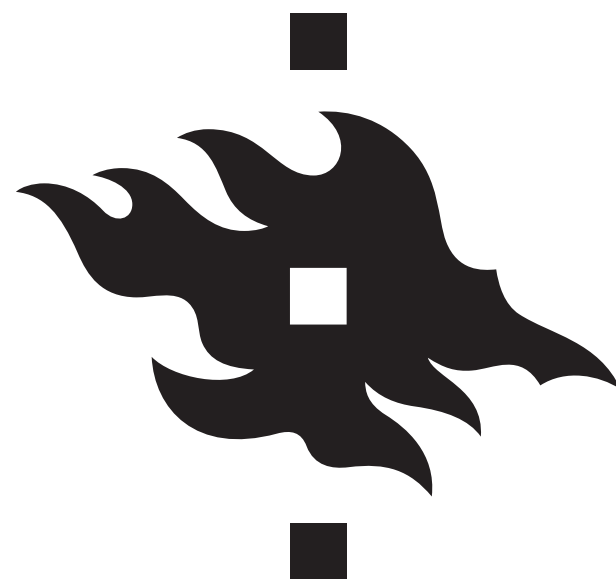


Genome Assembly

From theory to practice (and back)



UNIVERSITY OF HELSINKI

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Department of Computer Science



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Outline

1. The problem
2. Practical issues
3. Theoretical problem formulations
4. Practical genome assembly
 - Contig assembly
 - Scaffolding
 - Gap filling
5. And back: a more "practical" theoretical formulation

LECTURE
Theory + Abstract view



ASSIGNMENT
Practical + Hands-on view

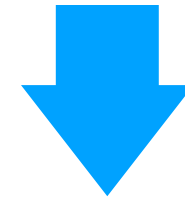
The problem

(A general description of the input and output)

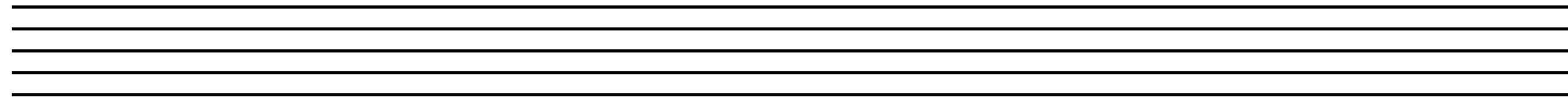
Short-read sequencing

(Third-generation sequencing)

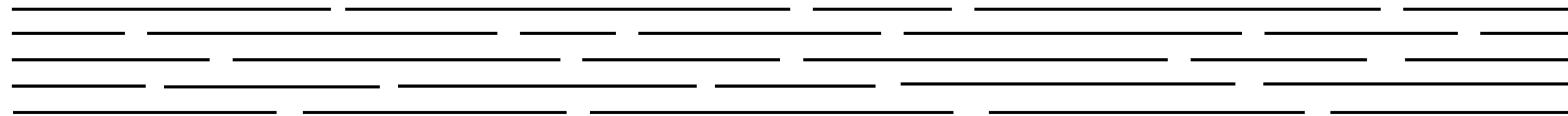
DNA



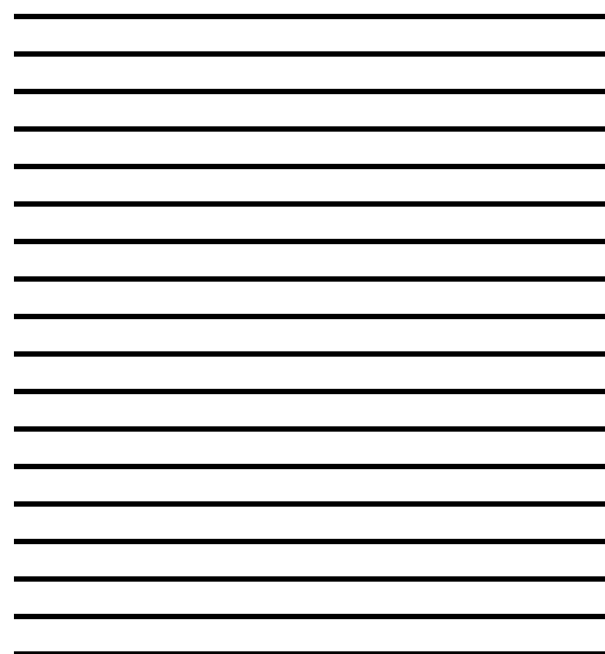
Amplification



Fragmentation

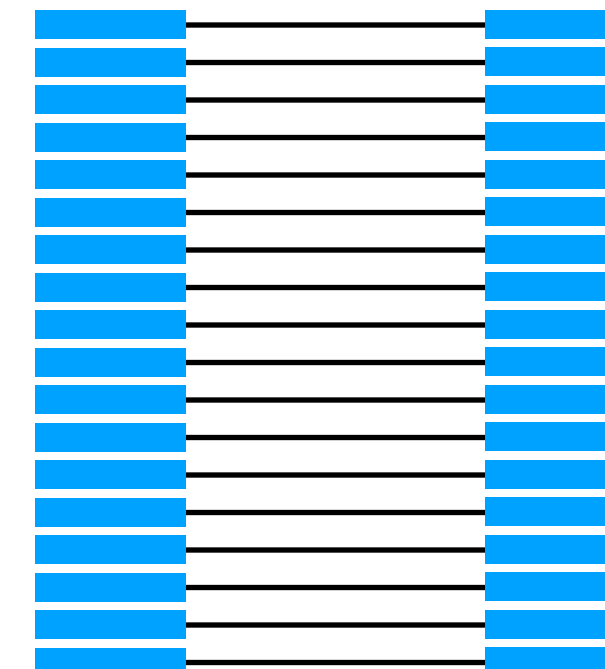
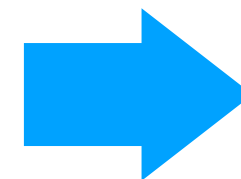


Size selection



length ~450

Sequencing



length 100 ~250 100

The assembly problem

INPUT: A collection of paired-end reads

OUTPUT: The genome from which they were sequenced

(we will see precise computational formulations later)

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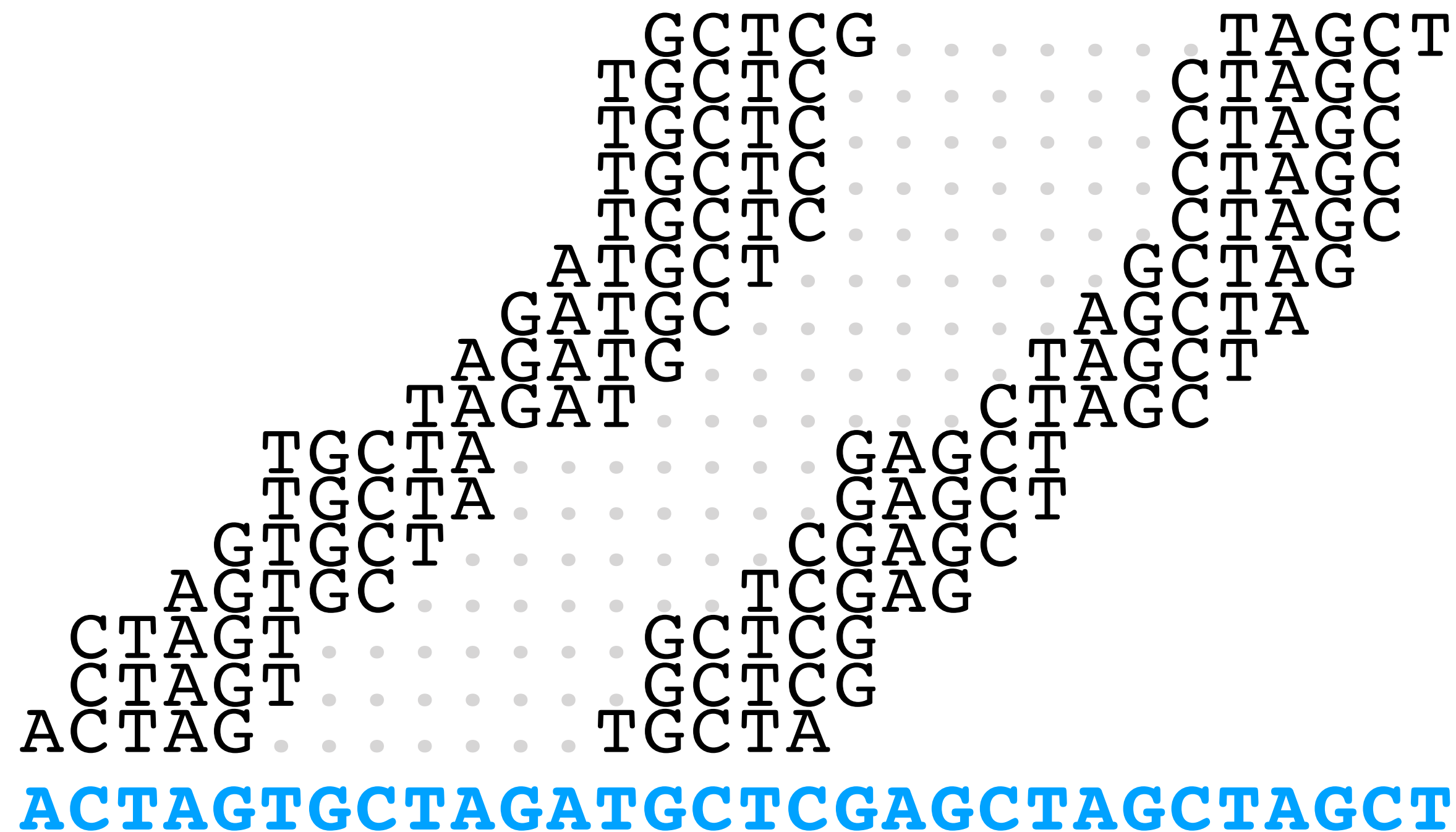
ACTAGTGCTAGATGCTCGAGCTAGCTAGCT

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INPUT: A collection of paired-end reads

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T	G	C	T	C	C	T	A	G	C
T	G	C	T	C	C	T	A	G	C
T	G	C	T	A	G	G	A	G	C
C	T	A	G	T	G	C	T	C	G
A	T	G	C	T	A	G	C	T	A
G	A	T	G	C	T	C	G	A	G
A	G	A	T	G	T	A	G	C	T
T	A	G	A	T	C	T	A	G	C
T	T	G	C	T	C	T	A	G	C
T	T	G	C	T	G	A	G	C	T
G	T	G	C	T	C	G	A	G	C
G	T	G	C	T	C	G	A	G	C
C	T	A	G	T	T	A	G	C	T
A	C	T	A	G	G	T	C	C	G
					T	G	C	T	A

???

Practical issues

(several of which are not covered in this course)

Repeats longer than read length

MAIN TOPIC

- If every substring of the genome of length = read length - 1 is unique, then assembly is trivial

AATTGAATTACACCAC

AATTG

ATTGA

TGAAT

TGAAT

GAATT

AATTT

ATTTA

TTTAC

TTACA

TACAC

ACACC

CACCA

ACCAC

read length - 1

Repeats longer than read length

MAIN TOPIC

- If every substring of the genome of length = read length - 1 is unique, then assembly is trivial
- Repeated substrings at least as long as the read length introduce ambiguity

Repeats longer than read length

MAIN TOPIC

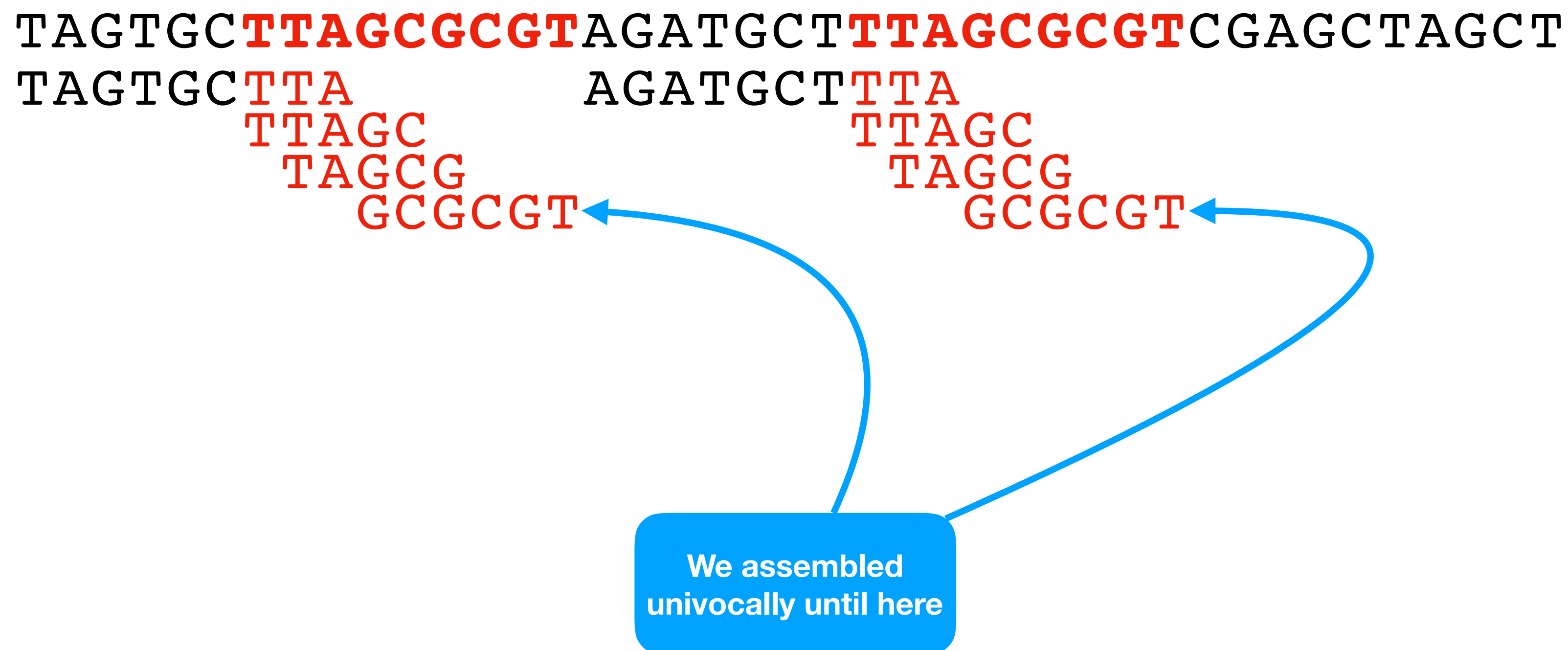
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- Repeated substrings at least as long as the read length introduce ambiguity

TAGTGC**TTAGCGCGT**AGATGCT**TTAGCGCGT**CGAGCTAGCT
TAGTGC**TTA**AGATGCT**TTA**
TAGC
TAGCG
GCGCGT

Repeats longer than read length

MAIN TOPIC

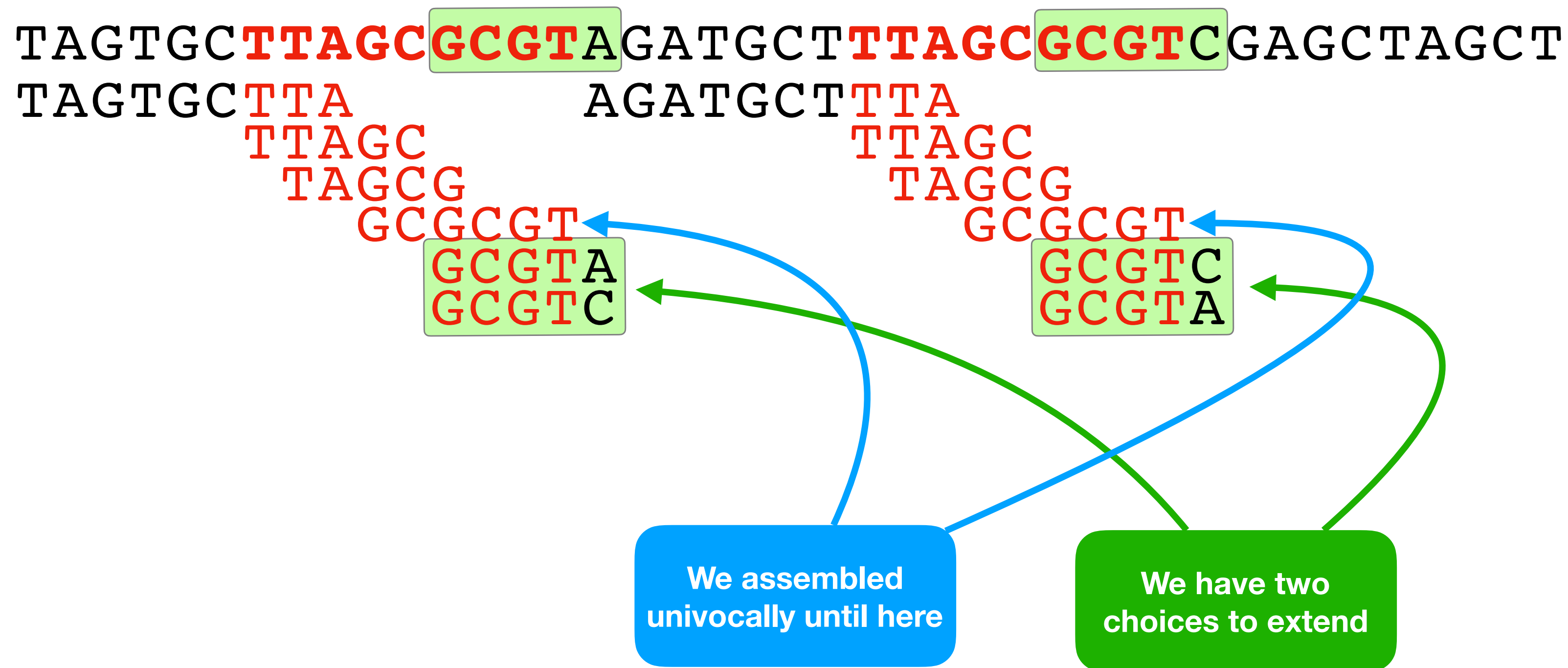
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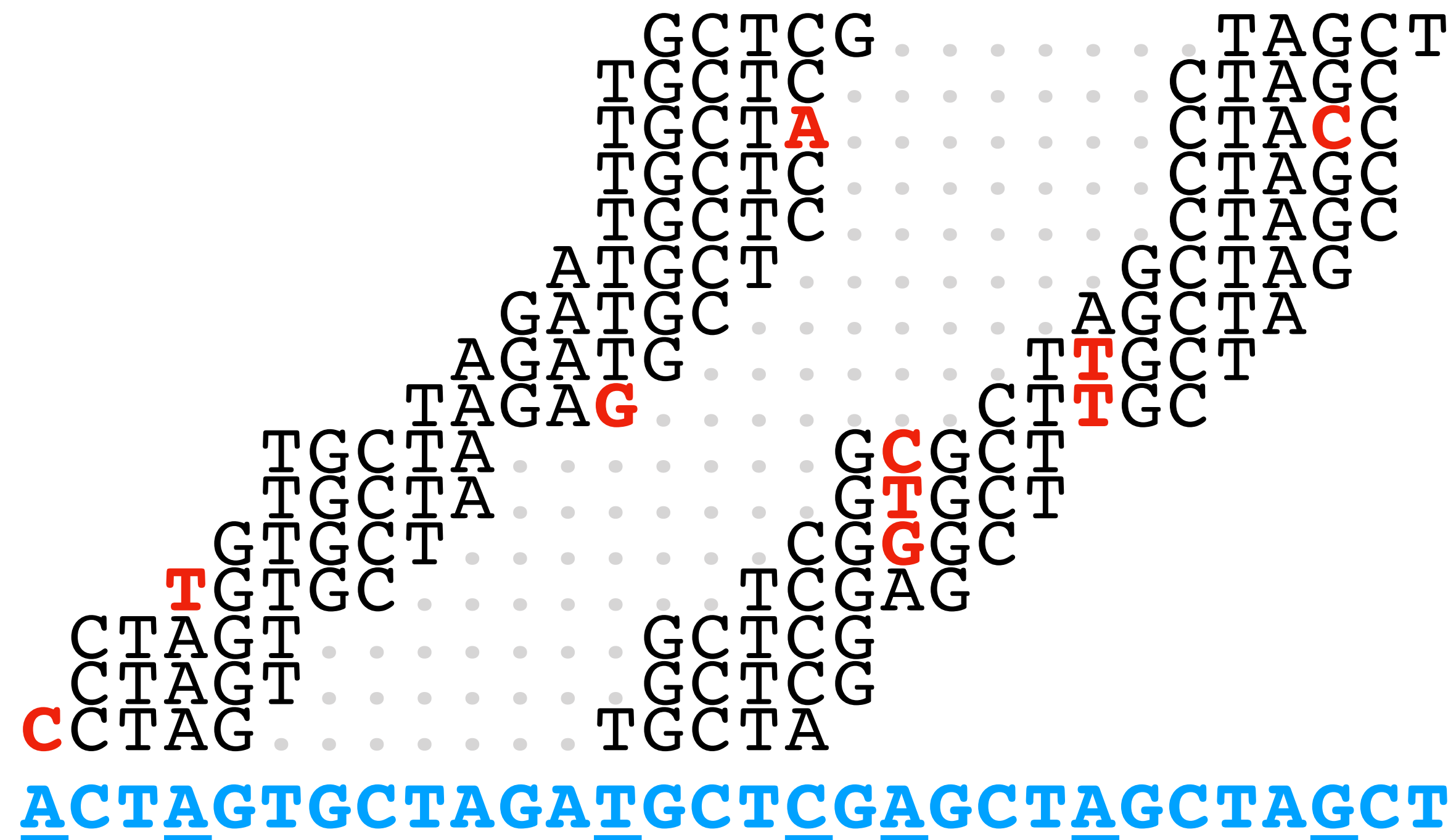
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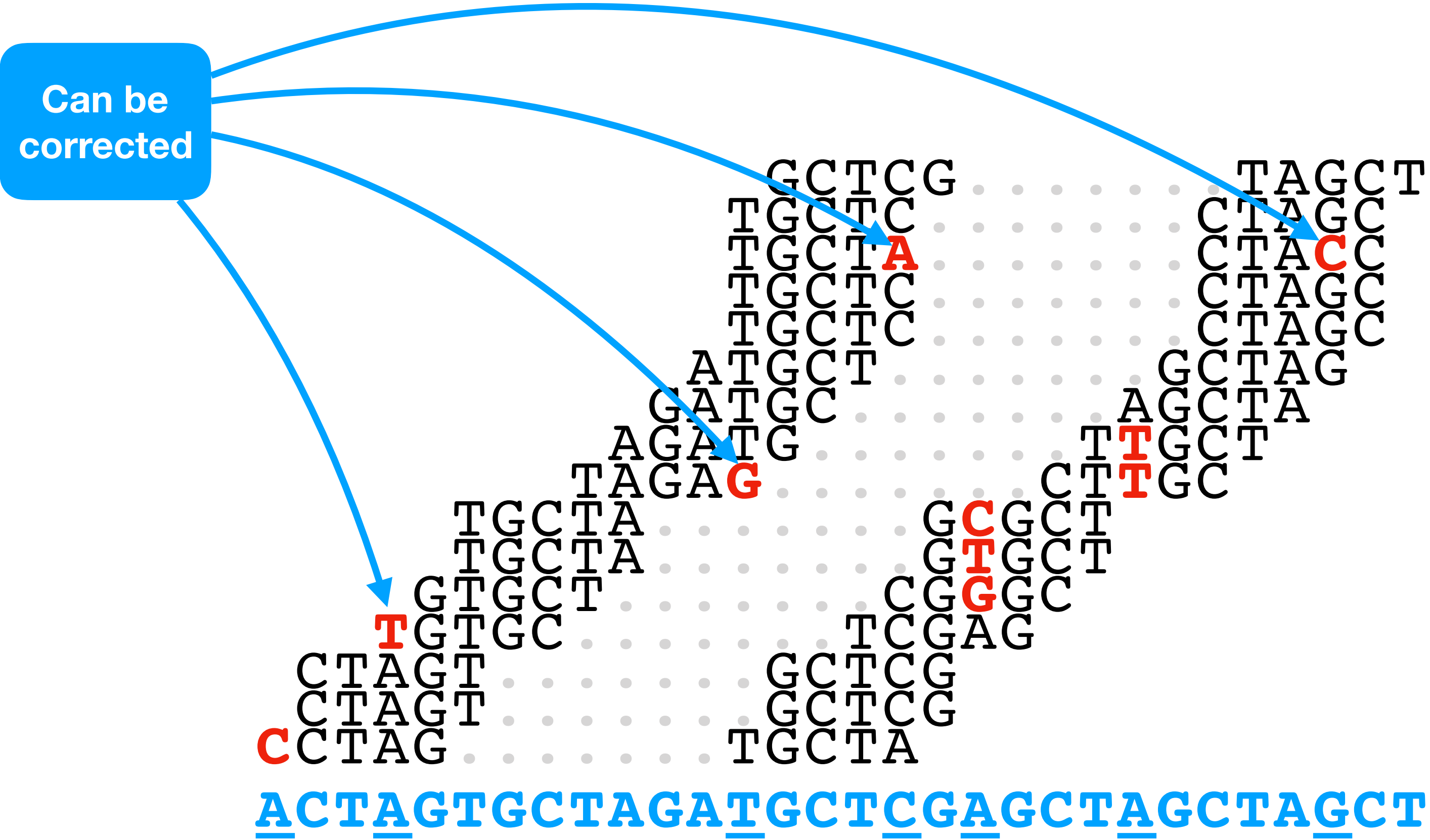
Sequencing errors

