Algorithms and Datastructures

Balanced Trees (AVL-Trees, (a,b)-Trees, Red-Black-Trees)

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Bioinformatics Group / Department of Computer Science

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Structure

Binary search tree:

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- Best case: $d \in O(\log n)$, keys are inserted randomly
- Worst case: $d \in O(n)$, keys are inserted in ascending / descending order $(20, 19, 18, \ldots)$

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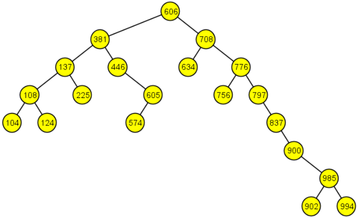

Figure: Binary search tree with random insert [Gna]

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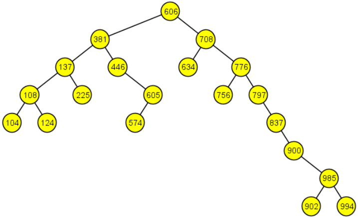

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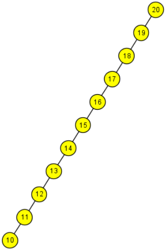

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- We rebalance the tree from time to time

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  - Used in C++ `std::map` and Java `SortedMap`

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- With that the height of the search tree is always $O(\log n)$
- We can perform all basic operations in $O(\log n)$

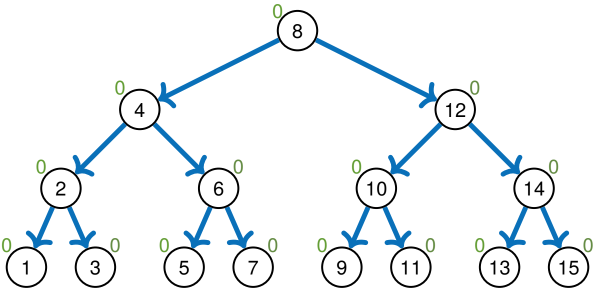


Figure: Example of an AVL-Tree

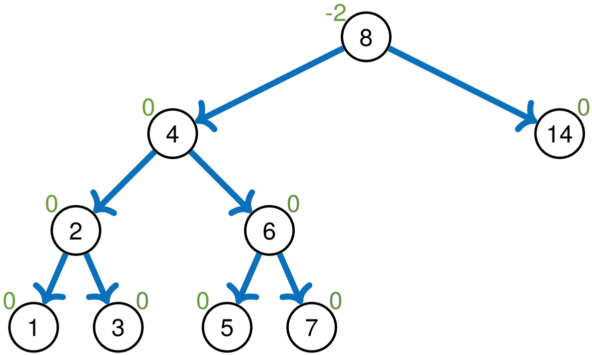


Figure: **Not** an AVL-Tree

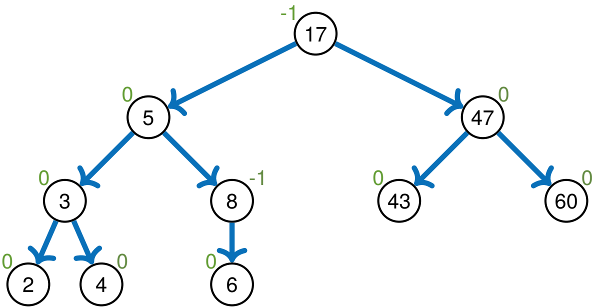


Figure: Another example of an AVL-Tree

Rotation:

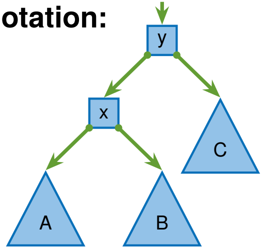
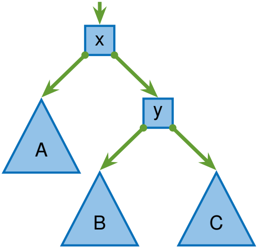


