

Parser

Goian Tudor – 933/2

<https://github.com/VaruTudor/Formal-Languages-and-Compiler-Design>

The Grammar has 5 fields:

N – non-terminals

E – terminals

P – productions

S – start symbol

isCFG – flag which checks if the grammar is context-free

For the production we use a dictionary, where for each key (symbol) there is a list of tuples (symbol, destination). Most methods are for file parsing except checkCFG(rules, N) which receives the rules and the set of non-terminals and getProductionsFor(nonTerminal) which receives a non-terminal symbol.

Below there are some test examples.

```
N = { S, A, B, C, D }
E = { (, ), +, *, int }
S = S
P = {
    S -> A B C,
    A -> ( S ),
    B -> * S | E,
    C -> + A | C
}
```

```
N = { S, A, B, C, D }
E = { (, ), +, *, int }
productions
[(['A B C', 1)], [(['( S )', 2)], [(['* S', 3)], [(['E', 4)], [(['+ A', 5)], [(['C', 6)], None]
productions for A
[(['( S )', 2)]
CFG
```

```

closure(itemList):
    For each item in the item list, if after dot there is a non-terminal, add
    it to the list of items and repeat the process.
    pseudocode:
        repeat
            for any  $[A \rightarrow \alpha.B\beta]$  in  $C$  do
                for any  $B \rightarrow \gamma$  in  $P$  do
                    if  $[B \rightarrow \gamma] \notin C$  then
                         $C = C \cup [B \rightarrow \gamma]$ 
                    end if
                end for
            end for
        until  $C$  stops changing
    input:
        itemList: a list of LR(0) items
    :return:
        the resulting state

goto(state, symbol):
    For each LR0 item in the state move the dot if the symbol follows it.
    Perform closure on the modified item.
    pseudocode:
        goto( $s, X$ ) = closure( $\{[A \rightarrow \alpha X \beta] \mid [A \rightarrow \alpha X \beta] \in s\}$ )
    input:
        state: the state to be checked
        symbol: the symbol to be checked
    output:
        closure of the updated items

computeCanonicalCollection():
    For each state  $s$  in the canonical collection, for each symbol  $X$  (in both
    terminals and non-terminals), check if goto( $s, X$ ) result is not an empty list
    nor exists already in the canonical collection and if so add it to the
    canonical collection
    pseudocode:
        repeat
            for any  $s$  in  $C$  do
                for any  $X$  in  $N \cup \Sigma$  do
                    if goto( $s, X$ )  $\neq \emptyset$  and goto( $s, X$ ) not in  $C$  then
                         $C = C \cup \text{goto}(s, X)$ 
                    end if
                end for
            end for
        until  $C$  stops changing

computeTableActions():
    For each state in the canonical collection add in the LR0 Table it's
    appropriate action

```

```
def buildInputStack(sequence):  
    Iterate and put each symbol found in the input list.  
    input:  
        sequence: a string containing symbols among it  
    output:  
        a list of symbols  
  
def parseSequence(sequence):  
    Performs the parsing algorithm using 3 stacks (input, working and output) handling each type of  
action for a state: shift, reduce or accept. The workingStack is a list considered a stack -> (meaning the  
top of the stack is the right most element), and the inputStack is also a list considered a stack <-  
(meaning the first element is the top of the stack)  
    input:  
        a string of symbols to be parsed  
    output:  
        outputStack
```

Bellow there is a test example computing the Canonical Collection for a grammar.

1) $S' \rightarrow S$
 $S \rightarrow aA$
 $A \rightarrow bA \mid c$

Canonical collection:

$$s_0 = \text{closure}(\{[S' \rightarrow \cdot S]\}) = \{[S' \rightarrow \cdot S], [S \rightarrow \cdot aA]\}$$

$$s_1 = \text{goto}(s_0, S) = \text{closure}(\{[S' \rightarrow S \cdot]\}) = \{[S' \rightarrow S \cdot]\}$$

$$s_2 = \text{goto}(s_0, a) = \text{closure}(\{[S \rightarrow a \cdot A]\}) = \{[S \rightarrow a \cdot A], [A \rightarrow \cdot bA], [A \rightarrow \cdot c]\}$$

$$\text{goto}(s_1, X) = \emptyset, \quad \forall X \in N \cup \Sigma$$

$$s_3 = \text{goto}(s_2, A) = \text{closure}(\{[S \rightarrow aA \cdot]\}) = \{[S \rightarrow aA \cdot]\}$$

$$s_4 = \text{goto}(s_2, b) = \text{closure}(\{[A \rightarrow b \cdot A]\}) = \{[A \rightarrow b \cdot A], [A \rightarrow \cdot bA], [A \rightarrow \cdot c]\}$$

$$\text{goto}(s_3, X) = \emptyset, \quad \forall X \in N \cup \Sigma$$

$$s_5 = \text{goto}(s_4, A) = \text{closure}(\{[A \rightarrow bA \cdot]\}) = \{[A \rightarrow bA \cdot]\}$$

$$s_6 = \text{goto}(s_4, b) = \text{closure}(\{[A \rightarrow b \cdot A]\})$$

$$s_7 = \text{goto}(s_4, c) = \text{closure}(\{[A \rightarrow c \cdot]\}) = \{[A \rightarrow c \cdot]\}$$

$$\text{goto}(s_2, c) = s_7$$

Non-terminals and Terminals:

$N = \{ S, A \}$

$E = \{ a, b, c \}$

Productions:

$\{ 'S': [('aA', 1)], 'A': [('bA', 2), ('c', 3)] \}$

Productions for A:

$[('bA', 2), ('c', 3)]$

The grammar is CFG

Canonical Collection:

$S_0([S' \rightarrow \cdot S, S \rightarrow \cdot aA])$

$S_1([S' \rightarrow S \cdot])$

$S_2([S \rightarrow a \cdot A, A \rightarrow \cdot bA, A \rightarrow \cdot c])$

$S_3([S \rightarrow aA \cdot])$

$S_4([A \rightarrow b \cdot A, A \rightarrow \cdot bA, A \rightarrow \cdot c])$

$S_5([A \rightarrow c \cdot])$

$S_6([A \rightarrow bA \cdot])$

Bellow there is a test example for parsing a sequence

LR(0) table:

	Action	a	b	c	S	A
0	shift	2			1	
1	acc					
2	shift		4	6		3
3	r1					
4	shift		4	6		5
5	r2					
6	r3					

$S \rightarrow S$
 $S \rightarrow aA$
 $A \rightarrow bA^2$
 $A \rightarrow c^3$

$S \Rightarrow aA \Rightarrow abA \Rightarrow abbA \Rightarrow abbc$
 (1) (2) (2) (3)

Parsing

rw=abbc

$\$0$ $\$0a2$ $\$0a2b4$ $\$0a2b4b4$ $\$0a2b4b4c6$ $\$0a2b4b4A5$ $\$0a2b4A5$ $\$0a2A5$ $\$0S1$ $\$acc$	$abbc\$$ $bbc\$$ $bc\$$ $c\$$ $\$$ $\$$ $\$$ $\$$ $\$$	$\overline{11}$ 3 $2, 3$ $2, 2, 3$ $1, 2, 2, 3$
--	--	---

Final output: ['1', '2', '2', '3']