Algorithms and Datastructures

Graphs, Depth-/Breadth-first Search, Graph-Connectivity

Albert-Ludwigs-Universität Freiburg

UNI FREIBURG

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Bioinformatics Group / Department of Computer Science

Algorithms and Datastructures, January 2017

Structure

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Feedback

- Exercises
- Lecture

Graphs

- Introduction
- Implementation
- Application example

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Feedback from the exercises

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The upcoming exercise sheet 12 and 13 will be merged together (finding largest connected component + Dijkstra)

Some people were asking for more solution sheets for the exercises

We are working on it.

Feedback from the lecture

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Code in the lecture will be a litte bit different from exercise sheet.

One person asked for additional explanations regarding proofs.

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Graphs - Overview:

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- Breadth first search (BFS)
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- Connected components of a graph

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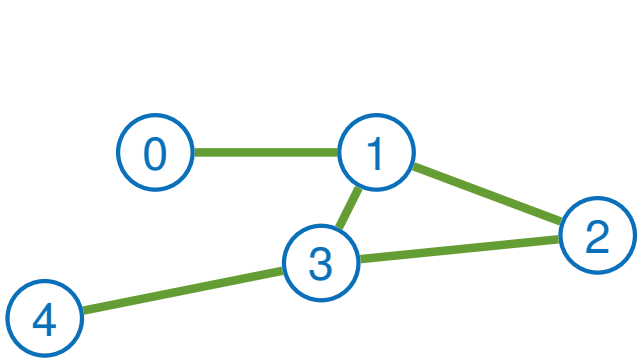

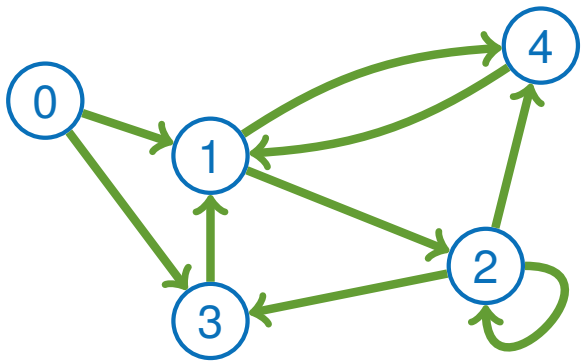

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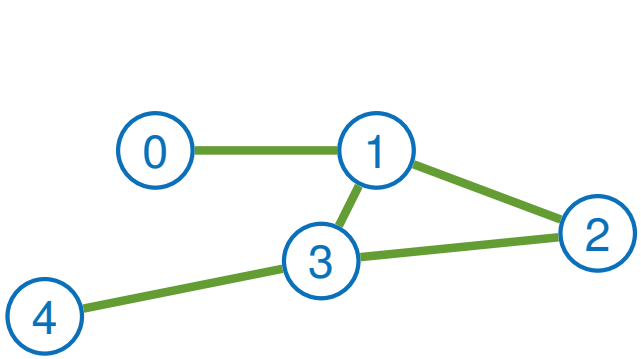

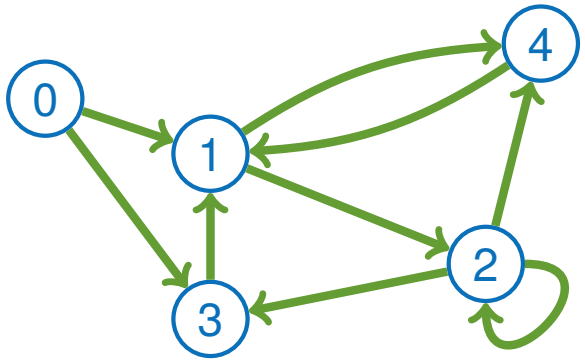

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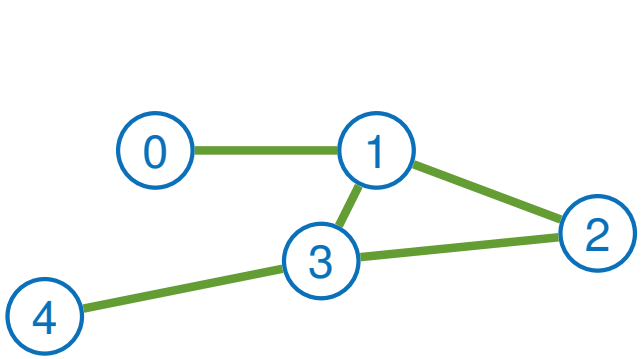

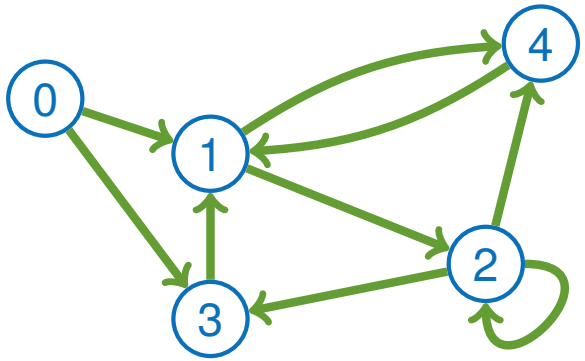

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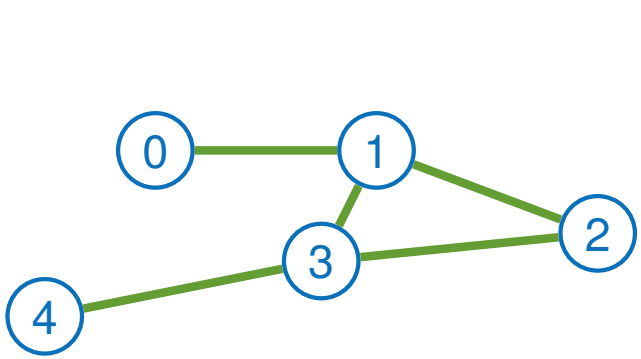

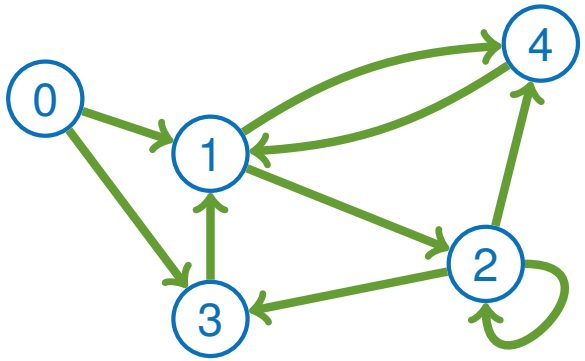

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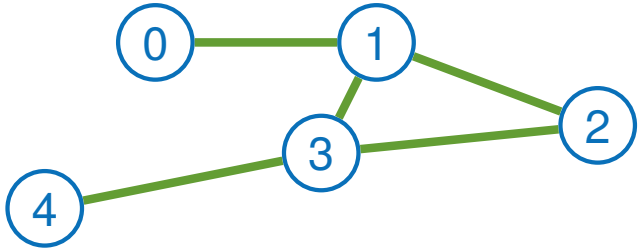

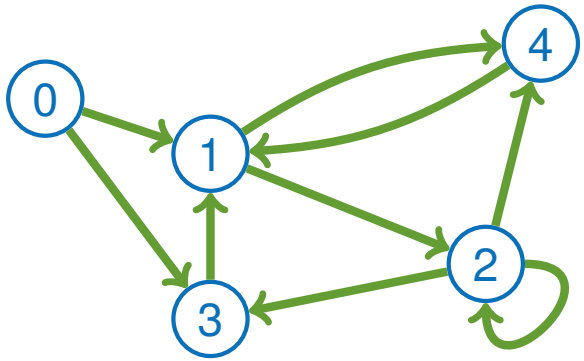

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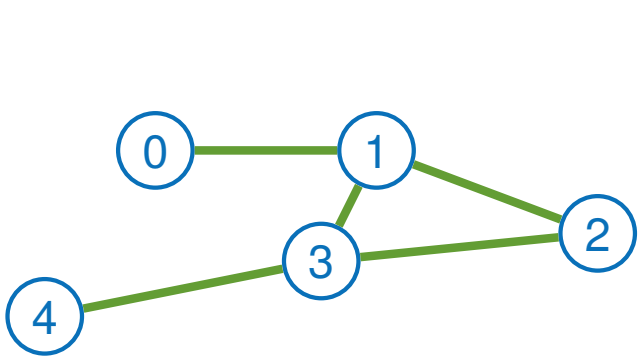

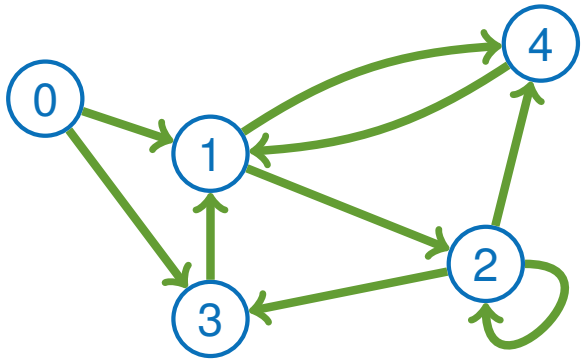

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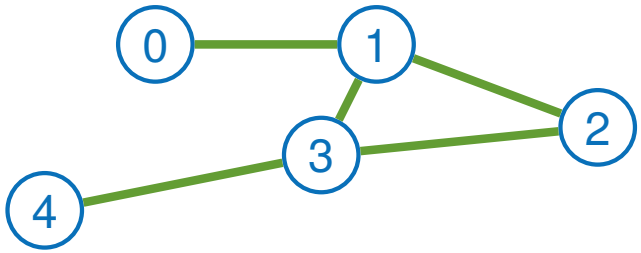

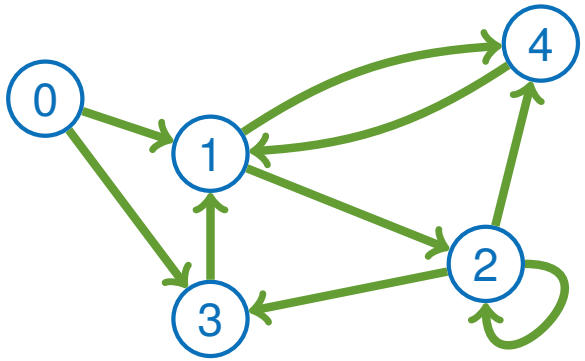

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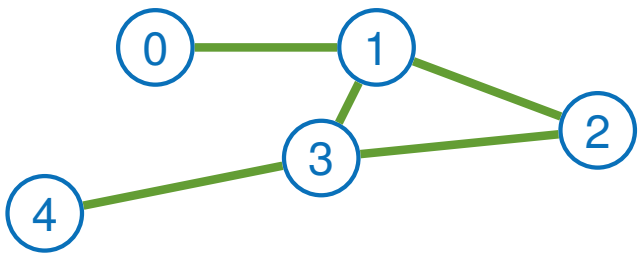

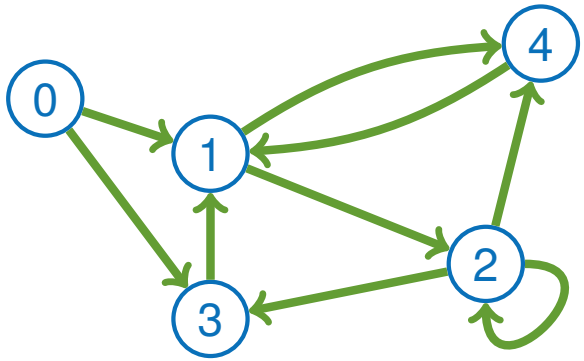

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- Each edge connects two vertices ($u, v \in V$)
  - Undirected edge: $e = \{u, v\}$ (set)
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- Self-loops are also possible: $e = (u, u)$ or $e = \{u, u\}$

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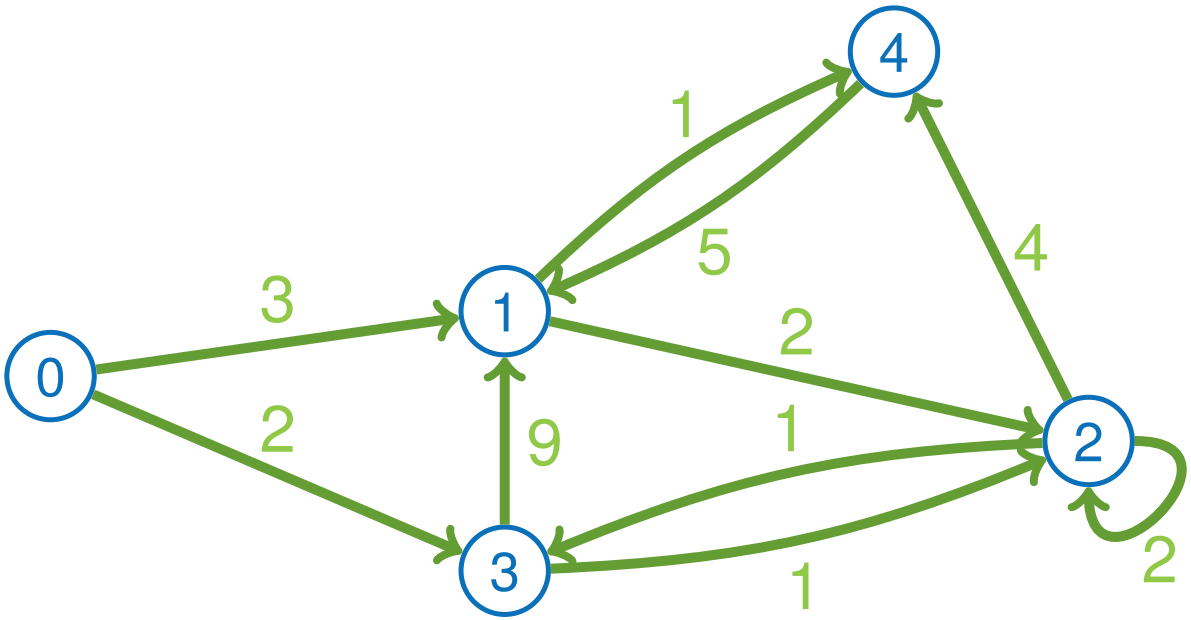

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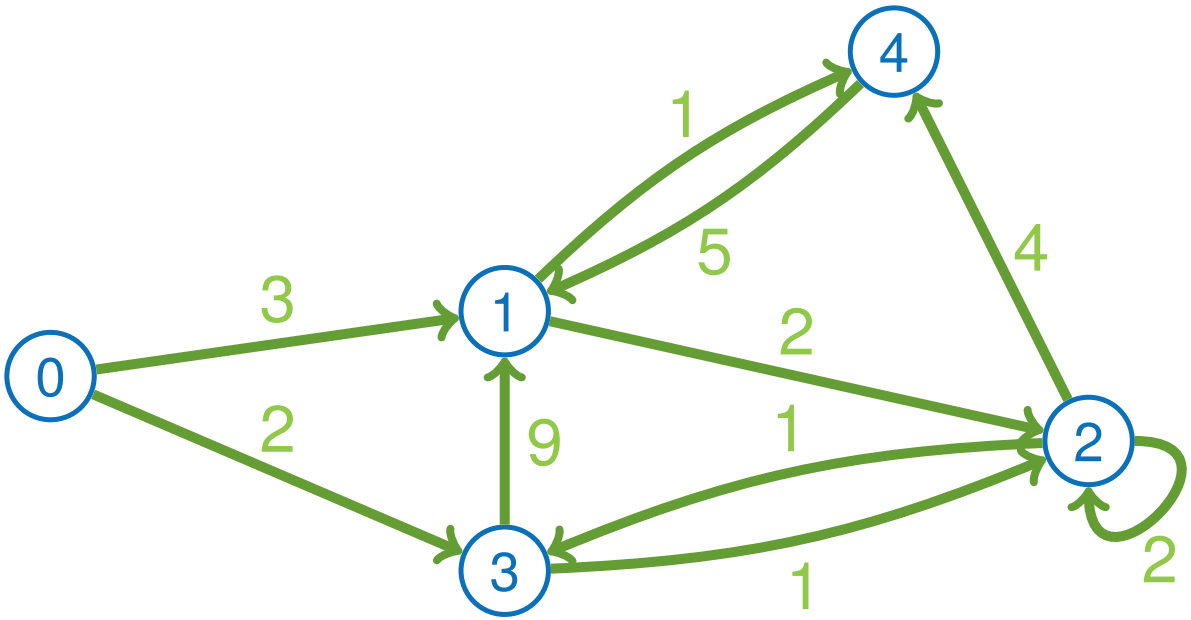

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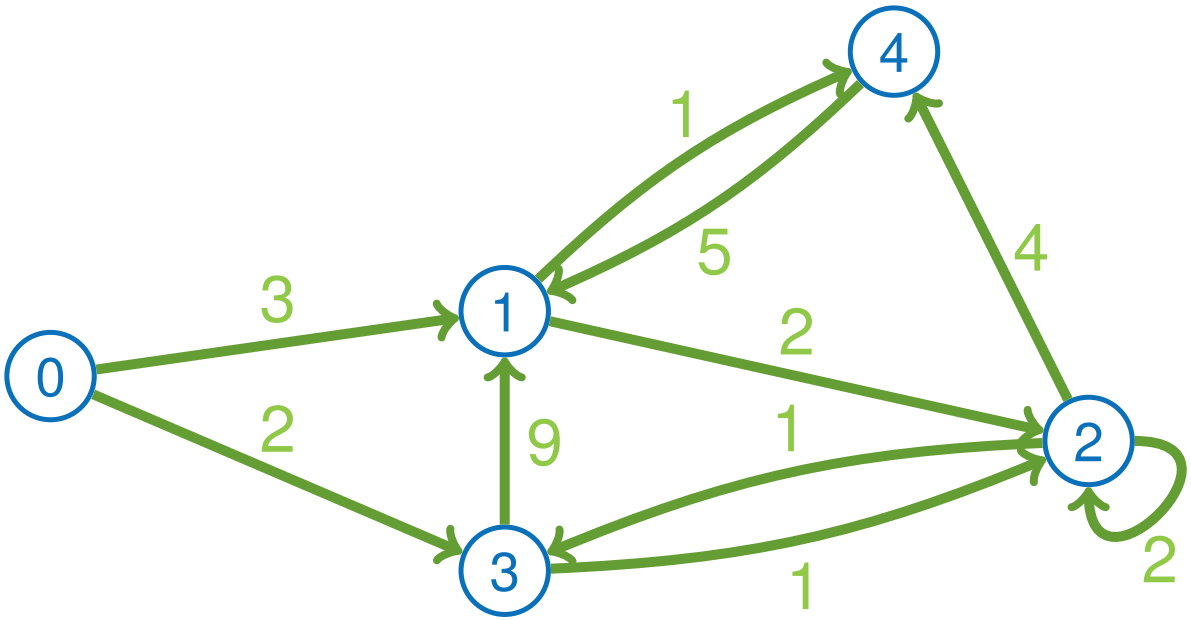

- Each edge is marked with a real number named weight
- The weight is also named length or cost of the edge depending on the application

Example: Road network

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- Intersections:
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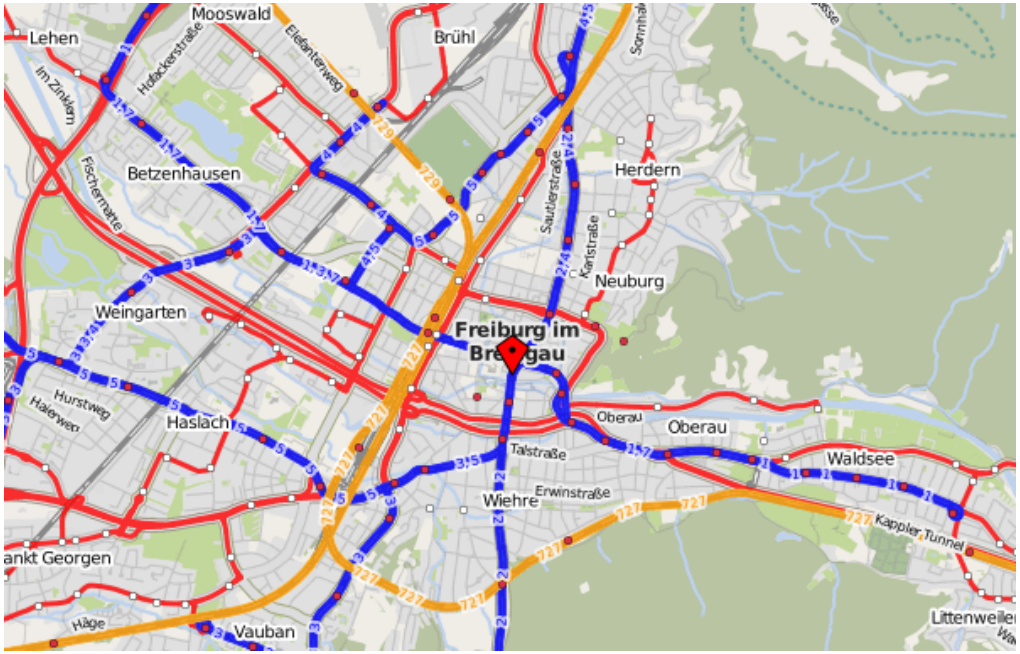

Abbildung: Map of Freiburg © OpenStreet-Map

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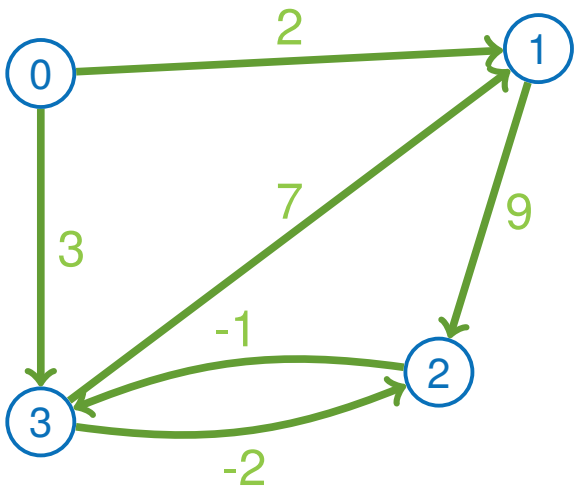

Abbildung: Weighted graph with $|V| = 4$, $|E| = 6$

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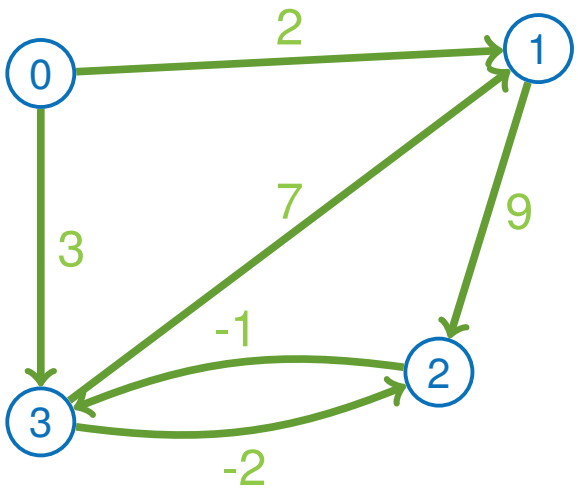

