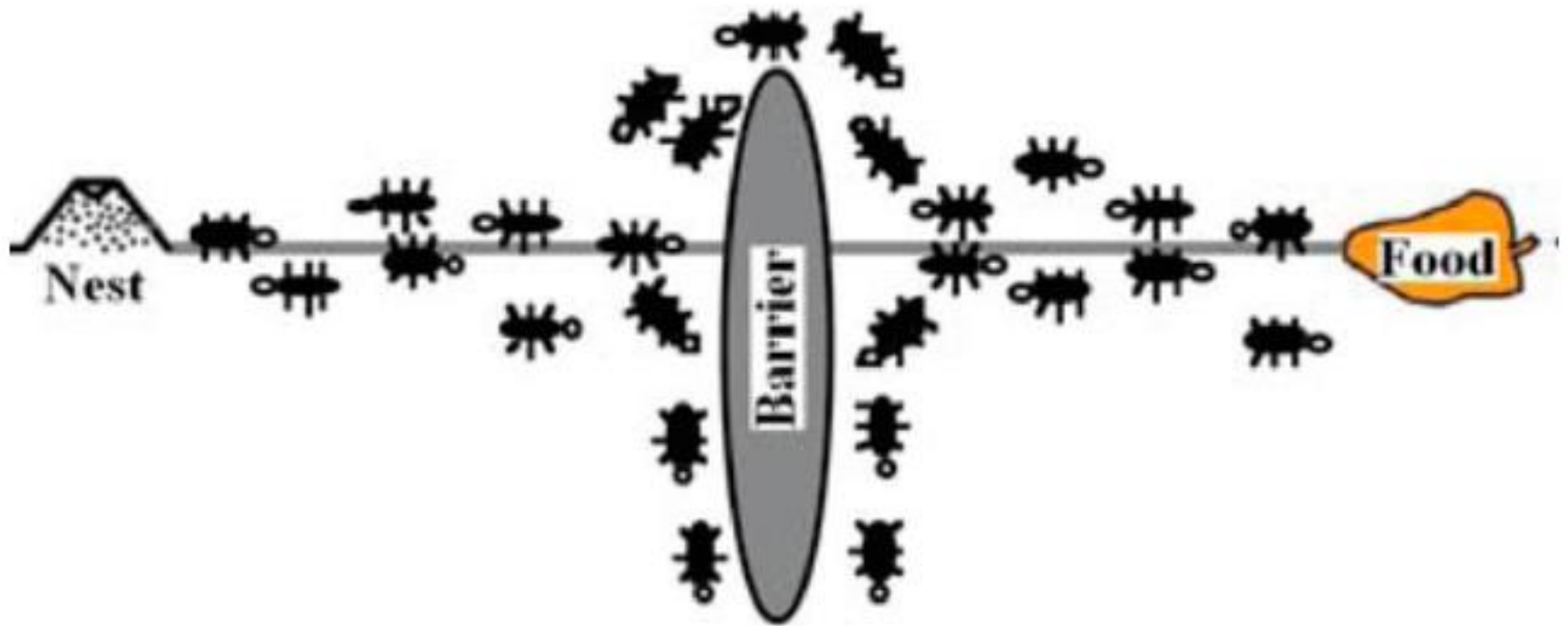


# Ant Algorithm





Initialize the pheromone matrix  $\tau$  for each pair of cities  
Place the  $m$  ants on  $n$  random cities  
**for**  $t = 1$  to  $nc$  **do**  
    **for**  $i = 1$  to  $n$  **do**  
        **for**  $k = 1$  to  $m$  **do**  
            Choose next city  $j$  according to the transition rule  
        **for**  $k = 1$  to  $m$  **do**  
            Calculate tour distance  $L_k$  for ant  $k$   
            **if** an improved tour is found **then**  
                Update  $T^*$  and  $L^*$   
    Update the pheromone matrix  $\tau$



Initialize  $T_{Global}^*$  {this data is shared, everything else is private}

**parallel region with  $nColonies$  threads**

Initialize the pheromone matrix  $\tau$  for each pair of cities

Place the  $m$  ants on  $n$  random cities

**for**  $t = 1$  to  $nc$  **do**

**for**  $i = 1$  to  $n$  **do**

**for**  $k = 1$  to  $m$  **do**

            Choose next city  $j$  according to the transition rule

**for**  $k = 1$  to  $m$  **do**

        Calculate tour distance  $L_k$  for ant  $k$

**if** an improved tour is found **then**

            Update  $T^*$  and  $L^*$

**if** this is an exchange cycle **then**

**if**  $L^* < L_{Global}^*$  **then**

                \*\*\*Critical section\*\*\*

**if**  $L^* < L_{Global}^*$  **then**

$T_{Global}^* = T^*$

                \*\*\*End critical section\*\*\*

                \*\*\*Synchronization barrier\*\*\*

$T^* = T_{Global}^*$

    Update the pheromone matrix  $\tau$



visited using a stochastic mechanism. An ant  $k$  at city  $i$  has not visited set of cities  $S_p$  then  $P_{ij}^k$  be the probability to visit edge  $k$  after edge  $i$ .

$$P_{ij}^k = \begin{cases} \frac{\tau_{ij}^\alpha \eta_{ij}^\beta}{\sum_{j \in S_p} \tau_{ij}^\alpha \eta_{ij}^\beta} & \text{if } j \in S_p \\ 0 & \text{otherwise} \end{cases} \quad (1)$$

$S_p$  represents the set of cities which has not been visited yet and to be visited again so that the probability of the ant visiting a city which has already visited becomes 0. Where  $\tau_{ij}$  is the pheromone content on the edge joining node  $i$  to  $j$ .  $\eta_{ij}$  represents the heuristic value which is inverse of the distance between the city  $i$  to  $j$ , which is given by:



$$\eta_{ij} = \frac{1}{d_{ij}}$$

Where  $d_{ij}$  is the distance between the city  $i$  to  $j$ .  $\alpha$  and  $\beta$  represents the dependency of probability on the pheromone content or the heuristic value respectively. Increasing the value of  $\alpha$  and  $\beta$  may vary the convergence of ACO.



