Algorithms and Datastructures

Cache Efficiency, Divide and Conquer

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Structure

Background:

- Up to now we always counted the number of operations
- Assuming this is a good measure for the runtime of a algorithm/tool
- Today we will see examples where this is not suitable

Example:

- We sum up all elements of a field a of size n in ...
  - natural order:

$$\mathrm{sum}(a) = a[1] + a[2] + \cdots + a[n]$$

  - random order:

$$\mathrm{sum}(a) = a[21] + a[5] + \cdots + a[8]$$

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Python:

```python
def init(size):
    """Creates the dataset."""

    # use system time as seed
    random.seed(None)

    # set linear order as accessor
    order = [a for a in range(0, size)]

    # init array with random data
    data = [random.random() for a in order]

    return (order, data)
```

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Python:

```python
def run(param):
    """Processes the dataset."""

    # unpack data
    (order, data) = param

    # init the sum value
    s = 0

    for index in order:
        s += data[index]

    return s
```

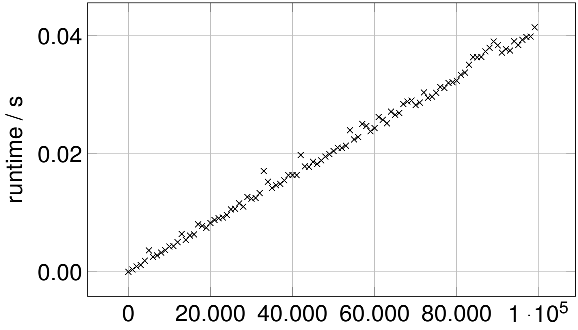

Figure: Summing elements in linear order

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Random Order - Python

```
def init(size):
    """Creates a randomly ordered dataset."""

    # use system time as seed
    random.seed(None)

    # set random order as accessor
    order = [a for a in range(0, size)]
    random.shuffle(order)

    # init array with random data
    data = [random.random() for a in order]

    return (order, data)
```

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Random Order

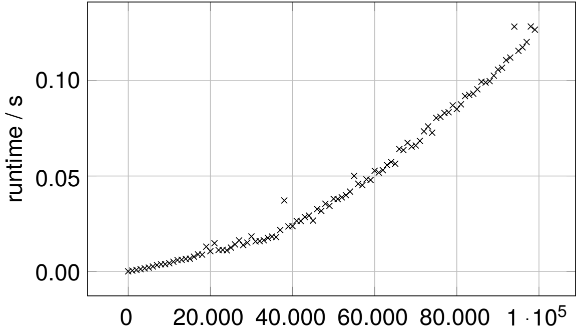
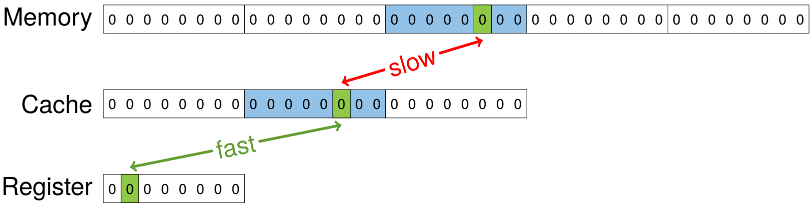


Figure: Summing elements in random order

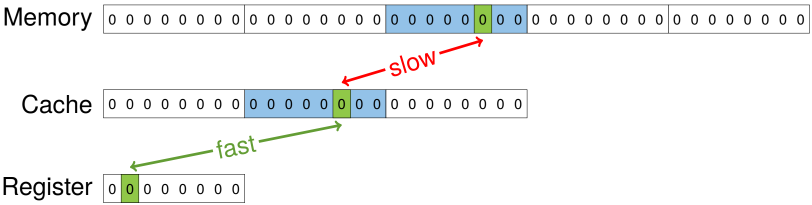
Conclusion:

- The number of operations are identical for both algorithms
- Accessing elements in random order takes a lot longer (Factor 10)
- The costs in terms of memory access are very different



Principle / organization:

- Accessing one byte of the main memory takes $\approx$ 100 ns
- Accessing one byte of (L1-)cache takes $\approx$ 1 ns
- Accessing one or more byte/s of main memory loads a whole block $\approx$ 100 B into the cache
- As long as this block is in the cache, it is not neccessary to access the memory for bytes of this block

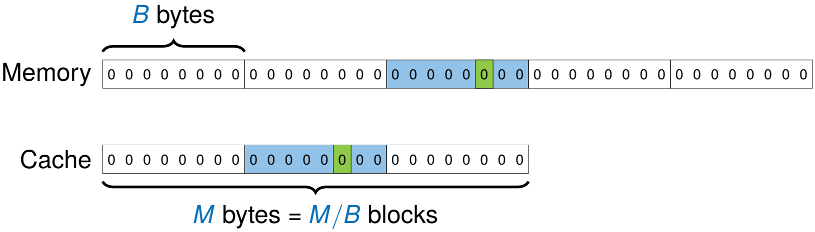


Cache organization:

- The (L1-)cache can hold multiple memory blocks
  - Cache lines ≈ 100 kB
- If the capacity is reached unused blocks are discarded
  - Least recently used (LRU)
  - Least frequently used (LFU)
  - First in first out (FIFO)
- Details of discarding are not the topic for today

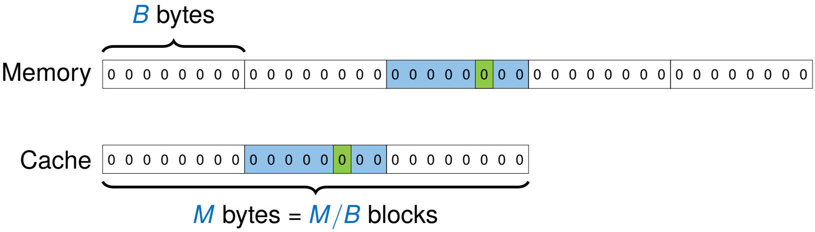
Cache Efficiency

Block Operations



Terminology:

- The system consists of slow and fast memory
- The slow memory is divided in blocks of size B
- The fast cache has size M an can store M/B blocks
- If data is not in fast memory, the corresponding block is loaded into the cache



Terminology:

- The program defines which blocks are held in the cache
- We use the number of block operations as runtime estimation
- We ignore runtime costs of cache accesses / management

Figure: Comparison good / bad locality

Accessing the cache B times:

- Best case: 1 block operation $\rightarrow$ good locality
- Worst case: B block operations $\rightarrow$ bad locality

Additional factors:

- The following settings change only a small constant factor in number of block operations
  - Partionining of the slow memory into blocks
  - Regardless of the block size: 1 Bytes or 4 Bytes or 8 Bytes

Note:

- If the input size is smaller than M we load the complete data chunk directly into the cache
- Cache handling is only interesting when the input size is greater than M

Typical values: (`Intel© i7-4770` Haswell, `WD© Blue` 2 TB)

- CPU L1 Cache: $B = 64\,\mathrm{B}$, $M = 4 \times (32\,\mathrm{kB} + 32\,\mathrm{kB})$
- CPU L2 Cache: $B = 64\,\mathrm{B}$, $M = 4 \times 256\,\mathrm{kB}$
- CPU L3 Cache: $B = 64\,\mathrm{B}$, $M = 8\,\mathrm{MB}$
- Disk Cache: $B = 64\,\mathrm{kB}$, $M = 64\,\mathrm{MB}$
  - Many operating systems use free system memory as disk cache

Terminology:

- Block loads on CPU-cache are called cache misses
- Block operations on disk-cache are called IOs (input / output operations)
- These also fall under the term cache efficiency or IO efficiency

Cache Efficiency

Block Operations - Linear Order

Example 1 - Linear order:

- We sum up all elements in natural order

$$\mathrm{sum}(a) = a[1] + a[2] + \cdots + a[n]$$

- The number of block operations is $\mathrm{ceil}\left(\frac{n}{B}\right)$

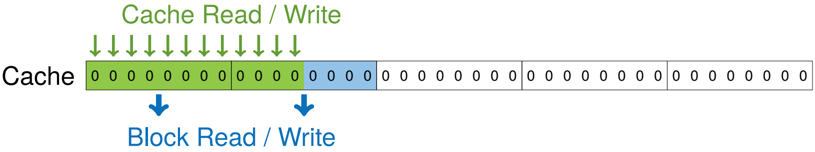


Figure: Good locality of sum operation

Example 2 - Random order:

- We sum up all elements in random order

$$\text{sum}(a) = a[21] + a[5] + \cdots + a[8]$$

- The number of block operations is n in the worst case
- This leads to a runtime factor difference of B

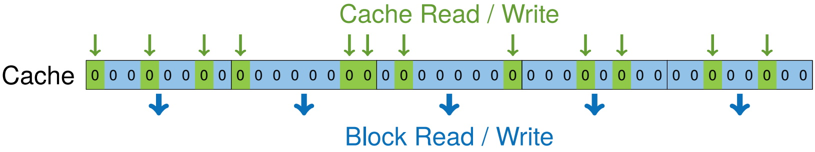


Figure: Bad locality of sum operation

Generally the factor is substantially $< B$

- Even with a random order we access 4 neighboring bytes at once per `int` (int32_t)
- The next element might already be loaded in the cache
- If not $n \gg M$ this might occur with a high probability

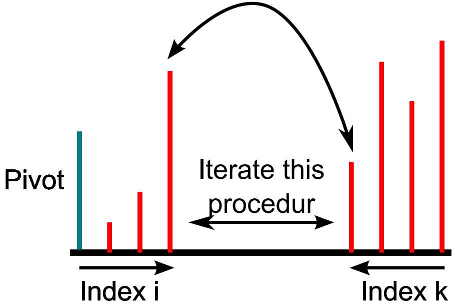
QuickSort:

- **Strategy:** Divide and conquer
- Divide the data into two parts where the “left” part contains all values $\leq$ those in the right part
- Choose one element (e.g the first one) as “pivot”-element
- Ideally both parts are the same size
- Both parts are sorted recursively

Figure: *QuickSort* with pivot-element

Idea of Quicksort

- **At start:** Pivot in first position, first re-arrange list such that left part contains small, right part larger elements
- Do required changes *in place*



- **End point:** k is left to left-most element greater than pivot

  swap position 0 (pivot) with k (smaller than pivot)

Python:

```
def quicksort(l, start, end):
  if (end - start) < 1:
    return

  i = start
  k = end
  piv = l[0]

  ...
```

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Block Operations - QuickSort - Python

```
def quicksort(l, start, end):
    ...

    while k > i:
        while l[i] <= piv and i <= end and k > i:
            i += 1
        while l[k] > piv and k >= start and k >= i:
            k -= 1

        if k > i: # swap elements
            (l[i], l[k]) = (l[k], l[i])

    (l[start], l[k]) = (l[k], l[start])
    quicksort(l, start, k - 1)
    quicksort(l, k + 1, end)
```

Number of operations for Quicksort:

- Let $T(n)$ be the runtime for the input size n

Assumptions:

- Fields are always separated perfectly in the middle
- n is a power of two and recursion depth is $k = \log_2 n$

$$
\begin{aligned}
T(n) &\leq \underbrace{A \cdot n}_{\text{splitting in two parts}} + \underbrace{2 \cdot T\left(\frac{n}{2}\right)}_{\text{recursive sort}} \\
&\leq A \cdot n + 2\left(A \cdot \frac{n}{2} + 2 \cdot T\left(\frac{n}{4}\right)\right) \\
&= 2A \cdot n + 4 \cdot T\left(\frac{n}{4}\right) \\
&\leq 3A \cdot n + 8 \cdot T\left(\frac{n}{8}\right) \\
&\leq \quad \cdots \\
&\leq k \cdot A \cdot n + 2^k \cdot T(1) \\
&= \log_2 n \cdot A \cdot n + n \cdot T(1) \\
&\leq \log_2 n \cdot A \cdot n + n \cdot A \in \mathcal{O}(n \log_2 n)
\end{aligned}
$$

