Mathematical and Logical Foundations of Computer Science

Lecture 5b - Propositional Logic (Natural Deduction & Sequent Calculus)

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(some slides were adapted from Rajesh Chitnis' slides)

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Where are we?

- Symbolic logic
- ► Propositional logic
- Predicate logic
- ► Constructive vs. Classical logic
- Type theory

Today

- Sequent Calculus vs. Natural Deduction
- Sequent Calculus proofs
- Natural Deduction proofs

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- Sequent Calculus proofs
- Natural Deduction proofs

Further reading

- Section 5 in "Proof and Types" https://www.paultaylor.eu/stable/prot.pdf
- ► Chapter 3 of http://leanprover.github.io/logic_and_proof/

Syntax:

$$P ::= a \mid P \wedge P \mid P \vee P \mid P \to P \mid \neg P$$

Syntax:

$$P ::= a \mid P \land P \mid P \lor P \mid P \to P \mid \neg P$$

Lower-case letters are atoms: p, q, r, etc.

Upper-case letters stand for any proposition: P, Q, R, etc.

Syntax:

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Two special atoms:

- ▶ T which stands for True
- which stands for False

Syntax:

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Two special atoms:

- ▶ T which stands for True
- which stands for False

We also introduced four connectives:

- $P \wedge Q$: we have a proof of both P and Q
- $P \vee Q$: we have a proof of at least one of P and Q
- ▶ $P \rightarrow Q$: if we have a proof of P then we have a proof of Q
- ▶ $\neg P$: stands for $P \rightarrow \bot$

How would you express these sentences in propositional logic?

• "if x > 2 then x > 1"

• "if x > 2 and x is even then x > 3"

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- "if x > 2 then x > 1"
 - atom p: "x > 2"
 - atom q: "x > 1"
 - ▶ proposition: $p \rightarrow q$
- "if x > 2 and x is even then x > 3"

How would you express these sentences in propositional logic?

- "if x > 2 then x > 1"
 - atom p: "x > 2"
 - atom q: "x > 1"
 - ▶ proposition: $p \rightarrow q$
- "if x > 2 and x is even then x > 3"
 - atom p: "x > 2"
 - ▶ atom q: "x is even"
 - ▶ atom r: "x > 3
 - proposition: $(p \land q) \rightarrow r$
 - we don't need parentheses, and can just write: $p \wedge q \rightarrow r$

Recap: Natural deduction vs. Sequent Calculus

2 deduction systems for propositional logic (don't mix their rules!)

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Natural Deduction

- "natural" style of constructing a proof
- start with the given premises
- repeatedly apply the given inference rules
- until you obtain the conclusion
- Can work both forwards and backwards
- "natural" doesn't mean there is a unique proof

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Sequent Calculus

- hypotheses are made explicit in a context
- instead of deriving proposition, we derive sequents
 - a sequent is of the form $\Gamma \vdash P$
 - where the environment/context Γ is a list of propositions
 - ▶ and *P* is a proposition
 - lacktriangle intuitively: P is true assuming that the formulas in Γ are true
- we typically go backward

Recap: Natural Deduction

Natural Deduction rules:

$$\frac{A}{A} \stackrel{\vdots}{\vdots} \qquad \frac{A}{B} \stackrel{\vdots}{} \qquad \frac{A}{B} \stackrel{\vdots}{} \qquad \frac{A}{B} \stackrel{\vdots}{} \qquad \frac{A \to B \quad A}{B} \quad [\to E]$$

$$\frac{A}{A} \stackrel{1}{\vdots} \qquad \frac{A}{A} \stackrel{1}{} \stackrel{\vdots}{\vdots} \qquad \frac{A}{A} \stackrel{1}{} \stackrel{1}{} \stackrel{\vdots}{} \qquad \frac{A}{A} \stackrel{1}{} \stackrel{1}{} \stackrel{1}{} \qquad \frac{A}{B} \stackrel{1}{} \stackrel{1}{} \stackrel{1}{} \qquad \frac{A}{B} \stackrel{1}{} \stackrel{1}{} \stackrel{1}{} \qquad \frac{A \to B \quad A}{B} \quad [\to E]$$

$$\frac{A}{A \lor B} \stackrel{1}{} \stackrel{1}{} \stackrel{1}{} \stackrel{1}{} \qquad \frac{A}{B} \stackrel{1}{} \stackrel{1}{} \stackrel{1}{} \qquad \frac{A \lor B \quad A \to C \quad B \to C}{C} \quad [\lor E]$$

$$\frac{A}{A \lor B} \stackrel{1}{} \stackrel{1}{} \stackrel{1}{} \stackrel{1}{} \qquad \frac{A \land B}{B} \quad [\land E_R] \qquad \frac{A \land B}{A} \quad [\land E_L]$$

Recap: Sequent Calculus

Sequence Calculus rules:

$$\begin{array}{lll} \frac{\Gamma \vdash A & \Gamma, B \vdash C}{\Gamma, A \to B \vdash C} & [\to L] & \frac{\Gamma, A \vdash B}{\Gamma \vdash A \to B} & [\to R] \\ \\ \frac{\Gamma \vdash A}{\Gamma, \neg A \vdash B} & [\neg L] & \frac{\Gamma, A \vdash \bot}{\Gamma \vdash \neg A} & [\neg R] \\ \\ \frac{\Gamma, A \vdash C & \Gamma, B \vdash C}{\Gamma, A \lor B \vdash C} & [\lor L] & \frac{\Gamma \vdash A}{\Gamma \vdash A \lor B} & [\lor R_1] & \frac{\Gamma \vdash A}{\Gamma \vdash B \lor A} & [\lor R_2] \\ \\ \frac{\Gamma, A, B \vdash C}{\Gamma, A \land B \vdash C} & [\land L] & \frac{\Gamma \vdash A & \Gamma \vdash B}{\Gamma \vdash A \land B} & [\land R] \\ \\ \frac{\Gamma}{A \vdash A} & [Id] & \frac{\Gamma \vdash B & \Gamma, B \vdash A}{\Gamma \vdash A} & [Cut] \\ \\ \frac{\Gamma, B, A, \Delta \vdash C}{\Gamma, A, B, \Delta \vdash C} & [X] & \frac{\Gamma \vdash B}{\Gamma, A \vdash B} & [W] & \frac{\Gamma, A, A \vdash B}{\Gamma, A \vdash B} & [C] \\ \end{array}$$

Recap: Sequent Calculus

In addition we allow using the following derived rules:

$$\begin{array}{ll} \frac{\Gamma_{1},\Gamma_{2}\vdash A & \Gamma_{1},B,\Gamma_{2}\vdash C}{\Gamma_{1},A\to B,\Gamma_{2}\vdash C} & [\to L] & \frac{\Gamma_{1},\Gamma_{2}\vdash A}{\Gamma_{1},\neg A,\Gamma_{2}\vdash B} & [\neg L] \\ \\ \frac{\Gamma_{1},A,\Gamma_{2}\vdash C & \Gamma_{1},B,\Gamma_{2}\vdash C}{\Gamma_{1},A\vee B,\Gamma_{2}\vdash C} & [\lor L] & \frac{\Gamma_{1},A,B,\Gamma_{2}\vdash C}{\Gamma_{1},A\wedge B,\Gamma_{2}\vdash C} & [\land L] \\ \\ \frac{\Gamma_{1},\Gamma_{2}\vdash B}{\Gamma_{1},A,\Gamma_{2}\vdash B} & [W] & \frac{\Gamma_{1},A,A,\Gamma_{2}\vdash B}{\Gamma_{1},A,\Gamma_{2}\vdash B} & [C] \\ \\ \hline \\ \frac{\Gamma_{1},A,\Gamma_{2}\vdash A}{\Gamma_{1},A,\Gamma_{2}\vdash A} & [Id] & \end{array}$$

Recap: Sequent Calculus

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All these **derived rules** can be proved/derived using the rules on the previous slide

Natural Deduction

Sequent Calculus

Natural Deduction Sequent Calculus

 $introduction/elimination\ rules \\ right/left\ rules$

Natural Deduction	Sequent Calculus
introduction/elimination rules	right/left rules
natural proofs	amenable to automation

Natural Deduction

Sequent Calculus

introduction/elimination rules

right/left rules

natural proofs

amenable to automation

$$\frac{A}{A}$$

$$\vdots$$

$$B$$

$$A \to B$$

$$1 \to I$$

$$\frac{\Gamma, A \vdash B}{\Gamma \vdash A \to B} \quad [\to R]$$

Natural Deduction

Sequent Calculus

introduction/elimination rules

right/left rules

natural proofs

amenable to automation

$$\frac{A}{A}$$

$$\vdots$$

$$B$$

$$A \to B$$

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$$\frac{\Gamma, A \vdash B}{\Gamma \vdash A \to B} \quad [\to R]$$

▶ in the Sequent Calculus the discharged hypothesis *A* is kept in the context!

