

PARALLELISM

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July 6, 2023



OUTLINE

- ► Introduction
- ► parallel Construct
- ► Worksharing Constructs
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- ► Data Clauses
- ► Synchronization Constructs

Introduction

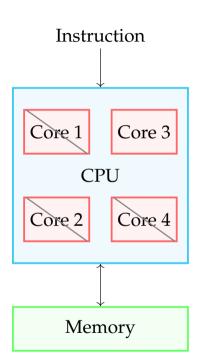
WHY OPENMP?

Consider this **serial** code. We want to run it on our *single* machine.

- Uses one available hardware core (the others are idle)
- ► Instructions performed one by one
- ► No concurrent access to memory

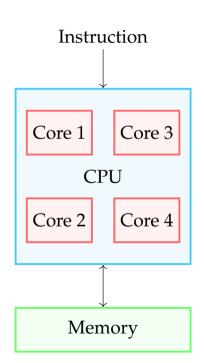
Waste of resources!

WHY OPENMP?



WHY OPENMP?

```
void add(double* x, double* y,
         double a, size t size)
    // Concurrently...
    // On core 1
    for (size t i = 0; i < size/4; i++)
        x[i] = a * x[i] + y[i];
    // On core 2
    for(size_t i = size/4; i < size/2; i++)</pre>
       x[i] = a * x[i] + y[i];
    // ...
```



NOTATION INTERLUDE

Core (Hardware)

- ► Hardware construct
- Single independent computing unit
- Reads and execute process instructions

Process (Software)

- Program under execution
- Might contain several concurrent execution flows (threads)
- Number of threads is dynamic during execution

Thread (Software)

- Smallest unit of execution
- All threads have access to process memory and system resources
- Several concurrent threads might improve performance (parallelism)
- Usually one thread is mapped to a single core

NOTATION INTERMEZZO

Difference between **process** and **thread**.



One process One thread



Multiple processes One thread



One process Multiple threads



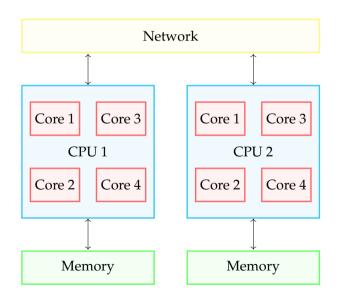


Multiple processes Multiple threads

WHY NOT MPI?

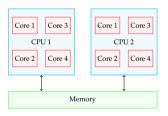
Shared Memory Core 1 Core 1 Core 3 Core 3 CPU 1 CPU 2 Core 2 Core 4 Core 2 Core 4 Memory

Distributed memory



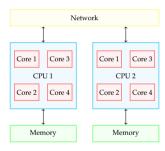
WHY NOT MPI?

Shared Memory



- Multiple processing units
- Shared memory

Distributed memory



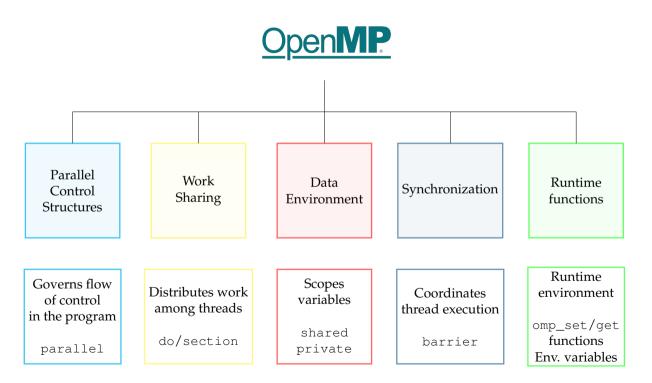
- ► More than one computing unit
- Each with own processing units and memory
- ► Interconnected via (hopefully fast) network

WHAT CAN OPENMP DO?

Parallelism paradigms supported by OpenMP standard.

- ► Threaded parallelism (multi-core, shared memory)
- ► Vectorized execution (SIMD)
- ▶ Offload execution on GPUs

WHAT IS OPENMP?



COMPILER DIRECTIVES

C/C++

```
#pragma omp construct [clauses]
{
    // Code to parallelize
}
```

Fortran

```
!$omp construct [clauses]
// Code to parallelize
!$omp end construct
```

OpenMP is **directive based**.

- ▶ Directives added to serial code
- Easy to port serial codes to multi-core CPU (and GPU)
- ► Compatible with C/C++ and **Fortran**

COMPILER FLAGS

Compiler flags to enable OpenMP.

Compilers	Fortran	C/C++
GNU	gfortran *.f90 -fopenmp	gcc/g++ *.c/cpp -fopenmp
Intel	ifort *.f90 -qopenmp	icc/icpc *.c/cpp -qopenmp
PGI	pgf90 *.f90 -mp	pgcc *.c/cpp -mp

Note

Pre-processor directives are ignored (without compiler errors) if the OpenMP compiler flag is missing. Thus remember to include this flag to compile with OpenMP parallism.

parallel construct

Question

What happens if I parallelize this simple code?

```
Output

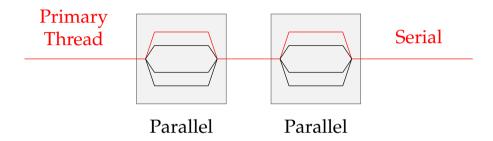
Hello World!

Hello World!

...
```

What's happening?