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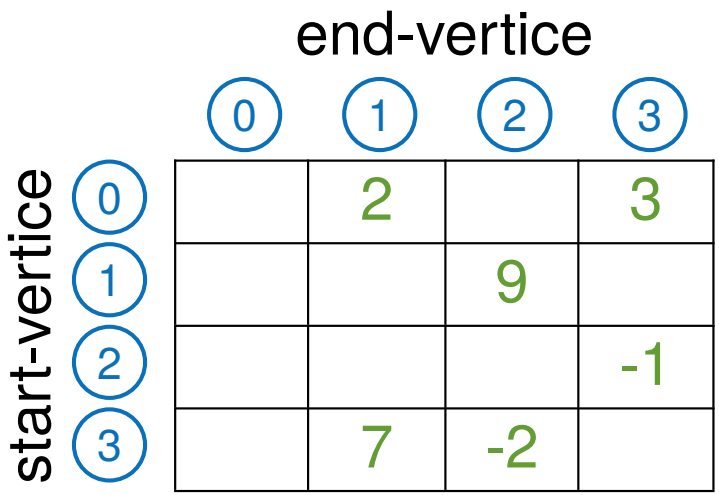

Abbildung: Adjacency matrix

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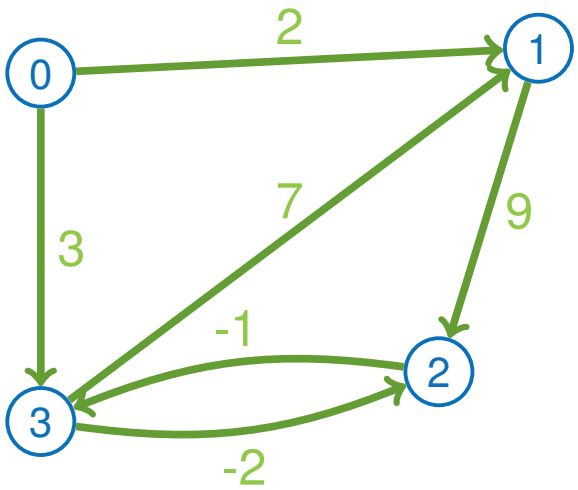

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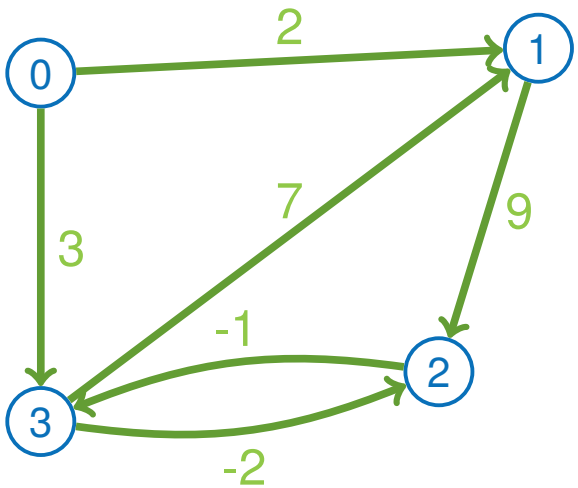

Abbildung: Weighted graph with $|V| = 4$, $|E| = 6$

start-vertice		
0	1, 2	3, 3
1	2, 9	
2	3, -1	
3	1, 7	2, -2

Abbildung: Adjacency list

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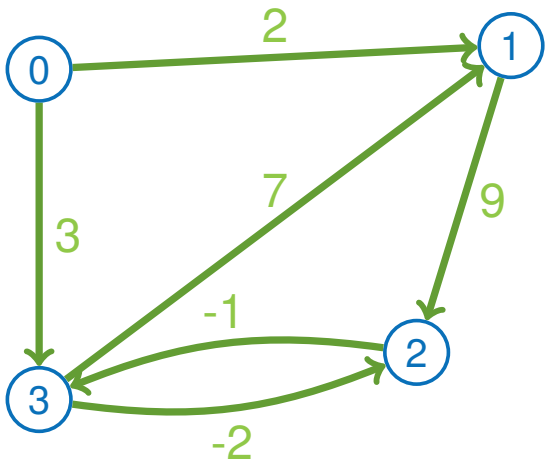

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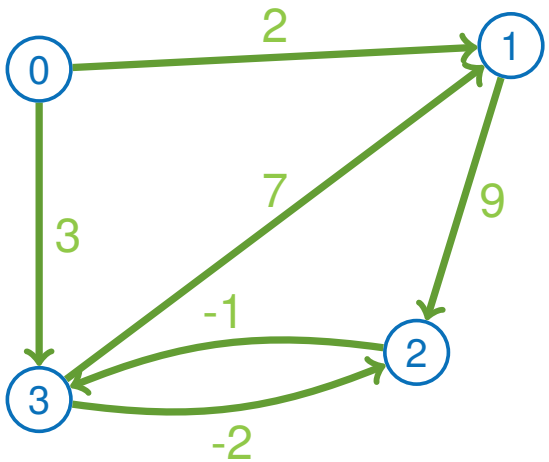

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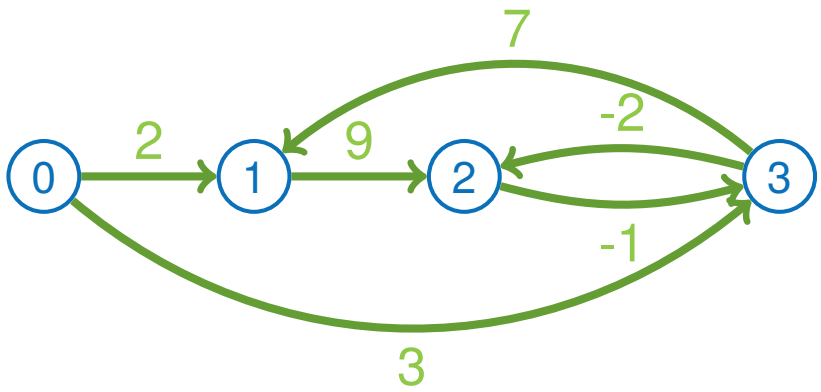

Abbildung: Same graph ordered by number - outer planar graph


```
class Graph:
    def __init__(self):
        self.vertices = []
        self.edges = []

    def addVertice(self, vert):
        self.vertices.append(vert)

    def addEdge(self, fromVert, toVert):
        self.edges.append((fromVert, toVert))

    ...
```

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```
...

