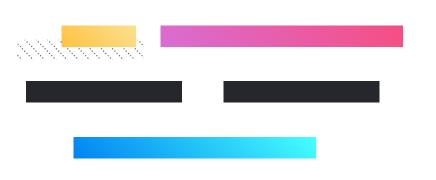
#### **Theory of Computing:**

#### 10. Turing Machine - 2



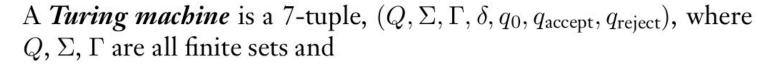
#### **Professor Imed Bouchrika**

National School of Artificial Intelligence imed.bouchrika@ensia.edu.dz

#### **Outline:**

- TM as
  - Language Acceptor/Recognizer
  - Transducer and Examples
  - Problem Solver
- Variations of Turing Machines
- Universal Turing Machine
- Algorithms and Church's Thesis
- Turing-Complete Systems

### Formalism of Turing Machine



- **1.** Q is the set of states,
- **2.**  $\Sigma$  is the input alphabet not containing the *blank symbol*  $\Box$ ,
- **3.**  $\Gamma$  is the tape alphabet, where  $\sqcup \in \Gamma$  and  $\Sigma \subseteq \Gamma$ ,
- **4.**  $\delta: Q \times \Gamma \longrightarrow Q \times \Gamma \times \{L, R\}$  is the transition function,
- 5.  $q_0 \in Q$  is the start state,
- **6.**  $q_{\text{accept}} \in Q$  is the accept state, and
- 7.  $q_{\text{reject}} \in Q$  is the reject state, where  $q_{\text{reject}} \neq q_{\text{accept}}$ .

### Classes of Languages for Turing Machine

- The collection of strings that M <u>accept</u>s is the language of M, or the language recognized by M
  - A language is called **Turing-recognizable** if some Turing machine recognizes it
  - Mainly: Accepting words that belong to the language.
  - For words not in the language:
    - Reject or Loop

### Classes of Languages for Turing Machine

- Turing-decidable language or simply decidable if some Turing machine decides it
  - Halts and Accepts for words in the language
  - Halts and Reject for words not in the language
- Every Decidable language is also recognizable.

### Classes of Languages for Turing Machine

#### • Terminologies:

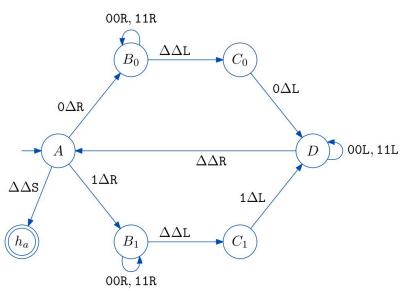
- Turing Recognizable is called a recursively enumerable language in some other textbooks.
- For turing decidable is called a recursive language

- What about the following language :
  - Even length palindromes
  - o Algorithm?
  - Turing Machine?

- What about the follo
  - Even length pali
  - Algorithm:

- 1. If first letter is:
  - 1.1. 0 → Remove and Skip to Blank at extremeRight, Move Left Remove 0
  - 1.2. 1 → Remove and Skip to Blank at extreme Right, Move Left Remove 1
  - 1.3. Blank → Accept
- Skip to Blank at extreme left.
- 3. Go to Step 1

- What about the following language:
  - Even length palindromes
  - Turing Machine



- What about the following language:
  - Odd-length palindromes
  - o Algorithm?
  - Turing Machine?

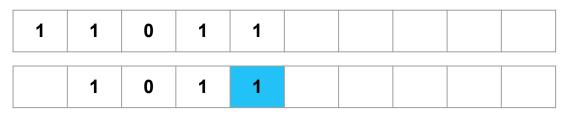


- Odd-length palindromes
- o Algorithm:

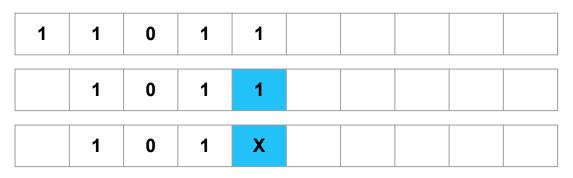
1	1	0	1	1			



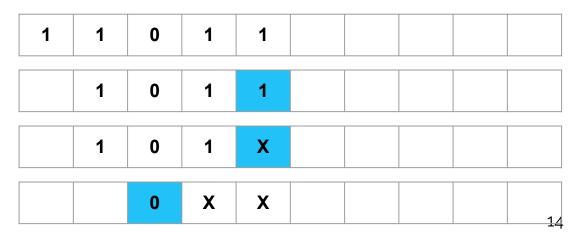
- Odd-length palindromes
- o Algorithm:



- What about the following language:
  - Odd-length palindromes
  - o Algorithm:



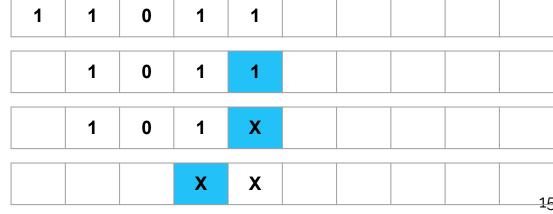
- What about the following language:
  - Odd-length palindromes
  - o Algorithm:



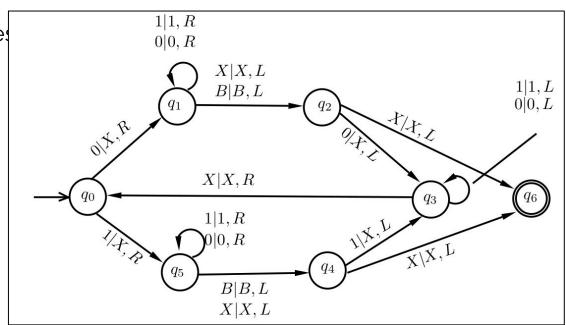
What about the following language:

If next transition takes directly to X, Just Accept

romes

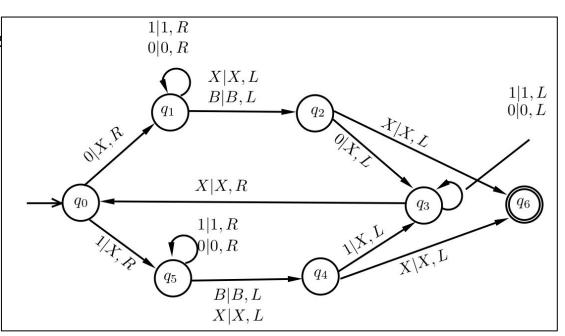


- What about the following language:
  - Odd-length palindromes
  - Turing Machine?



- What about the following language:
  - Odd-length palindromes
  - Turing Machine?

This Turing replaces even blanks at the left with X



- Given an input string composed of two words separated by the symbol # design the algorithm to check if they are equal
  - $\circ$  If Equal  $\rightarrow$  Accept
  - Otherwise → Reject

- Given an input string composed of two words separated by the symbol # design the algorithm to check if they are equal
  - $\circ$  If Equal  $\rightarrow$  Accept
  - Otherwise → Reject

- ? What's the algorithm?
  - Output Description 

    Output Description
  - How to inject the symbol # into the middle? or other approach?
    - Example:
      - 101001 ⇒ 101#001
      - 101001 ⇒ BABXXZ
      - ..?

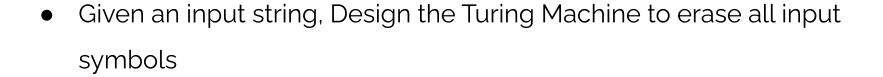
### **Turing Machine as Transducers**

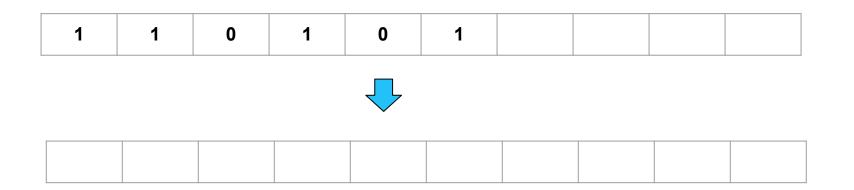
- Transducers are the devices used to convert one form of a signal into a different form.
- Transducer is a type of Turing Machine that is used to convert the given input into a specific output after the machine performs various read-writes.
- It doesn't accept or reject an input but performs series of operations to obtain the output right in the same tape and halts when finished.

#### **Turing Machine as Transducers**



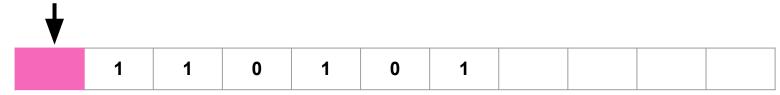
- Computation
  - Example:
    - Addition
- Operation on input strings and text processing
  - Example:
    - Removing special patterns from a string

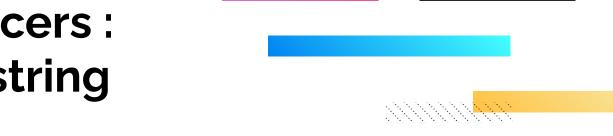




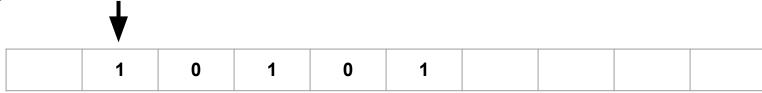


- Given an input string, Design TM to erase all input symbols
- Algorithm:





- Given an input string, Design Turing Machine to erase all input symbols
- Algorithm:





- Given an input string, Design Turing Machine to erase all input symbols
- Algorithm:

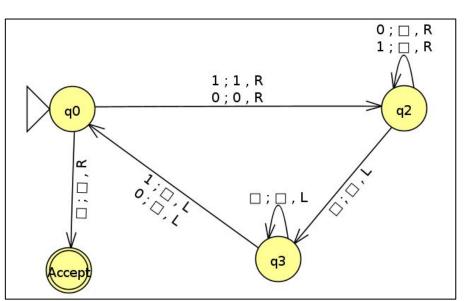


- Given an input string, Design Turing Machine to erase all input symbols
- Algorithm:

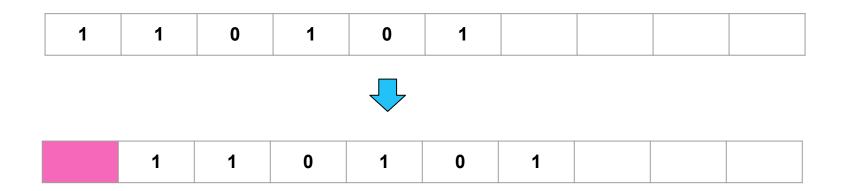


How to restore the head to its original position?
We skip left till we find what?