FORK-JOIN PARALLELISM



Fork:

- Primary thread spawns a team of secondary threads as needed
- ▶ Parallelism added incrementally until performance goals are met

Join:

- ▶ At the end of parallel region, thread team ends and only primary thread remains
- ▶ Primary thread execution is serial

What's happening?

- ► Threads **redundantly** execute code in the block
- ► Each thread evaluates std::cout
- ► Threads are synchronized at the end of the parallel region

What about the number of times Hello World! is printed?

NUMBER OF THREADS

Set the number of threads on the **terminal**.

- ► For the entire terminal session export OMP_NUM_THREADS=16; ./program
- ► For the single execution
 OMP_NUM_THREADS=16 ./program

Or in the code.

► In the OpenMP directive

```
#pragma omp parallel num_threads(16)
```

► With omp function (affects subsequent parallel regions)

```
void omp_set_num_threads(int num_threads)
```

Best practice

Rule of thumb: set the number of (software) threads = numbers of (hardware) cores.

Best practice

Avoid *magic numbers* in any code unless required. Prefer the OMP_NUM_THREADS option as bad things might happen: godbolt.org/z/K8v35dcnr.

Note

OpenMP thread number starts from 0 up to num_threads - 1.

```
#include <iostream>
#include <omp.h>
int main()
    // Note the header include above!
    omp_set_num_threads(4);
    #pragma omp parallel
    std::cout << "Hello World!"
               << std::endl;</pre>
    return 0;
```

Warning

Remember to include omp.h if OMP functions are used.

Run: OMP_NUM_THREADS=1 ./hello

```
#include <iostream>
#include <omp.h>
int main()
    omp_set_num_threads(2);
    // Note the num threads clause
    #pragma omp parallel num_threads(4)
    omp_set_num_threads(8);
    std::cout << "Hello World!"</pre>
               << std::endl;</pre>
    return 0;
```

Question

How many times Hello World! is printed in this case?

godbolt.org/z/xcTacrq1c

NUMBER OF THREADS HIERARCHY

```
Run: OMP NUM THREADS=1 ./hello
                                                         OMP NUM THREADS
#include <iostream>
#include <omp.h>
int main()
                                                                                  Hierarchy
                                                    omp_set_num_threads(int)
    omp_set_num_threads(2);
    // Note the num threads clause
    #pragma omp parallel num_threads(4)
    omp set num threads(8);
                                                        num threads clause
    std::cout << "Hello World!"</pre>
               << std::endl;</pre>
                                             godbolt.org/z/xcTacrq1c
    return 0;
```

OPENMP FUNCTIONS

most commonly used subset

Name	Result type	Purpose
omp_set_num_threads (int num_threads)	none	number of threads to be created for subsequent parallel region
omp_get_num_threads()	int	number of threads in currently executing region
omp_get_max_threads()	int	maximum number of threads that can be created for a subsequent parallel region
omp_get_thread_num()	int	thread number of calling thread (zero based) in currently executing region
<pre>omp_get_num_procs()</pre>	int	number of processors available
<pre>omp_get_wtime()</pre>	double	return wall clock time in seconds since some (fixed) time in the past
<pre>omp_get_wtick()</pre>	double	resolution of timer in seconds

Output

```
Hello World!Hello World!

Hello World!

Hello World!
```

Warning

std::cout seems to be accessed by threads non-orderly!

godbolt.org/z/5zx9GcnWv



Note

Atomic means that the operation is either completed or not. The other threads do not catch it in the middle of the operation.

```
#include <stdio.h>
#include <omp.h>
int main() {
    // Thread ID
    int tid, nthreads {4};
    omp set num threads (nthreads);
    #pragma omp parallel
    tid = omp_get_thread_num();
    nthreads=omp_get_num_threads();
    // Do some lengthy computation
    for (int i = 0; i < 1e5; i++) {};
    printf("Thread %d of %d.\n",
           tid, nthreads);
    return 0;
```

Output

```
Thread 1 of 4.
Thread 1 of 4.
Thread 3 of 4.
Thread 3 of 4.
```

godbolt.org/z/WfW6qY9aK

```
#include <stdio.h>
#include <omp.h>
int main() {
    // Thread ID
    int tid, nthreads {4};
    omp set num threads (nthreads);
    #pragma omp parallel private(tid)
    tid = omp_get_thread_num();
    nthreads=omp_get_num_threads();
    // Do some lengthy computation
    for (int i = 0; i < 1e5; i++) {};
    printf("Thread %d of %d.\n",
           tid, nthreads);
    return 0;
```

Possible solutions.

- ▶ private clause (see example)
- ▶ Define tid inside the parallel block

godbolt.org/z/K83q41fTv

EXERCISE 1



- ► Add a parallel in the correct place
- ► Get thread IDs in the parallel region
- ▶ Compile
- ► Run
- ► *Optional:* experiment with OMP_NUM_THREADS

Question

Do you get what you expected?

