ARCOS Group

uc3m Universidad Carlos III de Madrid

Lesson 3 (II)

Fundamentals of assembler programming

Computer Structure

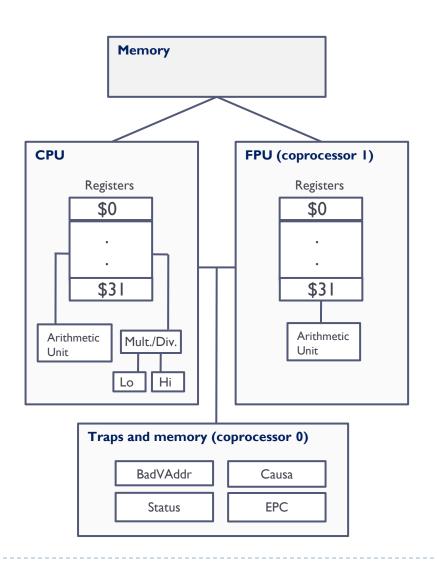
Bachelor in Computer Science and Engineering



Contents

- Basic concepts on assembly programming
- MIPS32 assembly language, memory model and data representation
- Instruction formats and addressing modes
- Procedure calls and stack convention

MIPS-32 Architecture



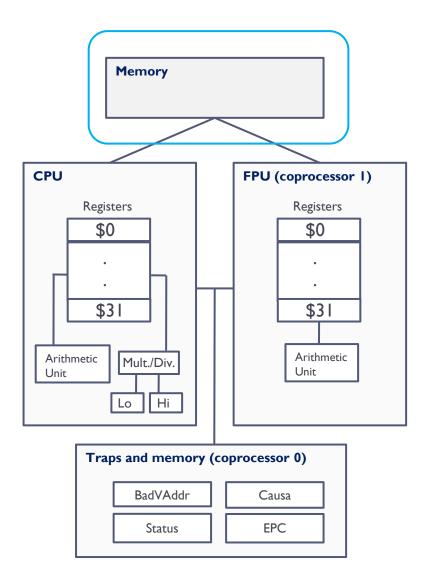
MIPS 32

- 32-bits processor
- RISC type
- CPU + auxiliar coprocessors

Coprocessor 0

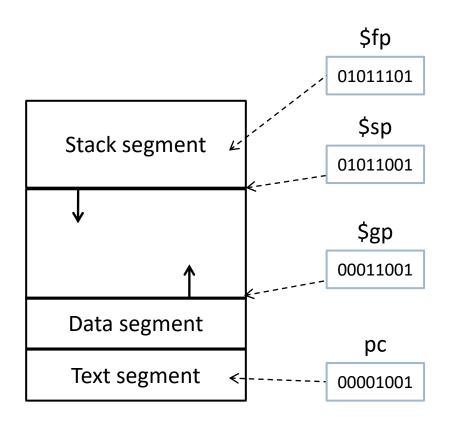
- exceptions, interrupts and virtual memory system
- Coprocessor I
 - FPU (Floating Point Unit)

MIPS-32 Architecture



- ▶ 32-bit memory addresses
- 4 GB addressable memory

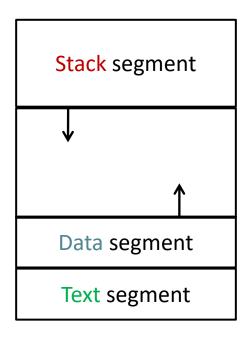
Memory layout for a process



5

- The memory space is divided in logic segments in order to organize the content:
 - Stack segment
 - Local variables
 - Function contexts
 - Data segments
 - Static data
 - Code segment (text)
 - Program code

Storing variables in memory



```
// global variables
int a;
main ()
   // local variables
   int b;
   // code (text)
   return a + b;
```

Register File (integers)

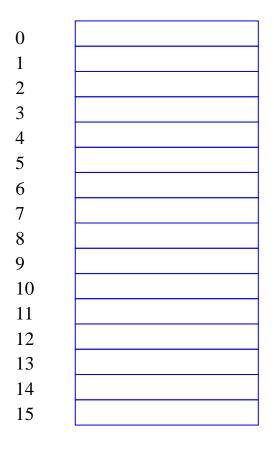
Symbolic name	Number	Usage
zero	0	Constant 0
at	I	Reserved for assembler
v0, v1	2, 3	Results of functions
a0,, a3	4,, 7	Function arguments
t0,, t7	8,, 15	Temporary (NO preserved across calls)
s0,, s7	16,, 23	Saved temporary (preserved across calls)
t8, t9	24, 25	Temporary (NO preserved across calls)
k0, k1	26, 27	Reserved for operating system
gp	28	Pointer to global area
sp	29	Stack pointer
fp	30	Frame pointer
ra	31	Return address (used by function calls)

▶ There are 32 registers

- Size: 4 bytes (I word)
- Used a \$ at the beginning

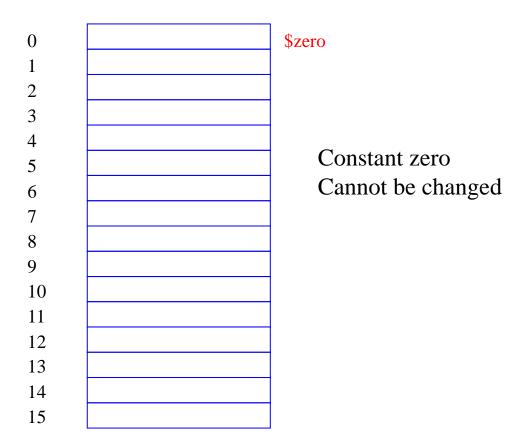
Use convention

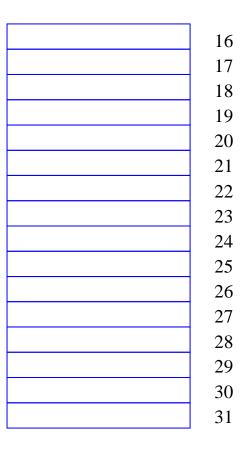
- Reserved
- Arguments
- Results
- Temporary
- Pointers



- 32 registers
 - □ 4 bytes of size (one word)
 - □ Name starts with \$ at the beginning
- Usage Convention
 - □ Reserved
 - Arguments
 - □ Results
 - □ Temporary
 - Pointers

,
16
17
18
19
20
21
22
23
24
25
26
27
28
29
30
31
_





0	\$2	zero			16
1					17
2		Temporary registers			18
3		p			19
4				,	20
5				,	21
6				,	22
7				,	23
8	\$	t0	\$t8	, <u>, , , , , , , , , , , , , , , , , , </u>	24
9	\$	t1	\$t9	,	25
10	\$	t2			26
11	\$	t3			27
12	\$	t4		,	28
13	\$	t5			29
14	\$	t6			30
15	\$	t7		,	31

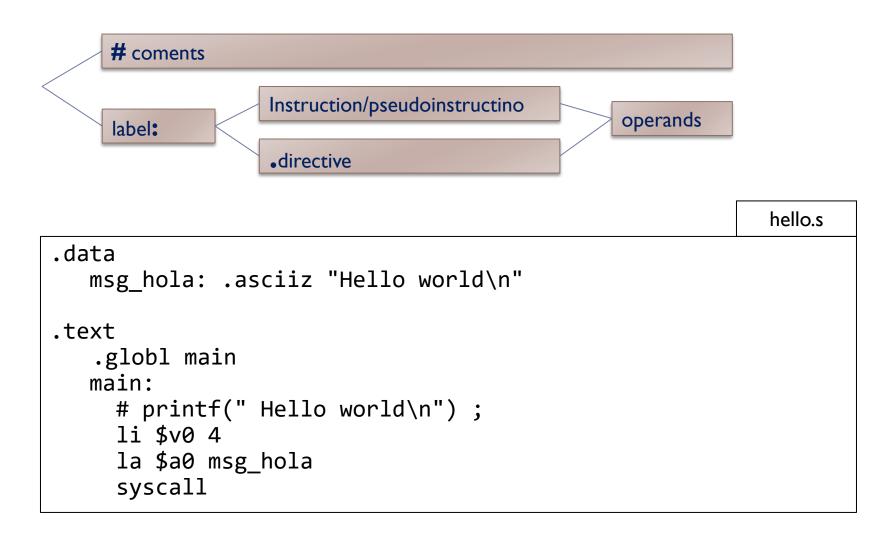
0	\$zero		\$s0	16
1			\$ s1	17
2		Preserved values	\$s2	18
3			\$s3	19
4			\$s4	20
5			\$s5	21
6			\$ s6	22
7			\$s7	23
8	\$t0		\$t8	24
9	\$t1		\$t9	25
10	\$t2			26
11	\$t3			27
12	\$t4			28
13	\$t5			29
14	\$t6			30
15	\$t7			31

0	\$zero	\$s0	16
1		\$ s1	17
2	\$v0	\$s2	18
3	\$v1	\$s3	19
4	\$a0	\$s4	20
5	\$a1	\$s5	21
6	\$a2	\$s6	22
7	\$a3	\$s7	23
8	\$tO	\$t8	24
9	\$t1	\$t9	25
10	\$t2		26
11	\$t3		27
12	\$t4 Arguments and		28
13	\$t5 subroutines support	\$sp	29
14	\$t6	\$fp	30
15	\$t7	\$ra	31

0	\$zero		\$s0	16
1	\$at		\$s1	17
2	\$v0		\$s2	18
3	\$v1		\$s3	19
4	\$a0		\$s4	20
5	\$a1		\$s5	21
6	\$a2		\$s6	22
7	\$a3		\$s7	23
8	\$t0		\$t8	24
9	\$t1		\$t9	25
10	\$t2		\$k0	26
11	\$t3		\$k1	27
12	\$t4	Others	\$gp	28
13	\$t5		\$sp	29
14	\$t6		\$fp	30
15	\$t7		\$ra	31

hello.s

```
.data
  msg_hola: .asciiz "hello world\n"
.text
   .glob1 main
  main:
    # printf("hello world\n") ;
    li $v0 4
             # syscall code: 4
    la $a0 msg_hola # address where msg_hola starts
    syscall
```



hello.s

```
.data
   msg hola: .asciiz "hello world\n"
.text
  .glob1 main
   main:
     # printf("hello world\n");
     li $v0 4
     la $a0 msg hola
     syscall
```

