Visualizing Dynamic Programming On Tree Decompositions

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- WHAT **is the motivation**
- WHO **benefits from visualization?**
- CHALLENGES **and solutions**
- WHAT **could be used otherwise?**
- OUTLOOK **and ideas**

Motivation

- DP-on-TD-algorithms can solve Model Counting and various combinatorial problems and are provable efficient at it
- Implementations of those are competing with modern solvers

- **But: those are fairly tedious to implement efficiently**

- Often bugs in the implementation
- Model counting is extremely space intensive (much more than SAT)

Contribution

This thesis created tdvisu as a tool that

- **integrates into existing implementations**
- **statically exports data from runs**
- **compiles simple DOT files and SVG graphics**

For further research it provides

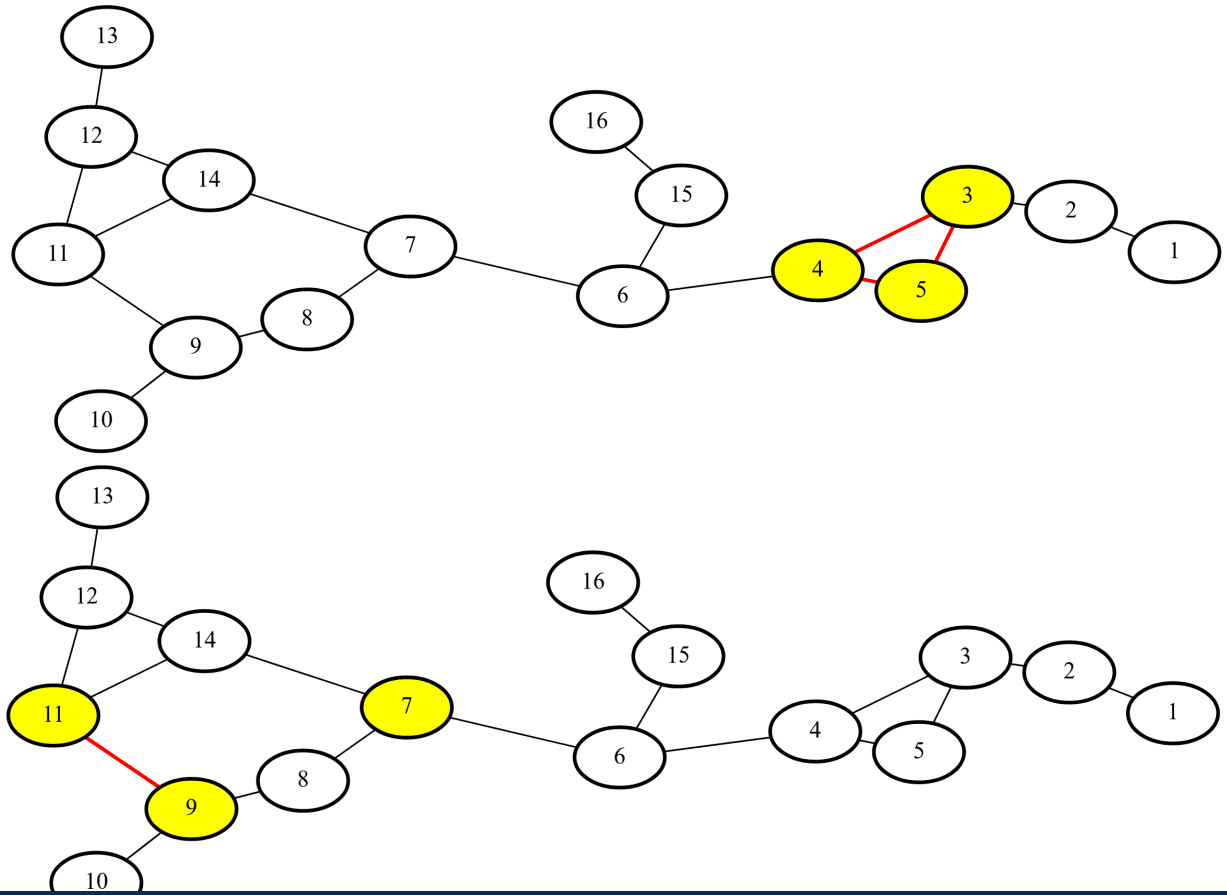
- **starting point for more complex investigations of**
  - **bug spotting**
  - **and fixing by using visualizations**

