

Image from: <http://www.kirrk.com/modularity/wp-content/uploads/2009/12/EncapsulatingDesign1.jpg>

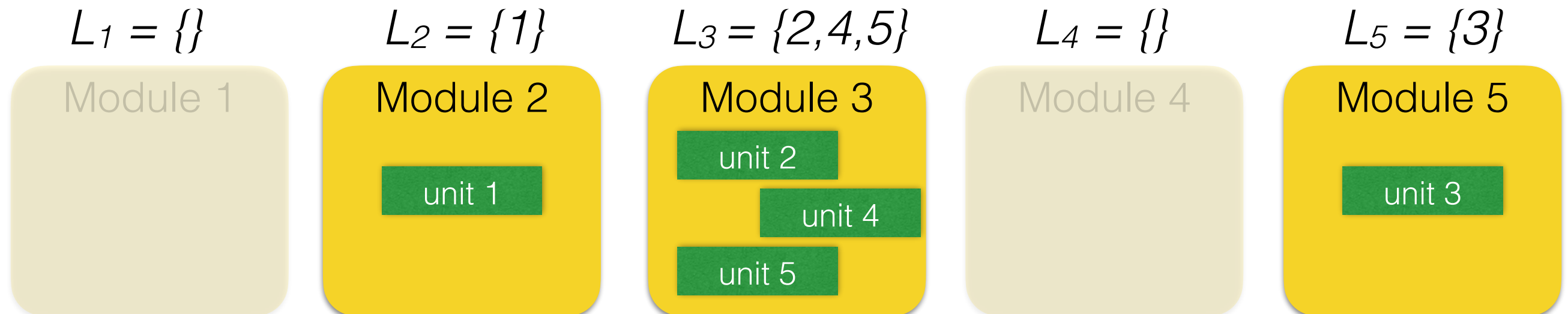
Example of Hill Climbing Application: Software Module Clustering (Algorithmic Design)

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Design Variable

Design variable: allocation of units into modules.

- Consider that we have N units, identified by natural numbers in $\{1, 2, \dots, N\}$.
- This means that we have at most N modules.
- Our design variable is a list L of N modules, where each module L_i , $i \in \{1, 2, \dots, N\}$, is a set containing a minimum of 0 and a maximum of N units.



Constraints and Objective Function

Constraints: N/A

Objective function: quality of modularisation (to be maximised).

$$\begin{array}{l} \text{Quality}(L) = \\ \text{(maximise)} \end{array} \sum_{\substack{i \in \{1,2,\dots,M\} \\ L_i \neq \{\}}} \text{Quality}(L_i)$$

$$\begin{array}{l} \text{Quality}(L_i) = \\ \text{(maximise)} \end{array} \frac{\# \text{IntraEdges}_i}{\# \text{IntraEdges}_i + 1/2 * \# \text{InterEdges}_i}$$

Problem Formulation

Hill-Climbing (assuming maximisation)

1. `current_solution` = generate initial solution randomly
2. Repeat:
 - 2.1 generate neighbour solutions (differ from current solution by a single element)
 - 2.2 `best_neighbour` = get highest quality neighbour of `current_solution`
 - 2.3 If `quality(best_neighbour) <= quality(current_solution)`
 - 2.3.1 Return `current_solution`
 - 2.4 `current_solution` = `best_neighbour`

Until a maximum number of iterations

Design variable —>
what is a candidate solution for us?

Objective —>
what is quality for us?

Are there any constraints that
need to be satisfied?

Designing Representation, Initialisation and Neighbourhood Operators

Hill-Climbing (assuming maximisation)