- Given an input string, Design Turing Machine to erase all input symbols
- Algorithm:



• Given an input string, design the algorithm to Shift all all symbols to the right by one cell.

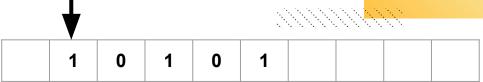


- Given an input string, design the algorithm to Shift all all symbols to the right by one cell.
- Algorithm :



0

- Given an input string, design the algorithm to Shift all all symbols to the right by one cell.
- Algorithm :



```
On Branch A:

If 1:

Keep it

Move Right On Branch A

If 0:

Invert to 1

Move Right on Branch B
```

- Given an input string, design the algorithm to Shift all all symbols to the right by one cell.
- Algorithm :



```
On Branch A:

If 1:

Keep it

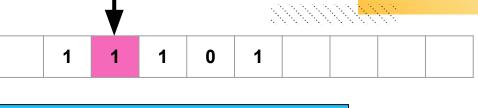
Move Right On Branch A

If 0:

Invert to 1

Move Right on Branch B
```

- Given an input string, design the algorithm to Shift all all symbols to the right by one cell.
- Algorithm :



```
On Branch B:

If 0:

Keep it

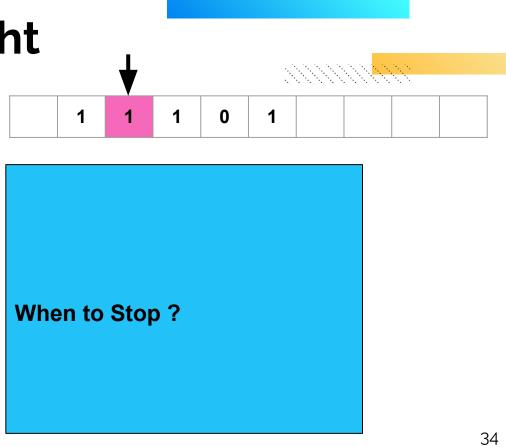
Move Right On Branch B

If 1:

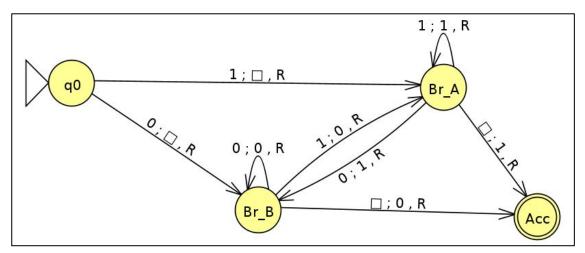
Invert to 0

Move Right on Branch A
```

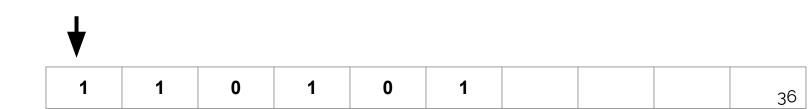
- Given an input string, design the algorithm to Shift all all symbols to the right by one cell.
- Algorithm :



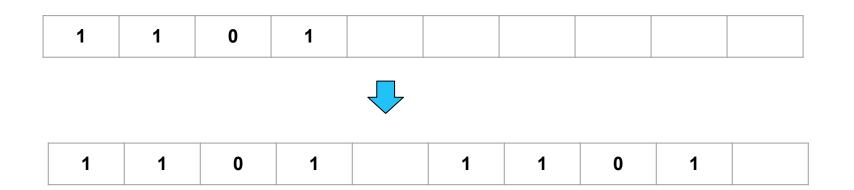
- Given an input string, design the algorithm to Shift all symbols to the right by one cell.
- Turing Machine :



- Given an input string, design the algorithm to Shift all all symbols to the right by one cell.
- Turing Machine:



 Given an input string, design the algorithm to copy the full string separated with a blank to the original string.



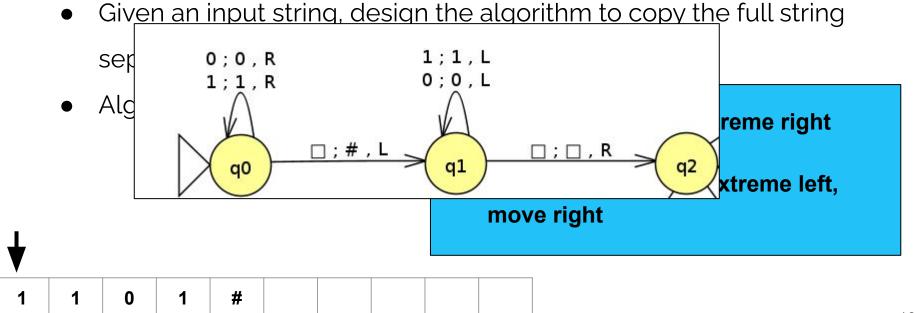
- Given an input string, design the algorithm to copy the full string separated with a blank to the original string.
- Algorithm?
- Turing Machine?

