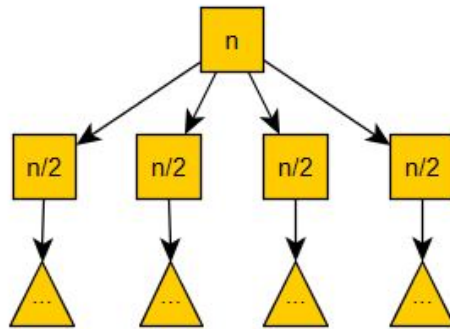


Assignment A4 – Recursion

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$$1 \quad T(n) = 4T\left(\frac{n}{2}\right) + n^2 \lg(n)$$



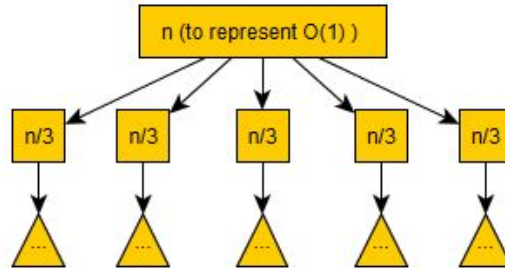
»> Apologies for some discrepancies in the diagrams «<

height of tree is $\log(n)$

$$\begin{aligned}
 T(n) &= \sum_{i=0}^{\log n} n^2 \log\left(\frac{n}{2^i}\right) \\
 &= n^2 \sum_{i=0}^{\log n} \log\left(\frac{n}{2^i}\right) \\
 &= n^2 \left[\sum_{i=0}^{\log n} \log n - \sum_{i=0}^{\log n} \log_2 i \right] \\
 &= n^2 \left[\log n * \log n - \sum_{i=0}^{\log n} i \right] \\
 &= n^2 \left[\log^2 n - \sum_{i=0}^{\log n} i \right] \\
 &= n^2 \left[\log^2 n - \frac{\log n (\log n + 1)}{2} \right] \\
 &= n^2 \left[\log^2 n - \frac{\log^2 n}{2} - \frac{\log n}{2} \right]
 \end{aligned}$$

$$T(n) = O(n^2 \log^2 n)$$

$$2 \quad T(n) = 5T\left(\frac{n}{3}\right) + O(1)$$



$$\begin{aligned} T(n) &= O(1) + 5\left(\frac{n}{3} + 5T\left(\frac{n}{6}\right)\right) \\ &= n + 5^2\left(5T\left(\frac{n}{12}\right) + \left(\frac{n}{6}\right)\right) + 5\left(\frac{n}{3}\right) \\ &= \dots \end{aligned}$$

$$T(n) = 5^k T\left(\frac{n}{3^k}\right) + n \sum_{i=0}^{k-1} \left(\frac{5^i}{3}\right)$$

$$n = 3^k \text{ for } k: k = \log_3 n$$

$$T(n) = 5^k T\left(\frac{3^k}{3^k}\right) + n \sum_{i=0}^{k-1} \left(\frac{5^i}{3}\right)$$

$$= n \log_3 5 + n \Theta\left(\frac{5^k}{3^k}\right)$$

$$= n^{\log_3 5} + n \Theta\left(\frac{5}{k}\right)$$

$$= n^{\log_3 5} + n \Theta(n^{\log_3 5})$$

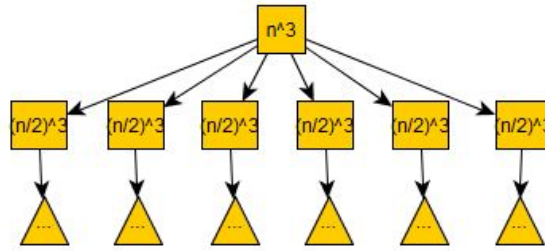
$$= \Theta(n^{\log_3 5})$$

$$3 \quad T(n) = 6T\left(\frac{n}{2}\right) + n^3$$

Subproblem size at $i = \frac{n}{2^i}$

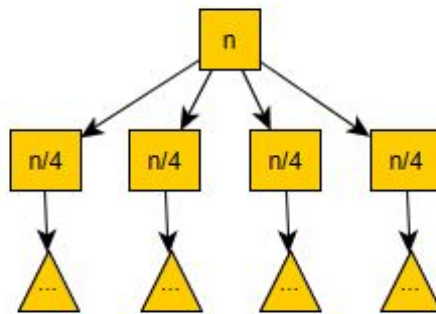
size = 1 when $2^i = n \mid i = \log_2 n$

