CH5 Algorithms(演算法)

Definition

• *An algorithm is an ordered (照

順序) set

of unambiguous (清楚),

executable (可

執行) steps that defines a terminating
(可終止) process.

- *Program
- Formal representation of an algorithm
 - *Process
- Activity of executing a

program

Ordered Set

• *Steps in an algorithm must

have a wellestablished structure in terms of the order in which its steps are executed

- *Each step must be an executable instruction
- Example: "Making a list of all the positive integers"

is not an executable instruction

*May involve more than one

thread (parallel algorithms)

Unambiguous Steps

• * During execution of an

algorithm, the information in the state of the process must be sufficient to determine uniquely and completely the actions required by each step

• *The execution of each step in an algorithm does not require creative skills. Rather, it requires only the ability to follow directions.

Terminating Process

- *All execution of an algorithm must lead to an end
- *There are, however, many meaningful application for non-terminating processes
- *Computer science seeks to distinguish both
- *Problems whose answers can
 be obtained

algorithmically

*Problems whose answers lie

beyond the capabilities of algorithmic systems

Algorithm and Its Representation

- *Like a story and a story book
- *Different language translation

of a textbook

*Example: converting temperature readings

from Celsius to Fahrenheit

- 1. F = (9/5)C + 32
- 2. Multiply the temperature

reading in Celsius by 9/5 and then add 32 to the product

• 3. Implemented by electronic

circuit

*Underlying algorithm is the same, only the representations differ

Level of Details

- *May cause problems in communicating algorithms
- *Example:
- * "Convert the Celsius

reading to its Fahrenheit equivalent"

• *This might suffice for

meteorologists

- *But a layperson would argue that this instruction is ambiguous
- *The problem is that the

algorithm is not represented in enough detail for the layperson

Algorithm Representation

- *Primitive
- *Set of building blocks from

which algorithm representations can be constructed

- *Programming language
- *Collection of primitives
- *Collection of rules stating

how the primitives can be combined to represent more complex ideas

Primitives

- *Syntax
- *Symbolic representation
- *Semantics
- *Concept represented

(meaning of the primitive)

Levels of Abstraction

- *Algorithm
- *Procedure to solve the

problem

*Often one of many

possibilities

- *Representation
- *Description of algorithm

sufficient to communicate it to the desired audience

*Always one of many

possibilities

Machine Instructions as **Primitives**

*Algorithm based on machine

instructions is suitable for machine execution

• *However, expressing

algorithms at this level is tedious

*Normally uses a collection of

higher level primitives, each being an abstract tool constructed from the low-level primitives provided in the machine's language

Pseudocode

• *Less formal, more intuitive

than the formalprogramming languages

• *A notation system in which

ideas can be expressed informally during the algorithm development process

• *A consistent, concise notation

for representing recurring semantic

structure

*Comparison with flow chart

The Point of Pseudocode

• *To communication the algorithm to the readers

• *The algorithm will later turn

into

Program

• *Also help the program

maintainer or developer to understand the program

Pseudocode Primitives

- *Assignment
- name <- expression
- *Conditional selection
- if condition then action
- *Repeated execution
- while condition do activity
- *Procedure
- procedure name (generic

names)

Basic Primitives

- *total <- price + tax
- *if (sales have decreased)

then (lower the price by 5%)

- *if (year is leap year)
 then (divide total by 366)
 else (divide total by 365)
- *while(tickets remain to be

sold) do

(sell a ticket)

Procedure Primitive

- *total <- price + tax
- *tax?
- *A procedure to calculate tax

Tax as a Procedure

Procedure tax
if (item is taxable)
then (if (price > limit)
then (return
price*0.1000)
else (return
price*0.0825)
)
else (return 0)

Nested Statements

*One statement within another

```
if (item is taxable)
then (if (price > limit)
then (return
price*0.1000)
else (return
price*0.0825)
)
else (return 0)
```

Indentations

*Easier to tell the levels of

```
nested statements
if (item is taxable)
then (if (price > limit)
then (return price*0.1000)
else (return price*0.0825)
)
else (return 0)
```

Structured (模組化) Program

- *Divide the long algorithm into smaller tasks
- *Write the smaller tasks as procedures
- *Call the procedures when needed
- *This helps the readers to understand the structure of the algorithm

if (customer credit is good) then (ProcessLoan) else (RejectApplication)

Algorithm Discovery

- *Development of a program consists of
- *Discovering the underlying algorithm
- *Representing the algorithm as a program
 *Algorithm discovery is usually

the more challenging step in the software

development process

*Requires finding a method of solving the problem

Problem Solving Steps

- 1. Understand the problem
- 2. Get an idea
- 3. Formulate the algorithm and represent
- it as a program
- 4. Evaluate the program
 - 1. For accuracy
- 2. For its potential as a tool for solving other problems

Not Yet Sure What to Do

- 1. Understand the problem
- 2. Get an idea
- 3. Formulate the algorithm and represent

it as a program

- 4. Evaluate the program
 - 1. For accuracy
- 2. For its potential as a tool for solving other problems

