# Algorithmns and Datastructures Graphs, Depth-/Breadth-first Search, Graph-Connectivity



Albert-Ludwigs-Universität Freiburg

#### Prof. Dr. Rolf Backofen

Bioinformatics Group / Department of Computer Science Algorithmns and Datastructures, January 2017

#### Structure



#### Feedback

Exercises Lecture

#### Graphs

Introduction Implementation Application example

#### Structure



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#### Feedback from the exercises



The upcoming exercise sheet 12 and 13 will be merged together (finding largest connected component + Dijkstra)

Some people were asking for more solution sheets for the exercises

We are working on it.

#### Feedback from the lecture



Code in the lecture will be a little bit different from exercise sheet.

One person asked for additional explanations regarding proofs.

#### Structure



Feedback Exercises

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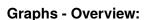
**Graphs - Overview:** 



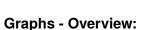
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#### **Graphs - Overview:**

 Besides arrays, lists and trees the most common datastructure (Trees are a special type of graph)



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- Breadth first search (BFS)

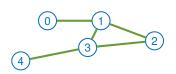


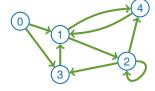
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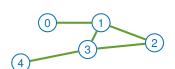
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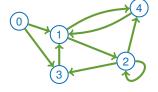
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- Connected components of a graph

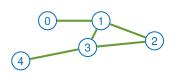


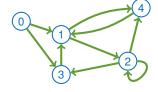




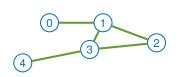


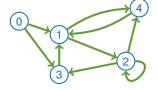
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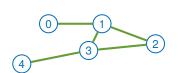


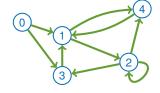
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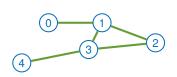


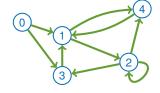
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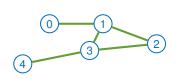


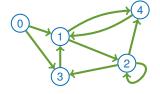
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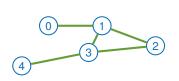


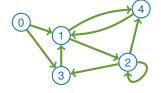
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  - Directed edge: e = (u, v) (tuple)
- Self-loops are also possible: e = (u, u) or  $e = \{u, u\}$



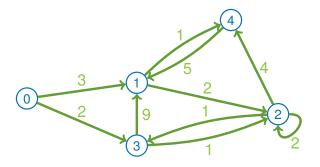
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# Weighted graph:



# UN:

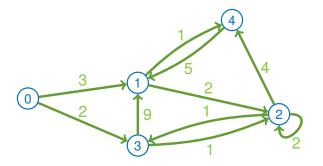
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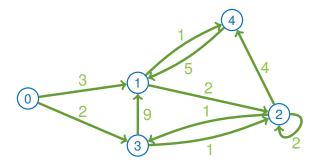
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#### Weighted graph:



Each edge is marked with a real number named weight

#### Weighted graph:



- Each edge is marked with a real number named weight
- The weight is also named length or cost of the edge depending on the application

# Graphs Introduction



Example: Road network



**Example:** Road network

Intersections:

vertices



**Example:** Road network

Intersections:

vertices

■ Roads: edges

#### **Example:** Road network

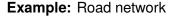
Intersections:

vertices

Roads: edges

Travel time:

costs of the edges



- Intersections: vertices
- Roads: edges
- Travel time: costs of the edges



Figure: Map of Freiburg © OpenStreetMap

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Two classic variants



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  - 1 Adjacency matrix with space consumption  $\Theta(|V|^2)$

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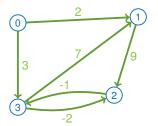


Figure: Weighted graph with

$$|V| = 4$$
,  $|E| = 6$ 

#### How to represent this graph computationally?

- Two classic variants
  - Adjacency matrix with space consumption  $\Theta(|V|^2)$

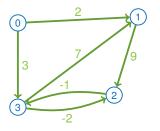


Figure: Weighted graph with |V| = 4, |E| = 6

|               | end-vertice |   |    |    |
|---------------|-------------|---|----|----|
|               | 0           | 1 | 2  | 3  |
| ice<br>0      |             | 2 |    | 3  |
| start-vertice |             |   | 9  |    |
| £ 2           |             |   |    | -1 |
| sta 3         |             | 7 | -2 |    |

Figure: Adjacency matrix

### Graphs

Implementation

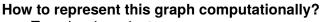
EIBURG

How to represent this graph computationally?

