reset by using an asynchronously asserted synchronously deaserted logic. To do it you can use the ResetCtrl.asyncAssertSyncDeassert function which will return you the filtred value.

Argument	Туре	Description
name		
input	Bool	Signal that should be filtered
clockDomain	ClockDomain	ClockDomain which will use the filtered value
inputPolarity	Polarity	HIGH/LOW (default=HIGH)
outputPolarity	Polarity	HIGH/LOW (default=clockDomain.config.resetActiveLevel)
bufferDepth	Int	Number of register stages used to avoid metastability (default=2)

There is also an ResetCtrl.asyncAssertSyncDeassertDrive version of tool which directly assign the clockDomain reset with the filtred value.

## 11.1.3 Special utilities

Syntax	Return	Description
LatencyAnalysis(paths : Node*)	Int	
		Return the shortest path,in therm of cycle, that travel through all nodes, from the first one to the last one

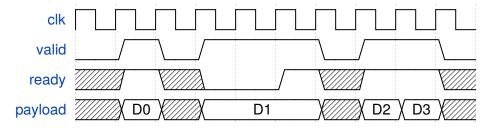
## 11.2 Stream

## 11.2.1 Specification

The Stream interface is a simple handshake protocol to carry payload.

It could be used for example to push and pop elements into a FIFO, send requests to a UART controller, etc.

Sig-	Type	Driver	Description	Don't
nal				care
				when
valid	Bool	Mas-	When high => payload present on the interface	
		ter		
ready	Bool	Slave	When low => transaction are not consumed by the slave	valid
				is
				low
pay-	T	Mas-	Content of the transaction	valid
load		ter		is
				low



There is some examples of usage in SpinalHDL:

```
class StreamFifo[T <: Data](dataType: T, depth: Int) extends Component {
  val io = new Bundle {
    val push = slave Stream (dataType)
    val pop = master Stream (dataType)
  }
  ...
}

class StreamArbiter[T <: Data](dataType: T,portCount: Int) extends Component {
  val io = new Bundle {
    val inputs = Vec(slave Stream (dataType),portCount)
    val output = master Stream (dataType)
  }
  ...
}</pre>
```

**Note:** Each slave can or can't allow the payload to change when valid is high and ready is low. For examples:

- An priority arbiter without lock logic can switch from one input to the other (which will change the payload).
- An UART controller could directly use the write port to drive UART pins and only consume the transaction at the end of the transmission. Be careful with that.

11.2. Stream 145

## 11.2.2 Semantics

When manually reading/driving the signals of a Stream keep in mind that:

- After being asserted, valid may only be deasserted once the current payload was acknowleged. This means valid can only toggle to 0 the cycle after a the slave did a read by asserting ready.
- In contrast to that ready may change at any time.
- A transfer is only done on cycles where both valid and ready are asserted.
- valid of a Stream must not depend on ready in a combinatorial way and any path between the two must be registered.
- It is recommended that valid does not depend on ready at all.

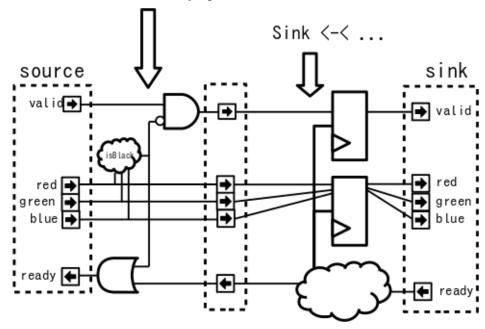
11.2. Stream 147

## 11.2.3 Functions

Syntax	Description	Re-	La-
Stream(type : Data)	Create a Stream of a given type	turn Stream[	tency
master/slave Stream(type : Data)	Create a Stream of a given type	Stream[	
muster/stave stream(type : Butu)		Stream	1
	Create a Stream of a given type		
	Initialized with corresponding in/out setup		
x.fire	Return True when a transaction is consumed	Bool	
	on the bus (valid && ready)		
x.isStall	Return True when a transaction is stall on the bus (valid &&! ready)	Bool	
x.queue(size:Int)	Return a Stream connected to x through a	Stream[	Т2
	FIFO	Stream[	Tl
	Determ a Character deise diberte	Stream	1,
x.m2sPipe()	Return a Stream drived by x		
x.stage()	through a register stage that cut valid/payload paths		
	Cost = (payload width $+ 1$ ) flop flop		
	The state of the s		
x.s2mPipe()		Stream[	ΤØ
	Return a Stream drived by x		
	ready paths is cut by a register stage		
	Cost = payload width * $(mux2 + 1 flip flop)$		
x.halfPipe()		Stream[	Π]
	Return a Stream drived by x		
	valid/ready/payload paths are cut by some		
	register		
	Cost = (payload width + 2) flip flop,		
	bandwidth divided by two		
	Connect y to x		0
x << y			
y >> x			
	Connect y to x through a m2sPipe		1
x <-< y			
y >-> x			
	Connect y to x through a s2mPipe		0
x < y</td <td></td> <td></td> <td></td>			
y >/> x			
V			
			1
x <-/< y	Connect y to x through		
y >/-> x	s2mPipe().m2sPipe()		
	Which imply no combinatorial path		
	between x and y		
x.haltWhen(cond : Bool)		Stream[	ТФ
140	Return a Stream connected to x		
148	Halted when cond is true	ler 11. L	ibraries
x.throwWhen(cond : Bool)		Stream[	T0

The following code will create this logic:

## source.throwWhen(source.payload.isBlack)



```
case class RGB(channelWidth : Int) extends Bundle{
  val red = UInt(channelWidth bit)
  val green = UInt(channelWidth bit)
  val blue = UInt(channelWidth bit)

  def isBlack : Bool = red === 0 && green === 0 && blue === 0
}

val source = Stream(RGB(8))
val sink = Stream(RGB(8))
sink <--< source.throwWhen(source.payload.isBlack)</pre>
```

## 11.2.4 Utils

There is many utils that you can use in your design in conjunction with the Stream bus, this chapter will document them.

### **StreamFifo**

On each stream you can call the .queue(size) to get a buffered stream. But you can also instantiate the FIFO component itself:

```
val streamA, streamB = Stream(Bits(8 bits))
//...
val myFifo = StreamFifo(
  dataType = Bits(8 bits),
  depth = 128
)
myFifo.io.push << streamA
myFifo.io.pop >> streamB
```

11.2. Stream 149

parameter name	Туре	Description
dataType	T	Payload data type
depth	Int	Size of the memory used to store elements

io	Type	Description
name		
push	Stream[T]	Used to push elements
pop	Stream[T]	Used to pop elements
flush	Bool	Used to remove all elements inside the FIFO
occu-	UInt of log2Up(depth + 1) bits	Indicate the internal memory occupancy
pancy		

## **StreamFifoCC**

You can instanciate the dual clock domain version of the fifo the following way:

```
val clockA = ClockDomain(???)
val clockB = ClockDomain(???)
val streamA, streamB = Stream(Bits(8 bits))
//...
val myFifo = StreamFifoCC(
  dataType = Bits(8 bits),
  depth = 128,
  pushClock = clockA,
  popClock = clockB
)
myFifo.io.push << streamA
myFifo.io.pop >> streamB
```

parameter name	Туре	Description
dataType	T	Payload data type
depth	Int	Size of the memory used to store elements
pushClock	ClockDomain	Clock domain used by the push side
popClock	ClockDomain	Clock domain used by the pop side

io	Туре	Description
name		
push	Stream[T]	Used to push elements
pop	Stream[T]	Used to pop elements
pushOc-	UInt of log2Up(depth + 1) bits	Indicate the internal memory occupancy (from the
cu-		push side perspective)
pancy		
popOc-	UInt of log2Up(depth + 1) bits	Indicate the internal memory occupancy (from the pop
cu-		side perspective)
pancy		

### StreamCCByToggle

Component that connects Streams across clock domains based on toggling signals.

This way of implementing a cross clock domain bridge is characterized by a small area usage but also a low bandwidth.

```
val clockA = ClockDomain(???)
val clockB = ClockDomain(???)
val streamA, streamB = Stream(Bits(8 bits))
//...
val bridge = StreamCCByToggle(
  dataType = Bits(8 bits),
  inputClock = clockA,
  outputClock = clockB
)
bridge.io.input << streamA
bridge.io.output >> streamB
```

parameter name	Туре	Description
dataType	T	Payload data type
inputClock	ClockDomain	Clock domain used by the push side
outputClock	ClockDomain	Clock domain used by the pop side

io name	Type	Description
input	Stream[T]	Used to push elements
output	Stream[T]	Used to pop elements

Alternatively you can also use a this shorter syntax which directly return you the cross clocked stream:

### **StreamArbiter**

When you have multiple Streams and you want to arbitrate them to drive a single one, you can use the StreamArbiterFactory.

11.2. Stream 151

Arbitration	Description
functions	
lowerFirst	Lower port have priority over higher port
roundRobin	Fair round robin arbitration
sequen-	
tialOrder	Could be used to retrieve transaction in a sequancial order First transaction should come from port zero, then from port one,

Lock func-	Description
tions	
noLock	The port selection could change every cycle, even if the transaction on the selected port is
	not consumed.
transaction-	The port selection is locked until the transaction on the selected port is consumed.
Lock	
fragmentLock	
	Could be used to arbitrate Stream[Flow[T]].
	In this mode, the port selection is locked until the selected port finish is burst (last=True).

Generation functions	Return
on(inputs : Seq[Stream[T]])	Stream[T]
onArgs(inputs : Stream[T]*)	Stream[T]

### **StreamJoin**

This utile takes multiple input streams and wait until all of them fire before letting all of them through.

```
val cmdJoin = Stream(Cmd())
cmdJoin.arbitrationFrom(StreamJoin.arg(cmdABuffer, cmdBBuffer))
```

### **StreamFork**

A StreamFork will clone each incoming data to all its output streams. If synchronous is true, all output streams will always fire together, which means that the stream will halt until all output streams are ready. If synchronous is false, output streams may be ready one at a time, at the cost of an additional flip flop (1 bit per output). The input stream will block until all output streams have processed each item regardlessly.

```
val inputStream = Stream(Bits(8 bits))
val (outputstream1, outputstream2) = StreamFork2(inputStream, synchronous=false)
```

or

```
val inputStream = Stream(Bits(8 bits))
val outputStreams = StreamFork(inputStream,portCount=2, synchronous=true)
```

## StreamDispatcherSequencial

This util take its input stream and routes it to outputCount stream in a sequential order.

```
val inputStream = Stream(Bits(8 bits))
val dispatchedStreams = StreamDispatcherSequencial(
  input = inputStream,
  outputCount = 3
)
```

## 11.3 Flow

## 11.3.1 Specification

The Flow interface is a simple valid/payload protocol which mean the slave can't halt the bus. It could be used, for example, to represent data coming from an UART controller, requests to write an on-chip memory, etc.

Sig-	Type	Driver	Description	Don't
nal				care
				when
valid	Bool	Mas-	When high => payload present on the interface	
		ter		
pay-	T	Mas-	Content of the transaction	valid
load		ter		is
				low

11.3. Flow 153

## 11.3.2 Functions

Syn-	Description	Re-	La-
tax		turn	tency
Flow(t	yp€reate a Flow of a given type	Flow[T	]
:			
Data)			
mas-		Flow[T	]
ter/slav	e Create a Flow of a given type		
Flow(t	/person a rion of a given type Initialized with corresponding in/out setup		
:	midalized with corresponding in/out setup		
Data)		-	
x.m2sF	ipe()	Flow[T	]] 1
	Return a Flow drived by x		
	through a register stage that cut valid/payload paths		
	Connect y to x		0
x <<			
у			
y >>			
X			
	Connect y to x through a m2sPipe		1
x <-<			
у			
y >->			
X			
x.throw	When(cond	Flow[T	] 0
:	Return a Flow connected to x		
Bool)			
	When cond is high, transaction are dropped		
v toRe	(Return a register which is loaded with payload when valid is high	Т	
A.WICE	pyretarn a register which is loaded with pay road when valid is high	1	

## 11.4 Fragment

## 11.4.1 Specification

The Fragment bundle is the concept of transmitting a "big" thing by using multiple "small" fragments. For examples:

- $\bullet \ \ A \ picture \ transmitted \ with \ width *height \ transaction \ on \ a \ {\tt Stream[Fragment[Pixel]]}$
- An UART packet received from an controller without flow control could be transmitted on a Flow[Fragment[Bits]]
- An AXI read burst could be carried by an Stream[Fragment[AxiReadResponse]]

Signals defined by the Fragment bundle are:

Signal	Type	Driver	Description
fragment	T	Master	The "payload" of the current transaction
last	Bool	Master	High when the fragment is the last of the current packet

As you can see with this specification and precedent example, the Fragment concept doesn't specify how transaction are transmitted (You can use Stream,Flow or any other communication protocol). It only add enough information (last) to know if the current transaction is the first one, the last one or one in the middle of a given packet.

**Note:** The protocol didn't carry a 'first' bit because it can be generated at any place by doing 'RegNextWhen(bus.last, bus.fire) init(True)'

## 11.4.2 Functions

For Stream[Fragment[T]] and Flow[Fragment[T]], following function are presents:

		Description
tax	turr	1
x.fi	rstBoo	l Return True when the next or the current transaction is/would be the first of a packet
x.ta	ilBoo	l Return True when the next or the current transaction is/would be not the first of a packet
x.is	FiBsto	l Return True when an transaction is present and is the first of a packet
x.is	T <b>a</b> bo	l Return True when an transaction is present and is the not the first/last of a packet
x.is	LBsoc	l Return True when an transaction is present and is the last of a packet

For Stream[Fragment[T]], following function are also accessible:

Syntax	Return	Description
x.insertHeader(header : T)	Stream[Fragment[T]]	Add the header to each packet on
		x and return the resulting bus

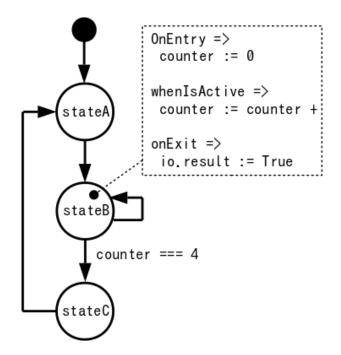
## 11.5 State machine

## 11.5.1 Introduction

In SpinalHDL you can define your state machine like in VHDL/Verilog, by using enumerations and switch cases statements. But in SpinalHDL you can also use a dedicated syntax.

The following state machine is implemented in following examples :

11.5. State machine 155



### Style A:

```
import spinal.lib.fsm._
class TopLevel extends Component {
 val io = new Bundle{
   val result = out Bool()
 val fsm = new StateMachine{
   val counter = Reg(UInt(8 bits)) init (0)
   io.result := False
   val stateA : State = new State with EntryPoint{
     whenIsActive (goto(stateB))
   val stateB : State = new State{
     onEntry(counter := 0)
     whenIsActive {
        counter := counter + 1
        when(counter === 4){
          goto(stateC)
     onExit(io.result := True)
   val stateC : State = new State{
     whenIsActive (goto(stateA))
   }
 }
```

## Style B:

```
import spinal.lib.fsm._
class TopLevel extends Component {
```

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```
val io = new Bundle{
    val result = out Bool()
  }
  val fsm = new StateMachine{
    val stateA = new State with EntryPoint
    val stateB = new State
    val stateC = new State
    val counter = Reg(UInt(8 bits)) init (0)
    io.result := False
    stateA
      .whenIsActive (goto(stateB))
    stateB
      .onEntry(counter := 0)
      .whenIsActive {
        counter := counter + 1
        when(counter === 4){
          goto(stateC)
        }
      }
      .onExit(io.result := True)
   stateC
      .whenIsActive (goto(stateA))
  }
}
```

## 11.5.2 StateMachine

StateMachine is the base class that will manage the logic of your FSM.

```
val myFsm = new StateMachine{
   // Here will come states definition
}
```

The StateMachine class also provide some utils :

Name	Return	Description
isAc-	Bool	Return True when the state machine is in the given state
tive(state)		
isEnter-	Bool	Return True when the state machine is entering the given state
ing(state)		

11.5. State machine

## 11.5.3 States

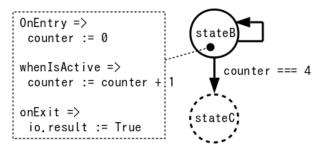
There is multiple kinds of states that you can use.

- State (the base one)
- StateDelay
- StateFsm
- StateParallelFsm

In each of them you have access the following utilities:

Name	Description
	yourStatements is executed the cycle before entering the state
onEntry{	
	urStatements
}	
	yourStatements is executed when the state machine will be in another state the next cycle
onExit{	
	on Ctatamanta
1, 1	urStatements
}	
	yourStatements is executed when the state machine is in the state
whenIsA	ctive{
yo	urStatements
}	
	vormetestaments is avacuated when the state mechine will be in the state the most evale
	yourStatements is executed when the state machine will be in the state the next cycle
whenIsN	ext{
yo	urStatements
}	
	State the state of the state machine by nextState
exit()	Set the state of the state machine to the boot one

For example, the following state could be defined in SpinalHDL by using the following syntax :



```
val stateB : State = new State{
  onEntry(counter := 0)
  whenIsActive {
    counter := counter + 1
    when(counter === 4) {
      goto(stateC)
    }
  }
  onExit(io.result := True)
}
```

You can also define your state as the entry point of the state machine by extends the EntryPoint trait.

```
val stateA: State = new State with EntryPoint {
  whenIsActive {
    goto(stateB)
  }
}
```

### **StateDelay**

StateDelay allow you to create a state which wait a fixed number of cycles before executing statments in your  $whenCompleted{...}$ . The standard way to write it is:

```
val stateG : State = new StateDelay(cyclesCount=40){
  whenCompleted{
    goto(stateH)
  }
}
```

But you can also write it like that:

```
val stateG : State = new StateDelay(40){whenCompleted(goto(stateH))}
```

### **StateFsm**

StateFsm Allow you to describe a state which contains a nested state machine. When the nested state machine is done, your statments in whenCompleted{...} are executed.

There is an example of StateFsm definition:

```
val stateC = new StateFsm(fsm=internalFsm()){
  whenCompleted{
    goto(stateD)
  }
}
```

As you can see in the precedent code, it use a internalFsm function to create the inner state machine. There is an example of definition bellow :

```
def internalFsm() = new StateMachine {
  val counter = Reg(UInt(8 bits)) init (0)

  val stateA: State = new State with EntryPoint {
    whenIsActive {
      goto(stateB)
```

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11.5. State machine 159

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In the precedent example, the exit() call will make the state machine jump to the boot state (a internal hidden state). This notify the StateFsm about the completion of the inner state machine.

#### **StateParallelFsm**

This state is able to handle multiple nested state machines. When all nested state machine are done, your statments in whenCompleted{...} are executed.

There is an example of declaration:

```
val stateD = new StateParallelFsm (internalFsmA(), internalFsmB()){
  whenCompleted{
    goto(stateE)
  }
}
```

### Notes about the entry state

By default, you have to specify the entry state of your state-machine, but between the reset and the first clock sampling, the state-machine will be in a bootState. It is only after the first clock sampling that your entry state become active.

This allow to properly the onEntry of your entry state, and also allow the implementation of the inner state-machines.

While it is usefull, it is also possible to bypass that feature and directly have your state-machine starting into a user state.

Do do so, you can use *makeInstantEntry*. This will return you the state which will be active directly after reset. Note, the onEntry of that state will only be called when it transition from another state to this state.

There is an example:

```
// State sequance : IDLE, STATE_A, STATE_B, ...
val fsm = new StateMachine{
  val IDLE = makeInstantEntry()
  val STATE_A, STATE_B, STATE_C = new State

  IDLE.whenIsActive(goto(STATE_A))
  STATE_A.whenIsActive(goto(STATE_B))
  STATE_B.whenIsActive(goto(STATE_C))
  STATE_C.whenIsActive(goto(STATE_B))
}
```

```
// State sequance : stateBoot, STATE_A, STATE_B, ...
val fsm = new StateMachine{
  val STATE_A, STATE_B, STATE_C = new State
  setEntry(STATE_A)

STATE_A.whenIsActive(goto(STATE_B))
  STATE_B.whenIsActive(goto(STATE_C))
  STATE_C.whenIsActive(goto(STATE_B))
}
```

## 11.6 VexRiscv (RV32IM CPU)

VexRiscv is an fpga friendly RISC-V ISA CPU implementation with following features :

- RV32IM instruction set
- Pipelined on 5 stages (Fetch, Decode, Execute, Memory, WriteBack)
- 1.44 DMIPS/Mhz when all features are enabled
- · Optimized for FPGA
- Optional MUL/DIV extension
- · Optional instruction and data caches
- Optional MMU
- Optional debug extension allowing eclipse debugging via an GDB >> openOCD >> JTAG connection
- Optional interrupts and exception handling with the Machine and the User mode from the riscv-privileged-v1.9.1 spec.
- Two implementation of shift instructions, Single cycle / shiftNumber cycles
- · Each stage could have bypass or interlock hazard logic
- FreeRTOS port https://github.com/Dolu1990/FreeRTOS-RISCV

Much more information there: https://github.com/SpinalHDL/VexRiscv

## 11.7 Bus Slave Factory

#### 11.7.1 Introduction

In many situation it's needed to implement a bus register bank. The BusSlaveFactory is a tool that provide an abstract and smooth way to define them.

To see capabilities of the tool, an simple example use the Apb3SlaveFactory variation to implement an *memory mapped UART*. There is also another example with an *Timer* which contain a memory mapping function.

You can find more documentation about the internal implementation of the BusSlaveFactory tool there

## 11.7.2 Functionality

Currently there is three implementation of the BusSlaveFactory tool: APB3, AXI-lite 3 and Avalon. Each implementation of that tool take as argument one instance of the corresponding bus and then offer following functions to map your hardware into the memory mapping:

Name	Re- turn	Description
busDataW- idth	Int	Return the data width of the bus
	e hit∩ffe	et)When the bus read the address, fill the response with that at bitOffset
·		seWhen the bus write the address, assign that with bus's data from bitOffset
	55,011011	Call doThat when a write transaction occur on address
on-	1.70	Can do that when a write transaction occur on address
Write(address)(	doinat)	
on-		Call doThat when a read transaction occur on address
Read(address)(	loThat)	
nonStop-		Permanently assign that by the bus write data from bitOffset
Write(that,bitO	ffset)	
readAnd-		Make that readable and writable at address and placed at bitOffset in the word
Write(that,addre	ess,bitOf	•
readMulti-	<u> </u>	,
Word(that,addre	(225	
Word (mar, addi)	033)	Create the memory mapping to read that from 'address'.
		If that is bigger than one word it extends the register on followings addresses
		in chart to engger than one word it entends the register on renewings that each
writeMulti-		
Word(that,addre	188)	Create the memory mapping to write that at 'address'.
		If that is bigger than one word it extends the register on followings addresses
		if that is bigger than one word it extends the register on followings addresses
240	Т	Charte a white only register of type data Type at address and placed at hit Office
cre-	_	Create a write only register of type dataType at address and placed at bitOffset
		didrehe, bit Odfset)
createRead-	T	Create a read write register of type dataType at address and placed at bitOffset
Write(dataType		
create-	Flow[T	Create a writable Flow register of type dataType at address and placed at
AndDrive-		bitOffset in the word
Flow(dataType,	address,t	itOffset)
		seDrive that with a register writable at address placed at bitOffset in the word
driveAn-		Drive that with a register writable and readable at address placed at bitOffset
dRead(that,add	ress hit∩t	
drive-	C38,011O1	
	1::00	Emit on that a transaction when a write happen at address by using data placed at
Flow(tnat,addre	SS,DITOII	selb)itOffset in the word
readStreamNon	Blocking	Read that and consume the transaction when a read happen at address.  valid <= validBitOffset bit
		payload <= payloadBitOffset+widthOf(payload) downto payloadBitOffset
address,		payroad + payroadziroinot + midii o 1 (payroad) do milio pay zodaz z co z zod
validBitC	offset,	
payloadB		
раугоаць	itonset)	
doBitsAccumul	ationAnd	Instanciate an internal register which at each cycle do : ClearOnRead(that, reg := reg   that
address,		Then when a read occur, the register is cleared. This register is readable at address and placed at bitOffset in the word
		and placed at DICOIISEC III die word
bitOffset)	)	
	_	

## 11.8 Fiber framework

Currently in developpement.

The Fiber to run the hardware elaboration in a out of order manner, a bit similarly to Makefile, where you can define rules and dependencies which will then be solved when you run a make command. It is very similar to the Scala Future feature.

Such framework complexify simple things but provide some strong feature for complex cases:

- You can define things before even knowing all their requirements, ex: instanciating a interruption controller, before knowing how many lines of interrupt you need
- Abstract/lazy/partial SoC architecture definition allowing the creation of SoC template for further specialisations
- Automatic requirements negotiation between multiple agents in a decentralized way, ex : between masters and slaves of a memory bus

The framework is mainly composed of:

- Handle[T], which can be used later to store a value of type T.
- handle.load which allow to set the value of a handle (will reschedule all tasks waiting on it)
- handle.get, which return the value of the given handle. Will block the task execution if that handle isn't loaded yet
- Handle{ code }, which fork a new task which will execute the given code. The result of that code will be loaded into the Handle
- soon(handle), which allow the current task to announce that soon it will load that handle with a value (used to track which handle will

Warning, this is realy not usual RTL description and aim large system generation. It is currently used as toplevel integration tool in SaxonSoC.

## 11.8.1 Simple dummy example

There is a simple example :