def toString(self):
    return '{'
        + ', '.join( \
            [str(len(self.vertices)), \
                str(len(self.edges))] \
            + ["(%s, %s)" % tup \
                for tup in self.edges]) \
        + '}'
```

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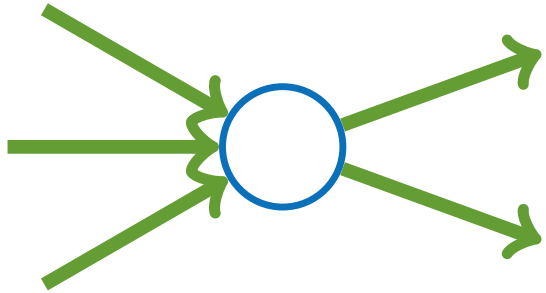

Abbildung: Vertex with in- / outdegree of 3 / 2

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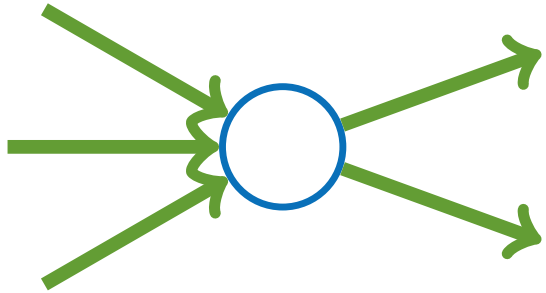

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- Indegree of a vertex u is the number of edge heads adjacent to the vertex

$$\deg^+(u) = |\{(v,u) : (v,u) \in E\}|$$

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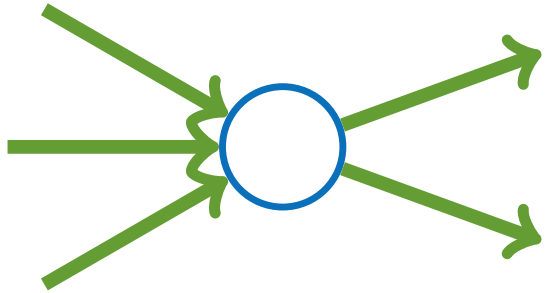

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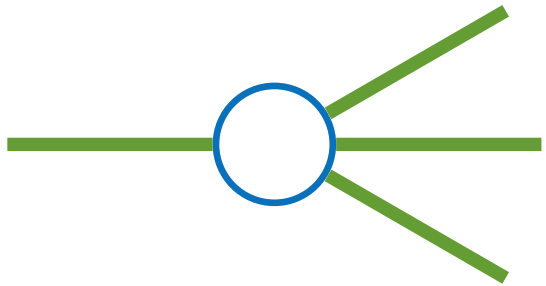

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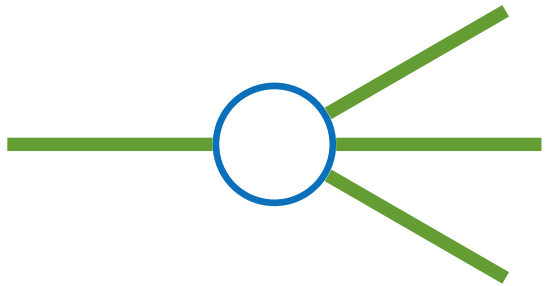

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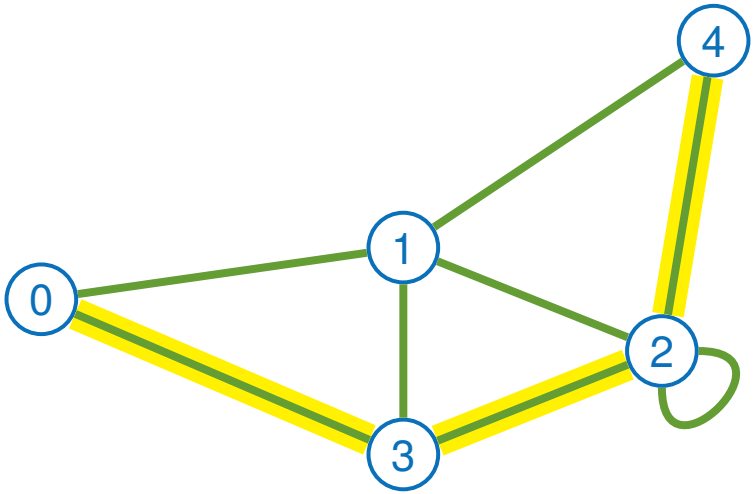
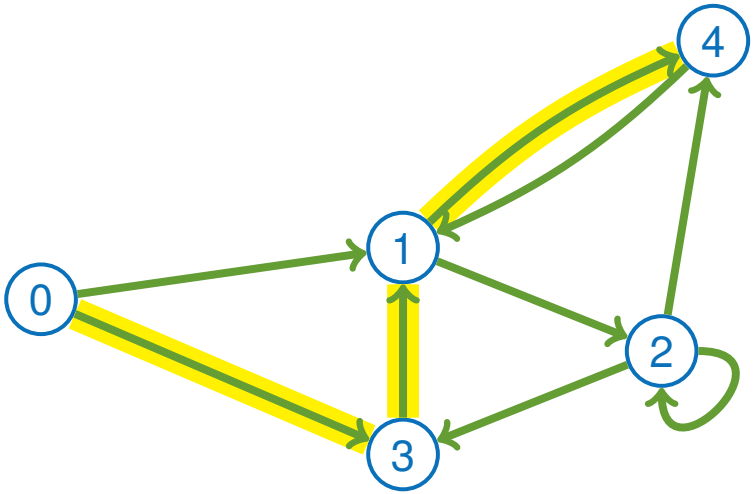


Abbildung: Undirected path of length 3

$P = (0, 3, 2, 4)$



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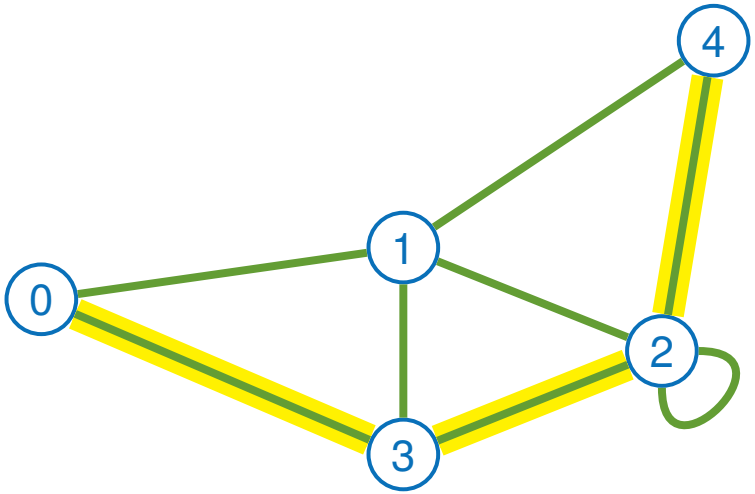

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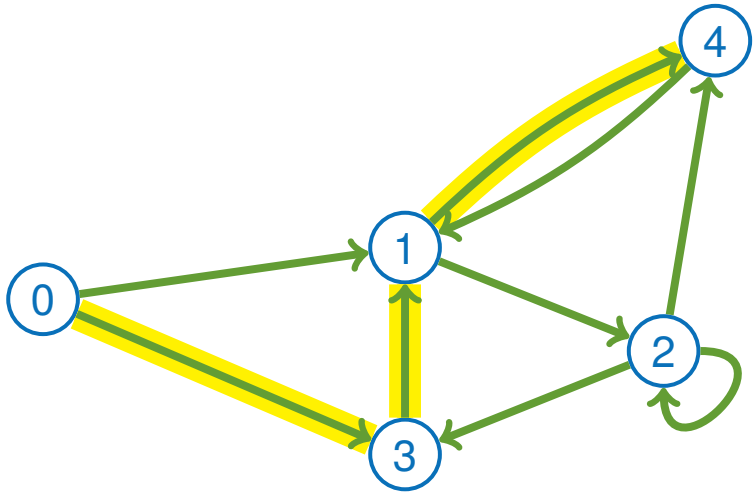

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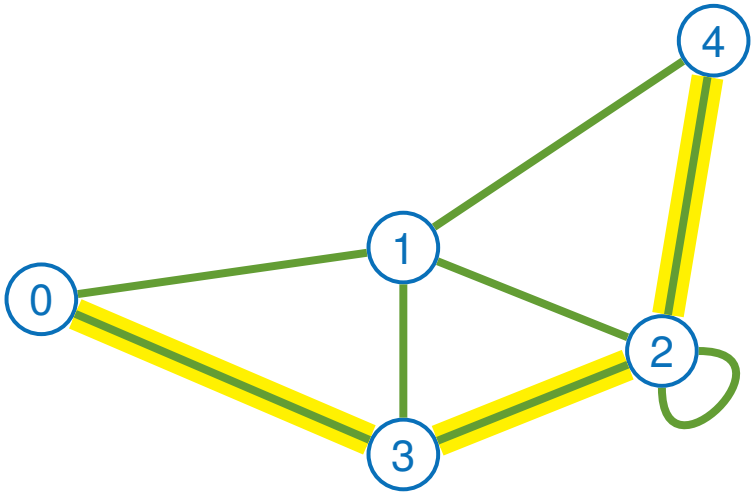

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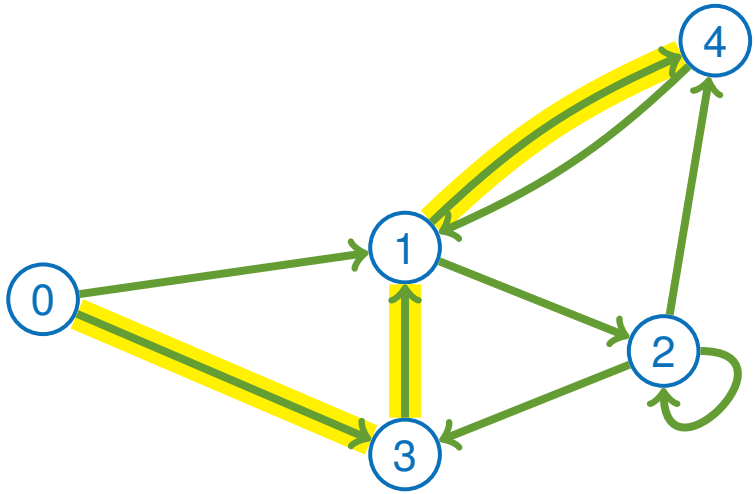
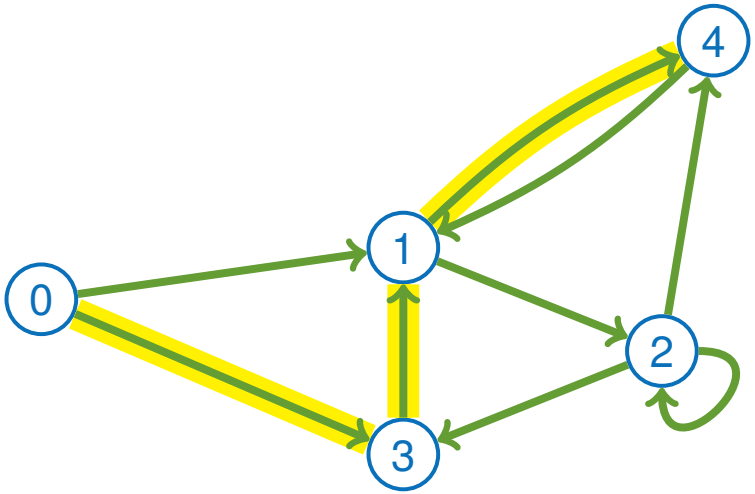


Abbildung: Directed path of length 3

$P = (0, 3, 1, 4)$

- A path of G is a sequence of edges $u_1, u_2, \ldots, u_i \in V$ with
  - Undirected graph: $\{u_1, u_2\}, \{u_2, u_3\}, \ldots, \{u_{i-1}, u_i\} \in E$
  - Directed graph: $(u_1, u_2), (u_2, u_3), \ldots, (u_{i-1}, u_i) \in E$

Paths in a graph: $G = (V, E)$



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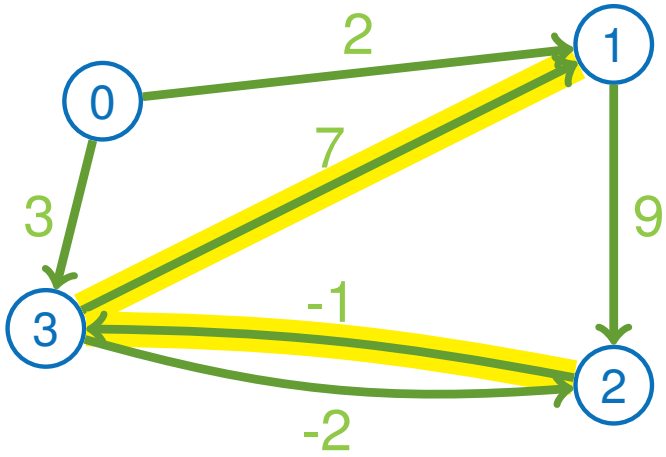

Abbildung: Weighted path with cost 6

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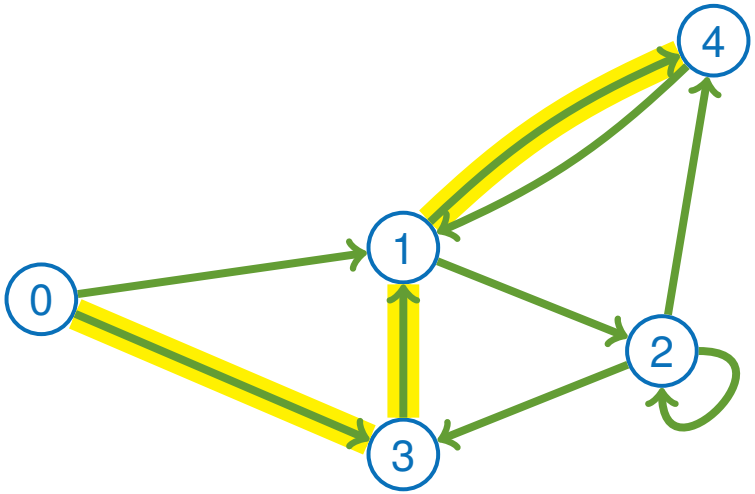

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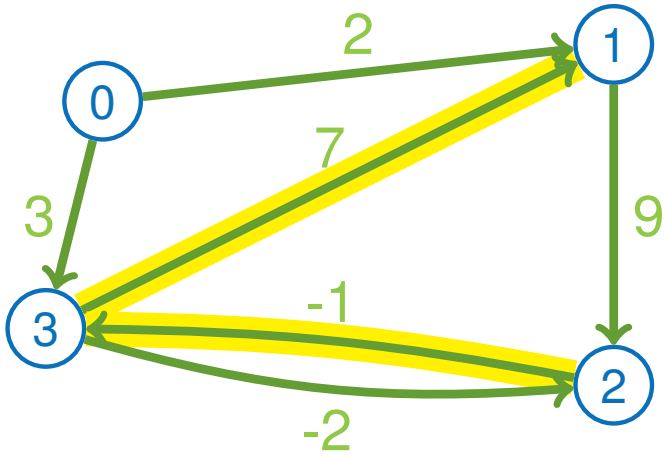

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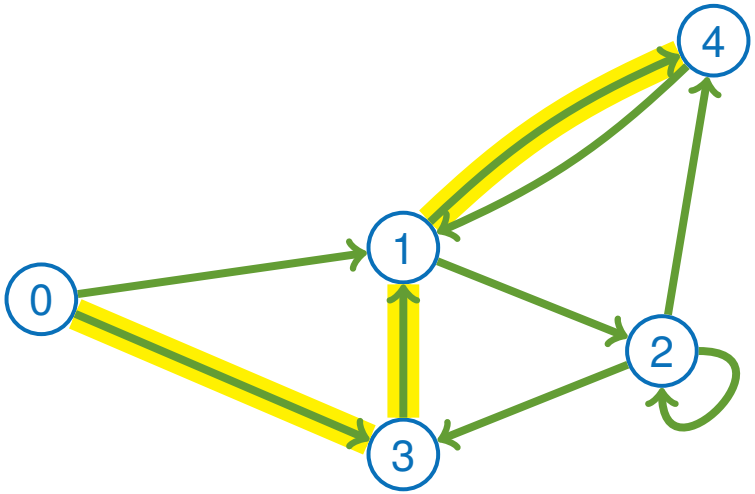

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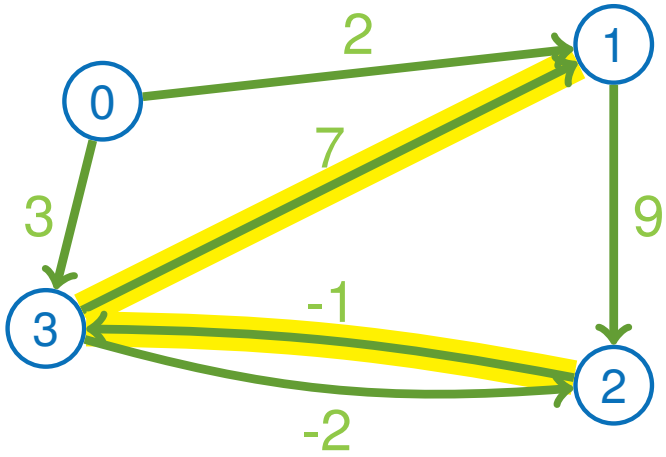

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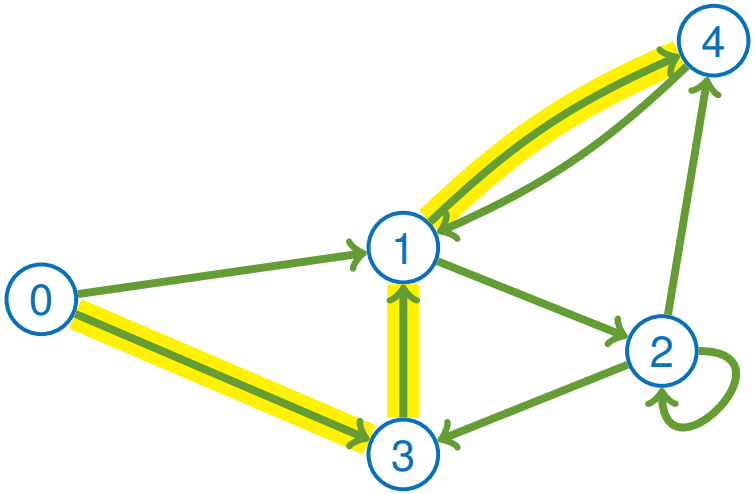

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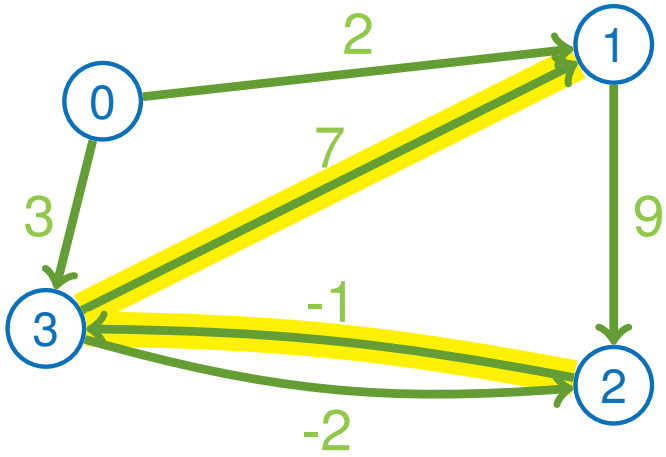

