

Math Week 1

Lecture 1 The natural numbers

The pitfalls of computer arithmetic

- $10^{10} = 1,1410,065,408$ (in C program)

1.1 The laws of arithmetic

- neutral elements $a + 0 = a$ $a \times 1 = a$
- commutativity $a + b = b + a$ $a \times b = b \times a$
- associativity $(a + b) + c = a + (b + c)$ $(a \times b) \times c = a \times (b \times c)$
- distributivity $a \times (b + c) = a \times b + a \times c$
- annihilation $a \times 0 = 0$

These actually hold for computer integers as well

1.2 Beyond equations

- additive cancellation $a + c = b + c \Rightarrow a = b$
- multiplicative cancellation $c \neq 0 \ \& \ a \times c = b \times c \Rightarrow a = b$

Additive cancellation also holds for computer integers

But multiplicative cancellation does not hold for computer integers

- Peano's Axioms
 1. 0 is a natural number
 2. If a is a natural number then so is $s(a)$
 3. A number of the form $s(a)$ is always different from 0
 4. If $s(a)$ and $s(b)$ are equal, then a and b are equal
 5. If $P(x)$ is a property of natural numbers that (ground case) holds of 0, and (inductive step) holds of $s(x)$ whenever it holds of x then P holds of all the natural numbers

The Axiom of Induction can be used to prove that some properties $P(a)$ are true for all natural numbers a

1.3 Place value systems

- Place value system
 - base $b > 0$ and digits $0, 1, \dots, b-1$
$$d_n \times b^n + d_{n-1} \times b^{n-1} + \dots + d_1 \times b^1 + d_0 \times b^0$$
or
$$d_n b^n + d_{n-1} b^{n-1} + \dots + d_1 b + d_0$$

1.4 Natural number representation in a computer

- The bit pattern as the digit representation of a number in base 2
- What happens when the result of an operation requires more than 32 digits (There are only 32 bits to represent binary digits)?
 - i. The excess digits are available for only a short moment in the cpu
 - ii. Java ignores them

Left associative: Read from left to right

Right associative: Read from right to left

Ambiguous: without either Precedence or Associativity

Lecture 2 The integers

2.1 The arithmetic laws of integers

- We only need to assume that for every integer a there is another integer denoted by $-a$ for which $a + (-a) = 0$ holds
- From this we can prove (additive) cancellation
- From this we can prove annihilation

- From this we can prove double negation: $\neg(\neg a) = a$
- From this we can prove minus times minus equals plus:
 $(-a) \times (-b) = a \times b$

2.2 Rings

The laws of rings

$a + 0 = a$	$a \times 1 = a$	(neutral elements)
$a + b = b + a$	$a \times b = b \times a$	(commutativity)
$a + (-a) = 0$		(additive inverse)
$(a + b) + c = a + (b + c)$	$(a \times b) \times c = a \times (b \times c)$	(associativity)
$a \times (b + c) = a \times b + a \times c$		(distributivity)

2.3 Integers in computers

- Java's int variables are based on 32-bit registers
- All calculations are done modulo 2^{32}
- The bit patterns from 100...000 to 111...111 are interpreted as **negative numbers**

2.4 Modulo arithmetic

- Computing "modulo m" can be done for any $m > 1$. We get the ring Z_m which has exactly m different elements
- Calculations in Z_m can be thought of in two different ways:
 - i. We can take the numbers from 0 to $m - 1$ as the standard members of Z_m , perform calculations with them as we would in Z , then reduce the result to an answer between 0 and $m - 1$ at the end. Example in Z_7

$$3 \times 5 = 15 \text{ in } Z$$

$$= 1 \text{ modulo } 7$$
 so in Z_7 we have $3 \times 5 = 1$
 - ii. Alternatively, we can do all calculations in Z and use $=$ for comparisons, instead of \equiv .
- Computer integers implement calculations in $Z_{2^{32}}$ and adopt the first approach internally, but when reporting the result back to the user, the numbers between 2^{31} and $2^{32} - 1$ are converted to negative numbers by subtracting 2^{32}