

OPERATING SYSTEMS: COMMUNICATION AND SYNCHRONIZATION AMONG PROCESSES



Threads and communication and synchronization mechanisms

To remember...

Before classes

Class

After class

Prepare the prerequisites.

Study the material associated with the **bibliography**:
slides alone are not enough.
Please ask questions (especially after study).

Exercising skills:

- ▶ Perform all **exercises**.
- ▶ Carrying out the **practice notebooks** and **the practical exercises** progressively.

Recommended reading

Base



1. Carretero 2020:
 1. Cap. 6
2. Carretero 2007:
 1. Cap. 6.1 and 6.2

Suggested



1. Tanenbaum 2006:
 1. (es) Chap. 5
 2. (en) Chap. 5
2. Stallings 2005:
 1. 5.1, 5.2 and 5.3
3. Silberschatz 2006:
 1. 6.1, 6.2, 6.5 and 6.6

Contents

- Introduction (definitions):
 - ▣ Concurrent processes.
 - ▣ Concurrency, communication and synchronization.
 - ▣ Critical section and Race conditions.
 - ▣ Mutual exclusion and critical section.
- Synchronization mechanisms (I):
 - ▣ Initial basic primitives.
 - ▣ Semaphores.
- Classic concurrency problems (I):
 - ▣ Producer-consumer
 - ▣ Reader-writers
- Synchronization mechanisms of threads (II)
 - ▣ Semaphores
 - System calls for semaphores.
 - Classic concurrency problems.
 - ▣ Mutex and condition variables
 - System calls for mutex.
 - Classic concurrency problems.
- Case study: concurrent server development

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POSIX Semaphores

- Synchronization mechanism for processes and/or threads on the same machine.

```
#include <semaphore.h>
```

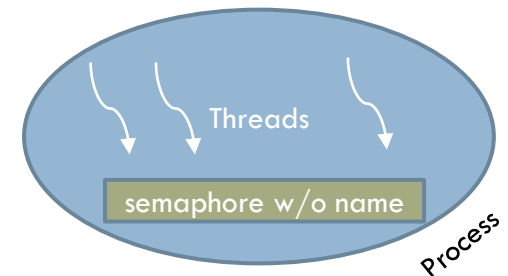
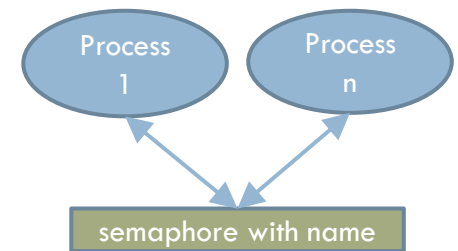
- Two types of POSIX semaphores:

- **Semaphores with name:**

- It can be used by different processes that know the name.
 - No shared memory required.
 - `sem_t *semaphore; // named`

- **Semaphores without name:**

- They can be used only by the processes that create them (and their threads) or by processes that have a shared memory area.
 - `sem_t semaphore; // no named`



POSIX Semaphores

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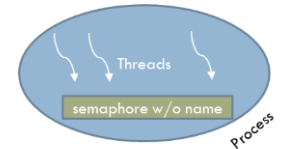
Alejandro Calderón Mateos 

- `int sem_init(sem_t *sem, int shared, int val);`

- Initializing an **unnamed** semaphore

- `int sem_destroy(sem_t *sem);`

- Destroy an **unnamed** semaphore



- `sem_t *sem_open(char *name, int flag, mode_t mode, int val);`

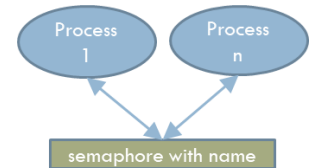
- Creates (or opens) a **named** semaphore.

- `int sem_close(sem_t *sem);`

- Closes a **named** semaphore.

- `int sem_unlink(char *name);`

- Remove a **named** semaphore.



- `int sem_wait(sem_t *sem);`

- Performs the **wait** operation on the semaphore.

- `int sem_trywait(sem_t *sem)`

- Attempts to **wait** but if it needs to block the process -> it does not block and gives -1

- `int sem_post(sem_t *sem);`

- Performs the **signal** operation on the semaphore.

Semaphores operations

```
sem_wait(s) {  
    s = s - 1;  
    if (s < 0) {  
        <Blocking  
        the process>  
    }  
}
```

```
sem_post(s) {  
    s = s + 1;  
    if (s <= 0)  
        <Unblock a  
        process blocked  
        by a wait  
        operation>  
    }  
}
```


Contents

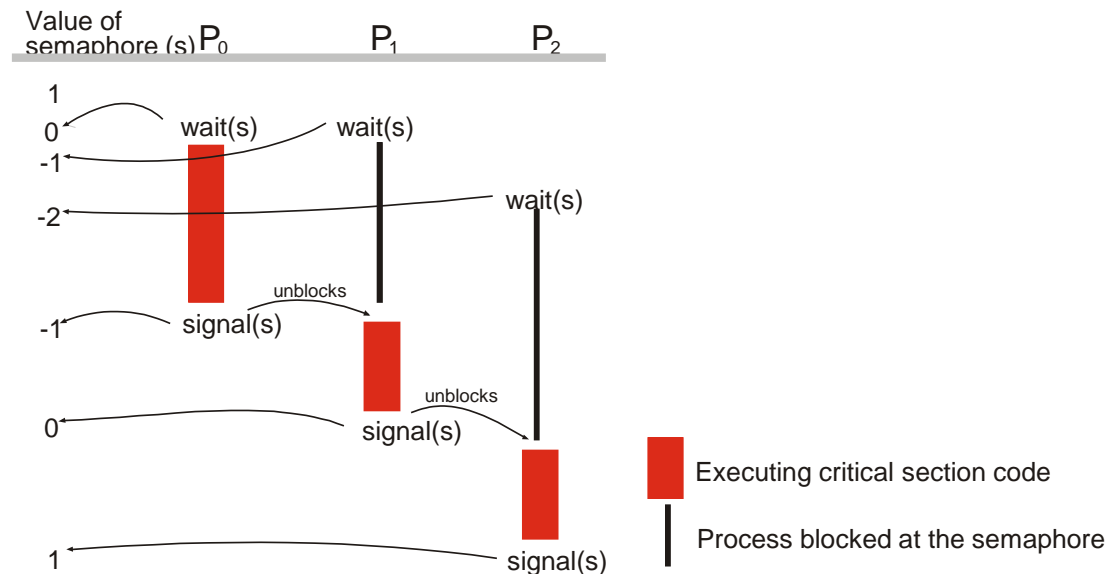
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 - ▣ Mutex and condition variables
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Critical sections with semaphores

