

## Problem type 1:

Last quiz you were asked to provide the recurrence that describes one of the backtracking problems from Labs 11/12. I will give you the recurrence to that problem below. **Now I want to know the evaluation order of the recurrence.** Specifically I want three things:

- The number of for loops needed to evaluate the recurrence.
- The order of each of those for loops ( $1 \rightarrow n$ ,  $n \rightarrow 1$ ,  $i \rightarrow n$ , etc. )
- The return value (which value/part of the array do we return)

```
A[n][n] = <Base Cases>
for <loop1 conditions>
    ▸ Fill in if needed
    for <loop2 conditions>
        ▸ Fill in if needed
        for <loop3 conditions>
            ▸ Fill in if needed
            Compute
        return ##
    ▸ Fill in
```

Not looking for full pseudocode. Just a basic idea of how to memorize the recurrence.

(See variants below)

### a. BYA & BYH

Given an array  $A[1..n]$  of integers, compute the length of a **longest alternating subsequence**: Let  $LAS^+(i, j)$  denote the length of the longest alternating subsequence of  $A[i..n]$  whose first element (if any) is larger than  $A[j]$  and whose second element (if any) is smaller than its first.

$$LAS^+(i, j) = \begin{cases} 0 & \text{if } i > n \\ LAS^+(i+1, j) & \text{if } i \leq n \text{ and } A[i] \leq A[j] \\ \max\{LAS^+(i+1, j), 1 + LAS^-(i+1, i)\} & \text{otherwise} \end{cases}$$

$$LAS^-(i, j) = \begin{cases} 0 & \text{if } i > n \\ LAS^-(i+1, j) & \text{if } i \leq n \text{ and } A[i] \geq A[j] \\ \max\{LAS^-(i+1, j), 1 + LAS^+(i+1, i)\} & \text{otherwise} \end{cases}$$

#### Solution:

- two for loops:
  - $i \leftarrow n$  down to 1
  - $j \leftarrow i - 1$  down to 1
- for  $i = 1 \rightarrow n$  return  $\max(LAS^+(i+1, 1), LAS^-(i+1, 1))$ . This problem is slightly harder than the others in the time required so we'll accept anything of the form:  $\max(LAS^+(n, n), LAS^-(n, n))$

## b. BYC &amp; BYE

Given an array  $A[1..n]$  of integers, compute the length of a **longest decreasing subsequence**. Let  $LDS(i, j)$  denote the length of the longest decreasing subsequence of  $A[i..n]$  where every element is smaller than  $A[j]$ .

$$LDS(i, j) = \begin{cases} 0 & \text{if } i > n \\ LDS(i+1, j) & \text{if } i \leq n \text{ and } A[j] \leq A[i] \\ \max\{LDS(i+1, j), 1 + LDS(i+1, i)\} & \text{otherwise} \end{cases}$$

**Solution:**

- two for loops:
  - $i \leftarrow n$  down to 1
  - $j \leftarrow i - 1$  down to 0
 We won't make off for  $\pm 1$  issues.
- return  $A[1, 0]$ .

## c. BYD &amp; BYG

Given an array  $A[1..n]$ , compute the length of a longest **palindrome** subsequence of  $A$ . Let  $LPS(i, j)$  denote the length of the longest palindrome subsequence of  $A[i..j]$ .

$$LPS(i, j) = \begin{cases} 0 & \text{if } i > j \\ 1 & \text{if } i = j \\ \max \left\{ \begin{array}{l} LPS(i+1, j) \\ LPS(i, j-1) \end{array} \right\} & \text{if } i < j \text{ and } A[i] \neq A[j] \\ \max \left\{ \begin{array}{l} 2 + LPS(i+1, j-1) \\ LPS(i+1, j) \\ LPS(i, j-1) \end{array} \right\} & \text{otherwise} \end{cases}$$

**Solution:**

- two for loops:
  - $i \leftarrow n$  down to 1
  - $j \leftarrow i + 1$  down to  $n$
 We won't make off for  $\pm 1$  issues.
- return  $A[1, n]$ .

d. **BYB & BYF**

Given an array  $A[1..n]$  of integers, compute the length of a longest **convex** subsequence of  $A$ . Let  $LCS(i, j)$  denote the length of the longest convex subsequence of  $A[i..n]$  whose first two elements are  $A[i]$  and  $A[j]$ .

$$LCS(i, j) = 1 + \max\{LCS(j, k) \mid j < k \leq n \text{ and } A[i] + A[k] > 2A[j]\}$$

**Solution:**

- three for loops:
  - $i \leftarrow n - 1$  down to 1
  - $j \leftarrow n$  down to  $i + 1$
  - $k \leftarrow j + 1$  to  $n$

We won't make off for  $\pm 1$  issues. Mainly want to see three for loops.

- return  $\max(A[1..n-1, i+1..n])$ . Will give full credit to anyone that take the max of the first two dimensions.

