

Algorithms and Datastructures

Runtime analysis Minsort / Heapsort, Induction

Albert-Ludwigs-Universität Freiburg



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Algorithms and Datastructures, October 2017

Runtime Example

Minsort

Basic Operations

Runtime analysis

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Heapsort

Introduction to Induction

Logarithms

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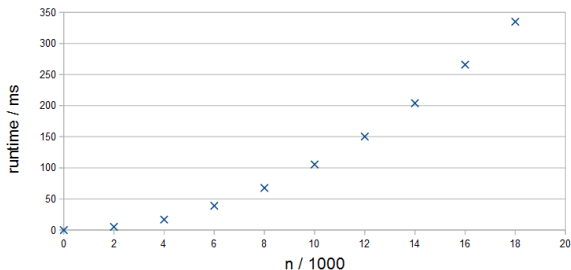
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- How can we say more precisely what is happening?

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 - Which kind of computer is the code executed on
 - What is running in the background
 - Which compiler is used to compile the code

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- **Problem:** The runtime is also depending on many other influences, especially:
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 - Which compiler is used to compile the code
- **Abstraction 1:** Analyze the number of basic operations, rather than analyzing the runtime

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Incomplete list of basic operations:

- Arithmetic operation, for example: $a + b$
- Assignment of variables, for example: $x = y$
- Function call, for example: *minsort(lst)*

Intuitive:

lines of code

Better:

lines of machine
code

Best:

process cycles

Important:

The actual runtime has to be roughly proportional to the number of operations.

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- **Abstraction 2:** We calculate the upper (lower) bound, rather than counting the operations exactly

Reason: Runtime is approximated by number of basic operations, but we can still infer:

- Upper bound
- Lower bound

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Reason: Runtime is approximated by number of basic operations, but we can still infer:

- Upper bound
 - Lower bound
- **Basic Assumption:**
 - n is size of the input data (i.e. array)
 - $T(n)$ number of operations for input n

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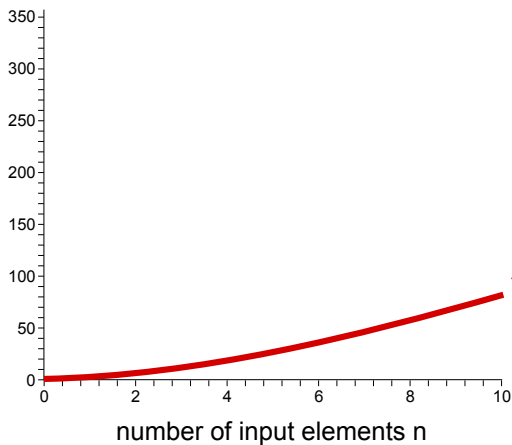
- **Observation:** The number of operations depends only on the size n of the array and not on the content!
- **Claim:** There are constants C_1 and C_2 such that:

$$C_1 \cdot n^2 \leq T(n) \leq C_2 \cdot n^2$$

- This is called “quadratic runtime” (due to n^2)

Runtime Example

Number of Steps
(runtime [s])

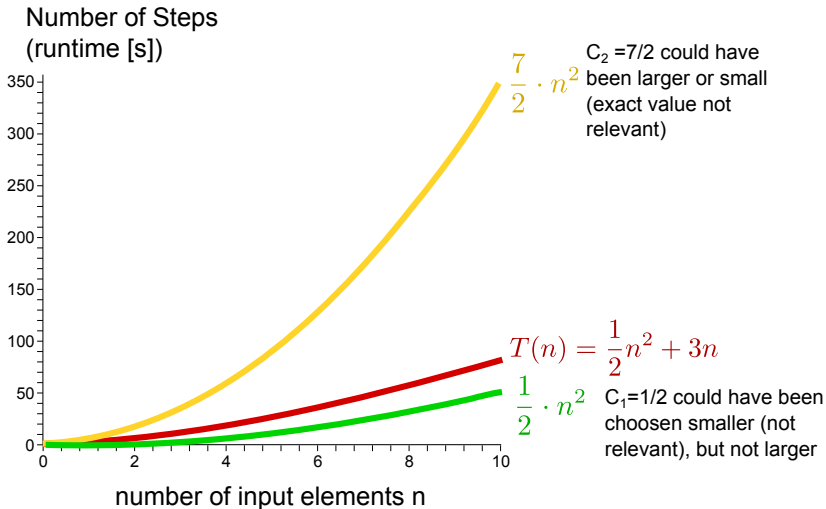


$$T(n) = \frac{1}{2}n^2 + 3n$$

Runtime Example



Runtime Example



We declare:

