

# Mathematical and Logical Foundations of Computer Science

## Lecture 11 - Predicate Logic (Syntax)

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(some slides were adapted from Rajesh Chitnis' slides)

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## Where are we?

- ▶ Symbolic logic
- ▶ Propositional logic
- ▶ **Predicate logic**
- ▶ Constructive vs. Classical logic
- ▶ Type theory

# Today

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## Further reading:

- ▶ Chapter 7 of  
[http://leanprover.github.io/logic\\_and\\_proof/](http://leanprover.github.io/logic_and_proof/)

## Recap: Propositional Logic

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**Four connectives:**

- ▶  $P \wedge Q$ : we have a proof of both  $P$  and  $Q$
- ▶  $P \vee Q$ : we have a proof of at least one of  $P$  and  $Q$
- ▶  $P \rightarrow Q$ : if we have a proof of  $P$  then we have a proof of  $Q$
- ▶  $\neg P$ : stands for  $P \rightarrow \perp$

## Recap: Proofs

**Natural Deduction**

**Sequent Calculus**

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introduction/elimination rules

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$$\frac{\begin{array}{c} \overline{\phantom{A}}^1 \\ A \\ \vdots \\ B \end{array}}{A \rightarrow B}^1 [\rightarrow I]$$

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$$\frac{\Gamma, A \vdash B}{\Gamma \vdash A \rightarrow B} [\rightarrow R]$$

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- ▶ Predicate logic allows us to reason about members of a (non-empty) domain

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- ▶ Socrates is one member of this domain
- ▶ Predicates are “being a man” and “being mortal”



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**Another example:** consider a database with 3 tables

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- ▶  $\exists y. \exists z. Student(y, x) \wedge Module(z, Math) \wedge Enroll(y, z)$

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- ▶ **precedence**: lower than the other connectives

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Examples of formulas in predicate logic

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**Notation:** We sometimes write  $p^k$  when we want to indicate that the predicate symbol  $p$  has arity  $k$

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The syntax of predicate logic is defined by the following grammar:

$$t ::= x \mid f(t, \dots, t)$$

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The scope of a quantifier extends as far right as possible. E.g.,  $P \wedge \forall x.p(x) \vee q(x)$  is read as  $P \wedge \forall x.(p(x) \vee q(x))$

# Examples

Consider the following domain and signature:

- ▶ Domain:  $\mathbb{N}$
- ▶ Functions:  $0, 1, 2, \dots$  (arity 0);  $+$  (arity 2)
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$$\forall x. \text{prime}(x) \rightarrow x = 2 \vee \text{odd}(x)$$
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$$\forall x. \text{even}(x) \rightarrow \exists y. \exists z. \text{prime}(y) \wedge \text{prime}(z) \wedge x = y + z$$
- ▶ There is no number greater than all numbers.  
$$\neg \exists x. \forall y. x > y$$
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$$\forall x. \exists y. y > x$$

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# Conclusion

## What did we cover today?

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## Next time?

- ▶ Predicate logic (Natural Deduction)