ASSIGNMENT



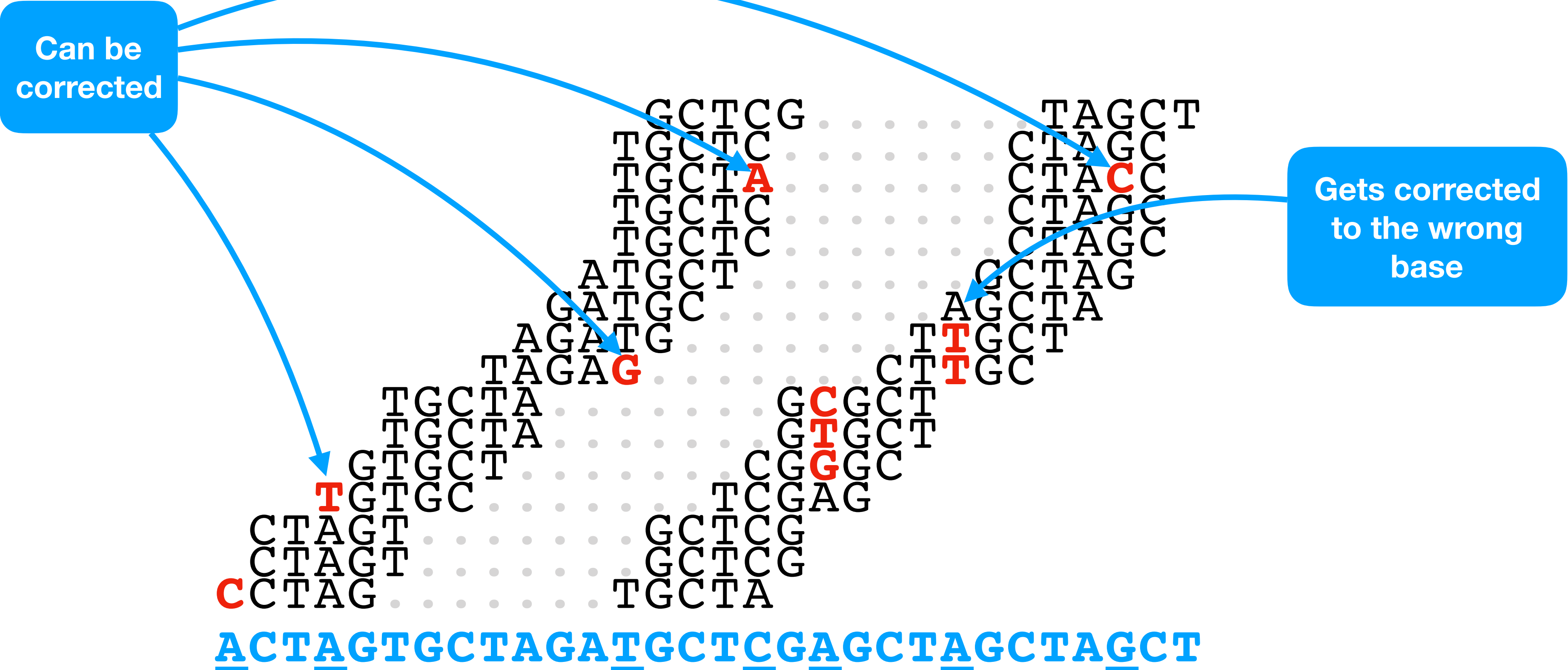
Sequencing errors

ASSIGNMENT



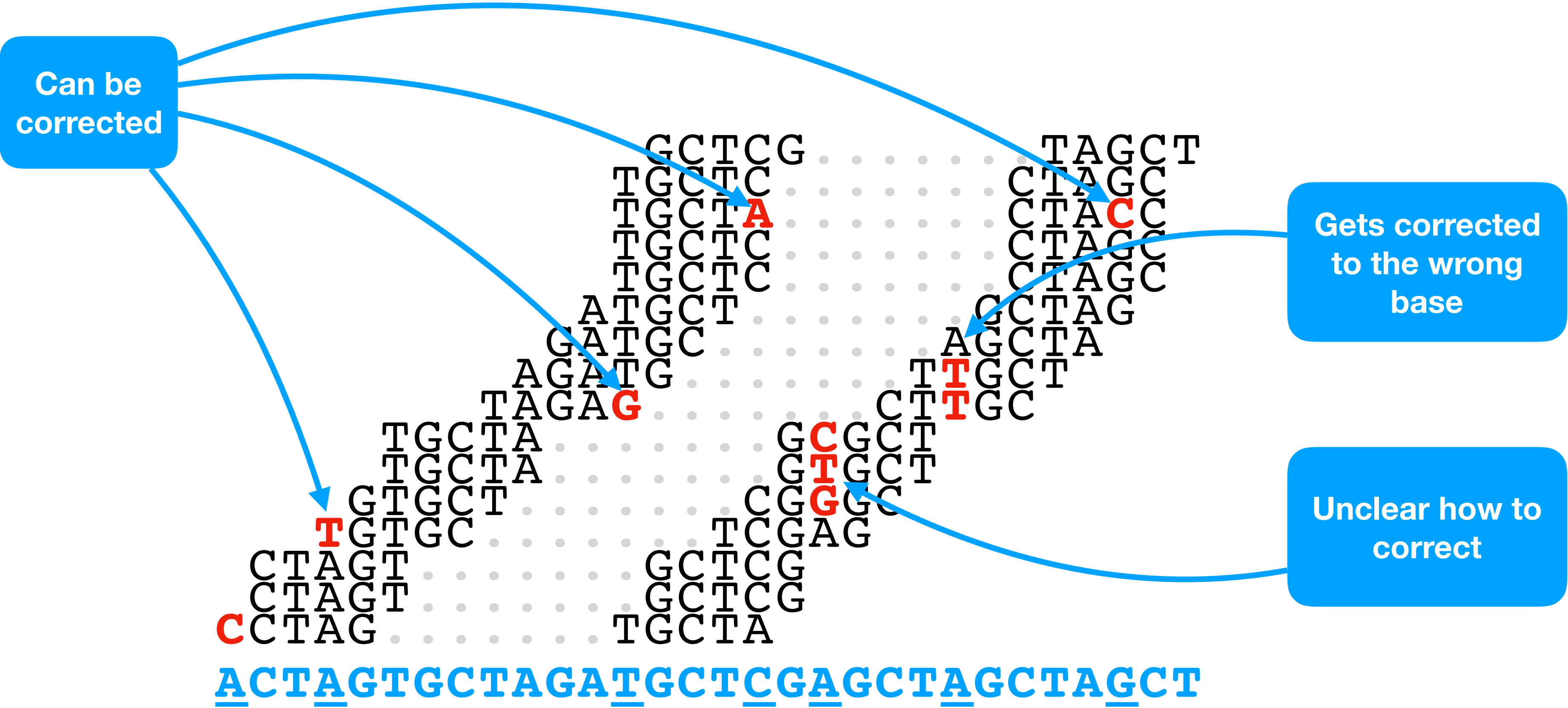
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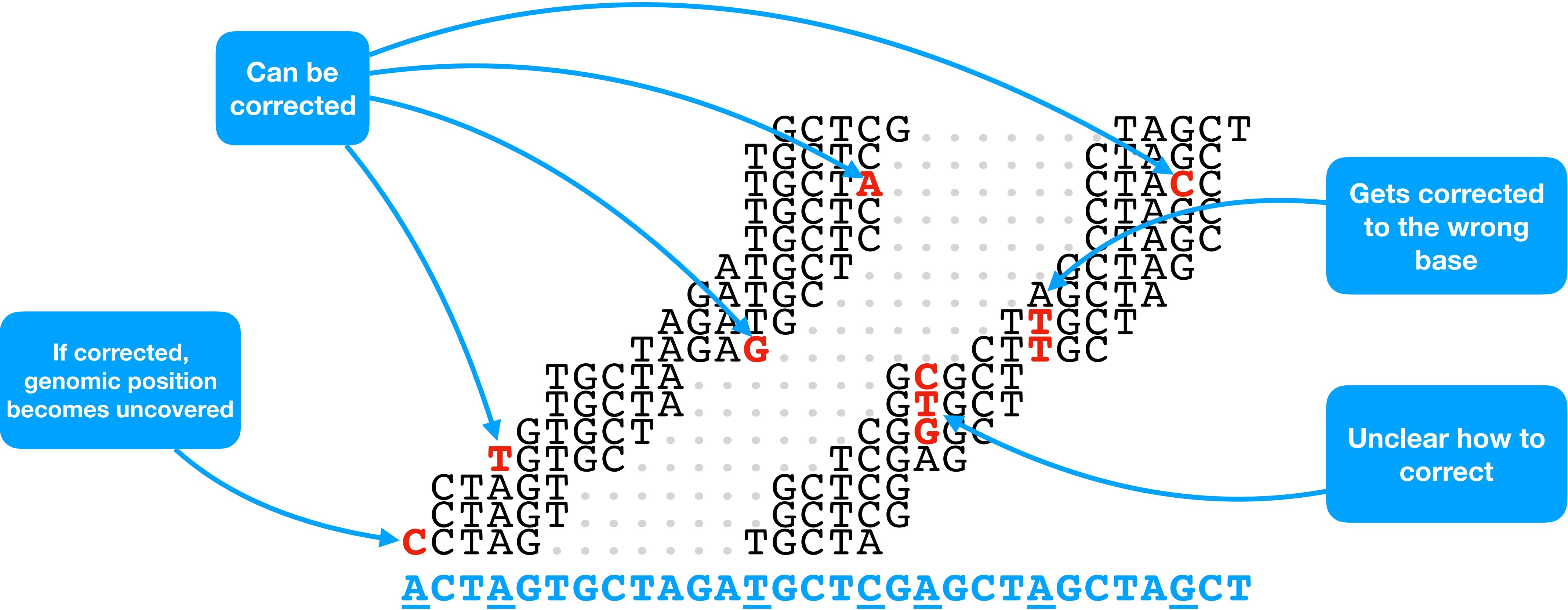
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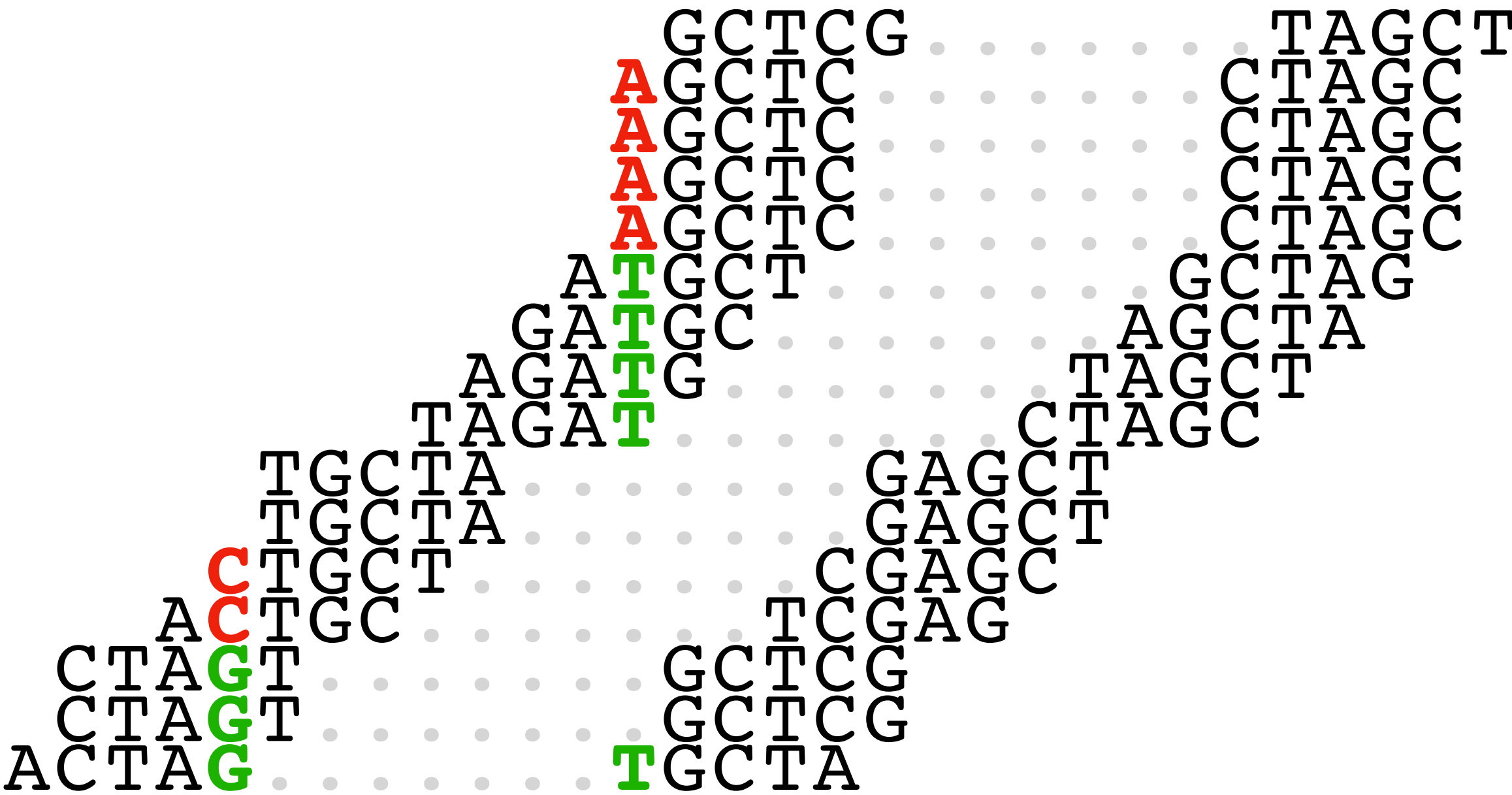
Sequencing errors

ASSIGNMENT



Ployploidy

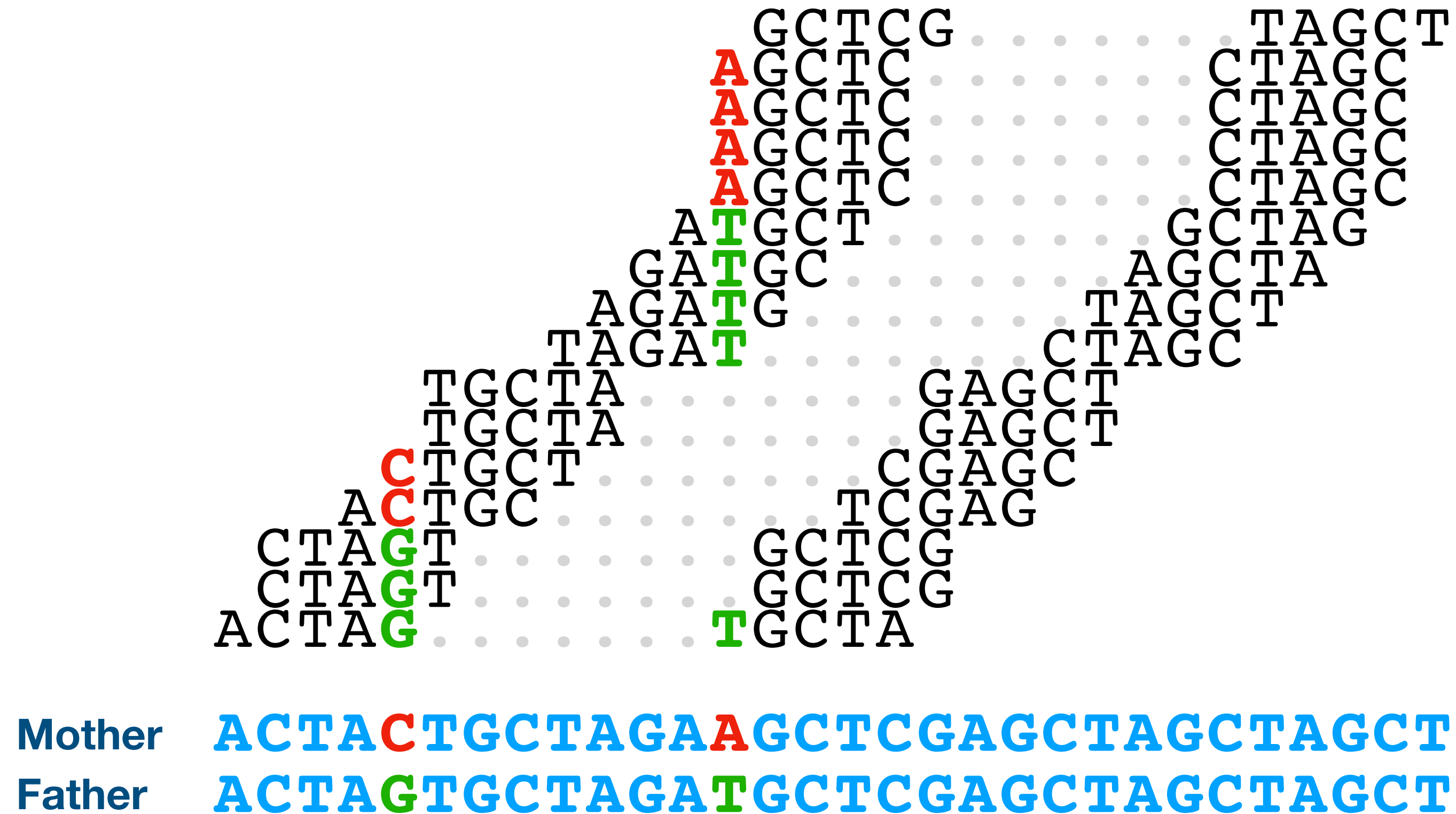
LATER IN THE LECTURE



Mother **ACTA****C****TGCTAGAA****A****GCTCGAGCTAGCTAGCT**
Father **ACTA****G****TGCTAGAT****G****GCTCGAGCTAGCTAGCT**

Polyploidy

LATER IN THE LECTURE



- Sequencing errors + polyploidy at the same time
- Phasing SNPs (**C** and **A** from same haplotype, **NOT** e.g. **C** and **T**) is a separate problem, called *haplotype assembly* or *haplotype phasing*

Unsequenced areas

CTAGT GCTCG
CTAGT GCTCG
ACTAG TGCTA

ACTAGT . . TAGATGCTCG . . CTAGCTAGCT

Non uniform paired-end distance

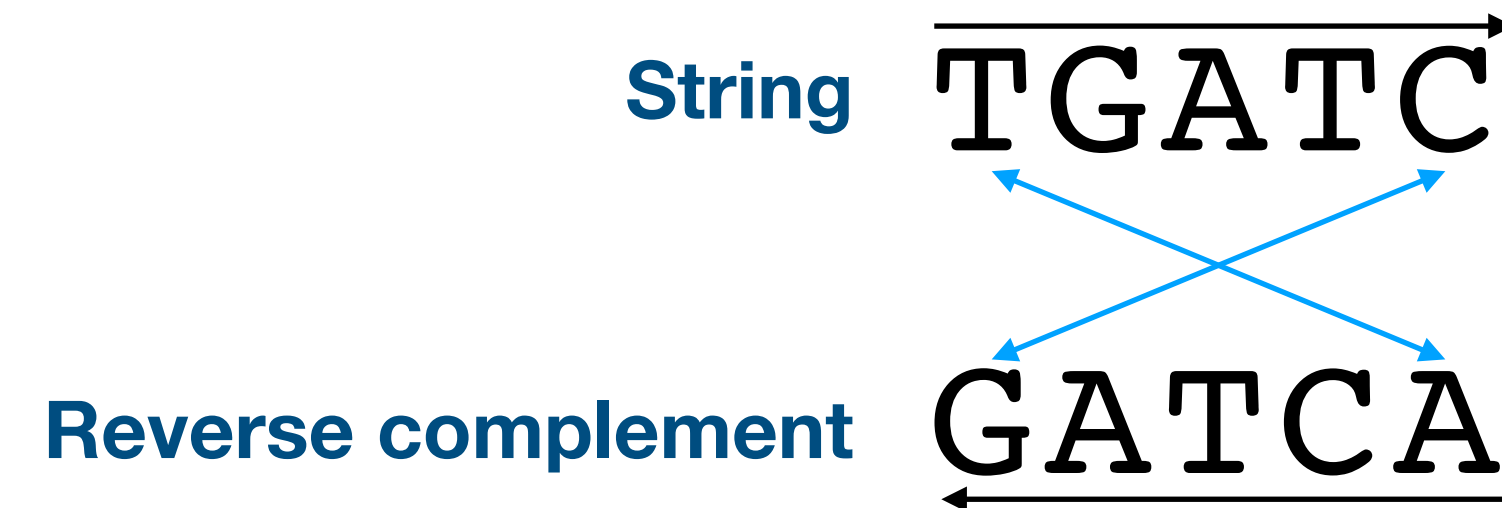
ACTAGTGCTAGATGCTCGAGCTAGCTAGCT

- Distance between each pair not known "exactly" from sequencer

Double-stranded DNA

NOT IN THE LECTURE

- Reads consist of strings and their reverse complements:



NOT IN THE LECTURE

-
- String TGATC
- Reverse complement GATCA



Double-stranded DNA

NOT IN THE LECTURE

- Reads consist of strings and their reverse complements:

String $\overrightarrow{\text{TGATC}}$

Reverse complement $\overleftarrow{\text{GATCA}}$

Diagram illustrating a double-stranded DNA structure. The top strand is a sequence of reads: CTAGT , AGTGT , TGCTA , GAGCT , GAGCT , CGAGC , TCGAG , GCTTC , CGAGC , GATCA , ACGAG . The bottom strand is the reverse complement of the top strand, shown as a sequence of reads: TGATC , GATCA , ACGAT , CTCGA , CTCGA , GCTTC , GCTTC , CGAGC , CGAGC , GATCA , TGATC . The two strands are connected by dots representing base pairing. The top strand is labeled "String" and the bottom strand is labeled "Reverse complement".

Diagram illustrating a double-stranded DNA structure. The top strand is a sequence of reads: TGCTA , GAGCT , GAGCT , CGAGC , TCGAG , GCTTC , CGAGC , GATCA , ACGAG . The bottom strand is the reverse complement of the top strand, shown as a sequence of reads: AGCTC , TAGCA , TAGCA , AGCAC , AGCAC , ACTAG , ACTAG , CTAGT . The two strands are connected by dots representing base pairing. The top strand is labeled "String" and the bottom strand is labeled "Reverse complement".

Large amount of data

NOT IN THE LECTURE

	Genome length	Total bases at 30x coverage	Size if each base takes 2 bits
E. coli	$4.6 \cdot 10^6$	$138 \cdot 10^6$	34 MBytes
Human	$3.2 \cdot 10^9$	$96 \cdot 10^9$	24 GBytes
Spruce	$25 \cdot 10^9$	$750 \cdot 10^9$	187.5 GBytes
Axolotl	$32 \cdot 10^9$	$960 \cdot 10^9$	240 GBytes

Theoretical problem formulations

("Classical" computational formulations of how to
obtain the output from the input)

Shortest superstring

INPUT: A collection of strings (the reads)

OUTPUT: A string S such that every given string is a substring of S (S is a *superstring*),
and S is shortest

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OUTPUT: A string S such that every given string is a substring of S (S is a *superstring*), and S is shortest

TAGA
ATAG
CATA
TCAT

Input

TCATAGA

Output S

Shortest superstring

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OUTPUT: A string S such that every given string is a substring of S (S is a *superstring*), and S is shortest

TAGA
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TAGA	TAGA
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CATA	CATA
TCAT	TCAT
	TCATAGA
Input	Output S

- NP-hard to compute (i.e. it cannot be solved efficiently)
- Not practical: it collapses repeats (main drawback)

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Input	Output S

- NP-hard to compute (i.e. it cannot be solved efficiently)
- Not practical: it collapses repeats (main drawback)

The genome [**TCATATATAGA**

The reads [TCAT
CATA
ATAT
TATA
ATAT
TATA
ATAG
TAGA

Shortest superstring

INPUT: A collection of strings (the reads)

OUTPUT: A string S such that every given string is a substring of S (S is a *superstring*), and S is shortest

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The genome	[TCATATATAGA	The shortest superstring S	[TCATATATAGA
The reads	[TCAT	The input	[TCAT
	CATA		CATA
	ATAT		ATAT
	TATA		TATA
	ATAT		ATAT
	TATA		TATA
	ATAG		ATAG
	TAGA		TAGA

Overlap graphs + Hamiltonian path

ACTAGACTAGACC

ACTA
CTAG
TAGA
AGAC
GACT
ACTA
CTAG
TAGA
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Overlap graphs + Hamiltonian path

ACTAGACTAGACC

ACTA
CTAG
TAGA
AGAC
GACT
ACTA
CTAG
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AGAC
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INPUT: Overlap graph of order t :

- Every read is a node
- Every suffix-prefix overlap of length $\geq t$ is an edge

OUTPUT: A path going through every node (i.e. read) exactly one (*Hamiltonian*)

Overlap graphs + Hamiltonian path

ACTAGACTAGACC
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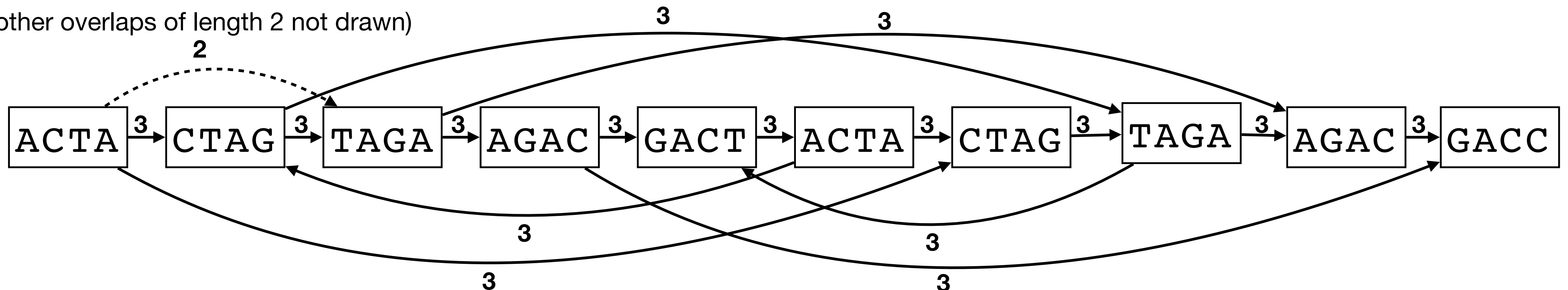
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Overlap graph of order 2:

(other overlaps of length 2 not drawn)



Overlap graphs + Hamiltonian path

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TAGA
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ACTA
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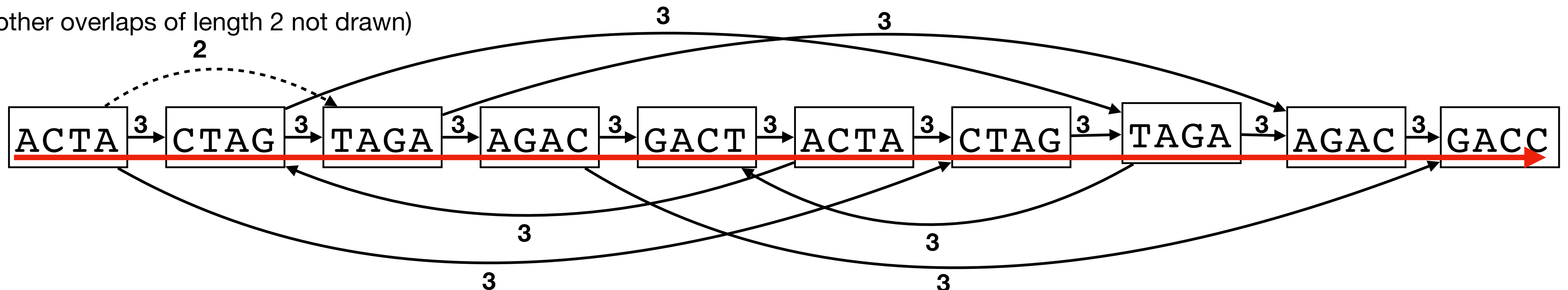
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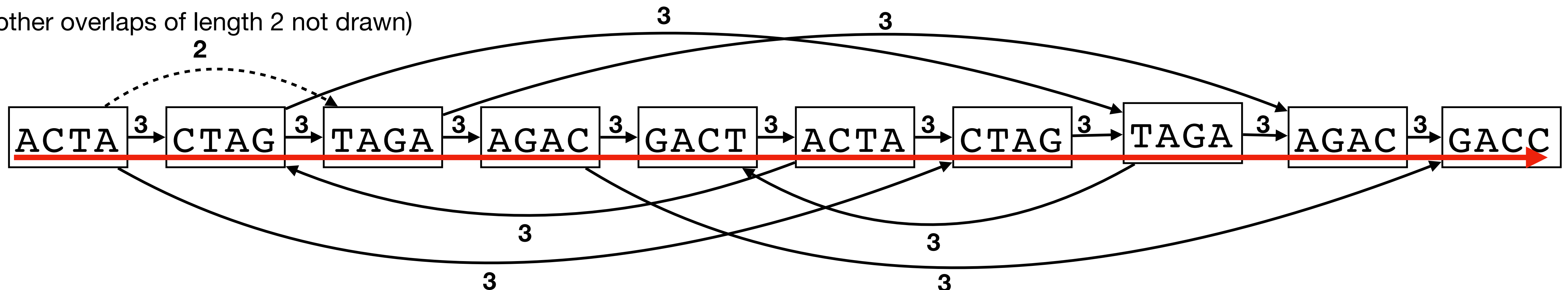
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OUTPUT: A path going through every node (i.e. read) exactly one (*Hamiltonian*)

- NP-hard to compute
- Not practical: usually graph has no Hamiltonian path (missing coverage, errors)

Overlap graph of order 2:

(other overlaps of length 2 not drawn)



De Bruijn graphs + Eulerian path



De Bruijn graphs + Eulerian path



INPUT: De Bruijn graph of order k :