- Given an input string, design the algorithm to copy the full string separated with a blank to the original string.
- Algorithm:

- 1. Go to the Blank at extreme right
- 2. Replace it with #
- Go back to Blank at extreme left, move right



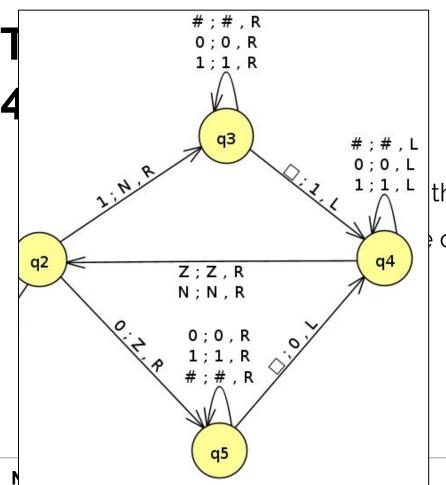
•							
1	1	0	1	#			



- Given an input string, design the algorithm to copy the full string separated with a blank to the original string.
- Algorithm:

- Replace 1 with N , Go to extreme right, write 1
- 2. Replace 0 with Z, Go to extreme right , write 0
- Skip until you find either N or Z, move Right





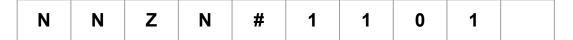
the algorithm to copy the full string

<del>priminal etrina</del>

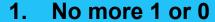
- 1. Replace 1 with N, Go to extreme right, write 1
- 2. Replace 0 with Z, Go to extreme right , write 0
- Skip until you find either N or Z, move Right

- Given an input string, design the algorithm to copy the full string separated with a blank to the original string.
- Algorithm:

- 1. No more 1 or 0
- 2. Invert:
  - a.  $N \rightarrow 1$
  - b.  $Z \rightarrow 0$
  - c. # → Blank



- Given an input string, design the algorithm to copy the full string separated with a blank to the original string.
- Algorithm:



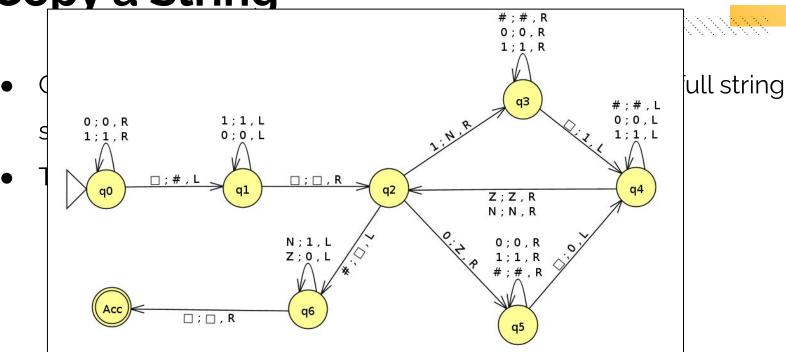
- 2. Invert:
  - a.  $N \rightarrow 1$
  - b.  $Z \rightarrow 0$
  - c. # → Blank



1	1	0	1	#	1	1	0	1	

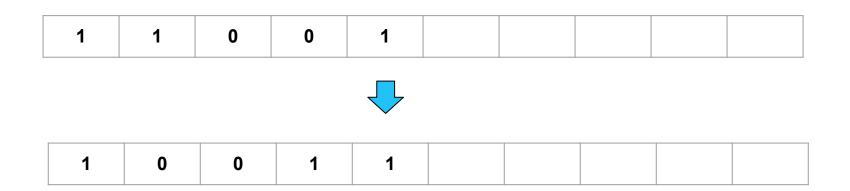
#### TM as Transducers:

4. Copy a String

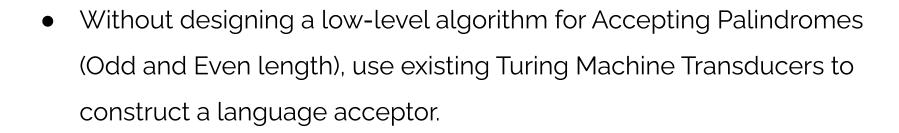


# TM as Transducers: 5. Reversing a String

Given an input string, reverse it or flip it.



#### **Combining Turing Machines**

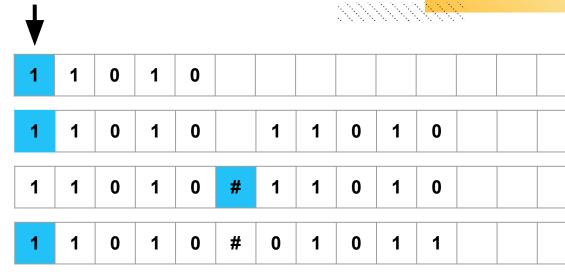


# **Combining Turing Machines**

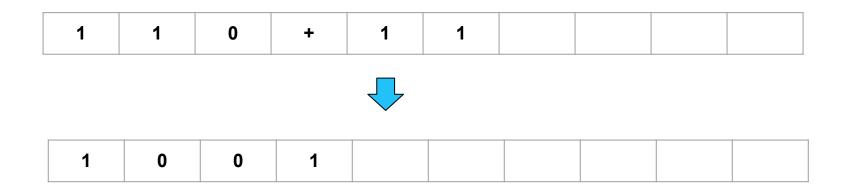
- Without designing a low-level algorithm for Accepting Palindromes, use existing Turing Machine Transducers to construct a language acceptor.
- Turing Machines to utilize:
  - Copy → Find\_Blank\_Replace\_# → Reverse → Equal\_Strings

