

Visualizing Dynamic Programming On Tree Decompositions

Martin Röbke Fakultät Informatik Technische Universität Dresden Germany

- ▶ WHAT is the motivation
- WHO benefits from visualization?
- CHALLENGES and solutions
- WHAT could be used otherwise?
- OUTLOOK and ideas



Motivation

- DP-on-TD-algorithms can solve various combinatorial problems like model counting
- Efficient (if the treewidth is small)
- ► Competitive with other modern solvers
- But tedious and hard to implement if done efficiently
- Often bugs in the implementation
- DP for model counting is extremely space demanding (much more than SAT)



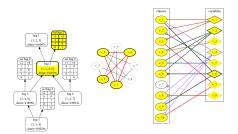
Ideas

- ► Inspect intermediate data during solving process
- ▶ Represent the input, tree decomposition and created solutions
- ▶ Lightweight but customizable output format



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Contribution

This thesis created tovisu as a tool that

- integrates into existing implementations
- statically exports data from runs
- compiles simple DOT files and SVG graphics

For further research it provides

- starting point for more complex investigations of
 - bug spotting
 - and fixing by using visualizations



Background

The algorithms of interest solve problems:

- NP-complete
- ▶ #P-complete problems instead of one solution we want to count all solutions

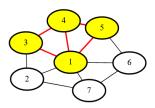


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Example of two snapshots of getting a minimal vertex cover via DP:



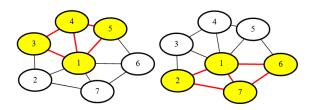


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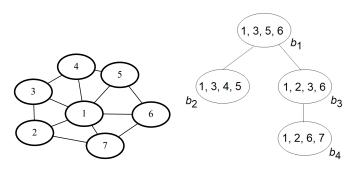
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Width of a TD is: size of largest bag -1 width = 3



Graphs for Boolean Formulas

► Example set of CNF-clauses:

$$\{c1 = \{v1, v3, \neg v4\}, c2 = \{\neg v1, v6\}, c3 = \{\neg v2, \neg v3, \neg v4\}, c4 = \{\neg v2, v6\}, c5 = \{\neg v3, \neg v4\}, c6 = \{\neg v3, v5\}, c7 = \{\neg v5, \neg v6\}, c8 = \{v5, v7\}\}$$

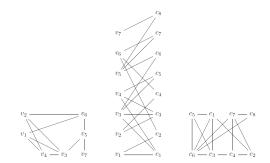
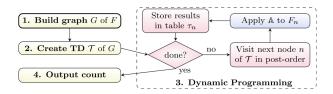


Figure: The primal (left), incidence (middle) and dual (right) graph



gpusat2 - Solving #SAT on GPU



Customized tree decompositions

Images: Markus Zisser. Solving the #SAT problem on the GPU with dynamic programming and OpenCL. Technische Universität Wien, 2018.



dpdb

Database templates in Python Generating SQL queries

- Create graph representation
- 2. Decompose graph
- 3. Solve sub-problems
- 4. Combine rows

```
 \begin{array}{ll} -\# \text{x-} \mathsf{Tab}\# : & \text{SELECT 1 AS cnt} \\ -\# \text{intr} \mathsf{Tab}\# : & \text{SELECT 1 AS val UNION ALL 0} \\ -\# \mathsf{plocalProbFilter}\# : (l_{1,1} \text{ OR } \ldots \text{ OR } l_{l,k_1}) \text{ AND } \ldots \text{ AND } (l_{n,1} \text{ OR } \ldots \text{ OR } l_{n,k_n}) \\ -\# \mathsf{aggrExp}\# : & \text{SUM}(\text{cnt}) \text{ AS cnt} \\ -\# \mathsf{extProj}\# : & \tau_1, \mathsf{cnt} * \ldots * \tau_\ell, \mathsf{cnt} \text{ AS cnt} \end{array}
```

(a) Problem #SAT

```
- #-Tab#: SELECT 0 AS card
- ##ortab#: SELECT 1 AS val WINDA ALL 0
- ##oraPob#iter#: ([n] 0.8 [n]) AND ... AND ([n] 0.8 [n])
- ##oraPob#iter#: ([n] 0.8 [n]) AND ... AND ([n] 0.8 [n])
- ##oraPob#: 71, card + ... + 72, card - Cf_{in}[\(\chi(n)\) ([n]) 1 0 7], [n]
- #foraphiter 1 71, card + ... + 72, card - Cf_{in}[\(\chi(n)\) ([n]) 1 0 7], [n]
- Cf_{in}[\(\chi(n)\) ([n]) 1 0 7], [n]
```

(b) Problem MinVC

- SAT and #SAT
- #o-Coloring
- Vertex cover

. . .

[&]quot;Exploiting Database Management Systems and Treewidth for Counting", Johannes Fichte et al., doi: 10.1007/978-3-030-39197-3-10.



