#### **ARCOS Group**

uc3m Universidad Carlos III de Madrid

## Lesson 3 (II)

Fundamentals of assembler programming

Computer Structure

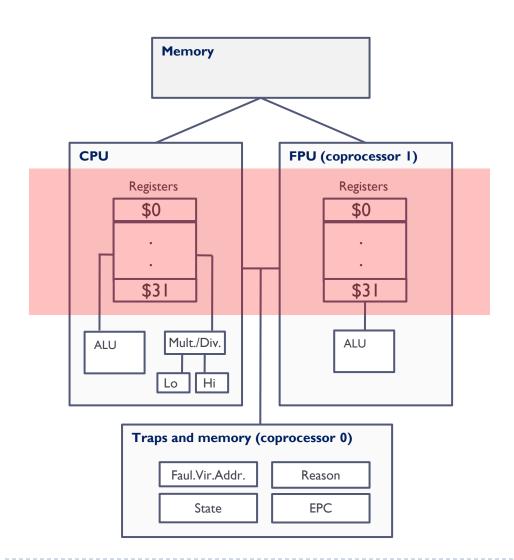
Bachelor in Computer Science and Engineering



#### Contents

- Basic concepts on assembly programming
- MIPS32 assembly language, memory model and data representation
  - ▶ MIPS architecture (II): registers and memory
  - Assembler directives
  - System services
  - Instructions (II): memory access
- Instruction formats and addressing modes
- Procedure calls and stack convention

#### MIPS-32 architecture

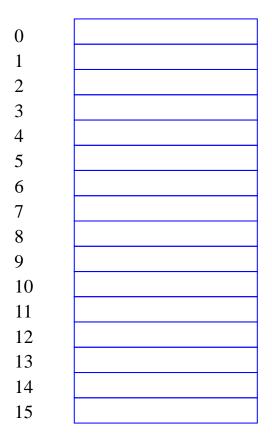


#### ▶ MIPS 32

- > 32 bits processor
- RISC type
- CPU + auxiliary coprocessors

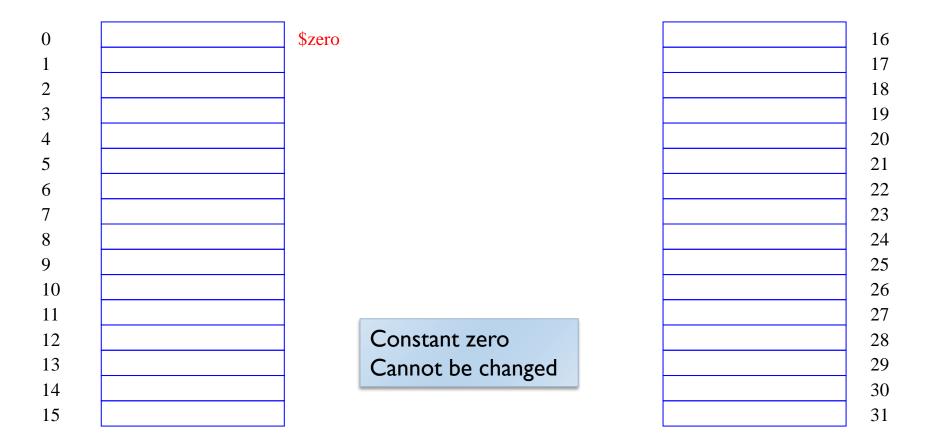
#### Coprocessor 0

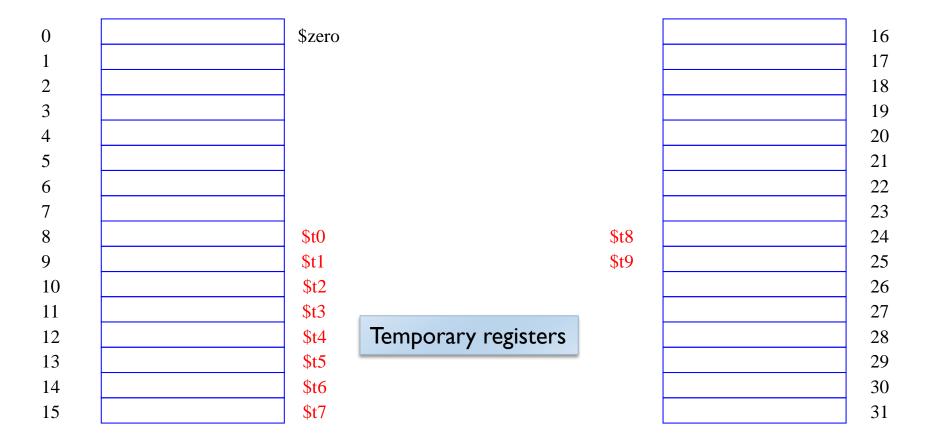
- exceptions, interrupts and virtual memory system
- Coprocessor I
  - FPU (floating point unit)

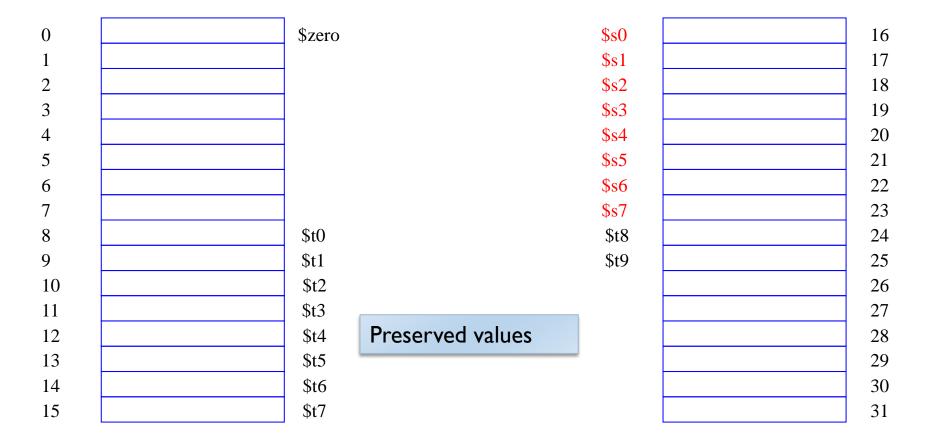


- 32 registers
  - □ 4 bytes of size (one word)
  - □ Name starts with \$ at the beginning
- Usage Convention
  - Reserved
  - Arguments
  - □ Results
  - □ Temporary
  - Pointers

16
17
18
19
20
21
22
23
24
25
26
27
28
29
 30
31







0	\$zero	\$s0	16
1		\$s1	17
2	\$v0	\$s2	18
3	\$v1	\$s3	19
4	\$a0	\$s4	20
5	\$a1	\$s5	21
6	\$a2	<b>\$</b> s6	22
7	\$a3	\$s7	23
8	\$tO	\$t8	24
9	\$t1	\$t9	25
10	\$t2		26
11	\$t3		27
12	\$t4 Arguments a	nd	28
13	\$t5 functions sup		29
14	\$t6	\$fp	30
15	\$t7	\$ra	31

0	\$zero		\$s0	16
1	\$at		<b>\$</b> s1	17
2	\$v0		\$s2	18
3	\$v1		\$s3	19
4	\$a0		\$s4	20
5	\$a1		\$s5	21
6	\$a2		<b>\$</b> s6	22
7	\$a3		\$s7	23
8	\$t0		\$t8	24
9	\$t1		\$t9	25
10	\$t2		\$k0	26
11	\$t3		<b>\$</b> k1	27
12	\$t4	Others	\$gp	28
13	\$t5		\$sp	29
14	\$t6		\$fp	30
15	\$t7		\$ra	31