- Every k -mer (substring of length k) in the reads is a ***single*** node
- Every $(k + 1)$ -mer is a ***different*** arc from its length- k prefix to its length- k suffix

ASSUMPTION: Every length- k interval of the genome appears exactly the same number of times in the reads (*uniform coverage*)

OUTPUT: A path going through every ***edge*** (i.e. k -mer) exactly one (*Eulerian*)

De Bruijn graphs + Eulerian path

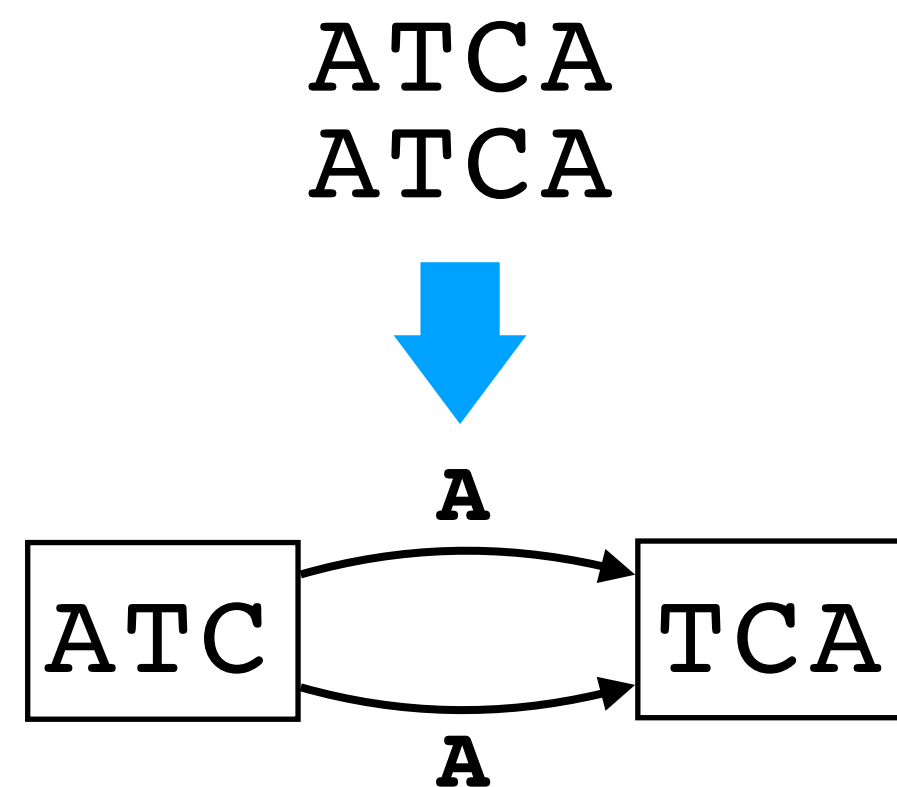


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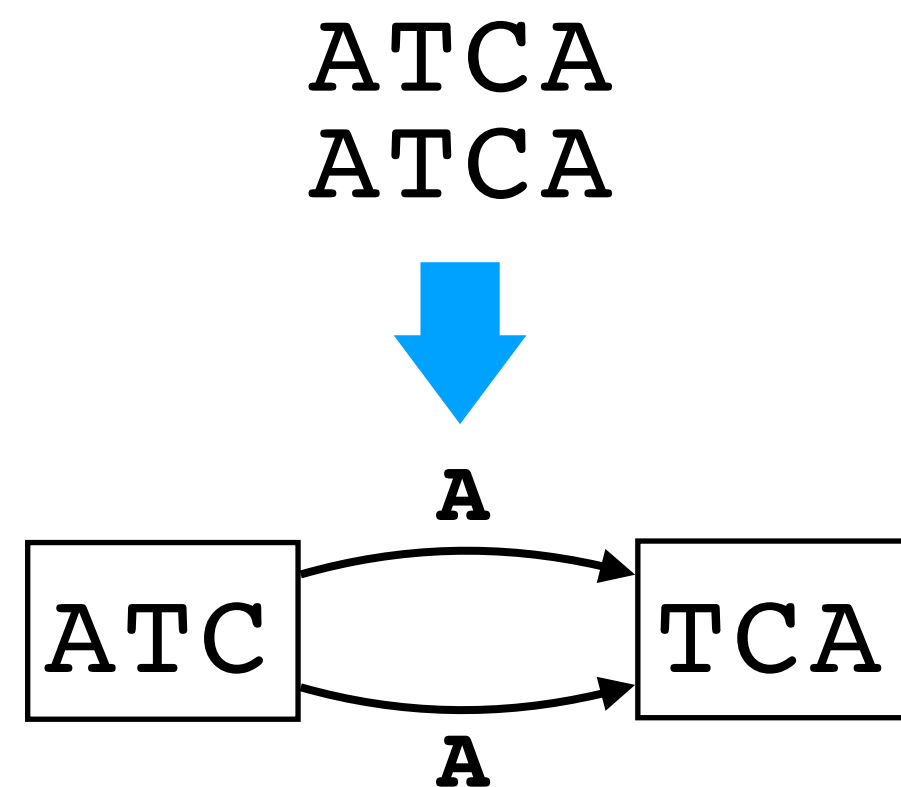
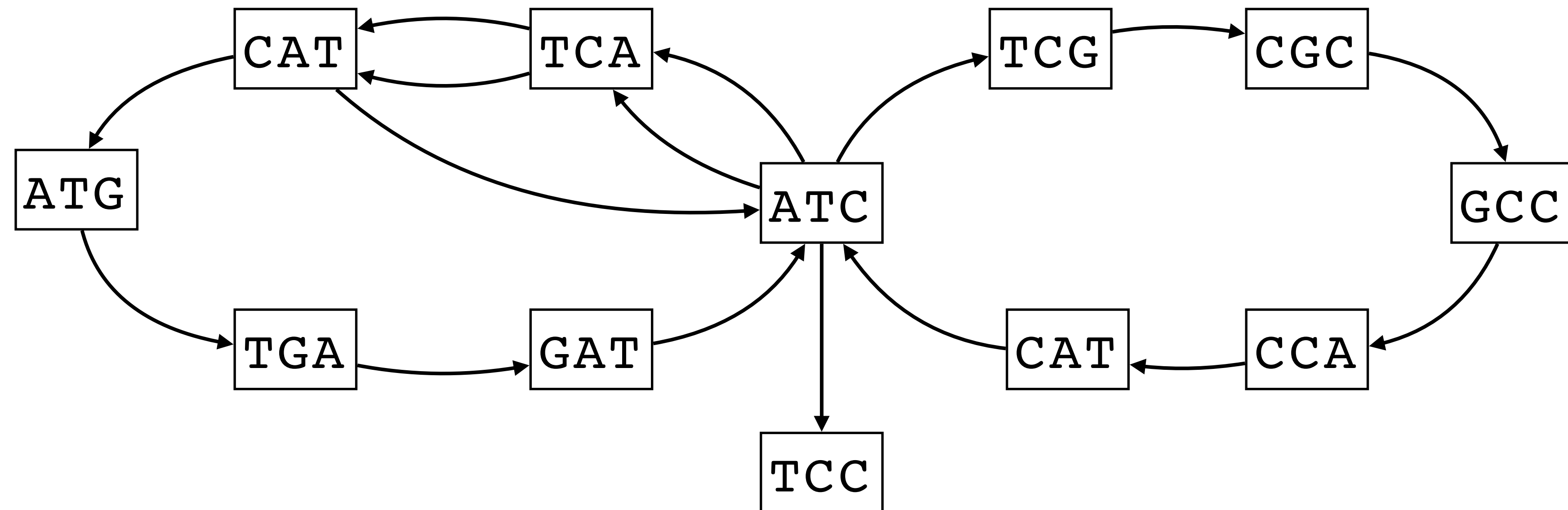
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De Bruijn graph of order 3:



De Bruijn graphs + Eulerian path



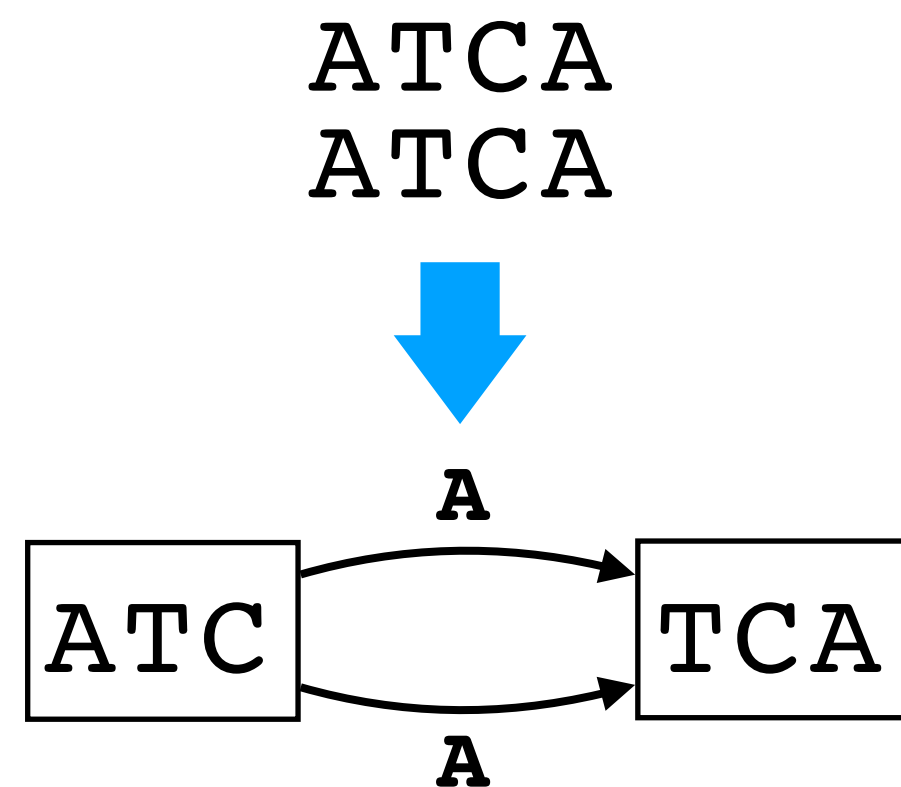
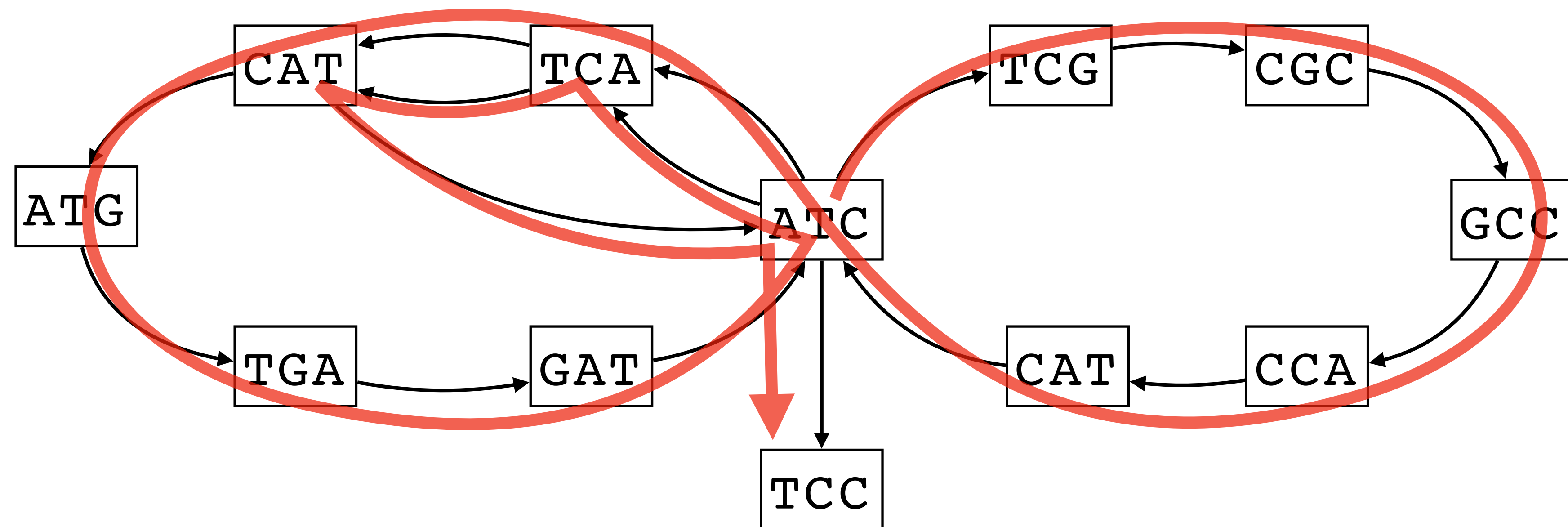
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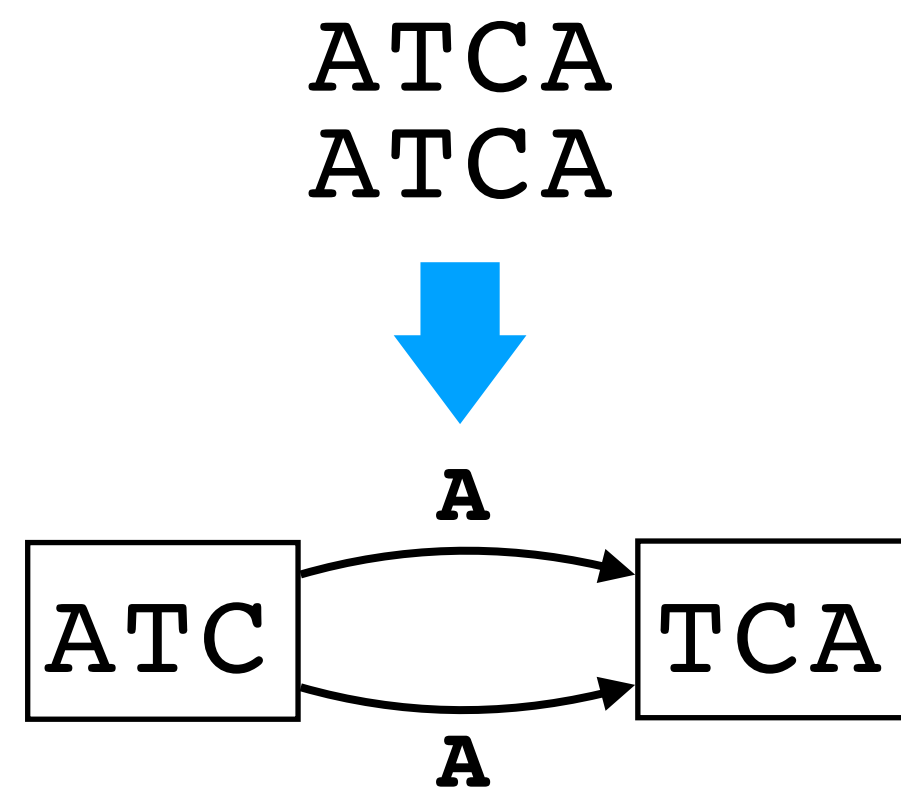
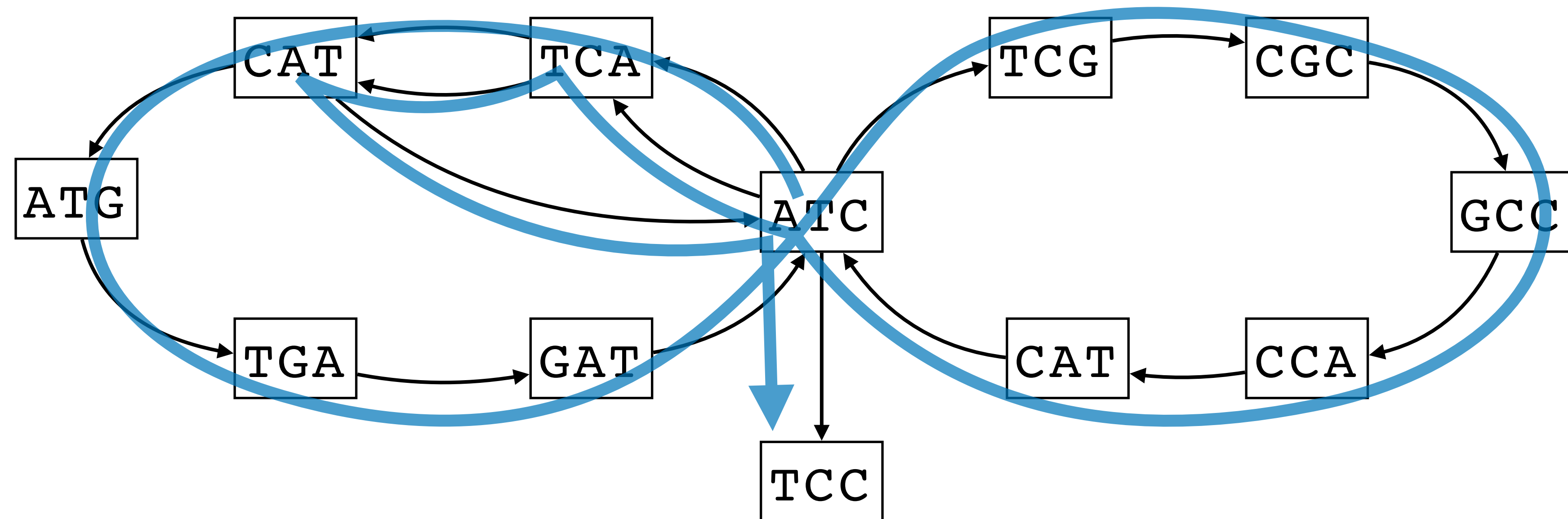
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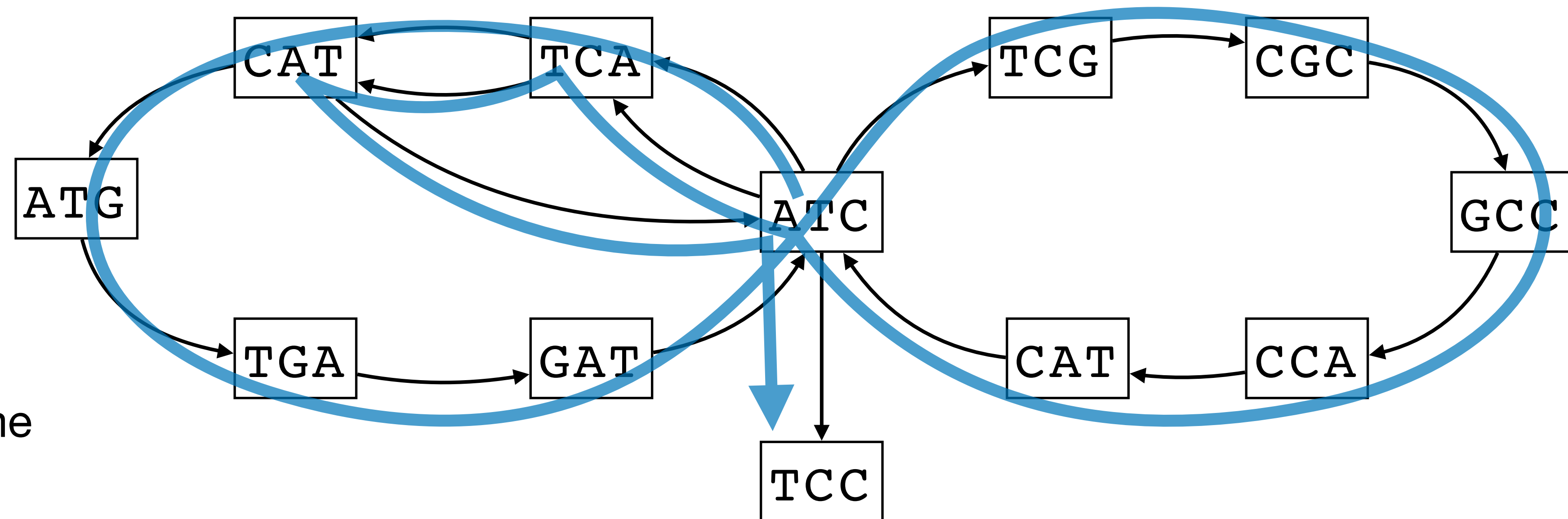
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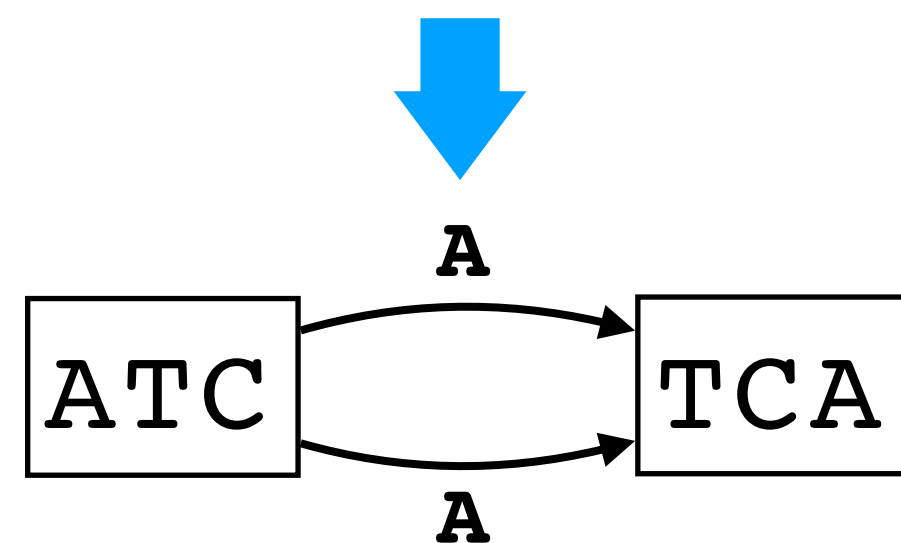
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De Bruijn graph of order 3:



ATCA
ATCA



- Can be solved in $O(|edges|)$ time
- Too restrictive assumption

Section summary

- Modeling the genome assembly problem evolved over time (and still does)

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- Modeling the genome assembly problem evolved over time (and still does)
- Computational complexity ranges from NP-hard to linear
- Not robust to practical issues, and hard to integrate them into the formulations
- Most importantly (even if the above are solved):
 - Many solutions, which one is the true genome?
 - We will see a different "theoretical" approach, closer to practice

AT END OF LECTURE

Practical genome assembly

(The sequence of algorithmic steps behind
"real" genome assemblers)

Contig assembly

Focus on single-end reads
Assembling a full genomic
sequence in one shot is
hopeless



Forget about the assembly
model (i.e. the problem
formulation)
Assemble only parts about
which we are sure (contigs →
contiguous sequences)

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- Assume an assembly graph (here de Bruijn with parallel edges collapsed into one)
- Focus on ***unitigs***: non-branching paths
- Unitig =_{def} "non-branching" path
- (Usually) Contig =_{def} unitig in a graph "corrected" for polyploidy

Unitigs

- Let $P = (v_1, v_2, \dots, v_{t-1}, v_t)$ be a path. We say that P is a ***unitig*** if either:

Unitigs

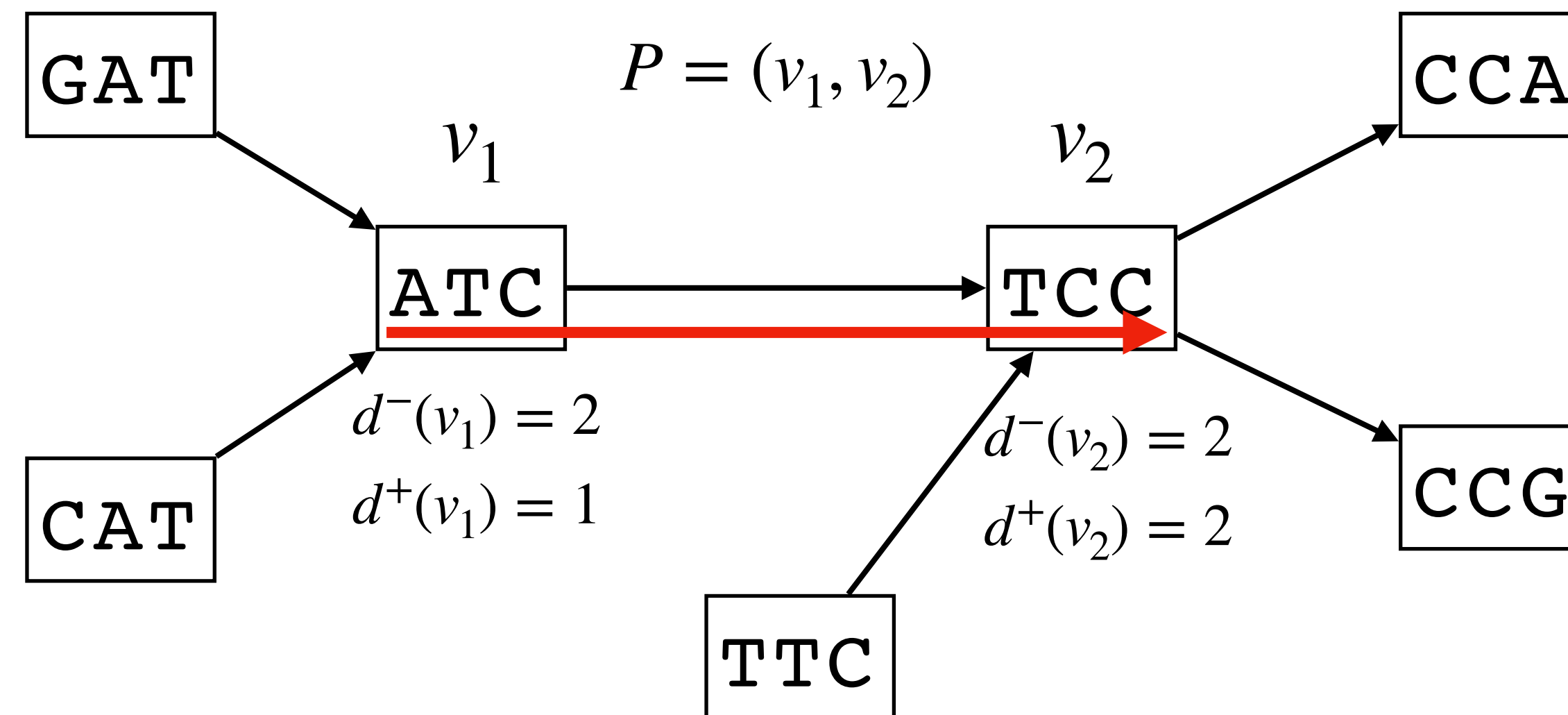
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(every internal node has exactly one in-neighbor and one out-neighbor)