EXERCISE 1 - PARALLEL

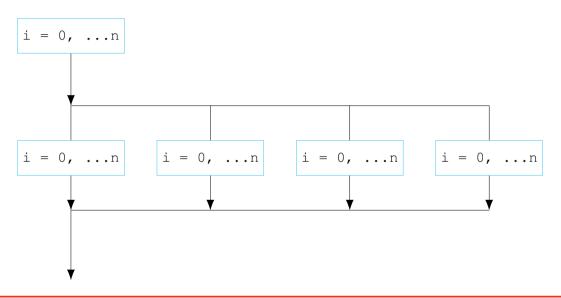
```
auto tick = omp_get_wtime();

#pragma omp parallel
{
    int tid = omp_get_thread_num();
    int nthreads{omp_get_num_threads()};

    for (std::size_t i = 0; i < n; i++) {
        c[i] = a[i] + b[i];
    }
}
auto tock = omp_get_wtime();</pre>
```

- ► Check the output: is it correct?
- ► How many times are we performing the sum on each element of c?

EXERCISE 1 - PARALLEL



Warning

parallel is executing the same code redundantly on every thread.

Worksharing constructs

EXERCISE 1 - MANUAL WORKSHARING



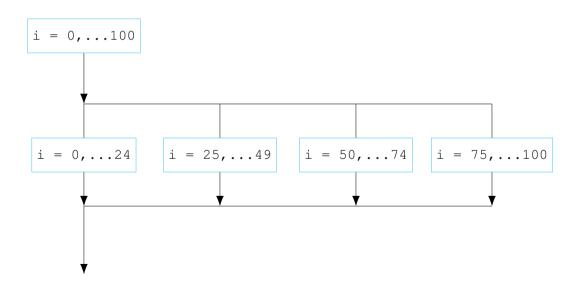
- ► Using the thread ID and the number of threads, subdivide the work between the threads (worksharing)
- ▶ Substitute the time functions with the appropriate omp functions
- ▶ Compile
- ► Run
- ► *Optional:* experiment with OMP_NUM_THREADS

Question

Do you get what you expected?

EXERCISE 1 - MANUAL WORKSHARING

```
for (std::size_t i = start; i <= end; i++) c[i] = a[i] + b[i];
```



Exercise 1 - Manual Worksharing

```
auto tick = omp get wtime();
                                            ► Check the output: is it correct?
#pragma omp parallel
                                            ► How many times are we performing
    int tid = omp_get_thread_num();
                                              the sum on each element of c?
    int nthreads = omp_get_num_threads();
    // Iteration start and end
    std::size t start = tid * n/nthreads;
    std::size t end = (tid+1) * n/nthreads - 1;
    for (std::size t i = start; i <= end; i++) {
        c[i] = a[i] + b[i];
auto tock = omp_get_wtime();
```

SPMD (SINGLE PROGRAM MULTIPLE DATA)

Single Program Multiple Data (SPDM) is the dominant style of parallel programming, where all processors use the same program, though each has its own data.

Worksharing in SPMD

- Workshare among a team of threads
- ▶ Region with no internal **barriers**, i.e., do not internally wait for other threads to continue
- ► **Implicit** barrier at the end
- ► Each thread must encounter the same worksharing regions and barriers

FOR CONSTRUCT

```
// Note the presence of parallel as well
#pragma omp parallel for
for (i = 0; i < n; i++)
{
    c[i] = a[i] + b[i];
}</pre>
```

- ► Team of threads is formed at parallel region
- ► Loop iterations are split among threads
- ► Explicit barrier at the end of the loop

Warning

Each loop iteration must be independent of the other iterations!

Note

private (i) is not needed, as loop iterators are private by default.

FOR CONSTRUCT

for construct to distribute work among threads

```
#pragma omp parallel {
    #pragma omp for {
        for (size_t i = 0; i < n; i++)
        {
            c[i] = a[i] + b[i];
        }
    }
}</pre>
```

▶ for construct in parallel region

```
#pragma omp parallel for
for (size_t i = 0; i < n; i++)
{
    c[i] = a[i] + b[i];
}</pre>
```

► Or equivalently for just a for

FOR CONSTRUCT

```
#pragma omp parallel {
    #pragma omp for {
                                                                   Team of threads is created (forked)
                                                fork
         for (i = 0; i < n; i++)
                                           distribute work
                                                                   Work is distributed among the team
                                               barrier
                                                                       Threads are synchronized
     /* Implicit barrier */
    #pragma omp for {
                                           distribute work
                                                                   Work is distributed among the team
          for (i = 0; i < n; i++)
                                                                       Threads are synchronized
                                               barrier
                                                join
                                                                  Team of threads is destroyed (joined)
     /* Implicit barrier */
```

EXERCISE 1 - FOR CONSTRUCT



- ► Switch from single parallel to parallel for
- ▶ Compile
- ► Run
- ► *Optional:* experiment with OMP_NUM_THREADS

Question

Is the workload correctly split among the threads?

NESTED LOOPS

```
for(size_t i {0}; i < n; i++) {
    for(size_t j {0}; j < n; j ++) {
        ...
}</pre>
```

What about this nested loops?

Note

Usual example: *matrix multiplication*.

$$C_{i,j} = \sum_{k=0}^{n} A_{i,k} B_{k,j}$$

with 3 loops (i, j, k).