- Runtime of operations: $T(n)$
- Number of Elements: n
- Constants: C_1 (lower bound), C_2 (upper bound)
$$C_1 \cdot n^2 \leq T(n) \leq C_2 \cdot n^2$$
- Number of operations in round i : T_i



Figure: *Minsort* at the iteration $i = 4$. We have to check $n - 3$ elements

Compares at each
iteration:



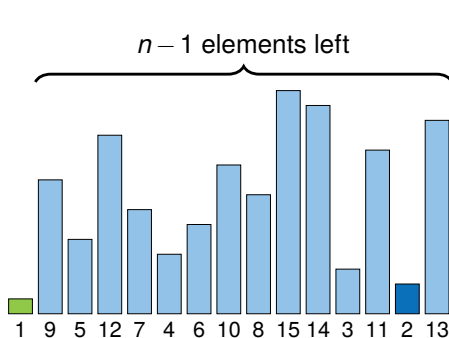
Figure: *Minsort* with start data

Compares at each iteration:

$$T_1 \leq C'_2 \cdot (n - 0)$$



Figure: *Minsort* at iteration $i = 1$



Compares at each iteration:

$$T_1 \leq C'_2 \cdot (n - 0)$$

$$T_2 \leq C'_2 \cdot (n - 1)$$

Figure: *Minsort* at iteration $i = 2$



Compares at each iteration:

$$T_1 \leq C'_2 \cdot (n - 0)$$

$$T_2 \leq C'_2 \cdot (n - 1)$$

$$T_3 \leq C'_2 \cdot (n - 2)$$

Figure: *Minsort* at iteration $i = 3$



Compares at each iteration:

$$T_1 \leq C'_2 \cdot (n - 0)$$

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$$T_3 \leq C'_2 \cdot (n - 2)$$

$$T_4 \leq C'_2 \cdot (n - 3)$$

Figure: Minsort at iteration $i = 4$



Figure: Minsort at iteration i

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$$T_{n-1} \leq C'_2 \cdot 2$$