# Register File (integers) summary

Symbolic name	Number	Usage
zero	0	Constant 0
at	I	Reserved for assembler
v0, v l	2, 3	Results of functions
a0,, a3	4,, 7	Function arguments
t0,, t7	8,, 15	Temporary (NO preserved across calls)
s0,, s7	16,, 23	Saved temporary (preserved across calls)
t8, t9	24, 25	Temporary (NO preserved across calls)
k0, k1	26, 27	Reserved for operating system
gp	28	Pointer to global area
sp	29	Stack pointer
fp	30	Frame pointer
ra	31	Return address (used by function calls)

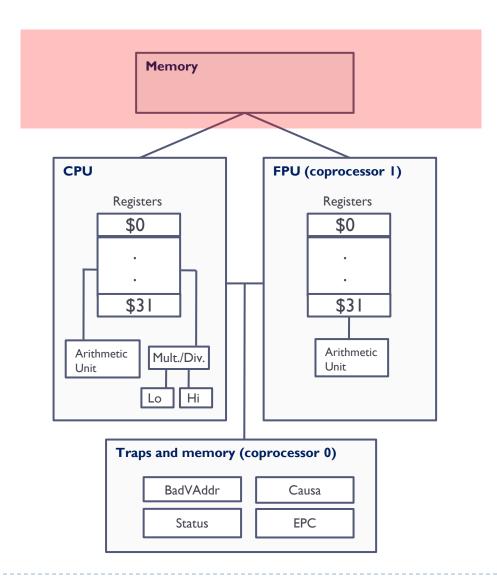
#### ▶ There are 32 registers

- Size: 4 bytes (I word)
- Used a \$ at the beginning

#### Use convention

- Reserved
- Arguments
- Results
- Temporary
- Pointers

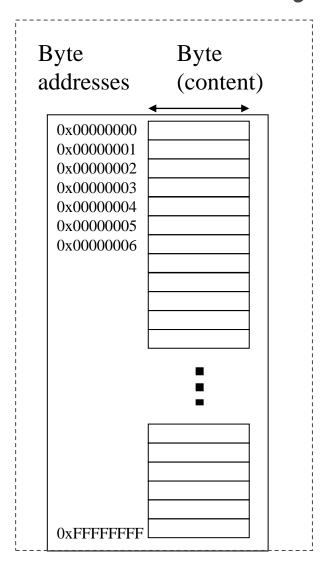
#### MIPS-32 Architecture



#### Main memory

- ▶ 32-bit memory addresses
- ▶ 4 GB addressable memory

#### MIPS32 memory model



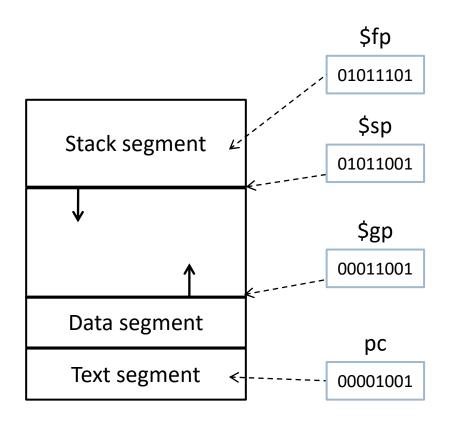
#### Memory is addressed at byte level:

- 32-bit addresses
- Content of each address: one byte
- Addressable space:  $2^{32}$  bytes = 4 GB

#### Access can be to:

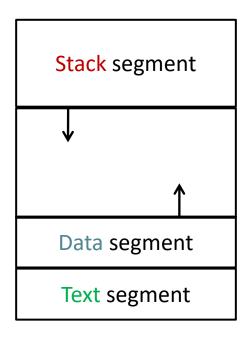
- Individual bytes
- Words (4 consecutive bytes)

## Memory layout for a process



- The memory space is divided in logic segments in order to organize the content:
  - Stack segment
    - Local variables
    - Function contexts
  - Data segments
    - Static data
  - Code segment (text)
    - Program code

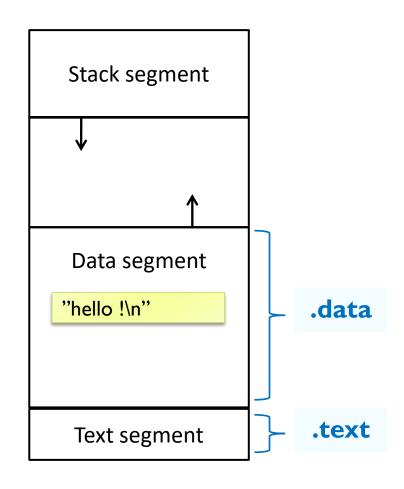
## Storing variables in memory



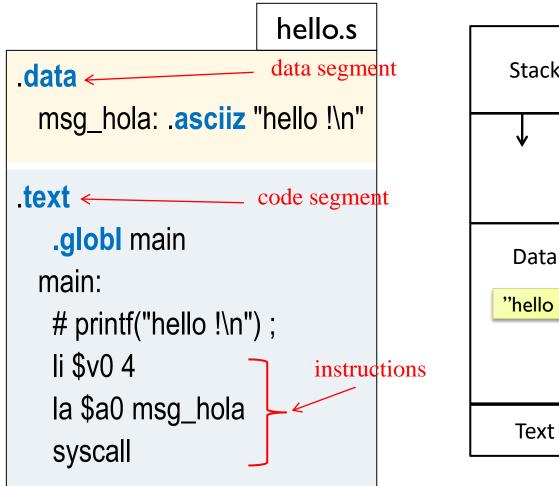
```
// global variables
int a;
main ()
   // local variables
   int b;
   // code (text)
   return a + b;
```

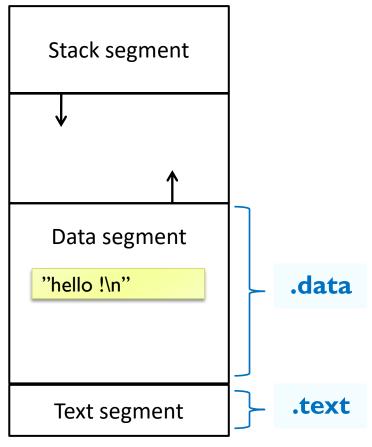
## Example: hello world...

```
hello.s
.data
 msg_hola: .asciiz "hello !\n"
text
   .globl main
 main:
  # printf("hello !\n");
   li $v0 4
   la $a0 msg_hola
   syscall
```