Natural Deduction

Sequent Calculus

introduction/elimination rules

right/left rules

natural proofs

amenable to automation

$$\frac{A}{A}$$

$$\vdots$$

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$$A \to B$$

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$$\frac{\Gamma, A \vdash B}{\Gamma \vdash A \to B} \quad [\to R]$$

- ▶ in the Sequent Calculus the discharged hypothesis *A* is kept in the context!
- all the available hypotheses are always kept in the context part of sequents

Natural Deduction

Sequent Calculus

introduction/elimination rules

right/left rules

natural proofs

amenable to automation

$$\begin{array}{c}
A \\
\vdots \\
B \\
\hline
A \to B
\end{array}$$
1 [\to I]

$$\frac{\Gamma, A \vdash B}{\Gamma \vdash A \to B} \quad [\to R]$$

- ▶ in the Sequent Calculus the discharged hypothesis *A* is kept in the context!
- all the available hypotheses are always kept in the context part of sequents
- ▶ a proposition provable in one system is provable in the other

Provide a Natural Deduction proof of $(A \land B) \rightarrow (B \land A)$

Provide a Natural Deduction proof of $(A \land B) \rightarrow (B \land A)$

Here is an example of a backward proof:

 $(A \land B) \to (B \land A)$

Provide a Natural Deduction proof of $(A \wedge B) \rightarrow (B \wedge A)$

$$\frac{\overline{A \wedge B}}{1} \qquad \overline{A \wedge B} \qquad 1$$

$$\frac{\overline{B \wedge A}}{(A \wedge B) \to (B \wedge A)} \qquad 1 \ [\to I]$$

Provide a Natural Deduction proof of $(A \land B) \rightarrow (B \land A)$

$$\frac{\overline{A \wedge B}}{B} \stackrel{1}{\underbrace{\frac{A \wedge B}{A}}} \stackrel{1}{\underbrace{\frac{A \wedge B}{A}}} \stackrel{[\wedge I]}{\underbrace{(A \wedge B) \rightarrow (B \wedge A)}} \stackrel{[\wedge I]}{1} \stackrel{[\to I]}{}$$

Provide a Natural Deduction proof of $(A \land B) \rightarrow (B \land A)$

$$\frac{\overline{A \wedge B}}{B} \stackrel{[\wedge E_R]}{=} \frac{\overline{A \wedge B}}{A} \stackrel{[\wedge I]}{=} \frac{B \wedge A}{(A \wedge B) \to (B \wedge A)} \stackrel{[\wedge I]}{=} 1$$

Provide a Natural Deduction proof of $(A \land B) \rightarrow (B \land A)$

$$\frac{\overline{A \wedge B}}{B} \stackrel{[\wedge E_R]}{=} \frac{\overline{A \wedge B}}{A} \stackrel{[\wedge E_L]}{=} \\
\frac{\overline{B \wedge A}}{(A \wedge B) \to (B \wedge A)} \stackrel{[\wedge I]}{=} \\$$

Provide a Natural Deduction proof of $(A \land B) \rightarrow (B \land A)$

Here is an example of a backward proof:

$$\frac{\overline{A \wedge B}}{B} \stackrel{[\wedge E_R]}{=} \frac{\overline{A \wedge B}}{A} \stackrel{[\wedge E_L]}{=} \frac{B \wedge A}{(A \wedge B) \to (B \wedge A)} \stackrel{[\wedge I]}{=} 1$$

How do we know where the introduced hypotheses ($A \wedge B$ above) will be used in the proof?

Provide a Natural Deduction proof of $(A \land B) \rightarrow (B \land A)$

Here is an example of a backward proof:

$$\frac{\overline{A \wedge B}}{B} \stackrel{1}{[\wedge E_R]} \frac{\overline{A \wedge B}}{A} \stackrel{[\wedge E_L]}{[\wedge I]} \\
\frac{\overline{B \wedge A}}{(A \wedge B) \to (B \wedge A)} \stackrel{1}{[\to I]}$$

How do we know where the introduced hypotheses $(A \wedge B \text{ above})$ will be used in the proof?

We typically don't so we can keep track of them on the side while doing the proof as follows:

Let us prove $(A \land$	$B) \to (B \land A)$	again:

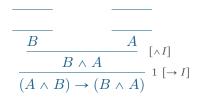
Let us prove $(A \land$	$B) \rightarrow$	$(B \wedge A)$	again:
		-	
		D) / F	
	$(A \land$	$B) \rightarrow (E$	$(\land A)$

Hypotheses:

Hypotheses:

▶ hypothesis 1: $A \wedge B$

Let us prove
$$(A \wedge B) \rightarrow (B \wedge A)$$
 again:



Hypotheses:

Let us prove $(A \wedge B) \rightarrow (B \wedge A)$ again:

$$\frac{\overline{A \wedge B}}{\underline{B}} \ [\wedge E_R] \ \overline{\frac{A}{A}} \ [\wedge I]$$

$$\frac{B \wedge A}{(A \wedge B) \to (B \wedge A)} \ 1 \ [\to I]$$

Hypotheses:

Let us prove $(A \wedge B) \rightarrow (B \wedge A)$ again:

$$\frac{\overline{A \wedge B}}{B} \stackrel{[\wedge E_R]}{=} \frac{\overline{A}}{A} \stackrel{[\wedge I]}{=} \frac{B \wedge A}{(A \wedge B) \to (B \wedge A)} \stackrel{[\wedge I]}{=} 1 \stackrel{[\to I]}{=}$$

Hypotheses:

Let us prove $(A \wedge B) \rightarrow (B \wedge A)$ again:

$$\frac{\overline{A \wedge B}}{B} \stackrel{[\wedge E_R]}{=} \frac{\overline{A \wedge B}}{A} \stackrel{[\wedge E_L]}{=} \frac{B \wedge A}{(A \wedge B) \to (B \wedge A)} \stackrel{[\wedge I]}{=} 1$$

Hypotheses:

Let us prove $(A \wedge B) \rightarrow (B \wedge A)$ again:

$$\frac{\overline{A \wedge B}}{B} \stackrel{1}{[\wedge E_R]} \frac{\overline{A \wedge B}}{A} \stackrel{[\wedge E_L]}{[\wedge I]}$$

$$\frac{B \wedge A}{(A \wedge B) \to (B \wedge A)} \stackrel{1}{[\to I]}$$

Hypotheses:

Let us prove $(A \wedge B) \rightarrow (B \wedge A)$ again:

$$\frac{\overline{A \wedge B}}{B} \stackrel{1}{[\wedge E_R]} \frac{\overline{A \wedge B}}{A} \stackrel{[\wedge E_L]}{[\wedge I]}$$

$$\frac{B \wedge A}{(A \wedge B) \to (B \wedge A)} \stackrel{1}{[\to I]}$$

Hypotheses:

▶ hypothesis 1: $A \wedge B$

This can be achieved using the Sequent Calculus!