NESTED LOOPS

```
#pragma omp parallel for
for(size_t i {0}; i < n; i++) {
    #pragma omp parallel for
    for(size_t j {0}; j < n; j ++) {
        ...
}</pre>
```

Warning

This compiles but do not work!

- Nested parallelism is disabled in OpenMP.
- ► The second pragma is effectively ignored.
- ► An *inner* team of one thread only is created, i.e., each inner loop iteration is processed by one thread
- ▶ But we create more overhead in the inner loop...

NESTED LOOPS: COLLAPSE

```
#pragma omp parallel for collapse(2)
for(size_t i {0}; i < n; i++) {
    for(size_t j {0}; j < n; j ++) {
        ...
    }
}</pre>
```

- ► Loops are collapsed into one for
- ► Iterations are shared between threads
- Each thread computes its assigned portion of iteration space
- ► Threads synchronize and join

Note

collapse might introduce additional scheduling overhead and increase computational time. Time your program after collapsing loops.

EXERCISE 2



- ▶ Parallelize the code in the appropriate regions
- ► Try without collapse
- ► Try with collapse (2)
- ► Check the results against changes in OMP_NUM_THREADS
- ▶ Optional: add another loop to perform the matrix multiplication. What does it change between collapse(2) and collapse(3)?

Question

Do you get the expected speed-up after parallelizing?

Explicit Worksharing constructs

MPMD (MULTIPLE PROGRAM MULTIPLE DATA)

Multiple Program Multiple Data (SPDM) is a style of parallel programming, where multiple autonomous processors simultaneously operate at least two independent programs.

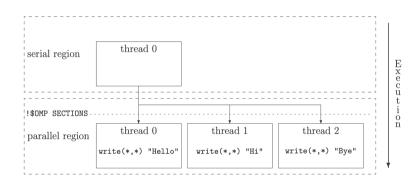
Examples

- Pre-determined independent work units
- ▶ Different functions in parallel
- ▶ Pipeline to overlap I/O with computations

SECTION CONSTRUCT

▶ section creates different tasks assigned to different threads in the parallel region. Examples: godbolt.org/z/6fxfde7z4 and godbolt.org/z/vz7ffb1oo

```
#pragma omp parallel sections
{
    #pragma omp section
    printf("Hello\n");
    #pragma omp section
    printf("Hi\n");
    #pragma omp section
    printf("Bye\n");
}
```



Warning

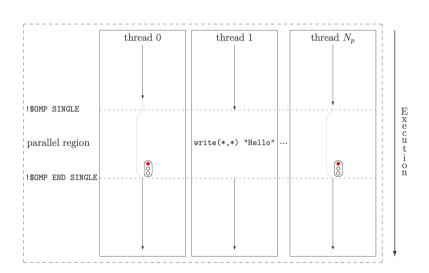
The variables defined in the parallel region are not available in the section. Example: godbolt.org/z/en9jr6acv

SINGLE CONSTRUCT

▶ single assigns a region exclusively to the first available thread. There is an **implicit** barrier.

Example: godbolt.org/z/ndhcGYba3

```
#pragma omp parallel
{
    #pragma omp single
    printf("Hello\n");
}
```



SINGLE CONSTRUCT

```
#include <cstdio>
#include <omp.h>
int main()
#pragma omp parallel num_threads(4)
    auto tid = omp get thread num();
    #pragma omp single
    printf("Hello from thread %d\n",
           tid);
    #pragma omp single
    printf("Bye from thread %d\n",
           tid);
    #pragma omp single nowait
    printf("Bye bye from thread %d\n",
           tid);
return 0;
```

Some outputs

Output

Hello from thread 1 Bye from thread 1 Bye bye from thread 2

Output

Hello from thread 3

Bye from thread 0

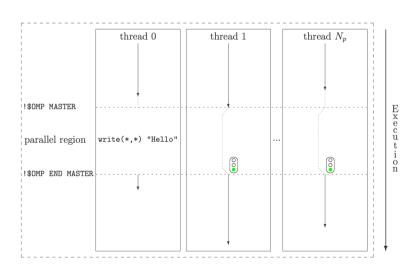
Bye bye from thread 3

godbolt.org/z/s6evTjeoe

MASTER CONSTRUCT

► master assigns a region exclusively to the master thread. There is **no implicit barrier**. Example: godbolt.org/z/aefz6Tnba

```
#pragma omp parallel
{
    #pragma omp master
    printf("Hello\n");
}
```



EXPLICIT WORK-SHARING CONSTRUCTS

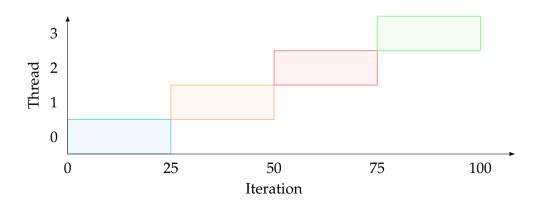
Name	Implicit barrier	Parallel region variables	Thread executing
section	Yes	No	First available
single	Yes	Yes	First available
master	No	Yes	Master

Scheduling directives

SCHEDULING

How OpenMP assign the iterations to each thread?

