Generating Random Variables

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Generating Random Variables

- Applications where random variables need to be generated
 - Simulations
 - Lotteries
 - Computer Games
- General strategy for generating an arbitrary random variable
 - Generate uniform random variables in the unit interval
 - Transform the uniform random variables to obtain the desired random variables

Generating Uniform Random Variables

X ~ U[a, b] has density function

$$f_X(x) = \begin{cases} \frac{1}{b-a} & \text{for } a \le x \le b \\ 0 & \text{otherwise} \end{cases}$$

The distribution function is

$$F_X(x) = \begin{cases} 0 & x < a \\ \frac{x-a}{b-a} & a \le x \le b \\ 1 & x > b \end{cases}$$

Y ~ U[0, 1] has distribution function

$$F_{Y}(x) = \begin{cases} 0 & x < 0 \\ x & 0 \le x \le 1 \\ 1 & x > 1 \end{cases}$$

- Given Y, can we generate X?
- (b-a)Y + a has the same distribution as $\mathcal{U}[a,b]$

Generating $\mathcal{U}[0, 1]$

- Computers can represent reals upto a finite precision
- Generate a random integer X from 0 to some positive integer m
- Generate the uniform random variable in [0, 1] as

$$U=\frac{X}{m}$$

The linear congruential method for generating integers from 0 to m

$$X_{n+1} = (aX_n + c) \bmod m, \quad n \ge 0$$

where m, a, c are integers called the modulus, multiplier and increment respectively. X_0 is called the starting value.

• For m = 10 and $X_0 = a = c = 7$, the sequence generated is

$$7, 6, 9, 0, 7, 6, 9, 0, \cdots$$

The linear congruential method is eventually periodic

Maximal Period Linear Congruential Generators

$$X_{n+1}=(aX_n+c) \bmod m, \quad n\geq 0$$

Theorem

The linear congruential sequence has period m if and only if

- c is relatively prime to m
- b = a 1 is a multiple of p, for every prime p dividing m
- b is a multiple of 4, if m is a multiple of 4.

Remarks

- Having maximal period is not a guarantee of randomness
- For a = c = 1, we have $X_{n+1} = (X_n + 1) \mod m$
- Additional tests are needed (see reference on last slide)

Generating a Bernoulli Random Variable

The probability mass function is given by

$$P[X = x] = \begin{cases} p & \text{if } x = 1\\ 1 - p & \text{if } x = 0 \end{cases}$$

where $0 \le p \le 1$

- Generate a uniform random variable U ~ U[0, 1]
- Generate the Bernoulli random variable by the following rule

$$X = \begin{cases} 1 & \text{if } U \le p \\ 0 & \text{if } U > p \end{cases}$$

How can we generate a binomial random variable?

The Inverse Transform Method

- Suppose we want to generate a random variable with distribution function F. Assume F is one-to-one.
- Generate a uniform random variable $U \sim \mathcal{U}[0, 1]$
- $X = F^{-1}(U)$ has the distribution function F

$$P(X \le x) = P(F^{-1}(U) \le x) = P(U \le F(x)) = F(x)$$

Example (Generating Exponential RVs)

X is an exponential RV with parameter $\lambda > 0$ if it has distribution function

$$F(x) = 1 - e^{-\lambda x}, \quad x \ge 0$$

How can it be generated?

Generating Discrete Random Variables

- Suppose we want to generate a discrete random variable X with distribution function F. F is usually not one-to-one.
- Let $x_1 \le x_2 \le x_3 \le \cdots$ be the values taken by X
- Generate a uniform random variable $U \sim \mathcal{U}[0, 1]$
- Generate X according to the rule

$$X = \begin{cases} x_1 & \text{if } 0 \le U \le F(x_1) \\ x_k & \text{if } F(x_{k-1}) < U \le F(x_k) \text{ for } k \ge 2 \end{cases}$$

Example (Generating Binomial RVs)

The probability mass function of a Binomial RV X with parameters n and p is

$$P[X = k] = \binom{n}{k} p^k (1 - p)^{n-k} \quad \text{if } 0 \le k \le n$$

How can it be generated?