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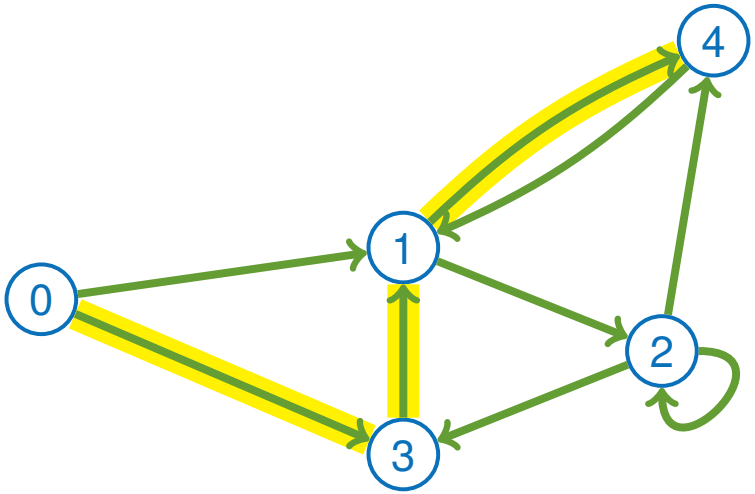

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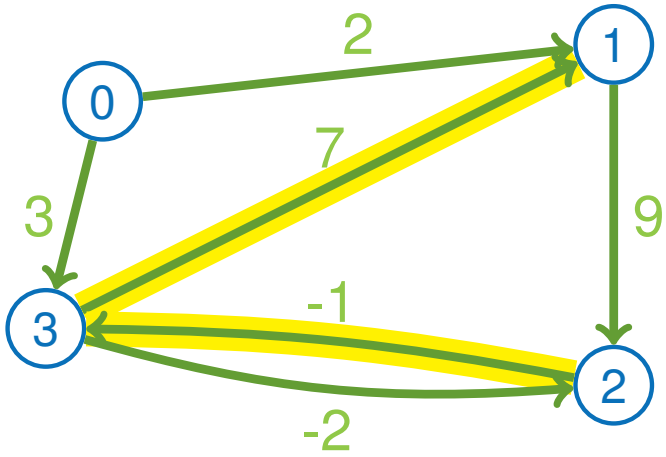

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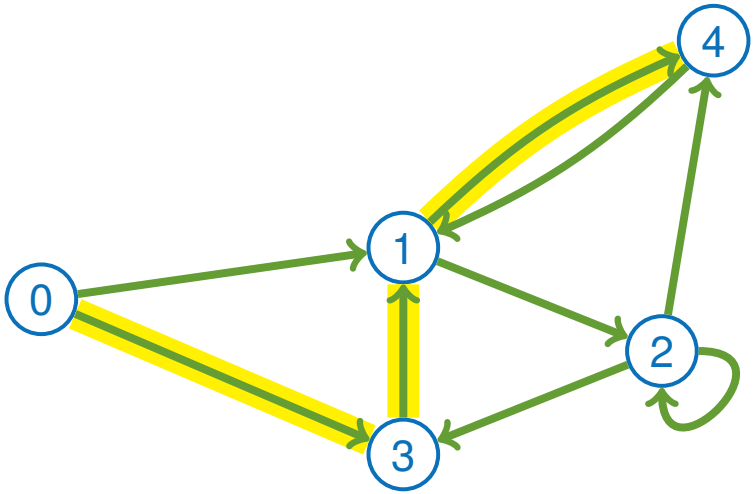

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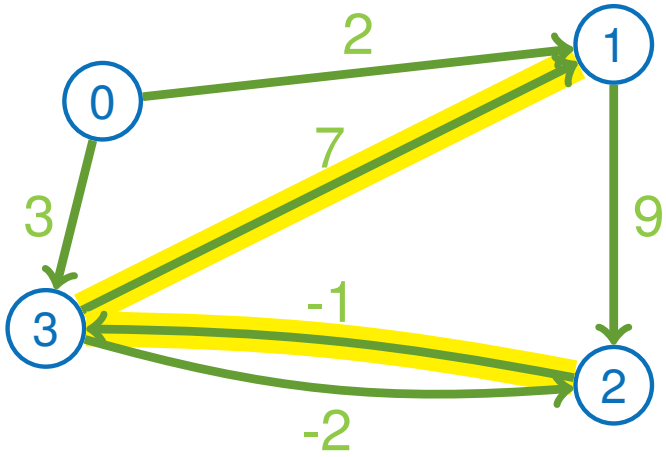

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Shortest path in a graph: $G = (V, E)$

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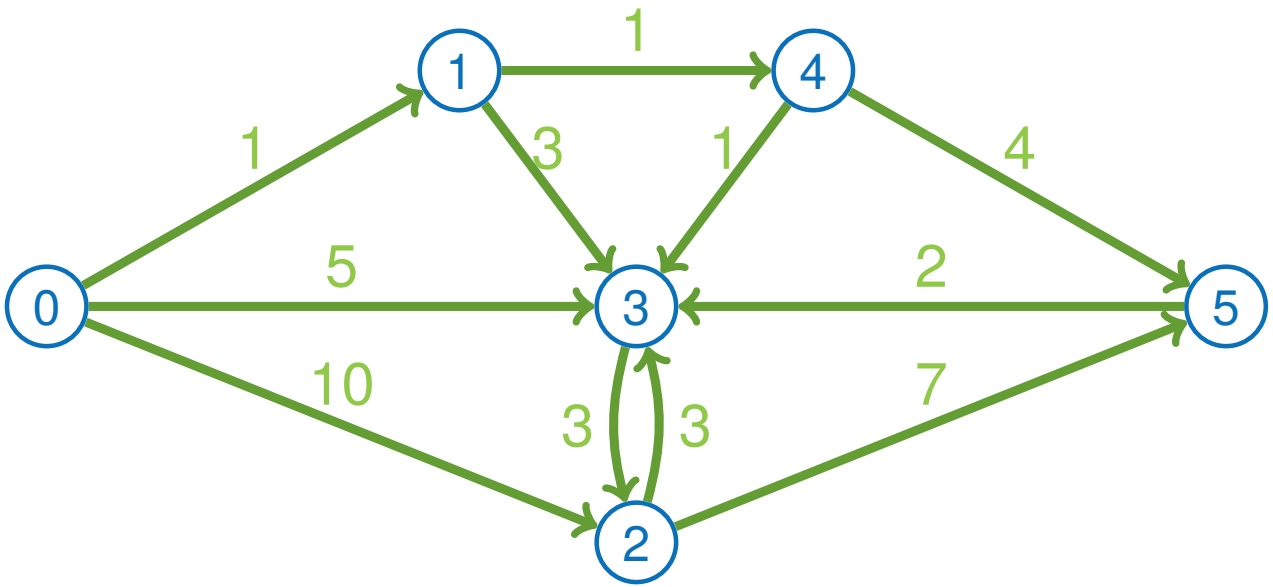

Abbildung: Shortest path from 0 to 2 with cost / distance $d(0,2) = ?$

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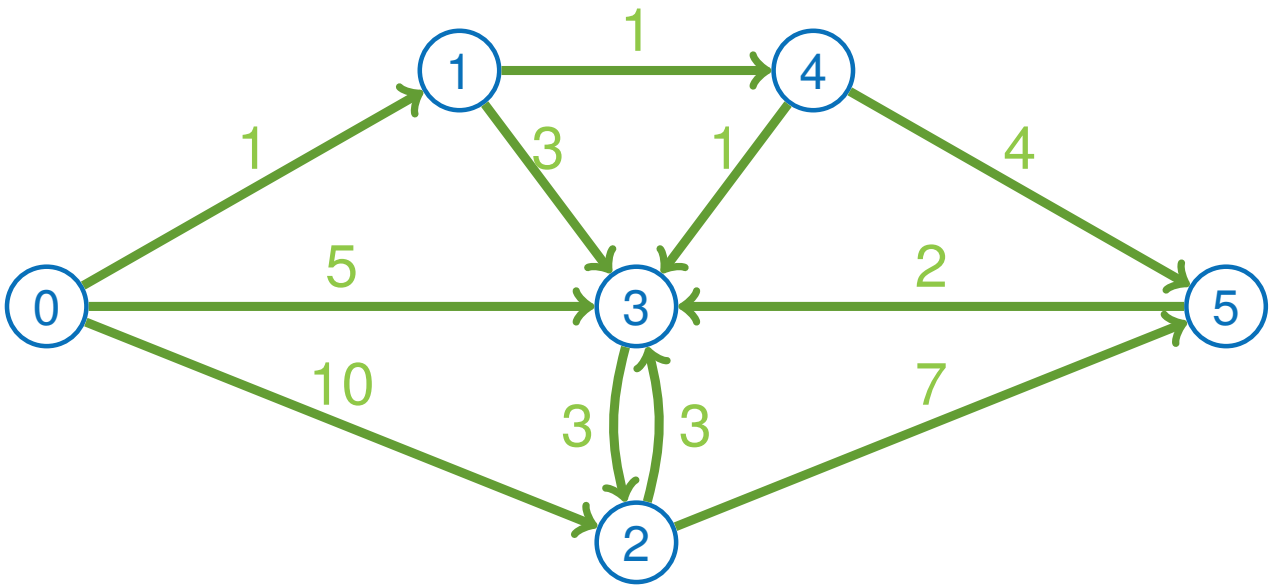

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- The shortest path between two vertices u, v is the path $P = (u, \ldots, v)$ with the shortest length $d(u, v)$ or lowest costs

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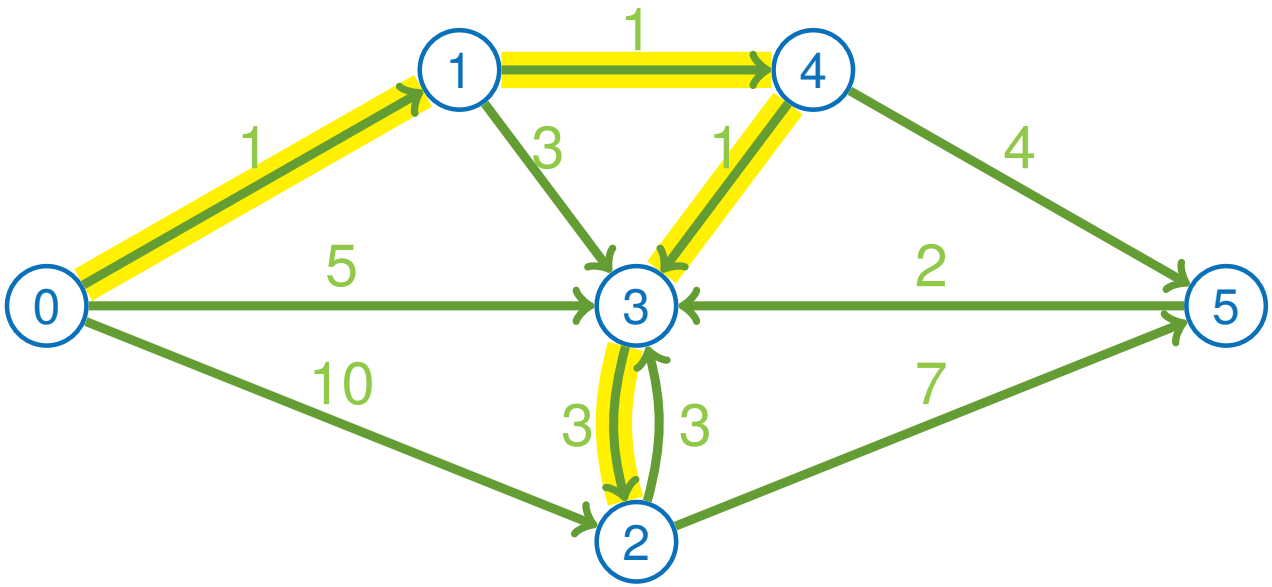

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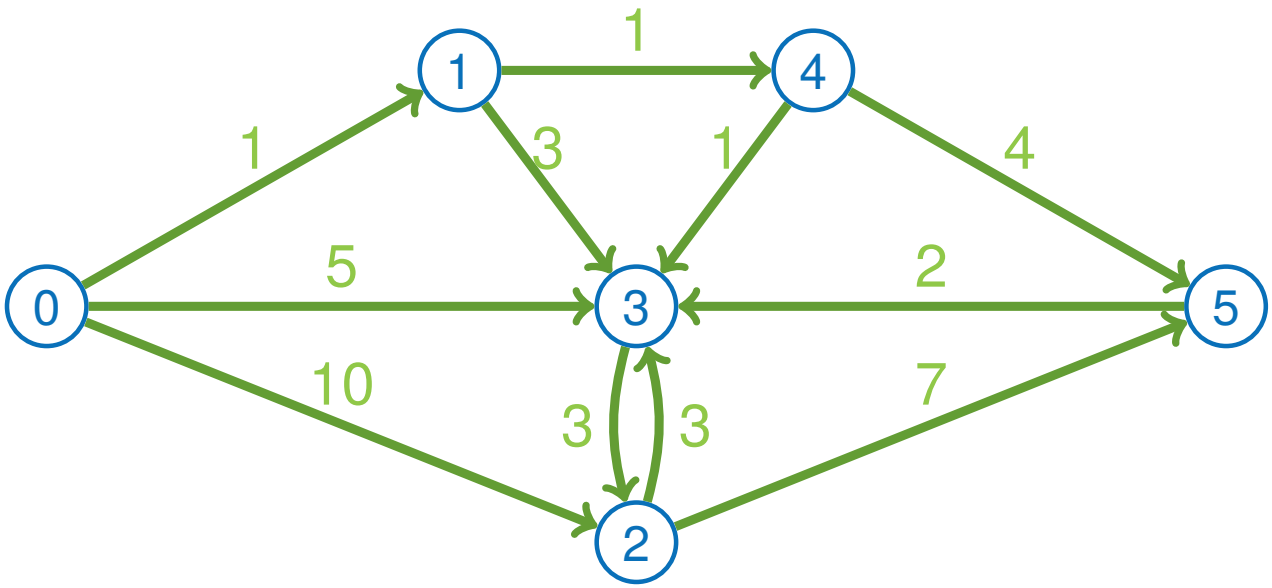

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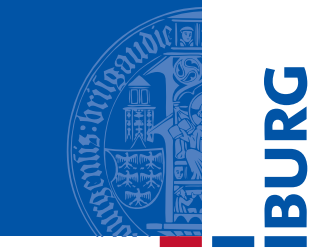

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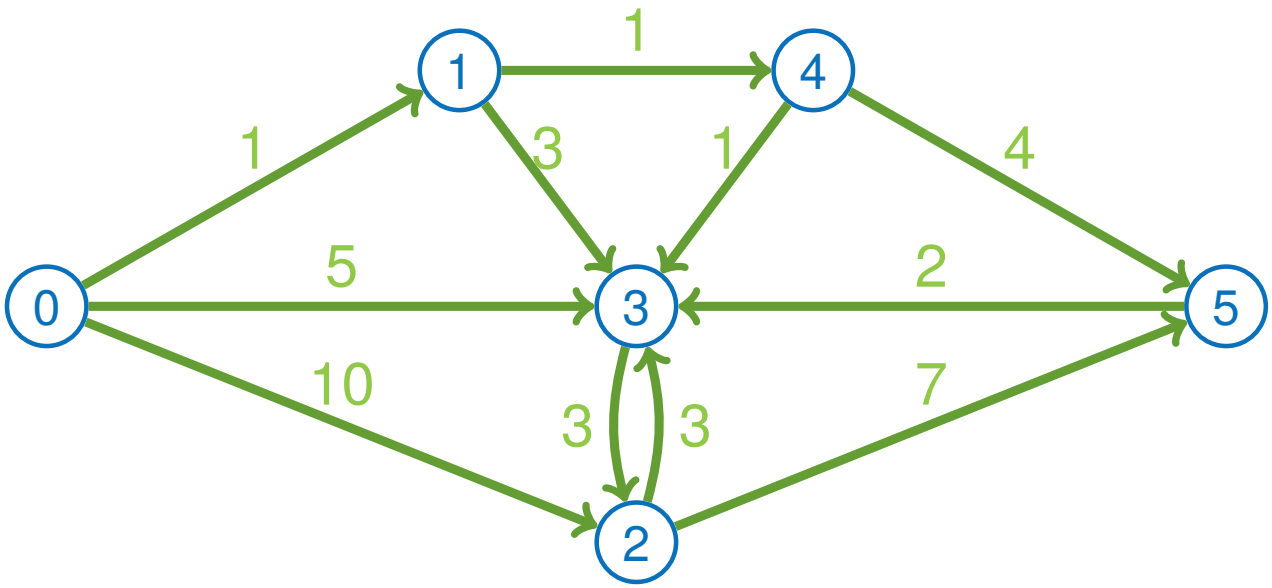

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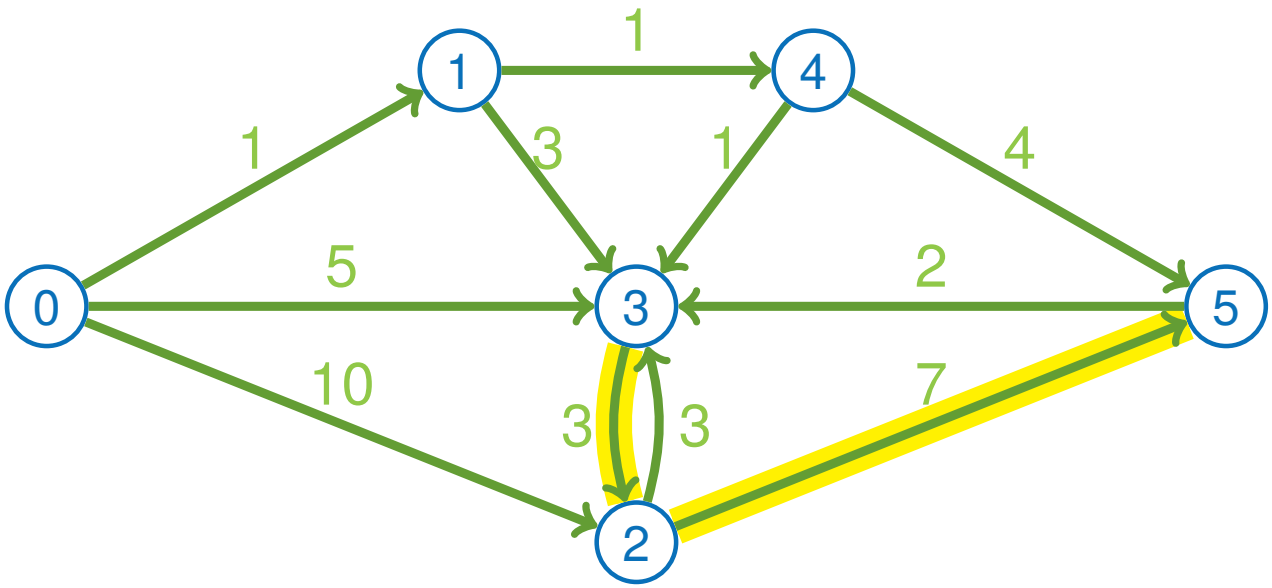

Abbildung: Diameter of graph is $d = 10$, $P = (3, 2, 5)$

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Connected components: $G = (V, E)$

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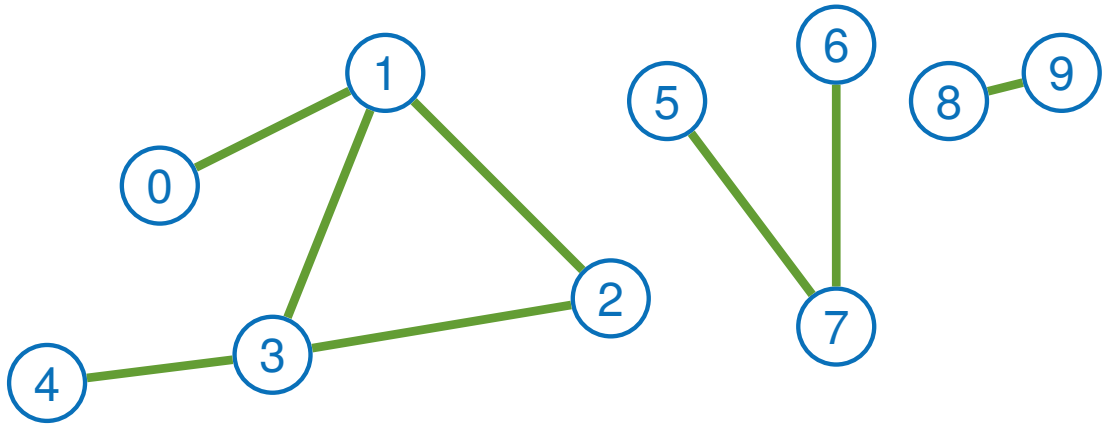
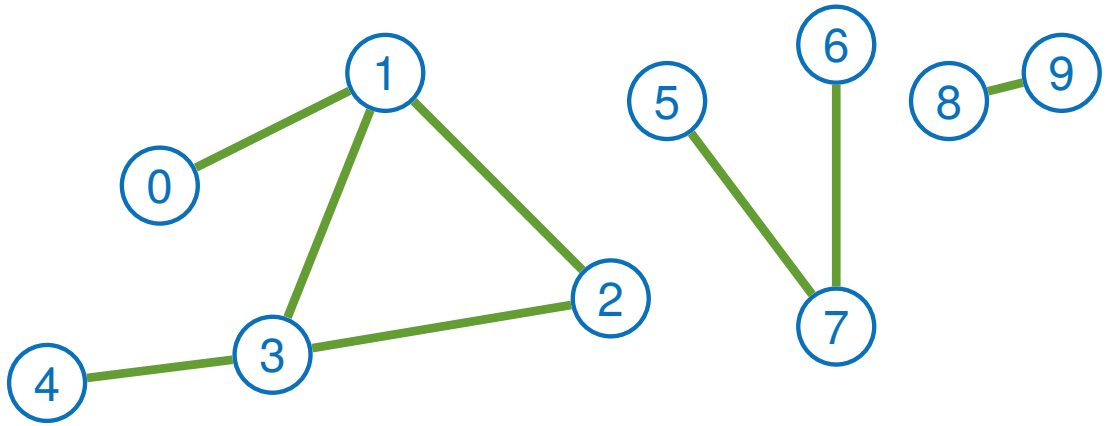


Abbildung: Three connected components

- Undirected graph:

Connected components: $G = (V, E)$



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  - All connected components are a partition of V

$$V = V_1 \cup \cdots \cup V_k$$

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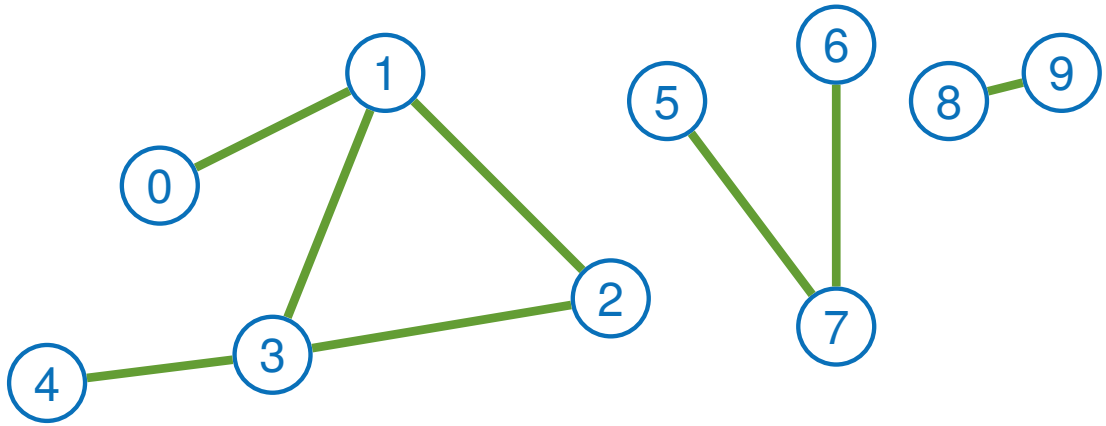

Abbildung: Three connected components

- Undirected graph:
  - All connected components are a partition of V

$$V = V_1 \cup \cdots \cup V_k$$

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- Breadth-first search: in sequence of the smallest distance to s
- Depth-first search: in sequence of the largest distance to s
- Not a problem on its own but is often used as subroutine of other algorithms

Idea:

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6. All connected nodes are now marked and in the same connected component as the start vertex s