# Combining Turing Machines

- Turing Machines to utilize:
  - Copy
  - Find\_Blank\_Replace\_#
  - Reverse
  - $\circ$  Equal\_Strings :  $\rightarrow$  Accept or  $\rightarrow$  reject



 Given an input string as binary number + binary number, Write the algorithm to do the addition



- Given an input string as binary number + binary number, Write the algorithm to do the addition
- Algorithm:



- Given an input string as binary number + binary number,
- Write the algorithm to do the addition



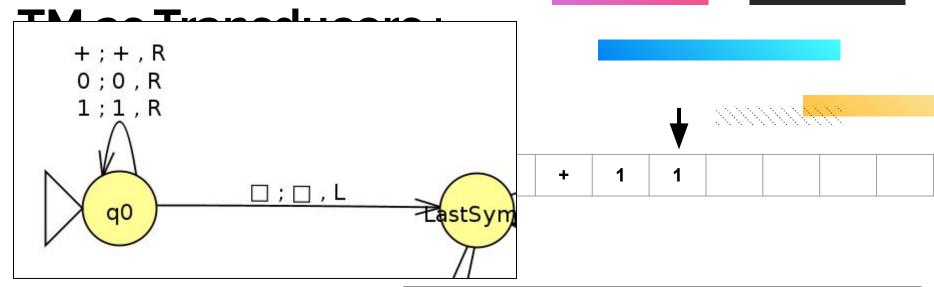
- Given an input string as binary number + binary number,
- Write the algorithm to do the addition



- 1. Go to blank at extreme right, move left one cell.
- 2. If:

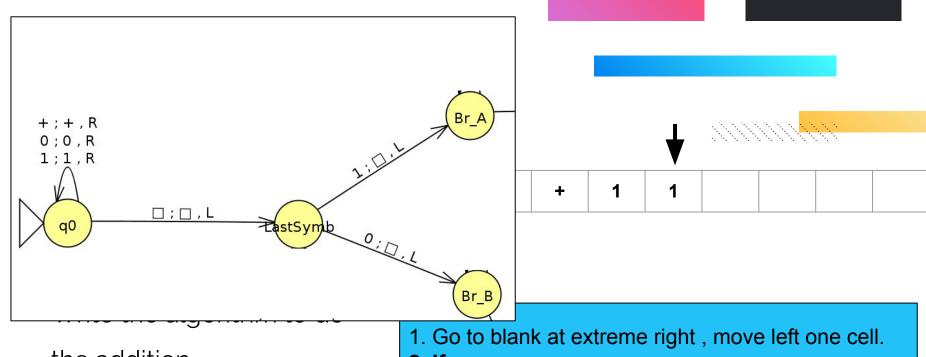
1 take a branch A 0 take branch B

3. Erase the Cell



 Write the algorithm to do the addition

- 1. Go to blank at extreme right, move left one cell.
- 2. If:
  - 1 take a branch A
    - 0 take branch B
- 3. Erase the Cell



the addition

2. If:

1 take a branch A 0 take branch B

3. Erase the Cell

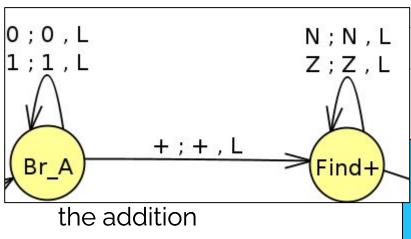
- Given an input string as binary number + binary number,
- Write the algorithm to do the addition



- Given an input string as binary number + binary number,
- Write the algorithm to do the addition



1. Skip Left 1 and 0 till we reach +:

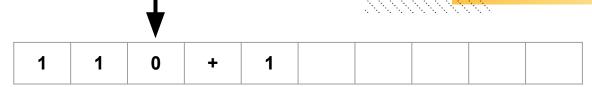


1. Skip Left 1 and 0 till we reach + :

0

2. We Keep skipping N and Z to the left

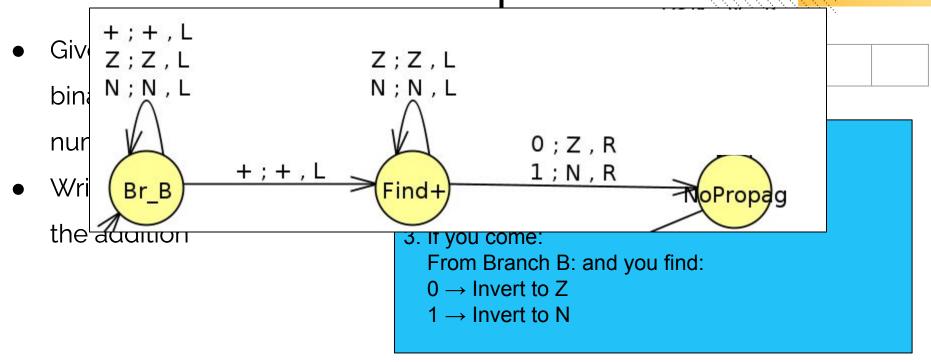
- Given an input string as binary number + binary number,
- Write the algorithm to do the addition