Box-Muller Method for Generating Gaussian RVs

- 1. Generate two independent uniform RVs U_1 and U_2 between 0 and 1
- 2. Let $V_1 = 2U_1 1$ and $V_2 = 2U_2 1$
- 3. Let $S = V_1^2 + V_2^2$.
- 4. If $S \ge 1$, go to Step 1
- 5. If S < 1, let

$$X_1 = V_1 \sqrt{\frac{-2 \ln S}{S}}, \ X_2 = V_2 \sqrt{\frac{-2 \ln S}{S}}$$

6. X_1 and X_2 are independent standard Gaussian random variables

Proof

- (V_1, V_2) represents a random point in the unit circle
- Let $V_1 = R \cos \Theta$ and $V_2 = R \sin \Theta$
- $\Theta \sim \mathcal{U}[0, 2\pi]$ and $R^2 = S \sim \mathcal{U}[0, 1]$. Θ and S are independent
- $X_1 = \sqrt{-2 \ln S} \cos \Theta$ and $X_2 = \sqrt{-2 \ln S} \sin \Theta$
- X_1, X_2 also are in polar coordinates with radius $R' = \sqrt{-2 \ln S}$ and angle Θ

Proof Continued

• The probability density function of R' is $f_R(r) = re^{-r^2/2}$

$$\Pr\left[R' \le r\right] = \Pr\left[\sqrt{-2\ln S} \le r\right] = \Pr\left[S \ge e^{-r^2/2}\right] = 1 - e^{-r^2/2}$$

• The joint probability distribution of X_1 and X_2 is given by

$$P(X_{1} \leq x_{1}, X_{2} \leq x_{2}) = \int_{\{(r,\theta) \mid r \cos \theta \leq x_{1}, r \sin \theta \leq x_{2}\}} \frac{1}{2\pi} e^{-\frac{r^{2}}{2}} r dr d\theta$$

$$= \frac{1}{2\pi} \int_{\{x \leq x_{1}, y \leq x_{2}\}} e^{-\frac{x^{2}+y^{2}}{2}} dx dy$$

$$= \frac{1}{\sqrt{2\pi}} \int_{-\infty}^{x_{1}} e^{-\frac{x^{2}}{2}} dx \cdot \frac{1}{\sqrt{2\pi}} \int_{-\infty}^{x_{2}} e^{-\frac{y^{2}}{2}} dy$$

 This proves that X₁ and X₂ are independent and have standard Gaussian distribution

Acceptance-Rejection Method

- Suppose we want to generate a random variable X having density f
- Suppose X is difficult to generate using the inversion method
- Suppose there is a random variable Y with density g which is easy to generate
- For some $c \in \mathbb{R}$, suppose f and g satisfy

$$\frac{f(y)}{cg(y)} \le 1$$
 for all y .

- Generate a uniform random variable U ~ U[0, 1]
- Generate the random variable Y
- If $U \le \frac{f(Y)}{cg(Y)}$, set X = Y. Otherwise, generate another pair (U, Y) and keep trying until the inequality is satisfied
- To show that the method is correct, we have to show that

$$P\left(Y \leq x \middle| U \leq \frac{f(Y)}{cg(Y)}\right) = F(x)$$

where
$$F(x) = \int_{-\infty}^{x} f(t) dt$$

Example of Acceptance-Rejection Method

 Suppose we want to generate a random variable X with probability density function

$$f(x) = 20x(1-x)^3, \quad 0 < x < 1$$

- We need a pdf g(x) such that $\frac{f(x)}{g(x)} \le c$ for some $c \in \mathbb{R}$
- Consider g(x) = 1 for 0 < x < 1

$$\frac{f(x)}{g(x)} = 20x(1-x)^3 \le 20 \cdot \frac{1}{4} \cdot \left(\frac{3}{4}\right)^3 = \frac{135}{64}$$

- Let $c = \frac{135}{64} \implies \frac{f(x)}{cg(x)} = \frac{256}{27}x(1-x)^3$
- X can now be generated as follows
 - 1. Generate $U \sim \mathcal{U}[0,1]$ and $Y \sim \mathcal{U}[0,1]$
 - 2. If $U < \frac{256}{37}Y(1-Y)^3$, set X = Y
 - 3. Otherwise, return to step 1

Reference

• Chapter 3, *The Art of Computer Programming,* Seminumerical Algorithms (Volume 2), Third Edition, Pearson Education, 1998.