- The marked vertices create a „spanning tree“ containing all reachable nodes

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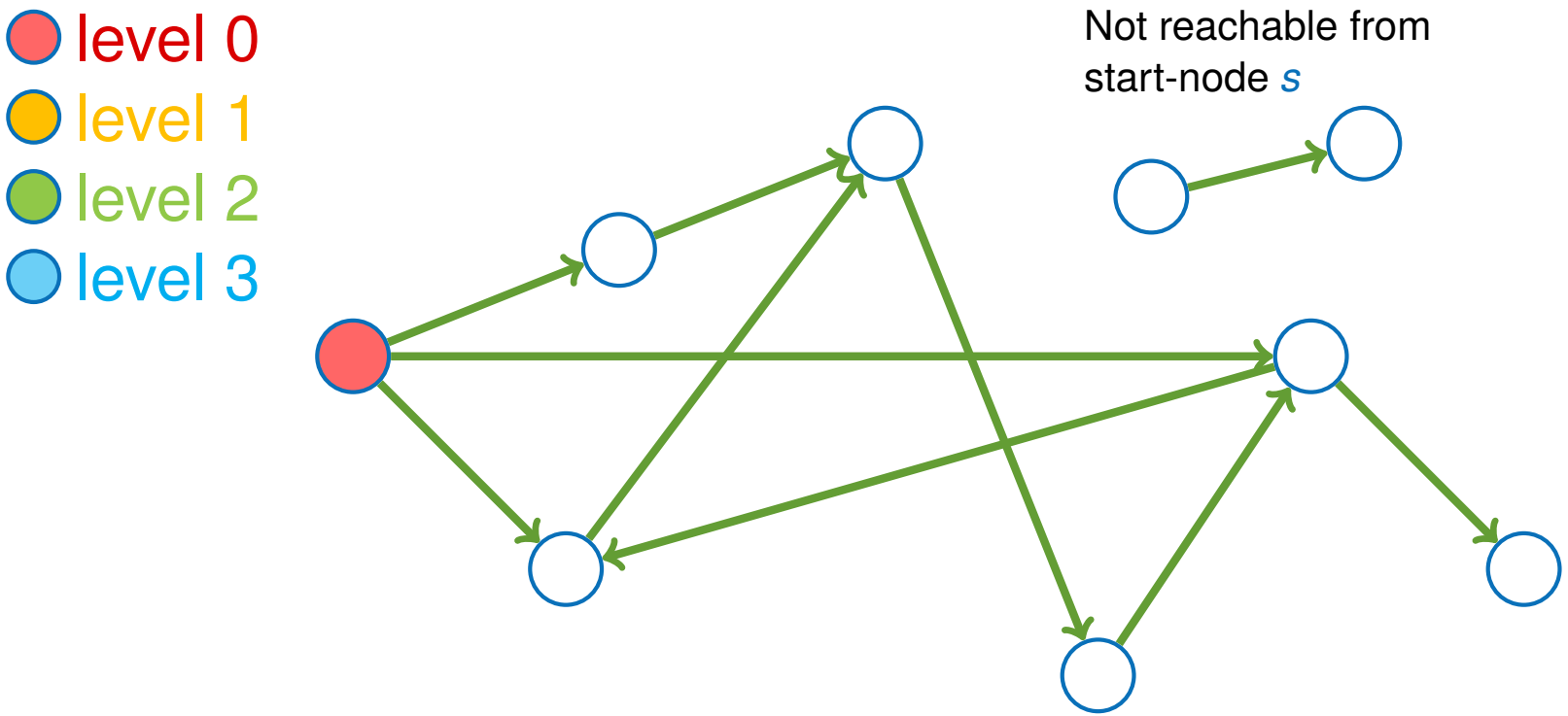

Abbildung: spanning tree of a breadth-first search

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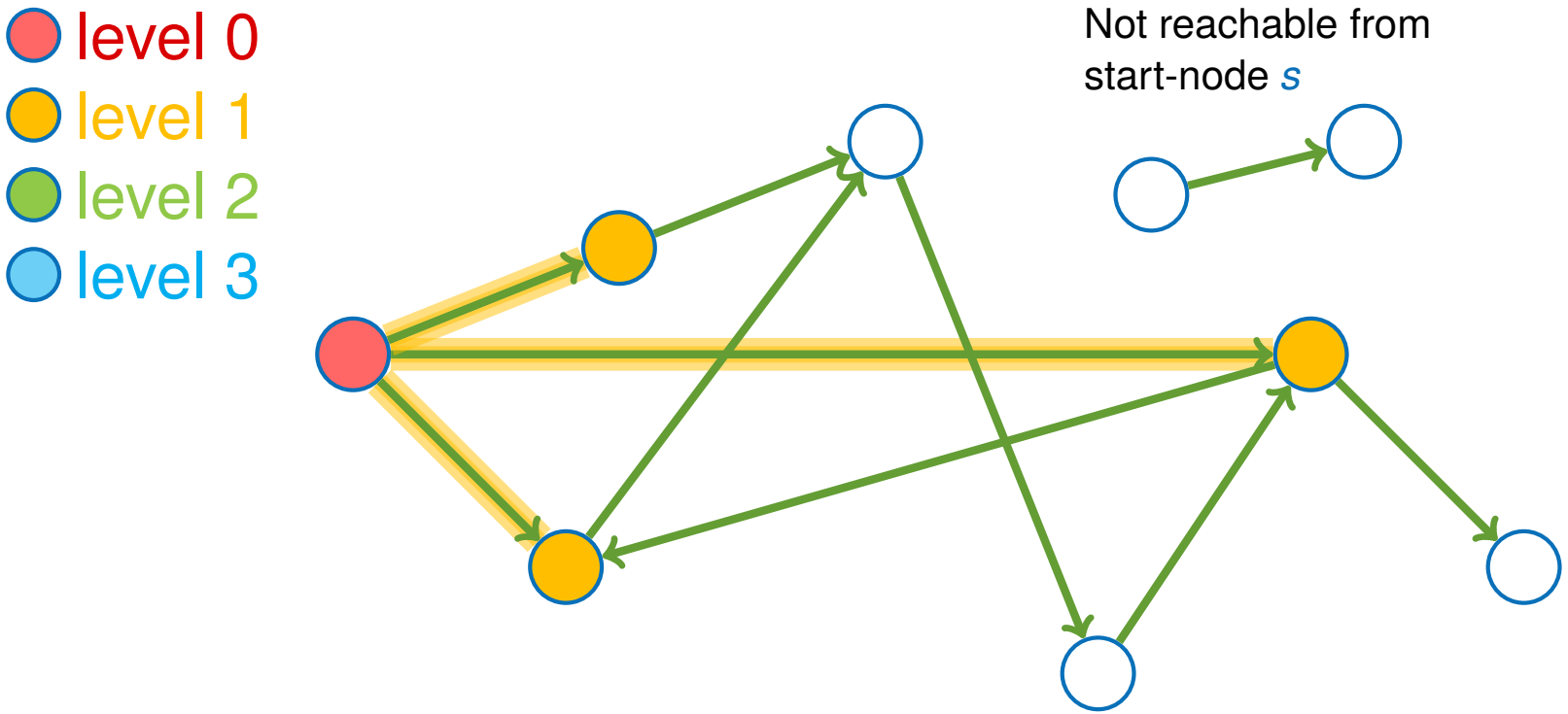

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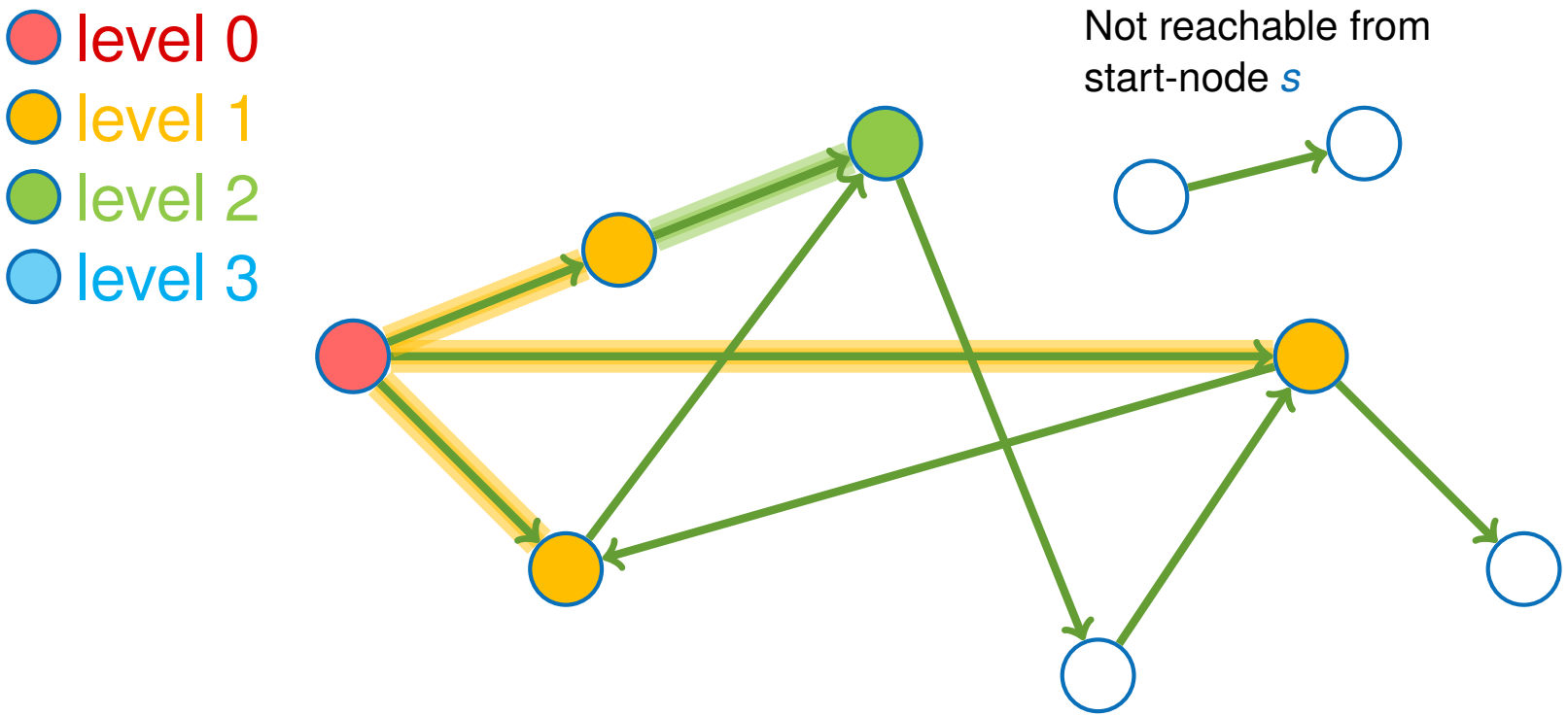

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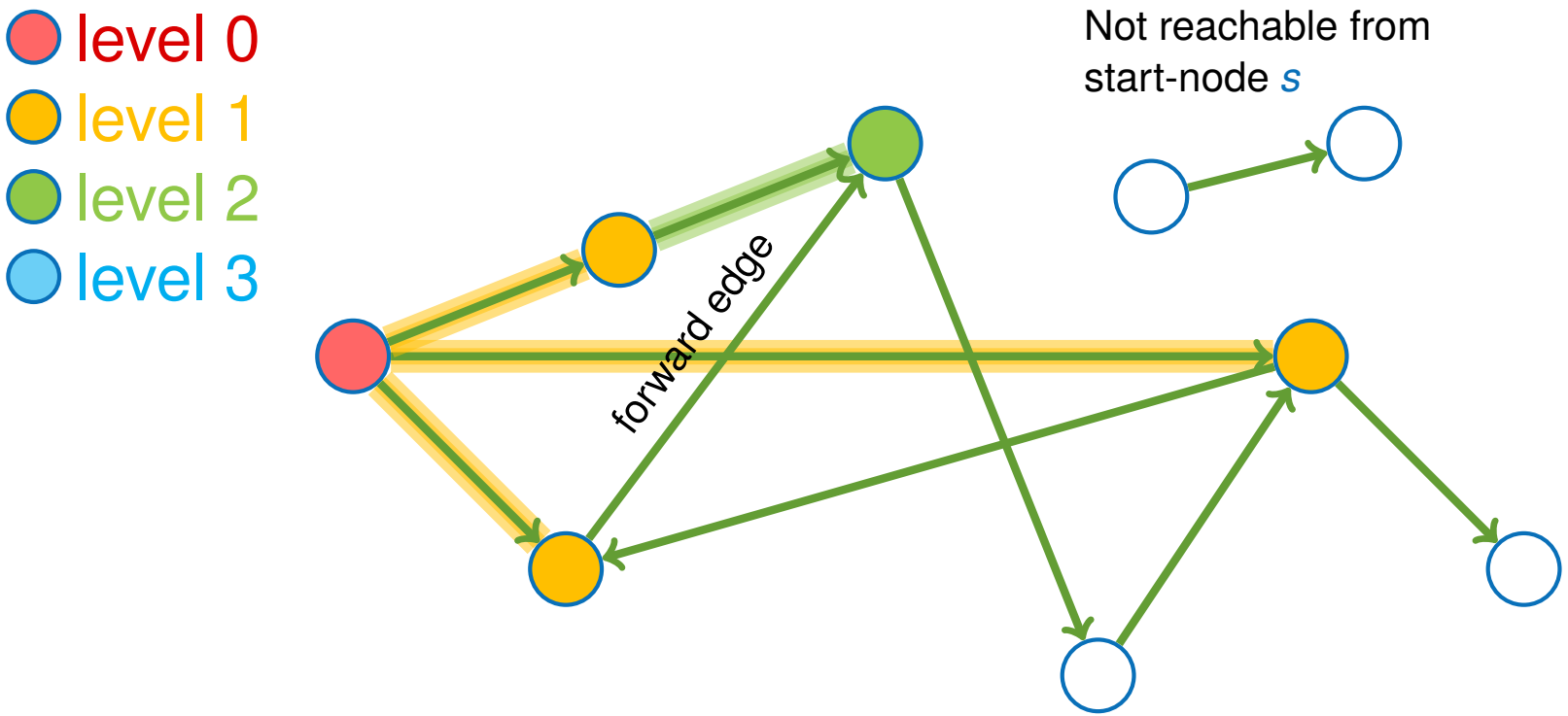

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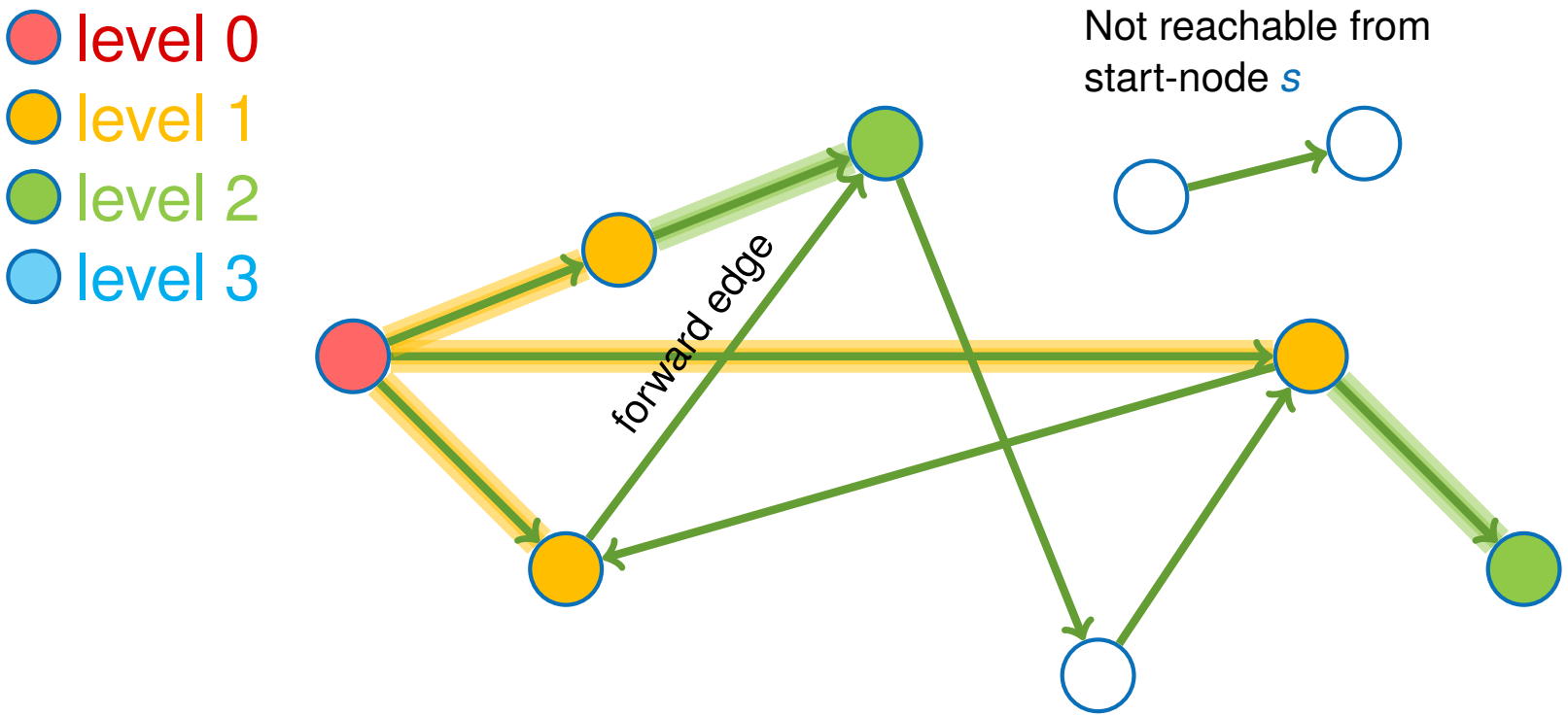

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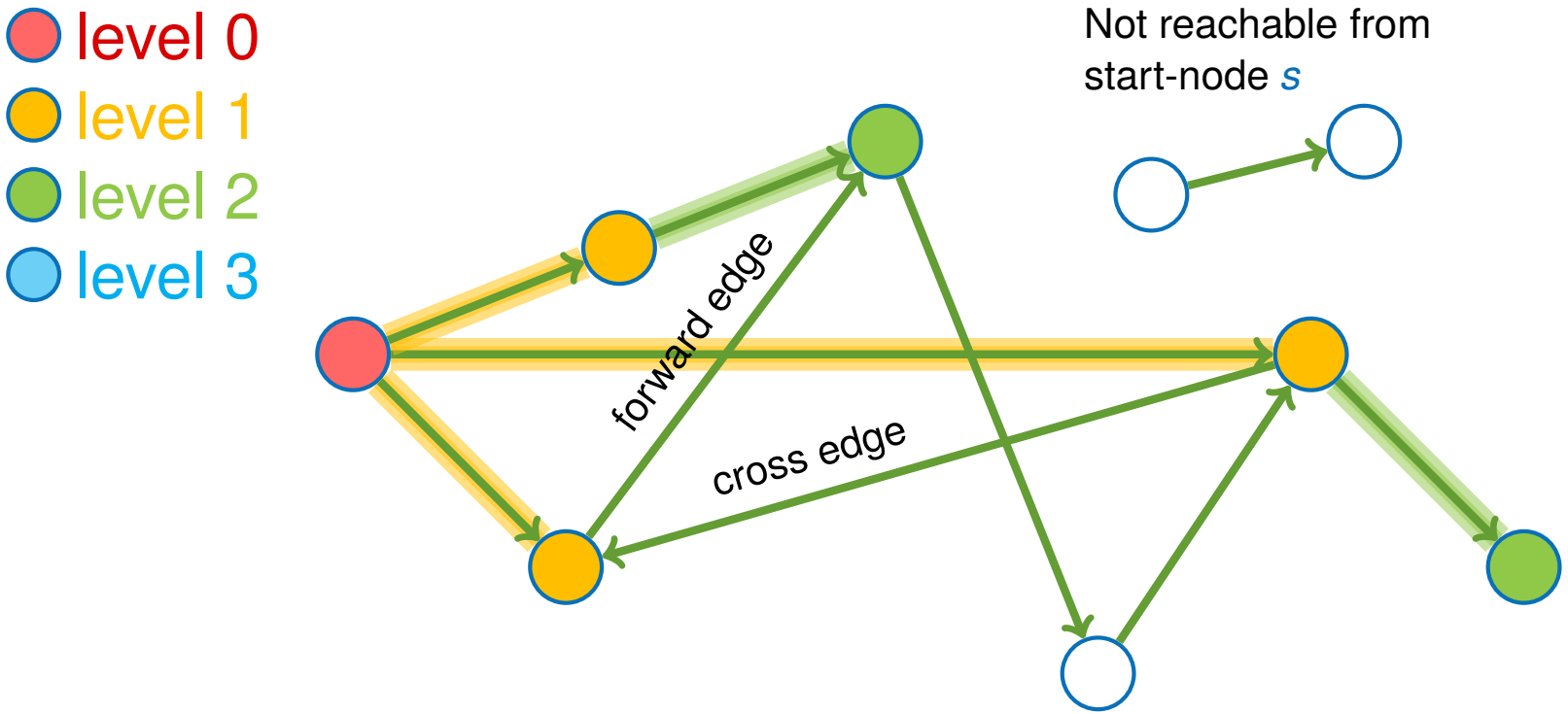

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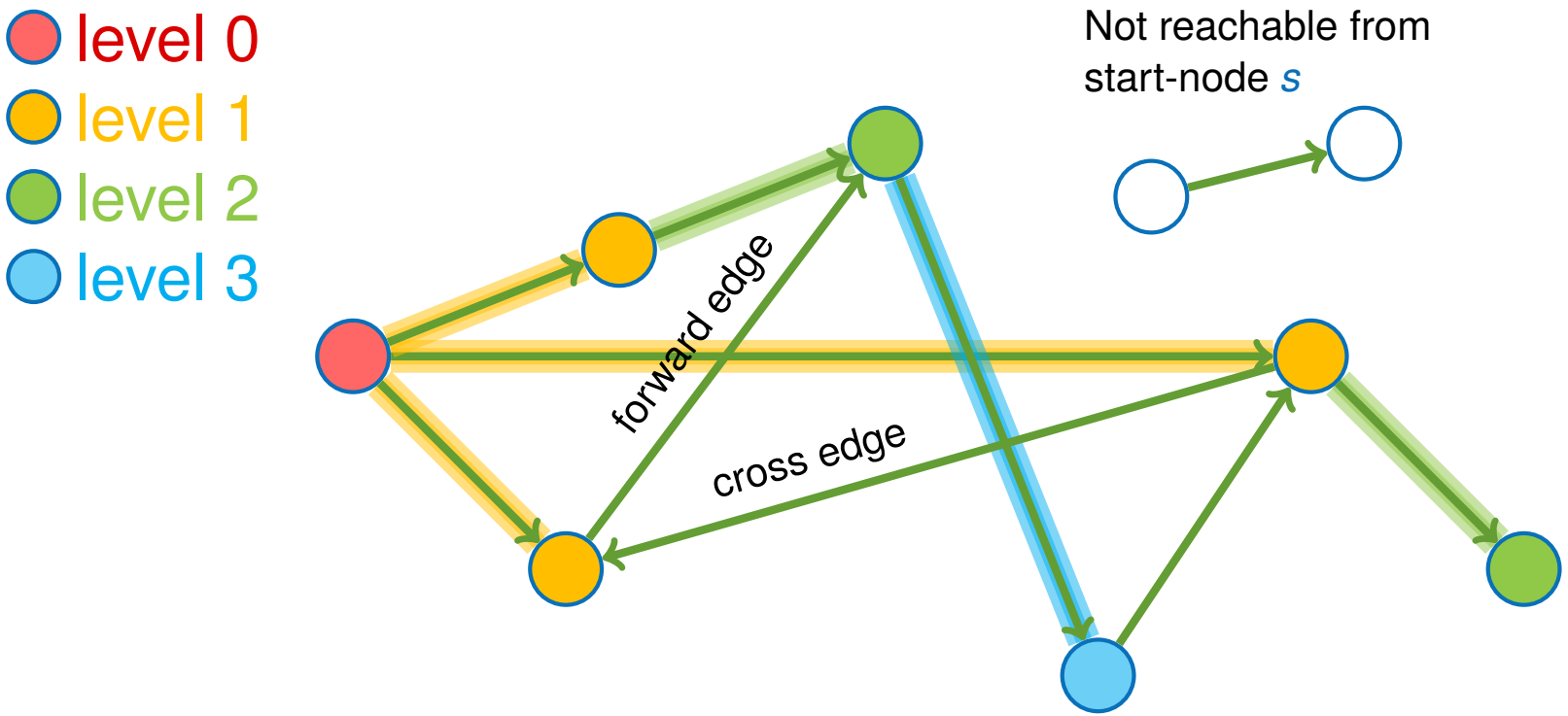

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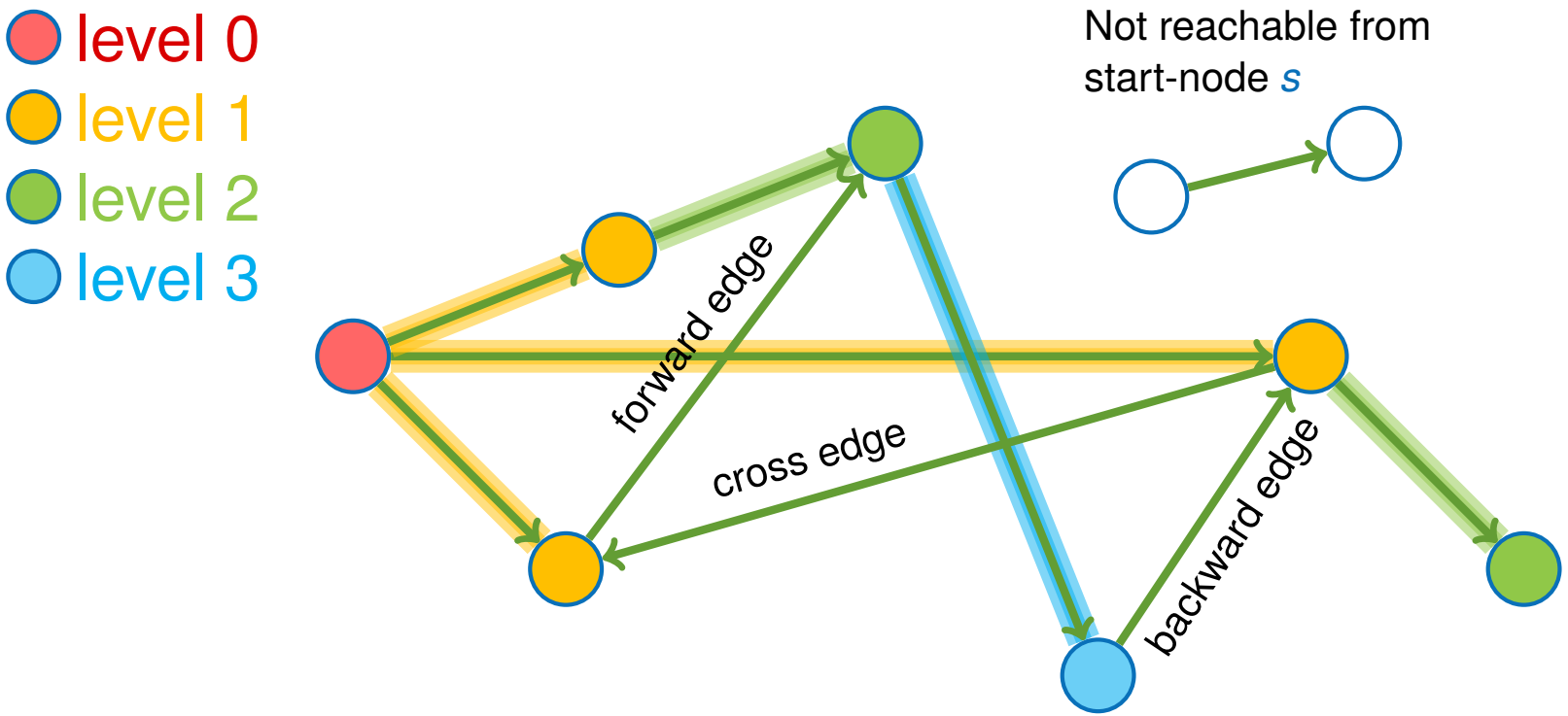

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4. If no unmarked connected vertex exists go one vertex back (reduce the recursion level by one)

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  - The numbers increase with path from the start vertex

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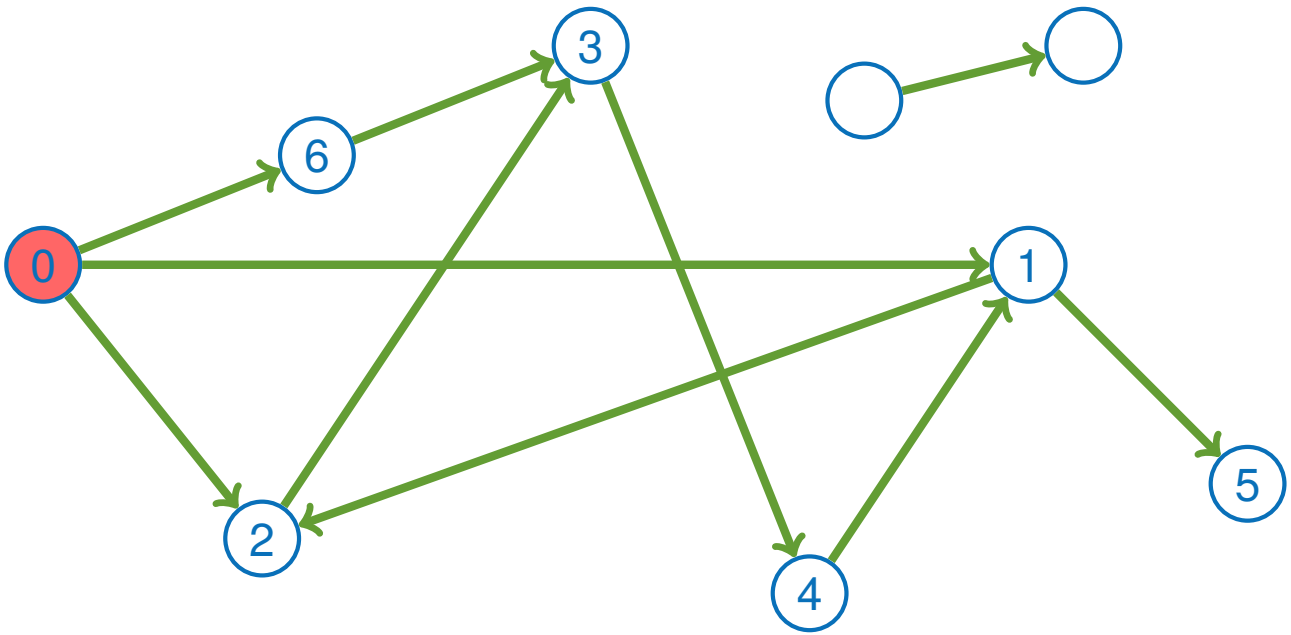
path 1
path 2
path 3



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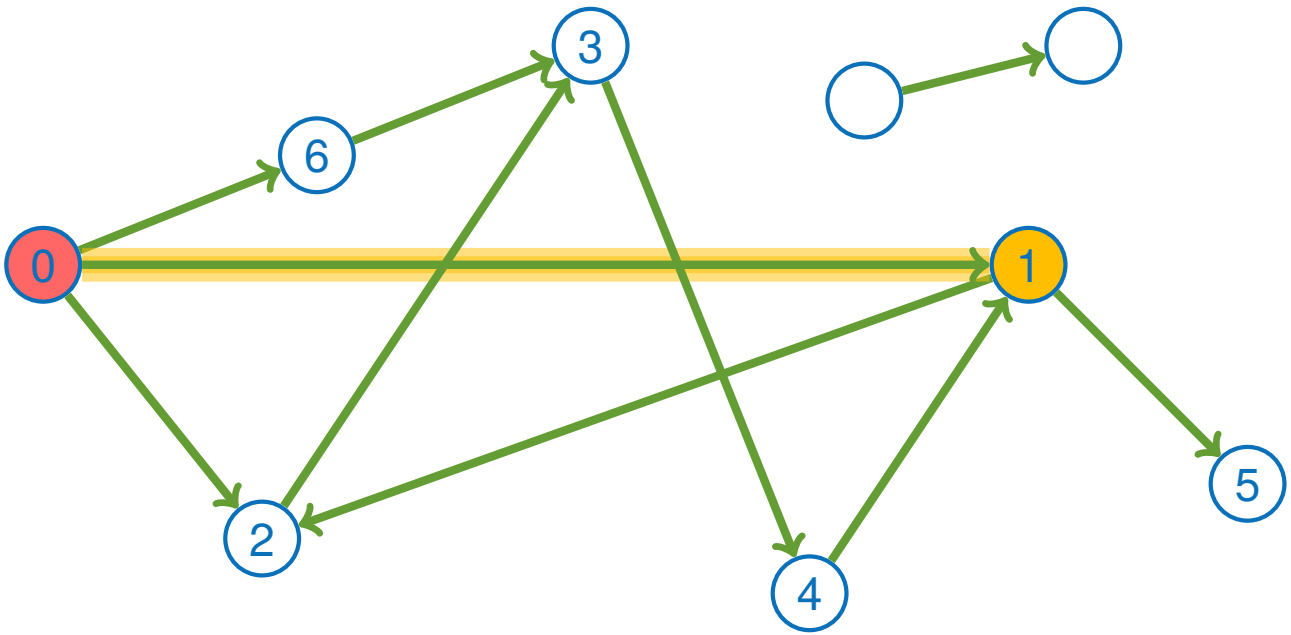

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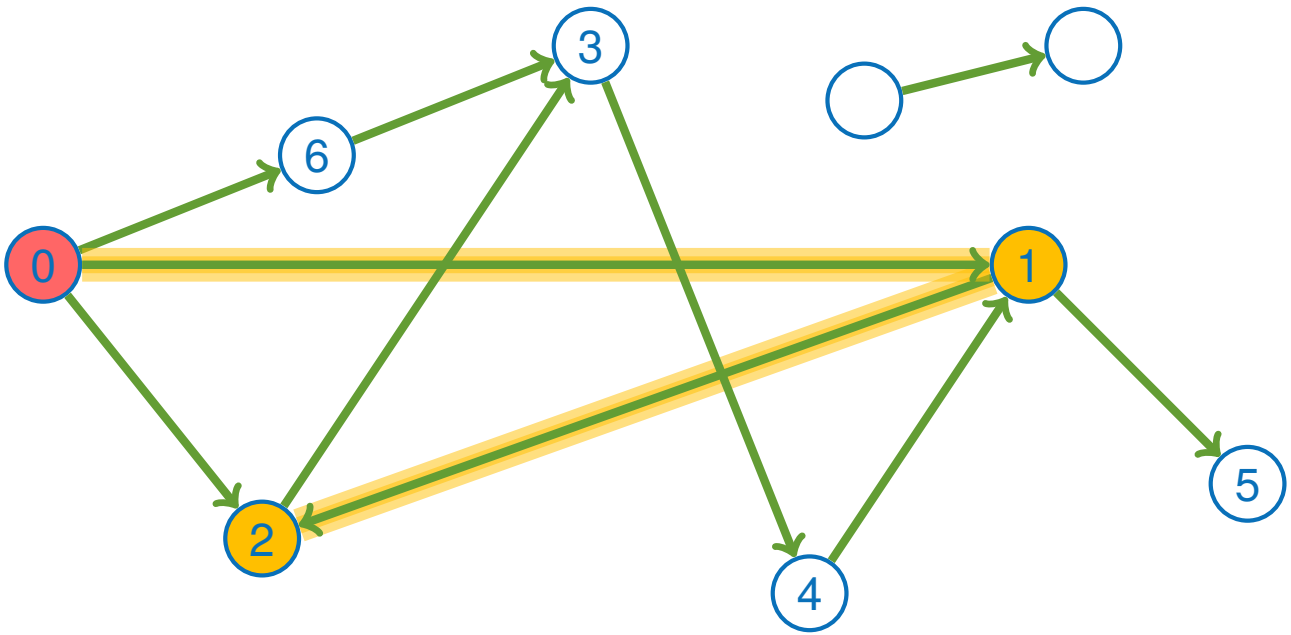

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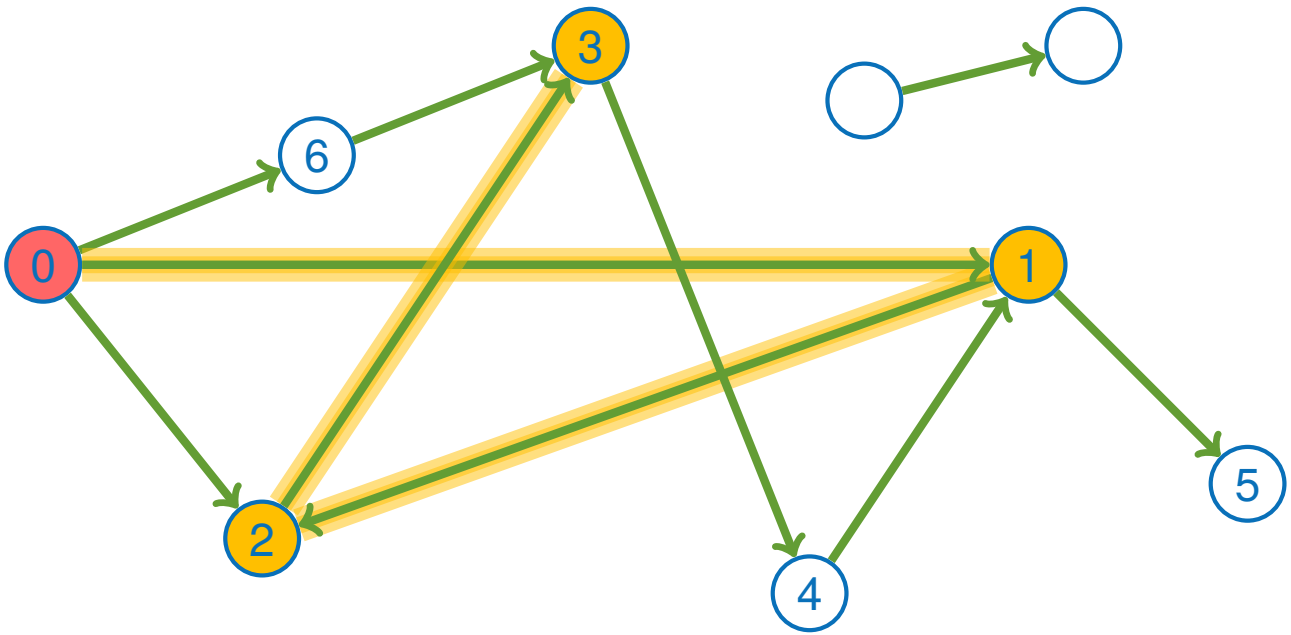

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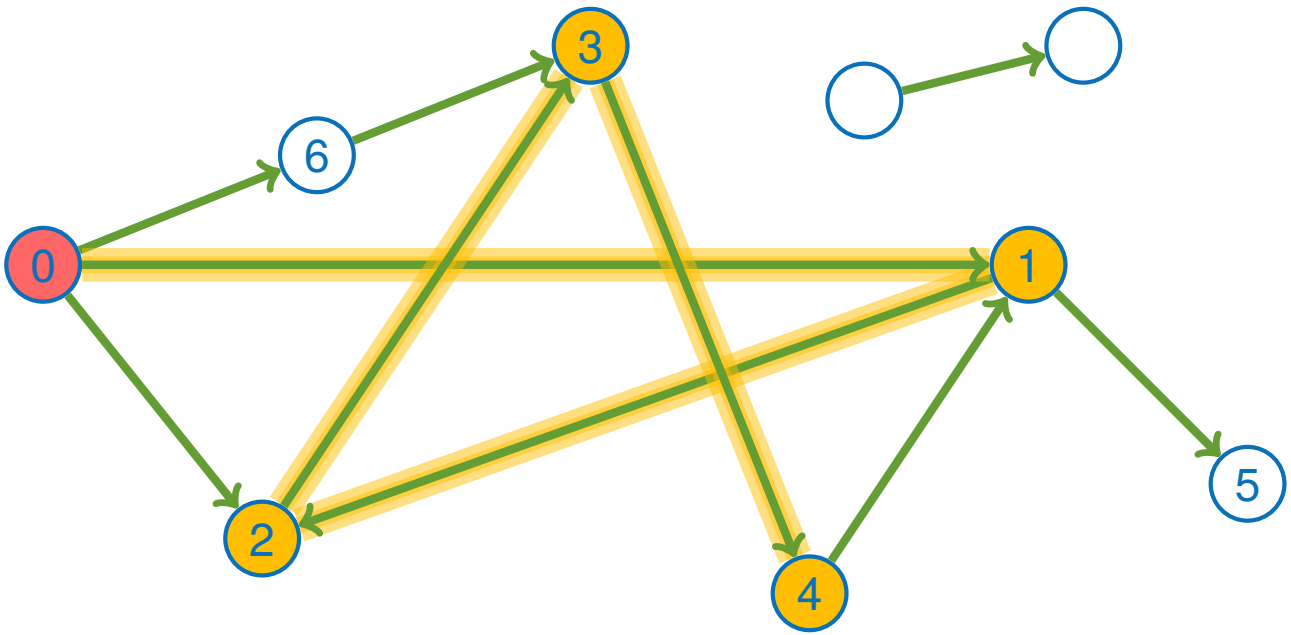

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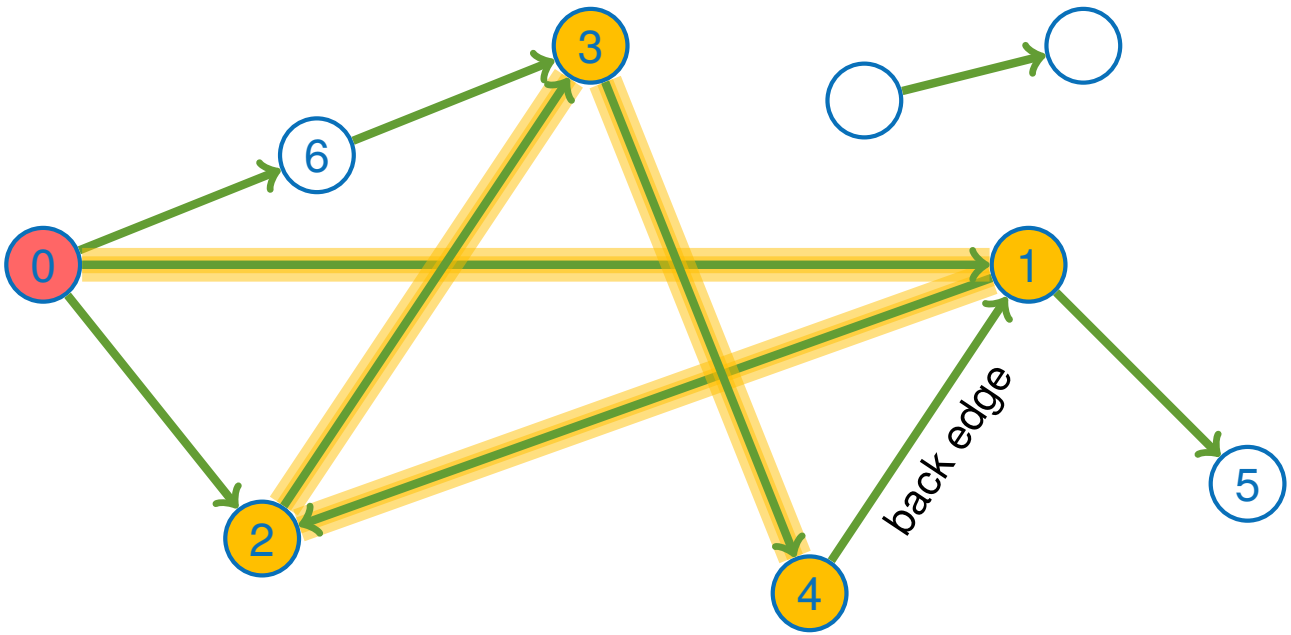

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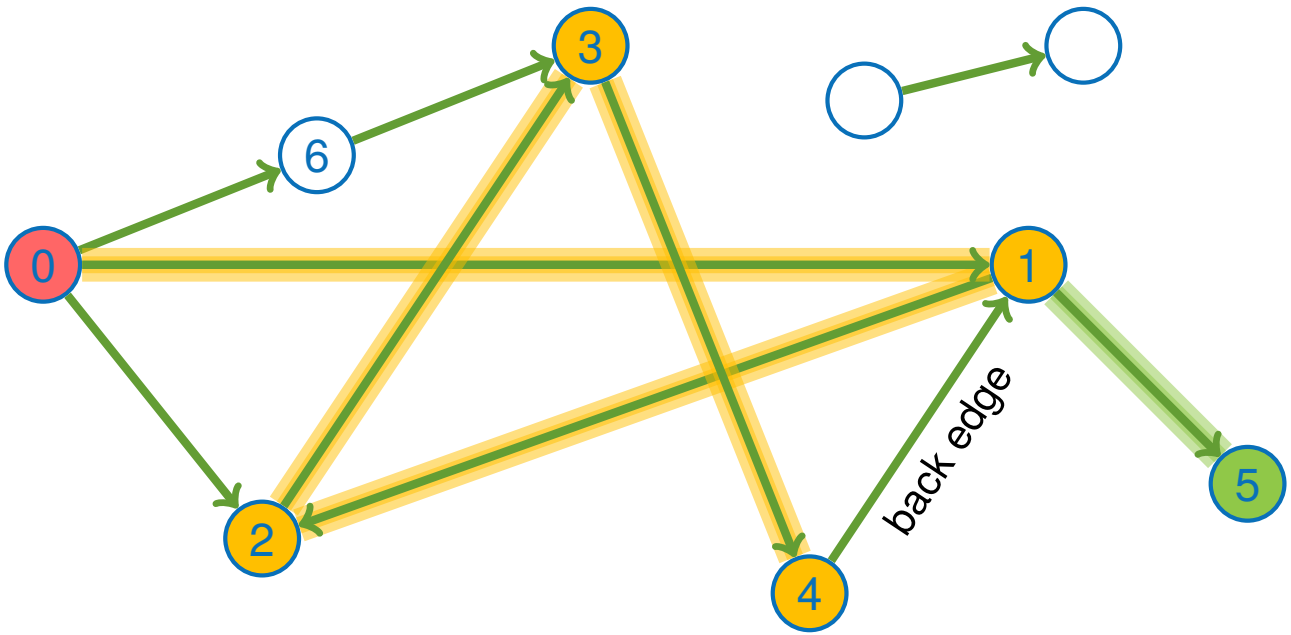

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path 2
path 3

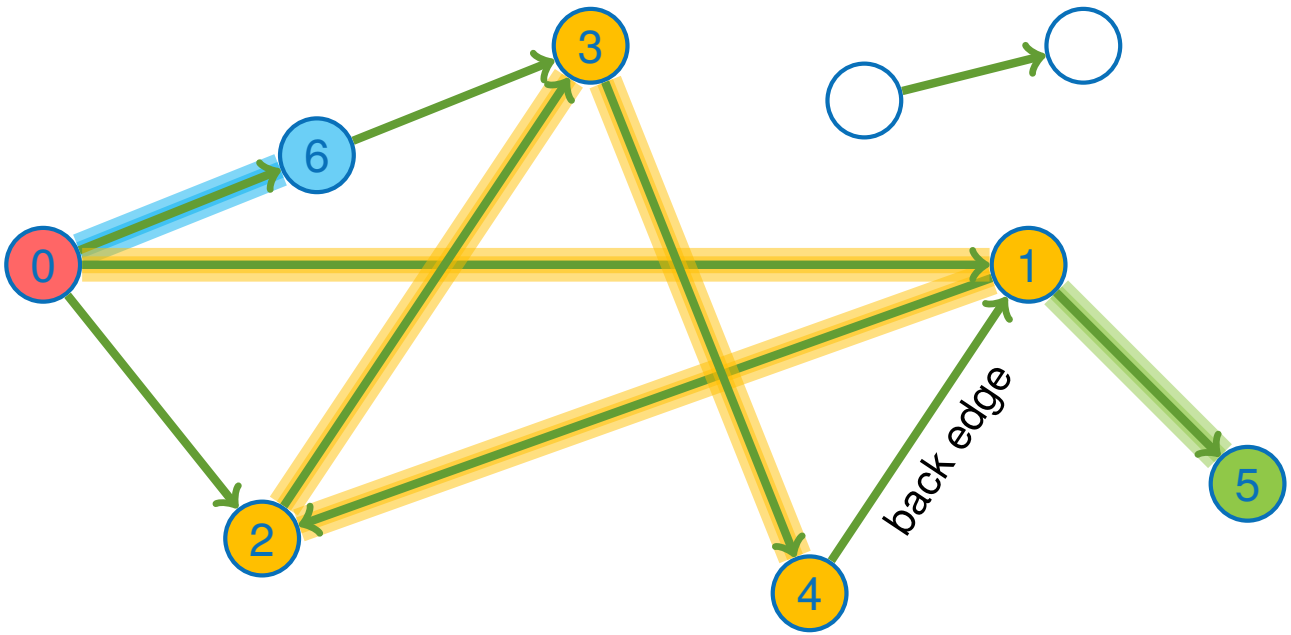

Abbildung: spanning tree of a depth-first search

- The marked vertices create a different spanning tree containing all reachable nodes

start-node
path 1
path 2
path 3

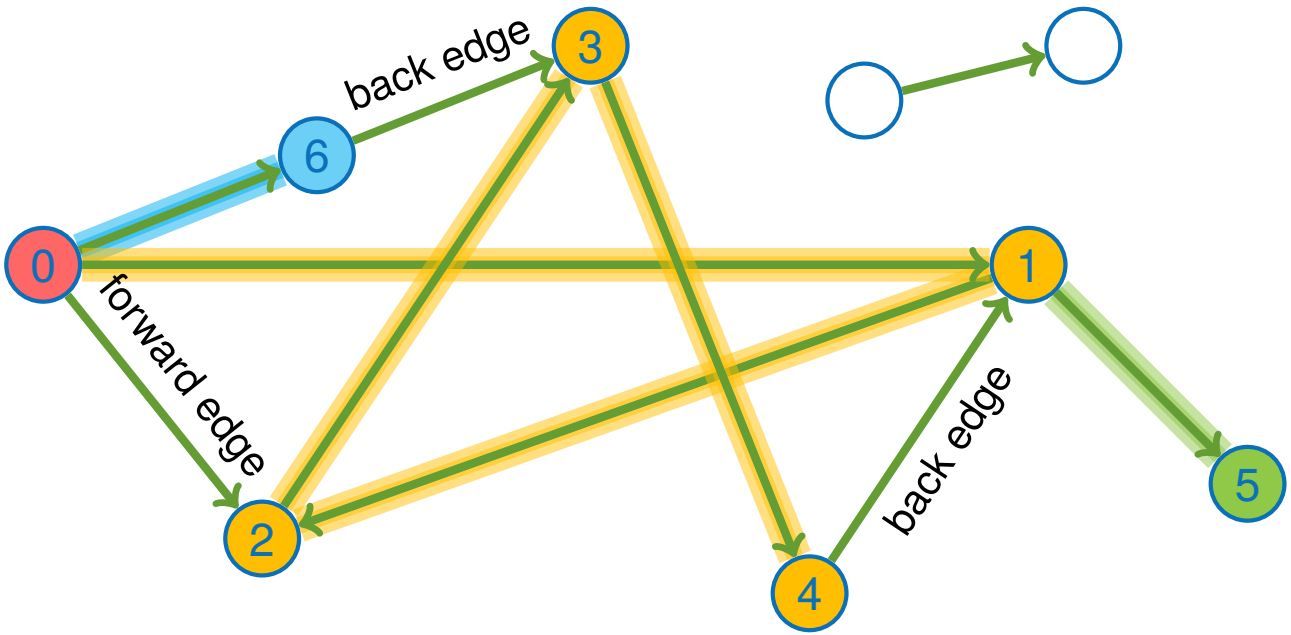

Abbildung: spanning tree of a depth-first search

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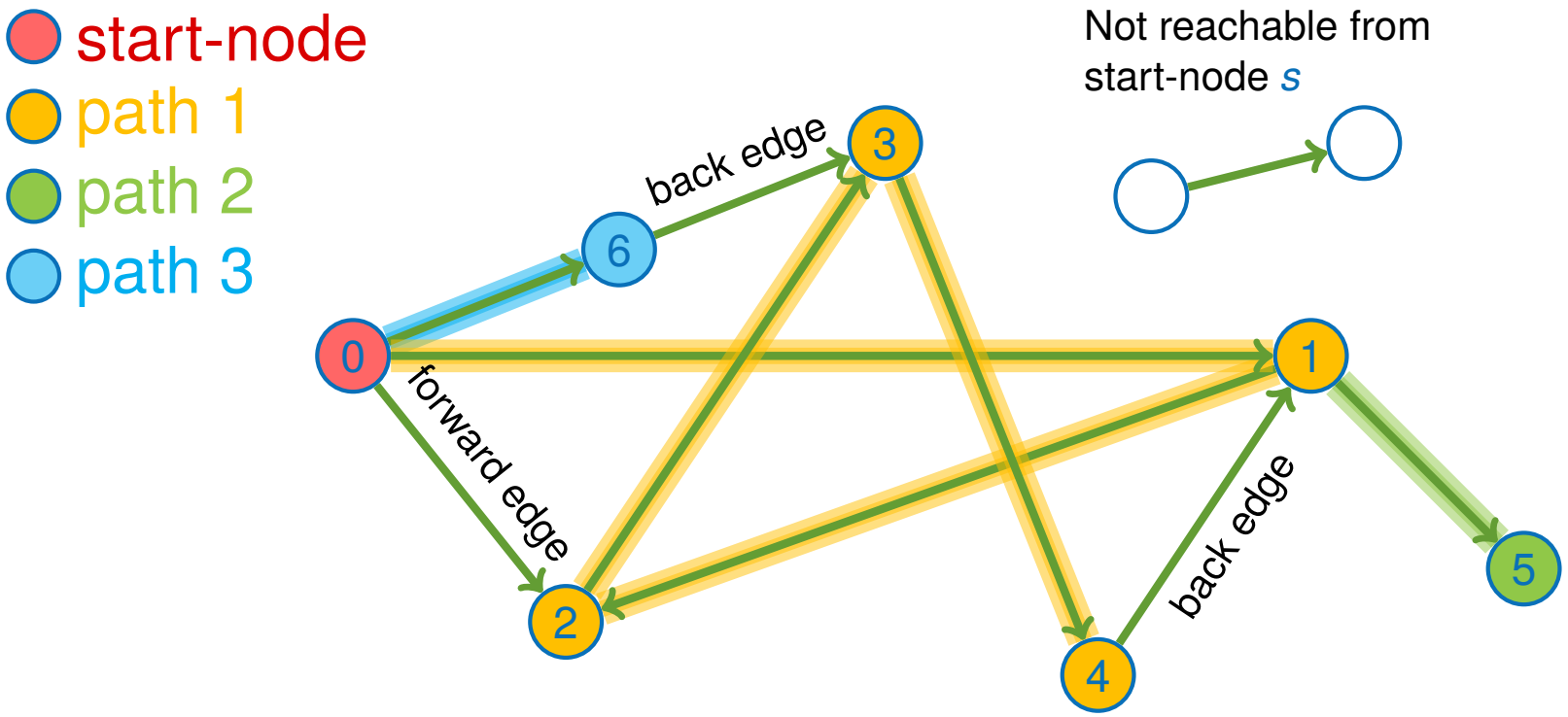

Abbildung: spanning tree of a depth-first search

Runtime complexity:

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- We get a runtime complexity of $\Theta(|V'| + |E'|)$