Background

The algorithms of interest solve problems of:

- **combinatorics (NP-problems)**
- **model-counting (#P-problems)**

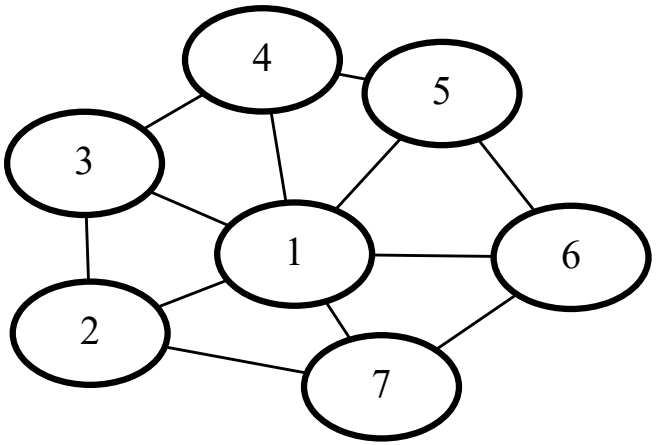
Modelcounting: Instead of **one** solution we want to count **all** solutions.

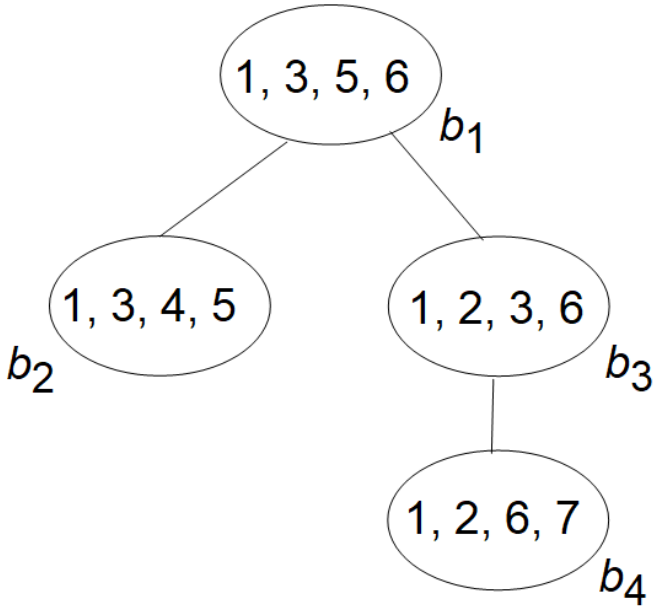


Tree Decomposition

Gives the DP algorithm a partial ordering for sub-problems.

1. **Each vertex must occur in some bag**
2. **For each edge, there is a bag containing both endpoints**
3. **Subgraph "restricted" to any vertex must be connected**





Width of a TD is: size of largest bag -1

width $=$ **3**

Graphs for Boolean Formulas

- **Example set of CNF-clauses:**

$\{c1 = \{v1, v3, \neg v4\}, c2 = \{\neg v1, v6\}, c3 = \{\neg v2, \neg v3, \neg v4\}, c4 = \{\neg v2, v6\}, c5 = \{\neg v3, \neg v4\}, c6 = \{\neg v3, v5\}, c7 = \{\neg v5, \neg v6\}, c8 = \{v5, v7\}\}$

