

Bachelor Thesis

Visualizing Dynamic Programming on Tree Decompositions

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Abstract

This thesis introduces a practical and lightweight implementation for visualizing dynamic programming on tree decompositions. The provided source code is located in the python-package `tdvisu` for the purpose of visualizing, teaching and analyzing the dynamic solving process.

It defines a JSON-format specification for portability and customization of the visualization combined in one human-readable file and two reference implementations in actual solvers. The implementation currently does not support hyper-graphs and assumes that each node in the tree decomposition has either one or two children. The visualization output consists by default of scalable-vector-graphics SVG, a very flexible text-based that can be compressed and modified very easily without loss of quality. The images are split up into different views on the current state of the tree decomposition for consecutive user-defined time steps showing the progress of the dynamic programming. As illustrations of the possibilities for an application smaller examples from the problem-types "`#SAT`" and "minimal vertex cover" are presented, as well as an example of a faulty tree-decomposition that occurred during development. Intended audience:

- Developer of dynamic programming on tree decompositions for debugging.
- Researcher of such algorithms for comparisons and visualizations.
- Teachers or students looking for automatic visualization of their examples and the dynamic programming.

As two reference implementations of dynamic programming on tree decompositions we chose the projects `GPUSAT` and `dpdb`.

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1. Introduction

Graphs are increasingly interesting in scientific work, as the applications of interconnected datasets grow. The use cases outlined by Neo4j [1], a provider of open source database tools for "labeled property graphs" and include fields like

- Network and Database Infrastructure
- Recommendation Engines
- Artificial Intelligence and Analytics

While many problems can be expressed using the language of graphs, there is still plenty of room to utilize graph The idea for this project comes from my supervisor Dr. Johannes Fichte, who works on and with many projects such as the ones listed above on solving monadic second order logic (MSOL [2]) problems using highly parallelized architectures like graphics processing units or state of the art databases. One early implementation is published in [3] where for different real world examples the results looked promising These projects are very competitive ???REF???? for solving even large instances of those problems. The source code for TDVisu is available under GPL3 license.

Graphviz is a open source graph visualization software that provides customizable visualization for directed and undirected graphs. The information processed by graphviz intro. mit motivation und related work, state of the art, advancements.

Visualization Pipeline

Stand Umsetzung, Tools: Slack, Trello, GitHub, Presentations

2. Background

This chapter provides the reader with a brief background for this work.

We begin with a description on SAT and #SAT as examples for a very general problem that can be described with monadic second order logic (MSOL). Furthermore the general case of MSOL will be described, as well as the *DIMACS*-file-format used in the projects. The following section describes Tree Decompositions (TDs) which are the basis for our visualization. Finally we shortly discuss Courcelle's Theorem [2] as a related method of solving these problems.

2.1. Boolean satisfiability problem

https://en.wikipedia.org/wiki/Boolean_satisfiability_problem

SAT was the first known NP-complete problem, shown by Stephen Cook at the University of Toronto in 1971 [4]

literal \equiv boolean variable v or its negation

clause \equiv finite set of literals, interpreted as the disjunction

unit \equiv clause with $|c|=1$

CNF formula \equiv set of clauses, interpreted as their conjunction

var \equiv set of variables contained in the clause or clause set C

assignment $\equiv \alpha: \text{var}(C) \rightarrow \{0,1\}$

satisfiedclause \equiv if $\exists v \in \text{var}(c), v \in c$ and $\alpha(v)=1$ or $\neg v \in c$ and $\alpha(v)=0$. Otherwise falsified

satisfiedform \equiv each clause in the formula is satisfied by assignment

Connection to graphs with [5]

SAT Handbook: Even finding a single solution can be a challenge for such problems; counting the number of solutions is much harder. Not surprisingly, the largest formulas we can solve for the model counting problem with state-of-the-art model counters are orders of magnitude smaller than the formulas we can solve with the best SAT solvers. Generally speaking, current exact counting methods can tackle problems with a couple of hundred variables, while approximate counting methods push this to around 1,000 variables.

