

## **Visualizing Dynamic Programming On Tree Decompositions**

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- REWORK AFTER CHAPTERS
- ▶ WHAT was the motivation
- ▶ WHAT could be used otherwise?
- WHO benefits from visualization?
- ▶ METHODOLOGY challenges and solutions
- ▶ WHAT could be developed next?





#### Motivation

- DP-on-TD-algorithms can solve Model Counting and various combinatorial problems
- Implementations of those are competing with modern solvers
- ▶ But: those are fairly hard to implement efficiently
- Practical debug output quickly becomes very large (GB)
- ▶ Finding the cause of the problem is a time consuming challenge

The B.T. probably focused too much on the convenience features, and not on the urgent need for better debugging and visualization needs for those algorithms.

## **Background**

The algorithms of interest solve problems of:

- combinatorics (NP-problems)
- model-counting (#P-problems)

Recent promising results for Projected Model Counting by Markus Hecher<sup>1</sup>.

<sup>&</sup>lt;sup>1</sup> Hecher M., Thier P., Woltran S. (2020) Taming High Treewidth with Abstraction, Nested Dynamic Programming, and Database Technology. In: Pulina L., Seidl M. (eds) Theory and Applications of Satisfiability Testing - SAT 2020. SAT 2020. SAT 2020. In: Computer Science, vol 12178. Springer, Cham. https://doi.org/10.1007/978-3-030-51825-7.25



#### **Tree Decompositions**

A tree decomposition is a tree obtained from an arbitrary graph s.t.

- 1. Each vertex must occur in some bag
- 2. For each edge, there is a bag containing both endpoints
- 3. <u>Connected</u>: Subgraph "restricted" to any vertex must be connected Graphic



## **Graphs for Boolean Formulas**

#### ► Example set of CNF-clauses:

$$\{c1 = \{v1, v3, \neg v4\}, c2 = \{\neg v1, v6\}, c3 = \{\neg v2, \neg v3, \neg v4\}, c4 = \{\neg v2, v6\}, c5 = \{\neg v3, \neg v4\}, c6 = \{\neg v3, v5\}, c7 = \{\neg v5, \neg v6\}, c8 = \{v5, v7\}\}$$



Figure: The primal (left), incidence (middle) and dual (right) graph



## gpusat2 - Solving on GPU



- Customized tree decompositions
- Adapted memory-management
- Improved precision handling



<sup>&</sup>lt;sup>1</sup> Images: Markus Zisser. Solving the #SAT problem on the GPU with dynamic programming and OpenCL. Technische Universität Wien. 2018.



### dpdb

# Database templates in Python Generating SQL queries

- Create graph representation
- 2. Decompose graph
- 3. Solve sub-problems
- 4. Combine rows
- SAT and #SAT
- #o-Coloring
- Vertex cover

. . .

```
- #sTab#:
                             SELECT 1 AS ont
– #intrTab#:
                             SELECT 1 AS val UNION ALL O
- #localProbFilter#: (l_{1,1} 	ext{ OR } ... 	ext{ OR } l_{1,k_1}) 	ext{ AND } ... 	ext{ AND } (l_{n,1} 	ext{ OR } ... 	ext{ OR } l_{n,k_n})
#aggrExp#:
                             SUM(cnt) AS cnt
- #extProj#:
                             \tau_1.cnt * ... * \tau_\ell.cnt AS cnt
                                       (a) Problem #SAT
- #εTab#:
                        SELECT 0 AS card
                        SELECT 1 AS val UNION ALL O
- #localProbFilter#: ([u<sub>1</sub>] OR [v<sub>1</sub>]) AND ... AND ([u<sub>n</sub>] OR [v<sub>n</sub>])
- #aggrExp#:
                        MIN(card) AS card
                        \tau_1.card + ... + \tau_\ell.card - (\Sigma_{i=1}^{\ell} | \chi(t_i) \cap \{a_1\} | - 1) *
                        \tau_1 \cdot [a_1] - \dots - (\Sigma_{i-1}^{\ell} | \chi(t_i) \cap \{a_k\} | -1) * \tau_1 \cdot [a_k]
```

(b) Problem MinVC



## Diff. debug information and problem types

Display user-defined strings where applicable Different ?????? Grafik bag, Grafik solution



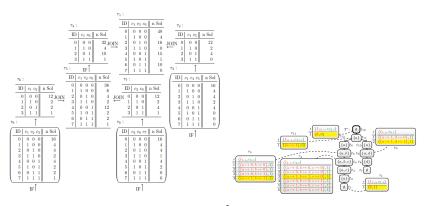
## Challenge2

- Wie robust ist die Datenverarbeitung in der Visu Restrictions in strings for ids - Was Gedanken bei der Visu waren



#### Visualization

Manually for one run



dpdb????2

<sup>&</sup>lt;sup>2</sup>"Solving #SAT on the GPU with Dynamic Programming and OpenCL", Diploma Markus Zisser 2018 Technische Universität Wien, p.33





© Gephi.org - a tool for data analysts and scientists keen to explore and understand graphs.3



Tulip - Better Visualization Through Research. 4



=¿ Related Work Schluss / Wiss Arbeiten -¿ Nicht speziell Angeschaut / Format aus Solvern extrahiert - kann trotzdem sehr generisch sein (dpdb speziell)

<sup>3</sup>https://aephi.org/

<sup>4</sup>https://tulip.labri.fr/TulipDrupal/

<sup>&</sup>lt;sup>5</sup>https://neo4i.com/developer/tools-graph-visualization/



MinVC example size 90 (expected 82)



1. Inspect visualization



- 1. Inspect visualization
- 2. Verify findings in solver (in this case dpdb)

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3. Cross reference with standalone tree-decomposition

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3. Cross reference with standalone tree-decomposition



4. Fix the root cause



#### Outlook

## Static → Dynamic

- Enrich the visualization with debugging info for each node
- cross reference the creation of rows in parent nodes
- For more advanced debugging tasks you may also need to revise the approach
- Bottlenecks of different architectures
- Preselection of interesting areas/data to visualize
- Utilizing graph databases for visualization and queries for debugging



### Final slide

I would like to thank you for your attention.

#### **Benchmark**

#### Performance of all three programs on #SAT instances:





#### Visualization

Manually for dpdb

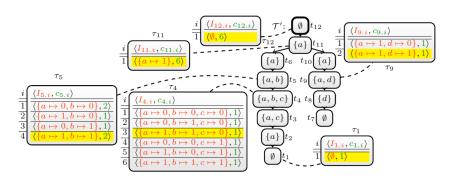


Figure: Handcrafted #SAT example-run from dpdb7

 $<sup>^{7}{\</sup>rm "Exploiting}$  Database Management Systems and Treewidth for Counting", Fichte, Hecher, Thier, Woltran



## **Bibliography**