1. Let $L = \{w \in \{\mathtt{0}, \mathtt{1}\}^* \mid w \text{ starts and ends with } \mathtt{0}\}$.

   (a) Construct an NFA for L with exactly three states.

   (b) Convert the NFA you just constructed into a DFA using the incremental subset construction. Draw the resulting DFA. Your DFA should have four states, all reachable from the start state.

   (c) Convert the DFA you constructed in part (b) into a regular expression using the state elimination algorithm.

   (d) Write a simpler regular expression for L.

2. (a) Convert the regular expression $(\mathtt{0}^*\mathtt{1} + \mathtt{01}^*)^*$ into an NFA using Thompson's algorithm.

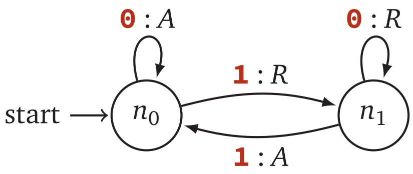
   (b) Convert the NFA you just constructed into a DFA using the incremental subset construction. Draw the resulting DFA. Your DFA should have four states, all reachable from the start state. (Some of these states are obviously equivalent, but keep them separate.)

   (c) Convert the DFA you constructed in part (b) into a regular expression using the state elimination algorithm. You should *not* get the same regular expression you started with.

   (d) **Think about later:** Find the smallest DFA that is equivalent to your DFA from part (b) and convert that smaller DFA into a regular expression using the state elimination algorithm. Again, you should *not* get the same regular expression you started with.

   (e) What is this language?

3. In a previous lab/homework we talked about a new machine called a *finite-state transducer* (FST). The special part thing about this type of machine is that it gives an output on the transition instead of the state that it is in. An example of a finite state transducer is as follows:



defined by the five tuple: $(\Sigma, \Gamma, Q, \delta, s)$. Let's constrain this machine (call is FST_{AR}) a bit and say the output alphabet consists of two signals: accept or reject ($\Gamma = \{A, R\}$). We say that $L(FST_{AR})$ represents the language consisting of all strings that end with a accept (A) output signal.

Prove that $L(FST_{AR})$ represents the class of regular languages.

4. Let $L = \{w \in \{\mathtt{0}, \mathtt{1}\}^* \mid \text{a } \mathtt{0} \text{ appears in some position } i \text{ of } w, \text{ and a } \mathtt{1} \text{ appears in position } i+2\}$.

   (a) Construct an NFA for L with exactly four states.

   (b) Convert the NFA you just constructed into a DFA using the incremental subset construction. Draw the resulting DFA. Your DFA should have eight states, all reachable from the start state.

(c) Convert the NFA you constructed in part (a) into a regular expression using the state elimination algorithm.

5. **Think about later:**

(a) Convert the regular expression $(\varepsilon + (\mathtt{0} + \mathtt{11})^* \mathtt{0})\mathtt{1}(\mathtt{11})^*$ into an NFA using Thompson's algorithm.

(b) Convert the NFA you just constructed into a DFA using the incremental subset construction. Draw the resulting DFA. Your DFA should have six states, all reachable from the start state. (Some of these states are obviously equivalent, but keep them separate.)

(c) Convert the DFA you constructed in part (b) into a regular expression using the state elimination algorithm. You should *not* get the same regular expression you started with.

(d) What is this language?