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- All vertices of V' are in the same connected component as our start vertex s
- This can only be improved by a constant factor

Structure

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Feedback
- Exercises
- Lecture

Graphs
- Introduction
- Implementation
- Application example

Application example

Image processing

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Application example

Image processing

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- Connected component labeling

Application example

Image processing

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- Connected component labeling
- Counting of objects in an image

Application example

Image processing

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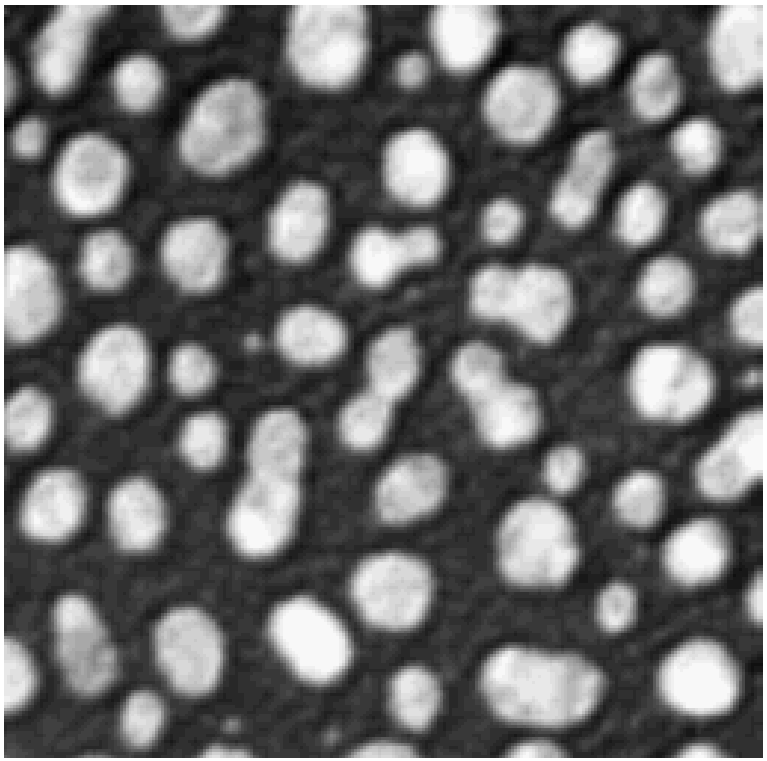
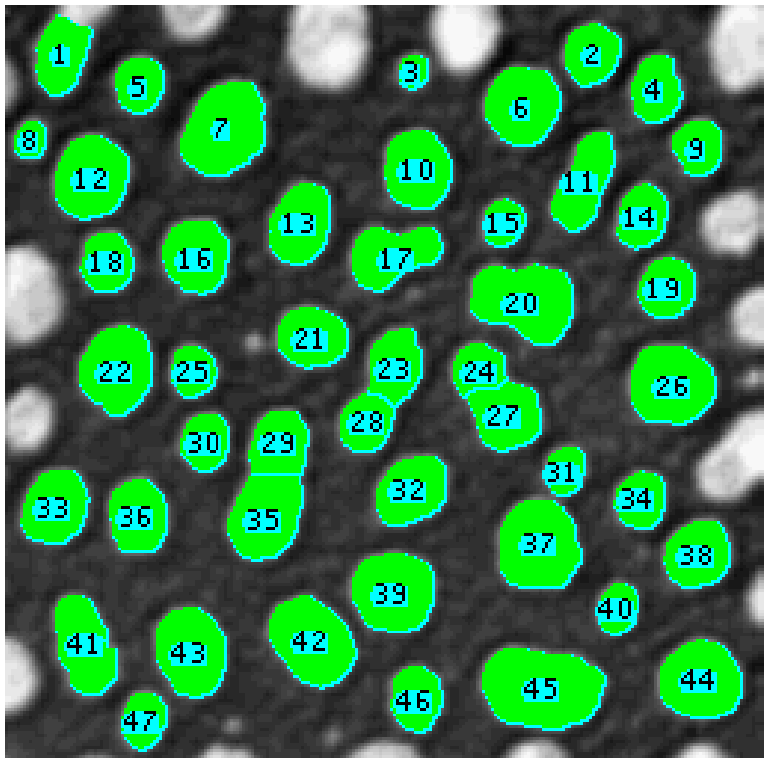
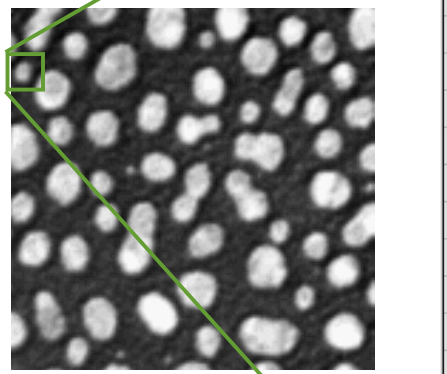
- Counting of objects in an image



$\Rightarrow$



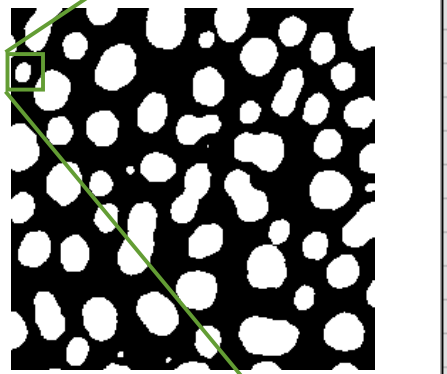
What's object, what's background?



	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U
35	104	80	56	40	16	16	8	16	16	24	32	32	32	32	32	32	32	32	24	24	16
36	80	64	48	32	16	16	16	24	32	40	40	40	40	40	40	40	32	32	24	24	24
37	56	48	32	24	8	16	16	32	40	48	48	48	40	40	40	40	32	32	24	24	24
38	40	32	24	24	16	32	48	64	72	80	80	72	56	56	48	48	40	40	32	32	32
39	16	16	16	24	24	48	72	88	104	112	112	96	72	64	56	48	40	40	40	40	40
40	16	16	24	40	56	88	120	128	136	144	144	120	96	88	72	56	48	48	40	40	40
41	8	16	24	56	80	120	160	168	168	168	168	144	120	104	80	64	48	48	40	40	32
42	16	32	40	80	112	144	176	176	176	176	168	152	128	112	88	64	48	40	32	32	24
43	24	40	56	96	136	160	184	184	176	176	168	152	136	112	88	64	40	32	24	24	16
44	40	56	80	112	152	168	184	184	176	176	168	152	136	112	80	64	40	32	16	16	16
45	48	72	96	128	160	176	184	184	176	176	168	152	136	104	72	56	32	24	8	16	16
46	48	72	96	136	168	176	192	192	184	184	176	160	136	104	72	56	32	24	16	24	32
47	48	72	96	136	168	184	192	192	192	192	184	160	136	104	72	48	24	24	16	32	48
48	48	72	96	128	168	184	200	200	200	192	184	160	128	96	64	48	24	32	32	56	72
49	48	72	88	128	160	184	200	200	200	192	184	152	120	88	56	40	24	32	40	72	96
50	48	64	80	112	136	160	176	176	176	168	160	136	104	80	48	40	32	40	56	88	128
51	48	64	72	96	112	128	144	152	152	144	136	112	88	64	40	40	32	48	64	112	152
52	48	56	64	80	88	104	112	112	120	112	104	88	72	56	32	32	32	64	88	128	168
53	40	48	48	56	64	72	72	80	80	80	72	64	48	40	24	32	32	72	104	144	184
54	48	48	48	48	48	56	56	56	64	56	56	48	40	32	24	40	48	88	128	160	200

Convert to black white using threshold:

```
value = 255 if value > 100 else 0
```



	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P
35	255	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
36	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
37	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
38	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
39	0	0	0	0	0	0	0	0	255	255	255	0	0	0	0	0
40	0	0	0	0	0	0	255	255	255	255	255	255	0	0	0	0
41	0	0	0	0	0	255	255	255	255	255	255	255	255	255	0	0
42	0	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0
43	0	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0
44	0	0	0	255	255	255	255	255	255	255	255	255	255	255	0	0
45	0	0	0	255	255	255	255	255	255	255	255	255	255	255	0	0
46	0	0	0	255	255	255	255	255	255	255	255	255	255	255	0	0
47	0	0	0	255	255	255	255	255	255	255	255	255	255	255	0	0
48	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0
49	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0
50	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0
51	0	0	0	0	255	255	255	255	255	255	255	255	0	0	0	0
52	0	0	0	0	0	255	255	255	255	255	255	0	0	0	0	0
53	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
54	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
55	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
56	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
57	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
58	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

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Interpret image as graph:

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Interpret image as graph:

- Each white pixel is a node

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- Edges between adjacent pixels (normally 4 or 8 neighbors)

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Interpret image as graph:

- Each white pixel is a node
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- Edges are not saved externally, algorithm works directly on array

Interpret image as graph:

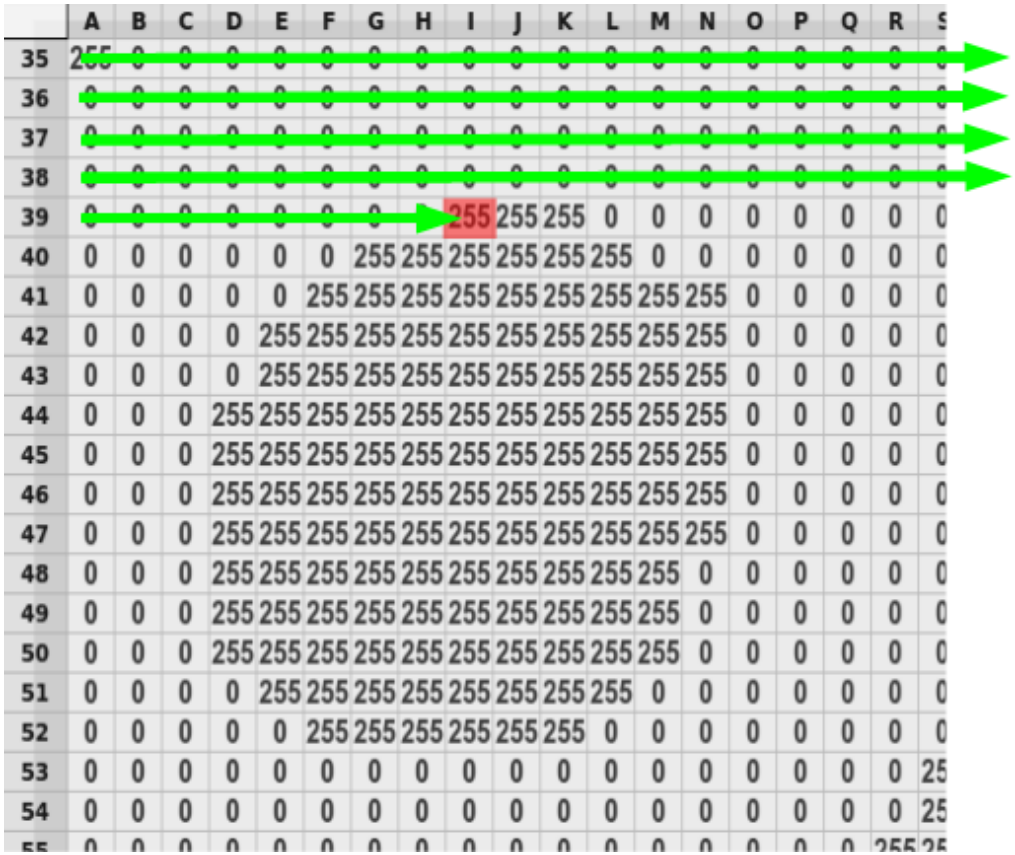
- Each white pixel is a node
- Edges between adjacent pixels (normally 4 or 8 neighbors)
- Edges are not saved externally, algorithm works directly on array
- Breadth- / depth-first search find all connected components (particles)

Find connected components:

Application example

Image processing

Find connected components:

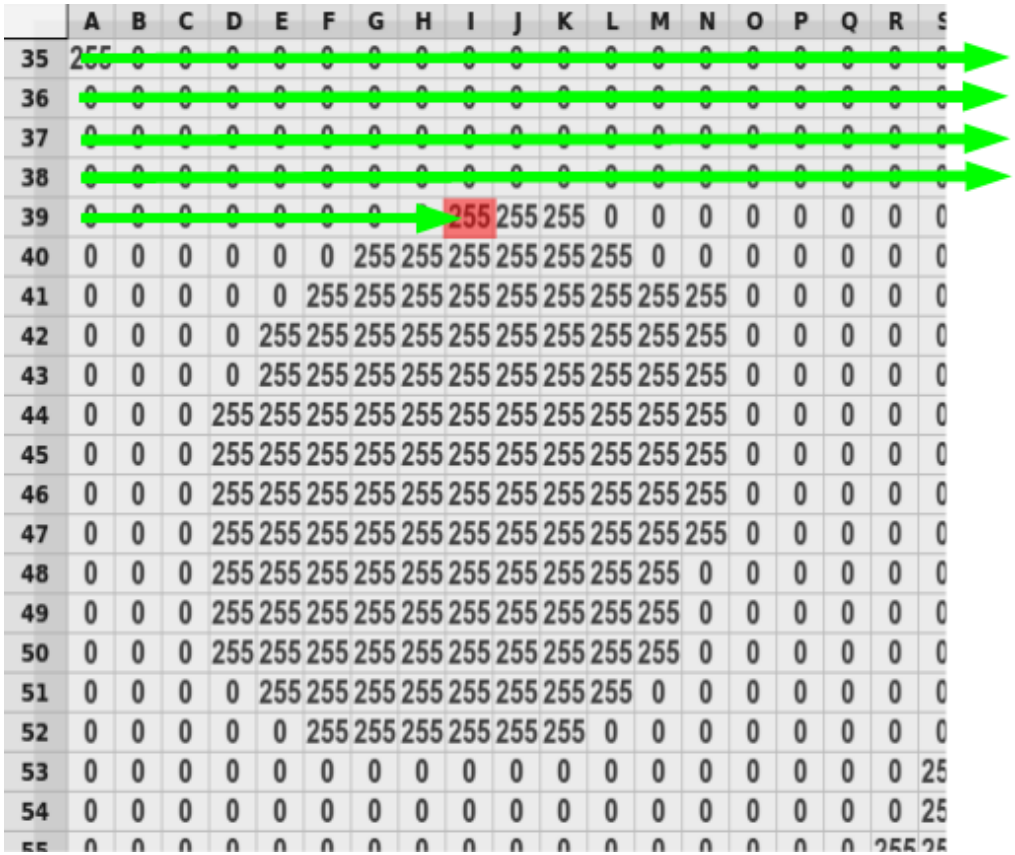


	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S
35	255	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
36	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
37	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
38	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
39	0	0	0	0	0	0	0	255	255	255	0	0	0	0	0	0	0	0	0
40	0	0	0	0	0	0	255	255	255	255	255	255	0	0	0	0	0	0	0
41	0	0	0	0	0	255	255	255	255	255	255	255	255	255	0	0	0	0	0
42	0	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
43	0	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
44	0	0	0	255	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
45	0	0	0	255	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
46	0	0	0	255	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
47	0	0	0	255	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
48	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0	0
49	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0	0
50	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0	0
51	0	0	0	0	255	255	255	255	255	255	255	255	0	0	0	0	0	0	0
52	0	0	0	0	0	255	255	255	255	255	255	0	0	0	0	0	0	0	0
53	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	25
54	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	25
55	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	255	25

Application example

Image processing

Find connected components:



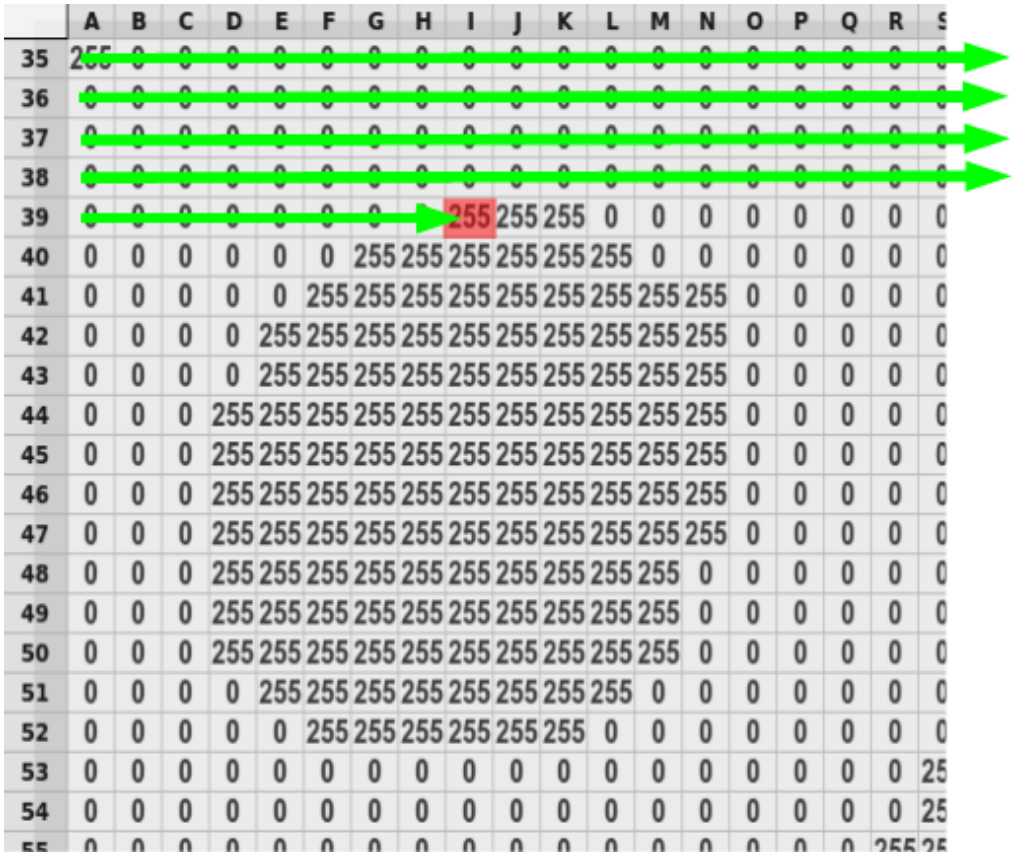
	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S
35	255	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
36	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
37	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
38	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
39	0	0	0	0	0	0	0	255	255	255	0	0	0	0	0	0	0	0	0
40	0	0	0	0	0	0	255	255	255	255	255	255	0	0	0	0	0	0	0
41	0	0	0	0	0	255	255	255	255	255	255	255	255	255	0	0	0	0	0
42	0	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
43	0	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
44	0	0	0	255	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
45	0	0	0	255	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
46	0	0	0	255	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
47	0	0	0	255	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
48	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0	0
49	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0	0
50	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0	0
51	0	0	0	0	255	255	255	255	255	255	255	255	0	0	0	0	0	0	0
52	0	0	0	0	0	255	255	255	255	255	255	0	0	0	0	0	0	0	0
53	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	25
54	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	25
55	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	255	25

- Search pixel-by-pixel for non-zero intensity

Application example

Image processing

Find connected components:



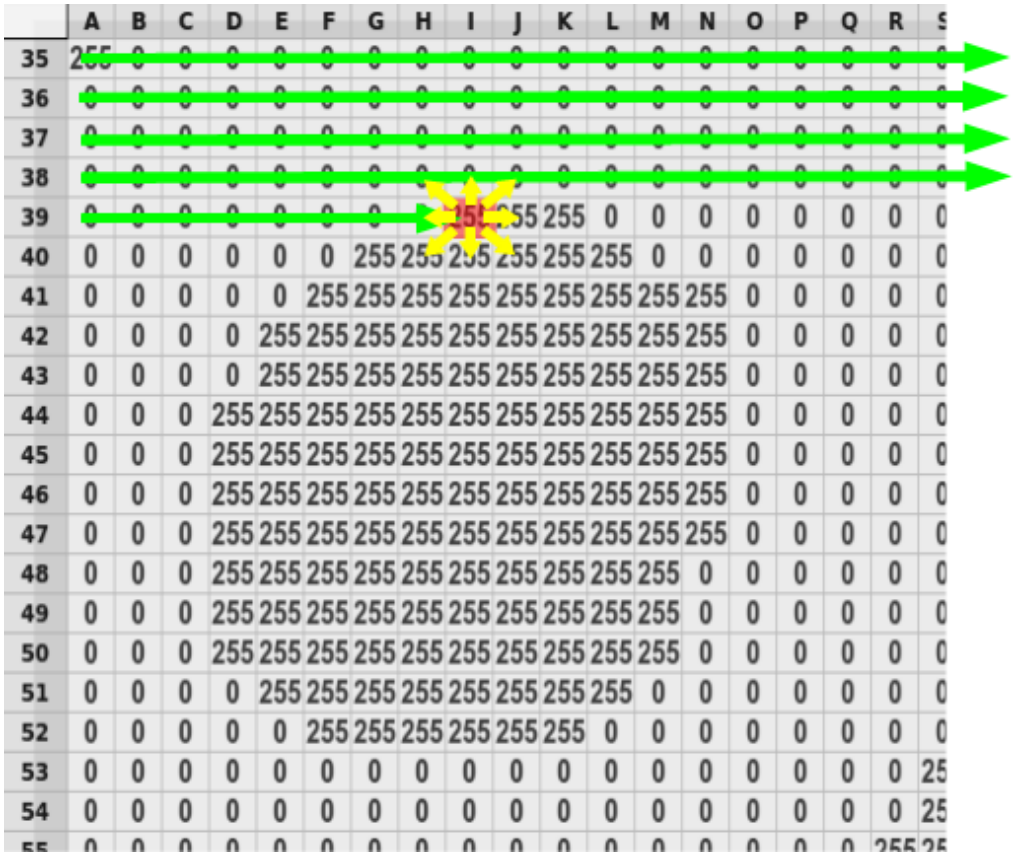
	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S
35	255	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
36	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
37	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
38	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
39	0	0	0	0	0	0	0	255	255	255	0	0	0	0	0	0	0	0	0
40	0	0	0	0	0	0	255	255	255	255	255	255	0	0	0	0	0	0	0
41	0	0	0	0	0	255	255	255	255	255	255	255	255	255	0	0	0	0	0
42	0	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
43	0	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
44	0	0	0	255	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
45	0	0	0	255	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
46	0	0	0	255	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
47	0	0	0	255	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
48	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0	0
49	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0	0
50	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0	0
51	0	0	0	0	255	255	255	255	255	255	255	255	0	0	0	0	0	0	0
52	0	0	0	0	0	255	255	255	255	255	255	0	0	0	0	0	0	0	0
53	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	25
54	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	25
55	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	255	25

- Search pixel-by-pixel for non-zero intensity
- Label found pixel as component 1

Application example

Image processing

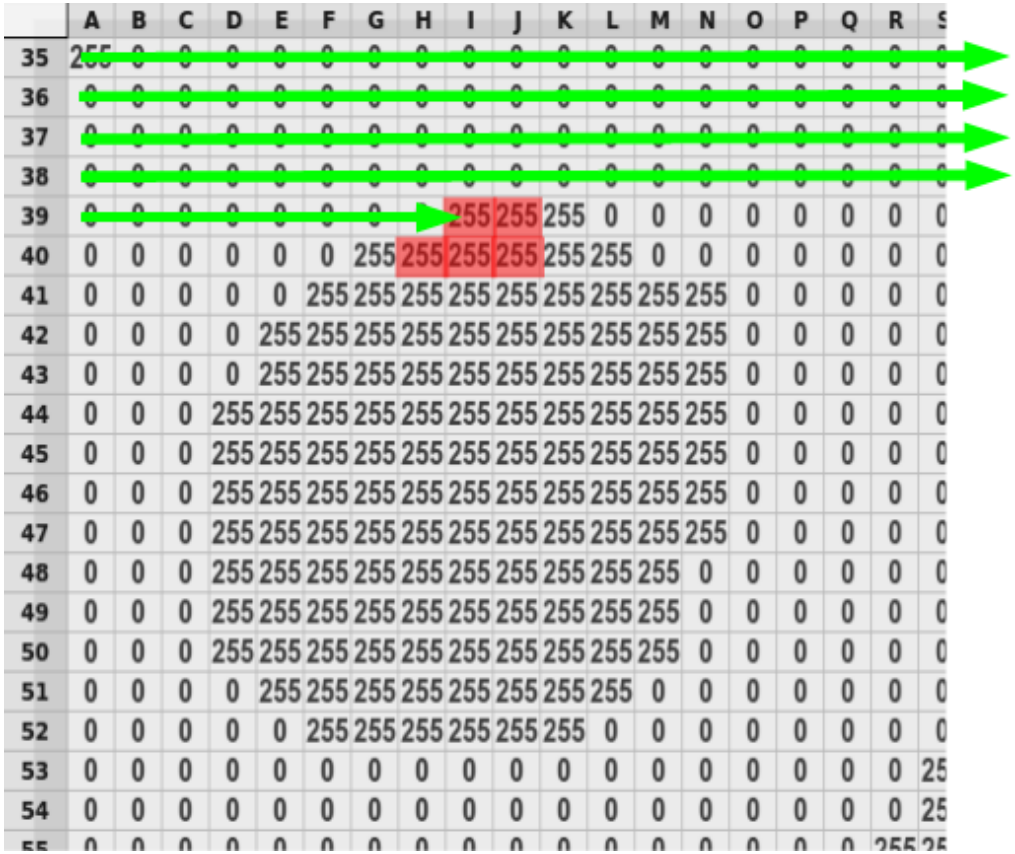
Find connected components:



	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S
35	255	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
36	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
37	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
38	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
39	0	0	0	0	0	0	0	255	255	255	0	0	0	0	0	0	0	0	0
40	0	0	0	0	0	0	255	255	255	255	255	255	0	0	0	0	0	0	0
41	0	0	0	0	0	255	255	255	255	255	255	255	255	255	0	0	0	0	0
42	0	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
43	0	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
44	0	0	0	255	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
45	0	0	0	255	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
46	0	0	0	255	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
47	0	0	0	255	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
48	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0	0
49	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0	0
50	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0	0
51	0	0	0	0	255	255	255	255	255	255	255	255	0	0	0	0	0	0	0
52	0	0	0	0	0	255	255	255	255	255	255	0	0	0	0	0	0	0	0
53	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	25
54	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	25
55	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	255	25

- Search pixel-by-pixel for non-zero intensity
- Label found pixel as component 1
- Check neighbors of all new labeled pixels

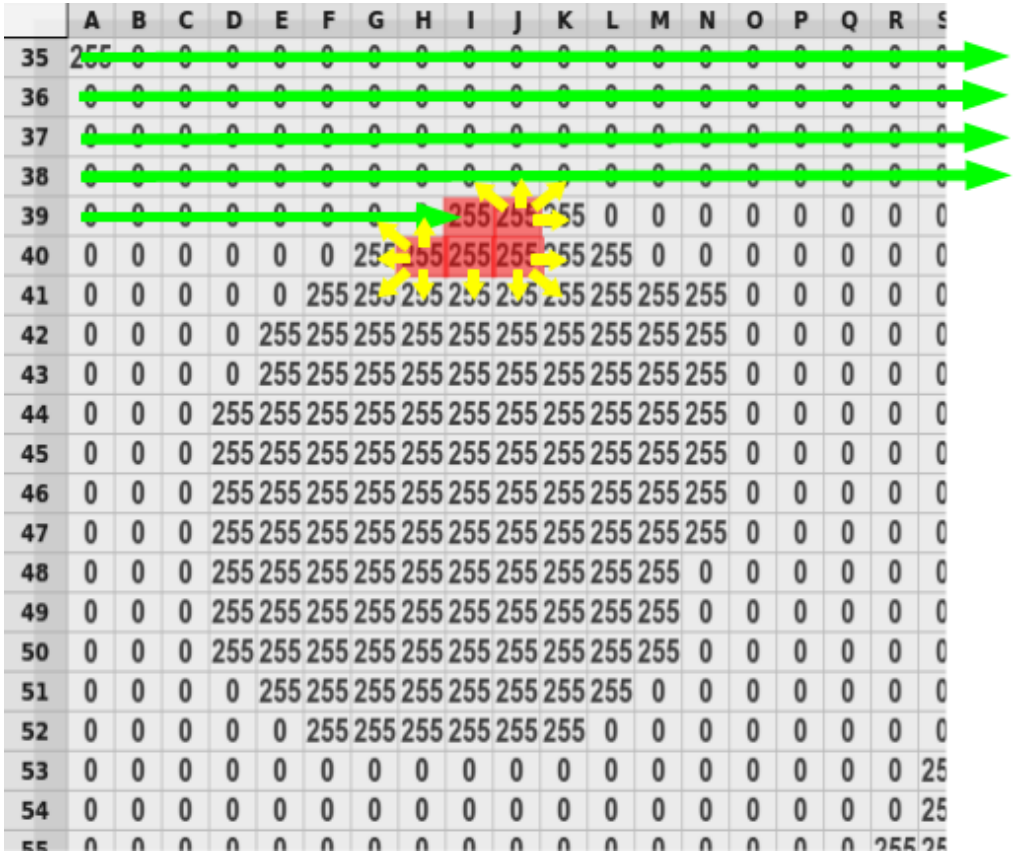
Find connected components:



	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S