After solution construction we have to update the pheromone accordingly, as follows:

$$\tau_{ij} \leftarrow (1 - \rho) \cdot \tau_{ij} + \sum_{k=1}^m \Delta \tau_{ij}^k$$

Where  $\rho$  is the evaporation rate,  $m$  is the number of ants, and  $\Delta \tau_{ij}^k$  is the quantity of pheromone laid on edge(i,j) by an ant  $k$ :

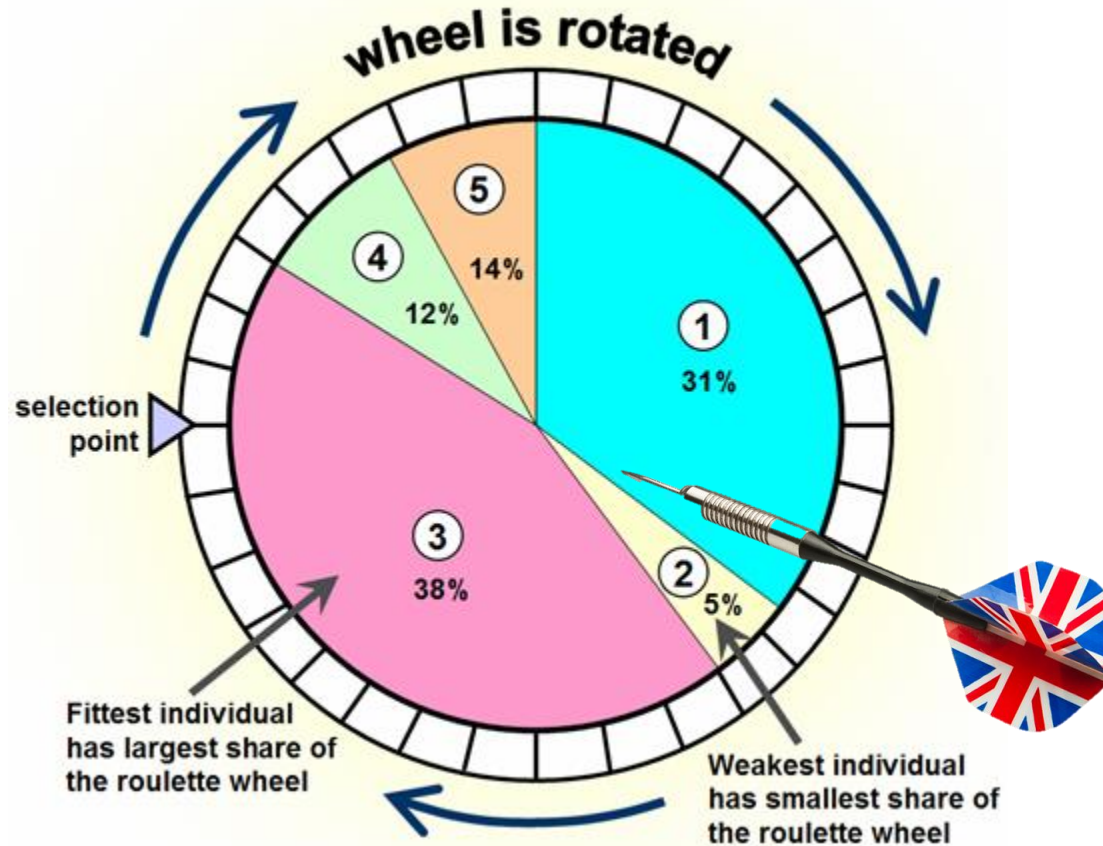
$$\Delta \tau_{ij}^k = \begin{cases} Q/L_k & \text{if ant } k \text{ uses edge } (i,j) \text{ in its tour,} \\ 0 & \text{otherwise} \end{cases}$$

Where  $Q$  is a constant and  $L_k$  is the length of the tour constructed by an ant  $k$ .





# Roulette Wheel Selection





# Printing the best tour

```
struct {  
    int cost;  
    int rank;  
} loc_data, global_data;  
  
loc_data.cost = Tour_cost(loc_best_tour);  
loc_data.rank = my_rank;  
  
MPI_Allreduce(&loc_data, &global_data, 1, MPI_2INT, MPI_MINLOC, comm);  
if (global_data.rank == 0) return;  /* 0 already has the best tour */  
if (my_rank == 0)  
    Receive best tour from process global_data.rank;  
else if (my_rank == global_data.rank)  
    Send best tour to process 0;
```



# Homework 6

- 使用MPI+OpenMP實作，每一台電腦各啟動一個process，每個process再fork出multi-thread