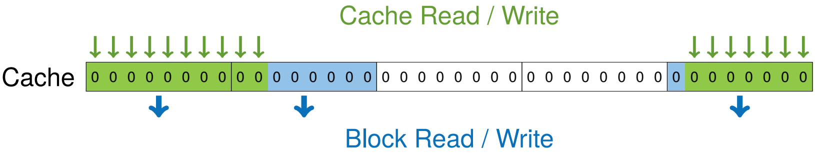

Figure: Locality of quicksort

- Let $IO(n)$ be the number of block operations for input size n
- Assumptions as before but recursion depth is $k = \log_2 \frac{n}{B}$

$$
\begin{aligned}
IO(n) &\leq \underbrace{A \cdot n/B}_{\text{splitting in two parts}} + \underbrace{2 \cdot IO(n/2)}_{\text{recursive sort}} \\
&\leq A \cdot n/B + 2(A \cdot n/2B + 2 \cdot IO(n/4)) \\
&\leq 2 \cdot A \cdot n/B + 4 \cdot IO(n/4) \\
&\leq 3 \cdot A \cdot n/B + 8 \cdot IO(n/8) \\
&\leq \cdots \\
&\leq k \cdot A \cdot n/B + 2^k \cdot IO(n/2^k) \\
&= \log_2(n/B) \cdot A \cdot (n/B) + n/B \cdot IO(B) \\
&\leq \log_2(n/B) \cdot A \cdot (n/B) + A \cdot n/B \in \mathcal{O}\left(\frac{n}{B} \cdot \log_2\left(\frac{n}{B}\right)\right)
\end{aligned}
$$

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Concept:

- Divide the problem into smaller subproblems
- Conquer the subproblems through recursive solving.
  If subproblems are small enough solve them directly
- Connect all solutions of the subproblems to a solution of the full problem
- Recursive application of the algorithm to ever smaller subproblems
- Direct solving of sufficently small subproblems

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- Function solve for solving a problem of size *n*

```
def solve(problem):
    if n < threshold:
        return solution # solve directly
    else:
        # divide problem into subproblems
        # P1, P2, ..., Pk with k>=2
        S1 = solve(P1)
        S2 = solve(P2)
        ...
        Sk = solve(Pk)

        # combine solutions
        return S1 + S2 + ... + Sk
```

Divide and Conquer:

- Can help with conceptual hard problems
- Solution of the trivial problems has to be known
- Dividing in subproblems has to be possible
- Combination of solutions has to be possible

Features:

- Realization of efficient solutions
  - If trivial solution is $\in O(1)$
  - And separation / combination of subproblems is $\in O(n)$
  - And the number of subproblems is limited
  - The runtime is $\in O(n \cdot \log n)$
- Suitable for parallel processing
  - Subproblems are independent of each other
  - Only needed input for each subproblem has to be known

Definition of the trivial case:

- Smaller subproblems are elegant and simple
- Otherwise the efficiency will be improved if relative big subproblems can be solved directly
- Recursion depth should not get too big (stack / memory overhead)

Division in subproblems:

- Choosing the number of subproblems and the concrete allocation can be demanding

Combination of solutions:

- Typically conceptional demanding

Example - Maximum Subtotal Input:

- Sequence X of n integers

Output:

- Maximum sum of related subsequence and its index boundary

Index	0	1	2	3	4	5	6	7	8	9
Value	31	-41	59	26	-53	58	97	-93	-23	84

Output: Sum: 187, Start: 2, End: 6

Application:

- Maximum profit of buying and selling shares

Figure: Stock value over time

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Naive solution (brute force)

```python
def maxSubArray(X):
    # Store sum, start, end
    result = (X[0], 0, 0)
    for i in range(0, len(X)):
        for j in range(i, len(X)):
            subSum = 0
            for k in range(i, j + 1):
                subSum += X[k]
            if result[0] < subSum:
                result = (subSum, i, j)
    return result
```

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Runtime - Upper bound

```
def maxSubArray(X):
    result = (X[0], 0, 0)
    # n loops -> O(n)
    for i in range(0, len(X)):
        # max n loops -> O(n)
        for j in range(i, len(X)):
            # max n loops -> O(n)
            subSum = sum(X[i:j+1])
            if result[0] < subSum: # O(1)
                result = (subSum, i, j)
    return result
```