35	255	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
36	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
37	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
38	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
39	0	0	0	0	0	0	0	255	255	255	0	0	0	0	0	0	0	0	0
40	0	0	0	0	0	0	255	255	255	255	255	255	0	0	0	0	0	0	0
41	0	0	0	0	0	255	255	255	255	255	255	255	255	255	0	0	0	0	0
42	0	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
43	0	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
44	0	0	0	255	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
45	0	0	0	255	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
46	0	0	0	255	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
47	0	0	0	255	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
48	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0	0
49	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0	0
50	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0	0
51	0	0	0	0	255	255	255	255	255	255	255	255	0	0	0	0	0	0	0
52	0	0	0	0	0	255	255	255	255	255	255	0	0	0	0	0	0	0	0
53	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	25
54	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	25
55	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	255	25

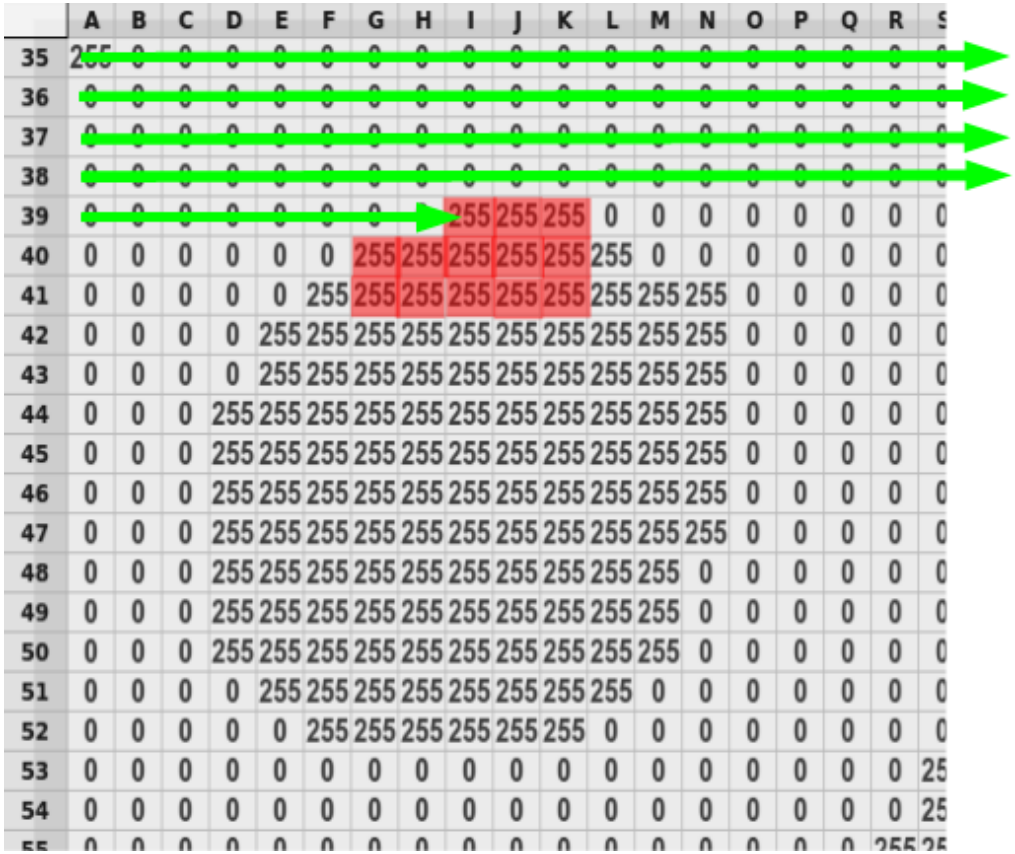
- Search pixel-by-pixel for non-zero intensity
- Label found pixel as component 1
- Check neighbors of all new labeled pixels
- Label non-zero pixels as component 1

Find connected components:



- Search pixel-by-pixel for non-zero intensity
- Label found pixel as component 1
- Check neighbors of all new labeled pixels
- Label non-zero pixels as component 1

Find connected components:



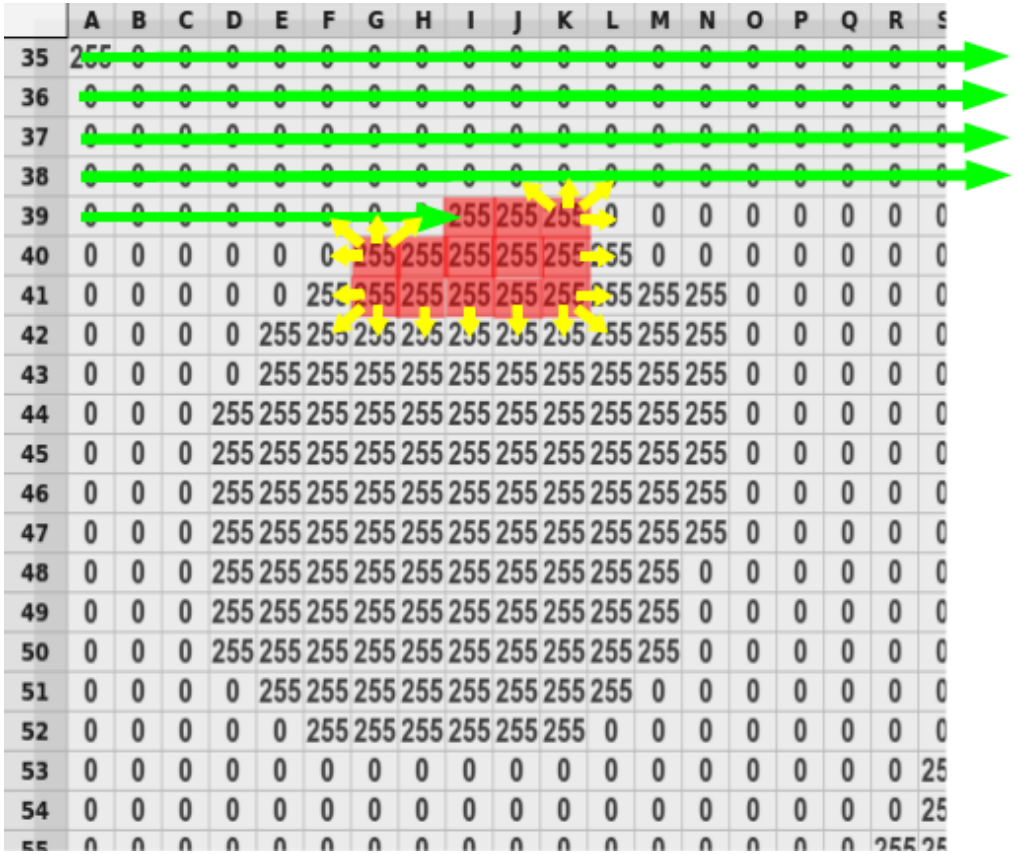
	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S
35	255	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
36	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
37	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
38	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
39	0	0	0	0	0	0	0	255	255	255	0	0	0	0	0	0	0	0	0
40	0	0	0	0	0	0	255	255	255	255	255	255	0	0	0	0	0	0	0
41	0	0	0	0	0	255	255	255	255	255	255	255	255	255	0	0	0	0	0
42	0	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
43	0	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
44	0	0	0	255	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
45	0	0	0	255	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
46	0	0	0	255	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
47	0	0	0	255	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
48	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0	0
49	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0	0
50	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0	0
51	0	0	0	0	255	255	255	255	255	255	255	255	0	0	0	0	0	0	0
52	0	0	0	0	0	255	255	255	255	255	255	0	0	0	0	0	0	0	0
53	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	25
54	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	25
55	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	255	25

- Search pixel-by-pixel for non-zero intensity
- Label found pixel as component 1
- Check neighbors of all new labeled pixels
- Label non-zero pixels as component 1

Application example

Image processing

Find connected components:

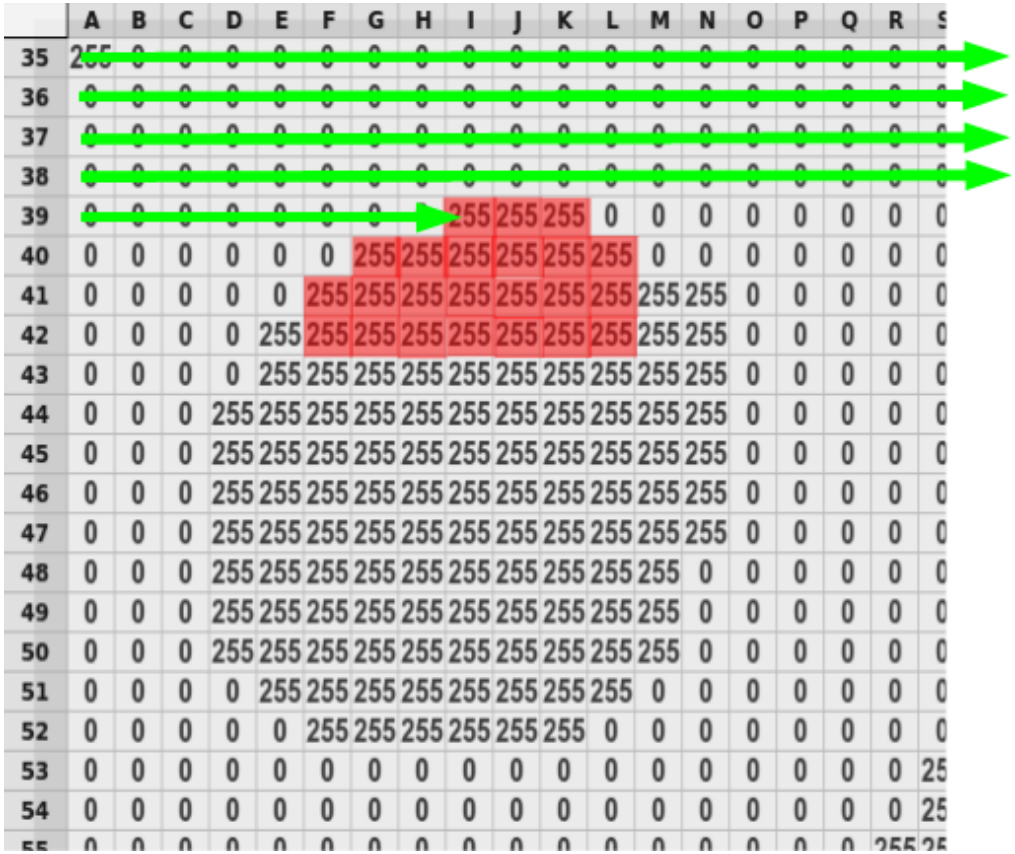


- Search pixel-by-pixel for non-zero intensity
- Label found pixel as component 1
- Check neighbors of all new labeled pixels
- Label non-zero pixels as component 1

Application example

Image processing

Find connected components:



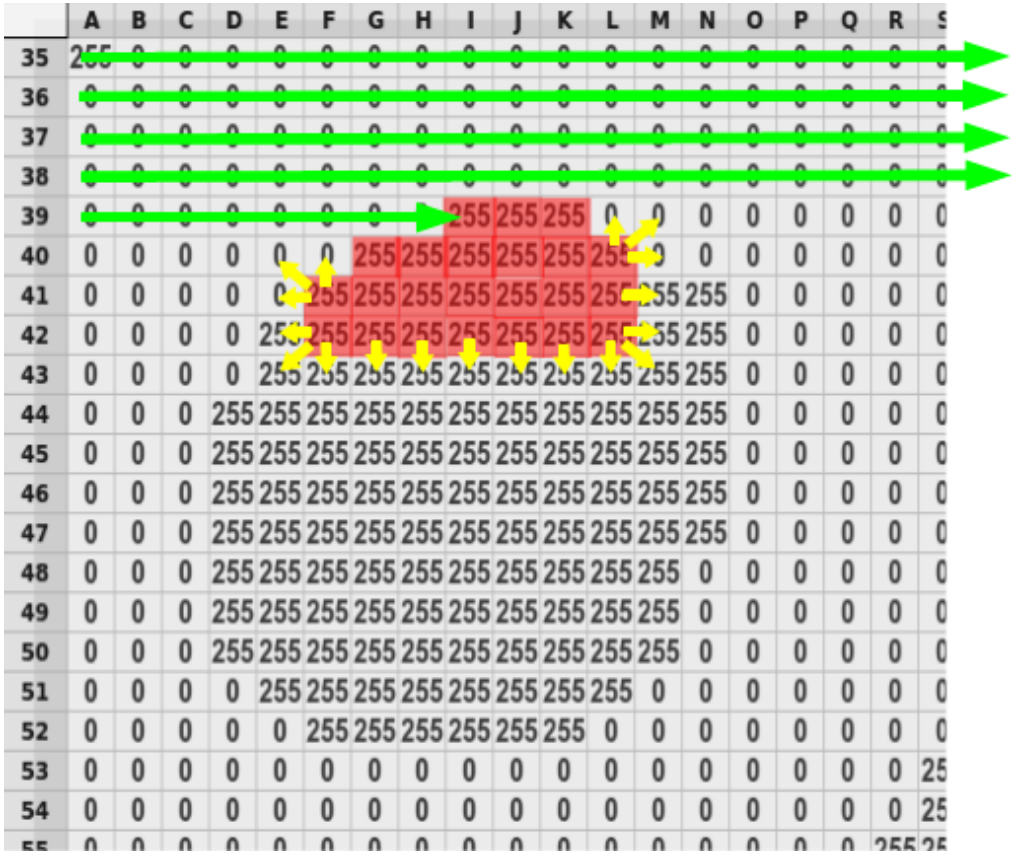
	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S
35	255	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
36	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
37	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
38	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
39	0	0	0	0	0	0	0	255	255	255	0	0	0	0	0	0	0	0	0
40	0	0	0	0	0	0	255	255	255	255	255	255	0	0	0	0	0	0	0
41	0	0	0	0	0	255	255	255	255	255	255	255	255	255	0	0	0	0	0
42	0	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
43	0	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
44	0	0	0	255	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
45	0	0	0	255	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
46	0	0	0	255	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
47	0	0	0	255	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
48	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0	0
49	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0	0
50	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0	0
51	0	0	0	0	255	255	255	255	255	255	255	255	0	0	0	0	0	0	0
52	0	0	0	0	0	255	255	255	255	255	255	0	0	0	0	0	0	0	0
53	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	25
54	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	25
55	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	255	25

- Search pixel-by-pixel for non-zero intensity
- Label found pixel as component 1
- Check neighbors of all new labeled pixels
- Label non-zero pixels as component 1

Application example

Image processing

Find connected components:

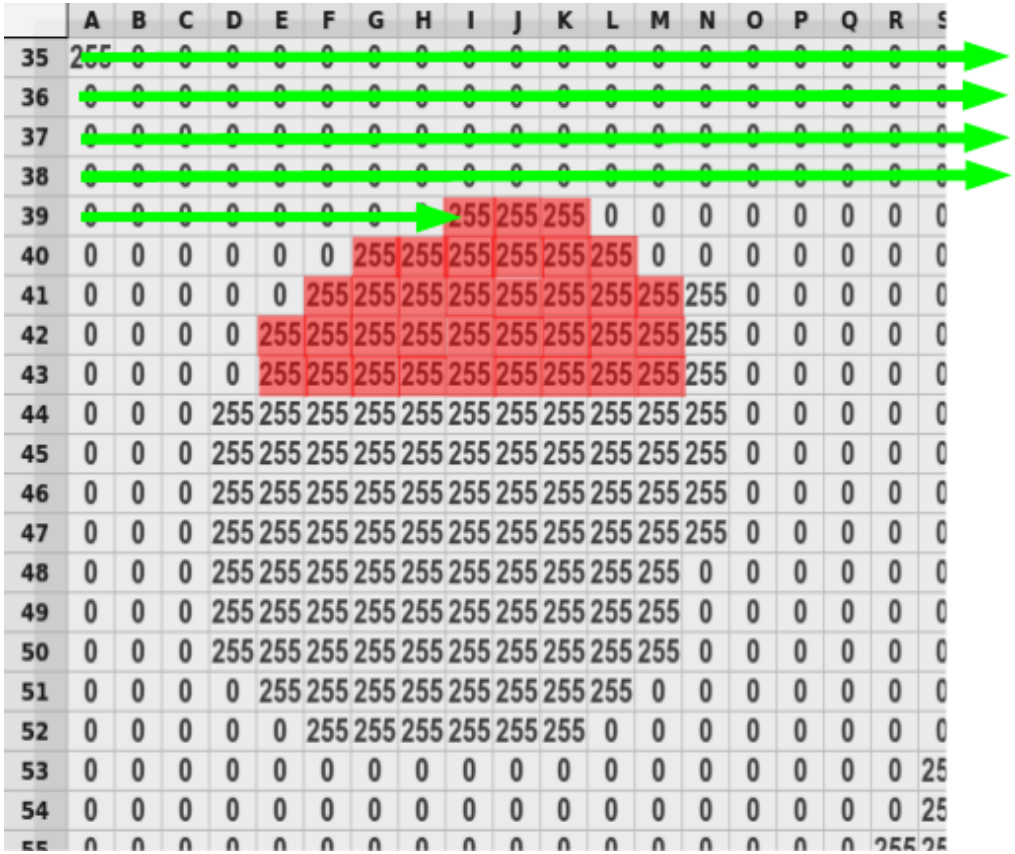


- Search pixel-by-pixel for non-zero intensity
- Label found pixel as component 1
- Check neighbors of all new labeled pixels
- Label non-zero pixels as component 1

Application example

Image processing

Find connected components:



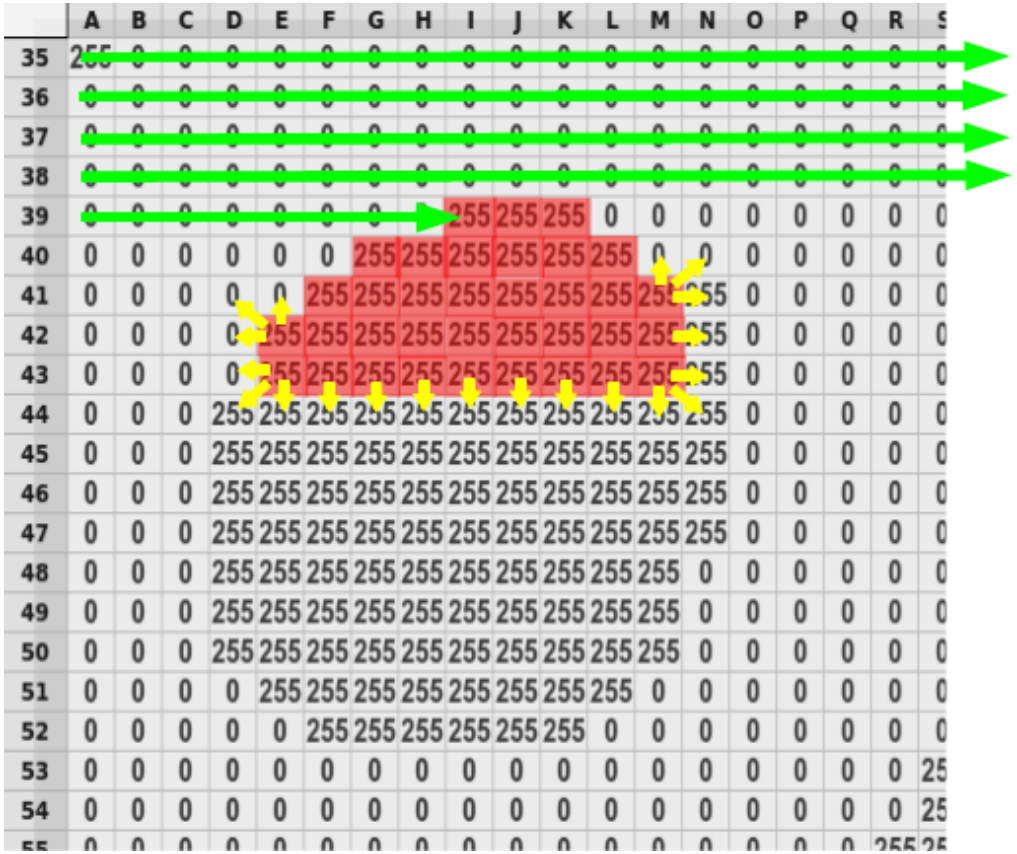
	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S
35	255	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
36	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
37	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
38	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
39	0	0	0	0	0	0	0	255	255	255	0	0	0	0	0	0	0	0	0
40	0	0	0	0	0	0	255	255	255	255	255	255	0	0	0	0	0	0	0
41	0	0	0	0	0	255	255	255	255	255	255	255	255	255	0	0	0	0	0
42	0	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
43	0	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
44	0	0	0	255	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
45	0	0	0	255	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
46	0	0	0	255	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
47	0	0	0	255	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0
48	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0	0
49	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0	0
50	0	0	0	255	255	255	255	255	255	255	255	255	255	0	0	0	0	0	0
51	0	0	0	0	255	255	255	255	255	255	255	255	0	0	0	0	0	0	0
52	0	0	0	0	0	255	255	255	255	255	255	0	0	0	0	0	0	0	0
53	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	25
54	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	25
55	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	255	25

- Search pixel-by-pixel for non-zero intensity
- Label found pixel as component 1
- Check neighbors of all new labeled pixels
- Label non-zero pixels as component 1

Application example

Image processing

Find connected components:

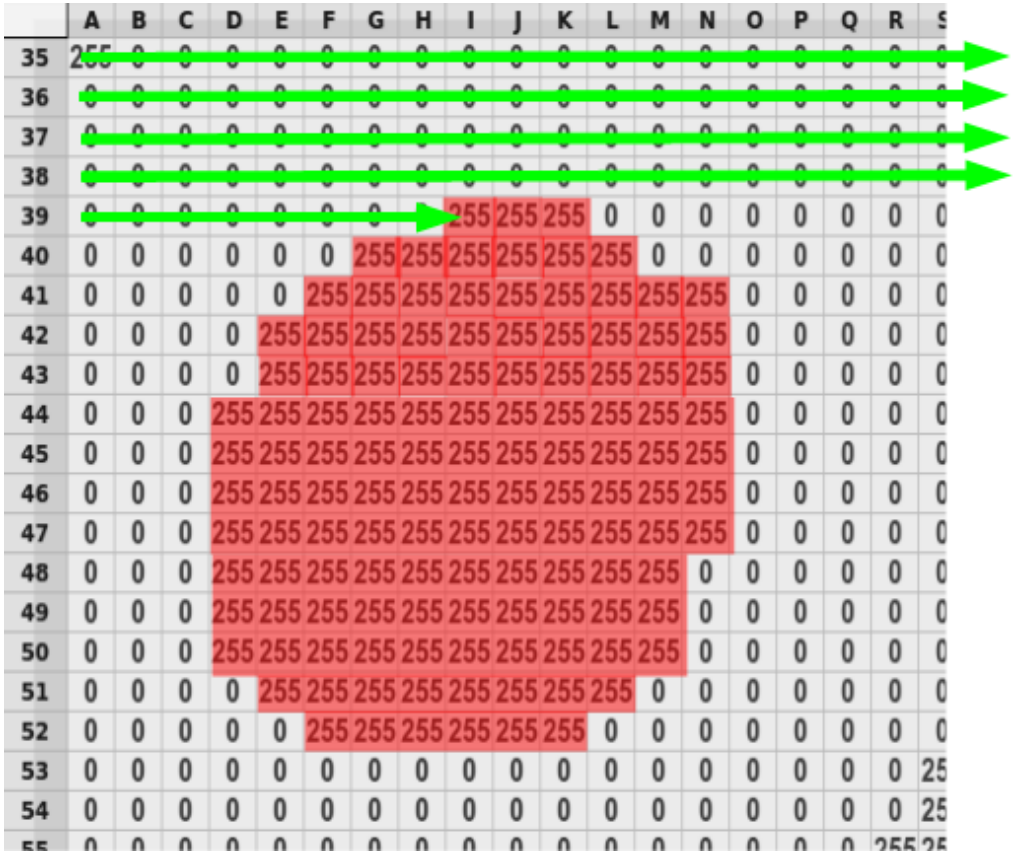


- Search pixel-by-pixel for non-zero intensity
- Label found pixel as component 1
- Check neighbors of all new labeled pixels
- Label non-zero pixels as component 1

Application example

Image processing

Find connected components:



- Search pixel-by-pixel for non-zero intensity
- Label found pixel as component 1
- Check neighbors of all new labeled pixels
- Label non-zero pixels as component 1

Find connected components:

- Search pixel-by-pixel for non-zero intensity
- Label found pixel as component 2
- ...

Result of connected component labeling:

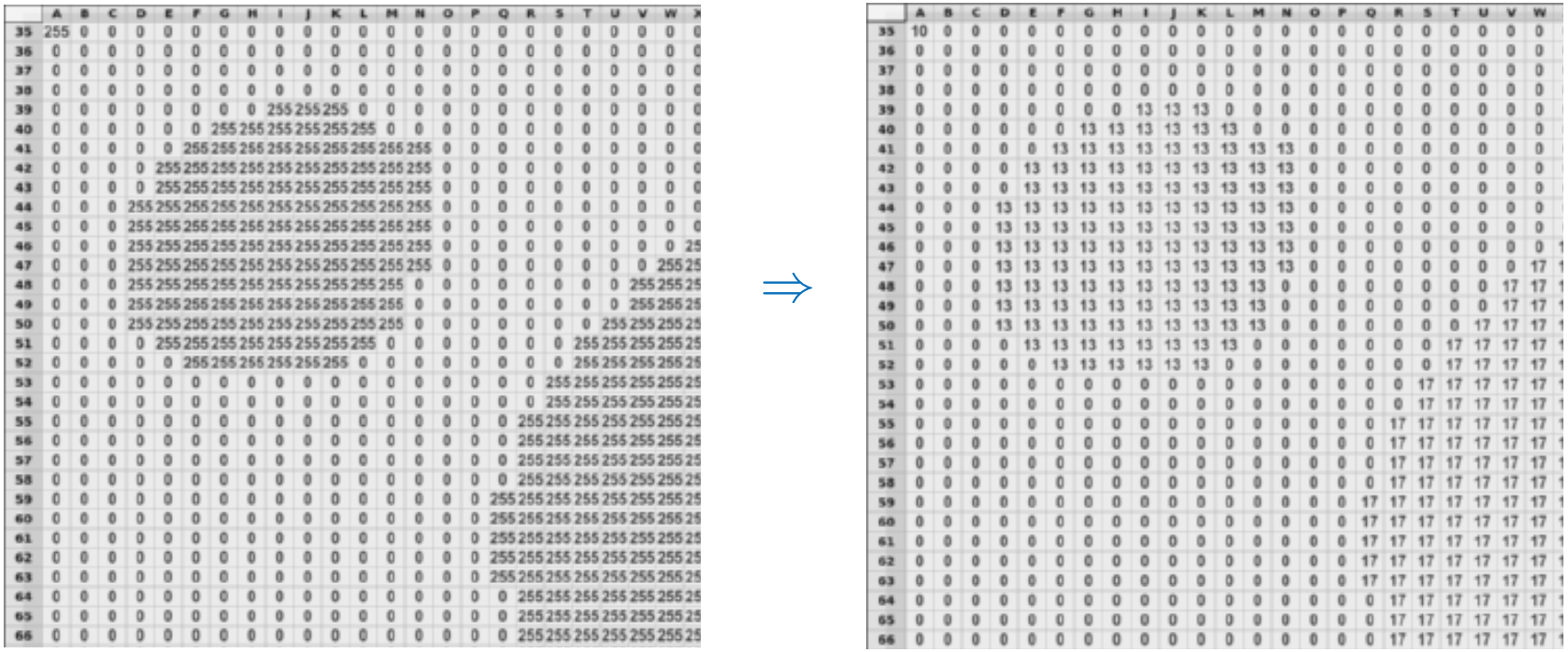


Abbildung: Result: particle indices instead of intensities

Further Literature

UNI FREIBURG

- **General**

[CRL01] Thomas H. Cormen, Ronald L. Rivest, and Charles E. Leiserson.
Introduction to Algorithms.
MIT Press, Cambridge, Mass, 2001.

[MS08] Kurt Mehlhorn and Peter Sanders.
Algorithms and data structures, 2008.
https://people.mpi-inf.mpg.de/~mehlhorn/ftp/Mehlhorn-Sanders-Toolbox.pdf.

Further Literature

UNI FREIBURG

- **Graph-Search**

  [Wika] Breadth-first search
  `https://en.wikipedia.org/wiki/Breadth-first_search`

  [Wikb] Depth-first search
  `https://en.wikipedia.org/wiki/Depth-first_search`

- **Graph-Connectivity**

  [Wik] Connectivity (graph theory)
  `https://en.wikipedia.org/wiki/Connectivity_(graph_theory)`