Height = $\log_2 n + 1 \mid$ Each level has 6^i blocks Level $i = 6^{\log_2 n} = n^{\log_2 6}$



$$\begin{aligned}
 T(n) &= \sum_{i=0}^{\log_2 n - 1} \left(\frac{6}{8}\right)^i n^3 + \Theta(n^{\log_2 6}) < \sum_{i=0}^{\infty} \left(\frac{6}{8}\right)^i n^3 + \Theta(n^{\log_2 6}) \\
 &= \frac{1}{1 - (\frac{6}{8})} n^3 + \Theta(n^{\log_2 6}) \\
 &= O(n^3)
 \end{aligned}$$

$$4 \quad T(n) = 4T\left(\frac{n}{4}\right) + n$$



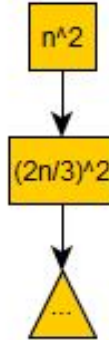
Let number of levels be from 0 extending through h with h being the final level

$$\frac{n}{4^h} = 1 \mid 4^h = n \mid h = \log_4 n$$

$$\begin{aligned}
 T(n) &= 4^h * T(1) + \sum_{i=0}^{h-1} n \\
 &= 4^{\log_4 n} T(1) + n \log_4 n \\
 &= n T(1) + n \log_4 n
 \end{aligned}$$

$$= \Theta(n \log n)$$

$$5 \quad T(n) = T\left(\frac{2n}{3}\right) + n^2$$



Let number of levels be from 0 extending through h with h being the final level

Subproblem size at $i = \frac{2^i n}{3^i}$

$$\frac{n}{(\frac{3}{2})^h} = 1 \mid h = \log_{3/2} n$$

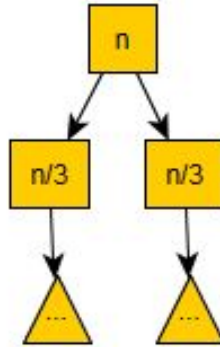
$$\begin{aligned} T(n) &= \sum_{i=0}^{h-1} \left(\frac{4}{9}\right)^i n^2 + T(1) < \sum_{i=0}^{\infty} \left(\frac{4}{9}\right)^i n^2 + T(1) \\ &= \frac{1}{1-\frac{4}{9}} n^2 + T(1) \\ &= \frac{9}{5} n^2 + T(1) \\ &= O(n^2) \end{aligned}$$

$$6 \quad T(n) = 2T\left(\frac{n}{3}\right) + n$$

Let number of levels be from 0 extending through h with h being the final level

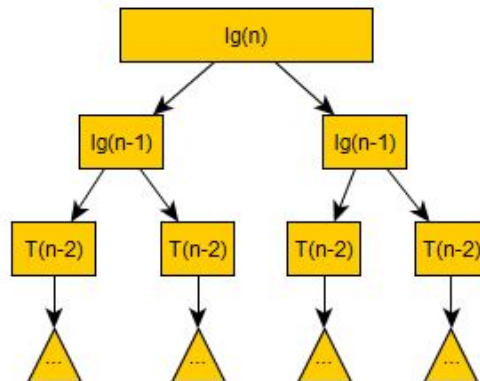
Subproblem size at $i = \frac{n}{3^i}$

$$\frac{n}{3^h} = 1 \mid 3^h = n \mid h = \log_3 n$$



$$\begin{aligned}
 T(n) &= 2^{\log_3 n} T(1) + \sum_{i=0}^{h-1} \left(\frac{2}{3}\right)^i n \\
 &= \Theta(n^{\log_3 2}) + \sum_{i=0}^{h-1} \left(\frac{2}{3}\right)^i n < \Theta(n^{\log_3 2}) + \sum_{i=0}^{\infty} \left(\frac{2}{3}\right)^i n \\
 &= \frac{1}{1-\frac{2}{3}} n + \Theta(n^{\log_3 2}) \\
 &= 3n + \Theta(n^{\log_3 2}) \\
 &= O(n)
 \end{aligned}$$

7 $T(n) = 2T(n-1) + \lg(n)$



$$\begin{aligned}
 T(n) &= \lg(n) + 2 \lg(n-1) + 2^2 \lg(n-2) + \dots + 2^{n-2} \lg(2) \\
 &\leq \lg(n)(1+2+2^2+2^3+\dots+2^{n-2})
 \end{aligned}$$

$$\leq (2^{n-1}-1)(\lg(n))$$

$$\leq O(2^n \lg(n))$$

$$= \Theta(\lg(n))$$

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