Exercise Sheet 1 - Solutions Symbolic Logic

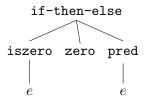
1. • The BNF for this language is defined as follows:

$$e ::= \mathsf{zero} \mid \mathsf{succ}(e) \mid \mathsf{pred}(e) \mid \mathsf{iszero}(e) \mid \mathsf{if}\ e\ \mathsf{then}\ e\ \mathsf{else}\ e \mid \mathsf{true} \mid \mathsf{false}$$

- No expression is ambiguous using this grammar.
- The expression that given an expression e, checks whether e is zero, and if it is returns zero, else returns e's predecessor, can be written as:

$$\verb|if iszero|(e) then zero else pred|(e)$$

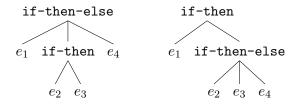
Its parse tree is



2. • The BNF for this extended language is defined as follows:

$$e := \mathsf{zero} \mid \mathsf{succ}(e) \mid \mathsf{pred}(e) \mid \mathsf{iszero}(e) \mid \mathsf{if} \ e \ \mathsf{then} \ e \ \mathsf{else} \ e \mid \mathsf{true} \mid \mathsf{false} \mid \mathsf{if} \ e \ \mathsf{then} \ e$$

- It now has ambiguities because expressions of the form (if e_1 then if e_2 then e_3 else e_4) could be (if e_1 then (if e_2 then e_3) else e_4) or (if e_1 then (if e_2 then e_3 else e_4)).
- The parse trees corresponding to the two above expressions are:



3. • Our BNF is now:

$$e$$
 ::= zero | succ (e) | pred (e) | iszero (e) | if e then e else e | true | false | if e then e eq ::= $e = e$

• An axiom schema capturing that "the expression that given an expression e, checks whether e is zero, and if it is returns zero, else returns e's predecessor" is equal to "e's predecessor" is

if
$$iszero(e)$$
 then $zero$ else $pred(e) = pred(e)$

This axiom schema makes use of one metavariable: e.

- Here are two instances of this axiom:
 - if iszero(zero) then zero else pred(zero) = pred(zero)
 - if iszero(succ(zero)) then zero else pred(succ(zero)) = pred(succ(zero))