# EIBURG

#### How to represent this graph computationally?

- Two classic variants
  - 2 Adjacency list / fields with space consumption  $\Theta(|V| + |E|)$



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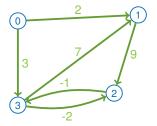
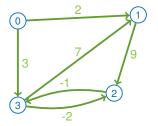


Figure: Weighted graph with

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- Two classic variants
  - 2 Adjacency list / fields with space consumption  $\Theta(|V| + |E|)$ 
    - Each list item stores the target vertice and the cost of the edge



| <u>8</u> (0)  | 1, 2  | 3, 3  |
|---------------|-------|-------|
| start-vertice | 2, 9  |       |
| £ 2           | 3, -1 |       |
| sta 3         | 1, 7  | 2, -2 |
|               |       | •     |

Figure: Weighted graph with

$$|V| = 4$$
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Figure: Adjacency list

# Graphs Implementation

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#### Graphs

Implementation



#### **Graph: Arrangement**

■ Graph is fully defined through the adjacency matrix / list

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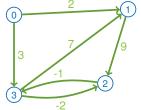


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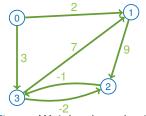


Figure: Weighted graph with |V| = 4, |E| = 6

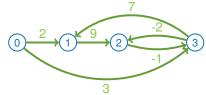


Figure: Same graph ordered by number - outer planar graph

```
class Graph:
    def init (self):
        self.vertices = []
        self.edges = []
    def addVertice(self, vert):
        self.vertices.append(vert)
    def addEdge(self, fromVert, toVert):
        self.edges.append((fromVert, toVert))
```



# Graphs Degrees (Valency)

**Degree of a vertex:** Directed graph: G = (V, E)



### Graphs

Degrees (Valency)



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Figure: Vertex with in- / outdegree of 3 / 2



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$$\deg^+(u) = |\{(v, u) : (v, u) \in E\}|$$

### Degrees (Valency)

**Degree of a vertex:** Directed graph: G = (V, E)



Figure: Vertex with in- / outdegree of 3 / 2

Indegree of a vertex u is the number of edge heads adjacent to the vertex

$$\deg^+(u) = |\{(v, u) : (v, u) \in E\}|$$

 Outdegree of a vertex u is the number of edge tails adjacent to the vertex

$$\deg^{-}(u) = |\{(u, v) : (u, v) \in E\}|$$

### Graphs Degrees (Valency)



FREE

**Degree of a vertex:** Undirected graph: G = (V, E)



Figure: Vertex with degree of 4

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# Graphs Paths





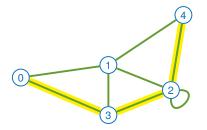


Figure: Undirected path of length 3 P = (0,3,2,4)

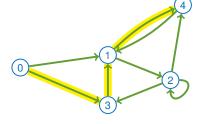


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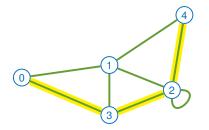


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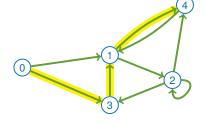


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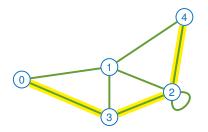


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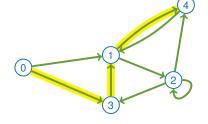


Figure: Directed path of length 3 P = (0, 3, 1, 4)

- A path of G is a sequence of edges  $u_1, u_2, ..., u_i \in V$  with
  - Undirected graph:  $\{u_1, u_2\}, \{u_2, u_3\}, \dots, \{u_{i-1}, u_i\} \in E$
  - Directed graph:  $(u_1, u_2), (u_2, u_3), \dots, (u_{i-1}, u_i) \in E$

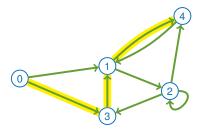


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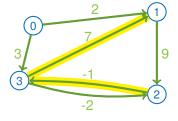


Figure: Weighted path with cost 6 P = (2,3,1)

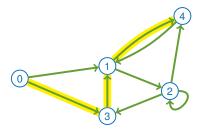


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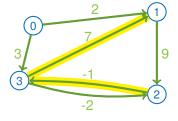