$$T_n \leq C'_2 \cdot 1$$

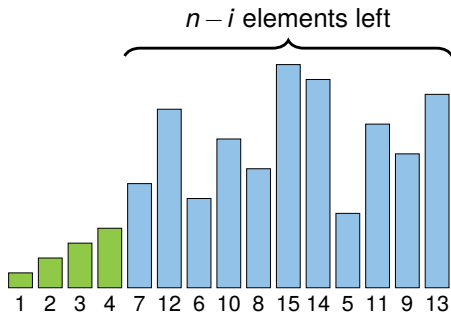


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$$T_n \leq C'_2 \cdot 1$$

$$T(n) = C'_2 \cdot (T_1 + \dots + T_n) \leq \sum_{i=1}^n (C'_2 \cdot i)$$

Alternative: Analyse the Code:

```
def minsort(elements):  
    for i in range(0, len(elements)-1):  
        minimum = i  
  
        for j in range(i+1, len(elements)):  
            if elements[j] < elements[minimum]:  
                minimum = j  
  
        if minimum != i:  
            elements[i], elements[minimum] = \  
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        } const.  
        } runtime  
  
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Diagram illustrating the runtime analysis of the Minsort algorithm:

- The inner loop (for j in range(i+1, len(elements))) is labeled "const. runtime".
- The inner loop is repeated $n-i-1$ times for each iteration of the outer loop.
- The outer loop (for i in range(0, len(elements)-1)) is repeated $n-1$ times.

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Remark: C'_2 is cost of comparison \Rightarrow assumed constant

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Excursion - Small Gauss Formula

Proof of lower bound: $C_1 \cdot n^2 \leq T(n)$

Like for the upper boundary there exists a C_1 . Summation analysis is the same

$$T(n) \geq \sum_{i=1}^{n-1} C'_1 \cdot (n-i)$$

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How do we get to n^2 ?

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Summarized:

$$\frac{C'_1}{4} \cdot n^2 \leq T(n) \leq C'_2 \cdot n^2$$

Quadratic runtime proven:

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- **Quadratic runtime = “big” problems unsolvable**

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Formal:

- Let $T(n)$ be the runtime for the *Heapsort* algorithm with n elements
- On the next pages we will proof $T(n) \leq C \cdot n \log_2 n$

Depth of a binary tree:

- **Depth d :** longest path through the tree
- Complete binary tree has $n = 2^d - 1$ nodes
- Example: $d = 4$
 $\Rightarrow n = 2^4 - 1 = 15$

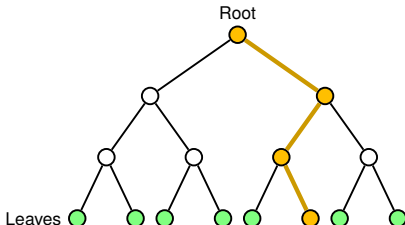


Figure: Binary tree with 15 nodes

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- If both has been proven, then $A(n)$ holds for all natural numbers n by **induction**

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Figure: Tree of depth 1 has 1 node

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A **complete** binary tree of depth d has $n(d) = 2^d - 1$ nodes

- **Induction basis:** Assumption holds for $d = 1$

Root



$$n(1) = 2^1 - 1 = 1$$

\Rightarrow correct ✓

Figure: Tree of depth 1 has 1 node

Induction - Example 1



Number of nodes $n(d)$ in a binary tree with depth d :

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Induction - Example 1



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- **Induction basis:** $n(1) = 2^1 - 1 = 2^1 - 1 = 1$ ✓

Induction - Example 1

Number of nodes $n(d)$ in a binary tree with depth d :

- **Induction assumption:** $n(d) = 2^d - 1$
- **Induction basis:** $n(1) = 2^1 - 1 = 2^0 - 1 = 1$ ✓
- **Induction step:** to show for $d := d + 1$



Figure: Binary tree with subtrees

Induction - Example 1

Number of nodes $n(d)$ in a binary tree with depth d :

- **Induction assumption:** $n(d) = 2^d - 1$
- **Induction basis:** $n(1) = 2^1 - 1 = 2^0 - 1 = 1$ ✓
- **Induction step:** to show for $d := d + 1$



$$n(d+1) = 2 \cdot n(d) + 1$$

Figure: Binary tree with subtrees

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$$\begin{aligned} n(d+1) &= 2 \cdot n(d) + 1 \\ &= 2 \cdot (2^d - 1) + 1 \end{aligned}$$

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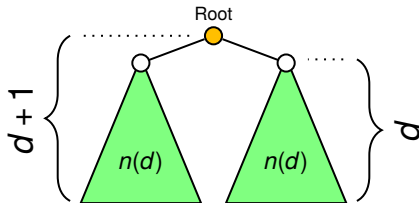


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⇒ **By induction:** $n(d) = 2^d - 1 \quad \forall n \in \mathbb{N} \quad \square$