## Example: hello world...

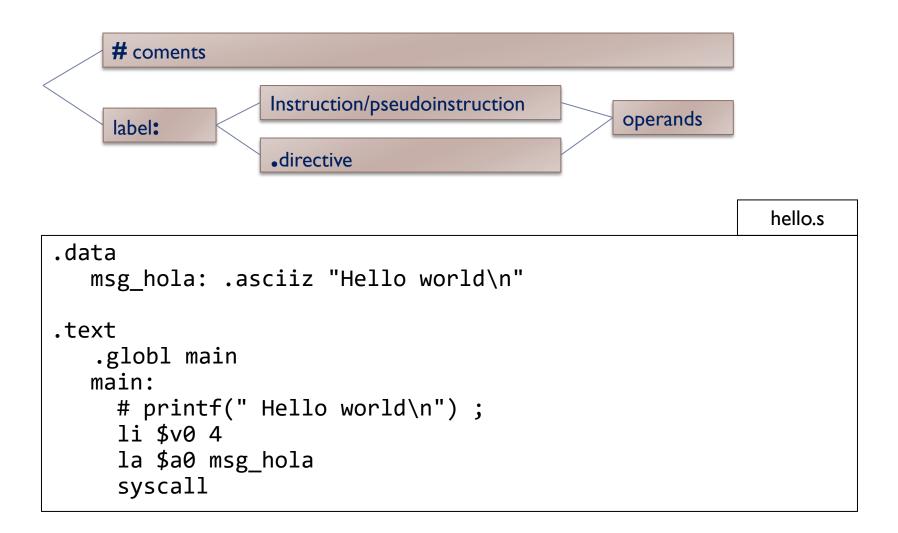




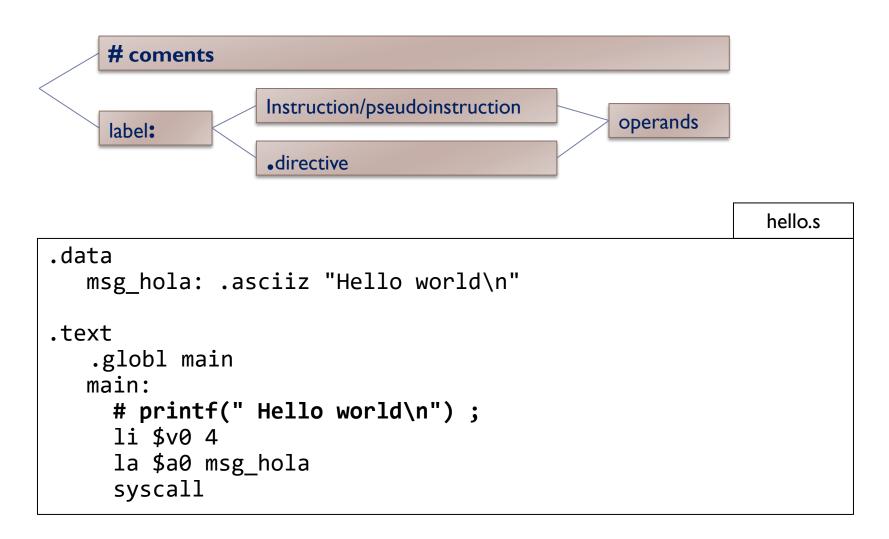
#### Contents

- Basic concepts on assembly programming
- MIPS32 assembly language, memory model and data representation
  - ▶ MIPS architecture (II): registers and memory
  - Assembler directives
  - System services
  - Instructions (II): memory access
- Instruction formats and addressing modes
- Procedure calls and stack convention

## Example: hello world...



# Assembler program: line comment with #



## Assembler program: labels:

```
# coments
                    Instruction/pseudoinstruction
                                                      operands
     label:
                     directive
                                                                     hello.s
.data
   msg_hola: ..asciiz "Hello world\n"
                                    msg_hola: represents the memory address
.text
                                    where the string (.asciiz) starts to be stored.
    .globl main
   main: <
     # printf(" Hello world\n") ;
      li $v0 4
                                    main: represent the memory address
      la $a0 msg hola
                                    where first instruction of main is stored
      syscall
```

# Assembler program: assembler directives (preprocessing)

```
# coments
                   Instruction/pseudoinstruction
                                                   operands
     label:
                   directive
                                                                hellos
.data
   msg hola: .asciiz "Hello world\n"
.text
   .globl main
   main:
     # printf(" Hello world\n") ;
     li $v0 4
     la $a0 msg_hola
     syscall
```

## Assembly: directives

Directives	Description	
.data	Next elements will go to the data segment	
.text	Next elements will go to the code segment	
.ascii "string value"	String definition without '\0' ending terminator	
.asciiz "string value"	String definition with '\0' ending terminator ('\0' = 0)	
.byte 1, 2, 3	Bytes stored in memory consecutively	
.half 300, 301, 302	Half-words stored in memory consecutively	
.word 800000, 800001	Words stored in memory consecutively	
.float 1.23, 2.13	Floats stored in memory consecutively	
.double 3.0e21	Doubles stored in memory consecutively	
.space 10	Allocates a space of 10 bytes in the current segment	
.extern label n	Declare that label is global of size n	
.globl label	Declare label as global	
.align n	Align next element to a address multiple of 2 <sup>n</sup>	

## Representation of basic data types (1/3)

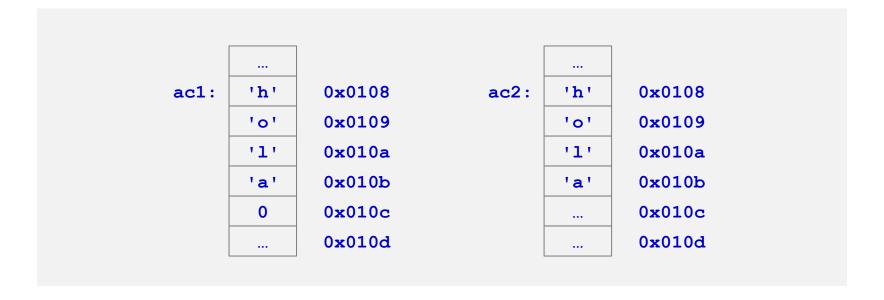
```
// boolean
bool t b1;
bool t b2 = false ;
// caracter
char c1;
char c2 = x';
// enteros
int res1;
int op1 = -10;
// coma flotante
float f0;
float f1 = 1.2;
double d2 = 3.0e10;
```

```
.data
# boolean
b1: .space 1
b2: .byte 0
               # 1 byte
# caracter
c1: .byte
c2: .byte 'x'
                     # 1 bytes
# enteros
res1: .space 4
op1: .word -10
                  # 4 bytes
# coma flotante
f0: .float
f1: .float 1.2 # 4 bytes
d2: .double 3.0e10 # 8 bytes
```

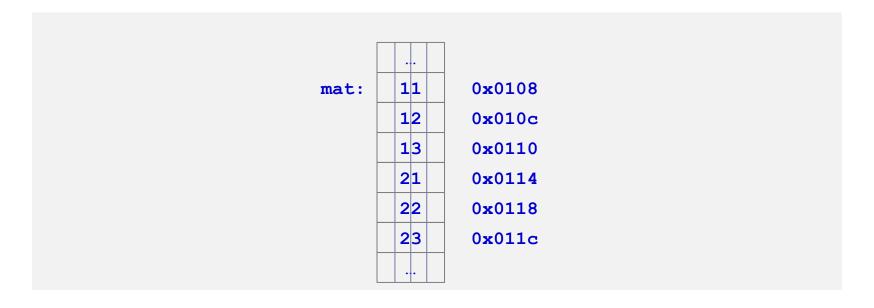
#### Representation of basic data types (2/3)

```
// strings
char c1[10] ;
char ac1[] = "hola" ;
```

```
# strings
c1: .space 10  # 10 byte
ac1: .asciiz "hola" # 5 bytes (!)
ac2: .ascii "hola" # 4 bytes
```



## Representation of basic data types (3/3)



#### Contents

- Basic concepts on assembly programming
- MIPS32 assembly language, memory model and data representation
  - ▶ MIPS architecture (II): registers and memory
  - Assembler directives
  - System services
  - Instructions (II): memory access
- Instruction formats and addressing modes
- Procedure calls and stack convention

## Example: hello world...

hello.s

```
.data
  msg_hola: .asciiz "hello world\n"
.text
   .glob1 main
  main:
    # printf("hello world\n") ;
    li $v0 4
             # syscall code: 4
    la $a0 msg_hola # address where msg_hola starts
    syscall
```

## System calls

- Many assembler simulators include a small "operating system"
  - The SPIM simulator provides 17 services.
- How to invoke:
  - Call code in register \$v0
  - Other arguments on specific records
  - Invocation by the syscall instruction

## System calls

Service	Call code (\$v0)	Arguments	Result
print_int	I	\$a0 = integer	
print_float	2	\$f12 = float	
print_double	3	\$f12 = double	
print_string	4	\$a0 = string	
read_int	5		integer in \$v0
read_float	6		float in \$f0
read_double	7		double in \$f0
read_string	8	\$a0 = buffer, \$a1 = length	
sbrk	9	\$a0 = amount	address in \$v0
exit	10		