- The semaphore must have an initial value of 1:

```
sem_wait(s); /* enter in the critical section */  
< critical section >  
sem_post(s); /* leave the critical section */
```

- Example:



Producer-consumer with bounded buffered Semaphores without name

```
/* buffer size*/
#define      MAX_BUFFER  1024
int buffer[MAX_BUFFER];    /* búfer común */
sem_t mutex;               /* critical section */
sem_t elements;            /* elements in buffer*/
sem_t huecos;              /* spaces in the buffer */
```

```
void Producer(void)
{
    int pos = 0;
    int dato, i;

    for (i=0; i<DATA_TO_PRODUCE; i++)
    {
        dato = i;
        sem_wait(&huecos);
        sem_wait(&mutex);
        buffer[pos] = i;
        pos = (pos + 1) % MAX_BUFFER;
        sem_post(&mutex);
        sem_post(&elements);
    }
    pthread_exit(0);
}
```

```
void Consumer ( void )
{
    int pos = 0;
    int dato, i;

    for (i=0; i<DATA_TO_PRODUCE; i++)
    {
        sem_wait(&elements);
        sem_wait(&mutex);
        dato = buffer[pos];
        pos = (pos + 1) % MAX_BUFFER;
        sem_post(&mutex);
        sem_post(&huecos);
        /* consume data */
    }
    pthread_exit(0);
}
```

Producer-consumer with bounded buffered Semaphores without name (1/4)

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Sistemas operativos: una visión aplicada

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```
//...
#include <semaphore.h>

/* buffer size */
#define MAX_BUFFER 1024
/* datos a producir */
#define DATA_TO_PRODUCE 100000

sem_t mutex; /* critical section */
sem_t elements; /* elements in the buffer */
sem_t huecos; /* spaces in the buffer */
int buffer[MAX_BUFFER]; /* shared buffer */

int main ( int argc, char*argv[] )
{
    pthread_t th1, th2;

    /* initialize the semaphores */
    sem_init(&mutex, 0, 1);
    sem_init(&elements, 0, 0);
    sem_init(&huecos, 0, MAX_BUFFER);

    /* create threads */
    pthread_create(&th1, NULL, Producer, NULL);
    pthread_create(&th2, NULL, Consumer, NULL);

    /* wait for its completion */
    pthread_join(th1, NULL);
    pthread_join(th2, NULL);

    sem_destroy(&mutex);
    sem_destroy(&elements);
    sem_destroy(&huecos);
    return (0);
}

void Producer(void)
{
    int pos = 0;
    int dato, i;

    for (i=0; i<DATA_TO_PRODUCE; i++)
    {
        dato = i;
        sem_wait(&huecos);
        sem_wait(&mutex);
        buffer[pos] = i;
        pos = (pos + 1) % MAX_BUFFER;
        sem_post(&mutex);
        sem_post(&elements);
    }
    pthread_exit(0);
}

void Consumer ( void )
{
    int pos = 0;
    int dato, i;

    for (i=0; i<DATA_TO_PRODUCE; i++)
    {
        sem_wait(&elements);
        sem_wait(&mutex);
        dato = buffer[pos];
        pos = (pos + 1) % MAX_BUFFER;
        sem_post(&mutex);
        sem_post(&huecos);
        /* consume data */
    }
    pthread_exit(0);
}
```

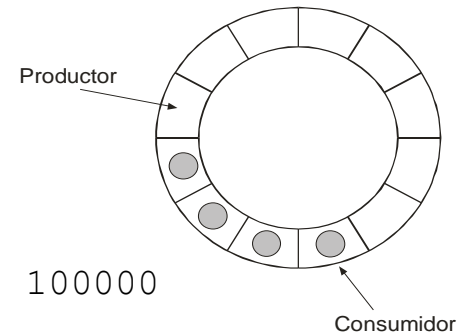
```
// ...
#include <semaphore.h>

/* buffer size */
#define MAX_BUFFER 1024
/* data to produce */
#define DATA_TO_PRODUCE 100000
```

```
sem_t mutex; /* critical section */
sem_t elements; /* eltos in the buffer */
sem_t huecos; /* spaces in the buffer */
int buffer[MAX_BUFFER]; /* shared buffer */
```

```
int main ( int argc, char *argv[] )
{
    pthread_t th1, th2;
```

```
/* initialize the semaphores */
sem_init(&mutex, 0, 1);
sem_init(&elements, 0, 0);
sem_init(&huecos, 0, MAX_BUFFER);
```



Producer-consumer with bounded buffered Semaphores without name (2/4)

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Sistemas operativos: una visión aplicada

Alejandro Calderón Mateos



```
//...
#include <semaphore.h>

/* buffer size */
#define MAX_BUFFER 1024
/* datos a producir */
#define DATA_TO_PRODUCE 100000

sem_t mutex; /* critical section */
sem_t elements; /* elements in the buffer */
sem_t huecos; /* spaces in the buffer */
int buffer[MAX_BUFFER]; /* shared buffer */

int main ( int argc, char*argv[] )
{
    pthread_t th1, th2;

    /* initialize the semaphores */
    sem_init(&mutex, 0, 1);
    sem_init(&elements, 0, 0);
    sem_init(&huecos, 0, MAX_BUFFER);

    /* create threads */
    pthread_create(&th1, NULL, Producer, NULL);
    pthread_create(&th2, NULL, Consumer, NULL);

    /* wait for its completion */
    pthread_join(th1, NULL);
    pthread_join(th2, NULL);

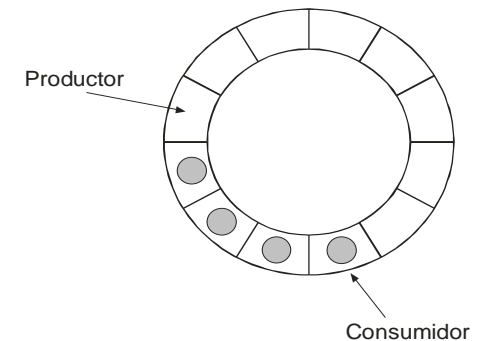
    sem_destroy(&mutex);
    sem_destroy(&huecos);
    sem_destroy(&elements);
    return (0);
}
```