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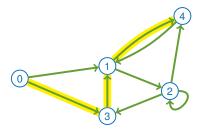


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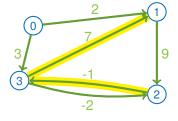
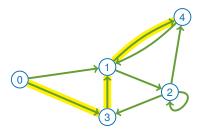


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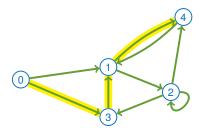


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Figure: Directed path of length 3 P = (0,3,1,4)

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The length of a path is: (also costs of a path)

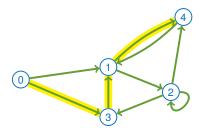


3 -1 -2 2

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## Graphs Paths



Shortest path in a graph: G = (V, E)



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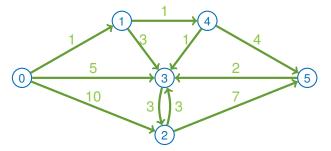


Figure: Shortest path from 0 to 2 with cost / distance d(0,2) = ?

#### Shortest path in a graph: G = (V, E)

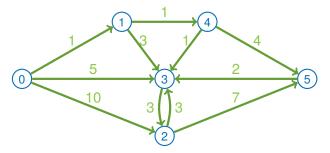


Figure: Shortest path from 0 to 2 with cost / distance d(0,2) = ?

The shortest path between two vertices u, v is the path P = (u, ..., v) with the shortest length d(u, v) or lowest costs



#### Shortest path in a graph: G = (V, E)

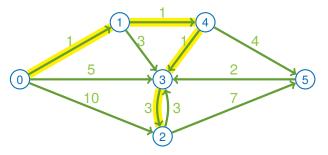


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# Graphs Paths



Diameter of a graph: G = (V, E)



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$$d = \max_{u,v \in V} d(u,v)$$

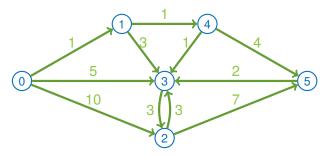


Figure: Diameter of graph is d = ?

Paths

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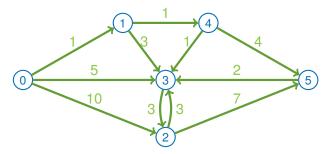


Figure: Diameter of graph is d = ?

The diameter of a graph is the length / the costs of the longest shortest path

Paths

# Diameter of a graph: G = (V, E)

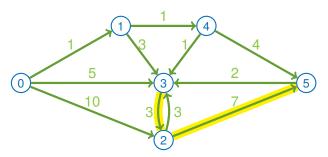


Figure: Diameter of graph is d = 10, P = (3, 2, 5)

The diameter of a graph is the length / the costs of the longest shortest path

# Graphs Connected Components

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Connected components: G = (V, E)

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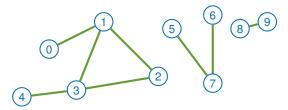


Figure: Three connected components

Undirected graph:

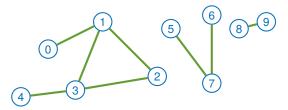


Figure: Three connected components

- Undirected graph:
  - All connected components are a partition of V

$$V = V_1 \cup \cdots \cup V_k$$

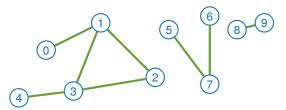


Figure: Three connected components

- Undirected graph:
  - All connected components are a partition of V

$$V = V_1 \cup \cdots \cup V_k$$

Two vertices u, v are in the same connected component if a path between u and v exists

# Graphs Connected Components



# Graphs Connected Components



Connected components: G = (V, E)

Directed graph:

- Directed graph:
  - Named strongly connected components



- Directed graph:
  - Named strongly connected components
  - Direction of edge has to be regarded



- Directed graph:
  - Named strongly connected components
  - Direction of edge has to be regarded
  - Not part of this lecture

Connected Components - Graph Exploration



**Graph Exploration:** (Informal definition)

Connected Components - Graph Exploration



**Graph Exploration:** (Informal definition)

■ Let G = (V, E) be a graph and  $s \in V$  a start vertex

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- We visit each reachable vertex connected to s

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- Breadth-first search: in sequence of the smallest distance to s

### **Graph Exploration:** (Informal definition)

- Let G = (V, E) be a graph and  $s \in V$  a start vertex
- We visit each reachable vertex connected to s
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- Depth-first search: in sequence of the largest distance to s
- Not a problem on its own but is often used as subroutine of other algorithms