```
import spinal.core.fiber._
// Create two empty Handles
val a, b = Handle[Int]

// Create a Handle which will be loaded asynchronously by the given body result
val calculator = Handle {
    a.get + b.get // .get will block until they are loaded
}

// Same as above
val printer = Handle {
    println(s"a + b = ${calculator.get}") // .get is blocking until the calculator...
    body is done
}

// Synchronously load a and b, this will unblock a.get and b.get
a.load(3)
b.load(4)
```

Its runtime will be:

- · create a and b
- fork the calculator task, but is blocked when executing a.get
- fork the printer task, but is blocked when executing calculator.get
- load a and b, which reschedule the calculator task (as it was waiting on a)
- calculator do its a + b sum, and load its Handle with that result, which reschedule the printer task
- printer task print its stuff
- · everything done

So, the main point of that example is to show that we kind of overcome the sequential execution of things, as a and b are loaded after the definition of the calculator.

## 11.8.2 Handle[T]

Handle[T] are a bit like scala's Future[T], they allow to talk about something before it is even existing, and wait on it.

### soon(handle)

In order to maintain a proper graph of dependencies between tasks and Handle, a task can specify in advance that it will load a given handle. This is very usefull in case of a generation starvation/deadlock for SpinalHDL to report accuratly where is the issue.

### 11.9 Bus

### 11.9.1 AHB-Lite3

## **Configuration and instanciation**

First each time you want to create a AHB-Lite3 bus, you will need a configuration object. This configuration object is an AhbLite3Config and has following arguments:

Parameter name	Type	Default	Description
addressWidth	Int		Width of HADDR (byte granularity)
dataWidth	Int		Width of HWDATA and HRDATA

There is in short how the AHB-Lite3 bus is defined in the SpinalHDL library :

```
case class AhbLite3(config: AhbLite3Config) extends Bundle with IMasterSlave{
    // Address and control
    val HADDR = UInt(config.addressWidth bits)
    val HSEL = Bool()
    val HREADY = Bool()
    val HWRITE = Bool()
```

(continues on next page)

11.9. Bus 165

(continued from previous page)

```
val HSIZE = Bits(3 bits)
val HBURST = Bits(4 bits)
val HPROT = Bits(4 bits)
val HTANS = Bits(2 bits)
val HMASTLOCK = Bool()

// Data
val HWDATA = Bits(config.dataWidth bits)
val HRDATA = Bits(config.dataWidth bits)

// Transfer response
val HREADYOUT = Bool()
val HRESP = Bool()

override def asMaster(): Unit = {
   out(HADDR,HWRITE,HSIZE,HBURST,HPROT,HTRANS,HMASTLOCK,HWDATA,HREADY,HSEL)
   in(HREADYOUT,HRESP,HRDATA)
}
```

There is a short example of usage:

```
val ahbConfig = AhbLite3Config(
  addressWidth = 12,
  dataWidth = 32
)
val ahbX = AhbLite3(ahbConfig)
val ahbY = AhbLite3(ahbConfig)
when(ahbY.HSEL){
  //...
}
```

### **Variations**

There is an AhbLite3Master variation. The only difference is the absence of the HREADYOUT signal. This variation should only be used by masters while the interconnect and slaves use AhbLite3.

## 11.9.2 Apb3

### Introduction

The AMBA3-APB bus is commonly used to interface low bandwidth peripherals.

### **Configuration and instanciation**

First each time you want to create a APB3 bus, you will need a configuration object. This configuration object is an Apb3Config and has following arguments:

Parameter name Type		Default	Description
addressWidth	Int		Width of PADDR (byte granularity)
dataWidth	Int		Width of PWDATA and PRDATA
selWidth	Int	1	With of PSEL
useSlaveError	Boolean	false	Specify the presence of PSLVERROR

There is in short how the APB3 bus is defined in the SpinalHDL library:

There is a short example of usage:

```
val apbConfig = Apb3Config(
  addressWidth = 12,
  dataWidth = 32
)
val apbX = Apb3(apbConfig)
val apbY = Apb3(apbConfig)
when(apbY.PENABLE){
   //...
}
```

## **Functions and operators**

Name	Return	Description
X >> Y		Connect X to Y. Address of Y could be smaller than the one of X
X << Y		Do the reverse of the >> operator

### 11.9.3 Axi4

### Introduction

The AXI4 is a high bandwidth bus defined by ARM.

11.9. Bus 167

### **Configuration and instanciation**

First each time you want to create a AXI4 bus, you will need a configuration object. This configuration object is an Axi4Config and has following arguments:

Note: useXXX specify if the bus has XXX signal present.

Parameter name	Type	Default
addressWidth	Int	
dataWidth	Int	
idWidth	Int	
userWidth	Int	
useId	Boolean	true
useRegion	Boolean	true
useBurst	Boolean	true
useLock	Boolean	true
useCache	Boolean	true
useSize	Boolean	true
useQos	Boolean	true
useLen	Boolean	true
useLast	Boolean	true
useResp	Boolean	true
useProt	Boolean	true
useStrb	Boolean	true
useUser	Boolean	false

There is in short how the AXI4 bus is defined in the SpinalHDL library:

```
case class Axi4(config: Axi4Config) extends Bundle with IMasterSlave{
  val aw = Stream(Axi4Aw(config))
  val w = Stream(Axi4W(config))
  val b = Stream(Axi4B(config))
  val ar = Stream(Axi4Ar(config))
  val r = Stream(Axi4R(config))

  override def asMaster(): Unit = {
    master(ar,aw,w)
    slave(r,b)
  }
}
```

There is a short example of usage:

```
val axiConfig = Axi4Config(
  addressWidth = 32,
  dataWidth = 32,
  idWidth = 4
)
val axiX = Axi4(axiConfig)
val axiY = Axi4(axiConfig)
when(axiY.aw.valid) {
    //...
}
```

### **Variations**

There is 3 other variation of the Axi4 bus:

Туре	Description
Axi4ReadOnly	Only AR and R channels are present
Axi4WriteOnly	Only AW, W and B channels are present
Axi4Shared	
	This variation is a library initiative.  It use 4 channels, W, B, R and also a new one which is named AWR.  The AWR channel can be used to transmit AR and AW transactions. To dissociate them, a signal write is present.  The advantage of this Axi4Shared variation is to use less area, especially in the
	interconnect.

### **Functions and operators**

Name	Return	Description
X >> Y		Connect X to Y. Able infer default values as specified in the AXI4 specification,
		and also to adapt some width in a safe manner.
X << Y		Do the reverse of the >> operator
X.toWriteOnly	Axi4WriteOnl	Return an Axi4WriteOnly bus drive by X
X.toReadOnly	Axi4ReadOnly	Return an Axi4ReadOnly bus drive by X

## 11.9.4 AvalonMM

### Introduction

The AvalonMM bus fit very well in FPGA. It is very flexible:

- Able of the same simplicity than APB
- Better for than AHB in many application that need bandwidth because AvalonMM has a mode that decouple read response from commands (reduce latency read latency impact).
- Less performance than AXI but use much less area (Read and write command use the same handshake channel. The master don't need to store address of pending request to avoid Read/Write hazard)

### Configuration and instanciation

The AvalonMM Bundle has a construction argument AvalonMMConfig. Because of the flexible nature of the Avalon bus, the AvalonMMConfig as many configuration elements. For more information the Avalon spec could be find there.

(continues on next page)

11.9. Bus 169

(continued from previous page)

```
useWaitRequestn : Boolean,
useReadDataValid : Boolean,
useBurstCount : Boolean,
//useEndOfPacket : Boolean,
addressUnits : AddressUnits = symbols,
burstCountUnits : AddressUnits = words,
burstOnBurstBoundariesOnly : Boolean = false,
constantBurstBehavior : Boolean = false,
holdTime : Int = 0,
linewrapBursts : Boolean = false,
maximumPendingReadTransactions : Int = 1,
maximumPendingWriteTransactions : Int = 0, // unlimited
readLatency : Int = 0,
readWaitTime : Int = 0,
setupTime : Int = 0,
writeWaitTime : Int = 0
)
```

This configuration class has also some functions:

Name	Return	Description
getReadOn-	AvalonMM-	Return a similar configuration but with all write feature disabled
lyConfig	Config	
getWriteOn-	AvalonMM-	Return a similar configuration but with all read feature disabled
lyConfig	Config	

This configuration companion object has also some functions to provide some AvalonMMConfig templates:

Name	Return	Description
fixed(addressWid	th <b>AdataWWMIM</b> i,readI	aRectuyn a simple configuration with fixed read timings
	Config	
pipelined(address	<b>WaidalhodaxaW</b> idth)	Return a configuration with variable latency read (readDataValid)
	Config	
bursted(addressW	/iaAthadlataWWithth,bu	rs <b>Redurnt Width</b> figuration with variable latency read and burst capabilities
	Config	

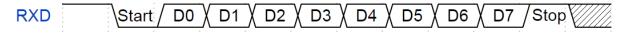
## 11.10 Com

### 11.10.1 UART

### Introduction

The UART protocol could be used, for instance, to emit and receive RS232 / RS485 frames.

There is an example of an 8 bits frame, with no parity and one stop bit :



### **Bus definition**

```
case class Uart() extends Bundle with IMasterSlave {
  val txd = Bool() // Used to emit frames
  val rxd = Bool() // Used to receive frames

  override def asMaster(): Unit = {
    out(txd)
    in(rxd)
  }
}
```

### **UartCtrl**

An Uart controller is implemented in the library. This controller has the specificity to use a sampling window to read the rxd pin and then to using an majority vote to filter its value.

Ю	di-	type	Description
name	rec-		
	tion		
con-	in	UartC-	Used to set the clock divider/parity/stop/data length of the controller
fig		trl-	
		Con-	
		fig	
write	slave	Stream	[Bits]am port used to request a frame transmission
read	mas-	Flow[B	its]ow port used to receive decoded frames
	ter		
uart	mas-	Uart	Interface to the real world
	ter		

The controller could be instantiated via an  ${\tt UartCtrlGenerics}$  configuration object :

11.10. Com 171

At-	type	Description
tribute		
dataW-	Int	Maximal number of bit inside a frame
idth-		
Max		
clock-	Int	Width of the internal clock divider
Di-		
vider-		
Width		
pre-	Int	Specify how many samplingTick are drop at the beginning of a UART baud
Sam-		
pling-		
Size		
sam-	Int	Specify how many samplingTick are used to sample rxd values in the middle of the UART
pling-		baud
Size		
post-	Int	Specify how many samplingTick are drop at the end of a UART baud
Sam-		
pling-		
Size		

## 11.10.2 USB device

### Introduction

There is a USB device controller in the SpinalHDL library. In a few bullet points it can be resumed to:

- Implemented to allow a CPU to configure and manage the endpoints
- A internal ram which store the endpoints states and transactions descriptors
- Up to 16 endpoints (for virtualy no price)
- Support USB host full speed (12Mbps)
- Test on linux using its own driver (https://github.com/SpinalHDL/linux/blob/dev/drivers/usb/gadget/udc/spinal\_udc.c)
- Bmb memory interace for the configuration
- Require a clock for the internal phy which is a multiple of 12 Mhz at least 48 Mhz
- The controller frequency is not restricted
- No external phy required

Linux gadget tested and functional:

- Serial connection
- Ethernet connection
- Mass storage (~8 Mbps on ArtyA7 linux)

### Deployments:

- https://github.com/SpinalHDL/SaxonSoc/tree/dev-0.3/bsp/digilent/ArtyA7SmpLinux
- https://github.com/SpinalHDL/SaxonSoc/tree/dev-0.3/bsp/radiona/ulx3s/smp

### **Architecture**

The controller is composed of:

- A few control registers
- A internal ram used to store the endpoint status, the transfer descriptors and the endpoint 0 SETUP data.

A linked list of descriptors for each endpoint in order to handle of the USB IN/OUT transactions and data.

The endpoint 0 manage the IN/OUT transactions like all the other endpoints but has some additional hardware to manage the SETUP transactions :

- Its linked list is cleared on each setup transactions
- The data of the SETUP transaction are stored in a fixed location (SETUP\_DATA)
- It has a specific interrupt flag for SETUP transactions

## **Registers**

## FRAME (0xFF00)

Name	Type	Bits	Description
usbFrameId	RO	31-0	Current usb frame id

## ADDRESS (0xFF04)

Name	Type	Bits	Description
usbAd-	WO	6-0	The device will only listen at tokens with the specified address This field is auto-
dressFilter			maticaly cleared on usb reset events

### **INTERRUPT (0xFF08)**

All bits of this register can be cleared by writing '1' in them.

Name	Type	Bits	Description
endpoints	RC	15-0	Raised when a enpoint generate a interrupt
reset	RC	16	Raised when a USB reset appeared
ep0Setup	RC	17	Raised when endpoint 0 receive a setup transaction
suspend	RC	18	Raised when a USB suspend appeared
resume	RC	19	Raised when a USB resume appeared
disconnect	RC	20	Raised when a USB disconnect appeared

## HALT (0xFF0C)

This register allow to place a single enpoint in a dormant state in order to ensure atomicity of CPU operations. The peripheral will return NAK if the given endpoint is addressed by the usb host.

Name	Type	Bits	Description
endpointId	WO	3-0	The endpoint you want to put in sleep
enable	WO	4	
effective en-	RO	5	After setting the enable, you need to wait for this bit to be set by the hardware
able			itself to ensure atomicity

11.10. Com 173

## CONFIG (0xFF10)

Name	Type	Bits	Description
pullupSet	SO	0	Write '1' to enable the USB device pullup on the dp pin
pullupClear	SO	1	
interruptEnableSet	SO	2	Write '1' to let the present and future interrupt happening
interruptEnableClear	SO	3	

## INFO (0xFF14)

- 1				Description
	ramSize	RO	3-0	The internal ram will have (1 << this) bytes

## **ENDPOINTS (0x0000 - 0x003F)**

The endpoints status are stored at the beginning of the internal ram over one 32 bits word each.

Name	Туре	Bits	Description	
enable	RW	0	If not set, the endpoint will ignore all the trafic	
stall	RW	1	If set, the endpoint will always return STALL status	
nack	RW	2	If set, the endpoint will always return NACK status	
dataPhase	RW	3	Specify the IN/OUT data PID used. '0' => DATA0. This field is also updated	
			by the controller.	
head	RW	15-4	Specify the current descriptor head (linked list). 0 => empty list, byte address =	
			this << 4	
isochronous	RW	16		
maxPacket-	RW	31-		
Size		22		

## **SETUP\_DATA** (0x0040 - 0x0047)

When endpoint 0 receive a SETUP transaction, the data of the transaction will be stored at that place.

# **Descriptors**

Descriptors allows to specify how a endpoint need to handle the data phase of IN/OUT transactions. They are stored in the internal ram, can be linked together via their linked lists and need to be aligned on 16 bytes boundaries

Name	Word Bits		Description	
offset	0	15-	Specify the current progress in the transfer (in byte)	
		0		
code	0	19-	$0xF \Rightarrow \text{in progress}, 0x0 \Rightarrow \text{success}$	
		16		
next	1	15-	Point the next descriptor $0 \Rightarrow$ nothing, byte address = this $<< 4$	
		4		
length	1	31-	Number of bytes allocated for the data field	
		16		
direc-	2	16	'0' => OUT, '1' => IN	
tion				
inter-	2	17	If set, the completion of the descriptor will generate a interrupt.	
rupt				
com-	2	18	Normaly, a descriptor completion only occure when a USB transfer is smaller than the	
ple-			maxPacketSize. But if this field is set, then when the descriptor become full is also a	
tionOn-			considered as a completion event. (offset == length)	
Full				
data1OnCo2mpletib9		ti <b>d</b> A	force the endpoint dataPhase to DATA1 on the completion of the descriptoo	
data				

# **Usage**

```
import spinal.core._
import spinal.core.sim._
import spinal lib bus bmb BmbParameter
import spinal.lib.com.usb.phy.UsbDevicePhyNative
import spinal.lib.com.usb.sim.UsbLsFsPhyAbstractIoAgent
import spinal.lib.com.usb.udc.{UsbDeviceCtrl, UsbDeviceCtrlParameter}
case class UsbDeviceTop() extends Component {
 val ctrlCd = ClockDomain.external("ctrlCd", frequency = FixedFrequency(100 MHz))
 val phyCd = ClockDomain.external("phyCd", frequency = FixedFrequency(48 MHz))
 val ctrl = ctrlCd on new UsbDeviceCtrl(
   p = UsbDeviceCtrlParameter(
     addressWidth = 14
    bmbParameter = BmbParameter(
      addressWidth = UsbDeviceCtrl.ctrlAddressWidth,
      dataWidth = 32,
      sourceWidth = 0,
      contextWidth = 0,
      lengthWidth = 2
   )
 )
 val phy = phyCd on new UsbDevicePhyNative(sim = true)
 ctrl.io.phy.cc(ctrlCd, phyCd) <> phy.io.ctrl
```

(continues on next page)

11.10. Com 175

```
val bmb = ctrl.io.ctrl.toIo()
val usb = phy.io.usb.toIo()
val power = phy.io.power.toIo()
val pullup = phy.io.pullup.toIo()
val interrupts = ctrl.io.interrupt.toIo()
}

object UsbDeviceGen extends App{
   SpinalVerilog(new UsbDeviceTop())
}
```

# 11.10.3 USB OHCI

#### Introduction

There is a USB OHCi controller (host) in the SpinalHDL library. In a few bullet points it can be resumed to:

- It follow the *OpenHCI Open Host Controller Interface Specification for USB* specification (OHCI).
- It is compatible with the upstream linux / uboot OHCI drivers already. (there is also a OHCI driver on tinyUSB)
- This provide USB host full speed and low speed capabilities (12Mbps and 1.5Mbps)
- · Tested on linux and uboot
- One controller can host multiple ports (up to 16)
- Bmb memory interface for DMA accesses
- Bmb memory interace for the configuration
- Require a clock for the internal phy which is a multiple of 12 Mhz at least 48 Mhz
- The controller frequency is not restricted
- No external phy required

#### Devices tested and functional:

- Mass storage (~8 Mbps on ArtyA7 linux)
- · Keyboard / Mouse
- Audio output
- Hub

#### Limitations:

- Some USB hub (had one so far) do not like having a full speed host with low speed devices attached.
- Some modern devices will not work on USB full speed (ex : Gbps ethernet adapter)
- Require memory coherency with the CPU (or the cpu need to flush his data cache in the driver)

#### Deployments:

- https://github.com/SpinalHDL/SaxonSoc/tree/dev-0.3/bsp/digilent/ArtyA7SmpLinux
- https://github.com/SpinalHDL/SaxonSoc/tree/dev-0.3/bsp/radiona/ulx3s/smp

#### Usage

```
import spinal.core._
import spinal.core.sim._
import spinal.lib.bus.bmb._
import spinal.lib.bus.bmb.sim._
import spinal.lib.bus.misc.SizeMapping
import spinal.lib.com.usb.ohci._
import spinal.lib.com.usb.phy.UsbHubLsFs.CtrlCc
import spinal.lib.com.usb.phy._
class Usb0hciTop(val p : Usb0hciParameter) extends Component {
 val ohci = UsbOhci(p, BmbParameter(
   addressWidth = 12,
   dataWidth = 32,
   sourceWidth = 0.
    contextWidth = 0,
    lengthWidth = 2
 ))
 val phyCd = ClockDomain.external("phyCd", frequency = FixedFrequency(48 MHz))
 val phy = phyCd(UsbLsFsPhy(p.portCount, sim=true))
 val phyCc = CtrlCc(p.portCount, ClockDomain.current, phyCd)
 phyCc.input <> ohci.io.phy
 phyCc.output <> phy.io.ctrl
 // propagate io signals
 val irq = ohci.io.interrupt.toIo
 val ctrl = ohci.io.ctrl.toIo
 val dma = ohci.io.dma.toIo
 val usb = phy.io.usb.toIo
 val management = phy.io.management.toIo
object UsbHostGen extends App{
 val p = UsbOhciParameter(
   noPowerSwitching = true,
   powerSwitchingMode = true,
   noOverCurrentProtection = true,
   powerOnToPowerGoodTime = 10,
   dataWidth = 64, //DMA data width, up to 128
   portsConfig = List.fill(4)(OhciPortParameter()) //4 Ports
 SpinalVerilog(new UsbOhciTop(p))
}
```

11.10. Com 177

# 11.11 IO

# 11.11.1 ReadableOpenDrain

#### ReadableOpenDrain

The ReadableOpenDrain bundle is defined as following:

Then, as a master, you can use the read signal to read the outside value and use the write to set the value that you want to drive on the output.

There is an example of usage:

```
val io = new Bundle{
  val dataBus = master(ReadableOpenDrain(Bits(32 bits)))
}
io.dataBus.write := 0x12345678
when(io.dataBus.read === 42){
}
```

# 11.11.2 TriState

# Introduction

Tri-state signals are weird to handle in many cases:

- They are not really kind of digital things
- And except for IO, they aren't used for digital design
- The tristate concept doesn't fit naturally in the SpinalHDL internal graph.

SpinalHDL provides two different abstractions for tristate signals. The TriState bundle and *Analog and inout* signals. Both serve different purposes:

- TriState should be used for most purposes, especially within a design. The bundle contains an additional signal to carry the current direction.
- Analog and inout should be used for drivers on the device boundary and in some other special cases. See the referenced documentation page for more details.

As stated above, the recommended approach is to use TriState within a design. On the top-level the TriState bundle is then assigned to an analog inout to get the synthesis tools to infer the correct I/O driver. This can be done automatically done via the *InOutWrapper* or manually if needed.

#### **TriState**

The TriState bundle is defined as following:

A master can use the read signal to read the outside value, the writeEnable to enable the output, and finally use write to set the value that is driven on the output.

There is an example of usage:

```
val io = new Bundle{
  val dataBus = master(TriState(Bits(32 bits)))
}
io.dataBus.writeEnable := True
io.dataBus.write := 0x12345678
when(io.dataBus.read === 42){
}
```

#### **TriStateArray**

In some case, you need to have the control over the output enable of each individual pin (Like for GPIO). In this range of cases, you can use the TriStateArray bundle.

It is defined as following:

```
case class TriStateArray(width : BitCount) extends Bundle with IMasterSlave{
  val read,write,writeEnable = Bits(width)

  override def asMaster(): Unit = {
    out(write,writeEnable)
    in(read)
  }
}
```

It is the same than the TriState bundle, except that the writeEnable is an Bits to control each output buffer.

There is an example of usage:

```
val io = new Bundle{
  val dataBus = master(TriStateArray(32 bits)
}
io.dataBus.writeEnable := 0x87654321
io.dataBus.write := 0x12345678
when(io.dataBus.read === 42){
}
```

11.11. IO 179

# 11.12 Graphics

#### 11.12.1 Colors

#### **RGB**

You can use an Rgb bundle to model colors in hardware. This Rgb bundle take as parameter an RgbConfig classes which specify the number of bits for each channels :

```
case class RgbConfig(rWidth : Int,gWidth : Int,bWidth : Int){
  def getWidth = rWidth + gWidth + bWidth
}

case class Rgb(c: RgbConfig) extends Bundle{
  val r = UInt(c.rWidth bits)
  val g = UInt(c.gWidth bits)
  val b = UInt(c.bWidth bits)
}
```

Those classes could be used as following:

```
val config = RgbConfig(5,6,5)
val color = Rgb(config)
color.r := 31
```

#### 11.12.2 VGA

#### **VGA** bus

An VGA bus definition is available via the Vga bundle.