```
void Producer(void)
{
    int pos = 0;
    int dato, i;

    for (i=0; i<DATA_TO_PRODUCE; i++)
    {
        dato = i;
        sem_wait(&huecos);
        sem_wait(&mutex);
        buffer[pos] = i;
        pos = (pos + 1) % MAX_BUFFER;
        sem_post(&mutex);
        sem_post(&elements);
    }
    pthread_exit(0);
}

void Consumer ( void )
{
    int pos = 0;
    int dato, i;

    for (i=0; i<DATA_TO_PRODUCE; i++)
    {
        sem_wait(&elements);
        sem_wait(&mutex);
        dato = buffer[pos];
        pos = (pos + 1) % MAX_BUFFER;
        sem_post(&mutex);
        sem_post(&huecos);
        /* Consume data */
    }
    pthread_exit(0);
}
```



```
/* create threads */
pthread_create(&th1, NULL, Producer, NULL);
pthread_create(&th2, NULL, Consumer, NULL);
```

```
/* wait for its completion */
pthread_join(th1, NULL);
pthread_join(th2, NULL);
```

```
sem_destroy(&mutex);
sem_destroy(&huecos);
sem_destroy(&elements);
return (0);
```

Producer-consumer with bounded buffered Semaphores without name (3/4)

```
//...
#include <semaphore.h>

/* buffer size */
#define MAX_BUFFER 1024
/* datos a producir */
#define DATA_TO_PRODUCE 100000

sem_t mutex; /* critical section */
sem_t elements; /* elements in the buffer */
sem_t huecos; /* spaces in the buffer */
int buffer[MAX_BUFFER]; /* shared buffer */

int main ( int argc, char*argv[] )
{
    pthread_t th1, th2;

    /* initialize the semaphores */
    sem_init(&mutex, 0, 1);
    sem_init(&elements, 0, 0);
    sem_init(&huecos, 0, MAX_BUFFER);

    /* create threads */
    pthread_create(&th1, NULL, Producer, NULL);
    pthread_create(&th2, NULL, Consumer, NULL);

    /* wait for its completion */
    pthread_join(th1, NULL);
    pthread_join(th2, NULL);

    sem_destroy(&mutex);
    sem_destroy(&huecos);
    sem_destroy(&elements);
    return (0);
}
```

```
void Producer(void)
{
    int pos = 0;
    int dato, i;

    for (i=0; i<DATA_TO_PRODUCE; i++)
    {
        dato = i;
        sem_wait(&huecos);
        sem_wait(&mutex);
        buffer[pos] = i;
        pos = (pos + 1) % MAX_BUFFER;
        sem_post(&mutex);
        sem_post(&elements);
    }
    pthread_exit(0);
}
```

```
void Consumer ( void )
{
    int pos = 0;
    int dato, i;

    for (i=0; i<DATA_TO_PRODUCE; i++)
    {
        sem_wait(&elements);
        sem_wait(&mutex);
        dato = buffer[pos];
        pos = (pos + 1) % MAX_BUFFER;
        sem_post(&mutex);
        sem_post(&huecos);
        /* Consume data */
    }
    pthread_exit(0);
}
```

```
void Producer(void)
```

```
{
```

```
    int pos = 0;
```

```
    int dato;
```

```
    int i;
```

```
    for (i=0; i<DATA_TO_PRODUCE; i++)
```

```
    {
```

```
        dato = i; /* produce... */
```

```
        sem_wait(&huecos);
```

```
        sem_wait(&mutex);
```

```
        buffer[pos] = i;
```

```
        pos = (pos + 1) % MAX_BUFFER;
```

```
        sem_post(&mutex);
```

```
        sem_post(&elements);
```

```
    }
```

```
    pthread_exit(0);
```

```
}
```

Producer-consumer with bounded buffered Semaphores without name (4/4)

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Sistemas operativos: una visión aplicada

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```
//...
#include <semaphore.h>

/* buffer size */
#define MAX_BUFFER 1024
/* datos a producir */
#define DATA_TO_PRODUCE 100000

sem_t mutex; /* critical section */
sem_t elements; /* elements in the buffer */
sem_t huecos; /* spaces in the buffer */
int buffer[MAX_BUFFER]; /* shared buffer */

int main ( int argc, char*argv[] )
{
    pthread_t th1, th2;

    /* initialize the semaphores */
    sem_init(&mutex, 0, 1);
    sem_init(&elements, 0, 0);
    sem_init(&huecos, 0, MAX_BUFFER);

    /* create threads */
    pthread_create(&th1, NULL, Producer, NULL);
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    /* wait for its completion */
    pthread_join(th1, NULL);
    pthread_join(th2, NULL);

    sem_destroy(&mutex);
    sem_destroy(&huecos);
    sem_destroy(&elements);
    return (0);
}
```

```
void Producer(void)
{
    int pos = 0;
    int dato, i;

    for (i=0; i<DATA_TO_PRODUCE; i++)
    {
        dato = i;
        sem_wait(&huecos);
        sem_wait(&mutex);
        buffer[pos] = i;
        pos = (pos + 1) % MAX_BUFFER;
        sem_post(&mutex);
        sem_post(&elements);
    }
    pthread_exit(0);
}
```

```
void Consumer ( void )
{
    int pos = 0;
    int dato, i;

    for (i=0; i<DATA_TO_PRODUCE; i++)
    {
        sem_wait(&elements);
        sem_wait(&mutex);
        dato = buffer[pos];
        pos = (pos + 1) % MAX_BUFFER;
        sem_post(&mutex);
        sem_post(&huecos);
        /* Consume data */
    }
    pthread_exit(0);
}
```