Connected Components - Breadth-First Search



Connected Components - Breadth-First Search



#### Idea:

We start with all vertices unmarked and mark visited vertices

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- We start with all vertices unmarked and mark visited vertices
- Mark the start vertex s (level 0)
- Mark all unmarked connected vertices (level 1)
- Mark all unmarked vertices connected to a level 1-vertex (level 2)
- 5 Iteratively mark reachable vertices for all levels
- All connected nodes are now marked and in the same connected component as the start vertex s

Connected Components - Breadth-First Search



#### Connected Components - Breadth-First Search



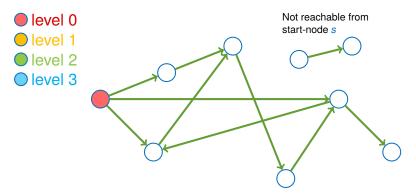


Figure: spanning tree of a breadth-first search

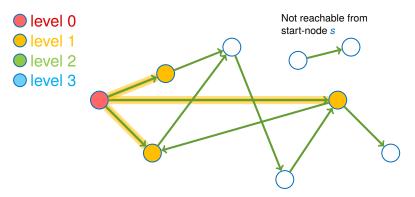


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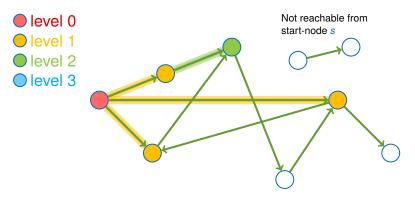


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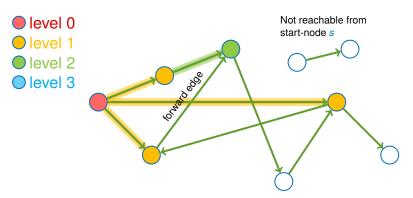


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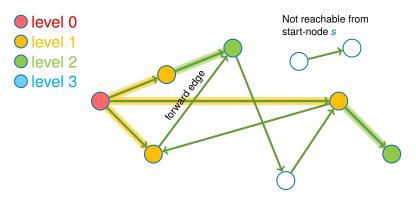


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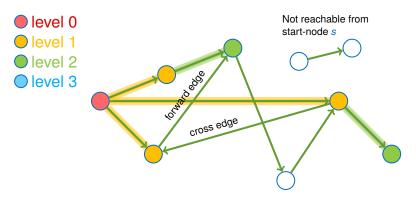


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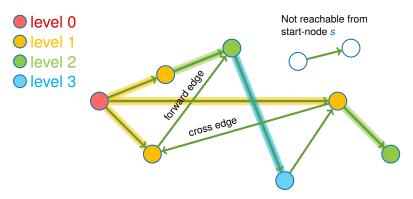


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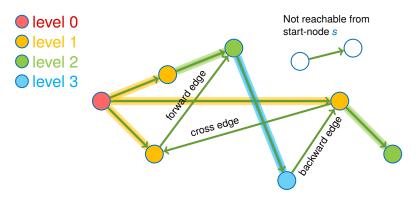


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Connected Components - Depth-First Search



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- Mark the start vertex s
- Pick an unmarked connected vertex and start a recursive depth-first search with the vertex as start vertex (continue on step 2)
- If no unmarked connected vertex exists go one vertex back (reduce the recursion level by one)

### Graphs

Connected Components - Depth-First Search



Search starts with long paths (searching with depth)

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- Marks like breadth-first search all connected vertices
- If the graph is acyclic we get a topological sorting
  - Each newly visited vertex gets marked by an increasing number
  - The numbers increase with path from the start vertex

# Graphs Connected Components - Depth-First Search



January 2017

### Graphs

#### Connected Components - Depth-First Search



- The marked vertices create a different spanning tree containing all reachable nodes
- start-node
- opath 1
- path 2
- opath 3

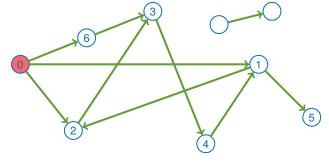


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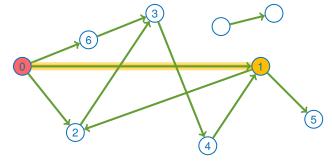