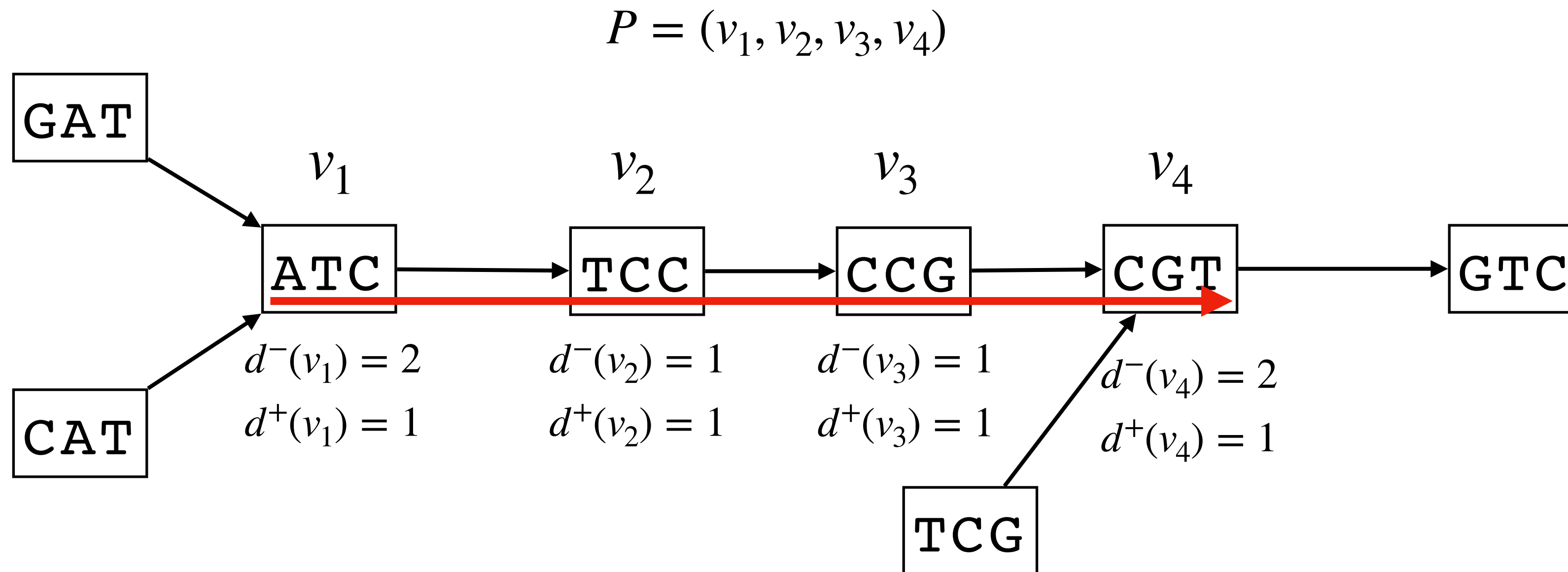
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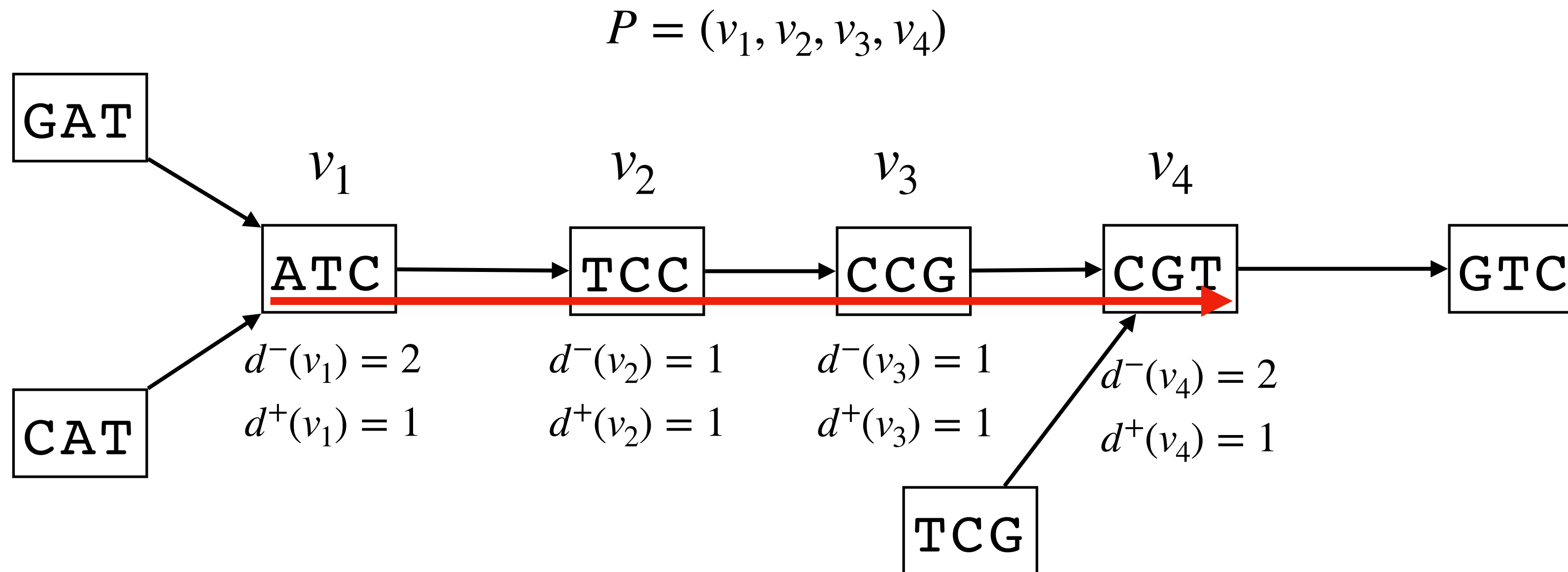
Unitigs

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 - $t = 2$, that is, $P = (v_1, v_2)$ is a single edge (we "trust" edges), or
 - for every $i \in \{2, \dots, t-1\}$, we have $d^-(v_i) = d^+(v_i) = 1$,
(every internal node has exactly one in-neighbor and one out-neighbor)



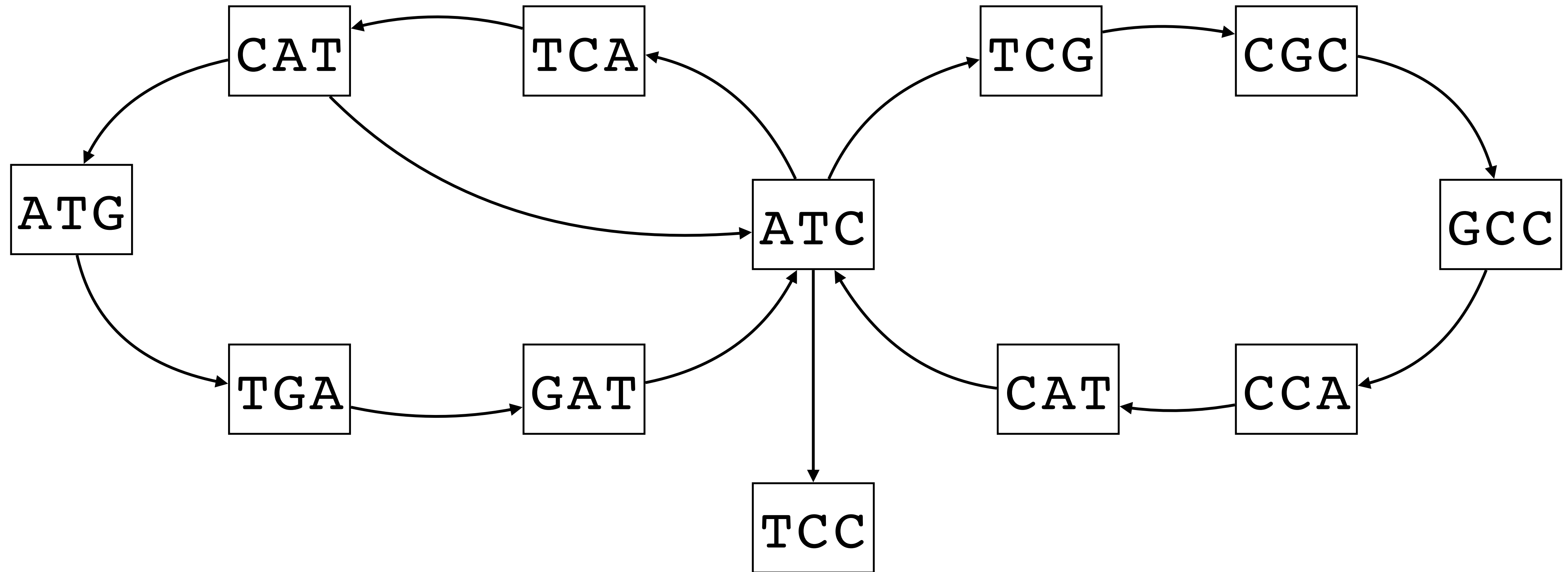
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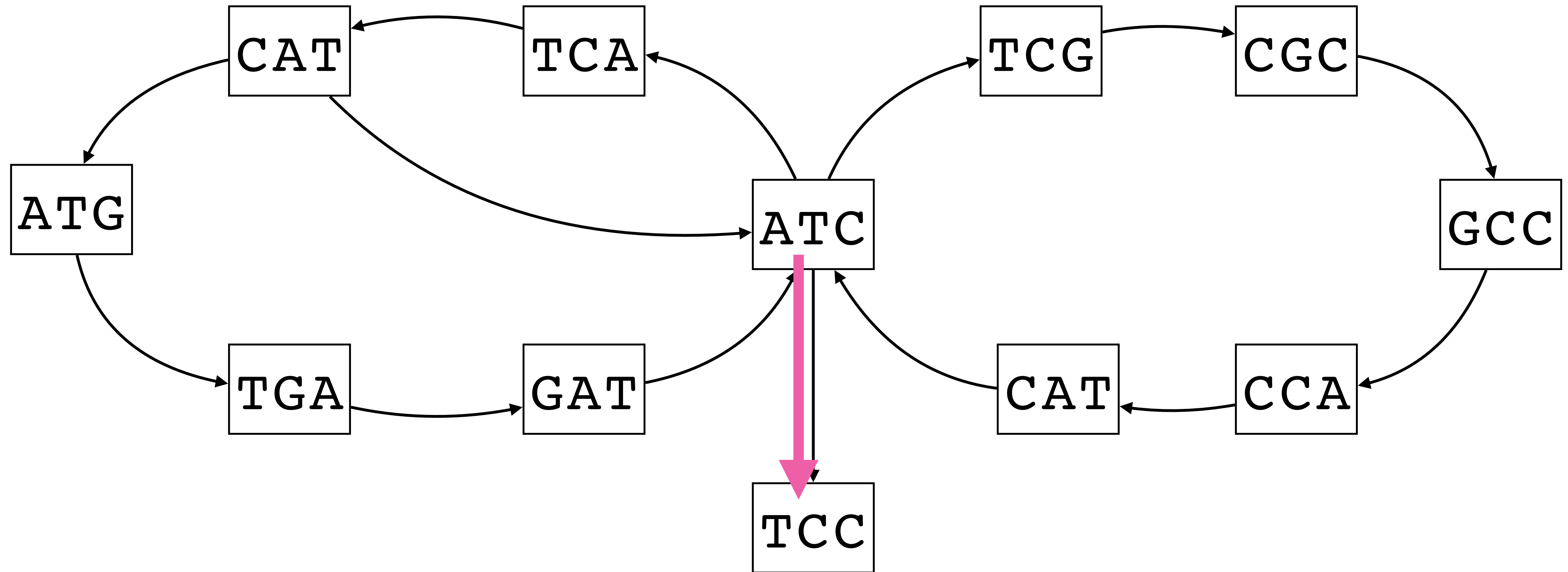


- We want maximal (longest) unitigs

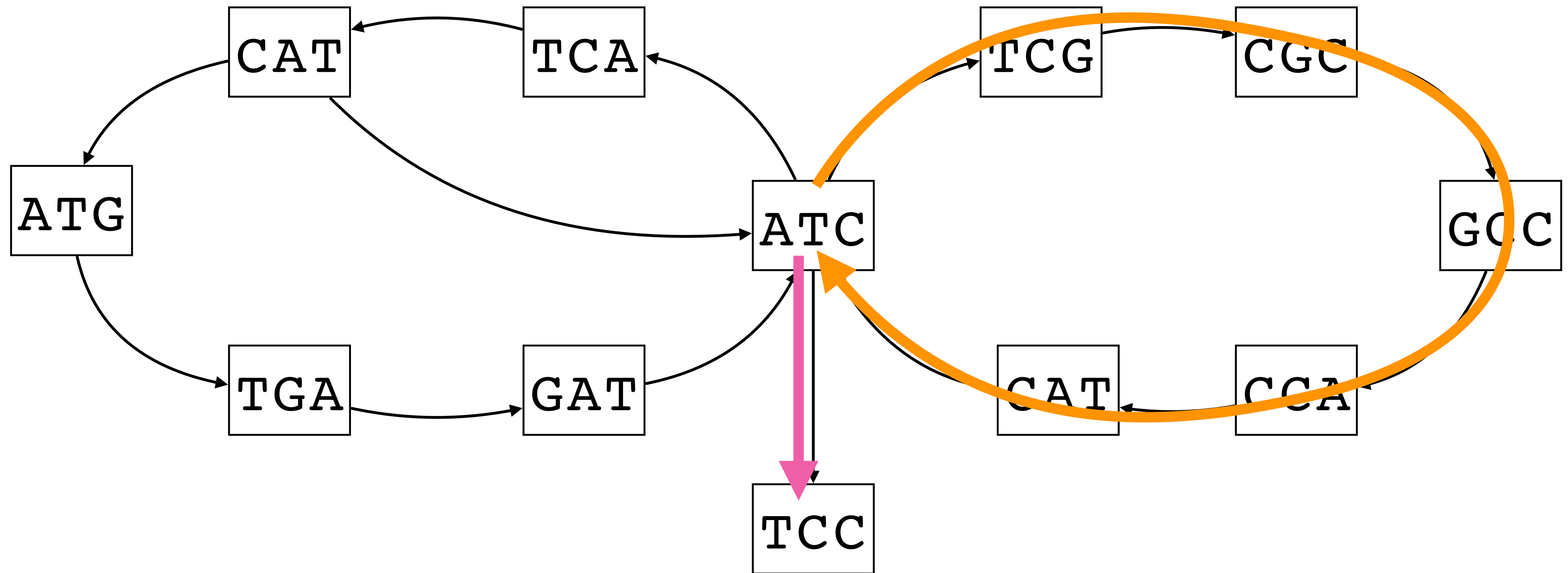
Example



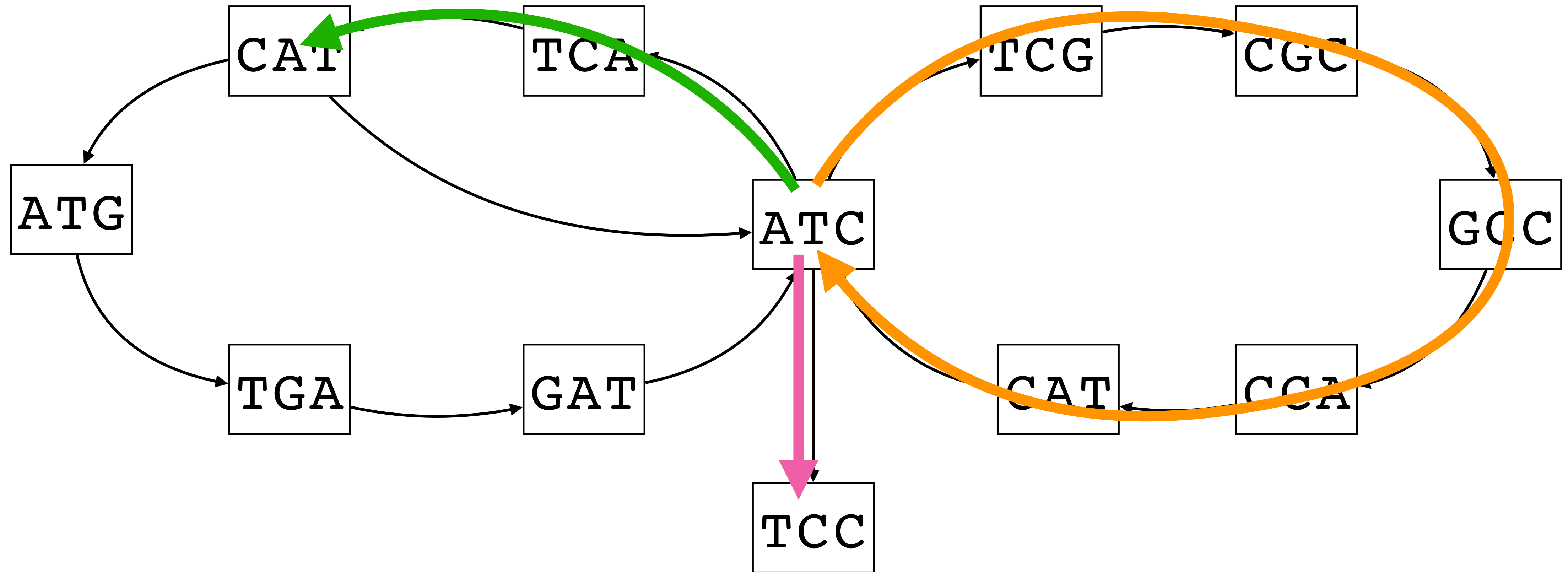
Example



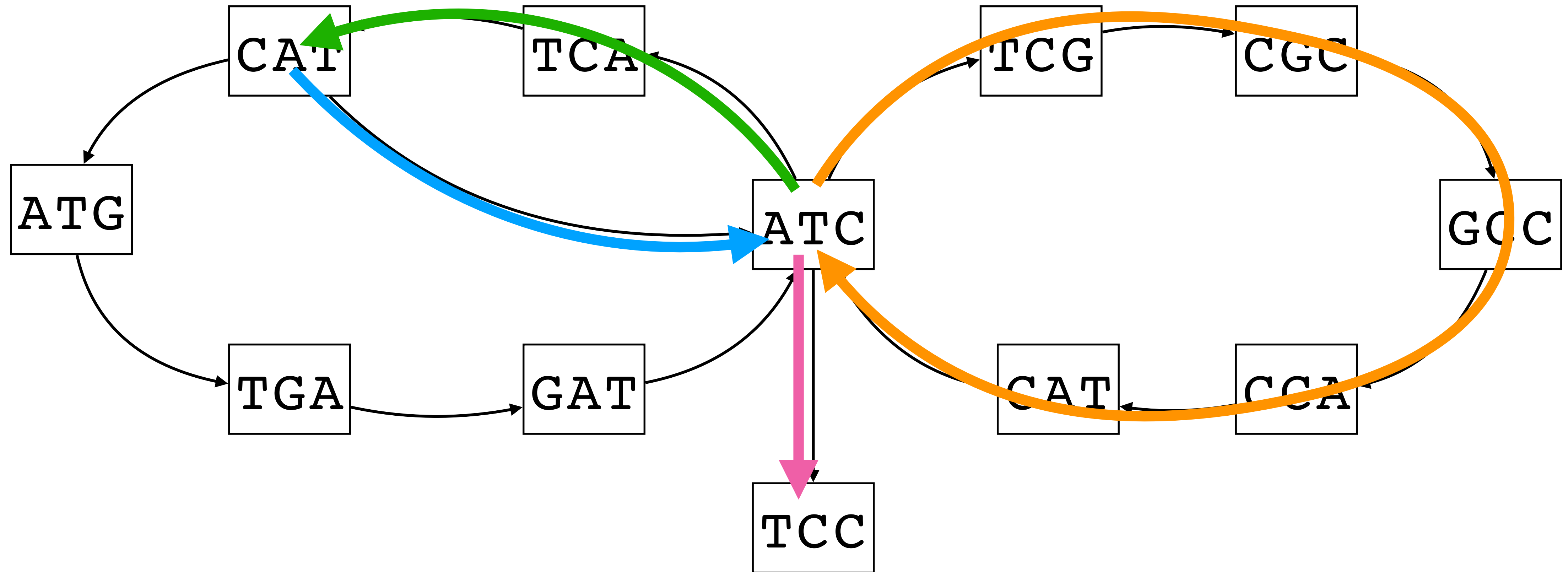
Example



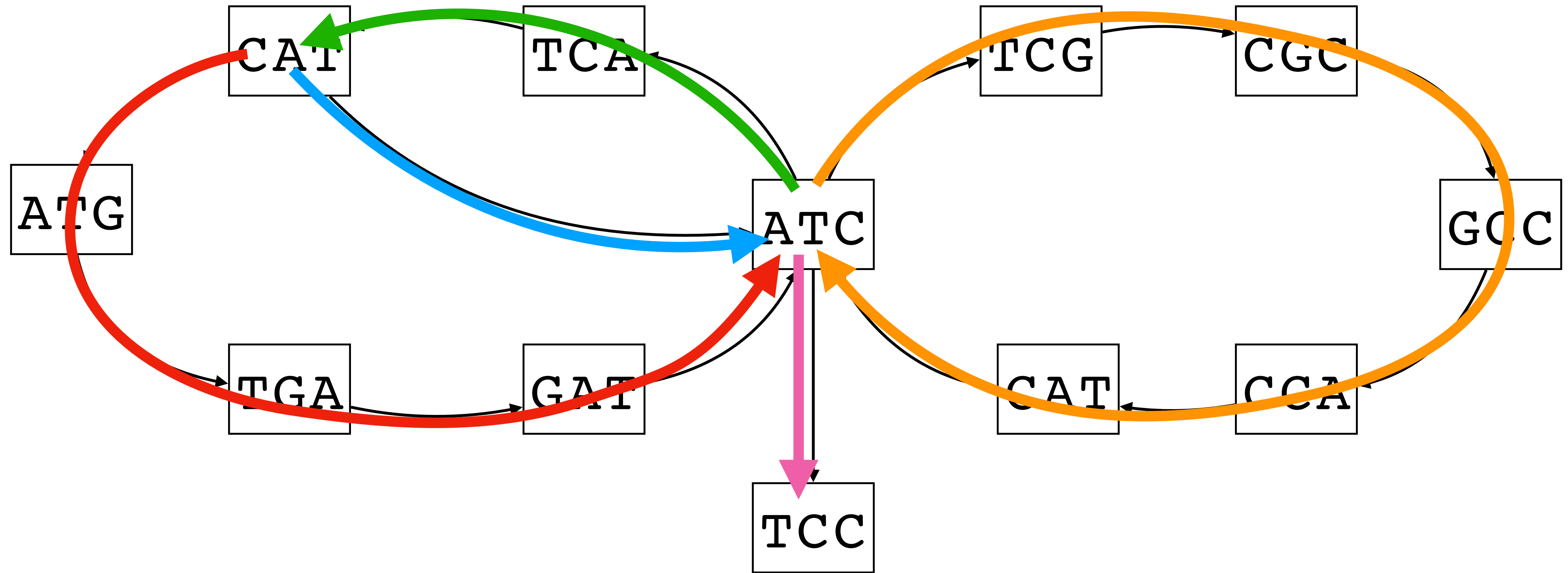
Example



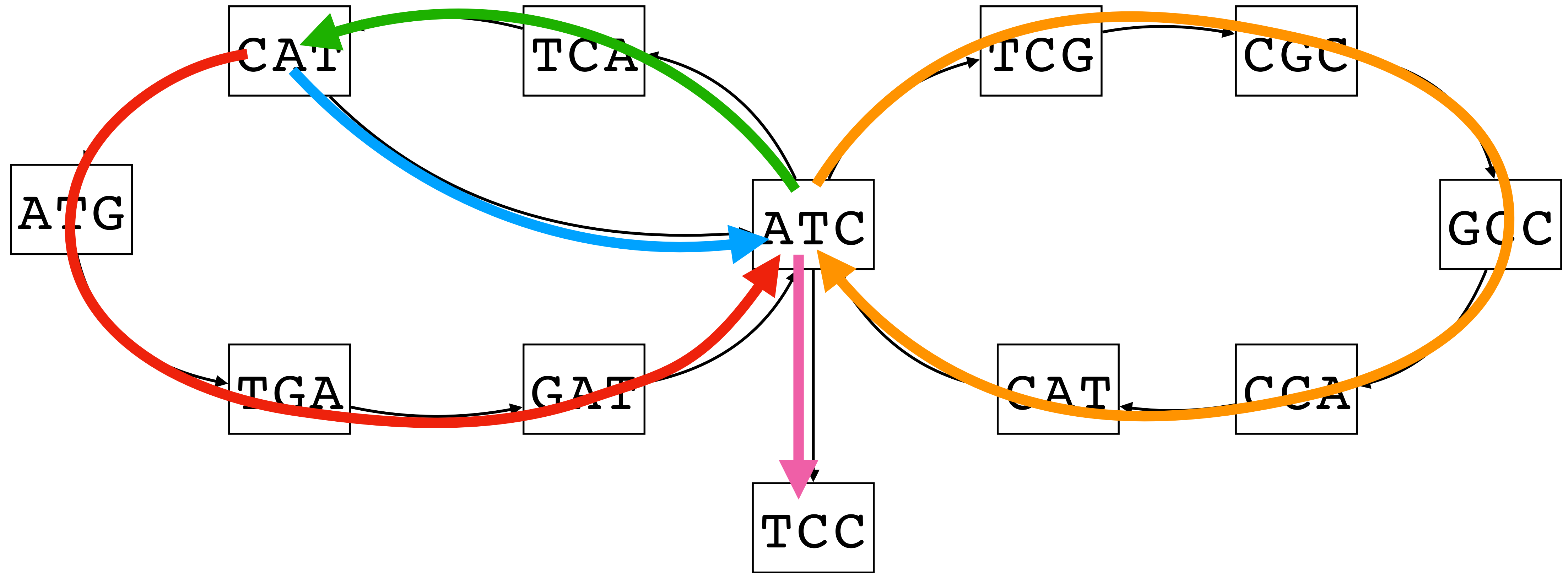
Example



Example



Example



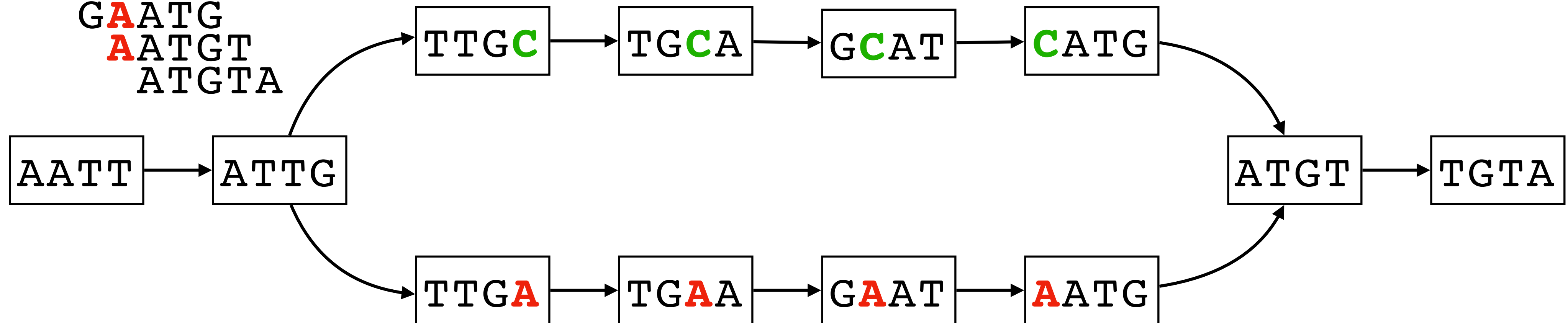
- Unitigs can be found in $O(|\text{edges}|)$ time

Bubble popping

ATGTA
CATGT
GCATG
TGCAT
TTGCA
ATTGC
AATTG
Mother **AATTGCATGTA**
Father **AATTGAATGTA**
AATTG
ATTG**A**
TTG**A**A
TG**A**AT
G**A**ATG
AATGT
ATGTA

