Systems Programming Devices

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Topics

I/O Subsystem

Introduction
Device Types
I/O Software
Accessing Devices

Device Drivers

Interface Implementation Device Access

Topics

I/O Subsystem

Introduction

Device Types I/O Software Accessing Devices

Device Drivers

Interface Implementation Device Access

I/O Devices

- ▶ O/S controls all I/O devices
- issues commands to devices
- catches interrupts
- handles errors
- provides interface

Device Controllers

- devices consist of:
 - mechanical components
 - electronic components: device controller
- ▶ O/S deals with controller
 - connected through a standard interface
 - SCSI, USB, Firewire, ...

Controller Registers

- ▶ CPU communicates with the controller through registers
- data register: for sending/receiving data
- control register: for sending commands to device
- status register: for getting/setting the state of device

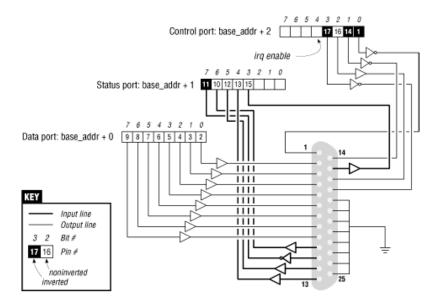
I/O Architecture

- ports: special address space for I/O
 - separate lines for I/O ports
 - special instructions for I/O
- memory-mapped: registers part of regular address space
 - directly-mapped: part of address space reserved for I/O
 - software-mapped: I/O space part of virtual memory

PC Parallel Interface

- ▶ parallel interface base addresses on a PC: 0x378, 0x278
- ports:
 - ▶ +0: bidirectional data register
 - ► +1: status register (read-only) online, out-of-paper, busy
 - ► +2: control register (write-only) enable/disable interrupts

Parallel Interface



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Device Types

- character devices
- block devices
- network interfaces
- clocks and timers

Character Devices

- a character device acts like a stream of characters
- arbitrary-sized data transfer
- ▶ not addressable: no seek operation

examples

- console, mouse
- sound card
- ► serial port, parallel port

Block Devices

- a block device can host a filesystem
- data transfer in fixed-size blocks
- each block has its own address
- read/write each block independently

example

disks

Device Type

▶ the device type is more the characteristic of the driver rather than the device itself

example: disk

- usually a block device
- ▶ it can also be used as a character device: tar

Topics

I/O Subsystem

I/O Software

Accessing Devices

Device Access

I/O Software

- blocking vs interrupt-driven
 - better for CPU to work interrupt-driven fashion
 - better for user-space programs to work in blocking fashion
 - easier to develop programs that work in blocking fashion
 - O/S makes interrupt-driven operations look blocking
- standardized interface
- uniform naming

I/O Software

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 - better for CPU to work interrupt-driven fashion
 - ▶ better for user-space programs to work in blocking fashion
 - easier to develop programs that work in blocking fashion
 - O/S makes interrupt-driven operations look blocking
- standardized interface
- uniform naming

- ▶ in Unix, every device has a device node
- under the /dev folder
- /dev/sda: first SCSI disk
- /dev/sdb: second SCSI disk
- /dev/sdb1: first partition of the second SCSI disk
- /dev/sdb2: second partition of the second SCSI disk
- /dev/parport0: first parallel port

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- device nodes have major and minor numbers
- major number identifies the driver
- minor number identifies the physical device
- ▶ all /dev/sd* devices have the same major number
- they all have different minor numbers
- (recently) major number alone doesn't identify driver
- ▶ major number + region of minor numbers

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- copy semantics: transfer the snapshot of data at the time of the I/O request
- scheduling: issue order may not be the best execution order
- buffering: adapt between different data transfer sizes
- caching
- spooling: deal with dedicated devices (e.g. printers)
 - a daemon for controlling the device
 - a spooling directory
- error handling

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Software Layers

- ► top-down:
- user-space applications
- device-independent software
- device drivers
- ▶ interrupt handlers

Device-Independent Software

- functions common to all devices
- uniform interface to user-level software
- device naming
- device protection
- provide device-independent block sizes
- buffering
- allocating and releasing dedicated devices
- error reporting