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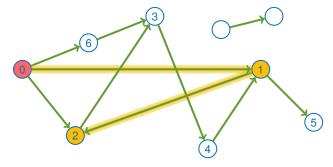


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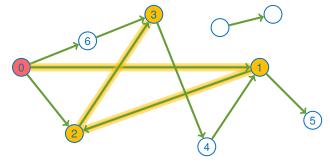


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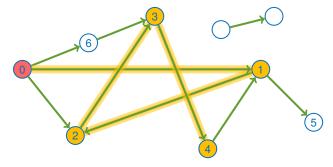


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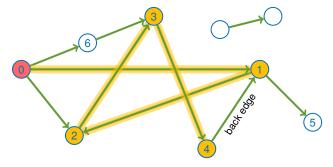


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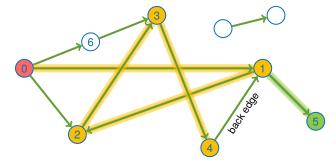


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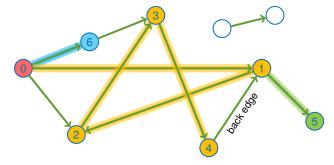


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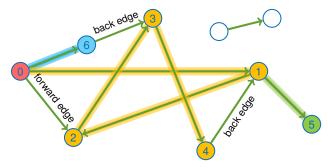


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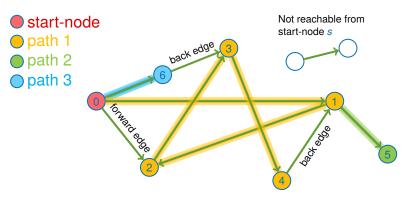


Figure: spanning tree of a depth-first search

### Graphs

Why is this called Breadth - and Depth First Search?



Constant costs for each visited vertex and edge

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- We get a runtime complexity of  $\Theta(|V'| + |E'|)$

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- Constant costs for each visited vertex and edge
- We get a runtime complexity of  $\Theta(|V'| + |E'|)$
- Let V' and E' be the reachable vertices and edges
- All vertices of V' are in the same connected component as our start vertex s
- This can only be improved by a constant factor

#### Structure



Feedback

Exercises

#### Graphs

Introduction Implementation

Application example

Image processing



Image processing



Connected component labeling

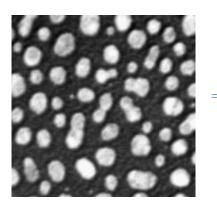
Image processing



- Connected component labeling
- Counting of objects in an image

Image processing

- Connected component labeling
- Counting of objects in an image



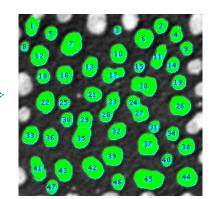
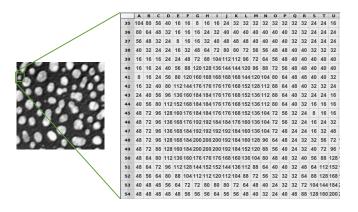


Image processing



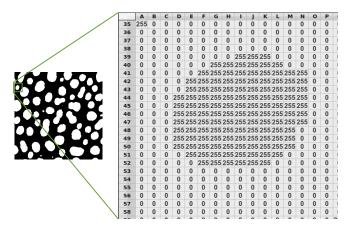
#### What's object, what's background?





#### Convert to black white using threshold:

value = 255 if value > 100 else 0



# Application example Image processing



Interpret image as graph:

# Application example Image processing



NE NE

#### Interpret image as graph:

■ Each white pixel is a node

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- Each white pixel is a node
- Edges between adjacent pixels (normally 4 or 8 neighbors)

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Image processing

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- Edges between adjacent pixels (normally 4 or 8 neighbors)
- Edges are not saved externally, algorithm works directly on array
- Breadth- / depth-first search find all connected components (particles)

Image processing



Image processing



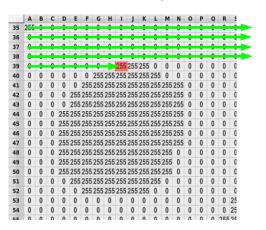
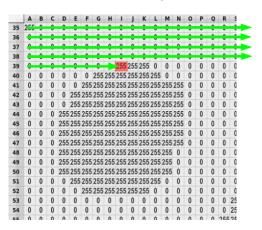