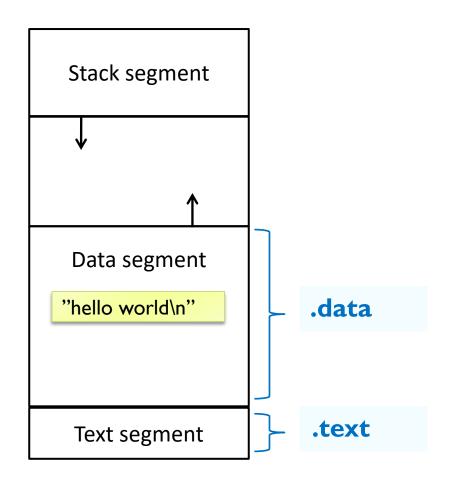
```
data segment
                                                    hello.s
.data <
   msg_hola: .asciiz "hello world\n"
                              msg_hola: represents the memory address
                              where the string starts to be stored.
.text ←
   .glob1 main
                                     code segment
   main:
      # printf("hello world\n") ;
      li $v0 4
      la $a0 msg hola
      syscall
```

hello.s

```
.data
   msg hola: .asciiz "hello world\n"
.text
                           label: represent the memory address
                           where first instruction of main is stored
   .glob1 main
   main
                                  comments
      # printf("hello world\n") ;
      li $v0 4
                                           instructions
      la $a0 msg_hola
      syscall
```

Assembler program: assembler directives (preprocessing)

```
hello.s
.data
 msg_hola: .asciiz "hello
   world\n"
.text
 .globl main
 main:
  # printf("hello world\n");
  li $v0 4
  la $a0 msg_hola
  syscall
```



Assembly: directives

Directives	Description	
.data	Next elements will go to the data segment	
.text	Next elements will go to the code segment	
.ascii "string value"	String definition without '\0' ending terminator	
.asciiz "string value"	String definition with '\0' ending terminator ('\0' = 0)	
.byte 1, 2, 3	Bytes stored in memory consecutively	
.half 300, 301, 302	Half-words stored in memory consecutively	
.word 800000, 800001	Words stored in memory consecutively	
.float 1.23, 2.13	Floats stored in memory consecutively	
.double 3.0e21	Doubles stored in memory consecutively	
.space 10	Allocates a space of 10 bytes in the current segment	
.extern label n	Declare that label is global of size n	
.globl label	Declare label as global	
.align n	Align next element to a address multiple of 2 ⁿ	

Static data definition

label (address) datatype (directive) value

```
.data
cadena :
         .asciiz_"Hola mundo\n"
                 # int i1=10
i1: .word 10
i2: .word -5
               # int i2=-5
i3: .half 300  # short i3=300
c1: .byte 100 # char c1=100
c2: .byte 'a' # char c2='a '
f1: .float 1.3e-4 # float f1=1.3e-4
d1: .double .001 # double d1=0.001
# int v[3] = \{ 0, -1, 0xfffffffff \}; int w[100];
v: .word 0, −1, 0 xffffffff
w: .word 400
```

System calls

- Many assembler simulators include a small "operating system"
 - The SPIM simulator provides 17 services.
- How to invoke:
 - Call code in register \$v0
 - Other arguments on specific records
 - Invocation by the syscall instruction

System calls

Service	Call code (\$v0)	Arguments	Result
print_int	I	\$a0 = integer	
print_float	2	\$f12 = float	
print_double	3	\$f12 = double	
print_string	4	\$a0 = string	
read_int	5		integer in \$v0
read_float	6		float in \$f0
read_double	7		double in \$f0
read_string	8	\$a0 = buffer, \$a1 = length	
sbrk	9	\$a0 = amount	address in \$v0
exit	10		

System calls

Service	Call code (\$v0)	Arguments	Result
print_char	П	\$a0 (ASCII code)	
read_char	12		\$v0 (ASCII code)
open	13	Equivalent to \$v0=open(\$a0, \$a1, \$a2)	file descriptor in \$v0
read	14	Equivalent to \$v0=read(\$a0, \$a1, \$a2)	read bytes in \$v0
write	15	Equivalent to \$v0=write(\$a0, \$a1, \$a2)	written bytes in \$v0
close	16	Equivalent to \$v0=close(\$a0)	0 in \$v0
exit2	17	End the program. Return the value stored in \$a0	

hello.s

.data

msg_hola: .asciiz "hello world\n"

.text

.glob1 main

main:

Service	Call code (\$v0)	Arguments
print_int	I	\$a0 = integer
print_float	2	\$f12 = float
print_double	3	\$f12 = double
print_string	4	\$a0 = string

```
# printf("hello world\n");
```

li \$v0 4 # syscall code: 4

la \$a0 msg_hola # address where msg_hola starts

syscall ← Operating system invocation instruction

Exercise

```
int valor ;
int valor ;

readInt(&valor) ;
valor = valor + 1 ;
printInt(valor) ;

. . .
```

service	code	arguments	results
print_int	I	\$a0 = integer	
print_float	2	\$f12 = float	
print_double	3	\$f12 = double	
print_string	4	\$a0 = string	
read_int	5		integer in \$v0
read_float	6		float in \$f0
read_double	7		double in \$f0
read_string	8	\$a0=buffer, \$a1=long.	
sbrk	9	\$a0=amount	address in \$v0
exit	10		

Exercise (solution)

```
int valor ;

readInt(&valor) ;

valor = valor + 1 ;
printInt(valor) ;

. . .
```

service	code	arguments	results
print_int	I	\$a0 = integer	
read_int	5		integer en \$v0

```
# readInt(&valor)
li $v0 5
syscall
sw $v0 valor
# valor = valor + 1
add $a0 $v0 1
sw $a0 valor
# printInt
li $v0 1
syscall
```

Instructions and pseudo-instructions

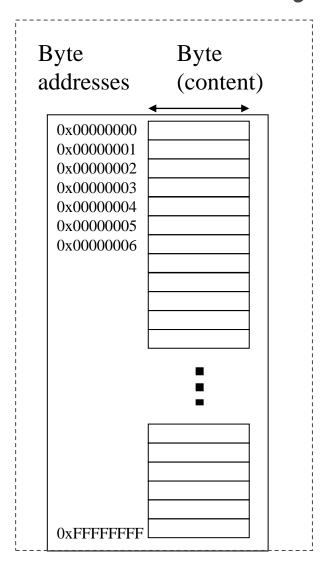
- ▶ There is an assembly instruction per machine instruction :
 - ▶ Each machine instruction occupies 32 bits in MIPS32
 - addi \$t1,\$t0,4
- A pseudo-instruction can be used in an assembler program and it corresponds to one or several assembly instructions:
 - ▶ E.g: li \$v0, 4 move \$t1,\$t0
- In the assembly process, they are replaced by the sequence of assembly instructions that perform the same functionality.
 - E.g.: ori \$v0, \$0, 4 replaces to: li \$v0, 4 addu \$t1, \$0, \$t2 replaces to: move \$t1, \$t2

Other examples of pseudo-instructions

- An assembler pseudoinstruction can correspond to several machine instructions.
 - ▶ li \$t1,0×00800010
 - □ It does not fit in 32 bits, but can be used as a pseudo-instruction.
 - □ It is equivalent to:

```
lui $t1,0x0080
ori $t1,$t1,0x0010
```

MIPS32 memory model



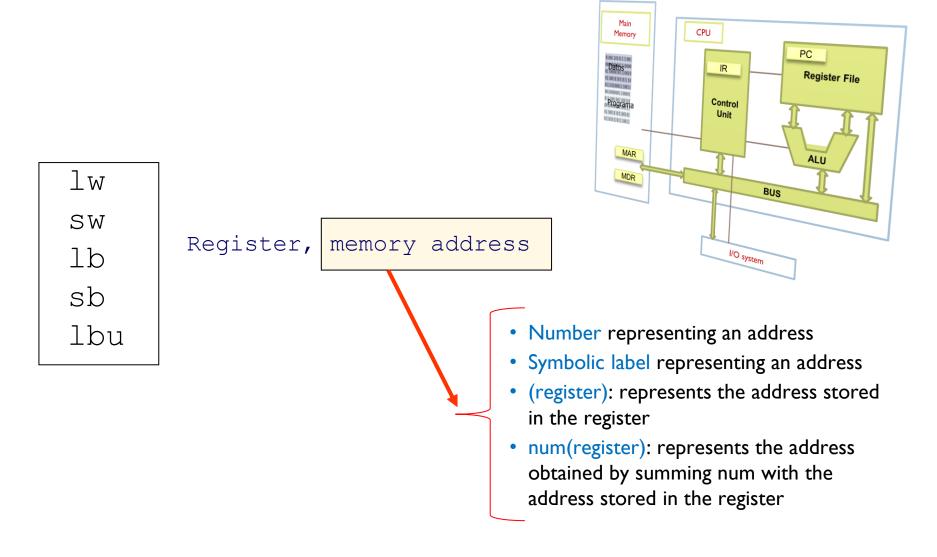
Memory is addressed by bytes:

- 32-bit addresses
- Content of each address: one byte
- Addressable space: 2^{32} bytes = 4 GB

Access can be to:

- Individual bytes
- Words (4 consecutive bytes)

Format of memory access instructions



Memory access instruction formats

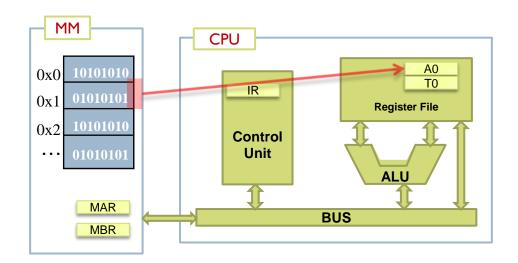
- ▶ lbu \$t0, 0x0F000002
- ▶ 1bu \$t0, labeled
 - Direct addressing. The byte stored in the memory location labeled is loaded into \$t0.
- ▶ lbu \$t0, (\$t|)
 - Indirect register addressing. The byte stored in the memory location stored in \$t1 is loaded in \$t0.
- ▶ lbu \$t0, 80(\$t|)
 - Relative addressing. The byte stored in the memory location obtained by adding the contents of \$t1 with 80 is loaded in \$t0.

Data transfer **bytes**

▶ Copies a byte from memory to a register or vice versa.