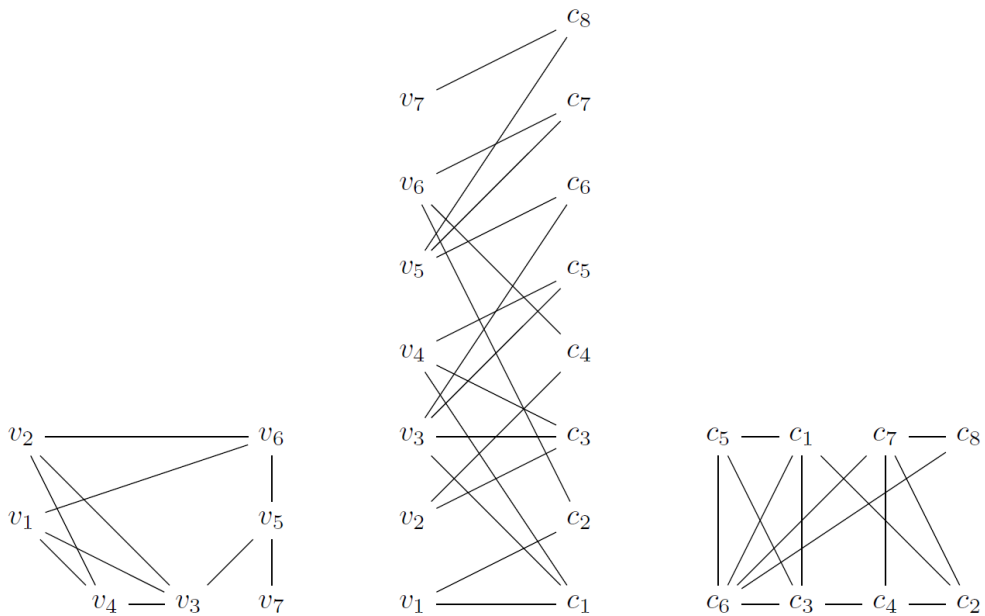
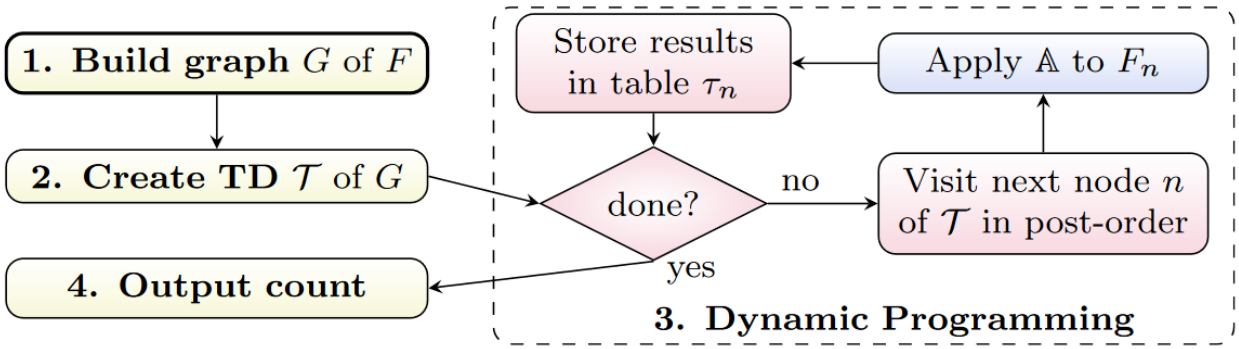


Figure: The primal (left), incidence (middle) and dual (right) graph

gpusat2 - Solving #SAT on GPU



- **Customized tree decompositions**
- **Adapted memory-management**
- **Improved precision handling**

[0] Images: Markus Zisser. *Solving the #SAT problem on the GPU with dynamic programming and OpenCL*. Technische Universität Wien, 2018.

dpdb

Database templates in Python
Generating SQL queries

1. **Create graph representation**
2. **Decompose graph**
3. **Solve sub-problems**
4. **Combine rows**

- #εTab#: `SELECT 1 AS cnt`
- #intrTab#: `SELECT 1 AS val UNION ALL 0`
- #localProbFilter#: $(l_{1,1}$ `OR ... OR` $l_{1,k_1})$ `AND ... AND` $(l_{n,1}$ `OR ... OR` $l_{n,k_n})$
- #aggrExp#: `SUM(cnt) AS cnt`
- #extProj#: τ_1`.cnt * ... *` τ_ℓ`.cnt AS cnt`

(a) Problem #SAT

- #εTab#: `SELECT 0 AS card`
- #intrTab#: `SELECT 1 AS val UNION ALL 0`
- #localProbFilter#: $([\![u_1]\!]$ `OR` $[\![v_1]\!])$ `AND ... AND` $([\![u_n]\!]$ `OR` $[\![v_n]\!])$
- #aggrExp#: `MIN(card) AS card`
- #extProj#: τ_1`.card + ... +` τ_ℓ`.card -` $(\Sigma_{i=1}^{\ell}|\chi(t_i) \cap \{a_1\}|$ `- 1) *` $\tau_1.[\![a_1]\!]$ `- ... -` $(\Sigma_{i=1}^{\ell}|\chi(t_i) \cap \{a_k\}|$ `- 1) *` $\tau_1.[\![a_k]\!]$