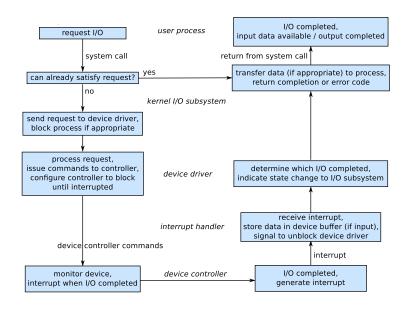
Device Drivers

- device-dependent code
- a driver for each device type
- accept request from device-independent software
- decide on sequence of controller operations

Interrupt Handlers

- interrupts hidden from rest of system
 - ▶ requesting process is blocked until I/O is completed
- ▶ when I/O is completed, interrupt occurs
 - process is made to unblock

I/O Life Cycle



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Accessing Devices

- directly: using ports or memory
- ▶ through device drivers: using the device driver interface

Direct Access

- ▶ input: inb, inw, inl
- ▶ output: outb, outw, outl
- ▶ get permission from O/S: ioperm system call

Direct Access

- ▶ input: inb, inw, inl
- output: outb, outw, outl
- ▶ get permission from O/S: ioperm system call

Direct Access Example

output to parallel interface

```
ioperm(0x378, 1, 255);
outb(0xff, 0x378);
```

Reading Material

- ► Silberschatz, 8/e
 - ► Chapter 13: I/O Systems

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Unix Device Driver Interface

- ▶ in Unix, the device driver interface is similar to the file interface
- ▶ open, close
- ▶ read, write

Device Specific Operations

- some operations are neither read nor write
- ioctl: issue command specific to device

device-specific operation examples

- ▶ eject CDROM
- make the speaker beep
- set communication parameters for modem

open:

- ► flags:
 - O_RDONLY O_WRONLY O_RDWR
 - 0_CREAT 0_APPEND
- mode: permissions
- returns: file descriptor

close:

```
int close(int fd);
```

▶ returns: success / failure status

read:

- returns: number of bytes read (x)
 - \triangleright x = count: successful completion
 - x = 0: end-of-file
 - ► *x* < 0: error
 - ightharpoonup 0 < x < count: partial transfer, retry remaining part

write:

- returns: number of bytes written (x)
 - $\rightarrow x = count$: successful completion
 - x = 0: end-of-file
 - ► *x* < 0: error
 - ightharpoonup 0 < x < count: partial transfer, retry remaining part

▶ ioctl:

parameter and return values depend on request

Device Access Example

output to parallel port

```
fd = open("/dev/parport0", O_WRONLY);
if (fd == -1)
{
    perror("cannot access device");
    exit(EXIT_FAILURE);
}
write(fd, buffer, len);
close(fd);
```

Device Specific Command Example

make the speaker beep

```
fd = open("/dev/console", O_RDWR);
status = ioctl(fd, KDMKTONE, 0x100011AA);
if (status == -1)
{
    perror("cannot generate beep");
    exit(EXIT_FAILURE);
}
close(fd);
```

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Implementing Device Drivers

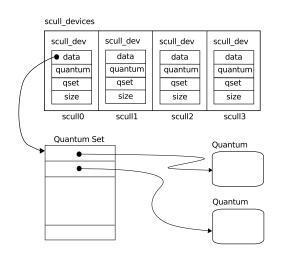
- ▶ implement system calls for device
- convert system calls to device specific I/O instructions

Device Driver Example

simplified scull

- use memory as device
 - ▶ /dev/scull0
 - ▶ /dev/scull1
- ▶ each device can hold data up to a limit
 - data persists during module's lifetime

Memory Layout



- each device has a quantum set
- each quantum contains the actual data
- memory is allocated as data is written

Global Definitions

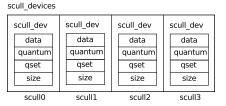
scull.h

```
#define SCULL_MAJOR 0
#define SCULL_NR_DEVS 4
#define SCULL_QUANTUM 4000
#define SCULL_QSET 1000
```

Module Parameters

```
int scull_major = SCULL_MAJOR;
int scull_minor = 0;
int scull_nr_devs = SCULL_NR_DEVS;
int scull_quantum = SCULL_QUANTUM;
int scull_qset = SCULL_QSET;
module_param(scull_major, int, S_IRUGO);
module_param(scull_minor, int, S_IRUGO);
module_param(scull_nr_devs, int, S_IRUGO);
module_param(scull_quantum, int, S_IRUGO);
module_param(scull_gset, int, S_IRUGO);
```

Data Structures



```
struct scull_dev {
    char **data;
    int quantum;
    int qset;
    unsigned long size;
    struct semaphore sem;
    struct cdev cdev;
};
struct scull_dev *scull_devices;
```