Examples:

- Memory to registerlb \$a0, dirlbu \$a0, dir
- Register to memory sb \$t0, dir





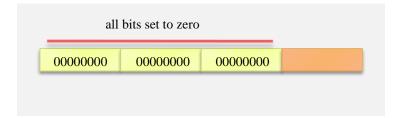


Data transfer

Sign extension

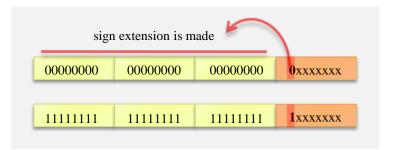
- There are two possibilities when transferring a byte from memory to register:
- A) Transfer without sign, for example: Ibu \$a0, dir



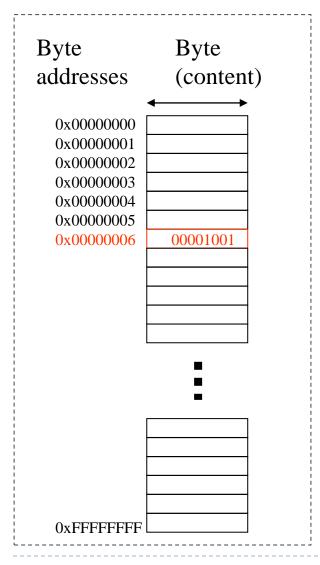


B) Transfer with sign, for example: Ib \$a0, dir





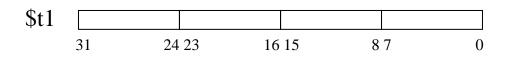
Access to bytes with 1b (load byte)



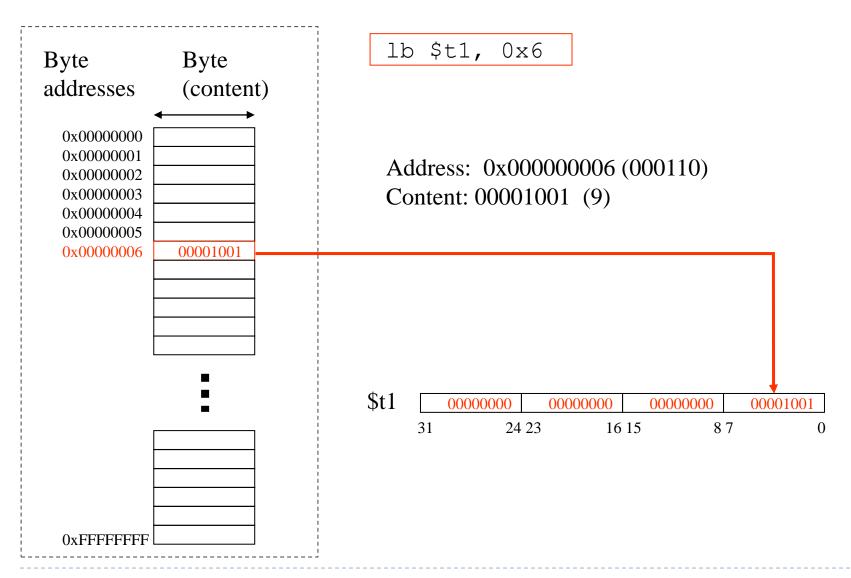
lb \$t1, 0x6

Address: 0x00000006 (000110)

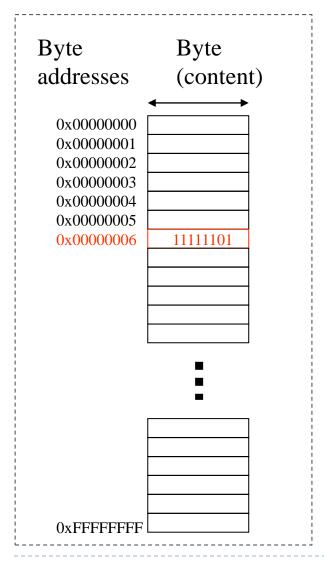
Content: 00001001 (9)



Access to bytes with lb



Access to bytes with lb



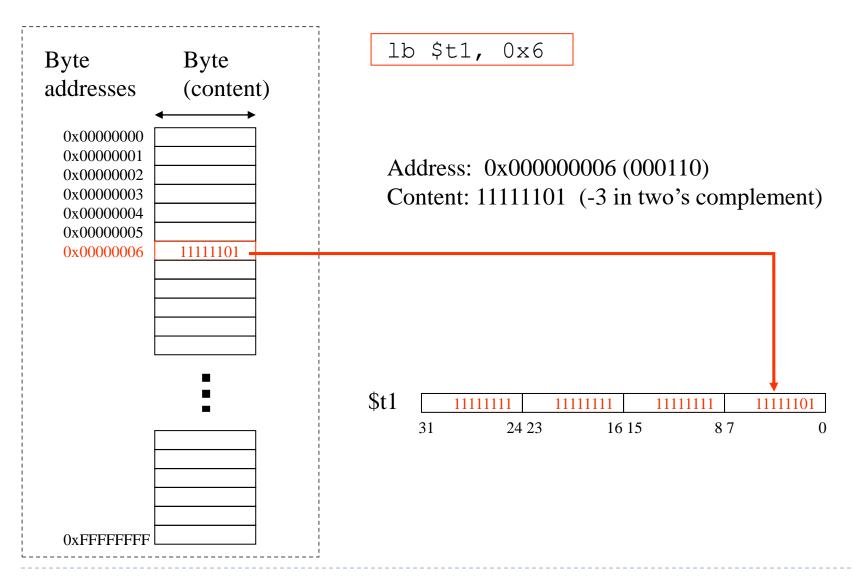
lb \$t1, 0x6

Address: 0x00000006 (000110)

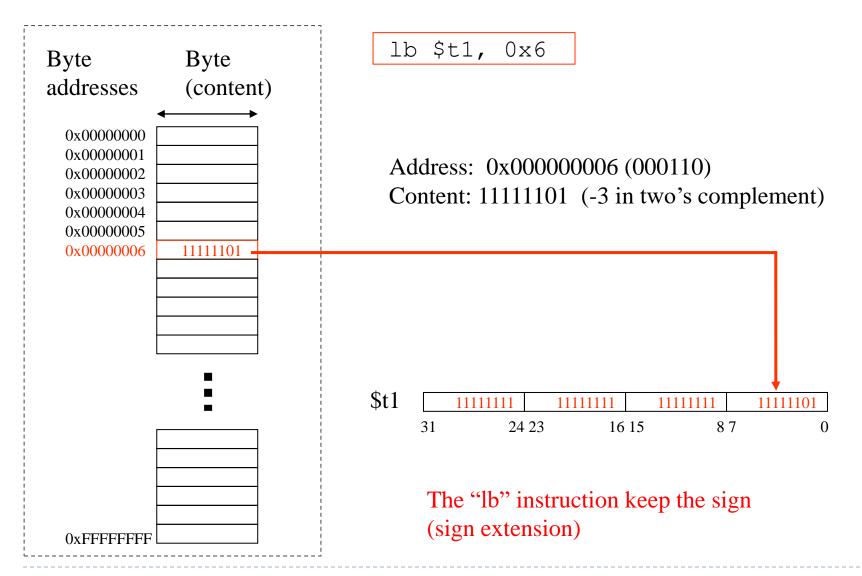
Content: 11111101 (-3 in two's complement)



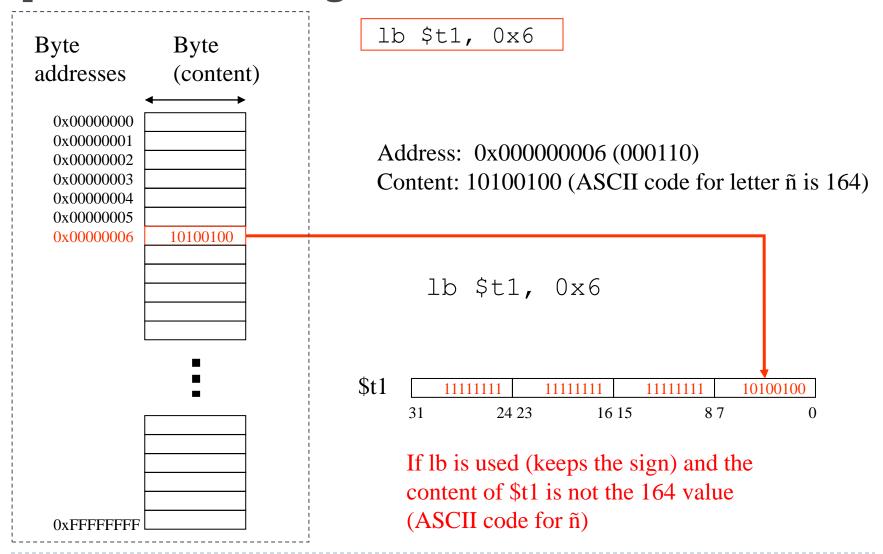
Access to bytes with lb



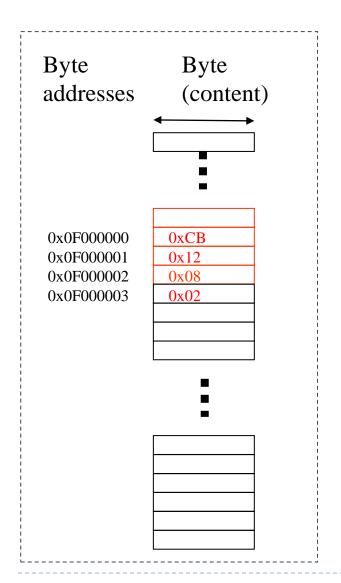
Access to bytes with lb



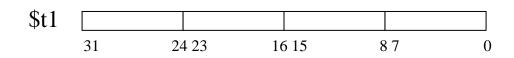
Access to bytes with 1b problems accessing characters



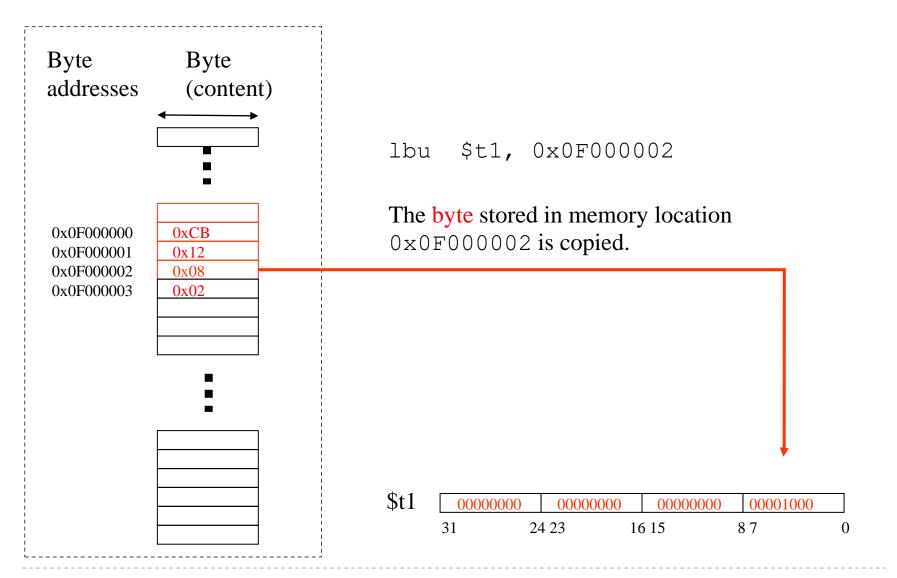
Access to bytes with lbu (load byte unsigned)