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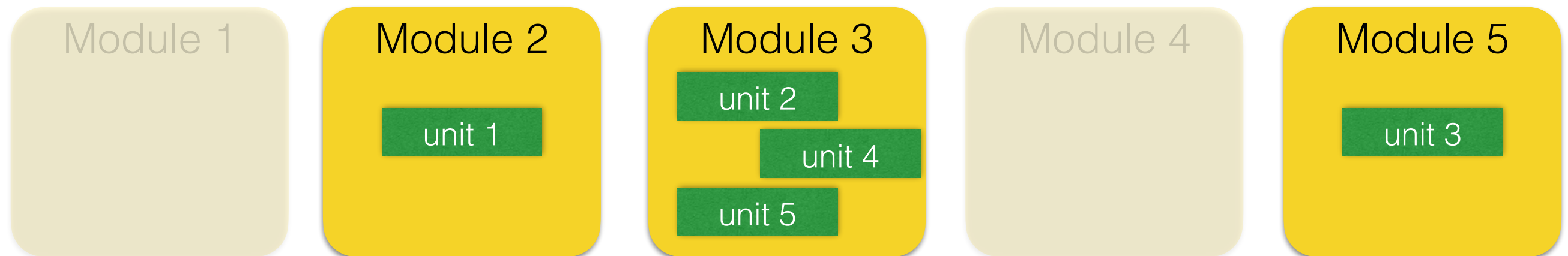
Until a maximum number of iterations

- Representation:
 - How to store the design variable.
 - E.g., boolean, integer or float variable or array.
- Initialisation:
 - Usually involve randomness.
- Neighbourhood operator:
 - How to generate neighbour solutions.

Representation

How to represent the design variable internally in the implementation?

- E.g., list of N modules, where each module is a list of integers in $\{1, 2, \dots, N\}$ identifying the existing units.

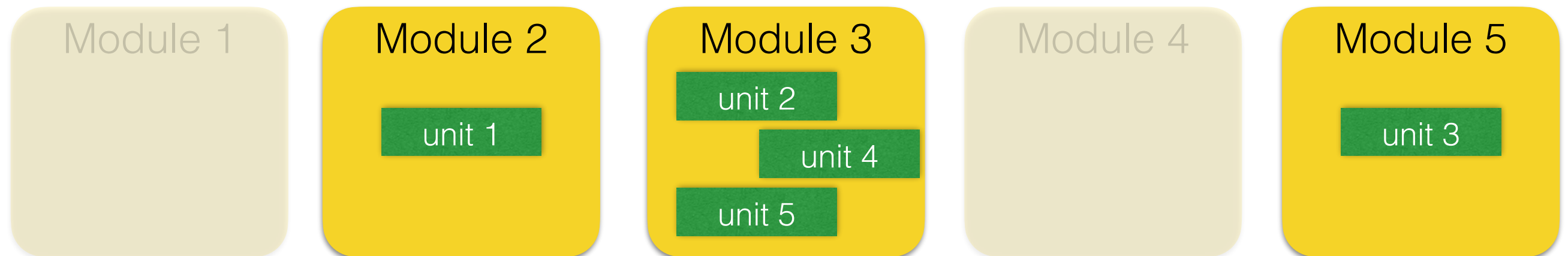


- E.g., if we have $N=5$, a possible allocation is $L = \{\{\}, \{1\}, \{2, 4, 5\}, \{\}, \{3\}\}$.

Representation

How to represent the design variable internally in the implementation?

- E.g., matrix $A_{N \times N}$, where $A_{ij} = 1$ if unit j is in module i , and 0 otherwise.



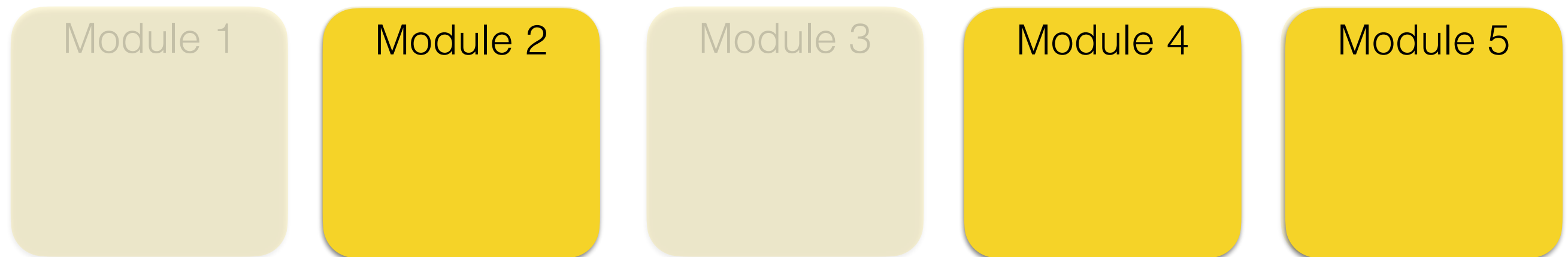
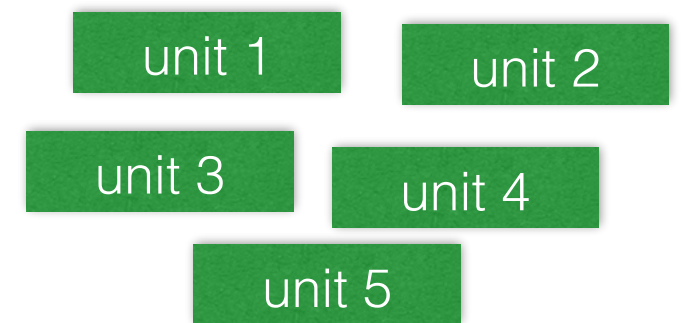
$$A = \begin{bmatrix} 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 \end{bmatrix}$$