Figure: Before rotating

⇒



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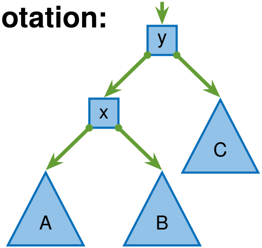

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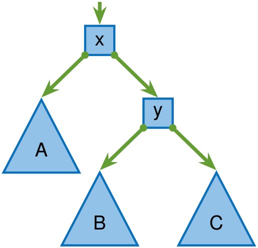

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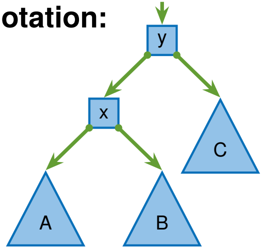

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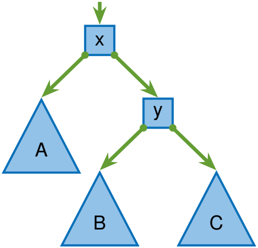

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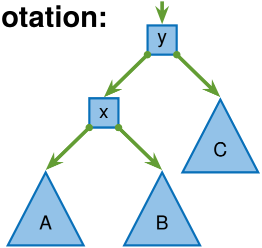

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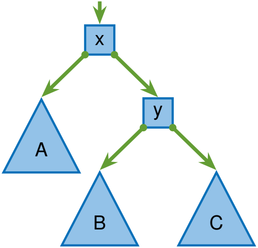

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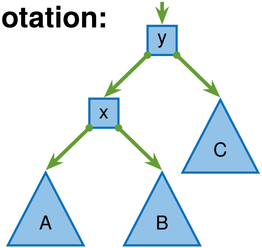

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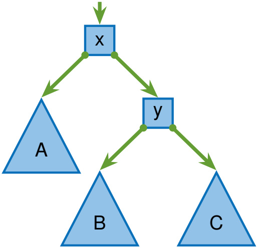

Figure: After rotating

- Central operation of rebalancing
- After rotation to the right:
  - Subtree A is a layer higher and subtree C a layer lower
  - The parent child relations between nodes x and y have been swapped

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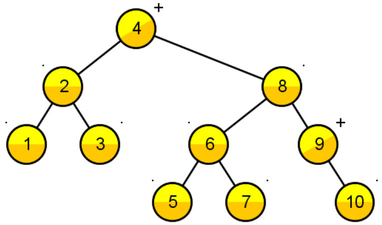

Figure: Inserting $1, \ldots, 10$ into an AVL-tree [Gna]

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- However not amortized update costs of $O(1)$
- Additional memory costs: We have to save a height difference for every node
- Better (and easier) to implement are (a,b)-trees

Structure

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- So we have space for elements on an `insert` and balance operation

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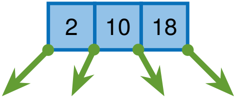

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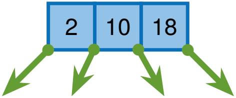

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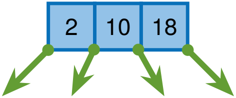

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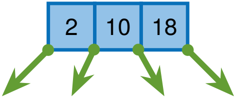

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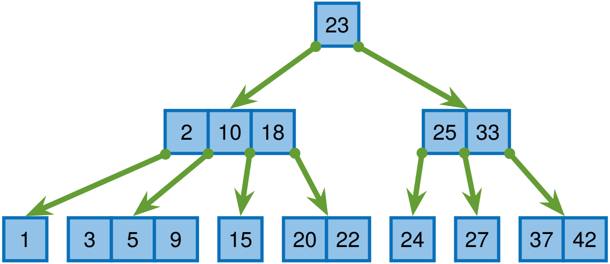
- Subtrees are located "between" the elements
- We require: $a \geq 2$ and $b \geq 2a-1$

(2,4)-Tree:



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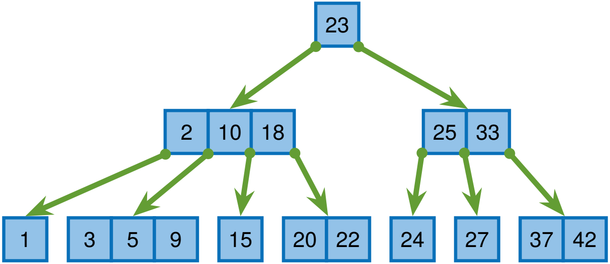

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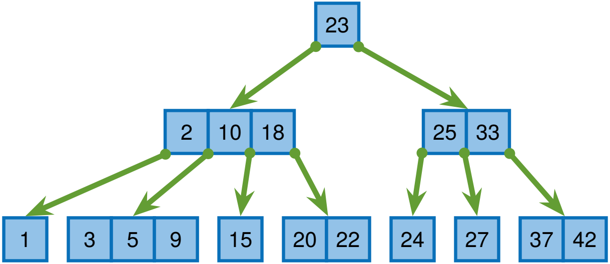

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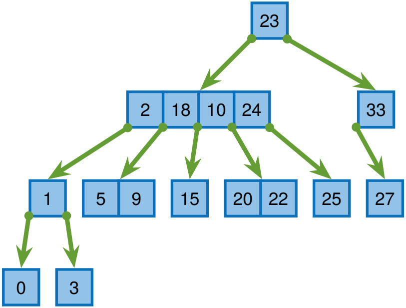
- (2,4)-tree with depth of 3
- Each node has between 2 and 4 children (1 to 3 elements)

Not an (2,4)-Tree:



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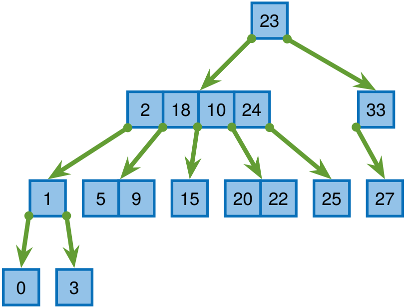

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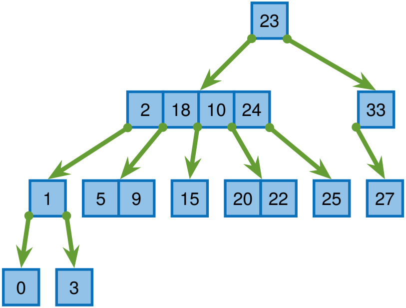

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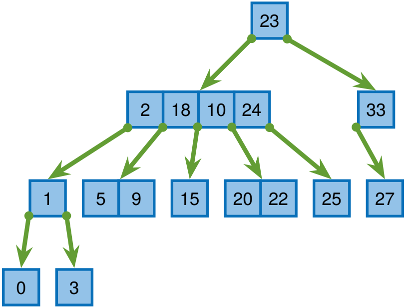

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- Invalid sorting
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- Leaves on different levels

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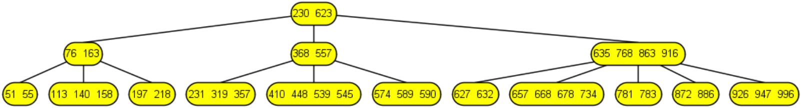

Figure: (3,5)-Tree [Gna]

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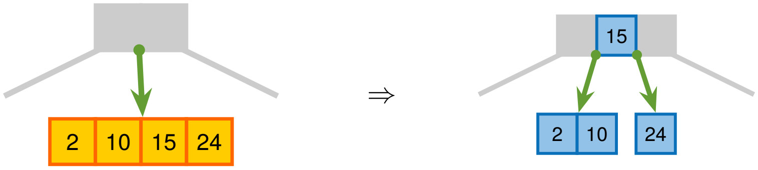

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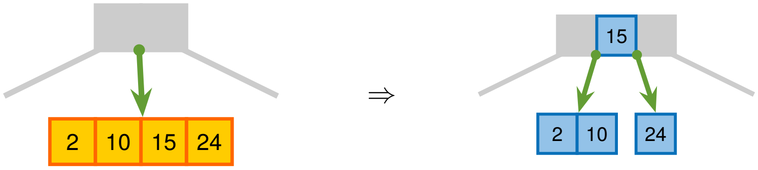