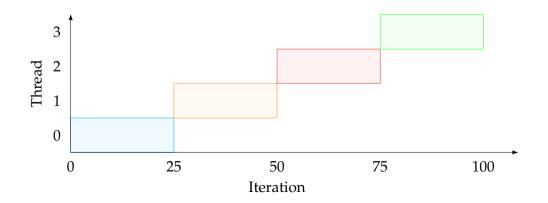
- ► Loop is splitted in chunks
- ▶ By default: static schedule with fixed chunk value



SCHEDULING: STATIC

```
#pragma omp parallel for schedule(static) for (i = 0; i < 100; i++) \{...\};
```

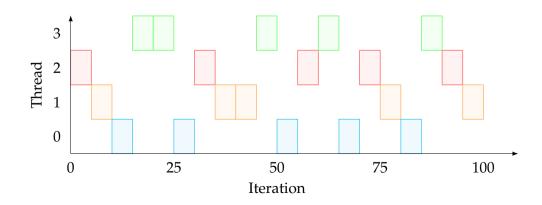
- ▶ Iterations divided into chunks
- Smallest overhead
- ► Specify chunk size with schedule (static, 25)
- ► Default chunk size divides the iterations equally



SCHEDULING: DYNAMIC

```
#pragma omp parallel for schedule(dynamic) for (i = 0; i < 100; i++) {...};
```

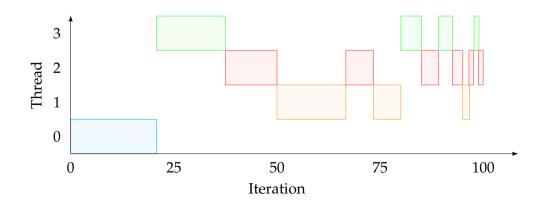
- ▶ Iterations divided into chunks
- ► Each thread requests and executes a chunk, until no chunks remains
- Useful for unbalanced workloads, e.g., when some threads complete work faster
- ▶ Default chunk size is 1



SCHEDULING: GUIDED

```
#pragma omp parallel for schedule(guided) for (i = 0; i < 100; i++) \{...\};
```

- ► Similar to dynamic, but chunk size decreases in time
- Useful for unbalanced workloads, e.g., when some threads complete work faster
- ► But mitigates overhead of dynamic by starting with large chunks



SCHEDULING

Schedule	Overhead	Default chunk	Use case
static	Small	Equally divided by threads	Small work unbalance
dynamic	Big	1	Unbalanced workload
guided	Medium	1	Unbalanced workload

Best practice

Use static for generic loops to avoid overhead. Change it when the loop/architecture is known to perform better with another.

Note

auto schedule is also available. The compiler tries to choose the best schedule.

EXERCISE 3



- ▶ Parallelize the code in the appropriate regions
- ► Add work sharing constructs
- ► Experiment with schedule clause
- ► Check the results against changes in OMP_NUM_THREADS

Question

What is the best schedule for this work?

Question

Do you get the expected speed-up after parallelizing?

AMDAHL'S LAW

Speedup
$$(n) = \frac{1}{(1-p) + p/n}$$

- ▶ *n*: the number of workers
- ightharpoonup 0 : fraction of parallelizable code

Example

- ightharpoonup n = 2 processors
- ightharpoonup p = 0.7 of the work is parallelizable

Theorical max Speedup = 1.43.

Warning

Thus sometime it might be better to revise the algorithm than blindly throw more hardware into the computation...

OPENMP OVERHEAD/SCALABILITY

OpenMP programs do not always scale well,

i.e., increasing OpenMP threads might not increase speedup.

- ► Considerable overhead involved (e.g., schedule)
- ▶ More threads competing for available memory bandwidth
- ► Cache mismanagement

Best practice

Parallelize most outer loop possible (in some cases even for small iterations)

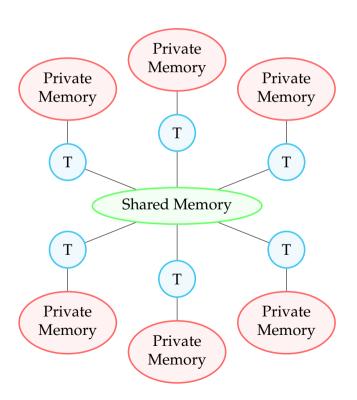
Best practice

Check the speedup against the overheads.

- ► Is the number of iterations large enough?
- ▶ Is the amount of work per iteration enough?

Data constructs

MEMORY MODEL



- ► **Shared memory**: available to all threads
- ► **Private memory**: available to only one thread

Warning

It is up to the programmer to choose the memory scope to variables.