TDVisu



Figure: TDVisu producing flexible and further processable formats



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TDVisu



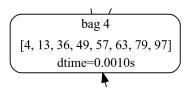
Figure: TDVisu producing flexible and further processable formats

JSON-Schema specifying:

timeline, tree decomposition incidence graph, general graph orientation, maximum lines, maximum columns, emphasis, svgjoin info

Graphics for TD

- Most cells are interpreted as strings.
- Extendable header and footer.





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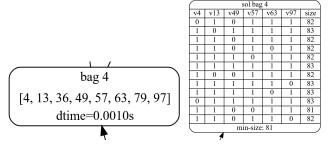


Figure: Bag node and solution node



Capabilities and Limitations

Integration of solvers via TDVisu.schema.json 1

Capabilities:

- Extracting basic (extendable) information (TD, solution nodes, order+time of processing...) from gpusat + dpdb
- Constructing and enriching the with solver information
- Adding multiple graphs for e.g. problems on Boolean formulas
- Providing a discrete timeline
- Parameters to control the layout and coloring of the data

Limitations:

- Can not further animate for example the origin of solutions
- Maneuvering in very large graphs is not very ergonomic with static content
- ▶ In the optimal case, a comprehensive test suite should be run prior to this.



Visualization in Action

MinVC example size 90 (expected 82)



1. Inspect visualization



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- 2. Verify findings in solver (in this case dpdb)

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3. Cross reference with standalone tree-decomposition



4. Fix the root cause



Related Work on the Algorithms

- J. Fichte, M. Hecher, M. Morak, S. Woltran (2018)
 "Exploiting Treewidth for Projected Model Counting and Its Limits."
- M. Hecher (2020)
 "Treewidth-aware Reductions of Normal ASP to SAT Is Normal ASP Harder than SAT after All?".
- M. Hecher, P. Thier, S. Woltran (2020)
 "Taming High Treewidth with Abstraction, Nested Dynamic Programming, and Database Technology." In: Pulina L., Seidl M. (eds) Theory and Applications of Satisfiability Testing - SAT 2020.



Related Work on Visualizations

- M.-C. Harre, Jelschen, Winter. "ELVIZ: A querybased approach to model visualization". (Jan. 2014)
- S. Diehl. "Software Visualization. Visualizing the Structure, Behaviour, and Evolution of Software." Springer, 2007.
- J. Daida et al. "Visualizing Tree Structures in Genetic Programming". (Mar. 2005)



Outlook

Static → Dynamic

Interesting Questions:

- Cross reference the creation of rows in parent nodes
- Enriching the visualization with more data for each node
- For more advanced debugging tasks you may also need to revise the approach
- Utilizing graph databases for debugging

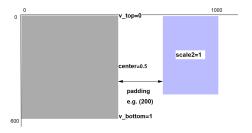


Final slide



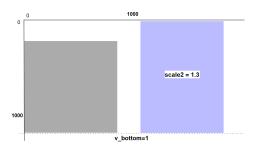
SVG-Join

▶ Joining single graphs for each time step

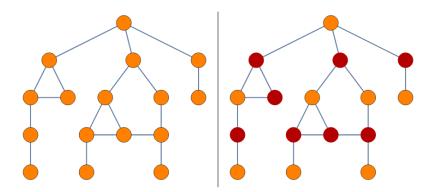


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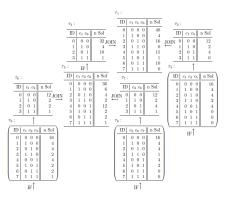
MinVC for example graph

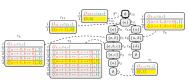




Visualization

Manually for one run





[&]quot;Exploiting Database Management Systems and Treewidth for Counting", Johannes Fichte et al. doi: 10.1007/978-3-030-39197-3-10.

[&]quot;Solving #SAT on the GPU with Dynamic Programming and OpenCL", Diploma Markus Zisser 2018 Technische Universität Wien, p.33











Gephi.org¹ Tulip ²



3 Vis.js



Sigma.js



vasturiano/3d-force-graph 4

With the diverse / large node labels and special layout the creation of a lightweight and customizable exchange format took precedence over the integration into special layout software.

https://gephi.org/ - Tool for data analysts and scientists keen to explore and understand graphs. tulip.labri.fr/TulipDrupal/ - Better Visualization Through Research.

https://neo4j.com/developer/tools-graph-visualization/



ELVIZ - Query based approach to software visualization

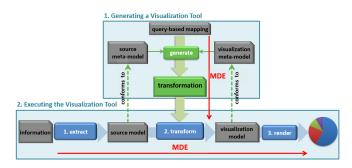


Figure: Overview of the ELVIZ-approach⁵

Fig 1 in: Marie-Christin Harre, J. Jelschen, A. Winter. "ELVIZ: A querybased approach to model visualization". In: Lecture Notes in Informatics (LNI), Proceedings - Series of the Gesellschaft fur Informatik (GI) (Jan. 2014), pp. 105–120.



Bibliography

See the citations in the thesis.