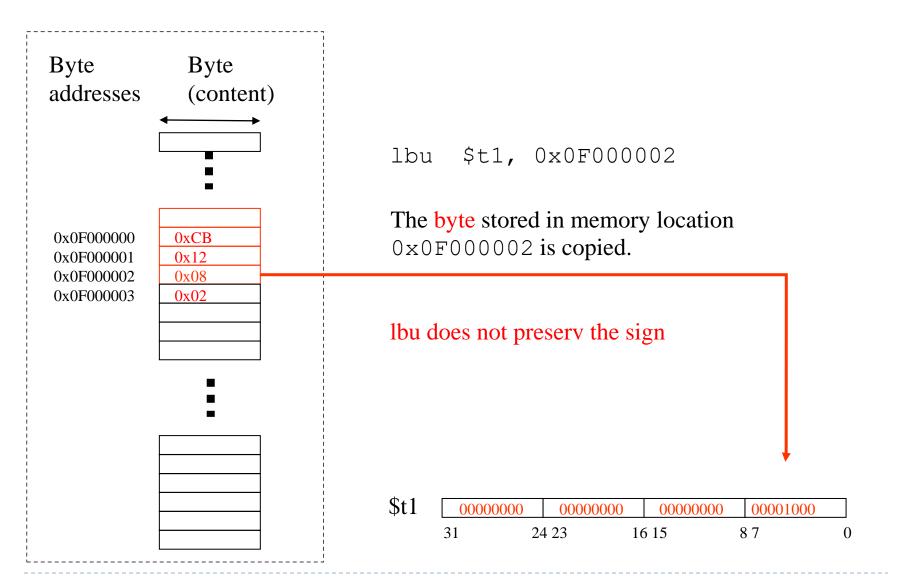
lbu \$t1, 0x0F000002



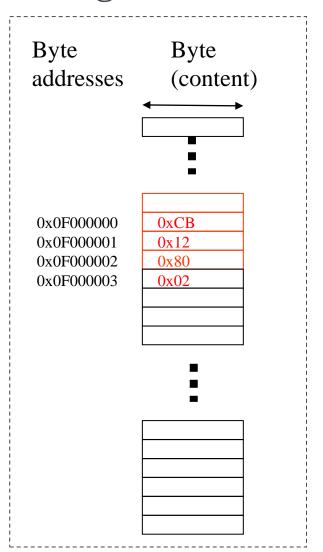
Access to bytes with lbu (load byte unsigned)



Access to bytes with lbu (load byte unsigned)



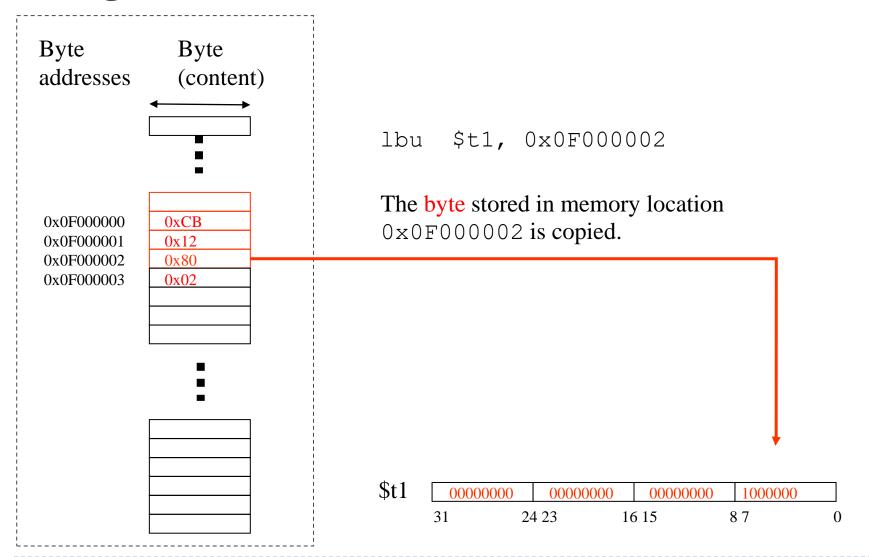
Access to bytes with lbu (load byte unsigned) No sign-extension

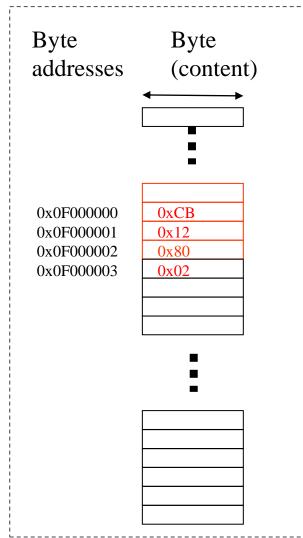


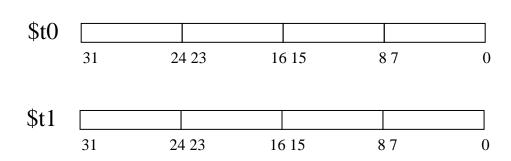
lbu \$t1, 0x0F000002

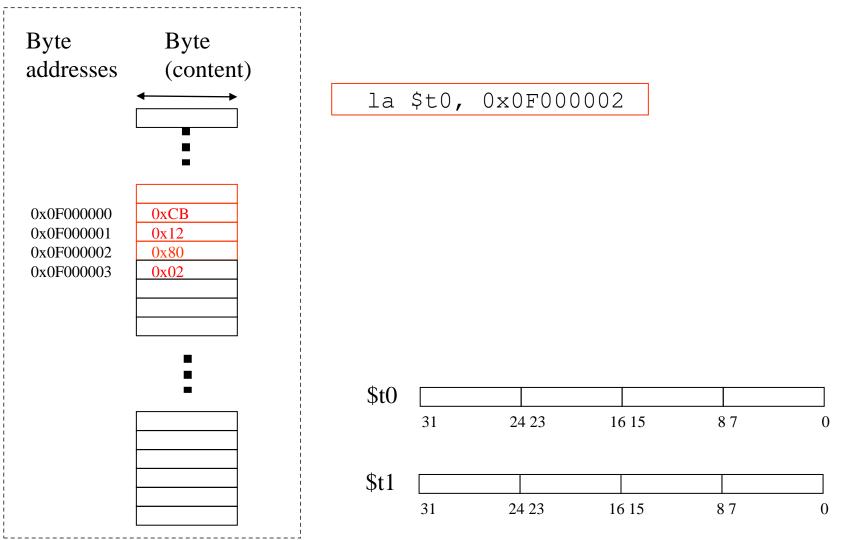


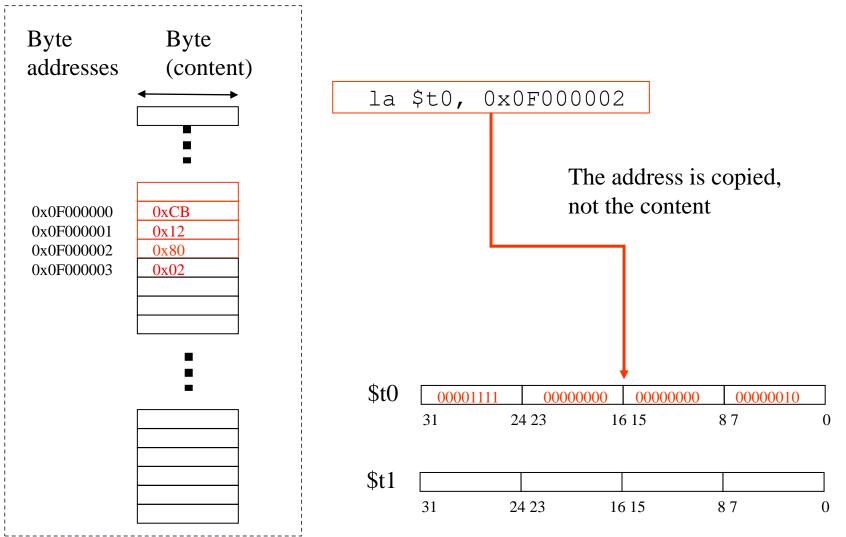
Access to bytes with lbu (load byte unsigned) No sign-extension

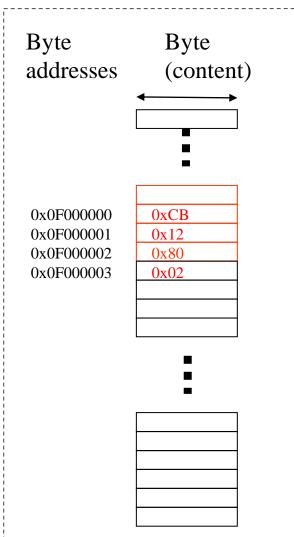




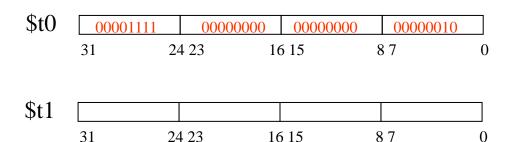


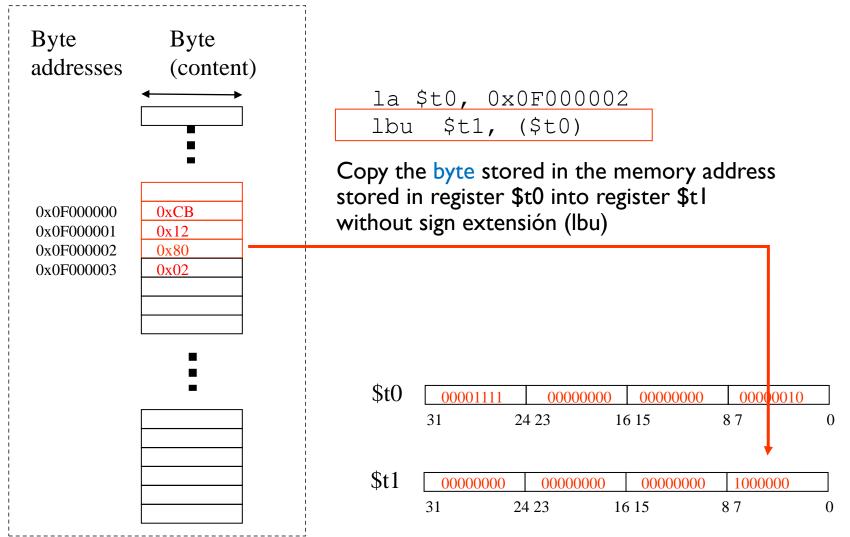






la \$t0, 0x0F000002 lbu \$t1, (\$t0)



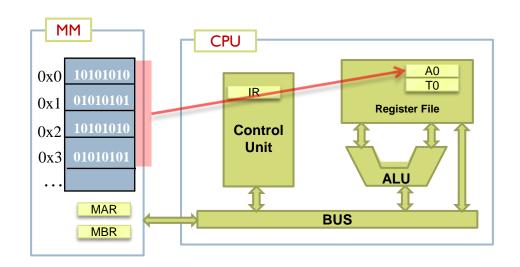


Data transfer words

▶ Copies a word from memory to a register or vice versa.