(b) Problem MinVC

- **SAT and #SAT**
- **#o-Coloring**
- **Vertex cover**

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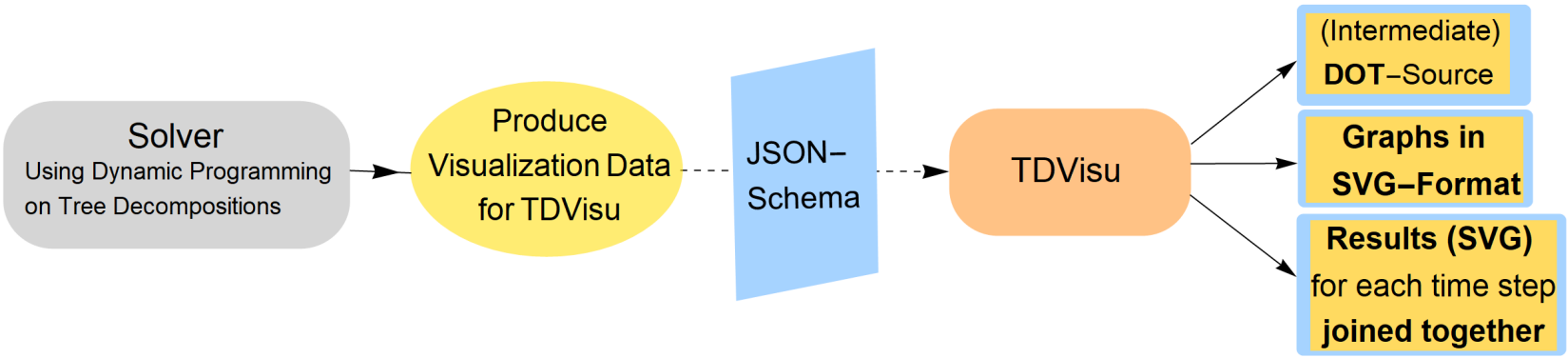
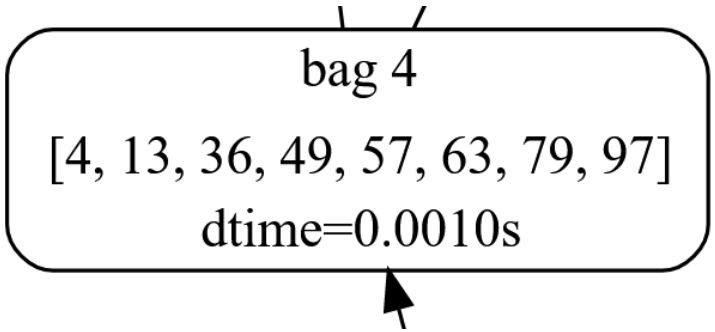
[0] "Exploiting Database Management Systems and Treewidth for Counting", Johannes Fichte et al. doi: 10.1007/978-3-030-39197-3_10.

Running tdvisu



TDVisu producing flexible and further processable formats



sol bag 4						
v4	v13	v49	v57	v63	v97	size
0	1	0	1	1	1	82
1	0	1	1	1	1	83
1	1	0	1	1	1	82
1	1	0	1	0	1	82
1	1	1	0	1	1	82
1	1	1	1	1	1	83
1	0	0	1	1	1	82
1	1	1	1	1	0	83
1	1	1	1	0	1	83
0	1	1	1	1	1	83
1	1	0	0	1	1	81
1	1	0	1	1	0	82
min-size: 81						

Challenge2

- Wie robust ist die Datenverarbeitung in der Visu Restrictions in strings for ids - Was Gedanken bei der Visu waren

Visualization in Action

MinVC example size 90 (expected 82)