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Heapsort has the following steps:

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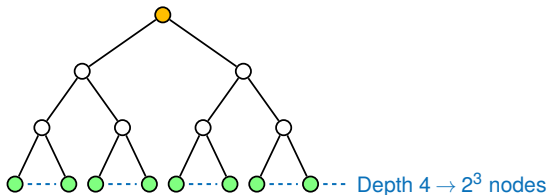
Heapsort has the following steps:

- **Initially:** heapify list of n elements
- **Then:** until all n elements are sorted
 - Remove root as minimal element
 - Move last leaf to root position
 - Repair heap by sifting

Runtime of heapify depends on depth d :



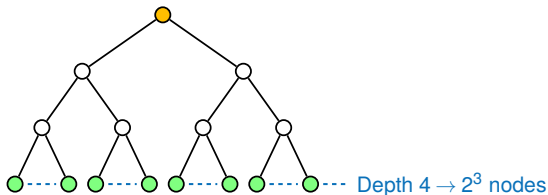
Runtime of heapify depends on depth d :



Runtime of heapify with depth of d :

- No costs at depth d with 2^{d-1} (or less) nodes

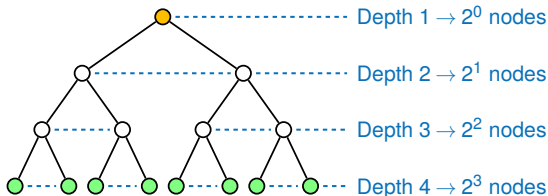
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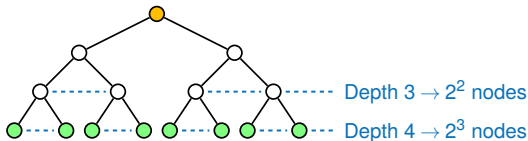
- No costs at depth d with 2^{d-1} (or less) nodes
- The cost for sifting with depth 1 is at most $1C$ per node
- In general: Sifting costs are linear with path length **and** number of nodes

Heapify total runtime:



| Depth | Nodes | Path length | Costs per node |
|-------|-----------|-------------|------------------|
| d | 2^{d-1} | 0 | $\leq C \cdot 0$ |

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In total:
$$T(d) \leq \sum_{i=1}^d (C \cdot (i-1) \cdot 2^{d-i})$$

Heapify total runtime:



| Depth | Nodes | Path length | Costs per node | Upper bound |
|-------|-----------|-------------|------------------|----------------------|
| d | 2^{d-1} | 0 | $\leq C \cdot 0$ | Standard Equation |
| $d-1$ | 2^{d-2} | 1 | $\leq C \cdot 1$ | |
| $d-2$ | 2^{d-3} | 2 | $\leq C \cdot 2$ | |
| $d-3$ | 2^{d-4} | 3 | $\leq C \cdot 3$ | |

In total:
$$T(d) \leq \sum_{i=1}^d \left(C \cdot (i-1) \cdot 2^{d-i} \right) \leq \sum_{i=1}^d \left(C \cdot i \cdot 2^{d-i} \right)$$

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| $d-1$ | 2^{d-2} | 1 | $\leq C \cdot 1$ | $\leq C \cdot 2$ |
| $d-2$ | 2^{d-3} | 2 | $\leq C \cdot 2$ | $\leq C \cdot 3$ |
| $d-3$ | 2^{d-4} | 3 | $\leq C \cdot 3$ | $\leq C \cdot 4$ |

In total:
$$T(d) \leq \sum_{i=1}^d \left(C \cdot (i-1) \cdot 2^{d-i} \right) \leq \sum_{i=1}^d \left(C \cdot i \cdot 2^{d-i} \right)$$

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- **Hence:** Resulting costs for heapify:

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- **However:** We want costs in relation to n



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Figure: Partial binary tree

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- $2^{d-1} - 1$ nodes in full tree till layer $d - 1$



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- A binary tree of depth d has $2^{d+1} - 1$ nodes **Why?**
- $2^{d-1} - 1$ nodes in full tree till layer $d - 1$
- At least 1 node in layer d



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- Equation multiplied by 2^2
 $\Rightarrow 2^{d-1} \cdot 2^2 \leq 2^2 \cdot n$

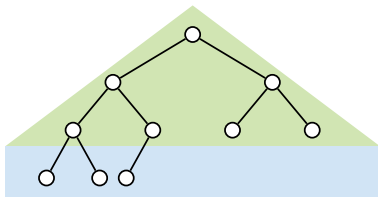


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- At least 1 node in layer d
- Equation multiplied by 2^2
 $\Rightarrow 2^{d-1} \cdot 2^2 \leq 2^2 \cdot n$
- Cost for heapify:
 $\Rightarrow T(n) \leq C \cdot 4 \cdot n$



Figure: Partial binary tree

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Introduction to Induction

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- We want to proof (induction assumption):

$$\underbrace{\sum_{i=1}^d (i \cdot 2^{d-i})}_{A(d)} \leq \underbrace{2^{d+1}}_{B(d)}$$

- We denote the left side with A , the right side with B

- **Induction basis:** $d := 1$:

$$A(d) \leq B(d)$$

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■ **Induction basis:** $d := 1$:

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$$\sum_{i=1}^1 (i \cdot 2^{1-i}) \leq 2^{1+1}$$

$$2^0 \leq 2^2 \quad \checkmark$$

Induction step: ($d := d + 1$):