## System calls

Service	Call code (\$v0)	Arguments	Result
print_char	П	\$a0 (ASCII code)	
read_char	12		\$v0 (ASCII code)
open	13	Equivalent to \$v0=open(\$a0, \$a1, \$a2)	file descriptor in \$v0
read	14	Equivalent to \$v0=read(\$a0, \$a1, \$a2)	read bytes in \$v0
write	15	Equivalent to \$v0=write(\$a0, \$a1, \$a2)	written bytes in \$v0
close	16	Equivalent to \$v0=close(\$a0)	0 in \$v0
exit2	17	End the program. Return the value stored in \$a0	

## Example: Hello world...

hello.s

#### .data

msg\_hola: .asciiz "hello world\n"

#### .text

.glob1 main

main:

Call code (\$v0)	Arguments
I	\$a0 = integer
2	\$f12 = float
3	\$f12 = double
4	\$a0 = string
	1

```
# printf("hello world\n");
```

li \$v0 4 # syscall code: 4

la \$a0 msg\_hola # address where msg\_hola starts

syscall ← Operating system invocation instruction

#### Exercise

```
int valor ;

int valor ;

readInt(&valor) ;

valor = valor + 1 ;
printInt(valor) ;

. . . .
```

service	code	arguments	results
print_int	ı	\$a0 = integer	
print_float	2	\$f12 = float	
print_double	3	\$f12 = double	
print_string	4	\$a0 = string	
read_int	5		integer in \$v0
read_float	6		float in \$f0
read_double	7		double in \$f0
read_string	8	\$a0=buffer, \$a I =long.	
sbrk	9	\$a0=amount	address in \$v0
exit	10		

## Exercise (solution)

```
int valor ;
int valor ;

readInt(&valor) ;
valor = valor + 1 ;
printInt(valor) ;

. . .
```

service	code	arguments	results
print_int	I	\$a0 = integer	
read_int	5		integer en <b>\$v0</b>

```
# readInt(&valor)
li $v0 5
syscall
sw $v0 valor
# valor = valor + 1
add $a0 $v0 1
sw $a0 valor
# printInt
li $v0 1
syscall
```

#### Contents

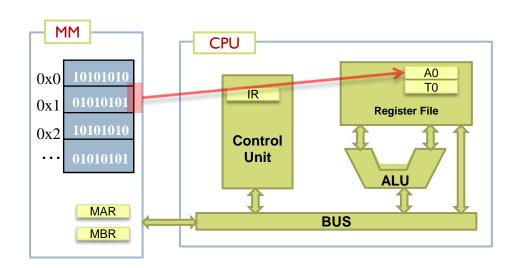
- Basic concepts on assembly programming
- MIPS32 assembly language, memory model and data representation
  - ▶ MIPS architecture (II): registers and memory
  - Assembler directives
  - System services
  - Instructions (II): memory access
- Instruction formats and addressing modes
- Procedure calls and stack convention

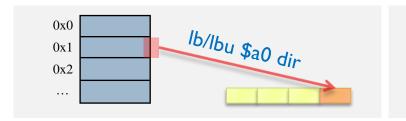
# Data transfer **bytes**

▶ Copies a byte from memory to a register or vice versa.

#### Examples:

- Memory to registerlb \$a0, dirlbu \$a0, dir
- Register to memorysb \$t0, dir



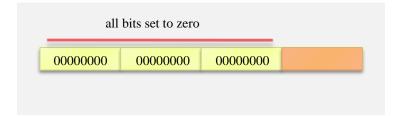




# Data transfer bytes, sign extension

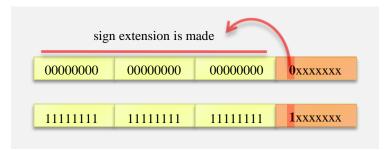
- There are two possibilities when transferring a byte from memory to register:
- A) Transfer without sign, for example: Ibu \$a0, dir





B) Transfer with sign, for example: Ib \$a0, dir





## Basic data types **booleans**

```
bool_t b1 = false;
bool_t b2 = true;
...
main ()
{
   b1 = true ;
...
}
```

```
.data
b1: .byte 0 # 1 byte
b2: .byte 1
. . .
.text
.globl main
 main: la $t0 b1
        li $t1 1
        sb $t1 ($t0)
```

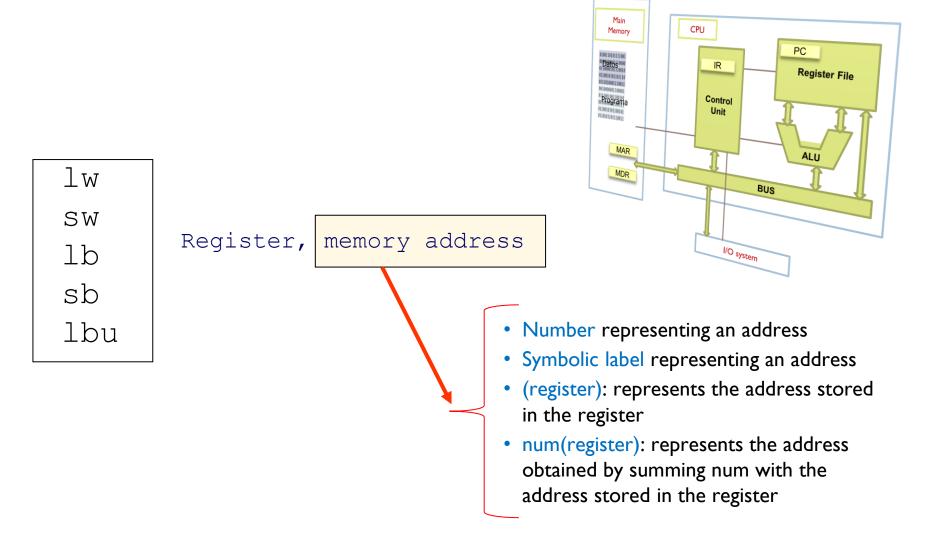
## Basic data types characters

```
char c1 ;
char c2 = 'a';
...

main ()
{
    c1 = c2;
...
}
```

```
.data
c1: .space 1 # 1 byte
c2: .byte 'a'
. . .
.text
.globl main
 main: la $t0 c1
        1bu $t1 c2
        sb $t1 ($t0)
```

#### Format of memory access instructions



### Memory access instruction formats

- ▶ lbu \$t0, 0x0F000002
- ▶ 1bu \$t0, labeled
  - Direct addressing. The byte stored in the memory location labeled is loaded into \$t0.
- ▶ lbu \$t0, (\$t|)
  - Indirect register addressing. The byte stored in the memory location stored in \$t1 is loaded in \$t0.
- ▶ lbu \$t0, 80(\$t|)
  - Relative addressing. The byte stored in the memory location obtained by adding the contents of \$t1 with 80 is loaded in \$t0.

### Instructions and pseudo-instructions

- ▶ There is an assembly instruction per machine instruction :
  - ▶ Each machine instruction occupies 32 bits in MIPS32
  - addi \$t1,\$t0,4
- A pseudo-instruction can be used in an assembler program and it corresponds to one or several assembly instructions:
  - E.g.: li \$v0, 4 move \$t1,\$t0
- In the assembly process, they are replaced by the sequence of assembly instructions that perform the same functionality.
  - E.g.: ori \$v0, \$0, 4 replaces to: li \$v0, 4 addu \$t1, \$0, \$t2 replaces to: move \$t1, \$t2

#### Other examples of pseudo-instructions

- An assembler pseudoinstruction can correspond to several machine instructions.
  - ▶ li \$t1,0×00800010
    - □ It does not fit in 32 bits but can be used as a pseudo-instruction.
    - ☐ It is equivalent to:

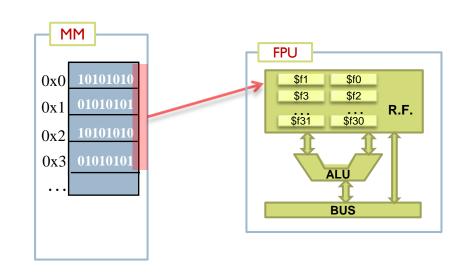
```
lui $t1,0x0080
ori $t1,$t1,0x0010
```

# Data transfer IEEE 754

▶ Copies a number from memory to a register or vice versa.