2.2. Monadic Second Order Logic

See also figure 2.

Explain graphs? Node, Edge

MSO graph properties are "fixed-parameter-tractable" with respect to clique-width and tree-width. <https://www.youtube.com/watch?v=hZI-wANH01w> 5th workshop on Graph Classes, Optimization, and Width Parameters (GROW 2011) 2011-10-28. MSO

counting (k-colorings) and optimizing (distance between two vertices...) functions. Interested in MSO logic over graphs.

Two types of MSO formulas *or logical graph representations*.

- MSO formulas
- MSO_2 formulas with edge quantification \equiv MSO formulas over incidence graphs
- $G = (\text{vertices}, \text{edges as binary relation})$
- $\text{INC}(G) = (\text{vertices and edges}, \text{Inc})$ for G undirected: $\text{Inc}(e, v)$ i.e. v is a vertex of edge e
- FPT for clique width
- FPT for tree-width

This can also be done for directed graphs!

Typical MSO_2 graph properties:

has a perfect matching has a Hamiltonian circuit spanning tree of degree ≤ 3

The expressions have the form: "There exists a **set of edges** that is..." can not be transferred into "set of vertices"

<https://youtu.be/Wyn3djrYg7c?t=1385> Bruno Courcelle: Recognizable sets of graphs: algebraic and logical aspects <https://library.cirm-math.fr/Record.htm?idlist=2&record=19276851124910940339> Recording during the thematic meeting: "Frontiers of reconnaissability" the April 29, 2014 at the Centre International de Rencontres Mathématiques (Marseille, France)

FPT for model checking: An algorithm is FPT if it takes time $f(k) \cdot n^c$ for some fixed function f and constant c . The size of the input is n . The value k is a parameter of the input. This algorithm is then usable for small values of k . usually tree-width and clique width.

2.3. DIMACS Format

Developed in 1993 at Rutgers University. DIMACS (the Center for Discrete Mathematics and Theoretical Computer Science)

Wolfram Language fully supports the DIMACS format for storing a single undirected graph. <https://reference.wolfram.com/language/ref/format/DIMACS.html>

DIMACS CNF: This format is used to define a Boolean expression, written in conjunctive normal form. <https://people.sc.fsu.edu/~jburkardt/data/cnf/cnf.html>

Other formats for WMC, different graphs...

Supported also in Maple <https://www.maplesoft.com/support/help/maple/view.aspx?path=Formats/CNF>

Examples in appendix

2.4. Tree Decomposition

[5]chapter 2.2 Tree decompositions were originally introduced by Robertson and Seymour [6] in 1984. A *tree decomposition* (TD) of a graph G is a pair (T, χ) . T is a tree and χ is a mapping which assigns each node $n \in V(T)$ a set $\chi(n) \subseteq V(G)$ called a *bag*. Then (T, χ) is a TD if the following conditions hold:

1. for each vertex $v(n) \in V(G)$ there is a node $n \in V(T)$ such that $v \in \chi(n)$
2. for each edge $(x, y) \in E(G)$ there is a node $n \in V(T)$ such that $x, y \in \chi(n)$
3. if $x, y, z \in V(T)$ and y lies on the path from x to z then $\chi(x) \cap \chi(z) \subseteq \chi(y)$.

The width $width(T)$ of a tree decomposition T is $\max_{n \in V(T)} (|\chi(n)|) - 1$. The tree width of a graph is the *minimal width* over all tree decompositions of the graph.

!!!Example (can take one from the visualizations) also in [7] page 169

2.5. Courcelle's Theorem

Every graph property definable in monadic second-order logic (MSO) is decidable in linear time on graphs of bounded tree-width.

Courcelle, Bruno (1990)¹

For all $k \in \mathbb{N}$ and MSO-sentences F is the decision problem for a given graph G , whether $G \models F$ is true, in time $2^{p(tw(G))} \cdot |G|$ with a polynom p decidable.