- **Idea:** Write down right hand formula and try to get $A(d)$ and $B(d)$ out of it

$$A(d) \leq B(d) \quad \Rightarrow \quad A(d+1) \leq B(d+1)$$

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$$A(d) \leq B(d) \quad \Rightarrow \quad A(d+1) \leq B(d+1)$$

$$\sum_{i=1}^{d+1} (i \cdot 2^{d+1-i}) \leq 2^{d+1+1}$$

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$$\sum_{i=1}^{d+1} (i \cdot 2^{d+1-i}) \leq 2^{d+1+1}$$

$$2 \cdot \sum_{i=1}^{d+1} (i \cdot 2^{d-i}) \leq 2 \cdot 2^{d+1}$$

\vdots

Induction step: ($d := d + 1$):

\vdots

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$$2 \cdot A(d) + (d+1) \leq 2 \cdot B(d)$$

■ **Problem:** Does not work but claim still holds

Working proof:

- Show a **little bit stronger** claim

$$\sum_{i=1}^d (i \cdot 2^{d-i}) \leq 2^{d+1} - d - 2 \leq 2^{d+1}$$

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- **Advantage:** Results in a stronger induction assumption
 \Rightarrow **exercise**

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- **Recall:** The depth and number of elements is decreasing

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- **Recall:** The depth and number of elements is decreasing
 - **Hence:** $T(n) \leq n \cdot (1 + \log_2 n) \cdot C$

Runtime of the other operations:

- Constant costs for taking out $n \times$ maximum
- Maximum of d steps repairing the heap n times
- Depth of heap at the start is $d \leq 1 + \log_2 n$ Why?

$$2^{d-1} \leq n \Rightarrow d-1 \leq \log_2 n \Rightarrow d \leq 1 + \log_2 n$$

- **Recall:** The depth and number of elements is decreasing
 - **Hence:** $T(n) \leq n \cdot (1 + \log_2 n) \cdot C$
 - We can reduce this to:

$$T(n) \leq 2 \cdot n \log_2 n \cdot C \quad (\text{holds for } n > 2)$$

Runtime costs:

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Runtime costs:

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- Remove: $T(n) \leq 2 \cdot n \log_2 n \cdot C$
- Total runtime: $T(n) \leq 6 \cdot n \log_2 n \cdot C$
- Constraints:
 - Upper bound: $C_2 \cdot n \log_2 n \geq T(n)$ (for $n \geq 2$)
 - Lower bound: $C_1 \cdot n \log_2 n \leq T(n)$ (for $n \geq 2$)

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 - $\Rightarrow C_1$ and C_2 are constant

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Logarithm to different bases:

$$\log_a n = \frac{\log_b n}{\log_b a} = \log_b n \cdot \frac{1}{\log_b a}$$

The only difference is a constant coefficient $\frac{1}{\log_b a}$

Examples:

- $\log_2 4 = \log_{10} 4 \cdot \frac{1}{\log_2 10} = 0.602 \dots \cdot 3.322 \dots = 2 \checkmark$
- $\log_{10} 1000 = \log_e 1000 \cdot \frac{1}{\log_e 10} = \ln 1000 \cdot \frac{1}{\ln 10} = 3 \checkmark$

Runtime of $n \log_2 n$:

- Assume we have constants C_1 and C_2 with

$$C_1 \cdot n \cdot \log_2 n \leq T(n) \leq C_2 \cdot n \cdot \log_2 n \quad \text{for } n \geq 2$$

Runtime of $n \log_2 n$:

- Assume we have constants C_1 and C_2 with

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Runtime of $n \log_2 n$:

- Assume we have constants C_1 and C_2 with

$$C_1 \cdot n \cdot \log_2 n \leq T(n) \leq C_2 \cdot n \cdot \log_2 n \quad \text{for } n \geq 2$$

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 - $C \cdot n \cdot \log_2 n = 10^{-9} \text{ s} \cdot 2^{30} \cdot 30 = 32 \text{ s}$
- **Runtime $n \log_2 n$ is nearly as good as linear!**

■ General for this Lecture

[CRL01] Thomas H. Cormen, Ronald L. Rivest, and Charles E. Leiserson.

Introduction to Algorithms.

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[MS08] Kurt Mehlhorn and Peter Sanders.

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<https://people.mpi-inf.mpg.de/~mehlhorn/ftp/Mehlhorn-Sanders-Toolbox.pdf>.

■ Mathematical Induction

[Wik] [Mathematical induction](https://en.wikipedia.org/wiki/Mathematical_induction)

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