- ▶ allocate I/O region
 - base address
 - number of ports
- register driver with the kernel
 - major and minor numbers
 - capabilities: file operations

driver registration: major and minor numbers

```
if (scull_major)
{
    devno = MKDEV(scull_major, scull_minor);
    result = register_chrdev_region(devno,
                scull_nr_devs, "scull");
else
      /* dynamic */
    result = alloc_chrdev_region(&devno,
        scull_minor, scull_nr_devs, "scull");
    scull_major = MAJOR(devno);
```

data structure allocation

File Operations

map system calls to functions: struct file_operations

```
struct file_operations scull_fops = {
    .open = scull_open,
    .release = scull_release,
    .read = scull_read,
    .write = scull_write,
    .llseek = scull_llseek,
    .ioctl = scull_ioctl,
};
```

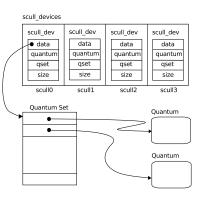
driver activation

```
for (i = 0; i < scull_nr_devs; i++)</pre>
{
    dev = &scull_devices[i];
    dev->quantum = scull_quantum;
    dev->gset = scull_gset;
    init_MUTEX(&dev->sem);
    devno = MKDEV(scull_major, scull_minor + i);
    cdev_init(&dev->cdev, &scull_fops);
    dev->cdev.owner = THIS_MODULE:
    dev->cdev.ops = &scull_fops;
    cdev_add(&dev->cdev, devno, 1);
```

Module Cleanup

```
dev_t devno = MKDEV(scull_major, scull_minor);
if (scull_devices)
{
    for (i = 0; i < scull_nr_devs; i++)</pre>
        scull_trim(scull_devices + i);
        cdev_del(&scull_devices[i].cdev);
    kfree(scull_devices);
unregister_chrdev_region(devno, scull_nr_devs);
```

Module Cleanup



data structure deallocation

```
if (dev->data)
    for (i = 0; i < dev->qset;
         i++)
        if (dev->data[i])
            kfree(dev->data[i]);
    kfree(dev->data):
dev->data = NULL;
dev->quantum = scull_quantum;
dev->qset = scull_qset;
dev->size = 0;
```

Kernel Data Structures

- a structure for each device node: struct inode
- a structure for each open file: struct file
 - ▶ f_mode: readable, writable, both
 - ► f_pos: current reading/writing position
 - ▶ f_flags
 - f_op: operations associated with the file
 - private_data: pointer to allocated data

Open

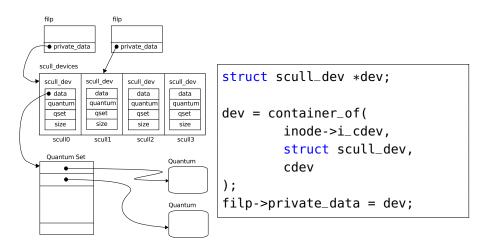
- ► identify actual device
- check for device-specific errors
- ▶ initialize device
- allocate and initialize data structures

Kernel System Call Interface

open system call:

kernel function to implement:

Open



Kernel System Call Interface

close system call:

```
int close(int fd);
```

kernel function to implement:

Kernel System Call Interface

write system call:

kernel function to implement:

Write

```
struct scull_dev *dev = filp->private_data;
  ssize_t retval = -ENOMEM:
  if (down_interruptible(&dev->sem))
      return - ERESTARTSYS;
 /* determine position */
 /* allocate quantum if necessary */
 /* copy from user space */
 /* update size */
out:
 up(&dev->sem);
  return retval;
```

Write

determine position

```
int quantum = dev->quantum, qset = dev->qset;
int s_pos, q_pos;
if (*f_pos >= quantum * qset)
{
    retval = 0;
    goto out;
s_{pos} = (long) *f_{pos} / quantum;
q_pos = (lonq) *f_pos % quantum;
```

allocate quantum if necessary

```
if (!dev->data)
{
    dev->data = kmalloc(gset * sizeof(char *),
                        GFP_KERNEL);
    if (!dev->data)
        qoto out;
    memset(dev->data, 0, qset * sizeof(char *));
if (!dev->data[s_pos])
{
    dev->data[s_pos] = kmalloc(quantum, GFP_KERNEL);
    if (!dev->data[s_pos])
        goto out;
```

copy from user space

```
/* adjust write amount */
if (count > quantum - q_pos)
    count = quantum - q_pos;
if (copy_from_user(dev->data[s_pos] + q_pos,
                   buf, count))
{
    retval = -EFAULT;
    goto out;
```

update size

```
*f_pos += count;
retval = count;

if (dev->size < *f_pos)
    dev->size = *f_pos;
```