Bubble popping

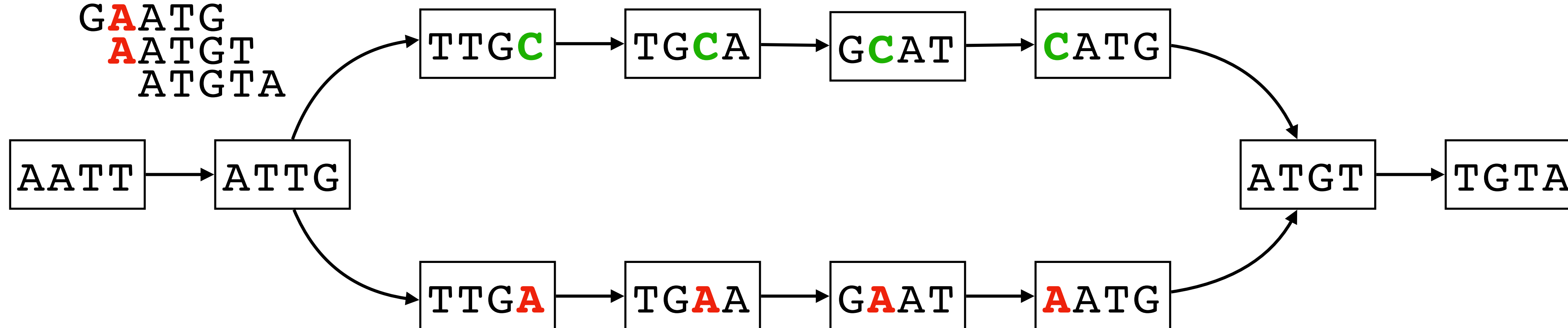
ATGTA
CATGT
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ATTGC
AATTG
Mother AATTGCATGTA
Father AATTGAATGTA
AATTG
ATTGA
TTGAA
TGAAAT
GAATG
AATGT
ATGTA



Bubble popping

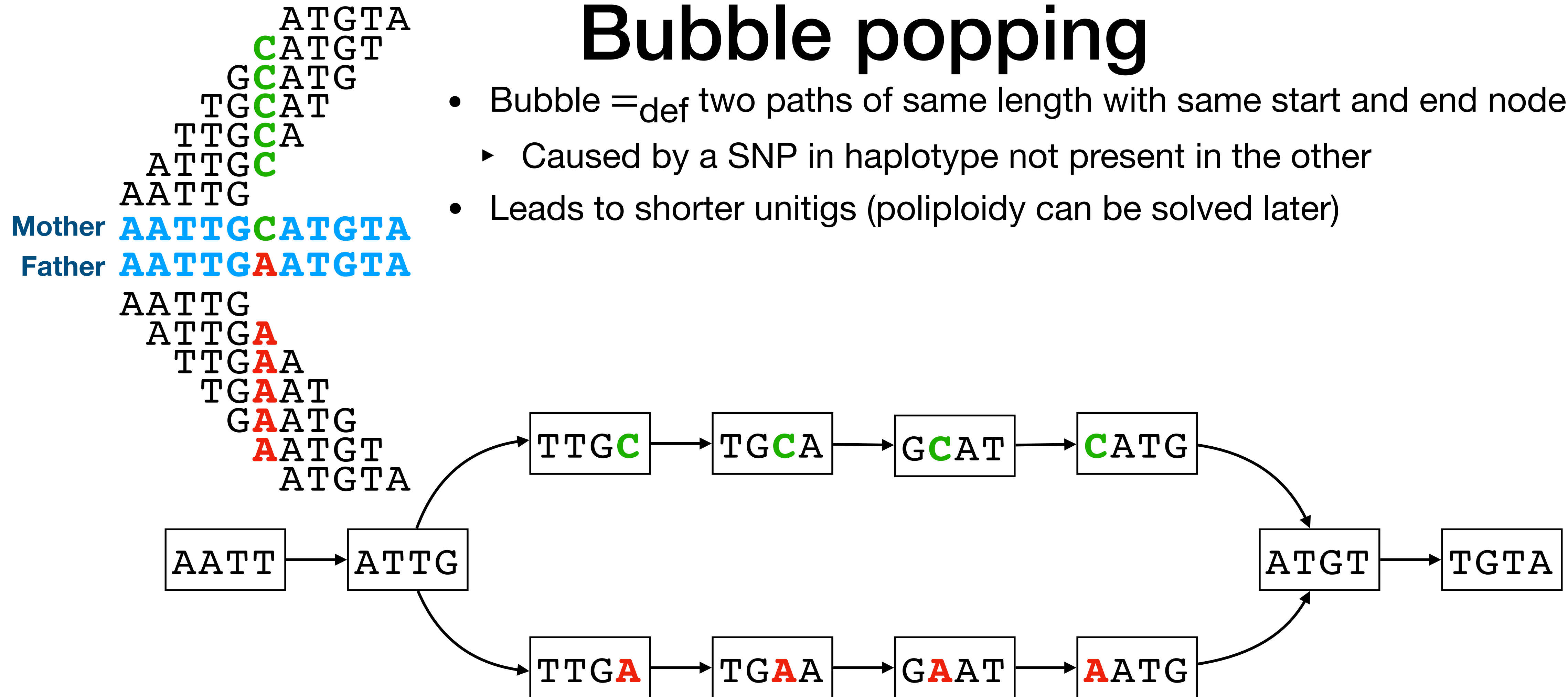
- Bubble =_{def} two paths of same length with same start and end node
 - Caused by a SNP in haplotype not present in the other

ATGTA
CATGT
GCATG
TGCA
TTGC
ATTG
AATTG
Mother **AATTGCATGTA**
Father **AATTGAATGTA**
AATTG
ATTG**A**
TTG**AA**
TG**AA**T
G**A**ATG
AATGT
ATGTA



Bubble popping

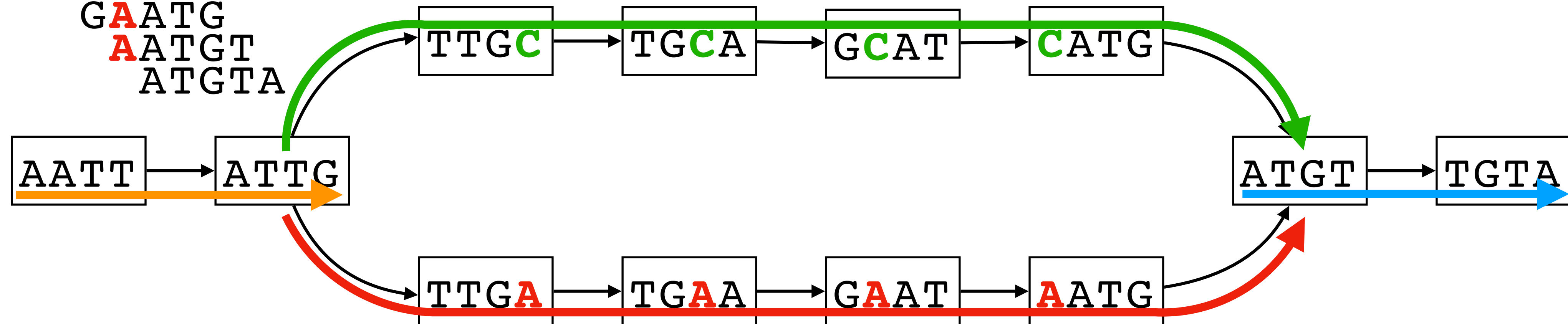
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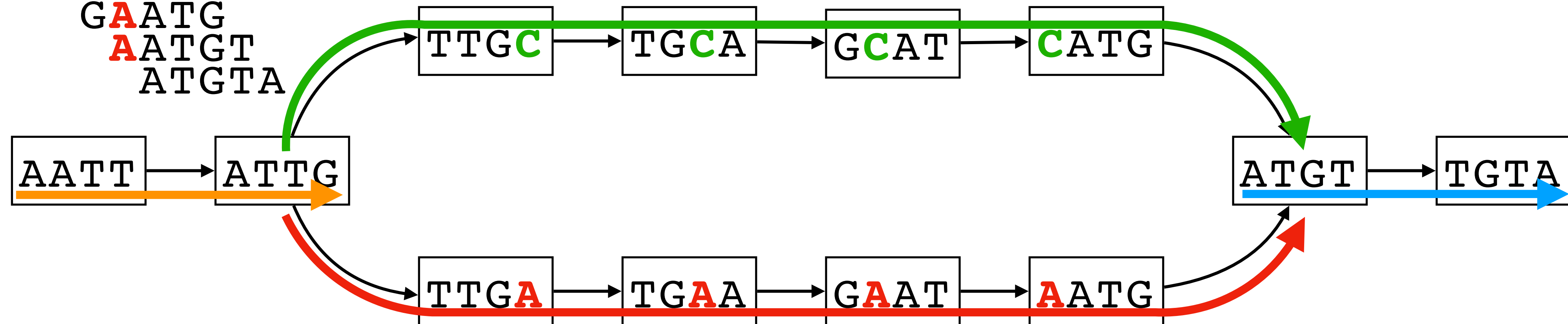
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- Generalizations possible (e.g. *superbubble*), handling more complex scenarios
 - Taku Onodera, Kunihiro Sadakane, Tetsuo Shibuya:
Detecting Superbubbles in Assembly Graphs. WABI 2013: 338-348

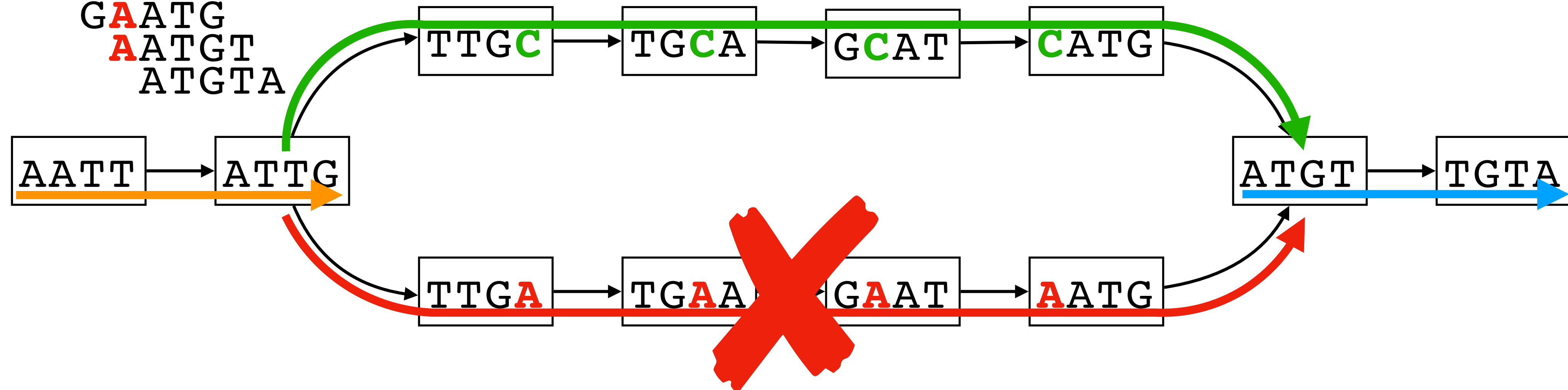
ATGTA
 CATGT
 GCATG
 TGCAT
 TTGCA
 ATTGC
 AATTG
 Mother **AATTGCATGTA**
 Father **AATTGAATGTA**
 AATTG
 ATTGA
 TTGAA
 TGAAT
 GAATG
 AATGT
 ATGTA



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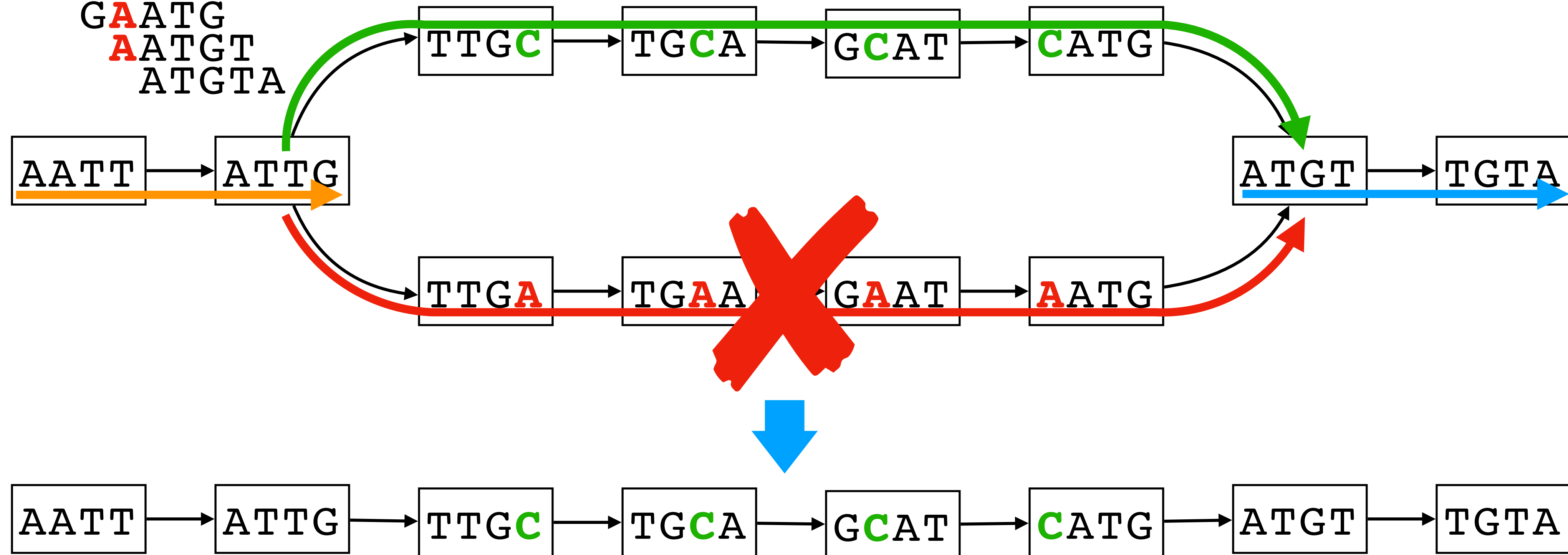
ATGTA
CATGT
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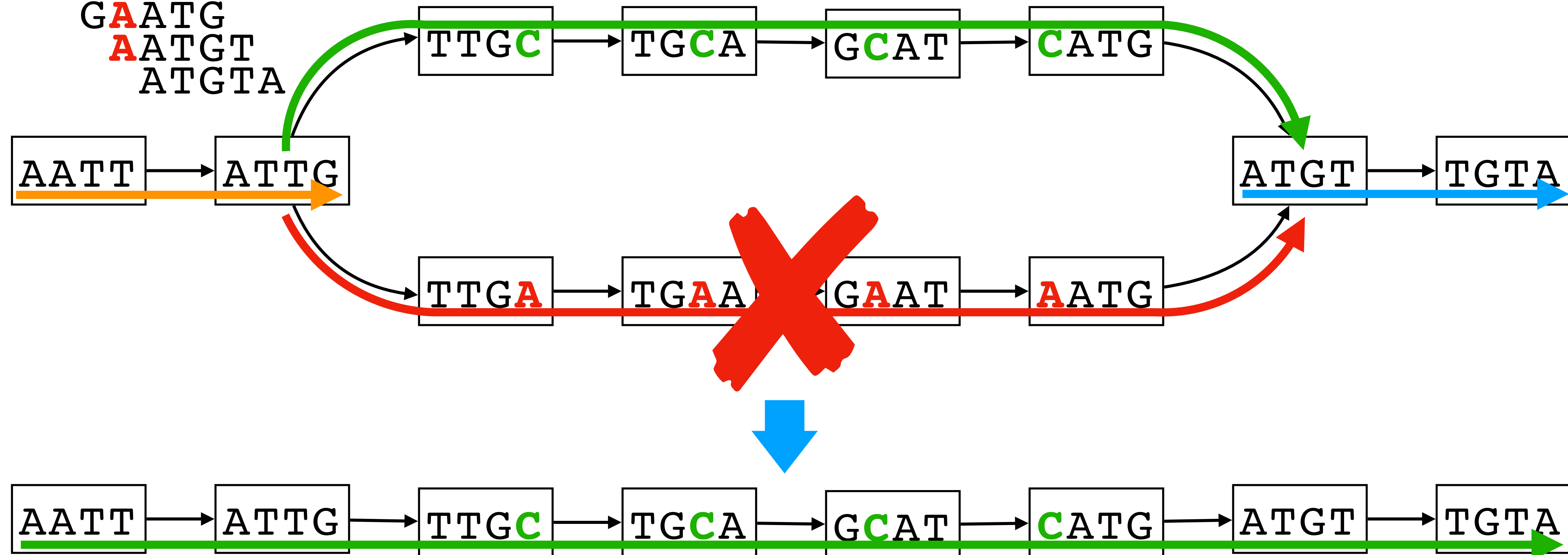
ATGTA
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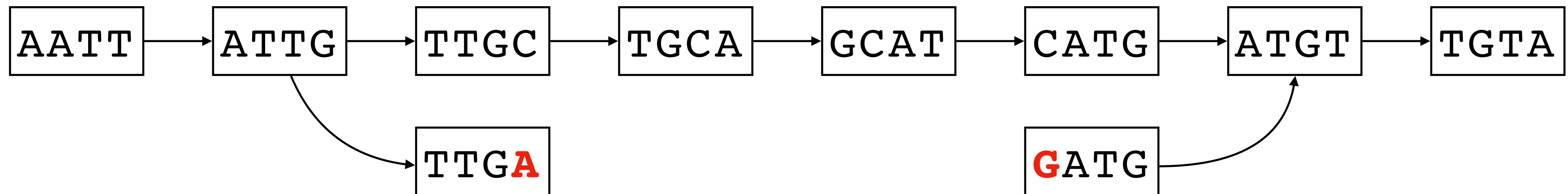


Tip removal

ATGTA
GATGT
GCATG
TGCAT
TTGCA
ATTGC
ATTGA
AATTG
AATTGCATGTA

Tip removal

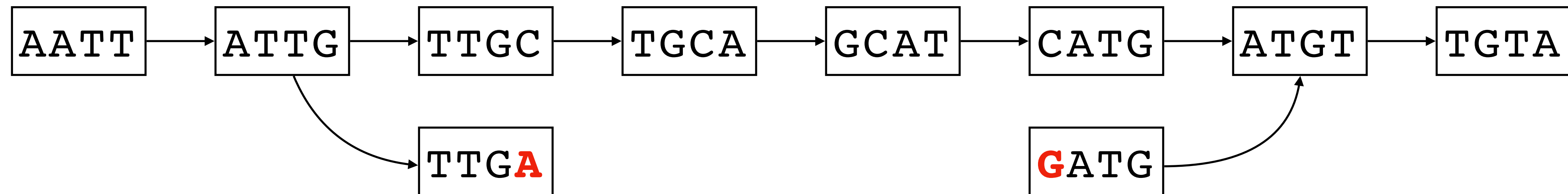
ATGTA
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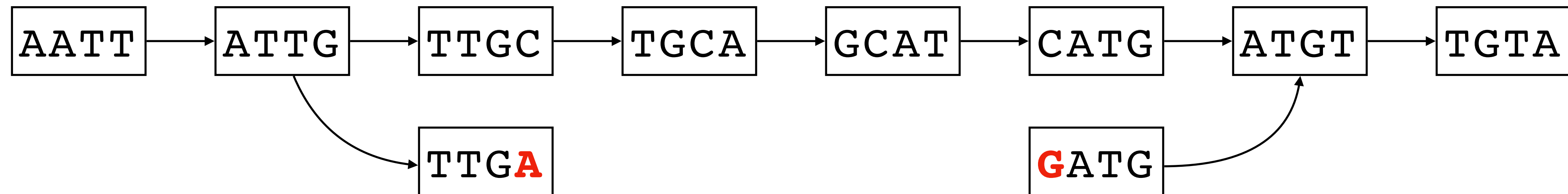
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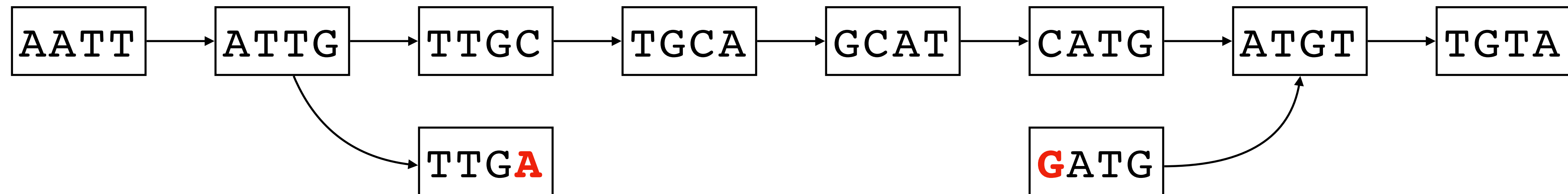
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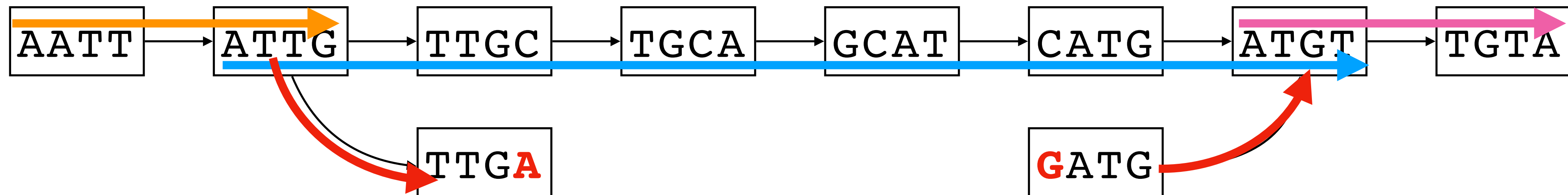
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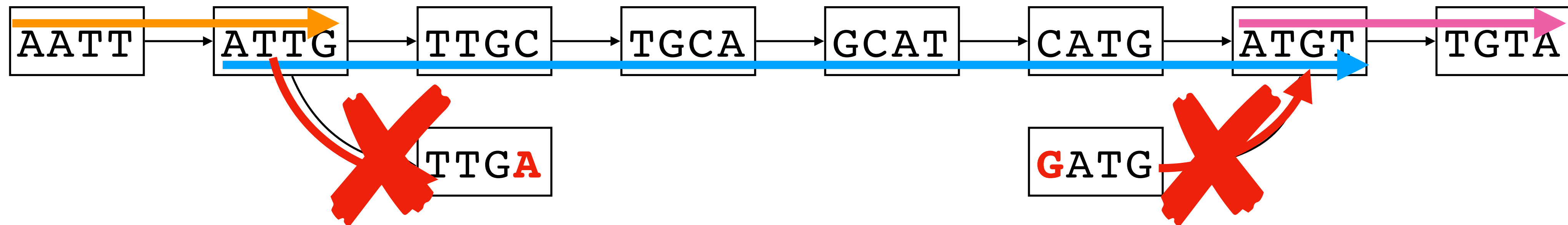
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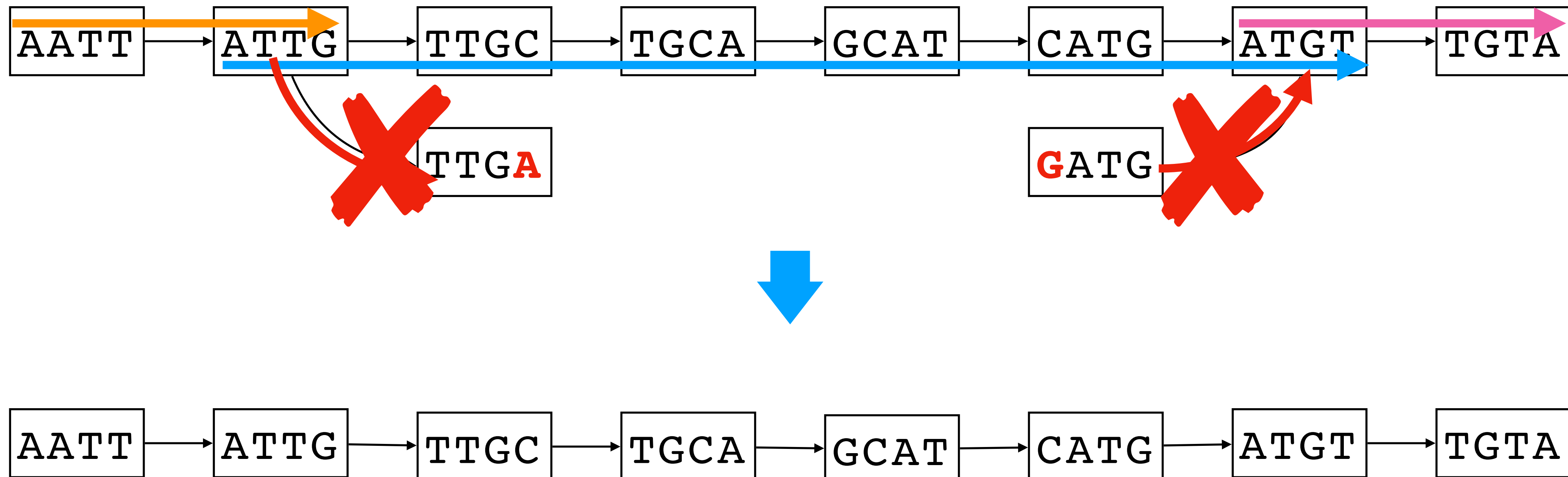
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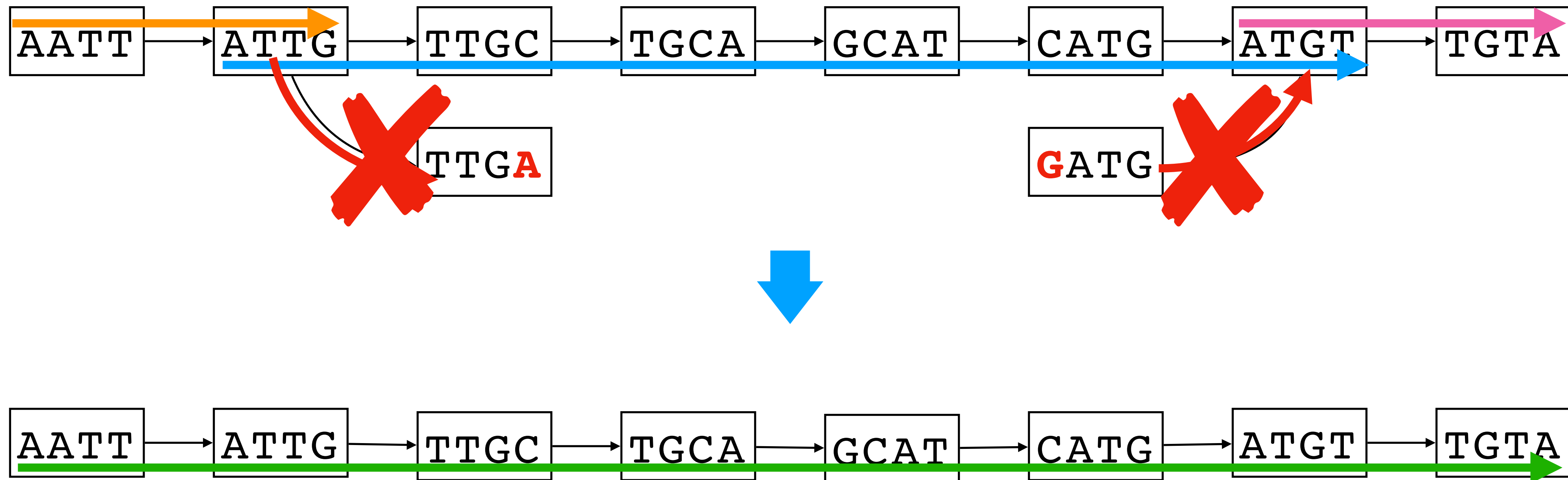
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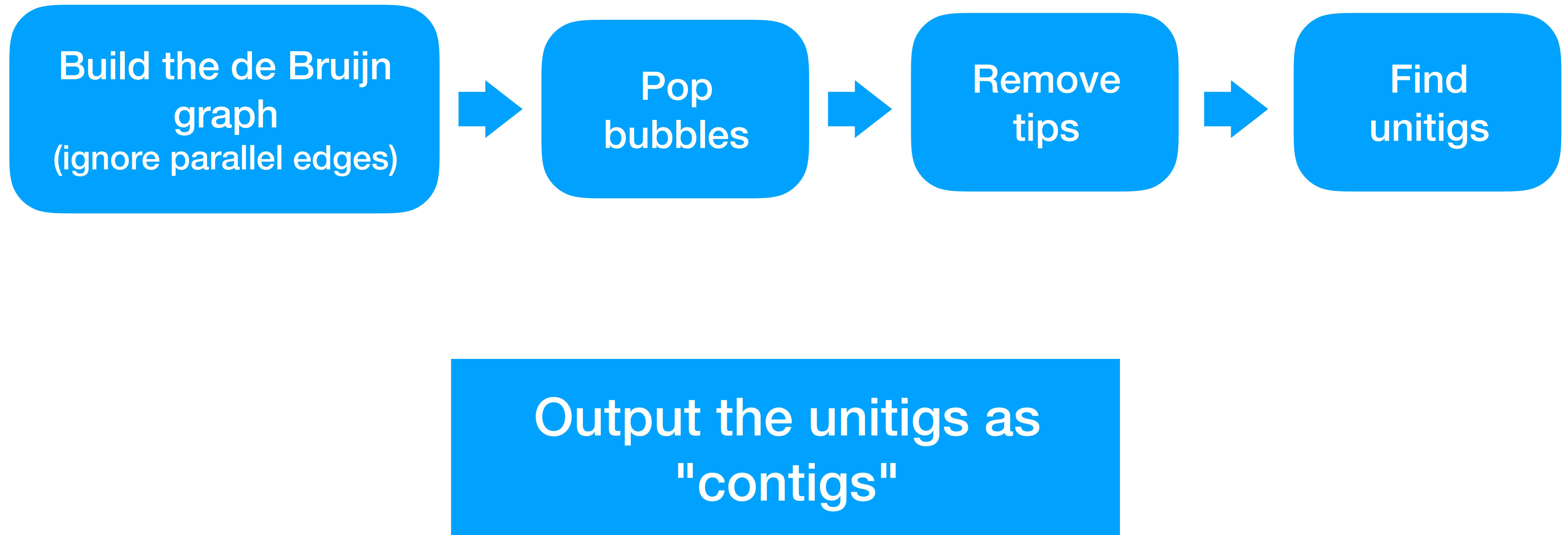
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Contigs assembly