```
case class Vga (rgbConfig: RgbConfig) extends Bundle with IMasterSlave{
  val vSync = Bool()
  val hSync = Bool()

val colorEn = Bool() //High when the frame is inside the color area
  val color = Rgb(rgbConfig)

override def asMaster() = this.asOutput()
}
```

# **VGA timings**

VGA timings could be modeled in hardware by using an VgaTimings bundle :

```
case class VgaTimingsHV(timingsWidth: Int) extends Bundle {
  val colorStart = UInt(timingsWidth bit)
  val colorEnd = UInt(timingsWidth bit)
  val syncStart = UInt(timingsWidth bit)
  val syncEnd = UInt(timingsWidth bit)
}

case class VgaTimings(timingsWidth: Int) extends Bundle {
  val h = VgaTimingsHV(timingsWidth)
```

```
val v = VgaTimingsHV(timingsWidth)

def setAs_h640_v480_r60 = ...
   def driveFrom(busCtrl : BusSlaveFactory, baseAddress : Int) = ...
}
```

#### **VGA** controller

An VGA controller is available. It's definition is the following:

```
case class VgaCtrl(rgbConfig: RgbConfig, timingsWidth: Int = 12) extends Component {
  val io = new Bundle {
    val softReset = in Bool()
    val timings = in(VgaTimings(timingsWidth))

  val frameStart = out Bool()
  val pixels = slave Stream (Rgb(rgbConfig))
  val vga = master(Vga(rgbConfig))

  val error = out Bool()
  }

// ...
}
```

frameStart is a signals that pulse one cycle at the beginning of each new frame. pixels is a stream of color used to feed the VGA interface when needed. error is high when a transaction on the pixels is needed, but nothing is present.

# 11.13 EDA

# 11.13.1 QSysify

# Introduction

QSysify is a tool which is able to generate a QSys IP (tcl script) from a SpinalHDL component by analysing its IO definition. It currently implement the following interfaces features :

- Master/Slave AvalonMM
- Master/Slave APB3
- · Clock domain input
- Reset output
- Interrupt input
- Conduit (Used in last resort)

11.13. EDA 181

#### **Example**

In the case of a UART controller:

```
case class AvalonMMUartCtrl(...) extends Component{
  val io = new Bundle{
    val bus = slave(AvalonMM(AvalonMMUartCtrl.getAvalonMMConfig))
    val uart = master(Uart())
  }

//...
}
```

The following main will generate the Verilog and the QSys TCL script with io.bus as an AvalonMM and io.uart as a conduit:

```
object AvalonMMUartCtrl{
  def main(args: Array[String]) {
    //Generate the Verilog
    val toplevel = SpinalVerilog(AvalonMMUartCtrl(UartCtrlMemoryMappedConfig(...))).
    →toplevel

    //Add some tags to the avalon bus to specify it's clock domain (information used...)
    →by QSysify)
    toplevel.io.bus addTag(ClockDomainTag(toplevel.clockDomain))

    //Generate the QSys IP (tcl script)
    QSysify(toplevel)
  }
}
```

# tags

Because QSys require some information that are not specified in the SpinalHDL hardware specification, some tags should be added to interface:

#### AvalonMM / APB3

```
io.bus addTag(ClockDomainTag(busClockDomain))
```

# Interrupt input

182

# **Reset output**

```
io.resetOutput addTag(ResetEmitterTag(resetOutputClockDomain))
```

# Adding new interface support

Basically, the QSysify tool can be setup with a list of interface emitter (as you can see here)

You can create your own emitter by creating a new class extending QSysifyInterfaceEmiter

# 11.13.2 QuartusFlow

#### Introduction

A compilation flow is an Altera-defined sequence of commands that use a combination of command-line executables. A full compilation flow launches all Compiler modules in sequence to synthesize, fit, analyze final timing, and generate a device programming file.

Tools in this file help you get rid of redundant Quartus GUI.

# For a single rtl file

The object spinal.lib.eda.altera.QuartusFlow can automatically report the used area and maximum frequency of a single rtl file.

# **Example**

```
val report = QuartusFlow(
    quartusPath="/eda/intelFPGA_lite/17.0/quartus/bin/",
    workspacePath="/home/spinalvm/tmp",
    toplevelPath="TopLevel.vhd",
    family="Cyclone V",
    device="5CSEMA5F31C6",
    frequencyTarget = 1 MHz
)
println(report)
```

The code above will create a new Quartus project with TopLevel.vhd.

Warning: This operation will remove the folder workspacePath!

**Note:** The family and device values are passed straight to the Quartus CLI as parameters. Please check the Quartus documentation for the correct value to use in your project.

11.13. EDA 183

#### Tip

To test a component that has too many pins, set them as VIRTUAL\_PIN.

```
val miaou: Vec[Flow[Bool]] = Vec(master(Flow(Bool())), 666)
miaou.addAttribute("altera_attribute", "-name VIRTUAL_PIN ON")
```

# For an existing project

The class spinal.lib.eda.altera.QuartusProject can automatically find configuration files in an existing project. Those are used for compilation and programming the device.

# **Example**

Specify the path that contains your project files like .qpf and .cdf.

```
val prj = new QuartusProject(
   quartusPath = "F:/intelFPGA_lite/20.1/quartus/bin64/",
   workspacePath = "G:/"
)
prj.compile()
prj.program() // automatically find Chain Description File of the project
```

**Important:** Remember to save the .cdf of your project before calling prj.program().

# 11.14 Misc

# 11.14.1 Plic Mapper

The PLIC Mapper defines the register generation and access for a PLIC (Platform Level Interrupt Controller.

# PlicMapper.apply

```
(bus: BusSlaveFactory, mapping: PlicMapping)(gateways : Seq[PlicGateway], targets :
Seq[PlicTarget])
```

args for PlicMapper:

- bus: bus to which this ctrl is attached
- mapping: a mapping configuration (see above)
- gateways: a sequence of PlicGateway (interrupt sources) to generate the bus access control
- targets: the sequence of PlicTarget (eg. multiple cores) to generate the bus access control

It follows the interface given by riscv: https://github.com/riscv/riscv-plic-spec/blob/master/riscv-plic.adoc

As of now, two memory mappings are available:

#### PlicMapping.sifive

Follows the SiFive PLIC mapping (eg. E31 core complex Manual), basically a full fledged PLIC

#### PlicMapping.light

This mapping generates a lighter PLIC, at the cost of some missing optional features:

- no reading the intrerrupt's priority
- no reading the interrupts's pending bit (must use the claim/complete mechanism)
- no reading the target's threshold

The rest of the registers & logic is generated.

# 11.15 Introduction

# 11.15.1 Introduction

The spinal.lib package goals are:

- Provide things that are commonly used in hardware design (FIFO, clock crossing bridges, useful functions)
- Provide simple peripherals (UART, JTAG, VGA, ..)
- Provide some bus definition (Avalon, AMBA, ..)
- Provide some methodology (Stream, Flow, Fragment)
- Provide some example to get the spirit of spinal
- Provide some tools and facilities (latency analyser, QSys converter, ...)

To use features introduced in followings chapter you need, in most of cases, to import spinal.lib.\_ in your sources.

#### **Important:**

This package is currently under construction. Documented features could be considered as stable.

Do not hesitate to use github for suggestions/bug/fixes/enhancements

11.15. Introduction 185

**CHAPTER** 

**TWELVE** 

# **SIMULATION**

# 12.1 Installation instructions

#### 12.1.1 Scala

To enable SpinalSim, the following lines have to be added in your build.sbt file:

```
fork := true
```

Also the following imports have to be added in testbenches sources:

```
import spinal.core._
import spinal.core.sim._
```

# 12.1.2 Backend-dependent installation instructions

# Setup and installation of GHDL

Even though GHDL is generally available in linux distributions package system, SpinalHDL depends on bugfixes of GHDL codebase that were added after the release of GHDL v0.37. Therefore it is reccomended to install GHDL from source. The C++ library boost-interprocess, which is contained in the libboost-dev package in debian-like distributions, has to be installed too. boost-interprocess is required to generate the shared memory communication interface.

#### Linux

```
sudo apt-get install build-essential libboost-dev git
sudo apt-get install gnat # Ada compiler used to buid GHDL
git clone https://github.com/ghdl/ghdl.git
cd ghdl
mkdir build
cd build
../configure
make
sudo make install
```

Also the openidk package that corresponds to your Java version has to be installed.

For more configuration options and Windows installation see https://ghdl.readthedocs.io/en/latest/getting/

#### Setup and installation of Icarus Verilog

In most recent linux distributions, a recent version of Icarus Verilog is generally available through the package system. The C++ library boost-interprocess, which is contained in the libboost-dev package in debian-like distributions, has to be installed too. boost-interprocess is required to generate the shared memory communication interface.

#### Linux

```
sudo apt-get install build-essential libboost-dev iverilog
```

Also the openjdk package that corresponds to your Java version has to be installed. Refer to https://iverilog.fandom.com/wiki/Installation Guide for more informations about Windows and installation from source.

#### Setup and installation of Verilator

SpinalSim + Verilator is supported on both Linux and Windows platforms.

#### Scala

Don't forget to add the following in your build.sbt file:

```
fork := true
```

And you will always need the following imports in your Scala testbench:

```
import spinal.core._
import spinal.core.sim._
```

#### Linux

You will also need a recent version of Verilator installed:

```
sudo apt-get install git make autoconf g++ flex bison # First time prerequisites
git clone http://git.veripool.org/git/verilator # Only first time
unsetenv VERILATOR_ROOT # For csh; ignore error if on bash
unset VERILATOR_ROOT # For bash
cd verilator
git pull # Make sure we're up-to-date
git checkout v4.040
autoconf # Create ./configure script
./configure
make -j$(nproc)
sudo make install
echo "DONE"
```

#### **Windows**

In order to get SpinalSim + Verilator working on Windows, you have to do the following:

- Install MSYS2
- Via MSYS2 get gcc/g++/verilator (for Verilator you can compile it from the sources)
- Add bin and usr\bin of MSYS2 into your windows PATH (ie: C:\msys64\usr\bin;C:\msys64\ mingw64\bin)
- Check that the JAVA\_HOME environnement variable point to the JDK installation folder (ie: C:\Program Files\Java\jdk-13.0.2)

Then you should be able to run SpinalSim + Verilator from your Scala project without having to use MSYS2 anymore.

From a fresh install of MSYS2 MinGW 64-bit, you will have to run the following commands inside the MSYS2 MinGW 64-bits shell (enter commands one by one):

# From the MinGW package manager

# From source

```
pacman -Syuu
# Close the MSYS2 shell once you're asked to
pacman -Syuu
pacman -S --needed base-devel mingw-w64-x86_64-toolchain \
                   git flex\
                   mingw-w64-x86_64-cmake
git clone http://git.veripool.org/git/verilator
unset VERILATOR ROOT
cd verilator
git pull
git checkout v4.040
autoconf
./configure
export CPLUS_INCLUDE_PATH=/usr/include:$CPLUS_INCLUDE_PATH
export PATH=/usr/bin/core_perl:$PATH
cp /usr/include/FlexLexer.h ./src
make -j$(nproc)
make install
```

```
echo "DONE"
# Add C:\msys64\usr\bin;C:\msys64\mingw64\bin to your Windows PATH
```

**Important:** Be sure that your PATH environnement variable is pointing to the JDK 1.8 and doesn't contain a JRE installation.

**Important:** Adding the MSYS2 bin folders into your windows PATH could potentially have some side effects. This is why it is safer to add them as the last elements of the PATH to reduce their priority.

# 12.2 Boot a simulation

#### 12.2.1 Introduction

There is an example hardware definition + testbench:

# 12.2.2 Configuration

SimConfig will return a default simulation configuration instance on which you can call multiple functions to configure your simulation:

Syntax	Description	
withWave	Enable simulation wave capture	
withConfig(SpinalConfig)	Specify the SpinalConfig that should be use to generate the hardware	
allOptimisation	Enable all the RTL compilation optimizations to reduce simulation time (will	
	increase compilation time)	
workspacePath(path)	Change the folder where the sim files are generated	
withVerilator	Use Verilator as simulation backend (default)	
withGhdl	Use GHDL as simulation backend	
withIVerilog	Use Icarus Verilog as simulation backend	

Then you can call the compile(rtl) function to compile the hardware and warm up the simulator. This function will return a SimCompiled instance.

On this SimCompiled instance you can run your simulation with the following functions:

Syntax	Description		
<pre>doSim[(simName[, seed])]{dut</pre>	Run the simulation until the main thread is done (doesn't wait on forked		
=>}	threads) or until all threads are stuck		
<pre>doSimUntilVoid[(simName[,</pre>	Run the simulation until all threads are done or stuck		
seed])]{dut =>}			

#### For example:

```
val spinalConfig = SpinalConfig(defaultClockDomainFrequency = FixedFrequency(10 MHz))
SimConfig
.withConfig(spinalConfig)
.withWave
.allOptimisation
.workspacePath("~/tmp")
.compile(new TopLevel)
.doSim { dut =>
    // Simulation code here
}
```

Note that by default, the simulation files will be placed into the simWorkspace/xxx folders. You can override the simWorkspace location by setting the SPINALSIM\_WORKSPACE environnement variable.

# 12.2.3 Running multiple tests on the same hardware

```
val compiled = SimConfig.withWave.compile(new Dut)

compiled.doSim("testA") { dut =>
    // Simulation code here
}

compiled.doSim("testB") { dut =>
    // Simulation code here
}
```

# 12.2.4 Throw Success or Failure of the simulation from a thread

At any moment during a simulation you can call simSuccess or simFailure to end it.

# 12.3 Accessing signals of the simulation

# 12.3.1 Read and write signals

Each interface signal of the toplevel can be read and written from Scala:

Syntax	Description			
Bool.toBoolean	Read a hardware Bool as a Scala Boolean value			
Bits/UInt/SInt.toInt	Read a hardware BitVector as a Scala Int value			
Bits/UInt/SInt.toLong	Read a hardware BitVector as a Scala Long value			
Bits/UInt/SInt.toBigInt	Read a hardware BitVector as a Scala BigInt value			
SpinalEnumCraft.toEnum	Read a hardware SpinalEnumCraft as a Scala			
	SpinalEnumElement value			
Bool #= Boolean	Assign a hardware Bool from an Scala Boolean			
Bits/UInt/SInt #= Int	Assign a hardware BitVector from a Scala Int			
Bits/UInt/SInt #= Long	Assign a hardware BitVector from a Scala Long			
Bits/UInt/SInt #= BigInt	Assign a hardware BitVector from a Scala BigInt			
SpinalEnumCraft #= SpinalEnumEle-	Assign a hardware SpinalEnumCraft from a Scala			
ment	SpinalEnumElement			

```
dut.io.a #= 42
dut.io.a #= 421
dut.io.a #= BigInt("101010", 2)
dut.io.a #= BigInt("0123456789ABCDEF", 16)
println(dut.io.b.toInt)
```

# 12.3.2 Accessing signals inside the component's hierarchy

To access signals which are inside the component's hierarchy, you have first to set the given signal as simPublic.

```
You can add this {\tt simPublic} tag directly in the hardware description:
```

```
object SimAccessSubSignal {
 import spinal.core.sim._
 class TopLevel extends Component {
   val counter = Reg(UInt(8 bits)) init(0) simPublic() // Here we add the simPublic_
→tag on the counter register to make it visible
    counter := counter + 1
 def main(args: Array[String]) {
    SimConfig.compile(new TopLevel).doSim{dut =>
      dut.clockDomain.forkStimulus(10)
      for(i <- 0 to 3) {
        dut.clockDomain.waitSampling()
        println(dut.counter.toInt)
      }
   }
 }
}
```

Or you can add it later, after having instantiated your toplevel for the simulation:

```
object SimAccessSubSignal {
  import spinal.core.sim._
  class TopLevel extends Component {
    val counter = Reg(UInt(8 bits)) init(0)
    counter := counter + 1
  }
```

```
def main(args: Array[String]) {
    SimConfig.compile {
        val dut = new TopLevel
            dut.counter.simPublic()
            dut
    }.doSim{dut =>
            dut.clockDomain.forkStimulus(10)

        for(i <- 0 to 3) {
            dut.clockDomain.waitSampling()
            println(dut.counter.toInt)
        }
    }
}</pre>
```

# 12.4 Clock domains

# 12.4.1 Stimulus API

Below is a list of ClockDomain stimulation functions:

ClockDomain stimulus	Description		
functions			
forkStimulus(period)	Fork a simulation process to generate the clockdomain stimulus (clock, reset,		
	softReset, clockEnable signals)		
forkSimSpeedPrinter(pri	t (Rerkizosi) mulation process which will periodically print the simulation speed in		
	kilo-cycles per real time second. printPeriod is in realtime seconds		
clockToggle()	Toggle the clock signal		
fallingEdge()	Clear the clock signal		
risingEdge()	Set the clock signal		
assertReset()	Set the reset signal to its active level		
deassertReset()	Set the reset signal to its inactive level		
assertClockEnable()	Set the clockEnable signal to its active level		
deassertClockEnable()	Set the clockEnable signal to its active level		
assertSoftReset()	Set the softReset signal to its active level		
deassertSoftReset()	Set the softReset signal to its active level		

12.4. Clock domains

# 12.4.2 Wait API

Below is a list of ClockDomain utilities that you can use to wait for a given event from the domain:

ClockDomain	Description			
wait functions				
waitSampling(	[dyailess(i)duht]ClockDomain makes a sampling, (active clock edge && deassertReset &&			
	assertClockEnable)			
waitRisingEdg	waitRisingEdge (Wait colors Sciental) rising edges on the clock; cycleCount defaults to 1 cycle if not other-			
	wise specified. Note, cyclesCount = 0 is legal, and the function is not sensitive to re-			
	set/softReset/clockEnable			
waitFallingEd	waitFallingEdge Safocy as was to this jagedge but for the falling edge			
waitActiveEdg	dge Sagge Lessadut Rilling Edge but for the edge level specified by the ClockDomainConfig			
waitRisingEdg	waitRisingEdge Warene (cvaditRisin) gEdge, but to exit, the boolean condition must be true when the rising			
	edge occurs			
waitFallingEdgeWinnere(wandtRtison) EdgeWhere, but for the falling edge				
waitActiveEdg	eWhavennee (consumbly ing Edge Where, but for the edge level specified by the			
	ClockDomainConfig			

**Warning:** All the functionalities of the wait API can only be called from inside of a thread, and not from a callback.

# 12.4.3 Callback API

Below is a list of ClockDomain utilities that you can use to wait for a given event from the domain:

ClockDomain	Description	
callback func-	·	
tions		
onNextSampling Execute the callback code only once on the next ClockDomain sample (active edge		
{ callback }	off + clock enable on)	
onSamplings	Execute the callback code each time the ClockDomain sample (active edge + reset off +	
{ callback }	clock enable on)	
onActiveEdges	Execute the callback code each time the ClockDomain clock generates its configured edge	
{ callback }		
onEdges {	Execute the callback code each time the ClockDomain clock generates a rising or falling	
callback }	edge	
onRisingEdges	Execute the callback code each time the ClockDomain clock generates a rising edge	
{ callback }		
onFallingEdge	s Execute the callback code each time the ClockDomain clock generates a falling edge	
{ callback }		

# 12.4.4 Default ClockDomain

You can access the default ClockDomain of your toplevel as shown below:

```
// Example of thread forking to generate a reset, and then toggling the clock each 5...
time units.
// dut.clockDomain refers to the implicit clock domain created during component...
instantiation.
fork {
  dut.clockDomain.assertReset()
```

```
dut.clockDomain.fallingEdge()
    sleep(10)
    while(true) {
        dut.clockDomain.clockToggle()
        sleep(5)
    }
}
```

Note that you can also directly fork a standard reset/clock process:

```
dut.clockDomain.forkStimulus(period = 10)
```

An example of how to wait for a rising edge on the clock:

```
dut.clockDomain.waitRisingEdge()
```

# 12.4.5 New ClockDomain

If your toplevel defines some clock and reset inputs which aren't directly integrated into their ClockDomain, you can define their corresponding ClockDomain directly in the testbench:

```
// In the testbench ClockDomain(dut.io.coreClk, dut.io.coreReset).forkStimulus(10)
```

# 12.5 Thread-full API

In SpinalSim, you can write your testbench by using multiple threads in a similar way to SystemVerilog, and a bit like VHDL/Verilog process/always blocks. This allows you to write concurrent tasks and control the simulation time using a fluent API.

# 12.5.1 Fork and join simulation threads

```
// Create a new thread
val myNewThread = fork {
    // New simulation thread body
}

// Wait until `myNewThread` is execution is done.
myNewThread.join()
```

# 12.5.2 Sleep and waitUntil

```
// Sleep 1000 units of time
sleep(1000)

// waitUntil the dut.io.a value is bigger than 42 before continuing
waitUntil(dut.io.a > 42)
```

12.5. Thread-full API 195

# 12.6 Thread-less API

There are some functions that you can use to avoid the need for threading, but which still allow you to control the flow of simulation time.

Threadless	Description
functions	
delayed(delay	Register the callback code to be called at a simulation time delay steps after the current
callback }	timestep.

The advantages of the delayed function over using a regular simulation thread + sleep are:

- Performance (no context switching)
- Memory usage (no native JVM thread memory allocation)

Some other thread-less functions related to ClockDomain objects are documented as part of the *Callback API*, and some others related with the delta-cycle execution process are documented as part of the *Sensitive API* 

# 12.7 Sensitive API

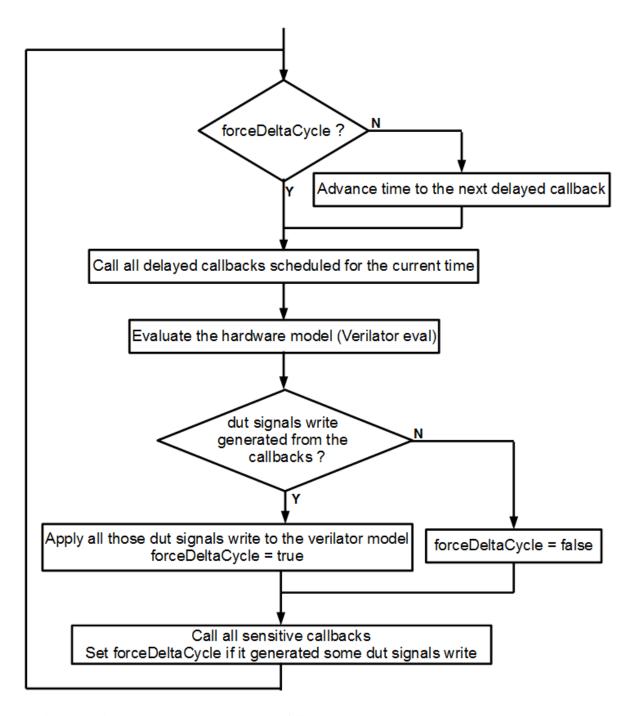
You can register callback functions to be called on each delta-cycle of the simulation:

Sensitive	Description		
functions			
forkSensitive Register the callback code to be called at each delta-cycle of the simulation			
{ callback }			
forkSensitive	forkSensitive Wikelgister the callback code to be called at each delta-cycle of the simulation, while the call-		
{ callback }	back return value is true (meaning it should be rescheduled for the next delta-cycle)		

# 12.8 Simulation engine

This page explains the internals of the simulation engine.

The simulation engine emulates an event-driven simulator (VHDL/Verilog like) by applying the following simulation loop on the top of the Verilator C++ simulation model:



At a low level, the simulation engine manages the following primitives:

- Sensitive callbacks, which allow users to call a function on each simulation delta cycle.
- Delayed callbacks, which allow users to call a function at a future simulation time.
- Simulation threads, which allow users to describe concurrent processes.
- *Command buffer*, which allows users to delay write access to the DUT (Device Under Test) until the end of the current delta cycle.

There are some practical uses of those primitives:

- Sensitive callbacks can be used to wake up a simulation thread when a given condition happens, like a rising edge on a clock.
- Delayed callbacks can be used to schedule stimuli, such as deasserting a reset after a given time, or toggling the clock.

- Both sensitive and delayed callbacks can be used to resume a simulation thread.
- A simulation thread can be used (for instance) to produce stimulus and check the DUT's output values.
- The command buffer's purpose is mainly to avoid all concurrency issues between the DUT and the testbench.

# 12.9 Examples

# 12.9.1 Asynchronous adder

This example creates a Component out of combinational logic that does some simple arithmetic on 3 operands.

The test bench performs the following steps 100 times:

- Initialize a, b, and c to random integers in the 0..255 range.
- Stimulate the DUT's matching a, b, c inputs.
- Wait 1 simulation timestep (to allow the inputs to propagate).
- Check for correct output.

```
import spinal.sim._
import spinal.core._
import spinal.core.sim._
import scala util Random
object SimAsynchronousExample {
 class Dut extends Component {
    val io = new Bundle {
     val a, b, c = in UInt (8 bits)
      val result = out UInt (8 bits)
    io.result := io.a + io.b - io.c
 }
 def main(args: Array[String]): Unit = {
    SimConfig.withWave.compile(new Dut).doSim{ dut =>
      var idx = 0
      while(idx < 100){</pre>
        val a, b, c = Random.nextInt(256)
        dut.io.a #= a
        dut.io.b #= b
        dut.io.c #= c
        sleep(1) // Sleep 1 simulation timestep
        assert(dut.io.result.toInt == ((a + b - c) & 0xFF))
        idx += 1
      }
    }
 }
}
```

#### 12.9.2 Dual clock fifo

This example creates a StreamFifoCC, which is designed for crossing clock domains, along with 3 simulation threads.

The threads handle:

- Management of the two clocks
- Pushing to the FIFO
- Popping from the FIFO

The FIFO push thread randomizes the inputs.

The FIFO pop thread handles checking the DUT's outputs against the reference model (an ordinary scala. collection.mutable.Queue instance).