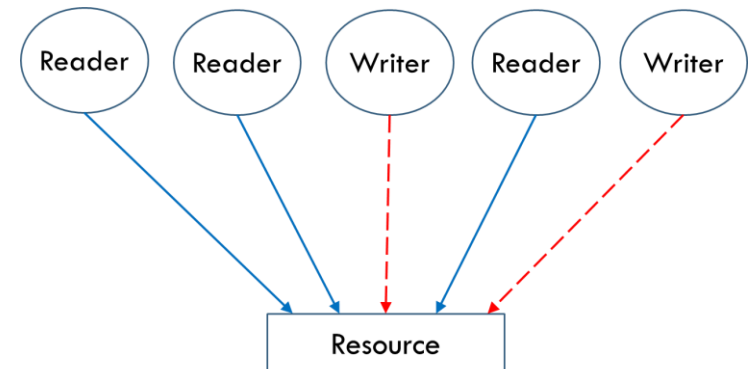
```
void Consumer ( void )
{
    int pos = 0;
    int dato;
    int i;

    for (i=0; i<DATA_TO_PRODUCE; i++)
    {
        sem_wait(&elements);
        sem_wait(&mutex);
        dato = buffer[pos];
        pos = (pos + 1) % MAX_BUFFER;
        sem_post(&mutex);
        sem_post(&huecos);
        /* consume data */
    }

    pthread_exit(0);
}
```

Reader-writer problem

- Problem that arises when you have:
 - A **shared storage** area.
 - **Multiple processes read** information.
 - **Multiple processes write** information.
- Conditions:
 - Any number of readers can read from the data zone concurrently: **multiple readers possible at the same time**.
 - **Only one writer** can modify the information **at a time**.
 - **During a writing no reader** can read



Readers and writers

Semaphores without name

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```
int dato = 5;          /* resource */
int n_readers = 0;     /* num. of readers */
sem_t sem_reader;      /* control access to n_readers */
sem_t mutex;           /* control access to dato */
```

```
void Reader(void)
{
    sem_wait(&sem_reader);
    n_readers = n_readers + 1;
    if (n_readers == 1)
        sem_wait(&mutex);
    sem_post(&sem_reader);

    printf("%d\n", dato);

    sem_wait(&sem_reader);
    n_readers = n_readers - 1;
    if (n_readers == 0)
        sem_post(&mutex);
    sem_post(&sem_reader);
    pthread_exit(0);
}
```

```
void Writer(void)
{
    sem_wait(&mutex);
    dato = dato + 2;
    sem_post(&mutex);

    pthread_exit(0);
}
```

Readers and writers

Semaphores without name

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```
int dato = 5;           /* resource */
int n_readers = 0;      /* num. of readers */
sem_t sem_reader;       /* control access to n_readers */
sem_t mutex;            /* control access to dato */

int main ( int argc, char *argv[] )
{
    pthread_t th1, th2, th3, th4;
    sem_init(&mutex, 0, 1);
    sem_init(&sem_reader, 0, 1);

    pthread_create(&th1, NULL, Reader, NULL);
    pthread_create(&th2, NULL, Writer, NULL);
    pthread_create(&th3, NULL, Reader, NULL);
    pthread_create(&th4, NULL, Writer, NULL);
    pthread_join(th1, NULL);
    pthread_join(th2, NULL);
    pthread_join(th3, NULL);
    pthread_join(th4, NULL);

    sem_destroy(&mutex);
    sem_destroy(&sem_reader);

    return 0;
}
```

Reader-writers (readers have priority)

Semaphores without name

REVIEW

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<http://faculty.juniata.edu/rhodes/os/ch5d.htm>

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SHARED MEMORY:

```
int nreaders; semaphore s_read=1; semaphore s_write=1;
```

READER:

```
for(;;) {  
    sem_wait(s_read);  
    nreaders++;  
    if (nreaders==1)  
        sem_wait(s_write);  
    sem_signal(s_read);  
  
    perform_read();  
  
    sem_wait(s_read);  
    nreaders--;  
    if (nreaders==0)  
        sem_signal(s_write);  
    sem_signal(s_read);  
}
```

WRITER:

```
for(;;) {  
    sem_wait(s_write);  
    perform_write();  
    sem_signal(s_write);  
}
```

Reader-writers (writers have priority)

Semaphores without name

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<https://computationstructures.org/lectures/synchronization/synchronization.html>
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SHARED MEMORY:

```
int nreaders, nescr = 0; semaphore lect, s_write = 1;
semaphore x, y, z = 1;
```

READER:

```
for(;;) {
    sem_wait(z);
    sem_wait(lect);
    sem_wait(x);
    nreaders++;
    if (nreaders==1)
        sem_wait(s_write);
    sem_signal(x);
    sem_signal(lect);
    sem_signal(z);
    // doReading();
    sem_wait(x);
    nreaders--;
    if (nreaders==0)
        sem_signal(s_write);
    sem_signal(x);
}
```

WRITER:

```
for(;;) {
    sem_wait(y);
    nescr++;
    if (nescr==1)
        sem_wait(lect);
    sem_signal(y);
    sem_wait(s_write);
    // doWriting();
    sem_signal(s_write);
    sem_wait(y);
    nescr--;
    if (nescr==0)
        sem_signal(s_readt);
    sem_signal(y);
}
```

Semaphores with name

Naming

- Allow synchronization of different processes without using shared memory.
- The name of a semaphore is a string of characters (with the same restrictions as a filename).
 - ▣ If the name (path) is relative, only the process that creates it and its children can access the semaphore.
 - ▣ If the name is absolute (starts with "/") the semaphore can be shared by any process that knows its name and has permissions.
- Common mechanism for creating semaphores shared by parents and children
 - ▣ "Unnamed" does not apply -> processes do NOT share memory.

Semaphores with name

Creation and use

□ To create:

```
sem_t *sem_open(char *name, int flag, mode_t mode, int val);
```

- Flag = O_CREAT will create it.
- Flag: O_CREAT | O_EXECL will create it if it does not exist. -1 in case it exists.
- Mode: access permissions;
- Val: initial value of the semaphore (≥ 0);

□ To use:

```
sem_t *sem_open(char *name, int flag);
```

- With flag 0. If it does not exist, it returns -1.

□ Important:

- All processes must know "name" and use the same name.