#### Examples:

- Memory to registerl.s \$f0 dir ll.d \$f2 dir2
- Register to memorys.s \$f0 dir ls.d \$f0 dir 2







### Basic data types float

```
float result;
float op1 = 100;
float op2 = 2.5
main ()
  result = op1 + op2;
```

```
.data
.align 2
   result: .word 0 # 4 bytes
   op1: .float 100
   op2:
           .float 2.5
.text
      .globl main
main: 1.s $f0 op1
       1.s $f1 op2
       add.s $f3 $f1 $f2
       s.s $f3 result
```

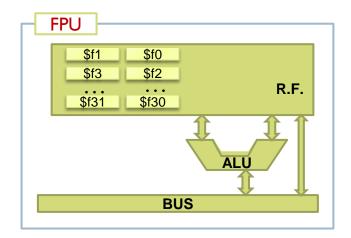
### Basic data types double

```
double result ;
double op1 = 100;
double op2 = -10.27;
main ()
  result = op1 * op2;
```

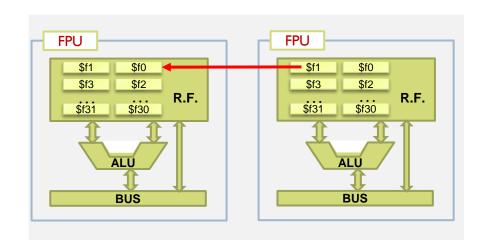
```
.data
.align 3
   result: .space 8
   op1: .double 100
   op2:
           .double -10.27
.text
      .globl main
main: 1.d $f0 op1 # ($f0,$f1)
      1.d $f2 op2 # ($f2,$f3)
      mul.d $f6 $f0 $f2
      s.d $f6 result
```

### Floating point. IEEE 754

- ▶ The coprocessor I has 32 registers of 32 bits (4 bytes) each.
  - It is possible to work with single or double precision
- Simple precision (32 bits):
  - From \$f0 to \$f3 I
  - E.g.: add.s \$f0 \$f1 \$f5 f0 = f1 + f5
  - Other operations:
    - add.s, sub.s, mul.s, div.s, abs.s
- Doble precision (64 bits):
  - Registers used in pairs
  - E.g.: add.d f0 f2 f8(f0,f1) = (f2,f3) + (f8,f9)
  - Other operations:
    - add.d, sub.d, mul.d, div.d, abs.d

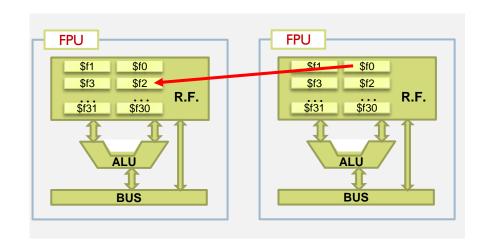


### Operations with registers (FPU, FPU)



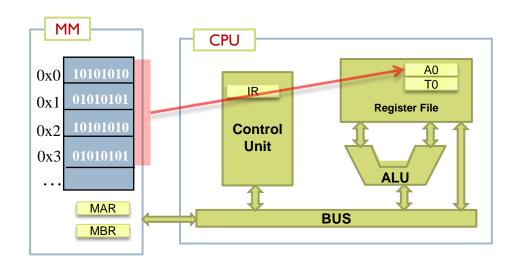
mov.d \$f0 \$f2

▶ (\$f0, \$f1) ← (\$f2, \$f3)



# Data transfer words

- ▶ Copies a word from memory to a register or vice versa.
- Examples:
  - Memory to registerlw \$a0 (\$t0)
  - Register to memorysw \$a0 (\$t0)







# Basic data types integers

```
int result;
int op1 = 100;
int op2 = -10;
main ()
  result = op1+op2;
```

```
.data
.align 2
result: .word 0 # 4 bytes
op1: .word 100
op2: .word -10
.text
.globl main
main: lw $t1 op1
       lw $t2 op2
       add $t3 $t1 $t2
       la $t4 result
       sw $t3 ($t4)
```

### Basic data types

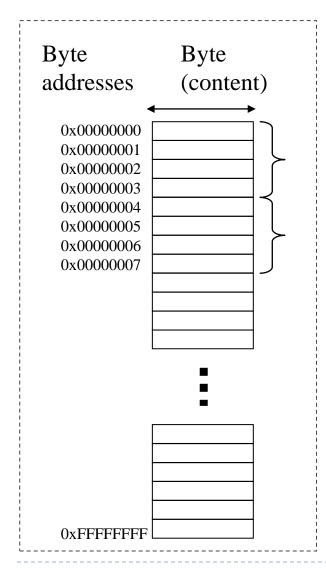
integers

global variable without initial value

```
int result;
int op1 = 100;
int op2 = -10;
     global variable with initial value
main ()
  result = op1+op2;
```

```
.data
.align 2
result: .word 0 # 4 bytes
        .word 100
op1:
         ?word -10
op2:
.text
.globl main
main: lw $t1 op1
        lw $t2 op2
        add $t3 $t1 $t2
        la $t4 result
        sw $t3 ($t4)
```

#### Accessing to words



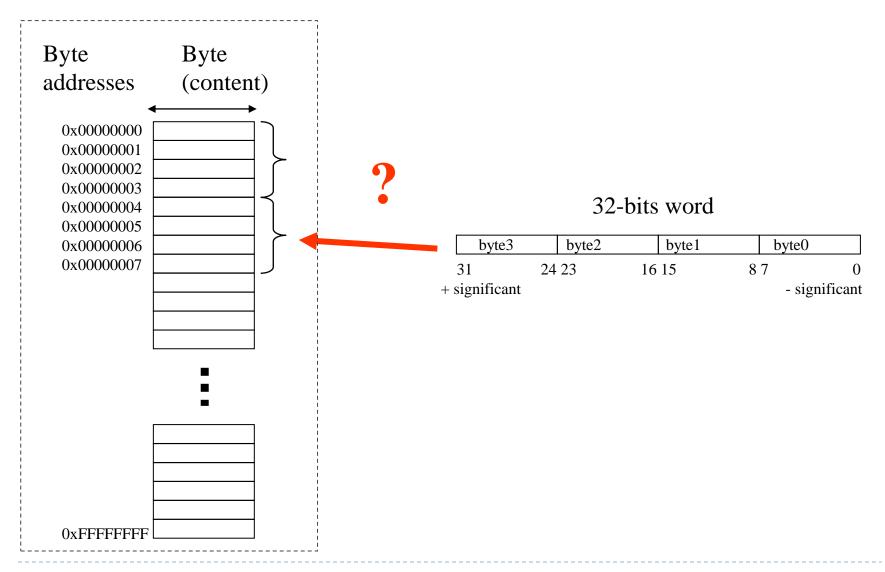
A word: 4 bytes in a 32-bits processor

Word stored starting at byte 0

Word stored starting at byte 4

Words (32 bits, 4 bytes) are stored using 4 consecutive memory locations, starting with the first position at an address multiple of 4

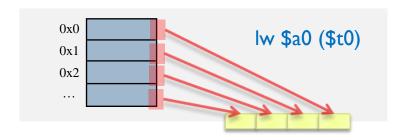
### Accessing to words



#### Data transfer byte order

▶ There are 2 types of byte order:

Little-endian ( 'small' address ends the word...)