```
import spinal.sim._
import spinal.core._
import spinal.core.sim._
import scala.collection.mutable.Queue
object SimStreamFifoCCExample {
 def main(args: Array[String]): Unit = {
   // Compile the Component for the simulator.
   val compiled = SimConfig.withWave.allOptimisation.compile(
      rtl = new StreamFifoCC(
        dataType = Bits(32 bits),
        depth = 32,
        pushClock = ClockDomain.external("clkA"),
        popClock = ClockDomain.external("clkB")
      )
   )
    // Run the simulation.
   compiled.doSimUntilVoid{dut =>
      val queueModel = mutable.Queue[Long]()
      // Fork a thread to manage the clock domains signals
      val clocksThread = fork {
        // Clear the clock domains' signals, to be sure the simulation captures their.
⊶first edges.
        dut.pushClock.fallingEdge()
        dut.popClock.fallingEdge()
        dut.pushClock.deassertReset()
        dut.popClock.deassertReset()
        sleep(0)
        // Do the resets.
        dut.pushClock.assertReset()
        dut.popClock.assertReset()
        sleep(10)
        dut.pushClock.deassertReset()
        dut.popClock.deassertReset()
        sleep(1)
        // Forever, randomly toggle one of the clocks.
```

(continues on next page)

12.9. Examples 199

```
// This will create asynchronous clocks without fixed frequencies.
        while(true) {
          if(Random.nextBoolean()) {
            dut.pushClock.clockToggle()
          } else {
            dut.popClock.clockToggle()
          sleep(1)
        }
      }
      // Push data randomly, and fill the queueModel with pushed transactions.
      val pushThread = fork {
        while(true) {
          dut.io.push.valid.randomize()
          dut.io.push.payload.randomize()
          dut.pushClock.waitSampling()
          if(dut.io.push.valid.toBoolean && dut.io.push.ready.toBoolean) {
            queueModel.enqueue(dut.io.push.payload.toLong)
        }
      }
      // Pop data randomly, and check that it match with the queueModel.
      val popThread = fork {
        for(i <- 0 until 100000) {</pre>
          dut.io.pop.ready.randomize()
          dut.popClock.waitSampling()
          if(dut.io.pop.valid.toBoolean && dut.io.pop.ready.toBoolean) {
            assert(dut.io.pop.payload.toLong == queueModel.dequeue())
        }
        simSuccess()
      }
    }
 }
}
```

# 12.9.3 Single clock fifo

This example creates a StreamFifo, and spawns 3 simulation threads. Unlike the *Dual clock fifo* example, this FIFO does not need complex clock management.

The 3 simulation threads handle:

- Managing the clock/reset
- Pushing to the FIFO
- Popping from the FIFO

The FIFO push thread randomizes the inputs.

The FIFO pop thread handles checking the DUT's outputs against the reference model (an ordinary scala. collection.mutable.Queue instance).

```
import spinal.sim._
import spinal.core._
```

```
import spinal.core.sim._
import scala.collection.mutable.Queue
object SimStreamFifoExample {
 def main(args: Array[String]): Unit = {
   // Compile the Component for the simulator.
   val compiled = SimConfig.withWave.allOptimisation.compile(
      rtl = new StreamFifo(
        dataType = Bits(32 bits),
        depth = 32
      )
   )
   // Run the simulation.
   compiled.doSimUntilVoid{dut =>
      val queueModel = mutable.Queue[Long]()
      dut.clockDomain.forkStimulus(period = 10)
      SimTimeout(1000000*10)
      // Push data randomly, and fill the queueModel with pushed transactions.
      val pushThread = fork {
        dut.io.push.valid #= false
        while(true) {
          dut.io.push.valid.randomize()
          dut.io.push.payload.randomize()
          dut.clockDomain.waitSampling()
          if(dut.io.push.valid.toBoolean && dut.io.push.ready.toBoolean) {
            queueModel.enqueue(dut.io.push.payload.toLong)
          }
        }
      }
      // Pop data randomly, and check that it match with the queueModel.
      val popThread = fork {
        dut.io.pop.ready #= true
        for(i <- 0 until 100000) {</pre>
          dut.io.pop.ready.randomize()
          dut.clockDomain.waitSampling()
          if(dut.io.pop.valid.toBoolean && dut.io.pop.ready.toBoolean) {
            assert(dut.io.pop.payload.toLong == queueModel.dequeue())
          }
        }
        simSuccess()
      }
   }
 }
}
```

12.9. Examples 201

# 12.9.4 Synchronous adder

This example creates a Component out of sequential logic that does some simple arithmetic on 3 operands.

The test bench performs the following steps 100 times:

- Initialize a, b, and c to random integers in the 0..255 range.
- Stimulate the DUT's matching a, b, c inputs.
- Wait until the simulation samples the DUT's signals again.
- Check for correct output.

The main difference between this example and the *Asynchronous adder* example is that this Component has to use forkStimulus to generate a clock signal, since it is using sequential logic internally.

```
import spinal.sim._
import spinal.core._
import spinal.core.sim._
import scala.util.Random
object SimSynchronousExample {
 class Dut extends Component {
   val io = new Bundle {
      val a, b, c = in UInt (8 bits)
      val result = out UInt (8 bits)
   io.result := RegNext(io.a + io.b - io.c) init(0)
 }
 def main(args: Array[String]): Unit = {
   SimConfig.withWave.compile(new Dut).doSim{ dut =>
      dut.clockDomain.forkStimulus(period = 10)
      var resultModel = 0
      for(idx <- 0 until 100){</pre>
        dut.io.a #= Random.nextInt(256)
        dut.io.b #= Random.nextInt(256)
        dut.io.c #= Random.nextInt(256)
        dut.clockDomain.waitSampling()
        assert(dut.io.result.toInt == resultModel)
        resultModel = (dut.io.a.toInt + dut.io.b.toInt - dut.io.c.toInt) & 0xFF
    }
 }
```

# 12.9.5 Uart decoder

```
// Fork a simulation process which will analyze the uartPin and print transmitted.
⇒bytes into the simulation terminal.
fork {
 // Wait until the design sets the uartPin to true (wait for the reset effect).
 waitUntil(uartPin.toBoolean == true)
 while(true) {
    waitUntil(uartPin.toBoolean == false)
    sleep(baudPeriod/2)
    assert(uartPin.toBoolean == false)
    sleep(baudPeriod)
    var buffer = 0
    for(bitId <- 0 to 7) {</pre>
      if(uartPin.toBoolean)
        buffer |= 1 << bitId</pre>
      sleep(baudPeriod)
    }
    assert(uartPin.toBoolean == true)
    print(buffer.toChar)
 }
}
```

# 12.9.6 Uart encoder

```
// Fork a simulation process which will get chars typed into the simulation terminal.
→and transmit them on the simulation uartPin.
fork{
 uartPin #= true
 while(true) {
    // System.in is the java equivalent of the C's stdin.
   if(System.in.available() != 0) {
     val buffer = System.in.read()
     uartPin #= false
     sleep(baudPeriod)
      for(bitId <- 0 to 7) {
        uartPin #= ((buffer >> bitId) & 1) != 0
        sleep(baudPeriod)
     uartPin #= true
      sleep(baudPeriod)
   } else {
      sleep(baudPeriod * 10) // Sleep a little while to avoid polling System.in too.
⇔often.
   }
 }
}
```

12.9. Examples 203

# 12.10 Introduction

As always, you can use your standard simulation tools to simulate the VHDL/Verilog generated by SpinalHDL. However, since SpinalHDL 1.0.0, the language integrates an API to write testbenches and test your hardware directly in Scala. This API provides the capabilities to read and write the DUT signals, fork and join simulation processes, sleep and wait until a given condition is reached. Therefore, using SpinalHDL's simulation API, it is easy to integrate testbenches with the most common Scala unit-test frameworks.

To be able to simulate user-defined components, SpinalHDL uses external HDL simulators as backend. Currently, three simulators are supported:

- Verilator
- GHDL (experimental, since SpinalHDL 1.4.1)
- Icarus Verilog (experimental, since SpinalHDL 1.4.1)

With external HDL simulators it is possible to directly test the generated HDL sources without increasing the SpinalHDL codebase complexity.

# 12.11 How SpinalHDL simulates the hardware with Verilator backend

- 1. Behind the scenes, SpinalHDL generates a Verilog equivalent hardware model of the DUT and then uses Verilator to convert it to a C++ cycle-accurate model.
- 2. The C++ model is compiled into a shared object (.so), which is bound to Scala via JNR-FFI.
- 3. The native Verilator API is abstracted by providing a simulation multi-threaded API.

# Advantages:

• Since the Verilator backend uses a compiled C++ simulation model, the simulation speed is fast compared to most of the other commercial and free simulators.

#### **Limitations:**

- Verilator accepts only synthesizable Verilog/System Verilog code. Therefore special care has to be taken when simulating Verilog blackbox components that may have non-synthesizable statements.
- VHDL blackboxes cannot be simulated.
- The simulation boot process is slow due to the necessity to compile and link the generated C++ model

# 12.12 How SpinalHDL simulates the hardware with GHDL/Icarus Verilog backend

- 1. Depending on the chosen simulator, SpinalHDL generates a Verilog or VHDL hardware model of the DUT.
- 2. The HDL model is loaded in the simulator.
- 3. The communication between the simulation and the JVM is established through shared memory. The commands are issued to the simulator using VPI.

#### Advantages:

- Both GHDL and Icarus Verilog can accept non-synthesizable HDL code.
- The simulation boot process is quite faster compared to Verilator.

# **Limitations:**

• GHDL accepts VHDL code only. Therefore only VHDL blackboxes can be used with this simulator.

- Icarus Verilog accepts Verilog code only. Therefore only Verilog blackboxes can be used with this simulator.
- The simulation speed is around one order of magnitude slower compared to Verilator.

Finally, as the native Verilator API is rather crude, SpinalHDL abstracts over it by providing both single and multi-threaded simulation APIs to help the user construct testbench implementations.

# 12.13 Performance

When a high-performance simulation is required, Verilator should be used as a backend. On a little SoC like Murax, an Intel® Core<sup>TM</sup> i7-4720HQ is capable of simulating 1.2 million clock cycles per second. However, when the DUT is simple and a maximum of few thousands clock cycles have to be simulated, using GHDL or Icarus Verilog could yield a better result, due to their lower simulation loading overhead.

12.13. Performance 205

**CHAPTER** 

# **THIRTEEN**

# **EXAMPLES**

# 13.1 Simple ones

# 13.1.1 APB3 definition

#### Introduction

This example will show the syntax to define an APB3 Bundle.

# **Specification**

The specification from ARM could be interpreted as follows:

Signal	Туре	Driver side	Comment
Name			
PADDR	UInt(addressWidth bits)	Master	Address in byte
PSEL	Bits(selWidth)	Master	One bit per slave
PENABLE	Bool	Master	
PWRITE	Bool	Master	
PWDATA	Bits(dataWidth bits)	Master	
PREADY	Bool	Slave	
PRDATA	Bits(dataWidth bits)	Slave	
PSLVER-	Bool	Slave	Optional
ROR			

# **Implementation**

This specification shows that the APB3 bus has multiple possible configurations. To represent that, we can define a configuration class in Scala:

```
case class Apb3Config(
  addressWidth : Int,
  dataWidth : Int,
  selWidth : Int = 1,
  useSlaveError : Boolean = true
)
```

Then we can define the APB3 Bundle which will be used to represent the bus in hardware:

```
val PENABLE = Bool()
val PREADY = Bool()
val PWRITE = Bool()
val PWDATA = Bits(config.dataWidth bit)
val PRDATA = Bits(config.dataWidth bit)
val PSLVERROR = if(config.useSlaveError) Bool else null

override def asMaster(): Unit = {
  out(PADDR,PSEL,PENABLE,PWRITE,PWDATA)
  in(PREADY,PRDATA)
  if(config.useSlaveError) in(PSLVERROR)
}
```

#### **Usage**

Here is a usage example of this definition:

```
val apbConfig = Apb3Config(
  addressWidth = 16,
  dataWidth = 32,
  selWidth = 1,
  useSlaveError = false
)

val io = new Bundle{
  val apb = slave(Apb3(apbConfig))
}

io.apb.PREADY := True
when(io.apb.PSEL(0) && io.apb.PENABLE){
    //...
}
```

# 13.1.2 Carry adder

This example defines a component with inputs a and b, and a result output. At any time, result will be the sum of a and b (combinatorial). This sum is manually done by a carry adder logic.

```
c \= (a & b) | (a & c) | (b & c);  //variable assignment
}

object CarryAdderProject {
  def main(args: Array[String]) {
    SpinalVhdl(new CarryAdder(4))
  }
}
```

# 13.1.3 Color summing

First let's define a Color Bundle with an addition operator.

```
case class Color(channelWidth: Int) extends Bundle {
 val r = UInt(channelWidth bits)
 val g = UInt(channelWidth bits)
 val b = UInt(channelWidth bits)
 def +(that: Color): Color = {
   val result = Color(channelWidth)
   result.r := this.r + that.r
   result.g := this.g + that.g
   result.b := this.b + that.b
   return result
 }
 def clear(): Color ={
   this.r := 0
   this.g := 0
    this.b := 0
    this
 }
}
```

Then let's define a component with a sources input which is a vector of colors, and a result output which is the sum of the sources input.

```
class ColorSumming(sourceCount: Int, channelWidth: Int) extends Component {
  val io = new Bundle {
    val sources = in Vec(Color(channelWidth), sourceCount)
    val result = out(Color(channelWidth))
  }

  var sum = Color(channelWidth)
  sum.clear()
  for (i <- 0 to sourceCount - 1) {
    sum \= sum + io.sources(i)
  }
  io.result := sum
}</pre>
```

13.1. Simple ones

#### 13.1.4 Counter with clear

This example defines a component with a clear input and a value output. Each clock cycle, the value output is incrementing, but when clear is high, value is cleared.

```
class Counter(width : Int) extends Component{
  val io = new Bundle{
    val clear = in Bool()
    val value = out UInt(width bits)
}

val register = Reg(UInt(width bits)) init(0)
register := register + 1
when(io.clear){
    register := 0
}
io.value := register
}
```

# 13.1.5 PLL BlackBox and reset controller

Let's imagine you want to define a TopLevel component which instantiates a PLL BlackBox, and create a new clock domain from it which will be used by your core logic. Let's also imagine that you want to adapt an external asynchronous reset into this core clock domain to a synchronous reset source.

The following imports will be used in code examples on this page:

```
import spinal.core._
import spinal.lib._
```

#### The PLL BlackBox definition

This is how to define the PLL BlackBox:

```
class PLL extends BlackBox{
  val io = new Bundle{
    val clkIn = in Bool()
    val clkOut = out Bool()
    val isLocked = out Bool()
}

noIoPrefix()
}
```

This will correspond to the following VHDL component:

```
component PLL is
  port(
    clkIn : in std_logic;
    clkOut : out std_logic;
    isLocked : out std_logic
  );
end component;
```

# **TopLevel definition**

This is how to define your TopLevel which instantiates the PLL, creates the new ClockDomain, and also adapts the asynchronous reset input to a synchronous reset:

```
class TopLevel extends Component{
 val io = new Bundle {
   val aReset = in Bool()
   val clk100Mhz = in Bool()
   val result = out UInt(4 bits)
 }
 // Create an Area to manage all clocks and reset things
 val clkCtrl = new Area {
   //Instanciate and drive the PLL
   val pll = new PLL
   pll.io.clkIn := io.clk100Mhz
   //Create a new clock domain named 'core'
   val coreClockDomain = ClockDomain.internal(
     name = "core",
      frequency = FixedFrequency(200 MHz) // This frequency specification can be used
                                           // by coreClockDomain users to do some_

→ calculations

    //Drive clock and reset signals of the coreClockDomain previously created
    coreClockDomain.clock := pll.io.clkOut
    coreClockDomain.reset := ResetCtrl.asyncAssertSyncDeassert(
      input = io.aReset || ! pll.io.isLocked,
      clockDomain = coreClockDomain
 }
  //Create a ClockingArea which will be under the effect of the clkCtrl.
→ coreClockDomain
 val core = new ClockingArea(clkCtrl.coreClockDomain){
   //Do your stuff which use coreClockDomain here
   val counter = Reg(UInt(4 bits)) init(0)
    counter := counter + 1
    io.result := counter
 }
```

# 13.1.6 RGB to gray

Let's imagine a component that converts an RGB color into a gray one, and then writes it into external memory.

io name	Direction	Description
clear	in	Clear all internal registers
r,g,b	in	Color inputs
wr	out	Memory write
address	out	Memory address, incrementing each cycle
data	out	Memory data, gray level

13.1. Simple ones 211

```
class RgbToGray extends Component{
  val io = new Bundle{
    val clear = in Bool()
    val r,g,b = in UInt(8 bits)
    val wr = out Bool()
    val address = out UInt(16 bits)
    val data = out UInt(8 bits)
  }
  def coef(value : UInt,by : Float) : UInt = (value * U((255*by).toInt,8 bits) >> 8)
  val gray = RegNext(
    coef(io.r, 0.3f) +
    coef(io.g, 0.4f) +
    coef(io.b, 0.3f)
  val address = CounterFreeRun(stateCount = 1 << 16)</pre>
  io.address := address
  io.wr := True
  io.data := gray
  when(io.clear){
    gray := 0
    address.clear()
    io.wr := False
  }
}
```

# 13.1.7 Sinus rom

Let's imagine that you want to generate a sine wave and also have a filtered version of it (which is completely useless in practical, but let's do it as an example).

Parameters name Type		Description
resolutionWidth	Int	Number of bits used to represent numbers
sampleCount	Int	Number of samples in a sine period

Ю	Di-	Type	Description
name	rec-		
	tion		
sin	out	SInt(resolutionWidth bits)	Output which plays the sine wave
sin-	out	SInt(resolutionWidth bits)	Output which plays the filtered version of the sine
Fil-			
tred			

So let's define the Component:

```
class TopLevel(resolutionWidth : Int, sampleCount : Int) extends Component {
  val io = new Bundle {
    val sin = out SInt(resolutionWidth bits)
    val sinFiltred = out SInt(resolutionWidth bits)
  }
}
```

```
// Here will come the logic implementation }
```

To play the sine wave on the sin output, you can define a ROM which contain all samples of a sine period (tt could be just a quarter, but let's do things by the simplest way).

Then you can read that ROM with an phase counter and this will generate your sine wave.

```
//Function used to generate the rom (later)
def sinTable = for(sampleIndex <- 0 until sampleCount) yield {
  val sinValue = Math.sin(2 * Math.PI * sampleIndex / sampleCount)
  S((sinValue * ((1<<resolutionWidth)/2-1)).toInt,resolutionWidth bits)
}
val rom = Mem(SInt(resolutionWidth bits),initialContent = sinTable)
val phase = Reg(UInt(log2Up(sampleCount) bits)) init(0)
phase := phase + 1

io.sin := rom.readSync(phase)</pre>
```

Then to generate sinFiltred, you can for example use a first order low pass filter implementation:

```
io.sinFiltred := RegNext(io.sinFiltred - (io.sinFiltred >> 5) + (io.sin >> 5))

io.sinFiltred := RegNext(io.sinFiltred - (io.sinFiltred >> 5) + (io.sin >> 5))

io.sinFiltred := RegNext(io.sinFiltred - (io.sinFiltred >> 5) + (io.sin >> 5))

io.sinFiltred := RegNext(io.sinFiltred - (io.sinFiltred >> 5) + (io.sin >> 5))

io.sinFiltred := RegNext(io.sinFiltred - (io.sinFiltred >> 5) + (io.sin >> 5))

io.sinFiltred := RegNext(io.sinFiltred - (io.sinFiltred >> 5) + (io.sin >> 5))

io.sinFiltred := RegNext(io.sinFiltred - (io.sinFiltred >> 5) + (io.sin >> 5))

io.sinFiltred := RegNext(io.sinFiltred - (io.sinFiltred >> 5) + (io.sin >> 5))

io.sinFiltred := RegNext(io.sinFiltred - (io.sinFiltred >> 5) + (io.sin >> 5))

io.sinFiltred := RegNext(io.sinFiltred - (io.sinFiltred >> 5) + (io.sin >> 5))

io.sinFiltred := RegNext(io.sinFiltred - (io.sinFiltred >> 5) + (io.sinFiltred >> 5)

io.sinFiltred := RegNext(io.sinFiltred - (io.sinFiltred >> 5) + (io.sinFiltred >> 5)

io.sinFiltred := RegNext(io.sinFiltred - (io.sinFiltred >> 5) + (io.sinFiltred >> 5)

io.sinFiltred := RegNext(io.sinFiltred - (io.sinFiltred >> 5) + (io.sinFiltred >> 5)

io.sinFiltred := RegNext(io.sinFiltred - (io.sinFiltred >> 5)

io.sinFiltred := RegNext(io.sinFiltred - (io.sinFiltred >> 5)

io.sinFiltred := RegNext(io.sinFiltred - (io.sinFiltred - (io.sinFiltred
```

Here is the complete code:

```
class TopLevel(resolutionWidth : Int,sampleCount : Int) extends Component {
  val io = new Bundle {
    val sin = out SInt(resolutionWidth bits)
    val sinFiltred = out SInt(resolutionWidth bits)
}

def sinTable = for(sampleIndex <- 0 until sampleCount) yield {
    val sinValue = Math.sin(2 * Math.PI * sampleIndex / sampleCount)
    S((sinValue * ((1<<resolutionWidth)/2-1)).toInt,resolutionWidth bits)
}

val rom = Mem(SInt(resolutionWidth bits),initialContent = sinTable)
    val phase = Reg(UInt(log2Up(sampleCount) bits)) init(0)
    phase := phase + 1

io.sin := rom.readSync(phase)
    io.sinFiltred := RegNext(io.sinFiltred - (io.sinFiltred >> 5) + (io.sin >> 5))
    init(0)
}
```

13.1. Simple ones 213

# 13.2 Intermediates ones

### 13.2.1 Fractal calculator

### Introduction

This example will show a simple implementation (without optimization) of a Mandelbrot fractal calculator by using data streams and fixed point calculations.

# **Specification**

The component will receive one Stream of pixel tasks (which contain the XY coordinates in the Mandelbrot space) and will produce one Stream of pixel results (which contain the number of iterations done for the corresponding task).

Let's specify the IO of our component:

Ю	Di-	Type	Description
Name	rec-		
	tion		
cmd	slave	Stream	ProvidekXY coordinates to process
rsp	mas-	Stream	Recire stuttation count needed for the corresponding cmd transaction
	ter		

Let's specify the PixelTask Bundle:

Element Name	Туре	Description
X	SFix	Coordinate in the Mandelbrot space
у	SFix	Coordinate in the Mandelbrot space

Let's specify the PixelResult Bundle:

Element Name	Туре	Description
iteration	UInt	Number of iterations required to solve the Mandelbrot coordinates

# **Elaboration parameters (Generics)**

Let's define the class that will provide construction parameters of our system:

Note: iterationType and fixType are functions that you can call to instantiate new signals. It's like a typedef in C.

### **Bundle definition**

```
case class PixelTask(g : PixelSolverGenerics) extends Bundle{
  val x,y = g.fixType
}

case class PixelResult(g : PixelSolverGenerics) extends Bundle{
  val iteration = g.iterationType
}
```

### **Component implementation**

And now the implementation. The one below is a very simple one without pipelining / multi-threading.

```
case class PixelSolver(g : PixelSolverGenerics) extends Component{
  val io = new Bundle{
    val cmd = slave Stream(PixelTask(g))
    val rsp = master Stream(PixelResult(g))
  import g._
  //Define states
          = Reg(fixType) init(0)
  val x,y
  val iteration = Reg(iterationType) init(0)
  //Do some shared calculation
  val xx = x*x
  val yy = y*y
  val xy = x*y
  //Apply default assignment
  io.cmd.ready := False
  io.rsp.valid := False
  io.rsp.iteration := iteration
  when(io.cmd.valid) {
    //Is the mandelbrot iteration done ?
    when(xx + yy >= 4.0 || iteration === iterationLimit) {
      io.rsp.valid := True
      when(io.rsp.ready){
       io.cmd.ready := True
        x := 0
        y := 0
        iteration := 0
      }
    } otherwise {
     x := (xx - yy + io.cmd.x).truncated
      y := (((xy) << 1) + io.cmd.y).truncated
      iteration := iteration + 1
    }
  }
}
```

# 13.2.2 **UART**

# **Specification**

This UART controller tutorial is based on this implementation.

This implementation is characterized by:

- ClockDivider/Parity/StopBit/DataLength configs are set by the component inputs.
- RXD input is filtered by using a sampling window of N samples and a majority vote.

Interfaces of this UartCtrl are:

Name	Type	Description
config	UartC-	Give all configurations to the controller
	trl-	
	Con-	
	fig	
write	Stream[	BRoft used by the system to give transmission order to the controller
read	Flow[Bi	ts ort used by the controller to notify the system about a successfully received frame
uart	Uart	Uart interface with rxd / txd

### **Data structures**

Before implementing the controller itself we need to define some data structures.

# **Controller construction parameters**

Nameyr	eDescription
dataWnt	Maximum number of data bits that could be sent using a single UART frame
idth-	
Max	
clockInt	Number of bits that the clock divider has
Di-	
vider-	
Width	
pre- Int	Number of samples to drop at the beginning of the sampling window
Sam-	
pling-	
Size	
sam-Int	Number of samples use at the middle of the window to get the filtered RXD value
pling-	
Size	
post-Int	Number of samples to drop at the end of the sampling window
Sam-	
pling-	
Size	

To make the implementation easier let's assume that preSamplingSize + samplingSize + postSamplingSize is always a power of two.

Instead of adding each construction parameters (generics) to UartCtrl one by one, we can group them inside a class that will be used as single parameter of UartCtrl.

#### **UART** bus

Let's define a UART bus without flow control.