(simplified to ignore some practical issues e.g. errors, reverse complements)



Scaffolding

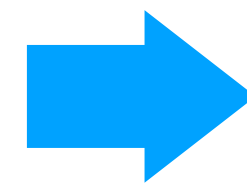
Bring in paired-end
information
Align reads to contigs



Chain (order) the contigs
Output chains of contigs
with "gaps" (NNNN...)
between them

Contigs

TCGATAGCTAAAA
AATTGT
ATAGAGATATTT
ATATCGCTAGA



Scaffolds

TCGATAGCTAAAA**NNNNNNNNNN**AATTGT**NNN**ATAGAGATATTT
ATATCGCTAGA

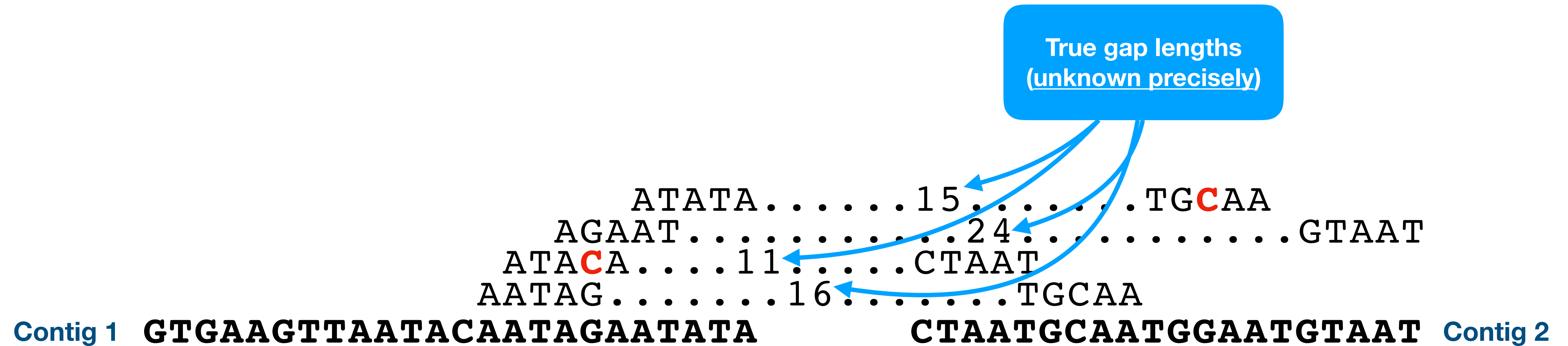
Align paired-end reads to contigs

```
          ATATA . . . . . 15 . . . . . TGCAA
        AGAAT . . . . . 24 . . . . . GTAAT
      ATACA . . . . 11 . . . . CTAAT
    AATAG . . . . . 16 . . . . . TGCAA
```

Contig 1 GTGAAGTTAATACAATAGAATATA CTAATGCAATGGAATGTAAT **Contig 2**

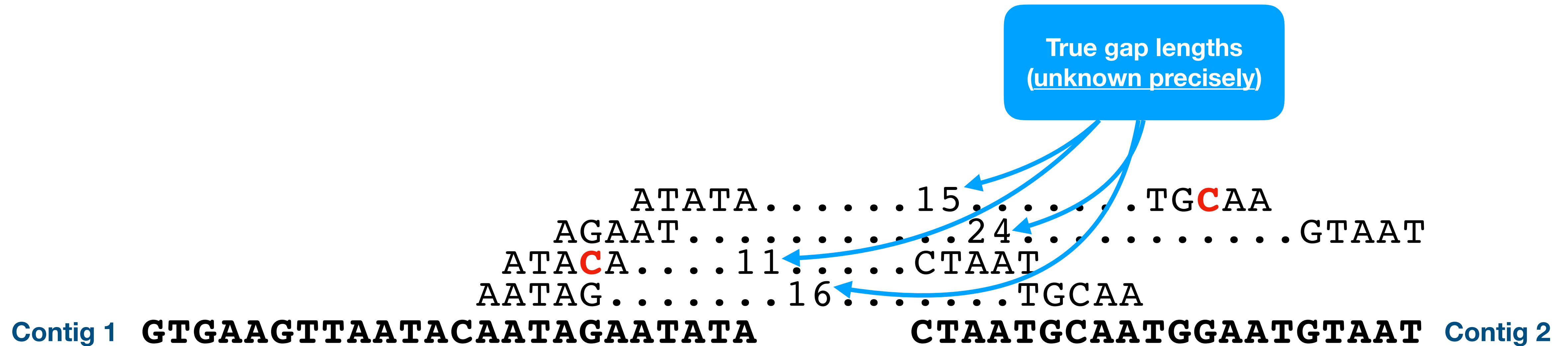
- Align paired-end reads to contigs, focus on read pairs aligning to different contigs

Align paired-end reads to contigs



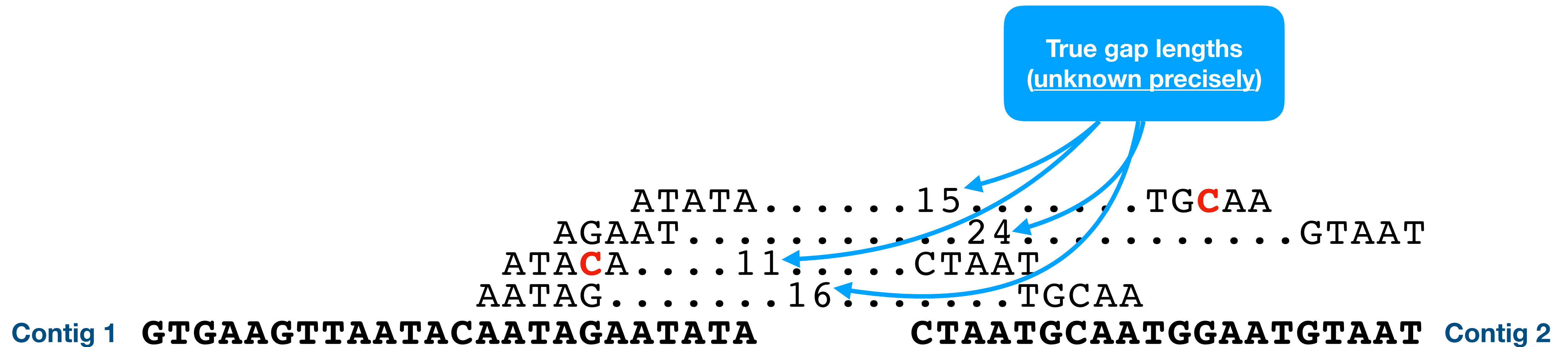
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Align paired-end reads to contigs



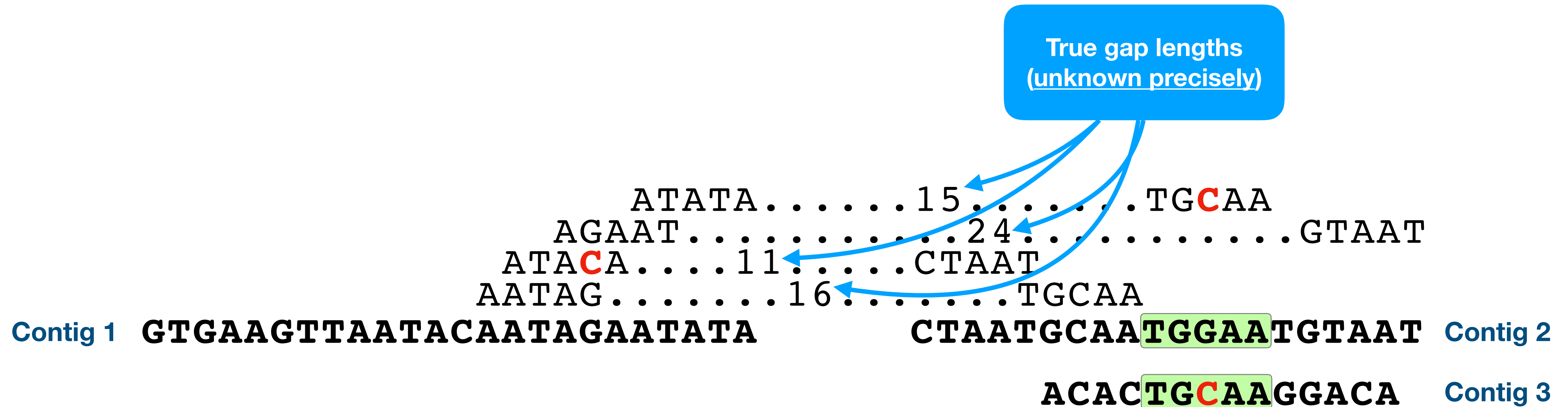
- Align paired-end reads to contigs, focus on read pairs aligning to different contigs
 - Even if we know two contigs are "consecutive" in a genome, it is not easy to estimate the gap length between them, see e.g.

Align paired-end reads to contigs



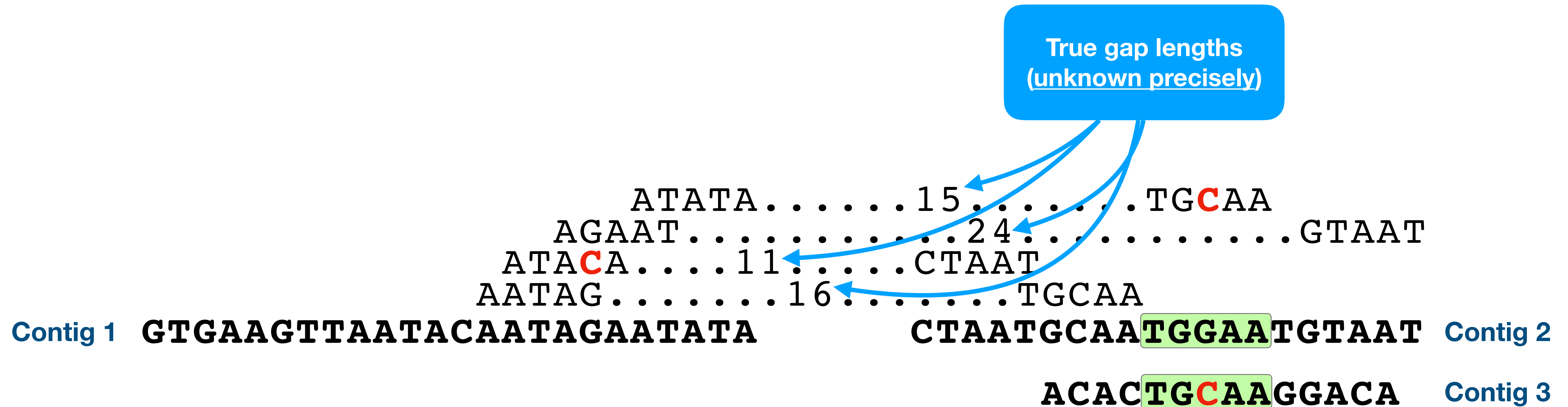
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Align paired-end reads to contigs



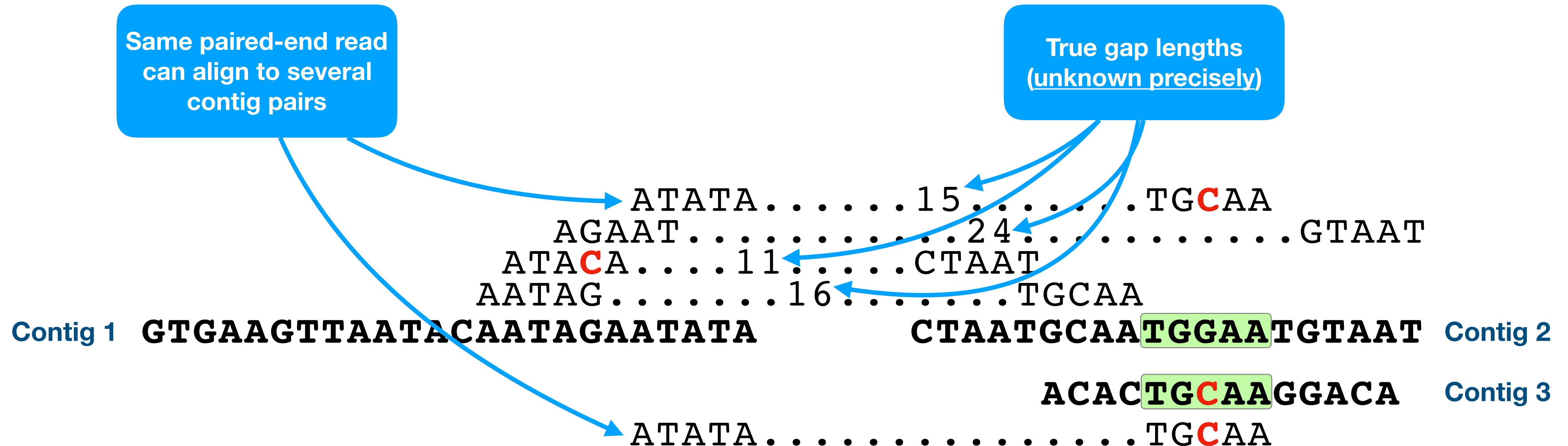
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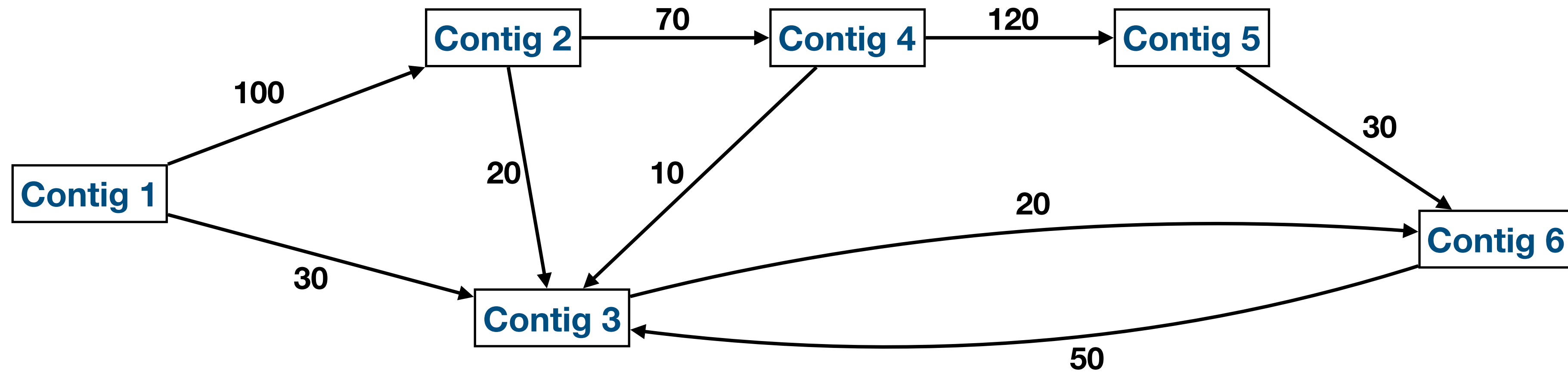
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Align paired-end reads to contigs

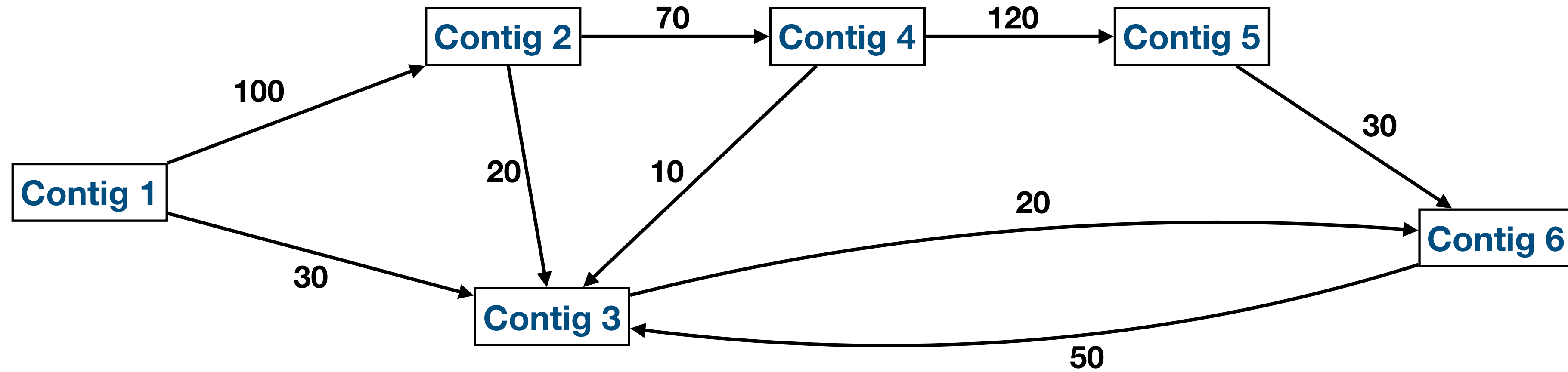


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Chaining (scaffolding) contigs

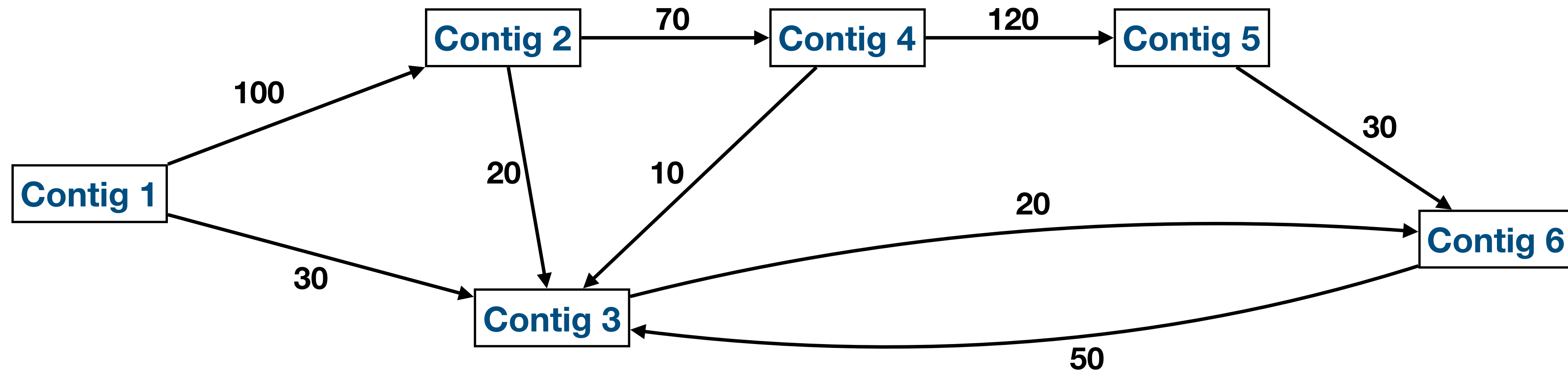


Chaining (scaffolding) contigs



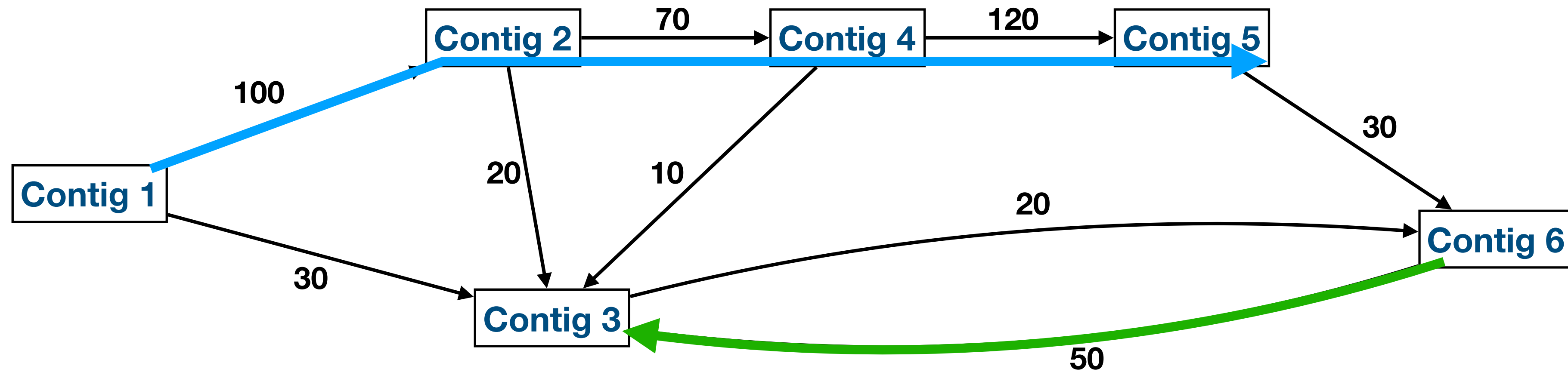
- Another contig assembly-like problem

Chaining (scaffolding) contigs



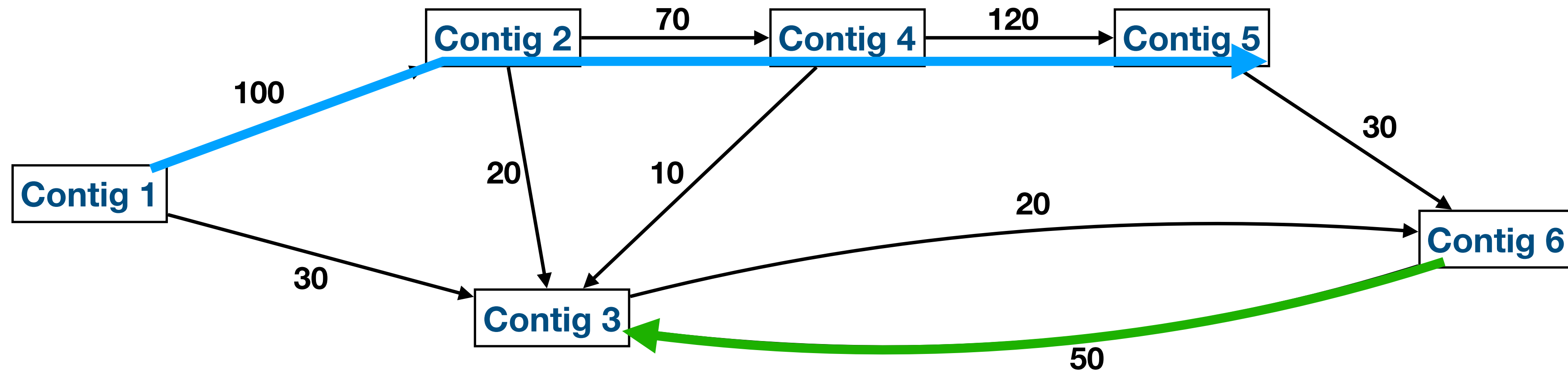
- Another contig assembly-like problem
- Now weights on edges (how much "evidence" there is): new problem formulations

Chaining (scaffolding) contigs



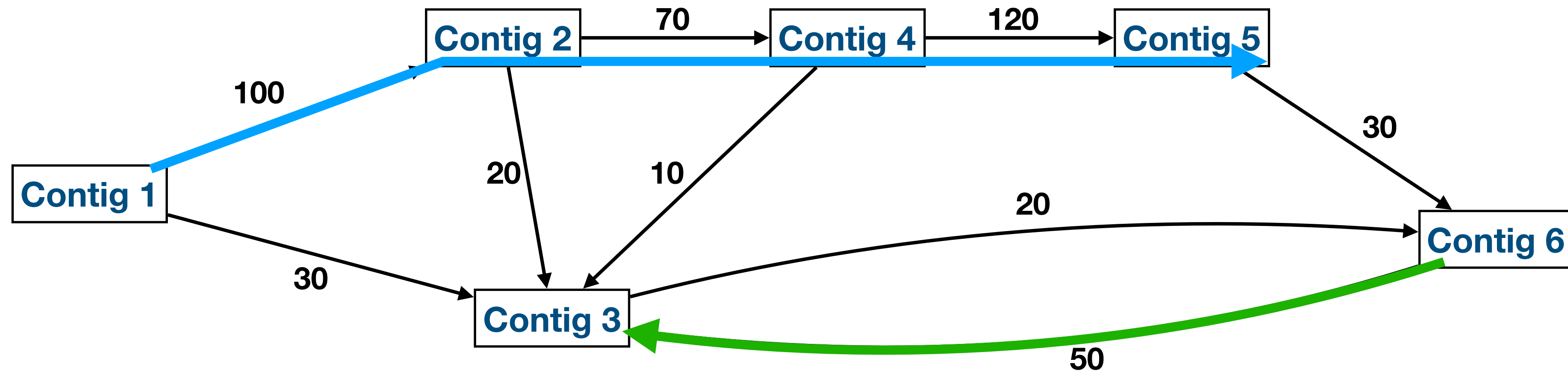
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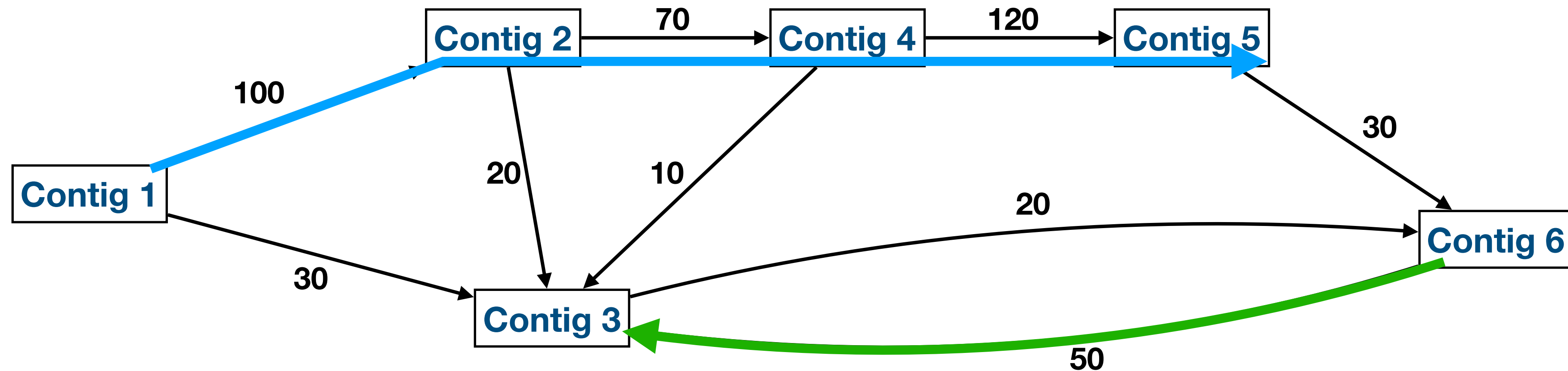
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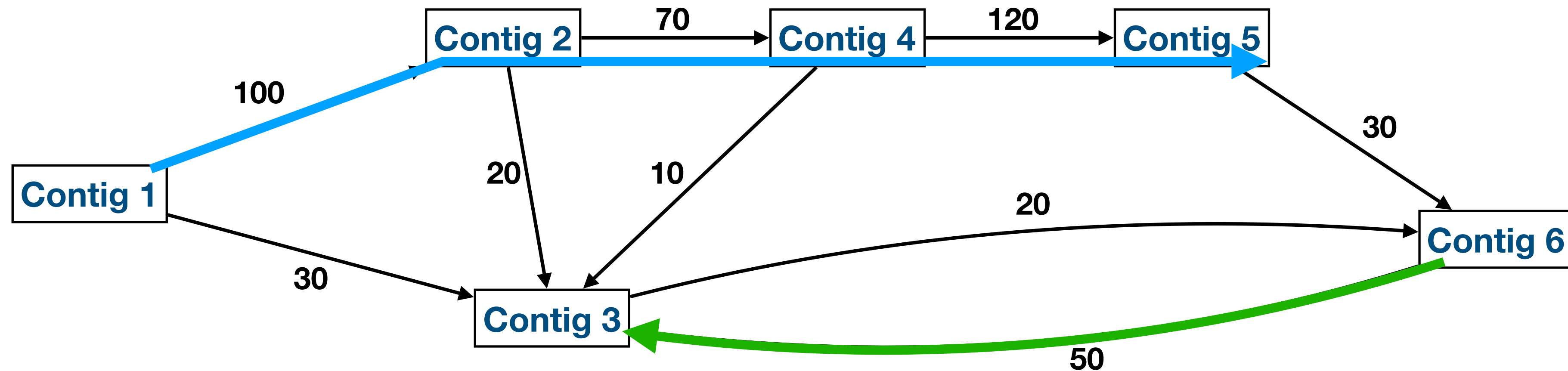
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Chaining (scaffolding) contigs