Readers and writers

semaphores with name

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```
int main ( int argc, char *argv[] )
{
    int i, n= 5; pid_t pid;

    /* Creates the named semaphore */
    if ((mutex=sem_open("/tmp/sem_1", O_CREAT, 0644, 1))== (sem_t *)-1)
        { perror("Cannot create a semaphore: "); exit(1); }
    if((sem_reader=sem_open("/tmp/sem_2", O_CREAT, 0644, 1))== (sem_t *)-1)
        { perror("Cannot create a semaphore: "); exit(1); }

    /* Creates the processes */
    for (i = 1; i< atoi(argv[1]); ++i)
    {
        pid = fork();
        if (pid ==-1)
            { perror("The process cannot be created: "); exit(-1);}
        if (pid==0)
            { reader(getpid()); break; }
        else writer(pid);
    }
    sem_close(mutex);
    sem_close(sem_reader);
    sem_unlink("/tmp/sem_1");
    sem_unlink("/tmp/sem_2");

    return 0;
}
```

Readers and writers semaphores with name

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```
int dato = 5;           /* resource */
int n_readers = 0;      /* num. of readers */
sem_t *sem_reader;
sem_t *mutex;
```

```
void reader (int pid)
{
    sem_wait(sem_reader);
    n_readers = n_readers + 1;
    if (n_readers == 1)
        sem_wait(mutex);
    sem_post(sem_reader);

    printf("reader %d  dato: %d\n",
           pid, dato);

    sem_wait(sem_reader);
    n_readers = n_readers - 1;
    if (n_readers == 0)
        sem_post(mutex);
    sem_post(sem_reader);
}
```

```
void writer (int pid)
{
    sem_wait(mutex);
    dato = dato + 2;

    printf("writer %d  dato: %d\n",
           pid, dato);

    sem_post(mutex);
}
```


Contents

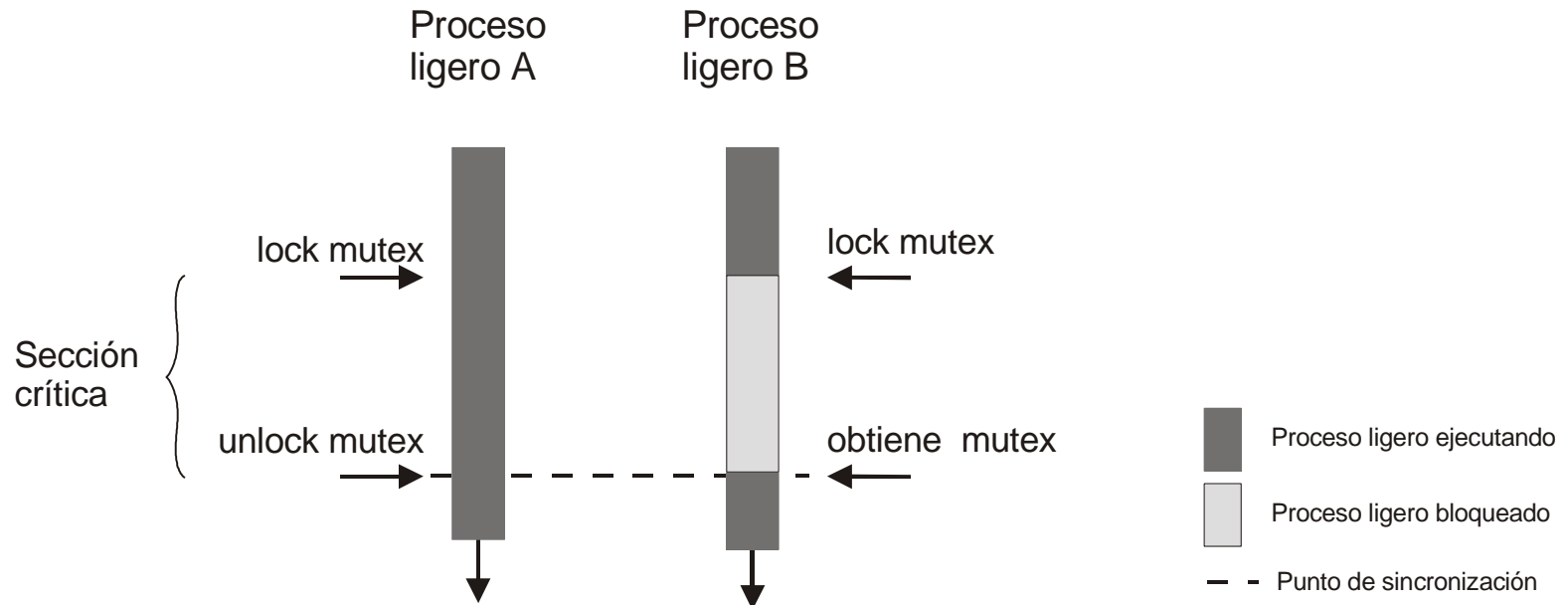
- Introduction (definitions):
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 - Classic concurrency problems.
 - ▣ **Mutex and condition variables**
 - System calls for mutex.
 - Classic concurrency problems.
- Case study: concurrent server development

Mutex and conditional variables

- A mutex is a synchronization mechanism suitable for threads.
- It is a binary semaphore with 2 atomic operations:
 - ▣ **lock(m)** Lock the mutex and if the mutex is already locked the process is suspended.
 - ▣ **unlock(m)** Unlocks the mutex and if there are suspended processes in the mutex then it unblock one.
- TIP: The `unlock` operation must be performed by the thread that executed `lock`