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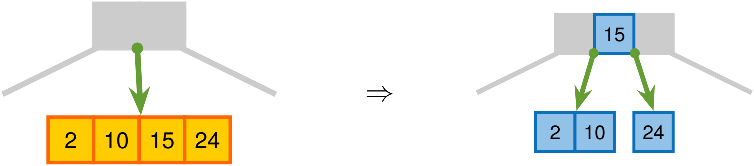

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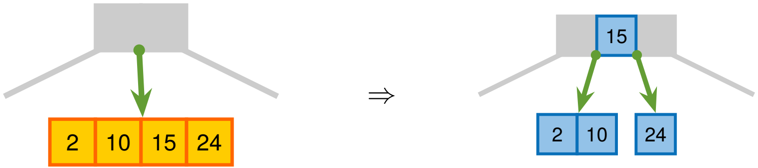

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- Thats why we have the limit $b \geq 2a - 1$

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- If we split the root node we create a new parent root node (The tree is now one level deeper)

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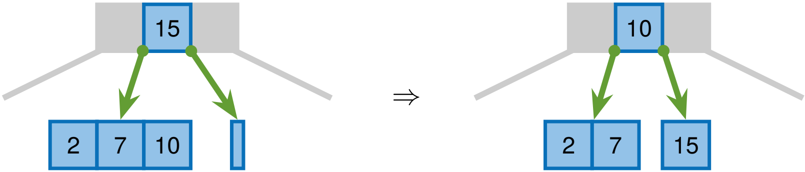

Figure: Borrow an element

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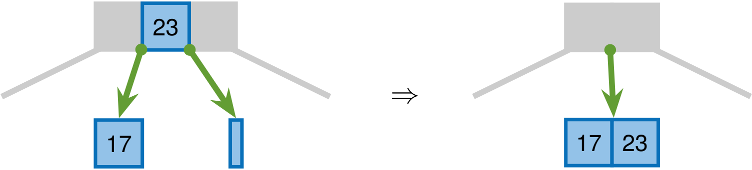

Figure: Merge two nodes

Removing an element: (`remove`)

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Removing an element: (`remove`)

- Now the parent node can be of degree $a - 1$
- We merge parent nodes the same way
- If the root has only a single child
  - Remove the root
  - Define sole child as new root
  - The tree shrinks by one level

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- `lookup` always takes $\Theta(d)$
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- **Worst case:** split or merge all nodes on path up to the root
- Therefore instead of $b \geq 2a - 1$ we need $b \geq 2a$

Counter example (2,3)-Tree:

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- Before executing `delete(11)`

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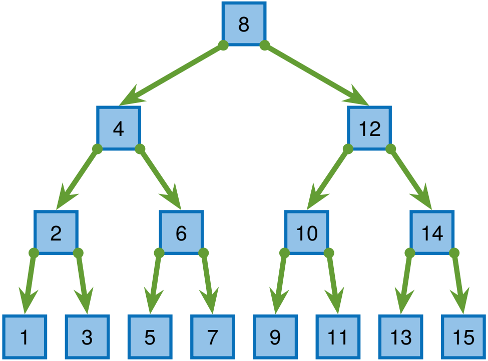

Figure: Normal (2,3)-Tree

Counter example (2,3)-Tree:

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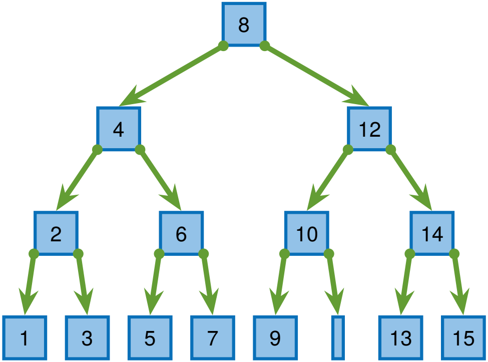