Provide a Sequent Calculus proof of $(A \land B) \rightarrow (B \land A)$

Provide a Sequent Calculus proof of $(A \wedge B) \rightarrow (B \wedge A)$

 $\vdash (A \land B) \to (B \land A)$

Provide a Sequent Calculus proof of $(A \wedge B) \rightarrow (B \wedge A)$

$$\frac{\overline{A \wedge B \vdash B \wedge A}}{\vdash (A \wedge B) \to (B \wedge A)} \ [\to R]$$

Provide a Sequent Calculus proof of $(A \wedge B) \rightarrow (B \wedge A)$

$$\frac{A, B \vdash B \land A}{A \land B \vdash B \land A} \ [\land L]}{\vdash (A \land B) \to (B \land A)} \ [\to R]$$

Provide a Sequent Calculus proof of $(A \land B) \rightarrow (B \land A)$

$$\frac{\overline{A,B \vdash B} \quad [Id] \quad \overline{A,B \vdash A} \quad [Id]}{\frac{A,B \vdash B \land A}{A \land B \vdash B \land A} \quad [\land L]}$$

$$\frac{|A \land B \vdash B \land A}{| \vdash (A \land B) \rightarrow (B \land A)} \quad [\rightarrow R]$$

Provide a Natural Deduction proof of

$$(A \to B) \to (C \to D) \to (A \lor C) \to (B \lor D)$$

Provide a Natural Deduction proof of $(A \to B) \to (C \to D) \to (A \lor C) \to (B \lor D)$ We will keep track of our hypotheses on the side **Hypotheses:**

 $(A \to B) \to (C \to D) \to (A \lor C) \to (B \lor D)$

Provide a Natural Deduction proof of

$$(A \to B) \to (C \to D) \to (A \lor C) \to (B \lor D)$$

We will keep track of our hypotheses on the side

	_	 _

$$\frac{(C \to D) \to (A \lor C) \to (B \lor D)}{(A \to B) \to (C \to D) \to (A \lor C) \to (B \lor D)} \ \ 1 \ [\to I]$$

Hypotheses:

• hyp. 1:
$$A \rightarrow B$$

Provide a Natural Deduction proof of

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We will keep track of our hypotheses on the side

$$\frac{\overline{(A \lor C) \to (B \lor D)}}{\overline{(C \to D) \to (A \lor C) \to (B \lor D)}} \stackrel{2 [\to I]}{= I}$$

$$\overline{(A \to B) \to (C \to D) \to (A \lor C) \to (B \lor D)} \stackrel{1 [\to I]}{=}$$

Hypotheses:

• hyp. 1:
$$A \rightarrow B$$

• hyp. 2:
$$C \rightarrow D$$

Provide a Natural Deduction proof of

$$(A \to B) \to (C \to D) \to (A \lor C) \to (B \lor D)$$

We will keep track of our hypotheses on the side

$$\frac{B \vee D}{(A \vee C) \to (B \vee D)} \xrightarrow{3 \ [\to I]} \frac{}{(C \to D) \to (A \vee C) \to (B \vee D)} \xrightarrow{2 \ [\to I]} \frac{}{(A \to B) \to (C \to D) \to (A \vee C) \to (B \vee D)} \xrightarrow{1 \ [\to I]} \frac{}{(A \to B) \to (C \to D) \to (A \vee C) \to (B \vee D)} \xrightarrow{1 \ [\to I]} \frac{}{(A \to B) \to (C \to D) \to (A \vee C) \to (B \vee D)} \xrightarrow{1 \ [\to I]} \frac{}{(A \to B) \to (C \to D) \to (A \vee C) \to (B \vee D)} \xrightarrow{1 \ [\to I]} \frac{}{(A \to B) \to (C \to D) \to (A \vee C) \to (B \vee D)} \xrightarrow{1 \ [\to I]} \frac{}{(A \to B) \to (C \to D) \to (A \vee C) \to (B \vee D)} \xrightarrow{1 \ [\to I]} \frac{}{(A \to B) \to (C \to D) \to (A \vee C) \to (B \vee D)} \xrightarrow{1 \ [\to I]} \frac{}{(A \to B) \to (C \to D) \to (A \vee C) \to (B \vee D)} \xrightarrow{1 \ [\to I]} \frac{}{(A \to B) \to (C \to D) \to (A \vee C) \to (B \vee D)} \xrightarrow{1 \ [\to I]} \frac{}{(A \to B) \to (C \to D) \to (A \vee C) \to (B \vee D)} \xrightarrow{1 \ [\to I]} \frac{}{(A \to B) \to (C \to D) \to (A \vee C) \to (B \vee D)} \xrightarrow{1 \ [\to I]} \frac{}{(A \to B) \to (C \to D) \to (A \vee C) \to (B \vee D)} \xrightarrow{1 \ [\to I]} \frac{}{(A \to B) \to (C \to D) \to (A \vee C) \to (B \vee D)} \xrightarrow{1 \ [\to I]} \frac{}{(A \to B) \to (C \to D) \to (A \vee C) \to (B \vee D)} \xrightarrow{1 \ [\to I]} \frac{}{(A \to B) \to (C \to D) \to (A \vee C) \to (B \vee D)} \xrightarrow{1 \ [\to I]} \frac{}{(A \to B) \to (C \to D) \to (A \vee C) \to (B \vee D)} \xrightarrow{1 \ [\to I]} \frac{}{(A \to B) \to (B \vee D)} \xrightarrow{1 \ [\to I]} \frac{}{(A \to B) \to (B \vee D)} \xrightarrow{1 \ [\to I]} \frac{}{(A \to B) \to (B \vee D)} \xrightarrow{1 \ [\to I]} \frac{}{(A \to B) \to (B \vee D)} \xrightarrow{1 \ [\to I]} \frac{}{(A \to B) \to (B \vee D)} \xrightarrow{1 \ [\to I]} \frac{}{(A \to B) \to (B \vee D)} \xrightarrow{1 \ [\to I]} \frac{}{(A \to B) \to (B \vee D)} \xrightarrow{1 \ [\to I]} \frac{}{(A \to B) \to (B \vee D)} \xrightarrow{1 \ [\to I]} \frac{}{(A \to B) \to (B \vee D)} \xrightarrow{1 \ [\to I]} \frac{}{(A \to B) \to (B \vee D)} \xrightarrow{1 \ [\to I]} \frac{}{(A \to B) \to (B \vee D)} \xrightarrow{1 \ [\to I]} \frac{}{(A \to B) \to (B \vee D)} \xrightarrow{1 \ [\to I]} \frac{}{(A \to B) \to (B \vee D)} \xrightarrow{1 \ [\to I]} \frac{}{(A \to B) \to (B \vee D)} \xrightarrow{1 \ [\to I]} \frac{}{(A \to B) \to (B \vee D)} \xrightarrow{1 \ [\to I]} \frac{}{(A \to B) \to (B \vee D)} \xrightarrow{1 \ [\to I]} \frac{}{(A \to B) \to (B \vee D)} \xrightarrow{1 \ [\to I]} \xrightarrow{1 \ [\to I]} \frac{}{(A \to B) \to (B \vee D)} \xrightarrow{1 \ [\to I]} \frac{}{(A \to B) \to (B \vee D)} \xrightarrow{1 \ [\to I]} \frac{}{(A \to B) \to (B \vee D)} \xrightarrow{1 \ [\to I]} \frac{}{(A \to B) \to (B \vee D)} \xrightarrow{1 \ [\to I]} \xrightarrow{1 \ [\to I]} \frac{}{(A \to B) \to (B \to B)} \xrightarrow{1 \ [\to I]} \frac{}{(A \to B) \to (B \to B)} \xrightarrow{1 \ [\to I]} \frac{}{(A \to B) \to (B \to B)} \xrightarrow{1 \ [\to I]} \xrightarrow{$$

Hypotheses:

• hyp. 1: $A \rightarrow B$

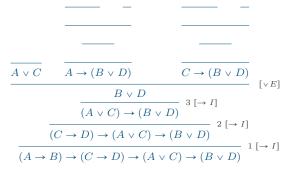
• hyp. 2: $C \rightarrow D$

■ hyp. 3: *A* ∨ *C*

Provide a Natural Deduction proof of

$$(A \to B) \to (C \to D) \to (A \lor C) \to (B \lor D)$$

We will keep track of our hypotheses on the side



Hypotheses:

• hyp. 1: $A \rightarrow B$ • hyp. 2: $C \rightarrow D$

■ hyp. 3: *A* ∨ *C*

Provide a Natural Deduction proof of

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$$\frac{\overline{A \vee C} \stackrel{3}{\longrightarrow} \overline{A \to (B \vee D)} \qquad \overline{C \to (B \vee D)}}{\overline{C \to (B \vee D)}} \quad [\vee E]$$

$$\frac{B \vee D}{(A \vee C) \to (B \vee D)} \stackrel{3 [\to I]}{\longrightarrow} \stackrel{2 [\to I]}{\longrightarrow} \stackrel{1 [\to I]}{\longrightarrow} \stackrel{1}{\longrightarrow} \stackrel{1}{\longrightarrow} \stackrel{I}{\longrightarrow} \stackrel{I}{$$

Hypotheses:

• hyp. 1: $A \rightarrow B$

• hyp. 2: $C \rightarrow D$

• hyp. 3: $A \lor C$

Provide a Natural Deduction proof of

$$(A \to B) \to (C \to D) \to (A \lor C) \to (B \lor D)$$

We will keep track of our hypotheses on the side

$$\frac{\overline{B \vee D}}{A \vee C} \stackrel{3}{\longrightarrow} \frac{\overline{B \vee D}}{A \to (B \vee D)} \stackrel{4 [\to I]}{\longrightarrow} \frac{\overline{C \to (B \vee D)}}{C \to (B \vee D)} \stackrel{[\vee E]}{\longrightarrow} \frac{\overline{B \vee D}}{(A \vee C) \to (B \vee D)} \stackrel{3 [\to I]}{\longrightarrow} \frac{\overline{C \to (B \vee D)}}{(C \to D) \to (A \vee C) \to (B \vee D)} \stackrel{1 [\to I]}{\longrightarrow} \frac{\overline{C \to (B \vee D)}}{(A \to B) \to (C \to D) \to (A \vee C) \to (B \vee D)} \stackrel{1 [\to I]}{\longrightarrow} \frac{\overline{C \to (B \vee D)}}{\longrightarrow} \stackrel{1 [\to I]}{\longrightarrow} \stackrel{1 [\to$$

Hypotheses:

• hyp. 1: $A \rightarrow B$

• hyp. 2: $C \rightarrow D$

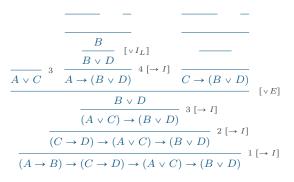
■ hyp. 3: *A* ∨ *C*

• hyp. 4: *A*

Provide a Natural Deduction proof of

$$(A \to B) \to (C \to D) \to (A \lor C) \to (B \lor D)$$

We will keep track of our hypotheses on the side



Hypotheses:

- hyp. 1: $A \rightarrow B$
- hyp. 2: $C \rightarrow D$
- hyp. 3: $A \vee C$
- hyp. 4: *A*

Provide a Natural Deduction proof of

$$(A \to B) \to (C \to D) \to (A \lor C) \to (B \lor D)$$

We will keep track of our hypotheses on the side

$$\frac{A \to B \qquad A}{B \qquad B \qquad [\to E]} \qquad -\frac{B}{B \lor D} \qquad [\lor I_L] \qquad -\frac{A \to C}{A \to (B \lor D)} \qquad 4 \ [\to I] \qquad C \to (B \lor D) \qquad [\lor E]$$

$$\frac{B \lor D}{(A \lor C) \to (B \lor D)} \qquad 3 \ [\to I]$$

$$\frac{(C \to D) \to (A \lor C) \to (B \lor D)}{(A \to B) \to (C \to D) \to (A \lor C) \to (B \lor D)} \qquad 1 \ [\to I]$$

Hypotheses:

• hyp. 1: $A \rightarrow B$

• hyp. 2: $C \rightarrow D$

■ hyp. 3: *A* ∨ *C*

■ hyp. 4: *A*

Provide a Natural Deduction proof of

$$(A \to B) \to (C \to D) \to (A \lor C) \to (B \lor D)$$

We will keep track of our hypotheses on the side

Hypotheses:

• hyp. 1: $A \rightarrow B$

• hyp. 2: $C \rightarrow D$

■ hyp. 3: *A* ∨ *C*

• hyp. 4: *A*

Provide a Natural Deduction proof of

$$(A \to B) \to (C \to D) \to (A \lor C) \to (B \lor D)$$

We will keep track of our hypotheses on the side

$$\frac{A \to B}{A} \xrightarrow{I} \xrightarrow{A} \xrightarrow{I} \xrightarrow{E} \xrightarrow{I}$$

$$\frac{B}{B \lor D} \xrightarrow{[\lor I_L]} \xrightarrow{C} \xrightarrow{C} \xrightarrow{(B \lor D)}$$

$$\frac{B \lor D}{(A \lor C) \to (B \lor D)} \xrightarrow{3 \ [\to I]} \xrightarrow{[\lor E]}$$

$$\frac{B \lor D}{(C \to D) \to (A \lor C) \to (B \lor D)} \xrightarrow{2 \ [\to I]}$$

$$\frac{A \to B}{(A \to B) \to (C \to D) \to (A \lor C) \to (B \lor D)} \xrightarrow{1 \ [\to I]}$$

Hypotheses:

• hyp. 1: $A \rightarrow B$

• hyp. 2: $C \rightarrow D$

■ hyp. 3: *A* ∨ *C*

• hyp. 4: *A*

Provide a Natural Deduction proof of

$$(A \to B) \to (C \to D) \to (A \lor C) \to (B \lor D)$$

We will keep track of our hypotheses on the side

Hypotheses:

- hyp. 1: $A \rightarrow B$
- hyp. 2: $C \rightarrow D$
- hyp. 3: $A \lor C$
- hyp. 4: A
- hyp. 5: *C*

Provide a Natural Deduction proof of

$$(A \to B) \to (C \to D) \to (A \lor C) \to (B \lor D)$$

Provide a Natural Deduction proof of

$$(A \to B) \to (C \to D) \to (A \lor C) \to (B \lor D)$$

$$\frac{\overline{A \to B} \quad \stackrel{1}{A} \quad \stackrel{4}{A} \quad \overline{C \to D} \quad \overline{C}}{\overline{B} \quad [\to E]} \qquad \begin{array}{c} \overline{B} \quad [\to E] \\ \overline{B} \quad [\lor I_L] \\ \overline{A \lor C} \end{array} \quad \begin{array}{c} \overline{B} \quad [\lor I_L] \\ \overline{A \lor C} \end{array} \quad \begin{array}{c} \overline{B} \quad [\lor I_L] \\ \overline{A \to (B \lor D)} \end{array} \quad \begin{array}{c} [\lor I_R] \\ \overline{C \to (B \lor D)} \end{array} \quad \begin{array}{c} 5 \ [\to I] \\ \overline{C \to (B \lor D)} \end{array} \quad \begin{array}{c} \bullet \quad \text{hyp. 1: } A \to B \\ \bullet \quad \text{hyp. 2: } C \to D \\ \hline \bullet \quad \text{hyp. 3: } A \lor C \\ \hline \bullet \quad [\lor E] \\ \hline \hline \bullet \quad (A \lor C) \to (B \lor D) \end{array} \quad \begin{array}{c} B \lor D \\ \hline \hline \bullet \quad (C \to D) \to (A \lor C) \to (B \lor D) \end{array} \quad \begin{array}{c} \bullet \quad [\lor E] \\ \hline \bullet \quad \text{hyp. 4: } A \\ \hline \bullet \quad \text{hyp. 5: } C \end{array}$$

