

Problem type 1:

Please provide a algorithm for the problem below. Any algorithms given should be computationally efficient (please no brute-forcing). If needed, you may refer to the *Explore* algorithm (below) from lecture as a black box:

```
EXPLORE( $G, u$ ):
  Visited[1.. $n$ ]  $\leftarrow$  false
  Add  $u$  to ToExplore and to  $S$ 
  Visited[ $u$ ]  $\leftarrow$  true
  While (ToExplore is non-empty)
    Remove node  $x$  from ToExplore
    for each edge  $xy$  in  $Adj(x)$ 
      if (Visited[ $y$ ] = false)
        Visited[ $y$ ]  $\leftarrow$  true
        Add  $y$  to ToExplore
        Add  $y$  to  $S$ 
  return  $S$ 
```

(See variants below)

a. **BYG**

G is a directed graph and I want to know if node u can reach node v .

b. **BYD**

G is a directed graph and I want to find all nodes that u can reach.

c. **BYA**

G is a directed graph and I want to find all nodes that can reach u .

d. **BYH**

G is a directed graph and I want to find all nodes in u 's strong connected component.

Problem type 2:

Answer the following problem:

(See variants below)

a. **BYC**

What type of *directed*-graph has only one strongly connected component?

b. **BYF**

Assuming a directed graph with n nodes, how many edges do you need to make the graph have a single strongly connected component.

c. **BYE**

Assuming a *undirected* graph with n nodes, how many edges do you need to make the graph connected.

d. **BYB**

Assuming a *directed* graph with n nodes, what type of graph would have the most number of cycles of size n . How many cycles would this graph have.