- *drawback*: still expensive ($2^{p(tw(G))}$, $2^{2^{(\#Q)}}$, large constants)

The workflow then looks like we see in figure 1.

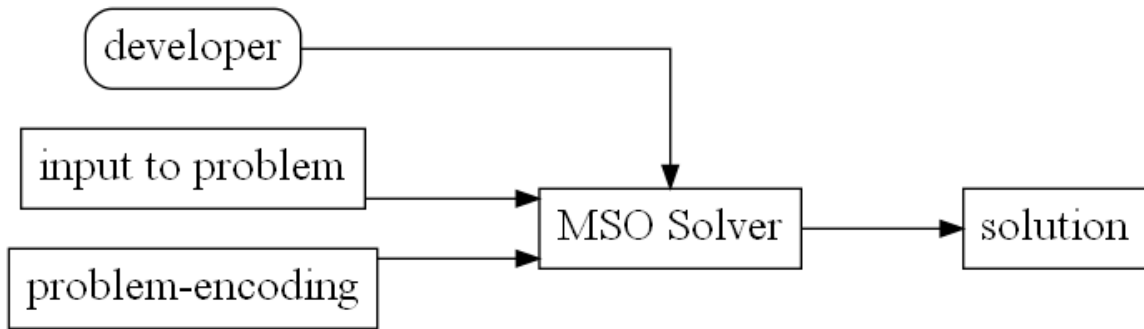


Figure 1: Implementation of the theorem

¹Courcelle, Bruno "The monadic second-order logic of graphs. I. Recognizable sets of finite graphs", Information and Computation, 85 (1990) no. 1: 12-75

3. Concept

We did setup a board with <https://trello.com> to track the progress and gather information connected to the project. Discussion and problems were mostly discussed with a dedicated <https://slack.com/intl/en-de/> in the *Collaborations Parameterized Algorithms and Complexity* team.

Research: language (python - explain) graph-construction (graphviz vs networkX), examples (diploma at first).

My previous experience connected to this work mainly stemmed from these two very good courses:

- Visualization with python from the lecture "Computational Physics" by Prof. Dr. A. Bäcker, chair of computational physics, TU Dresden 2016
- algorithms and various manipulations on graphs from the lecture "Graph Data Management and Analytics" by Hannes Voigt. [8]

The nodes in the dot-language are *labeled*, so creating a node takes one string identifier and can additionally be provided a string label. Valid examples for IDs include: a, b, A1, node1. The complete abstract grammar for DOT can be viewed at the DOT language.

It supports directed (*digraph* with edges indicated by '*-i*') an undirected (*graph* with edges indicated by '*-*') graphs. The visualizations presented here are constructed as undirected graphs, but would be easily extendable to directed representations since almost all operations keep the order of edge-endpoints given as input.

Another concept utilized were the sub-graphs and clusters available in DOT. To get a well structured (bipartite) incidence graph, each partition is placed in an individual cluster and sorted by node-label to easier find single nodes in potentially large clusters.

4. My Visualization Project

Pip, requirements files, virtual environments and how to pick quality Python libraries:
https://realpython.com/products/managing-python-dependencies/?utm_source=drip&utm_medium=email&utm_campaign=productfooter&utm_content=mpd

The first stages were in <https://github.com/VaeterchenFrost/gpusat-VISU> and the first releases of the source code outsourced to <https://github.com/VaeterchenFrost/tdvisu>

The objective of this project was/is to support the visualization mainly to document and improve the development efforts of dynamic programming on tree decompositions.

The tree decompositions in every tested application were provided by the utility <https://github.com/mabseher/htd> (small but efficient C++ library for computing (customized) tree and hypertree decompositions). Files / Classes / Methods Current perspective Checking with www.deepcode.ai

4.1. Integration in GPUSAT

To study and improve the handling of the C++ program was chronologically the first task I did experiment with. Getting the program up and running proved to be more difficult than we envisioned due to a probable bug in the driver when running OpenCL Drivers from CUDA on the Windows OS.

