Topics to discuss

Master Theorem / Master Method Some More Examples.

Master Theorem or Master Method: General Form, $T(n) = aT(\frac{n}{b}) + f(n)$, where $f(n) = \theta(n^k \log^k n)$ So, it become, $T(n) = aT(\frac{n}{b}) + \theta(n^k \log^p n)$ where, az1, b>1, k≥0 and p is real number. Case I: If a > bk, then T(n) = O(n loga) Ly a=b? a) If b>-1 then $T(n)=\theta(n\log a\log n)$ b) If b=-1 then $T(n)=\theta(n\log a\log n)$ c) If b<-1 then $T(n)=\theta(n\log a\log n)$ case II: If a=bk

Case III: If
$$a < b^k$$

a) If $b \ge 0$ then $T(n) = \theta(n^k \log^{p_n})$

b) If $b < 0$ then $T(n) = \theta(n^k)$

Solve $T(n) = 4T(\frac{n}{2}) + n^2$ using Master Theorems Solution: $T(n) = aT(\frac{n}{b}) + \Theta(n^k \log^p n)$ So, a=4, b=2, $f(n)=\theta(n^2\log^n n)$ and, K=2 and b=0 bK = 2 = 4 => a=bK & \$>-1 case-II (a) T(n) = of (n log b log pth) $= \theta \left(n \log^{\frac{4}{2}} \log^{\frac{4}{1}} n \right)$ = O (N login)

Solve
$$T(n) = 2T(\frac{n}{2}) + \frac{n}{\log n}$$
 using Master theorem.
Solution: $T(n) = a T(\frac{n}{b}) + \theta(n^k \log^k n)$
 $a = 2$, $b = 2$, $f(n) = \theta(n^k \log^k n)$
sw, $K = 1$, $p = -1$
Now, $a = 2$
 $b^k = 2^l = 2$
 $= 7$ $a = b^k$ and $b = -1 = 7$ Case II (b)
 $T(n) = \theta(n \log \log n)$
 $= \theta(n^k \log \log n)$
 $= \theta(n \log \log n)$

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