Critical sections with mutex

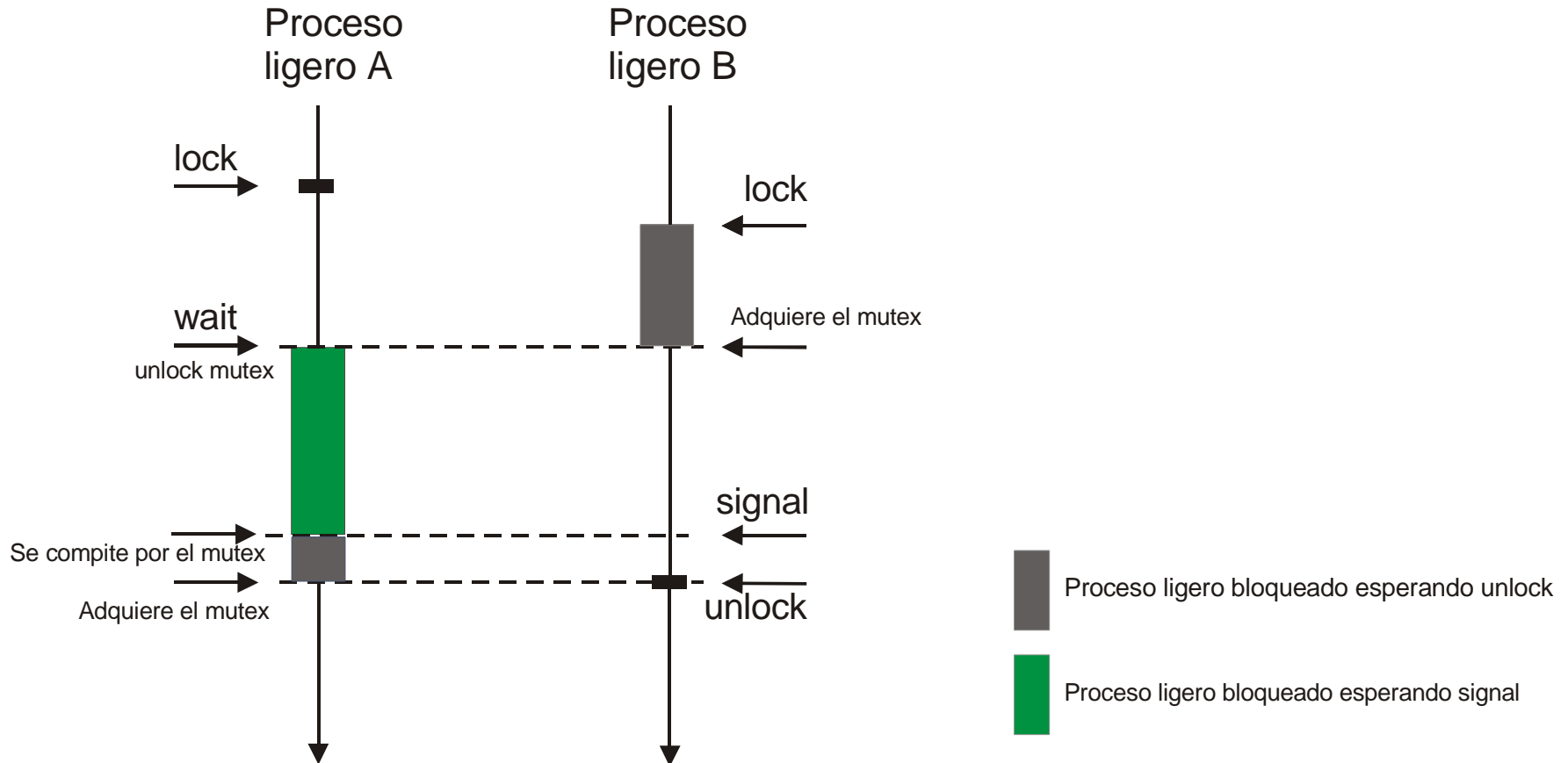
```
lock (m);    /* enter in the critical section */  
    < critical section >  
unlock (m); /* leave the critical section */
```



Conditional variables

- Synchronization variables associated with a mutex
- Two atomic operations:
 - **wait:** Blocks the thread that executes it and ejects it from the mutex
 - **signal:** Unblocks one or more suspended processes in the conditional variable and the waking process competes again to lock the mutex
- Convenient to execute between lock and unlock

Conditional variables



Use of mutex and conditional variables

□ Thread A

```
lock(mutex); /* access to the resource */
```

```
check the data structures;
```

```
while (busy resource)
```

```
    wait(condition, mutex);
```

```
mark the resource as busy;
```

```
unlock(mutex);
```

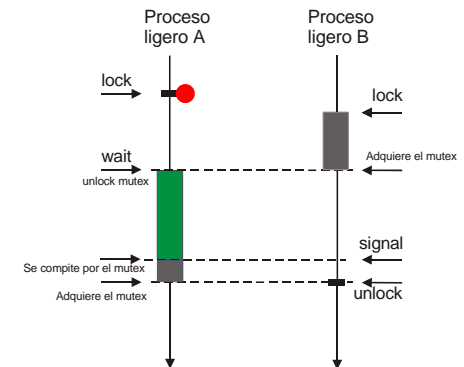
□ Thread B

```
lock(mutex); /* access to the resource */
```

```
mark the resource as free;
```

```
signal(condition);
```

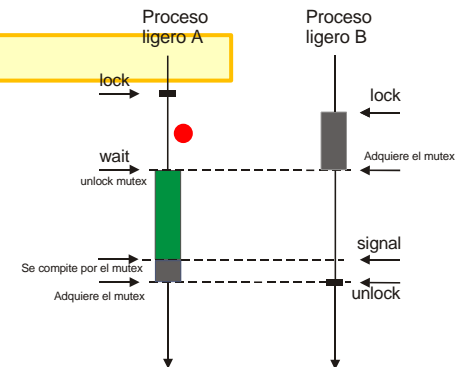
```
unlock(mutex);
```



Use of mutex and conditional variables

□ Thread A

```
lock(mutex); /* access to the resource */  
check the data structures;  
while (busy resource)  
    wait(condition, mutex);  
mark the resource as busy;  
unlock(mutex);
```



□ Thread B

```
lock(mutex); /* access to the resource */  
mark the resource as free;  
signal(condition);  
unlock(mutex);
```

Use of mutex and conditional variables

□ Thread A

```
lock(mutex); /* access to the resource */  
check the data structures;  
while (busy resource)
```

```
wait(condition, mutex);
```

```
mark the resource as busy;
```

```
unlock(mutex);
```

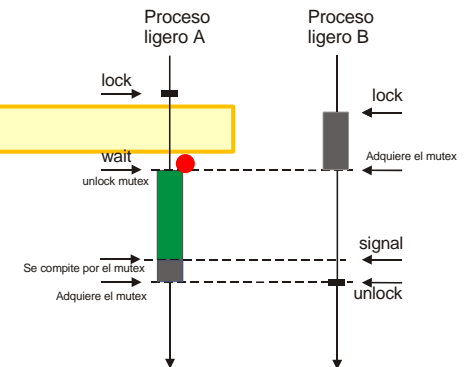
□ Thread B

```
lock(mutex); /* access to the resource */
```

```
mark the resource as free;
```

```
signal(condition);
```

```
unlock(mutex);
```