Impact on performance: Utilizing small classes and streams we tried to keep the impact on performance during the solving process low. However running on the same thread especially for larger problem instances

Some non-functional changes made to the source were

- a) some adjustments to satisfy the local compiler (shortening kernel string, replacing '*and*' with '&&')
- b) some explicit casts
- c) fixed some documentation of cmd line arguments

necessary The functional changes were:

1. allow tabs instead of spaces in input files
2. output more information about the hardware used (*device_query*)
3. add verbose output globally toggled by a flag
4. decide on https://en.wikipedia.org/wiki/DOT_%28graph_description_language%29 as the intermediate format for storing graphs.
5. start collecting information that might be needed for a visualization *graphfile* for saving the decomposition graph
6. add *graphout* to solver flow to get automated insight into the structure of a run.
7. add function *solutiontable* to extract tables of variable assignments as a string
8. add labels to each node (bags and solutions at this point)
9. encapsulating the previous functions into *gpusat::Graphoutput* class and instantiate it in main.
10. using the class functionality in the Solver
11. changing enum to the scoped enum class for better encapsulation and strongly typed.
12. with inlining *gpusatutils* the current functionality of creating raw dot was completed.

13. Experimented with neo4j [1], but found the visualization in particular not that presentable. The functionality to create cypher-queries for the SAT formula with primal, incidence and dual graph is still present in the class.
14. The functionality to create the cypher-query for the graph of the tree-decomposition was added.
15. Rename *visualisierung* into *visualization*
16. Using JsonCPP for processing and formatting json objects in C++
17. Setting format to BasedOnStyle: LLVM, UseTab: Never, IndentWidth: 4, TabWidth: 4, ColumnLimit: 0
18. Creating a *Grid* class for efficiently storing unsigned integer values in a two-dimensional structure based on ideas from this thread
19. create guide for remote development with *Visual Code*
20. Converted intermediate string operations to string-streams instead of files
21. Updated README.md
22. Added Doxygen for docbook, html, latex

Programm <https://github.com/VaeterchenFrost/GPUSAT>

Differences: <https://github.com/daajoe/GPUSAT/compare/master...VaeterchenFrost:master> Commits 142 Files changed 94 Nagoya talk: Graphs for performance are Ordered by used time per algorithm - gpusat quite good

Working with cmake remotly. ssh @(sg1.)dbai.tuwien.ac.at CPU branch wasn't working. Only AMD/Nvidia graphics with respective flags.

First steps: Connect ssh mkdir Ba cd Ba git clone <https://github.com/mabseher/htd.git> cd htd/ cmake . make -B -j1 make test (100pct tests passed, 0 tests failed out of 38 Total Test time (real) = 0.81 sec)

git clone <https://github.com/VaeterchenFrost/GPUSAT.git>

Configuration manual with the include options in CMakeLists.txt <https://github.com/VaeterchenFrost/GPUSAT/blob/master/CMakeLists.txt> or with help from <https://marketplace.visualstudio.com/items?itemName=ms-vscode.cmake-tools> to set up for the (potentially remote) environment.

Usage: ./gpusat [OPTIONS]

Options: -s,-seed INT number used to initialize the pseudorandom number generator -f,-formula TEXT path to the file containing the sat formula -d,-decomposition TEXT path to the file containing the tree decomposition -CPU run the solver on a cpu -NVIDIA run the solver on an NVIDIA device -AMD run the solver on an AMD device -weighted use weighted model count -noExp don't use extended exponents -v,-verbose print additional program information -p,-nopreprocess skips the preprocessing step for

debugging and visualisation-purposes -w,-combineWidth INT=20 maximum width to combine bags of the decomposition -g,-graph TEXT filename for saving the decomposition graph -visufile TEXT filename for saving the visualization file

4.1.1. Class Graphoutput

First steps in automatically visualizing the solving process with its tree decomposition. Outputs a .dot graph !!!EXAMPLE!!! with the bags and their solution nodes.

