Algorithms and Datastructures Runtime analysis Minsort / Heapsort, Induction

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Albert-Ludwigs-Universität Freiburg

Prof. Dr. Rolf Backofen

Bioinformatics Group / Department of Computer Science Algorithms and Datastructures, October 2018

Structure



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Forum

Checkstyle

Unit Tests

Version management

Jenkins

Sorting

Minsort

Heapsort

Topics of the Lecture:

- Algorithms and Data Structures
- **Algorithm** Solving of complex computional problems

Topics of the Lecture:

- Algorithms and Data Structures
 Efficient data handling and processing
 ... for problems that occur in practical any larger program / project
- Algorithm

 Solving of complex computional problems
- **Datastructure**

 Representation of data on computer

Example 1: Sorting



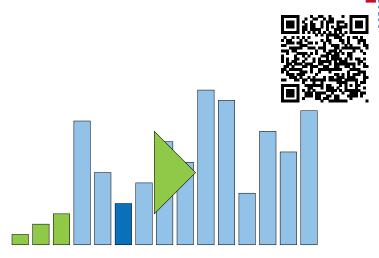


Figure: Sorting with *Minsort*



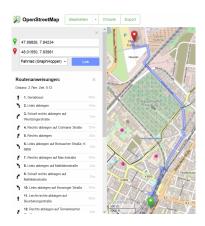


Figure: Navigationplan © OpenStreetMap

■ **Datastructures:** How to represent the map as data?



Figure: Navigationplan © OpenStreetMap

- Datastructures: How to represent the map as data?
- **Algorithms:** How to find the shortest / fastest way?



Figure: Navigationplan © OpenStreetMap

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Content of the Lecture 1 / 2



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■ Most of you had a lecture on basic progamming ... performance was not an issue



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Content of the Lecture 1 / 2



- Most of you had a lecture on basic progamming ... performance was not an issue
- Here it is going to be:
 - How fast is our program?
 - 2 How can we make it faster?
 - 3 How can we proof that it will always be that fast?
- Important issues:
 - Most of the time: application runtime
 - Sometimes also: resource / space consumption



Algorithms:



Algorithms:

- Sorting
- Dynamic Arrays
- Associative Arrays
- Hashing
- Edit distance

- Priority Queue
- Linked Lists
- Pathfinding / Dijkstra Algorithm
- Search Trees



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- Sorting
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Mathematics:

Content of the Lecture 2 / 2



Algorithms:

- Sorting
- Dynamic Arrays
- Associative Arrays
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- Priority Queue
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- Pathfinding / Dijkstra Algorithm
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Mathematics:

- Runtime analysis
- Ø-Notation

Proof of correctness



■ ... you should be able to understand the joke

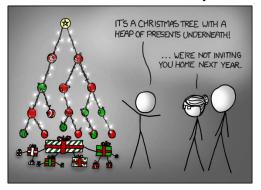


Figure: Comic @ xkcd/835

After the lecture ...



■ ... you should be able to understand the joke

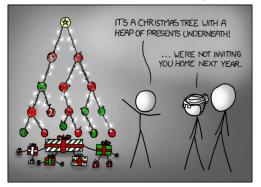


Figure: Comic @ xkcd/835

■ Hopefully your parents will still invite you

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Homepage:

- Exercise sheets
- Lectures
- Materials

Link to Homepage

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Lecture:

- Tuesday, 12:00 14:00, HS 00 006, Build. 082
- Recordings of the lecture will be uploaded to the webpage

Exercises:

- One exercise sheet per week
- Submission / Correction / Assistance online
- Tutorial: (if needed)Wednesday, 13:00-14:00 HS 00 006, Build. 082

Exam:

■ Planned: Sa. 23th March 2019, 10:00-12:00, Build. 101, Lec. theater 026 & 036



■ 80% practical, 20% theoretical



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■ 50% of all points from the exercise sheets are needed

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Exam:

- 50% of all points from the exercise sheets are needed
- Content of exam: whole lecture and all exercises



■ Tutors: Tim Maffenbeier, Till Steinmann, Tobias Faller



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- Deadline: ESE: 1 week, IEMS: none



Post questions into the forum (link later)



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Exercises:

- Post questions into the forum (link later)
- Submission via "commit" through svn and Daphne
- Feedback one week after deadline through "update" (svn)
- Unit test / checkstyle via Jenkins



Exercises - Points:

- Practical:
 - 60% functionality
 - 20% tests
 - 20% documentation, Checkstyle, etc.
 - Program is not running \Rightarrow 0 points

Exercises - Points:

- Practical:
 - 60% functionality
 - 20% tests
 - 20% documentation, Checkstyle, etc.
 - Program is not running \Rightarrow 0 points
- Theoretical (mathematical proof):
 - 40% general idea / approach
 - 60% clean / complete

Effort:

- 4 ECTS (ESE), 6 ECTS (IEMS)
- 120 / 180 working hours per semester
- 14 Lectures each 6h / 8h + exam
- 4h / 6h per exercise sheet (one per week)

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Daphne:

- Provides the following information:
 - Name / contact information of your tutor
 - Download of / info needed for exercise sheets
 - Collected points of all exercise sheets
 - Links to:
 - Coding standards
 - 2 Build system
 - 3 The other systems
- Link: Daphne

Forum:

- Please don't hesitate to ask if something is unclear
- Ask in the forum and not separate. Others might also be interested in the answer
- The tutors or the coordinators will reply as soon as possible
- Link: Forum

Checkstyle / Linting (flake8):

■ Installation: python3 -m pip install flake8

■ Check file: **python3** -m flake8 path/to/files/*.py

■ Link: flake8

Unit Tests



Why unit tests?

A non-trivial method without a unit test is probably wrong

Unit Tests



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- Simplifies debugging



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Unit test checks desired output for a given input



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- At least one typical input

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What is a good unit test?

- Unit test checks desired output for a given input
- At least one typical input
- At least one critical case
 E.g. double occurrence of a value in sorting



Testing (doctest):

```
def subOne(n):
    """Subtracts 1 from n
    >>> subOne(5)
    >>> subOne(3)
    . . .
    return n-2
if __name__ == "__main__":
    print("2 minus 1: %d" % subOne(2))
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doctest

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- Tests are contained in docstrings
- Module doctest runs them
- Run check with: python3 -m doctest path/to/files/*.py -v

Version management (subversion):

- Keeps a history of code changes
- Initialize / update directory: svn checkout <URL>
- Add files / folders: **svn** add <file> --all
- Create snapshot: svn commit -m "<Your Message>" Data is uploaded to Jenkins automatically
- Link: Subversion

Jenkins:

- Provides our build system
- You can check if your uploded code runs
 - Especially whether all unit test pass
 - And if checkstyle (flake8) is statisfied
- Will be shown in the first exercise
- Link: Jenkins

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Problem:

- Input: *n* elements $x_1, ..., x_n$
- Transitive operator "<" which returns true if the left value is smaller than the right one
 - Transitivity: x < y, $y < z \rightarrow x < z$
- Output: x_1, \ldots, x_n sorted with operator

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Example

Input: 14, 4, 32, 19, 8, 44, 65

Output:



Why do we need sorting?

- Nearly every program needs a sorting algorithm
- Examples:
 - Index of a search engine
 - Listing filesystem in explorer / finder
 - (Music) library
 - Highscore list

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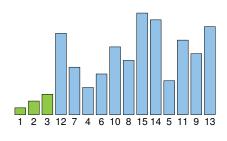
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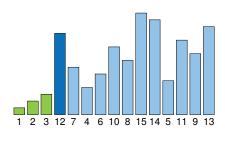
- Find the minimum and switch the value with the first position
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. . . .



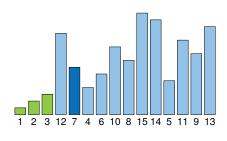
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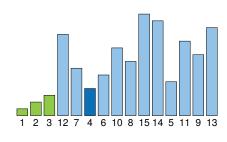


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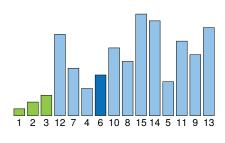


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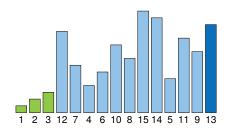


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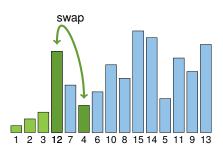
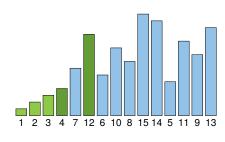


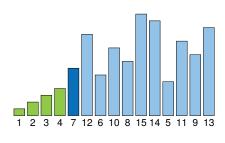
Figure: Minsort

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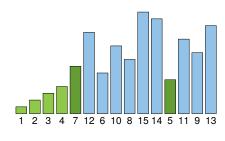


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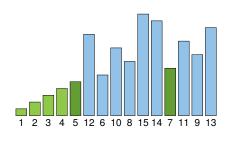


Figure: Minsort

Minsort in Python:

```
def minsort(lst):
    for i in range (0, len(lst)-1):
        minimum = i
        for j in range(i+1, len(lst)):
             if lst[j] < lst[minimum]:</pre>
                 minimum = i
        if minimum != i:
             [st[i], [st[minimum] = \]
                 Ist[minimum], Ist[i]
    return 1st
```

MinSort - Runtime



How long does our program run?

We test it for different input sizes



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We test it for different input sizes

Table: Run	time for <i>Minsort</i>
n	Runtime / ms
2×10^3	5.24
4×10^3	16.92
6×10^3	39.11
8×10^3	67.80
10×10^3	105.50
12×10^3	150.38
14×10^3	204.00
16×10^3	265.98
18×10^{3}	334 94

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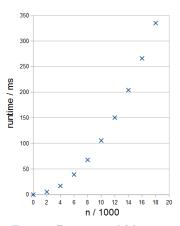


Figure: Runtime of *Minsort*



Runtime analysis:

- Minsort runtime depicted in a diagram
 - That is what you should do in the first exercise sheet

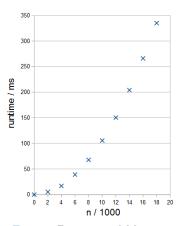


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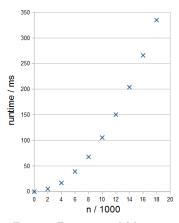


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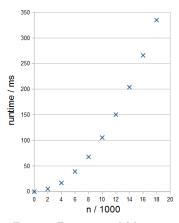


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- Better structure for finding the smallest element quicker

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Binary heap:

- Preferably a complete binary tree
- Heap property: Each child is smaller (larger) than the parent element



Min heap:

■ **Heap property:** Each child is smaller (larger) than the parent element

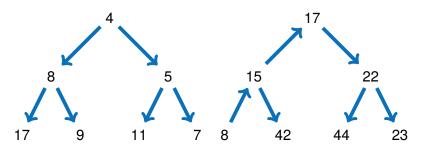


Figure: Valid min heap

Figure: Invalid min heap



Min heap:

- **Heap property:** Each child is smaller (larger) than the parent element
- A valid heap fulfills the property at each node

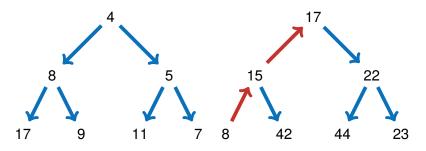


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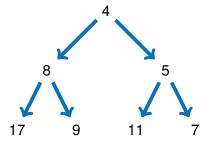
Figure: Invalid min heap

Heapsort - Algorithm 3 / 10



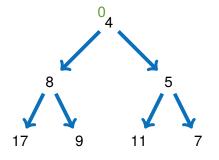
How to save the heap?

- We number all nodes from top to bottom and left to right starting at 0
 - The children of node i are 2i + 1 and 2i + 2
 - The parent node of node *i* is floor $\left(\frac{i-1}{2}\right)$



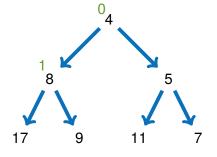


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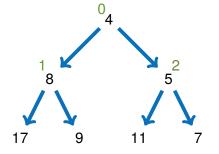
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0	1			
4	8			

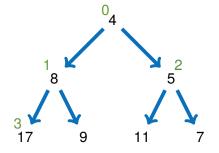


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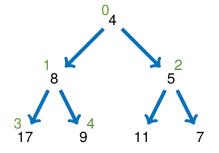
0	1	2		
4	8	5		

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0	1	2	3	
4	8	5	17	

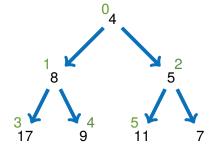
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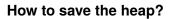
0	1	2	3	4	
4	8	5	17	9	



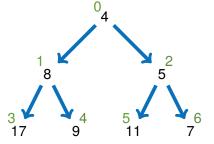
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4	8	5	17	9	11	



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0	1	2	3	4	5	6
4	8	5	17	9	11	7





Remove the smallest element (root node)



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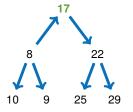


Figure: Repairing a min heap