```
case class Uart() extends Bundle with IMasterSlave {
  val txd = Bool()
  val rxd = Bool()

  override def asMaster(): Unit = {
    out(txd)
    in(rxd)
  }
}
```

# **UART** configuration enums

Let's define parity and stop bit enumerations.

```
object UartParityType extends SpinalEnum(sequancial) {
  val NONE, EVEN, ODD = newElement()
}
object UartStopType extends SpinalEnum(sequancial) {
  val ONE, TWO = newElement()
  def toBitCount(that : T) : UInt = (that === ONE) ? U"0" | U"1"
}
```

# **UartCtrl configuration Bundles**

Let's define Bundles that will be used as IO elements to setup UartCtrl.

```
case class UartCtrlFrameConfig(g: UartCtrlGenerics) extends Bundle {
  val dataLength = UInt(log2Up(g.dataWidthMax) bit) //Bit count = dataLength + 1
  val stop = UartStopType()
  val parity = UartParityType()
}

case class UartCtrlConfig(g: UartCtrlGenerics) extends Bundle {
  val frame = UartCtrlFrameConfig(g)
  val clockDivider = UInt (g.clockDividerWidth bit) //see UartCtrlGenerics.
```

# **Implementation**

In UartCtrl, 3 things will be instantiated:

- One clock divider that generates a tick pulse at the UART RX sampling rate.
- One UartCtrlTx Component
- One UartCtrlRx Component

### **UartCtrlTx**

The interfaces of this Component are the following:

Name	Type	Description
con-	UartC-	Contains data bit width count and party/stop bits configurations
figFram	e trl-	
	Frame-	
	Con-	
	fig	
sam-	Bool	Time reference that pulses rxSamplePerBit times per UART baud
plingTic	k	
write	Stream[	BRoft used by the system to give transmission orders to the controller
txd	Bool	UART txd pin

Let's define the enumeration that will be used to store the state of UartCtrlTx:

```
object UartCtrlTxState extends SpinalEnum {
  val IDLE, START, DATA, PARITY, STOP = newElement()
}
```

Let's define the skeleton of UartCtrlTx:

```
class UartCtrlTx(g : UartCtrlGenerics) extends Component {
  import g._

val io = new Bundle {
  val configFrame = in(UartCtrlFrameConfig(g))
  val samplingTick = in Bool()
  val write = slave Stream (Bits(dataWidthMax bit))
  val txd = out Bool
  }

// Provide one clockDivider.tick each rxSamplePerBit pulses of io.samplingTick
  // Used by the stateMachine as a baud rate time reference
  val clockDivider = new Area {
    val counter = Reg(UInt(log2Up(rxSamplePerBit) bits)) init(0)
```

```
val tick = False
 }
 // Count up each clockDivider.tick, used by the state machine to count up data bits.
\rightarrowand stop bits
 val tickCounter = new Area {
   val value = Reg(UInt(Math.max(dataWidthMax, 2) bit))
   def reset() = value := 0
 }
 val stateMachine = new Area {
   import UartCtrlTxState._
   val state = RegInit(IDLE)
   val parity = Reg(Bool)
   val txd = True
   switch(state) {
   }
 }
 io.txd := RegNext(stateMachine.txd) init(True)
```

And here is the complete implementation:

```
class UartCtrlTx(g : UartCtrlGenerics) extends Component {
 import g._
 val io = new Bundle {
   val configFrame = in(UartCtrlFrameConfig(g))
   val samplingTick = in Bool
   val write = slave Stream (Bits(dataWidthMax bit))
   val txd
                   = out Bool
 // Provide one clockDivider.tick each rxSamplePerBit pulse of io.samplingTick
 // Used by the stateMachine as a baud rate time reference
 val clockDivider = new Area {
   val counter = Reg(UInt(log2Up(rxSamplePerBit) bits)) init(0)
   val tick = False
   when(io.samplingTick) {
     counter := counter - 1
     tick := counter === 0
   }
 }
 // Count up each clockDivider.tick, used by the state machine to count up data bits.
→and stop bits
 val tickCounter = new Area {
   val value = Reg(UInt(Math.max(dataWidthMax, 2) bit))
   def reset() = value := 0
```

```
when(clockDivider.tick) {
    value := value + 1
}
val stateMachine = new Area {
  import UartCtrlTxState._
  val state = RegInit(IDLE)
  val parity = Reg(Bool)
  val txd = True
  when(clockDivider.tick) {
   parity := parity ^ txd
  }
  io.write.ready := False
  switch(state) {
    is(IDLE){
      when(io.write.valid && clockDivider.tick){
        state := START
      }
    }
    is(START) {
      txd := False
      when(clockDivider.tick) {
        state := DATA
        parity := io.configFrame.parity === UartParityType.ODD
        tickCounter.reset()
      }
    }
    is(DATA) {
      txd := io.write.payload(tickCounter.value)
      when(clockDivider.tick) {
        when(tickCounter.value === io.configFrame.dataLength) {
          io.write.ready := True
          tickCounter.reset()
          when(io.configFrame.parity === UartParityType.NONE) {
            state := STOP
          } otherwise {
            state := PARITY
        }
      }
    }
    is(PARITY) {
      txd := parity
      when(clockDivider.tick) {
        state := STOP
        tickCounter.reset()
      }
    }
    is(STOP) {
      when(clockDivider.tick) {
        when(tickCounter.value === toBitCount(io.configFrame.stop)) {
          state := io.write.valid ? START | IDLE
```

```
}
}

io.txd := RegNext(stateMachine.txd, True)
}
```

## **UartCtrlRx**

The interfaces of this Component are the following:

Name	Type	Description
con-	UartC-	Contains data bit width and party/stop bits configurations
figFram	e trl-	
	Frame-	
	Con-	
	fig	
sam-	Bool	Time reference that pulses rxSamplePerBit times per UART baud
plingTic	k	
read	Flow[Bi	ts ort used by the controller to notify the system about a successfully received frame
rxd	Bool	UART rxd pin, not synchronized with the current clock domain

Let's define the enumeration that will be used to store the state of UartCtrlTx:

```
object UartCtrlRxState extends SpinalEnum {
  val IDLE, START, DATA, PARITY, STOP = newElement()
}
```

Let's define the skeleton of the UartCtrlRx:

```
class UartCtrlRx(g : UartCtrlGenerics) extends Component {
 import g._
 val io = new Bundle {
   val configFrame = in(UartCtrlFrameConfig(g))
   val samplingTick = in Bool
              = master Flow (Bits(dataWidthMax bit))
   val read
   val rxd
                   = in Bool
 }
 // Implement the rxd sampling with a majority vote over samplingSize bits
 // Provide a new sampler.value each time sampler.tick is high
 val sampler = new Area {
   val syncroniser = BufferCC(io.rxd)
   val samples
                = History(that=syncroniser,when=io.samplingTick,
→length=samplingSize)
   val value = RegNext(MajorityVote(samples))
   val tick
                   = RegNext(io.samplingTick)
 }
 // Provide a bitTimer.tick each rxSamplePerBit
 // reset() can be called to recenter the counter over a start bit.
 val bitTimer = new Area {
```

```
val counter = Reg(UInt(log2Up(rxSamplePerBit) bit))
   def reset() = counter := preSamplingSize + (samplingSize - 1) / 2 - 1)
   val tick = False
 }
 // Provide bitCounter.value that count up each bitTimer.tick, Used by the state.
→machine to count data bits and stop bits
 // reset() can be called to reset it to zero
 val bitCounter = new Area {
   val value = Reg(UInt(Math.max(dataWidthMax, 2) bit))
   def reset() = value := 0
 }
 val stateMachine = new Area {
   import UartCtrlRxState._
   val state = RegInit(IDLE)
   val parity = Reg(Bool)
   val shifter = Reg(io.read.payload)
   switch(state) {
   }
 }
```

And here is the complete implementation:

```
class UartCtrlRx(g : UartCtrlGenerics) extends Component {
 import g._
 val io = new Bundle {
   val configFrame = in(UartCtrlFrameConfig(g))
   val samplingTick = in Bool
   val read = master Flow (Bits(dataWidthMax bit))
   val rxd
                   = in Bool
 // Implement the rxd sampling with a majority vote over samplingSize bits
 // Provide a new sampler.value each time sampler.tick is high
 val sampler = new Area {
   val syncroniser = BufferCC(io.rxd)
   val samples
                = History(that=syncroniser,when=io.samplingTick,
→length=samplingSize)
   val value = RegNext(MajorityVote(samples))
   val tick
                   = RegNext(io.samplingTick)
 }
 // Provide a bitTimer.tick each rxSamplePerBit
 // reset() can be called to recenter the counter over a start bit.
 val bitTimer = new Area {
   val counter = Reg(UInt(log2Up(rxSamplePerBit) bit))
   def reset() = counter := preSamplingSize + (samplingSize - 1) / 2 - 1
   val tick = False
   when(sampler.tick) {
```

```
counter := counter - 1
     when(counter === 0) {
       tick := True
   }
 }
 // Provide bitCounter.value that count up each bitTimer.tick, Used by the state.
→machine to count data bits and stop bits
 // reset() can be called to reset it to zero
 val bitCounter = new Area {
   val value = Reg(UInt(Math.max(dataWidthMax, 2) bit))
   def reset() = value := 0
   when(bitTimer.tick) {
     value := value + 1
   }
 }
 val stateMachine = new Area {
   import UartCtrlRxState._
   val state = RegInit(IDLE)
   val parity = Reg(Bool)
   val shifter = Reg(io.read.payload)
   //Parity calculation
   when(bitTimer.tick) {
     parity := parity ^ sampler.value
   io.read.valid := False
   switch(state) {
     is(IDLE) {
       when(sampler.value === False) {
         state := START
         bitTimer.reset()
       }
     }
     is(START) {
       when(bitTimer.tick) {
         state := DATA
         bitCounter.reset()
         parity := io.configFrame.parity === UartParityType.ODD
         when(sampler.value === True) {
           state := IDLE
         }
       }
     }
     is(DATA) {
       when(bitTimer.tick) {
         shifter(bitCounter.value) := sampler.value
         when(bitCounter.value === io.configFrame.dataLength) {
           bitCounter.reset()
           when(io.configFrame.parity === UartParityType.NONE) {
             state := STOP
```

```
} otherwise {
              state := PARITY
        }
      }
      is(PARITY) {
        when(bitTimer.tick) {
          state := STOP
          bitCounter.reset()
          when(parity =/= sampler.value) {
            state := IDLE
          }
        }
      }
      is(STOP) {
        when(bitTimer.tick) {
          when(!sampler.value) {
            state := IDLE
          }.elsewhen(bitCounter.value === toBitCount(io.configFrame.stop)) {
            state := IDLE
            io.read.valid := True
        }
      }
    }
 io.read.payload := stateMachine.shifter
}
```

# **UartCtrl**

Let's write UartCtrl that instantiates the UartCtrlRx and UartCtrlTx parts, generate the clock divider logic, and connect them to each other.

```
class UartCtrl(g : UartCtrlGenerics = UartCtrlGenerics()) extends Component {
 val io = new Bundle {
   val config = in(UartCtrlConfig(g))
   val write = slave(Stream(Bits(g.dataWidthMax bit)))
   val read = master(Flow(Bits(g.dataWidthMax bit)))
   val uart = master(Uart())
 }
 val tx = new UartCtrlTx(g)
 val rx = new UartCtrlRx(g)
 //Clock divider used by RX and TX
 val clockDivider = new Area {
   val counter = Reg(UInt(g.clockDividerWidth bits)) init(0)
   val tick = counter === 0
   counter := counter - 1
   when(tick) {
     counter := io.config.clockDivider
```

```
tx.io.samplingTick := clockDivider.tick
rx.io.samplingTick := clockDivider.tick

tx.io.configFrame := io.config.frame
rx.io.configFrame := io.config.frame

tx.io.write << io.write
rx.io.read >> io.read

io.uart.txd <> tx.io.txd
io.uart.rxd <> rx.io.rxd
}
```

# Simple usage

To synthesize a UartCtrl as 115200-N-8-1:

```
val uartCtrl: UartCtrl = UartCtrl(
  config = UartCtrlInitConfig(
    baudrate = 115200,
    dataLength = 7,  // 8 bits
    parity = UartParityType.NONE,
    stop = UartStopType.ONE
  )
)
```

If you are using txd pin only:

```
uartCtrl.io.uart.rxd := True  // High is the idle state for UART
txd := uartCtrl.io.uart.txd
```

On the contrary, if you are using rxd pin only:

```
val uartCtrl: UartCtrl = UartCtrl(
  config = UartCtrlInitConfig(
    baudrate = 115200,
    dataLength = 7,  // 8 bits
    parity = UartParityType.NONE,
    stop = UartStopType.ONE
    ),
    readonly = true
)
```

### **Example with test bench**

Here is a top level example that does the followings things:

- Instantiate UartCtrl and set its configuration to 921600 baud/s, no parity, 1 stop bit.
- Each time a byte is received from the UART, it writes it on the leds output.
- Every 2000 cycles, it sends the switches input value to the UART.

```
class UartCtrlUsageExample extends Component{
 val io = new Bundle{
   val uart = master(Uart())
   val switchs = in Bits(8 bits)
   val leds = out Bits(8 bits)
 }
 val uartCtrl = new UartCtrl()
 uartCtrl.io.config.setClockDivider(921600)
 uartCtrl.io.config.frame.dataLength := 7 //8 bits
 uartCtrl.io.config.frame.parity := UartParityType.NONE
 uartCtrl.io.config.frame.stop := UartStopType.ONE
 uartCtrl.io.uart <> io.uart
 //Assign io.led with a register loaded each time a byte is received
 io.leds := uartCtrl.io.read.toReg()
 //Write the value of switch on the uart each 2000 cycles
 val write = Stream(Bits(8 bits))
 write.valid := CounterFreeRun(2000).willOverflow
 write.payload := io.switchs
 write >-> uartCtrl.io.write
}
object UartCtrlUsageExample{
 def main(args: Array[String]) {
    SpinalVhdl(new UartCtrlUsageExample,
→defaultClockDomainFrequency=FixedFrequency(50e6))
}
```

The following example is just a "mad one" but if you want to send a 0x55 header before sending the value of switches, you can replace the write generator of the preceding example by:

```
val write = Stream(Fragment(Bits(8 bits)))
write.valid := CounterFreeRun(4000).willOverflow
write.fragment := io.switchs
write.last := True
write.stage().insertHeader(0x55).toStreamOfFragment >> uartCtrl.io.write
```

Here you can get a simple VHDL testbench for this small UartCtrlUsageExample.

# **Bonus: Having fun with Stream**

If you want to queue data received from the UART:

```
val uartCtrl = new UartCtrl()
val queuedReads = uartCtrl.io.read.toStream.queue(16)
```

If you want to add a queue on the write interface and do some flow control:

```
val uartCtrl = new UartCtrl()
val writeCmd = Stream(Bits(8 bits))
val stopIt = Bool
writeCmd.queue(16).haltWhen(stopIt) >> uartCtrl.io.write
```

# 13.2.3 VGA

#### Introduction

VGA interfaces are becoming an endangered species, but implementing a VGA controller is still a good exercise.

An explanation about the VGA protocol can be found here.

This VGA controller tutorial is based on this implementation.

#### **Data structures**

Before implementing the controller itself we need to define some data structures.

### **RGB** color

First, we need a three channel color structure (Red, Green, Blue). This data structure will be used to feed the controller with pixels and also will be used by the VGA bus.

```
case class RgbConfig(rWidth : Int,gWidth : Int,bWidth : Int){
  def getWidth = rWidth + gWidth + bWidth
}

case class Rgb(c: RgbConfig) extends Bundle{
  val r = UInt(c.rWidth bit)
  val g = UInt(c.gWidth bit)
  val b = UInt(c.bWidth bit)
}
```

### **VGA** bus

io	Driver	Description
name		
vSync	mas-	Vertical synchronization, indicate the beginning of a new frame
	ter	
hSync	mas-	Horizontal synchronization, indicate the beginning of a new line
	ter	
col-	mas-	High when the interface is in the visible part
orEn	ter	
color	mas-	Carry the color, don't care when colorEn is low
	ter	

```
case class Vga (rgbConfig: RgbConfig) extends Bundle with IMasterSlave{
  val vSync = Bool()
  val hSync = Bool()

  val colorEn = Bool()
  val color = Rgb(rgbConfig)

  override def asMaster() : Unit = this.asOutput()
}
```

This Vga Bundle uses the IMasterSlave trait, which allows you to create master/slave VGA interfaces using the following:

```
master(Vga(...))
slave(Vga(...))
```

# **VGA timings**

The VGA interface is driven by using 8 different timings. Here is one simple example of a Bundle that is able to carry them.

```
case class VgaTimings(timingsWidth: Int) extends Bundle {
  val hSyncStart = UInt(timingsWidth bits)
  val hSyncEnd = UInt(timingsWidth bits)
  val hColorStart = UInt(timingsWidth bits)
  val hColorEnd = UInt(timingsWidth bits)
  val vSyncStart = UInt(timingsWidth bits)
  val vSyncEnd = UInt(timingsWidth bits)
  val vColorStart = UInt(timingsWidth bits)
  val vColorEnd = UInt(timingsWidth bits)
}
```

But this not a very good way to specify it because it is redundant for vertical and horizontal timings.

Let's write it in a clearer way:

```
case class VgaTimingsHV(timingsWidth: Int) extends Bundle {
  val colorStart = UInt(timingsWidth bit)
  val colorEnd = UInt(timingsWidth bit)
  val syncStart = UInt(timingsWidth bit)
  val syncEnd = UInt(timingsWidth bit)
}

case class VgaTimings(timingsWidth: Int) extends Bundle {
  val h = VgaTimingsHV(timingsWidth)
  val v = VgaTimingsHV(timingsWidth)
}
```

Then we could add some some functions to set these timings for specific resolutions and frame rates:

```
case class VgaTimingsHV(timingsWidth: Int) extends Bundle {
  val colorStart = UInt(timingsWidth bit)
  val colorEnd = UInt(timingsWidth bit)
  val syncStart = UInt(timingsWidth bit)
  val syncEnd = UInt(timingsWidth bit)
}
```

```
case class VgaTimings(timingsWidth: Int) extends Bundle {
 val h = VgaTimingsHV(timingsWidth)
 val v = VgaTimingsHV(timingsWidth)
 def setAs_h640_v480_r60: Unit = {
   h.syncStart := 96 - 1
   h.syncEnd := 800 - 1
   h.colorStart := 96 + 16 - 1
   h.colorEnd := 800 - 48 - 1
   v.syncStart := 2 - 1
   v.syncEnd := 525 - 1
   v.colorStart := 2 + 10 - 1
   v.colorEnd := 525 - 33 - 1
 }
 def setAs_h64_v64_r60: Unit = {
   h.syncStart := 96 - 1
   h.syncEnd := 800 - 1
   h.colorStart := 96 + 16 - 1 + 288
   h.colorEnd := 800 - 48 - 1 - 288
   v.syncStart := 2 - 1
   v.syncEnd := 525 - 1
   v.colorStart := 2 + 10 - 1 + 208
   v.colorEnd := 525 - 33 - 1 - 208
 }
}
```

# **VGA Controller**

### **Specification**

io	Di-	Description
name	rec-	
	tion	
soft-	in	Reset internal counters and keep the VGA interface inactive
Reset		
tim-	in	Specify VGA horizontal and vertical timings
ings		
pixels	slave	Stream of RGB colors that feeds the VGA controller
error	out	High when the pixels stream is too slow
frameS-	out	High when a new frame starts
tart		
vga	mas-	VGA interface
	ter	

The controller does not integrate any pixel buffering. It directly takes them from the pixels Stream and puts them on the vga.color out at the right time. If pixels is not valid then error becomes high for one cycle.

# Component and io definition

Let's define a new VgaCtrl Component, which takes as RgbConfig and timingsWidth as parameters. Let's give the bit width a default value of 12.

```
class VgaCtrl(rgbConfig: RgbConfig, timingsWidth: Int = 12) extends Component {
  val io = new Bundle {
    val softReset = in Bool
    val timings = in(VgaTimings(timingsWidth))
    val pixels = slave Stream (Rgb(rgbConfig))

  val error = out Bool
  val frameStart = out Bool
  val vga = master(Vga(rgbConfig))
  }
  ...
}
```

# Horizontal and vertical logic

The logic that generates horizontal and vertical synchronization signals is quite the same. It kind of resembles ~PWM~. The horizontal one counts up each cycle, while the vertical one use the horizontal syncronization signal as to increment.

Let's define HVArea, which represents one ~PWM~ and then instantiate it two times: one for both horizontal and vertical syncronization.

```
class VgaCtrl(rqbConfig: RgbConfig, timingsWidth: Int = 12) extends Component {
 val io = new Bundle {...}
 case class HVArea(timingsHV: VgaTimingsHV, enable: Bool) extends Area {
   val counter = Reg(UInt(timingsWidth bit)) init(0)
   val syncStart = counter === timingsHV.syncStart
   val syncEnd = counter === timingsHV.syncEnd
   val colorStart = counter === timingsHV.colorStart
   val colorEnd = counter === timingsHV.colorEnd
   when(enable) {
      counter := counter + 1
      when(syncEnd) {
        counter := 0
      }
   }
               = RegInit(False) setWhen(syncStart) clearWhen(syncEnd)
   val colorEn = RegInit(False) setWhen(colorStart) clearWhen(colorEnd)
   when(io.softReset) {
      counter := 0
      sync
           := False
      colorEn := False
   }
 }
 val h = HVArea(io.timings.h, True)
 val v = HVArea(io.timings.v, h.syncEnd)
}
```

As you can see, it's done by using Area. This is to avoid the creation of a new Component which would have been much more verbose.

#### Interconnections

Now that we have timing generators for horizontal and vertical synchronization, we need to drive the outputs.

```
class VgaCtrl(rgbConfig: RgbConfig, timingsWidth: Int = 12) extends Component {
  val io = new Bundle {...}

  case class HVArea(timingsHV: VgaTimingsHV, enable: Bool) extends Area {...}
  val h = HVArea(io.timings.h, True)
  val v = HVArea(io.timings.v, h.syncEnd)

val colorEn = h.colorEn && v.colorEn
  io.pixels.ready := colorEn
  io.error := colorEn &&! io.pixels.valid

io.frameStart := v.syncEnd

io.vga.hSync := h.sync
  io.vga.vSync := v.sync
  io.vga.colorEn := colorEn
  io.vga.colorEn := colorEn
  io.vga.colorEn := colorEn
  io.vga.color := io.pixels.payload
}
```

#### **Bonus**

The VgaCtrl that was defined above is generic (not application specific). We can imagine a case where the system provides a Stream of Fragment of RGB, which means the system transmits pixels between start/end of picture indications.

In this case we can automatically manage the softReset input by asserting it when an error occurs, then wait for the end of the current pixels picture to deassert error.

Let's add a function to VgaCtrl that can be called from the parent component to feed VgaCtrl by using this Stream of Fragment of RGB.

```
class VgaCtrl(rgbConfig: RgbConfig, timingsWidth: Int = 12) extends Component {
    ...
    def feedWith(that : Stream[Fragment[Rgb]]): Unit ={
        io.pixels << that.toStreamOfFragment

        val error = RegInit(False)
        when(io.error) {
            error := True
        }
        when(that.isLast) {
            error := False
        }
        io.softReset := error
        when(error) {
            that.ready := True
        }
    }
}</pre>
```

# 13.3 Advanced ones

### 13.3.1 JTAG TAP

#### Introduction

**Important:** The goal of this page is to show the implementation of a JTAG TAP (a slave) by a non-conventional way.

### **Important:**

This implementation is not a simple one, it mix object oriented programming, abstract interfaces decoupling, hardware generation and hardware description.

Of course a simple JTAG TAP implementation could be done only with a simple hardware description, but the goal here is really to going forward and creating an very reusable and extensible JTAG TAP generator

Important: This page will not explains how JTAG work. A good tutorial could be find there.

One big difference between commonly used HDL and Spinal, is the fact that SpinalHDL allow you to define hardware generators/builders. It's very different than describing hardware. Let's take a look into the example bellow because the difference between generate/build/describing could seem "playing with word" or could be interpreted differently.

The example bellow is a JTAG TAP which allow the JTAG master to read switchs/keys inputs and write leds outputs. This TAP could also be recognized by a master by using the UID 0x87654321.

```
class SimpleJtagTap extends Component {
 val io = new Bundle {
   val jtag
               = slave(Jtag())
   val switchs = in Bits(8 bit)
   val keys
             = in Bits(4 bit)
   val leds
               = out Bits(8 bit)
 }
 val tap = new JtagTap(io.jtag, 8)
 val idcodeArea = tap.idcode(B"x87654321") (instructionId=4)
 val switchsArea = tap.read(io.switchs)
                                            (instructionId=5)
 val keysArea = tap.read(io.keys)
                                            (instructionId=6)
 val ledsArea = tap.write(io.leds)
                                            (instructionId=7)
```

As you can see, a JtagTap is created but then some Generator/Builder functions (idcode,read,write) are called to create each JTAG instruction. This is what i call "Hardware generator/builder", then these Generator/Builder are used by the user to describing an hardware. And there is the point, in commonly HDL you can only describe your hardware, which imply many donkey job.

This JTAG TAP tutorial is based on this implementation.

#### JTAG bus

First we need to define a JTAG bus bundle.