SHARED CLAUSE

```
#include <stdio.h>
#include <omp.h>
int main()
    // Thread ID
    int x \{3\};
    omp set num threads (8);
    #pragma omp parallel shared(x)
        auto tid = omp_get_thread_num();
        x += 1;
        printf("Thread %d has x: %d.\n",
               tid, x);
    return 0;
```

shared variables have shared memory among all threads.

```
Output

Thread 0 has x: 4.

Thread 3 has x: 5.

Thread 1 has x: 6.

Thread 2 has x: 7.
```

Note

Variables are shared by default, i.e., the shared clause is redundant in this example.

godbolt.org/z/TahqMEv1E

PRIVATE CLAUSE

```
#include <stdio.h>
#include <omp.h>
int main()
    // Thread ID
    int x \{3\};
    omp set num threads (3);
    #pragma omp parallel for private(x)
    for(size_t i = 0; i < 9; i++)
        auto tid = omp_get_thread_num();
        x = i:
        printf("Thread %d has x: %d.\n",
               tid, x);
    printf("Final x: %d.\n", x);
    return 0;
```

private variables are **uninitialized** copies of the global ones, and visible to its thread.

```
Output

Thread 0 has x: 4.
Thread 0 has x: 0.
Thread 0 has x: 1.
Thread 0 has x: 2.
Thread 2 has x: 6.
Thread 2 has x: 7.
Thread 2 has x: 8.
Thread 1 has x: 3.
Thread 1 has x: 4.
Thread 1 has x: 5.
Final x: 3.
```

godbolt.org/z/ncYG1n88z

PRIVATE CLAUSE

```
#include <stdio.h>
#include <omp.h>
int main()
    // Thread ID
    int x \{3\};
    omp set num threads (3);
    #pragma omp parallel for private(x)
    for (size t i = 0; i < 9; i++)
        auto tid = omp_get_thread_num();
        x += i; // Instead of =
        printf("Thread %d has x: %d.\n",
               tid, x);
    printf("Final x: %d.\n", x);
    return 0;
```

Warning

Be careful about the initialization!

```
Output

Thread 1 has x: 5.

Thread 1 has x: 9.

Thread 1 has x: 14.

Thread 2 has x: 6.

Thread 2 has x: 13.

Thread 2 has x: 21.

Thread 0 has x: -608901888.

Thread 0 has x: -608901887.

Thread 0 has x: -608901885.

Final x: 3.
```

godbolt.org/z/xvb4bfs45

FIRSTPRIVATE CLAUSE

```
#include <stdio.h>
#include <omp.h>
int main()
    // Thread ID
    int x \{3\};
    omp set num threads (3);
    #pragma omp parallel for \
            firstprivate(x)
    for (size_t i = 0; i < 3; i++) {
        auto tid = omp_get_thread_num();
        printf("Thread %d has x: %d.\n",
               tid, x);
        x = i;
        printf("Thread %d has x: %d.\n",
               tid, x);
    printf("Final x: %d.\n", x);
    return 0;
```

firstprivate variables are copies of the global ones, visible to its thread, and initialized to the global variable value.

```
Output

Thread 0 has x: 3.

Thread 0 has x: 0.

Thread 2 has x: 3.

Thread 2 has x: 2.

Thread 1 has x: 3.

Thread 1 has x: 1.

Final x: 3.
```

Best practice

Prefer firstprivate to avoid unwanted uninitialized variables.

godbolt.org/z/3jv6K84br

LASTPRIVATE CLAUSE

```
#include <stdio.h>
#include <omp.h>
int main()
    // Thread ID
    int x \{3\};
    omp set num threads (3);
    #pragma omp parallel for lastprivate(x)
    for (size_t i = 0; i < 3; i++)
        auto tid = omp_get_thread_num();
        x = i:
        printf("Thread %d has x: %d.\n",
               tid, x);
    printf("Final x: %d.\n", x);
    return 0;
```

lastprivate variables are **uninitialized** copies of the global ones, visible to its thread. Once the parallel region is left, it is copied back to the value of the last iteration.

```
Output

Thread 0 has x: 0.
Thread 1 has x: 1.
Thread 2 has x: 2.
Final x: 2.
```

godbolt.org/z/MeMa757x6

DATA CLAUSES

Name	Sharing Policy	Initialization	Outside Value
shared	All threads	Global value	Last modified
private	One thread	No	Pre-parallel region
firstprivate	One thread	Global value	Pre-parallel region
lastprivate	One thread	No	Last iteration

Best practice

Prefer firstprivate over private to avoid uninitialized variable issues.

Note

Default data sharing policy:

- ▶ shared if defined outside of the parallel region
- ▶ private if defined inside the parallel region or loop iterators when using the for construct.

DEFAULT CLAUSE

```
// Some variables with any datatype T
T \times, y
// All variables are shared
#pragma omp parallel default(shared) \
                     private(x)
// All variables are private
#pragma omp parallel default(private) \
                      shared(v)
// Force to specify explicitly the
// data sharing policy
#pragma omp parallel default(none) \
                     private(x)
                      shared(v)
```

default changes the default data policy behaviour.

Warning

Using none, the programmer is forced to specify the data sharing policy for all variables used in the parallel region.

Best practice

Always use default (none) and explicity specify the data sharing policy for all variables used in the parallel region. [clang-tidy]

Synchronization Constructs

EXAMPLE

```
#include <iostream>
#include <omp.h>
int main()
    int sum {0}, n {10000};
    auto start = omp_get_wtime();
    for (size t i = 1; i < n; i++) {
        sum += i;
    auto end = omp_get_wtime();
    std::cout << "Sum: " << sum
              << " (compare with "
               << n*(n-1)/2 << ")."
              << std::endl;</pre>
    std::cout << "Runtime: "</pre>
              << end-start << std::endl; godbolt.org/z/TqaM1nczG</pre>
    return 0;
```

Output

Sum: 49995000 (compare with 49995000). Runtime: 4.1008e-05

Ouestion

Can you parallelize this loop?