Image processing



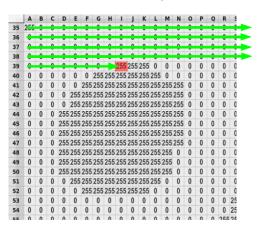
#### Find connected components:



Search pixel-by-pixel for non-zero intensity

Image processing

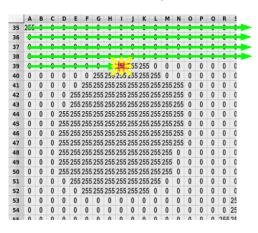




- Search pixel-by-pixel for non-zero intensity
- Label found pixel as component 1

Image processing

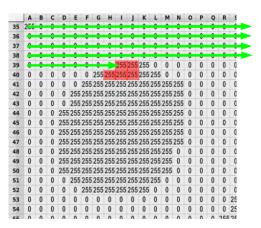




- Search pixel-by-pixel for non-zero intensity
- Label found pixel as component 1
- Check neighbors of all new labeled pixels

William Willia

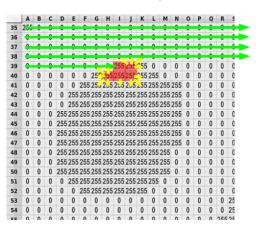
#### Image processing



- Search pixel-by-pixel for non-zero intensity
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Image processing

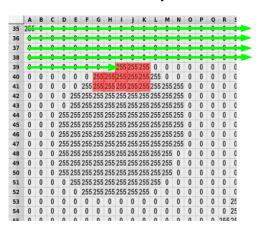




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Image processing

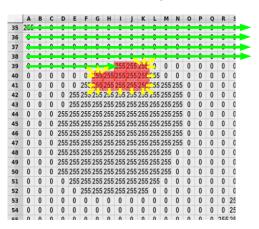




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Image processing

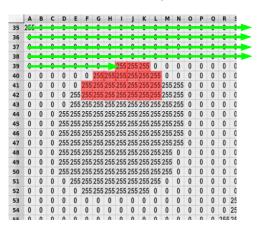




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Image processing

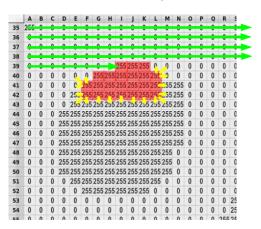




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Image processing

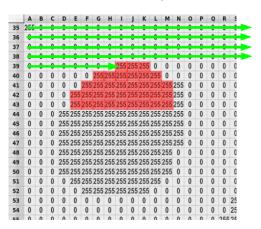




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Image processing

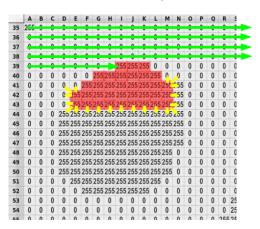




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Image processing

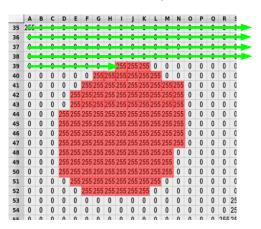




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Image processing

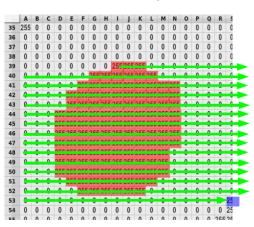




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Image processing





- Search pixel-by-pixel for non-zero intensity
- Label found pixel as component 2
- ...

#### Result of connected component labeling:

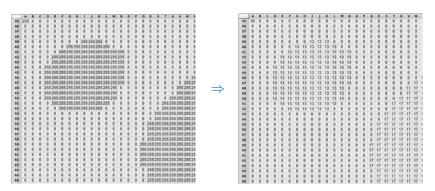


Figure: Result: particle indices instead of intensities

#### General

- [CRL01] Thomas H. Cormen, Ronald L. Rivest, and Charles E. Leiserson. Introduction to Algorithms. MIT Press, Cambridge, Mass, 2001.
- Kurt Mehlhorn and Peter Sanders. [MS08] Algorithms and data structures, 2008. https://people.mpi-inf.mpg.de/~mehlhorn/

ftp/Mehlhorn-Sanders-Toolbox.pdf.

#### ■ Graph-Search

#### Graph-Connectivity

```
[Wik] Connectivity (graph theory)
    https://en.wikipedia.org/wiki/Connectivity_
    (graph_theory)
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