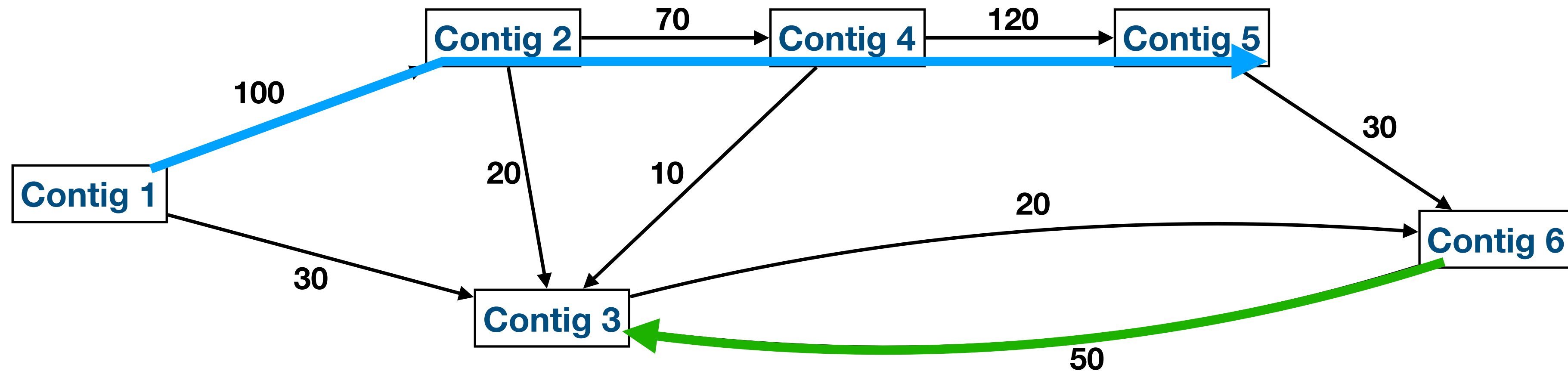
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Chaining (scaffolding) contigs



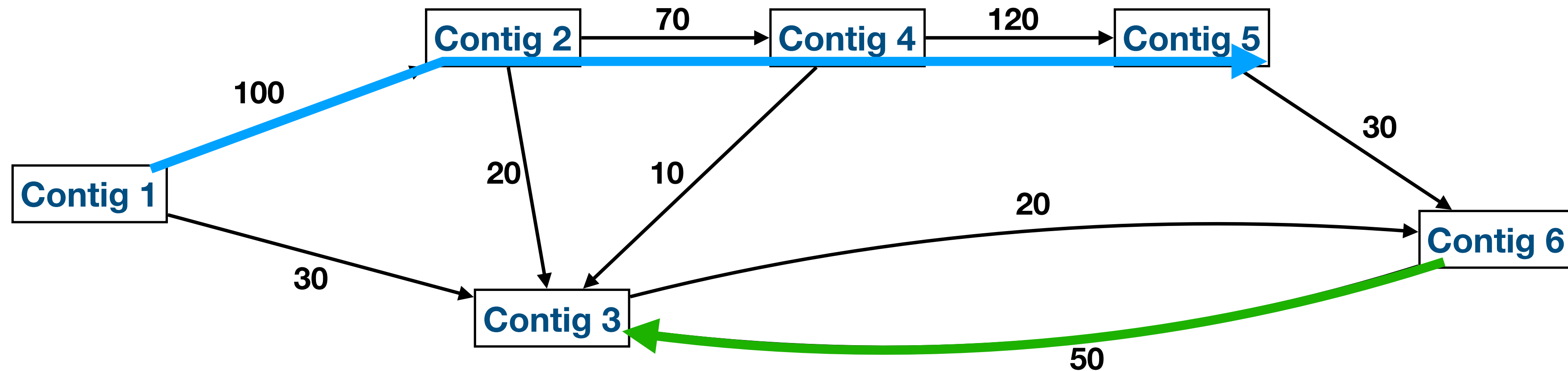
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Chaining (scaffolding) contigs



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 - Igor Mandric, Alex Zelikovsky:
ScaffMatch: Scaffolding Algorithm Based on Maximum Weight Matching. RECOMB 2015: 222-223

Gap filling

Scaffolds contain gap
length estimates
(number of Ns)

Bring back all reads



Find filling paths from the
assembly graph

Output the scaffolds in
which some gaps are "filled"

TCGATAGCTAAAANNNNNNNNNNAATTGTNNNATAGAGATATT

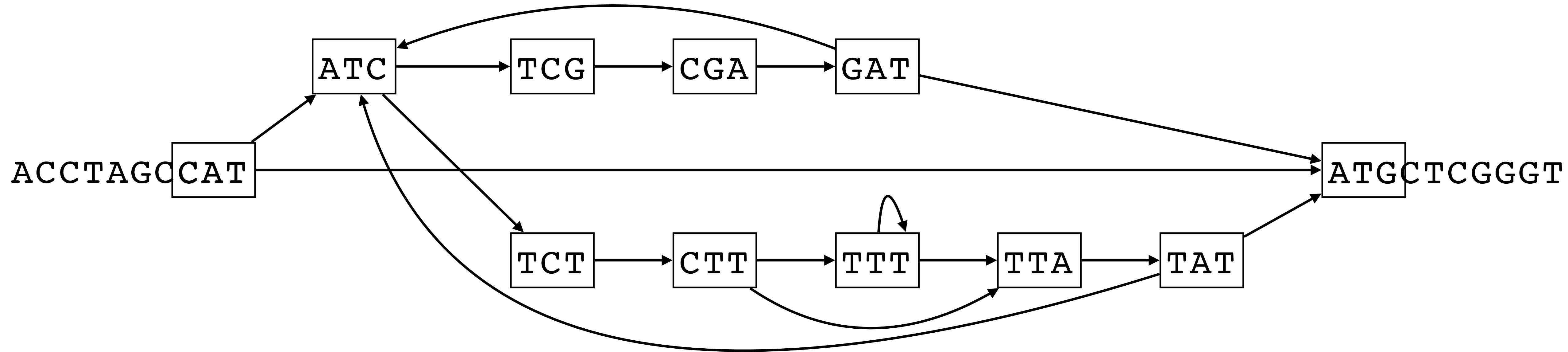


TCGATAGCTAAAATGCCGTTCGGAATTGTNNNATAGAGATATT

Find path of given length

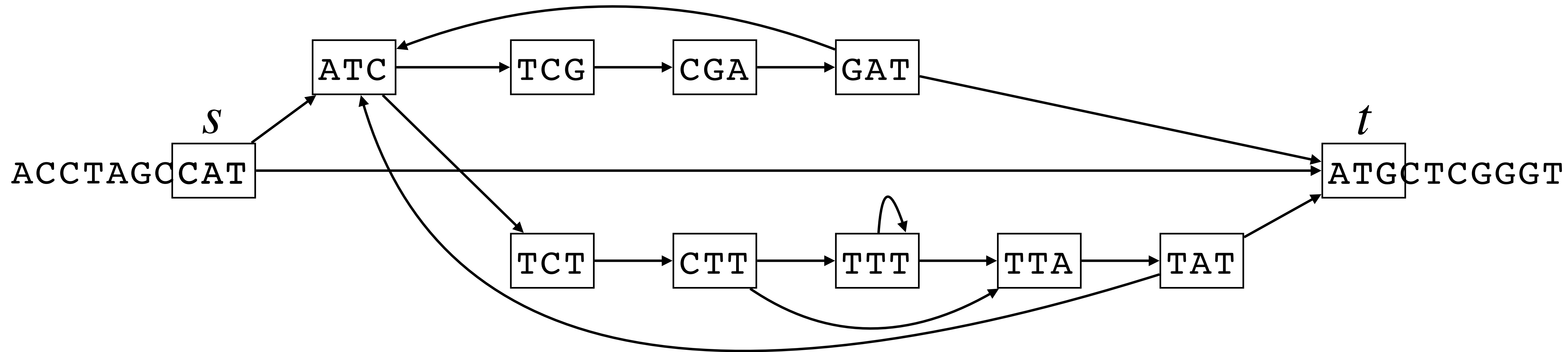


Find path of given length



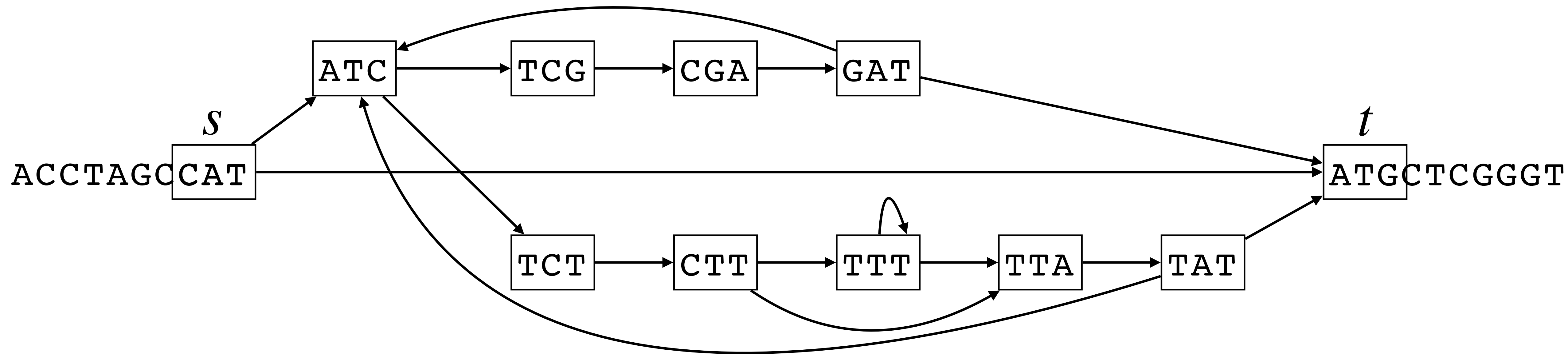
- Build the de Bruijn graph of all reads

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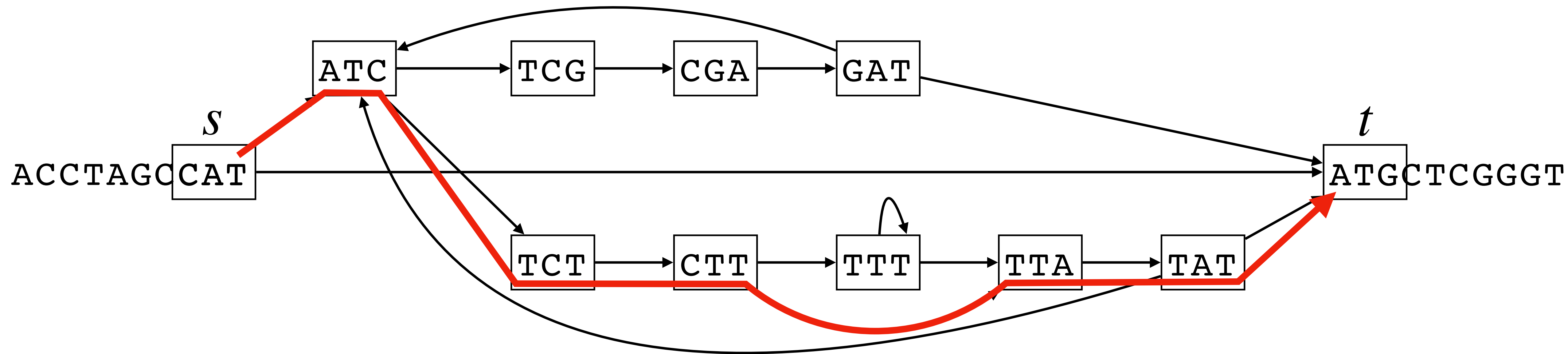
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Find path of given length



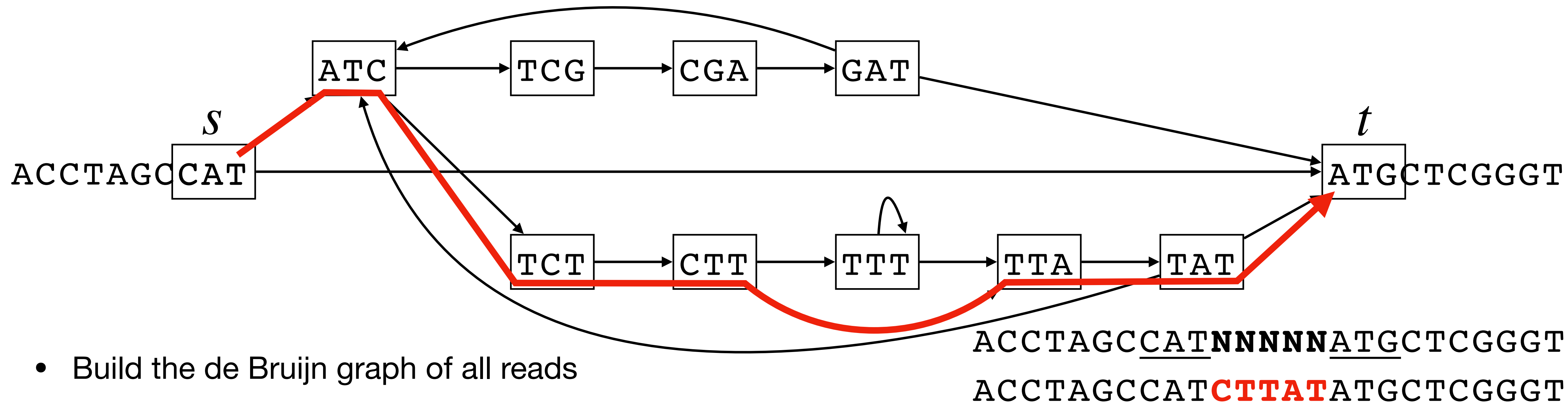
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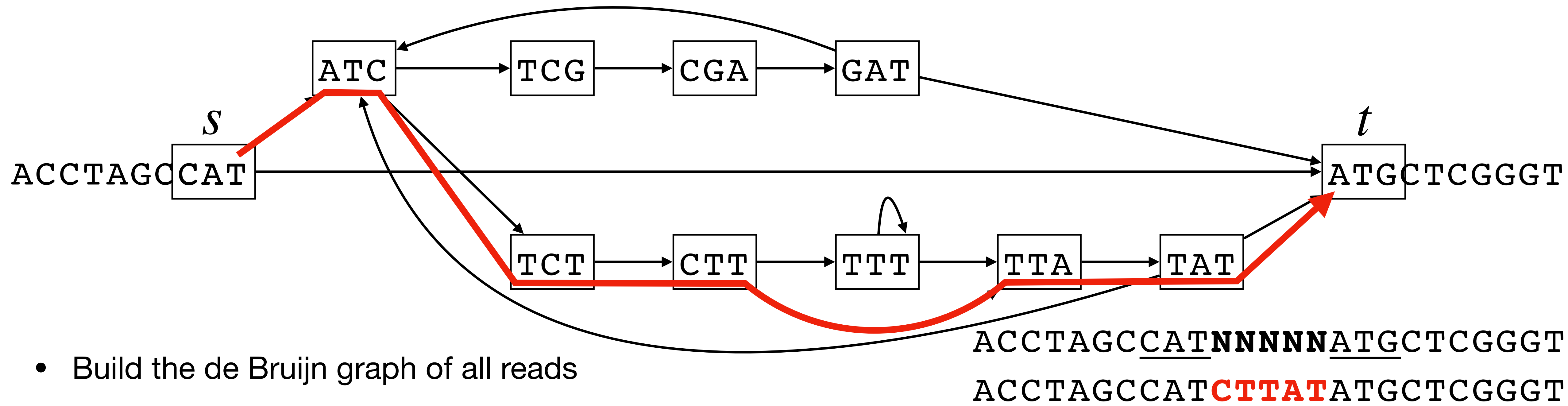
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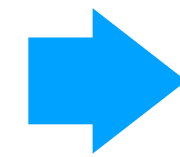
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- Can be solved by dynamic programming in time $O(d \mid \text{edges} \mid)$
 - Leena Salmela, Kristoffer Sahlin, Veli Mäkinen, Alexandru I. Tomescu:
Gap Filling as Exact Path Length Problem. RECOMB 2015: 281-292

Section summary

Contig assembly

Forget about the assembly model (i.e. the problem formulation)

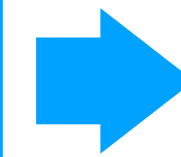
Assemble only parts about which we are sure (contigs → contiguous sequences)



Scaffolding

Chain (order) the contigs

Output chains of contigs with "gaps" (NNNN...) between them



Gap filling

Find filling paths from the assembly graph

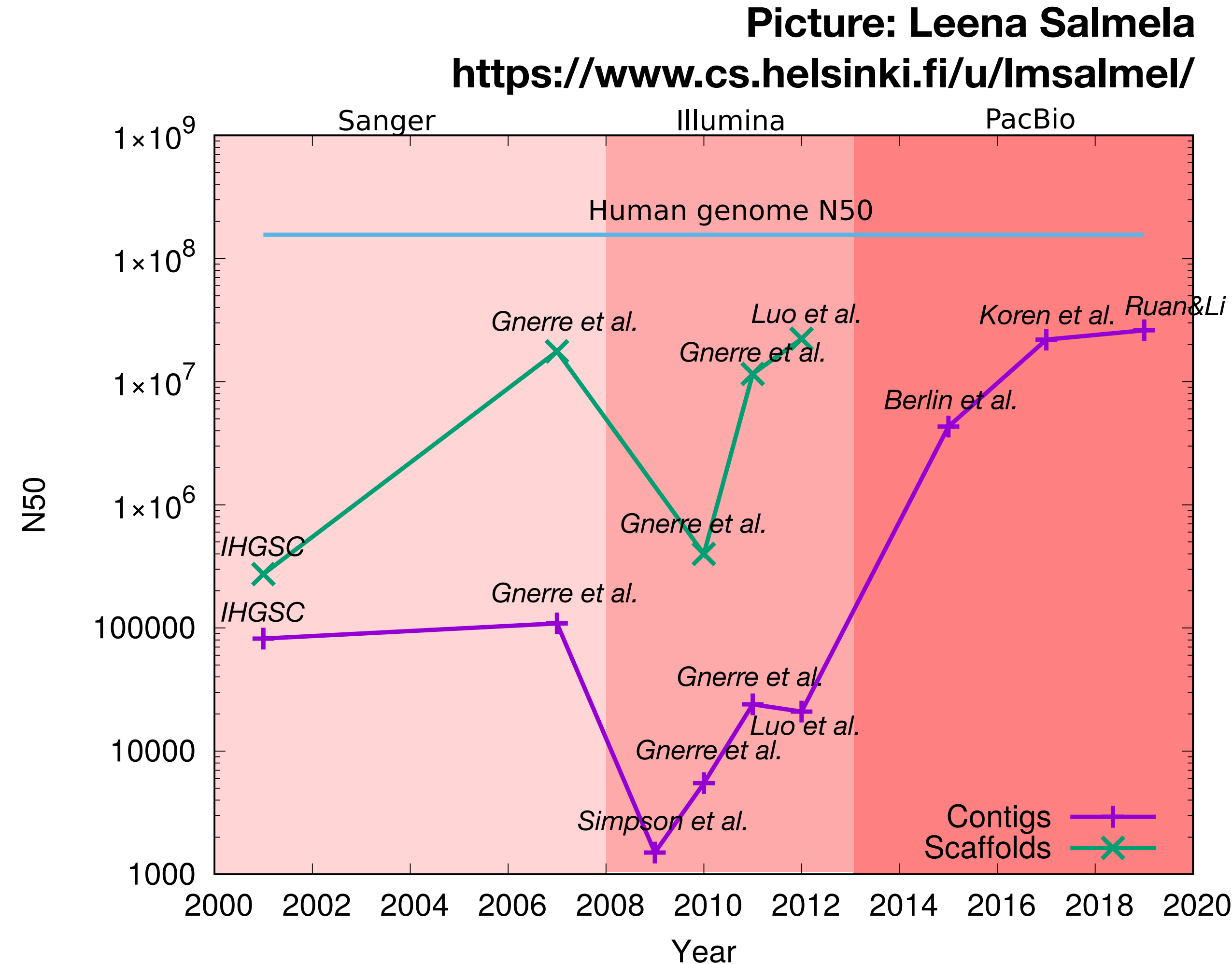
Output the scaffolds in which some gaps are "filled"

- A natural decomposition into subproblems based on the available paired-end information
- One can improve each step individually, thus improving the overall result

Long-read sequencing

(Third-generation sequencing)

- No paired-end reads (focus is on contig assembly)
- Higher error rate: 15% compared to 0.1% for short reads
 - Still developing: accurate PacBio HiFi reads
- No "clear" best strategy
- Short reads still relevant for some scenarios (e.g. metagenomic sequencing)



N50 measure → ASSIGNMENT

A more "practical" theoretical formulation

(A principled approach to contig assembly)

Back to contig assembly

- Goal: obtain sequences that are "guaranteed" to occur in the genome

Back to contig assembly

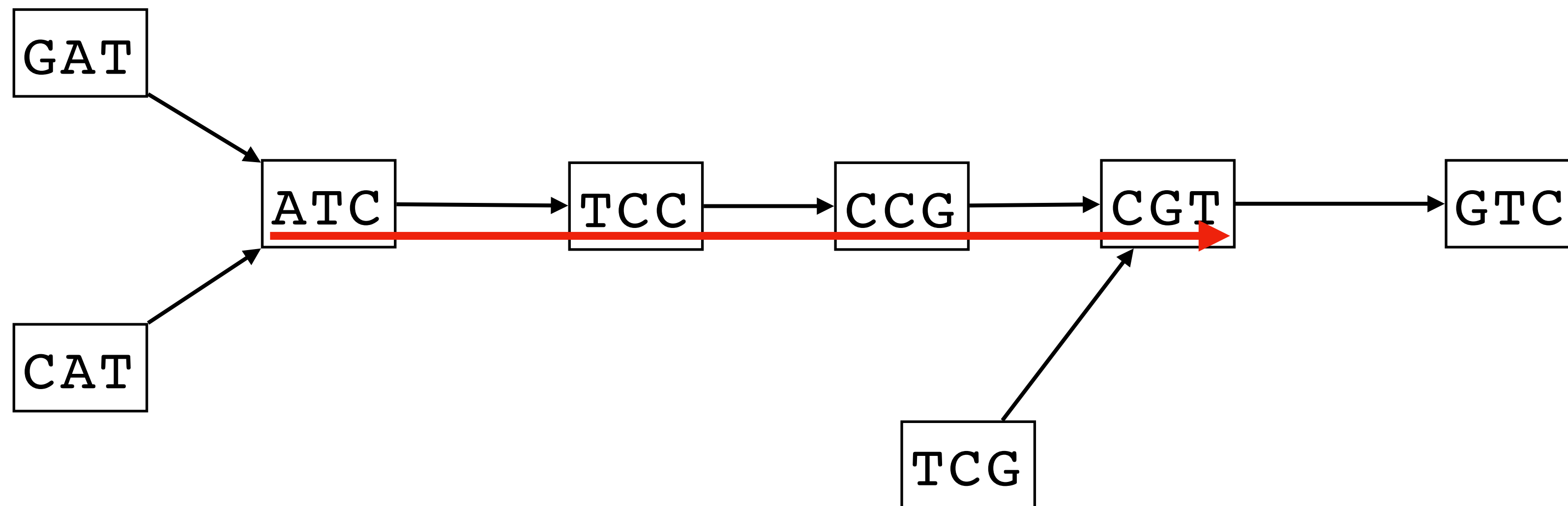
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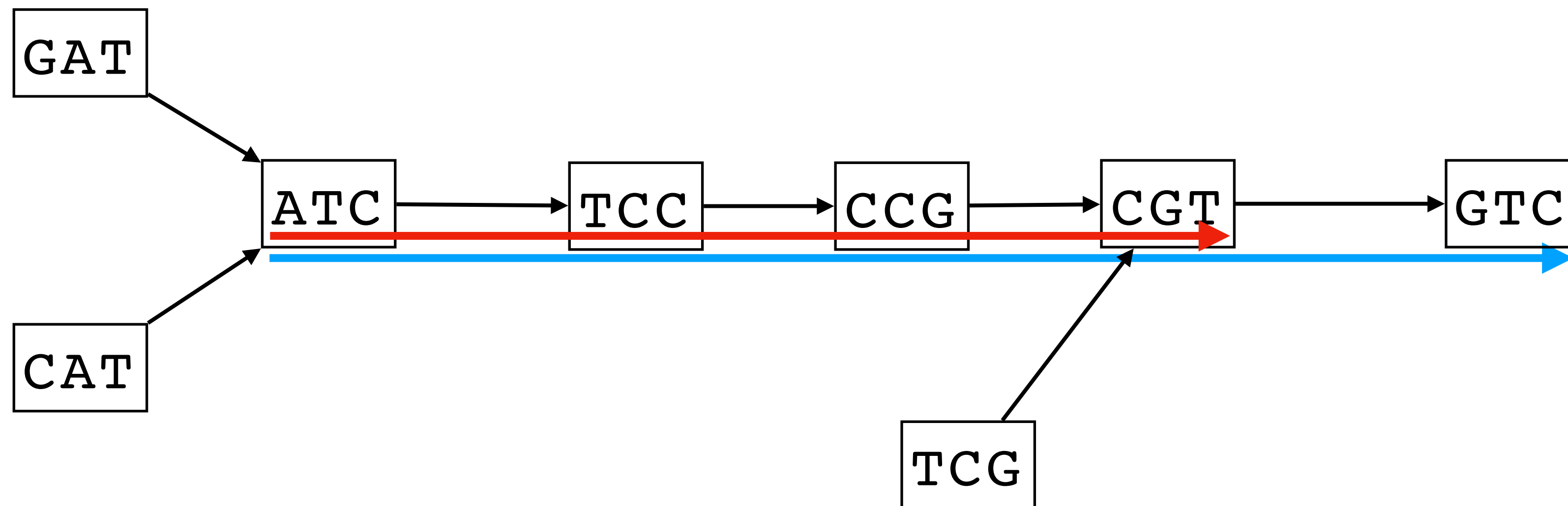
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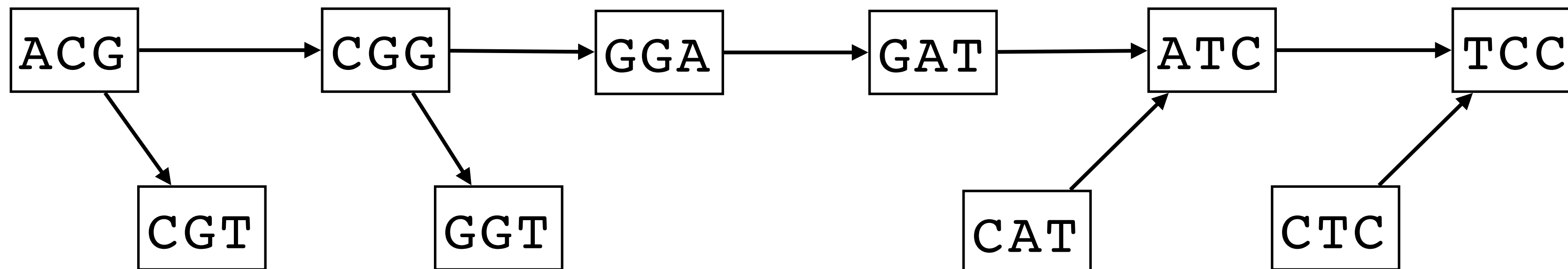