**AMD** 

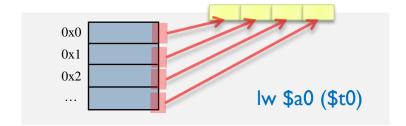


Big-endian

('big' address ends the word...)



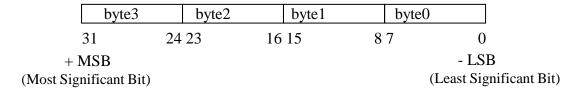
(bi-endian)

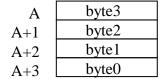




### Storing words in memory

32-bit word

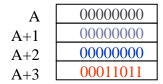




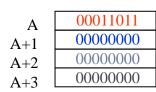
BigEndian

A	byte0
A+1	byte1
A+2	byte2
A+3	byte3

LittleEndian



BigEndian

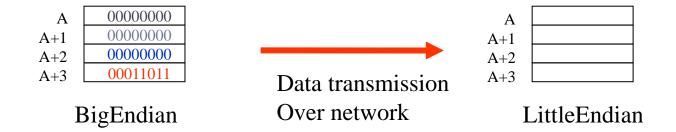


LittleEndian

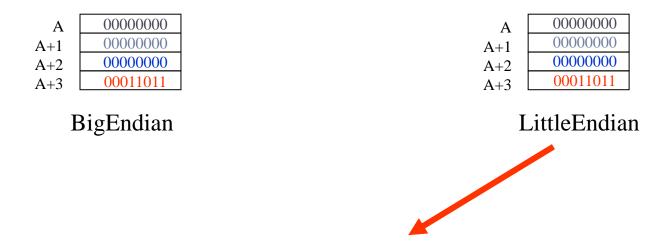
# Communication problems in computers with different architectures



# Communication problems in computers with different architectures



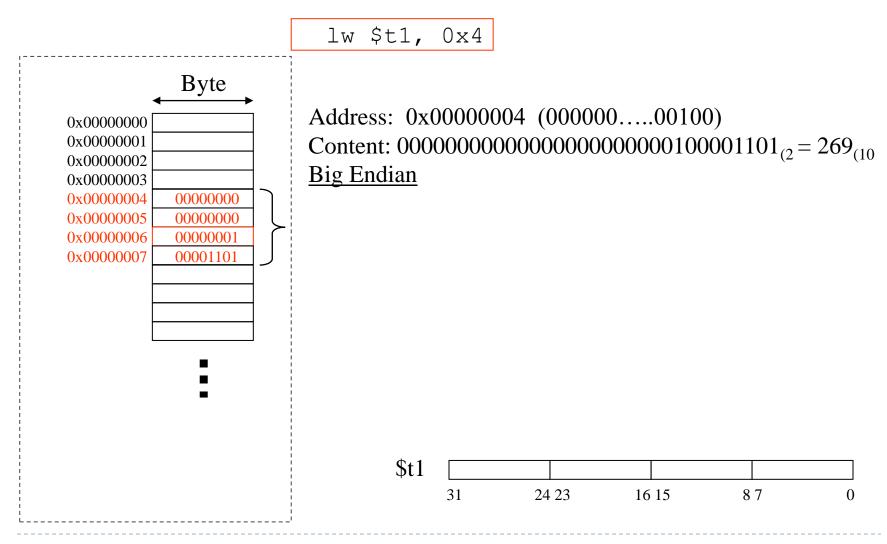
# Communication problems in computers with different architectures

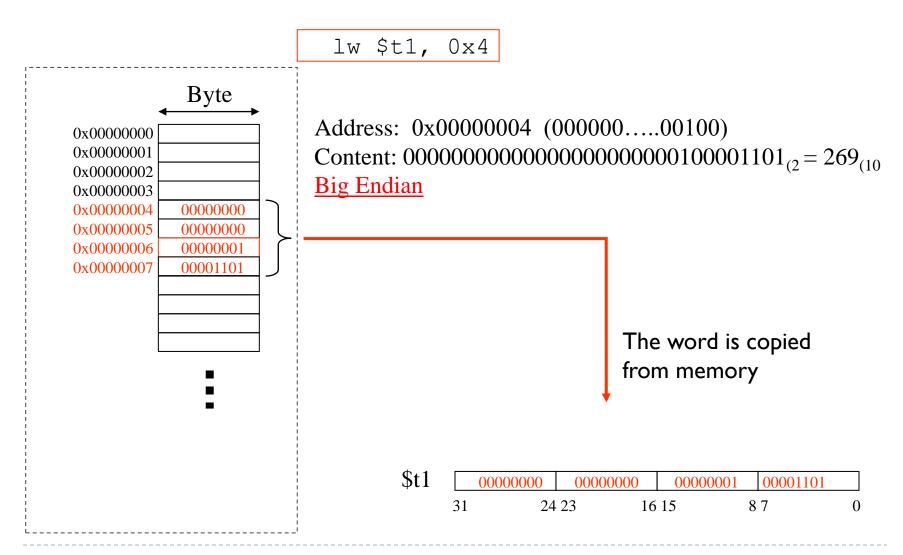


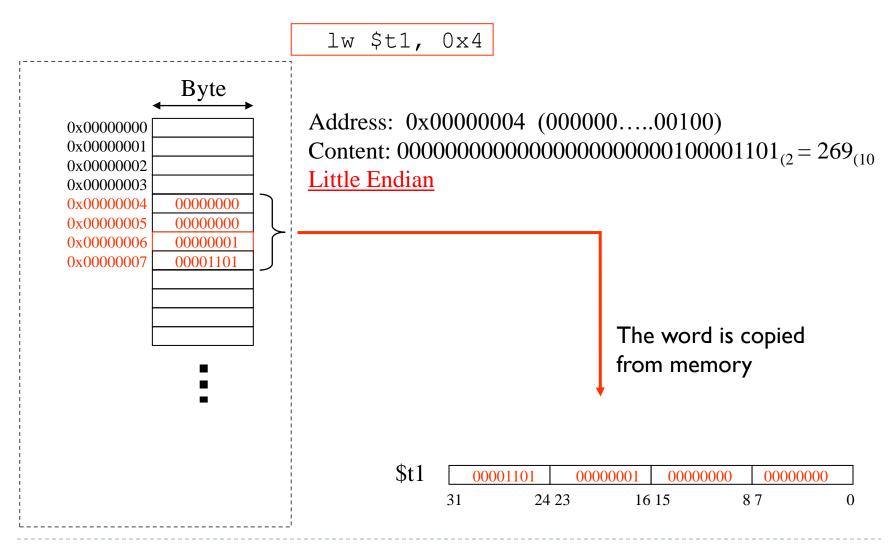
### Example

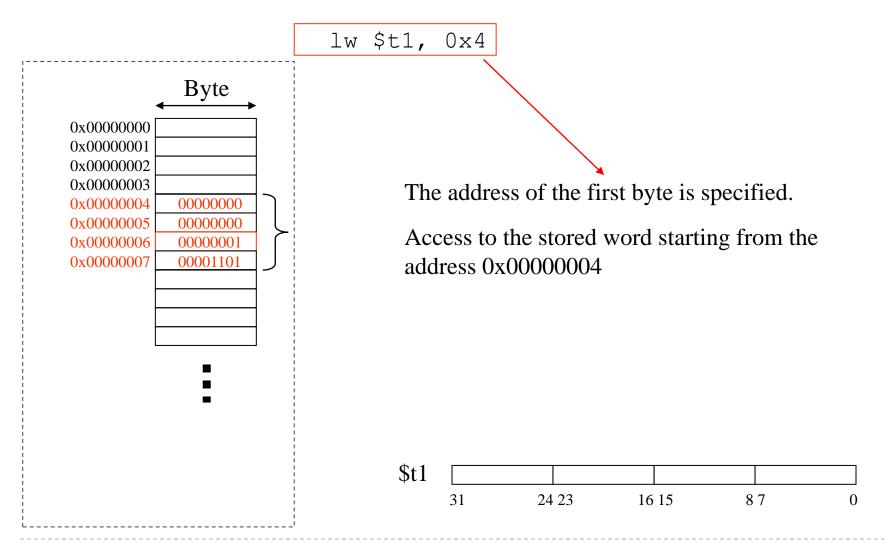
endian.s

```
.data
 b1: .byte 0x00, 0x11, 0x22, 0x33
.text
.globl main
main:
   lw $t0 b1
```