Initialisation

E.g.: place each unit into a randomly picked module.

For each unit $u \in \{1, \dots, N\}$

Add u to a module L_i , where $i \sim U\{1, N\}$

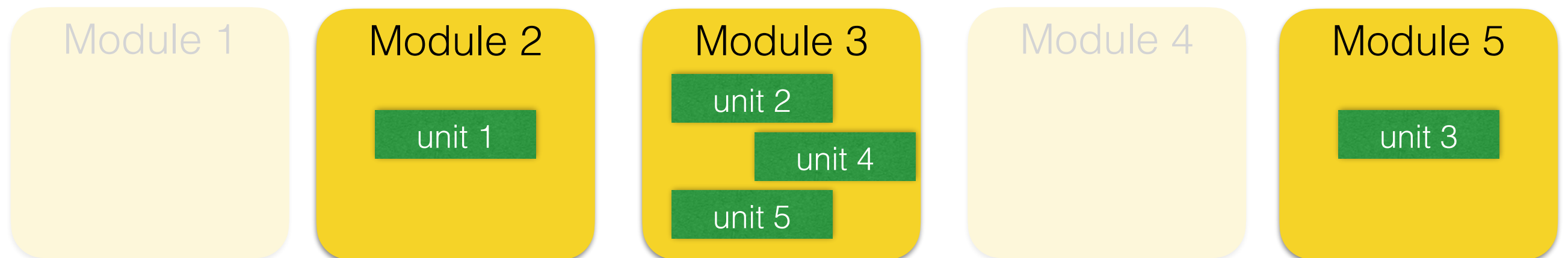


$$L = \{\{\}, \{1, 3\}, \{2\}, \{4\}, \{5\}\}$$

Neighbourhood Operator

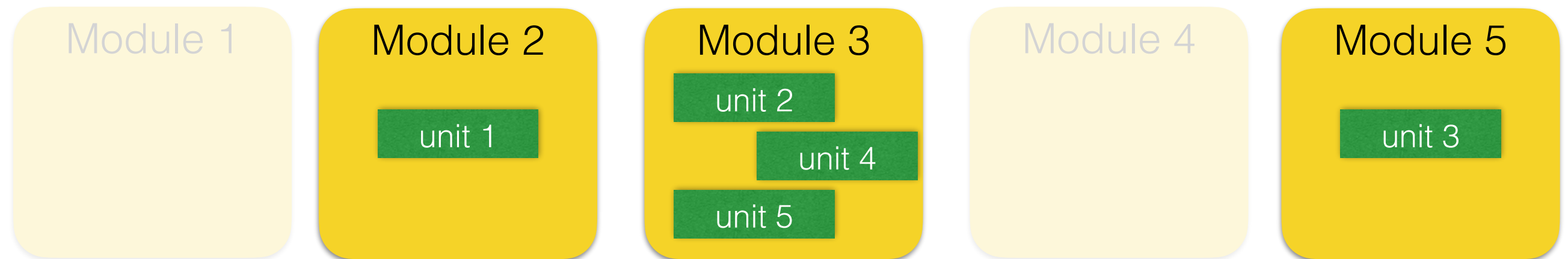
- What would be a possible neighbourhood operator for the software clustering problem?
- A neighbour in the software module clustering problem would be a solution where a single unit moves from one module to another. E.g.:

$$L = \{\{\}, \{1\}, \{2,4,5\}, \{\}, \{3\}\} \longrightarrow L = \{\{\}, \{1,5\}, \{2,4\}, \{\}, \{3\}\}$$