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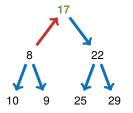


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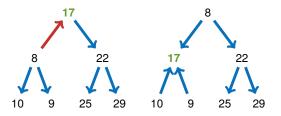


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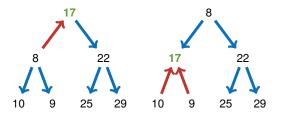


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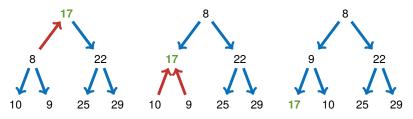


Figure: Repairing a min heap

- Organize the *n* elements as heap
- While the heap still contains elements
 - Take the smallest element
 - Move the last node to the root
 - Repair the heap as described

- Organize the *n* elements as heap
- While the heap still contains elements
 - Take the smallest element
 - Move the last node to the root
 - Repair the heap as described
- Output: 4

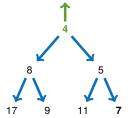


Figure: One iteration of Heapsort

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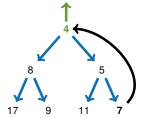


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HeapSort - Algorithm 5 / 10

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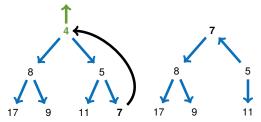


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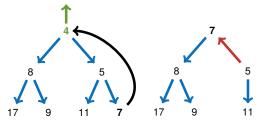


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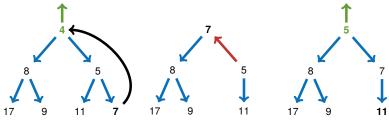


Figure: One iteration of Heapsort



Heapsort:

- Organize the n elements as heap
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 - Move the last node to the root
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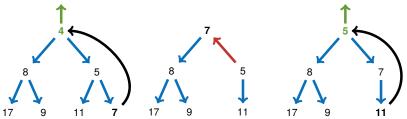


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- Interpret the array as binary heap where the heap property is not yet satisfied
- We repair the heap from bottom up (in layers) with sift



Table: Input in array

0	1	2	3	4	5	6
11	7	8	3	2	5	4

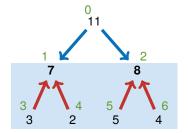
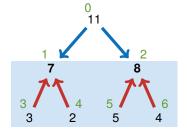


Figure: Heapify lower layer



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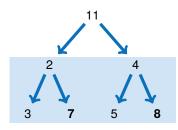


Figure: Heapify lower layer

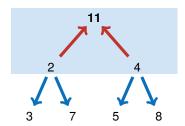
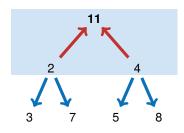


Figure: Heapify upper layer



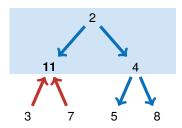


Figure: Heapify upper layer

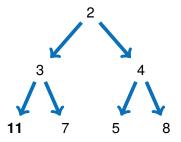


Figure: Resulting heap



Finding the minimum is intuitive:

- Minsort: Iterate through all non-sorted elements
- Heapsort: Finding the minimum is trivial (concept)
 Just take the root of the heap

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- **Heapsort:** Finding the minimum is trivial (concept)

 Just take the root of the heap

Removing the minimum in Heapsort:

- Repair the heap and restore the heap property
 - We don't have to repair the whole heap
- More of this in the next lecture

■ Course literature

[CRL01] Thomas H. Cormen, Ronald L. Rivest, and Charles E. Leiserson. Introduction to Algorithms. MIT Press, Cambridge, Mass, 2001.

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[Wikb] Wikipedia - Selectionsort

https://de.wikipedia.org/wiki/Selectionsort

Further Literature



Subversion

[Apa] Apache Subversion

https://subversion.apache.org/