Visualization in Action

1. **Inspect visualization**
2. **Verify findings in solver (in this case dpdb)**

public.p3_td_node_59/logicsem/postgres@PostgreSQL 12

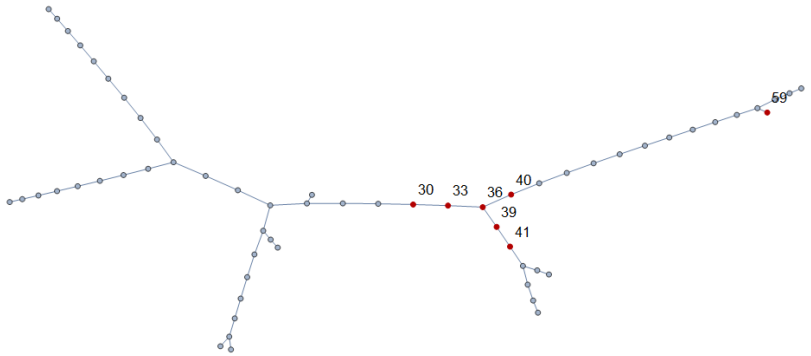
Query Editor | Query History

```
1  SELECT * FROM public.p3_td_node_59
2
```

Data Output | Explain | Messages | Notifications

	v16 boolean	v23 boolean	v40 boolean	v41 boolean	v52 boolean	v55 boolean	v60 boolean	v61 boolean	v66 boolean	v84 boolean	v99 boolean	v100 boolean	size integer
1	[null]	[null]	[null]	[null]	[null]	[null]	[null]	[null]	[null]	[null]	[null]	[null]	8

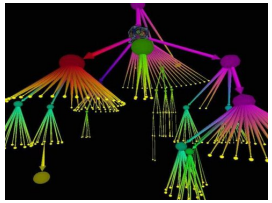
3. **Cross reference with standalone tree-decomposition**



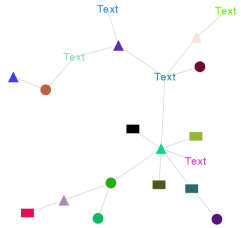
4. **Fix the root cause**

Related Work

- Fichte, Johannes & Hecher, Markus & Morak, Michael & Woltran, Stefan. (2018). Exploiting Treewidth for Projected Model Counting and Its Limits. 10.1007/978-3-319-94144-8_11.
- Hecher, Markus. (2020). Treewidth-aware Reductions of Normal ASP to SAT - Is Normal ASP Harder than SAT after All?. 485-495. 10.24963/kr.2020/49.
- Hecher M., Thier P., Woltran S. (2020)
  Taming High Treewidth with Abstraction, Nested Dynamic Programming, and Database Technology. In: Pulina L., Seidl M. (eds) Theory and Applications of Satisfiability Testing - SAT 2020. SAT 2020. Lecture Notes in Computer Science, vol 12178. Springer, Cham. https://doi.org/10.1007/978-3-030-51825-7_25



Gephi.org[1] Tulip [2]



[3] Vis.js Sigma.js vasturiano/3d-force-graph [4]

With the diverse / large node labels and special layout
the creation of a lightweight and customizable exchange format
took precedence over the integration into special layout software.

[1] https://gephi.org/ - Tool for data analysts and scientists keen to explore and understand graphs.
[2] tulip.labri.fr/TulipDrupal/ - Better Visualization Through Research.
[3] https://neo4j.com/developer/tools-graph-visualization/
[4] https://github.com/vasturiano/3d-force-graph

Outlook

Static $\rightarrow$ Dynamic

Interesting Questions :