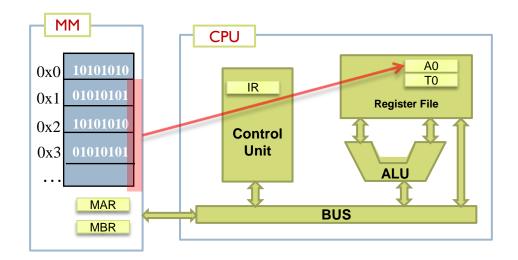




# Data transfer alignment and access size

#### Peculiarities:

- Alignment of elements in memory
- Default access size



### Data alignment

#### In general:

A data of K bytes is aligned when the address D used to access this data fulfills the condition:

 $D \mod K = 0$ 

#### Data alignment implies:

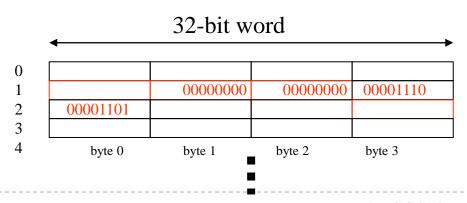
- Data of 2 bytes are stored in even addresses
- Data of 4 bytes are stored in addresses multiple of 4
- Data of 8 bytes (double) are stored in addresses multiple of 8

### Data alignment

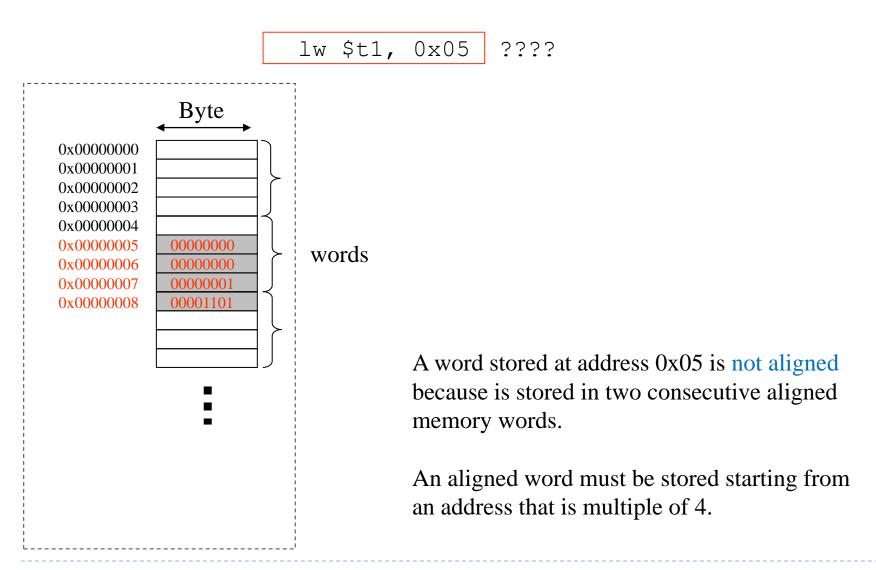
.data ▶ The alignment requires the address to be a .byte 0x0f multiple of the word size: .word 10 31 23 15. 0 This word is aligned, next ones are not 08 address 12 16 20 24

### Data alignment

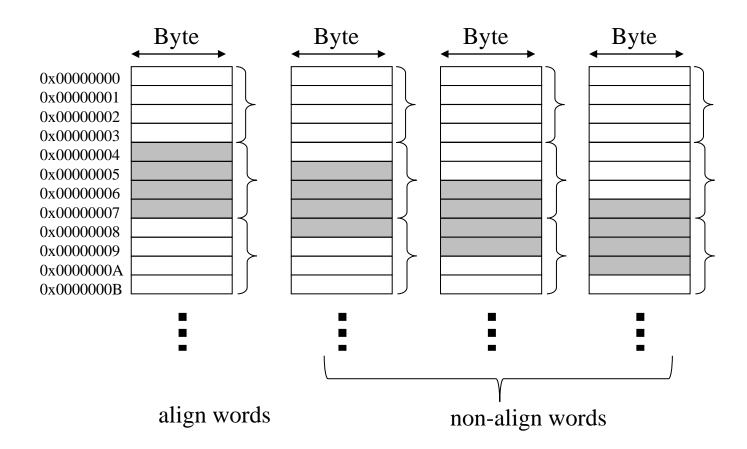
- Many computers does not allow the access to not aligned data:
  - Goal: reduce the number of memory accesses
  - Compilers assign addresses aligned to variables
- Some processors, such as Intel models, allow the access to not aligned data:
  - Non-aligned data needs several memory access



### Non-aligned data

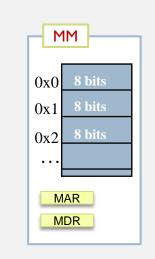


### Non-aligned data



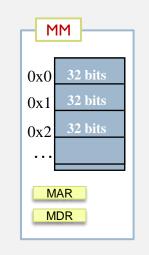
### Word-level or byte-level addressing

- The main memory is similar to a large one-dimensional vector of items.
- ▶ A memory address is the index of one item in the vector.
- There are two types of addressing:
  - Byte addressing



- Each memory element is 1 byte
- Transferring a word means transferring 4 bytes (in a 32-bit CPU)

Word addressing



- Each memory element is a word
- b means transferring one word and keeping one byte.

### Summary

- The instructions and data of a program must be loaded in memory for the execution (process)
- All data and instructions are stored in memory so all have an associated memory address where is stored
- ▶ In a 32-bit computer such as MIPS 32:
  - Registers have 32 bits
  - Memory can store bytes (8 bits)
    - ▶ Instructions: memory → register: lb, lbu
    - ▶ Instructions: register → memory: sb
  - Memory can store words (32 bits)
    - ▶ Instructions: memory → register: 1w
    - ▶ Instructions: register → memory: SW

# Assembly compound data types: vector, matrix and structs

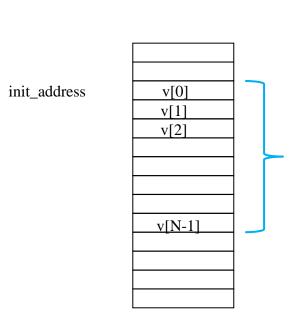
- Vector:
   consecutives elements of the same type indexed by position
  - String
- Matrix: two dimensional indexed elements of same type
- Structs:
   consecutive elements of same/different types indexed by name

# Compound data types Arrays

- Collection of data ítems stored consecutively in memory
- The address of the j element can be computed as:

init\_address + j \* p

Where **p** is the size of each item



# Compound data types Arrays

```
int vec[5] ;
...
main ()
{
    vec[4] = 8;
}
```

```
.data
.align 2 # next item aligned to 4
vec: .space 20 # 5 items * 4 bytes/item
.text
.globl main
main:
       la $t1 vec
       li $t2 8
       sw $t2 16($t1)
```

# Compound data types Arrays

```
int vec[5] ;
...
main ()
{
    vec[4] = 8;
}
```

```
.data
.align 2 # next item aligned to 4
vec: .space 20 # 5 items * 4 bytes/item
.text
.globl main
main:
       li $t0 16
       la $t1 vec
       add $t3, $t1, $t0
       li $t2 8
       sw $t2, ($t3)
```

# Compound data types Arrays

```
int vec[5] ;
...
main ()
{
    vec[4] = 8;
}
```

```
.data
.align 2 # next item align to 4
vec: .space 20 # 5 items * 4 bytes/item
.text
main:
       li $t2 8
       li $t1 16
        sw $t2 vec($t1)
        . . .
```

### Exercise

- Let V be an array of integer elements
  - V represents the initial address of the array
- ▶ What is the address of the V[5] item?
- Which are the instruction to load in register \$t0 the value of v[5]?