Provide a Natural Deduction proof of

$$(A \to B) \to (C \to D) \to (A \lor C) \to (B \lor D)$$

$$\frac{\overline{A \to B} \quad \overline{A} \quad \overline{A} \quad \overline{C \to D} \quad \overline{C} \quad \overline{C}}{\overline{B} \quad \overline{B} \quad \overline{C} \quad \overline{D} \quad \overline{C}} \quad \overline{C} \quad \overline{C}} \quad \overline{C} \quad \overline{C$$

Provide a Natural Deduction proof of

$$(A \to B) \to (C \to D) \to (A \lor C) \to (B \lor D)$$

$$\frac{\overline{A \to B} \quad \overline{A} \quad \overline{A}}{\overline{B} \quad \overline{A}} \quad \overline{[\to E]} \quad \overline{C \to D} \quad \overline{C} \quad \overline{C} \quad \overline{D} \quad \overline{[\to E]} \quad \overline{B} \quad \overline{B \vee D} \quad \overline{[\to E]} \quad \overline{B \vee D} \quad \overline{[\to E]} \quad \overline{B \vee D} \quad \overline{[\to E]} \quad \overline{B \vee D} \quad \overline{S} \quad \overline{[\to I]} \quad \overline{B \vee D} \quad \overline{S} \quad \overline{[\to I]} \quad \overline{D} \quad \overline{S} \quad \overline{S}$$

Provide a Natural Deduction proof of

$$(A \to B) \to (C \to D) \to (A \lor C) \to (B \lor D)$$

We will keep track of our hypotheses on the side

$$\frac{\overline{A \to B} \quad \overline{A} \quad \overline{A}}{\overline{B} \quad \overline{A}} \quad \overline{[\to E]} \quad \overline{C \to D} \quad \overline{C} \quad \overline{C} \quad \overline{D} \quad \overline{[\to E]} \quad \overline{B} \quad \overline{B \vee D} \quad \overline{[\to E]} \quad \overline{B \vee D} \quad \overline{[\to E]} \quad \overline{B \vee D} \quad \overline{[\to E]} \quad \overline{B \vee D} \quad \overline{S} \quad \overline{[\to I]} \quad \overline{B \vee D} \quad \overline{S} \quad \overline{[\to I]} \quad \overline{D} \quad \overline{S} \quad \overline{S}$$

If an hypothesis is introduced in a branch, make sure you don't use it in another branch (e.g. 4 cannot be used in the far right branch)

Provide a Natural Deduction proof of

$$(A \to B) \to (C \to D) \to (A \lor C) \to (B \lor D)$$

$$\frac{\overline{A \to B} \quad \overline{A} \quad \overline{A}}{\overline{B} \quad \overline{A}} \quad \overline{[\to E]} \quad \overline{C \to D} \quad \overline{C} \quad \overline{C} \quad \overline{D} \quad \overline{[\to E]} \quad \overline{B} \quad \overline{B \vee D} \quad \overline{[\to E]} \quad \overline{B \vee D} \quad \overline{[\to E]} \quad \overline{B \vee D} \quad \overline{[\to E]} \quad \overline{B \vee D} \quad \overline{S} \quad \overline{[\to I]} \quad \overline{B \vee D} \quad \overline{S} \quad \overline{[\to I]} \quad \overline{D} \quad \overline{S} \quad \overline{S}$$

- If an hypothesis is introduced in a branch, make sure you don't use it in another branch (e.g. 4 cannot be used in the far right branch)
- This is enforced by sequents in the Sequent Calculus

$$(A \to B) \to (C \to D) \to (A \lor C) \to (B \lor D)$$

Provide a Sequent Calculus proof of

$$(A \to B) \to (C \to D) \to (A \lor C) \to (B \lor D)$$

 $\vdash (A \to B) \to (C \to D) \to (A \lor C) \to (B \lor D)$

$$(A \to B) \to (C \to D) \to (A \lor C) \to (B \lor D)$$

$$\frac{A \to B \vdash (C \to D) \to (A \lor C) \to (B \lor D)}{\vdash (A \to B) \to (C \to D) \to (A \lor C) \to (B \lor D)} \quad [\to R]$$

$$(A \to B) \to (C \to D) \to (A \lor C) \to (B \lor D)$$

$$\cfrac{\cfrac{A \to B, C \to D \vdash (A \lor C) \to (B \lor D)}{A \to B \vdash (C \to D) \to (A \lor C) \to (B \lor D)}}{[\to R]}$$

$$[\to R]$$

$$[\to R]$$

$$(A \to B) \to (C \to D) \to (A \lor C) \to (B \lor D)$$

$$\frac{A \to B, C \to D, A \lor C \vdash B \lor D}{A \to B, C \to D \vdash (A \lor C) \to (B \lor D)} \quad [\to R]$$

$$\frac{A \to B, C \to D \vdash (A \lor C) \to (B \lor D)}{A \to B \vdash (C \to D) \to (A \lor C) \to (B \lor D)} \quad [\to R]$$

$$\vdash (A \to B) \to (C \to D) \to (A \lor C) \to (B \lor D)$$

$$(A \to B) \to (C \to D) \to (A \lor C) \to (B \lor D)$$

$$\frac{A \to B, C \to D, A \vdash B \lor D}{A \to B, C \to D, C \vdash B \lor D} \qquad [\lor L]$$

$$\frac{A \to B, C \to D, A \lor C \vdash B \lor D}{A \to B, C \to D \vdash (A \lor C) \to (B \lor D)} \qquad [\to R]$$

$$\frac{A \to B, C \to D \vdash (A \lor C) \to (B \lor D)}{A \to B \vdash (C \to D) \to (A \lor C) \to (B \lor D)} \qquad [\to R]$$

$$(A \to B) \to (C \to D) \to (A \lor C) \to (B \lor D)$$

$$\frac{A \to B, C \to D, A \vdash B \lor D}{A \to B, C \to D, A \vdash B \lor D} [\to L] \qquad A \to B, C \to D, C \vdash B \lor D$$

$$\frac{A \to B, C \to D, A \vdash B \lor D}{A \to B, C \to D, A \lor C \vdash B \lor D} [\to R]$$

$$\frac{A \to B, C \to D, A \lor C \vdash B \lor D}{A \to B, C \to D, C \vdash B \lor D} [\to R]$$

$$\frac{A \to B, C \to D, A \lor C \vdash B \lor D}{A \to B, C \to D, C \vdash B \lor D} [\to R]$$

$$\frac{A \to B, C \to D, A \lor C \vdash B \lor D}{A \to B, C \to D, C \vdash B \lor D} [\to R]$$

$$\frac{A \to B, C \to D, A \lor C \vdash B \lor D}{A \to B, C \to D, C \vdash B \lor D} [\to R]$$

$$(A \to B) \to (C \to D) \to (A \lor C) \to (B \lor D)$$

$$\frac{C \to D, A \vdash A}{A \to B, C \to D, A \vdash B \lor D} \xrightarrow{[\to L]} \frac{A \to B, C \to D, A \vdash B \lor D}{A \to B, C \to D, C \vdash B \lor D} \xrightarrow{[\lor L]} \frac{A \to B, C \to D, A \lor C \vdash B \lor D}{A \to B, C \to D, A \lor C \vdash B \lor D} \xrightarrow{[\to R]} \frac{A \to B, C \to D \vdash (A \lor C) \to (B \lor D)}{A \to B \vdash (C \to D) \to (A \lor C) \to (B \lor D)} \xrightarrow{[\to R]} \xrightarrow{[\to R]} \frac{A \to B \vdash (C \to D) \to (A \lor C) \to (B \lor D)}{[\to R]}$$

$$(A \to B) \to (C \to D) \to (A \lor C) \to (B \lor D)$$

$$(A \to B) \to (C \to D) \to (A \lor C) \to (B \lor D)$$

$$\frac{D}{C \to D, A \vdash A} \quad [Id] \quad \frac{D}{B, C \to D, A \vdash B} \quad [Id] \quad D$$

$$\frac{A \to B, C \to D, A \vdash B \lor D}{A \to B, C \to D, A \vdash B \lor D} \quad [\to L] \quad D$$

$$\frac{A \to B, C \to D, A \vdash B \lor D}{A \to B, C \to D, C \vdash B \lor D} \quad [\lor L]$$

$$\frac{A \to B, C \to D, A \lor C \vdash B \lor D}{A \to B, C \to D, C \vdash B \lor D} \quad [\to R]$$

$$\frac{A \to B, C \to D \vdash (A \lor C) \to (B \lor D)}{A \to B \vdash (C \to D) \to (A \lor C) \to (B \lor D)} \quad [\to R]$$

$$\frac{A \to B \vdash (C \to D) \to (A \lor C) \to (B \lor D)}{A \to B \vdash (C \to D) \to (A \lor C) \to (B \lor D)} \quad [\to R]$$

$$(A \to B) \to (C \to D) \to (A \lor C) \to (B \lor D)$$

$$(A \to B) \to (C \to D) \to (A \lor C) \to (B \lor D)$$

$$\frac{C \to D, A \vdash A}{A \to B, C \to D, A \vdash B \lor D} \begin{bmatrix} Id] \\ B, C \to D, A \vdash B \lor D \end{bmatrix} \begin{bmatrix} \lor R_1 \end{bmatrix} \xrightarrow{A \to B, C \vdash C} \begin{bmatrix} Id] \xrightarrow{A \to B, D, C \vdash B \lor D} \\ A \to B, C \to D, A \vdash B \lor D \end{bmatrix} \begin{bmatrix} \lor L \end{bmatrix} \xrightarrow{A \to B, C \to D, C \vdash B \lor D} \begin{bmatrix} \lor L \end{bmatrix} \begin{bmatrix} A \to B, C \to D, C \vdash B \lor D \\ \hline A \to B, C \to D, C \vdash B \lor D \end{bmatrix} \begin{bmatrix} \lor L \end{bmatrix} \xrightarrow{A \to B, C \to D \vdash (A \lor C) \to (B \lor D)} \begin{bmatrix} \vdash R \end{bmatrix} \begin{bmatrix} \bot R \end{bmatrix} \begin{bmatrix} \bot A \to B \vdash (C \to D) \to (A \lor C) \to (B \lor D) \end{bmatrix} \begin{bmatrix} \bot B \end{bmatrix}$$

Provide a Sequent Calculus proof of

$$(A \to B) \to (C \to D) \to (A \lor C) \to (B \lor D)$$

$$\frac{C \to D, A \vdash A}{A \to B, C \to D, A \vdash B \lor D} \begin{bmatrix} Id \\ B, C \to D, A \vdash B \lor D \end{bmatrix} \begin{bmatrix} \lor R_1 \\ \to L \end{bmatrix} \xrightarrow{A \to B, C \vdash C} \begin{bmatrix} Id \end{bmatrix} \xrightarrow{A \to B, D, C \vdash D} \begin{bmatrix} \lor R_2 \\ A \to B, D, C \vdash B \lor D \end{bmatrix} \begin{bmatrix} \lor R_2 \\ \to L \end{bmatrix}$$

$$\frac{A \to B, C \to D, A \vdash B \lor D}{A \to B, C \to D, C \vdash B \lor D} \begin{bmatrix} \lor L \end{bmatrix} \begin{bmatrix} \to L \end{bmatrix}$$

$$\frac{A \to B, C \to D, A \lor C \vdash B \lor D}{A \to B, C \to D \vdash (A \lor C) \to (B \lor D)} \begin{bmatrix} \to R \end{bmatrix}$$

$$A \to B \vdash (C \to D) \to (A \lor C) \to (B \lor D)$$

 $\vdash (A \to B) \to (C \to D) \to (A \lor C) \to (B \lor D)$

Provide a Sequent Calculus proof of

$$(A \to B) \to (C \to D) \to (A \lor C) \to (B \lor D)$$

$$\frac{C \to D, A \vdash A}{A \to B, C \to D, A \vdash B \lor D} \begin{bmatrix} Id \end{bmatrix} \begin{bmatrix}$$

 $\vdash (A \to B) \to (C \to D) \to (A \lor C) \to (B \lor D)$

Provide a Natural Deduction proof of $(B \to C \to \neg A) \to A \to \neg (B \land C)$

Provide a Natural Deduction proof of $(B \to C \to \neg A) \to A \to \neg (B \land C)$

We will keep track of our hypotheses on the side

			_
	-		
	$(B \rightarrow C \rightarrow$	$\rightarrow \neg A) \rightarrow A \rightarrow \neg$	$(B \land C)$

Hypotheses:

Provide a Natural Deduction proof of $(B \to C \to \neg A) \to A \to \neg (B \land C)$

We will keep track of our hypotheses on the side

• hyp. 1: $B \to C \to \neg A$

$$\frac{\overline{A \to \neg (B \land C)}}{(B \to C \to \neg A) \to A \to \neg (B \land C)} \ ^1 \ [\to I]$$

Provide a Natural Deduction proof of $(B \to C \to \neg A) \to A \to \neg (B \land C)$

We will keep track of our hypotheses on the side

		_

Hypotheses:

- hyp. 1: $B \to C \to \neg A$
- hyp. 2: *A*

$$\frac{\overline{-(B \wedge C)}}{A \to \neg(B \wedge C)} \ ^{2} [\to I]$$

$$\frac{(B \to C \to \neg A) \to A \to \neg(B \wedge C)}{} \ ^{1} [\to I]$$

Provide a Natural Deduction proof of

$$(B \to C \to \neg A) \to A \to \neg (B \land C)$$

We will keep track of our hypotheses on the side

	 	 _	
		 _	
			-

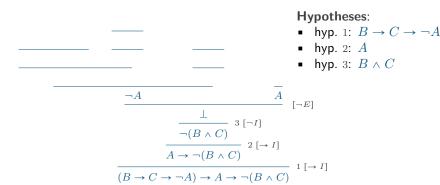
Hypotheses:

- hyp. 1: $B \to C \to \neg A$
- hyp. 2: *A*
- hyp. 3: *B* ∧ *C*

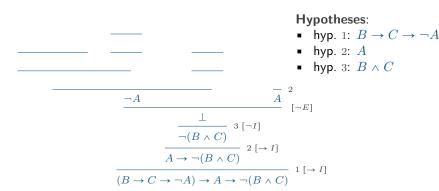
$$\frac{\frac{\bot}{\neg (B \land C)} \quad 3 \, [\neg I]}{\frac{A \rightarrow \neg (B \land C)}{A \rightarrow \neg (B \land C)}} \quad 2 \, [\rightarrow I]}$$

$$\frac{(B \rightarrow C \rightarrow \neg A) \rightarrow A \rightarrow \neg (B \land C)}{} \quad 1 \, [\rightarrow I]$$

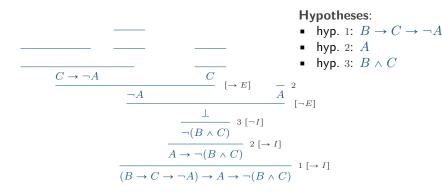
Provide a Natural Deduction proof of $(B \to C \to \neg A) \to A \to \neg (B \land C)$



Provide a Natural Deduction proof of $(B \to C \to \neg A) \to A \to \neg (B \land C)$

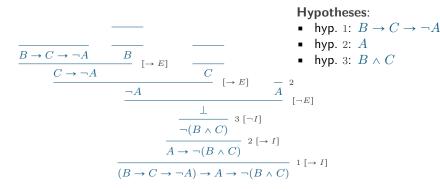


Provide a Natural Deduction proof of $(B \to C \to \neg A) \to A \to \neg (B \land C)$



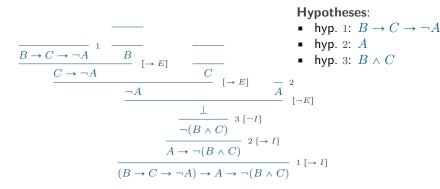
Provide a Natural Deduction proof of

$$(B \to C \to \neg A) \to A \to \neg (B \land C)$$



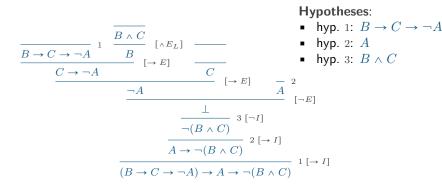
Provide a Natural Deduction proof of

$$(B \to C \to \neg A) \to A \to \neg (B \land C)$$



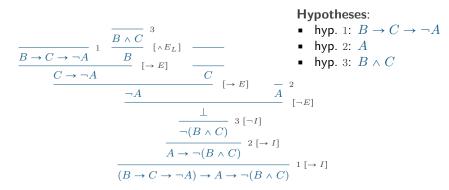
Provide a Natural Deduction proof of

$$(B \to C \to \neg A) \to A \to \neg (B \land C)$$



Provide a Natural Deduction proof of

$$(B \to C \to \neg A) \to A \to \neg (B \land C)$$



Provide a Natural Deduction proof of

$$(B \to C \to \neg A) \to A \to \neg (B \land C)$$

$$\frac{B \to C \to \neg A}{2} \stackrel{1}{\longrightarrow} \frac{B \land C}{B} \stackrel{[\land E_L]}{\longrightarrow} \frac{B \land C}{C} \stackrel{[\land E_R]}{\longrightarrow} \frac{B \land C}{C} \stackrel{[\land E_R]}{\longrightarrow} \frac{A}{A} \stackrel{\bullet}{\longrightarrow} \frac{1}{A} \stackrel{\bullet}{\longrightarrow} \frac{1$$

Provide a Natural Deduction proof of $(B \to C \to \neg A) \to A \to \neg (B \land C)$

$$\frac{B \rightarrow C \rightarrow \neg A}{B \land C} \stackrel{\text{$| A \land C |}}{= B \land C} \stackrel{\text{$| A \land C |}}{= A} \stackrel{\text{$| A \land C |}$$

Provide a Sequent Calculus proof of $(B \to C \to \neg A) \to A \to \neg (B \land C)$

		lculus pro $A \to \neg (B$		
Here is	*	- (-		
_				
			-	

$$\vdash (B \to C \to \neg A) \to A \to \neg (B \land C)$$

Provide a Sequent Calculus proof of $(B \to C \to \neg A) \to A \to \neg (B \land C)$

Here is a proof:

$$\frac{\overline{B \to C \to \neg A \vdash A \to \neg (B \land C)}}{\vdash (B \to C \to \neg A) \to A \to \neg (B \land C)} \ \ [\to R]$$

Provide a Sequent Calculus proof of $(B \to C \to \neg A) \to A \to \neg (B \land C)$

$$\frac{\overline{B \to C \to \neg A, A \vdash \neg (B \land C)}}{\overline{B \to C \to \neg A \vdash A \to \neg (B \land C)}} \xrightarrow{[\to R]} \overline{+ (B \to C \to \neg A) \to A \to \neg (B \land C)}$$

Provide a Sequent Calculus proof of $(B \to C \to \neg A) \to A \to \neg (B \land C)$ Here is a proof:

$$\frac{\overline{B \to C \to \neg A, A, B \land C \vdash \bot}}{\overline{B \to C \to \neg A, A \vdash \neg (B \land C)}} [\neg R]}$$

$$\frac{\overline{B \to C \to \neg A, A \vdash \neg (B \land C)}}{\overline{B \to C \to \neg A \vdash A \to \neg (B \land C)}} [\to R]$$

$$\vdash (B \to C \to \neg A) \to A \to \neg (B \land C)}$$

Provide a Sequent Calculus proof of $(B \to C \to \neg A) \to A \to \neg (B \land C)$ Here is a proof:

 $\frac{B \to C \to \neg A, A, B, C \vdash \bot}{B \to C \to \neg A, A, B \land C \vdash \bot} [\land L]$ $\frac{B \to C \to \neg A, A, B \land C \vdash \bot}{B \to C \to \neg A, A \vdash \neg (B \land C)} [\to R]$ $\frac{B \to C \to \neg A \vdash A \to \neg (B \land C)}{\vdash (B \to C \to \neg A) \to A \to \neg (B \land C)} [\to R]$

Provide a Sequent Calculus proof of $(B \to C \to \neg A) \to A \to \neg (B \land C)$

$$\frac{A, B, C \vdash B}{C \to \neg A, A, B, C \vdash \bot} \xrightarrow{[\to L]} [\to L]$$

$$\frac{B \to C \to \neg A, A, B, C \vdash \bot}{B \to C \to \neg A, A, B \land C \vdash \bot} \xrightarrow{[\neg R]} [\to R]$$

$$\frac{B \to C \to \neg A, A, B \land C \vdash \bot}{B \to C \to \neg A, A \vdash \neg (B \land C)} \xrightarrow{[\to R]} [\to R]$$

$$\frac{A, B, C \vdash B}{A, C \to \neg A, A, B, C \vdash \bot} \xrightarrow{[\to L]} [\to L]$$

$$\frac{A, B, C \vdash B}{B \to C \to \neg A, A, B, C \vdash \bot} \xrightarrow{[\land L]} [\to L]$$

Provide a Sequent Calculus proof of $(B \to C \to \neg A) \to A \to \neg (B \land C)$

$$\frac{A,B,C \vdash B}{B \to C \to \neg A,A,B,C \vdash \bot} \xrightarrow{[\land L]} \xrightarrow{[\land L]} \xrightarrow{B \to C \to \neg A,A,B,C \vdash \bot} \xrightarrow{[\land L]} \xrightarrow{B \to C \to \neg A,A,B \land C \vdash \bot} \xrightarrow{[\land L]} \xrightarrow{[\land R]} \xrightarrow{B \to C \to \neg A,A \vdash \neg (B \land C)} \xrightarrow{[\to R]} \xrightarrow{[\to R]} \xrightarrow{[\vdash (B \to C \to \neg A) \to A \to \neg (B \land C)} \xrightarrow{[\vdash R]}$$

Provide a Sequent Calculus proof of $(B \to C \to \neg A) \to A \to \neg (B \land C)$

$$\frac{A,B,C \vdash C}{A,B,C \vdash B} \xrightarrow{\neg A,A,B,C \vdash \bot} [\rightarrow L]$$

$$\frac{B \to C \to \neg A,A,B,C \vdash \bot}{C \to \neg A,A,B,C \vdash \bot} [\rightarrow L]$$

$$\frac{B \to C \to \neg A,A,B,C \vdash \bot}{B \to C \to \neg A,A,B \land C \vdash \bot} [\rightarrow L]$$

$$\frac{B \to C \to \neg A,A,B \land C \vdash \bot}{B \to C \to \neg A,A \vdash \neg (B \land C)} [\rightarrow R]$$

$$\frac{B \to C \to \neg A \vdash A \to \neg (B \land C)}{[\rightarrow R]} [\rightarrow R]$$

$$\vdash (B \to C \to \neg A) \to A \to \neg (B \land C)$$

Provide a Sequent Calculus proof of $(B \to C \to \neg A) \to A \to \neg (B \land C)$

$$\frac{A,B,C \vdash B}{A,B,C \vdash C} \stackrel{[Id]}{\longrightarrow} \frac{A,B,C \vdash C}{\neg A,A,B,C \vdash \bot} \stackrel{[\rightarrow L]}{\longrightarrow} \frac{A,B,C \vdash \bot}{[\rightarrow L]} \stackrel{[\rightarrow L]}{\longrightarrow} \frac{A,B,C \vdash \bot}{[\rightarrow L$$

Provide a Sequent Calculus proof of $(B \to C \to \neg A) \to A \to \neg (B \land C)$

$$\frac{A,B,C \vdash B}{A,B,C \vdash B} \xrightarrow{[Id]} \frac{A,B,C \vdash A}{\neg A,A,B,C \vdash \bot} \xrightarrow{[\neg L]} \xrightarrow{[\neg L]}$$

$$\frac{B \to C \to \neg A,A,B,C \vdash \bot}{B \to C \to \neg A,A,B,C \vdash \bot} \xrightarrow{[\neg R]} \xrightarrow{[\neg R]} \xrightarrow{[\neg R]} \xrightarrow{[\neg R]}$$

$$\frac{B \to C \to \neg A,A,B,C \vdash \bot}{B \to C \to \neg A,A \vdash \neg (B \land C)} \xrightarrow{[\neg R]} \xrightarrow{[$$

Provide a Sequent Calculus proof of

$$(B \to C \to \neg A) \to A \to \neg (B \land C)$$

e is a proof:
$$\frac{A,B,C \vdash C}{A,B,C \vdash B} \stackrel{[Id]}{=} \frac{\overline{A,B,C \vdash A}}{C \to \neg A,A,B,C \vdash \bot} \stackrel{[Id]}{=} \stackrel{[\neg L]}{=} \\ \frac{B \to C \to \neg A,A,B,C \vdash \bot}{B \to C \to \neg A,A,B,C \vdash \bot} \stackrel{[\land L]}{=} \\ \overline{B \to C \to \neg A,A,B \land C \vdash \bot} \stackrel{[\land L]}{=} \\ \overline{B \to C \to \neg A,A \vdash \neg (B \land C)} \stackrel{[\neg R]}{=} \\ \overline{B \to C \to \neg A \vdash A \to \neg (B \land C)} \stackrel{[\to R]}{=} \\ \overline{B \to C \to \neg A \vdash A \to \neg (B \land C)} \stackrel{[\to R]}{=} \\ \overline{B \to C \to \neg A \vdash A \to \neg (B \land C)} \stackrel{[\to R]}{=} \\ \overline{B \to C \to \neg A \vdash A \to \neg (B \land C)} \stackrel{[\to R]}{=} \\ \overline{B \to C \to \neg A \vdash A \to \neg (B \land C)} \stackrel{[\to R]}{=} \\ \overline{B \to C \to \neg A \vdash A \to \neg (B \land C)} \stackrel{[\to R]}{=} \\ \overline{B \to C \to \neg A \vdash A \to \neg (B \land C)} \stackrel{[\to R]}{=} \\ \overline{B \to C \to \neg A \vdash A \to \neg (B \land C)} \stackrel{[\to R]}{=} \\ \overline{B \to C \to \neg A \vdash A \to \neg (B \land C)} \stackrel{[\to R]}{=} \\ \overline{B \to C \to \neg A \vdash A \to \neg (B \land C)} \stackrel{[\to R]}{=} \\ \overline{B \to C \to \neg A \vdash A \to \neg (B \land C)} \stackrel{[\to R]}{=} \\ \overline{B \to C \to \neg A \vdash A \to \neg (B \land C)} \stackrel{[\to R]}{=} \\ \overline{B \to C \to \neg A \vdash A \to \neg (B \land C)} \stackrel{[\to R]}{=} \\ \overline{B \to C \to \neg A \vdash A \to \neg (B \land C)} \stackrel{[\to R]}{=} \\ \overline{B \to C \to \neg A \vdash A \to \neg (B \land C)} \stackrel{[\to R]}{=} \\ \overline{B \to C \to \neg A \vdash A \to \neg (B \land C)} \stackrel{[\to R]}{=} \\ \overline{B \to C \to \neg A \vdash A \to \neg (B \land C)} \stackrel{[\to R]}{=} \\ \overline{B \to C \to \neg A \vdash A \to \neg (B \land C)} \stackrel{[\to R]}{=} \\ \overline{B \to C \to \neg A \vdash A \to \neg (B \land C)} \stackrel{[\to R]}{=} \\ \overline{B \to C \to \neg A \vdash A \to \neg (B \land C)} \stackrel{[\to R]}{=} \\ \overline{B \to C \to \neg A \vdash A \to \neg (B \land C)} \stackrel{[\to R]}{=} \\ \overline{B \to C \to \neg A \vdash A \to \neg (B \land C)} \stackrel{[\to R]}{=} \\ \overline{B \to C \to \neg A \vdash A \to \neg (B \land C)} \stackrel{[\to R]}{=} \\ \overline{B \to C \to \neg A \vdash A \to \neg (B \to C)} \stackrel{[\to R]}{=} \\ \overline{B \to C \to \neg A \vdash A \to \neg (B \to C)} \stackrel{[\to R]}{=} \\ \overline{B \to C \to \neg A \vdash A \to \neg (B \to C)} \stackrel{[\to R]}{=} \\ \overline{B \to C \to \neg A \vdash A \to \neg (B \to C)} \stackrel{[\to R]}{=} \\ \overline{B \to C \to \neg A \vdash A \to \neg (B \to C)} \stackrel{[\to R]}{=} \\ \overline{B \to C \to \neg A \vdash A \to \neg (B \to C)} \stackrel{[\to R]}{=} \\ \overline{B \to C \to \neg A \vdash A \to \neg (B \to C)} \stackrel{[\to R]}{=} \\ \overline{B \to C \to \neg A \vdash A \to \neg (B \to C)} \stackrel{[\to R]}{=} \\ \overline{B \to C \to \neg A \vdash A \to \neg (B \to C)} \stackrel{[\to R]}{=} \\ \overline{B \to C \to \neg A \vdash A \to \neg (B \to C)} \stackrel{[\to R]}{=} \\ \overline{B \to C \to \neg A \vdash A \to \neg (B \to C)} \stackrel{[\to R]}{=} \\ \overline{B \to C \to \neg A \vdash A \to \neg (B \to C)} \stackrel{[\to R]}{=} \\ \overline{B \to C \to \neg A \vdash A \to \neg (B \to C)} \stackrel{[\to R]}{=} \\ \overline{B \to C \to \neg A \to \neg (B \to C)} \stackrel{[\to R]}{=} \\ \overline{B \to C \to \neg A \to \neg (B \to C)} \stackrel{[\to R]}{=} \\ \overline{B \to C \to \neg A \to \neg (B \to C)} \stackrel{[\to R]}{=} \\ \overline{B \to C \to \neg A \to \neg (B \to C)} \stackrel{[\to R]}{=} \\ \overline{B \to$$

Provide a Sequent Calculus proof of $(B \to C \to \neg A) \to A \to \neg (B \land C)$

Here is a proof:

e is a proof:
$$\frac{\overline{A,B,C\vdash C}}{A,B,C\vdash B} \stackrel{[Id]}{=} \frac{\overline{A,B,C\vdash A}}{C \to \neg A,A,B,C\vdash \bot} \stackrel{[Id]}{=} \frac{}{\neg L]} \stackrel{[\neg L]}{=} \frac{}{[\neg R]} \stackrel{[\neg R]}{=} \stackrel{[\neg R]}{=} \frac{}{[\neg R]} \stackrel{[\neg R]}{=} \stackrel{[\neg R]}{=} \frac{}{[\neg R]} \stackrel{[\neg R]}{=} \stackrel{[\neg R]}{=} \frac{}{[\neg R]} \stackrel{[\neg R]}{=} \stackrel{[\neg R]}{=}$$

Note that compared to the Natural Deduction proof, we only have to eliminate $B \wedge C$ once here

▶ sequents are useful to keep track of available hypotheses

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Conclusion

What did we cover today?

- Sequent Calculus vs. Natural Deduction
- Sequent Calculus proofs
- Natural Deduction proofs

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Further reading

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- ► Chapter 3 of http://leanprover.github.io/logic_and_proof/

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What did we cover today?

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Next time?

Classical reasoning