Neighbourhood

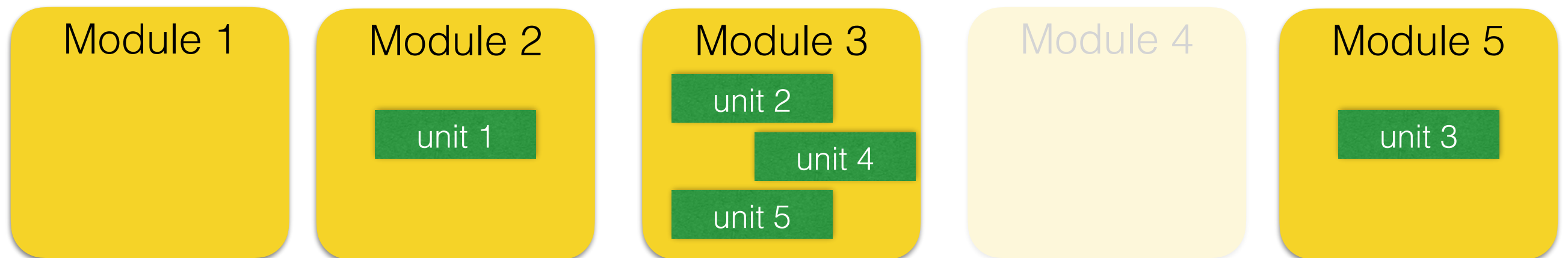
- Real world problems will frequently have more than two neighbours for each candidate solution.
- How many neighbours do we have for the candidate solution below, if we allow for equivalent neighbours?



5 units * 4 possible modules to move to = 20

Neighbourhood

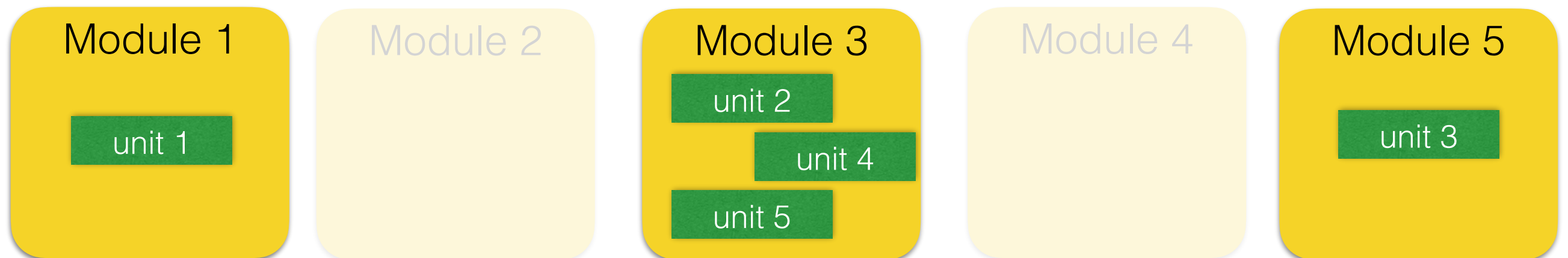
- How many neighbours do we have for the candidate solution below, if we allow for equivalent neighbours?



Some neighbours will be equivalent.
Duplicates could be eliminated.

Neighbourhood

- How many neighbours do we have for the candidate solution below, if we allow for equivalent neighbours?



For $i \in \{1, \dots, N\}$ // module

For $j \in \{1, \dots, \text{size}(L_i)\}$ // unit within module

For $i' \in \{1, \dots, N\} \setminus i$ // another module

$L' = \text{clone of } L$

Move unit L'_{ij} to module $L'_{i'}$

Yield L' as a neighbour

Hill Climbing

Hill-Climbing (assuming maximisation)

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- Until a maximum number of iterations

Simulated Annealing would also require a representation, initialisation procedure, and neighbourhood operator to solve a problem.

Summary

- Software Module Clustering problem formulation.
- Representation, initialisation and neighbourhood operators.

Next

- Application of Simulated Annealing.