Davis Putnam ATP

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Davis Putnam Algorithm

- Negate Conclusion
- Convert each statement to CNF
- Turn CNFs into clauses
- Run DP algorithm

boolean Satisfiable(S)

begin

```
if S = {} return true;
```

```
if S = \{\{\}\} return false
```

```
select L \in lit(S);
```

```
return Satisfiable(SL) || Satisfiable(SL');
```

end

Add-ons

- Subsumption Elimination
 - [A,B] subsumes [A,B,C]
 - remove the subsumed statemetrs
- Pure-Literal Elimination
 - L ∈ S AND ~L ∉ S
 - remove all statements containing L

- Unit Literal
 - [L] ∈ S
 - don't branch
- Tautological Elimination

Input Parsing

- Two formats:
 - "Classic"
 - natural language (kinda)
 - "New"

Formats

"Classic" format:

A implies (N or Q)

not(N or not A)

A implies Q

"New" format:

 $A \rightarrow (N \vee Q)$

~(N v ~A)

A -> Q

CNF

- Putting in CNF is Hard...
- Two approaches:
 - Regular Expressions
 - Wolfram-Alpha

$$A < -> (N \lor Q)$$

$$\sim$$
(N \wedge (Q \vee \sim P) \wedge J) $<->$ A

Regular Expressions

- Fast
- Only for simple expressions
- A
 - re.match('(\w)\$', line, re.l)
- A <-> ~B
 - $m = re.match('(\w+) xnor NOT (\w+)$', line, re.l)$

Wolfram Alpha

- Slow...
- Any complexity

- Free Developer API
 - 2000 queries/month
 - Simple HTTP requests
 - Returns XML

Demo!