Upper bound:

- Three interleaved loops
- Each loop with runtime $O(n)$
- Algorithm runtime of $O(n^3)$

Lower bound:

Table: Operations

i	Additions	j
$\frac{n}{3} \in O(n)$	$\frac{n}{3} \in O(n)$	$\frac{n}{3} \in O(n)$

- We iterate at least $\frac{n}{3}$ values for i
- For each i we iterate at least $\frac{n}{3}$ values for j
- For each j we have at least $\frac{n}{3}$ additions
- We need at least $T(n) = (\frac{n}{3})^3 \in \Omega(n^3)$ steps

Runtime:

- With $T(n) \in O(n^3)$ and $T(n) \in \Omega(n^3)$ we know:

$$T(n) \in \Theta(n^3)$$

- It is hard to solve the problem in a worse way ...

Current approach:

- Calculating the sum for range from i to j with loop

$$S_{i,j} = X[i] + X[i+1] + \dots + X[j]$$

Better approach:

- Incremental sum instead of loop

$$\begin{aligned} S_{i,j+1} &= X[i] + X[i+1] + \dots + X[j] + X[j+1] \\ S_{i,j+1} &= S_{i,j} + X[j+1] \in O(1) \quad \text{instead of} \quad \in O(n) \end{aligned}$$

Better solution:

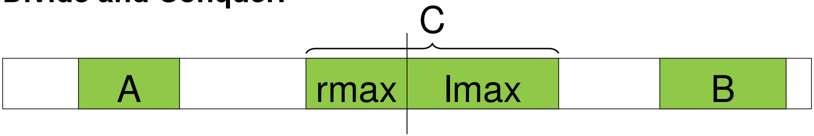
```
def maxSubArray(X):
    result = (X[0], 0, 0)
    # n loops -> O(n)
    for i in range(0, len(X)):
        subSum = 0
        # max n loops -> O(n)
        for j in range(i, len(X)):
            subSum += X[j] # O(1)
            if result[0] < subSum: # O(1)
                result = (subSum, i, j)
    return result
```

- Runtime $\in O(n^2)$

Divide and Conquer

Example - Maximum Subtotal

Divide and Conquer:



Divide and Conquer Idea to solve:

- Split the sequence in the middle
- Solve left half of the problem
- Solve right half and combine both solutions into one
- Maximum might be located in left half (A) or right half (B)
- Problem: Maximum can overlap the split
- To solve this case we have to calculate rmax and lmax
- The overall solution is the maximum of A, B and C

Principle - Divide and Conquer:

- Small problems are solved directly: $n = 1 \Rightarrow \text{max} = X[0]$
- Bigger problems are partitioned into two subproblems and recursively solved. Subsolutions A and B are returned
- To determine subsolution C, rmax and lmax for the subproblems are computed
- The overall solution is the maximum of A, B and C

```python
def maxSubArray(X, i, j):
    if i == j: # trivial case
        return (X[i], i, i)

    # recursive subsolutions for A, B
    m = (i + j) / 2
    A = maxSubArray(X, i, m)
    B = maxSubArray(X, m + 1, j)

    # rmax and lmax for cornercase C
    C1, C2 = rmax(X, i, m), lmax(X, m + 1, j)
    C = (C1[0] + C2[0], C1[1], C2[1])

    # compute solution from results A, B, C
    return max([A, B, C], key=lambda i: i[0])
```

Further Literature

- **General**

[CRL01] Thomas H. Cormen, Ronald L. Rivest, and Charles E. Leiserson.
Introduction to Algorithms.
MIT Press, Cambridge, Mass, 2001.

[MS08] Kurt Mehlhorn and Peter Sanders.
Algorithms and data structures, 2008.
https://people.mpi-inf.mpg.de/~mehlhorn/ftp/Mehlhorn-Sanders-Toolbox.pdf.

Further Literature

- **Caching**

[Wik] Cache
https://en.wikipedia.org/wiki/Cache