- Let V be an array of integer elements
  - V represents the initial address of the array
- ▶ What is the address of the V[5] item?
  - V + 5\*4
- Which are the instruction to load in register \$t0 the value of v[5]?
  - ▶ li \$t1, 20
  - h lw \$t0, v(\$t1)

# Compound data types **String**

```
    Array of bytes
```

• '\0' ends string

```
char c1 ;
char c2 = 'h' ;
char *ac1 = "hola" ;
...

main ()
{
    printf("%s",ac1) ;
...
}
```

```
.data
c1: .space 1
                       # 1 byte
c2: .byte 'h'
ac1: .asciiz "hola"
.text
.globl main
main:
        li $v0 4
         la $a0 ac1
         syscall
```

### String layout in memory

```
// strings
char c1[10] ;
char ac1[] = "hola" ;
```

```
# strings
c1: .space 10  # 10 byte
ac1: .asciiz "hola" # 5 bytes (!)
ac2: .ascii "hola" # 4 bytes
```

```
ac1:
           'h'
                      0 \times 0108
                                             ac2:
                                                         \h'
                                                                    0 \times 0108
           10'
                      0 \times 0109
                                                         10'
                                                                    0 \times 0109
           111
                      0x010a
                                                         11'
                                                                    0x010a
           \a'
                      0 \times 010b
                                                         'a'
                                                                    0 \times 010 b
            0
                      0 \times 010c
                                                                    0x010c
                      0x010d
                                                                    0x010d
```

# Compound data types String length

```
char c1;
char c2 = 'h';
char *ac1 = "hola" ;
char *c;
main ()
  c = ac1; int 1 = 0;
  while (c[1] != NULL) {
        1++;
  printf("%d", 1);
```

# Compound data types String length

```
char c1;
char c2 = 'h';
char *ac1 = "hola" ;
char *c;
main ()
  c = ac1; int 1 = 0;
  while (c[1] != NULL) {
        1++;
  printf("%d", 1);
```

```
.data
c1: .space 1 # 1 byte
c2: .byte 'h'
ac1: .asciiz "hola"
.align 2
c: .word 0 # pointer => address
.text
.globl main
main: la $t0, ac1
         li $a0, 0
         lbu $t1, ($t0)
        begz $t1, fin
  buc:
         addi $t0, $t0, 1
         addi $a0, $a0, 1
         lbu $t1, ($t0)
         b buc
  fin: li $v0 1
         syscall
```

## Arrays and strings

### Review (in general) :

▶ sw \$t0, 4(\$s3) # M[\$s3+4] \$t0

#### Exercise

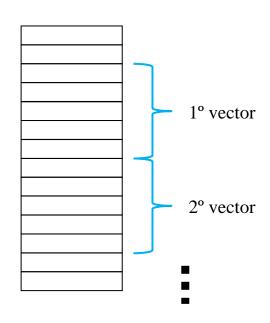
- Write a program that:
  - ▶ Calculate the number of occurrences of a char in a string
    - String address stored in \$a0
    - Char to look for in \$a1
    - Result must be stored in \$v0

## Compound data types Matrix

- A matrix m x n consists of m vectors (m rows) of length n
- Usually stored by rows
- ▶ The element  $a_{ij}$  is stored in the address:

$$init_address + (i \cdot n + j) \times p$$

where p is the size of each item



## Compound data types Matrix

```
.data
.align 2  # next item align to 4
Vec: .space 20 # 5 item * 4 bytes/item
mat: .word 11, 12, 13
     .word 21, 22, 23
.text
.globl main
 main:
        lw $t1 mat+0
         lw $t2 mat+12
         add $t3 $t1 $t2
         sw $t3 mat+4
```

## Example: integer matrix

```
int vec[5];
int mat[2][3] = \{\{11,12,13\},
                  {21,22,23}};
main ()
 mat[1][2] = mat[1][1] +
  mat[2][1];
```

```
.data
align 2
Vec: .space 20 # 5 items*4 bytes/item
mat: .word 11, 12, 13
     .word 21, 22, 23
. . .
.text
.qlobl main
 main:
         lw $t1 mat+0
         lw $t2 mat+12
         add $t3 $t1 $t2
         sw $t3 mat+4
```

### Exercise

```
// global variables
char v1;
int v2 ;
float v3 = 3.14 ;
char v4 = "ec" ;
int v5[] = { 20, 22 } ;
```

```
// global variables
char v1;
int v2 ;
float v3 = 3.14 ;
char v4 = "ec" ;
int v5[] = { 20, 22 } ;
```

```
.data
v1: .byte 0
.align 2
v2: .space 4
v3: .float 3.14
v4: .asciiz "ec"
.align 2
v5: .word 20, 22
```

**v1**: ? ? ? 0x0100 0x0101 0x0102 0x0103

```
.data
v1: .byte 0
.align 2
v2: .space 4
v3: .float 3.14
v4: .asciiz "ec"
.align 2
v5: .word 20, 22
```

		1
<b>v</b> 1:	0	0x0100
	?	0x0101
	?	0x0102
	?	0x0103
<b>v</b> 2:	0	0x0104
	0	0x0105
	0	0x0106
	0	0x0107
<b>v</b> 3:	(3.14)	0x0108
	(3.14)	0x0109
	(3.14)	0x010A
	(3.14)	0x010B
<b>v4</b> :	\e'	0x010C
	\c'	0x010D
	0	0x010E
		0x010F
<b>v</b> 5:	(20)	0x0110
	(20)	0x0111
	(20)	0x0112
	(20)	

```
.data
v1: .byte 0
.align 2
v2: .space 4
v3: .float 3.14
v4: .asciiz "ec"
.align 2
v5: .word 20, 22
```

## Tips

- Do not program directly in assembler
  - Better to first do the design in DFD, Java/C/Pascal...
  - Gradually translate the design to assembler.
- Sufficiently comment the code and data
  - By line or by group of lines comment which part of the design implements.
- Test with enough test cases
  - Test that the final program works properly to the given specifications.

### Exercise

- Write an assembly program that:
  - ▶ Load the value -3.141516 in register \$f0
  - Obtain the exponent and mantissa values stored in the register
     \$f0 (IEEE 754 format)
    - Display the sign
    - Display the exponent
    - Display the mantissa

```
.data
  newline: .asciiz "\n"
.text
.qlobl main
main:
         li.s $f0, -3.141516
         # print value
         mov.s $f12, $f0
         li $v0, 2
         syscall
         la $a0, newline
         li $v0, 4
         syscall
         # copy to processor
         mfc1 $t0, $f12
```

```
li $s0, 0x80000000
                      #sign
and $a0, $t0, $s0
srl $a0, $a0, 31
li $v0, 1
syscall
la $a0, newline
li $v0, 4
syscall
li $s0, 0x7F800000
                      #exponent
and $a0, $t0, $s0
srl $a0, $a0, 23
li $v0, 1
syscall
la $a0, newline
li $v0, 4
syscall
                      #mantissa
li $s0, 0x007FFFFF
and $a0, $t0, $s0
li $v0, 1
syscall
jr $ra
```