EXAMPLE

```
#include <iostream>
#include <omp.h>
int main()
    int sum {0}, n {10000};
    auto start = omp get wtime();
    #pragma omp parallel for
    for(size_t i = 1; i < n; i++) {
        sum += i;
    auto end = omp_get_wtime();
    std::cout << "Sum: " << sum
              << " (compare with "
               << n*(n-1)/2 << ")."
               << std::endl;</pre>
    std::cout << "Runtime: "</pre>
              << end-start << std::endl; godbolt.org/z/MaMEoq6Kh</pre>
    return 0;
```

Output

Sum: 36949316 (compare with 49995000). Runtime: 9.6305e-05

Warning

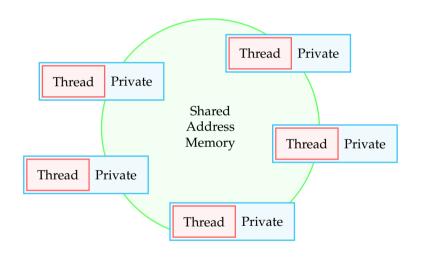
Again: data race!

But in this case, private is not the best choice...

Ouestion

What happens using clang instead?

DATA RACE



- Data race occurs when two or more threads access the same memory location
- Synchronization to assure legal order of operations

Warning

It is up to the programmer to choose the correct synchronization method.

DATA RACE SOLUTIONS

Mutual Exclusion (Mutex)

Define block of code executed by one thread only at a time.

Constructs

- ▶ critical
- ► atomic
- ▶ barrier

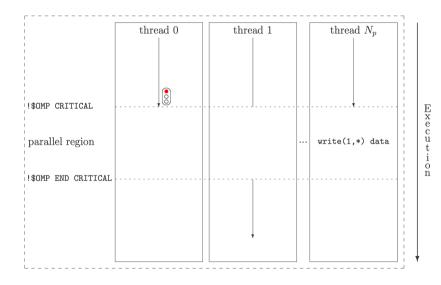
Reduction

Perform global operation on quantity from different threads.

Operations

- ► Sum, subtraction, multiplication
- ► Logical operations: or, and, ...
- Maximum, minimum

CRITICAL CONSTRUCT



critical prevents multiple threads from accessing a section of the code at the same time.

- ▶ **Pro**: prevents data race
- ► Con: other threads wait their turn in idle status

CRITICAL CONSTRUCT

```
int sum {0}, n {10000};
auto start = omp_get_wtime();
#pragma omp parallel for
for(size_t i = 1; i < n; i++) {
    #pragma omp critical
    sum += i;
}
auto end = omp_get_wtime();</pre>
```

critical prevents multiple threads from accessing a section of the code at the same time.

Output

Sum: 49995000 (compare with 49995000).

Runtime: 0.000228857

godbolt.org/z/e8aq3MWj7

ATOMIC CONSTRUCT

```
int sum {0}, n {10000};
auto start = omp_get_wtime();
#pragma omp parallel for
for(size_t i = 1; i < n; i++) {
    #pragma omp atomic
    sum += i;
}
auto end = omp_get_wtime();</pre>
```

atomic guarantees mutually exclusive access to a specific memory location (variable).

Output

Sum: 49995000 (compare with 49995000).

Runtime: 0.000147173

godbolt.org/z/zq4b7fKr3

ATOMIC CONSTRUCT LIMITATIONS

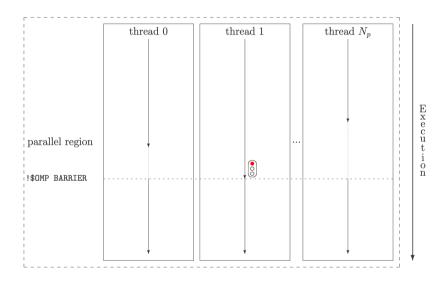
atomic guarantees exclusive access to a variable for the following operations.

- ightharpoonup Read: v = x
- ightharpoonup Write: x = v
- ► Update: x++, x-, -x, +=, *= ...
- ► Capture: v = x++,...

Limitations

- ▶ x and v must be scalar
- ► No operator overload
- ► No complex expressions

BARRIER CONSTRUCT



barrier prevents threads from continuing the execution until all threads reach the it.

- ► **Pro**: prevents data race
- ► Con: other threads wait in idle status

Best practice

Use barrier as little as possible! If you have too many barriers it might be good to consider alternatives or rethink the algorithm.

BARRIER CONSTRUCT

```
#include <cstdio>
#include <omp.h>
int main()
    #pragma omp parallel
        auto tid {omp_get_thread_num()};
        printf("Thread %d here \n", tid);
        #pragma omp barrier
        printf("Thread %d passed the "
        "barrier \n", tid);
    return 0;
```

barrier prevents threads from continuing the execution until all threads reach the it.