- 1. Skip Left 1 and 0 till we reach +:
- 2. We Keep skipping N and Z to the left
- 3. If you come:

From Branch B: and you find:

- $0 \rightarrow Invert to Z$
- $1 \rightarrow Invert to N$



- Given an input string as binary number + binary number,
- Write the algorithm to do the addition



- 1. Skip Left 1 and 0 till we reach +:
- 2. We Keep skipping N and Z to the left
- 3. If you come:

From Branch **A**: and you find:

- $0 \rightarrow Invert to N$
- $1 \rightarrow Invert to Z ??$

- Given an input string as binary number + binary number,
- Write the algorithm to do the addition



- 1. Skip Left 1 and 0 till we reach +:
- 2. We Keep skipping N and Z to the left
- 3. If you come:

From Branch **A** : and you find:

- $0 \rightarrow Invert to N$
- $1 \rightarrow Invert to Z and$

keep looping left:

$$1 \rightarrow 0$$

 $0 \rightarrow 1$  and break from Loop

- Given an input string as binary number + binary number.
- Write the algorithm to do the addition



- 1. Skip Left 1 and 0 till we reach +:
- 2. We Keep skipping N and Z to the left
- 3. If you come:

From Branch **A**: and you find:

- $0 \rightarrow Invert to N$
- 1 → Invert to Z and

keep looping left:

$$1 \rightarrow 0$$

0 or blank→ 1 and break from Loop

3. Go to Blank at Extreme right and repeat

- Given an input string as binary number + binary number.
- Write the algorithm to do the addition



- 1. Skip Left 1 and 0 till we reach +:
- 2. We Keep skipping N and Z to the left
- 3. If you come:

From Branch **A**: and you find:

- $0 \rightarrow Invert to N$
- $1 \rightarrow Invert to Z and$

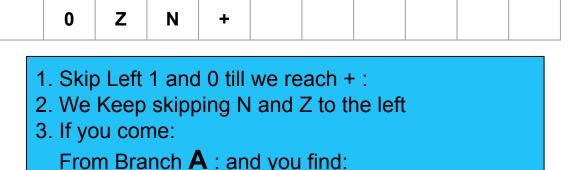
keep looping left:

$$1 \rightarrow 0$$

0 or blank→ 1 and break from Loop

3. Go to Blank at Extreme right and repeat

- Given an input string as binary number + binary number.
- Write the algorithm to do the addition



0 or blank  $\rightarrow$  1 and break from Loop

5

 $0 \rightarrow Invert to N$ 

 $1 \rightarrow Invert to Z and$ 

keep looping left:

3. Go to Blank at Extreme right and repeat

 $1 \rightarrow 0$ 

- Given an input string as binary number + binary number,
- Write the algorithm to do the addition



- 1. Skip Left 1 and 0 till we reach +:
- 2. We Keep skipping N and Z to the left
- 3. If you come:

From Branch **A**: and you find:

- $0 \rightarrow Invert to N$
- $1 \rightarrow Invert to Z and$

keep looping left:

$$1 \rightarrow 0$$

0 or blank → 1 and break from Loop

3. Go to Blank at Extreme right and repeat

- Given an input string as binary number + binary number,
- Write the algorithm to do the addition



- Given an input string as binary number + binary number,
- Write the algorithm to do the addition

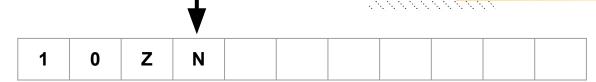


- 1. If + is on the extreme right.
  - 1. Delete the + sign
  - 2. Skip to extreme LEFT and Invert:

$$N \rightarrow 1$$

$$Z \rightarrow 0$$

- Given an input string as binary number + binary number,
- Write the algorithm to do the addition

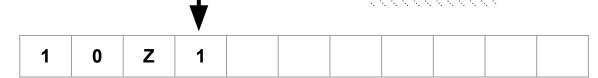


- 1. If + is on the extreme right.
  - 1. Delete the + sign
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$$N \rightarrow 1$$

$$Z \rightarrow 0$$

- Given an input string as binary number + binary number,
- Write the algorithm to do the addition

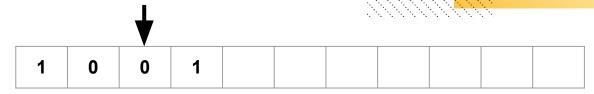


- 1. If + is on the extreme right.
  - 1. Delete the + sign
  - 2. Skip to extreme LEFT and Invert:

$$N \rightarrow 1$$

$$Z \rightarrow 0$$

- Given an input string as binary number + binary number,
- Write the algorithm to do the addition



- 1. If + is on the extreme right.
  - 1. Delete the + sign
  - 2. Skip to extreme LEFT and Invert:

$$N \rightarrow 1$$

$$Z \rightarrow 0$$

- Given an input string as binary number + binary number,
- Write the algorithm to do the addition