Figure: (2,3)-Tree - Delete step 1

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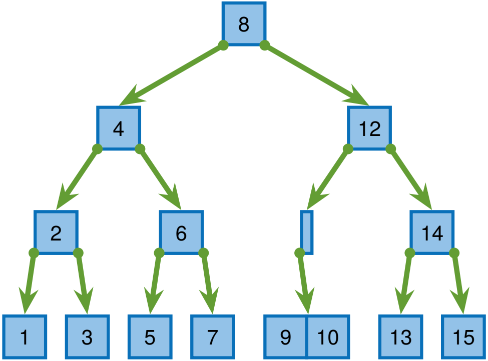

Figure: (2,3)-Tree - Delete step 2

Counter example (2,3)-Tree:

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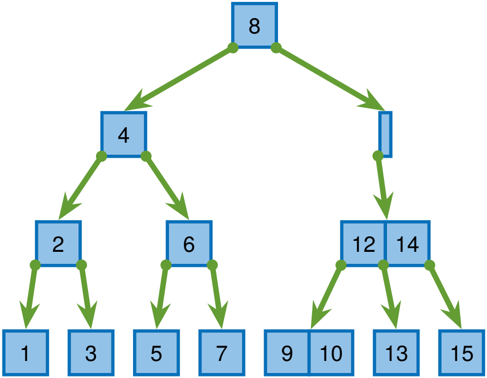

Figure: (2,3)-Tree - Delete step 3

Counter example (2,3)-Tree:

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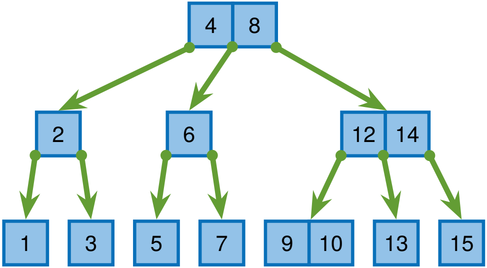

Figure: (2,3)-Tree - Delete step 4

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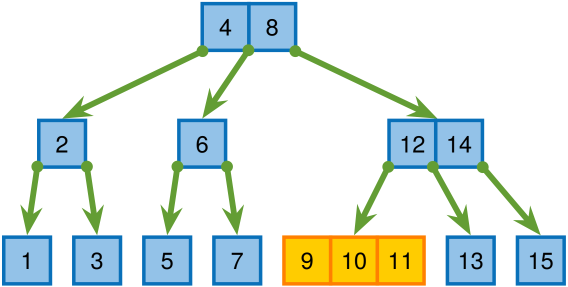

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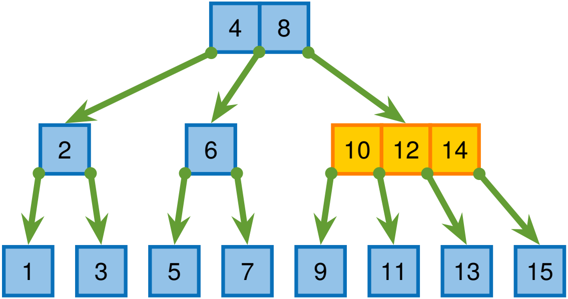

Figure: (2,3)-Tree - Insert step 2

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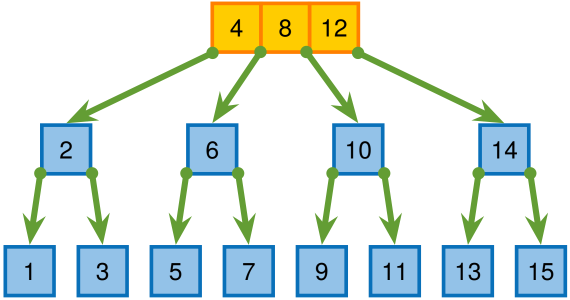

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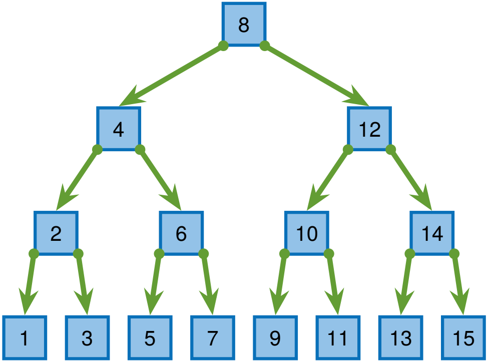

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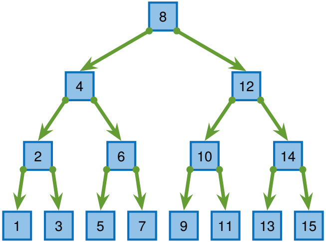
- We are exactly where we started



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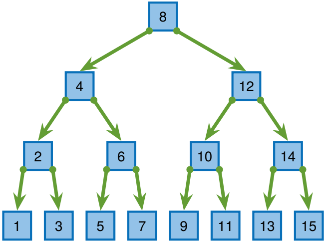

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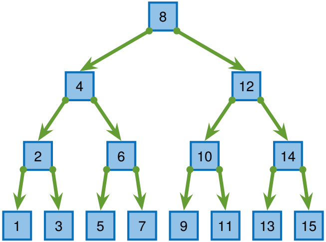
- We need $b \geq 2a$ instead of $b \geq 2a - 1$



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$\Rightarrow$ **Nodes of degree 3 are stable**
Neither an `insert` nor a `remove` operation trigger rebalancing operations

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(2,4)-Tree:

- **Idea:**
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- Like with dynamic arrays:
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  - If we overallocate clever we have an amortized runtime of $O(1)$

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- Empty tree has 0 nodes: $\Phi = 0$

Example:

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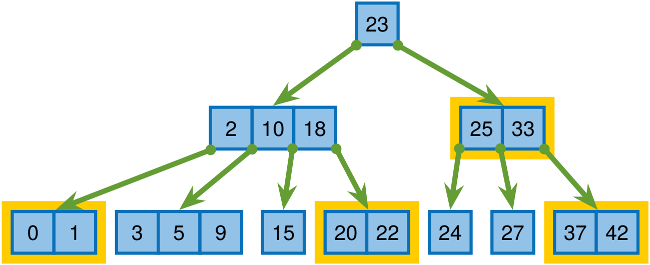

Figure: Tree with potential $\Phi = 4$

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- Each operation has an amortitzed cost of $O(1)$ summing up to $O(n)$ in total

Case 1: *i-th* operation is an `insert` operation on a full node

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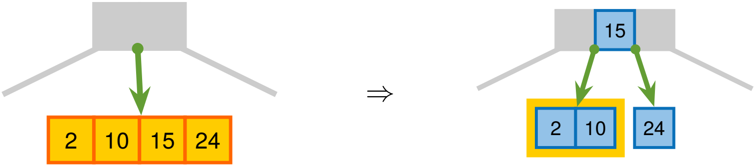

Figure: Splitting a node on `insert`

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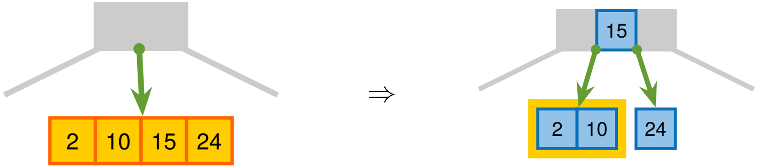

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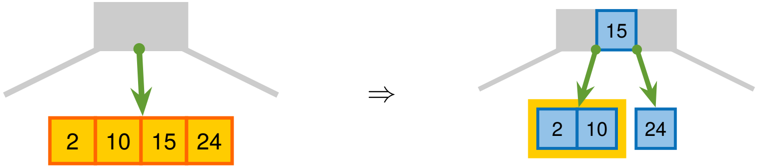

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- Each splitted node creates a node of degree 3
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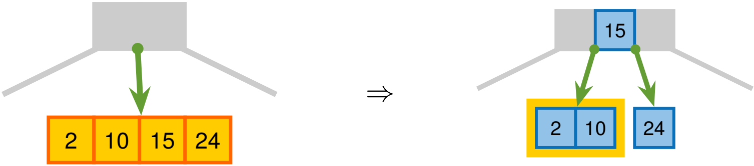

Figure: Splitting a node on `insert`

- Each splitted node creates a node of degree 3
- The parent node receives an element from the splitted node
- If the parent node is also full we have to split it too

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\begin{aligned}
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\Rightarrow m &\leq \Phi_i - \Phi_{i-1} + 1
\end{aligned}
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$$\Phi_i \geq \Phi_{i-1} + m - 1$$
$$\Rightarrow m \leq \Phi_i - \Phi_{i-1} + 1$$

Costs: $c_i \leq A \cdot m + B$

$$\Rightarrow c_i \leq A \cdot (\Phi_i - \Phi_{i-1} + 1) + B$$
$$c_i \leq A \cdot (\Phi_i - \Phi_{i-1}) + \underbrace{A + B}_{B'}$$

Case 2: *i-th* operation is an `remove` operation

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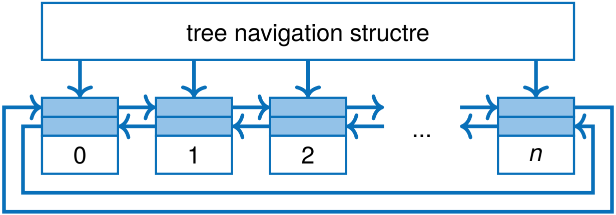

Figure: Tree with doubly linked list

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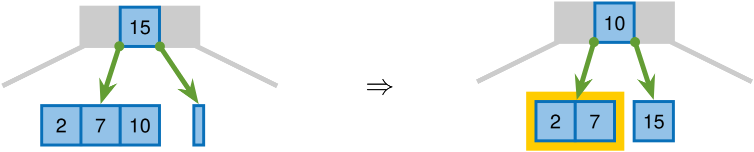

Figure: Case 2.1.1: Borrow an element

Case 2: *i-th* operation is an `remove` operation

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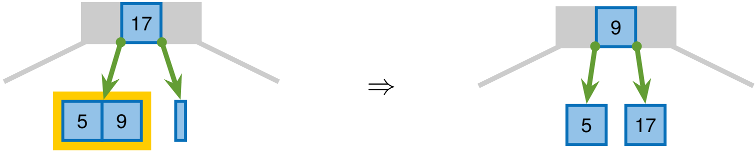

Figure: Case 2.1.2: Borrow an element

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- **Case 2.2:** Merging two node

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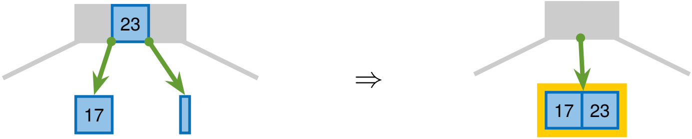

Figure: Merging two nodes

- Potential rises by one

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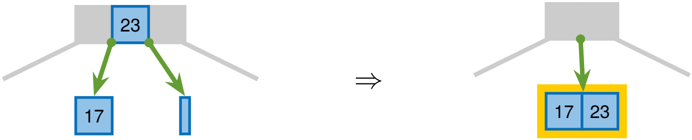

Figure: Merging two nodes

- Potential rises by one
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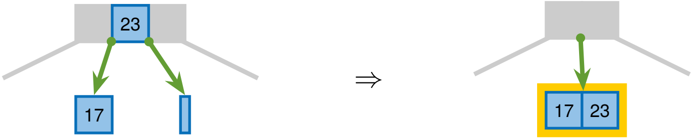

Figure: Merging two nodes

- Potential rises by one
- Parent node has one element less after the operation
- This operation propagates upwards until a node of degree > 2 or a node of degree 2, which can borrow from a neighbour

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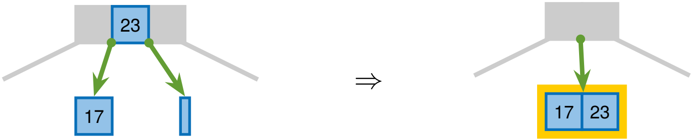
- **Case 2.2:** Merging two node



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- **Case 2.2:** Merging two node

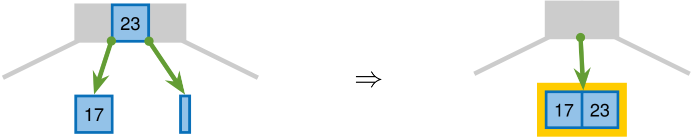

Figure: Merging two nodes

- The potential rises by *m*

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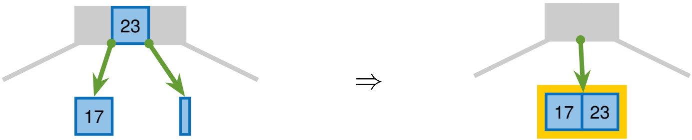

Figure: Merging two nodes

- The potential rises by m
- If the "stop-node" is of degree 2 then the potential eventually goes down by one

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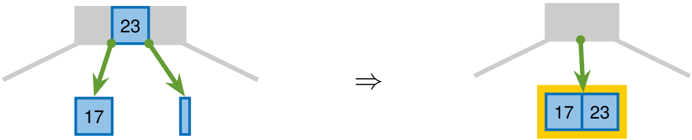

Figure: Merging two nodes

- The potential rises by *m*
- If the “stop-node” is of degree 2 then the potential eventually goes down by one
- Same costs as `insert`

Lemma:

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- With that we can conclude:

$$\sum_{i=0}^{n} c_i \in O(n)$$

Proof:

$$\sum_{i=0}^{n} c_i \leq \underbrace{A \cdot (\Phi_1 - \Phi_0) + B}_{\leq c_1} + \underbrace{A \cdot (\Phi_2 - \Phi_1) + B}_{\leq c_2} + \cdots + \underbrace{A \cdot (\Phi_n - \Phi_{n-1}) + B}_{\leq c_n}$$

$$= A \cdot (\Phi_n - \Phi_0) + B \cdot n \qquad | \text{ telescope sum}$$

$$= A \cdot \Phi_n + B \cdot n \qquad | \text{ we start with an empty tree}$$

$$< A \cdot n + B \cdot n \in O(n) \qquad | \text{ number of degree 3 nodes } < \text{ number of nodes}$$

Structure

Red-Black Tree:

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- Binary tree with red and black nodes

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- Binary tree with red and black nodes
- Number of black nodes on path to leaves is equal
- Can be interpreted as (2,4)-tree (also named 2-3-4-tree)
- Each (2,4)-tree-node is a small red-black-tree with a black root node

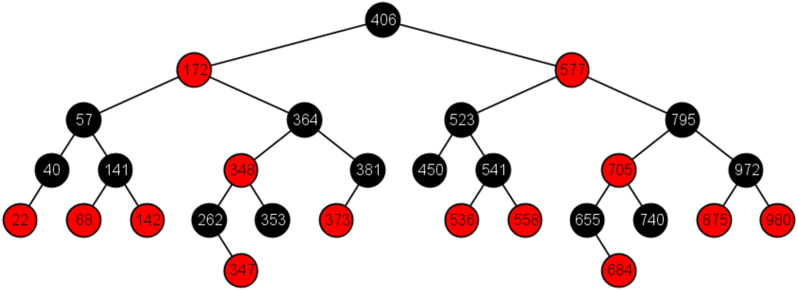

Figure: Example of an red-black-tree [Gna]

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