```
case class Jtag() extends Bundle with IMasterSlave {
  val tms = Bool()
  val tdi = Bool()
  val tdo = Bool()

  override def asMaster() : Unit = {
    out(tdi, tms)
    in(tdo)
  }
}
```

As you can see this bus don't contain the TCK pin because it will be provided by the clock domain.

#### JTAG state machine

Let's define the JTAG state machine as explained here

```
object JtagState extends SpinalEnum {
 val RESET, IDLE,
      IR_SELECT, IR_CAPTURE, IR_SHIFT, IR_EXIT1, IR_PAUSE, IR_EXIT2, IR_UPDATE,
      DR_SELECT, DR_CAPTURE, DR_SHIFT, DR_EXIT1, DR_PAUSE, DR_EXIT2, DR_UPDATE =_
→newElement()
}
class JtagFsm(jtag: Jtag) extends Area {
 import JtagState._
 val stateNext = JtagState()
 val state = RegNext(stateNext) randBoot()
 stateNext := state.mux(
    default
              -> (jtag.tms ? RESET
                                      | IDLE),
                                                          //RESET
               -> (jtag.tms ? DR_SELECT | IDLE),
   TDLE
   IR_SELECT -> (jtag.tms ? RESET | IR_CAPTURE),
    IR_CAPTURE -> (jtag.tms ? IR_EXIT1 | IR_SHIFT),
             -> (jtag.tms ? IR_EXIT1 | IR_SHIFT),
    IR_SHIFT
    IR_EXIT1
              -> (jtag.tms ? IR_UPDATE | IR_PAUSE),
    IR_PAUSE
              -> (jtag.tms ? IR_EXIT2 | IR_PAUSE),
             -> (jtag.tms ? IR_UPDATE | IR_SHIFT),
   IR_EXIT2
   IR_UPDATE -> (jtag.tms ? DR_SELECT | IDLE),
   DR_SELECT -> (jtag.tms ? IR_SELECT | DR_CAPTURE),
   DR_CAPTURE -> (jtag.tms ? DR_EXIT1 | DR_SHIFT),
             -> (jtag.tms ? DR_EXIT1 | DR_SHIFT),
   DR_SHIFT
             -> (jtag.tms ? DR_UPDATE | DR_PAUSE),
   DR_EXIT1
              -> (jtag.tms ? DR_EXIT2 | DR_PAUSE),
   DR_PAUSE
   DR_EXIT2
              -> (jtag.tms ? DR_UPDATE | DR_SHIFT),
   DR_UPDATE -> (jtag.tms ? DR_SELECT | IDLE)
 )
}
```

**Note:** The randBoot() on state make it initialized with a random state. It's only for simulation purpose.

13.3. Advanced ones 233

# **JTAG TAP**

Let's implement the core of the JTAG TAP, without any instruction, just the base manage the instruction register (IR) and the bypass.

```
class JtagTap(val jtag: Jtag, instructionWidth: Int) extends Area{
 val fsm = new JtagFsm(jtag)
 val instruction = Reg(Bits(instructionWidth bit))
 val instructionShift = Reg(Bits(instructionWidth bit))
 val bypass = Reg(Bool)
 jtag.tdo := bypass
 switch(fsm.state) {
    is(JtagState.IR_CAPTURE) {
      instructionShift := instruction
    is(JtagState.IR_SHIFT) {
      instructionShift := (jtag.tdi ## instructionShift) >> 1
      jtag.tdo := instructionShift.lsb
    is(JtagState.IR_UPDATE) {
      instruction := instructionShift
    is(JtagState.DR_SHIFT) {
      bypass := jtag.tdi
 }
}
```

# **Jtag instructions**

Now that the JTAG TAP core is done, we can think about how to implement JTAG instructions by an reusable way.

# **JTAG TAP class interface**

First we need to define how an instruction could interact with the JTAG TAP core. We could of course directly take the JtagTap area, but it's not very nice because is some situation the JTAG TAP core is provided by another IP (Altera virtual JTAG for example).

So let's define a simple and abstract interface between the JTAG TAP core and instructions:

```
trait JtagTapAccess {
  def getTdi : Bool()
  def getTms : Bool()
  def setTdo(value : Bool) : Unit

  def getState : JtagState.T
  def getInstruction() : Bits
  def setInstruction(value : Bits) : Unit
}
```

Then let's the JtagTap implement this abstract interface:

```
//JtagTapAccess impl
override def getTdi: Bool = jtag.tdi
override def setTdo(value: Bool): Unit = jtag.tdo := value
override def getTms: Bool = jtag.tms

override def getState: JtagState.T = fsm.state
override def getInstruction(): Bits = instruction
override def setInstruction(value: Bits): Unit = instruction := value
}
```

#### **Base class**

Let's define a useful base class for JTAG instruction that provide some callback (doCapture/doShift/doUpdate/doReset) depending the selected instruction and the state of the JTAG TAP:

```
class JtagInstruction(tap: JtagTapAccess,val instructionId: Bits) extends Area {
 def doCapture(): Unit = {}
 def doShift(): Unit = {}
 def doUpdate(): Unit = {}
 def doReset(): Unit = {}
 val instructionHit = tap.getInstruction === instructionId
 Component.current.addPrePopTask(() => {
   when(instructionHit) {
      when(tap.getState === JtagState.DR_CAPTURE) {
        doCapture()
      when(tap.getState === JtagState.DR_SHIFT) {
        doShift()
      }
      when(tap.getState === JtagState.DR_UPDATE) {
        doUpdate()
      }
   when(tap.getState === JtagState.RESET) {
      doReset()
    }
 })
}
```

### Note:

About the Component.current.addPrePopTask(...):

This allow you to call the given code at the end of the current component construction. Because of object oriented nature of JtagInstruction, doCapture, doShift, doUpdate and doReset should not be called before children classes construction (because children classes will use it as a callback to do some logic)

13.3. Advanced ones 235

# **Read instruction**

Let's implement an instruction that allow the JTAG to read a signal.

#### Write instruction

Let's implement an instruction that allow the JTAG to write a register (and also read its current value).

```
class JtagInstructionWrite[T <: Data](data: T) (tap: JtagTapAccess,instructionId:_
__Bits) extends JtagInstruction(tap,instructionId) {
  val shifter,store = Reg(Bits(data.getBitsWidth bit))

  override def doCapture(): Unit = {
    shifter := store
  }
  override def doShift(): Unit = {
    shifter := (tap.getTdi ## shifter) >> 1
    tap.setTdo(shifter.lsb)
  }
  override def doUpdate(): Unit = {
    store := shifter
  }
  data.assignFromBits(store)
}
```

# **Idcode** instruction

Let's implement the instruction that return a idcode to the JTAG and also, when a reset occur, set the instruction register (IR) to it own instructionId.

```
class JtagInstructionIdcode[T <: Data](value: Bits)(tap: JtagTapAccess,__
instructionId: Bits)extends JtagInstruction(tap,instructionId) {
  val shifter = Reg(Bits(32 bit))

  override def doShift(): Unit = {
    shifter := (tap.getTdi ## shifter) >> 1
    tap.setTdo(shifter.lsb)
  }

  override def doReset(): Unit = {
```

```
shifter := value
  tap.setInstruction(instructionId)
  }
}
```

# **User friendly wrapper**

Let's add some user friendly function to the JtagTapAccess to make instructions instantiation easier .

```
trait JtagTapAccess {
    ...

def idcode(value: Bits)(instructionId: Bits) =
    new JtagInstructionIdcode(value)(this,instructionId)

def read[T <: Data](data: T)(instructionId: Bits) =
    new JtagInstructionRead(data)(this,instructionId)

def write[T <: Data](data: T, cleanUpdate: Boolean = true, readable: Boolean = true)(instructionId: Bits) =
    new JtagInstructionWrite[T](data,cleanUpdate,readable)(this,instructionId)
}</pre>
```

### **Usage demonstration**

And there we are, we can now very easily create an application specific JTAG TAP without having to write any logic or any interconnections.

```
class SimpleJtagTap extends Component {
 val io = new Bundle {
   val jtag = slave(Jtag())
   val switchs = in Bits(8 bit)
   val keys = in Bits(4 bit)
   val leds = out Bits(8 bit)
 }
 val tap = new JtagTap(io.jtag, 8)
 val idcodeArea = tap.idcode(B"x87654321") (instructionId=4)
 val switchsArea = tap.read(io.switchs)
                                           (instructionId=5)
 val keysArea = tap.read(io.keys)
                                           (instructionId=6)
                                           (instructionId=7)
 val ledsArea = tap.write(io.leds)
}
```

This way of doing things (Generating hardware) could also be applied to, for example, generating an APB/AHB/AXI bus slave.

13.3. Advanced ones 237

# 13.3.2 Memory mapped UART

### Introduction

This example will take the <code>UartCtrl</code> component implemented in the previous <code>example</code> to create a memory mapped UART controller.

# **Specification**

The implementation will be based on the APB3 bus with a RX FIFO.

Here is the register mapping table:

Name	Туре	Access	Address	Description
clockDi-	UInt	RW	0	Set the UartCtrl clock divider
vider				
frame	UartCtrl-	RW	4	Set the dataLength, the parity and the stop bit configuration
	Frame-			
	Config			
writeCmd	Bits	W	8	Send a write command to UartCtrl
write-	Bool	R	8	Bit 0 => zero when a new writeCmd can be sent
Busy				
read	Bool /	R	12	
	Bits			Dita 7 dayyata 0 x ay maylaad
				Bits 7 downto 0 => rx payload
				Bit 31 => rx payload valid

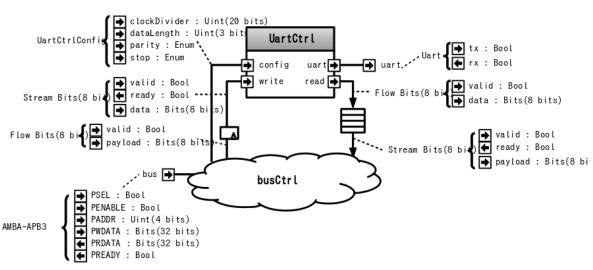
# **Implementation**

For this implementation, the Apb3SlaveFactory tool will be used. It allows you to define a APB3 slave with a nice syntax. You can find the documentation of this tool *there*.

First, we just need to define the Apb3Config that will be used for the controller. It is defined in a Scala object as a function to be able to get it from everywhere.

```
object Apb3UartCtrl{
  def getApb3Config = Apb3Config(
    addressWidth = 4,
    dataWidth = 32
  )
}
```

Then we can define a Apb3UartCtrl component which instantiates a UartCtrl and creates the memory mapping logic between it and the APB3 bus:



```
class Apb3UartCtrl(uartCtrlConfig : UartCtrlGenerics, rxFifoDepth : Int) extends_
→Component{
 val io = new Bundle{
   val bus = slave(Apb3(Apb3UartCtrl.getApb3Config))
   val uart = master(Uart())
 // Instanciate an simple uart controller
 val uartCtrl = new UartCtrl(uartCtrlConfig)
 io.uart <> uartCtrl.io.uart
 // Create an instance of the Apb3SlaveFactory that will then be used as a slave.
→ factory drived by io.bus
 val busCtrl = Apb3SlaveFactory(io.bus)
 // Ask the busCtrl to create a readable/writable register at the address 0
 // and drive uartCtrl.io.config.clockDivider with this register
 busCtrl.driveAndRead(uartCtrl.io.config.clockDivider,address = 0)
 // Do the same thing than above but for uartCtrl.io.config.frame at the address 4
 busCtrl.driveAndRead(uartCtrl.io.config.frame,address = 4)
 // Ask the busCtrl to create a writable Flow[Bits] (valid/payload) at the address 8.
 // Then convert it into a stream and connect it to the uartCtrl.io.write by using.
→an register stage (>->)
 busCtrl.createAndDriveFlow(Bits(uartCtrlConfig.dataWidthMax bits),address = 8).
→toStream >-> uartCtrl.io.write
 // To avoid losing writes commands between the Flow to Stream transformation just.
→above.
 // make the occupancy of the uartCtrl.io.write readable at address 8
 busCtrl.read(uartCtrl.io.write.valid,address = 8)
 // Take uartCtrl.io.read, convert it into a Stream, then connect it to the input of.
→a FIFO of 64 elements
 // Then make the output of the FIFO readable at the address 12 by using a non.
→blocking protocol
 // (Bit 7 downto 0 => read data <br>> Bit 31 => read data valid )
 busCtrl.readStreamNonBlocking(uartCtrl.io.read.toStream.queue(rxFifoDepth),
                                address = 12, validBitOffset = 31, payloadBitOffset = __
```

}

# **Important:**

Yes, that's all it takes. It's also synthesizable.

The Apb3SlaveFactory tool is not something hard-coded into the SpinalHDL compiler. It's something implemented with SpinalHDL regular hardware description syntax.

# 13.3.3 Pinesec

Remember to add it

# 13.3.4 Timer

### Introduction

A timer module is probably one of the most basic pieces of hardware. But even for a timer, there are some interesting things that you can do with SpinalHDL. This example will define a simple timer component which integrates a bus bridging utile.

### **Timer**

So let's start with the Timer component.

# **Specification**

The Timer component will have a single construction parameter:

Parameter Name	Туре	Description	
width	Int	Specify the bit width of the timer counter	ì

And also some inputs/outputs:

IO Name	Direction	Туре	Description
tick	in	Bool	When tick is True, the timer count up until limit.
clear	in	Bool	When tick is True, the timer is set to zero. clear has
			priority over tick.
limit	in	UInt(width bits)	When the timer value is equal to limit, the tick in-
			put is inhibited.
full	out	Bool	full is high when the timer value is equal to limit
			and tick is high.
value	out	UInt(width bits)	Wire out the timer counter value.

# **Implementation**

```
case class Timer(width : Int) extends Component{
 val io = new Bundle{
                = in Bool()
    val tick
   val clear = in Bool()
val limit = in UInt(width bits)
    val full = out Bool()
    val value = out UInt(width bits)
 }
 val counter = Reg(UInt(width bits))
 when(io.tick && !io.full){
    counter := counter + 1
 when(io.clear){
    counter := 0
 }
 io.full := counter === io.limit && io.tick
  io.value := counter
```

# **Bridging function**

Now we can start with the main purpose of this example: defining a bus bridging function. To do that we will use two techniques:

- Using the BusSlaveFactory tool documented *here*
- Defining a function inside the Timer component which can be called from the parent component to drive the Timer's IO in an abstract way.

# **Specification**

This bridging function will take the following parameters:

Pa-	Туре	Description
ram-		
eter		
Name		
busCtrl	Bus-	The BusSlaveFactory instance that will be used by the function to create the bridging
	Slave-	logic.
	Fac-	
	tory	
baseAd-	Big-	The base address where the bridging logic should be mapped.
dress	Int	
ticks	Seq[Boo	plA list of Bool sources that can be used as a tick signal.
clears	Seq[Boo	plA list of Bool sources that can be used as a clear signal.

The register mapping assumes that the bus system is 32 bits wide:

13.3. Advanced ones 241

Name	Ac-	Width	Ad-	Bit	Description
	cess		dress	off-	
			off-	set	
			set		
tick-	RW	len(tic	kØ)	0	Each ticks bool can be actived if the corresponding ticksEnable bit
sEn-					is high.
able					
clears	E <b>iR</b> -W	len(cl	eafors)	16	Each clears bool can be actived if the corresponding clearsEnable
able					bit is high.
limit	RW	width	4	0	
					Access the limit value of the timer component.
					When this register is written to, the timer is cleared.
value		width		0	Access the value of the timer.
clear	W		8		When this register is written to, it clears the timer.

### **Implementation**

Let's add this bridging function inside the Timer component.

```
case class Timer(width : Int) extends Component{
 val io = new Bundle{
   val tick = in Bool()
   val clear = in Bool()
   val limit = in UInt(width bits)
   val full = out Bool()
   val value = out UInt(width bits)
 // Logic previously defined
 // ....
 // The function prototype uses Scala currying funcName(arg1,arg2)(arg3,arg3)
 // which allow to call the function with a nice syntax later
 // This function also returns an area, which allows to keep names of inner signals.
→in the generated VHDL/Verilog.
 def driveFrom(busCtrl : BusSlaveFactory, baseAddress : BigInt)(ticks : Seq[Bool],

¬clears : Seq[Bool]) = new Area {
   //Address 0 => clear/tick masks + bus
   val ticksEnable = busCtrl.createReadWrite(Bits(ticks.length bits),baseAddress +__
\rightarrow 0.0) init(0)
   val clearsEnable = busCtrl.createReadWrite(Bits(clears.length bits),baseAddress +__
\rightarrow 0.16) init(0)
   val busClearing = False
   io.clear := (clearsEnable & clears.asBits).orR | busClearing
   io.tick := (ticksEnable & ticks.asBits ).orR
    //Address 4 => read/write limit (+ auto clear)
   busCtrl.driveAndRead(io.limit,baseAddress + 4)
   busClearing setWhen(busCtrl.isWriting(baseAddress + 4))
    //Address 8 => read timer value / write => clear timer value
   busCtrl.read(io.value,baseAddress + 8)
```

```
busClearing setWhen(busCtrl.isWriting(baseAddress + 8))
}
```

# **Usage**

Here is some demonstration code which is very close to the one used in the Pinsec SoC timer module. Basically it instantiates following elements:

- One 16 bit prescaler
- · One 32 bit timer
- Three 16 bit timers

Then by using an Apb3SlaveFactory and functions defined inside the Timers, it creates bridging logic between the APB3 bus and all instantiated components.

```
val io = new Bundle{
 val apb = Apb3(ApbConfig(addressWidth = 8, dataWidth = 32))
 val interrupt = in Bool()
 val external = new Bundle{
   val tick = Bool()
   val clear = Bool()
 }
}
//Prescaler is very similar to the timer, it mainly integrates a piece of auto reload.
val prescaler = Prescaler(width = 16)
val timerA = Timer(width = 32)
val timerB, timerC, timerD = Timer(width = 16)
val busCtrl = Apb3SlaveFactory(io.apb)
val prescalerBridge = prescaler.driveFrom(busCtrl,0x00)
val timerABridge = timerA.driveFrom(busCtrl,0x40)(
 // The first element is True, which allows you to have a mode where the timer is.
→always counting up.
 ticks = List(True, prescaler.io.overflow),
 // By looping the timer full to the clears, it allows you to create an autoreload.
⊶mode.
 clears = List(timerA.io.full)
)
val timerBBridge = timerB.driveFrom(busCtrl,0x50)(
 //The external.tick could allow to create an impulsion counter mode
 ticks = List(True, prescaler.io.overflow, io.external.tick),
 //external.clear could allow to create an timeout mode.
 clears = List(timerB.io.full, io.external.clear)
val timerCBridge = timerC.driveFrom(busCtrl,0x60)(
 ticks = List(True, prescaler.io.overflow, io.external.tick),
  clears = List(timerC.io.full, io.external.clear)
```

```
val timerDBridge = timerD.driveFrom(busCtrl,0x70)(
   ticks = List(True, prescaler.io.overflow, io.external.tick),
   clears = List(timerD.io.full, io.external.clear)
)

val interruptCtrl = InterruptCtrl(4)
val interruptCtrlBridge = interruptCtrl.driveFrom(busCtrl,0x10)
interruptCtrl.io.inputs(0) := timerA.io.full
interruptCtrl.io.inputs(1) := timerB.io.full
interruptCtrl.io.inputs(2) := timerC.io.full
interruptCtrl.io.inputs(3) := timerD.io.full
io.interrupt := interruptCtrl.io.pendings.orR
```

# 13.4 Introduction

Examples are split into three kinds:

- Simple examples that could be used to get used to the basics of SpinalHDL.
- Intermediate examples which implement components by using a traditional approach.
- Advanced examples which go further than traditional HDL by using object-oriented programming, functional programming, and meta-hardware description.

They are all accessible in the sidebar under the corresponding sections.

**Important:** The SpinalHDL workshop contains many labs with their solutions. See here.

Note: You can also find a list of repositories using SpinalHDL there

# 13.5 Getting started

All examples assume that you have the following imports on the top of your scala file:

```
import spinal.core._
import spinal.lib._
```

To generate VHDL for a given component, you can place the following at the bottom of your scala file:

```
object MyMainObject {
  def main(args: Array[String]) {
    SpinalVhdl(new TheComponentThatIWantToGenerate(constructionArguments)) //Or
    SpinalVerilog
  }
}
```

# **FOURTEEN**

# **LEGACY**

# 14.1 RiscV

**Warning:** This page document the first RISC-V cpu iteration done in SpinalHDL. The second iteration of this CPU is available there and already offer better perforance/area/features.

# 14.1.1 Features

#### RISC-V CPU

- Pipelined on 5 stages (Fetch Decode Execute0 Execute1 WriteBack)
- Multiple branch prediction modes : (disable, static or dynamic)
- Data path parameterizable between fully bypassed to fully interlocked

# Extensions

- One cycle multiplication
- 34 cycle division
- Iterative shifter (N shift -> N cycles)
- Single cycle shifter
- Interruption controller
- Debugging module (with JTAG bridge, openOCD port and GDB)
- Instruction cache with wrapped burst memory interface, one way
- Data cache with instructions to evict/flush the whole cache or a given address, one way

# Performance/Area (on cyclone II)

- small core -> 846 LE, 0.6 DMIPS/Mhz
- debug module (without JTAG) -> 240 LE
- JTAG Avalon master -> 238 LE
- big core with MUL/DIV/Full shifter/I\$/Interrupt/Debug -> 2200 LE, 1.15 DMIPS/Mhz, at least 100 Mhz (with default synthesis option)

# 14.1.2 Base FPGA project

You can find a DE1-SOC project which integrate two instance of the CPU with MUL/DIV/Full shifter/I\$/Interrupt/Debug there :

https://drive.google.com/drive/folders/0B-CqLXDTaMbKNkktb2k3T3lzcUk?usp=sharing

CPU/JTAG/VGA IP are pre-generated. Quartus Prime: 15.1.

# 14.1.3 How to generate the CPU VHDL

**Warning:** This avalon version of the CPU isn't present in recent releases of SpinalHDL. Please considerate the VexRiscv instead.

# 14.1.4 How to debug

You can find the openOCD fork there:

https://github.com/Dolu1990/openocd\_riscv

An example target configuration file could be find there:

https://github.com/Dolu1990/openocd\_riscv/blob/riscv\_spinal/tcl/target/riscv\_spinal.cfg

Then you can use the RISCV GDB.

### 14.1.5 Todo

- Documentation
- Optimise instruction/data caches FMax by moving line hit condition forward into combinatorial paths.

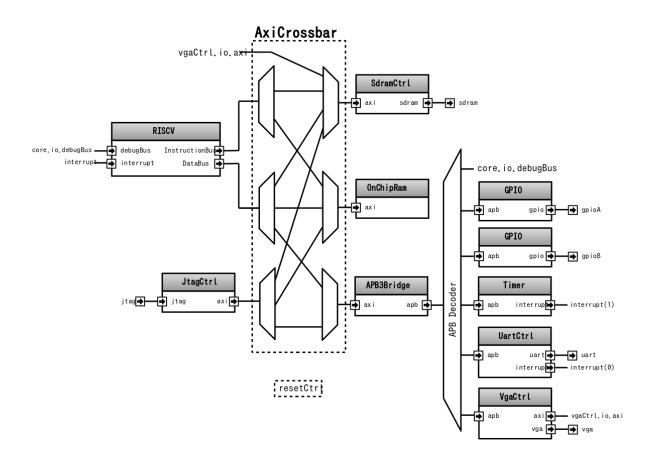
Contact spinalhdl@gmail.com for more information

# 14.2 pinsec

### 14.2.1 Hardware

### Introduction

There is the Pinsec toplevel hardware diagram:



#### **RISCV**

The RISCV is a 5 stage pipelined CPU with following features :

- · Instruction cache
- Single cycle Barrel shifter
- Single cycle MUL, 34 cycle DIV
- Interruption support
- Dynamic branch prediction
- Debug port

## AXI4

As previously said, Pinsec integrate an AXI4 bus fabric. AXI4 is not the easiest bus on the Earth but has many advantages like :

- A flexible topology
- High bandwidth potential
- Potential out of order request completion
- Easy methods to meets clocks timings
- · Standard used by many IP
- An hand-shaking methodology that fit with SpinalHDL Stream.

From an Area utilization perspective, AXI4 is for sure not the lightest solution, but some techniques could dramatically reduce that issue:

- Using Read-Only/Write-Only AXI4 variations where it's possible
- Introducing an Axi4-Shared variation where a new ARW channel is introduced to replace AR and AW channels. This solution divide resources usage by two for the address decoding and the address arbitration.
- Depending the interconnect implementation, if masters doesn't use the R/B channels ready, this path will be removed until each slaves at synthesis, which relax timings.
- As the AXI4 spec suggest, the interconnect can expand the transactions ID by aggregating the corresponding input port id. This allow the interconnect to have an infinite number of pending request and also to support out of order completion with a negligible area cost (transaction id expand).

The Pinsec interconnect doesn't introduce latency cycles.

#### APB3

In Pinsec, all peripherals implement an APB3 bus to be interfaced. The APB3 choice was motivated by following reasons:

- Very simple bus (no burst)
- Use very few resources
- Standard used by many IP

#### Generate the RTL

To generate the RTL, you have multiple solutions:

You can download the SpinalHDL source code, and then run:

```
sbt "project SpinalHDL-lib" "run-main spinal.lib.soc.pinsec.Pinsec"
```

Or you can create your own main into your own SBT project and then run it:

```
import spinal.lib.soc.pinsec._

object PinsecMain{
  def main(args: Array[String]) {
    SpinalVhdl(new Pinsec(100 MHz))
    SpinalVerilog(new Pinsec(100 MHz))
  }
}
```

**Note:** Currently, only the verilog version was tested in simulation and in FPGA because the last release of GHDL is not compatible with cocotb.

## 14.2.2 SoC toplevel (Pinsec)

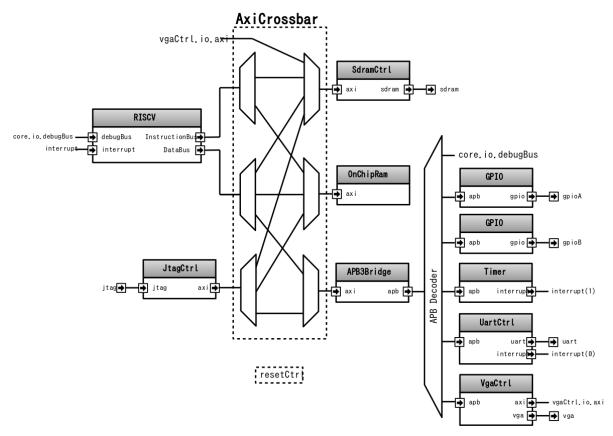
#### Introduction

Pinsec is a little SoC designed for FPGA. It is available in the SpinalHDL library and some documentation could be find *there* 

Its toplevel implementation is an interesting example, because it mix some design pattern that make it very easy to modify. Adding a new master or a new peripheral to the bus fabric could be done in the seconde.

This toplevel implementation could be consulted there: https://github.com/SpinalHDL/SpinalHDL/blob/master/lib/src/main/scala/spinal/lib/soc/pinsec/Pinsec.scala

There is the Pinsec toplevel hardware diagram:



## **Defining all IO**

```
val io = new Bundle{
  //Clocks / reset
  val asyncReset = in Bool()
  val axiClk = in Bool()
  val vgaClk
               = in Bool()
  //Main components IO
  val jtag = slave(Jtag())
               = master(SdramInterface(IS42x320D.layout))
  val sdram
 //Peripherals IO
  val gpioA = master(TriStateArray(32 bits)) //Each pin has it's individual_
→output enable control
 val gpioB = master(TriStateArray(32 bits))
val uart = master(Uart())
                = master(Vga(RgbConfig(5,6,5)))
  val vga
}
```

#### **Clock and resets**

Pinsec has three clocks inputs:

- axiClock
- vgaClock
- jtag.tck

And one reset input:

• asyncReset

Which will finally give 5 ClockDomain (clock/reset couple):

Name	Clock	Description
resetCtrlClock-	axi-	Used by the reset controller, Flops of this clock domain are initialized by the FPGA
Domain	Clock	bitstream
axiClockDo-	axi-	Used by all component connected to the AXI and the APB interconnect
main	Clock	
coreClockDo-	axi-	The only difference with the axiClockDomain, is the fact that the reset could also
main	Clock	be asserted by the debug module
vgaClockDo-	vga-	Used by the VGA controller backend as a pixel clock
main	Clock	
jtagClockDo-	jtag.tck	Used to clock the frontend of the JTAG controller
main		

#### Reset controller

First we need to define the reset controller clock domain, which has no reset wire, but use the FPGA bitstream loading to setup flipflops.

```
val resetCtrlClockDomain = ClockDomain(
  clock = io.axiClk,
  config = ClockDomainConfig(
    resetKind = BOOT
  )
)
```

Then we can define a simple reset controller under this clock domain.

```
val resetCtrl = new ClockingArea(resetCtrlClockDomain) {
  val axiResetUnbuffered = False
  val coreResetUnbuffered = False

  //Implement an counter to keep the reset axiResetOrder high 64 cycles
  // Also this counter will automaticly do a reset when the system boot.
  val axiResetCounter = Reg(UInt(6 bits)) init(0)
  when(axiResetCounter =/= U(axiResetCounter.range -> true)){
    axiResetCounter := axiResetCounter + 1
    axiResetUnbuffered := True
  }
  when(BufferCC(io.asyncReset)){
    axiResetCounter := 0
  }

  //When an axiResetOrder happen, the core reset will as well
  when(axiResetUnbuffered){
```

```
coreResetUnbuffered := True
}

//Create all reset used later in the design
val axiReset = RegNext(axiResetUnbuffered)
val coreReset = RegNext(coreResetUnbuffered)
val vgaReset = BufferCC(axiResetUnbuffered)
}
```

## Systems clock domains

Now that the reset controller is implemented, we can define clock domain for all part of Pinsec:

```
val axiClockDomain = ClockDomain(
        = io.axiClk,
 clock
 reset
           = resetCtrl.axiReset,
 frequency = FixedFrequency(50 MHz) //The frequency information is used by the SDRAM_
\hookrightarrow controller
)
val coreClockDomain = ClockDomain(
 clock = io.axiClk,
 reset = resetCtrl.coreReset
)
val vgaClockDomain = ClockDomain(
 clock = io.vgaClk,
 reset = resetCtrl.vgaReset
val jtagClockDomain = ClockDomain(
 clock = io.jtag.tck
```

Also all the core system of Pinsec will be defined into a axi clocked area:

```
val axi = new ClockingArea(axiClockDomain) {
   //Here will come the rest of Pinsec
}
```

## **Main components**

Pinsec is constituted mainly by 4 main components:

- One RISCV CPU
- One SDRAM controller
- One on chip memory
- One JTAG controller

#### **RISCV CPU**

The RISCV CPU used in Pinsec as many parametrization possibilities:

```
val core = coreClockDomain {
 val coreConfig = CoreConfig(
   pcWidth = 32,
    addrWidth = 32.
    startAddress = 0x000000000,
   regFileReadyKind = sync,
   branchPrediction = dynamic,
   bypassExecute0 = true,
   bypassExecute1 = true,
   bypassWriteBack = true,
   bypassWriteBackBuffer = true,
   collapseBubble = false,
    fastFetchCmdPcCalculation = true,
    dynamicBranchPredictorCacheSizeLog2 = 7
 )
 //The CPU has a systems of plugin which allow to add new feature into the core.
 //Those extension are not directly implemented into the core, but are kind of ...
→additive logic patch defined in a separated area.
 coreConfig.add(new MulExtension)
 coreConfig.add(new DivExtension)
 coreConfig.add(new BarrelShifterFullExtension)
 val iCacheConfig = InstructionCacheConfig(
    cacheSize =4096,
   bytePerLine =32,
   wayCount = 1, //Can only be one for the moment
   wrappedMemAccess = true,
   addressWidth = 32,
   cpuDataWidth = 32,
   memDataWidth = 32
 )
 //There is the instanciation of the CPU by using all those construction parameters
 new RiscvAxi4(
   coreConfig = coreConfig,
    iCacheConfig = iCacheConfig,
    dCacheConfig = null,
   debug = true,
    interruptCount = 2
 )
}
```

## On chip RAM

The instanciation of the AXI4 on chip RAM is very simple.

In fact it's not an AXI4 but an Axi4Shared, which mean that a ARW channel replace the AR and AW ones. This solution use less area while being fully interoperable with full AXI4.

```
val ram = Axi4SharedOnChipRam(
  dataWidth = 32,
  byteCount = 4 KiB,
  idWidth = 4  //Specify the AXI4 ID width.
)
```

## **SDRAM** controller

First you need to define the layout and timings of your SDRAM device. On the DE1-SOC, the SDRAM device is an IS42x320D one.

```
object IS42x320D {
 def layout = SdramLayout(
   bankWidth = 2,
   columnWidth = 10,
             = 13,
   rowWidth
    dataWidth
               = 16
 def timingGrade7 = SdramTimings(
   bootRefreshCount =
   tPOW
                    = 100 \text{ us,}
                   = 64 ms,
   tREF
   tRC
                    = 60 ns,
    tRFC
                    = 60 ns,
    tRAS
                    = 37 ns,
   tRP
                    = 15 ns,
   tRCD
                    = 15 ns,
    cMRD
                    = 2,
    tWR
                    = 10 ns,
    cWR
                       1
 )
```

Then you can used those definition to parametrize the SDRAM controller instantiation.

```
val sdramCtrl = Axi4SharedSdramCtrl(
  axiDataWidth = 32,
  axiIdWidth = 4,
  layout = IS42x320D.layout,
  timing = IS42x320D.timingGrade7,
  CAS = 3
)
```

#### **JTAG** controller

The JTAG controller could be used to access memories and debug the CPU from an PC.

```
val jtagCtrl = JtagAxi4SharedDebugger(SystemDebuggerConfig(
  memAddressWidth = 32,
  memDataWidth = 32,
  remoteCmdWidth = 1,
  jtagClockDomain = jtagClockDomain
))
```

## **Peripherals**

Pinsec integrate some peripherals:

- GPIO
- Timer
- UART
- VGA

#### **GPIO**

```
val gpioACtrl = Apb3Gpio(
   gpioWidth = 32
)

val gpioBCtrl = Apb3Gpio(
   gpioWidth = 32
)
```

#### **Timer**

The Pinsec timer module is constituted of:

- One prescaler
- One 32 bits timer
- Three 16 bits timers

All of them are packed into the PinsecTimerCtrl component.

```
val timerCtrl = PinsecTimerCtrl()
```

#### **UART** controller

First we need to define a configuration for our UART controller:

Then we can use it to instantiate the UART controller

```
val uartCtrl = Apb3UartCtrl(uartCtrlConfig)
```

#### **VGA** controller

First we need to define a configuration for our VGA controller:

Then we can use it to instantiate the VGA controller

```
val vgaCtrl = Axi4VgaCtrl(vgaCtrlConfig)
```

#### **Bus interconnects**

There is three interconnections components:

- AXI4 crossbar
- AXI4 to APB3 bridge
- · APB3 decoder

#### **AXI4 to APB3 bridge**

This bridge will be used to connect low bandwidth peripherals to the AXI crossbar.

```
val apbBridge = Axi4SharedToApb3Bridge(
  addressWidth = 20,
  dataWidth = 32,
  idWidth = 4
)
```

#### **AXI4** crossbar

The AXI4 crossbar that interconnect AXI4 masters and slaves together is generated by using an factory. The concept of this factory is to create it, then call many function on it to configure it, and finally call the build function to ask the factory to generate the corresponding hardware:

First you need to populate slaves interfaces:

```
// Slave -> (base address, size) ,

axiCrossbar.addSlaves(
  ram.io.axi -> (0x00000000L, 4 KiB),
   sdramCtrl.io.axi -> (0x40000000L, 64 MiB),
   apbBridge.io.axi -> (0xF0000000L, 1 MiB)
)
```

Then you need to populate interconnections between slaves and masters:

Then to reduce combinatorial path length and have a good design FMax, you can ask the factory to insert pipelining stages between itself a given master or slave :

#### Note:

halfPipe / >> / << / >/-> in the following code are provided by the Stream bus library.

Some documentation could be find *there*. In short, it's just some pipelining and interconnection stuff.

```
//Pipeline the connection between the crossbar and the apbBridge.io.axi
axiCrossbar.addPipelining(apbBridge.io.axi,(crossbar,bridge) => {
  crossbar.sharedCmd.halfPipe() >> bridge.sharedCmd
  crossbar.writeData.halfPipe() >> bridge.writeData
  crossbar.writeRsp << bridge.writeRsp</pre>
```

#### **APB3** decoder

The interconnection between the APB3 bridge and all peripherals is done via an APB3Decoder:

```
val apbDecoder = Apb3Decoder(
   master = apbBridge.io.apb,
   slaves = List(
      gpioACtrl.io.apb -> (0x00000, 4 KiB),
      gpioBCtrl.io.apb -> (0x01000, 4 KiB),
      uartCtrl.io.apb -> (0x10000, 4 KiB),
      timerCtrl.io.apb -> (0x20000, 4 KiB),
      vgaCtrl.io.apb -> (0x30000, 4 KiB),
      core.io.debugBus -> (0xF0000, 4 KiB)
)
```

#### Misc

To connect all toplevel IO to components, the following code is required:

```
io.gpioA <> axi.gpioACtrl.io.gpio
io.gpioB <> axi.gpioBCtrl.io.gpio
io.jtag <> axi.jtagCtrl.io.jtag
io.uart <> axi.uartCtrl.io.uart
io.sdram <> axi.sdramCtrl.io.sdram
io.vga <> axi.vgaCtrl.io.vga
```

And finally some connections between components are required like interrupts and core debug module resets

```
core.io.interrupt(0) := uartCtrl.io.interrupt
core.io.interrupt(1) := timerCtrl.io.interrupt

core.io.debugResetIn := resetCtrl.axiReset
when(core.io.debugResetOut){
   resetCtrl.coreResetUnbuffered := True
}
```

## 14.2.3 Introduction

**Note:** This page document the SoC implemented with the first RISC-V cpu iteration done in SpinalHDL. The second iteration of this SoC (and CPU) is available there and offer better perforance/area/features.

## Introduction

Pinsec is the name of a little FPGA SoC fully written in SpinalHDL. Goals of this project are multiple:

- Prove that SpinalHDL is a viable HDL alternative in non-trivial projects.
- Show advantage of SpinalHDL meta-hardware description capabilities in a concrete project.
- Provide a fully open source SoC.

Pinsec has followings hardware features:

- AXI4 interconnect for high speed busses
- APB3 interconnect for peripherals
- RISCV CPU with instruction cache, MUL/DIV extension and interrupt controller
- JTAG bridge to load binaries and debug the CPU
- SDRAM SDR controller
- On chip ram
- One UART controller
- One VGA controller
- Some timer module
- Some GPIO

The toplevel code explanation could be find there

#### **Board support**

A DE1-SOC FPGA project can be find there with some demo binaries.

#### 14.2.4 Software

#### **RISCV** tool-chain

Binaries executed by the CPU can be defined in ASM/C/C++ and compiled by the GCC RISCV fork. Also, to load binaries and debug the CPU, an OpenOCD fork and RISCV GDB can be used.

RISCV tools: https://github.com/riscv/riscv-wiki/wiki/RISC-V-Software-Status

OpenOCD fork: https://github.com/Dolu1990/openocd\_riscv Software examples: https://github.com/Dolu1990/pinsecSoftware

## OpenOCD/GDB/Eclipse configuration

About the OpenOCD fork, there is the configuration file that could be used to connect the Pinsec SoC: https://github.com/Dolu1990/openocd\_riscv/blob/riscv\_spinal/tcl/target/riscv\_spinal.cfg

There is an example of arguments used to run the OpenOCD tool:

```
openocd -f ../tcl/interface/ftdi/ft2232h_breakout.cfg -f ../tcl/target/riscv_spinal. \rightarrow cfg -d 3
```

To debug with eclipse, you will need the Zylin plugin and then create an "Zynlin embedded debug (native)".

Initialize commands:

```
target remote localhost:3333
monitor reset halt
load
```

Run commands:

continue

# **DEVELOPERS AREA**

# 15.1 Bus Slave Factory Implementation

## 15.1.1 Introduction

This page will document the implementation of the BusSlaveFactory tool and one of those variant. You can get more information about the functionality of that tool *there*.

## 15.1.2 Specification

The class diagram is the following:



 $\label{thm:continuous} The \ Bus Slave Factory\ abstract\ class\ define\ minimum\ requirements\ that\ each\ implementation\ of\ it\ should\ provide\ .$ 

Name	Description	
busDataWidth	Return the data width of the bus	
read(that,address,	band the bus read the address, fill the response with that at bitOffset	
write(that,address	s, MACDETISETHE bus write the address, assign that with bus's data from bitOffset	
on-	Call doThat when a write transaction occur on address	
Write(address)(de	oThat)	
on-	Call doThat when a read transaction occur on address	
Read(address)(doThat)		
nonStop-	Permanently assign that by the bus write data from bitOffset	
Write(that,bitOffset)		

By using them the  ${\tt BusSlaveFactory}$  should also be able to provide many utilities:

Name	Re-	Description
1.4 1	turn	M1 11 1 211 1 211 1 1 1 1 1 1 1 1 1 1 1
readAnd-	1 : 00	Make that readable and writable at address and placed at bitOffset in the word
Write(that,addre	ess,bitOff	(set)
readMulti-		
Word(that,addre	SS)	Create the memory mapping to read that from 'address'.:
		If that is bigger than one word it extends the register on followings addresses
writeMulti-		
Word(that,addre	ss)	Create the memory mapping to write that at 'address'.:
		If that is bigger than one word it extends the register on followings addresses
cre-	T	Create a write only register of type dataType at address and placed at bitOffset
ateWriteOnly(da	ataType,a	adidréhe, bit (Odfset)
createRead-	T	Create a read write register of type dataType at address and placed at bitOffset
Write(dataType,	address,	pit Office tyord
create-	Flow[T	Create a writable Flow register of type dataType at address and placed at
AndDrive-		bitOffset in the word
Flow(dataType,		
	ss,bitOff	selDrive that with a register writable at address placed at bitOffset in the word
driveAn-		Drive that with a register writable and readable at address placed at bitOffset
dRead(that,addr	ess,bitOt	
drive-		Emit on that a transaction when a write happen at address by using data placed at
Flow(that,addre	ss,bitOff	selb)itOffset in the word
		Read that and consume the transaction when a read happen at address. valid <= validBitOffset bit
validBitOffset,p	ayloadBi	t@axload <= payloadBitOffset+widthOf(payload) downto payloadBitOffset
doBitsAccumul	ationAnd	Instanciate an internal register which at each cycle do : Clear OnRead reg := reg   that
(that,address,bit	Offset)	Then when a read occur, the register is cleared. This register is readable at address and placed at bitOffset in the word

About BusSlaveFactoryDelayed, it's still an abstract class, but it capture each primitives (BusSlaveFactoryElement) calls into a data-model. This datamodel is one list that contain all primitives, but also a HashMap that link each address used to a list of primitives that are using it. Then when they all are collected (at the end of the current component), it do a callback that should be implemented by classes that extends it. The implementation of this callback should implement the hardware corresponding to all primitives collected.

## 15.1.3 Implementation

## **BusSlaveFactory**

Let's describe primitives abstract function:

Then let's operate the magic to implement all utile based on them:

```
trait BusSlaveFactory extends Area{
 //...
 def readAndWrite(that : Data,
                   address: BigInt,
                   bitOffset : Int = 0): Unit = {
   write(that,address,bitOffset)
   read(that,address,bitOffset)
 }
 def drive(that : Data,
            address : BigInt,
            bitOffset : Int = 0) : Unit = {
   val reg = Reg(that)
   write(reg,address,bitOffset)
   that := reg
 def driveAndRead(that : Data,
                   address : BigInt,
                   bitOffset : Int = 0) : Unit = {
   val reg = Reg(that)
   write(reg,address,bitOffset)
   read(reg,address,bitOffset)
    that := reg
 }
 def driveFlow[T <: Data](that : Flow[T],</pre>
                           address: BigInt,
                           bitOffset : Int = 0) : Unit = {
```

```
that.valid := False
   onWrite(address){
     that.valid := True
   nonStopWrite(that.payload,bitOffset)
 def createReadWrite[T <: Data](dataType: T,</pre>
                                 address: BigInt,
                                 bitOffset : Int = 0): T = {
   val reg = Reg(dataType)
   write(reg,address,bitOffset)
   read(reg,address,bit0ffset)
   rea
 }
 def createAndDriveFlow[T <: Data](dataType : T,</pre>
                                 address: BigInt,
                                 bitOffset : Int = 0) : Flow[T] = {
   val flow = Flow(dataType)
   driveFlow(flow,address,bitOffset)
   flow
 }
 def doBitsAccumulationAndClearOnRead(
                                           that: Bits,
                                           address : BigInt,
                                           bitOffset : Int = 0): Unit = {
   assert(that.getWidth <= busDataWidth)</pre>
   val reg = Reg(that)
   reg := reg | that
   read(reg,address,bitOffset)
   onRead(address){
     reg := that
 }
 def readStreamNonBlocking[T <: Data] (that : Stream[T],</pre>
                                         address: BigInt,
                                         validBitOffset : Int,
                                        payloadBitOffset : Int) : Unit = {
   that.ready := False
   onRead(address){
     that.ready := True
   read(that.valid ,address,validBitOffset)
   read(that.payload,address,payloadBitOffset)
 def readMultiWord(that : Data,
               address : BigInt) : Unit = {
   val wordCount = (widthOf(that) - 1) / busDataWidth + 1
   val valueBits = that.asBits.resize(wordCount*busDataWidth)
   val words = (0 until wordCount).map(id => valueBits(id * busDataWidth ,__
→busDataWidth bit))
   for (wordId <- (0 until wordCount)) {</pre>
     read(words(wordId), address + wordId*busDataWidth/8)
```

```
}
 }
 def writeMultiWord(that : Data,
                 address : BigInt) : Unit = {
   val wordCount = (widthOf(that) - 1) / busDataWidth + 1
    for (wordId <- (0 until wordCount)) {</pre>
      write(
        that = new DataWrapper{
          override def getBitsWidth: Int =
            Math.min(busDataWidth, widthOf(that) - wordId * busDataWidth)
          override def assignFromBits(value : Bits): Unit = {
            that.assignFromBits(
                      = value.resized,
              bits
              offset = wordId * busDataWidth,
              bitCount = getBitsWidth bits)
        },address = address + wordId * busDataWidth / 8,0
   }
 }
}
```

#### **BusSlaveFactoryDelayed**

Let's implement classes that will be used to store primitives:

```
trait BusSlaveFactoryElement
// Ask to make `that` readable when a access is done on `address`.
// bitOffset specify where `that` is placed on the answer
case class BusSlaveFactoryRead(that : Data,
                               address : BigInt,
                               bitOffset : Int) extends BusSlaveFactoryElement
// Ask to make `that` writable when a access is done on `address`.
// bitOffset specify where `that` get bits from the request
case class BusSlaveFactoryWrite(that : Data,
                                address : BigInt,
                                bitOffset : Int) extends BusSlaveFactoryElement
// Ask to execute `doThat` when a write access is done on `address`
case class BusSlaveFactoryOnWrite(address : BigInt,
                                  doThat : () => Unit) extends BusSlaveFactoryElement
// Ask to execute `doThat` when a read access is done on `address`
case class BusSlaveFactoryOnRead( address : BigInt,
                                  doThat : () => Unit) extends BusSlaveFactoryElement
// Ask to constantly drive `that` with the data bus
// bitOffset specify where `that` get bits from the request
case class BusSlaveFactoryNonStopWrite(that : Data,
                                       bitOffset : Int) extends BusSlaveFactoryElement
```

Then let's implement the BusSlaveFactoryDelayed itself:

```
trait BusSlaveFactoryDelayed extends BusSlaveFactory{
 // elements is an array of all BusSlaveFactoryElement requested
 val elements = ArrayBuffer[BusSlaveFactoryElement]()
 // elementsPerAddress is more structured than elements, it group all.
→BusSlaveFactoryElement per requested addresses
 val elementsPerAddress = collection.mutable.HashMap[BigInt,
→ArrayBuffer[BusSlaveFactoryElement]]()
 private def addAddressableElement(e : BusSlaveFactoryElement,address : BigInt) = {
    elements += e
    elementsPerAddress.getOrElseUpdate(address,_
→ArrayBuffer[BusSlaveFactoryElement]()) += e
 override def read(that : Data,
          address : BigInt,
          bitOffset : Int = 0) : Unit = {
   assert(bitOffset + that.getBitsWidth <= busDataWidth)</pre>
   addAddressableElement(BusSlaveFactoryRead(that,address,bitOffset),address)
 }
 override def write(that : Data,
           address : BigInt,
           bitOffset : Int = 0) : Unit = {
   assert(bitOffset + that.getBitsWidth <= busDataWidth)</pre>
    addAddressableElement(BusSlaveFactoryWrite(that,address,bitOffset),address)
 }
 def onWrite(address : BigInt)(doThat : => Unit) : Unit = {
    addAddressableElement(BusSlaveFactoryOnWrite(address,() => doThat),address)
 def onRead (address : BigInt)(doThat : => Unit) : Unit = {
   addAddressableElement(BusSlaveFactoryOnRead(address,() => doThat),address)
 def nonStopWrite( that : Data,
                    bitOffset : Int = 0) : Unit = {
   assert(bitOffset + that.getBitsWidth <= busDataWidth)</pre>
   elements += BusSlaveFactoryNonStopWrite(that,bitOffset)
 //This is the only thing that should be implement by class that extends.
→ BusSlaveFactoryDelayed
 def build() : Unit
 component.addPrePopTask(() => build())
```

#### **AvalonMMSlaveFactory**

First let's implement the companion object that provide the compatible AvalonMM configuration object that correspond to the following table :

Pin name	Туре	Description
read	Bool High one cycle to produce a read request	
write	Bool	High one cycle to produce a write request
address	UInt(addressWidth bits)	Byte granularity but word aligned
writeData	Bits(dataWidth bits)	
readDataValid	Bool	High to respond a read command
readData	Bool(dataWidth bits)	Valid when readDataValid is high

```
object AvalonMMSlaveFactory{
 def getAvalonConfig( addressWidth : Int,
                       dataWidth : Int) = {
    AvalonMMConfig.pipelined(
                              //Create a simple pipelined configuration of the
→ Avalon Bus
      addressWidth = addressWidth.
      dataWidth = dataWidth
                               //Change some parameters of the configuration
   ).copy(
     useByteEnable = false,
      useWaitRequestn = false
   )
 }
 def apply(bus : AvalonMM) = new AvalonMMSlaveFactory(bus)
}
```

Then, let's implement the AvalonMMSlaveFactory itself.