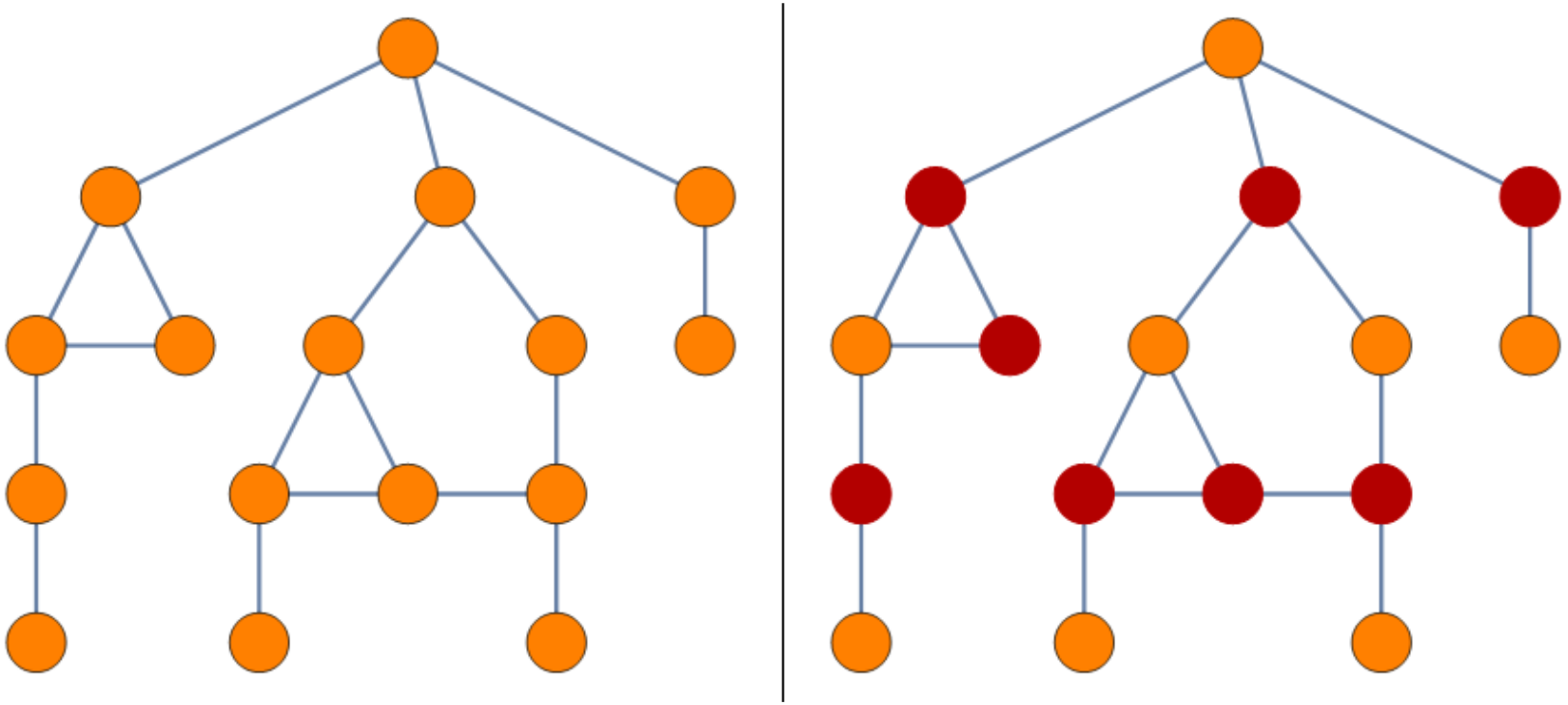
- **Utilizing graph databases for visualization and queries for debugging**
- **Enrich the visualization with debugging info for each node**
- **Cross reference the creation of rows in parent nodes**
- **For more advanced debugging tasks you may also need to revise the approach**

Final slide

Bibliography

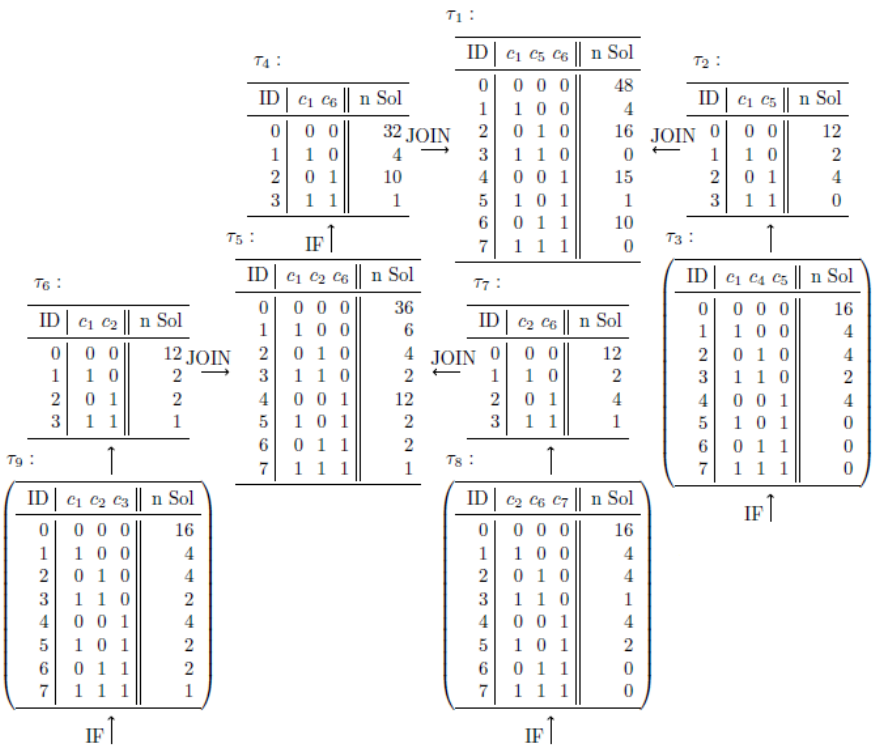
See the citations in the thesis.