Difficulties

- *Understanding the problem
- *There are complicated

problems and easy problems

• *A complete understanding of

the problem before proposing any solutions is somewhat idealistic *Get an idea

• *Take the 'Algorithm'

course

• *Mysterious inspiration

Getting a Foot in the Door

- *Work the problem backwards
 - *Solve for an example and

then generalize

• *Solve an easier related problem

*Relax some of the

problem

constraints

- *Divide and conquer
- *Stepwise refinement
- *top-down methodology
- *Popular technique

because it produces modular programs

- *Solve easy pieces of the problem first
 - *bottom up methodology

Work the Problem Backwards

- *Simplify the problem
- *Build up a scenario for the simplified problem
- *Try to solve this scenario
- *Generalize the special solution

to general scenarios

- *Consider a more general problem
- *Repeat the process

Example Problem

- *Person A is assigned the task of determining the ages of B's three children.
- *B tells A that the product of the children's ages is X.
- *A replies that another clue is required.
- *B tells A the sum of the children's ages Y.
- *A replies that another clue is needed.
- *B tells A that the oldest child plays the piano.
- *A tells B the ages of the three children.
- *Come out with an algorithm for person A

Divide and Conquer Concept

• *Not trying to conquer an entire

task at once

- *First view the problem at hand in terms of several subproblems
- *Approach the overall solution in terms of steps, each of which is easier to solve than the entire original problem

Divide and Conquer

- *Steps be decomposed into smaller steps
- *These smaller steps be broken into still smaller ones

• *Until the entire problem has

- been reduced to a collection of easily solved subproblems
- *Solve from the small subproblems and gradually have it all.

Iterative Structures

- *Used in describing algorithmic process
- *A collection of instructions is repeated in a looping manner
 Search Problem
- *Search a list for the occurrence of a

particular target value

- *If the value is in the list, we consider the search a success; otherwise we consider it a failure
- *Assume that the list is sorted according to some rule for ordering its entries while vs. repeat Structure
- In repeat structure the loop's body is always performed at least once (posttest loop)

• While in while structure, the

body is never executed if the termination is satisfied the first time it is tested (pretest loop)

A Sort Problem

- *Sort a list of names into alphabetical order
 - *The constraint is to sort the

list "within itself"

- *In other words, sort by shuffling its entries as opposed to moving the list to another location
 - *Typical in computer

applications to use the storage space efficiently

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Recursive Structures

Involves repeating the set of instructions as a subtask of itself

•

An example is in processing incoming telephone calls using the call-waiting (來電等候) feature •

- 1.An incomplete telephone conversation is set aside while another incoming call is processed 2.Two conversations are performed
- 3.But not in a one-after-the-other manner as in the loop structure 4.Instead one is performed within the other

Recursion

- 1.Execution is performed in which each stage of repetition is as a subtask of the previous stage
- .
- 2.Example: divide-and-conquer in binary search

Characteristics of Recursion

- 1.Existence of multiple copies of itself (or multiple activations of the program)
- 2.At any given time only one is actively progressing
- 3. Each of the others waits for

another activation to terminate before it can continue

Recursive Control

Also involves

- 1.Initialization
- 2. Modification
- 3.Test for termination (degenerative case)

•

Test for degenerative case

- 1.Before requesting further activations
- 2.If not met, assigns another activation to solve a revised problem that is closer to the termination condition

•

Similar to a loop control

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Software Efficiency

- 1.Measured as number of instructions executed
- 2.notation for efficiency classes
- 3.Best, worst, and average case

Big-theta Notation

Identification of the shape of the graph representing the resources required with respect to the size of the input data

1. Normally based on the worst-case analysis

2.Insertion sort: $\Theta(n2)$

3.Binary search: $\Theta(\lg n)$

(n2-n)/2

lg2(n+1) - 1

Formal Definition

 Θ (n2): complexity is kn2+o(n2)

 $1.f(n)/n2 \cdot k, n \cdot \infty$

•

o(n2): functions grow slower than n2

 $1.f(n)/n2 \cdot 0, n \cdot \infty$

 $(n2-n)/2 = \frac{1}{2} n2 - \frac{1}{2} n$ o(n2)

Problem Solving Steps

- 1.Understand the problem
- 2.Get an idea
- 3. Formulate the algorithm and represent it as a program
- 4.Evaluate the program
- 1. For its potential as a tool for solving other problems (the faster, the better)
 - 2.For accuracy

Software Verification

- 1.Evaluate the accuracy of the solution
- 2. This is not easy
- 3. The programmer often does not know whether the solution is accurate (enough)
- 4. Example: Traveler's gold chain

Ways to Level the Confidence

For perfect confidence

- 1.Prove the correctness of a algorithm
- 2. Application of formal logic to prove the correctness of a program

For high confidence

- 1.Exhaustive tests
- 2. Application specific test generation

For some confidence

- 1.Program verification
- 2. Assertions

Program Conditions

Preconditions

1. Conditions satisfied at the beginning of the program execution

Propagations

1. The next step is to consider how the consequences of the preconditions propagate through the program

Program Verification

Assertions

- 1.Statements that can be established at various points in the program
- 2.To check for rightful

precondition & propagation

Proof of correctness to some degree

- 1.Establish a collection of assertions
- 2.If all assertions pass
- 3. The program is correct to some degree