Examples:

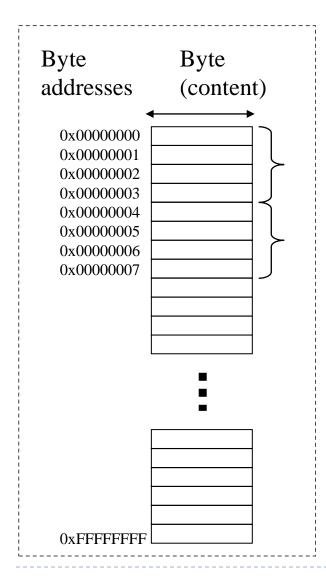
- Memory to registerlw \$a0 (\$t0)
- Register to memorysw \$a0 (\$t0)







Accessing to words



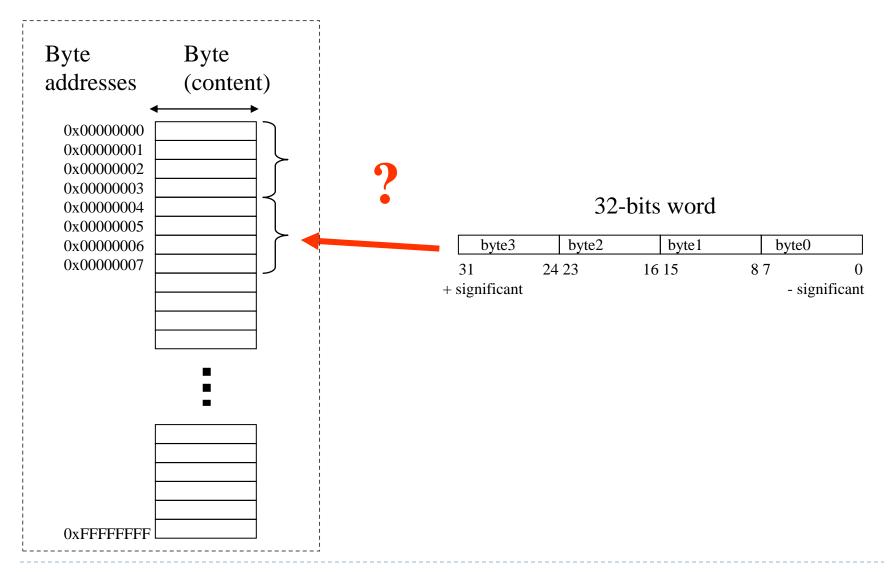
A word: 4 bytes in a 32-bits processor

Word stored starting at byte 0

Word stored starting at byte 4

Las palabras (32 bits, 4 bytes) se almacenan utilizando cuatro posiciones consecutivas de memoria, comenzando la primera posición en una dirección múltiplo de 4

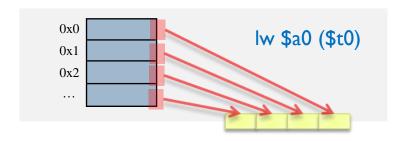
Accessing to words



Data transfer byte order

▶ There are 2 types of byte order:

Little-endian ('small' address ends the word...)



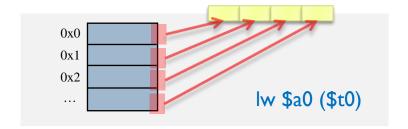
(intel)

AMD

Big-endian ('big' address ends the word...)



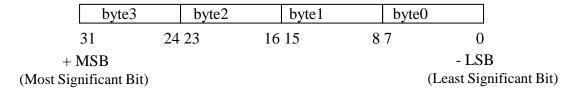
(bi-endian)





Storing words in memory

32-bit word



Α	byte3
A+1	byte2
A+2	byte1
A+3	byte0

BigEndian

byte0
byte1
byte2
byte3

LittleEndian



BigEndian

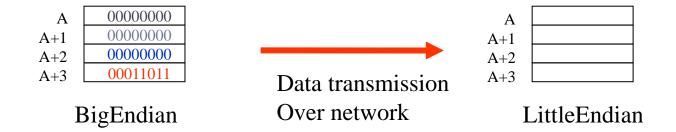
A	00011011
A+1	00000000
A+2	00000000
A+3	00000000

LittleEndian

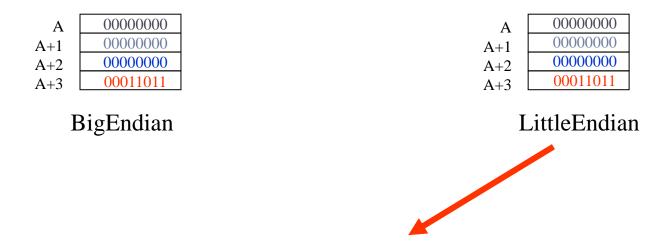
Communication problems in computers with different architectures



Communication problems in computers with different architectures



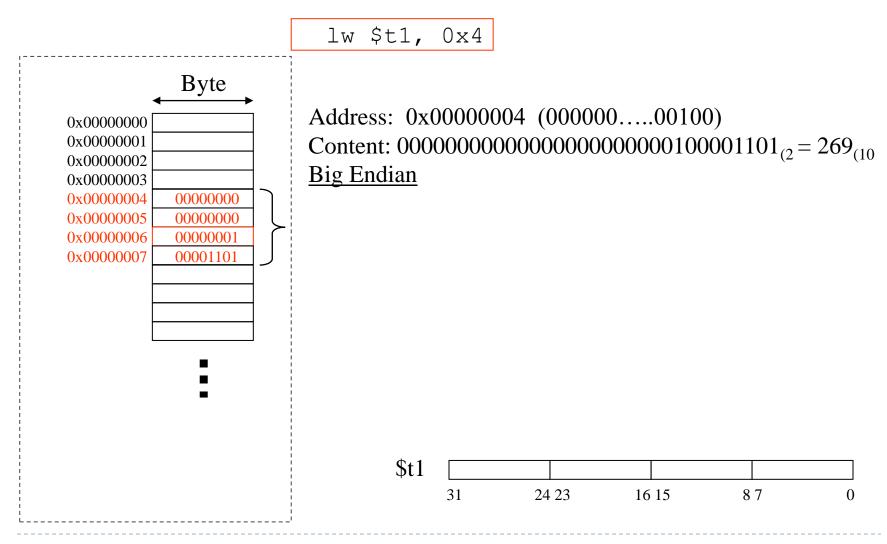
Communication problems in computers with different architectures

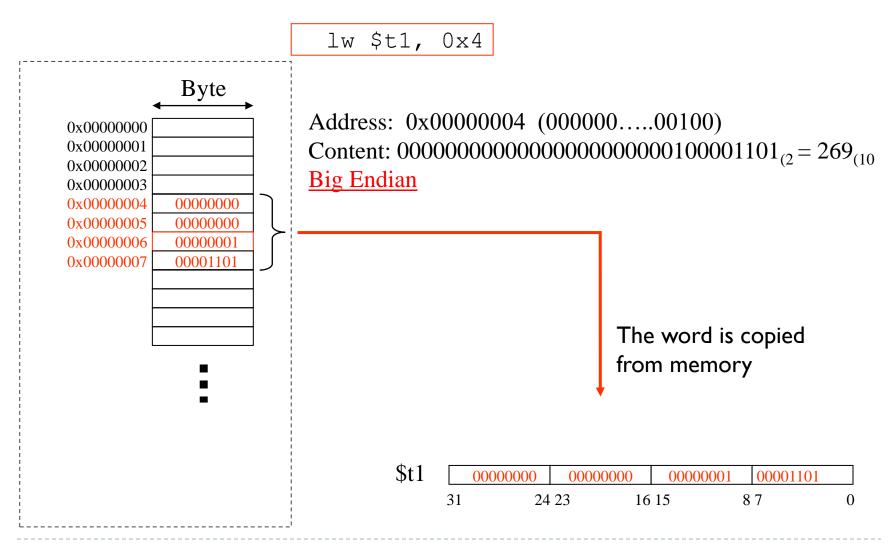


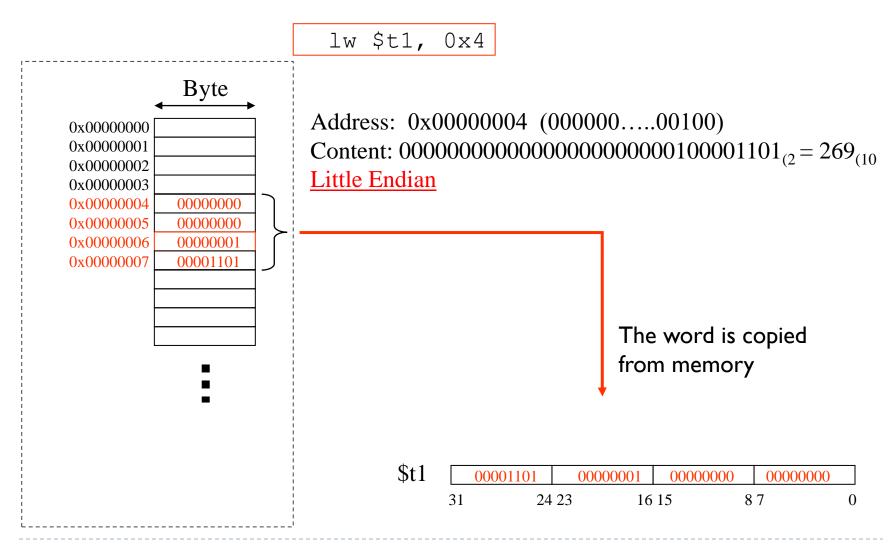
Example

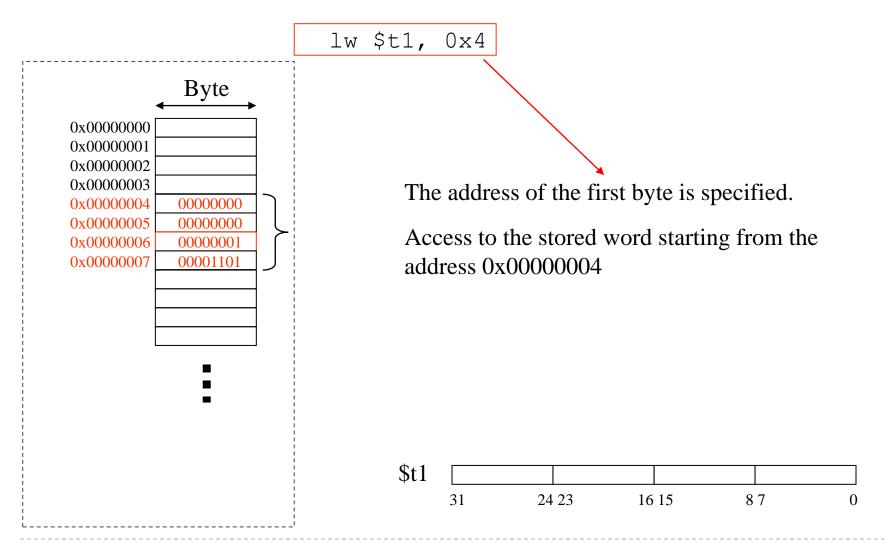
endian.s

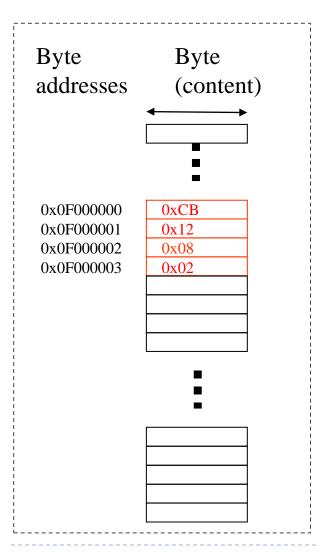
```
.data
 b1: .byte 0x00, 0x11, 0x22, 0x33
.text
.globl main
main:
   lw $t0 b1
```



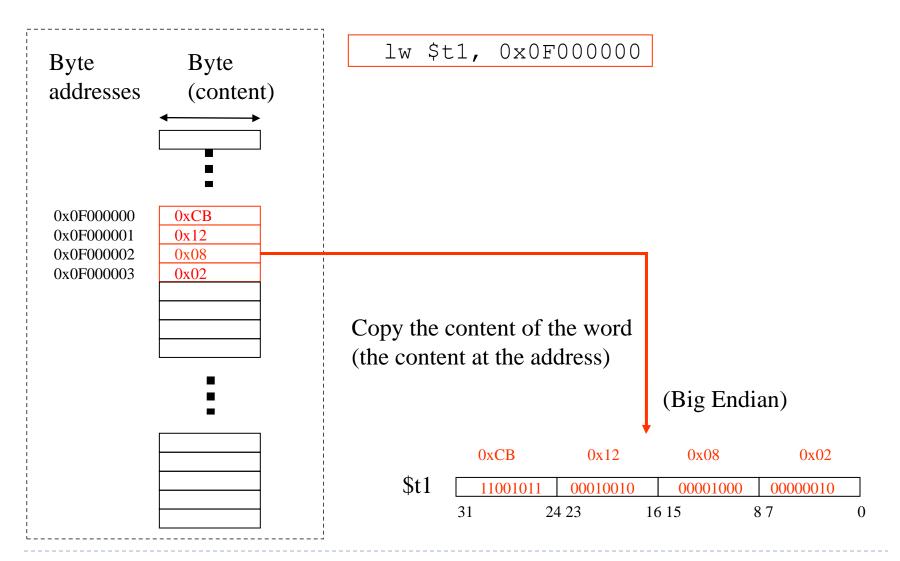


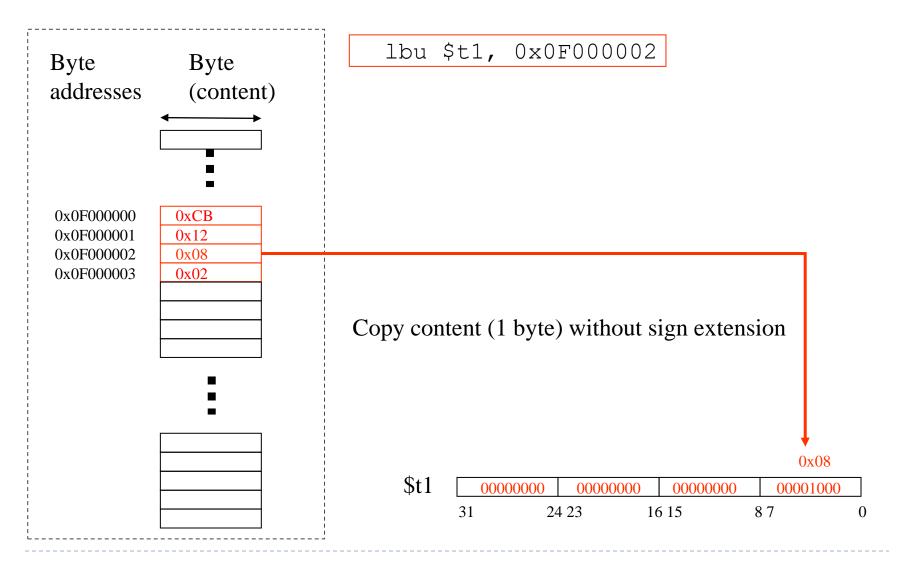


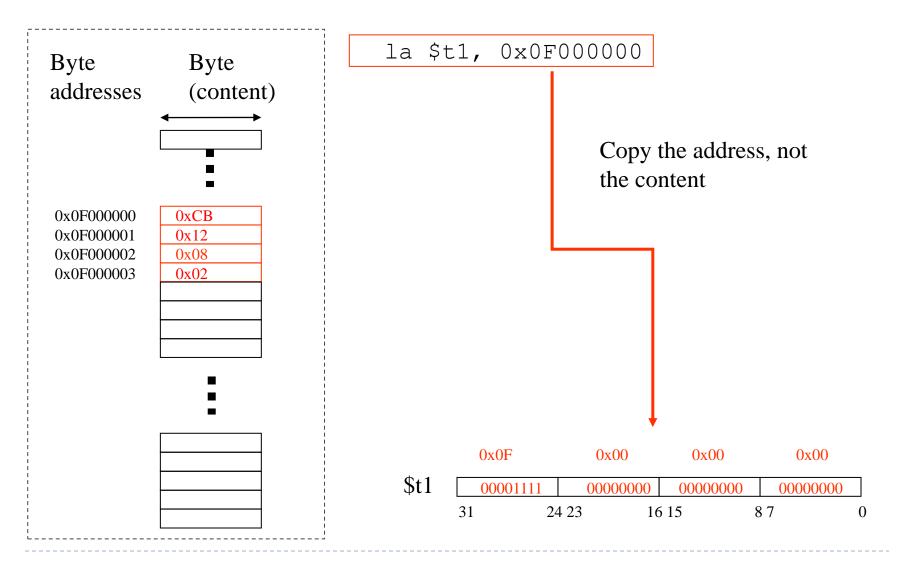




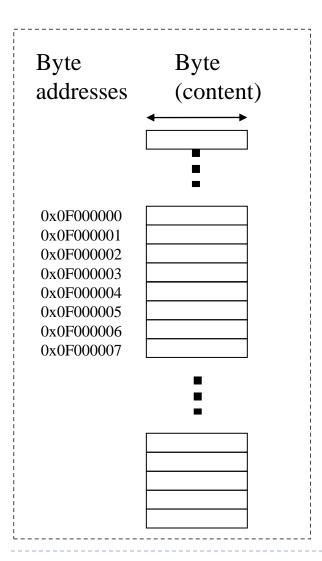
lw \$t1, 0x0F000000





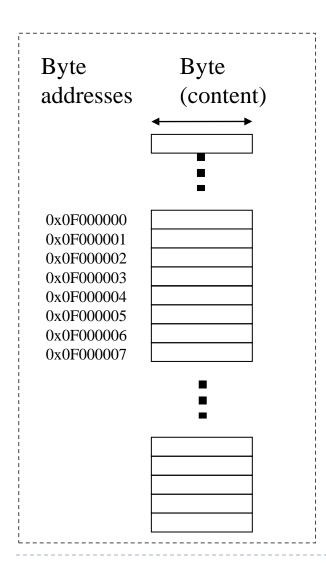


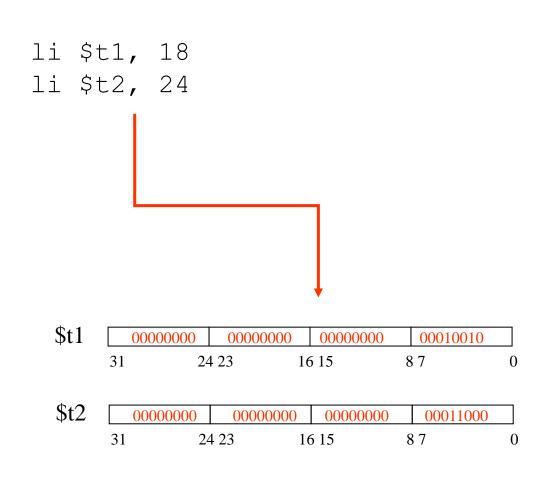
Example



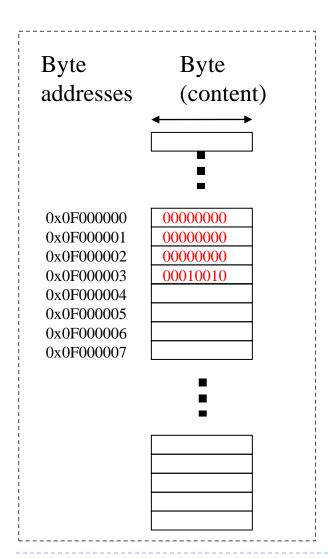
li \$t1, 18 li \$t2, 24

Example



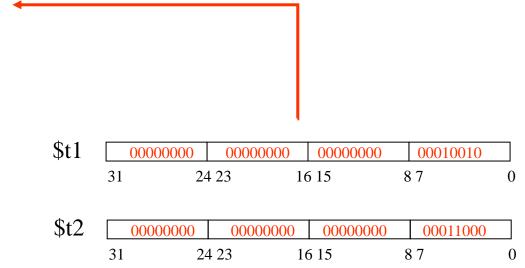


Write word in memory

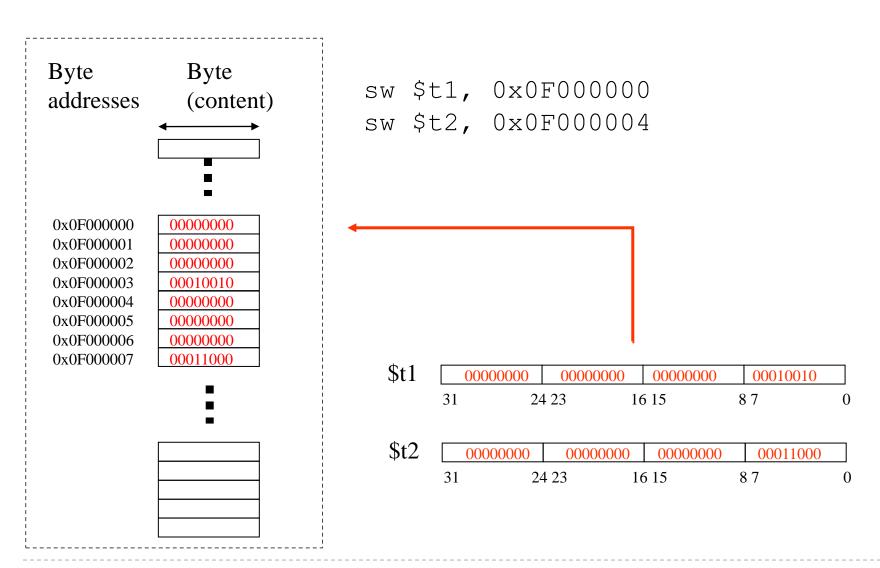


sw \$t1, 0x0F000000

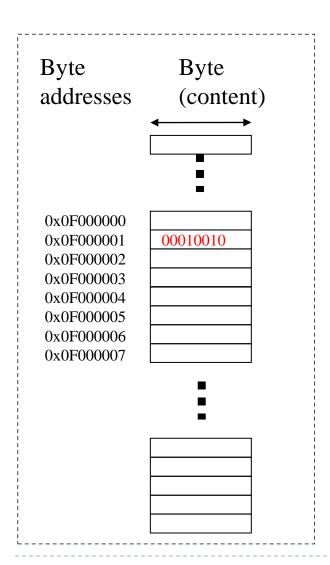
Write the content of a register into memory (the full word value stored in the register)



Write word in memory

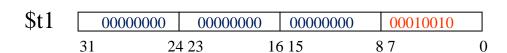


Write byte in memory



sb \$t1, 0x0F000001

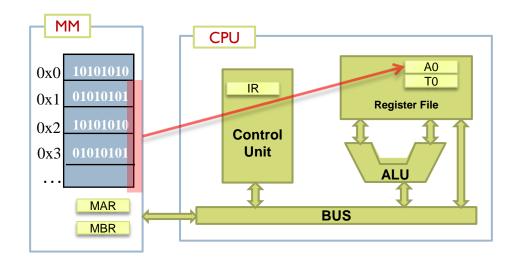
Write the less significant byte of register \$t1 in memory



Data transfer alignment and access size

Peculiarities:

- Alignment of elements in memory
- Default access size



Data alignment

In general:

A data of K bytes is aligned when the address D used to access this data fulfills the condition:

 $D \mod K = 0$

Data alignment implies:

- Data of 2 bytes are stored in even addresses
- Data of 4 bytes are stored in addresses multiple of 4
- Data of 8 bytes (double) are stored in addresses multiple of 8

Data alignment

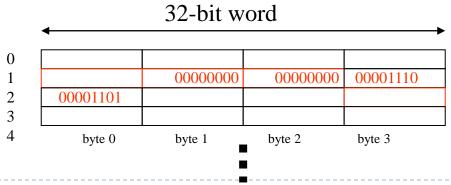
75

.data ▶ The alignment requires the address to be a .byte 0x0f multiple of the word size: .word 10 31 23 15. 0 This word is aligned, next ones are not 08 address 12 16 20 24

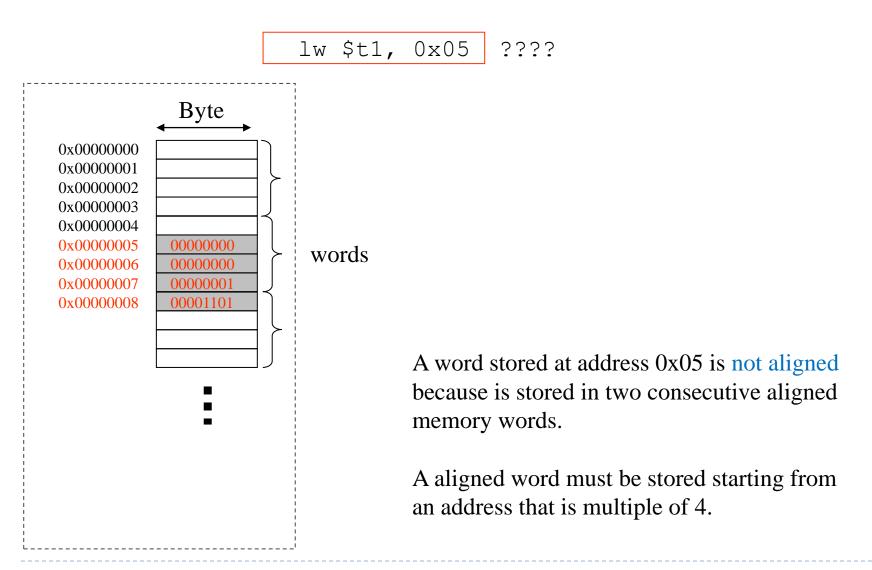
Data alignment

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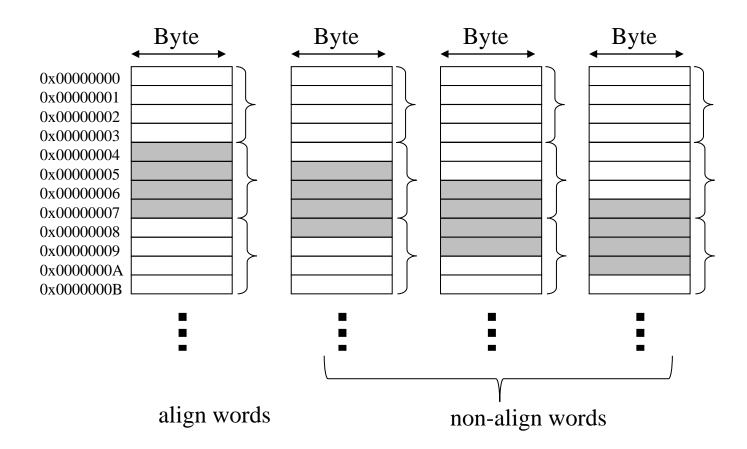
- Many computers does not allow the access to not aligned data:
 - Goal: reduce the number of memory accesses
 - Compilers assign addresses aligned to variables
- Some processors, such as Intel models, allow the access to not aligned data:
 - Non-aligned data needs several memory access



Non-aligned data

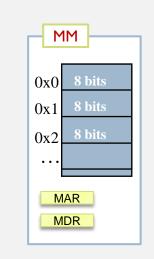


Non-aligned data



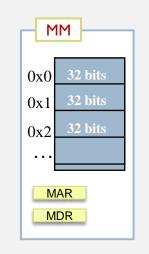
Word-level or byte-level addressing

- ▶ The main memory is similar to a large one-dimensional vector of items.
- A memory address is the index of one item in the vector.
- There are two types of addressing:
 - Byte addressing



- Each memory element is 1 byte
- Transferring a word means transferring 4 bytes (in a 32-bit CPU)

Word addressing



- Each memory element is a word
- b means transferring one word and keeping one byte.

Summary

- The instructions and data of a program must be loaded in memory for the execution (process)
- All data and instructions are stored in memory so all have an associated memory address where is stored
- ▶ In a 32-bit computer such as MIPS 32:
 - Registers have 32 bits
 - Memory can store bytes (8 bits)
 - ▶ Instructions: memory → register: lb, lbu
 - ▶ Instructions: register → memory: sb
 - Memory can store words (32 bits)
 - ▶ Instructions: memory → register: 1w
 - ▶ Instructions: register → memory: SW

Format of the memory access instructions **summary**

lw
sw
lb
sb
lbu

• Number that represent an address
• Symbolic label that represents the
associated address
• (register): address is stored in the

the address stored in the register

num(register): represent the address that is obtained by adding num with

register

Memory access instruction formats summary

- ▶ lbu \$t0, 0x0F000002
- ▶ 1bu \$t0, labeled
 - Direct addressing. The byte stored in the memory location labeled is loaded into \$t0.
- ▶ lbu \$t0, (\$t|)
 - Indirect register addressing. The byte stored in the memory location stored in \$t1 is loaded in \$t0.
- ▶ lbu \$t0, 80(\$t|)
 - Relative addressing. The byte stored in the memory location obtained by adding the contents of \$t1 with 80 is loaded in \$t0.

Instructions to write in memory summary

- > sw \$t0, 0x0F000000
 - ► Copy the word stored in \$t0 in the address 0x0F00000
- sb \$t0, 0x0F000000
 - ► Copy the (least significant) byte stored in \$t0 in the address 0x0F00000

Assembly data types

Basic

- Booleans
- Characters
- Integers
- Decimals (float/double)

Compound

- Vector
- String
- Matrix
- Others... (struct)

Basic data types **booleans**

```
bool_t b1 = false;
bool_t b2 = true;
....
main ()
{
    b1 = true ;
....
}
```

```
.data
b1: .byte 0 # 1 byte
b2: .byte 1
. . .
.text
.globl main
 main: la $t0 b1
        li $t1 1
        sb $t1 ($t0)
```

Basic data types characters

```
char c1 ;
char c2 = 'a';
...
main ()
{
    c1 = c2;
...
}
```

```
.data
c1: .space 1 # 1 byte
c2: .byte 'a'
. . .
.text
.globl main
 main: la $t0 c1
       lbu $t1 c2
        sb $t1 ($t0)
```

Basic data types Integers

```
int result;
int op1 = 100;
int op2 = -10;
main ()
  result = op1+op2;
```

```
.data
.align 2
result: .word 0 # 4 bytes
op1: .word 100
op2: .word -10
.text
.globl main
main: lw $t1 op1
       lw $t2 op2
       add $t3 $t1 $t2
       la $t4 result
       sw $t3 ($t4)
```

Basic data types

Integers

global variable without initial value

```
int result;
int op1 = 100;
int op2 = -10;
     global variable with initial value
main ()
  result = op1+op2;
```

```
.data
.align 2
result: .word 0 # 4 bytes
         .word 100
op1:
         ?word -10
op2:
.text
.globl main
main: lw $t1 op1
        lw $t2 op2
        add $t3 $t1 $t2
        la $t4 result
        sw $t3 ($t4)
```

Exercise

Write in MIPS-32 assembly a fragment of code with the same functionality that:

```
int b;
int a = 100;
int c = 5;
int d;
main ()
{
   d = 80;
   b = -(a+b*c+a);
}
```

Assuming that a, b, c and d are variables stored in memory

Basic data types float

```
float result;
float op1 = 100;
float op2 = 2.5
main ()
  result = op1 + op2;
```

```
.data
.align 2
   result: .word 0 # 4 bytes
   op1: .float 100
   op2:
           .float 2.5
.text
      .globl main
main: 1.s $f0 op1
       1.s $f1 op2
       add.s $f3 $f1 $f2
       s.s $f3 result
```

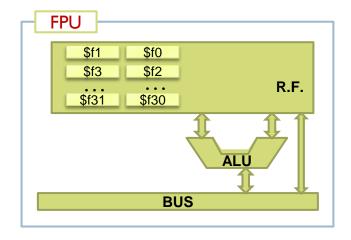
Basic data types double

```
double result ;
double op1 = 100;
double op2 = -10.27;
main ()
  result = op1 * op2;
```

```
.data
.align 3
   result: .space 8
   op1: .double 100
   op2:
           .double -10.27
.text
      .globl main
main: 1.d $f0 op1 # ($f0,$f1)
      1.d $f2 op2 # ($f2,$f3)
      mul.d $f6 $f0 $f2
      s.d $f6 result
```

Floating point. IEEE 754

- ▶ The coprocessor I has 32 registers of 32 bits (4 bytes) each.
 - It is possible to work with single or double precision
- Simple precision (32 bits):
 - From \$f0 to \$f3 I
 - E.g.: add.s \$f0 \$f1 \$f5 f0 = f1 + f5
 - Other operations:
 - add.s, sub.s, mul.s, div.s, abs.s
- Doble precision (64 bits):
 - Registers used in pairs
 - E.g.: add.d f0 f2 f8(f0,f1) = (f2,f3) + (f8,f9)
 - Other operations:
 - add.d, sub.d, mul.d, div.d, abs.d

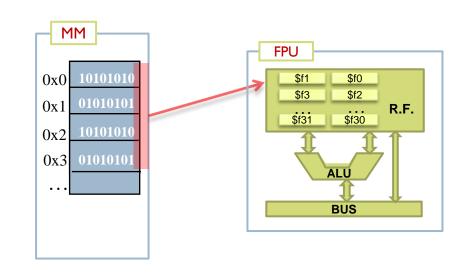


Data transfer IEEE 754

▶ Copies a number from memory to a register or vice versa.

Examples:

- Memory to registerl.s \$f0 dir ll.d \$f2 dir2
- Register to memorys.s \$f0 dir ls.d \$f0 dir 2

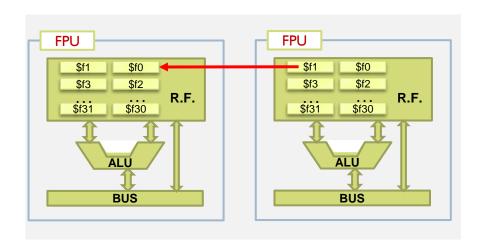






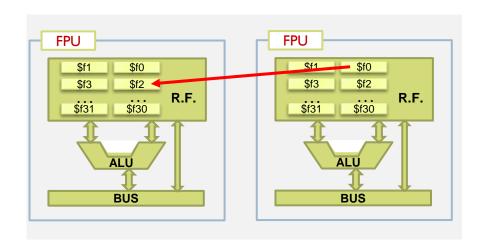
Operations with registers (FPU, FPU)





Operations with registers (FPU, FPU)

mov.d \$f0 \$f2

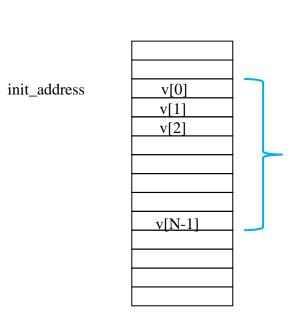


 $(\$f0, \$f1) \leftarrow (\$f2, \$f3)$

- Collection of data ítems stored consecutively in memory
- The address of the j element can be computed as:

init_address + j * p

Where **p** is the size of each item



```
int vec[5] ;
...
main ()
{
    vec[4] = 8;
}
```

```
.data
.align 2 # next item aligned to 4
vec: .space 20 # 5 items * 4 bytes/item
.text
.globl main
main:
       la $t1 vec
       li $t2 8
       sw $t2 16($t1)
```

```
int vec[5] ;
...
main ()
{
    vec[4] = 8;
}
```

```
.data
.align 2 # next item aligned to 4
vec: .space 20 # 5 items * 4 bytes/item
.text
.globl main
main:
       li $t0 16
       la $t1 vec
       add $t3, $t1, $t0
       li $t2 8
       sw $t2, ($t3)
```

```
int vec[5] ;
...
main ()
{
    vec[4] = 8;
}
```

```
.data
.align 2 # next item align to 4
vec: .space 20 # 5 items * 4 bytes/item
.text
main:
       li $t2 8
       li $t1 16
        sw $t2 vec($t1)
        . . .
```

Exercise

- ▶ Let V be an array of integer elements
 - V represents the initial address of the array
- ▶ What is the address of the V[5] item?
- Which are the instruction to load in register \$t0 the value of v[5]?

- Let V be an array of integer elements
 - V represents the initial address of the array
- ▶ What is the address of the V[5] item?
 - V + 5*4
- Which are the instruction to load in register \$t0 the value of v[5]?
 - ▶ li \$t1, 20
 - ► Iw \$t0, v(\$t1)

Compound data types **String**

```
    Array of bytes
```

• '\0' ends string

```
char c1 ;
char c2 = 'h' ;
char *ac1 = "hola" ;
...

main ()
{
   printf("%s",ac1) ;
...
}
```

```
.data
c1: .space 1
                       # 1 byte
c2: .byte 'h'
ac1: .asciiz "hola"
.text
.globl main
main:
        li $v0 4
         la $a0 ac1
         syscall
```

String layout in memory

```
// strings
char c1[10] ;
char ac1[] = "hola" ;
```

```
# strings
c1: .space 10  # 10 byte
ac1: .asciiz "hola" # 5 bytes (!)
ac2: .ascii "hola" # 4 bytes
```

```
ac1:
           'h'
                      0 \times 0108
                                             ac2:
                                                         \h'
                                                                    0 \times 0108
           10'
                      0 \times 0109
                                                         10'
                                                                    0 \times 0109
           111
                      0x010a
                                                         11'
                                                                    0x010a
           \a'
                      0 \times 010b
                                                         'a'
                                                                    0 \times 010 b
            0
                      0 \times 010c
                                                                    0x010c
                      0x010d
                                                                    0x010d
```

Exercise

```
// global variables
char v1;
int v2;
float v3 = 3.14;
char v4[10];
char v5 = "ec";
int v6[] = { 20, 22 };
```

```
// variables globales
char v1;
int v2 ;
float v3 = 3.14 ;
char v4 = "ec" ;
int v5[] = { 20, 22 } ;
```

```
.data
v1: .byte 0
.align 2
v2: .space 4
v3: .float 3.14
v4: .asciiz "ec"
.align 2
v5: .word 20, 22
```

v1: ? ? ? 0x0100 0x0101 0x0102 0x0103

```
.data
v1: .byte 0
.align 2
v2: .space 4
v3: .float 3.14
v4: .asciiz "ec"
.align 2
v5: .word 20, 22
```

v1 :	0	0x0100
	?	0x0101
	?	0x0102
	?	0x0103
v 2:	0	0x0104
	0	0x0105
	0	0x0106
	0	0x0107
v 3:	(3.14)	0x0108
	(3.14)	0x0109
	(3.14)	0x010A
	(3.14)	0x010B
v4 :	\e'	0x010C
	\c'	0x010D
	0	0x010E
		0x010F
v 5:	(20)	0x0110
	(20)	0x0111
	(20)	0x0112
	(20)	

```
.data
v1: .byte 0
.align 2
v2: .space 4
v3: .float 3.14
v4: .asciiz "ec"
.align 2
v5: .word 20, 22
```

Compound data types String length

```
char c1;
char c2 = 'h';
char *ac1 = "hola" ;
char *c;
main ()
  c = ac1; int 1 = 0;
  while (c[1] != NULL) {
        1++;
  printf("%d", 1);
```

Compound data types String length

```
char c1;
char c2 = 'h';
char *ac1 = "hola" ;
char *c;
main ()
  c = ac1; int 1 = 0;
  while (c[1] != NULL) {
        1++;
  printf("%d", 1);
```

```
.data
c1: .space 1 # 1 byte
c2: .byte 'h'
ac1: .asciiz "hola"
.align 2
c: .word 0 # pointer => address
.text
.globl main
main: la $t0, ac1
         li $a0, 0
         lbu $t1, ($t0)
        begz $t1, fin
  buc:
         addi $t0, $t0, 1
         addi $a0, $a0, 1
         lbu $t1, ($t0)
         b buc
  fin: li $v0 1
         syscall
```

Arrays and strings

Review (in general) :

▶ sw \$t0, 4(\$s3) # M[\$s3+4] \$t0

Exercise

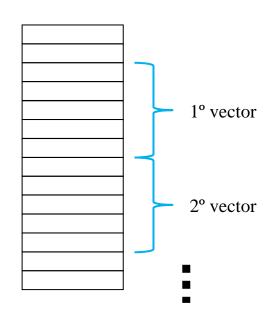
- Write a program that:
 - ▶ Calculate the number of occurrences of a char in a string
 - String address stored in \$a0
 - Char to look for in \$a1
 - Result must be stored in \$v0

Compound data types Matrix

- A matrix m x n consists of m vectors (m rows) of length n
- Usually stored by rows
- The element a_{ij} is stored in the address:

$$init_address + (i \cdot n + j) \times p$$

where p is the size of each item



Compound data types Matrix

```
.data
.align 2  # next item align to 4
Vec: .space 20 # 5 item * 4 bytes/item
mat: .word 11, 12, 13
     .word 21, 22, 23
.text
.globl main
 main:
        lw $t1 mat+0
         lw $t2 mat+12
         add $t3 $t1 $t2
         sw $t3 mat+4
```

Example: integer matrix

```
int vec[5];
int mat[2][3] = \{\{11,12,13\},
                  {21,22,23}};
main ()
 mat[1][2] = mat[1][1] +
  mat[2][1];
```

```
.data
align 2
Vec: .space 20 # 5 items*4 bytes/item
mat: .word 11, 12, 13
     .word 21, 22, 23
. . .
.text
.qlobl main
 main:
         lw $t1 mat+0
         lw $t2 mat+12
         add $t3 $t1 $t2
         sw $t3 mat+4
```

Tips

- Do not program directly in assembler
 - ▶ Better to first do the design in DFD, Java/C/Pascal...
 - Gradually translate the design to assembler.
- Sufficiently comment the code and data
 - By line or by group of lines comment which part of the design implements.
- ▶ Test with enough test cases
 - Test that the final program works properly to the given specifications.

Exercise

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- Write an assembly program that:
 - ▶ Load the value -3.141516 in register \$f0
 - Obtain the exponent and mantissa values stored in the register
 \$f0 (IEEE 754 format)
 - Display the sign
 - Display the exponent
 - Display the mantissa

```
.data
  newline: .asciiz "\n"
.text
.qlobl main
main:
         li.s $f0, -3.141516
         # print value
         mov.s $f12, $f0
         li $v0, 2
         syscall
         la $a0, newline
         li $v0, 4
         syscall
         # copy to processor
         mfc1 $t0, $f12
```

```
li $s0, 0x80000000
                      #sign
and $a0, $t0, $s0
srl $a0, $a0, 31
li $v0, 1
syscall
la $a0, newline
li $v0, 4
syscall
li $s0, 0x7F800000
                      #exponent
and $a0, $t0, $s0
srl $a0, $a0, 23
li $v0, 1
syscall
la $a0, newline
li $v0, 4
syscall
                      #mantissa
li $s0, 0x007FFFFF
and $a0, $t0, $s0
li $v0, 1
syscall
jr $ra
```