

Week 1

Thursday, June 18, 2020 3:57 PM

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Knowledge-based Agents :

Agents that reason by operating on internal representations of knowledge

Sentence :

An assertion about the world in a knowledge representation language

Model :

Assignment of a truth value to every propositional symbol (a "possible world")

Knowledge base :

A set of sentences known by a knowledge-based agent

Entailment ($a \models b$) :

In every model a is true b is also true

Inference :

The process of deriving a new sentence from old ones

Usage of inference : Does $A \models B$?

Example for inference algorithm : Model Checking

Propositional logic

Propositional symbols :

Representation of a fact about the world

Logical connectives :

Not \neg

And \wedge

Or \vee

implication \rightarrow

biconditional \leftrightarrow

Truth table for :

Implication :

P	Q	$P \rightarrow Q$
true	true	true
true	false	false
false	true	true
false	false	true

Biconditional :

P	Q	$P \leftrightarrow Q$
true	true	true
true	false	false
false	false	true
false	true	false

Knowledge Engineering

To solve a problem using propositional logic :

- 1- start by asking what propositional symbols you are going to need
- 2- using the symbols create logical sentences (knowledge)
- 3- using the sentences create the KB

Model Checking

To determine if $KB \models a$:

- Enumerate all possible models
- if in every model where KB is true , a is true
Then KB entail a
- Otherwise KB does not entail a

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Inference Rules

:: Set of rulers that take knowledge that already exists and translate it to a new form of knowledge

Modus Ponens :

- o If we have :
 - $A \rightarrow B$
 - A
- o Then we have :
 - B

And Elimination :

- o If we have :
 - $A \wedge B$
- o Then we have :
 - A
 - B

Double Negation Elimination :

- o If we have :
 - $\neg(\neg A)$
- o Then we have :
 - A

Implication Elimination :

- o If we have :
 - $A \rightarrow B$
- o Then we have :
 - $\neg A \vee B$

Biconditional Elimination :

- o If we have :
 - $A \leftrightarrow B$
- o Then we have :
 - $(A \rightarrow B) \wedge (B \rightarrow A)$

De Morgan's law :

- o If we have :
 - $\neg(A \wedge B)$
- o Then we have :
 - $\neg(A) \vee \neg(B)$
- o The reverse is also right **
 - If : $\neg(A \vee B)$ then we have $\neg(A) \wedge \neg(B)$

Distributive Property :

- o $A \wedge (B \vee C)$ give us : $(A \wedge B) \vee (A \wedge C)$
- o $A \vee (B \wedge C)$ give us : $(A \vee B) \wedge (A \vee C)$

Theorem Proving

Initial State : starting Knowledge base

Actions : inference rules

Transition model : New KB after inference

Goal test : check statement we are trying to prove

Path cost function : number of steps in proof

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Inference by resolution

To determine if $KB \models A$:

Check if $(KB \wedge \neg(A))$ is a contradiction?

if so, then $KB \models A$

Otherwise, no entailment

And we do the check by :

Convert $(KB \wedge \neg(A))$ to conjunction Normal form(CNF)

Keep checking if we can use resolution to produce new clause

if we ever produce the empty clause (\square)

Then we have a contradiction and

$KB \models A$

Otherwise, if we can not add new clause

Then : no entailment

Clause :

A disjunction of literals
 $P \vee Q \vee R$

Conjunction Normal form (CNF):

Logical sentence that is a conjunction of clauses
 $(A \vee B \vee C) \wedge (D \vee E \vee F) \wedge P$

Conversion to CNF

- Eliminate biconditionals
 - o Turn $(A \leftrightarrow B)$ into $(A \rightarrow B) \wedge (B \rightarrow A)$
- Eliminate implications
 - o Turn $(A \rightarrow B)$ into $\text{not}(A) \vee B$
- Move not inward using De Morgan's Law
 - o Turn $\text{not}(A \wedge B)$ into $\text{not}(A) \vee \text{not}(B)$
- Use distribution law to distribute \vee whenever possible
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Resolution rule

If we have :
 $P \vee Q$
 $\text{Not}(p)$
Then we have :
 Q

If we have :
 $P \vee Q$
 $\text{Not}(p) \vee R$
Then we have :
 $Q \vee R$

If we have :
 A
 $\text{Not}(A)$
Then we have :
 $()[\text{false}]$

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First-Order Logic

We have :

Constant Symbols : represent objects

Predicate Symbols : relations or functions (properties that might hold true or false)

Quantifiers

We have :

Universal Quantification : something is going to be true for all value of a variable

Existential Quantification : something going to be true for some value of a variable (at least one)