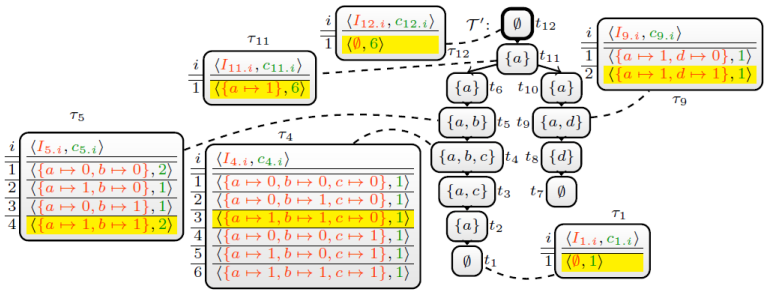
MinVC for example graph



Visualization

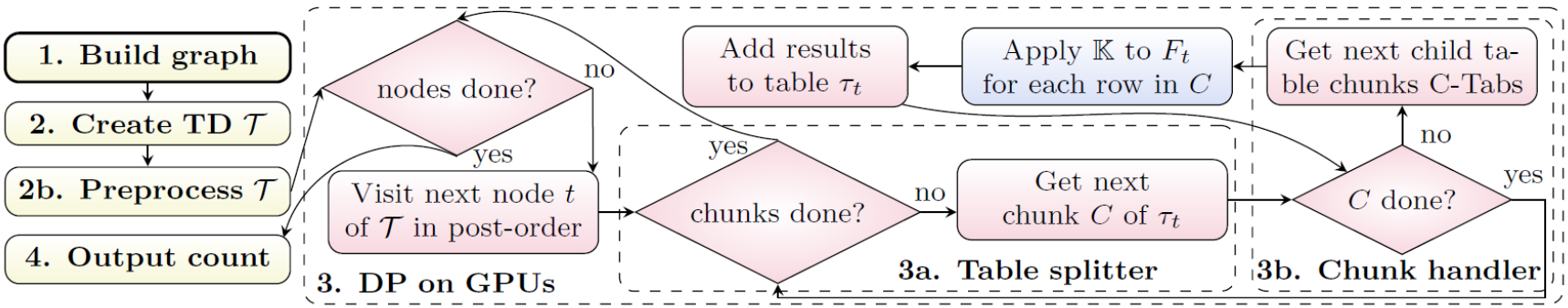
Manually for one run





[4] "Exploiting Database Management Systems and Treewidth for Counting", Johannes Fichte et al. doi: 10.1007/978-3-030-39197-3_10.

[4] "Solving #SAT on the GPU with Dynamic Programming and OpenCL", Diploma Markus Zisser 2018 Technische Universität Wien, p.33


1. Build graph
2. Create TD $\mathcal{T}$
2b. Preprocess $\mathcal{T}$
4. Output count
nodes done?
no
yes
Visit next node t of $\mathcal{T}$ in post-order
3. DP on GPUs
yes
chunks done?
no
Get next chunk C of τ_t
3a. Table splitter
Add results to table τ_t
Apply $\mathbb{K}$ to F_t for each row in C
Get next child table chunks C-Tabs
no
C done?
yes
3b. Chunk handler

Benchmark

Performance of all three programs on #SAT instances: