

Mathematical and Logical Foundations of Computer Science – Rules

University of Birmingham

November 24, 2021

1 Propositional Logic

1.1 Constructive Natural Deduction

$$\begin{array}{c}
 \frac{}{\perp} \quad [\perp E] \quad \quad \quad \frac{}{\top} \quad [\top I] \quad \quad \quad \frac{\overline{A}^1 \quad \vdots \quad B}{A \rightarrow B}^1 \quad [\rightarrow I] \quad \quad \quad \frac{A \rightarrow B \quad A}{B} \quad [\rightarrow E] \\
 \\
 \frac{\overline{A}^1 \quad \vdots \quad \perp}{\neg A}^1 \quad [\neg I] \quad \quad \quad \frac{\neg A \quad A}{\perp} \quad [\neg E] \\
 \\
 \frac{A}{A \vee B} \quad [\vee I_L] \quad \quad \quad \frac{A}{B \vee A} \quad [\vee I_R] \quad \quad \quad \frac{A \vee B \quad A \rightarrow C \quad B \rightarrow C}{C} \quad [\vee E] \\
 \\
 \frac{A \quad B}{A \wedge B} \quad [\wedge I] \quad \quad \quad \frac{A \wedge B}{B} \quad [\wedge E_R] \quad \quad \quad \frac{A \wedge B}{A} \quad [\wedge E_L]
 \end{array}$$

1.2 Constructive Sequent Calculus (Derived Rules)

$$\begin{array}{c}
 \frac{\Gamma_1, \Gamma_2 \vdash A \quad \Gamma_1, B, \Gamma_2 \vdash C}{\Gamma_1, A \rightarrow B, \Gamma_2 \vdash C} \quad [\rightarrow L] \quad \quad \quad \frac{\Gamma, A \vdash B}{\Gamma \vdash A \rightarrow B} \quad [\rightarrow R] \\
 \\
 \frac{\Gamma_1, \Gamma_2 \vdash A}{\Gamma_1, \neg A, \Gamma_2 \vdash B} \quad [\neg L] \quad \quad \quad \frac{\Gamma, A \vdash \perp}{\Gamma \vdash \neg A} \quad [\neg R] \\
 \\
 \frac{\Gamma_1, A, \Gamma_2 \vdash C \quad \Gamma_1, B, \Gamma_2 \vdash C}{\Gamma_1, A \vee B, \Gamma_2 \vdash C} \quad [\vee L] \quad \quad \quad \frac{\Gamma \vdash A}{\Gamma \vdash A \vee B} \quad [\vee R_1] \quad \quad \quad \frac{\Gamma \vdash A}{\Gamma \vdash B \vee A} \quad [\vee R_2] \\
 \\
 \frac{\Gamma_1, A, B, \Gamma_2 \vdash C}{\Gamma_1, A \wedge B, \Gamma_2 \vdash C} \quad [\wedge L] \quad \quad \quad \frac{\Gamma \vdash A \quad \Gamma \vdash B}{\Gamma \vdash A \wedge B} \quad [\wedge R] \\
 \\
 \frac{}{\Gamma_1, A, \Gamma_2 \vdash A} \quad [Id] \quad \quad \quad \frac{\Gamma \vdash B \quad B, \Gamma \vdash A}{\Gamma \vdash A} \quad [Cut] \\
 \\
 \frac{\Gamma_1, B, A, \Gamma_2 \vdash C}{\Gamma_1, A, B, \Gamma_2 \vdash C} \quad [X] \quad \quad \quad \frac{\Gamma_1, \Gamma_2 \vdash B}{\Gamma_1, A, \Gamma_2 \vdash B} \quad [W] \quad \quad \quad \frac{\Gamma_1, A, A, \Gamma_2 \vdash B}{\Gamma_1, A, \Gamma_2 \vdash B} \quad [C]
 \end{array}$$

1.3 Classical Natural Deduction

It includes all the Constructive Natural Deduction rules, plus:

$$\frac{}{A \vee \neg A} \quad [LEM] \quad \quad \quad \frac{\neg \neg A}{A} \quad [DNE]$$

1.4 Classical Sequent Calculus – 1st version

It includes all the Constructive Sequent Calculus rules, plus:

$$\frac{}{\Gamma \vdash A \vee \neg A} \quad [LEM] \qquad \frac{\Gamma \vdash \neg \neg A}{\Gamma \vdash A} \quad [DNE]$$