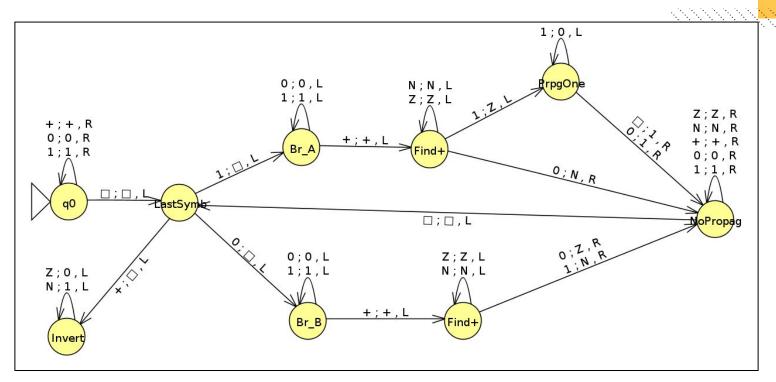


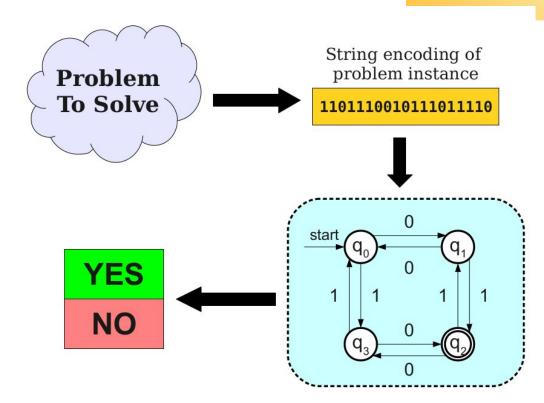
- 1. If + is on the extreme right.
  - 1. Delete the + sign
  - 2. Skip to extreme LEFT and Invert:

$$N \rightarrow 1$$

$$Z \rightarrow 0$$

## TM as Transducers : 6. Addition of Binaries

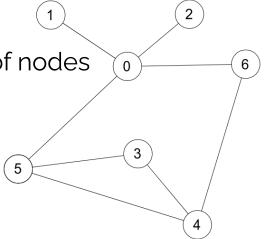




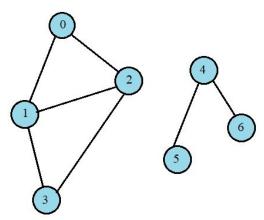


The problem: is the graph connected?

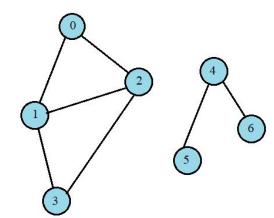
There is a link or path between any pair of nodes



- Given a graph with some nodes and links between the nodes
- The problem: is the graph connected?
  - There is a link or path between any pair of nodes.



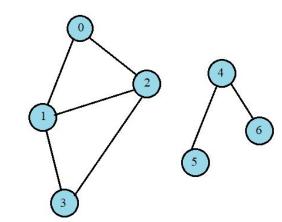
- Given a graph with some nodes and links between the nodes
- The problem: is the graph connected?
- How to construct Turing Machine to answer such Question?



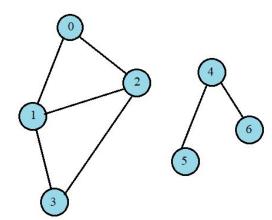
- Given a graph with some nodes and links between the nodes
- The problem: is the graph connected?
- How to construct Turing Machine to answer such

Question?

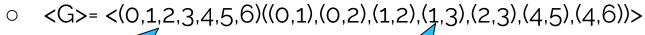
Important : FIRST how to encode the problem ?



- Given a graph G
  - Its encoding can denoted as: <G>
  - $\circ$  <G>= ?

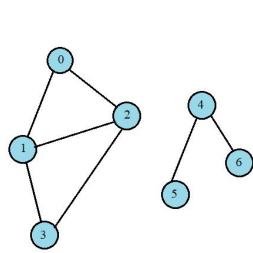


- Given a graph G
  - Its encoding can denoted as: <G>



List of nodes

List of edges



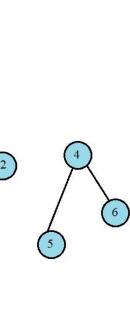
Given a graph G

Its enco

<G>= <(( any use any format:

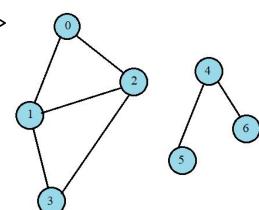
Unary (1111), binary (0101), decimal (1,2,3), hexadecimal (A) depending on the alphabet.

For the representation of numbers, feel free to

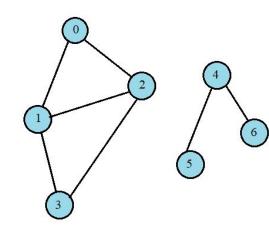


- Given a graph G
  - Its encoding can denoted as : <G>
  - $\circ$  <G>= <(0,1,2,3,4,5,6)((0,1),(0,2),(1,2),(1,3),(2,3),(4,5),(4,6))>
- Language of the connected graphs:

L = { <G> | G is a connected graph }



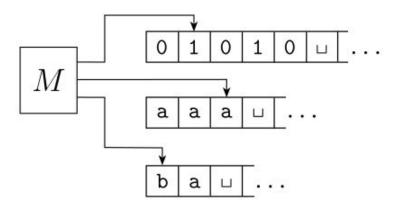
- Construction of the Turing Machine
- Better to use Subroutines:
  - Verify the encoding format
  - Verify that there is no repetition in the first part
  - Check that edges contain only existing nodes
  - Check the links between nodes