Additionally two functions to generate a Neo4J Cypher ***Link*** Query with:

- one graph representing the SAT formula and queries to construct incidence, dual and primal representations. output as satFile = "cypherSatFormula.txt"
- a graph representing the tree decomposition of the primal graph with it's bags containing variables. output as tdFile = "cypherTreedec.txt"

4.1.2. Class SolverVisualization

Umsetzung Special fork connecting the Visualization directly to the Solver. Utilizing the c++ library JsonCpp <https://github.com/open-source-parsers/jsoncpp> version 1.9.2

Beispiel

4.2. Integration in dpdb

Programm The integration of the API with dpdb was easier to implement after the solving process instead of hooking into the solving, provided that all necessary information got persisted in the database. The integration is included in the TDVisu project as a separate python file. This can be accomplished using the arguments

- `-store-formula` as a problem specific option for Sat and SharpSat
- `-gr-file GR_FILE` for problems like VertexCover with graph input

when calling `dpdb.py`

The integration for SharpSat was the first implemented in the file *construct_dpdb_visu.py* with the main parameters

- **problemnumber** the problem-id to select in the database
- `-twfile TWFILE` tw-file containing the edges of the graph
- `-outfile OUTFILE` file to write the output to, default `'dbjson%d.json'`
- `-pretty` pretty-print the JSON
- `-inter-nodes` Calculate and animate the shortest path between successive bags in the order of evaluation. To accomplish this task, an efficient implementation of the bidirectional Dijkstra's algorithm [9] based on the implementation by NetworkX [10] in `bidirectional_dijkstra`. in our case the weight function is a constant, as the edges have no weight associated with them.

Umsetzung Beispiel

5. Application and Images

6. Summary and Outline

What is achieved? What worked good, what bad?

References

- [1] “Graph database use cases and solutions.” <http://neo4j.com/use-cases>, Aug. 2016. [Online; accessed 11-June-2020].
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A. Images

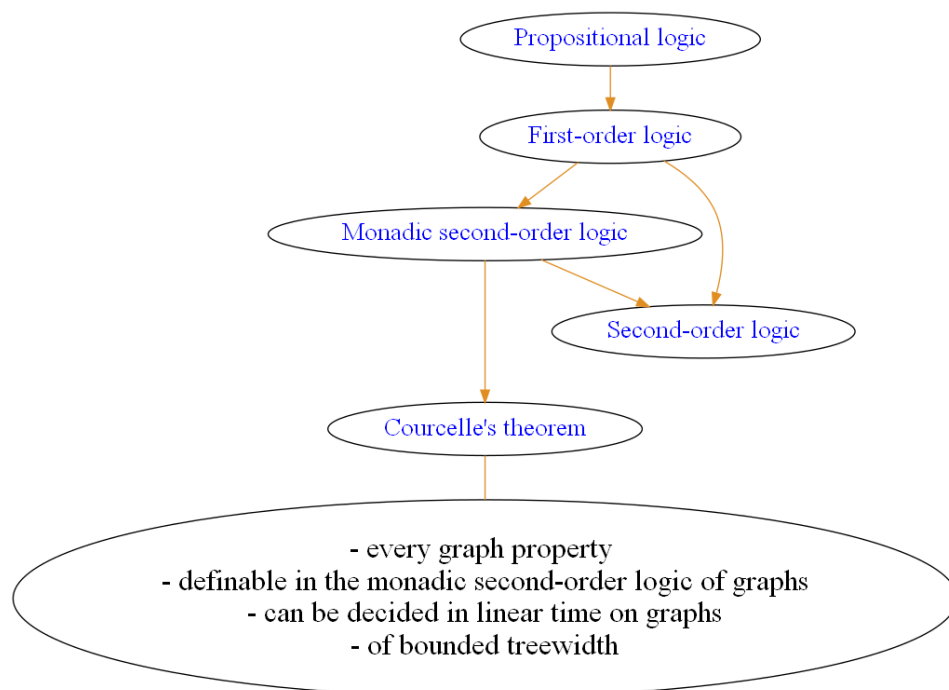


Figure 2: From propositional logic to monadic second order logic and Courcelle's Theorem