Output without barrier

```
Thread 0 here
Thread 0 passed the barrier
Thread 1 here
Thread 1 passed the barrier
```

Output with barrier

```
Thread 0 here
Thread 1 here
Thread 0 passed the barrier
Thread 1 passed the barrier
```

godbolt.org/z/cYrrYYE84

DATA RACE SOLUTIONS

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Reduction

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Operations

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- Maximum, minimum

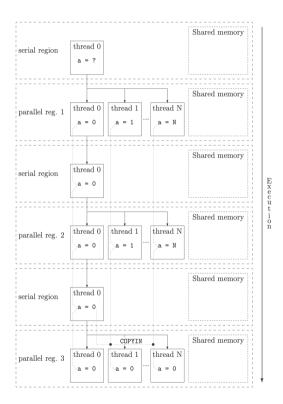
EXAMPLE

```
#include <iostream>
#include <omp.h>
int main()
    int sum {0}, n {10000};
    auto start = omp_get_wtime();
    #pragma omp parallel for
    for (size_t i = 1; i < n; i++) {
        sum += i;
    auto end = omp_get_wtime();
    std::cout << "Sum: " << sum
              << " (compare with "
               << n*(n-1)/2 << ")."
              << std::endl;
    std::cout << "Runtime: "</pre>
              << end-start << std::endl;</pre>
    return 0;
```

We want a way to gather contributions of the sum from different threads.

- ► Manual ... No ...
- ▶ reduction clause

REDUCTION CLAUSE



reduction combines values on a single accumulation variable.

Syntax: reduction(operation:var)

Operations

+	▶ .neqv
> -	▶ .max
*	▶ .min
▶ .and	▶ .iand
▶ .or	▶ .ior
▶ .eqv	▶ .ieor

REDUCTION CLAUSE

```
#include <iostream>
#include <omp.h>
int main()
                                              Output
                                              Sum: 49995000 (compare with
    int sum {0}, n {10000};
                                              49995000).
                                              Runtime: 7.1988e-05
    auto start = omp get wtime();
    #pragma omp parallel for reduction(+:sum)
    for (size_t i = 1; i < n; i++) {
                                             godbolt.org/z/45xKrsqox
        sum += i;
    auto end = omp_get_wtime();
    std::cout << "Sum: " << sum
              << " (compare with "
              << n*(n-1)/2 << ")."
              << std::endl;
    std::cout << "Runtime: "</pre>
              << end-start << std::endl;</pre>
    return 0;
```

NOWAIT CLAUSE

```
#pragma omp parallel
    #pragma omp for
    for(size_t i=0; i<10; i++) {
        c();
    d();
#pragma omp parallel
    #pragma omp for nowait
    for(size t i=0; i<10; i++) {
        c();
    d();
```

Without nowait, threads wait for completion before evaluating d() Example: godbolt.org/z/qbadoqG19

▶ With nowait, threads can evaluate d() as soon as they exit the for Example: godbolt.org/z/8hs4fhaTd

EXERCISE 4



Computation of π using

$$\pi = \int_0^1 \frac{4}{1+x^2} \, dx \cong \sum_{i=0}^N \frac{4}{1+x^2} \, \Delta x \, .$$

- ▶ Parallelize the code in the appropriate regions
- ▶ Try to add critical and atomic to solve the data race
- ► Then try with reduction
- Run and time the code
- ► Check the results against changes in OMP_NUM_THREADS

EXERCISE 4

Implementation	Runtime (s)
Serial	0.14
critical	10.60
atomic	8.34
reduction	0.04

Best practice

Use reduction when possible, especially with SPMD algorithms.

Take-Home Messages

TAKE-HOME MESSAGES

- ▶ OpenMP: to parallelize on shared memory systems (or hybrid with MPI)
- ▶ OpenMP works with threads: master thread is forked in threads when a parallel region is reached; the control is given back to the master at the end of the parallel region
- ▶ **OpenMP** consists in directives (code), run-time routines (in omp.h) and environmental variables (shell)

TAKE-HOME MESSAGES

Constructs/Clauses

- ► To create teams of threads
 - parallel
- ► To share work among threads
 - for (SPMD)
 - section (MPMD)
 - single (MPMD)
 - master (MPMD)
- ► To prevent conflicts
 - critical
 - atomic
 - barrier
- ► Data environment clauses
 - private
 - firstprivate
 - lastprivate
 - reduction

API functions

- ► To manage threads
 - omp_set_num_threads(int num_threads)
 - omp_get_num_threads()
 - omp_get_thread_num()
 - omp_ge_num_procs()
- ► To time the code
 - omp_get_wtime()
 - omp_get_wtick()

TAKE-HOME MESSAGES

Best practice

Prefer OpenMP over MPI on shared memory systems: the overhead of memory transfer in MPI is huge. Or an hybrid approach for both shared and distributed memory systems.

Warning

Be aware of the overheads of OpenMP (scheduling, barriers, ...). Always time (or better profile) your code.

TAKE-HOME MESSAGES: C++

Warning

Be careful when using std::cout, std:cin or similar, due to data race. This is another reason a good debugger might help!

Warning

Check if some advanced functionalities of C++ are compatible with the constructs used. For instance, avoid brace initialization for the internal loop iterator variables, e.g. **avoid**

```
#pragma omp parallel for
for(size_t i {0}; i < N; i++) { ... };</pre>
```

RESOURCES

- ► Essential: Official OpenMP Reference Guides
- ► Tim Mattson's OpenMP Tutorial
- ► OpenMP Compilers and Tools Official List
- ► And obviously: Stackoverflow (with OpenMP tag)

Thank you!