Use of mutex and conditional variables

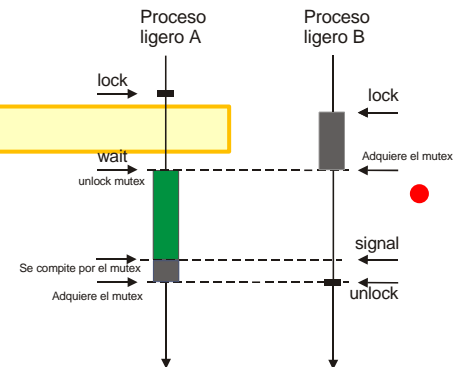
□ Thread A

```
lock(mutex); /* access to the resource */  
check the data structures;  
while (busy resource)
```

```
wait(condition, mutex);
```

```
mark the resource as busy;
```

```
unlock(mutex);
```



□ Thread B

```
lock(mutex); /* access to the resource */
```

```
mark the resource as free;
```

```
signal(condition);
```

```
unlock(mutex);
```

Use of mutex and conditional variables

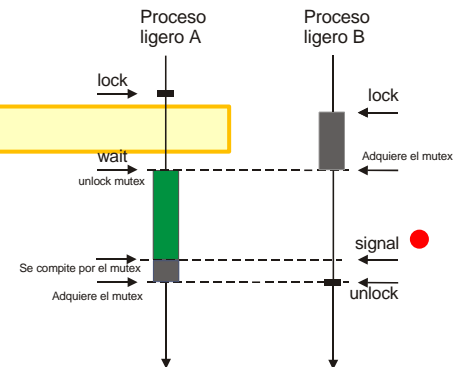
□ Thread A

```
lock(mutex); /* access to the resource */  
check the data structures;  
while (busy resource)
```

```
wait(condition, mutex);
```

```
mark the resource as busy;
```

```
unlock(mutex);
```



□ Thread B

```
lock(mutex); /* access to the resource */  
mark the resource as free;
```

```
signal(condition);
```

```
unlock(mutex);
```

Use of mutex and conditional variables

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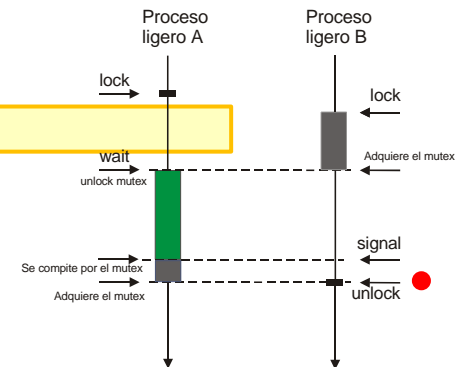
□ Thread A

```
lock(mutex); /* access to the resource */  
check the data structures;  
while (busy resource)
```

```
wait(condition, mutex);
```

```
mark the resource as busy;
```

```
unlock(mutex);
```



□ Thread B

```
lock(mutex); /* access to the resource */  
mark the resource as free;
```

```
signal(condition);
```

```
unlock(mutex);
```

Use of mutex and conditional variables

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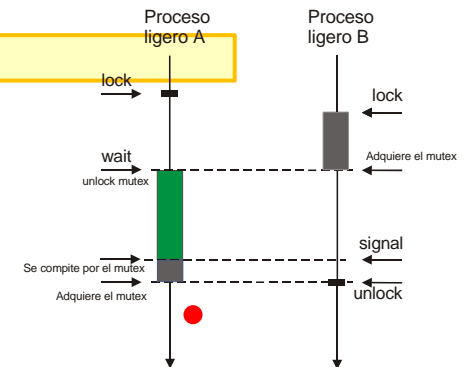
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□ Thread A

```
lock(mutex); /* access to the resource */  
check the data structures;  
while (busy resource)  
    wait(condition, mutex);  
mark the resource as busy;  
unlock(mutex);
```

□ Thread B

```
lock(mutex); /* access to the resource */  
mark the resource as free;  
signal(condition);  
unlock(mutex);
```



Use of mutex and conditional variables

□ Thread A

```
lock(mutex); /* access to the resource */  
check the data structures;  
while (busy resource)  
    wait(condition, mutex);  
mark the resource as busy;  
unlock(mutex);
```

Important to use
while to re-
evaluate condition

□ Thread B

```
lock(mutex); /* access to the resource */  
mark the resource as free;  
signal(condition);  
unlock(mutex);
```

Use of mutex and conditional variables

□ Thread A

```
lock(mutex); /* access to the resource */  
check the data structures;  
while (busy resource)  
    wait(condition, mutex);  
mark the resource as busy;  
unlock(mutex);
```

□ Thread B

```
lock(mutex); /* access to the resource */  
mark the resource as free;  
signal(condition);  
unlock(mutex);
```

Use of mutex and conditional variables

□ Thread A

```
lock(mutex); /* access to the resource */  
check the data structures;  
while (busy resource)  
    wait(condition, mutex);  
mark the resource as busy;  
unlock(mutex);
```

□ Thread B

```
lock(mutex); /* access to the resource */  
mark the resource as free;  
signal(condition);  
unlock(mutex);
```

Use of mutex and conditional variables

□ Thread A

```
lock(mutex); /* access to the resource */
```

```
check the data structures;
```

```
while (busy resource)
```

```
    wait(condition, mutex);
```

```
mark the resource as busy;
```

```
unlock(mutex);
```

- A signal before the **wait** is "lost".
- Herefore, the Boolean condition (free/busy resource) is important.

□ Thread B

```
lock(mutex); /* access to the resource */
```

```
mark the resource as free;
```

```
signal(condition);
```

```
unlock(mutex);
```


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POSIX Services

```
int pthread_mutex_init ( pthread_mutex_t *mutex,  
                        pthread_mutexattr_t * attr );
```

- ▣ Initializes a mutex.

```
int pthread_mutex_destroy ( pthread_mutex_t *mutex );
```

- ▣ Destroy a mutex.

```
int pthread_mutex_lock ( pthread_mutex_t *mutex );
```

- ▣ Try to obtain the mutex.
- ▣ locks the thread if the mutex is acquired by another thread.

```
int pthread_mutex_unlock ( pthread_mutex_t *mutex );
```

- ▣ Unlock the mutex.

POSIX Services

```
int pthread_cond_init ( pthread_cond_t*cond,  
                        pthread_condattr_t*attr );
```

- ▣ Initializes a conditional variable.

```
int pthread_cond_destroy ( pthread_cond_t *cond );
```

- ▣ Destroys a conditional variable.

```
int pthread_cond_signal ( pthread_cond_t *cond );
```

- ▣ At least one of the threads suspended in `cond` is reactivated.
- ▣ No effect if there is no thread waiting (different from semaphores).

```
int pthread_cond_wait ( pthread_cond_t*cond,  
                        pthread_mutex_t*mutex );
```

- ▣ At the same time suspends the calling thread and releases `mutex`.
- ▣ When another thread calls `..._cond_signal` on `cond` the thread wakes up and competes again for the `mutex`.