## Variations of Turing Machine

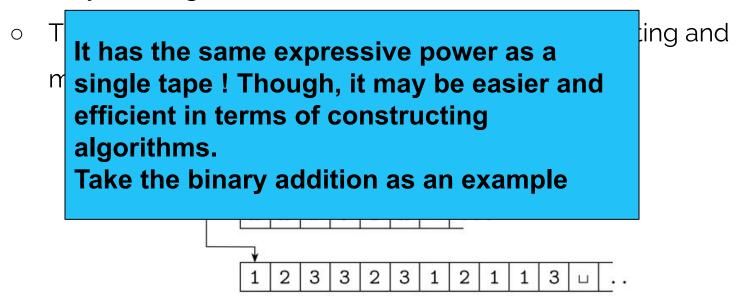
#### Multitape Turing Machine:

 The control can access multiple tapes reading, writing and moving the head of each tape at each step.



## Variations of Turing Machine

Multitape Turing Machine:



## Variations of Turing Machine

#### Nondeterministic Turing Machine:

- Where at a given state, there can be many possible transitions for the same symbol being read.
- The machine would keep track of all actions in the same way as an NFA.
- Every non-deterministic TM, there is an equivalent deterministic
   Turing machine → Nondeterministic TM has the same expressive
   power as a normal turing machine.

 You write a programming code in new Language X, it does some computation and return some input.

- You write a programming code in new Language X, it does some computation and return some input.
- You need to compile and execute the program written in X by compiler written in which language?

- Given any Turing Machine M to conduct some computation, there is the Universal Turing Machine to simulate or execute M for any given input w
- The universal Turing Machine:
  - Halt iff M halts on input w.
  - o If M is a deciding/semi-deciding machine, then
    - If M accepts, accept.
    - If M rejects, reject.
  - If M computes a function, then U (<M, w>) must equal M (w)

We can construct a universal TM that accepts the language
 L = {<M, w> | M is a TM and w ∈ L(M)}

### **Church-Turing Thesis**

- Turing Machines are an abstract model of computation, their purpose is to define in a mathematical way what problems are theoretically computable and which are not.
- In 1900, mathematician David Hilbert identified 23 mathematical problems and posed them as a challenge for the coming century.
- Hilbert's tenth problem was to devise an algorithm that tests whether
  a polynomial has an integral root. (Integers to be assigned to the polynomial
  variables to reach a value of zero)

### **Church-Turing Thesis**

- Hilbert did not use the term algorithm but rather "a process according to which it can be determined by a finite number of operations."
  - He assumed that such an algorithm must exist—someone need only find it.
  - But: it is algorithmically unsolvable.
- Church-Turing thesis provides the definition of algorithm necessary to resolve Hilbert's tenth problem

### **Church-Turing Thesis**

• Church-Turing thesis. There is an "effective procedure" for a problem if and only if there is a TM for the problem.

# Turing-Complete Systems



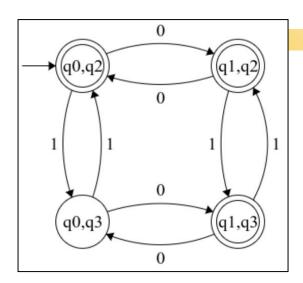
# Turing-Complete Systems

- Are there models of computation more powerful than Turing machines?
  - We do not know if there are more powerful models.
  - However, there are many computational models equivalent in power to TM's. They are called **Turing-complete systems**.
- A system of data-manipulation rules is said to be complete if it can be used to simulate any Turing machine

# Turing-Complete Systems

- Modern computers
- Church's lambda calculus.
- Gödel's µ-recursive functions
- Lindenmayer systems or L-systems.
- ..

Convert the following NFA to RegEx



```
Produce the context free grammar for the following language: The number of b+2 = the number of a. The number of a is more than b but with strictly only two letters. Example of words in the language: aa, aaba, baaa, baaaab,....
```

Let D =  $\{xy | x, y \in \{0,1\}* \text{ and } |x| = |y| \text{ but } x \neq y\}$ . Show that D is a context-free language.

Use the pumping lemma to show that the following languages are not context free.  $\{0^n\ 1^n\ 0^n\ 1^n\ |\ n\ \ge\ 0\}$ 

Use the pumping lemma to show that the following languages are not context free  $\{0^n\ \#0^{2n}\ \#0^{3n}\ |\ n\ge 0\}$ 



Prove that  $L = \{a^n \mid n \text{ is prime}\}\$ is not CFL.

```
Over alphabet {0,1}, Produce the RegExs for the languages:
1. L whose words do not contain the substring 101
Easy and non-brainy way -> DFA for the complement -> convert to RegEx
Hard way:
Enumerate possible strings that we need to accept ...
11111 → 1*
00000 \rightarrow 0*
1111000000 \rightarrow 1*0*
00000111111 \rightarrow 0*1*
001110000111000 \rightarrow \{000*,111*\}*
011111000001 → ...
Look for more possible words and later
      optimize
```

Over alphabet  $\{0,1\}$ , Produce the RegExs for the languages : L which does not contain the string 101