1.5 Classical Sequent Calculus – 2nd version (Derived Rules)

$$\begin{array}{c} \frac{\Gamma_1, \Gamma_2 \vdash A, \Delta_1 \quad \Gamma_1, B, \Gamma_2 \vdash \Delta_2}{\Gamma_1, A \rightarrow B, \Gamma_2 \vdash \Delta_1, \Delta_2} \quad [\rightarrow L] \quad \frac{\Gamma, A \vdash \Delta_1, B, \Delta_2}{\Gamma \vdash \Delta_1, A \rightarrow B, \Delta_2} \quad [\rightarrow R] \quad \frac{\Gamma_1, \Gamma_2 \vdash A, \Delta}{\Gamma_1, \neg A, \Gamma_2 \vdash \Delta} \quad [\neg L] \\ \\ \frac{\Gamma_1, A, \Gamma_3 \vdash \Delta_1 \quad \Gamma_2, B, \Gamma_4 \vdash \Delta_2}{\Gamma_1, \Gamma_2, A \vee B, \Gamma_3, \Gamma_4 \vdash \Delta_1, \Delta_2} \quad [\vee L] \quad \frac{\Gamma \vdash \Delta_1, A, B, \Delta_2}{\Gamma \vdash \Delta_1, A \vee B, \Delta_2} \quad [\vee R] \quad \frac{\Gamma, A \vdash \Delta_1, \Delta_2}{\Gamma \vdash \Delta_1, \neg A, \Delta_2} \quad [\neg R] \\ \\ \frac{\Gamma_1, A, B, \Gamma_2 \vdash \Delta}{\Gamma_1, A \wedge B, \Gamma_2 \vdash \Delta} \quad [\wedge L] \quad \frac{\Gamma_1 \vdash \Delta_1, A, \Delta_3 \quad \Gamma_2 \vdash \Delta_2, B, \Delta_4}{\Gamma_1, \Gamma_2 \vdash \Delta_1, \Delta_2, A \wedge B, \Delta_3, \Delta_4} \quad [\wedge R] \quad \frac{}{\Gamma_1, A, \Gamma_2 \vdash \Delta_1, A, \Delta_2} \quad [Id] \\ \\ \frac{\Gamma_1 \vdash B, \Delta_1 \quad \Gamma_2, B \vdash \Delta_2}{\Gamma_1, \Gamma_2 \vdash \Delta_1, \Delta_2} \quad [Cut] \quad \frac{\Gamma_1, B, A, \Gamma_2 \vdash \Delta}{\Gamma_1, A, B, \Gamma_2 \vdash \Delta} \quad [X_L] \quad \frac{\Gamma \vdash \Delta_1, B, A, \Delta_2}{\Gamma \vdash \Delta_1, A, B, \Delta_2} \quad [X_R] \\ \\ \frac{\Gamma_1, \Gamma_2 \vdash \Delta}{\Gamma_1, A, \Gamma_2 \vdash \Delta} \quad [W_L] \quad \frac{\Gamma_1, A, A, \Gamma_2 \vdash \Delta}{\Gamma_1, A, \Gamma_2 \vdash \Delta} \quad [C_L] \quad \frac{\Gamma \vdash \Delta_1, \Delta_2}{\Gamma \vdash \Delta_1, A, \Delta_2} \quad [W_R] \quad \frac{\Gamma \vdash \Delta_1, A, A, \Delta_2}{\Gamma \vdash \Delta_1, A, \Delta_2} \quad [C_R] \end{array}$$

2 Predicate Logic

2.1 Natural Deduction

The Natural Deduction rules for Predicate Logic include all Proposition Logic rules plus the following rules:

$$\frac{P[x \setminus y]}{\forall x. P} \quad [\forall I] \quad \frac{\forall x. P}{P[x \setminus t]} \quad [\forall E] \quad \frac{P[x \setminus t]}{\exists x. P} \quad [\exists I] \quad \frac{\exists x. P \quad \overline{P[x \setminus y]}^1 \quad \vdots \quad Q}{Q} \quad [\exists E]$$

Side conditions:

- for $[\forall I]$: y must not be free in any not-yet-discharged hypothesis or in $\forall x. P$
- for $[\forall E]$: $\mathbf{fv}(t)$ must not clash with $\mathbf{bv}(P)$
- for $[\exists I]$: $\mathbf{fv}(t)$ must not clash with $\mathbf{bv}(P)$
- for $[\exists E]$: y must not be free in Q or in not-yet-discharged hypotheses or in $\exists x. P$

2.2 Sequent Calculus (Derived Rules)

The Sequent Calculus rules for Predicate Logic include all Propositional Logic rules plus the following rules:

$$\frac{\Gamma \vdash P[x \setminus y]}{\Gamma \vdash \forall x. P} \quad [\forall R] \quad \frac{\Gamma_1, P[x \setminus t], \Gamma_2 \vdash Q}{\Gamma_1, \forall x. P, \Gamma_2 \vdash Q} \quad [\forall L] \quad \frac{\Gamma \vdash P[x \setminus t]}{\Gamma \vdash \exists x. P} \quad [\exists R] \quad \frac{\Gamma_1, P[x \setminus y], \Gamma_2 \vdash Q}{\Gamma_1, \exists x. P, \Gamma_2 \vdash Q} \quad [\exists L]$$

Side conditions:

- for $[\forall R]$: y must not be free in Γ or $\forall x. P$
- for $[\forall L]$: $\mathbf{fv}(t)$ must not clash with $\mathbf{bv}(P)$
- for $[\exists R]$: $\mathbf{fv}(t)$ must not clash with $\mathbf{bv}(P)$
- for $[\exists L]$: y must not be free in Γ_1, Γ_2, Q , or $\exists x. P$