

# Mathematical and Logical Foundations of Computer Science

## Week 2 - Additional Examples

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# Propositions

Upper case letters stand for variables, while lower case letters stand for atomic propositions. We will often use variables rather than specific atomic propositions.

Here are examples of propositions:

- ▶  $P \rightarrow Q \rightarrow R$ 
  - ▶ it stands for  $P \rightarrow (Q \rightarrow R)$
  - ▶ because  $\rightarrow$  is right-associative
  - ▶ this is different from  $(P \rightarrow Q) \rightarrow R$
  - ▶ it uses 2 variables  $P$  and  $Q$
- ▶  $p \rightarrow q \rightarrow r$ 
  - ▶ it stands for  $p \rightarrow (q \rightarrow r)$
  - ▶ because  $\rightarrow$  is right-associative
  - ▶ this is different from  $(p \rightarrow q) \rightarrow r$
  - ▶ it uses 2 atoms  $p$  and  $q$

# Propositions

Here are further examples of propositions:

- ▶  $P \wedge Q \rightarrow P \vee Q$ 
  - ▶ it stands for  $(P \wedge Q) \rightarrow (P \vee Q)$
  - ▶ because  $\wedge$  and  $\vee$  have precedence over  $\rightarrow$
  - ▶ this is different from  $P \wedge (Q \rightarrow P \vee Q)$
  - ▶ it uses 2 variables  $P$  and  $Q$
- ▶  $\neg\neg P \rightarrow P$ 
  - ▶ it stands for  $(\neg(\neg P)) \rightarrow P$
  - ▶ because  $\neg$  has precedence over  $\rightarrow$
  - ▶ this is different from  $\neg\neg(P \rightarrow P)$
  - ▶ it uses 1 variable  $P$

# Propositions

Those are not propositions:

- ▶  $P \wedge \wedge Q$ 
  - ▶ this is not derivable from the syntax of propositional logic
  - ▶ it is not possible to have two  $\wedge$  next to each other
- ▶  $P \rightarrow \wedge Q$ 
  - ▶ this is not derivable from the syntax of propositional logic
  - ▶ it is not possible to have a  $\wedge$  next to a  $\rightarrow$

## Natural Deduction Proofs

Here is a Natural Deduction proof of  $A \vee B \vdash C \rightarrow (A \vee B) \wedge C$

$$\frac{\frac{A \vee B \quad \overline{C}^1}{(A \vee B) \wedge C} [\wedge I]}{C \rightarrow (A \vee B) \wedge C}^1 [\rightarrow I]$$

Note that  $C \rightarrow (A \vee B) \wedge C$  is read as  $C \rightarrow ((A \vee B) \wedge C)$

# Natural Deduction Proofs

Here is a Natural Deduction proof of  $Q \vdash (Q \rightarrow R) \rightarrow R$

$$\frac{\frac{\overline{Q \rightarrow R}^1 \quad Q}{R} [\rightarrow E]}{(Q \rightarrow R) \rightarrow R}^1 [\rightarrow I]$$

# Natural Deduction Proofs

Here is a Natural Deduction proof of  $\vdash \neg(A \wedge B) \rightarrow A \rightarrow \neg B$

$$\frac{\frac{\frac{}{\neg(A \wedge B)} 1 \quad \frac{\frac{A}{\neg A} 2 \quad \frac{B}{\neg B} 3}{A \wedge B} [\wedge I]}{\perp} [\neg E] \quad \frac{}{\neg B} 3 [\neg I]}{A \rightarrow \neg B} 2 [\rightarrow I] \quad \frac{}{\neg(A \wedge B) \rightarrow A \rightarrow \neg B} 1 [\rightarrow I]$$

Note that  $\neg(A \wedge B) \rightarrow A \rightarrow \neg B$  is read as  $\neg(A \wedge B) \rightarrow (A \rightarrow \neg B)$