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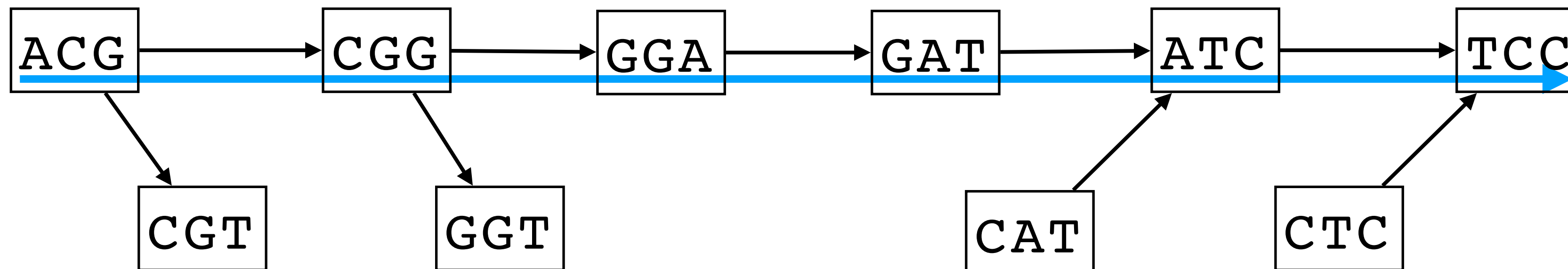
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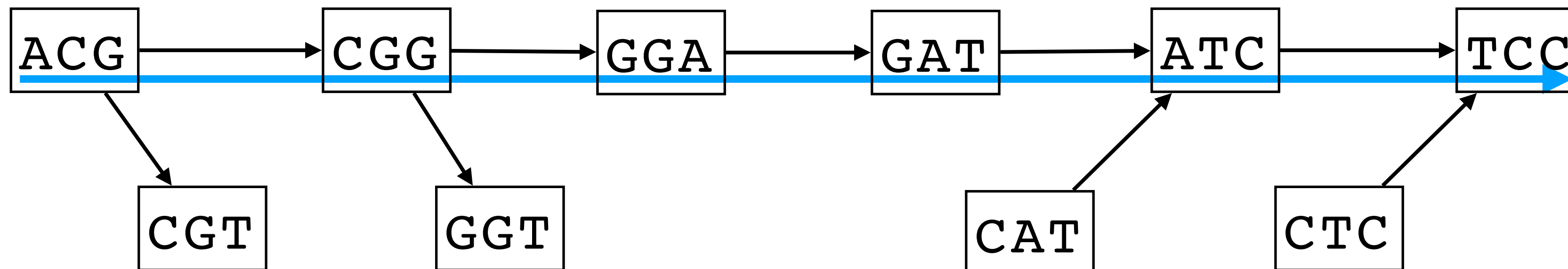
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- Assume that the "genome assembly solution" is a circular walk covering every edge at least once (walk can repeat nodes)
 - ▶ Trivial to find one, exponential to find all
 - ▶ Makes sense for single circular chromosomes (i.e. most bacteria), full coverage, no errors

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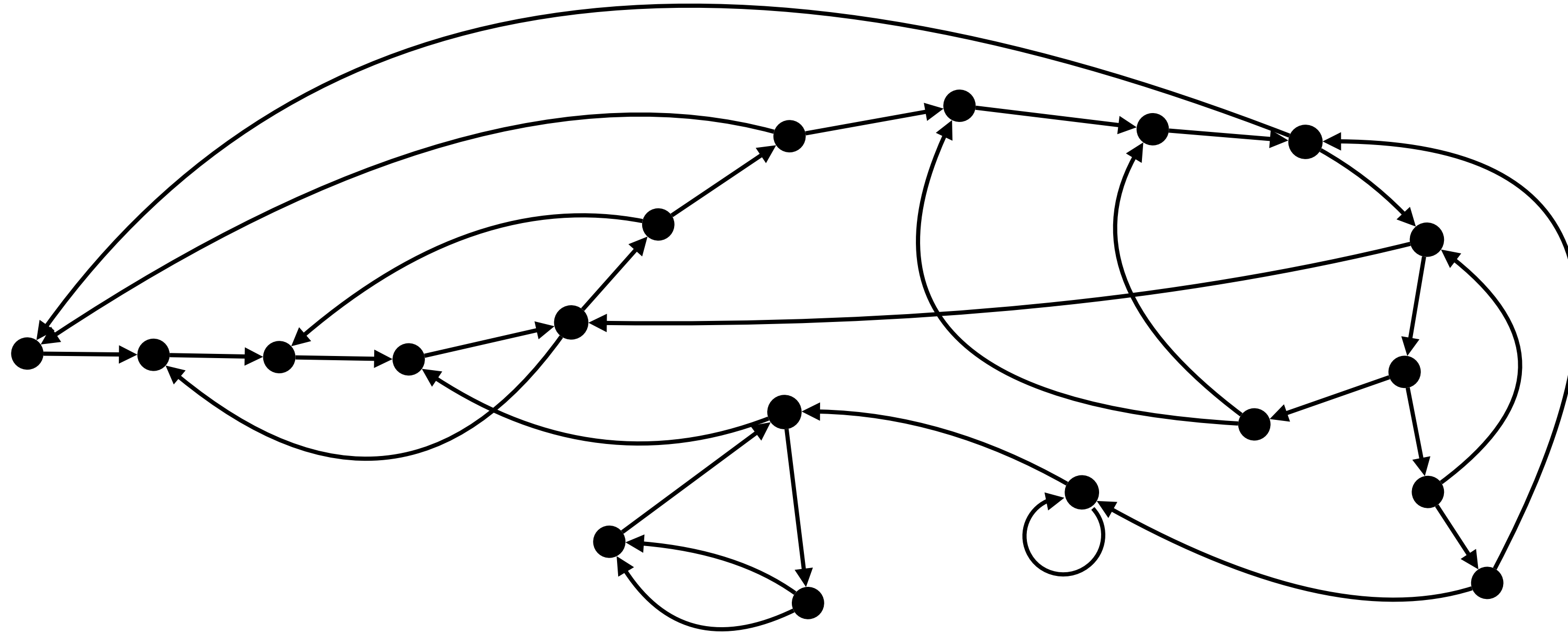
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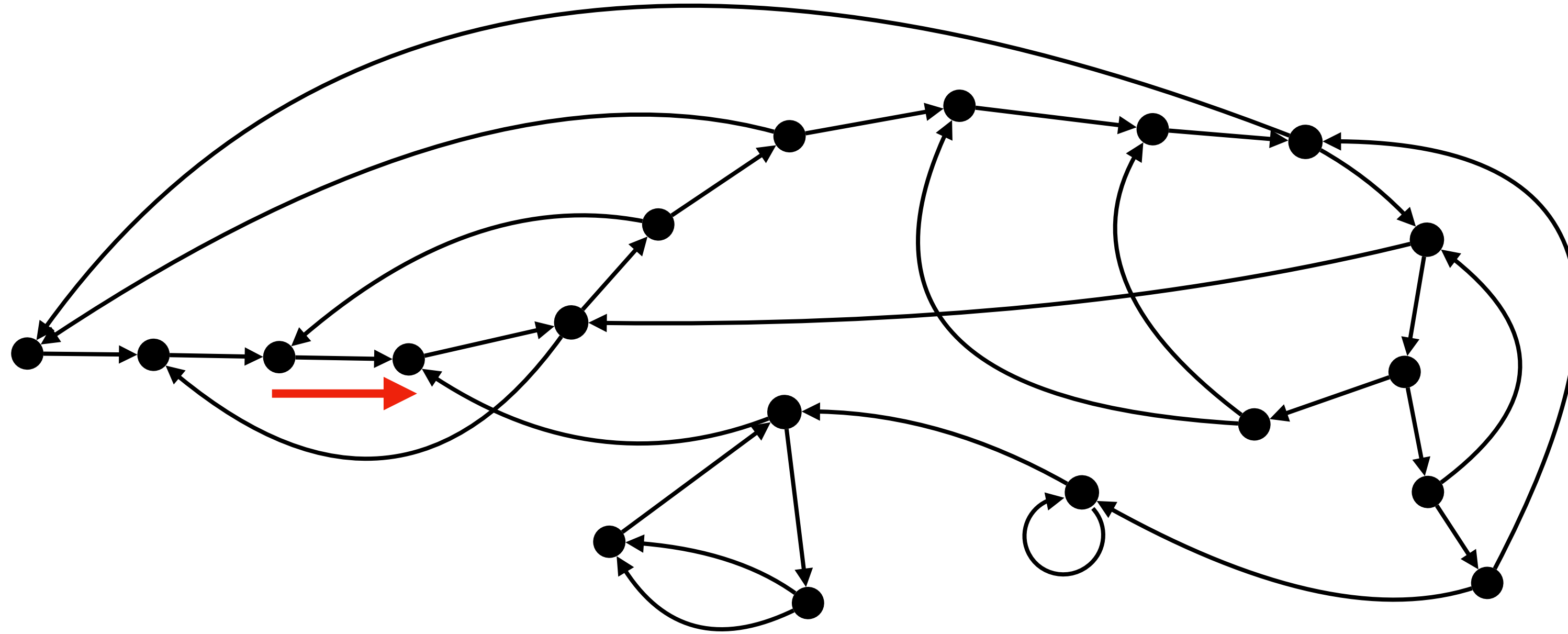
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- Can be adapted to deal with practical issues (ongoing work)
 - ▶ **Subprojects available as Master thesis topics**

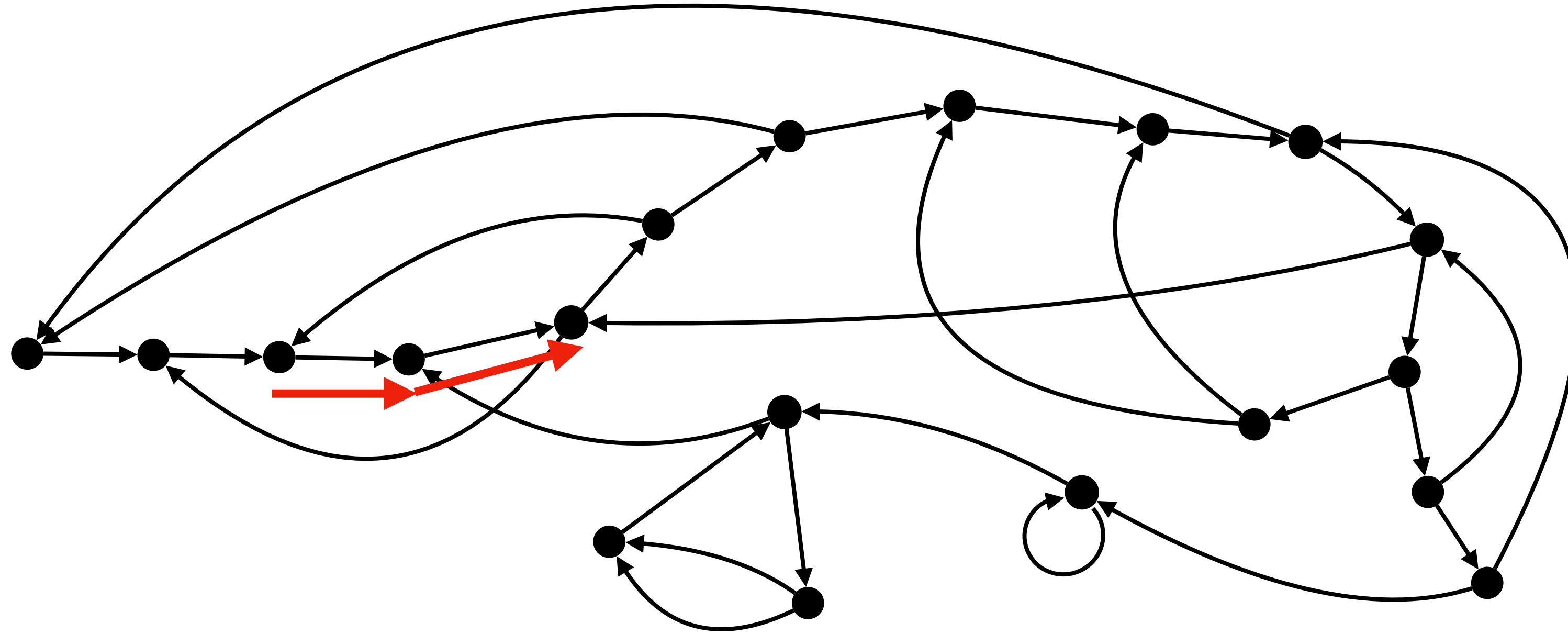
Example + Results on perfect data



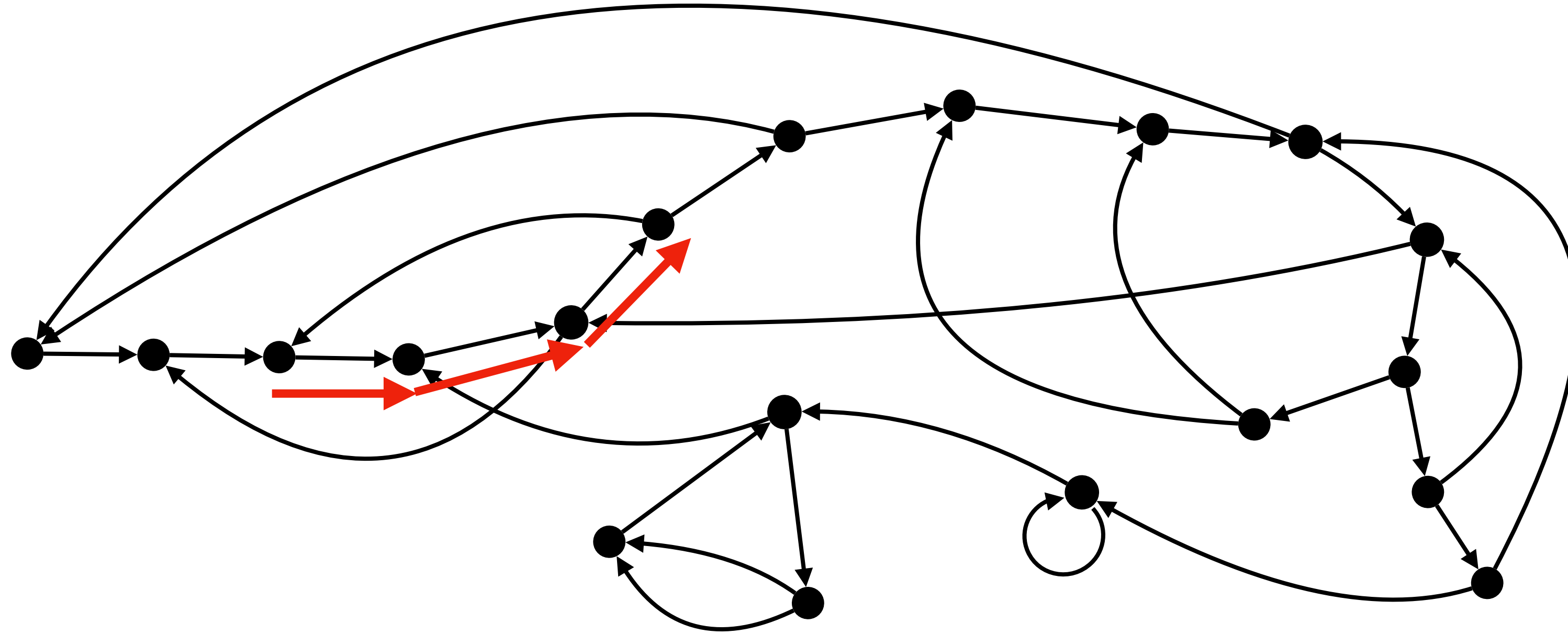
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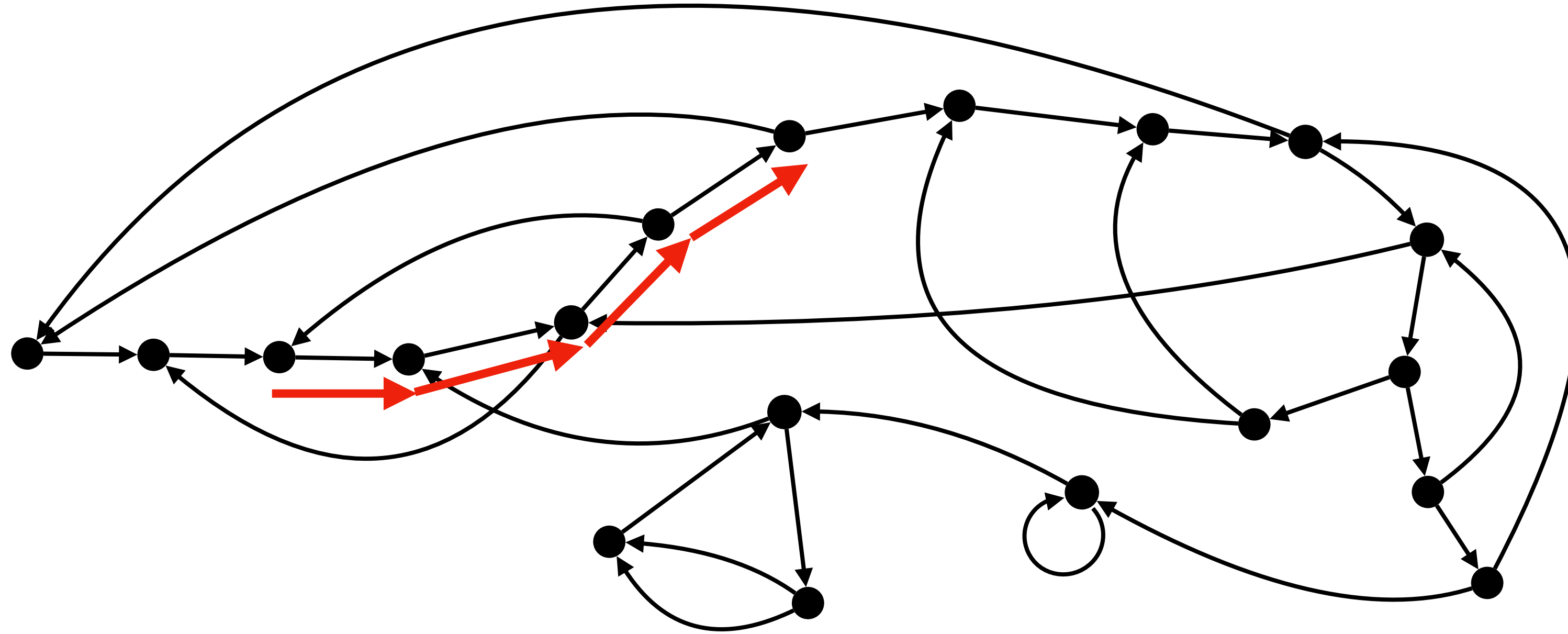
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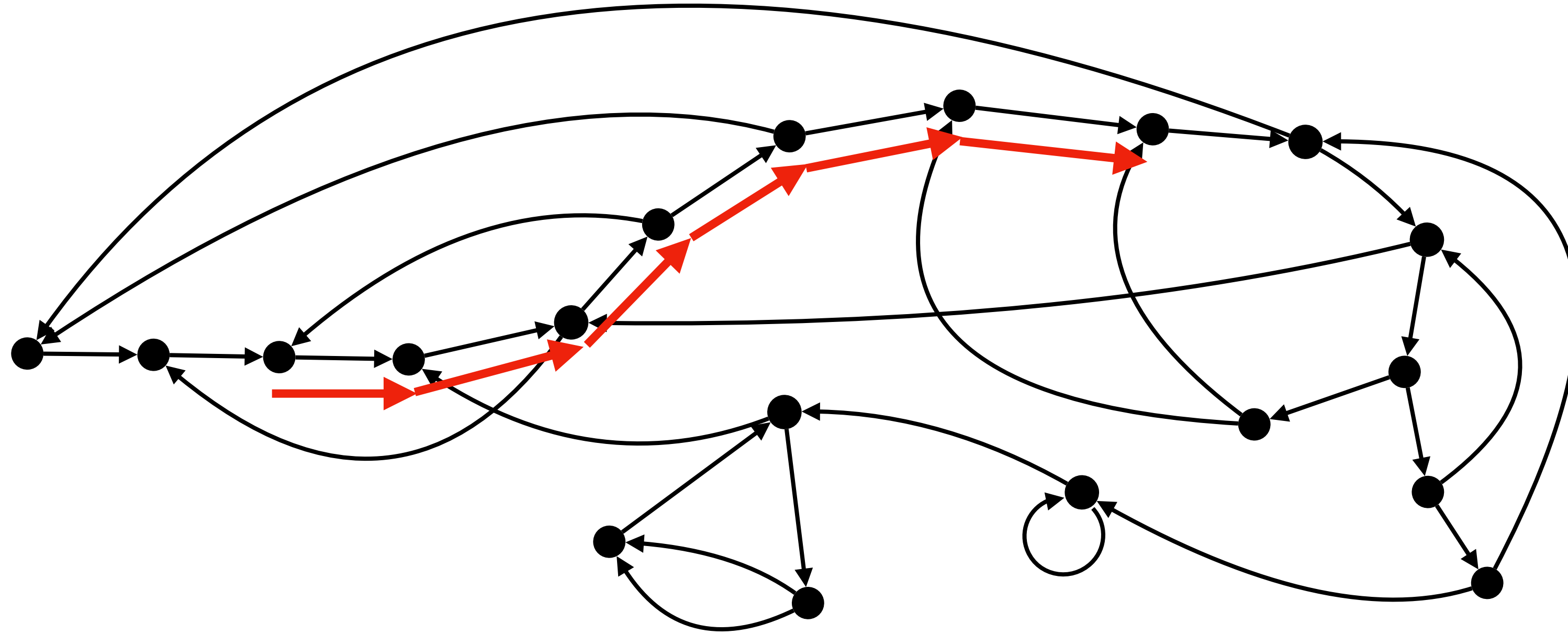
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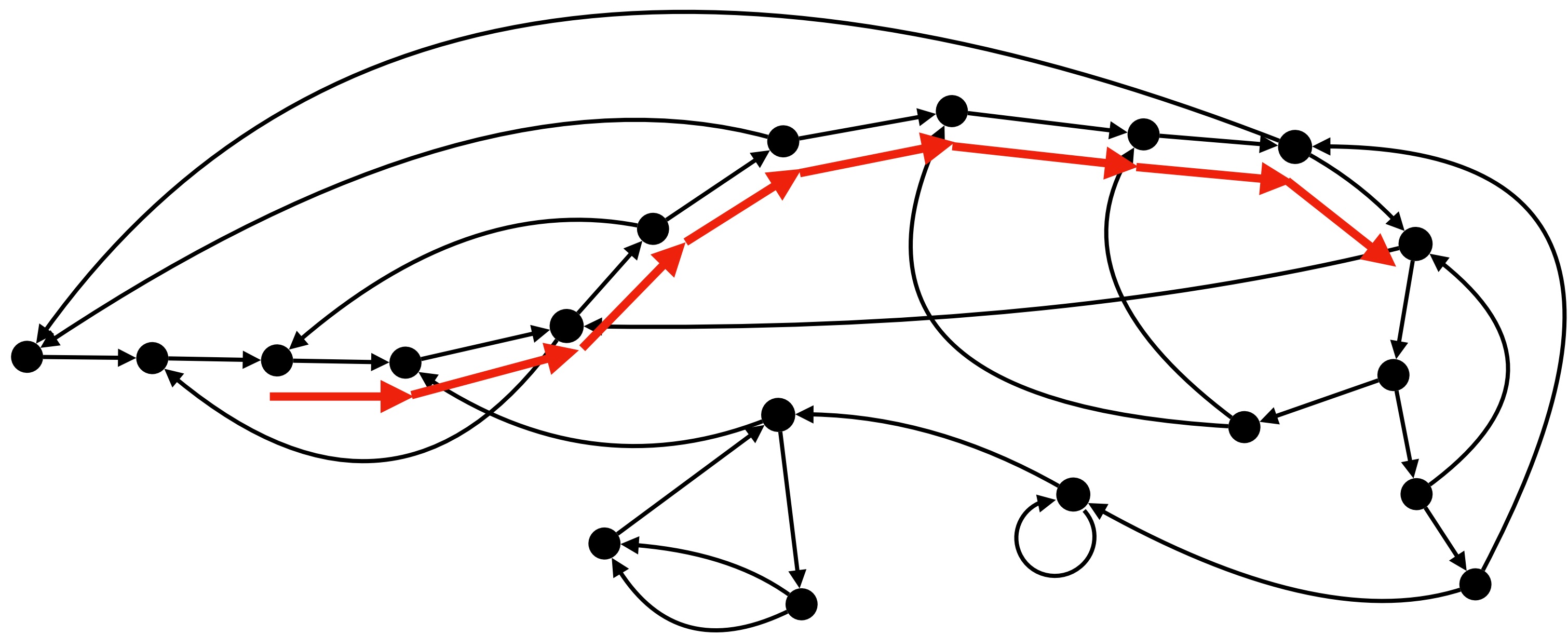
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chr10, length 135M	#Strings	Avg. length	Avg. #SNPs / string
unitigs	260K	546	26
omnitigs	158K (-40%)	887 (+62%)	41 (+58%)

Section summary

Theory is important,
but more so when
it is motivated by practice