Kernel System Call Interface

read system call:

kernel function to implement:

Read

```
struct scull_dev *dev = filp->private_data;
  ssize_t retval = 0;
  if (down_interruptible(&dev->sem))
      return -ERESTARTSYS;
  /* determine position */
  /* copy to user space */
out:
  up(&dev->sem);
  return retval;
```

Read

determine position

```
int quantum = dev->quantum;
int s_pos, q_pos;
if (*f_pos >= dev->size)
   goto out;
if (*f_pos + count > dev->size)
    count = dev->size - *f_pos;
s_{pos} = (long) *f_{pos} / quantum;
q_pos = (long) *f_pos % quantum;
if (dev->data == NULL || ! dev->data[s_pos])
    qoto out;
```

Reading from the Device

copy to user space

```
/* adjust read amount */
if (count > quantum - q_pos)
    count = quantum - q_pos;
if (copy_to_user(buf, dev->data[s_pos] + q_pos, count))
    retval = -EFAULT:
    qoto out;
}
*f_pos += count;
retval = count:
```

Kernel System Call Interface

lseek system call:

kernel function to implement:

Seek

calculate new position

```
switch(whence)
{
    case 0: /* SEEK_SET */
        newpos = off;
        break:
    case 1: /* SEEK_CUR */
        newpos = filp->f_pos + off;
        break:
    case 2: /* SEEK_END */
        newpos = dev->size + off;
        break:
    default: /* can't happen */
        return -EINVAL;
```

Seek

set new position

```
if (newpos < 0)
    return -EINVAL;
filp->f_pos = newpos;
return newpos;
```

Kernel System Call Interface

ioctl system call:

kernel function to implement:

- ► SCULL_IOCRESET: assign default values to quantum set size and quantum size
- ► SCULL_IOCSQUANTUM: set quantum size from pointer
- ► SCULL_IOCTQUANTUM: (tell) set quantum size from value
- ► SCULL_IOCGQUANTUM: get quantum size to pointer
- ► SCULL_IOCQQUANTUM: (query) return quantum size
- ► SCULL_IOCXQUANTUM: (exchange) set + get
- ► SCULL_IOCHQUANTUM: (shift) tell + query
- similar operations for quantum set size

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Device Operations

```
switch(cmd)
{
    case SCULL_IOCRESET:
        scull_quantum = SCULL_QUANTUM;
        scull_qset = SCULL_QSET;
        break;

    /* other cases */
}
```

Device Operations

setting quantum size

```
case SCULL_IOCSQUANTUM:
    if (! capable (CAP_SYS_ADMIN))
        return - EPERM;
    retval = __get_user(scull_quantum,
                         (int __user *) arg);
    break:
case SCULL_IOCTQUANTUM:
    if (! capable (CAP_SYS_ADMIN))
        return - EPERM;
    scull_quantum = arg;
    break;
```

Device Operations

getting quantum size

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Device Driver Example

short: read/write I/O ports

- each device node accesses a different port:
 - ► /dev/short0: port at base
 - ▶ /dev/short1: port at base+1
- module parameters:
 - major number (default dynamic)
 - base address (default 0x378)

Region Allocation

module initializion

module cleanup

```
release_region(short_base, SHORT_NR_PORTS);
```

Read

```
int retval = count;
int minor = iminor(filp->f_dentry->d_inode);
unsigned long port = short_base + (minor & 0x0f);
unsigned char *kbuf, *ptr;
kbuf = kmalloc(count, GFP_KERNEL);
if (!kbuf)
    return - ENOMEM;
/* do the I/O */
kfree(kbuf);
return retval;
```

Read

do the I/O

```
ptr = kbuf;
while (count--)
{
    *(ptr++) = inb(port);
    rmb();
}
if ((retval > 0) && copy_to_user(buf, kbuf, retval))
    retval = -EFAULT;
```

```
if (copy_from_user(kbuf, buf, count))
    return -EFAULT;
ptr = kbuf;
while (count--)
{
    outb(*(ptr++), port);
    wmb();
}
```

Reading Material

- ► Corbet-Rubini-Hartman, 3/e
 - ► Chapter 3: Char Drivers
 - ► Chapter 9: Communicating with Hardware