POSIX Services

```
int pthread_cond_broadcast ( pthread_cond_t *cond );
```

- ▣ All suspended threads in the conditional variable `cond` are reactivated.
- ▣ It has no effect if no thread is waiting.

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Producer-consumer with mutex

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Alejandro Calderón Mateos 

```
int main ( int argc, char *argv[] )
{
    pthread_t th1, th2;

    pthread_mutex_init(&mutex, NULL);
    pthread_cond_init(&no_full, NULL);
    pthread_cond_init(&no_empty, NULL);

    pthread_create(&th1, NULL, Producer, NULL);
    pthread_create(&th2, NULL, Consumer, NULL);
    pthread_join(th1, NULL);
    pthread_join(th2, NULL);

    pthread_mutex_destroy(&mutex);
    pthread_cond_destroy(&no_full);
    pthread_cond_destroy(&no_empty);
    return 0;
}
```

Producer-consumer with mutex

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Alejandro Calderón Mateos 

```
#define MAX_BUFFER      1024    /* buffer size */
#define DATA_TO_PRODUCE 100000 /* data to be produced */

pthread_mutex_t mutex;          /* shared buffer access mutex */
pthread_cond_t no_full;         /* controls the filling of the buffer */
pthread_cond_t no_empty;       /* controls the emptying of the buffer */
int n_elements;                /* number of elements in the buffer */
int buffer[MAX_BUFFER];        /* common buffer */
```

```
void Producer ( void )
{
    int dato, i ,pos = 0;

    for(i=0; i<DATA_TO_PRODUCE; i++)
    {
        dato = i;
        pthread_mutex_lock(&mutex);
        while (n_elements == MAX_BUFFER)
            pthread_cond_wait(&no_full,
                               &mutex);

        buffer[pos] = i;
        pos = (pos + 1) % MAX_BUFFER;
        n_elements ++;
        pthread_cond_signal(&no_empty);
        pthread_mutex_unlock(&mutex);
    }
    pthread_exit(0);
}
```

```
void Consumer(void)
{
    int dato, i ,pos = 0;

    for(i=0; i<DATA_TO_PRODUCE; i++)
    {
        pthread_mutex_lock(&mutex);
        while (n_elements == 0)
            pthread_cond_wait(&no_empty,
                               &mutex);

        dato = buffer[pos];
        pos = (pos + 1) % MAX_BUFFER;
        n_elements --;
        pthread_cond_signal(&no_full);
        pthread_mutex_unlock(&mutex);
        printf("Consume %d \n", dato);
    }
    pthread_exit(0);
}
```

Producer-consumer with mutex

```
#define MAX_BUFFER      1024    /* buffer size */
#define DATA_TO_PRODUCE 100000 /* data to be produced */

pthread_mutex_t mutex;          /* shared buffer access mutex */
pthread_cond_t no_full;         /* controls the filling of the buffer */
pthread_cond_t no_empty;       /* controls the emptying of the buffer */
int n_elements;                /* number of elements in the buffer */
int buffer[MAX_BUFFER];        /* common buffer */
```

```
void Producer ( void )
{
    int dato, i ,pos = 0;

    for(i=0; i<DATA_TO_PRODUCE; i++)
    {
        dato = i;
        pthread_mutex_lock(&mutex);
        while (n_elements == MAX_BUFFER)
            pthread_cond_wait(&no_full,
                               &mutex);
        buffer[pos] = i;
        pos = (pos + 1) % MAX_BUFFER;
        n_elements ++;
        pthread_cond_signal(&no_empty);
        pthread_mutex_unlock(&mutex);
    }
    pthread_exit(0);
}
```

```
void Consumer(void)
{
    int dato, i ,pos = 0;

    for(i=0; i<DATA_TO_PRODUCE; i++)
    {
        pthread_mutex_lock(&mutex);
        while (n_elements == 0)
            pthread_cond_wait(&no_empty,
                               &mutex);
        dato = buffer[pos];
        pos = (pos + 1) % MAX_BUFFER;
        n_elements --;
        pthread_cond_signal(&no_full);
        pthread_mutex_unlock(&mutex);
        printf("Consume %d \n", dato);
    }
    pthread_exit(0);
}
```


Reader-writers with mutex

```
int main ( int argc, char *argv[] )
{
    pthread_t th1, th2, th3, th4;

    pthread_mutex_init(&mutex, NULL);
    pthread_mutex_init(&mutex_readers, NULL);

    pthread_create(&th1, NULL, Reader, NULL);
    pthread_create(&th2, NULL, Writer, NULL);
    pthread_create(&th3, NULL, Reader, NULL);
    pthread_create(&th4, NULL, Writer, NULL);
    pthread_join(th1, NULL);
    pthread_join(th2, NULL);
    pthread_join(th3, NULL);
    pthread_join(th4, NULL);

    pthread_mutex_destroy(&mutex);
    pthread_mutex_destroy(&mutex_readers);
    return 0;
}
```

Reader-writers with mutex

```
int dato = 5; /* resource */
int n_readers = 0; /* number of readers*/
pthread_mutex_t mutex; /* control access to dato */
pthread_mutex_t mutex_readers; /* control access to n_readers */
```

```
void *Reader( void *arg )
{
    pthread_mutex_lock(&mutex_readers);
    n_readers++;
    if (n_readers == 1)
        pthread_mutex_lock(&mutex);
    pthread_mutex_unlock(&mutex_readers);

    /* read dato */
    printf("%d\n", dato);

    pthread_mutex_lock(&mutex_readers);
    n_readers--;
    if (n_readers == 0)
        pthread_mutex_unlock(&mutex);
    pthread_mutex_unlock(&mutex_readers);

    pthread_exit(0);
}
```

```
void *Writer ( void *arg )
{
    pthread_mutex_lock(&mutex);

    /* modify the resource */
    dato = dato + 2;

    pthread_mutex_unlock(&mutex);

    pthread_exit(0);
}
```

OPERATING SYSTEMS: COMMUNICATION AND SYNCHRONIZATION AMONG PROCESSES



Threads and communication and synchronization mechanisms