```
class AvalonMMSlaveFactory(bus : AvalonMM) extends BusSlaveFactoryDelayed{
 assert(bus.c == AvalonMMSlaveFactory.getAvalonConfig(bus.c.addressWidth,bus.c.
→dataWidth))
 val readAtCmd = Flow(Bits(bus.c.dataWidth bits))
 val readAtRsp = readAtCmd.stage()
 bus.readDataValid := readAtRsp.valid
 bus.readData := readAtRsp.payload
 readAtCmd.valid := bus.read
 readAtCmd.payload := 0
 override def build(): Unit = {
    for(element <- elements) element match {</pre>
      case element : BusSlaveFactoryNonStopWrite =>
        element.that.assignFromBits(bus.writeData(element.bitOffset, element.that.
→getBitsWidth bits))
     case _ =>
    for((address, jobs) <- elementsPerAddress){</pre>
      when(bus.address === address){
        when(bus.write){
          for(element <- jobs) element match{</pre>
            case element : BusSlaveFactoryWrite => {
```

```
element.that.assignFromBits(bus.writeData(element.bitOffset, element.
→that.getBitsWidth bits))
            case element : BusSlaveFactoryOnWrite => element.doThat()
            case _ =>
          }
        }
        when(bus.read){
          for(element <- jobs) element match{</pre>
            case element : BusSlaveFactoryRead => {
              readAtCmd.payload(element.bitOffset, element.that.getBitsWidth bits) := __
⇔element.that.asBits
            case element : BusSlaveFactoryOnRead => element.doThat()
          }
        }
      }
    }
 }
 override def busDataWidth: Int = bus.c.dataWidth
}
```

## 15.1.4 Conclusion

That's all, you can check one example that use this Apb3SlaveFactory to create an Apb3UartCtrl there.

If you want to add the support of a new memory bus, it's very simple you just need to implement another variation of the BusSlaveFactoryDelayed trait. The Apb3SlaveFactory is probably a good starting point:D

## 15.2 How to HACK this documentation

If you want to add your page to this documentation you need to add your source file in the appropriate section. I opted to create a structure that resample the various section of the documentation, this is not strictly necessary, but for clarity sake, highly encourage.

This documentation uses a recursive index tree: every folder have a special index.rst files that tell sphinx witch file, and in what order put it in the documentation tree.

## 15.2.1 Title convention

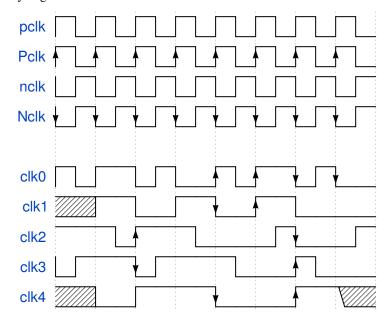
Sphinx is very smart, the document structure is deduced from how you use non alphanumerical characters (like: = - `: ' "  $\sim$  ^ \_ \* + # < >), you only need to be consistent. Still, for consistency sakes we use this progression:

- = over and underline for section titles
- = underline for titles
- - underline for paragraph
- · ^ for subparagraph

## 15.2.2 Wavedrom integration

This documentation makes use of the sphinxcontrib-wavedrom plugin, So you can specify a timing diagram, or a register description with the WaveJSON syntax like so:

and you get:



**Note:** if you want the Wavedrom diagram to be present in the pdf export, you need to use the "non relaxed" JSON dialect. long story short, no javascript code and use " around key value (Eg. "name").

you can describe register mapping with the same syntax:

```
{"reg":[
    {"bits": 8, "name": "things"},
    {"bits": 2, "name": "stuff" },
    {"bits": 6}
],
    "config": { "bits":16,"lanes":1 }
}
```

```
15 10 9 8 7 0 stuff things
```

#### 15.2.3 New section

if you want to add a new section you need to specify in the top index, the index file of the new section. I suggest to name the folder like the section name, but is not required; Sphinx will take the name of the section from the title of the index file.

## example

I want to document the new feature in SpinalHDL, and I want to create a section for it; let's call it Cheese So I need to create a folder named Cheese (name is not important), and in it create a index file like:

```
cheese
=====

.. toctree::
:glob:
introduction
*
```

**Note:** The .. toctree:: directive accept some parameters, in this case :glob: makes so you can use the \* to include all the remaining files.

Note: The file path is relative to the index file, if you want to specify the absolute path, you need to prepend /

**Note:** introduction.rst will be always the first on the list because it's specified in the index file. Other files will be included in alphabetical order.

Now I can add the introduction.rst and other files like cheddar.rst, stilton.rst, etc.

The only thing remaining to do is to add cheese to the top index file like so:

```
Welcome to SpinalHDL's documentation!
.. toctree::
   :maxdepth: 2
   :titlesonly:
  rst/About SpinalHDL/index
  rst/Getting Started/index
  rst/Data types/index
   rst/Structuring/index
   rst/Semantic/index
   rst/Sequential logic/index
   rst/Design errors/index
   rst/Other language features/index
  rst/Libraries/index
  rst/Simulation/index
   rst/Examples/index
   rst/Legacy/index
```

rst/Developers area/index
rst/Cheese/index

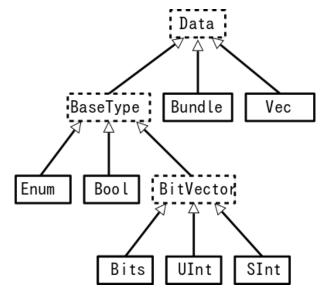
that's it, now you can add all you want in cheese and all pages will show up in the documentation.

# **15.3 Types**

## 15.3.1 Introduction

The language provides 5 base types and 2 composite types that can be used.

- Base types: Bool, Bits, UInt for unsigned integers, SInt for signed integers, Enum.
- Composite types : Bundle, Vec.



Those types and their usage (with examples) are explained hereafter.

About the fixed point support it's documented there

## 15.3.2 Bool

This is the standard boolean type that correspond to a bit.

## **Declaration**

The syntax to declare such as value is as follows:

Syntax	Description	Return
Bool()	Create a Bool	Bool
True	Create a Bool assigned with true	Bool
False	Create a Bool assigned with false	Bool
Bool(value : Boolean)	Create a Bool assigned with a Scala Boolean	Bool

Using this type into SpinalHDL yields:

15.3. Types 271

## **Operators**

The following operators are available for the Bool type

Operator	Description	Return
		type
!x	Logical NOT	Bool
	Logical AND	Bool
x && y		
x & y		
	Logical OR	Bool
II		
x    y		
x   y		
x ^ y	Logical XOR	Bool
x.set[()]	Set x to True	
x.clear[()]	Set x to False	
x.rise[()]	Return True when x was low at the last cycle and is now high	Bool
x.rise(initAt :	Same as x.rise but with a reset value	Bool
Bool)		
x.fall[()]	Return True when x was high at the last cycle and is now low	Bool
x.fall(initAt : Bool)	Same as x.fall but with a reset value	Bool
x.setWhen(cond)	Set x when cond is True	Bool
x.clearWhen(cond)	Clear x when cond is True	Bool

# 15.3.3 The BitVector family - (Bits, UInt, SInt)

BitVector is a family of types for storing multiple bits of information in a single value. This type has three subtypes that can be used to model different behaviours:

Bits do not convey any sign information whereas the UInt (unsigned integer) and SInt (signed integer) provide the required operations to compute correct results if signed / unsigned arithmetics is used.

## **Declaration syntax**

Syntax	Description	Re-	
		turn	
Bits/UInt/SInt [()]	Create a BitVector, bits count is inferred	Bits/UIn	t/SInt
Bits/UInt/SInt(x bits)	Create a BitVector with x bits	Bits/UIn	t/SInt
B/U/S(value : Int[,width : BitCount])	Create a BitVector assigned with 'value'	Bits/UIn	t/SInt
B/U/S"[[size']base]value"	Create a BitVector assigned with 'value'	Bits/UIn	t/SInt
B/U/S([x bits], element,)	Create a BitVector assigned with the value spec-	Bits/UIn	t/SInt
	ified by elements (see bellow table)		

Elements could be defined as follows:

Element syntax	Description
x : Int -> y : Boolean/Bool	Set bit x with y
x : Range -> y : Boolean/Bool	Set each bits in range x with y
x : Range -> y : T	Set bits in range x with y
x : Range -> y : String	
	Set bits in range x with y The string format follow same rules than B/U/S"xyz" one
x : Range -> y : T	Set bits in range x with y
default -> y : Boolean/Bool	
	Set all unconnected bits with the y value.
	This feature could only be use to do assignments without the U/B/S prefix

You can define a Range values

Range syntax	Description	Width
(x downto y)	[x:y] x >= y	x-y+1
(x to y)	$[x:y] x \le y$	y-x+1
(x until y)	[x:y[ x < y	у-х

```
val myUInt = UInt(8 bits)
myUInt := U(2,8 bits)
myUInt := U(2)
myUInt := U"0000_0101" // Base per default is binary => 5
myUInt := U"h1A"
                        // Base could be x (base 16)
                                        h (base 16)
                        //
                        //
                                         d (base 10)
                                         o (base 8)
                        //
                                         b (base 2)
myUInt := U"8'h1A"
                        // You can use scala Int as literal value
myUInt := 2
val myBool := myUInt === U(7 \rightarrow true, (6 downto 0) \rightarrow false)
val myBool := myUInt === U(myUInt.range -> true)
//For assignment purposes, you can omit the B/U/S, which also alow the use of the
→[default -> ???] feature
myUInt := (default -> true)
                                                   //Assign myUInt with "11111111"
                                                   //Assign myUInt with "11111111"
myUInt := (myUInt.range -> true)
myUInt := (7 -> true, default -> false)
                                                  //Assign myUInt with "10000000"
myUInt := ((4 downto 1) -> true,default -> false) //Assign myUInt with "00011110"
```

15.3. Types 273

## **Operators**

Operator	Description	Return
~X	Bitwise NOT	T(w(x) bits)
x & y	Bitwise AND	$T(\max(w(x), w(y))$
		bits)
x   y	Bitwise OR	$T(\max(w(x), w(y))$
		bits)
x ^ y	Bitwise XOR	$T(\max(w(x), w(y)))$
		bits)
x(y)	Readbit, y: Int/UInt	Bool
x(hi,lo)	Read bitfield, hi: Int, lo: Int	T(hi-lo+1 bits)
x(offset,width)	Read bitfield, offset: UInt, width: Int	T(width bits)
x(y) := z	Assign bits, y: Int/UInt	Bool
x(hi,lo) := z	Assign bitfield, hi: Int, lo: Int	T(hi-lo+1 bits)
x(offset, width) := z	Assign bitfield, offset: UInt, width: Int	T(width bits)
x.msb	Return the most significant bit	Bool
x.lsb	Return the least significant bit	Bool
x.range	Return the range (x.high downto 0)	Range
x.high	Return the upper bound of the type x	Int
x.xorR	XOR all bits of x	Bool
x.orR	OR all bits of x	Bool
x.andR	AND all bits of x	Bool
x.clearAll[()]	Clear all bits	T
x.setAll[()]	Set all bits	T
x.setAllTo(value : Boolean)	Set all bits to the given Boolean value	
x.setAllTo(value : Bool)	Set all bits to the given Bool value	
x.asBools	Cast into a array of Bool	Vec(Bool,width(x))

## **Masked comparison**

Some time you need to check equality between a BitVector and a bits constant that contain hole (don't care values).

There is an example about how to do that:

```
val myBits = Bits(8 bits)
val itMatch = myBits === M"00--10--"
```

# 15.3.4 Bits

Operator	Description	Return
x >> y	Logical shift right, y: Int	T(w(x) - y  bits)
x >> y	Logical shift right, y: UInt	T(w(x) bits)
x << y	Logical shift left, y: Int	T(w(x) + y  bits)
x << y	Logical shift left, y: UInt	T(w(x) + max(y) bits)
x.rotateLeft(y)	Logical left rotation, y: UInt	T(w(x))
x.resize(y)	Return a resized copy of x, filled with zero, y: Int	T(y bits)

# 15.3.5 UInt, SInt

Operator	Description	Return
x + y	Addition	$T(\max(w(x), w(y) \text{ bits})$
x - y	Subtraction	$T(\max(w(x), w(y) \text{ bits})$
x * y	Multiplication	T(w(x) + w(y) bits)
x > y	Greater than	Bool
x >= y	Greater than or equal	Bool
x < y	Less than	Bool
x <= y	Less than or equal	Bool
x >> y	Arithmetic shift right, y: Int	T(w(x) - y  bits)
x >> y	Arithmetic shift right, y: UInt	T(w(x) bits)
x << y	Arithmetic shift left, y: Int	T(w(x) + y  bits)
x << y	Arithmetic shift left, y: UInt	T(w(x) + max(y) bits)
x.resize(y)	Return an arithmetic resized copy of x, y: Int	T(y bits)

# 15.3.6 Bool, Bits, Ulnt, SInt

Operator	Description	Return
x.asBits	Binary cast in Bits	Bits(w(x) bits)
x.asUInt	Binary cast in UInt	UInt(w(x) bits)
x.asSInt	Binary cast in SInt	SInt(w(x) bits)

## 15.3.7 Vec

Declaration	Description
Vec(type : Data, size : Int)	Create a vector of size time the given type
Vec(x,y,)	
	Create a vector where indexes point to given elements. this construct support mixed element width

Operator	Description	Return
x(y)	Read element y, y: Int/UInt	T
x(y) := z	Assign element y with z, y: Int/UInt	

15.3. Types 275

#### 15.3.8 Bundle

Bundles could be used to model data structure line buses and interfaces.

All attributes that extends Data (Bool, Bits, UInt, ...) that are defined inside the bundle are considered as part of the bundle.

## Simple example (RGB/VGA)

The following example show an RGB bundle definition with some internal function.

```
case class RGB(channelWidth : Int) extends Bundle{
  val red = UInt(channelWidth bits)
  val green = UInt(channelWidth bits)
  val blue = UInt(channelWidth bits)

def isBlack : Bool = red === 0 && green === 0 && blue === 0
  def isWhite : Bool = {
   val max = U((channelWidth-1 downto 0) -> true)
   return red === max && green === max && blue === max
  }
}
```

Then you can also incorporate a Bundle inside Bundle as deeply as you want:

```
case class VGA(channelWidth : Int) extends Bundle{
  val hsync = Bool
  val vsync = Bool
  val color = RGB(channelWidth)
}
```

And finaly instanciate your Bundles inside the hardware:

If you want to specify your bundle as an input or an output of a Component, you have to do it by the following way :

#### Interface example (APB)

If you want to define an interface, let's imagine an APB interface, you can also use bundles :

```
class APB(addressWidth: Int,
         dataWidth: Int,
         selWidth : Int,
         useSlaveError : Boolean) extends Bundle {
 val PADDR
                = UInt(addressWidth bit)
 val PSEL
                = Bits(selWidth bits)
 val PENABLE
                = Bool
 val PREADY
               = Bool
 val PWRITE
               = Bool
 val PWDATA = Bits(dataWidth bit)
 val PRDATA = Bits(dataWidth bit)
 val PSLVERROR = if(useSlaveError) Bool() else null //This wire is created only_
→when useSlaveError is true
// Example of usage :
val bus = APB(addressWidth = 8,
             dataWidth = 32,
             selWidth = 4.
             useSlaveError = false)
```

One good practice is to group all construction parameters inside a configuration class. This could make the parametrization much easier later in your components, especially if you have to reuse the same configuration at multiple places. Also if one time you need to add another construction parameter, you will only have to add it into the configuration class and everywhere this one is instantiated:

```
case class APBConfig(addressWidth: Int,
                    dataWidth: Int,
                     selWidth : Int,
                    useSlaveError : Boolean)
class APB(val config: APBConfig) extends Bundle { //[val] config, make the_
val PADDR = UInt(config.addressWidth bit)
 val PSEL
                = Bits(config.selWidth bits)
 val PENABLE = Bool
 val PREADY = Bool
 val PWRITE = Bool
             = Bits(config.dataWidth bit)
= Bits(config.dataWidth bit)
 val PWDATA
 val PRDATA
 val PSLVERROR = if(config.useSlaveError) Bool() else null
// Example of usage
val apbConfig = APBConfig(addressWidth = 8,dataWidth = 32,selWidth = 4,useSlaveError_
\rightarrow= false)
val busA = APB(apbConfig)
val busB = APB(apbConfig)
```

Then at some points, you will probably need to use the APB bus as master or as slave interface of some components. To do that you can define some functions :

15.3. Types 277

```
import spinal.core._
case class APBConfig(addressWidth: Int,
                     dataWidth: Int,
                     selWidth : Int,
                     useSlaveError : Boolean)
class APB(val config: APBConfig) extends Bundle {
 val PADDR
                = UInt(config.addressWidth bit)
 val PSEL
                = Bits(config.selWidth bits)
 val PENABLE = Bool
 val PREADY = Bool
 val PWRITE = Bool
               = Bits(config.dataWidth bit)
 val PWDATA
 val PRDATA = Bits(config.dataWidth bit)
 val PSLVERROR = if(config.useSlaveError) Bool() else null
 def asMaster(): this.type = {
   out (PADDR, PSEL, PENABLE, PWRITE, PWDATA)
   in(PREADY,PRDATA)
   if(config.useSlaveError) in(PSLVERROR)
    this
 }
 def asSlave(): this.type = this.asMaster().flip() //Flip reverse all in out_
\rightarrow configuration.
// Example of usage
val apbConfig = APBConfig(addressWidth = 8,dataWidth = 32,selWidth = 4,useSlaveError_
\rightarrow= false)
val io = new Bundle{
 val masterBus = APB(apbConfig).asMaster()
 val slaveBus = APB(apbConfig).asSlave()
```

Then to make that better, the spinal.lib integrate a small master slave utile named IMasterSlave. When a bundle extends IMasterSlave, it should implement/override the asMaster function. It give you the ability to setup a master or a slave interface by a smoother way:

There is an example of an APB bus that implement this IMasterSlave:

```
class APB(val config: APBConfig) extends Bundle with IMasterSlave {
 val PADDR
             = UInt(addressWidth bit)
 val PSEL
                = Bits(selWidth bits)
 val PENABLE = Bool
 val PREADY = Bool
 val PWRITE
              = Bool
 val PWDATA = Bits(dataWidth bit)
 val PRDATA = Bits(dataWidth bit)
 val PSLVERROR = if(useSlaveError) Bool() else null //This wire is created only.
→when useSlaveError is true
 override def asMaster() : Unit = {
   out(PADDR, PSEL, PENABLE, PWRITE, PWDATA)
   in(PREADY,PRDATA)
   if(useSlaveError) in(PSLVERROR)
 //The asSlave is by default the flipped version of asMaster.
```

## 15.3.9 Enum

SpinalHDL support enumeration with some encodings:

En-	Bit	Description
cod-	width	
ing		
native		Use the VHDL enumeration system, this is the default encoding
bina-	log2Up	state:Chitatto store states in declaration order (value from 0 to n-1)
rySe-		
quan-		
cial		
bina-	state-	Use Bits to store state. Each bit correspond to one state
ry-	Count	
One-		
Hot		

Define a enumeration type:

```
object UartCtrlTxState extends SpinalEnum { // Or_
    SpinalEnum(defaultEncoding=encodingOfYouChoice)
    val sIdle, sStart, sData, sParity, sStop = newElement()
}
```

Instantiate a enumeration signal and assign it:

```
val stateNext = UartCtrlTxState() // Or UartCtrlTxState(encoding=encodingOfYouChoice)
stateNext := UartCtrlTxState.sIdle

//You can also import the enumeration to have the visibility on its elements
import UartCtrlTxState._
stateNext := sIdle
```

15.3. Types 279

# 15.3.10 Data (Bool, Bits, Ulnt, SInt, Enum, Bundle, Vec)

All hardware types extends the Data class, which mean that all of them provide following operators :

Operator	Description	Return
x === y	Equality	Bool
x =/= y	Inequality	Bool
x.getWidth	Return bitcount	Int
x ## y	Concatenate, x->high, y->low	Bits(width(x) + width(y))
		bits)
Cat(x)	Concatenate list, first element on lsb, x : Ar-	Bits(sumOfWidth bits)
	ray[Data]	
Mux(cond,x,y)	if cond? x: y	$T(\max(w(x), w(y) \text{ bits})$
x.asBits	Cast in Bits	Bits(width(x) bits)
x.assignFromBits(bits)	Assign from Bits	
x.assignFromBits(bits,hi,lo)	Assign bitfield, hi: Int, lo: Int	T(hi-lo+1 bits)
x.assignFromBits(bits,offset,width)Assign bitfield, offset: UInt, width: Int		T(width bits)
x.getZero	Get equivalent type assigned with zero	T

# 15.3.11 Literals as signal declaration

Literals are generally use as a constant value. But you can also use them to do two things in a single one :

• Define a wire which is assigned with a constant value

There is an example:

## WELCOME TO SPINALHDL'S DOCUMENTATION!

# 16.1 Site purpose and structure

This site presents the *SpinalHDL* language and how to use it on concrete examples.

If you are learning the language from scratch, this presentation is probably a good starting point.

# 16.2 What is SpinalHDL?

SpinalHDL is an open source high-level hardware description language. It can be used as an alternative to VHDL or Verilog and has several advantages over them.

Also, SpinalHDL is not an HLS approach. Its goal is not to push something abstract into flip-flops and gates, but by using simple elements (flip-flops, gates, if / case statments) create a new abstraction level and help the designer to reuse their code and not write the same thing over and over again.

**Note:** SpinalHDL is *fully interoperable* with standard VHDL/Verilog-based EDA tools (simulators and synthetizers) as the output generated by the toolchain could be VHDL or Verilog. It also enables mixed designs where SpinalHDL components inter-operate with VHDL or Verilog IPs.

## 16.2.1 Advantages of using SpinalHDL over VHDL / Verilog

As SpinalHDL is based on a high-level language, it provides several advantages to improve your hardware coding:

- 1. No more endless wiring Create and connect complex buses like AXI in one single line.
- 2. Evolving capabilities Create your own bus definitions and abstraction layers.
- 3. **Reduce code size** By a high factor, especially for wiring. This enables you to have a better overview of your code base, increase your productivity and create fewer headaches.
- 4. **Free and user friendly IDE** Thanks to Scala tools for auto-completion, error highlighting, navigation shortcuts, and many others.
- 5. **Powerful and easy type conversions** Bidirectional translation between any data type and bits. Useful when loading a complex data structure from a CPU interface.
- 6. Loop detection Tools check that there are no combinatorial loops / latches.
- 7. Clock domain safety The tools inform you that there are no unintentional clock domain crossings.
- 8. **Generic design** There are no restrictions to the genericity of your hardware description by using Scala constructs.

## **16.2.2 License**

SpinalHDL uses two licenses, one for spinal.core, and one for spinal.lib.

spinal.core (the compiler) is under the LGPL license, which could be summarized with following statements:

- You can make money with your SpinalHDL description and its generated RTL.
- You don't have to share your SpinalHDL description and its generated RTL.
- There are no fees and no royalties.
- If your make improvements to the SpinalHDL core, please share your modifications to make the tool better for everybody.

spinal.lib (a general purpose library of components/tools/interfaces) is under the permissive MIT license.

# 16.3 Getting started

Want to try it for yourself? Then jump to the getting started section and have fun!

## 16.4 Links

SpinalHDL repository: https://github.com/SpinalHDL/SpinalHDL
A short show case (PDF): motivation.pdf
Presentation of the language (PDF): presentation.pdf
SBT base project: https://github.com/SpinalHDL/SpinalTemplateSbt
Jupyter bootcamp: https://github.com/SpinalHDL/Spinal-bootcamp
Workshop: https://github.com/SpinalHDL/SpinalWorkshop
VexRiscv CPU and SoC: https://github.com/SpinalHDL/VexRiscv
StackOverflow (tag: SpinalHDL): StackOverflow
Google group:

https://groups.google.com/forum/#! forum/spinalhdl-hardware-description-language



16.4. Links 283